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## Cubic Spline Interpolation on Nested Polygon Triangulations

### Oleg Davydov, Günther Nürnberger, and Frank Zeilfelder

Abstract. We develop an algorithm for constructing Lagrange and Hermite interpolation sets for spaces of cubic  $C^1$ -splines on general classes of triangulations built up of nested polygons whose vertices are connected by line segments. Additional assumptions on the triangulation are significantly reduced compared to the special class given in [4]. Simultaneously, we have to determine the dimension of these spaces, which is not known in general. We also discuss the numerical aspects of the method.

#### §1. Introduction

In contrast to univariate splines, it is a non-trivial problem to construct even one single set of interpolation points for bivariate spline spaces. Such interpolation sets for  $S_q^r(\Delta)$ , the space of splines of degree q and smoothness r, were constructed for crosscut-partitions  $\Delta$  (see the survey [9] and the references therein). For general triangulations  $\Delta$ , interpolation sets were constructed for  $S_q^1(\Delta)$ ,  $q \geq 4$  in [3].

The case q = 3 is much more complicated given that not even the dimension of  $S_3^1(\Delta)$  is known for arbitrary triangulations  $\Delta$ . It is an open question whether the dimension of  $S_3^1(\Delta)$  is equal to Schumaker's lower bound [12]. The aim of this paper is to investigate interpolation by  $S_3^1(\Delta)$  for general classes of triangulations  $\Delta$  consisting of nested polygons whose vertices are connected by line segments. Following a general principle of locally choosing interpolation points for  $S_3^1(\Delta)$  by passing from triangle to triangle, we describe an inductive method for constructing point sets that admit unique Lagrange (respectively Hermite) interpolation by  $S_3^1(\Delta)$  under certain assumptions on  $\Delta$ . Moreover, we prove that the dimension of these spaces is equal to Schumaker's lower bound.

In this way we obtain a class of triangulations  $\triangle$  which is significantly larger than the special class described in [4]. Moreover, the methods of proof in this paper are different from those in [4]. It is important to note that

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triangulations of this type can be constructed starting from any given points in the plane, see [11].

The numerical examples (with up to 100,000 interpolation points) show that in order to obtain good approximations, it is desirable to subdivide some of the triangles. Our method of constructing interpolation points also works for these modified triangulations.

We note that our interpolation method can be used for the construction of smooth surfaces without involving any derivative data. For scattered data fitting, the needed Lagrange data are approximately computed by local methods. In contrast to the finite element methods for cubic splines, we do not need to subdivide all triangles by a Clough-Tocher split or use derivatives.

#### §2. Preliminaries

Let  $\triangle$  be a regular triangulation of a simply connected polygonal domain  $\Omega$  in  $\mathbb{R}^2$ . We denote by  $S_3^1(\triangle) = \{s \in C^1(\Omega) : s |_T \in \Pi_3, T \in \triangle\}$  the space of bivariate splines of degree 3 and smoothness 1 (with respect to  $\triangle$ ). Here  $\Pi_3 = \operatorname{span} \{x^{\nu}y^{\mu} : \nu, \mu \ge 0, \nu + \mu \le 3\}$  denotes the space of bivariate polynomials of total degree 3.

We investigate the following interpolation problem. Construct a set  $\{z_1, \ldots, z_N\}$  in  $\Omega$ , where  $N = \dim S_3^1 = (\Delta)$ , such that for each function  $f \in C(\Omega)$ , a unique spline  $s \in S_3^1(\Delta)$  exists such that  $s(z_i) = f(z_i)$ ,  $i = 1, \ldots, N$ . Such a set  $\{z_1, \ldots, z_N\}$  is called a Lagrange interpolation set for  $S_3^1(\Delta)$ . If also partial derivatives of f are involved, then we speak of a Hermite interpolation set for  $S_3^1(\Delta)$ .

In contrast to [4], we will use *Bernstein-Bézier techniques* [2,5]. Given a spline  $s \in S_3^1(\Delta)$ , we consider the following representation of the polynomial pieces  $p = s|_T \in \Pi_3$  on the triangle  $T \in \Delta$  with vertices  $v_1, v_2, v_3$ ,

$$p(x,y) = \sum_{\nu+\mu+\sigma=3} a_{\nu,\mu,\sigma}^{[T]} \frac{3!}{\nu!\mu!\sigma!} \Phi_1^{\nu}(x,y) \Phi_2^{\mu}(x,y) \Phi_3^{\sigma}(x,y), \ (x,y) \in T,$$
(1)

where  $\Phi_l \in \Pi_1$ , l = 1, 2, 3, is uniquely defined by  $\Phi_l(v_k) = \delta_{k,l}$ , k = 1, 2, 3. This representation of p is called the Bernstein-Bézier representation of p, the real numbers  $a_{\nu,\mu,\sigma}^{[T]}$  are called the Bernstein-Bézier coefficients of p, and  $\Phi_l(x, y)$ , l = 1, 2, 3, are the barycentric coordinates (w.r.t. T) of  $(x, y) \in T$ .

**Definition 1.** A set  $A \subset \{(\nu, \mu, \sigma, T) : \nu + \mu + \sigma = 3, T \in \Delta\}$  is called an admissible set for  $S_3^1(\Delta)$  if for every choice of coefficients  $a_{\nu,\mu,\sigma}^{[T]}, (\nu, \mu, \sigma, T) \in A$ , a unique spline  $s \in S_3^1(\Delta)$  exists with these coefficients in the above Bernstein-Bézier representation.

The above Bernstein-Bézier form can be used to express smoothness conditions of polynomial pieces on adjacent triangles  $T_1, T_2$  with vertices  $v_1, v_2, v_3$ , respectively  $v_1, v_2, v_4$  (cf. [2,5]).

**Theorem 2.** Let s be a piecewise cubic polynomial function defined on  $T_1 \cup T_2$ . Then  $s \in S_3^1(\{T_1, T_2\})$  iff  $a_{\nu,\mu,0}^{[T_2]} = a_{\nu,\mu,0}^{[T_1]}$ ,  $\nu + \mu = 3$ , and  $a_{\nu,\mu,1}^{[T_2]} = a_{\nu+1,\mu,0}^{[T_1]} \Phi_1(v_4) + a_{\nu,\mu+1,0}^{[T_1]} \Phi_2(v_4) + a_{\nu,\mu,1}^{[T_1]} \Phi_3(v_4)$ ,  $\nu + \mu = 2$ .

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#### Cubic Spline Interpolation

For later use, we also mention here the following relations between the Bernstein-Bézier coefficients of a cubic polynomial p in the representation (1) and its partial derivatives at  $v_1$  in direction of a unit vector parallel to the edge  $e = [v_1, v_2]$ , denoted by  $\frac{\partial}{\partial e}$ .

$$a_{3,0,0}^{[T]} = p(v_1), \quad a_{2,1,0}^{[T]} = p(v_1) + \frac{1}{3} \frac{\partial p(v_1)}{\partial e} ||v_1 - v_2||_2,$$

$$a_{1,2,0}^{[T]} = p(v_1) + \frac{2}{3} \frac{\partial p(v_1)}{\partial e} ||v_1 - v_2||_2 + \frac{1}{6} \frac{\partial^2 p(v_1)}{\partial e^2} ||v_1 - v_2||_2^2, \quad (2)$$

$$\frac{\partial p(v_1)}{\partial e} = \frac{3(a_{2,1,0}^{[T]} - a_{3,0,0}^{[T]})}{||v_1 - v_2||_2}, \quad \frac{\partial^2 p(v_1)}{\partial e^2} = \frac{6(a_{1,2,0}^{[T]} - 2a_{2,1,0}^{[T]} + a_{3,0,0}^{[T]})}{||v_1 - v_2||_2^2}.$$

#### §3. Main Results

In this section, we state our main results on  $S_3^1(\Delta)$ , where  $\Delta$  consists of nested polygons whose vertices are connected by line segments. We first define this class of triangulations. Then, we determine the dimension and construct interpolation sets for the corresponding spline space. Moreover, we show that this dimension is equal to Schumaker's lower bound [12]. Finally, we discuss a property of  $\Delta$  which is essential for the local construction of interpolation points.

First, we describe triangulations of nested polygons and decompose the domain into finitely many subsets needed in our construction of interpolation points.

**Triangulations of nested polygons.** We consider the following general type of triangulation  $\triangle$ . Let  $P_0, P_1, \ldots, P_k$  be a sequence of closed simple polygonal lines, and let  $\Omega_{\mu}$  be the closed (not necessarily convex) bounded polygon with boundary  $P_{\mu}$ . Suppose that the polygons  $\Omega_{\mu}$  are nested, i.e.,  $\Omega_{\mu-1} \subset \Omega_{\mu}$ ,  $\mu = 0, \ldots, k$ . The vertices of  $\Delta$  are the vertices of  $P_{\mu}, \mu = 0, \ldots, k$ , and one vertex inside  $P_0$ . The edges of  $\triangle$  are the edges of  $P_{\mu}$ ,  $\mu = 0, \ldots, k$ , and additional line segments connecting the vertices of  $P_{\mu}$  with the vertices of  $P_{\mu+1}, \mu = 0, \ldots, k-1$ . The resulting triangulation  $\triangle$  of  $\Omega := \Omega_k$  does not have vertices in the interior of  $\Omega_{\mu+1} \setminus \Omega_{\mu}$ ,  $\mu = 0, \ldots, k-1$ , and does not have edges connecting two vertices of  $P_{\mu}$  other than the edges of  $P_{\mu}$ , see Figure 1. **Decomposition of the domain.** We decompose the domain  $\Omega$  into finitely many sets  $V_0 \subset V_1 \subset \cdots \subset V_m = \Omega$ , where each set  $V_i$ , is the union of closed triangles of  $\triangle$ ,  $i = 0, \ldots, m$ . Let  $V_0$  be an arbitrary closed triangle of  $\triangle$  in  $\Omega_0$ . We define the sets  $V_1 \subset \cdots \subset V_m$  by induction. Assuming  $V_{i-1}$  is defined, we choose a vertex  $v_i$  of  $\triangle$  such that there exists at least one triangle of  $\triangle$ with vertex  $v_i$  and a common edge with  $V_{i-1}$ . Let  $T_{i,1}, \ldots, T_{i,n_i}, n_i \ge 1$ , be all such triangles. We set  $V_i = V_{i-1} \cup \overline{T}_{i,1} \cup \cdots \cup \overline{T}_{i,n_i}$ , and denote by  $\Delta_i = \{T \in \Delta : T \subset V_i\}$  the subtriangulation which corresponds to the set  $V_i$ .

The vertices  $v_i$ , i = 1, ..., m, are chosen as follows. After choosing  $V_0$  to be an arbitrary closed triangle of  $\Delta$  in  $\Omega_0$ , we pass through the vertices of  $P_0$  in clockwise order by applying the above rule. (It is clear that the choice of



Fig. 1. Triangulation of nested polygons.

these vertices is unique after fixing the first vertex.) Now, we assume that we have passed through the vertices of  $P_{\mu-1}$ . We fix a vertex  $w_{\mu}$  of  $P_{\mu}$  that is connected with at least two vertices of  $P_{\mu+1}$ . Then w.r.t. clockwise order, we choose the first vertex of  $P_{\mu}$  greater than  $w_{\mu}$  which is connected with at least two vertices of  $P_{\mu-1}$ . Then we pass through the vertices of  $P_{\mu}$  in clockwise order until  $w_{\mu}^-$ , and pass through the vertices of  $P_{\mu}$  in anticlockwise order until  $w_{\mu}^+$  by applying the above rule. (Here  $w_{\mu}^+$  denotes the vertex next to  $w_{\mu}$  in clockwise order and  $w_{\mu}^-$  denotes the vertex next to  $w_{\mu}$  in anticlockwise order.) Finally, we choose the vertex  $w_{\mu}$ . (It is clear that after fixing  $w_{\mu}$ , the choice of the vertices on  $P_{\mu}$  is unique.)

The construction of an admissible set for  $S_3^1(\Delta)$  and the choice of interpolation points depend on the following properties of the triangulation  $\Delta$ .

**Definition 3.** (1) An interior edge e with vertex v of the triangulation  $\triangle$  is called degenerate at v if the edges with vertex v adjacent to e lie on a line. (2) An interior vertex v of  $\triangle$  is called singular if v is a vertex of exactly four edges and these edges lie on two lines. (3) An interior vertex v of  $\triangle$  on the boundary of a given subtriangulation  $\triangle'$  of  $\triangle$  is called semi-singular of type 1 w.r.t.  $\triangle'$  if exactly one edge with endpoint v is not contained in  $\triangle'$  and this edge is degenerate at v. (4) An interior vertex v of  $\triangle$  on the boundary of a given subtriangulation  $\triangle'$  of  $\triangle$  is called semi-singular of type 2 w.r.t.  $\triangle'$  if exactly two edges with endpoint v are not contained in  $\triangle'$  and these edges are degenerate at v. (5) A vertex v of  $\triangle$  is called semi-singular w.r.t.  $\triangle'$  if v satisfies (3) or (4).

In the following, we construct an admissible set and interpolation sets for  $S_3^1(\Delta)$ , where  $\Delta$  is a nested-polygon triangulation.

Construction of an admissible set. First, we choose  $\mathcal{A}_0 = \{(\nu, \mu, \sigma, V_0) : \nu + \mu + \sigma = 3\}$  and then, proceeding by induction, we successively add admissible points on  $V_i \setminus V_{i-1}$ ,  $i = 1, \ldots, m$ . Assuming that an admissible set  $\mathcal{A}_{i-1}$  on  $V_{i-1}$  has been constructed, we choose admissible points on  $V_i \setminus V_{i-1}$  as follows. By the above decomposition of  $\Omega$ ,  $V_i \setminus V_{i-1}$  is the union of consecutive triangles  $T_{i,1}, \ldots, T_{i,n_i}$  with vertex  $v_i$  and common edges with  $V_{i-1}$ . We denote the consecutive endpoints of these edges by  $v_{i,0}, v_{i,1}, \ldots, v_{i,n_i}$ , and the piecewise polynomials in the representation (1) on  $T_{i,j}$  by  $p_{i,j} \in \Pi_3$ ,  $j = 1, \ldots n_i$ , where the vertices of  $T_{i,j}$  are ordered as follows:  $v_i, v_{i,j}, v_{i,j+1}$ . Furthermore, we denote by  $e_{i,j}$  the edges  $[v_{i,j}, v_i], j = 0, \ldots, n_i$ .

We need the following properties of the subtriangulation  $\Delta_i = \{T \in \Delta : T \subset V_i\}$  at the vertices  $v_{i,0}, \ldots, v_{i,n_i}$ :

- (a)  $e_{i,j}$  is non-degenerate at  $v_{i,j}$ ,
- (b)  $v_{i,j}$  is semi-singular w.r.t.  $\triangle_i$ . (This latter property is only relevant if  $v_{i,j}$  lies on the boundary of  $\triangle_i$ , *i.e.*, for  $j \in \{0, n_i\}$ .)

For  $j \in \{1, \ldots, n_i - 1\}$ , we set  $c_{i,j} = 1$  if (a) holds, and  $c_{i,j} = 0$  otherwise. For  $j \in \{0, n_i\}$ , we set  $c_{i,j} = 1$  if both (a) and (b) hold, and  $c_{i,j} = 0$  otherwise. Moreover, we set  $c_i = \sum_{j=0}^{n_i} c_{i,j}$ , and assume  $c_i \leq 3$ ,  $i = 1, \ldots, n$ .

Now, we construct the following admissible points on  $V_i \setminus V_{i-1}$ . If  $c_i = 3$ , then no point is chosen. If  $c_i = 2$ , then we choose  $(3, 0, 0, T_{i,1})$ . If  $c_i = 1$ , then we choose  $(3, 0, 0, T_{i,1})$  and  $(2, 0, 1, T_{i,j})$ , where  $e_{i,j}$  is an edge with  $c_{i,j} = 0$ . If  $c_i = 0$ , then we choose  $(3, 0, 0, T_{i,1})$ ,  $(2, 0, 1, T_{i,1})$  and  $(2, 1, 0, T_{i,1})$ . The admissible set  $\mathcal{A}_i$  on  $V_i$  is obtained by adding these points to  $\mathcal{A}_{i-1}$ .

**Construction of interpolation sets.** We choose interpolation points in  $V_0$  and then in  $V_i \setminus V_{i-1}$ ,  $i = 1, \ldots, m$ , successively. In the first step, we choose 10 different points in  $V_0$  (respectively 10 Hermite interpolation conditions) which admit unique Lagrange interpolation (respectively Hermite interpolation) by the space  $\Pi_3$ . For example, for Lagrange interpolation, we may choose four parallel line segments  $l_{\nu}$  in  $V_0$  and  $\nu$  different points on each  $l_{\nu}$ ,  $\nu = 1, 2, 3, 4$ . Assuming that the interpolation points in  $V_{i-1}$  have already been chosen, we proceed to  $V_i \setminus V_{i-1}$  as follows.

For Lagrange interpolation, we choose the following points in  $V_i \setminus V_{i-1}$ . If  $c_i = 3$ , then no point is chosen. If  $c_i = 2$ , then we choose  $v_i$ . If  $c_i = 1$ , then we choose  $v_i$  and one further point on some edge  $e_{i,j}$  with  $c_{i,j} = 0$ . If  $c_i = 0$ , then we choose  $v_i$  and two further points on two different edges.

For Hermite interpolation, we require the following interpolation conditions for  $s \in S_3^1(\Delta)$  at the vertex  $v_i$ . If  $c_i = 3$ , then no interpolation condition is required at  $v_i$ . If  $c_i = 2$ , then we require  $s(v_i) = f(v_i)$ . If  $c_i = 1$ , then we require  $s(v_i) = f(v_i)$  and  $\frac{\partial s}{\partial e_{i,j}}(v_i) = \frac{\partial f}{\partial e_{i,j}}(v_i)$ , where  $e_{i,j}$  is some edge with  $c_{i,j} = 0$ . If  $c_i = 0$ , then we require  $s(v_i) = f(v_i)$ ,  $\frac{\partial s}{\partial x}(v_i) = \frac{\partial f}{\partial x}(v_i)$  and  $\frac{\partial s}{\partial y}(v_i) = \frac{\partial f}{\partial y}(v_i)$ .

By the above construction, we obtain a set of points for Lagrange interpolation respectively a set of Hermite interpolation conditions.

**Theorem 4.** Let  $\triangle$  be a triangulation of nested polygons. If for all  $i \in \{1, \ldots, m\}$ ,  $c_i \leq 3$  and no vertex  $v_i$  is simultanously semi-singular (of type 2) w.r.t  $\triangle_i$  and non-singular, then a unique spline in  $S_3^1(\triangle)$  exists which satisfies the above Lagrange (respectively Hermite) interpolation conditions. In particular, the total number of interpolation conditions is equal to the dimension of  $S_3^1(\triangle)$ .

**Proof:** First, we prove that the set constructed above is an admissible set for  $S_3^1(\Delta)$ . To this end, we show by induction that  $\mathcal{A}_i$  is an admissible set for  $S_3^1(\Delta)|_{\Delta_i} = \{s|_{\Delta_i} : s \in S_3^1(\Delta)\}$ . This is clear for i = 0. Now, we assume that  $\mathcal{A}_{i-1}$  is an admissible set for  $S_3^1(\Delta)|_{\Delta_{i-1}}$ , where  $i \in \{1, \ldots, m\}$ , and consider  $V_i$ . For simplicity, we omit here the index i for  $v_i, v_{i,j}, e_{i,j}, p_{i,j}, T_{i,j}$  and  $n_i$ . It follows from the induction hypothesis and Theorem 2 that the coefficients  $a_{\nu,3-\nu-\sigma,\sigma}^{[T_j]}, \sigma = 0, \ldots, 3-\nu, \nu = 0, 1$ , of  $p_j \in \Pi_3, j = 1, \ldots, n$ , on  $T_j$ , are uniquely determined. Moreover, if  $c_{i,j} = 1$  for some  $j \in \{1, \ldots, n-1\}$ , then it follows from Theorem 2 that the coefficient  $a_{2,0,1}^{[T_j]}$  is uniquely determined.

In the following, we show that if  $c_{i,0} = 1$ , then the coefficient  $a_{2,1,0}^{[T_1]}$  is uniquely determined. Let us consider the case where  $v_0$  is semi-singular of type 2 w.r.t.  $\Delta_i$ . (The case that  $v_0$  is semi-singular of type 1 w.r.t.  $\Delta_i$  is analogous.) We denote by  $\tilde{T}_l \in \Delta$ ,  $l = 1, \ldots, 3$ , the triangles with vertex  $v_0$ not contained in  $\Delta_i$  in anticlockwise order, and by  $\tilde{e}_l$  the common edge of  $\tilde{T}_l$ and  $\tilde{T}_{l+1}$ , l = 1, 2. Since  $\tilde{T}_3$  has a common edge with  $\Delta_{i-1}$ , it follows from Theorem 2 that the coefficient  $\tilde{a}_{1,1,1}^{[\tilde{T}_3]}$  of  $\tilde{p}_3 \in \Pi_3$  on  $\tilde{T}_3$  is uniquely determined. Moreover, since  $\tilde{e}_2$  and  $\tilde{e}_1$  are degenerate at  $v_0$ , the coefficients  $\tilde{a}_{1,1,1}^{[\tilde{T}_i]}$  of  $\tilde{p}_l \in \Pi_3$ on  $\tilde{T}_l$ , l = 1, 2, are uniquely determined. Since  $e_0$  is non-degenerate at  $v_0$ , it follows from Theorem 2 that the coefficient  $a_{2,1,0}^{[T_1]}$  is uniquely determined. We note that since  $\Delta$  is a nested-polygon triangulation, at least two triangles with vertex  $v_0$  not contained in  $\Delta_i$  exist. Therefore, if  $c_{i,0} = 0$ , then the coefficient  $a_{2,1,0}^{[T_1]}$  is not yet determined.

Analogously as above, it can be shown that the coefficient  $a_{2,0,1}^{[T_n]}$  is uniquely determined if  $c_{i,n} = 1$ . Otherwise, this coefficient is not yet determined.

Now, we consider the vertex v. The arguments below will show that we may assume that v is an interior point of  $\triangle$ . We denote by  $T_{n+l} \in \triangle$ , l = $1, \ldots, r, r \geq 3$ , the triangles with vertex v not contained in  $\Delta_i$  in anticlockwise order. Moreover, let the piecewise polynomials  $p_{n+l} \in \Pi_3$ ,  $l = 1, \ldots, r$ , on  $T_{n+l}$  in the representation (1) be given such that the first barycentric coordinate always corresponds to v. The above arguments show that exactly  $c_i \leq 3$  coefficients of the set  $\mathcal{C}_1 = \{a_{\nu,3-\nu-\sigma,\sigma}^{[T_i]}: \sigma = 0,\ldots,3-\nu, \nu =$ 2,3, l = 1, ..., n + r are uniquely determined. On the other hand, we construct  $3 - c_i$  additional admissible points from  $C_1$  on  $V_i \setminus V_{i-1}$ . Now, it follows from the  $C^1$ -property at v and Theorem 2 that all coefficients from  $\mathcal{C}_1$  are uniquely determined. By our method of passing through the vertices of  $\triangle$ , v is not semi-singular of type 1 w.r.t.  $\Delta_i$ . In particular, if  $v = w_{\mu}$  for some  $\mu \in \{0, \ldots, k\}$ . Moreover, by assumption v can be semi-singular of type 2 w.r.t.  $\Delta_i$  only if v is singular. In this case, we have r = 3, and it follows from Theorem 3.3 in [13] that the coefficient  $a_{1,1,1}^{[T_3]}$  is uniquely determined. Otherwise, if  $r \ge 4$ , then for some  $l \in \{1, \ldots, r-1\}$  one common edge of  $T_{n+l}$  and  $T_{n+l+1}$  is non-degenerate at v, and we can also proceed with our arguments.

Since all relevant differentiability conditions at the edges with endpoint v, respectively  $v_j$ , were involved, the above shows that  $\mathcal{A}_i$  is an admissible set

for  $S_3^1(\triangle)|_{\triangle_i}$ . Thus, the set  $\mathcal{A}_m$  is an admissible set for  $S_3^1(\triangle)$ .

Therefore, the cardinality of  $\mathcal{A}_m$  is equal to the dimension of  $S_3^1(\triangle)$ . By construction, it is evident that the number of Lagrange interpolation points, respectively the number of Hermite interpolation conditions coincides with this cardinality.

By an inductive argument, it follows from (2) that the Hermite interpolation conditions at v determine the Bernstein-Bézier coefficients of the admissible points chosen on  $V_i \setminus V_{i-1}$ . Analogously, the Lagrange interpolation conditions uniquely determine the interpolating spline on the edges of  $V_i \setminus V_{i-1}$ . Therefore, the interpolating spline is uniquely determined on all of  $V_i \setminus V_{i-1}$ . This completes the proof of Theorem 4.  $\Box$ 

For arbitrary triangulations, Schumaker [12] gave the following lower bound  $L(\Delta)$  for the dimension of  $S_3^1(\Delta)$ ,

$$L(\triangle) = 3V_B(\triangle) + 2V_I(\triangle) + \sigma(\triangle) + 1.$$
(3)

Here,  $V_B(\triangle)$  is the number of boundary vertices of  $\triangle$ ,  $V_I(\triangle)$  is the number of interior vertices of  $\triangle$  and  $\sigma(\triangle)$  is the number of singular vertices of  $\triangle$ . For bounds on the dimension of bivariate spline spaces see also Manni [6].

**Theorem 5.** If a triangulation  $\triangle$  of nested polygons satisfies the hypotheses of Theorem 4, then the dimension of  $S_3^1(\triangle)$  is equal to  $L(\triangle)$ .

**Proof:** We have to show that the cardinality of  $\mathcal{A}_m$  is equal to  $L(\Delta)$ . We prove this by induction. We set  $\mathcal{S}(\Delta_0) = \emptyset$  and for  $i \in \{1, \ldots, m\}$ , we denote by  $\mathcal{S}(\Delta_i)$  the set of boundary vertices w of  $\Delta_i$  such that  $w = v_{l,0}$  and  $c_{l,0} = 1$  (respectively  $w = v_{l,n_l}$  and  $c_{l,n_l} = 1$ ) for some  $l \in \{1, \ldots, i\}$ . Moreover, let  $\tilde{\sigma}_i$  be the cardinality of  $\mathcal{S}(\Delta_i)$  and  $a_i$  be the cardinality of  $\mathcal{A}_i$ . We will show that

$$L(\Delta_i) = a_i + \tilde{\sigma}_i, \qquad i = 0, \dots, m.$$
(4)

This is evident for i = 0. We assume that  $L(\Delta_{i-1}) = a_{i-1} + \tilde{\sigma}_{i-1}$  for some  $i \in \{1, \ldots, m\}$  and consider  $V_i$ . We have  $V_B(\Delta_i) = V_B(\Delta_{i-1}) - n_i + 2$ ,  $V_I(\Delta_i) = V_I(\Delta_{i-1}) + n_i - 1$ ,  $\sigma(\Delta_i) = \sigma(\Delta_{i-1}) + \gamma_i$ , where  $\gamma_i$  is the number of singular vertices from the set  $\{v_{i,j} : j = 1, \ldots, n_i - 1\}$ . Since  $a_i = a_{i-1} + 3 - c_i$ , it follows from the induction hypothesis and some elementary computations that

$$L(\Delta_i) = a_i + \tilde{\sigma}_{i-1} + c_i + \gamma_i - n_i + 1.$$

By our method of passing through the vertices of  $\triangle$ , it is evident that if  $v_{i,0} = v_{i-1} \in \mathcal{S}(\triangle_i)$ , then  $v_{i,0} \notin \mathcal{S}(\triangle_{i-1})$ . In the following, we show that if  $v_{i,n_i} \in \mathcal{S}(\triangle_i)$ , then  $v_{i,n_i} \notin \mathcal{S}(\triangle_{i-1})$ . First, let us assume that  $v_{i,n_i} = v_{l,0}$  for some  $l \in \{1, \ldots, i-1\}$ . If  $v_{i,n_i}$  is semi-singular of type 2 w.r.t.  $\triangle_i$ , then at least three edges of  $\triangle$  not contained in  $\triangle_l$  are attached to  $v_{l,0}$ . Hence,  $c_{l,0} = 0$ . If  $v_{i,n_i}$  is semi-singular of type 1 w.r.t.  $\triangle_i$ , then the edge  $e_{i,n_i}$  is non-degenerate at  $v_{i,n_i}$ , since  $c_{i,n_i} = 1$ . Therefore,  $v_{l,0}$  is not semi-singular of type 2 w.r.t.  $\triangle_l$ . Again,  $c_{l,0} = 0$  holds. The remaining case  $v_{i,n_i} = v_{i-1,n_{i-1}}$ , where  $n_i = 1$ , follows by the same arguments.

Now, we show for  $j \in \{1, \ldots, n_i - 1\}$  that every non-singular vertex  $v_{i,j}$  such that  $e_{i,j}$  is degenerate at  $v_{i,j}$  lies in  $\mathcal{S}(\Delta_{i-1})$ . First, we consider the case j = 1. Set  $v_{i_1} = v_{i,1}$  and let  $\tilde{e}_0$  be the edge that connects  $v_{i,0}$  and  $v_{i,1}$ . We have to consider two cases.

**Case 1.** (The vertices  $v_{i,1}$  and  $v_{i-2}$  are connected by an edge e.) If  $\tilde{e}_0$  is non-degenerate at  $v_{i,1}$  then  $c_{i-1,n_{i-1}} = 1$ . (In this case  $v_{i,1}$  is semi-singular of type 1 w.r.t.  $\Delta_{i-1}$ .) Otherwise, since  $v_{i,1}$  is non-singular, e is non-degenerate at  $v_{i,1}$ . Thus,  $c_{i-2,n_{i-2}} = 1$ . (In this case  $v_{i,1}$  is semi-singular of type 2 w.r.t.  $\Delta_{i-2}$ .) We note that  $v_{i,1}$  is not semi-singular w.r.t.  $\Delta_{i,1+1}$ , since at least three edges of  $\Delta$  not contained in  $\Delta_{i,1+1}$  are attached to  $v_{i,1}$ .

**Case 2.** (The vertices  $v_{i,1}$  and  $v_{i-2}$  are not connected by an edge.) If  $\tilde{e}_0$  is non-degenerate at  $v_{i,1}$  then we also have  $c_{i-1,n_{i-1}} = 1$ . (In this case  $v_{i,1}$  is semi-singular of type 1 w.r.t.  $\Delta_{i-1}$ .) We note that  $v_{i,1}$  is not semi-singular w.r.t.  $\Delta_{i_1+1}$ , since  $\tilde{e}_0$  is non-degenerate at  $v_{i,1}$ . Otherwise, let e be the edge that connects  $v_{i,1}$  with  $v_{i_1+1}$ . Since  $v_{i,1}$  is non-singular, e is non-degenerate at  $v_{i,1}$ . Thus,  $c_{i_1+1,0} = 1$ . (In this case  $v_{i,1}$  is semi-singular of type 2 w.r.t.  $\Delta_{i_1+1}$ .) We note that in this case  $v_{i,1}$  is semi-singular of type 1 w.r.t.  $\Delta_{i-1}$ , but  $c_{i-1,n_{i-1}} = 0$ .

Now, we consider the remaining case  $j \in \{2, \ldots, n_i - 1\}$ . Set  $v_{i_j} = v_{i,j}$ and let e be the edge that connects  $v_{i,j}$  with  $v_{i_j+1}$ . Since  $v_{i,j}$  is non-singular, it follows that  $v_{i,j}$  is not semi-singular of type 2 w.r.t. to  $\Delta_{i_j}$ . Therefore, e is non-degenerate at  $v_{i,j}$ . Hence,  $c_{i_j+1,0} = 1$ . (In this case  $v_{i,j}$  is semi-singular of type 1 w.r.t.  $\Delta_{i_j+1}$ .) We note that in the case  $j \in \{2, \ldots, n_i - 1\}$ , by our method of passing through the vertices of  $\Delta$ , the value  $c_{i_j-1,n_{i_j-1}}$  is not influenced by the geometrical properties of  $\Delta$  at  $v_{i,j}$ .

The above proof now implies  $\tilde{\sigma}_i = \tilde{\sigma}_{i-1} + c_i + \gamma_i - n_i + 1$ , and therefore, (4) holds. Since  $\tilde{\sigma}_m = 0$ , we get  $L(\Delta) = a_m$ . This proves the theorem.  $\Box$ 

In Theorem 4 we assume that for all  $i \in \{1, \ldots, m\}$ , no vertex  $v_i$  is simultaneously semi-singular (of type 2) w.r.t.  $\Delta_i$  and non-singular. In the following, we show that this assumption is essential for the local construction of interpolation points.

**Example 6.** Let  $v = v_i = (0, 0)$ ,  $v_5 = v_0 = v_{i,0} = (\gamma, 0)$ ,  $\gamma < 0$ ,  $v_1 = v_{i,1} = (\tau, m\tau)$ ,  $\tau < 0$ , m > 0,  $v_2 = v_{i,2} = (0, \delta)$ ,  $\delta < 0$ , and set  $v_3 = (\alpha, 0)$ ,  $\alpha > 0$ ,  $v_4 = (0, \beta)$ ,  $\beta > 0$ . Let v be connected with  $v_3$  and  $v_4$  and  $v_{l-1}$  be connected with  $v_l$ ,  $l = 1, \ldots, 5$ . Then v is simultaneously semi-singular (of type 2) w.r.t.  $\Delta_i$  and non-singular. Furthermore, we denote by  $T_l$  the triangle with vertices  $v, v_{l-1}, v_l$  and by  $p_l \in \Pi_3$  the polynomial pieces on  $T_l$ ,  $l = 1, \ldots, 5$ , in the representation (1). We consider the set  $C_2 = \{a_{\nu,3-\nu-\sigma,\sigma}^{[T_1]}, \sigma = 0, \ldots, 3-\nu, \nu = 1, \ldots, 3, l = 1, \ldots, 5\}$ . For  $C^1$ -splines, it follows from Theorem 3.3 in [13] that each subset of  $C_2$  that uniquely determines all coefficients of  $C_2$  has cardinality 8 and contains the coefficients  $a_{1,0,2}^{[T_1]}$ , l = 3, 4. By the proof of Theorem 4,

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the coefficients  $a_{1,2-\sigma,\sigma}^{[T_1]}$ ,  $\sigma = 0, 1, 2$ , and  $a_{1,2-\sigma,\sigma}^{[T_2]}$ ,  $\sigma = 1, 2$ , are uniquely determined. If  $e_{i,1}$  is non-degenerate at  $v_1$ , then in addition  $a_{2,0,1}^{[T_1]}$  is uniquely determined. Otherwise, this coefficient is not determined. Hence, if  $e_{i,1}$  is nondegenerate at  $v_1$ , then we have to choose exactly one additional coefficient to determine all coefficients of  $C_2$ , and otherwise, we have to choose exactly two additional coefficients. We claim that in the latter case every choice of exactly two additional coefficients from the set  $\{a_{3,0,0}^{[T_1]}, a_{2,0,1}^{[T_1]}, a_{2,0,1}^{[T_2]}\}$  fails to determine all coefficients of  $C_2$ .

**Proof:** Suppose that  $e_{i,1}$  is degenerate at  $v_1$  and choose, for example,  $a_{2,1,0}^{[T_1]}$  and  $a_{2,0,1}^{[T_1]}$ . For simplicity, we set  $a_1 = a_{2,1,0}^{[T_3]}$ ,  $a_2 = a_{3,0,0}^{[T_3]}$ ,  $a_3 = a_{2,0,1}^{[T_4]}$ ,  $a_4 = a_{1,1,1}^{[T_4]}$ ,  $a_5 = a_{2,1,0}^{[T_4]}$ ,  $a_6 = a_{1,1,1}^{[T_3]}$ , and assume that the remaining coefficients in  $C_2$  are zero. By Theorem 2,

$$a_{3} = (\frac{\tau - \gamma}{m\tau\gamma}\beta + 1)a_{2}, \quad a_{4} = (1 - \frac{\alpha}{\gamma})a_{3}, \quad a_{1} = (1 - \frac{\delta}{\beta})a_{2} + \frac{\delta}{\beta}a_{3},$$
  
$$a_{6} = (1 - \frac{\delta}{\beta})a_{5} + \frac{\delta}{\beta}a_{4}, \quad a_{5} = (1 - \frac{\alpha}{\gamma})a_{2}, \quad 0 = ((-\tau)(\frac{m}{\delta} + \frac{1}{\alpha}) + 1)a_{1} + \frac{\tau}{\alpha}a_{6}.$$

Eliminating  $a_j$ ,  $j \in \{3, 4, 5\}$ , yields  $a_1 = (1 + \delta \frac{\tau - \gamma}{m \tau \gamma})a_2$ ,  $a_6 = (1 - \frac{\alpha}{\gamma})(1 + \delta \frac{\tau - \gamma}{m \tau \gamma})a_2$ . By some elementary computations, we obtain for the determinant D of the corresponding system

$$D = rac{(-1)( au(m\gamma+\delta)-\delta\gamma)^2}{m au\delta\gamma^2},$$

and it is easy to verify that D = 0 iff  $e_{i,1}$  is degenerate at  $v_1$ . Other choices of exactly two additional coefficients from the set  $\{a_{3,0,0}^{[T_1]}, a_{2,1,0}^{[T_1]}, a_{2,0,1}^{[T_2]}, a_{2,0,1}^{[T_2]}\}$  can be examined in the same way, which proves our claim.  $\Box$ 

Note that if  $e_{i,1}$  is non-degenerate at  $v_1$ , then every choice of exactly one additional coefficient in the set  $\{a_{3,0,0}^{[T_1]}, a_{2,1,0}^{[T_2]}\}$  determines all coefficients in  $C_2$ .

We finally discuss some numerical aspects of our scheme. A method for constructing nested polygon triangulations  $\triangle$  of given points in the plane which satisfy the conditions of Theorem 4 was developed in [11]. Our numerical tests show that in order to obtain good approximations, it is necessary to subdivide some of the triangles (for details see [10,11]). Meanwhile, we have computed such examples with a high number of interpolation conditions. We only mention here that, for example, Lagrange respectively Hermite interpolation of Franke's test function by cubic  $C^1$ -splines with 118,822 interpolation conditions yields an error of 4.66902 \* 10<sup>-6</sup> in the uniform norm.

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