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Public reporting burden for this collection o gathering and mantaring the data needed without of inducation, enduding suggest	I information is estimated to average 1 hour per response, in , and completing and reviewing the collection of information, lons for reducing this burden, to Washington Headquarters 2	cluding the time for reviewing instruction Send comment regarding this burden a envices. Directorate for information Op Acenvork Reduction Project (0704-018)	is, searching existing data sources, stimates or any other aspect of this erations and Reports, 1215 Jefferson)), Washington, DC 20503.
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4. TITLE AND SUBTITLE			DING NUMBERS
Course Development in	Parallel and Distributed Database	DA1	AH04-95-1-025
6. AUTHOR(S) Jamal R. Als	abbagh and Carolyn Winston		
7. PERFORMING ORGANIZATION	NAMES(S) AND ADDRESS(ES)	8. PER REP	ORMING ORGANIZATION
Crombling State Un	iversity		
Department of Matr Carver Hall, Room	and Computer Science		
Grambling, Louisia	na 71245		NOODING LUONITODING
9. SPONSORING / MONITORING	AGENCY NAME(S) AND ADDRESS(ES)	10. SP AG	ONSORING / MONITORING ENCY REPORT NUMBER
U.S. Army Research Office P.O. Box 12211 Research Triangle Park, N	e C 27709-2211	ARO	34157.27-MA-I
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COURSE DEVELOPMENT IN PARALLEL AND DISTRIBUTED **DATABASE MANAGEMENT SYSTEMS**

Jamal Alsabbagh and Carolyn Winston

ADSRC - Grambling State University

ASRC annual meeting, Orlando, Fl.

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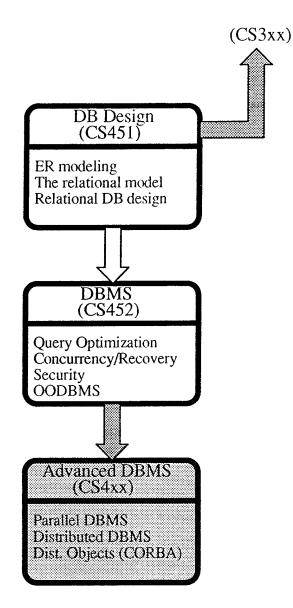
OUTLINE

- Context and Plan for This Work
- Overview of Parallel Database Systems

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CURRENT AND PROPOSED DATABASE COURSES (AT GSU)

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Classification of Parallel Database Systems

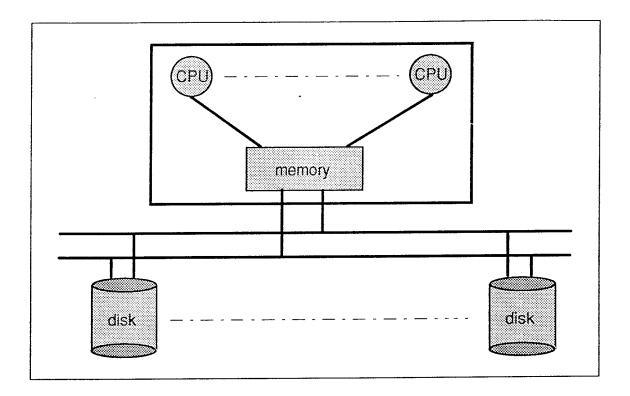
• Shared Memory (SM); also called Shared Everything (SE)

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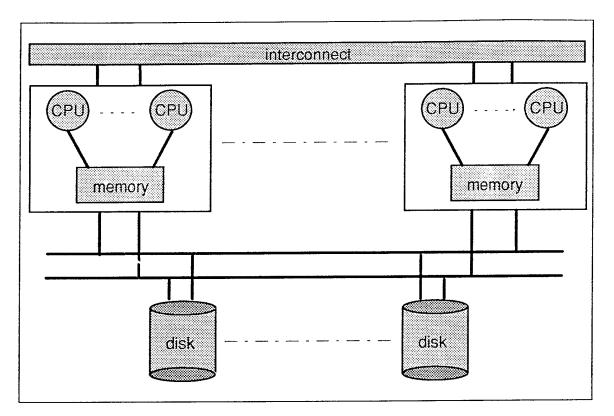
- Shared Disk (SD)
- Shared Nothing (SN)

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Shared-Memory (SM) Parallel Database Systems



- All processors directly access, in a symmetrical fashion, all main memory and disks.
- In general, the operating system allocates processors to processes.
- Processors have local caches to reduce network traffic, but loading/flushing caches can degrade performance.
- Hardware-specific solutions are required to ensure coherence among the caches (e. g. processors continuously snoop the shared bus to see if their cached data is required elsewhere.)
- Typical Hardware: IBM 3090, IBM 370, Bull DBS8, Encore, Sequent Symmetry

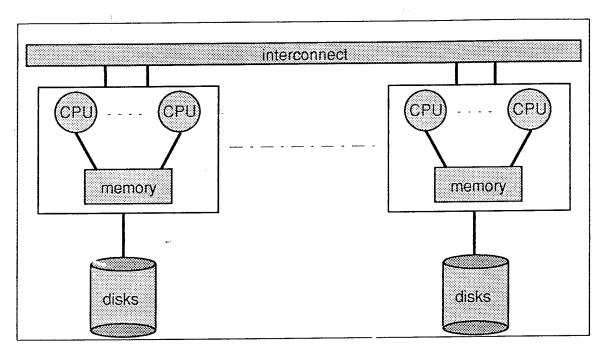


Shared-Disk (SD) Parallel Database Systems

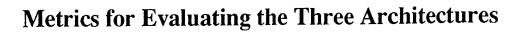
- Each node has its own memory and may itself be an SMP box.
- The nodes share the same disks logically (and may be physically too).
- System (or application) software must ensure the coherency among multiple copies of disk pages requested by different nodes.
- A query or update by a node requires it to:
 - 1. Transmit, to all other nodes, an intention to query/update the database.
 - 2. If the required page is currently being updated by any other node, then wait until it is released.
 - 3. Read or receive the required page.
 - 4. perform the query or update.
- Typical Hardware: DEC VM Cluster, SUN Sparc 1000 cluster

#4 - 1

Shared-Nothing (SN) Parallel Database Systems



- Each node has its own local memory and its own local disks.
- The database is partitioned across the nodes, thus allowing I/O parallelization
- Each node acts as a server for its local data.
- An update by a processor requires:
 - 1. Transmit a request for update to the relevant server.
 - 2. The server performs the update, locally.
 - 3. The server acknowledges the success back to the requester.
- A query by a processor requires:
 - 1. Transmit the query to the relevant server.
 - 2. The server performs the query locally.
 - 3. The server sends the query result to the requester.
- Typical Hardware: AT&T 3600, IBM SP2, nCUBE, VAXcluster



	Metric	Explanation	
1	Price	Using commodity hardware reduces system cost.	
2	Throughput	Inter-query parallelism improves throughput.	
3	Response Time	Intra-query parallelism improves response time.	
4	Speedup	Ideally, twice the hardware should solve the problem in half the time.	
5	Scaleup	Ideally, twice the hardware should solve twice the problem in the same time.	
6	Startup Cost	Preparing a query for parallel execution is an overhead.	
7	Interference	Processors slow each other when competing for shared resources.	
8	Load Balancing	Ideally, all the processors should be working concurrently.	
9	Comm. Overheads	Ideally, sub-problems of one problem should require least communication.	
10	Data Availability	It is desirable to be able to tolerate failure of some nodes.	
11	Portability	Porting centralized DBMS software should be relatively easy.	

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