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AN ADAPTIVE ARQ STRATEGY FOR
PACKET SWITCHING DATA COMMUNICATION
NETWORKS

By

Kurtis J. Guth

June 1989

Thesis Advisor:

Tri T. Ha

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An Adaptive ARQ Strategy for
Packet Switching Data Communication Networks

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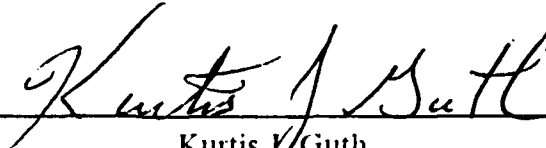
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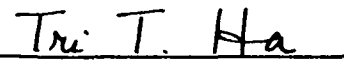
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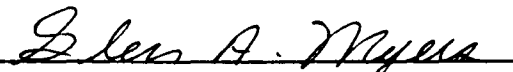
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ABSTRACT

Automatic Repeat reQuest (ARQ) techniques are often used by a packet switching data communication network to provide an error-free communications link between stations. The ARQ technique ensures consistent data quality under varying link conditions. Unfortunately, the information throughput is link dependent and as the noise or interference on the link increases, the throughput decreases. In an effort to improve the throughput on a short range, RF, packet switching data communications network, an adaptive ARQ strategy applied to Stop-and-Wait (SW) protocols was developed. To provide a system designer with flexibility, different adaptive strategies for different system and link conditions were developed. Examples of information transfer between two stations using the adaptive SW protocol are presented. A simulation to compute the throughput efficiency of several adaptive SW protocols was performed. A comparison of the throughput efficiencies of the simulated adaptive SW protocol with the non-adaptive SW protocol showed good gains could be achieved using the adaptive strategy when the networks are subject to high channel bit error rates.

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TABLE OF CONTENTS

I. INTRODUCTION	1
II. THE DATA LINK CONTROL (DLC) LAYER AND PROTOCOLS	2
A. DLC LAYER	2
B. DLC PROTOCOLS	4
III. NETWORK CONFIGURATIONS AND PROTOCOL STRUCTURES	7
A. CONFIGURATIONS AND DATA TRANSFER MODES	7
B. STRUCTURE	8
IV. ADAPTIVE ARQ STRATEGY DEVELOPMENT	14
A. GENERAL CONSIDERATIONS	14
1. Throughput Efficiency	14
2. Signal Fading and Noise	14
B. ADAPTIVE STRUCTURE	15
1. Strategy Structure	15
2. Two-Level Adaptive Strategy Structure	17
3. Multi-level Adaptive Strategy Structure	20
C. ADAPTIVE STRATEGY IMPLEMENTATION	22
D. ADAPTIVE STOP-AND-WAIT (SW) PROTOCOL	23
1. Adaptive SW Operations and Example Notation	23
2. An Example of One-Way Information Transfer	26
3. An Example of Two-Way Information Transfer	28
V. ADAPTIVE SW SIMULATION	32
A. INTRODUCTION	32
B. SIMULATION DEVELOPMENT	32
1. Model Design	32
2. Simulation Program	33
C. SIMULATION RESULTS	34

VI. CONCLUSIONS	37
APPENDIX A. SIMULATION INPUT FILES	38
APPENDIX B. SIMULATION FLOW DIAGRAMS	58
LIST OF REFERENCES	64
BIBLIOGRAPHY	65
INITIAL DISTRIBUTION LIST	66

LIST OF TABLES

Table 1. ADAPTIVE SW STRUCTURE REPRESENTATION AND STATE VECTOR	26
Table 2. ADAPTIVE SW STRUCTURE AND STATE VECTOR	29

LIST OF FIGURES

Figure 1. OSI network architecture and data communications network	3
Figure 2. Frame structure [From Ref. 2]	4
Figure 3. Supervisory and unnumbered frame format	9
Figure 4. Information frame format	9
Figure 5. Flag field	10
Figure 6. Control field format	11
Figure 7. Adaptive strategy structure and state vector	16
Figure 8. Alternative adaptive strategy structure design	18
Figure 9. Two-level adaptive strategy	19
Figure 10. Multi-level adaptive strategy	21
Figure 11. Verification procedure and packet length modification	24
Figure 12. Adaptive SW - an example of one-way information transfer	27
Figure 13. Adaptive SW - an example of two-way information transfer	30
Figure 14. Simulation throughput efficiency vs probability of bit error	36

I. INTRODUCTION

In this thesis an adaptive *Automatic Repeat reQuest* (ARQ) strategy for a packet switching network is developed. The adaptive technique is to be applied to a Stop-and-Wait (SW) protocol. The SW protocol was chosen because the adaptive techniques designed are useful for radio frequency (RF) systems transferring data over short distances. In the short range RF systems, the propagation delay is very short, and the inefficient use of bandwidth resulting from a transmitter waiting for a response is not a significant factor. Also, RF systems often operate in a high noise environment and are affected by signal fading or even jamming. The throughput of all ARQ systems is very low under high noise conditions. These high noise conditions will often result in a large number of retransmissions which increases the delays for the data transfer and further justifies the use of adaptive SW type protocols.

The adaptive strategies developed in this thesis all vary the length of the message packets by reducing or increasing the information contained in the packet. These increases or decreases in the packet length are made based on the number of consecutive successful transmissions or consecutive retransmissions of a packet. By keeping track of the consecutive packet transmissions and using this information to base the decision of when to adapt the packet length, no additional hardware should be needed to upgrade an SW protocol to an adaptive SW protocol.

Chapters 2 and 3 discuss the basic requirements to be met by an ARQ protocol and the structure used to meet these requirements. Chapter 4 develops the adaptive strategies and provides examples of their implementation. Chapter 5 describes some simulations of several adaptive SW strategies and computes the throughput efficiencies for each type.

II. THE DATA LINK CONTROL (DLC) LAYER AND PROTOCOLS

A. DLC LAYER

The International Standards Organization (ISO) has developed the Open System Interconnection (OSI) network architecture composed of a hierarchy of seven layers or modules [Ref. 1,2]. Each layer performs certain functions and provides certain services to the layers above. Each function performed by a data communications terminal or a station on a data communications network is typically assigned to a particular layer. The actual operation of transferring data or messages by a station is the responsibility of the first three layers. Figure 1 shows the OSI network architecture and a block diagram representing a communications network between n stations.

The third layer, the network layer, and the second layer, the data link control (DLC) layer, provide the upper layers with a virtual link for messages with the other stations on the communications network. This means the layers above the network and DLC layer only provide a message to be transferred to another station on the network, and expect the network and DLC layers to perform the necessary functions to ensure its proper transfer. The network layer works closely with the DLC layer to break the message into packets for transmission. The DLC layer is responsible for providing the upper layers with an "error-free" communications link. The DLC layer uses an unreliable bit pipe provided by the first layer, the physical layer, to transmit the actual bits through the communications link. The physical layer contains the transmission equipment used to provide the interface between stations on the data communications network. This layer is only concerned with transferring bits from station to station and does not perform any bit error detection to determine the accuracy of the bit transmission. Therefore, the reception of a correct bit is "unreliable" and must be checked by the DLC layer before accepted. The DLC layer uses a specified set of procedures to transmit the information and correct the bit errors that occur on the unreliable bit pipe.

To establish the "error-free" communications link, the network and DLC layers break the message into packets and then add some overhead control bits to each packet. As shown in Figure 2, a packet with the overhead control bits added is called a frame. The DLC layer is responsible for establishing a bit sequence to identify the frame starting and ending points. The bit sequence must be unique and no inadvertent occurrences of the sequences can be transmitted prior to sending the bit sequence marking the end

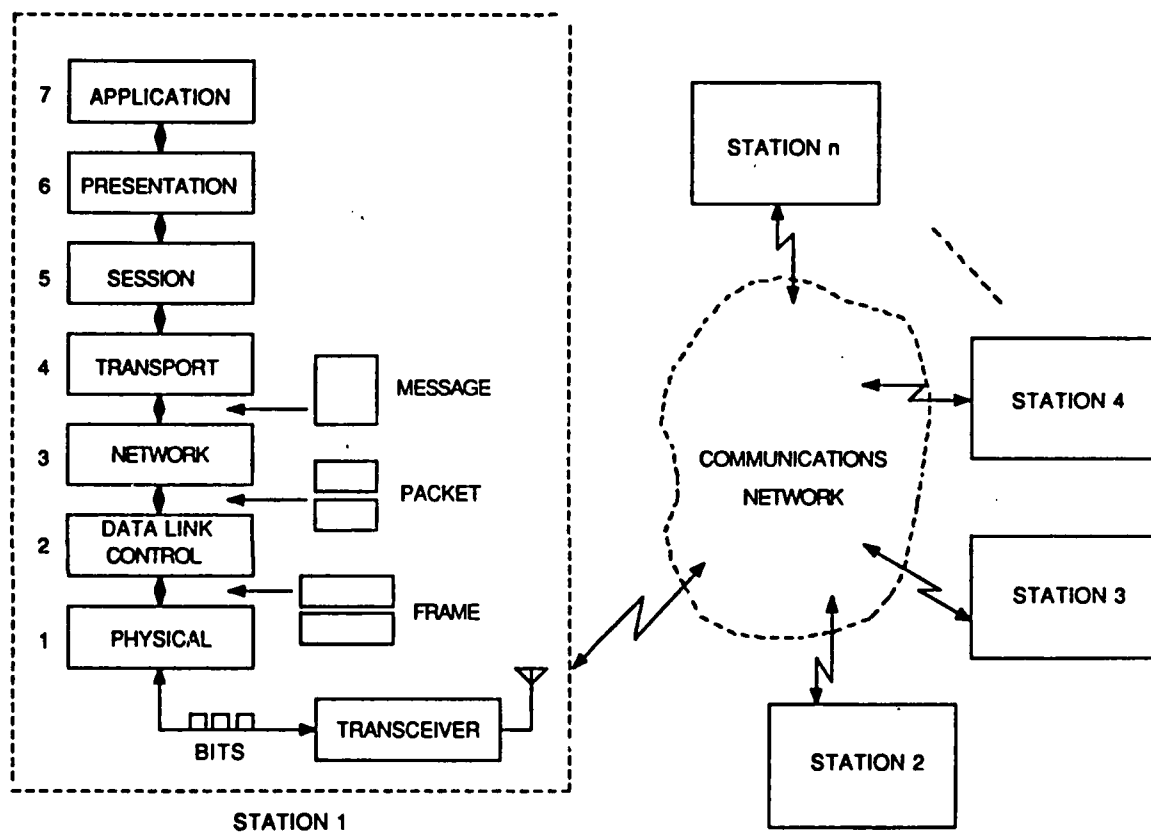


Figure 1. OSI network architecture and data communications network

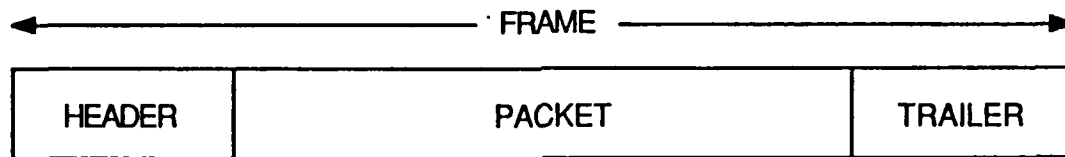


Figure 2. Frame structure [From Ref. 2]

of the frame. The DLC layer must account for the frames that have been transmitted and the acknowledgements identifying receipt from the receiving station. Additionally, the DLC layer is responsible for ensuring the frames are received in the proper order so the message can be reassembled and passed in the correct form to the layer above. Frames which acquire errors during transmission over the unreliable bit pipe must be corrected by the DLC. Finally, the DLC layer must be able to establish and control the communications link between the stations.

B. DLC PROTOCOLS

A protocol is a set of specific procedures for data transfer that the DLC layers use to obtain error free communications when utilizing the unreliable bit pipe. The protocols or sets of procedures must be designed to allow the DLC layers to identify and correct all possible transmission errors. One technique used to correct these transmission errors is called automatic repeat request (ARQ). This technique draws on the ability of the DLC to detect frame errors which occur during transmission and request the sender to retransmit the frame. To implement these ARQ protocols, the DLC establishes a specific format for the frames and transfer procedures to handle the data transfer.

One type of ARQ protocol is called the Stop-and-Wait (SW) protocol. In the SW protocol, a single frame is transmitted and the transmitter stops and waits for a response from the receiving station. If the frame is received without errors, then an acknowledgement (ACK) will be returned from the receiver and the next frame is transmitted. If an error is encountered, then a request for retransmission (NAK) will be returned and the same frame is retransmitted. If no response is received, the sender only waits a cer-

tain amount of time referred to as a *time-out* and then the frame is retransmitted. The SW protocol is effective for data communications when the round trip propagation delay time for a packet is short relative to the frame transmission time. This protocol becomes inefficient in using the available bandwidth when the round trip propagation time is long, such as in satellite communications, or when stopping the transmission introduces any additional delays associated with resynchronization of the transmission equipment.

To make better use of the bandwidth on these channels with long propagation delays, continuous or "pipelining" protocols are used. A pipelining type protocol allows the transmitter to send frames continuously without stopping and waiting for a response for each frame. There is a limit on how many frames can be sent without having received a response from the receiver for the first of the unacknowledged frames transmitted. This limit is referred to as a "sliding window" and is determined by the system capabilities and the propagation delay. For the most efficient operation, the window should be greater when the propagation delay is longer. [Ref. 2]

Go-Back-N (GBN) is a pipelining protocol which allows a transmitter to send up to N-1 frames before stopping and waiting for a response from the receiver. The receiver, in contrast to the transmitter, has a receive window equal to one frame. The receiver is looking to receive a specific frame and once that frame is received, it immediately looks for the next one. Typically, the receiver discards a frame received with errors. When a later frame arrives correctly and is out of sequence, or not the frame the receiver was expecting, the transmitter is notified that an error has occurred. When the transmitter learns that one of the frames was received with errors, it backs up and begins to transmit again starting with that frame. The beginning of the window *slides* up to the frame being retransmitted, so this frame is the first of N-1 possible frames in the new window. All the frames that were transmitted following the frame received in error are also retransmitted. If the protocol is operating in an environment which causes several consecutive packets to incur errors during transmission, a large number of frame retransmissions will result. Under these conditions, the transmission delays can be quite large and the protocol efficiency quite low.

Selective Repeat (SR) is another pipelining protocol. Similar to GBN, SR protocol uses a sliding window to allow N-1 frames to be sent without waiting for a response. However, SR protocol uses a different method of error correction than GBN. For the SR protocol, only frames received in error are retransmitted. Since the DLC must still provide the message in the proper form to the layer above, frames received without errors must be stored in a buffer while waiting for the retransmitted frames. This selective

repeat feature adds complexity to the system and additional memory requirements which may be undesirable for some systems.

A short range, radio frequency (RF) network will often operate in a noisy environment with probability of bit error P_b rates in the range of 10^{-2} to 10^{-5} . The short distance for this RF network means the propagation delay is very small. When the frame processing time can also be kept small relative to the frame transmission time, the average delay in transferring a packet using a SW protocol is almost the same as using a GBN protocol. Considering the similar delay time and the fact that SW is a simple, easy to implement protocol with low storage requirements, then the SW protocol is often preferable for this type of RF network. [Ref. 3]

III. NETWORK CONFIGURATIONS AND PROTOCOL STRUCTURES

A. CONFIGURATIONS AND DATA TRANSFER MODES

Networks are configured in many different ways and require flexible data transfer modes and procedures that can handle a varied range of applications. Networks are generally configured as multi-point or point-to-point systems and operate in a full-duplex or half-duplex mode. Depending on the configuration and type of stations operating on the network, different types of data transfer modes may be used. For the multi-point configuration where only one station may be considered the primary station communicating with outlying secondary stations, the normal response mode (NRM) is often the best suited data transfer mode. In the NRM, a primary station controls the data transfer link and informs each secondary station on the network when it is allowed to transfer data. This mode is particularly useful in a polled operation where each secondary station has only the limited capability of responding to a poll and then transmitting or receiving data when directed. The primary station directs all the interaction between itself and the outlying stations, relays all the transmissions between stations, and performs all error recovery procedures on the link. This mode does have the disadvantage of high overhead associated with the polling operation. [Ref. 2]

To take advantage of secondary stations with slightly more capability, the asynchronous response mode (ARM) can be used to limit the polling overhead required. In the ARM, the secondary stations are capable of initiating responses without the strict direction from the primary station as in the NRM. In the ARM, there is still a primary station controlling the data link between itself and the remote secondary stations. The primary station is typically still superior in capability to the secondary stations and when communication with a secondary is desired, the secondary is activated. Once a secondary is activated, the data transfer between the primary and the secondary flows freely back and forth without any polling or direction required. The advantage of the ARM is the ability to communicate in an asynchronous manner between the primary and the activated secondary stations in the multi-point configuration, without the polling operation. [Ref. 4]

For point-to-point communications when stations have comparable capabilities to initiate data transfer and control the communication link, the asynchronous balanced mode (ABM) is a very effective mode of data transfer. In the ABM, the stations act as

both the primary and the secondary stations during the data transfer. Therefore, stations operating in the ABM are called combined stations [Ref. 4].

In a radio frequency (RF) network, it is often desirable for stations to communicate over a full-duplex, point-to-point communications link. Both stations will often be of comparable capabilities and will operate as combined stations, so the data transfer is conducted in the ABM. For short range RF networks, the ABM is most efficient. This is the network of interest in this research.

B. STRUCTURE

A protocol establishes a specific frame structure and transfer procedures for the proper error-free transmission of information and control of the communications link. The protocol uses three different types of frames to perform the data transfer and the communications link control:

- Information frame (I frame).
- Supervisory frame (S frame).
- Unnumbered frame (U frame).

Information or I frames are used to transfer the data and can be of variable length depending on the length of the message packet. An I frame can also be used to return the appropriate acknowledgment for a frame received from the other station. Supervisory or S frames are used to control the data flow on the communications link and to provide message acknowledgments or rejections. They also can be used to recover from conditions causing errors. Unnumbered or U frames provide the necessary commands and responses to establish and terminate the communications session and to specify the various parameters and restrictions to be used during the session. Depending on the type of frame, two different formats can be used. The first format type, used for the S or U frames when no information is to be transmitted, is shown in Figure 3. The second format type, used for I frames when information is to be transmitted, is shown in Figure 4.

The starting and ending bits of both frame formats are delineated by the flag field. The flag field shown in Figure 5 is a unique 8-bit sequence from which the receiver can acquire frame synchronization. The bit sequence is a "0" bit followed by six "1" bits and another "0" bit. To ensure that the flag sequence is unique and not reproduced anywhere else in the frame, the transmitter uses a "bit stuffing" technique. In bit stuffing, the transmitter inserts a "0" bit after five consecutive "1" bits in all fields after the be-

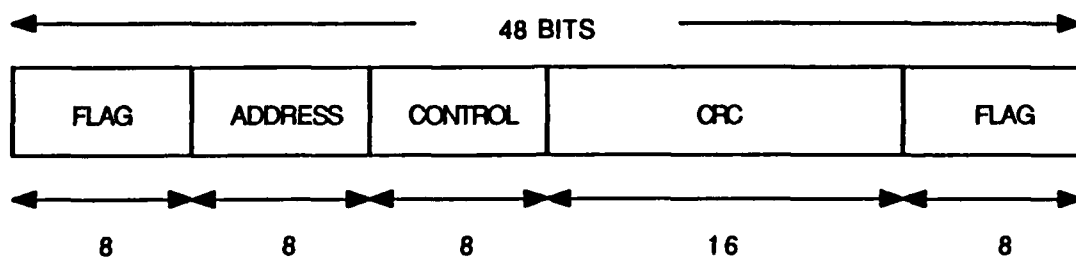


Figure 3. Supervisory and unnumbered frame format

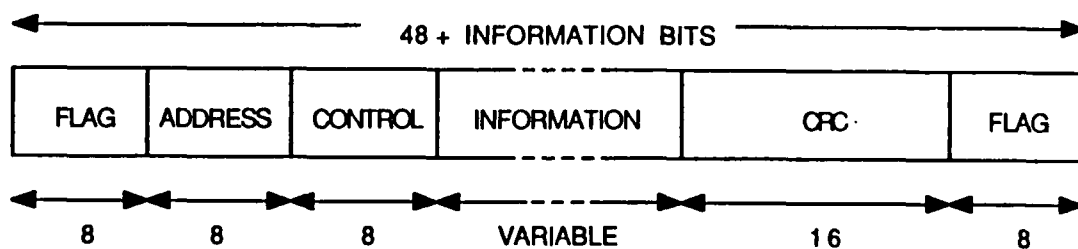


Figure 4. Information frame format

gining flag and before the ending flag. The receiver, to "destuff" the frame, simply removes the "0" bit after five consecutive "1" bits for all bits received between the flag sequences. [Ref. 1]

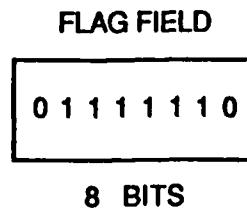


Figure 5. Flag field

The address field is typically eight bits or a byte in length and used to identify the transmitter or the receiver of the frame. For conditions where more than 256 addresses are required, the address field can be extended by any number of additional bytes. By convention, when using the extended addressing, all bytes preceding the final byte of the address field have the first bit set to "0" and the first bit of the final byte is set to "1" [Ref. 4].

In systems with multi-point configurations, the address field is used to identify the station receiving the frame. In point-to-point configurations, where the stations are combined stations, the address field is used to determine if the frame being transmitted contains a command or a response. All frames contain either a command or a response and are one of the three types of frames: I, S, or U frames. When a local station transmits a frame with the remote station specified in the address, that frame contains a command. When a station uses its own address, then the frame contains a response.

The control field is another eight bit field and this field actually identifies the type and function of the frame. The first bit of the control field is used to distinguish between an I frame and an S or U frame. The control field format is specified in Figure 6. If a frame is determined not to be an I frame, the second bit of the control field is used to distinguish between an S and a U frame.

An I frame is identified by a "0" in the first bit position of the control field. The next three bit positions of the I frame control field compose the send sequence number N(s) which identifies the frame. Only an I frame contains a send sequence number. The Poll/Final (P/F) is the fifth bit position of the control field in all the frame types and used to generate a checkpointing system. In the checkpointing system, the P/F bit is P

CONTROL FIELD FOR	CONTROL FIELD BITS							
	1	2	3	4	5	6	7	8
INFORMATION TRANSFER FRAME	0	N(s)			P/F	N(r)		
SUPERVISORY FRAME	1	0	TYPE		P/F	N(r)		
UNNUMBERED FRAME	1	1	TYPE		P/F	TYPE		

Figure 6. Control field format

for frames containing commands and F for the related response frame. When a station sends a frame containing a command and has the P bit set to "1", the remote station can only send a frame containing a response and must set the F bit to "1". The final three bit positions of an I frame and an S frame compose a receive sequence number N(r). By convention, the N(r) is the sequence number of the next expected frame and acknowledges receipt of all the previous frames with sequence numbers less than or equal to N(r)-1 [Ref. 2].

An S frame is identified by a "1" and a "0" in the first and second bit positions of the control field. The third and fourth bits of the control field encode the four *types* of S frames used for data flow control:

- Receive Ready (RR) indicates that the station is ready to receive I frames and is used to acknowledge I frames received when the station does not have information to transfer.
- Receive Not Ready (RNR) indicates that the station is unable to receive I frames at this time and acknowledges the I frames that have been received.
- REject (REJ) is used to indicate that a transmission error has occurred and that retransmission is to begin with the frame indicated in the N(r) field.

- Selective REject (SREJ) indicates a frame has been received in error and retransmission of this frame is requested.

The use of RR and RNR are required for a system to provide proper data flow control. REJ and SREJ are typically optional and best suited for pipelining protocols to enhance error recovery. The final three bit positions contain the receive sequence number $N(r)$ which enable a station to use an S frame to acknowledge receipt of $N(r)-1$ frames when it has no data to transfer.

A U frame has the first bit of the control field set to "1" and is distinguished from the S frame by setting the second bit to "1". The next two and the last three bit positions of the control field are used to encode the commands and responses required for data link control. To initiate a communications session, a station will use a U frame SET Mode (SETM) command. The SETM command is used to represent the set of mode setting commands which specifies the mode of communications and the modulus used for the sequence numbering [Ref. 2]. The U frame response Disconnected Mode (DM) to the SETM command is used when a station is temporarily unable to participate in a communications session. The response to the SETM command when a station is ready for a communications session is the Unnumbered Acknowledgment (UA) response. The UA response is used also for acknowledging other U frame commands. To terminate a communications session, the U frame DISConnect (DISC) command is used. To recover from a frame error condition which can not be solved by retransmitting the frame, the FRaMe Reject (FRMR) response is used. A more detailed description of the types of U frames available can be found in Ref. 4.

The information field only exists in an I frame. The information field contains the information or the message packet to be transferred across the link and is typically of variable length. Since the information is composed of bits occurring in any order with no particular pattern, the bit stuffing technique must be used to ensure a bit sequence resembling a flag sequence is not inadvertently transmitted in the information field.

The Cyclic Redundancy Check (CRC) field of the frame is a 16-bit checksum used to detect bit errors between the flags of the frame. To determine the bit sequence in the CRC field, the transmitting station performs a long division on the frame's data bits (excluding the flag field) by a generator polynomial. The remainder resulting from this long division then becomes the bit sequence for the CRC field. The generator polynomials used to create this checksum are typically one of two types of polynomials:

- CRC-CCITT polynomial $x^{16} + x^{12} + x^5 + 1$.

- CRC-16 polynomial $x^{16} + x^{15} + x^2 + 1$.

The receiver, to determine if a frame has suffered any transmission errors, performs a similar long division using the same generator polynomial as the transmitter on all the fields of the frame between the flags (including the CRC field). If a frame is received without any transmission errors, the remainder of the long division will equal zero and the frame is accepted. If a non-zero remainder results from the long division, then a transmission error has occurred and the frame is not accepted. If a zero remainder is undesirable to satisfy the CRC, other methods of calculating the checksum can be used which results in the receiver obtaining a remainder when performing the division on the received frame. The receiver then compares this remainder to the bit sequence in the CRC field to determine if a transmission error has occurred. If the remainder and the bit sequence in the CRC field match, then the frame is accepted. If they do not match, then the frame has errors and is rejected. [Ref. 2]

IV. ADAPTIVE ARQ STRATEGY DEVELOPMENT

A. GENERAL CONSIDERATIONS

1. Throughput Efficiency

The effectiveness of information transfer of the typical ARQ protocol largely depends on the bit error rate of the channel. For a low and generally stable channel bit error rate, the optimum packet size for the maximum throughput efficiency can be determined. Throughput efficiency is defined as the amount of actual information transferred relative to the total number of bits transmitted. Since every packet transmitted has the same amount of overhead required by the ARQ protocol, the larger the information field, the better the packet efficiency. Unfortunately, as the packet size increases the probability of an error occurring during transmission increases correspondingly. Generally, the maximum throughput efficiency is realized by utilizing the optimum packet length for the link conditions and minimizing the number of retransmissions required. For long term changes in the channel bit error rate, the optimum packet lengths can be modified manually to maximize the throughput efficiency. If the channel bit error rate increases quickly, the typical ARQ strategy will continue to retransmit packets which were received in error, at the packet length determined to be optimum for the average channel bit error rate. For these typical ARQ strategies the throughput efficiency suffers greatly during the periods of operation with high bit error rate. [Ref. 5]

2. Signal Fading and Noise

The channel bit error rate does not often remain constant for a radio frequency (RF) system, particularly for a system affected by signal fading or subject to jamming. Destructive interference of the RF signal as a result of multi-path interference or jamming may result in an abrupt, large increase in the channel bit error rate. This condition may only last for several packet transmissions followed by an equally abrupt decrease in the bit error rate. A less destructive degradation of the channel, such as an increase in noise level, may result in a less abrupt and severe increase in the bit error rate. This noisy degradation of the channel may last longer than the quick destructive interference causing an increase of the bit error rate to persist for a large number of packet transmissions. Since the bit error rate does not remain constant over the channel and actually may vary widely, an ARQ strategy which adapts the packet length for a changing bit error rate would increase the throughput efficiency over one that does not.

Different types of signal fading conditions and channel noise dictate that different adaptive strategies be employed. Since many different channel conditions exist, developing a different strategy for each condition is nearly impossible and certainly impractical. Therefore, two strategies with some optional design variations were developed in an attempt to counteract the effects of signal fading and increased channel noise conditions.

B. ADAPTIVE STRUCTURE

1. Strategy Structure

An adaptive ARQ strategy decreases the length of the packets by decreasing the amount of data transferred in the information field. This reduction of data occurs for packets which have already failed a certain number of retransmission attempts indicating an increase in the bit error rate on the forward channel. Also, the lengths of packets not yet transmitted for the first time may be increased (up to a maximum size allowed), if the previous packets are transmitted without error, indicating that the forward channel bit error rate has decreased.

The structure for the adaptive strategies can be viewed as consisting of various levels where each level equates to a packet length as shown in Figure 7. The top level of the structure is associated with the largest information field and is typically the starting point for a communications session. Each level below the top level has an information field length associated with it which is less than the largest information field.

Each successive level is connected by a series of steps. The number of steps between different levels do not have to be the same. Once the structure is established, the number of steps are fixed for that structure. These steps connecting the levels determines how many retransmission attempts at a certain length are made prior to the modification of the packet length to the new length associated with the new level. A step up to the next step or level of the structure is made for each successful packet transmission until the top level is reached. For each packet which suffers an error in transmission, a step down is made until the bottom level is reached.

In Figure 7, the set of steps between two levels have a packet length equal to the length associated with the upper level. In this figure, to go from level $L(0)$ to level $L(1)$ requires four steps. This means that four consecutive attempts are made to transfer a packet with length equal to the maximum length associated with $L(0)$. If four consecutive attempts fail to transfer the packet error-free, then level $L(1)$ is reached and the packet length is modified prior to the next retransmission attempt.

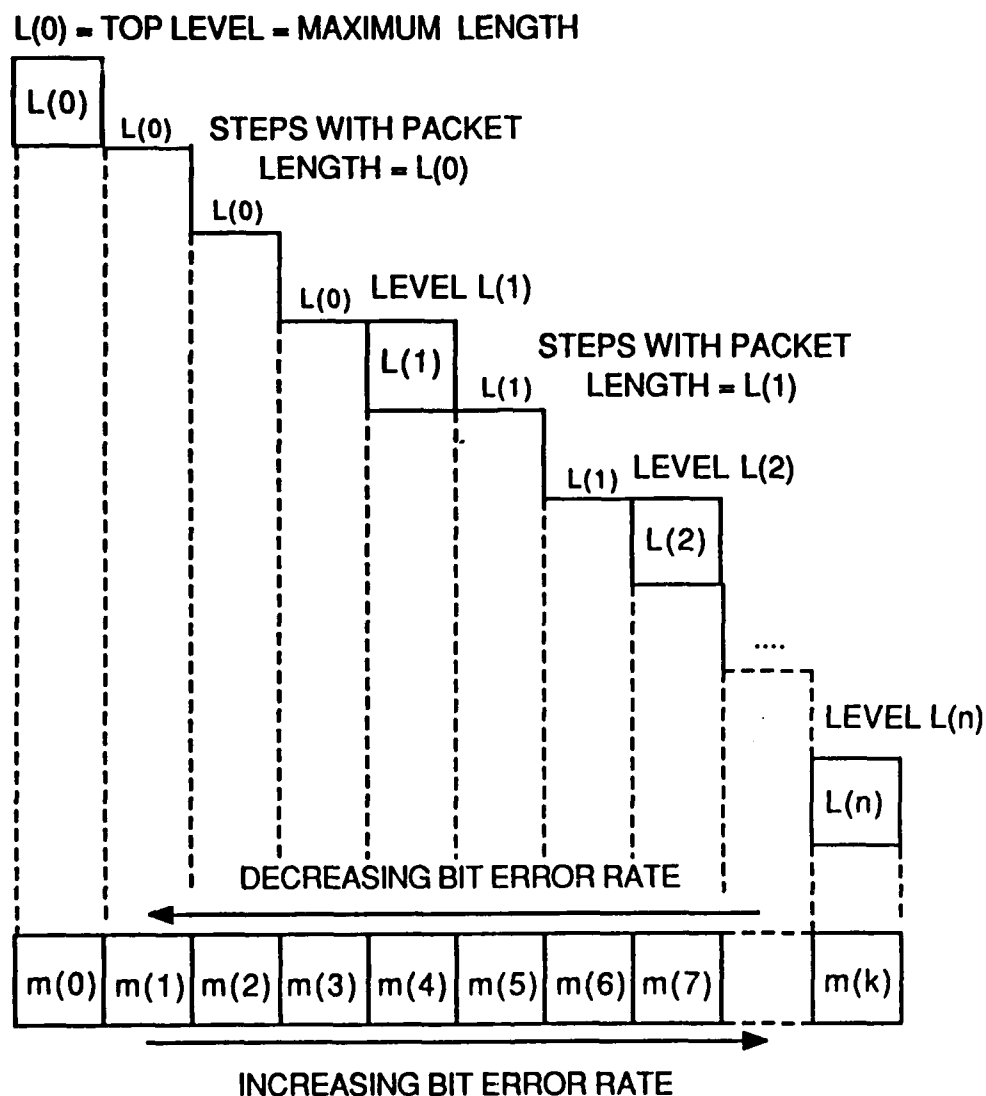


Figure 7. Adaptive strategy structure and state vector

The adaptive structure can be represented by a state vector m . Each of the elements or the states of the state vector $m(i)$, $i = 0, 1, \dots, k$, correspond to either one of the levels or the steps of the adaptive structure. The number of states ($k + 1$) in the state vector is equal to the number of levels and steps in the adaptive structure. Moving left or right to a new state in the vector corresponds to stepping up or down the adaptive structure respectively. The value of each state $m(i)$ is the packet length associated with the corresponding level or step of the structure. Since some levels and steps have the same associated packet length, the values of some states will be the same. For example,

as shown in Figure 7, when the current state is $m(3)$, then the packet is transmitted at the length associated with level $L(0)$. If the packet is received in error, then the state is updated to $m(4)$ and the length is modified to that associated with $L(1)$. [Ref. 6]

This state vector will be used by the station software to implement the adaptive strategy. The packet length to be transmitted by the station is determined by the value of the current state of the vector. The ends of the state vector correspond to the levels associated with the maximum and minimum packet lengths of the structure. When the state at the left-hand end of the vector is reached, this corresponds to reaching the top of the adaptive structure and the system state will remain at this vector element until a frame suffers an error during transmission. Likewise, when the right-hand end of the vector is reached, the state corresponds to the bottom of the structure and the frame is retransmitted at the minimum length associated with this state until it is successfully transferred.

An example of one alternative method of designing the steps between levels is shown in Figure 8. In this structure, a different set of steps or path is provided when ascending from one level to another. This strategy may be useful in the situations where additional packet transmissions at the length associated with the lower level are desired prior to modifying the packet lengths to the longer length of the level above. The steps ascending to the next level can be variable in number and may intersect any of the descending steps or go all the way up to the next level. The state vector m for this structure is made up of different segments. The states $m(i)$ in each segment of the state vector correspond to one level and the steps which ascend and descend from that level, which all have the same associated packet length. As in the previous vector, moving left or right corresponds to stepping up and down the structure. For this vector, the arrows indicate which is the next state when moving from one segment to another. This method of transition between levels provides the system designer with added flexibility to optimize the system throughput efficiency.

2. Two-Level Adaptive Strategy Structure

When the signal is subjected to fading or jamming on the forward channel, there is a correspondingly large increase in the bit error rate which occurs quickly relative to the time required for packet transmission. Under these conditions, the packet length must be modified substantially to gain an increase in the information throughput. The type of adaptive strategy to counteract the conditions imposed by this large changing bit error rate is a *two-level* adaptive strategy, shown in Figure 9.

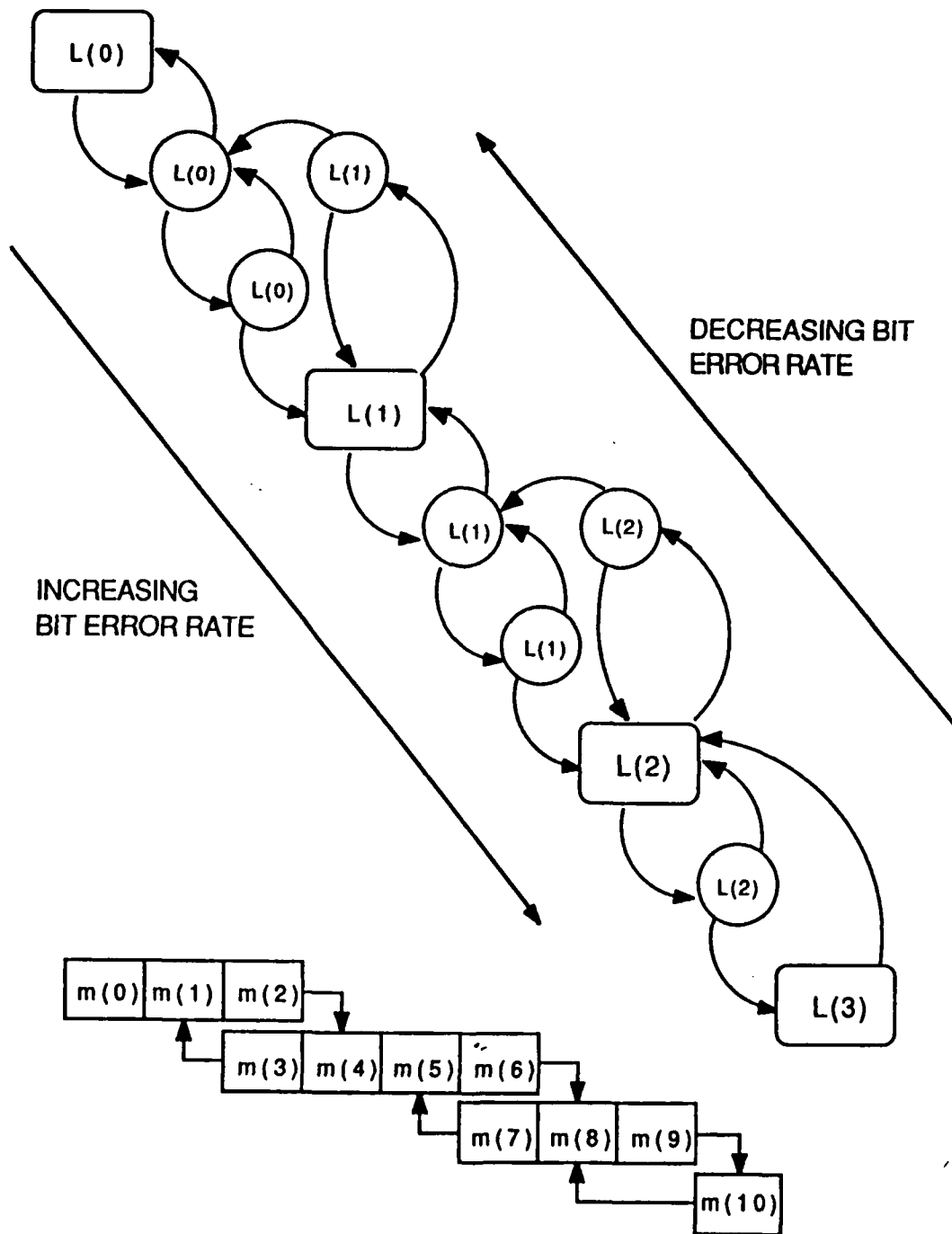
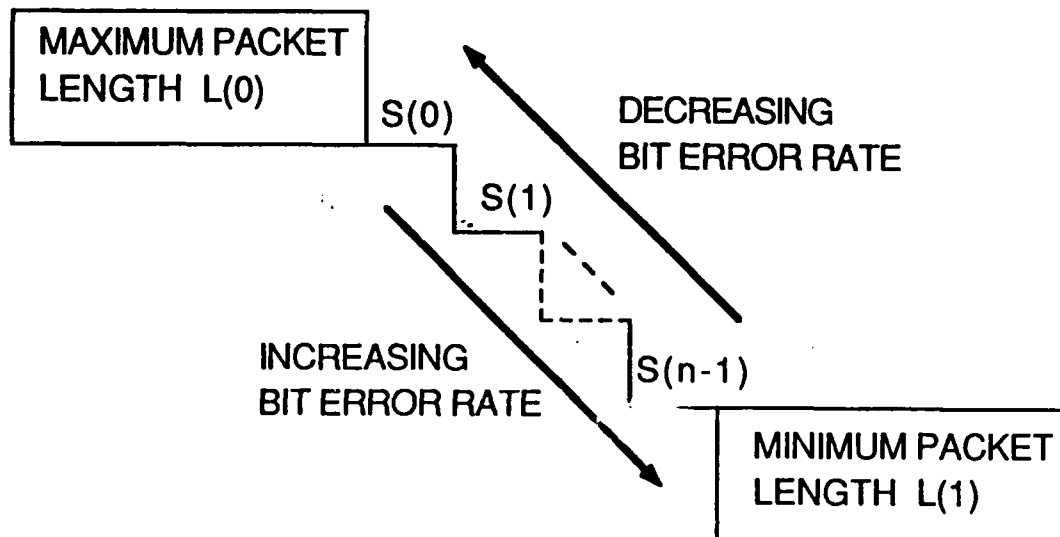


Figure 8. Alternative adaptive strategy structure design



$S(0), S(1), \dots, S(n-1)$ = STEPS (NUMBER OF TRANSMISSION ATTEMPTS TO CHANGE LEVELS)

Figure 9. Two level adaptive strategy

The structure of the two-level strategy, as the name implies, transmits only packets with two lengths. The top level is the maximum packet length and would typically be the length determined to be optimum for the average channel bit error rate of the system employed. The packet length associated with the lower level is the length determined to be optimum for the system during periods of operation under the high bit error rate conditions which occur most frequently. These channel bit error rates are usually not known for a system and must typically be determined from empirical testing of the system and results accumulated for long periods of time.

The steps connecting the levels which determine the number of retransmission attempts or the number of consecutive, successful packet transmissions prior to a length modification are completely variable. Using the alternative structure design, the number of steps to descend to the lower level can be specified to be different than the number of steps to ascend to the top level. This variability allows a system designer to specify how quickly the system will react to a perceived increase or decrease in the forward channel bit error rate. Determining the number of steps between the levels which optimizes the system throughput efficiency, similar to the level determination, is very system dependent. System testing under as many conditions as possible would be required to find the number of steps which optimizes the system throughput.

3. Multi-level Adaptive Strategy Structure

A different strategy is needed in the cases where varying noise, which may be modeled as Gaussian noise, causes an improvement or degradation of the channel quality resulting in a forward channel bit error rate which changes accordingly. The changing bit error rate may vary slowly relative to the time required for packet transmission and in smaller amounts than in the signal fading case. For this type of changing bit error rate condition the adaptive strategy must vary the packet length in smaller increments than the previous strategy in an attempt to optimize the packet length for the channel conditions. The type of adaptive strategy shown in Figure 10 is a *multi-level* adaptive strategy and will step or modify the packet length to various lengths associated with each level ranging from the maximum to the minimum packet length allowed.

Similar to the two-level strategy, the top level is associated with the maximum packet length allowed. This packet length is typically chosen to optimize the system operating with the most frequently occurring bit error rate. Because this strategy targets the cases where the changes in the bit error rate are less severe than in the two-level strategy, the information field lengths of each lower level are typically no less than one-half the length of the level above. The number of levels and steps in the multi-level

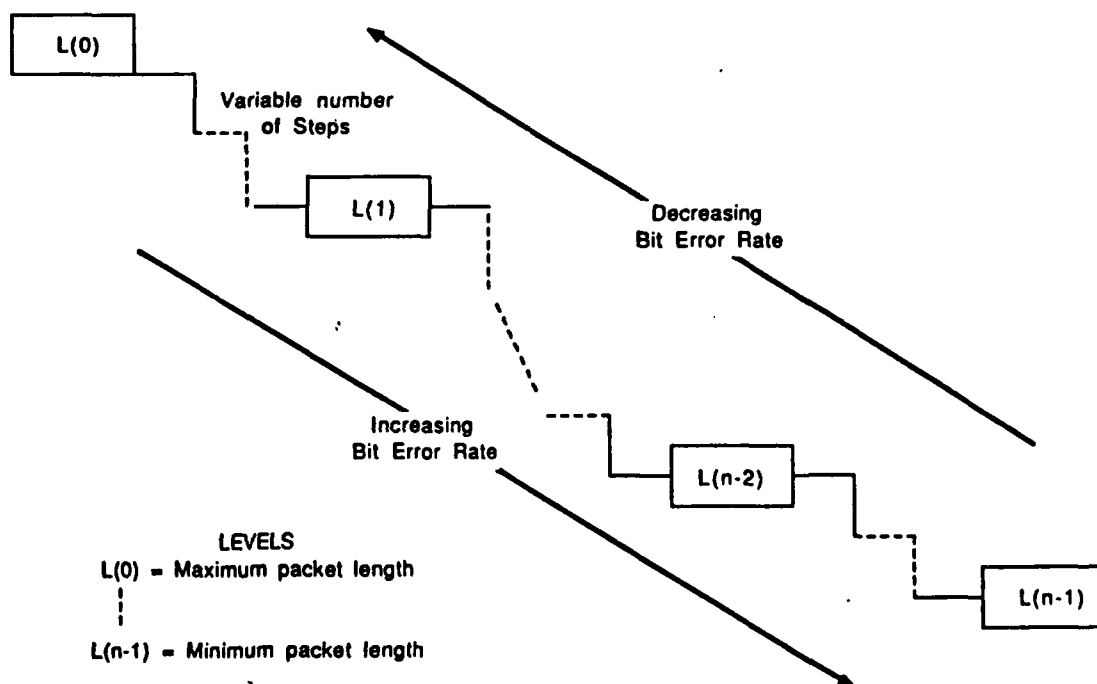


Figure 10. Multi-level adaptive strategy

strategy are completely variable. As the channel bit error rate increases, after a specified number of failed retransmission attempts of a packet with length greater than the minimum length allowed, the packet length is changed to the next lower level and retransmitted. This procedure is continued until the packet is successfully transmitted or the lowest level is reached. Once a specified number of packets at a certain level have been successfully transmitted, the packet length of the next packet to be transmitted is increased to the next level. As long as packets are transmitted successfully, stepping up the packet length to the next level continues until the highest level is reached and the packet length is the maximum for the system.

C. ADAPTIVE STRATEGY IMPLEMENTATION

The main feature of the adaptive strategies is the modification of the length of the message packet after a specified number of consecutive retransmission attempts or successful transmissions. All ARQ protocols require that the transmitting station store all transmitted packets until their successful transfer is acknowledged by the receiving station. Since the transmitted information is stored until acknowledged, it can be broken up into smaller packets to be retransmitted at lengths associated with the new level of the adaptive structure. For the case where the bit error rate decreases and the packet length is to be increased, the number of information bits in the information field is simply increased accordingly.

The addition of the adaptive feature is designed to be strictly a software modification which adds only some additional software overhead requirements to the transmitting station. The transmitting station must know what information has been successfully transferred, the information that is attempting to be transferred, and the beginning of the next block to be transferred. In addition, the number of attempted transfers at the current packet length, as well as the number of consecutive packets transmitted successfully, must be accounted for. The addition of a state vector for the two-level or multi-level strategy can take care of most of the additional packet accounting required. The state vector is used by the transmitting station software to implement the adaptive strategy and identifies the proper length of the packets to be transmitted at any particular time. The adaptive strategy can be designed to be as complex or as simple as the system designer determines necessary to optimize the system throughput.

The state vector is always updated when an acknowledgment for receipt or rejection of a transmitted packet is received. For an accepted packet, the state is updated by moving one position to the left in the vector and the next packet of information is as-

sembled into a frame at the length specified by this new state and transmitted. If a frame is rejected, the state is updated by moving to the right and the new state specifies if the packet length is to be modified. When modification of the packet length is indicated prior to retransmitting the frame, the local station must verify that the receive sequence number $N(r)$ of the remote station is the same as the send sequence number $N(s)$ of the frame that the local station is attempting to transfer. If the received frame from the remote station which contained the negative acknowledgment (NAK) was received without error, then the verification is complete. But, if the frame which contained the NAK is received with errors or not received at all, the verification procedure can be accomplished through the use of a P/F exchange. A flow diagram of the verification procedure is shown in Figure 11. The local station transmits a supervisory command frame with the P set to "1", which forces the remote station to send a supervisory response frame with the F bit set to "1". The remote station also provides its current $N(r)$ value in this response frame. Then if a modification of the packet length is required, the frame is reassembled with the new packet length, a new CRC is performed and the checksum is put in the CRC field. The frame is then retransmitted and the data transfer process continues.

D. ADAPTIVE STOP-AND-WAIT (SW) PROTOCOL

1. Adaptive SW Operations and Example Notation

Point-to-point communications using adaptive SW protocols operate in a very similar manner to typical communications with SW protocols [Ref. 1]. The sessions are initiated in the same manner and when operating in the asynchronous balanced mode (ABM), the information can flow in both directions [Ref. 2]. Since a SW protocol can be thought of as a GBN protocol with a send window equal to one, the modulus used for the sequence numbering sets $N(s)$ and $N(r)$ would be modulo 2. Acknowledgments (ACK) for accepted frames and negative acknowledgments (NAK) for rejected frames are returned by the $N(r)$ value in the control field of the frame transmitted from the remote station. Error recovery from the frames that are lost and consequently never acknowledged, or discarded due to errors received in transmission, is handled by the time-out function. The minimum length of time set for the time-out function is typically equal to two times the propagation delay plus the transmission time required for the largest frame plus a small amount of time for processing delay and a small safety margin to help eliminate premature time-outs [Ref. 1].

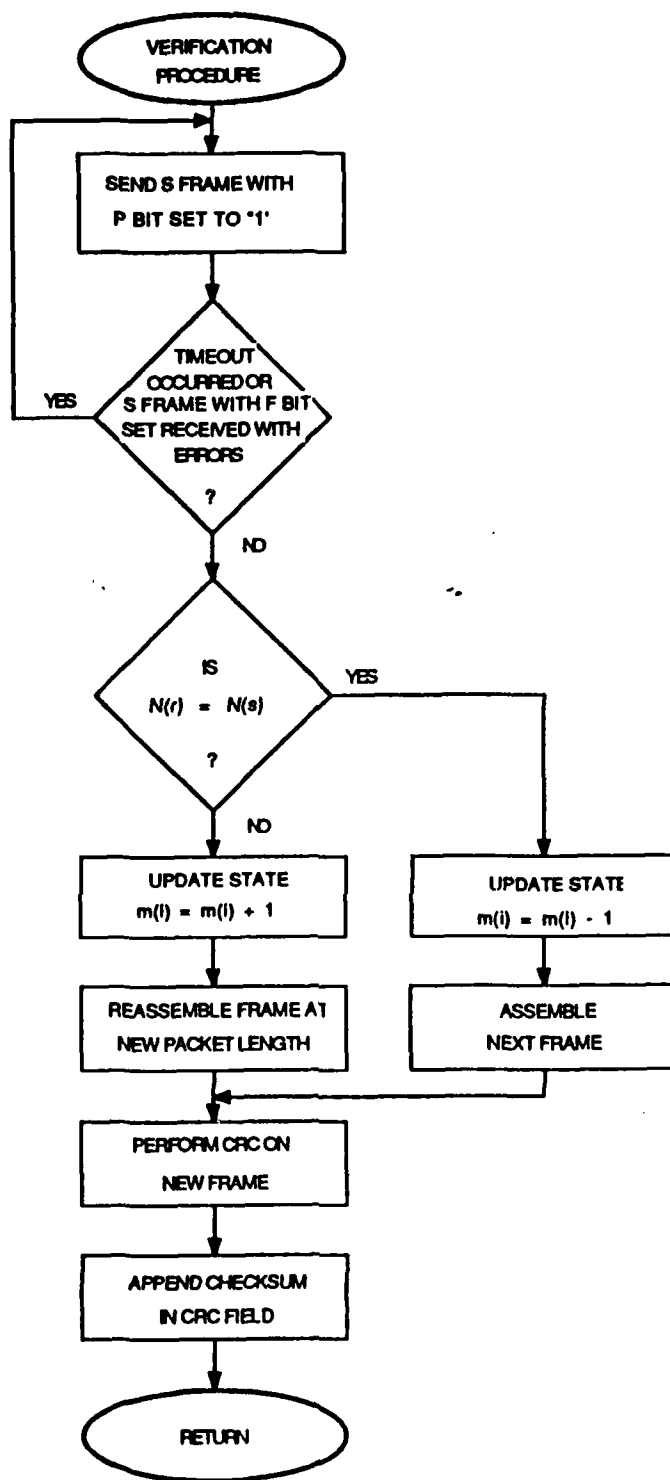


Figure 11. Verification procedure and packet length modification

Figure 12 and Figure 13 illustrate some typical operations for point-to-point communications sessions between two stations A and B. These examples are for short range, RF packet switching networks which may be subject to noise and signal fading so the probability P_b of bit error may be different for the channel from A to B than from B to A. Both stations are combined stations so the mode used for the data transfer is ABM. In these examples, I frames are represented by a single "I" symbol of various lengths for each frame. The maximum length for each frame is specified by the current state $m(i)$ of the state vector m . The S and U frames are specified by the short, double "II" symbols and are all transmitted at a length of 48 bits. The length of the frame symbols represents the time required to transmit the frame, including the small amount of time used to process the frame. The slanted lines indicate the amount of propagation delay in the transmission of the frames. These lines are only slightly angled which reflects how small the propagation delay is relative to the frame transmission time for the short range system. The disruptions in some of the propagation lines, illustrated by a rotated "Z" symbol, indicate that the packet has received at least one error in transmission.

Often times a quick exchange of S or U frames is desired for a station to verify that the frame it is attempting to send is the same as the frame that the remote station is expecting to receive. This exchange can best be handled by using the P/F bit of the control field which initiates the checkpointing mechanism. To force the remote station to respond with either an S or U frame, the I frames are restricted to only containing commands. This restriction never allows an I frame to be used in responding to a received frame with the P bit set to "1".

The notation used in the examples of protocol operations to describe the frames is as follows: [Ref. 4]

L T N(s) N(r) P/F.

- **L** represents the address of the frame. For frames which contain commands, the address is specified as the station receiving the frame. For frames which contain responses, the address is the station sending the frame.
- **T** represents the abbreviation of the frame function. For I frames T is "I" and for S and U frames it is the specific function the frame performs (e.g., RR, RNR, SETM, UA, DISC, etc.).
- **N(s)** is the send sequence number and will only be used for I frames.
- **N(r)** is the receive sequence number and is present for all I and S frames.

- P/F represents either the P or F bit and is only present when the frame sets this bit to "1".

2. An Example of One-Way Information Transfer

Figure 12 illustrates an example of a point-to-point communications session between stations A and B with information being transferred only from A to B. The protocol established for these stations is multi-level adaptive SW, with each station using a four-level structure specified in Table 1. The adaptive structure is implemented by use of the state vector specified in row three of Table 1.

Table 1. ADAPTIVE SW STRUCTURE REPRESENTATION AND STATE VECTOR

Adaptive Structure Levels	L(0)	L(0)	L(0)	L(1)	L(2)	L(3)
Max Information Field Length (bits)	1000	1000	1000	500	250	125
State Vector	m(0)	m(1)	m(2)	m(3)	m(4)	m(5)

Station A initiates the communication session by sending a set mode command frame with the P bit set to "1". B is ready to receive data and responds with an Un-numbered Acknowledgement (UA) response frame. Using its own address in the frame, B is indicating that this is a response frame and the F bit is set to "1" in response to the P bit in the command frame from A. Upon transfer of the UA frame, B sets its N(s) and N(r) values to zero, initializes its current state to m(0), and enters the information transfer state. When A receives the UA response from B, it performs the same initializations and commences transferring I frames.

The stations commence transferring frames according to the adaptive SW protocol which requires they stop after each transmitted frame and wait for a response. Since B does not have any data to transfer to A during this session, B provides a response for each frame received utilizing a supervisory frame and the appropriate N(r) value. Essentially, B provides an acknowledgement (ACK) for each frame received without errors and a negative acknowledgement (NAK) for each frame received with errors. For the system errors caused by transmitted frames that are lost and not received, recovery is performed by activating the time-out function.

A transmits the first I frame BI00, at the length specified by the current state m(0). Since this frame is received at B with errors being detected, B responds with a

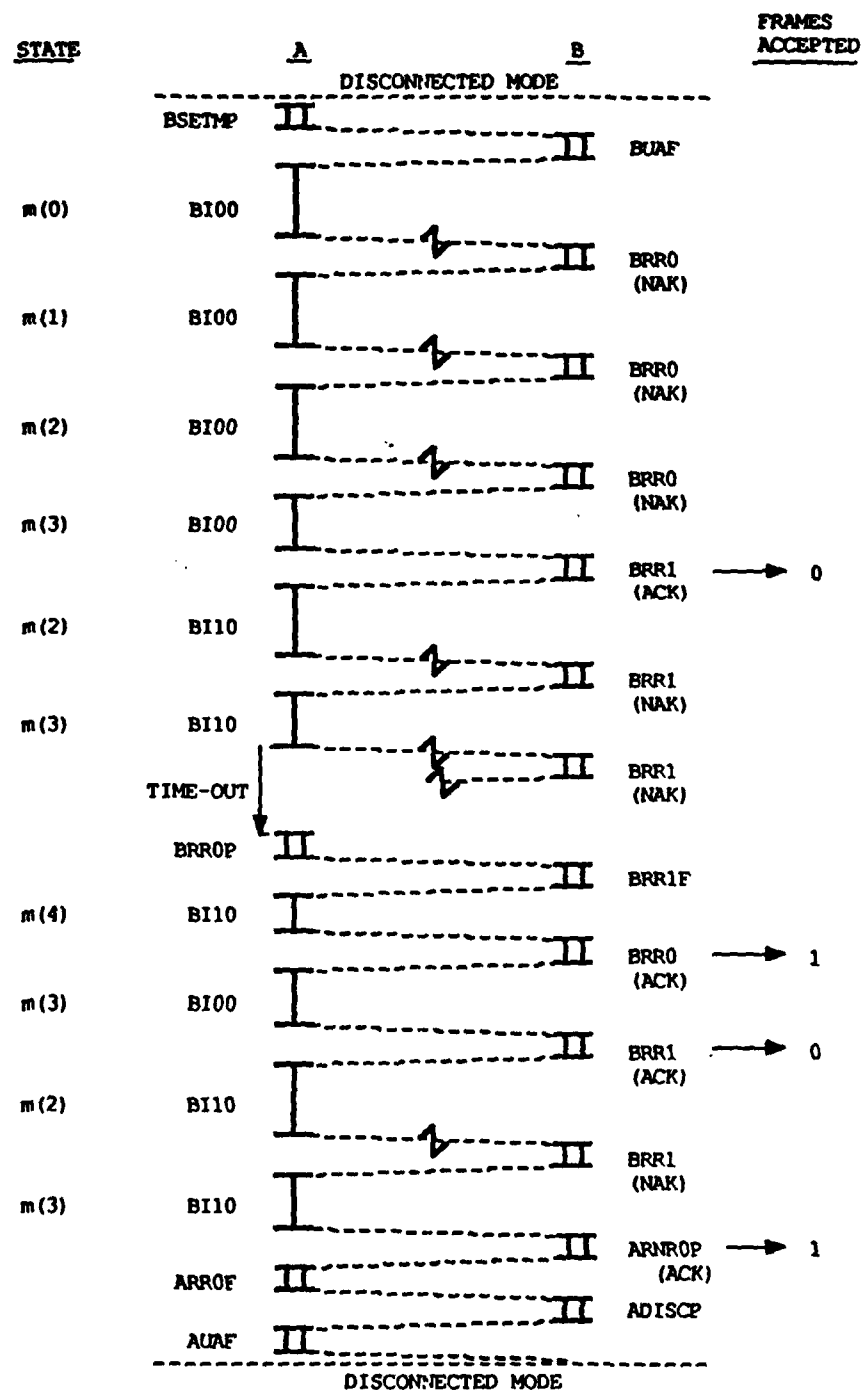


Figure 12. Adaptive SW - an example of one-way information transfer

NAK, by setting the $N(r)$ value equal to $N(s)$ value of the frame just received. Upon receipt of frame BRR0, A updates its state by moving one state to the right. The current state is now $m(1)$ which specifies that the frame is to be retransmitted using the same maximum packet length constraint. The attempted data transfer continues and after the third NAK is received by A, an updating of the current state to $m(3)$ specifies a modification to the frame length prior to retransmitting the frame. To modify the length, A divides the information field into two blocks, places the first block into the information field, performs a CRC on the new frame, appends the new checksum, and transmits BI00. The second block of information is now the next block of information waiting to be transferred.

After the successful transfer of BI00 and the ACK response is received from B, A updates its state to $m(2)$. This state specifies that the maximum frame length is increased to the next level. Therefore A assembles and transmits frame BI10 at the length associated with level $L(0)$. Frame BI10 is received at B with errors and B responds with a NAK. After updating its state to $m(3)$, A must modify the frame length to the length associated with level $L(1)$ in the same manner as before. A then retransmits the modified frame BI10. The frame BI10 again suffers transmission errors, so B responds with a NAK. Since the NAK sent by B is never received at A, the time-out function at A expires and a recovery action is initiated. In this case, A sends a supervisory type command frame BRR0P, with the P bit set to "1". This ensures the remote station B is still active and updates the $N(r)$ value. B responds with BRR1F, with the F bit set to "1" in response to the set P bit of the previous frame.

Data transfer from A to B continues with the length of the I frames increasing and decreasing as specified by the state vector. After A has successfully transferred frame BI10, station B is interested in ending the communications session. B sends the command ARNR0P with the P bit set, indicating that B is not ready to receive any additional data from A. A responds with ARR0F indicating receipt of the command frame. B then initiates the disconnect procedures with station A. When A responds with the UA response, both stations enter the disconnect mode.

3. An Example of Two-Way Information Transfer

Figure 13 illustrates another example of a point-to-point communications session between stations A and B. In this example, the data link has already been established and both stations have information to transfer. Whenever both stations have information to transfer, the ACK(s) and NAK(s) for the received frames are returned with or "piggybacked" on the transmitted I frames. This case requires both local stations

to receive, without errors, an I frame transmitted by the remote station to determine if the previous frame transmitted by the local station was accepted. A local station will discard frames received with errors and uses the time-out function to initiate a retransmission. For this situation, setting the optimal length of the time-out function is very important for maximizing the information throughput.

In this adaptive SW example, both stations A and B are using multi-level strategies with the adaptive structure represented by Table 2. As in the previous example, the state vector shown in the third row of Table 2 is used to implement the adaptive strategy.

Table 2. ADAPTIVE SW STRUCTURE AND STATE VECTOR

Adaptive Structure Levels	L(0)	L(0)	L(1)	L(2)	L(3)
Max Information Field Length (bits)	1000	1000	500	250	125
State Vector	m(0)	m(1)	m(2)	m(3)	m(4)

A sends frame BI00 at the frame length specified by the state m(0). Since the frame arrives at B suffering from errors which occurred during transmission, B waits for the time-out to expire before sending AI00. The frame AI00 sent by B provides a piggybacked negative acknowledgment displayed by the N(r) value remaining equal to "0". As B transmits the last bit of frame AI00, the time-out function is reset and started. Upon receipt of AI00, A performs the CRC on the frame and accepts the frame since no errors are detected. A then updates its state to m(1) and retransmits the previous frame at the same length. B detects errors in the received frame BI01 and again discards the frame. Since B is unaware that its previous frame AI00 was accepted, after the time-out expires, B updates its state and retransmits the frame.

A detects errors in the received frame AI00, and discards the frame. After the time-out expires, A updates its state to m(2) and determines that a packet length change is required for the next retransmission. Since the N(r) value at B has not been verified to be the same as the N(s) value at A following the last transmission of this packet, A transmits supervisory command frame BRR1P. This supervisory frame contains a command and with the P bit set to "1", forces B to respond. Since an I frame is restricted to containing a command only, B must use a S or U frame with the F bit set to "1" for the response. Limiting this checkpointing exchange to the 48 bit S or U frames provides

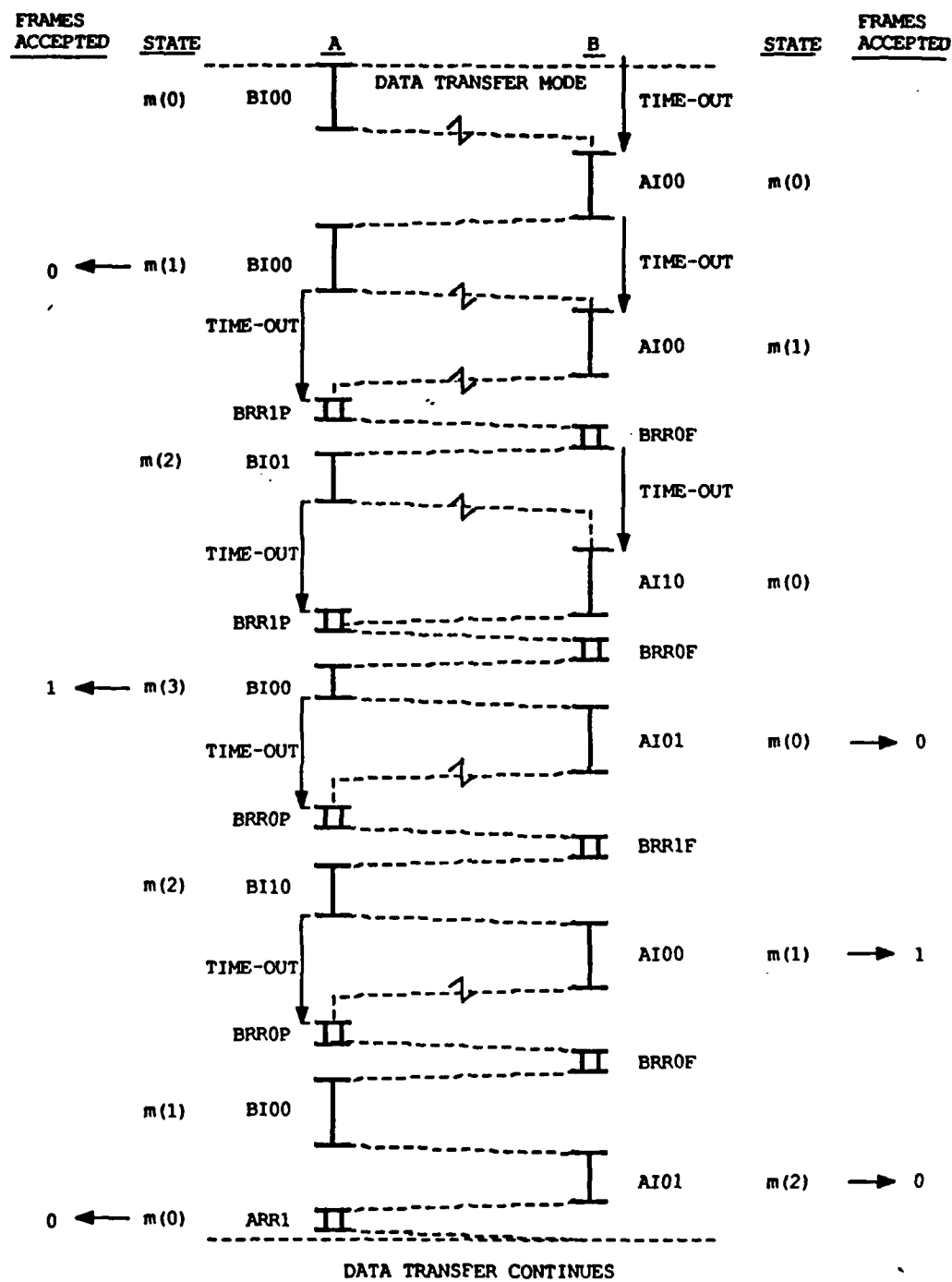


Figure 13. Adaptive SW - an example of two-way information transfer

high probability that the frames will be received without errors and the amount of time for the exchange is very short. When A receives the response frame BRR0F from B, the frame length is modified to the length specified by state $m(2)$ using the same length modification procedures as described in the previous example. A then retransmits this modified frame at the new length and resets the time-out function.

Since the modified frame from A is received with errors, B discards the frame and waits for the time-out to expire before sending frame AI10, A's time-out expires prior to receiving frame AI10 from B and this causes A to initiate the necessary verification procedure prior to modifying the frame length for retransmission. When A receives the error-free frame AI10 from B, A updates the current state to $m(3)$ and retransmits frame BI00 at the new frame length. B responds to the command frame BRR1P from A and prepares to retransmit the previous frame after receipt of the next frame. B then receives frame BI00 from A which indicates B's previous frame was accepted. B then updates the current state to $m(0)$ and transmits frame AI01. The I frame AI01 is discarded by A once errors are detected and upon expiration of the time-out, A sends a supervisory command to determine B's $N(r)$ value. When B responds with BRR1F, A updates its state to $m(2)$ and transmits frame BI10.

The transfer of data continues and when a station no longer has any I frames to transfer, it responds to received frames from the remote station with supervisory frames and the appropriate $N(r)$. If any new information should arrive and require transmission, it is put into I frames with lengths specified by the current state of the state vector and given the next $N(s)$. If both stations run out of I frames to transfer, then one station will initiate the disconnect procedure as in the previous example. Once both stations acknowledge the disconnect commands, then they enter the disconnect mode.

V. ADAPTIVE SW SIMULATION

A. INTRODUCTION

The methods available to determine the throughput efficiency of any ARQ protocols, short of using or building an actual packet switching data communications network and collecting empirical data, are through analysis and computing numerical results or through simulation. The difficulty in performing the analysis of the system operating with the adaptive protocol is in being able to mathematically or statistically model the system properly. Often, to model all states and processes of the system successfully, certain assumptions and preconditions must be made. Many times the assumptions pose no restrictions on the results and the analysis results closely resemble actual system performance.

Simulation of the adaptive protocol provides another method of determining the throughput efficiency of the system. To simulate the adaptive protocol operating on a system, the system operation and parameters must be reproduced. In addition, the modeling of the statistical nature of the channel allows the simulation to reflect the actual efficiency of the system. As in analysis, assumptions and preconditions are made to perform the simulation and to the extent that these conditions are accurate, the simulation results may be indicative of actual system performance.

B. SIMULATION DEVELOPMENT

1. Model Design

The simulation model was designed to produce results from which the throughput efficiency of the system operating with one of the adaptive SW strategies could be computed. The model was not developed to show that the adaptive protocols worked correctly. The simulation modeled a point-to-point, RF, data communications system with only two stations. The model is for a system operating over a short distance where the propagation delay is very short relative to the time required for packet transmission. The bit rate for the data transfer is 4800 bits per second. The computer systems modeled in the simulation are capable of processing the packets with almost negligible delay and they are assumed to be dedicated to the data communications network. The RF channel connecting the stations is modeled to be a noisy channel, and the probability of bit error for the channel can be specified. To model the adaptive SW protocol, the packet sizes

had to be controllable and the overhead associated with each packet held constant. The conditions used for the system in developing the simulation model are:

- Only one station transfers data.
- The station transmitting data is always saturated with information.
- The ACK(s) and NAK(s) are only 48 bits and always received without errors.
- Delays associated with the processing time are considered negligible and not used for the throughput efficiency calculations.
- The bit errors which occur during transmission are considered independent.

These conditions were applied to the model to provide results which could be used to calculate the throughput efficiency.

2. Simulation Program

The simulation for the data communications system was performed using the PC-NETWORK II.5 simulation software package from the CACI Products Company [Ref. 7]. This simulation package is intended to model computer system configurations and local area networks. Because the package has the capability to model the instructions required for the local area network, modeling some packet switching features was possible.

To simulate the adaptive SW protocol, two stations A and B, called processing elements (PE)s, were created. Within these PE(s) the instructions for sending the various sized frames, receiving data frames or acknowledgments, and sending ACK(s) or NAK(s) were developed. To simulate the channel connecting stations A and B, a transfer device (TD) was created. The parameters specified in the creation of the TD established the bit rate, the overhead associated with the parity checks for each word, and the control overhead for each frame. These TD parameters established the transmission time required for the data and control frames.

The software which simulates the ARQ adaptive SW protocol operating on the two PE(s) is in the form of software modules. A module contains a list of executable instructions to perform the various functions on a specified PE. Modules were created for each size packet required by the adaptive structure. Other modules were used to decide what size packet to send based on the current value of the state vector. Each module has a set of preconditions which must be satisfied prior to its execution. Through the setting of the preconditions for the modules, the software logic for the adaptive protocol could be implemented.

The module preconditions used to implement the adaptive protocol are described as message and semaphore based preconditions. The message preconditions require that a specified message be received prior to that condition being satisfied. The semaphore preconditions require that a set of semaphore values be checked and satisfied before execution of the module begins.

Semaphores are programming elements which can be used to "flag" certain events or conditions, or used as counters. Semaphores have a count associated with them which is incremented when the semaphore is set and decremented when reset. Semaphores are set or reset through the execution of the appropriate semaphore instruction. Through the use of semaphores, updating the state, which implements stepping up or down the adaptive structure, can be simulated.

The TD provided in the simulation package did not have any capability of simulating a noisy channel which would cause some frames to be received with errors. To simulate this noisy channel condition, a transmitter would send randomly a number of I frames which contain errors. This number of transmitted I frames which have errors would be a percentage of the total number of frames sent. The percentage represents the probability of a frame error P_e for a specified P_b and a frame length. Using the assumption that all bit errors are independent, the probability of a frame error P_e for a certain P_b of the channel is determined from

$$P_e = 1 - (1 - P_b)^{l+h} \quad (1)$$

where l specifies the information field size of the frame and h specifies the frame overhead. As can be seen from Equation (1), the P_e of an I frame decreases as the frame length decreases which means a lower percentage of I frames containing errors are transmitted. Since the percentages used to select the I frames which contain errors are set for a simulation run, the channel P_b had to also be set and not varied while performing the simulation.

C. SIMULATION RESULTS

The adaptive SW protocol for several different adaptive structures were simulated using the simulation model and program described in the previous section. In all simulations, the configuration was point-to-point and operated under the aforementioned assumptions.

These simulations provide results that are used to calculate the gain or loss in a station's information throughput efficiency ρ when using an adaptive SW strategy versus

a non-adaptive SW protocol. The throughput efficiency for the simulation is defined as the number of information bits a station transfers divided by the total number it transmits. The simulation program used the output reports provided by PC-Network II.5 to record the total number of frames transmitted, the number of frames received without errors, and the number of frames received with errors.

To calculate the station throughput efficiency for a particular adaptive strategy, the number of information bits transferred is first computed. This result is obtained by summing the product of the number of frames successfully transferred for each different length times its respective frame efficiency times the number of bits in each frame. Frame efficiency is the ratio of the number of information bits to the total number of bits in a frame. Then the total number of bits transmitted by a station is computed. This number is obtained by summing the product of the number of frames transmitted for each different length times the number of bits in each frame. The throughput efficiency is obtained from the quotient of the number of bits transferred divided by the number of bits transmitted. The throughput efficiency results computed as described above for three different adaptive SW strategies and a non-adaptive SW protocol are shown in Figure 14.

Three adaptive strategies were used in the simulation and all three showed a higher throughput efficiency than the non-adaptive SW when the P_b was greater than 1×10^{-3} . The three adaptive strategies were different and the relative merit of each can be seen from the results. The results indicate that all the adaptive strategies are very similar in throughput efficiency when the channel P_b is less than about 1×10^{-3} . For data transfer operations when the P_b gets larger than 1×10^{-3} , the multi-level strategy with the state vector that was a single segment representing a structure with the same ascending and descending steps between levels, provided the best throughput efficiency. This strategy is referred to as "Multi-level Strategy 1" in Figure 14. The multi-level strategy with the state vector composed of multiple segments representing a structure with different ascending steps than descending steps between levels, indicated a ρ value consistently less than or equal to the multi-level strategy with a single segment state vector. This strategy is referred to as "Multi-level Strategy 2" in Figure 14. The two-level strategy indicated the worst ρ performance of the adaptive strategies, but as the P_b got very high, approaching 1×10^{-2} , the efficiency curve appeared to level off. This leveling off might indicate that the two-level adaptive strategy might be quite effective for large swings of the channel P_b .

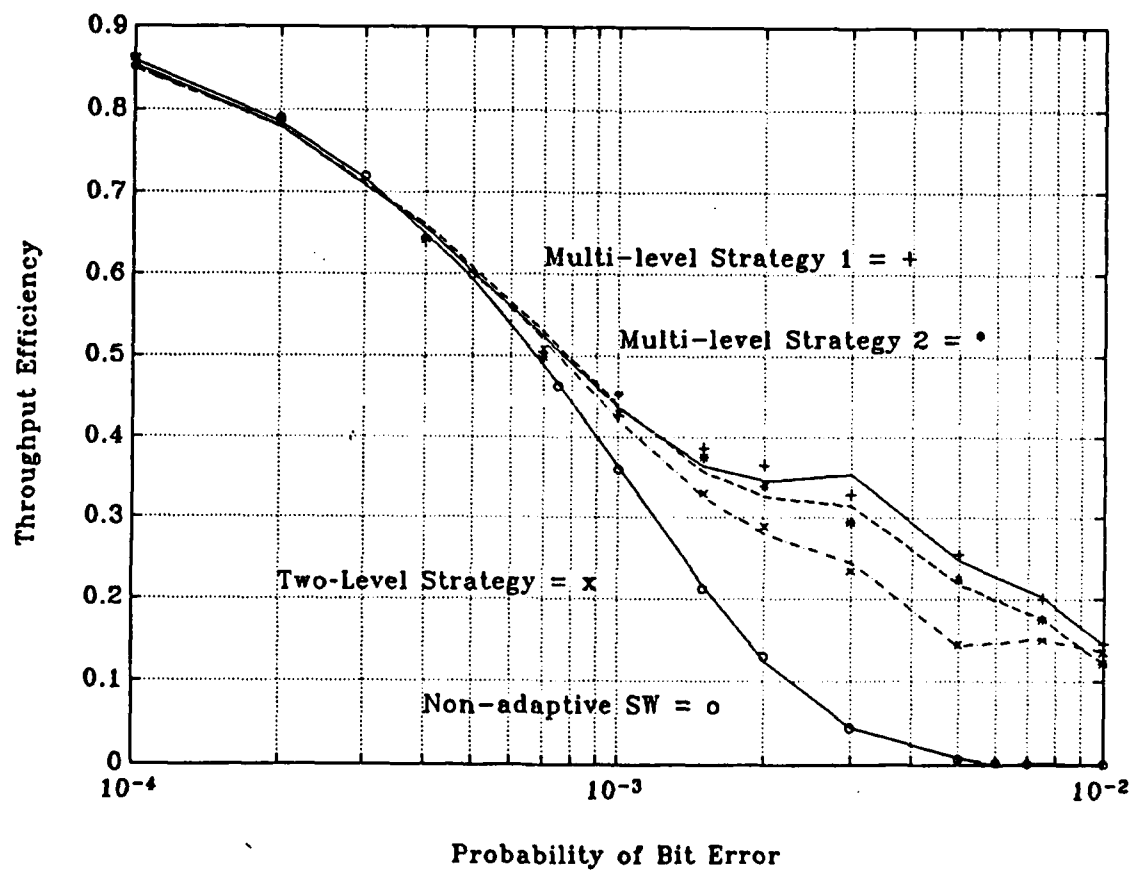


Figure 14. Simulation throughput efficiency vs probability of bit error P_b

VI. CONCLUSIONS

Adaptive SW protocols used for short range, RF networks performing packet switching data communications, where the propagation delays are short relative to the packet transmission time, can increase the information throughput efficiency over a typical SW protocol. As shown by the simulation results this improvement is particularly evident for a communication system operating in an environment with high channel P_b . The addition of the adaptive features to an existing SW protocol should be a fairly inexpensive modification as it is designed to be only a software modification with no additional hardware requirements.

Several different adaptive strategies were presented; all adapt the packet length as the channel P_b increases or decreases. To determine how many levels or steps to use in an adaptive structure for a system would require an analysis of the noise or interference of the specific system or empirical data which reflects system conditions. The different strategies were presented to provide a system designer with the ability to optimize a specific system for the particular noise or fading conditions experienced.

The simulation program used to model the adaptive SW protocol can possibly be modified to simulate other ARQ protocols. This would allow comparisons with other protocols in an attempt to further optimize the system.

APPENDIX A. SIMULATION INPUT FILES

CACI NETWORK II.5 RELEASE 4.01

4 Level Adaptive SW strategy with multi-segment state vector - Probability of Bit Error is 1×10^{-4} .

This is the input file generated on the CACI NETWORK II.5 editor NETIN, used to run the simulation on the NETWORK program.

```
1  * 4 Level Adaptive SW strategy
2  * NETIN RELEASE 4.01
3
4  ***** GLOBAL FLAGS
5  GLOBAL FLAGS =
6    RANDOMIZER = 3
7    BATCH = NO
8
9  ***** PROCESSING ELEMENTS - SYS. PE. SET
10  HARDWARE TYPE = PROCESSING
11    NAME = STATION A
12    BASIC CYCLE TIME = 1.000000 MICROSEC
13    INPUT CONTROLLER = YES
14    MESSAGE LIST SIZE = 640000.0
15    LOSE OVERFLOW MESSAGES = NO
16    INSTRUCTION REPERTOIRE =
17    INSTRUCTION TYPE = MESSAGE
18    NAME ; SEND GOOD DATA FRAME L(0)
19    MESSAGE ; GOOD DATA
20    LENGTH ; 1000 BITS
21    DESTINATION PROCESSOR ; STATION B
22    QUEUE FLAG ; YES
```

23 RESUME FLAG ; YES
24 ALLOWABLE BUSSES ;
25 CHANNEL 1
26 NAME ; SEND BAD DATA FRAME L(0)
27 MESSAGE ; BAD DATA
28 LENGTH ; 1000 BITS
29 DESTINATION PROCESSOR ; STATION B
30 QUEUE FLAG ; YES
31 RESUME FLAG ; YES
32 ALLOWABLE BUSSES ;
33 CHANNEL 1
34 NAME ; SEND GOOD DATA FRAME L(1)
35 MESSAGE ; GOOD DATA
36 LENGTH ; 500 BITS
37 DESTINATION PROCESSOR ; STATION B
38 QUEUE FLAG ; YES
39 RESUME FLAG ; YES
40 ALLOWABLE BUSSES ;
41 CHANNEL 1
42 NAME ; SEND BAD DATA FRAME L(1)
43 MESSAGE ; BAD DATA
44 LENGTH ; 500 BITS
45 DESTINATION PROCESSOR ; STATION B
46 QUEUE FLAG ; YES
47 RESUME FLAG ; YES
48 ALLOWABLE BUSSES ;
49 CHANNEL 1
50 NAME ; SEND GOOD DATA FRAME L(2)
51 MESSAGE ; GOOD DATA
52 LENGTH ; 250 BITS
53 DESTINATION PROCESSOR ; STATION B
54 QUEUE FLAG ; YES
55 RESUME FLAG ; YES
56 ALLOWABLE BUSSES ;

57 CHANNEL 1
 58 NAME ; SEND BAD DATA FRAME L(2)
 59 MESSAGE ; BAD DATA
 60 LENGTH ; 250 BITS
 61 DESTINATION PROCESSOR ; STATION B
 62 QUEUE FLAG ; YES
 63 RESUME FLAG ; YES
 64 ALLOWABLE BUSSES ;
 65 CHANNEL 1
 66 NAME ; SEND GOOD DATA FRAME L(3)
 67 MESSAGE ; GOOD DATA
 68 LENGTH ; 125 BITS
 69 DESTINATION PROCESSOR ; STATION B
 70 QUEUE FLAG ; YES
 71 RESUME FLAG ; YES
 72 ALLOWABLE BUSSES ;
 73 CHANNEL 1
 74 NAME ; SEND BAD DATA FRAME L(3)
 75 MESSAGE ; BAD DATA
 76 LENGTH ; 125 BITS
 77 DESTINATION PROCESSOR ; STATION B
 78 QUEUE FLAG ; YES
 79 RESUME FLAG ; YES
 80 ALLOWABLE BUSSES ;
 81 CHANNEL 1
 82 NAME ; START DATA TRANSFER
 83 MESSAGE ; DATA
 84 LENGTH ; 1 BITS
 85 DESTINATION PROCESSOR ; STATION A
 86 QUEUE FLAG ; YES
 87 RESUME FLAG ; YES
 88 INSTRUCTION TYPE = SEMAPHORE
 89 NAME ; A SET MESSAGE TRANSFERRED
 90 SEMAPHORE ; A MESSAGE TRANSFERRED

91 SET/RESET FLAG ; SET
92 NAME ; A RESET MESSAGE TRANSFERRED
93 SEMAPHORE ; A MESSAGE TRANSFERRED
94 SET/RESET FLAG ; RESET
95 NAME ; A SET LOCNTR
96 SEMAPHORE ; A LOCNTR
97 SET/RESET FLAG ; SET
98 NAME ; A RESET LOCNTR
99 SEMAPHORE ; A LOCNTR
100 SET/RESET FLAG ; RESET
101 NAME ; A SET L1GOOD
102 SEMAPHORE ; A L1GOOD
103 SET/RESET FLAG ; SET
104 NAME ; A RESET L1GOOD
105 SEMAPHORE ; A L1GOOD
106 SET/RESET FLAG ; RESET
107 NAME ; A SET L1BAD
108 SEMAPHORE ; A L1BAD
109 SET/RESET FLAG ; SET
110 NAME ; A RESET L1BAD
111 SEMAPHORE ; A L1BAD
112 SET/RESET FLAG ; RESET
113 NAME ; A SET L2GOOD
114 SEMAPHORE ; A L2GOOD
115 SET/RESET FLAG ; SET
116 NAME ; A RESET L2GOOD
117 SEMAPHORE ; A L2GOOD
118 SET/RESET FLAG ; RESET
119 NAME ; A SET L2BAD
120 SEMAPHORE ; A L2BAD
121 SET/RESET FLAG ; SET
122 NAME ; A RESET L2BAD
123 SEMAPHORE ; A L2BAD
124 SET/RESET FLAG ; RESET

```

125      NAME ; A SET L3GOOD
126      SEMAPHORE ; A L3GOOD
127      SET/RESET FLAG ; SET
128      NAME ; A RESET L3GOOD
129      SEMAPHORE ; A L3GOOD
130      SET/RESET FLAG ; RESET
131  NAME = STATION B
132      BASIC CYCLE TIME =      1.000000 MICROSEC
133      INPUT CONTROLLER = YES
134      MESSAGE LIST SIZE = 640000.0
135      LOSE OVERFLOW MESSAGES = NO
136      INSTRUCTION REPERTOIRE =
137      INSTRUCTION TYPE = MESSAGE
138      NAME ; SEND GOOD DATA FRAME L(0)
139      MESSAGE ; GOOD DATA
140      LENGTH ; 1000 BITS
141      DESTINATION PROCESSOR ; STATION A
142      QUEUE FLAG ; YES
143      RESUME FLAG ; YES
144      ALLOWABLE BUSSES ;
145      CHANNEL 1
146      NAME ; SEND BAD DATA FRAME L(0)
147      MESSAGE ; BAD DATA
148      LENGTH ; 1000 BITS
149      DESTINATION PROCESSOR ; STATION A
150      QUEUE FLAG ; YES
151      RESUME FLAG ; YES
152      ALLOWABLE BUSSES ;
153      CHANNEL 1
154      NAME ; SEND GOOD DATA FRAME L(1)
155      MESSAGE ; GOOD DATA
156      LENGTH ; 500 BITS
157      DESTINATION PROCESSOR ; STATION A
158      QUEUE FLAG ; YES

```

159 RESUME FLAG ; YES
160 ALLOWABLE BUSSES ;
161 CHANNEL 1
162 NAME ; SEND BAD DATA FRAME L(1)
163 MESSAGE ; BAD DATA
164 LENGTH ; 500 BITS
165 DESTINATION PROCESSOR ; STATION A
166 QUEUE FLAG ; YES
167 RESUME FLAG ; YES
168 ALLOWABLE BUSSES ;
169 CHANNEL 1
170 INSTRUCTION TYPE = PROCESSING
171 NAME ; DATA FRAME PROCESSING
172 TIME ; 1100 CYCLES
173 NAME ; START DATA TRANSFER
174 TIME ; 1 CYCLES
175 INSTRUCTION TYPE = SEMAPHORE
176 NAME ; SET MESSAGE TRANSFERRED
177 SEMAPHORE ; MESSAGE TRANSFERRED
178 SET/RESET FLAG ; SET
179 NAME ; RESET MESSAGE TRANSFERRED
180 SEMAPHORE ; MESSAGE TRANSFERRED
181 SET/RESET FLAG ; RESET
182 NAME ; SET LOCNTR
183 SEMAPHORE ; LOCNTR
184 SET/RESET FLAG ; SET
185 NAME ; RESET LOCNTR
186 SEMAPHORE ; LOCNTR
187 SET/RESET FLAG ; RESET
188 NAME ; SET L1GOOD
189 SEMAPHORE ; L1GOOD
190 SET/RESET FLAG ; SET
191 NAME ; RESET L1GOOD
192 SEMAPHORE ; L1GOOD

```

193         SET/RESET FLAG ; RESET
194         NAME ; SET L1BAD
195         SEMAPHORE ; L1BAD
196         SET/RESET FLAG ; SET
197         NAME ; RESET L1BAD
198         SEMAPHORE ; L1BAD
199         SET/RESET FLAG ; RESET
200
201 ***** BUSSES - SYS.BUS.SET
202 HARDWARE TYPE = DATA TRANSFER
203     NAME = CHANNEL 1
204     CYCLE TIME =      208.333000 MICROSEC
205     BITS PER CYCLE =      1
206     CYCLES PER WORD =      7
207     WORDS PER BLOCK =    125
208     WORD OVERHEAD TIME =    208.333000 MICROSEC
209     BLOCK OVERHEAD TIME =   9999.990000 MICROSEC
210     BUS CONNECTIONS =
211         STATION A
212         STATION B
213
214 ***** MODULES - SYS.MODULE.SET
215 SOFTWARE TYPE = MODULE
216     NAME = A TRANSFER GOOD DATA L(0)
217     PRIORITY =      0
218     INTERRUPTABILITY FLAG = NO
219     CONCURRENT EXECUTION = NO
220     START TIME = 0.0
221     ORED PREDECESSOR LIST =
222         PROCESSOR 0
223     REQUIRED SEMAPHORE STATUS =
224         WAIT FOR ; A MESSAGE TRANSFERRED
225         TO BE ; RESET
226     INSTRUCTION LIST =

```

```

227     EXECUTE A TOTAL OF ; 1 SEND GOOD DATA FRAME L(0)
228     EXECUTE A TOTAL OF ; 1 A SET MESSAGE TRANSFERRED
229     EXECUTE A TOTAL OF ; 1 A RESET LOCNTR
230     ANDED SUCCESSORS =
231     CHAIN TO ; PROCESSOR 0
232     WITH ITERATIONS THEN CHAIN COUNT OF ;      0
233     NAME = PROCESSOR 0
234     PRIORITY =      1
235     INTERRUPTABILITY FLAG = NO
236     CONCURRENT EXECUTION = NO
237     START TIME = 0.0
238     STATISTICAL SUCCESSOR STREAM = 21
239     ORED PREDECESSOR LIST =
240     A TRANSFER GOOD DATA L(0)
241     A TRANSFER BAD DATA L(0)
242     A START
243     A COUNTER RESETTER 1
244     REQUIRED MESSAGES =
245     DATA*
246     REQUIRED SEMAPHORE STATUS =
247     CHAIN IF ; A LOCNTR
248     IS ; < 3
249     STATISTICAL SUCCESSORS =
250     CHOOSE AS SUCCESSOR ; 82.00 % A TRANSFER GOOD DATA L(0)
251     CHOOSE AS SUCCESSOR ; 18.00 % A TRANSFER BAD DATA L(0)
252     INSTRUCTION LIST =
253     EXECUTE A TOTAL OF ; 1 A RESET MESSAGE TRANSFERRED
254     NAME = A TRANSFER BAD DATA L(0)
255     PRIORITY =      0
256     INTERRUPTABILITY FLAG = NO
257     CONCURRENT EXECUTION = NO
258     ANDED PREDECESSOR LIST =
259     PROCESSOR 0
260     REQUIRED SEMAPHORE STATUS =

```

```

261      WAIT FOR ; A MESSAGE TRANSFERRED
262      TO BE ; RESET
263      INSTRUCTION LIST =
264          EXECUTE A TOTAL OF ; 1 SEND BAD DATA FRAME L(0)
265          EXECUTE A TOTAL OF ; 1 A SET MESSAGE TRANSFERRED
266          EXECUTE A TOTAL OF ; 1 A SET LOCNTR
267      ANDED SUCCESSORS =
268          CHAIN TO ; PROCESSOR 0
269          WITH ITERATIONS THEN CHAIN COUNT OF ;      0
270          CHAIN TO ; PROCESSOR 1
271          WITH ITERATIONS THEN CHAIN COUNT OF ;      0
272      NAME = A START
273          PRIORITY =      9
274          INTERRUPTABILITY FLAG = NO
275          CONCURRENT EXECUTION = NO
276          START TIME = 0.0
277          ALLOWED PROCESSORS =
278              STATION A
279          INSTRUCTION LIST =
280              EXECUTE A TOTAL OF ; 1 START DATA TRANSFER
281          ANDED SUCCESSORS =
282              CHAIN TO ; PROCESSOR 0
283              WITH ITERATIONS THEN CHAIN COUNT OF ;      0
284      NAME = A TRANSFER GOOD DATA L(1)
285          PRIORITY =      0
286          INTERRUPTABILITY FLAG = NO
287          CONCURRENT EXECUTION = NO
288          STATISTICAL SUCCESSOR STREAM =      1
289          ANDED PREDECESSOR LIST =
290              PROCESSOR 1
291          REQUIRED SEMAPHORE STATUS =
292              WAIT FOR ; A MESSAGE TRANSFERRED
293              TO BE ; RESET
294          INSTRUCTION LIST =

```

```

295     EXECUTE A TOTAL OF ; 1 A SET L1GOOD
296     EXECUTE A TOTAL OF ; 1 SEND GOOD DATA FRAME L(1)
297     EXECUTE A TOTAL OF ; 1 A SET MESSAGE TRANSFERRED
298     EXECUTE A TOTAL OF ; 1 A RESET L1BAD
299     ANDED SUCCESSORS =
300     CHAIN TO ; A COUNTER RESETTER 1
301     WITH ITERATIONS THEN CHAIN COUNT OF ;      0
302     CHAIN TO ; PROCESSOR 1
303     WITH ITERATIONS THEN CHAIN COUNT OF ;      0
304     NAME = PROCESSOR 1
305     PRIORITY =      4
306     INTERRUPTABILITY FLAG = NO
307     CONCURRENT EXECUTION = NO
308     STATISTICAL SUCCESSOR STREAM = 35
309     ORED PREDECESSOR LIST =
310     A TRANSFER BAD DATA L(0)
311     A TRANSFER GOOD DATA L(1)
312     A TRANSFER BAD DATA L(1)
313     A COUNTER RESETTER 2
314     REQUIRED MESSAGES =
315     DATA*
316     REQUIRED SEMAPHORE STATUS =
317     CHAIN IF ; A LOCNTR
318     IS ; > 2
319     CHAIN IF ; A L1BAD
320     IS ; < 3
321     CHAIN IF ; A L1GOOD
322     IS ; < 3
323     STATISTICAL SUCCESSORS =
324     CHOOSE AS SUCCESSOR ; 90.00 % A TRANSFER GOOD DATA L(1)
325     CHOOSE AS SUCCESSOR ; 10.00 % A TRANSFER BAD DATA L(1)
326     INSTRUCTION LIST =
327     EXECUTE A TOTAL OF ; 1 A RESET MESSAGE TRANSFERRED
328     NAME = A TRANSFER BAD DATA L(1)

```

```

329     PRIORITY =      0
330     INTERRUPTABILITY FLAG = NO
331     CONCURRENT EXECUTION = NO
332     STATISTICAL SUCCESSOR STREAM =  1
333     ANDED PREDECESSOR LIST =
334         PROCESSOR 1
335     REQUIRED SEMAPHORE STATUS =
336         WAIT FOR ; A MESSAGE TRANSFERRED
337         TO BE ; RESET
338     INSTRUCTION LIST =
339         EXECUTE A TOTAL OF ; 1 SEND BAD DATA FRAME L(1)
340         EXECUTE A TOTAL OF ; 1 A SET MESSAGE TRANSFERRED
341         EXECUTE A TOTAL OF ; 1 A RESET L1GOOD
342         EXECUTE A TOTAL OF ; 1 A SET L1BAD
343     ANDED SUCCESSORS =
344         CHAIN TO ; PROCESSOR 1
345         WITH ITERATIONS THEN CHAIN COUNT OF ;      0
346         CHAIN TO ; PROCESSOR 2
347         WITH ITERATIONS THEN CHAIN COUNT OF ;      0
348     NAME = A COUNTER RESETTER 1
349     PRIORITY =      5
350     INTERRUPTABILITY FLAG = NO
351     CONCURRENT EXECUTION = NO
352     ORED PREDECESSOR LIST =
353         A TRANSFER GOOD DATA L(1)
354     REQUIRED SEMAPHORE STATUS =
355         CHAIN IF ; A LOCNTR
356         IS ; > 2
357         CHAIN IF ; A L1GOOD
358         IS ; > 2
359     INSTRUCTION LIST =
360         EXECUTE A TOTAL OF ; 3 A RESET LOCNTR
361         EXECUTE A TOTAL OF ; 3 A RESET L1GOOD
362     ANDED SUCCESSORS =

```

```

363     CHAIN TO ; PROCESSOR 0
364     WITH ITERATIONS THEN CHAIN COUNT OF ;      0
365     NAME = A COUNTER RESETTER 2
366     PRIORITY =      5
367     INTERRUPTABILITY FLAG = NO
368     CONCURRENT EXECUTION = NO
369     ORED PREDECESSOR LIST =
370         A TRANSFER GOOD DATA L(2)
371     REQUIRED SEMAPHORE STATUS =
372     CHAIN IF ; A L1BAD
373     IS ; > 2
374     CHAIN IF ; A L2GOOD
375     IS ; > 1
376     INSTRUCTION LIST =
377         EXECUTE A TOTAL OF ; 3 A RESET L1BAD
378         EXECUTE A TOTAL OF ; 2 A RESET L2GOOD
379     ANDED SUCCESSORS =
380     CHAIN TO ; PROCESSOR 1
381     WITH ITERATIONS THEN CHAIN COUNT OF ;      0
382     NAME = B TRANSFER GOOD DATA L(0)
383     PRIORITY =      0
384     INTERRUPTABILITY FLAG = NO
385     CONCURRENT EXECUTION = NO
386     STATISTICAL SUCCESSOR STREAM =    7
387     ANDED PREDECESSOR LIST =
388         B PROCESSOR 0
389     REQUIRED SEMAPHORE STATUS =
390     WAIT FOR ; MESSAGE TRANSFERRED
391     TO BE ; RESET
392     INSTRUCTION LIST =
393         EXECUTE A TOTAL OF ; 1 SEND GOOD DATA FRAME L(0)
394         EXECUTE A TOTAL OF ; 1 SET MESSAGE TRANSFERRED
395         EXECUTE A TOTAL OF ; 1 RESET LOCNTR
396     ANDED SUCCESSORS =

```

```

397     CHAIN TO ; B PROCESSOR 0
398     WITH ITERATIONS THEN CHAIN COUNT OF ;      0
399     NAME = B PROCESSOR 0
400     PRIORITY =      1
401     INTERRUPTABILITY FLAG = NO
402     CONCURRENT EXECUTION = NO
403     ORED PREDECESSOR LIST =
404         B TRANSFER GOOD DATA L(0)
405         B TRANSFER BAD DATA L(0)
406         B START
407         B COUNTER RESETTER
408     REQUIRED MESSAGES =
409         DATA*
410     REQUIRED SEMAPHORE STATUS =
411         CHAIN IF ; LOCNTR
412         IS ; < 3
413     STATISTICAL SUCCESSORS =
414         CHOOSE AS SUCCESSOR ; 82.00 % B TRANSFER GOOD DATA L(0)
415         CHOOSE AS SUCCESSOR ; 18.00 % B TRANSFER BAD DATA L(0)
416     INSTRUCTION LIST =
417         EXECUTE A TOTAL OF ; 1 RESET MESSAGE TRANSFERRED
418     NAME = B TRANSFER BAD DATA L(0)
419     PRIORITY =      0
420     INTERRUPTABILITY FLAG = NO
421     CONCURRENT EXECUTION = NO
422     STATISTICAL SUCCESSOR STREAM =      5
423     ANDED PREDECESSOR LIST =
424         B PROCESSOR 0
425     REQUIRED SEMAPHORE STATUS =
426         WAIT FOR ; MESSAGE TRANSFERRED
427         TO BE ; RESET
428     INSTRUCTION LIST =
429         EXECUTE A TOTAL OF ; 1 SEND BAD DATA FRAME L(0)
430         EXECUTE A TOTAL OF ; 1 SET MESSAGE TRANSFERRED

```

```

431     EXECUTE A TOTAL OF ; 1 SET LOCNTR
432     ANDED SUCCESSORS =
433     CHAIN TO ; B PROCESSOR 0
434     WITH ITERATIONS THEN CHAIN COUNT OF ;      0
435     CHAIN TO ; B PROCESSOR 1
436     WITH ITERATIONS THEN CHAIN COUNT OF ;      0
437     NAME = B START
438     PRIORITY =      9
439     INTERRUPTABILITY FLAG = NO
440     CONCURRENT EXECUTION = NO
441     START TIME = 0.0
442     ALLOWED PROCESSORS =
443     STATION B
444     INSTRUCTION LIST =
445     EXECUTE A TOTAL OF ; 1 START DATA TRANSFER
446     ANDED SUCCESSORS =
447     CHAIN TO ; B PROCESSOR 0
448     WITH ITERATIONS THEN CHAIN COUNT OF ;      0
449     NAME = B TRANSFER GOOD DATA L(1)
450     PRIORITY =      0
451     INTERRUPTABILITY FLAG = NO
452     CONCURRENT EXECUTION = NO
453     STATISTICAL SUCCESSOR STREAM =    3
454     ANDED PREDECESSOR LIST =
455     B PROCESSOR 1
456     REQUIRED SEMAPHORE STATUS =
457     WAIT FOR ; MESSAGE TRANSFERRED
458     TO BE ; RESET
459     INSTRUCTION LIST =
460     EXECUTE A TOTAL OF ; 1 SEND GOOD DATA FRAME L(1)
461     EXECUTE A TOTAL OF ; 1 SET MESSAGE TRANSFERRED
462     EXECUTE A TOTAL OF ; 1 SET L1GOOD
463     ANDED SUCCESSORS =
464     CHAIN TO ; B PROCESSOR 1

```

```

465     WITH ITERATIONS THEN CHAIN COUNT OF ;      0
466     CHAIN TO ; B COUNTER RESETTER
467     WITH ITERATIONS THEN CHAIN COUNT OF ;      0
468 NAME = B PROCESSOR 1
469     PRIORITY =      1
470     INTERRUPTABILITY FLAG = NO
471     CONCURRENT EXECUTION = NO
472     ORED PREDECESSOR LIST =
473         B TRANSFER BAD DATA L(0)
474         B TRANSFER GOOD DATA L(1)
475         B TRANSFER BAD DATA L(1)
476     REQUIRED MESSAGES =
477         DATA*
478     REQUIRED SEMAPHORE STATUS =
479         CHAIN IF ; LOCNTR
480         IS ; > 2
481         CHAIN IF ; L1GOOD
482         IS ; < 3
483     STATISTICAL SUCCESSORS =
484         CHOOSE AS SUCCESSOR ; 90.00 % B TRANSFER GOOD DATA L(1)
485         CHOOSE AS SUCCESSOR ; 10.00 % B TRANSFER BAD DATA L(1)
486     INSTRUCTION LIST =
487         EXECUTE A TOTAL OF ; 1 RESET MESSAGE TRANSFERRED
488 NAME = B TRANSFER BAD DATA L(1)
489     PRIORITY =      0
490     INTERRUPTABILITY FLAG = NO
491     CONCURRENT EXECUTION = NO
492     STATISTICAL SUCCESSOR STREAM = 9
493     ANDED PREDECESSOR LIST =
494         B PROCESSOR 1
495     REQUIRED SEMAPHORE STATUS =
496         WAIT FOR ; MESSAGE TRANSFERRED
497         TO BE ; RESET
498     INSTRUCTION LIST =

```

```

499      EXECUTE A TOTAL OF ; 1 SEND BAD DATA FRAME L(1)
500      EXECUTE A TOTAL OF ; 1 SET MESSAGE TRANSFERRED
501      EXECUTE A TOTAL OF ; 1 RESET L1GOOD
502      ANDED SUCCESSORS =
503      CHAIN TO ; B PROCESSOR 1
504      WITH ITERATIONS THEN CHAIN COUNT OF ;      0
505      NAME = B COUNTER RESETTER
506      PRIORITY =      5
507      INTERRUPTABILITY FLAG = NO
508      CONCURRENT EXECUTION = NO
509      ORED PREDECESSOR LIST =
510      B TRANSFER GOOD DATA L(1)
511      REQUIRED SEMAPHORE STATUS =
512      CHAIN IF ; LOCNTR
513      IS ; > 2
514      CHAIN IF ; L1GOOD
515      IS ; > 2
516      INSTRUCTION LIST =
517      EXECUTE A TOTAL OF ; 3 RESET LOCNTR
518      EXECUTE A TOTAL OF ; 3 RESET L1GOOD
519      ANDED SUCCESSORS =
520      CHAIN TO ; B PROCESSOR 0
521      WITH ITERATIONS THEN CHAIN COUNT OF ;      0
522      NAME = A TRANSFER GOOD DATA L(2)
523      PRIORITY =      0
524      INTERRUPTABILITY FLAG = NO
525      CONCURRENT EXECUTION = NO
526      STATISTICAL SUCCESSOR STREAM =      9
527      ANDED PREDECESSOR LIST =
528      PROCESSOR 2
529      REQUIRED SEMAPHORE STATUS =
530      WAIT FOR ; A MESSAGE TRANSFERRED
531      TO BE ; RESET
532      INSTRUCTION LIST =

```

```

533     EXECUTE A TOTAL OF ; 1 SEND GOOD DATA FRAME L(2)
534     EXECUTE A TOTAL OF ; 1 A SET MESSAGE TRANSFERRED
535     EXECUTE A TOTAL OF ; 1 A SET L2GOOD
536     EXECUTE A TOTAL OF ; 1 A RESET L2BAD
537     ANDED SUCCESSORS =
538     CHAIN TO ; A COUNTER RESETTER 2
539     WITH ITERATIONS THEN CHAIN COUNT OF ;      0
540     CHAIN TO ; PROCESSOR 2
541     WITH ITERATIONS THEN CHAIN COUNT OF ;      0
542     NAME = A TRANSFER BAD DATA L(2)
543     PRIORITY =      0
544     INTERRUPTABILITY FLAG = NO
545     CONCURRENT EXECUTION = NO
546     STATISTICAL SUCCESSOR STREAM =      8
547     ANDED PREDECESSOR LIST =
548     PROCESSOR 2
549     REQUIRED SEMAPHORE STATUS =
550     WAIT FOR ; A MESSAGE TRANSFERRED
551     TO BE ; RESET
552     INSTRUCTION LIST =
553     EXECUTE A TOTAL OF ; 1 SEND BAD DATA FRAME L(2)
554     EXECUTE A TOTAL OF ; 1 A SET MESSAGE TRANSFERRED
555     EXECUTE A TOTAL OF ; 1 A RESET L2GOOD
556     EXECUTE A TOTAL OF ; 1 A SET L2BAD
557     ANDED SUCCESSORS =
558     CHAIN TO ; PROCESSOR 2
559     WITH ITERATIONS THEN CHAIN COUNT OF ;      0
560     CHAIN TO ; PROCESSOR 3
561     WITH ITERATIONS THEN CHAIN COUNT OF ;      0
562     NAME = A COUNTER RESETTER 3
563     PRIORITY =      9
564     INTERRUPTABILITY FLAG = NO
565     CONCURRENT EXECUTION = NO
566     ORED PREDECESSOR LIST =

```

```

567      A TRANSFER GOOD DATA L(3)
568  REQUIRED SEMAPHORE STATUS =
569      CHAIN IF ; A L2BAD
570      IS ; > 1
571      CHAIN IF ; A L3GOOD
572      IS ; > 0
573  INSTRUCTION LIST =
574      EXECUTE A TOTAL OF ; 2 A RESET L2BAD
575      EXECUTE A TOTAL OF ; 1 A RESET L3GOOD
576  ANDED SUCCESSORS =
577      CHAIN TO ; PROCESSOR 2
578      WITH ITERATIONS THEN CHAIN COUNT OF ;      0
579  NAME = PROCESSOR 2
580      PRIORITY =      1
581  INTERRUPTABILITY FLAG = NO
582  CONCURRENT EXECUTION = NO
583  STATISTICAL SUCCESSOR STREAM = 12
584  ORED PREDECESSOR LIST =
585      A TRANSFER BAD DATA L(1)
586      A TRANSFER GOOD DATA L(2)
587      A TRANSFER BAD DATA L(2)
588      A COUNTER RESETTER 3
589  REQUIRED MESSAGES =
590      DATA*
591  REQUIRED SEMAPHORE STATUS =
592      CHAIN IF ; A L1BAD
593      IS ; > 2
594      CHAIN IF ; A L2GOOD
595      IS ; < 2
596      CHAIN IF ; A L2BAD
597      IS ; < 2
598  STATISTICAL SUCCESSORS =
599      CHOOSE AS SUCCESSOR ; 95.00 % A TRANSFER GOOD DATA L(2)
600      CHOOSE AS SUCCESSOR ; 5.00 % A TRANSFER BAD DATA L(2)

```

```

601     INSTRUCTION LIST =
602         EXECUTE A TOTAL OF ; 1 A RESET MESSAGE TRANSFERRED
603     NAME = PROCESSOR 3
604         PRIORITY =      1
605         INTERRUPTABILITY FLAG = NO
606         CONCURRENT EXECUTION = NO
607         STATISTICAL SUCCESSOR STREAM = 45
608     ORED PREDECESSOR LIST =
609         A TRANSFER BAD DATA L(2)
610         A TRANSFER GOOD DATA L(3)
611         A TRANSFER BAD DATA L(3)
612     REQUIRED MESSAGES =
613         DATA*
614     REQUIRED SEMAPHORE STATUS =
615         CHAIN IF ; A L2BAD
616         IS ; > 1
617         CHAIN IF ; A L3GOOD
618         IS ; < 1
619     STATISTICAL SUCCESSORS =
620         CHOOSE AS SUCCESSOR ; 97.50 % A TRANSFER GOOD DATA L(3)
621         CHOOSE AS SUCCESSOR ; 2.50 % A TRANSFER BAD DATA L(3)
622     INSTRUCTION LIST =
623         EXECUTE A TOTAL OF ; 1 A RESET MESSAGE TRANSFERRED
624     NAME = A TRANSFER GOOD DATA L(3)
625         PRIORITY =      0
626         INTERRUPTABILITY FLAG = NO
627         CONCURRENT EXECUTION = NO
628     ANDED PREDECESSOR LIST =
629         PROCESSOR 3
630     REQUIRED SEMAPHORE STATUS =
631         WAIT FOR ; A MESSAGE TRANSFERRED
632         TO BE ; RESET
633     INSTRUCTION LIST =
634         EXECUTE A TOTAL OF ; 1 SEND GOOD DATA FRAME L(3)

```

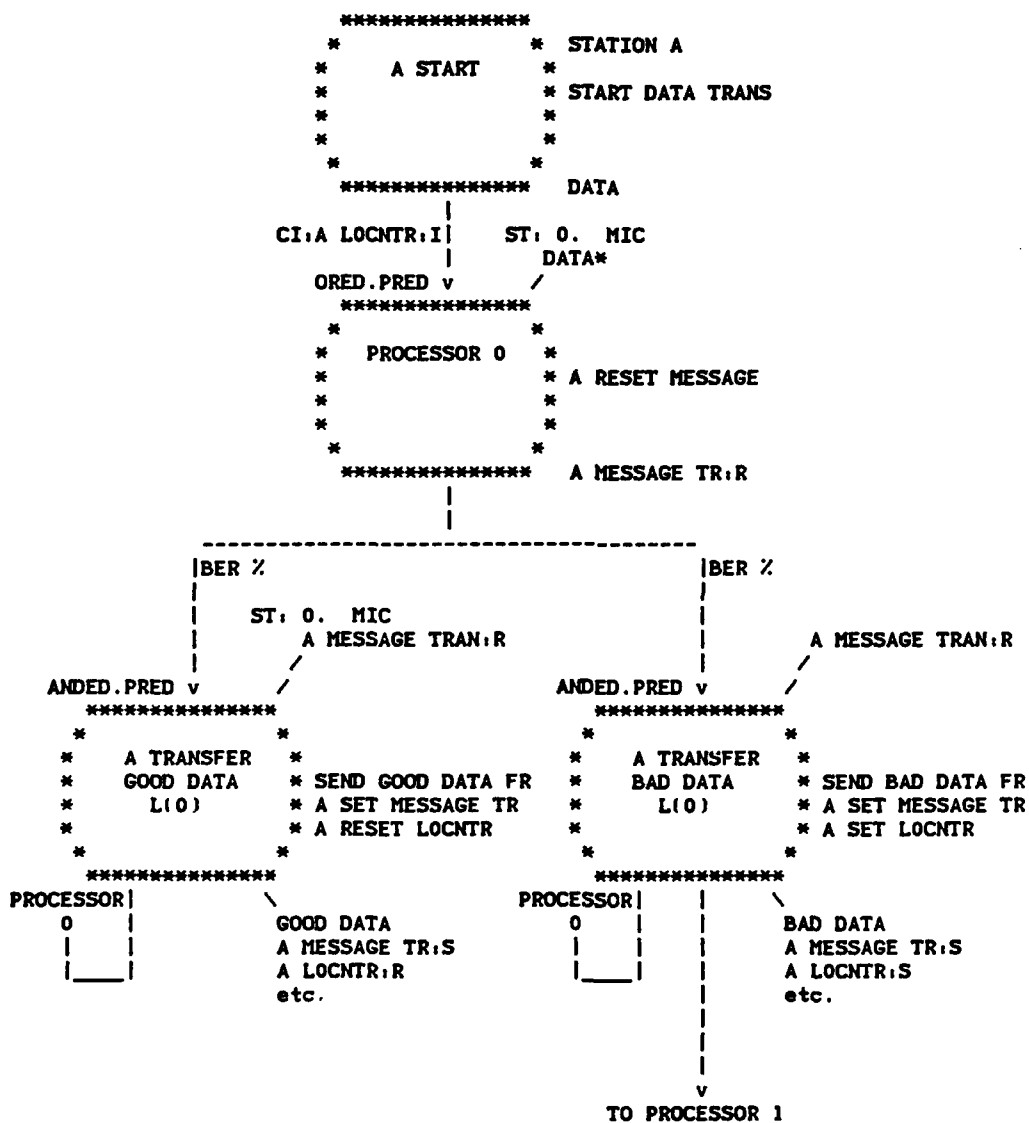
635 EXECUTE A TOTAL OF ; 1 A SET MESSAGE TRANSFERRED
 636 EXECUTE A TOTAL OF ; 1 A SET L3GOOD
 637 ANDED SUCCESSORS =
 638 CHAIN TO ; A COUNTER RESETTER 3
 639 WITH ITERATIONS THEN CHAIN COUNT OF ; 0
 640 CHAIN TO ; PROCESSOR 3
 641 WITH ITERATIONS THEN CHAIN COUNT OF ; 0
 642 NAME = A TRANSFER BAD DATA L(3)
 643 PRIORITY = 0
 644 INTERRUPTABILITY FLAG = NO
 645 CONCURRENT EXECUTION = NO
 646 ANDED PREDECESSOR LIST =
 647 PROCESSOR 3
 648 REQUIRED SEMAPHORE STATUS =
 649 WAIT FOR ; A MESSAGE TRANSFERRED
 650 TO BE ; RESET
 651 INSTRUCTION LIST =
 652 EXECUTE A TOTAL OF ; 1 SEND BAD DATA FRAME L(3)
 653 EXECUTE A TOTAL OF ; 1 A SET MESSAGE TRANSFERRED
 654 EXECUTE A TOTAL OF ; 1 A RESET L3GOOD
 655 ANDED SUCCESSORS =
 656 CHAIN TO ; PROCESSOR 3
 657 WITH ITERATIONS THEN CHAIN COUNT OF ; 0

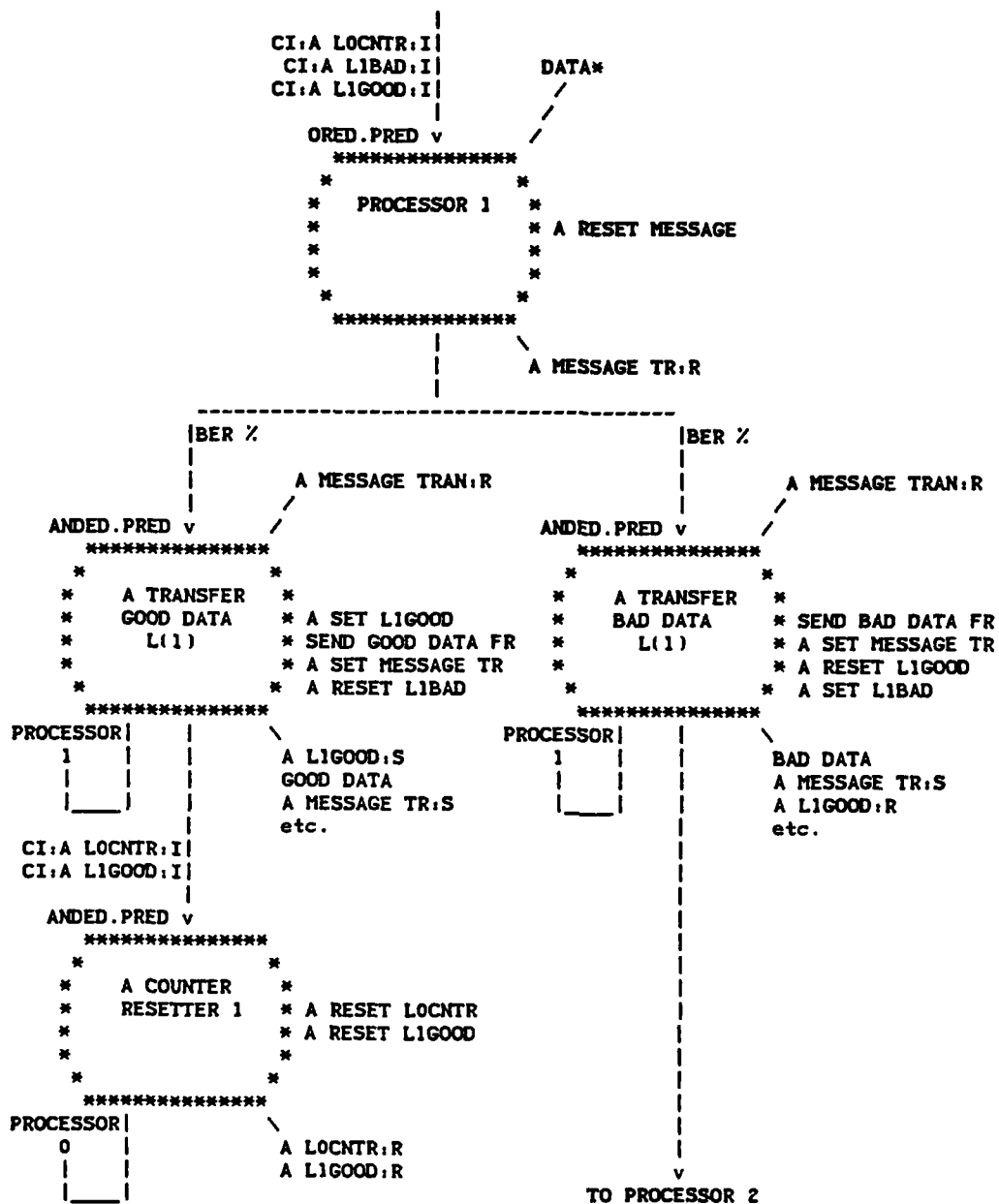
APPENDIX B. SIMULATION FLOW DIAGRAMS

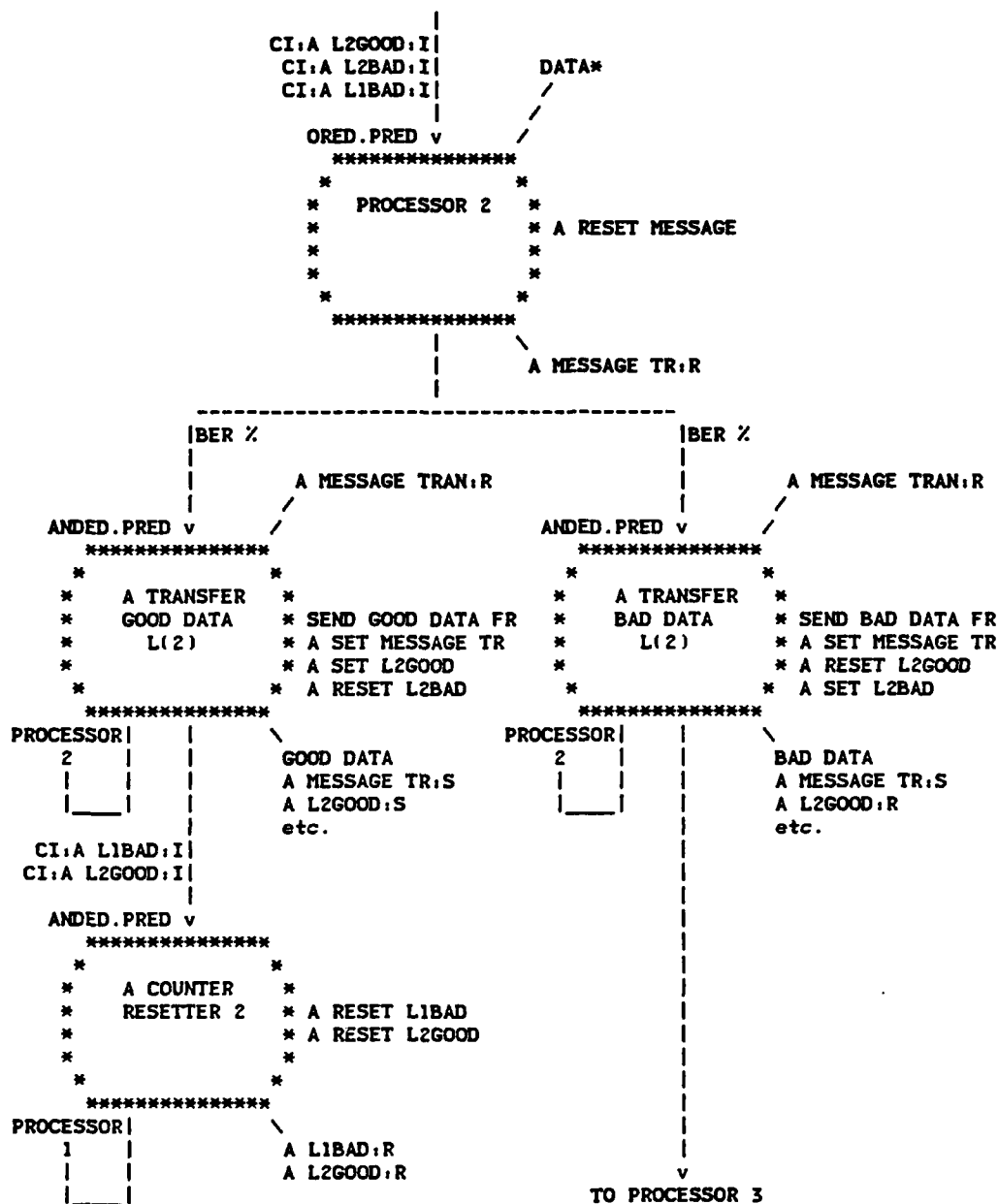
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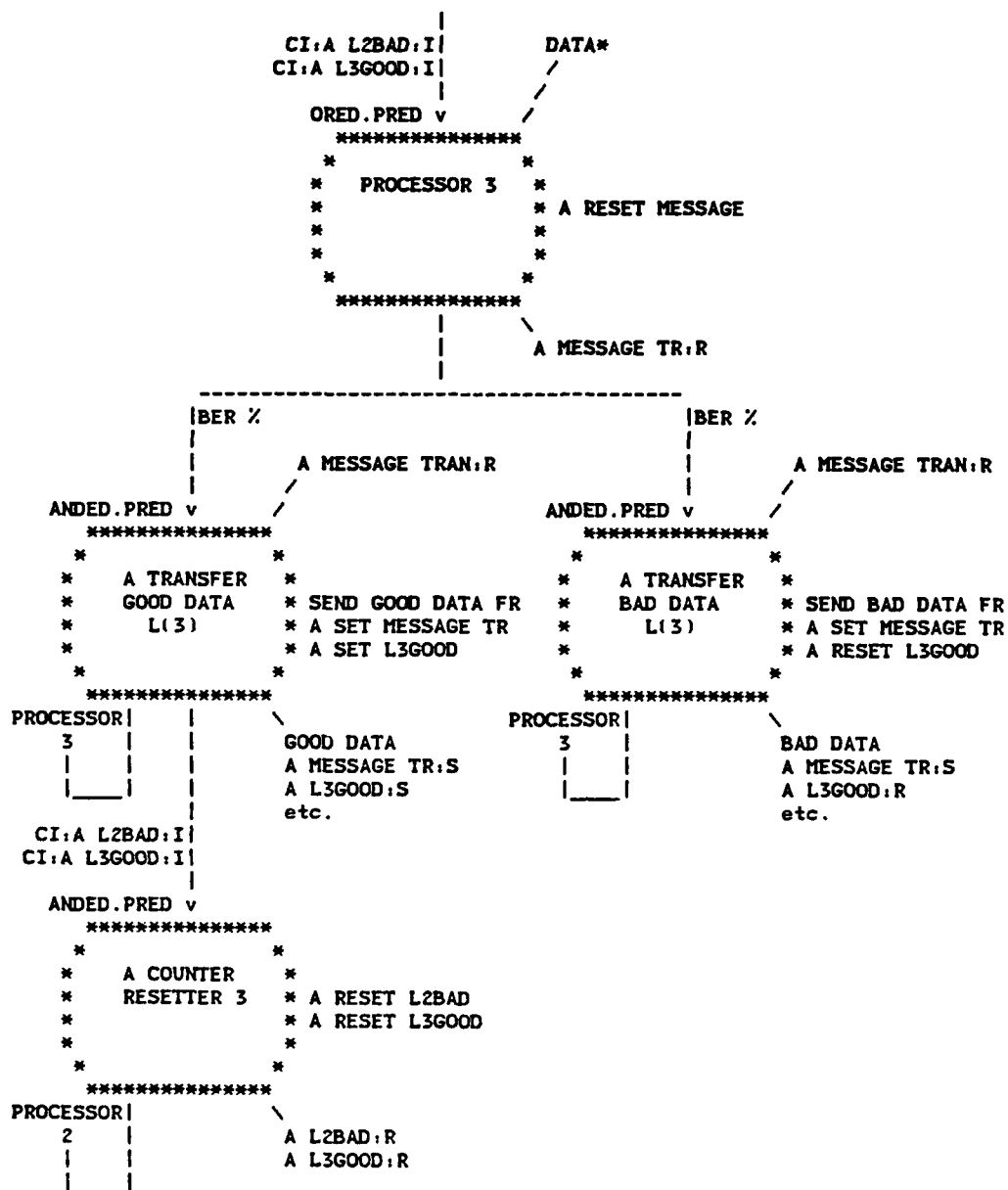
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STATION A - MULTI LEVEL ADAPTIVE ARQ STRATEGY

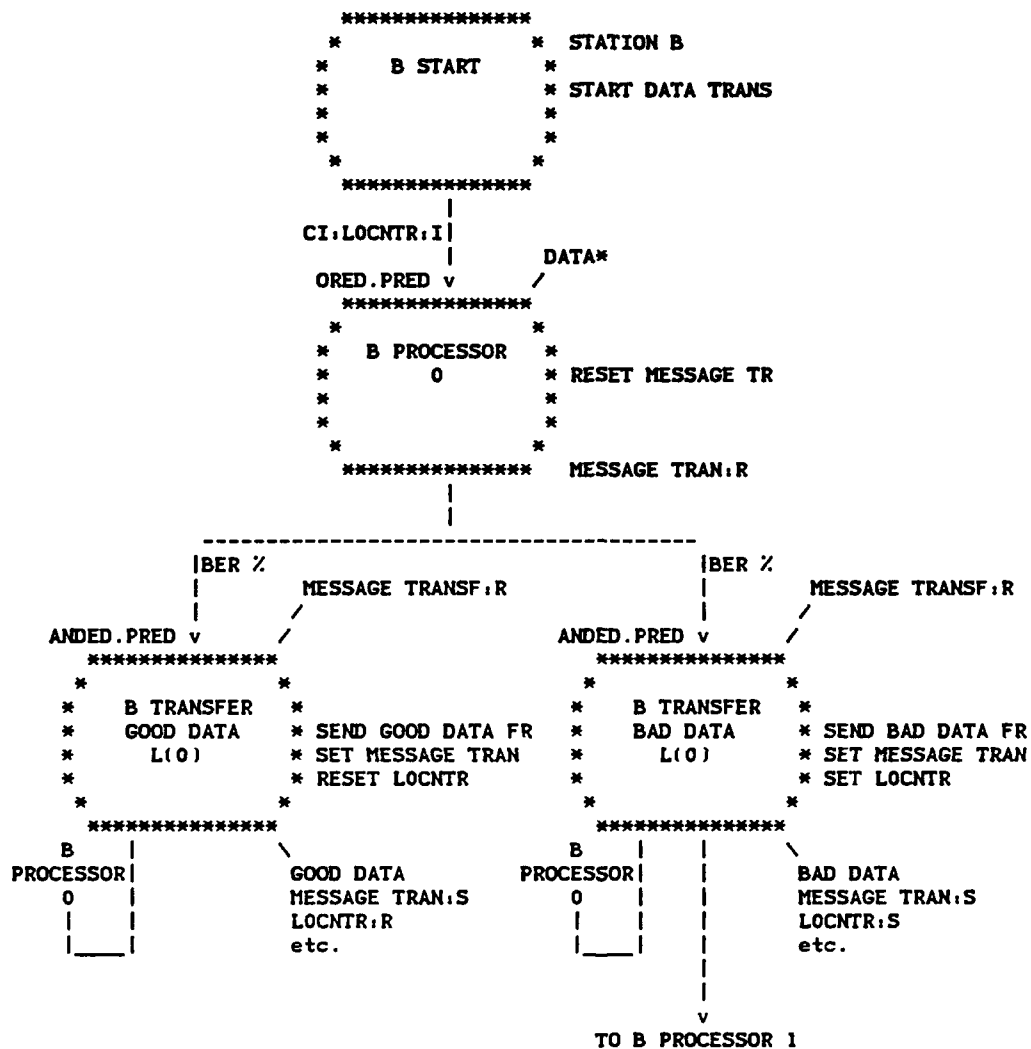


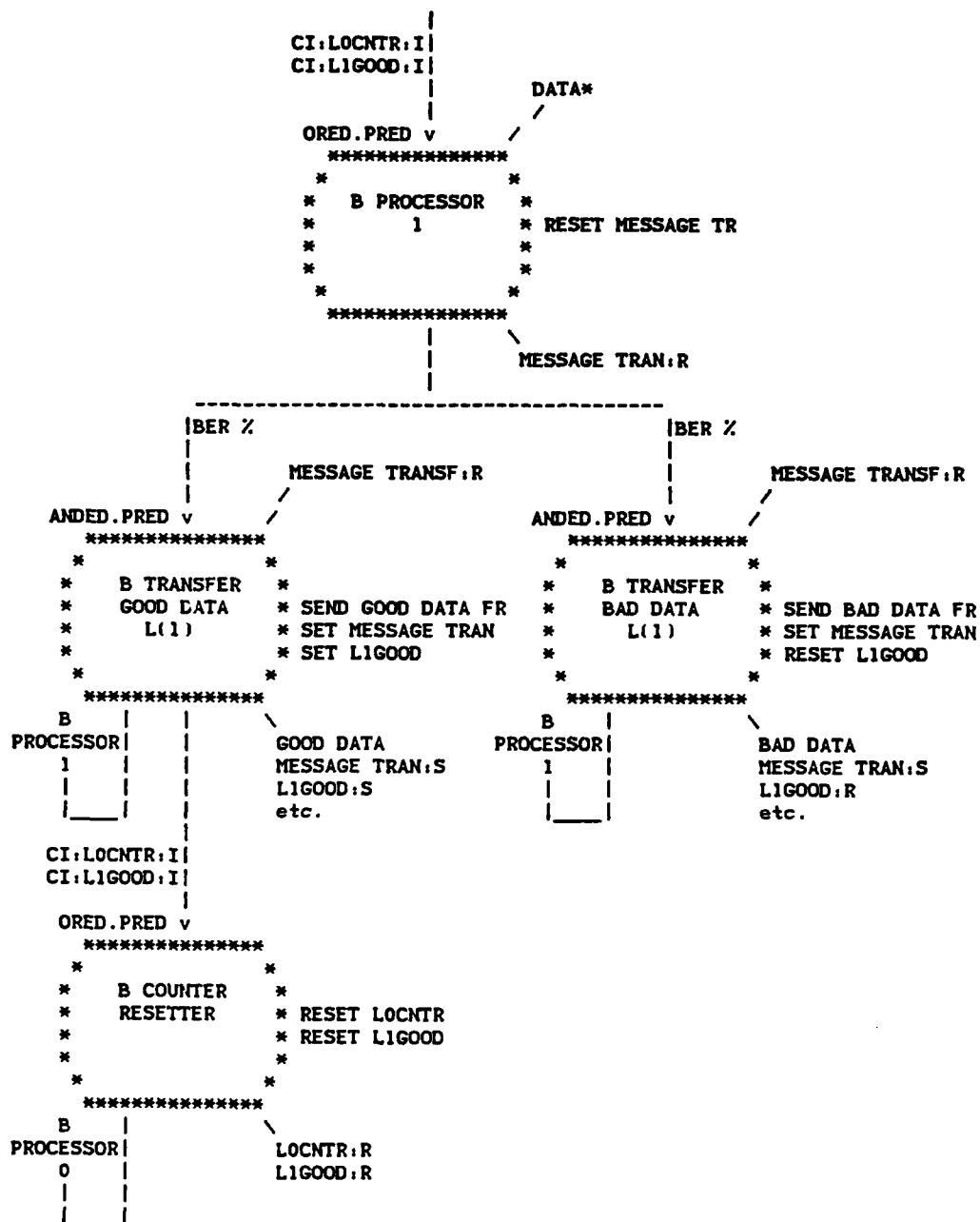






STATION B - TWO LEVEL ADAPTIVE ARQ STRATEGY





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