NAVAL POSTGRADUATE SCHOOL Monterey, California



19980417 151

Luqi

THESIS

DTIC QUALITY INSFECTED 4

A SYNTAX DIRECTED EDITOR FOR THE COMPUTER

AIDED PROTOTYPING SYSTEM

by

Charles A. Mock

September 1997

Thesis Advisor:

Approved for public release; distribution is unlimited.

REPORT DOCUMENTATION PAGE

Form Approved OMB No. 0704-0188

Public reporting burden for this collection of information is estimated to average 1 hour per response, including the time for reviewing instruction, searching existing data sources, gathering and maintaining the data needed, and completing and reviewing the collection of information. Send comments regarding this burden estimate or any other aspect of this collection of information, including suggestions for reducing this burden, to Washington headquarters Services, Directorate for Information Operations and Reports, 1215 Jefferson Davis Highway, Suite 1204, Arlington, VA 22202-4302, and to the Office of Management and Budget, Paperwork Reduction Project (0704-0188) Washington DC 20503.

1. AGENCY USE ONLY (Leave blank)	2. REPORT DATE September 1997	3. REPOR Master's	RT TYPE AND DATES COVERED Thesis	
4. TITLE AND SUBTITLE A SYNTAX DIRECTED EDITOR FOR PROTOTYPING SYSTEM.	THE COMPUTER AIDED		5. FUNDING NUMBERS	
6. AUTHOR(S) Mock, Charles A.				
7. PERFORMING ORGANIZATION NAME(S) A Naval Postgraduate School Monterey, CA 93943-5000	ND ADDRESS(ES)		8. PERFORMING ORGANIZATION REPORT NUMBER	
9. SPONSORING / MONITORING AGENCY NAME(S) AND ADDRESS(ES)			10. SPONSORING / MONITORING AGENCY REPORT NUMBER	
11. SUPPLEMENTARY NOTES			I	
The views expressed in this thesis are those	se of the author and do not re	flect the offici	al policy or position of the	
Department of Defense or the U.S. Govern	nment.		··· ·	
12a. DISTRIBUTION / AVAILABILITY STATEM	IENT		12b. DISTRIBUTION CODE	
Approved for public release; distribution i	s unlimited.			
13. ABSTRACT (maximum 200 word	ds)			
The Computer Aided Prototypin	g System (CAPS) is an integ	rated set of so	oftware engineering tools developed at	
the Naval Postgraduate School (NPS). It	is designed to support rapid p	prototyping of	f real-time systems. CAPS consists of	
four major subcomponents; the graphics/	text editor, the user interface	e, the softwar	re database system, and the execution	
support system. Reports from users of	CAPS, particularly novices,	indicated th	at the clumsy and unintuitive multi-	
windowed graphics/text editor present in	the system hampered the use	of the tool se	t. This thesis presents the substitution	
and integration of an efficient and user-fri	endly syntax directed editor i	nto CAPS. 7	The new syntax directed editor consists	

of a package of seven Ada95 parsers that recognize the elements of the Prototype System Description Language (PSDL) and an enhanced C\Motif based graphics editor. These modules combine the functionality of all the windows of the graphics/text editor into one window, using pop-up boxes and menus to guide the designer in providing the proper

occurring between modules and the internal consistency was being maintained at the inter-language interfaces. The result

During integration, particular attention was paid to ensuring the proper manipulation of data was

19. SECURITY

ABSTRACT

Unclassified

CLASSIFICATION OF

NSN 7540-01-280-5500

17. SECURITY

Unclassified

14. SUBJECT TERMS

CLASSIFICATION OF REPORT

information.

is a faster, intuitive, and more efficient designer interface.

18. SECURITY

Unclassified

CLASSIFICATION OF THIS PAGE

Syntax Directed Editor, Computer Aided Prototyping

UL

15. NUMBER OF PAGES

• <u>177</u> 16. PRICE CODE

20. LIMITATION OF

ABSTRACT

Standard Form 298 (Rev. 2-89) Prescribed by ANSI Std. 239-18

.

Approved for public release; distribution is unlimited

A SYNTAX DIRECTED EDITOR FOR THE COMPUTER AIDED PROTOTYPING SYSTEM

Charles A. Mock Captain, United States Marine Corps B.A., University of Florida, 1989

Submitted in partial fulfillment of the requirements for the degree of

MASTER OF SCIENCE IN COMPUTER SCIENCE

from the

NAVAL POSTGRADUATE SCHOOL September 1997

Author:

Charles A. Mock

Approved by:

Luqi, Thesis Advisor

Michael J. Holden, Second Reader 1L

Ted Lewis, Chairman Department of Computer Science

.

ABSTRACT

The Computer Aided Prototyping System (CAPS) is an integrated set of software engineering tools developed at the Naval Postgraduate School (NPS). It is designed to support rapid prototyping of real-time systems. CAPS consists of four major subcomponents; the graphics/text editor, the user interface, the software database system, and the execution support system. Reports from users of CAPS, particularly novices, indicated that the clumsy and unintuitive multi-windowed graphics/text editor present in the system hampered the use of the tool set. This thesis presents the substitution and integration of an efficient and user-friendly syntax directed editor into CAPS. The new syntax directed editor consists of a package of seven Ada95 parsers that recognize the elements of the Prototype System Description Language (PSDL) and an enhanced C'Motif based graphics editor. These modules combine the functionality of all the windows of the graphics/text editor into one window, using pop-up boxes and menus to guide the designer in providing the proper information. During integration, particular attention was paid to ensuring the proper manipulation of data was occurring between modules and the internal consistency was being maintained at the inter-language interfaces. The result is a faster, intuitive, and more efficient designer interface.

vi

I.	INTRO	DUCTION	.1
	А.	GOALS	.1
	B.	BACKGROUND AND MOTIVATION	.2
	С.	RESEARCH QUESTIONS	.3
	D.	ORGANIZATION	.4
Π.	BACK	TROUND	.5
	A	INTRODUCTION	5
		1 The Computer Revolution and The Software Crisis	6
		2 The Ranid Prototyning Solution	7
	в	SOFTWARE DEVELOPMENT SYSTEMS	. / 8
	D.	1 The Waterfall Model	.0 8
		2 Ranid Prototyning	10
		3 The Computer Aided Prototyning System (CAPS)	13
		4 Prototype System Design I anguage (PSDI)	16
	C	THE DESIGNER INTERFACE	20
	C.	1 Granhical Editor	.20
		 Oraphicar Lantor Syntax Directed Editor 	.21
		2. Sylitax Difference Lanoi	.22
		5. FOLL paising	.23
	Л		.20
ттт			.27
ш.	ARCHI	DECIDED ADAL SOL	.29
	· A.	SDECTED A TION	.29
	D. С		.32
	U. D		.54
137			.40
1.	INFLC		.41
	А.	1 Introduction	.41
		1. Introduction	.41
		2. Environment	.42
		5. Languages	.42
		4. Editor	.44
		5. Miscellancous Tools	.45
		 version Control 1001 I avial Applyment 	.40
		7. Lexical Analyzers	.4/
	ъ		.4/
	С		.48
	U. D		.48
37	ע. תוסדות		.50
۷.	KESUL		.33
	A. D		.33
١л	D. CIDAN		.38
¥1.	SOIMINT		.39
	А. Ъ		.39
	<u></u> . С		.01
در ۸	U. VICINZ		.61
		A. BUUNCE CODE (FBUL I I FEB)	.03
	LINDIA		./1
		C. BUTCE CODE (FOUL EDITOR)	.109
- AF			147

APPENDIX E: TEST DATA	153
LIST OF REFERENCES	163
INITIAL DISTRIBUTION LIST	

ACKNOWLEDGEMENTS

This research was possible due to the efforts of many people. Most notably is that of my thesis advisor, Professor Luqi and the members of the CAPS group. Professor Berzins' mentorship was truly instrumental in my understanding of the internal workings and interaction of the graphics editor and the syntax directed editor. The friendship and infallible cheerfulness of Professor Ying Jing made the long and arduous hours of coding bearable. I am also deeply indebted to Commander Mike Holden. Undoubtedly in the top 5% of instructors at this institute, he always places the student first. For his invaluable contributions to the organization and content of this paper, at a high cost to himself, I will always be grateful.

I would also like to thank the students and faculty members at the Naval Postgraduate School. I have found the educational experience to be a positive one and this is only possible with the outstanding faculty and peers with which this institution is blessed.

Finally, I am grateful to my wife Juli for her patience and understanding throughout this research. Without her support, none of this would have been possible.

I. INTRODUCTION

A. GOALS

The goals of this thesis are to produce an intuitive and efficient syntax directed graphical editor and document its integration into the Naval Postgraduate School's Computer Aided Prototyping System (CAPS).

CAPS operates using the Prototype System Design Language (PSDL), a design language created to present software designers with a tool for creating rapid prototypes. A syntax directed editor (SDE) for PSDL is a tool for editing PSDL programs, called prototypes. The editor checks each element of the PSDL language entered into the prototype by the designer for errors. This is done using a set of parsers designed to analyze text for specific words or strings of letters, known as tokens. If the tokens are correct for the language, the designer may proceed with the design, if not, the designer is alerted to the error and directed to the position of the mistake. In this way, an SDE can assist the designer in creating prototypes in a quick and effective manner.

CAPS is an integrated set of tools designed to allow designers to rapidly produce a prototype matching user specified requirements. Rapid prototyping allows the designer to clarify the needs and desires of the user early in the software development cycle. This eliminates costly errors that normally are not discovered until later, when significant effort is required for correction. The intent in inserting an updated designer interface is to compress the time necessary to create a working prototype by making CAPS easier to use. A successful integration of an enhanced syntax directed editor would make CAPS

more accessible to novices and allow experienced users to minimize time lost correcting errors during the prototyping cycle. This, in turn, will speed up the software development cycle.

B. BACKGROUND AND MOTIVATION

There is a problem with the creation of software today. It is costing the nation billions of dollars in failed projects and wasted efforts. Often the problem stems from a faulty design process. One solution to this problem, CAPS, is presently hampered by a poorly designed interface. Correcting this defect is one step toward solving the greater issue.

The well-know truism, "time is money", is especially relevant in the design of software. Certainly, the longer a project takes the higher the cost for labor and materials. The truth of this is quite obvious. The desire to reduce costs and meet deadlines invariably results in constricted timelines for identifying requirements, outlining specifications, and creating a solid design. The rush to implementation allows a product to be swiftly pushed out the door, but a cost must then be paid in user satisfaction and maintenance.

Time is money. As much as eighty percent of the cost involved in software production is maintenance [Ref. 4]. Yet that is not the least of the problem. In 1995, IEEE reported that for the Department of Defense (DOD) \$42 billion in software programs were canceled due to overruns in time and budget or they simply did not perform as required [Ref. 4]. The waste of billions of dollars is not confined to the DOD. Other prominent debacles include the automated baggage system at the Denver airport and the IRS integrated data network. While expensive, these systems are cheap compared to the seven billion dollar loss involved in the destruction of the Ariane rocket [Ref. 14] and the death of 225 passengers in the 6 August 1997 crash of a Boeing 747 on Guam [Ref. 16], both attributed to software errors.

Few of the problems with modern software systems involve poor coding. A NASA study performed during the Galileo project found that a full 98% of all software errors are traceable to incorrect requirements, specifications, or design [Ref. 4]. The advanced state of modern integrated development suites such as Aonix ObjectAda 7.1, Microsoft Visual Studio '97 and Borland C++ 5.0, combined with rigorous testing, have nearly eliminated implementation errors. Unfortunately, few tools exist for properly creating a design that faithfully represents the user's desire. CAPS is one of them.

Designed for rapid prototyping of real time systems, CAPS is aimed at reducing software design errors by presenting the user with successively more accurate representations of the final product. In order to do this effectively, CAPS itself must be fast and easy to use. That requires an easily learned, intuitive interface. That is the rationale for this thesis.

C. RESEARCH OBJECTIVES

The objectives for this thesis include the following specific objectives:

• Create parsers that correctly interpret the seven main elements of the Prototype System Design Language (PSDL).

- Automate coding of large segments of the project using existing UNIX environment tools and software engineering procedures
- Integrate the new parsers into a previously designed graphical editor.
- Integrate the combined graphical editor/parser package into the main flow of the CAPS environment.
- Document the flow of control and the design of data structures into a maintenance handbook to improve maintenance of this new enhanced version of CAPS.

D. ORGANIZATION

Chapter II provides a general overview of the workings of CAPS, PSDL and the need for and uses of rapid prototyping. Chapter III examines the processes and structures created during the analysis of the proposed evolution of CAPS for requirements, the specification of those requirements, and the creation of a design for the modifications to the CAPS code, with emphasis on the requirements for the appropriate interfaces, data manipulation, and improvements in storage management. Chapter IV presents the implementation phase of the project as well as testing and evaluation of the new version. In Chapter V the results are presented, with discussion on the utility of the improved interface and maintainability of CAPS. Chapter VI summarizes the thesis.

II. BACKGROUND

A. INTRODUCTION

The creation of a modern graphical editor for CAPS is a project that has been envisioned since 1991. By 1997 it became imperative that this goal become a reality for two major reasons.

The first is the state of software development. There is a software development crisis looming over the information age [Ref 8]. The rapid increase in the speed of computer hardware and the use of computer technology has not been matched by a corresponding improvement in software design. Software design has not only failed to keep pace, but has increasingly been the reason for late projects, exceeding budgets, and often complete failure of projects. This amounts to what was coined in 1968 as 'The Software Crisis'. This crisis has as its root cause the problem of complexity (brought about by sheer length of programs) combined with a poor control over how each line of code affects the overall system [Ref. 17]. These problems can often be corrected by using an alternative software development model. In many cases, rapid prototyping is a better method for designing software than current methodologies and CAPS is a powerful tool that could formalize the creation of software, if it achieves acceptance on a wide scale. Acceptance of CAPS on a wide scale requires a modern user-friendly interface.

The second reason is the participation of the Software Engineering Group in the DOD's Technology Transfer Program. The Department of Defense sponsors billions of dollars in advanced research yearly. In order to enable the various research laboratories

within the DOD to reap the maximum benefit from this effort, the Office of Technology Transition (OTT) was established. OTT serves as a clearinghouse for coordinating and facilitating the transition of such technologies and technological advancements within the DOD to other military organizations and the private sector. CAPS has been distributed throughout the DOD as a result of this program, and is in use by a large number of organizations. A number of these supporters have requested that improvements to the interface and in storage management be included with an enhanced version. In order for the enhanced version to be available for pending projects, its creation was necessary as soon as possible.

1. The Computer Revolution And The Software Crisis

The Information Age is here and is expanding at an extraordinary rate. The burgeoning use of computers and information systems in nearly every aspect of daily life is undeniable. Microprocessors and software programs inhabit our automobiles, videocassette recorders, coffee makers, and calculators. Fifty-one million Americans regularly access the Internet [Ref. 1] and it is the stated goal of the United States Government to place every classroom in the nation online. Riding the wave of computing power created by the realization of Moore's Law [Ref. 15], the use and scope of automated systems are increasing at a geometric rate.

Unfortunately, the same cannot be said for the state of software development. While microprocessor speed continues to double every 18 months, the creation of software remains labor intensive. Success is heavily dependent upon the skill and

experience of the individuals involved in production. Improvement is linear, if at all. Errors are common. The costs for labor are high, with starting salaries over \$50,000 common for recent four-year computer science graduates [Ref 13]. Complicating this equation is the continuing decline in the number of skilled professionals capable of producing high quality software. Over the previous decade, the number of degrees awarded for computer science in the U.S. has dropped an astounding 42% [Ref. 13].

2. The Rapid Prototyping Solution

With demand for faster, better, bigger, and more complex software accelerated by the explosion of information technology and automated systems, the cost of increasingly scarce software expertise has risen in tandem. It has become incumbent upon software manufacture's to maximize the productivity of software designers, engineers, and programmers in order to remain competitive and meet demand. This is where the utility of integrated rapid-prototyping environments such as CAPS becomes important. By enhancing the effectiveness in identifying and specifying requirements, CAPS decreases errors, rework, and reduces the overall time to complete a software project. However, for CAPS to be seriously considered as a viable alternative solution it must be easily understood and used. It is the intent of this thesis to make CAPS a more attractive alternative to other software development systems by creating an intuitive, user friendly interface.

B. SOFTWARE DEVELOPMENT SYSTEMS

1. The Waterfall Model

Problems exist with the predominant software development methodology, the Waterfall Model (Figure 1.). The Waterfall Model is a linear plan for the production of



Figure 1. Waterfall Methodology for the software life cycle.

software. The user and the designer analyze the users needs and produce a set of requirements that are used to create the design for the software system. The program is implemented, tested, and then returned to the user for validation.

Unfortunately, most errors occurring in software are requirements based. This stems from the imprecise communication that often occurs between the end user, who understands the problem domain but not software design, and the software designer, who understands software development but usually possesses little knowledge of the problem domain. The user, who is the ultimate arbiter of what is correct in the software, is unable to validate a system effectively until a working version is produced. At this point, a large portion of the development time and budget has been expended, leaving little time or



Figure 2. Relative cost for error correction over the software life cycle. [Ref. 4] money for corrective action. Furthermore, an invalid requirement at this stage has affected large portions of the implementation code, resulting in a 100 fold increase in the effort necessary to correct the mistake over detecting the same error before implementation had taken place (Figure 2).

2. Rapid Prototyping

One workable answer to the problem is rapid prototyping. As illustrated by Figure 3, rapid prototyping does not rely on validation of a finished product. As in the Waterfall Model, the users initial goals are analyzed and a set of requirements created. These are used to create a partially functional prototype that can be demonstrated to the user. The user then assists in identifying faults in the design of the system very early on,



Figure 3. Rapid Prototyping Software Life Cycle.

before valuable time and money are poured into propagating those faults through the entire system.

The payoff for using rapid prototyping is evident. Requirements are often changed repeatedly by users who are initially unsure of the functionality required, leading to wasted effort and frustration. Rapid prototyping obtains user feedback and solidifies requirements early, providing a clear path for future progress and reducing friction between user and designer.

In traditional software development, specifications evolved from the user requirements are presented in written form or a formal specification language. Neither of these methods are suitable to users with little or no knowledge of software design. Misunderstandings result. Rapid prototyping allows the user to perform validation on executable specifications. For example, if a user requires the system to initiate the firing sequence of a rocket motor when a certain air speed is reached, a prototype of the system, complete with simulations of the missile hardware, can be completed and demonstrated to the user long before a physical prototype could be delivered and at a fraction of the cost. The user could then evaluate the prototype's performance based on whether the rocket's motor was signaled to ignite at the appropriate air speed without actually having to launch the missile. This interaction between the user and actual working prototype results in a clearer understanding by both the user and designer of what the user truly desires.

The risk of failure in software development is considerable. Colonel Chadwick, the commander of the Marine Corps Tactical Systems Support Activity stated that software projects with a budget in excess of \$100 million had a 100% failure rate [Ref. 18]. Unfortunately, this is not an isolated phenomenon. In 1995, DOD wide, only 16% of software programs were completed on time and on budget. Furthermore, an astounding 31% were cancelled entirely (Figure. 4). Rapid prototyping can help resolve this problem. Large and complex systems can be modeled quickly, concentrating on the design of the software architecture while leaving the details for later implementation. This allows the user and designer to perform risk assessments and feasibility studies using the prototype. This is preferable to waiting for delivery of the beta version of the total system, when time and budget are nearly exhausted and changes are much more expensive.

•. • ·



Figure 4. DOD software delivery statistics for 1995 [Ref. 4].

While the advantages of prototyping are many, the prototyping process must be fast in order to be effective. The designer has limited time in which to create the final product and the sooner the prototype receives final validation from the user, the sooner it can be forwarded to the programming team for optimized implementation. This implies automation. Automation can generate code faster and more accurately than any human programmer as well as provide organization to help with decision support. Database software can automate searching through a system to find code that meets specific design requirements. Allowing software to perform tedious and time consuming tasks can be used to speed up nearly every aspect of the project.

3. The Computer Aided Prototyping System (CAPS)

CAPS supplies the automation that rapid prototyping requires. Hosting an array of tools to assist in producing a prototype in minimal time, CAPS is a completely integrated software design environment. The major areas of automation support provided by CAPS are designer interface, software databases, execution support, and an evolution control system. Combined, these tools allow the designer to perform all the necessary functions required by the validation cycle of the rapid prototyping model (Figure 5).

The designer interface contains the syntax directed editors and consists of:

Editing tools for the Prototype System Design Language (PSDL).
 Consisting of the Graphical Editor and the PSDL Editor, this is the portion of the CAPS system concerned with creating and modifying prototype



Figure 5. CAPS prototype creation flow diagram.

designs. It is also the major subject of this thesis.

- An Ada programming language editor. CAPS currently uses the Verdix Ada83 editor for creating and modifying Ada modules for use in the CAPS system.
- A requirements editor. The requirements editor allows for creating, updating, and tracing the accomplishment of user requirements over the life of the prototype.
- An interface editor. Currently the interface editor is The Transportable Applications Environment (TAE) Project, a state-of-the-art user interface development and management system. The interface editor allows the designer to rapidly create a graphic user interface from which to display and input data into the prototype.

These tools are included to assist the designer in rapidly entering information pertinent to the creation of the prototype. Speed is important. For the selection of editors, a premium was placed on propagation of appropriate constraints, preventing syntax errors, and providing for robust error diagnostics. Propagation of constraints is a necessity to ensure that the final product performs within limits initially set by the user. Prevention of errors and robust diagnostics serves a dual purpose. In the early stages of prototype design, the detection and elimination of errors is much less costly than correcting the same error later in the development cycle. The reduction in the number of errors in the design also serves to reduce the time required to create a prototype, allowing a working prototype to be demonstrated to the user quicker.

The Database consists of two sections:

- The Design database. The design database provides for the storage of prototype development data. This allows a designer to search the database for previous prototypes that may have components that fit functionality in the present design. Allowing reuse of prototype components not only speeds up the process, but reduces errors as completed prototypes have been tested repeatedly during the rapid prototyping cycle.
- The Software database. Just as the design database allows reuse of existing prototypes, the software database lets the designer search for reusable Ada and PSDL components. Additionally, since PSDL specifications are formalized, they can be used as the index for a database query that returns a component that is a very accurate match to what the designer needs.

Execution support provides compilers and other tools necessary to convert the PSDL code and graphs into a working real-time prototype. Execution support allows the designer to make an executable prototype a reality in four easy steps.

These are:

- The Translator translates PSDL language to Ada code.
- The Scheduler creates schedules for execution of a real-time prototype with corresponding Ada code.

- The Compiler compiles Ada code into executable form, presently Sun Ada version 1.1.
- The Execution Shell- opens a shell in which to run a prototype.

4. Prototype System Design Language (PSDL).

PSDL is the design language that supports CAPS. Standard programming languages are optimized to support efficient execution of a program. They consist of complex algorithms and data structures for creating small, fast executables. They are quite detailed and require a well-trained programmer in order to be used efficiently. Design languages are quite different. Execution is not the prime motivation. The objective is the creation of an efficient and easily understood design. Because of this, design languages are more expressive and simpler to use. Furthermore, where programming languages provide for comments as an adjunct to the main functionality, design languages have the recording of goals and justifications as a prime objective.

PSDL achieves these goals. Based on a simple grammar and a graph, PSDL provides for easy understanding of the language while yielding a powerful ability to create prototypes.

The most basic building blocks of PSDL are operators and streams. They represent the two things that can be done with data: manipulation or relocation. Operators are functions (or state machines) and streams are data flows.



Operators have three manifestations. Ordinary operators, represented by circles in

Figure 6. Graphical Interface for CAPS.

the example of the graphical editor shown (Figure 6) perform operations on the data from the streams. Composite operators, designated by a circle with a double ring, represent a sub-graph. Sub-graphs are standard PSDL graphs, and are represented by composite operators for the sake of clarity. A graph would rapidly become cluttered and confusing without the ability to compress functional areas into composite operators. Decomposition of functionality is easily achieved using this construct, as composite operators can contain other composite operators, allowing decomposition to continue until only simple atomic functions remain. Terminators represent sources of input external to the design. Terminators are simulations of external systems, are not required for the prototype, and

and the second second

will not be in the delivered software. The third kind of operator is the type operators used to represent the functionality of abstract data types defined by the user.

A stream is a communications link connecting two or more operators and is represented in the GE by a line connecting operators. Operators are producers and consumers. A stream requires at least one operator, producer or consumer. There are two types of streams available in PSDL, dataflow streams and sampled streams. The dataflow stream works on the first in, first out principle. Data values are never lost or replicated. Sampled streams model continuous data input. Only the most recent data value is used by the consumer with all others being lost.

Triggers are control constructs that act as decision points on the firing of operators as a result of receiving data through a stream. Triggers can be set to fire in the event that some or all of a set of data is present. Execution guards are the constructs within the operator in which triggers operate. If no execution guard is present within an operator, the default causes the operator to always execute. Execution guards and triggers can be thought of as implementing an if-then-else programming control structure.

Operators can also be fired through the use of a timer. Timers are software stopwatches that are declared within a composite operator and hold a time expression that, while running, is continuously updated to record the passage of real time. Timers have three control constraints that affect the value held by the timer. They are START, STOP, and RESET. Start obviously causes the timer to run. Stop freezes the current value. Reset returns the current value to zero. Like triggers, timers, and execution guards, other output guards exist to assist the user in controlling the flow of a program. Output guards allow conditional transmission of data to an output stream. Output guards are assigned by stream and are normally used to filter output to the remainder of the program. This allows a multiply threaded stream to pass data selectively rather than broadcasting to all.

PSDL is a real time prototyping language. As such, timing constraints are a necessary part of the language. The most important of PSDL's timing constraints is the maximum execution time (MET). The MET is the maximum amount of time that the operator is allowed to complete its activities. The addition of a MET constraint to an operator defines that operator as time critical and thus subject to timing constraints. Two types of time critical operators, periodic and sporadic, are used within PSDL.

Timers trigger periodic operators. A value respresenting the period of the operator is set to a time value. When the timer reaches that value, the operator is allowed to fire. It is not necessarily required to fire immediately. A second value, finished within (FW), is provided to allow flexibility in scheduling. After the operator is triggered by the period, the program can delay execution of the operator as long as the operator is completely executed prior to the arrival of the FW time.

Sporadic operators are triggered by the arrival of data. These operators are executed whenever data arrives subject to the minimum calling period (MCP) of the operator. The MCP is a time value. If a sporadic operator executes, it cannot be executed again until the time indicated by the MCP has elapsed, regardless of the receipt of data. This allows the designer to sample data at intervals, and prevent responses to every byte of data sent to the operator. Sporadic operators also require a maximum response time (MRT). The MRT is the maximum amount of time that can elapse between the arrival of data and the completion of the operation. This constraint provides the designer with a mechanism to ensure that critical work is completed within a known amount of time.

Operators also contain data elements that assist human users in understanding the purpose of that component. The description element allows the designer to place a natural language description of what the operator does within the structure of the PSDL code. This acts much like a comment would in a programming language. Keywords are also employed by PSDL. The keyword component allows the designer to place natural language identifiers within the prototype in order to assist future designers in evaluating the component for reuse. Keywords are used by the database as indices for queries.

Composite operators contain one other construct of interest, the graph description. PSDL provides for a component that describes the design graph of a decomposed composite operator to the graphical editor. The graphical editor translates this component in order to display the design within the editor window. A complex structure, the graph description contains coordinates, colors, fonts, and all visual components of child operators and streams.

C. THE DESIGNER INTERFACE

The graphical editor and the syntax directed editor are two representations of the same PSDL code. The graphical editor displays a truncated version of the code based on

the objects comprising a PSDL graphic. The syntax directed editor displays the hidden information that gives the prototype depth. When CAPS was originally designed, the use of graphical interfaces was a rarity and they were often poorly designed. Designers also prefered to directly access the PSDL code if they wished. These decisions are no longer valid. Graphical tools are much better and abstraction in software design is realized to be essential for truly understanding a system. There is no inherent reason for this separation of PSDL code into two separate interfaces.

1. The Graphical Editor (GE).

The graphical editor is the portion of the designer interface with which the designer creates the graphical representation of the prototype (Figure 6). The GE is simple in design and execution, allowing novice operators to begin creating real-time prototypes using the basic stream and operator objects with little formal training. With the use of the Quickstart manual for CAPS, a novice can create a simple prototype within minutes using only the graphical interface. Although a powerful concept and tool, the CAPS graphical interface is unable to deliver the necessary functionality required by more complex prototypes. For this, CAPS provides a syntax directed PSDL editor that works in conjunction with the graphical interface to produce a working design.

· · · · · · · ·

2. Syntax Directed Editor.

The CAPS syntax directed editor (SDE) displays the text version of the PSDL prototype – the part that the translator converts to programming language code (Figure 7). The SDE allows the designer to perform detailed, advanced PSDL programming. It is the principal interface for the designer, allowing the creation, editing and viewing of CAPS designs. PSDL language code generated by creating a design is modified within the SDE. The SDE also provides access to a wide variety of tools, such as the ability to

psdl_editor:c3i_system.psdl CAPS-Cnds File Fdd View Tools Onlines Sharkum Taxt	- teri
Read /tmp_mnt/n/sun55/work/mock/.caps/c3l_system/1.1/c3l_system.psdl	eth
of input sensor data) END IMPLEMENTATION ADA analyze_sensor_data	
IND OPERATOR c3i_system	
WARNINOS, ERRORS AND ALERTS:	
This Is A Root Operator 	
SPECIFICATION DESCRIPTION (This module implements a simplified version of a generic C3I workstation.)	
INCLEMENTATION GRAPH see graph viewer for details	
DATA STREAM comms_email : filename, comms_add_track : add_track_tuple, tdd_filter : set_track_filter, out_tracks : track_tuple,	
<pre>weapons_emrep : weapon_status_report, input_link_message : filename, initiate_trans : initiate_transmission_sequence, terminate_trans : BOULTAW, terd mission corrected : emission control</pre>	
tcd_network_setup : emissions_control_command, tcd_archive_setup : network_setup, tcd_transmit_command : transmit_command, position_data : ownship navigation info.	
sensor_add_track : add_track_tuple,	•

Figure 7. PSDL Editor.

query the software database, that assist the designer in creating a prototype.

The separation of these two editors is also the major deficiency in the design of CAPS release 1.0. As early as 1990 users of CAPS were underscoring the need for a syntax directed editor combined with the graphical editor for generating PSDL code [Ref. 2]. The necessity of writing PSDL language code in the SDE brings with it the entire training overhead associated with programming languages. Furthermore, since the modification of the design in the GE almost always requires a corresponding change in the SDE, the division of the two editors into separate windows has no practical purpose to justify the cumbersome environment. Merging these two windows and their underlying functionality is the major objective of this project.

3. PSDL parsing.

A syntax directed editor assists the designer in creating a prototype by automatically checking the PSDL language code for appropriate structure. For the purposes of this project, PSDL has seven major constructs that may be input into the graphical editor. Each of these constructs must be checked for syntax errors. Needless to say, the descriptions of the seven parsing elements are truncated. The full specification of the meanings of each can be found in Appendix D.

The seven constructs are:

• Expressions: Expressions represent a wide range of tokens within PSDL. An expression can be a string, an integer literal, Boolean expression, a type

operator identifier, time value, data stream, timer or a combination of the above. Example: X > 2.

 Initial expressions: Initial expressions are exactly the same as an expression except references to time values, timers, and data streams are excluded. Example: Y = 3.

• Output guards: A set of logical statements initiated by the OUTPUT token, output guard statements consist of a list of identifiers followed by IF, then an expression and ending with a requirements trace. For example:

OUTPUT temperature IF x<2 req1

In this example temperature is a data stream. When the stream is triggered, if x is less than two, then requirement #1 is fulfilled.

• Timer operations: Timer operations start with a timer_op token. This consists of one of three values - Start, Stop, and Reset. The timer_op token is followed by the identifier (name) of the timer performing the operation. This is then followed by the IF token, an expression, and a requirements trace. This works in the same manner as the output guard, with the exception that the expression is evaluated when the timer performs the action represented by the timer op token. Example:

START TIMER

timer1 -

• Operator specifications: A sequence of attributes and requirements traces comprise the body of this structure, starting with the OPERATOR token and

.

an identifier, followed by the SPECIFICATION token and concluding with an END token. An attribute is simply a characteristic of an operator, such as an input, output, state or an exception. Example:

OPERATOR Producer_1 SPECIFICATION GENERIC G1 : FLOAT OUTPUT DA : Missing_Info MAXIMUM EXECUTION TIME 0 MS END

• Type specifications: This element begins with TYPE token and an identifier followed by a SPECIFICATION token and ends with an END token, in the same manner as the operator specification except the interior is more complex. This component describes the attributes of a user-defined type. Generic declarations are parsed first, followed by a set of operator specifications as described above. A final list of functions available to the type completes the structure. Example:

```
TYPE STACK
SPECIFICATION
   GENERIC
     types : private
   type_2 : public
   OPERATOR PUSH
   SPECIFICATION
       INPUT
         I : INTEGER
       INPUT
         S : STACK
       OUTPUT
          S : STACK
   END
    OPERATOR POP
    SPECIFICATION
```

```
INPUT
           S : STACK
         OUTPUT
          I : INTEGER
        OUTPUT
          S : STACK
    END
    OPERATOR Empty
    SPECIFICATION
        OUTPUT
          dummy : STACK
    END
    KEYWORDS
      stack, adt
    DESCRIPTION
      {This is a generic stack adt}
    AXIOMS
      \{push (s,x) = s::x\}
END
```

• Exception guards: Exception guards begin with the EXCEPTIONS token followed by an identifier. The IF token is next followed by an expression and a requirements trace. With the exception of substituting the EXCEPTIONS token for the OUTPUT token, the exception guard and output guards are identical.

4. **Previous Work.**

The grammar necessary for creation of the seven parsers that were created in this thesis was already a part of the original CAPS code created by Professor Luqi. This code was combined in the Aflex and Ayacc files for parsing PSDL code as a whole language. The work necessary for converting these files into useful material for this project consisted of identifying and separating the syntax elements of a particular PSDL component from the overall grammar. This greatly simplified this portion of the project.
Originally, a commercial graphical editor had been selected for use with the new version of CAPS. A powerful tool, it also provided a very intuitive user interface. Unfortunately, licensing restrictions connected to this product would have seriously hampered the free use and distribution of the completed system. As an alternative, a graphical editor created as a class project by NPS graduate students, under the guidance of Professor Mantak Shing , was selected. Although not as robust and extensive as the commercial version, enhancements added by Ken Moeller, a graduate, have made it a viable and cost effective replacement for the commercial graphical editor.

D. SUMMARY

Rapid prototyping is a viable solution that is available now that can alleviate many of the problems presently experienced by the software industry. CAPS is an excellent system with which to create prototypes. An effort to improve the efficiency with which a software prototype designer can use CAPS by simplifying and enhancing the interface is valuable to the Naval Postgraduate School, the DOD, and the software development community as a whole. An improved interface allows prototypes to be developed more accurately and in less time. The reduction in time to create a prototype version creates a corresponding reduction in time to produce a user-validated design. An improved interface reduces errors; again saving time spent correcting implemented code.

The interface is essential in creating an effective CAPS environment. An effective CAPS environment is important in extending the use of rapid prototyping. Extending the use of rapid prototyping will reduce the amount of late, poor, or simply

unusable software being produced. This is a small step on a road that can save the DOD and the nation billions of dollars.

(a) A second of the second se second seco

III. ARCHITECTURE OF THE PSDL EDITOR

The Software Development Method is normally divided into five distinct phases: Requirements Analysis, Functional Specification, Architectural Design, Implementation, and Evolution and Repair [Ref. 5]. This chapter will record the first three phases as they occurred during the project.

A. **REQUIREMENTS ANALYSIS**

Requirements analysis begins with the receipt of an initial problem statement from the customer. The CAPS team possessed an advantage in this respect, being the customer and the designer. The initial problem statement was quickly developed as:

The purpose of the proposed software system is to provide a more efficient and user friendly interface for the CAPS environment.

The goals of requirements analysis are to define the purpose of the proposed system and to determine the constraints on development [Ref. 5]. To achieve this, the initial problem statement is reviewed in the context of the environment. Specific requirements and constraints for the combined graphical editor are then developed.

Three specific requirements for the system were developed:

• Create seven PSDL parsers to be integrated into an existing graphical interface that will allow syntax directed editing of PSDL components from within the graphical editor. The editors must check one of the previously denoted PSDL constructs for accuracy, and return an error location in the event that a syntax error is discovered.

- Retain all the functionality of the present SDE and graphical interface. It is necessary to retain all functionality in order to ensure that prototypes currently developed under CAPS 1.1 are usable in the new version being developed.
- Remove the memory leak present in the current version. The failure to
 reclaim memory following the destruction of temporary data structures causes
 the previous version of CAPS to slowly drain resources from the overall
 system. Over time, if CAPS continues to run, all memory resources are
 exhausted and the system fails. The new system will rigorously recycle
 memory freed when data structures are no longer necessary and are destroyed.

Constraints on development must also be addressed during requirements analysis. Constraints are limits which the designer is required to respect in the construction of the proposed system. There are many ways in which constraints can be documented. For the CAPS project, constraints were divided into five categories: Resources, Performance, Environment, Form, and Method.

1. Resources.

Time to complete the project was limited. The Software Engineering Group at NPS participates in the DOD technology transfer program. As such, a number of organizations were dependent upon the next release of CAPS to accomplish their own objectives. Although the original delivery date was set for 17 July 1997, a more realistic assessment rescheduled the delivery of the product on 12 August 1997.

Manpower was also limited. The CAPS team consisted of seven members. Furthermore, only four of the seven were available for full time work on the project. Given that the project resulted in the modification of 65,000 lines of code, having only seven people provided a limited pool of man-hours from which to draw.

Budget considerations were minor in comparision to a comparable software effort within the private sector. Over the course of the project 3024 man hours were expended or about one and a half man-years. At current labor costs for software programmers/designers this would have represented an investment of approximately \$100,000. Other budget costs were inconsequential.

2. Performance.

Constraints on performance were limited to insistence that memory storage management be rigidly controlled to correct the memory leak in the original release.

3. Environment.

The targeted system was to use SunOS 4.1.3 operating on a Sparc 10 workstation The system will have Motif X.11 available. No other environmental constraints were imposed.

4. Form.

The system would be implemented in Ada95 to assist in maintainability, and the Graphical Editor (GE) developed and enhanced by NPS graduate students would be incorporated into the design.

an an an an the the

5. Methods.

There were no constraints assigned in regard to methods.

B. SPECIFICATION

The goal of the functional specification stage is to define precisely the external interfaces of the proposed system [Ref. 5]. The process of integrating the graphical editor (GE) into CAPS and integrating a PSDL syntax directed editor into the GE creates an unusual situation where interfaces that normally would be considered internal are external for our purposes. Three interfaces need to be considered (Figure 8).

The first interface for specification is between the CAPS prototyping menu and the PSDL editor. The action taking place consists of passing the name of the PSDL file to the PSDL editor. Passing of data occurs in only one direction. One argument is



Figure 8. PSDL Editor Architecture.

•• .

required. It must be a string with a length of at least six characters, the last five of which match the string ".psdl".

The second interface is also between the PSDL editor and the CAPS Prototype, except in the opposite direction, and requiring different arguments. This is not a direct call but rather an implicit obligation on the part of the PSDL Editor to supply useable PSDL code for translation. As such, the specification for this interface required that the PSDL Editor save a valid PSDL prototype to the file called when the editor was initialized. To accomplish this interface, a filename must be supplied that is exactly the same six character string as in the first interface. It must also have a valid PSDL *prototype*, an abstract data type defined in the editor, to save to the file.

The third interface is between the PSDL Editor and the GE. This interface passes data in both directions, with the editor feeding the GE the graph description of the prototype and the GE returning that description along with the users choice for the next action to be performed and any error messages that may be necessary. The transfer of information in this interface is complicated by the difference in languages between the two modules. Fortunately, data types exist within the PSDL Editor that fit the needs of the interface. The graph can be passed using an in/out parameter *Graph_Desc_Node*, the action will use an out *Action* type, and the error messages will also be an out parameter of the type *error_msg*. The *Action* and *error_msg* types are simply enumerations and are readily converted to structures of the C programming language. The *Graph_Desc_Node* is a specially constructed data type consisting of C compatible structures designed for passing graph data.

The final interface joins the GE and the parsers that make up the SDE. Again, a cross language exchange of data is necessary and again, data flows in both directions. The data that needs to be exchanged consists of a character string that requires parsing for correctness, a Boolean variable designating whether the string passes the parsing test, and integers designating the line, column, and length of the first error discovered. As the C to Ada library of functions and available data types is not as extensive as the reverse, the interface used C pointers., which can be readily converted in Ada95, providing access to each of the five data elements being passed.

C. DESIGN

Design encompasses the decomposition of the system into software modules.



Figure 9. Functional Decomposition of the PSDL Editor

Many of the modules were already in place and only required modification, which was extensive in some cases. However, a description of the improved PSDL editor is in order. As an aid to understanding, Figures 9 - 13 show the PSDL editor being used to model and prototype the PSDL editor.

At the top most level, the PSDL editor is expected to perform two functions (Figure 9). Upon initiation the program receives a PSDL filename as an argument from the command line of the operating system. The *Open* function takes that filename and searches the default prototype directory in search of the named file. If it exists, it is opened and the prototype name contained therein is passed to the next function,



Figure 10. Functional Decomposition for Edit_Program

edit_program. If the file is not found, a new file is created using the given filename and an empty prototype is forwarded.

Edit_program receives the prototype name, manipulates its contents and then returns the modified program to file. Edit_program, because of complexity of the tasks assigned, must be decomposed further (Figure 10).

When initialized *Edit_program* begins by assigning the prototype an unique identification number. This allows the system to recognize the prototype uniquely even



Figure 11. Functional Decomposition of Edit_Operator

if the user eventually creates other components using the same name. After assigning the identity, *read_prototype* reads the entire PSDL program from the file into a skeleton prototype data structure. The prototype is then subjected to semantic checks to ensure that the program is in a useable format and not corrupted in any way. The verified prototype is then ready to be edited by the GE. This is performed within the substructure of *edit_operator*, which will be detailed below. Upon completion, *edit_operator* will have received the modified prototype and an action code indicating the user preference for further use of the system. *Update_action* tests the action returned and performs the necessary action, such as *save_to_file* or *revert*. *Reinvoke* is then initiated, testing whether further action is required. If true, the program loops back to *semantic_check* and performs the cycle again. If false, the *Close* function closes the opened file and returns control to the operating system.

Edit_operator performs only three functions, but they are vital to the interface with the GE (Figure 11). First, user defined types and operators are separated within the operator list. This is necessary because types and operators display difference characteristics and must be manipulated in different manners in order to present the GE with a useable graph. The separation results in two lists of operators and types. These lists and a blank graph data stucture are passed to *make_graph_desc*. Make_graph_desc takes the PSDL data presented and manipulated it to form a graph description. The graph description is a data structure that represents all the information contained in the prototype, converted to use C language conventions. This allows the edit_graph interlanguage function call to import the data for use within the GE. Upon completion of user action, the GE returns the graph description to *edit_operator* along with the user's next requested action.



Figure 12. Functional Decomposition of Save_To_Disk

Save_to_disk, the most typical of users actions, also needs to be reviewed. It consists of two functions, update_prototype and save_prototype (Figure 12). Both functions do what their names suggest, updating the PSDL prototype (as opposed to the C oriented graph description) and saving the prototype to the file. Of the two, update_prototype deserves closer inspection.

Update_prototype is the function that unwraps the changes performed within the GE and converts them into PSDL code that is useable by the remainder of CAPS. It



Figure 13. Functional Decomposition of Update Prototype

consists of six subfunctions (Figure 13). The first function is *build_edited_types*. The second is *build_edited_operators*. These two functions review the graph description returned by *edit_graph* and determine which types and operators have been changed. The PSDL structure is then modified to reflect the changes. Upon completion of these updates, the children of modified composite operators must be updated, as well as any grandchildren, etc. *Modify_children* performs this task using a recursive structure that checks for changes rippling through to each operator from the root down. After the children are checked for changes, the type operations are updated. At this stage, all

modifications performed within the GE will have been implemented in PSDL structure. Unfortunately, the GE will not provide data on operators that were unchanged. This requires that a function be created to check the original prototype against the one being built and add any unmodified operators back into the prototype. The final action in the reconversion of the graph description is *synthesize_inferred_parts*. This function evaluates changes made by the GE that would effect other portions of the PSDL code indirectly, such as the setting of terminator flags, type definitions, and exception declarations.

A final note on the design of the seven PSDL parsers is in order. The parsers were designed to accept a string for an argument, and return an error message and error location. All are C structures. This allows each individual parser to be called directely from the appropriate place in the graphical editor.

D. SUMMARY

This chapter detailed the first three of the five phases (Requirements Analysis, Functional Specification, Architectural Design, Implementation, and Evolution and Repair) of the software development methodology for the enhanced PSDL Editor. The next chapter deals with its implementation, and the fifth phase, Evolution and Repair, will be discussed in Chapter VI.

IV. IMPLEMENTATION

Implementation is the fourth phase of the software development cycle. In this chapter we will cover the aspects of implemention as they applied to the enhanced PSDL Editor project. The tools used, the approach to attacking the problem, and the process by which the project was completed will be covered. Repair and Evolution (R&E) is the final phase of the software development cycle but is properly outside the scope of this paper. Testing and evaluation are normally considered as part of this phase and were a very important part of this project and will be covered in the place of R&E.

A. TOOLS AND ENVIRONMENT.

Tools can make or break a software development effort. The CAPS group attempted to assemble a comprehensive suite of tools that would meet our requirements at a minimal expense.

1. Introduction.

When a project expands to larger than a few hundred lines of code; extensive hand coding cannot be done. Hand coding introduces errors and becomes an unacceptable risk to the robustness of the final product. The proper line of attack is to use pre-existing tools in order to automate large portions of the implementation process.

2. Environment.

The environment for the project was standard for work in the software engineering lab at the Naval Postgraduate School. The project code was created on a Sun Micro Systems Sparc 10 workstation networked within the software engineering laboratory subnet, which in turn is linked to the computer science department network. Keeping the project confined to a subnet allowed work to continue at all times, even when the larger department net was offline for backups.

The operating system installed on the test system was SunOS version 4.1.3. Chosen for proven reliability and known compatibility with the present CAPS system, version 4.1.3 is the dominant operating system within the NPS computer science department, ensuring widespread compatibility of newer CAPS versions with the department at large. One account was used for the entire project with each member of the team having access.

3. Languages.

Programming languages for the project were chosen in deference to the constraints applied to the project by previous work.

Ada95 was selected as the main language for working within the main SDE modules. The original version of CAPS had been implemented using Ada83 in accordance with the DOD Ada mandate being enforced at the time. Ada95 is a logical choice as a successor for that reason alone. However, Ada95 has advantages that would have resulted in its selection regardless. Like CAPS, Ada95 is designed for software

engineering. Ada95 includes all of Ada83's discipline enhancing aspects, such as strong typing and information hiding, which made Ada83 the language of choice for large and complex projects [Ref 9]. In addition, Ada95 is now object oriented, adding flexibility to robustness. Using Ada reduces the cost of initial development of software and significantly reduces the risks involved in program development. Maintenance costs are reduced. Ada is especially useful for projects that involve a large number of personnel. Although this project used only seven team members, the nature of work at NPS has shown that working groups expand and contract. Ada's readability supports easy understanding of code without recourse to extensive commenting. This is an essential feature for long term projects, and CAPS, now operational for nine years, is a long-term project.

GNAT (GNU/NYU Ada Translator) Ada95 version 3.09 became the main compiler. A perfectly serviceable compiler, GNAT was the first compiler to implement the Ada95 design [Ref. 7]. It is certified. It was also documented as supporting the Unbounded_Strings abstract datatype, which contained the key functionality for enhancing storage management, a key requirement. Finally, cost is always a consideration, and GNAT provided the features needed by the project at no cost.

The choice to use the C programming language was also determined by the requirements of the project. The existing graphical editor to be integrated into the CAPS program was created using C. Although regrettable from the standardization and maintenance points of view, the original creators of the graphical editor did have good and sufficient reason for choosing C in that the X.11 Motif function calls necessary for

building a windowing environment were all in that language. Fortunately, experience in the C language is widespread and a number of team members were quite proficient in its use.

The C compiler chosen for the project was "gcc", the gnu C compiler, which is included as a standard part of SunOS. The reasons for its acceptance were the familiarity of the compiler to the team members and, as with the GNAT Ada95 compiler, it was available at no cost.

"Awk", though called a programming language by it's creators (and from whose initials it gets its name), is really a pattern matching language [Ref. 6]. Awk was included among the languages needed because early in the design phase, it became apparent that considerable amounts of hand reworking of code segments could be avoided if certain recognizable patterns could be replaced automatically. Awk performs tasks like this using very few lines of code, making it an excellent addition to the arsenal of tools available.

Make is not in itself a programming language, but is a tool that uses Unix shell commands to create powerful scripts for building executable code. Make was chosen because of this ability and the lack of any viable alternative.

4. Editor

The main text editor employed by the project team was vi, the standard visual line editor provided with the operating system at no cost. If employed, another text editor would have required installation on the project system, eating into project time and

administrative resources. Knowledge of vi's command set was widespread within the team, even if the knowledge was shallow in some cases. In addition, vi's command set is easily combined with Unix shell commands to create extremely powerful searching and replacing routines. This capability in itself was enough to guarantee inclusion in the project tool set.

5. Miscellaneous Tools

"Sed", the Unix stream editor, acts in much the same way as Awk, but possesses a simpler grammar. It was chosen in order to perform quick substitutions of text. As part of the Unix tool set, it also possesses the advantages of availability and cost.

"Gen" is a code manipulation tool for creating loops within the Ada structure that would normally be considered illegal by the Ada Reference Manual. For example, consider the following code segment:

For child_vertex: op_id in op_id_set_pkg.scan(vertices(g)) loop

Normally this would not be allowed because there is no way for the compiler to ascertain the number of child_vertex that will be discovered using the scan function. Gen solves this problem by unrolling the loop, generating new Ada code in the process. The advantage of using this tool over hand coding the affected areas was significant

"Gnatchop" is a text search tool for by the the GNAT compiler that scans a file for Ada package bodies and specifications. Gnatchop then divides the file into corresponding Ada files. Names are assigned according to the convention that the file is named the same as the package contained with ".adb" appended for the package body and

".ads" for the specification. Gnatchop was recognized as necessary to automate processes. Gen, Aflex, and Ayacc were all developed for Ada83 and produce files that require splitting before compilation. Gnatchop not only helped us to produce Ada95 packages, but allowed the use of sophisticated makefile scripts to automate the process.

6. Version Control Tool

The Revision Control System (RCS) was chosen as the archiving system for the project. RCS tracks changes in files from one version to the next and saves only the changes made. This delivers a substantial reduction in the amount of magnetic storage necessary to backup a multiple version library for a large project. In addition, RCS allows adding comments and descriptions to archived versions for easy recognition in the event retrieval is necessary. Beyond the standard need for a backup system for the code being produced, the use of RCS by the implementation team was necessary for project control. SunOS 4.1.3, while an excellent operating system, has no mechanisms to prevent simultaneous editing of a file by two or more programmers.

The implementation team used a process of mirrored public and private directories for holding the latest working version and the current version being modified. The latest working version was to be exported to the public directories following a full archive with RCS. The private directories would hold the code being created and modified. This process possessed two important advantages. Testing could be conducted on a stable version of the system in parallel with the creation of the next version, and in

46

the event that changes to the code created compiler errors that could not be corrected immediately, a working version of the system was still available for demonstrations.

7. Lexical Analyzers

Aflex and Ayacc are the Ada versions of the popular C tools, lex and yacc. Aflex is a lexical analyzer/generator that accepts a file containing the definitions of lexical element to be recognized and returns a file with a ".a" suffix that contains a lexical scanner recognizing tokens corresponding to those elements. Ayacc accepts a set of tokens, such as those created by Aflex, and a set of rules provided by the designer. Together, the tokens and rules make a grammar. Ayacc generates Ada code for a parser to recognize that grammar. Both tools have been used for years and are well known. For a project such as creating parsers for a syntax directed editor, ignoring the benefits provided by these tools would have cost the team days of hand coding.

8. Debuggers

Two debuggers were employed for this project. Both debuggers were from GNAT and were designed for the C and Ada languages. The reasons for the choice of these debuggers were cost, compatibility, and availability. Availability was the key factor in the choice of the Ada debugger. It is one of the more disheartening aspects of working with Ada95 at the current time that an effective debugger, public domain or commercial, is extremely difficult to find for the SunOS environment.

B. APPROACH

The CAPS team considered a number of factors when planning how to approach the implementation of the project. Automation was of prime importance to the group. Limited time in which to complete the effort and a shortage of available man-hours required that the maximum value of automated assistance be achieved.

The experience of senior members of the team in creating sophisticated makefiles was primary importance to achieving maximum automaton. A quick look at some of the code within the PSDL Editor Makefile will serve to demonstrate this. Eight of the major packages are created using Gen, one of them being modified by Sed in the process. A short script is available to run a test case directly following a compilation. Two scripts were created to allow for easy export of private versions to public directories, and for entering a current version into the RCS.

The other process used to enhance the effectiveness of the group was the early division of the four main designers into teams of two. The newer members of the group were then able to learn the system from experienced personnel while providing a second set of eyes to spot syntax and logic errors. The investment in time up front undoubtedly paid off in the end in reduced errors, team spirit, and knowledge learned.

C. PROCESS

The implementation process was performed at multiple sites over a period stretching from 9 June to 11 August 1997. While the majority of the effort was performed within the CAPS Laboratory and the Software Engineering Laboratory at NPS, one team member was located in Lancaster, CA. Furthermore, at various times, other members were working from as far away as San Diego, CA and Brazil. This dispersion over time and distance complicated the problems of project management, especially in the initial phases and made the use of the RCS indispensable to success.

The parser creation project was the first part of the implementation phase undertaken. It was an elementary undertaking that involved the analyzing of an existing Aflex grammar for PSDL and removing the components that were not applicable to the parser at hand. This part of the project was very automated, taking advantage of the properties Awk, Sed, Aflex, and Ayacc to achieve success.

Testing of the parsers outside of the GE proved them to be accurate. Later testing within the GE found them to be less so. A default property within the GE caused streams not assigned a specific type to be listed as "?" in the graph description. The "?" had no meaning to the parsers and locating the error became a difficult process since the offending symbol was propagated to the graph description from within a window that is not normally parsed for errors.

The second phase of implementation involved the creation of the PSDL editor interface to the GE. This was a project of considerable detail. Implementation was simply a time consuming task occasionally punctuated with challenging problems. The most serious of these problems involved the GNAT compiler, version 3.09, and the attempt to improve storage management through use of Ada.Unbounded_Strings, which was new in Ada95. Repeated system crashes ultimately forced concentration on discovering what appeared to be a particularly difficult fault in the code. Extensive debugging and testing found no code error, but rather an instability in the Unbounded_Strings package. This forced the abandonment of large segments of code created to make use of that functionality.

Software engineering princples were enforced during the implementation process and a number of poor programming practices performed during the original implementation of CAPS were corrected. There was a strong reliance on large omnibus utility packages in the original code. This clear violation of the principles of object oriented programming and information hiding was removed in all cases but one. This was not a trivial project. The break up of the GE utitilities package alone created 13 new packages. The one package not broken down is the PSDL concrete types package. This massive file would have required far more man-hours than was available to reduce to its component pieces. It is therefore slated for replacement during the next evolution of the system.

D. TESTING AND EVALUATION

Testing occurred concurrently with integration. This was done to achieve maximum benefit from working in parallel. Frequently during a project, programmers lose productive hours waiting for another team to complete testing of a required section of code. Our concurrent approach avoided this loss. It also ensured that the implementers received accurate and timely feedback on the viability of code modules, essential in preventing the propagation of errors throughout the code structure. Testing was conducted in cycles. When a version was exported to the public directories, testing was performed and the results returned to the programming team. While the testing was being performed on a version, the implementers were correcting problems identified during the previous test cycle. During the course of integrating the SDE into CAPS, twelve testing cycles were completed, uncovering 75 errors that were then corrected. The improved version of CAPS was not released until three error-free testing cycles had been completed -- testing cycles without an error that could result in an abnormal termination of the program. While continued testing with real users was necessary to thoroughly test the design, the system had to be robust enough to allow work to proceed.

V. RESULTS



A. GRAPHICAL INTERFACE

Figure 14. Improved Graphical Interface.

The primary objective of this project was the merging of the SDE and the Graphical Editor into one efficient module. That has been accomplished (Figure 14) and, without entering into detail appropriate to a user manual, a short display of the new features is in order. The user now can access the entire functionality for software prototype design through a single window. Although concentrating the functionality in one window is a major benefit, the true value of the new GE becomes apparent when the prototype designer requires information. All the most frequently used information groups are placed at the designer's fingertips. This is the true benefit as it allows the designer to maintain focus on the task at hand.

Buttons for operators, streams and terminators still exist and perform the same functions. The properties and select buttons have been removed. Simply clicking on that object using the left mouse button does selection of a graph component. Accessing the

·.	Operator	Proper	ty	· · · · · · · · · · · · · · · · · · ·
	in de p			
		_		
Name: User_1	nterface		Operatio	e 12 (* 1
			rendere.	
Implementation Lang	nade:	ece .		
Iriggers Unprot	ected .			
IF Co	ndition	TRUE		
Reartin	H Ry II	la de la composición de la composición Composición de la composición de la comp		
	<u></u>			
Tiaing:	rindic 👘	1	4.002	
HET:		f	8	1
			Kequineo Byr	
rer100: 1000	<u>. ns</u>		Required By-	
Finish	ins.	- . f	Required By	
Rivering, P				
Utput Buards				
Exception Suar	- age de la company			ð., j
Exception List				······································
Liner ups	STAR	T TIMER I	1 IF (a ≻= 90	
Keyacrids	-Informal 1	esc	Formal De	sc III (
		285		
	<u>a Cance</u>		<u>HELP</u>	

Figure 15. Operator Property Box.

properties of an operator or stream is now performed by clicking on that object using the right mouse button. This brings up the operator property box (Figure 15) or the stream property box (Figure 16). The old editor allowed only the modification of the Maximum Execution Time and the period from the properties box. Now available for modification are requirements, guards, triggers, implementation languages, descriptions and keywords. The more extensive variety of fields reflects not an increase in complexity of the editor, but a decrease in the requirement for hand coded PSDL components. Previously, the designer would be required to enter these attributes directly into the PSDL code in the PSDL SDE, with a corresponding loss of focus on the flow of the design and a significant increase in the possibility of creating an error. The new GE does not increase the functionality of the CAPS system in this case, but it makes the functionality that always existed much more accessible. Attributes that were invisible to the novice before are now readily available.





Figure 16. Stream Property Box.

PARENT SPEC, TIMERS and GRAPH DESC.

The first allows access to the timers tool box (Figure 17). The timers tool box enables the designer to view and modify timers from within the GE at the touch of a button. The tool box itself contains a suite of options for operating on timers. Prior to



Figure 17. Timers Tool Box.

this modification, timers were created and edited within the PSDL editor, with a corresponding loss in focus from the task at hand.

The parent specification box (Figure 18) displays and allows editing of information pertaining to a parent composite operator. When working on a decomposition graph, it is frequently necessary to access the data of the parent operator in order to accurately set up the child graph. This used to entail extensive searching for the appropriate information within the PSDL code. Now it is available, and modifiable, directly from the interface.

The graph description button on the new GE's toolbar provides data on the PSDL graph. The PSDL graph contains a wealth of information pertaining to the operation of

the prototype. Previously, the designer was required to continually check the PSDL editor in order to gather the necessary information about the PSDL graph Time is saved and attention to the project at hand is retained.

]	?rototype	Specifica	ation	
View on Edit Prototype	Specification	• • • • •		7. H. S
DESCIOR sharped toat	MB	19 11 12 12 1		
SPECIFICATION		4 7 .		
STATES data_size:	integer INITI	ALLY 1	- in .	
STRIES data per lo	Integer INITI	AIYA		
STATES start_stop:	boolean INIT	IALLY FALSE		
EXCEPTIONS e4				
EALEPTIDIG 68	10. 10. 10 A. 1			
	•			
1	·····	******	*******	Channess 2
		<u></u>		
l ≡ K		incel		lelp
Participation of the second				



Lastly, the GE provides further two functions to assist the designer. At the bottom of the editor window (Figure 14.) is an information window and a button labeled "CHECK". The information window simply informs the designer in the event that the graph has changed since the last save. If it has, then the designer receives a gentle reminder to save the prototype. The "CHECK" button is an error checker. If an error is created that the parsers detect, the label will change to "ERROR MSGS". This is to alert the users that an mistake has been made in the PSDL code. This can occur without the syntax parsers alerting the designer when editing data directly on the graph. If the button is depressed while in "CHECK" mode, the SDE will scan the PSDL program for errors and return the result. In effect, the check button is a miniature compiler for PSDL that generates a report of correctness.

B. TECHNOLOGY TRANSFER

The Software Engineering Group at the Naval Postgraduate School is a full participant in the DOD Technology Transfer Program. CAPS version 1.1 has been in use by various DOD agencies since its creation in 1988. In accordance with this, on 12 August 1997, the CAPS version 1.2 with the enhanced syntax directed graphical editor was presented to members of the Army Research Office, the Army Research Lab and the Naval Research and Development activity for further testing and improvement. When follow up testing is completed, the new version will be disseminated to the various agencies within the DOD that support the rapid prototyping paradigm and use or support CAPS.

Transfer of CAPS technology within the DOD has resulted in the successful application of rapid prototyping and the CAPS environment to numerous projects. Currently the Navy's SmartShip and the Distributed Computing Networks programs are evaluating the use of CAPS in support of DOD research.

VI. SUMMARY AND RECOMMENDATIONS

A. SUMMARY

The integration of an improved SDE into the CAPS environment will decrease the overall time necessary to produce a software prototype and can substantially reduce the number of errors per thousand line of code. With proper use, a similar decrease can be achieved in both the time and development costs associated with software development through the use of CAPS in the early stages of a project. Four of the five original objectives were completed successfuly.

1. Creation of parsers that correctly interpret the seven main elements of PSDL.

The first objective completed, the creation of parsers for PSDL, was the simplest portion of the project. Extensive code reuse contributed to the accomplishment of this goal. Always a prime tenet of sound software engineering, the reuse of code in this case significantly reduced the time necessary to create the parsers, by reusing over 3000 lines of validated and tested code. The parsers themselves work perfectly and the error-locating mechanism properly presents the position and size of any faulty code the parsers discover.

and the second second second

2. Use existing UNIX environment tools and software engineering procedures to automate coding of large segments of the project.

The use of Awk, Aflex, Ayacc, and gen significantly reduced the amount of hand coding necessary to create the parsers needed. However, it was the creative use of the Unix make utility that allowed large portions of code creation to be automated.

3. Integrate the new parsers into a previously designed graphical editor.

Integration of the parsers into the Graphical Editor was completed and required relatively little effort.

4. Integrate the combined graphical editor/parser package into the main flow of the CAPS environment.

This goal required the most extensive work with the system. The multiple levels at which data was transferred to and from the graphical editor as well as the language conversion issues tended to make straightforward solutions impractical. Success was due to hard work and rigorous testing.

5. Document flow control and design of data structures into a maintenance handbook to improve the maintenance of the enhanced version.

Although an admirable goal and one that would have greatly enhanced the ability of future prototype designers to understand the nuances of the code for the SDE, time constraints forced the postponement of this objective. The creation of a maintenance handbook will be left for future work.

B. LESSONS LEARNED

Planning ahead is a key to success in software development and this project was not an exception. Time not spent properly specifying requirements and outlining the design of the new system will be spent correcting errors in the testing and evaluation phase of development.

Proper use of Unix Makefiles to automate portions of a program was an enormous time saver and resulted in more accurate code. Effective use of makefiles can only be accomplished with thorough knowledge of the tools available and how to link them together, in effect creating an assembly line for the creation of code, much like that used in manufacturing. The power of an integrated tool set was demonstrated during this project. Even though only linked by shell commands and the make tool the capabilities were immense. This concept truly requires more emphasis in computer science curricula.

C. RECOMMENDATIONS FOR FUTURE RESEARCH

CAPS is a versatile and powerful designing environment that can be used by the DOD now. However, the system could benefit from research in a number of areas. The DOD is moving toward Microsoft WindowsNTTM as the standard for information systems. It is essential that a version of CAPS that operates under WindowsNT is

at the second second

produced soon. The graphical editor in CAPS would benefit from a new type of component, a composite stream, that would decrease the complexity of large designs. The graphical editor would also benefit from an auto-redraw function to assist the user in removing clutter. Implementing a CAPS prototype in CAPS would greatly improve the ability of the system to grow and evolve over time. Lastly, the PSDL parsers function admirably on PSDL components, but there is no mechanism in place to enforce strong typing within streams. Examining each of these areas in detail:

• Windows NT[™] was first introduced to the network computer community in 1993. In 1995, NT was the operating system of choice on 36% of the network computers sold. By 1997 that percentage had increased to almost half [Ref 12]. Also in 1997, the DOD, and the U.S. Navy in particular through the Information Technology Initiative for the 21st Century (IT-21), opted for the WindowsNT operating system as the standard toward which the military information system will march. The time when we must count on a Unix compatible system being available to supporting agencies for the installation of CAPS is coming to an end. Therefore, it is critical that a IntelPC/Windows compatible version of CAPS be developed. Although this is a very large task and planning for the eventuality cannot begin too soon, the increasing usage of Microsoft Foundation Classes and their Ada95 equivalents will enable significant use of preexisting software components for the creation of a new graphical interface.

• As a prototype increases in complexity, the number of data streams linking operators with the graph also increases. While operators can be aggregated into a composite operator and decomposed in a separate graph, there is no such option to use
with streams. Indeed, composite operators often require a large number of data streams to feed the sub-functions imbedded in the lower levels. The proliferation of streams makes the graph unwieldy to create and manipulate, and in no way enhances the designer's overall knowledge of the design. The creation of a composite stream seems to be the obvious solution. Incorporating all streams between two operators, a composite stream would be subject to decomposition like an operator and would contain it's own graph containing the two terminating operators and all the streams linking them. This addition to CAPS would vastly enhance the functionality of the system especially at the higher levels of abstraction of the prototype design.

• In the absence of a composite stream, redrawing a prototype containing numerous streams is an arduous and time-consuming task. Users of CAPS are sometimes forced to create designs in a separate drawing tool in order to organize the components prior to drawing them within the graphical editor simply to avoid the time required to redraw streams. An automatic redraw tool that can create a neat and understandable output with a minimum of crossed streams would be of inestimable value to the system.

• CAPS was designed and is still evolving using the outmoded Waterfall methodology. There is also an unfortunate shortage of design information on the CAPS system. The combination of these two factors makes understanding the workings of the environment extraordinarily difficult for a designer viewing the code for the first time. Because of this state of affairs, evolutions to CAPS are necessarily major events requiring a substantial commitment of time and money. The creation of a CAPS prototype by reverse engineering the CAPS system would greatly relieve this problem. Although a

63

substantial project, the ability of the CAPS system to decompose the problem at successively finer levels of detail would allow the project to be spread over several theses. Alternately, it could be done as an ongoing class project in an advanced software-engineering course. This would have the added advantage of teaching new students the utility of the tool while improving it. Upon completion, CAPS would benefit from all the evolutionary advantages that accrue to projects using rapid prototyping.

• The parsers installed in the Graphical Editor perform quite well in detecting errors with PSDL components. However, they do not check for type conformity within operators or streams. For example, if the designer decides that a stream type will be an integer, the system still allows the designer to enter 1.2, a float, as the initial state variable. This will undoubtedly result in problems when create problems during compilation of generated Ada code. Enforcing strong typing for data types within the PSDL structure should be a minor fix and will have a significant impact upon the robustness of the system.

• The parsers installed in the Graphical Editor check the designer's input and, if an error is present, show the point where the error occurs. This is a functional and adequate system. However, a much more robust method of ensuring that the designer provides the necessary information in the correct format is by using buttons to display templates for the various PSDL constructs. The designer can be prompted to choose from different variations of a component. This would result in a minimal amount of hand entered data, with a corresponding drop in the number of hand entered errors.

APPENDIX A: SELECTED SOURCE CODE (PSDL TYPE)

```
--Makefile for PSDL_TYPE
INCLUDE FLAGS = \setminus
  -I./GENERIC TYPES \
  -I./INSTANTIATIONS
GEN ADA = \setminus
  psdl io.adb \
  old psdl io.adb \
  psdl program pkg.adb \
  pre_expander_pkg.adb \
  psdl component pkg.adb \
  psdl concrete type pkg.adb \
  psdl_graph_pkg.adb \
  parser.adb
SPARSERS = \setminus
exception guard io.adb \
  exception io.adb \
  expression io.adb \
  op id io.adb \
  output guard_io.adb \
  parser.adb \
  timer op guard io.adb \
  type name io.adb
EXP YACC = \setminus
  expression_io_goto.ads \
  expression io shift reduce.ads \
  expression_io_tokens.ads
TYPE NAME YACC = \setminus
  type_name_io_goto.ads \
  type name io shift reduce.ads \
  type_name_io_tokens.ads
OP ID YACC = \setminus
  op_id_io_goto.ads \
  op id io_shift_reduce.ads \
  op id io_tokens.ads
OUTPUT GUARD YACC = \setminus
  output guard io goto.ads \
  output guard io shift reduce.ads \
```

```
output_guard_io_tokens.ads
```

```
EXCEPTION_GUARD YACC = \setminus
  exception_guard_io_goto.ads \
  exception_guard_io_shift reduce.ads \
  exception_guard_io_tokens.ads
EXCEPTION YACC = \setminus
  exception io goto.ads \
  exception io shift reduce.ads \
  exception io tokens.ads
TIMER OP GUARD YACC = \setminus
  timer_op_guard_io_goto.ads \
  timer_op_guard io shift reduce.ads \
  timer_op guard io tokens.ads
AYACC = \setminus
  parser goto.ads \
  parser shift reduce.ads \
  parser_tokens.ads \
  parser.g
AFLEX = \setminus
  parser_lex_dfa.ads \
  parser_lex_dfa.adb \
  parser_lex_io.ads \
  parser lex io.adb \
  parser lex.ads \
  parser lex.adb
all: $(GEN ADA) $(SPARSERS) $(AFLEX)
       (cd I* ; make gen)
       (cd GE* ; make gen)
      gnatmake -g -c $(INCLUDE FLAGS) psdl io.adb
gen: $(GEN ADA)
      (cd I* ; make gen)
      (cd GE* ; make gen)
parsers: $(SPARSERS) $(AFLEX)
psdl io.adb: psdl io.g
      gen < psdl io.g > psdl io.adb
pre_expander_pkg.adb: pre_expander_pkg.g
      gen < pre_expander_pkg.g > pre_expander_pkg.adb
old_psdl_io.adb: old_psdl_io.g
      gen < old_psdl_io.g > old_psdl_io.adb
psdl_program_pkg.adb: psdl_program_pkg.g
      gen < psdl_program_pkg.g > psdl_program_pkg.adb
```

psdl component_pkg.adb: psdl_component_pkg.g gen < psdl_component_pkg.g > psdl_component_pkg.adb psdl concrete type pkg.adb: psdl concrete type pkg.g gen < psdl_concrete_type_pkg.g > psdl_concrete_type_pkg.adb psdl_graph_pkg.adb: psdl_graph_pkg.g gen < psdl_graph_pkg.g > psdl_graph_pkg.adb \$(AFLEX): parser lex.l aflex -s parser lex.l gnatchop -w yylex.adb /bin/rm yylex.* \$(AYACC): parser.y ayacc parser.y debug =\> off verbose =\> on mv yyparse.adb parser.g /bin/rm yyparse.* \$(EXP YACC) expression_io.g: expression_io.y ayacc expression io.y debug =\> off verbose =\> on mv yyparse.adb expression io.g /bin/rm yyparse.* \$(TYPE NAME_YACC) type_name_io.g: type_name_io.y ayacc type name io.y debug = > off verbose = > on mv yyparse.adb type name io.g /bin/rm yyparse.* \$(OP_ID_YACC) op_id_io.g: op_id_io.y ayacc op_id_io.y debug =\> off verbose =\> on mv yyparse.adb op_id_io.g /bin/rm yyparse.* \$(OUTPUT_GUARD_YACC) output_guard io.g: output_guard_io.y ayacc output_guard_io.y debug =\> off verbose =\> on mv yyparse.adb output guard io.g /bin/rm yyparse.* \$(EXCEPTION_GUARD_YACC) exception_guard_io.g: exception_guard_io.y ayacc exception guard io.y debug = > off verbose = > on mv yyparse.adb exception guard io.g /bin/rm yyparse.* \$(EXCEPTION_YACC) exception_io.g: exception_io.y ayacc exception_io.y debug = > off verbose = > on mv yyparse.adb exception io.g /bin/rm yyparse.* \$(TIMER_OP_GUARD_YACC) timer_op_guard_io.g: timer_op_guard_io.y ayacc timer op guard io.y debug =\> off verbose =\> on mv yyparse.adb timer_op_guard_io.g /bin/rm yyparse.*

expression_io.adb: expression_io.g \$(AYACC) \$(AFLEX) gen < expression_io.g > expression_io.a gnatchop -w expression_io.a @sed -f expression io.sed expression io.adb > tmpfile @mv tmpfile expression io.adb /bin/rm expression io.a type_name_io.adb: type_name_io.g \$(AYACC) \$(AFLEX) gen < type_name_io.g > type_name_io.a gnatchop -w type name io.a @sed -f type name io.sed type name io.adb > tmpfile @mv tmpfile type_name io.adb /bin/rm type_name_io.a op_id_io.adb: op_id_io.g \$(AYACC) \$(AFLEX) gen < op id io.g > op_id_io.a gnatchop -w op id io.a @sed -f op_id_io.sed op id_io.adb > tmpfile @mv tmpfile op_id_io.adb /bin/rm op id io.a output_guard_io.adb: output_guard_io.g \$(AYACC) \$(AFLEX) gen < output_guard_io.g > output_guard io.a gnatchop -w output_guard io.a @sed -f output_guard_io.sed output_guard io.adb > tmpfile @mv tmpfile output_guard io.adb /bin/rm output guard io.a exception guard io.adb: exception guard_io.g \$(AYACC) \$(AFLEX) gen < exception guard io.g > exception guard io.a gnatchop -w exception_guard_io.a @sed -f exception_guard_io.sed exception_guard_io.adb > tmpfile @mv tmpfile exception_guard_io.adb /bin/rm exception_guard_io.a exception io.adb: exception_io.g \$(AYACC) \$(AFLEX) gen < exception io.g > exception io.a gnatchop -w exception io.a @sed -f exception io.sed exception io.adb > tmpfile @mv tmpfile exception io.adb /bin/rm exception_io.a timer_op_guard_io.adb: timer_op_guard_io.g \$(AYACC) \$(AFLEX) gen < timer_op_guard io.g > timer op_guard io.a gnatchop -w timer_op_guard_io.a @sed -f timer_op_guard_io.sed timer_op_guard_io.adb > tmpfile @mv tmpfile timer_op_guard_io.adb /bin/rm timer_op_guard_io.a parser.adb: parser.g \$(AYACC) \$(AFLEX) gen < parser.g > parser.a gnatchop -w parser.a

68

/bin/rm parser.a

```
# currently not used, old my pkg
# add to GEN ADA if used
psdl ops pkg.adb:
      gen < psdl_ops_pkg.g > psdl_ops_pkg.adb
ci: gen parsers
      ci files -tRCS/desc *.[Cgly] *.ad[sb] Makefile
      sleep 1
      touch *.adb
      touch *.ali *.o [IG]*/*.ali [IG]*/*.o
      (cd GENERIC TYPES ; make ci)
      (cd INSTANTIATIONS ; make ci)
test: all
      gnatmake -g -IINSTANTIATIONS -IGENERIC TYPES test.adb
      ./test > test.out
test2: all
      gnatmake -g -IINSTANTIATIONS -IGENERIC TYPES test2.adb
      ./test2 > test2.out
pre expander: all
      gnatmake -g -IINSTANTIATIONS -IGENERIC TYPES pre expander.adb
      ./pre_expander < autopilot.psdl > pre_expander.out
test3: all $(SPARSERS)
      gnatmake -g -IINSTANTIATIONS -IGENERIC TYPES test3.adb
      ./test3 > test3.out
test_text: all
      gnatmake -g -IINSTANTIATIONS -IGENERIC_TYPES test_text.adb
      ./test_text > test_text.out
```

APPENDIX B: SELECTED SOURCE CODE (PARSERS)

```
#
# Makefile for PSDL Parsers:
#
     . Expression
#
      . Initial Expression
#
      . Output Guard
#
     . Timer op Guard
     . Exception Guard
#
#
     . Exception List
#
     . Type Spec
#
      . Operator Spec
#
      . Opid
#
# CAPS Lab.
# Jun/1997
#
SOURCES = parse expression.y parse init expression.y
parse_output_guard.y parse_timer_op_guard.y parse_exception_guard.y
parse exception list.y parse type spec.y parse oper spec.y parse opid.y
parse_lex.l parser_pkg.adb parser pkg.ads
OBJ = yyparse.ads parse tokens.ads
COBJ = test.o test exp.o test init exp.o test output guard.o
test timer op quard.o test exception quard.o test exception list.o
test type spec.o test oper spec.o test opid.o
OBJ1 = parse expression goto.ads parse expression shift reduce.ads
OBJ2 = parse init expression goto.ads
parse init expression shift reduce.ads
OBJ3 = parse output guard goto.ads parse output guard shift reduce.ads
OBJ4 = parse_timer_op_guard_goto.ads
parse timer op guard shift reduce.ads
OBJ5 = parse exception guard goto.ads
parse_exception_guard_shift_reduce.ads
OBJ6 = parse exception list goto.ads
parse exception list shift reduce.ads
OBJ7 = parse_type_spec_goto.ads parse_type_spec_shift_reduce.ads
OBJ8 = parse_oper_spec_goto.ads parse_oper_spec_shift_reduce.ads
OBJ9 = parse opid goto.ads parse opid shift reduce.ads
SCRIPTS= Makefile parse.awk parse io.awk
PARSER= yyparse.adb yylex.adb
parsers: parser pkg.ali
test: parser pkg.ali ${COBJ}
      gnatbind -n parser pkg.ali
```

gnatlink -o test parser pkg.ali \${COBJ} parser_pkg.ali: parser_pkg.adb \${PARSER} \${SCRIPTS} \${SOURCES} \ \${OBJ1} \${OBJ2} \${OBJ3} \${OBJ4} \${OBJ5} \${OBJ6} \${OBJ7} \${OBJ8} \${OBJ9} \${OBJ} gnatmake parser pkg.adb \$(OBJ1): parse expression.y ayacc parse_expression.y debug =\> off verbose =\> on gnatchop -w yyparse.adb \$(OBJ2): parse_init expression.y ayacc parse_init_expression.y debug =\> off verbose =\> on gnatchop -w yyparse.adb \$(OBJ3): parse output guard.y ayacc parse_output guard.y debug = > off verbose = > on gnatchop -w yyparse.adb \$(OBJ4): parse timer op guard.y ayacc parse_timer_op_guard.y debug =\> off verbose =\> on gnatchop -w yyparse.adb \$(OBJ5): parse exception_guard.y ayacc parse_exception guard.y debug = > off verbose = > on gnatchop -w yyparse.adb \$(OBJ6): parse exception list.y ayacc parse_exception_list.y debug = > off verbose = > on gnatchop -w yyparse.adb \$(OBJ7): parse_type_spec.y ayacc parse_type_spec.y debug =\> off verbose =\> on gnatchop -w yyparse.adb \$(OBJ8): parse_oper_spec.y ayacc parse_oper_spec.y debug = > off verbose = > on gnatchop -w yyparse.adb \$(OBJ9): parse_opid.y ayacc parse_opid.y debug = > off verbose = > on gnatchop -w yyparse.adb @sed -f expression.sed parser_pkg-parse_expression.adb > tmpfile @mv tmpfile parser_pkg-parse_expression.adb @sed -f init_expression.sed parser pkg-parse_init_expression.adb > tmpfile @mv tmpfile parser_pkg-parse_init_expression.adb @sed -f output_guard.sed parser_pkg-parse_output_guard.adb > tmpfile @mv tmpfile parser pkg-parse_output_guard.adb @sed -f timer_op_guard.sed parser_pkg-parse_timer_op_guard.adb > tmpfile @mv tmpfile parser_pkg-parse_timer op guard.adb

```
@sed -f exception_guard.sed parser_pkg-parse_exception_guard.adb
> tmpfile
    @mv tmpfile parser_pkg-parse_exception_guard.adb
    @sed -f exception_list.sed parser_pkg-parse_exception_list.adb >
tmpfile
    @mv tmpfile parser_pkg-parse_exception_list.adb
    @sed -f type_spec.sed parser_pkg-parse_type_spec.adb > tmpfile
    @mv tmpfile parser_pkg-parse_oper_spec.adb
    @sed -f oper_spec.sed parser_pkg-parse_oper_spec.adb > tmpfile
    @mv tmpfile parser_pkg-parse_oper_spec.adb
    @sed -f opid.sed parser_pkg-parse_opid.adb > tmpfile
    @mv tmpfile parser_pkg-parse_opid.adb > tmpfile
    @mv tmpfile parser_pkg-parse_opid.adb
    @yes | rm tmpfile
    .IGNORE:
```

ci:

ci files -tRCS/desc \${SOURCES} \${SCRIPTS} \${OBJ}

yylex.ads yylex.adb: parse_lex.l parse.awk parse_io.awk
 aflex -s parse_lex.l
 nawk -f parse.awk yylex.adb > yylex.patched
 nawk -f parse_io.awk parse_lex_io.adb > temp
 mv temp parse_lex_io.adb
 gnatchop -w yylex.patched

-- Alex file for Parser --************** _____ _____ -- This file is the Aflex input file for PSDL grammar, -- and defines the lexical tokens for the ayacc psdl parser. -- Definitions of lexical classes [0-9] Digit Int {Digit}+ Letter [a-zA-Z] Alpha ({Letter}|{Digit}) Blank [$t\n$] -- Text is anything but '{' and '}' Text [^{}] -- StrLit is anything but '"' and '\' OR a '\' followed by '"' or a '\' StrLit [^"\\]|[\\]["\\] Quote ["]

응응

.

__****************

axioms AXIOMS	ł	return	(AXTOMS TOKEN) · }				
by{Blank}+all BY{Blank}+ALL	ł	return	(BY ALL TOKEN):				
by{Blank}+some BY{Blank}+SOME	ł	return	(BY SOME TOKEN):				
control CONTROL	ł	return	(CONTROL TOKEN) : }				
constraints CONSTRAINTS	ì	return	(CONSTRAINTS TOKEN) : }				
data DATA	i	return	(DATA TOKEN): }				
stream STREAM	ì	return	(STREAM TOKEN): }				
description DESCRIPTION	ì	return	(DESCRIPTION TOKEN) : }				
edge EDGE	Ì	return	(EDGE TOKEN); }				
end END	ł	return	(END TOKEN); }				
exceptions EXCEPTIONS	ł	return	(EXCEPTIONS TOKEN) : }				
exception EXCEPTION	ł	return	(EXCEPTION TOKEN); }				
finish FINISH	ł	return	(FINISH TOKEN); }				
within WITHIN	ł	return	(WITHIN TOKEN); }				
generic GENERIC	ł	return	(GENERIC TOKEN); }				
graph GRAPH	ł	return	(GRAPH TOKEN): }				
hours HOURS	ì	return	(HOURS TOKEN) : }				
if IF	ì	return	(IF TOKEN); }				
implementation IMPLEMENTATION	ł	return	(IMPLEMENTATION TOKEN):				
}ly INITIALLY	ł	return	(INITIALLY TOKEN): }				
input	{	return	(INPUT TOKEN):)				
ks KEYWORDS	ì	return	(KEYWORDS TOKEN) : }				
m MAXIMUM	;	return	(MAXIMUM TOKEN): }				
eon EXECUTION	Ì	return	(EXECUTION TOKEN): }				
tME	{	return	(TIME TOKEN): }				
rse RESPONSE	{	return	(RESPONSE TOKEN) :)				
microsec MICROSEC microseconds MICROSECONDS { return (MICROSEC TOKEN); }							

minimum | MINIMUM { return (MINIMUM TOKEN); } calling{Blank}+period|CALLING{Blank}+PERIOD { return (CALL PERIOD TOKEN); } min | MIN | minutes | MINUTES { return (MIN TOKEN); } { return (MS TOKEN); } ms |MS | milliseconds | MILLISECONDS operator | OPERATOR { return (OPERATOR TOKEN); } output | OUTPUT { return (OUTPUT TOKEN); } { return (PERIOD TOKEN); } period PERIOD { return (PROPERTY TOKEN); } property | PROPERTY required{Blank}+by|REQUIRED{Blank}+BY { return (REQ_BY_TOKEN); } reset{Blank}+timer|RESET{Blank}+TIMER { return (RESET TOKEN); } { return (SEC TOKEN); } sec|SEC|seconds|SECONDS specification | SPECIFICATION { return (SPECIFICATION TOKEN); } start{Blank}+timer|START{Blank}+TIMER { return (START_TOKEN); } { return (STATES TOKEN); } states | STATES stop{Blank}+timer|STOP{Blank}+TIMER { return (STOP TOKEN); } timer | TIMER { return (TIMER TOKEN); } { return (TRIGGERED TOKEN); } triggered | TRIGGERED { return (TYPE TOKEN); } type | TYPE vertex | VERTEX { return (VERTEX TOKEN); } "?" { return ('?'); } "{" { return ('{'); } "}" { return ('}'); } "!"|"@"|"#"|"\$"|"%"|"^"|"~"|"`" { return (ILLEGAL_TOKEN); } "(" { return ('('); } ")" { return (')'); } "[" { return ('['); } ייךיי { return (']'); } 11 <u>.</u> 11 { return (':'); } "," { return (','); } ท่า { return ('.'); } " | " { return ('|'); } "->" { return (ARROW); } "and" | "AND" { set token value((token category => psdl id cat, psdl id value => convert("AND"))); return (AND TOKEN); } "or"|"OR" { set token value((token category => psdl id cat, psdl id value => convert("OR"))); return (OR TOKEN); } "xor" | "XOR" { set_token_value((token_category => psdl_id_cat, psdl id value => convert("XOR"))); return (XOR TOKEN); } ">=" { set_token value((token category => psdl id cat, psdl id value => convert(">="))); return (GREATER_THAN_OR_EQUAL); } "<=" { set_token_value((token_category => psdl_id_cat, psdl id value => convert("<=")));</pre> return (LESS THAN OR EQUAL); } "/="|"~=" { set token value((token category => psdl id cat, psdl id value => convert("/="))); return (INEQUALITY); } "-" { set_token_value((token category => psdl id cat,

		<pre>psdl_id_value => convert("=")));</pre>		
		return ('='); }		
"+"	{	<pre>set_token_value((token_category => psdl id cat,</pre>		
		<pre>psdl_id_value => convert("+")));</pre>		
		return ('+'); }		
<u>n_</u> n	{	<pre>set_token_value((token_category => psdl_id cat,</pre>		
		<pre>psdl_id_value => convert("-")));</pre>		
		return ('-'); }		
u x n	{	<pre>set_token_value((token_category => psdl_id_cat,</pre>		
		<pre>psdl_id_value => convert("*")));</pre>		
		return ('*'); }		
	{	<pre>set_token_value((token_category => psdl_id_cat,</pre>		
		<pre>psdi_id_value => convert("/")));</pre>		
" C "	ſ	set taken value (/taken astanana) it is		
u .	ı	sec_coken_value((coken_category => psdi_id_cat,		
		return ('c') ·)		
">"	ŧ	set token value/(token gategory -> nadl id at		
	·	psdl id value => convert(">"))).		
		return ('>'): }		
"<"	{	set token value((token category => psdl id cat		
		<pre>psdl id value => convert("<")));</pre>		
		return ('<'); }		
"mod" "MOD"	{	<pre>set_token_value((token category => psdl id cat,</pre>		
		<pre>psdl_id_value => convert("MOD")));</pre>		
		return (MOD_TOKEN); }		
"rem" "REM"	{	<pre>set_token_value((token_category => psdl_id_cat,</pre>		
		<pre>psdl_id_value => convert("REM")));</pre>		
****	r	return (REM_TOKEN); }		
	ı	set_token_value((token_category => psdl_id_cat,		
		return (EXP TOKEN) · }		
"abs" "ABS"	ł	set token value((token category => psdl id cat		
	•	<pre>psdl id value => convert("ABS"))):</pre>		
		return (ABS TOKEN); }		
"not" "NOT"	{	<pre>set_token_value((token category => psdl id cat,</pre>		
•		<pre>psdl_id_value => convert("NOT")));</pre>		
		return (NOT_TOKEN); }		
true TRUE	{	<pre>set_token_value((token_category => psdl_id_cat,</pre>		
		<pre>psdl_id_value => convert("TRUE")));</pre>		
falcolENICE	,	return (TRUE); }		
Tarsellarse	t	<pre>set_token_value((token_category => psdl_id_cat,</pre>		
		<pre>psdl_id_value => convert("FALSE")));</pre>		
		Tecurn (EMDE); }		
{Letter}{Alpha}*		{ set token value((token category => nsd) id cat		
		psdl id value => convert(vvtext)));		
		return (IDENTIFIER); }		
{Quote}{StrLit}*{Quo	te	<pre>{ set_token_value((token_category => text cat,</pre>		
		<pre>text_value => convert(yytext));</pre>		
		<pre>return (STRING_LITERAL); }</pre>		

{Int}	{	<pre>set_token_value((token_category => integer_cat, integer_value => string_to_integer(yytext))); return (INTEGER_LITERAL); }</pre>
{Int}"."{Int}		<pre>{ set_token_value((token_category => text_cat,</pre>
"{"{Text}*"}"		<pre>{ set_token_value((token_category => text_cat,</pre>
[\n]		<pre>{ increment_line_number; }</pre>
[\t]		{ null; } ignore spaces and tabs
88		

```
with parser_tokens; use parser_tokens;
with psdl_concrete_type_pkg; use psdl_concrete_type_pkg;
with psdl_id_pkg; use psdl_id_pkg;
with text_io;
package parser_lex is
-- Global variables used by the lexical analyser
line_number: positive := 1;
num_errors: natural := 0;
```

```
-- Initialization procedure to allow multiple calls from the parser procedure initialize_yylex;
```

```
-- Line numbers to be used for error messages during parsing function current_line return positive;
```

```
-- Lexical analyzer function generated by aflex function yylex return token;
```

```
-- Procedure that increments line_numbers procedure increment_line_number;
```

```
end parser_lex;
```

```
with text_io; use text_io;
package body parser_lex is
-- Externally visible subprograms
procedure initialize_yylex is
begin
line_number := 1; -- reset line_number to 1
yy_init := true; -- tell yylex to reinitialize itself
end initialize_yylex;
```

function current_line return positive is

```
begin
      return(line number);
    end current_line;
  procedure increment line number is
    begin
      line number := line_number + 1;
    end increment_line number;
  -- Functions and subprograms used in actions associated with tokens
  package nat_io is new integer_io(natural);
  function string_to_integer (digit_string: string) return natural is
    digit, value : natural := 0;
  begin
    for i in 1 .. digit_string'length loop
      case digit string(\overline{i}) is
        when '0' => digit := 0;
        when '1' => digit := 1;
        when '2' => digit := 2;
        when '3' => digit := 3;
        when '4' => digit := 4;
        when '5' => digit := 5;
        when '6' => digit := 6;
        when '7' => digit := 7;
        when '8' => digit := 8;
        when '9' => digit := 9;
        when others => return value;
      end case;
      value := (10 * value) + digit;
    end loop;
    return value;
  end string_to_integer;
  procedure set_token_value(value: yystype) is
    -- Set the value of a token, like "$$ := value" in ayacc.
 begin
    yylval := value;
  end set token value;
 -- The generated yylex function goes here.
##
```

end parser lex;

```
__********************
-- Avacc file for Parser
---*******************
-- $Header: /work/sde/PSDL TYPE.NEW/RCS/parser.y,
-- v 1.19 1997/08/12 21:35:47 sde Exp sde $
-- This file is the ayacc source file for PSDL parser.
-- token declarations section
%token '?' '{' '}' ILLEGAL TOKEN '(' ')' ',' '[' ']' ':' '.' '|'
%token ARROW
%token TRUE FALSE
%token AXIOMS TOKEN
%token BY ALL TOKEN REQ BY TOKEN BY SOME TOKEN
%token CALL PERIOD TOKEN CONTROL TOKEN
%token CONSTRAINTS TOKEN
%token DATA TOKEN DESCRIPTION TOKEN
%token EDGE_TOKEN END_TOKEN EXCEPTIONS_TOKEN
%token EXCEPTION TOKEN EXECUTION TOKEN
%token FINISH TOKEN
%token GENERIC TOKEN GRAPH TOKEN
%token HOURS TOKEN
%token IF TOKEN IMPLEMENTATION TOKEN
%token INITIALLY TOKEN INPUT TOKEN
%token KEYWORDS TOKEN
%token MAXIMUM TOKEN MINIMUM TOKEN
%token MICROSEC_TOKEN
%token MIN_TOKEN MS_TOKEN MOD_TOKEN
%token NOT TOKEN
%token OPERATOR TOKEN OR TOKEN OUTPUT TOKEN
%token PERIOD TOKEN PROPERTY TOKEN
%token RESET TOKEN RESPONSE TOKEN
%token SEC TOKEN SPECIFICATION TOKEN
%token START_TOKEN STATES TOKEN STOP TOKEN
%token STREAM TOKEN
%token TIME TOKEN
%token TIMER_TOKEN TRIGGERED TOKEN TYPE TOKEN
Stoken VERTEX TOKEN
%token WITHIN TOKEN
%token IDENTIFIER
%token INTEGER_LITERAL REAL_LITERAL
%token STRING LITERAL
Stoken TEXT TOKEN
-- operator precedences
-- left means group and evaluate from the left
%left AND TOKEN OR TOKEN XOR TOKEN LOGICAL OPERATOR
$left '<' '>' '=' GREATER THAN OR EQUAL LESS THAN OR EQUAL INEQUALITY
RELATIONAL OPERATOR
```

```
%left '+' '-' '&' BINARY ADDING OPERATOR
 %left UNARY ADDING OPERATOR
 %left '*' '/' MOD_TOKEN REM TOKEN MULTIPLYING_OPERATOR
 %left EXP_TOKEN ABS_TOKEN NOT_TOKEN HIGHEST_PRECEDENCE_OPERATOR
%start start_symbol
-- this is an artificial start symbol, for initialization
-- declaration of the value type for the parser stack.
%with psdl_concrete_type_pkg, expression_pkg, psdl_id_pkg;
%use psdl_concrete_type_pkg, expression_pkg, psdl_id_pkg;
 {
    type token category_type is (integer_cat,
                                  text_cat,
                                  psdl id cat,
                                  psdl id sequence cat,
                                  op id cat,
                                  operator name cat,
                                  opt arg cat,
                                  type_name_cat,
                          type_decl cat,
                                  timer op id cat,
                                  expression cat,
                                  expression seq cat,
                                  property map cat,
                                  no value cat);
    type yystype (token_category: token_category_type := no value cat)
is
      record
         case token category is
           -- lexical token attributes:
           when integer cat =>
             integer value: integer;
           when text cat =>
             text value: text;
           -- grammar psdl id attributes:
           when psdl id cat =>
             psdl id value: psdl id;
           when psdl id sequence cat =>
             psdl_id_sequence_value: psdl_id_sequence;
           when op id cat =>
             op_id_value: op id;
           when operator name cat =>
             type_name_part, op_name_part: psdl_id;
           when opt_arg_cat =>
             input_value, output_value: psdl id sequence;
           when type name cat =>
             type_name_value: type_name;
           when type decl cat =>
             type_decl_value: type_declaration;
```

```
when timer op id cat =>
             timer_op_id_value: timer_op_id;
           when expression cat =>
             expression value: expression;
           when expression seq cat =>
             expression_seq_value: expression_sequence;
           when property map cat =>
             property map value: init map;
           when no_value_cat => null;
         end case;
       end record;
}
응응
start symbol
            { the program := empty psdl program; }
        :
          psdl
        ;
        : psdl component
psdl
            { if member(name(the component), the program)
              then yyerror("Component redefined: "
                            & convert(name(the_component)));
              else bind(name(the_component),
                         the component, the program);
              end if; }
        ł
        ;
component
        : data_type
        | operator
        ;
data_type
        : TYPE TOKEN IDENTIFIER
            { the_operation_map := empty_operation_map;
              is_specification := true; }
          type_spec
            { is specification := false; }
          type_impl
            { -- Construct the psdl type using global variables.
              build_psdl_type($2.psdl_id_value,
                               the ada name,
                               the imp lang,
                               the model,
                               the data structure,
                               the_operation map,
                               the type gen par,
                               the keywords,
                               the_description,
                               the axioms,
                               is_atomic_type,
```

```
81
```

;

```
type_spec
          SPECIFICATION_TOKEN optional_generic_param optional_type_decl
          op_spec_list functionality END TOKEN
        ï
optional generic param
        : GENERIC TOKEN
            { the_type_decl := empty_type_declaration; }
          list_of_type_decl
            { the_type_gen_par := the_type_decl; }
            { the type gen par := empty type declaration; }
optional_type_decl
            { the_type_decl := empty_type_declaration; }
        :
          list of type decl
            { the model := the type decl; }
            { the model := empty_type_declaration; }
        T
        ;
op_spec list
        : op_spec_list OPERATOR_TOKEN IDENTIFIER operator spec
             { build_psdl_operator($3.psdl id value,
                                    to ada id($3.psdl id value),
                                   the imp lang,
                                    the gen par,
                                   the_gen_par_rb,
                                   the keywords,
                                   the description,
                                   the axioms,
                                   the input,
                                   the output,
                                   the state,
                                   the initial expression map,
                                   the exceptions,
                                   the specified met,
                                   the input rb,
                                   the output rb,
                                   the_state_rb,
                                   the exception rb,
                                   the spec_met_rb,
                                   is_atomic => true,
                                   the_opr => the operator);
               bind_operation ($3.psdl id value,
                               the operator,
                               the_operation_map); }
        1
        ;
```

```
operator
        : OPERATOR TOKEN IDENTIFIER
            { is_specification := true; }
          operator_spec
            { is specification := false; }
          operator impl
            { -- construct the psdl operator using the global variables
             build psdl operator($2.psdl id value,
                                  the_ada_name,
                                   the_imp_lang,
                                   the_gen_par,
                                  the gen par rb,
                                   the keywords,
                                   the description,
                                  the axioms,
                                  the input,
                                  the_output,
                                  the state,
                                  the initial expression map,
                                  the exceptions,
                                  the specified met,
                                  the input rb,
                                  the output rb,
                                  the state rb,
                                  the exception rb,
                                  the spec met rb,
                                  the graph,
                                  the streams,
                                  the timers,
                                  the trigger map,
                                  the_exec_guard,
                                  the_out_guard,
                                  the_excep_trigger,
                                  the timer_op,
                                  the per,
                                  the fw;
                                  the_mcp,
                                  the mrt,
                                  the impl desc,
                                  the eg rb,
                                  the per rb,
                                  the fw rb,
                                  the mcp rb,
                                  the mrt rb,
                                  the o rb,
                                  the e rb,
                                  the reset rb,
                                  the_start_rb,
                                  the_stop_rb,
                                  is atomic operator,
                                  the_component); }
        ;
```

```
operator spec
        : SPECIFICATION TOKEN
             { -- Initialize the variables used to
              -- build an operator spec.
              the_gen_par := empty_type_declaration;
              the gen par rb := empty;
              the input rb := empty;
              the output rb := empty;
              the state rb := empty;
              the exception rb := empty;
              the spec met rb := empty;
              the_input := empty_type_declaration;
              the output := empty type_declaration;
              the_state := empty_type_declaration;
              expression_sequence_pkg.empty(the_init_exp_seq);
              the_initial_expression_map := empty_init_map;
            the_exceptions := empty;
              the_specified_met := undefined time; }
          interface
            { bind_initial_state(the_state, the_init_exp_seq,
                                  the_initial_expression_map); }
          functionality END TOKEN
        ;
interface
        : interface attribute reqmts trace
            { bind_spec_rb(the attribute type,
                            $2.psdl id sequence value,
                            $3.psdl_id_sequence_value); }
        1
        ;
attribute
        : GENERIC TOKEN
            { the_type_decl := the gen par; }
          list of type decl
            \{ $$ := $3;
              the_gen_par := the type decl;
              the_attribute_type := gen_par; }
        INPUT TOKEN
            { the type decl := the input; }
          list_of_type_decl
            { $$ := $3;
              the input := the_type_decl;
              the attribute type := input; }
        | OUTPUT TOKEN
            { the_type_decl := the_output; }
          list_of_type_decl
{ $$ := $3;
              the_output := the_type_decl;
              the_attribute_type := output; }
        STATES TOKEN
            { the_states_token_line := current line;
            -- For error messages.
```

```
the states token := convert(yytext);
              -- For error messages.
              the type decl := the state; }
          list of type decl
            { the state := the type decl; }
          INITIALLY_TOKEN initial_expression_list
            { $$ := $3;
              expression sequence pkg.append(the init exp seq,
                                              $6.expression seq value,
                                              the init exp seq);
              the attribute type := state; }
        | EXCEPTIONS TOKEN id list
            { $$ := $2;
              psdl_id_set_pkg.union(the_exceptions,
                                     to set($2.psdl id sequence value),
                                    the exceptions );
              the attribute type := exc; }
        MAXIMUM TOKEN EXECUTION TOKEN TIME TOKEN time
            { $$ := ( token category => psdl id sequence cat,
                      psdl id sequence value => empty );
              if the specified met = undefined time or
                 $4. integer value < the specified met
              then the specified met := $4.integer value;
              end if;
              the attribute type := met; }
              -- Time is converted into millisec .
        ;
  -- Initialization of the type decl is
  -- done by the callers of this rule.
list_of_type_decl
        : list_of_type_decl ',' type_decl
            { $$ := ( token_category => psdl_id_sequence_cat,
                      psdl_id_sequence_value =>
                 psdl id sequence pkg.append($1.psdl id sequence value,
                      $3.psdl id sequence value )); }
        | type decl
           { $$$ := $1; }
type decl
        : id list ':' type_name
            \{ $$ := $1;
              bind_type_declaration($1.psdl id sequence value,
                                    $3.type name value,
                            the type decl); }
        ;
type name
        : IDENTIFIER
            { -- Save the previous value of the type decl.
            -- Needed because the list of type decl below
            -- might contain nested type declarations.
            $$ := (token_category => type_decl_cat,
```

```
type_decl_value => the_type_decl);
               the_type_decl := empty_type_declaration; }
           '[' list_of_type_decl ']'
             { the type name :=
                 create(name => $1.psdl_id value,
                        formals => $4.psdl id sequence value,
                        gen_par => the_type decl);
             -- Now restore the previous value saved above.
             the_type_decl := $2.type decl value;
               $$ := (token_category => type_name_cat,
                      type_name_value => the type_name); }
        | IDENTIFIER
             { the_type_name :=
                create(name => $1.psdl id value,
                        formals => psdl_id_sequence_pkg.empty,
                        gen_par => empty_type_declaration);
              $$ := (token_category => type_name_cat,
                      type name value => the_type_name); }
        ř
id list
        : id list ',' IDENTIFIER
            { $$ := ( token_category => psdl_id_sequence_cat,
                      psdl_id sequence value =>
                         add($3.psdl id value,
                             $1.psdl_id_sequence_value )); }
        | IDENTIFIER
            { $$ := ( token_category => psdl_id_sequence_cat,
                        psdl id sequence_value => add($1.psdl_id_value,
                        empty) ); }
        ;
reqmts_trace
        : REQ_BY TOKEN id list
          { $$ := $2; }
        Ŧ
          { $$ := ( token_category => psdl_id_sequence_cat,
                    psdl_id_sequence_value => empty ); }
        ;
functionality
        : keywords informal desc formal desc
        ;
keywords
        : KEYWORDS_TOKEN id list
            { the_keywords := to_set($2.psdl_id_sequence_value); }
        L
            { the keywords := empty; }
        ;
informal desc
        : DESCRIPTION TOKEN TEXT TOKEN
            { if is specification then
                 the_description := remove braces($2.text_value);
```

```
else the impl desc := remove braces($2.text value);
              end if; }
        I
            { if is specification then
                  the_description := empty;
              else the impl desc := empty;
              end if; }
        ;
formal desc
        : axioms TOKEN TEXT TOKEN
            { the axioms:= remove braces($2.text value); }
            { the axioms:= empty; }
        ;
type impl
        : IMPLEMENTATION_TOKEN IDENTIFIER IDENTIFIER END TOKEN
            { is_atomic_type := true;
              the_imp_lang := $2.psdl_id_value;
              the ada name := to ada id($3.psdl id value); }
        | IMPLEMENTATION_TOKEN type_name op_impl_list END_TOKEN
            { is atomic type := false;
              the data_structure := $2.type_name_value; }
        ;
op impl list
        : op impl list OPERATOR TOKEN IDENTIFIER operator impl
            { -- add implementation part to the operator in the
operation map
              add op impl to op map($3.psdl id value,
                                     the ada name,
                                     is atomic operator,
                                     the_operation_map,
                                     the_graph,
                                     the streams,
                                     the timers,
                                     the trigger map,
                                     the exec_guard,
                                     the out guard,
                                     the_excep_trigger,
                                     the_timer_op,
                                     the_per,
                                     the fw,
                                     the mcp,
                                     the mrt,
                                     the impl desc ); }
        1
        ;
operator impl
        : IMPLEMENTATION TOKEN IDENTIFIER IDENTIFIER END TOKEN
            { is atomic operator := true;
              the_imp_lang := $2.psdl_id_value;
```

```
87
```

```
the_ada_name := to_ada_id($3.psdl_id_value); }
         | IMPLEMENTATION_TOKEN psdl_impl END_TOKEN
             { is_atomic_operator := false; }
psdl impl
         : data flow_diagram streams timers control_constraints
informal desc
         :
data flow diagram
        : { the_graph := empty_psdl graph; }
          GRAPH_TOKEN vertex_list edge_list
         ;
                 -- Time is the maximum execution time.
vertex list
        : vertex_list VERTEX_TOKEN op_id optional_time graph_properties
            { the_graph := psdl_graph_pkg.add_vertex($3.op_id_value,
                                                       the graph,
                                                       $4.integer value,
$5.property_map_value); }
        ;
                -- Time is the latency.
edge list
        : edge_list EDGE_TOKEN IDENTIFIER
          optional_time op_id ARROW op_id graph_properties
            { the_graph := psdl_graph_pkg.add_edge($5.op_id_value,
                                                    $7.op_id_value,
                                                    $3.psdl_id_value,
                                                    the graph,
                                                    $4. integer_value,
$8.property_map_value); }
        1
        ;
graph_properties
        : graph_properties PROPERTY_TOKEN IDENTIFIER '=' expression
            { bind($3.psdl_id_value, $5.expression_value,
             $1.property map value);
              $$ := ( token_category => property map cat,
                      property_map_value => $1.property_map_value ); }
        1
            { $$ := ( token_category => property_map_cat,
                      property_map_value => empty_init_map ); }
        ;
op id
        : operator name opt arg
            { $$ := ( token_category => op_id_cat,
```

```
op_id_value =>
                      ( type name => $1.type name part,
                        operation name => $1.op name part,
                        inputs => $2.input value,
                        outputs => $2.output_value )); }
        ;
operator name
        : IDENTIFIER '.' IDENTIFIER
            { $$ := ( token_category => operator_name_cat,
                      type_name_part => $1.psdl_id_value,
                      op_name_part => $3.psdl_id_value ); }
        | IDENTIFIER
            { $$ := ( token category => operator name cat,
                      type_name_part => empty,
                      op_name_part => $1.psdl_id_value ); }
        ;
opt_arg
        : '(' optional id list '|' optional_id_list ')'
            { $$ := ( token category => opt arg cat,
                      input_value => $2.psdl_id_sequence_value,
                      output value => $4.psdl id sequence value ); }
            { $$ := ( token category => opt arg cat,
        input_value => empty,
                      output value => empty ); }
        ;
optional id list
        : id_list { $$ := $1; }
        | { $$ := ( token_category => psdl id sequence cat,
                   psdl id sequence value => empty ); }
        ï
optional_time
        : ':' time
            { $$ := (token_category => integer_cat,
                     integer_value => $2.integer_value); }
            { $$:= (token_category => integer_cat,
        L
                    integer_value => undefined_time); }
        ;
streams
        : DATA TOKEN STREAM TOKEN
            { the_type_decl := empty_type_declaration; }
         list of type decl
            { the streams := the type decl; }
        1
            { the_streams := empty type declaration; }
        ;
              _____
       -- The order of id's is not important, so
```

```
-- we use psdl id set as the data structure
        -- to store the timers.
         _____
timers
        : TIMER TOKEN id list
            { the timers := to_set($2.psdl_id_sequence_value); }
            { the timers := empty; }
        L
        ;
control constraints
        : CONTROL TOKEN CONSTRAINTS TOKEN
            { the_trigger_map := empty_trigger_map;
              the_per := empty_timing_map;
              the fw := empty timing map;
              the_mcp := empty timing map;
              the_mrt := empty timing map;
              the_exec_guard := empty_exec_guard_map;
              the out guard := empty_out guard map;
              the_excep_trigger := empty_excep_trigger_map;
              the_timer_op := empty_timer_op map;
              the eg rb := empty;
              the per rb := empty;
              the fw rb := empty;
              the_mcp_rb := empty;
              the mrt rb := empty;
              the o_srbm := empty;
              the e_srbm := empty;
              the reset srbm := empty;
              the_start_srbm := empty;
              the stop_srbm := empty;
              the_o_rb := empty;
              the e rb := empty;
              the_reset_rb := empty;
              the start rb := empty;
              the_stop_rb := empty; }
          constraints
        ;
constraints
        : constraints OPERATOR TOKEN op id
            { the operator id := $3.op_id_value;
              the_timer_op_set := timer op_set pkg.empty; }
          opt_trigger opt_period opt_finish_within
          opt_mcp opt_mrt constraint options
            { bind(the_operator_id, the_o_srbm, the o rb);
              bind(the_operator_id, the_e_srbm, the_e_rb);
              bind(the_operator_id, the_reset_srbm, the reset rb);
              bind(the_operator_id, the start srbm, the start rb);
              bind(the_operator_id, the_stop_srbm, the_stop_rb); }
        1
        ;
constraint options
        : constraint_options OUTPUT TOKEN
```

```
id list IF TOKEN expression regmts trace
            { the output id.op := the operator id;
              for id: psdl id in
                  psdl id sequence pkg.scan($3.psdl id sequence value)
              loop
                  the output id.stream := id;
                  bind(the output id, $5.expression value,
                      the out guard);
                  bind(id, $6.psdl id sequence value, the o srbm);
              end loop; }
        | constraint options EXCEPTION TOKEN IDENTIFIER
          opt if predicate regmts trace
            { the_excep_id.op := the_operator_id;
              the excep id.excep := $3.psdl id value;
              bind (the excep id, $4.expression value,
                  the_excep_trigger);
              bind($3.psdl id value, $5.psdl id sequence value,
                  the e srbm); }
        | constraint options timer op IDENTIFIER
          opt if predicate reqmts trace
            { the timer op record.op id := $2.timer op id value;
              the timer op record.timer id := $3.psdl id value;
              the timer op record.guard := $4.expression value;
              timer_op_set_pkg.add (the_timer_op_record,
                                     the_timer_op_set);
              bind(the_operator id, the timer op set, the timer op);
              case the timer op record.op id is
                when t reset =>
                  bind($3.psdl_id_value, $5.psdl_id_sequence_value,
the reset srbm);
                when t start =>
                  bind($3.psdl id value, $5.psdl id sequence value,
the_start_srbm);
                when t stop =>
                  bind($3.psdl_id_value, $5.psdl_id_sequence_value,
the_stop_srbm);
                when t none => null;
                  -- This case is impossible but the compiler can't
recognize that.
              end case; }
        1
        ;
opt_trigger
        : TRIGGERED_TOKEN trigger opt if predicate reqmts trace
          { bind(the_operator id, $3.expression value, the exec guard);
            bind(the_operator_id, $4.psdl id sequence value,
the eg rb); }
        1
        ;
trigger
```

: BY ALL TOKEN id list

```
{ the_trigger.tt := by all;
               the_trigger.streams := to_set($2.psdl_id_sequence_value);
               bind(the_operator_id, the_trigger, the_trigger_map); }
         BY SOME TOKEN id list
             { the trigger. tt := by some;
               the_trigger.streams := to set($2.psdl_id_sequence_value);
               bind(the_operator id, the_trigger, the_trigger_map); }
             { the trigger.tt := by_none;
         L
               the trigger.streams := empty;
              bind(the_operator_id, the_trigger, the_trigger_map); }
        ;
opt period
         : PERIOD_TOKEN time reqmts_trace
             { bind(the_operator_id, $2.integer value, the per);
              bind(the_operator_id, $3.psdl id sequence value,
the per rb); }
         L
        ;
opt_finish_within
        : FINISH_TOKEN WITHIN_TOKEN time reqmts_trace
             { bind(the_operator_id, $3.integer_value, the_fw);
              bind(the_operator_id, $4.psdl id sequence value,
              the fw rb); }
        1
        ;
opt mcp
        : MINIMUM_TOKEN CALL PERIOD_TOKEN time reqmts trace
             { bind(the_operator_id, $3.integer value, the_mcp);
              bind(the_operator_id, $4.psdl id sequence value,
              the_mcp_rb); }
        1
        ;
opt_mrt
        : max resp_time time reqmts_trace
            { bind(the_operator_id, $2.integer_value, the_mrt);
              bind(the_operator id, $3.psdl id sequence value,
              the_mrt rb); }
        L
        ;
max resp time
        : MAXIMUM TOKEN RESPONSE_TOKEN TIME_TOKEN
        ;
timer op
        : RESET TOKEN
            { $$ := (token_category => timer_op_id_cat,
                     timer_op_id_value => t_reset); }
        | START TOKEN
```

```
{ $$ := (token category => timer op id cat,
                   timer op id value => t start); }
       | STOP TOKEN
           { $$ := (token category => timer_op_id_cat,
                   timer op id value => t stop); }
       ;
opt_if_predicate
       : IF_TOKEN expression
         { $$ := (token category => expression cat,
                 expression value => $2.expression value); }
       { $$ := (token_category => expression_cat,
                 expression_value => true_expression); }
       î
           -- The expression sequence
      -- is used by procedure bind initial state together with
      -- the states map to construct the init map.
      _____
                                                    _____
initial expression list
       : initial expression list ',' initial expression
           { $$ := (token category => expression seq cat,
                   expression seq value =>
                     expression sequence pkg.add($3.expression value,
$1.expression seq value )); }
       | initial expression
           { $$ := (token_category => expression seq_cat,
                   expression_seq_value =>
                     expression_sequence_pkg.add($1.expression_value,
                                              empty exp seq )); }
       ;
      -- There is one and only one initial state(initial expression)
      -- for each state variable. This production returns one
      -- expression to the parent rule corresponding to one state.
      -- This is done by using the internal stack ($$ convention).
      _____
initial expression
       : TRUE
           { $$ := (token category => expression cat,
                   expression value => true expression); }
       FALSE
           { $$ := (token_category => expression_cat,
                   expression_value => false expression); }
       | INTEGER_LITERAL
```

```
{ $$ := (token category => expression cat,
              expression value =>
              create integer_literal($1.integer_value)); }
| REAL LITERAL
    { \bar{\$}\$ := (token_category => expression_cat,
              expression_value =>
              create real literal($1.text_value)); }
| STRING LITERAL
    { $$ := (token category => expression cat,
             expression value =>
              create_string literal($1.text value)); }
| IDENTIFIER
    { $$ := (token category => expression cat,
             expression value =>
             create identifier($1.psdl id_value)); }
| type name '.' IDENTIFIER
    { $$ := (token_category => expression_cat,
             expression value =>
             create function_call($1.type_name_value,
                                   psdl id($3.psdl id value),
                                   empty exp seq)); }
| type name '.' IDENTIFIER '(' initial_expression_list ')'
    { $$ := (token_category => expression_cat,
             expression value =>
             create_function call($1.type name_value,
                                   psdl id($3.psdl id value),
                                   $5.expression seq value));}
| '(' initial_expression ')'
    { $$ := (token_category => expression_cat,
             expression_value => $2.expression value) ;
| initial_expression log_op initial_expression
  %prec logical_operator
    { $$ := (token category => expression cat,
             expression value =>
             create_binary_op ($1.expression value,
                                $2.psdl id value,
                                $3.expression value));
| initial_expression rel_op initial_expression
  %prec relational_operator
    { $$ := (token_category => expression_cat,
             expression value =>
             create_binary_op ($1.expression value,
                                $2.psdl id value,
                                $3.expression value));
'-' initial_expression
   %prec unary adding operator
    { $$ := (token_category => expression_cat,
             expression value =>
               create_unary_op(convert("-"),
                               $2.expression value )); }
| '+' initial expression
```

}

}

```
%prec unary_adding_operator
            { $$ := (token category => expression cat,
                     expression_value =>
                        create_unary_op(convert("+"),
                                        $2.expression value ));}
        | initial_expression bin_add_op_initial_expression
          %prec binary_adding_operator
            { $$ := (token category => expression cat,
                     expression_value =>
                     create_binary_op ($1.expression value,
                                        $2.psdl id value,
                                        $3.expression value)); }
        | initial_expression bin_mul_op initial_expression
          %prec multiplying_operator
            { $$ := (token category => expression cat,
                     expression value =>
                     create_binary_op ($1.expression_value,
                                        $2.psdl id value,
                                        $3.expression value)); }
        | initial expression EXP TOKEN initial expression
          %prec highest_precedence_operator
            { $$ := (token category => expression cat,
                     expression value =>
                     create_binary_op ($1.expression_value,
                                        convert("**"),
                                        $3.expression value));
        | NOT_TOKEN initial_expression
          %prec highest_precedence_operator
            { $$ := (token category => expression cat,
                     expression_value =>
                       create_unary_op(convert("NOT"),
                                        $2.expression_value )); }
        | ABS_TOKEN initial expression
          %prec highest precedence operator
            { $$ := (token_category => expression_cat,
                     expression_value =>
                       create_unary_op(convert("ABS"),
                                        $2.expression_value )); }
        | '?' { $$ := (token_category => expression_cat,
                     expression value => undefined expression ); }
        ;
log op
        : AND TOKEN
            { $$ := (token_category => psdl_id_cat,
                     psdl id value => convert("AND") ); }
        | OR TOKEN
            { $$ := (token_category => psdl_id_cat,
                     psdl id value => convert("OR") ); }
        | XOR TOKEN
```

{ \$\$:= (token_category => psdl id cat,

;

}

psdl id value => convert("XOR")); }

rel_op : '<' { \$\$:= (token_category => psdl_id_cat, psdl_id_value => convert("<")); }</pre> 1 '>' { \$\$:= (token_category => psdl_id_cat, psdl_id_value => convert(">")); } | '=' { \$\$:= (token_category => psdl_id_cat, psdl_id_value => convert("=")); } GREATER THAN_OR_EQUAL { \$\$:= (token_category => psdl_id_cat, psdl_id_value => convert(">=")); } | LESS THAN OR EQUAL { \$\$:= (token_category => psdl_id_cat, psdl_id_value => convert("<=")); }</pre> | INEQUALITY { \$\$:= (token_category => psdl_id_cat, psdl_id_value => convert("/=")); } ; bin_add_op : '+' { \$\$:= (token_category => psdl_id_cat, psdl_id_value => convert("+")); } 1 '-' { \$\$:= (token_category => psdl_id_cat, psdl_id_value => convert("-")); } 1 '&' { \$\$:= (token_category => psdl_id_cat, psdl_id_value => convert("&")); } ; bin_mul_op : '*' { \$\$:= (token_category => psdl_id_cat, psdl_id_value => convert("*")); } 1 1/1 { \$\$:= (token_category => psdl_id_cat, psdl_id_value => convert("/")); } | MOD TOKEN { \$\$:= (token_category => psdl_id_cat, psdl_id_value => convert("MOD")); } | REM TOKEN { \$\$:= (token_category => psdl_id_cat, psdl_id_value => convert("REM")); } ;

time

```
| time number MS TOKEN
            { $$ := (token category => integer cat,
                     integer value => $1.integer value); }
        | time number SEC TOKEN
            { $$ := (token category => integer cat,
                     integer value => $1.integer value * 1000); }
        | time number MIN TOKEN
            { $$ := (token category => integer cat,
                     integer_value => $1.integer_value * 60000); }
        | time number HOURS TOKEN
            { $$ := (token category => integer cat,
                     integer value => $1.integer value * 3600000); }
        ;
time number
        : INTEGER LITERAL
            { $$ := (token category => integer_cat,
                     integer value => $1.integer value); }
        ;
expression list
        : expression_list ',' expression
            { $$ := (token_category => expression_seq_cat,
                     expression_seq_value =>
                     expression sequence pkg.add($3.expression value,
$1.expression_seq_value )); }
        | expression
            { $$ := (token category => expression seq cat,
                     expression seq value =>
                     expression sequence pkg.add($1.expression value,
                                                  empty exp seq )); }
        ;
-- Expressions can appear in guards appearing in control constraints.
-- These guards can be associated with triggering conditions, or
-- conditional outputs, conditional exceptions, or conditional timer
-- operations. Similar to initial expression, except that tim e values
-- and references to timers and data streams are allowed.
expression
        : TRUE
            { $$ := (token_category => expression_cat,
                     expression_value => true_expression); }
        | FALSE
            { $$ := (token_category => expression_cat,
                     expression value => false expression); }
        | INTEGER LITERAL
            { $$ := (token_category => expression cat,
                     expression value =>
                     create_integer_literal($1.integer_value)); }
        | REAL LITERAL
```

```
{ $$ := (token_category => expression_cat,
              expression value =>
              create_real_literal($1.text_value)); }
| STRING LITERAL
    { $$ := (token category => expression cat,
              expression value =>
              create_string_literal($1.text_value)); }
| IDENTIFIER
    { $$ := (token_category => expression cat,
             expression_value =>
             create_identifier($1.psdl_id_value)); }
 -- The only difference from the initial expression
| time
    { $$ := (token_category => expression_cat,
             expression value =>
    create time literal (natural($1.integer_value)));}
| type name '.' IDENTIFIER
    { $$ := (token category => expression cat,
             expression value =>
             create function_call($1.type_name_value,
                                   psdl_id($3.psdl id value),
                                   empty exp seq)); }
| type_name '.' IDENTIFIER '(' expression_list ')'
    { $$ := (token_category => expression_cat,
             expression value =>
             create_function_call($1.type_name_value,
                                  psdl_id($3.psdl_id value),
                                   $5.expression_seq_value));}
| '(' expression ')'
    { $$ := (token category => expression cat,
             expression_value => $2.expression value);
              }
| expression log_op expression
                                  %prec logical operator
    { $$ := (token_category => expression_cat,
             expression_value =>
             create_binary_op ($1.expression_value,
                               $2.psdl_id_value,
                               $3.expression_value)); }
| expression rel_op expression %prec relational_operator
    { $$ := (token_category => expression_cat,
             expression value =>
             create_binary_op ($1.expression_value,
                               $2.psdl_id_value,
                               $3.expression value)); }
| '-' expression
                               %prec unary_adding_operator
    { $$ := (token_category => expression_cat,
             expression value =>
             create_unary_op
                              (convert("-"),
                               $2.expression_value)); }
| '+' expression
                                %prec unary adding operator
   { $$ := (token_category => expression_cat,
             expression_value =>
             create_unary op (convert("+"),
                               $2.expression value)); }
```
```
| expression bin add op expression
         %prec binary adding operator
           { $$ := (token category => expression cat,
                    expression_value =>
                    create binary op ($1.expression value,
                                    $2.psdl id value,
                                    $3.expression value)); }
       | expression bin_mul_op expression
         %prec multiplying operator
           { $$ := (token category => expression cat,
                    expression value =>
                    create binary op ($1.expression value,
                                    $2.psdl id value,
                                    $3.expression value)); }
       | expression EXP_TOKEN expression
         %prec highest_precedence_operator
           { $$ := (token_category => expression_cat,
                    expression value =>
                    create binary op ($1.expression value,
                                    convert("**"),
                                    $3.expression value)); }
       | NOT TOKEN expression
                                    %prec
highest precedence operator
           { $$ := (token category => expression cat,
                   expression value =>
                   create_unary_op (convert("NOT"),
                                    $2.expression value)); }
       | ABS TOKEN expression
                                    %prec
highest precedence operator
           { $$ := (token category => expression cat,
                   expression value =>
                   create_unary_op (convert("ABS"),
                                    $2.expression value)); }
       | '?' { $$ := (token_category => expression_cat,
                   expression_value => undefined_expression ); }
       ;
응용
package spec parser
        with psdl_program pkg; use psdl program pkg;
 with text io; use text io;
package parser is
 procedure get(item: in out psdl program);
 procedure get(file: in file type; item: in out psdl program);
 syntax error: exception;
  semantic error: exception;
```

end parser;

___ package body parser _____ with parser_tokens; use parser tokens; with parser goto; use parser goto; with parser shift_reduce; use parser shift_reduce; with parser_lex; use parser_lex; with parser_lex_dfa; use parser_lex_dfa; with psdl_component_pkg; use psdl_component_pkg; with psdl_concrete_type_pkg; use psdl_concrete_type_pkg; with psdl_id_pkg; use psdl_id_pkg; with ada_id_pkg; use ada_id_pkg; with psdl_graph_pkg; use psdl_graph_pkg; with expression_pkg; use expression pkg; with spec req map pkg; use spec req map pkg; with cc_req_map_pkg; use cc_req_map_pkg; with cc req map map pkg; use cc req map map pkg; with psdl_io; use psdl_io; package body parser is subtype exp_seq is expression_sequence_pkg.sequence; function empty_exp_seq return expression_sequence renames expression_sequence_pkg.empty; -- Returns an empty expression sequence. type attribute_type is (gen_par, input, output, state, exc, met); -- global variables used by the parser. the_program: psdl program; the_component: psdl_component; the operator: operator; the atomic type: atomic_type; the atomic operator: atomic operator; the_composite_type: composite type; the composite operator: composite operator; the ada name: ada id; the_imp_lang: psdl id; the gen par: type declaration; the_type_gen_par: type_declaration; the keywords: psdl id set; the description: text; the axioms: text; the_output_id: output id; -- a temporary variable to hold output_id to construct out_guard map the excep id: excep id; -- a temporary variable to hold excep_id to construct excep_trigger map the_model: type declaration;

```
the operation map: operation map;
  the data structure: type name;
  the input: type declaration;
  the output: type declaration;
  the state: type declaration;
  the states token line: natural;
  the states token: text;
  the initial expression map: init map;
  the exceptions: psdl id set;
  the specified met: millisec;
  the graph: psdl graph;
  the streams: type declaration;
  the timers: psdl id set;
  the trigger map: trigger map;
  the exec guard: exec guard map;
  the out guard: out guard map;
  the excep_trigger: excep_trigger_map;
  the timer op: timer op map;
  the_per: timing_map;
  the fw: timing map;
  the mcp: timing map;
  the mrt: timing map;
  the operator id: op id;
    -- is used for storing the operator id's in control constraints
part
  is atomic type: boolean;
    -- true if the psdl component is an atomic type.
  is atomic operator: boolean;
   -- true if the psdl component is an atomic operator.
  is specification: boolean;
   -- True if the current unit is a psdl specification part.
  the init exp seq: exp seq;
   -- Holds the initial expressions for all state variables in an
operator spec.
  the_type_name: type_name;
  the type decl: type declaration;
    -- Used to hold an inherited/synthesized attribtue pair.
  the trigger: trigger;
  the timer op record: timer op;
  the_timer_op_set: timer_op_set;
  the impl desc: text;
  the_attribute_type: attribute_type;
  the_gen_par_rb: spec_req_map;
  the input rb: spec req map;
  the output rb: spec req map;
  the state rb: spec req map;
  the exception rb: spec req map;
  the spec met rb: psdl id sequence;
  the eg rb: cc req map;
  the per rb: cc req map;
  the fw rb: cc req map;
  the_mcp_rb: cc_req_map;
```

```
the_mrt_rb: cc_req_map;
the_o_srbm: spec_req_map;
the_e_srbm: spec_req_map;
the_reset_srbm: spec_req_map;
the start srbm: spec req map;
the_stop_srbm: spec_req_map;
the_o_rb: cc req map map;
the e rb: cc req map map;
the reset rb: cc req map map;
the_start_rb: cc_req_map_map;
the_stop_rb: cc_req_map_map;
has_syntax_error: boolean := false;
   -- procedure initialize_state_variables
______
procedure initialize_state_variables is
begin
 yyval := (token_category => no_value_cat);
end initialize_state_variables;
-- procedure yyparse
procedure yyparse; -- Body is automatically generated.
-- procedure yyerror
procedure yyerror(s: in string := "syntax error";
             err_line: natural := current_line;
             err_token: text := convert(yytext)) is
 space: integer;
begin
 new line;
 put line(standard_error,
        "line" & integer'image(err_line) & ": "
        & convert(err_token));
 space := integer'image(err_line)'length +
        integer(convert(err token)'length) + 5;
 for i in 1 .. space loop put(standard error, "-"); end loop;
 put_line(standard error, "^ " & s);
 has_syntax_error := true;
end yyerror;
        -- given a string of characters corresponding to a natural number,
-- returns the natural value
function convert_to_natural(string_digit: string) return natural is
 digit, value : natural := 0;
```

```
begin
   for i in 1 .. string_digit'length loop
    case string_digit(i) is
      when '0' => digit := 0;
      when '1' \Rightarrow digit := 1;
      when '2' => digit := 2;
      when '3' => digit := 3;
      when '4' => digit := 4;
      when '5' => digit := 5;
      when '6' => digit := 6;
      when '7' => digit := 7;
      when '8' \Rightarrow digit := 8;
      when '9' => digit := 9;
      when others => return value;
    end case;
    value := (10 * value) + digit;
   end loop;
   return value;
 end convert to natural;
 ___
                     procedure get
 --
 -- reads the standard input, parses it, and creates the
 -- psdl adt.
 procedure get(item: in out psdl program) is
begin
 initialize_state_variables;
 initialize_yylex;
 yyparse;
 if has_syntax_error then
    assign(item, empty_psdl_program);
    raise syntax_error;
 end if;
 assign(item, the_program);
end get;
 procedure get
 -- reads the psdl source file from a file,
 -- parses it, and creates the psdl adt.
 _______
                                   procedure get(file: in file type;
            item: in out psdl program ) is
 begin
   set input(file);
   get(item);
   set input(standard input);
 end get;
 ___
                  procedure bind_type declaration
```

```
103
```

```
____
  --bind each id in id the id
  --set to the type name
  --return temp_type_decl
  procedure bind_type_declaration(i_s: in psdl_id_sequence; tn: in
type_name;
                              td: in out type_declaration) is
  begin
    for id : psdl_id in psdl_id_sequence_pkg.scan(i_s) loop
      bind(id, tn, td);
    end loop;
  end bind type declaration;
  ___
                     procedure bind initial state
  ___
  -- Bind each id in the state map domain
  -- set to the type name initial expression
 procedure bind_initial_state(state: in type_declaration; init_seq: in
exp seq;
                         init_exp_map: in out init_map) is
   i: natural := 0;
 begin
-- Added by Dave Dampier 20 April 1994 to
-- eliminate use of the M4 Macros,
-- and adopt use of the new generator processor for loops.
   for id : psdl_id, td : type_name in
type_declaration pkg.scan(state) loop
     i := i + 1;
     if i > expression_sequence_pkg.length(init_seq) then
       yyerror("semantic error - some states are not initialized.",
              the_states_token_line, the_states_token);
       raise semantic error;
     else bind(id,
              expression_sequence_pkg.fetch(init_seq, i),
              init_exp map);
     end if;
   end loop;
-- End of Added Code.
-- Also eliminated old M4 code.
   if i < expression_sequence_pkg.length(init_seq) then</pre>
     yyerror("semantic error -
             there are more initializations than the states",
             the states_token_line, the_states_token);
     raise semantic error;
   end if;
 end bind initial state;
 ___
                    procedure bind spec rb
```

```
-- Bind each id in the ids to the required-by sequence
  -- in the appropriate spec requirements map.
                                                            _____
  procedure bind spec rb(a t: attribute type; ids, reqs :
psdl id sequence) is
  begin
    case a t is
      when gen par =>
        for id : psdl_id in psdl_id_sequence_pkg.scan(ids) loop
          bind(id, reqs, the gen par rb);
        end loop;
      when input =>
        for id : psdl id in psdl id sequence pkg.scan(ids) loop
          bind(id, reqs, the input rb);
        end loop;
      when output =>
        for id : psdl id in psdl id sequence pkg.scan(ids) loop
          bind(id, reqs, the_output rb);
        end loop;
      when state =>
        for id : psdl id in psdl id sequence pkg.scan(ids) loop
          bind(id, reqs, the state rb);
        end loop;
      when exc =>
        for id : psdl id in psdl id sequence pkg.scan(ids) loop
          bind(id, reqs, the exception rb);
        end loop;
      when met =>
        -- This case is different because the specified-met required-by
        -- applies to the entire operator.
        -- Normally there should be at most one specified met,
        -- but if there is more than one, the requirements traces are
        -- combined.
        for id : psdl_id in psdl id sequence pkg.scan(reqs) loop
          if not member(id, the spec_met_rb) then
             add(id, the_spec_met_rb);
          end if;
        end loop;
    end case;
  end bind_spec_rb;
  function remove braces(t: text) return text is
    s: string := to_string(t);
    len: natural := length(t);
 begin
    return to text(s(2 .. len-1));
  end remove braces;
  -- Generated body of yyparse goes here.
##%procedure parse
end parser;
__********************
-- Parser Tokens
```

```
105
```

---******************

```
with
     Psdl_Concrete_Type_Pkg, Expression_Pkg, Psdl_Id_Pkg;
use
      Psdl_Concrete_Type_Pkg, Expression_Pkg, Psdl_Id_Pkg;
package Parser Tokens is
    type token_category_type is (integer cat,
                                  text cat,
                                  psdl id cat,
                                  psdl_id_sequence_cat,
                                  op_id_cat,
                                  operator name cat,
                                  opt arg cat,
                                  type name cat,
                          type decl cat,
                                  timer op id cat,
                                  expression cat,
                                  expression seq_cat,
                                  property map cat,
                                  no_value_cat);
    type yystype (token_category: token_category_type := no_value_cat)
is
      record
         case token_category is
           -- lexical token attributes:
           when integer cat =>
             integer value: integer;
           when text cat =>
             text value: text;
           -- grammar psdl_id attributes:
           when psdl id cat =>
             psdl_id_value: psdl_id;
           when psdl_id_sequence_cat =>
             psdl id sequence_value: psdl_id_sequence;
           when op id cat =>
             op_id value: op id;
           when operator name cat =>
             type_name_part, op_name_part: psdl_id;
           when opt arg cat =>
             input_value, output_value: psdl_id_sequence;
           when type name cat =>
             type_name_value: type name;
           when type decl cat =>
             type decl_value: type_declaration;
         when timer op id cat =>
             timer_op_id_value: timer_op_id;
          when expression cat =>
             expression_value: expression;
          when expression seq cat =>
             expression_seq_value: expression_sequence;
          when property map cat =>
            property_map_value: init map;
```

```
when no value cat => null;
     end case;
   end record;
YYLVal, YYVal : YYSType;
type Token is
    (End_Of_Input, Error, '?', '{',
     '}', Illegal_Token, '(',
     ')', ',', '[<sup>-</sup>,
']', ':', '.',
     '|', Arrow, True,
     False, Axioms Token, By All Token,
     Req By Token, By Some Token, Call Period Token,
     Control_Token, Constraints Token, Data Token,
     Description Token, Edge Token, End Token,
     Exceptions Token, Exception Token, Execution Token,
     Finish Token, Generic Token, Graph Token,
     Hours Token, If Token, Implementation Token,
     Initially_Token, Input_Token, Keywords_Token,
     Maximum_Token, Minimum_Token, Microsec_Token,
     Min Token, Ms Token, Mod Token;
     Not_Token, Operator_Token, Or_Token,
     Output_Token, Period_Token, Property_Token,
     Reset_Token, Response_Token, Sec_Token,
     Specification Token, Start Token, States Token,
     Stop Token, Stream Token, Time Token,
     Timer_Token, Triggered_Token, Type_Token,
     Vertex Token, Within Token, Identifier,
     Integer Literal, Real Literal, String Literal,
     Text_Token, And_Token, Xor_Token,
     Logical_Operator, '<', '>',
     '=', Greater_Than_Or_Equal, Less_Than_Or_Equal,
     Inequality, Relational_Operator, '+',
     '-', '&', Binary_Adding_Operator,
     Unary Adding_Operator, '*', '/',
     Rem_Token, Multiplying_Operator, Exp_Token,
    Abs Token, Highest Precedence Operator );
```

Syntax_Error : exception;

end Parser_Tokens;

```
__*****************
-- AWK script for Parser
--*****************
/procedure YY\_USER\_ACTION is/ {
 # keep this line and the "begin" that follows.
 print $0
 getline
 print $0
 # replace the body of the procedure with our code
 print " increment_column_number(yytext'length);"
 # get the "null;" and discard it
 getline
 next
 }
 #
 # otherwise pass the line through
 #
 { print $0 }
```

APPENDIX C: SELECTED SOURCE CODE (PSDL EDITOR)

```
# Makefile for PSDL Editor
SHELL = /usr/bin/csh
C DIR = ./C Code
INCLUDE FLAGS = \setminus
  -I. \
  -I (PT DIR) \setminus
  -I$(PT_DIR)/GENERIC TYPES \
  -I$ (PT DIR) / INSTANTIATIONS \
  -I$(PARSER_DIR)
LIBS = -1Xm - 1Xt - 1Xext - 1X11 - 1m - 1g++ - 1gcc
PSDL TYPE = \setminus
  $(PT_DIR)/psdl_io.ali
GE OBJECTS = \setminus
  $(GE_DIR)/graph_editor.o \
  $(GE DIR)/operator object.o \
  $(GE DIR)/stream object.o \
  $(GE DIR)/spline object.o \
  $(GE DIR)/graph object list.o \
  $(GE DIR)/font table.0 \
  $(GE DIR)/graph object.o \
  $(GE DIR)/setcursor.o \
  $(GE DIR)/gettopshell.o \
  $(GE DIR)/postpopup.o \
  $(GE DIR)/build option.o \
  $(GE DIR)/timer tool.o \
  $(GE DIR)/action area.o \
  $(GE_DIR)/warning.o \
  $(GE DIR)/ge utilities.o \
  $(GE DIR)/stream property menu.o \
  $(GE DIR)/operator property menu.o \
  $(GE DIR)/windows.o \
  $(GE DIR)/get unique id.o \
  $(GE DIR)/ge utilities debug.o \
  $(GE DIR)/report errors.o
PARSERS = $(PARSER DIR)/parser pkg.ali
LOCAL OBJECTS = \setminus
  $(C DIR)/main.o \
  $(C_DIR)/sde_structure.o \
  $(C_DIR)/sde_globals.o \
```

\$(C_DIR)/ge_support.o

109

```
GENERATED ADA = \setminus
  analysis_pkg.adb \
  at_list_pkg.adb \
  ge utilities.adb \
  id_list pkg.adb \
  io_utilities.adb \
  op_list_pkg.adb \
  psdl utilities pkg.adb \
  st_list_pkg.adb \
#
 editor_io_pkg.adb \
#
   editor_io_pkg-utilities.adb
SOURCES = \setminus
  action node pkg.ads \
  action_node_pkg.adb \
  analysis_pkg.ads \
  analysis_pkg.adb \
  at_list_pkg.ads \
  at list pkg.adb \
  c_boolean pkg.ads \
  c boolean pkg.adb \
  c int pkg.ads \
  c_string_pkg.ads \
  c_string_pkg.adb \
  c_unsigned int pkg.ads \
  duration_type_pkg.ads \
 error_msgs_pkg.ads \
 error_msgs_pkg.adb \
 ge_action pkg.ads \
 ge_action_pkg.adb \
 ge_interface_pkg.ads \
 ge_op_id_pkg.ads \
 ge_op_id_pkg.adb \
 ge_operator_pkg.ads \
 ge_operator_pkg.adb \
 ge_trigger_type_pkg.ads \
 ge_trigger_type_pkg.adb \
 ge_utilities.ads \
 ge_utilities.adb \
 graph_desc_pkg.ads \
 graph desc pkg.adb \
 id list pkg.ads \
 id_list_pkg.adb \
 io utilities.ads \
 io utilities.adb \
 op_list_pkg.ads \
 op_list_pkg.adb \
 property names pkg.ads \
 property_names pkg.adb \
 psdl editor.ads \
 psdl_editor.adb \
 psdl_editor_pkg.ads \
```

```
psdl editor pkg.adb \
  psdl utilities pkg.ads \
  psdl utilities pkg.adb \
  spline ptr pkg.ads \
  spline_ptr_pkg.adb \
  st id pkg.ads \
  st_list_pkg.ads \
  st_list pkg.adb \
  stream pkg.ads \
  stream pkg.adb \
  time trigger type pkg.ads \
  time trigger type pkg.adb \
  time_unit_type_pkg.ads \
  time unit type pkg.adb \
  unique id pkg.ads
SCRIPTS = \setminus
  Makefile \
  gen.sed
.SUFFIXES:
all:
      (cd $(C DIR) ; make all)
      (cd $(PARSER DIR) ; make parsers)
      (cd $(GE DIR) ; make ge)
      (cd $(PT DIR) ; make gen parsers)
      make generated sources
      make psdl_editor
psdl_editor::
      gnatmake -g -o psdl editor $(INCLUDE FLAGS) psdl editor.adb \
               -bargs -n $(PARSERS) \
               -largs $(GE_OBJECTS) $(LOCAL OBJECTS) $(LIBS)
generated_sources: $(GENERATED ADA) $(PARSERS)
at_list pkg.adb: at list pkg.g
      gen < at_list_pkg.g > tmp
      sed -f gen.sed tmp > at list pkg.adb
      /bin/rm tmp
id list_pkg.adb: id_list pkg.g
      gen < id_list_pkg.g > id_list_pkg.adb
io utilities.adb: io utilities.g
      gen < io_utilities.g > io_utilities.adb
op list pkg.adb: op list pkg.g
      gen < op_list_pkg.g > op_list_pkg.adb
st list pkg.adb: st list pkg.g
      gen < st_list_pkg.g > st_list pkg.adb
```

```
111
```

```
psdl_utilities_pkg.adb: psdl_utilities_pkg.g
      gen < psdl_utilities_pkg.g > psdl_utilities_pkg.adb
analysis_pkg.adb: analysis_pkg.g
      gen < analysis_pkg.g > analysis_pkg.adb
ge_utilities.adb: ge utilities.g
      gen < ge_utilities.g > ge_utilities.adb
ci: generated sources
      ci_files -tRCS/desc *.[Cgly] *.ad[sb] $(SCRIPTS)
      (cd C_Code ; make ci)
ci_export: generated sources
      ci_files -tRCS/desc *.[Cgly] *.ad[sb] $(SCRIPTS)
      (cd C_Code ; make ci)
      sleep 1
      touch *.ali *.o
edit: all
      cp psdl_editor ../CAPS.RELEASE.1.2/bin
      (cd TEST ; psdl_editor test.psdl ; psdl_editor autopilot.psdl)
clean:
      rm *.ali *.o
.IGNORE:
# ignore nonzero exit codes below.
xref:
```

```
gnatf -e -f -x6 $(INCLUDE_FLAGS) *.ad*
```

This is the main procedure of the psdl editor.
It is called from a C main program to enable
the loader to find all the code.
The problem is that gnathind does not allow multiple
Ada root programs with an Ada main program.
Since we have Ada calling C with calls
back to Ada, and there are no direct links between
the Ada driver at the top and the Ada parsers at the
bottom, the binder needs multiple roots to enable
the loader to find all the Ada modules.

procedure psdl_editor; pragma Export (C, psdl_editor, "psdl_editor");

```
-- This is the main program of the PSDL_EDITOR
---
-- Usage:
      psdl editor <psdl-input-filename>
___
-- Change:
--
with Ada.Command Line;
 with Ada.Exceptions; use Ada.Exceptions;
 with Text_Io; use Text Io;
 with psdl editor pkg; use psdl editor pkg;
procedure psdl editor is
 PSDL File: File Type;
begin
 if Ada.Command_Line.Argument Count = 1 then
    declare
             -- right number of arguments, get the file name.
      filename: string := Ada.Command Line.Argument(1);
    begin
      if (filename'length >= 6) and then
         filename(filename'last-4 .. filename'last) = ".psdl"
      then -- File name is ok, open the file and edit it.
         begin
           Open(PSDL File, In File, filename);
         exception
          when name error =>
            begin
              -- The file does not exist, try to create one.
              Create(PSDL_File, In_File, filename);
            exception
              when others =>
                PUT_LINE(STANDARD_ERROR,
                        "psdl editor: couldn't open or create "
                           & filename);
            end:
          when others => PUT LINE(STANDARD ERROR,
                                 "psdl_editor: couldn't open "
                           & filename);
         end;
         -- Got a PSDL file, edit and update it under user control.
         edit_program(PSDL_File => PSDL File,
                    prototype name =>
                     filename(1 .. filename'last-5));
      else PUT_LINE(STANDARD_ERROR, "error: bad file name");
          PUT LINE (STANDARD ERROR,
                   "usage: psdl_editor prototype_name.psdl");
      end if;
    end;
 else
```

```
PUT LINE (STANDARD ERROR, "psdl editor:
                                  error, wrong number of arguments");
       PUT LINE (STANDARD ERROR, "usage:
                                 psdl editor prototype name.psdl");
-- begin debugging patch
       PUT LINE(STANDARD ERROR, "using file test.psdl");
     declare
               -- right number of arguments, get the file name.
       filename: string := "test.psdl";
     begin
       if (filename'length >= 6) and then
          filename(filename'last-4 .. filename'last) = ".psdl"
       then -- File name is ok, open the file and edit it.
          begin
            Open(PSDL File, In File, filename);
          exception
            when name_error =>
              begin
                -- The file does not exist, try to create one.
                Create(PSDL File, In File, filename);
              exception
                when others =>
                  PUT LINE (STANDARD ERROR,
                           "psdl editor: couldn't open or create "
                            & filename);
              end;
            when others => PUT LINE(STANDARD ERROR,
                                     "psdl editor: couldn't open "
                                      & filename);
          end;
          -- Got a PSDL file, edit and update it under user control.
          edit program(PSDL File => PSDL File,
                       prototype name =>
                       filename(1... filename'last-5));
       else PUT LINE(STANDARD ERROR, "error: bad file name");
            PUT LINE (STANDARD ERROR, "usage:
                     psdl_editor prototype_name.psdl");
       end if;
     end;
-- end debugging patch
  end if;
exception
 when the exception: others =>
    PUT LINE(STANDARD ERROR, "Internal error: unexpected exception " &
                             Exception Name(the exception));
end psdl editor;
```

----********************** -- GE UTILITIES SPECIFICATION --**** with psdl_program_pkg; use psdl_program_pkg; with psdl_component_pkg; use psdl_component_pkg; with GE_Operator_pkg; use GE_Operator_pkg; with st_list_pkg; use st_list_pkg; with op_list_pkg; use op_list_pkg; with psdl_graph_pkg; use psdl_graph_pkg; package ge utilities is procedure add_edges(streams: in St_List; ops: in Op_List; g: in out psdl_graph); -- Includes adding edge properties. procedure modify_child(child: in GE_Operator; edges: in St_List; prototype: in psdl_program; edited_prototype: in out psdl_program); procedure modify_child_type_op(child: in GE_Operator; edges: in St List; prototype: in psdl_program; edited prototype: in out psdl program); procedure modify_child_operator(child: in GE_Operator; edges: in St List; prototype: in psdl program; edited_prototype: in out psdl program); procedure update_type_operation_names(current_op: in composite operator; edited op: in out composite_operator); end ge utilities;

```
__****************************
-- GE UTILITIES BODY
__***************************
  with stream pkg; use stream pkg;
  with psdl concrete type pkg; use psdl concrete type pkg;
  with psdl_utilities_pkg; use psdl_utilities_pkg;
  with psdl id pkg; use psdl id pkg;
  with ada_id_pkg; use ada_id_pkg;
  with expression pkg; use expression pkg;
  with spec req map_pkg; use spec_req_map_pkg;
  with time unit type pkg; use time unit type pkg;
  with property names pkg; use property names pkg;
  with id list pkg; use id list pkg;
  with spline ptr pkg; use spline ptr pkg;
  with substitution map pkg; use substitution map pkg;
  with vertex substitution map pkg; use vertex substitution map pkg;
  with text io; use text io;
  with psdl io; use psdl io;
package body ge utilities is
  not found: exception;
  procedure add edge(S: in STREAM; ops: in Op List; g: in out
psdl graph) is

    Includes adding edge properties.

    source op, sink op: GE Operator;
    source, sink : op_id;
    str: psdl id := to psdl id(label(S));
    lat : millisec;
 begin
    if not Is Deleted(S) then
       -- find an op_id to the sending vertex
       source := find op id(from(S), ops);
       -- find an op id to the receiving vertex
       sink := find_op_id(to(S), ops);
       -- create a new edge and assign properties
       lat := Latency(S);
       g := add_edge(source, sink, str, g, lat);
       set_property(source, sink, str, id p,
               create integer literal(Integer(Id(S))), g);
       set property(source, sink, str, label font p,
               create_integer_literal(Label_Font(S)), g);
       set_property(source, sink, str, label x offset p,
               create integer literal(Label X Offset(S)), g);
       set property(source, sink, str, label_y_offset_p,
               create_integer_literal(Label Y Offset(S)), g);
       set_property(source, sink, str, latency font p,
               create_integer_literal(Latency_Font(S)), g);
       set_property(source, sink, str, latency unit p,
               create_integer_literal(to_integer(Latency Unit(S))), g);
       set property(source, sink, str, latency_x_offset_p,
```

```
create_integer_literal(Latency X Offset(S)), g);
       set property(source, sink, str, latency y offset p,
                     create_integer_literal(Latency_Y_Offset(S)), g);
       set_property(source, sink, str, spline_p,
                     create_string_literal(to_text(Arc(S))), g);
    end if;
  end add edge;
  procedure add_edges(streams: in St_List; ops: in Op_List; g: in out
psdl graph) is
    -- Includes adding edge properties.
    L: St_List := streams;
    current edge: STREAM;
  begin
    while not St List Is Null(L) loop
      current edge := ST(L);
      add edge(current_edge, ops, g);
      L := Next(L);
    end loop;
  end add_edges;
  procedure modify_child(child: in GE_Operator; edges: in St_List;
                          prototype: in psdl program;
                          edited_prototype: in out psdl_program) is
    -- Create operator specs for new children and
    -- update operator specs for modified children.
  begin
   if not Is_Deleted(child) then
      if is_type_operation(child) then
         modify_child_type_op(child, edges,
                               prototype, edited prototype);
      else modify_child_operator(child, edges,
                                 prototype, edited prototype);
      end if;
   end if;
  end modify_child;
  procedure modify_child_type_op(child: in GE_Operator; edges: in
St List;
                                 prototype: in psdl program;
                                 edited prototype: in out psdl program)
is
    -- Create operator specs for new children and
    -- update operator specs for modified children.
    child_vertex: op_id := vertex_id(child);
    child name: psdl id := base name(child vertex);
    child_type_name: psdl_id := child_vertex.type_name;
    child_type: data_type;
    child_type_ops: operation map;
   child op: operator;
   exc_list: psdl id set;
   exc_rb: spec_req_map;
```

begin

```
child type := fetch(edited prototype, child type name);
    if child type = null component then
       -- The type has not been defined yet, create one.
       child type := make atomic type(child name);
       add(child_type, edited_prototype);
    end if;
    child type ops := operations(child type);
    if (Is New(child) or Is Modified(child)) or else
       not member(child name, child type ops)
    then
      to_exceptions(to_string(Exception List(child)), exc list,
exc rb);
      build psdl operator(
        c name => child name,
        c_a_name => to_ada_id(child name),
        imp lang => Impl Lang(child),
                                         -- GE does not supply this
        g_par => empty_type_declaration,
        gen par rb => spec req map pkg.empty, --GE doesn't supply this
        kwr => to psdl id set(Keyword List(child)),
        i desc => Informal Desc(child),
        f desc => Formal_Desc(child),
        inp => inputs(Id(child), edges),
        otp => outputs(Id(child), edges),
        st => empty type declaration,
        i_exp_map => empty_init_map,
        excps => exc list,
        s met => MET(child),
        input_rb => spec_req map pkg.empty, -- GE does not supply this
        output_rb => spec_req_map_pkg.empty, -- GE doesn't supply this
        state_rb => spec_req_map_pkg.empty, -- GE does not supply this
        excep rb => exc rb,
        smet rb => MET Reqmts(child),
        is atomic => not Is Composite(child),
        the opr => child op);
     bind operation(child_name, child_op, child_type_ops);
      set_operations(child_type, child type ops);
   elsif Is Modified(child) then
      child_op := fetch(child_type_ops, child_name);
      set_name(child_op, child_name);
      to_exceptions(to_string(Exception List(child)),
                              exc list, exc rb);
      if not Is Composite(child) then
         set_implementation_language(child op, Impl Lang(child));
      end if;
      set keywords (child op, to psdl id set (Keyword List (child)));
      set informal description(child op, Informal Desc(child));
      set axioms(child op, Formal Desc(child));
      set inputs(child op, inputs(Id(child), edges));
      set_outputs(child_op, outputs(Id(child), edges));
      set_exceptions(child_op, exc list);
      set specified met(child op, MET(child));
      set_specified_met_reqs(child op, MET Reqmts(child));
      set exception reqs(child_op, exc_rb);
```

else

```
child_op := fetch(child_type_ops, child_name);
      set_inputs(child_op, inputs(Id(child), edges));
      set_outputs(child_op, outputs(Id(child), edges));
    end if;
  end modify_child_type op;
  procedure modify_child_operator(child: in GE_Operator;
                                  edges: in St List;
                                  prototype: in psdl program;
                                  edited prototype: in out
psdl program) is
    -- Create operator specs for new children and
    -- update operator specs for modified children.
   child_name: psdl_id := name(child);
   child op: operator;
   exc_list: psdl_id_set;
   exc_rb: spec_req_map;
 begin
   if Is New(child) then
     to exceptions(to_string(Exception List(child)),
                    exc_list, exc rb);
     build psdl operator(
        c name => child name,
        c_a_name => to ada id(child name),
        imp lang => Impl Lang(child),
        g_par => empty_type_declaration, -- GE does not supply this
        gen_par_rb => spec_req_map_pkg.empty, --GE doesn't supply this
       kwr => to_psdl id_set(Keyword List(child)),
        i desc => Informal_Desc(child),
       f desc => Formal Desc(child),
       inp => inputs(Id(child), edges),
       otp => outputs(Id(child), edges),
        st => empty_type_declaration,
        i_exp_map => empty init map,
       excps => exc list,
       s met => MET(child),
       input_rb => spec_req_map_pkg.empty, -- GE does not supply this
       output_rb => spec_req_map_pkg.empty, -- GE doesn't supply this
       state_rb => spec_req_map_pkg.empty, -- GE does not supply this
       excep rb => exc_rb,
       smet_rb => MET_Reqmts(child),
       is atomic => not Is_Composite(child),
       the_opr => child_op);
     add(child op, edited_prototype);
   elsif Is Modified(child) then
     assign(child_op, find(suffix(child_name), prototype));
     -- Find the old operator based on the op_num,
     -- the name may have changed.
     -- Use assign to make a copy so recycle will be safe.
     set_name(child_op, child_name);
     to exceptions(to_string(Exception_List(child)),
                   exc_list, exc rb);
     if not Is Composite(child) then
        set_implementation_language(child_op, Impl_Lang(child));
```

```
120
```

```
end if;
      set keywords (child op, to psdl id set (Keyword List (child)));
      set informal description(child op, Informal Desc(child));
      set_axioms(child_op, Formal_Desc(child));
      set inputs(child op, inputs(Id(child), edges));
      set outputs(child op, outputs(Id(child), edges));
      set exceptions(child op, exc list);
      set specified met(child op, MET(child));
      set specified met reqs(child op, MET Reqmts(child));
      set_exception_reqs(child_op, exc_rb);
      add(child op, edited prototype);
    else
      assign(child op, find(suffix(child name), prototype));
        -- Use assign to make a copy so recycle will be safe.
      set_name(child_op, child_name);
      set inputs(child op, inputs(Id(child), edges));
      set_outputs(child_op, outputs(Id(child), edges));
        -- The inputs and outputs are derived from the graph,
        -- so they could have changed
        -- even if the explicit attributes of the child did not.
      add(child op, edited prototype);
    end if;
  end modify child operator;
  function find edge(edge id: expression; cg: psdl graph) return edge
is
    e_id: expression;
  begin
    for e: edge in edge_set_pkg.scan(edges(cg)) loop
        e id := get property(e.source, e.sink, e.stream name,
                             id_p, cg);
        if eq(e id, edge id) then return e; end if;
    end loop;
    raise not found;
  end find edge;
  -- will not work if different arcs of the same
  -- stream are renamed in different ways.
  function make_renaming(cg, eg: psdl_graph) return substitution_map is
    edge id: expression;
    original edge: edge;
    result: substitution map := empty;
 begin
    for e: edge in edge set pkg.scan(edges(eg)) loop
        edge_id := get property(e.source, e.sink, e.stream name,
                                id p, eg);
        begin
          original edge := find edge(edge id, cg);
        exception
          when not found => original edge := e;
        end;
        bind(original_edge.stream_name, e.stream_name, result);
    end loop;
    return result;
```

```
121
```

```
end make_renaming;
 procedure update_type_operation_names(current_op: in
composite_operator;
                                        edited op: in out
composite_operator) is
    cg: psdl_graph := graph(current op);
    eg: psdl graph := graph(edited op);
   id_renaming: substitution_map := make_renaming(cg, eg);
   vertex_renaming: vertex_substitution_map := empty;
   new_name: op id;
 begin
   for v: op_id in op_id_set_pkg.scan(vertices(eg)) loop
       if is type op (v) then
          new_name := transform_vertex(v, id_renaming);
          bind(v, new_name, vertex_renaming);
       end if;
   end loop;
   rename_vertices(edited_op, vertex_renaming);
 end update_type_operation_names;
```

```
end ge_utilities;
```

with psdl_program_pkg; use psdl_program_pkg; with psdl_component_pkg; use psdl_component_pkg; with psdl_graph_pkg; use psdl_graph_pkg; with psdl_concrete_type_pkg; use psdl_concrete_type_pkg; with psdl_id_pkg; use psdl_id_pkg; with expression_pkg; use expression_pkg; with St_List_pkg; use St_List_pkg; with spec_req_map_pkg; use spec_req_map_pkg; with ge_op_id_pkg; use ge_op_id_pkg; with C_Boolean_pkg; use C_Boolean_pkg;

-- The following are needed in the body
-- but can't go there because of gen.
with Ada.characters.handling; use Ada.characters.handling;
with Ada.strings; use Ada.strings;
with Ada.strings.fixed; use Ada.strings.fixed;
with Ada.strings.maps; use Ada.strings.maps;
with Ada.strings.maps.constants; use Ada.strings.maps.constants;

package PSDL Utilities Pkg is

function get_is_terminator(op_name: psdl_id; parent: operator) return
boolean;

-- returns true if the named operator is a terminator bubble

-- in the graph of the parent operator.

-- If there are several nodes with the given operator name,

-- uses the properties of the first one it finds.

-- of the operator

function Extract_Implementation_Id(c: in psdl_component) return Text; -- returns the implementation id of the atomic operator, -- e.g. Ada, TAE, C, etc

function Extract Vertex Ids(c: in psdl component) return op id set;

function Extract_MRT_Reqmts_Ids(name : in op id;

123

c : in psdl component) return psdl id set; function Extract_MCP_Reqmts_Ids(name : in op id; c : in psdl component) return psdl id set; function Extract_Trigger_Reqmts_Ids(name : in op_id; c : in psdl component) return psdl id set; function Extract_Edge_Set(c: in psdl_component) return edge_set; function Get_PSDL_O_GENERICS_Text(c : in psdl_component) return Text; -- the output string consists of all generic parameter declarations -- in the spec of the operator c function Get_Psdl_Types_Text(prog : in psdl_program) return Text; -- prog is a psdl program that contains only user-defined types function Get_Op_Spec_Text(o : in operator) return Text; -- gets the specification of the operator as a text string. function Get_Expression_Text(e : in expression) return Text; -- white characters are not allowed in the output -- string except when they -- are blanks inside a string literal, i.e. -- between two matching quotes function Get_Psdl_Id_Sequence_Text(s : in psdl_id_sequence) return Text; -- the output string consists of a sequence of -- non-white characters which -- correspons to a list of ids separated by commas function Get_Output_Guards_Text(name : in op_id; c : in psdl component) return text; -- the output string consists of all -- the output guards and requirements -- traces associated with the vertex "name" -- in the control constraint of c function Get_Exception_Triggers_Text(name : in op_id; c : in psdl component) return text; -- the output string consists of all the -- exception triggers and requirements -- traces associated with the vertex -- "name" in the control constraint of c function Get_Exception_List_Text(name : in op id; c : in psdl component) return text;

function Get Timer Operations Text(name : in op id; c : in psdl component) return text; -- the output string consists of all the timer -- operations and requirements -- traces associated with the vertex "name" -- in the control constraint of c function Get_Type_Name_Text(name : in type_name) return text; -- the output string represents the type_name as -- a sequence of non-white characters procedure separate_types(prototype: in psdl_program; root name: in psdl_id; types: in out psdl_program); -- Separates the operators from the types. procedure create root name(prototype name: in string; prototype: in out psdl program; root name: out psdl id); -- Produces the name of the root node if there is one, -- otherwise constructs one from the prototype name -- and creates a corresponding root node. procedure check_suffixes (prototype: in out psdl program); -- Check for suffixes and generates them if not there. function to psdl id set(s: psdl id sequence) return psdl id set; -- converts the sequence to a set. function to type name(s: c string) return type name; -- converts the string to a psdl type name. function to op id(label: text) return op id; -- converts the label and suffixes to a psdl op id. function to expression(s: c string) return expression; -- converts the string to a psdl expression. function to_operator(spec: text) return operator; -- returns an atomic operator with the given name -- and psdl specification. procedure to exceptions (s: in string; exc list: out psdl id set; exc rb: out spec req map); procedure to out guard map(s: in string; og: out out guard map; ogrb: out spec req map); procedure to_exception_guard_map(s: string; eg: out excep trigger map; egrb: out spec_req_map);

procedure to_timer_op_set(s: in string; timer_op: out timer op set;

resetrb, startrb, stoprb: in out spec_req_map);

end PSDL_Utilities_Pkg;

```
-- PSDL UTILITIES PKG BODY
with raw_text_file_pkg; use raw_text_file pkg;
 with text io; use text io;
 with io utilities; use io utilities;
 with psdl io; use psdl io;
 with expression io; use expression io;
 with type name io; use type name io;
 with op id io; use op id io;
 with output guard io; use output guard io;
 with exception guard io; use exception guard io;
 with exception io; use exception io;
 with timer op guard io; use timer op guard io;
 with Unique Id Pkg; use Unique Id Pkg;
 with property_names_pkg; use property_names_pkg;
 with cc_req_map_pkg; use cc_req_map_pkg;
 with substitution_map_pkg; use substitution_map_pkg;
 with vertex_substitution_map_pkg; use vertex_substitution_map_pkg;
package body PSDL Utilities Pkg is
 function get is terminator(op name: psdl id; parent: operator) return
boolean is
   -- returns true if the named operator is a terminator bubble
   -- in the graph of the parent operator.
   -- If there are several nodes with the given operator name,
   -- uses the properties of the first one it finds.
   g: psdl graph;
 begin
   g := graph(parent);
   -- Find the graph vertex corresponding to the given operator name.
   for v: op_id in op_id_set_pkg,scan(vertices(g)) loop
       if eq(base_name(v), op_name) then -- found it.
          return eq(get property(v, is terminator p, g),
          true_expression);
       end if;
   end loop;
   -- Should never get here.
   put_line(standard_error, "get_is_terminator: node name "
            & convert(op name));
   put line(standard error, " not found in the graph of "
            & convert(name(parent)));
   return false;
 exception
   when others =>
     put line(standard error,
              "get is terminator: unexpected exception");
     return false;
 end get is terminator;
```

```
function Extract_Input_Ids(c : in psdl_component) return psdl_id_set
is
    -- returns the set of psdl_ids which are names of input streams to
the operator
    input_ids : psdl_id_set;
    td : type declaration;
  begin
    td := inputs(c);
    assign(input_ids, map_domain(td));
    return input_ids;
  end Extract Input Ids;
  function Extract_Output_Ids(c : in psdl_component) return psdl_id_set
is
    -- returns the set of psdl ids which are names of output
    -- streams to the operator
    output ids : psdl id set;
    td : type_declaration;
  begin
    td := outputs(c);
    assign(output_ids, map_domain(td));
    return output ids;
  end Extract_Output_Ids;
 function Extract_State_Ids(c : in psdl_component) return psdl id set
is
    -- returns the set of psdl_ids which are names of state
   -- streams declared in the spec
   -- of the operator
   state ids : psdl id set;
    td : type declaration;
 begin
   td := states(c);
   assign(state_ids, map_domain(td));
   return state_ids;
 end Extract State Ids;
 function Extract_Input_Reqmts_Ids(name : in psdl_id;
                                    c : in psdl component)
   return psdl id set is
   -- returns the set of reqmts associated with the name
   -- input stream of the operator
   td : psdl_id sequence;
   input_reqmts_ids : psdl id set;
 begin
   td := input_reqs(c, name);
   return to_psdl_id_set(td);
 end Extract_Input_Reqmts_Ids;
 function Extract_Output_Reqmts_Ids(name : in psdl_id;
                                     c : in psdl component)
   return psdl id set is
   --returns the set of reqmts associated with
   -- the name output stream of the operator
```

```
td : psdl id sequence;
 begin
    td := output_reqs(c, name);
    return to psdl id set(td);
  end Extract Output Reqmts Ids;
  function Extract State Reqmts Ids (name : in psdl id;
                                     c : in psdl component)
    return psdl id set is
    -- returns the set of reqmts associated with the name
    -- state stream of the operator
    td : psdl_id_sequence;
 begin
    td := state reqs(c, name);
    return to psdl id set(td);
  end Extract State Reqmts Ids;
  function Extract Exception Reqmts Ids (name : in psdl id;
                                        c : in psdl component)
    return psdl id set is
    -- returns the set of reqmts associated with the name
    -- exception of the operator
    td : psdl id sequence;
 begin
    td := exception reqs(c, name);
    return to psdl id set(td);
  end Extract Exception Reqmts Ids;
  function Extract Met Reqmts Ids(c : in psdl component) return
psdl id set is
    td : psdl_id_sequence;
 begin
   td := specified maximum execution time reqs(c);
    return to psdl id set(td);
  end Extract_Met_Reqmts_Ids;
  function Extract Implementation Id(c: in psdl component) return Text
is
    -- returns the implementation id of the
    -- atomic operator, e.g. Ada, TAE, C, etc
 begin
    return to text(implementation language(c));
  end Extract Implementation Id;
  function Extract Vertex_Ids(c: in psdl_component) return op_id_set is
 begin
    return vertices(graph(c));
  end Extract Vertex Ids;
  function Extract_Edge_Set(c: in psdl component) return edge set is
 begin
    return edges(graph(c));
  end Extract Edge Set;
```

```
function Extract_Period_Reqmts_Ids(name : in op_id;
                                      c : in psdl component)
    return psdl id set is
    period reqmts ids : psdl id sequence;
  begin
    period reqmts ids := fetch(period reqs_map(c), name);
    return to psdl_id_set(period_reqmts_ids);
  end Extract_Period_Reqmts Ids;
  function Extract_FW Reqmts_Ids(name : in op id;
                                  c : in psdl_component)
    return psdl id set is
    fw_reqmts_ids : psdl_id_sequence;
  begin
    fw_reqmts_ids := fetch(finish_within_reqs_map(c), name);
    return to_psdl_id_set(fw_reqmts ids);
  end Extract_FW_Reqmts_Ids;
  function Extract MCP Reqmts Ids(name : in op id;
                                   c : in psdl component)
    return psdl id set is
    mcp_reqmts_ids : psdl_id_sequence;
  begin
    mcp_reqmts ids := fetch(minimum_calling_period_reqs_map(c), name);
    return to_psdl_id_set(mcp_reqmts_ids);
  end Extract MCP Reqmts Ids;
  function Extract MRT Reqmts_Ids(name : in op_id;
                                   c : in psdl component)
    return psdl id set is
   mrt reqmts_ids : psdl_id_sequence;
  begin
   mrt_reqmts_ids := fetch(maximum_response_time_reqs_map(c), name);
    return to psdl_id_set(mrt_reqmts_ids);
  end Extract_MRT_Reqmts_Ids;
  function Extract_Trigger_Reqmts_Ids(name : in op_id;
                                      c : in psdl component)
    return psdl id set is
   trigger_reqmts_ids : psdl_id_sequence;
 begin
   trigger_reqmts_ids := fetch(execution guard reqs_map(c), name);
    return to_psdl_id_set(trigger_reqmts ids);
  end Extract_Trigger_Reqmts Ids;
  function Get_PSDL_O_GENERICS_Text(c : in psdl_component) return Text
is
    -- the output string consists of all generic parameter declarations
   -- in the spec of the operator c
 begin
   -- not finished, probably not needed.
   return empty;
 end Get_PSDL_O_GENERICS Text;
```

```
function Get Expression Text(e : in expression) return Text is
    -- white characters are not allowed in the output
    -- string except when they
    -- are blanks inside a string literal, i.e.
    -- between two matching quotes
    out f: text file;
    in f: raw text file;
    result: text;
 begin
    -- Write the expression to a temporary file
    temporary_text_file(out_f);
    Set Output(out f);
    put expression(e);
    Set Output(Standard Output);
    Close(out f);
    -- Read it in as a text string from the temporary file
   temporary raw text file(in f);
    result := get(in f);
    raw text file pkg.Delete(in f);
    remove last char(result);
    return result;
 exception
    when others =>
      PUT LINE (STANDARD ERROR, "Get Expression Text: io error");
      return empty;
  end Get Expression Text;
  function Get_Psdl_Id_Sequence_Text(s : in psdl_id_sequence) return
Text is
    -- the output string consists of a
    -- sequence of non-white characters which
    -- correspons to a list of ids separated by commas
    out f: text file;
    in f: raw text file;
    result: text;
 begin
    -- Write the id sequence to a temporary file
    temporary_text_file(out_f);
    Set Output(out_f);
   put_id_seq(s);
    Set Output(Standard Output);
   Close(out f);
    -- Read it in as a text string from the temporary file
    temporary raw text file(in f);
    result := get(in f);
    remove last char(result);
    raw text file pkg.Delete(in f);
    return result;
```

```
exception
    when others =>
      PUT LINE(STANDARD ERROR, "error: bad file name");
    return result;
  end Get Psdl Id Sequence Text;
  function Get_Output_Guards_Text(name : in op id; c : in
psdl component) return text is
    -- the output string consists of
    -- all the output guards and requirements
    -- traces associated with the vertex
    -- "name" in the control constraint
    -- of c
    out_f: text_file;
    in f: raw text file;
    result: text;
  begin
    -- Write the id sequence to a temporary file
    temporary text file(out f);
    Set Output(out f);
    put output guard(name, c);
    Set_Output(Standard_Output);
    Close(out f);
    -- Read it in as a text string from the temporary file
    temporary_raw_text_file(in_f);
    result := get(in f);
    remove last char(result);
    raw_text_file_pkg.Delete(in_f);
    return result;
  exception
    when others =>
      PUT_LINE(STANDARD_ERROR, "error: bad file name");
    return result;
  end Get_Output Guards Text;
  function Get_Exception_Triggers_Text(name : in op_id; c : in
psdl_component) return text is
    -- the output string consists of all the
    -- exception triggers and requirements
    -- traces associated with the vertex
    -- "name" in the control constraint
    -- of c
    out f: text file;
    in f: raw text file;
    result: text;
  begin
    -- Write the id sequence to a temporary file
    temporary_text_file(out f);
    Set_Output(out f);
   put_excep trigger(name, c);
    Set_Output(Standard Output);
    Close(out_f);
```

```
132
```

```
-- Read it in as a text string from the temporary file
    temporary_raw text file(in f);
    result := get(in f);
    remove last char(result);
    raw_text_file_pkg.Delete(in f);
    return result;
  exception
    when others =>
      PUT_LINE(STANDARD_ERROR, "error: bad file name");
    return result;
  end Get_Exception_Triggers_Text;
  function Get_Exception_List_Text(name : in op_id; c : in
psdl_component) return text is
    -- the output string consists of all the
    -- exception triggers and requirements
    -- traces associated with the vertex
    -- "name" in the control constraint
    -- of c
    out_f: text file;
    in f: raw text file;
    result: text;
  begin
    -- Write the id sequence to a temporary file
    temporary text file(out f);
    Set_Output(out f);
    put id set(exceptions(c), "EXCEPTIONS", exception reqs map(c));
    Set Output(Standard Output);
    Close(out f);
    -- Read it in as a text string from the temporary file
    temporary_raw_text_file(in f);
    result := get(in_f);
    raw_text_file_pkg.Delete(in_f);
    remove last char(result);
    return result;
  exception
    when others =>
      PUT_LINE(STANDARD_ERROR, "error: bad file name");
    return result;
  end Get Exception List Text;
  function Get Timer_Operations_Text(name : in op_id;
                                     c : in psdl component)
    return text is
    -- the output string consists of all the
    -- timer operations and requirements
    -- traces associated with the vertex "name"
    -- in the control constraint
    -- of c
    out f: text file;
    in f: raw text file;
    result: text;
```

```
begin
  -- Write the id sequence to a temporary file
  temporary_text_file(out f);
  Set_Output(out_f);
  put timer op(name, c);
  Set_Output(Standard Output);
  Close(out f);
  -- Read it in as a text string from the temporary file
  temporary_raw_text_file(in_f);
  result := get(in_f);
  remove_last_char(result);
  raw_text_file_pkg.Delete(in_f);
  return result;
exception
  when others =>
    PUT LINE(STANDARD_ERROR, "error: bad file name");
  return result;
end Get_Timer_Operations_Text;
function Get_Type_Name_Text(name : in type_name) return text is
  -- the output string represents the type name
  -- as a sequence of non-white characters
  out_f: text file;
  in_f: raw text file;
  result: text;
begin
  -- Write the id sequence to a temporary file
  temporary_text file(out f);
  Set Output(out_f);
  put type_name(name);
  Set_Output(Standard Output);
  Close(out_f);
  -- Read it in as a text string from the temporary file
  temporary raw_text file(in_f);
  result := get(in f);
  remove last_char(result);
  raw_text_file_pkg.Delete(in f);
  return result;
exception
 when others =>
   PUT_LINE(STANDARD_ERROR, "error: bad file name");
  return result;
end Get_Type_Name_Text;
function Get_Op_Spec_Text(o : in operator) return Text is
 -- gets the specification of the operator as a text string.
 out f: text_file;
 in f: raw text file;
 result: text;
```
```
begin
    -- Write the id sequence to a temporary file
    temporary_text_file(out_f);
    Set Output(out f);
    put component spec(o);
    Set Output(Standard_Output);
    Close(out f);
    -- Read it in as a text string from the temporary file
    temporary raw text file(in f);
    result := get(in f);
    remove last_char(result);
    raw text file pkg.Delete(in f);
    return result;
  exception
    when others =>
      PUT LINE(STANDARD ERROR, "error: bad file name");
    return result;
  end Get Op_Spec_Text;
  function Get Psdl Types Text(prog : in psdl program) return Text is
    -- prog is a psdl_program that contains only user-defined types
    out f: text file;
    in f: raw text file;
    result: text;
  begin
    -- Write the id sequence to a temporary file
    temporary text file(out f);
    Set_Output(out_f);
    put(out f, prog);
    Set Output(Standard Output);
    Close(out_f);
    -- Read it in as a text string from the temporary file
    temporary raw_text_file(in f);
    result := get(in f);
    remove_last_char(result);
    raw_text_file_pkg.Delete(in_f);
    return result;
  exception
    when others =>
      PUT LINE(STANDARD ERROR, "error: bad file name");
    return result;
  end Get_Psdl_Types_Text;
  procedure separate types (prototype: in psdl program; root name: in
psdl id;
                           types: in out psdl program) is
    -- Separates the operators from the types.
    procedure bind(name: in psdl_id; module: in psdl_component;
program: in out psdl program)
      renames psdl_program_map_pkg.bind;
      -- Make sure we use the internal bind operation.
      -- Means ops and types do not have valid parent pointers,
```

```
-- and that get_definition will not
    -- work for data type operations.
begin
  for id : psdl_id, c : psdl_component in
      psdl_program_map_pkg.scan(prototype) loop
      if component_category(c) = psdl type then
         bind(id, c, types);
      end if;
  end loop;
end separate_types;
procedure create_root_name(prototype_name: in string;
                           prototype: in out psdl program;
                           root name: out psdl id) is
  -- Produces the name of the root node if there is one,
  -- otherwise constructs one from the prototype name
  -- and creates a corresponding root node.
  obsolete roots: id set := id set pkg.empty;
  root, op: operator;
  old root name: psdl_id;
begin
  root_name := find root(prototype);
 begin -- check root name
    if prefix(root_name) /= prototype_name then
       root_name := convert(prototype_name, Get_Unique_Id);
       root := fetch(prototype, root_name);
       set_name(root, root name);
    end if;
  exception
   when constraint error => -- The root does not
                              -- have an op_num suffix.
       root_name := convert(prototype_name, Get_Unique_Id);
       root := fetch(prototype, root name);
       set_name(root, root_name);
  end;
exception
 when no root =>
    root_name := convert(prototype_name, Get_Unique Id);
   root := make_composite_operator(root_name);
    add(root, prototype);
 when multiple roots =>
   root_name := convert(prototype_name, Get_Unique_Id);
   old root_name := to_psdl_id(prototype_name);
   root := fetch(prototype, old_root name);
   if root = null component then root :=
              make composite operator(root_name);
   else set_name(root, root_name);
   end if;
   for id : psdl_id, c : psdl_component in
                      psdl_program_map_pkg.scan(prototype) loop
        if component_category(c) = psdl_operator and
          parent(c) = null component then
           id_set_pkg.add(id, obsolete roots);
           op := c;
```

```
-- recycle(op);
          end if;
      end loop;
      remove(obsolete roots, prototype);
      add(root, prototype);
  end create root name;
  function make vertex renaming(g: psdl graph;
                                 renaming: in substitution map)
            return vertex substitution map is
    result: vertex substitution map := empty;
    op num: ge op id;
    new id: psdl id;
    new vertex: op_id;
 begin
    for child vertex: op id in op id set pkg.scan(vertices(g)) loop
        if not eq(child_vertex, external) then
           if member(child vertex.operation name, renaming) then
              -- We have an old style vertex, the op id operation name
              -- matches the psdl id in the original definition.
              new vertex := child vertex;
              -- get the renamed operation name.
              new id := fetch(renaming, child vertex.operation name);
              -- add the v num.
              op_num := ge_op_id(Get Unique Id);
              new_id := to_psdl_id(to_text(new_id), op_num);
              -- install the transformed operation name.
              new vertex.operation name := new id;
              bind(child vertex, new vertex, result);
           elsif
               member(to psdl id(prefix(child vertex.operation name)),
                     renaming) then
              -- Somehow, we got a new-style graph
              -- with vertex numbers but
              -- an old-style operator definition without op num's.
              new vertex := child vertex;
              -- get the renamed operation name.
              new id := fetch(renaming,
to psdl_id(prefix(child vertex.operation name)));
              -- add the v num.
              op_num := suffix(child_vertex.operation_name);
              new_id := to_psdl_id(to_text(new_id), op_num);
              -- install the transformed operation name.
              new vertex.operation name := new id;
              bind(child vertex, new vertex, result);
           end if;
        end if;
    end loop;
    return result;
  end make vertex renaming;
  function member(id: psdl_id; vertices: op id set) return boolean is
 begin
    for oid: op_id in op_id_set_pkg.scan(vertices) loop
```

```
137
```

```
if eq(oid.operation_name, id) then return true; end if;
    end loop;
    return false;
  end member;
  procedure check suffixes (prototype: in out psdl program) is
    op_num: ge_op_id;
    old id, new id: psdl_id;
    renaming: substitution map := empty;
    co: composite operator;
    result: psdl_program := empty_psdl_program;
  begin
    -- Find the bad names, generate names with suffixes, construct a
renaming map,
    -- change the names of the components, and bind them into the
result map.
    for id : psdl_id, c : psdl_component in
psdl_program_map_pkg.scan(prototype) loop
        if component_category(c) = psdl_operator then
           begin
             old id := name(c);
             op_num := suffix(old id);
             -- If this works the suffix exists.
             if parent(c) /= null_component and then
                member(old_id, vertices(graph(parent(c))))
             then
-- We have an old-style graph and definition, which happened
-- to have a numeric suffix. To prevent the original numeric suffix
-- from disappearing from sight, we add another invisble suffix.
                loop
                  op_num := ge_op_id(Get_Unique_Id);
                  new_id := to_psdl_id(to_text(old_id), op_num);
                  exit when not member(new_id, result);
                end loop;
                set_name(c, new id);
                add(c, result);
                bind(old id, new id, renaming);
             else add(c, result);
             end if;
           exception
             when constraint_error =>
               -- The component does not have an op_num suffix.
               -- We need to create and install a suffix.
               loop
                 op_num := ge_op_id(Get_Unique_Id);
                 new_id := to_psdl_id(to_text(old_id), op_num);
                 exit when not member (new id, result);
               end loop;
               set_name(c, new_id);
               add(c, result);
              bind(old_id, new_id, renaming);
           end;
       else add(c, result);
       end if;
```

```
end loop;
    -- Now apply the renaming substitution to
    -- the graphs and control constriants
    -- of all the composite operators in the result map.
    for id : psdl id, c : psdl component in
        psdl program map pkg.scan(result) loop
        if component_category(c) = psdl_operator and then
           component_granularity(c) = composite
        then co := c;
             rename_vertices(co, make_vertex_renaming(graph(co),
                             renaming));
        end if;
    end loop;
    assign(prototype, result);
 end check suffixes;
  function to psdl_id_set(s: psdl_id_sequence) return psdl_id_set is
    -- converts the sequence to a set.
    result: psdl id set := empty;
 begin
    for id: psdl id in psdl id sequence pkg.scan(s) loop
        add(id, result);
   end loop;
    return result;
 end to psdl id set;
 function find name(op spec: text) return psdl id is
   str: string := to string(op_spec);
   len: natural := length(op spec);
   id char set: character set := Alphanumeric set or To Set(' ');
   first, last: natural;
 begin
   Find Token(str, id_char_set, inside, first, last);
   -- OPERATOR keyword
   Find_Token(str(last+1 .. len), id char set, inside, first, last);
   -- The name
   return to_psdl_id(head(str(first .. last), 1 + last - first));
 end find name;
 function to operator(spec: text) return operator is
   -- returns an atomic operator with the given name and psdl
specification.
   f: text file;
   prog : psdl program;
   name: psdl id := find name(spec);
 begin
   temporary text file(f);
   put(f, to string(spec));
   new line(f);
   put_line(f, "IMPLEMENTATION ADA " & to string(name) & " END");
   reset(f, input);
   get(f, prog);
   delete(f);
   return fetch(prog, find root(prog));
```

```
exception
  when others =>
    PUT_LINE(STANDARD_ERROR, "to operator: I/O error");
    return make_atomic_operator(name);
end to operator;
function to_type_name(s: c_string) return type_name is
  -- converts the string to a psdl type name.
  f: text file;
  result: type name;
begin
  temporary text file(f);
  put(f, value(s));
  reset(f, input);
  get(f, result);
  delete(f);
  return result;
exception
  when others =>
    PUT_LINE(STANDARD_ERROR, "to_type_name: I/O error");
    return null type;
end to type name;
function to_op_id(label: text) return op_id is
  -- converts the string to a psdl op_id.
  f: text_file;
  result: op id;
begin
  -- convert the label to an op_id
  temporary text file(f);
 put(f, to_string(label));
  reset(f, input);
  get(f, result);
  delete(f);
  return result;
exception
  when others =>
    PUT_LINE(STANDARD_ERROR, "to_op id: I/O error");
    return empty;
end to op id;
function is_blank(s: string) return boolean is
begin
  for i in s'range loop
      if not is control(s(i)) then return false; end if;
 end loop;
  return true;
end is blank;
```

function to_expression(s: c_string) return expression is
 -- converts the string to a psdl expression.

```
f: text_file;
    result: expression;
    str: string := value(s);
  begin
    if is blank(str) then return undefined expression; end if;
    temporary text file(f);
    put(f, str);
    reset(f, input);
    get(f, result);
    delete(f);
    return result;
  exception
    when others =>
      PUT LINE(STANDARD_ERROR, "to expression: I/O error");
      return undefined expression;
  end to expression;
  procedure to exceptions(s: in string; exc list: out psdl id set;
                          exc_rb: out spec_req_map) is
    f: text file;
  begin
    temporary text file(f);
    put(f, s);
    reset(f, input);
    get(f, exc_list, exc_rb);
    delete(f);
  exception
    when others =>
      PUT LINE(STANDARD ERROR, "to exceptions: I/O error");
      exc list := empty;
      exc rb := empty;
  end to_exceptions;
  procedure to out guard map(s: in string; og: out out guard map;
                             ogrb: out spec req map) is
    f: text_file;
  begin
    temporary text file(f);
    put(f, s);
    reset(f, input);
    get(f, og, ogrb);
    delete(f);
  exception
    when others =>
      PUT LINE (STANDARD ERROR, "to out guard map: I/O error");
      og := empty out guard map;
      ogrb := empty;
  end to out guard map;
procedure to exception_guard_map(s: string; eg: out excep trigger map;
                                    egrb: out spec_req_map) is
    f: text file;
  begin
    temporary text_file(f);
```

```
put(f, s);
  reset(f, input);
  get(f, eg, egrb);
  delete(f);
exception
  when others =>
    PUT_LINE(STANDARD_ERROR, "to_exception_guard_map: I/O error");
     eg := empty excep_trigger_map;
     egrb := empty;
end to exception guard map;
procedure to_timer_op_set(s: in string; timer_op: out timer_op_set;
                           resetrb, startrb, stoprb:
                           in out spec req map) is
  f: text file;
begin
  temporary text file(f);
  put(f, s);
  reset(f, input);
  get(f, timer_op, resetrb, startrb, stoprb);
  delete(f);
exception
  when others =>
    PUT LINE (STANDARD ERROR, "to timer op set: I/O error");
    timer op := empty;
    resetrb := empty;
    startrb := empty;
    stoprb := empty;
end to_timer_op_set;
procedure add_output_guards(id : in op id;
                             og : in out_guard_map;
                             o_guard: in out out guard map) is
  tid: output id;
begin
  for oid: output_id,
        e: expression in out_guard_map_pkg.scan(og) loop
    tid := oid;
      tid.op := id;
      bind(tid, e, o_guard);
  end loop;
end add_output_guards;
procedure add_exception_guards(id : in op id;
                            eg : in excep trigger map;
                            e_guard: in out excep_trigger_map) is
  tid: excep id;
begin
  for eid: excep_id,
        e: expression in excep_trigger_map_pkg.scan(eg) loop
      tid := eid;
      tid.op := id;
      bind(tid, e, e_guard);
  end loop;
```

end add_exception_guards;

end PSDL_Utilities_Pkg;

APPENDIX D: PSDL GRAMMAR

```
-- $Header: $
ક્રક
start_symbol
        : psdl
        ;
psdl
      : psdl component
        ;
component
        : data type
        | operator
        ;
data_type
        : TYPE_TOKEN IDENTIFIER
          type_spec
          type_impl
                              •
        ;
                                              · · · ·
type spec
        :
          SPECIFICATION_TOKEN optional generic param optional_type_decl
          op_spec_list functionality END_TOKEN
        ;
optional_generic_param
        : GENERIC TOKEN
          list_of_type_decl
        I
        ;
optional_type_decl
        :
          list_of_type_decl
        L
        ;
op_spec_list
        : op_spec_list OPERATOR_TOKEN IDENTIFIER operator_spec
        L
        ;
operator
        : OPERATOR_TOKEN IDENTIFIER
          operator spec
          operator_impl
        ;
```

operator_spec : SPECIFICATION_TOKEN interface functionality END TOKEN ; interface : interface attribute reqmts trace ï attribute : GENERIC TOKEN list_of_type_decl INPUT TOKEN list_of_type_decl | OUTPUT TOKEN list of type decl | STATES TOKEN list_of_type_decl INITIALLY TOKEN initial expression list | EXCEPTIONS_TOKEN id_list | MAXIMUM_TOKEN EXECUTION_TOKEN TIME_TOKEN time ; -- Initialization of the_type_decl is done by the callers of this rule. list_of_type_decl : list_of_type_decl ',' type decl | type decl ; type decl : id_list ':' type_name ; type_name : IDENTIFIER '[' list_of_type_decl ']' | IDENTIFIER ; id list : id list ',' IDENTIFIER | IDENTIFIER ; reqmts_trace -- Ignored in this version. : REQ_BY_TOKEN id list ;

functionality

```
: keywords informal desc formal desc
        ;
keywords
        : KEYWORDS_TOKEN id_list
        1
        ;
informal_desc
        : DESCRIPTION_TOKEN TEXT TOKEN
        1
formal_desc
        : axioms TOKEN TEXT TOKEN
        L
        ;
type_impl
        : IMPLEMENTATION TOKEN ADA TOKEN IDENTIFIER END TOKEN
        | IMPLEMENTATION TOKEN type name op impl list END TOKEN
        ;
op impl list
        : op_impl_list OPERATOR_TOKEN IDENTIFIER operator_impl
        I
        ;
operator_impl
        : IMPLEMENTATION TOKEN ADA TOKEN IDENTIFIER END TOKEN
        | IMPLEMENTATION_TOKEN psdl_impl END_TOKEN
        ;
psdl impl
        : data_flow_diagram
          streams
          timers
          control_constraints
          informal_desc
        ;
data flow diagram
        :
          GRAPH_TOKEN vertex_list edge list
        ;
                -- Time is the maximum execution time.
vertex list
        : vertex_list VERTEX_TOKEN op id optional time graph properties
        Ι
        ;
                -- Time is the latency.
```

```
edge_list
        : edge_list EDGE_TOKEN IDENTIFIER
          optional_time op_id ARROW op_id graph properties
        ;
graph properties
        : graph_properties PROPERTY_TOKEN IDENTIFIER '=' expression
        1
        î
op_id
        : operator_name opt_arg
        ;
operator_name
        : IDENTIFIER '.' IDENTIFIER
        | IDENTIFIER
        ;
opt arg
        : '(' optional_id_list '|' optional_id_list ')'
        ;
optional id list
        : id list
        ;
optional_time
   : ':' time
        L
        ï
streams
        : DATA_TOKEN STREAM TOKEN
         list_of_type_decl
        1
        ;
                                   -- The order of id's is not important, so we use psdl_id_set
-- as the data structure to store the timers.
              _____
                                               _____
timers
       : TIMER_TOKEN id_list
       T
       ;
control_constraints
       : CONTROL_TOKEN CONSTRAINTS_TOKEN
```

```
148
```

```
constraints
        ;
constraints
        : constraints OPERATOR TOKEN op id
          opt_trigger opt_period opt finish within
          opt mcp opt mrt constraint options
        | OPERATOR_TOKEN op_id
          opt trigger opt period opt finish within
          opt mcp opt mrt constraint options
        ;
constraint options
        : constraint_options OUTPUT_TOKEN
          id_list IF_TOKEN expression reqmts trace
        | constraint options EXCEPTION TOKEN IDENTIFIER
          opt if predicate reqmts trace
        | constraint options timer op IDENTIFIER
          opt if predicate reqmts trace
        ;
opt_trigger
        : TRIGGERED_TOKEN trigger opt_if_predicate reqmts trace
        E
        ;
trigger
        : BY ALL TOKEN id list
        | BY SOME TOKEN id list
        ;
opt period
        : PERIOD TOKEN time reqmts trace
        t
        ;
opt finish within
        : FINISH_TOKEN WITHIN_TOKEN time reqmts_trace
        1
        ;
opt_mcp
        : MINIMUM_TOKEN CALL_PERIOD_TOKEN time reqmts_trace
        L
        ;
opt mrt
        : max_resp_time time reqmts_trace
        1
        ;
```

max_resp_time

: MAXIMUM_TOKEN RESPONSE_TOKEN TIME_TOKEN ; timer_op : RESET TOKEN START TOKEN STOP TOKEN ; opt_if predicate : IF TOKEN expression 1 ; -- The expression sequence -- is used by procedure bind initial state together with -- the states map to construct the init map. _____ initial expression list : initial_expression_list ',' initial_expression | initial expression ; ____ -- There is one and only one initial state(initial expression) -- for each state variable. This production returns one -- expression to the parent rule corresponding to one state. -- This is done by using the internal stack (\$\$ convention). _____ initial expression : TRUE | FALSE | INTEGER LITERAL REAL LITERAL | STRING LITERAL | IDENTIFIER | type_name '.' IDENTIFIER | type_name '.' IDENTIFIER '(' initial_expression_list ')' | '(' initial_expression ')' | initial_expression log op initial expression %prec logical operator | initial_expression rel_op initial expression %prec relational operator | '-' initial_expression %prec unary_adding_operator | '+' initial expression %prec unary adding operator | initial_expression_bin_add_op_initial_expression %prec binary_adding operator | initial expression bin_mul_op initial_expression %prec multiplying operator

```
| initial expression EXP TOKEN initial expression
          %prec highest_precedence_operator
        | NOT_TOKEN initial_expression
          %prec highest precedence operator
        | ABS TOKEN initial expression
          %prec highest_precedence_operator
        ;
log_op
        : AND_TOKEN
        | OR_TOKEN
        | XOR TOKEN
        ;
rel_op
        : '<'
        | '>'
        j '='
        | GREATER THAN OR EQUAL
        LESS THAN OR EQUAL
        | INEQUALITY
        ;
bin_add_op
        : '+'
        1 '-'
        1 '&'
        ;
bin_mul_op
       : '*'
        1 1/1
        | MOD_TOKEN
        | REM_TOKEN
        ;
time
        : time number MICROSEC TOKEN
        | time_number MS_TOKEN
        | time_number SEC_TOKEN
        | time_number MIN_TOKEN
        | time number HOURS TOKEN
        ;
time_number
        : INTEGER LITERAL
        ;
expression list
        : expression list ',' expression
        | expression
        ;
```

```
151
```

```
-- Expressions can appear in guards appearing in control constraints.
-- These guards can be associated with triggering conditions, or
-- conditional outputs, conditional exceptions, or conditional timer
-- operations. Similar to initial expression, except that tim e values
-- and references to timers and data streams are allowed.
                               ~------
expression
       : TRUE
       | FALSE
       | INTEGER LITERAL
       | REAL LITERAL
       STRING LITERAL
        | IDENTIFIER
        -- The only difference from the initial expression
       | time
       | type_name '.' IDENTIFIER
       type_name '.' IDENTIFIER '(' expression_list ')'
       | '(' expression ')'
       | expression log_op expression %prec logical_operator
       | expression rel_op expression %prec relational_operator
       | '-' expression
                                      %prec unary adding operator
       | '+' expression
                                       %prec unary_adding_operator
       | expression bin_add_op expression
         %prec binary adding operator
       | expression bin_mul_op expression
         %prec multiplying operator
       | expression EXP_TOKEN expression
         %prec highest_precedence_operator
       | NOT TOKEN expression
         %prec highest_precedence_operator
       | ABS_TOKEN expression
         %prec highest_precedence_operator
```

응용

;

APPENDIX E: TEST DATA

TYPE STACK SPECIFICATION GENERIC types : private type 2 : public OPERATOR PUSH SPECIFICATION INPUT I : INTEGER INPUT S : STACK OUTPUT S : STACK END OPERATOR POP SPECIFICATION INPUT S : STACK OUTPUT I : INTEGER OUTPUT S : STACK END OPERATOR Empty SPECIFICATION OUTPUT dummy : STACK END KEYWORDS stack, adt DESCRIPTION {This is a generic stack adt} AXIOMS $\{push (s,x) = s::x\}$ END IMPLEMENTATION ADA STACK END **OPERATOR Compute 8** SPECIFICATION INPUT X1 : INTEGER INPUT X2 : INTEGER

```
INPUT
```

```
X3 : INTEGER
    OUTPUT
      X1 : INTEGER
    OUTPUT
      X2 : INTEGER
    OUTPUT
      X3 : INTEGER
    OUTPUT
      DC : INTEGER
    OUTPUT
      DX : FLOAT
    OUTPUT
      DY : BOOLEAN
    EXCEPTIONS
      E1
    EXCEPTIONS
      E2
    KEYWORDS
      software, bubbles
    DESCRIPTION
      {This is an atomic bubble.
       Line 2.
       Line 3. }
    AXIOMS
      {P = NP}
END
IMPLEMENTATION ADA Compute_8
END
OPERATOR Consumer 4
SPECIFICATION
    INPUT
      DB : INTEGER
    INPUT
      DC : INTEGER
    EXCEPTIONS
      E1
    EXCEPTIONS
      E2
   MAXIMUM EXECUTION TIME 0 MS
END
IMPLEMENTATION ADA Consumer 4
END
OPERATOR OBSOLETE ROOT 0
SPECIFICATION
END
IMPLEMENTATION ADA OBSOLETE_ROOT_0
```

END

OPERATOR OP A 135

SPECIFICATION INPUT DX : FLOAT END IMPLEMENTATION ADA OP A 135 END OPERATOR OP B 136 SPECIFICATION INPUT DY : BOOLEAN END IMPLEMENTATION ADA OP_B_136 END OPERATOR Process Data 5 SPECIFICATION INPUT DA : Missing Info OUTPUT DB : INTEGER OUTPUT DC : INTEGER OUTPUT DX : FLOAT OUTPUT DY : BOOLEAN STATES X2 : INTEGER INITIALLY 2 STATES X1 : INTEGER INITIALLY 1 STATES S : STACK INITIALLY EMPTY EXCEPTIONS E1 EXCEPTIONS E2 EXCEPTIONS E3 MAXIMUM EXECUTION TIME 150 MS KEYWORDS Compute, Composite, Parent DESCRIPTION {This is an composite bubble. } AXIOMS {P != NP, believe it or not }

```
END
IMPLEMENTATION
 GRAPH
    VERTEX STACK.POP_14(S | DB, S) : 120 MS
      PROPERTY x = 456
      PROPERTY y = 110
      PROPERTY radius = 30
     PROPERTY color = 62
      PROPERTY label font = 0
      PROPERTY label x offset = - 25
     PROPERTY label y offset = 84
     PROPERTY met font = 0
     PROPERTY met_x_offset = 35
     PROPERTY met_y_offset = - 10
     PROPERTY is terminator = FALSE
   VERTEX STACK.PUSH 13(DA, S | S)
     PROPERTY x = 210
     PROPERTY y = 110
     PROPERTY radius = 30
     PROPERTY color = 62
     PROPERTY label font = 0
     PROPERTY label_x_offset = -43
     PROPERTY label_y_offset = 82
     PROPERTY met font = 0
     PROPERTY met x offset = 244
     PROPERTY met_y_offset = 78
     PROPERTY is terminator = FALSE
   VERTEX Compute 8 19
     PROPERTY x = 330
     PROPERTY y = 341
     PROPERTY radius = 30
     PROPERTY color = 62
     PROPERTY label font = 0
     PROPERTY label x offset = 3
     PROPERTY label_y_offset = 34
     PROPERTY met font = 0
     PROPERTY met x offset = 298
     PROPERTY met_y_offset = 205
     PROPERTY is_terminator = FALSE
   EDGE S
     STACK.POP 14(S | DB, S) \rightarrow
     STACK.PUSH 13(DA, S | S)
       PROPERTY id = 33
       PROPERTY label font = 0
       PROPERTY label_x_offset = - 4
       PROPERTY label_y_offset = 5
       PROPERTY latency_font = 0
       PROPERTY latency x offset = 318
       PROPERTY latency_y_offset = 182
       PROPERTY spline = "362 71 "
```

```
156
```

```
EDGE DB
  STACK.POP 14(S | DB, S) ->
  EXTERNAL
    PROPERTY id = 32
    PROPERTY label font = 0
    PROPERTY label_x_offset = 3
    PROPERTY label_y_offset = 1
    PROPERTY latency font = 0
    PROPERTY latency x offset = 292
    PROPERTY latency_y_offset = 243
    PROPERTY spline = "593 102 675 137 "
EDGE DA
  EXTERNAL ->
  STACK.PUSH 13(DA, S | S)
    PROPERTY id = 31
    PROPERTY label font = 0
    PROPERTY label_x_offset = - 19
    PROPERTY label_y_offset = - 17
    PROPERTY latency_font = 0
    PROPERTY latency x offset = 131
    PROPERTY latency_y_offset = 123
    PROPERTY spline = "46 137 142 107 "
EDGE X1
  Compute 8 19 ->
  Compute 8 19
    PROPERTY id = 51
    PROPERTY label font = 0
    PROPERTY label x offset = 15
    PROPERTY label_y_offset = 24
    PROPERTY latency font = 0
    PROPERTY latency x offset = 124
    PROPERTY latency y offset = 249
PROPERTY spline = "437 402 403 443 "
EDGE X2
  Compute 8 19 ->
  Compute 8 19
    PROPERTY id = 51
    PROPERTY label font = 0
    PROPERTY label_x_offset = - 29
    PROPERTY label_y_offset = 18
    PROPERTY latency_font = 0
    PROPERTY latency x offset = 224
    PROPERTY latency y offset = 149
PROPERTY spline = "287 407 318 444 "
EDGE S
  STACK.PUSH 13(DA, S | S) ->
  STACK.POP 14(S | DB, S)
    PROPERTY id = 34
    PROPERTY label font = 0
    PROPERTY label_x_offset = -1
```

```
157
```

```
PROPERTY label_y_offset = 12
    PROPERTY latency_font = 0
    PROPERTY latency_x_offset = 154
    PROPERTY latency_y_offset = 196
    PROPERTY spline = "365 185 "
EDGE X3
  Compute 8 19 ->
  Compute 8 19
    PROPERTY id = 50
    PROPERTY label font = 0.
    PROPERTY label x offset =
                               - 10
    PROPERTY label_y_offset = - 10
    PROPERTY latency_font = 0
    PROPERTY latency_x_offset = 124
    PROPERTY latency_y_offset = 149
    PROPERTY spline = "330 299 383 297 "
EDGE DC
  Compute_8_19 ->
  EXTERNAL
    PROPERTY id = 2909
    PROPERTY label_font = 2
    PROPERTY label_x_offset = 0
    PROPERTY label_y_offset = 0
    PROPERTY latency_font = 2
    PROPERTY latency_x_offset = 0
    PROPERTY latency_y_offset = 15
    PROPERTY spline = "439 309 496 266 "
EDGE DX
  Compute 8 19 ->
  EXTERNAL
    PROPERTY id = 2910
    PROPERTY label font = 2
    PROPERTY label x offset = 0
    PROPERTY label_y_offset = 0
    PROPERTY latency_font = 2
    PROPERTY latency_x_offset = 0
    PROPERTY latency_y_offset = 15
    PROPERTY spline = "482 360 574 351 "
EDGE DY
  Compute 8 19 ->
  EXTERNAL
    PROPERTY id = 2911
    PROPERTY label font = 2
    PROPERTY label x offset = 0
    PROPERTY label_y_offset = 0
    PROPERTY latency font = 2
    PROPERTY latency x offset = 0
    PROPERTY latency_y_offset = 15
    PROPERTY spline = "477 400 566 423 "
```

```
DATA STREAM
    X3 : INTEGER
  CONTROL CONSTRAINTS
      OPERATOR STACK.POP 14(S | DB, S)
        MINIMUM CALLING PERIOD 100 MS
        MAXIMUM RESPONSE TIME 4500 MICROSEC
        OUTPUT
          S
          IF S /= STACK.EMPTY
        OUTPUT
          DB
          IF DB > 0
        EXCEPTION E3
          IF S = STACK.EMPTY
        EXCEPTION E4
          IF DB < 0
        START TIMER
           Timer1
        RESET TIMER
          Timer2
      OPERATOR STACK. PUSH_13(DA, S | S)
        OUTPUT
          S
          IF S /= STACK.EMPTY
        EXCEPTION E3
          IF S = STACK.EMPTY
        START TIMER
           Timer1
      OPERATOR Compute_8_19
END
OPERATOR Producer 1
SPECIFICATION
    GENERIC
      G1 : FLOAT
    OUTPUT
     DA : Missing_Info
    MAXIMUM EXECUTION TIME 0 MS
END
IMPLEMENTATION ADA Producer 1
END
OPERATOR objects 2
SPECIFICATION
END
IMPLEMENTATION
  GRAPH
    VERTEX Process_Data_5_20 : 150 MS
```

```
159
```

```
PROPERTY x = 183
  PROPERTY y = 139
  PROPERTY radius = 30
  PROPERTY color = 62
  PROPERTY label_font = 0
  PROPERTY label_x_offset = 5
  PROPERTY label_y_offset = 44
  PROPERTY met_font = 0
  PROPERTY met x offset = 8
  PROPERTY met_y offset = -12
  PROPERTY is terminator = FALSE
VERTEX Producer 1 12 : 0 MS
  PROPERTY x = 14
  PROPERTY y = 79
  PROPERTY radius = 30
  PROPERTY color = 62
  PROPERTY label font = 0
  PROPERTY label_x_offset = 20
  PROPERTY label_y_offset = 39
  PROPERTY met font = 0
  PROPERTY met x offset = 63
  PROPERTY met_y_offset = _ - 5
  PROPERTY is terminator = TRUE
VERTEX Consumer 4 11 : 0 MS
  PROPERTY x = 353
  PROPERTY y = 205
  PROPERTY radius = 30
  PROPERTY color = 62
  PROPERTY label font = 0
  PROPERTY label x offset = 14
  PROPERTY label_y_offset = 33
  PROPERTY met_font = 0
  PROPERTY met_x_offset = 63 .
  PROPERTY met y offset = -5
  PROPERTY is terminator = TRUE
VERTEX OP A 135 129
  PROPERTY x = 114
  PROPERTY y = 306
  PROPERTY radius = 30
  PROPERTY color = 62
  PROPERTY label font = 2
  PROPERTY label x offset = 28
  PROPERTY label y offset = 40
  PROPERTY met font = 2
  PROPERTY met x offset = 144
  PROPERTY met_y_offset = 306
  PROPERTY is terminator = FALSE
VERTEX OP B 136 131
  PROPERTY x = 427
  PROPERTY y = 77
```

```
PROPERTY radius = 30
  PROPERTY color = 62
  PROPERTY label font = 2
  PROPERTY label_x_offset = 20
  PROPERTY label y_offset = 40
  PROPERTY met font = 2
  PROPERTY met x offset = 457
  PROPERTY met_y_offset = 77
  PROPERTY is terminator = FALSE
EDGE DA
  Producer_1_{12} \rightarrow
  Process Data 5 20
    PROPERTY id = 30
    PROPERTY label font = 0
    PROPERTY label x offset = 6
    PROPERTY label y_offset = -8
    PROPERTY latency_font = 0
    PROPERTY latency x offset = 124
    PROPERTY latency y offset = 149
PROPERTY spline = "143 140 "
EDGE DB
  Process Data 5 20 ->
  Consumer 4 11
    PROPERTY id = 40
    PROPERTY label font = 0
    PROPERTY label x offset = 1
    PROPERTY label_y_offset = - 6
    PROPERTY latency_font = 0
    PROPERTY latency_x_offset = 269
    PROPERTY latency y offset = 204
PROPERTY spline = "297 199 "
EDGE DC
  Process Data 5 20 ->
  Consumer 4 11
    PROPERTY id = 42
    PROPERTY label font = 0
    PROPERTY label x_offset = -8
    PROPERTY label_y_offset =
                                 - 11
    PROPERTY latency_font = 0
    PROPERTY latency x offset = 0
    PROPERTY latency_y_offset = 16
    PROPERTY spline = "302 358 "
EDGE DX
  Process_Data_5_20 ->
  OP A 135 129
    PROPERTY id = 130
    PROPERTY label font = 2
    PROPERTY label x offset = -27
    PROPERTY label_y_offset = - 6
    PROPERTY latency_font = 2
```

```
PROPERTY latency_x_offset = 184
      PROPERTY latency_y_offset = 237
      PROPERTY spline = "169 250 "
  EDGE DY
    Process_Data_5_20 ->
    OP B 136 131
      PROPERTY id = 134
      PROPERTY label font = 2
      PROPERTY label_x_offset = - 17
      PROPERTY label_y_offset = - 15
      PROPERTY latency_font = 2
     PROPERTY latency_x_offset = 288
      PROPERTY latency_y_offset = 165
      PROPERTY spline = "335 138 "
DATA STREAM
 DY : BOOLEAN,
  DX : FLOAT,
  DC : INTEGER,
  DB : INTEGER,
  DA : Missing Info
CONTROL CONSTRAINTS
    OPERATOR Process_Data_5_20
   OPERATOR Producer_1_12
     PERIOD 4000 MS
     FINISH WITHIN 100 MS
   OPERATOR Consumer_4_11
     TRIGGERED BY ALL
       DB
        IF DB > 0
   OPERATOR OP_A_135_129
   OPERATOR OP_B_136_131
```

END

LIST OF REFERENCES

- 1. San Jose Mercury News, Section C, Tech Ticker, September 6, 1997.
- 2. Anderson, Steven E., Functional Specification for a Generic C3I Workstation, Master's Thesis, Naval Postgraduate School, Monterey, California, September 90.
- 3. CAPS Interface Manual, Software Engineering Group, Naval Postgraduate School, Monterey, California.
- 4. CAPS Executive Briefing, Software Engineering Group, Naval Postgraduate School, Monterey, California.
- 5. Berzins, V., Luqi, Software Engineering with Abstractions, Addison Wesley, 1991.
- 6. Aho, A., Kernighan, B., Weinberger, P., *The Awk Programming Language*, Addison Wesley, 1988.
- 7. Feldman, M., Koffman, E., Ada95: Problem Solving and Program Design, Addison Wesley, 1991.
- 8. Booch, G., Bryan, D., *Software Engineering with Ada*, Benjamin/Cummings, Redwood City, California, 1995.
- 9. Barnes, J., Programming in Ada95, Addison Wesley, 1996.
- 10. Luqi, "Software Evolution Through Rapid Prototyping", IEEE Transactions on Software Engineering, May 1989.
 - 11. Luqi, Berzins, V., Yeh, R., "A Prototyping Language for Real-Time Software," IEEE Transactions on Software Engineering, October 1988.
 - 12. San Jose Mercury News, Section E, Can Silicon Graphics Evolve, June 2, 1997.
 - 13. San Jose Mercury News, Section A, Help Wanted: Tech Grads, February 2, 1997.
 - 14. The New York Times, Cybertimes, 1 Teensy Little Bug, 1 Humongous Crash, December 1, 1996.
 - 15. Biggerstaff, T., "Moore's Law: Change or Die!", IEEE Software, Vol 13, No. 1, January 1996.
 - 16. San Jose Mercury News, Section A, Radar System Failed in Crash, Software Glitch Prevented Warning, August 11, 1997.
 - 17. J.A. McDermid, ``Safety-critical software: a vignette", Software Engineering Journal, Vol. 8, No. 1, pp. 2-3, 1993

18. Chadwick, C., in personal conversation with the author, June 26, 1996.

INITIAL DISTRIBUTION LIST

1.	Defense Technical Information Center 8725 John J. Kingman Road, Ste 0944 Ft. Belvoir, VA 22060-6218	2
2.	Dudley Knox Library Naval Postgraduate School 411 Dyer Rd. Monterey, CA 93943-5101	2
3.	Director, Training and Education MCCDC, Code C46 1019 Elliot Road Quantico, VA 22134-5027	1
4.	Director, Marine Corps Research Cente MCCDC, Code C40RC 2040 Broadway Street Quantico, VA 22134-5107	er2
5.	Director, Studies and Analysis Division MCCDC, Code C45 3300 Russell Road Quantico, VA 22134-5130	1
6.	Marine Corps Representative Naval Postgraduate School Code 037, Bldg. 234, HA-220 Monterey, CA 93940	

1

7.	Cmdr Michael J. Holden, Code CS/Hm1 Computer Science Department Nava! Postgraduate School Monterey, CA 93943
8.	Marine Corps Tactical Systems Support Activity
9.	Dr. Luqi, Code CS/Lq