#### **Research Review** 2017

## Automated Assurance of Security Policy Enforcement (AASPE)

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#### **Executive Summary**

- 1. Relevance for the DoD warfighter Safety-critical systems are no longer closed but connected, thus exposed to security threats
- 2. Relevance to state-of-the-art in software engineering or cybersecurity Avionics industry embraces model-based architecture-centric virtual integration for safety
- 3. Expected DoD practice improvements

Continuous assurance through automated architecture security analysis to complement code level security analysis

- Analysis of security policy specifications for vulnerabilities
- Analysis of security enforcement runtime architecture for vulnerabilities
- Generation of runtime system configurations from verified security models

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# Security Challenges as Safety-Critical Systems Become Connected







Bad Example: Automotive



Lack of physical and logical isolation within system leads to costly rework and recalls

*More than secure code and external firewalls!* Security policy in form of acceptable command and information flows and isolation requirements

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# A Model-Based Analysis and Generation Approach

#### **Modeling Tool: AADL**

Architecture Analysis & Design Language Widely used standard (SAE International)

- Designed for embedded software systems
- Includes support for timing, performance, safety analysis

#### **Extensions:**

#### Model security policy and enforcement

- Security levels/domains
- Trust/verification
- Encryption
- Authentication
- Physical and logical exposure
- Concurrency



## **Security Policies and Their Enforcement**





Mission System Security Policy

#### Security policy vulnerabilities: Analyze Information Flows Examples: Verify secrets stay secret, and Sensors can't send commands



Security enforcement vulnerabilities: Analyze Deployment Mechanisms Example: Hi and low-security channels shouldn't coexist on unpartitioned hardware



#### **Research Connection:**

Apply *Multiple Independent Levels of Security* (MILS) framework (confidentiality) to system security (integrity)

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# **Security Analysis Techniques and Tools**

- o. Consistency in security policy specification and enforcement
- 1. Model-Based Attack Impact Analysis (AIA) tool
- 2. Model-Based Attack Tree Analysis (ATA) tool
- 3. Generation of security configuration files
  - Model-based auto-configuration of certified kernel (seL4/CAmkES) security policy



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## **Using Security Assurance Techniques and Tools**

- 1. Specify security policy as verifiable requirements
- 2. Formalize verification activities
- 3. Automate execution of verification plans



MILS-R0: Components sharing a bus should have the same security level.
 MILS-R1: Inter-communicating components should have the same security level.
 MILS-R2: Processes with different security levels use isolated memory regions.
 MILS-R3: Components associated with identical processing resources share the same security level.
 MILS-R4: Threads inside the same process share the same security levels.

CWE-131 Incorrect calculation of buffer size.

CWE-311 Missing encryption of sensitive data.

CWE-805 Buffer Access with Incorrect Length Value.

# Extension to Architecture-Led Incremental System Assurance (ALISA) workbench

- System case JeepSecurityCase: (S94 F9 T0 E0 tbd0 EL0 TS0)
- Image: Model JeepSecurityCaseJeepSecurityPlan(integration.attack)
  - V Claim MILS\_R5(integration.attack): MILS\_R5: All non-verified
  - Claim CWE131(integration.attack): CWE131: Incorrect calc
    Evidence vaCWE131a (203 ms): check connections for c
    - ▶ ✓ Evidence vaCWE131b (252 ms): Check that timing requ
  - V Claim CWE311(integration.attack): CWE311: Missing Encry
  - V Claim CWE805(integration.attack): CWE805: Buffer Access
  - Subsystem cellular: (S4 F2 T0 E0 tbd0 EL0 TS0)
    - III Claim MILS\_R0(cellular): MILS\_R0: Components sharing
    - II Claim MILS\_R1(cellular): R1: Components with different
    - V Claim MILS\_R5(cellular): MILS\_R5: All non-verified com
    - V Claim CWE311(cellular): CWE311: Missing Encryption o
    - V Claim CWE805(cellular): CWE805: Buffer Access with In
    - V Claim MILS\_R6(cellular): R6: All communication that an
    - Subsystem internet: (S5 F1 T0 E0 tbd0 EL0 TS0)
    - V Claim MILS\_R0(internet): MILS\_R0: Components sharing
    - El Claim MILS\_R1(internet): R1: Components with differen
    - V Claim MILS\_R5(internet): MILS\_R5: All non-verified com
    - V Claim CWE311(internet): CWE311: Missing Encryption c
    - V Claim CWE805(internet): CWE805: Buffer Access with In
    - V Claim MILS\_R6(internet): R6: All communication that ar
  - Subsystem router\_cel: (S3 F0 T0 E0 tbd0 EL0 TS0)
  - Subsystem car: (S62 F6 T0 E0 tbd0 EL0 TS0)
  - Subsystem attacker\_cel: (S5 F0 T0 E0 tbd0 EL0 TS0)
  - Subsystem attacker\_wifi: (S5 F0 T0 E0 tbd0 EL0 TS0)
  - Subsystem attacker\_internet: (S5 F0 T0 E0 tbd0 EL0 TS0)

### **AASPE** Results

Code and Examples on GitHub (https://github.com/cmu-sei/AASPE)

Tools

- Security policy and enforcement verification on AADL models
- Graphical attack impact and attack tree analysis tools
- Generation of attack impact graphs and attack trees from AADL models with security annotations
- Generation of seL4 configuration files from AADL based specifications
- Example models
  - Automotive: Jeep, Prius
  - Aircraft model
  - Drone case study

Proposal for AADL Security Annex standard

Papers/report on security analysis, security assurance workflow

Proposal for an integrated safety and security engineering approach

## **Summary and Future Work**

Where We Started: DARPA High-Assurance Cyber Military Systems (HACMS) program successfully demonstrated AADL-based verification and generation for reducing vulnerabilities in unmanned drones.

What We Did: We demonstrated the feasibility of improving security assurance through architecture modeling and analysis of vulnerabilities in security policy specification and enforcement.

What's Next for the Community: Use of SAE International AADL standard offers transition path through the Open Source AADL Tool Environment (OSATE). We will advance the proposed Security Annex to AADL towards approval.

What's Next for the SEI: Develop an integrated approach to safety and security engineering approach to mission critical systems funded as new three-year SEI line project.

## **Contact Information**

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