## NAVAL POSTGRADUATE SCHOOL Monterey, California

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## **THESIS**

PRELIMINARY FLIGHT SOFTWARE SPECIFICATION FOR THE PETITE AMATEUR NAVY SATELLITE (PANSAT)

by

Teresa Owen Ford

March, 1994

Thesis Advisor: Second Reader: Douglas J. Fouts Frederick W. Terman

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# Preliminary Flight Software Specification For the Petite Amateur Navy Satellite (PANSAT)

by

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Lieutenant, United States Navy
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Submitted in partial fulfillment of the requirements for the degree of

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#### **ABSTRACT**

PANSAT is a small, spread-spectrum, communications satellite under design at the Naval Postgraduate School. It will support a store and forward bulletin board system for use by the amateur radio community. The flight software is responsible for the autonomous telemetry collection and hardware control operations of the satellite, communications and file transfer protocols allowing access to the bulletin board system, and command interpretation and response to ground control commands.

In this thesis, the complete flight software architecture and module interfaces are specified using the Estelle Formal Description Technique. The module bodies dealing with communications and file transfer protocols are specified in detail in Estelle. The current design goals for the remainder of the flight software modules are discussed. Appendices include the preliminary flight software specification itself, a data flow diagram interpretation of the specification, and a summary of the Estelle syntax used.

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## SYMBOL TABLE

SYMBOL	WHERE USED	MEANING
bold	thesis and specification	A Pascal or Estelle reserved word.
italics	thesis and specification	A specification constant, including a member of an enumerated type.
ALL_CAPS	thesis and specification	Module name or state name.
Beginning_Caps	thesis and specification	Name of a user defined type.
'single_quotes'	thesis	Name of a variable or record field in the specification, or the contents of a variable or record field.
0xNN	thesis and specification	A hexadecimal number. "N" stands for the digits 0 through F.
uchar	thesis and specification (primitive data type)	"Unsigned character": 8 bits of binary data, a 1 byte unsigned integer, or 1 ascii character.
uint	thesis and specification (primitive data type)	"Unsigned integer": 16 bits of binary data, a 2 byte unsigned integer, or 2 ascii characters.
ulong	thesis and specification (primitive data type)	"Unsigned long integer": 32 bits of binary data, a 4 byte unsigned integer, or 4 ascii characters.
int	thesis and specification (primitive data type)	"Integer": a signed integer.

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#### I. INTRODUCTION

#### A. PANSAT

The Petite Amateur Navy Satellite (PANSAT) is a small, experimental, communications satellite currently being designed and constructed at the Naval Postgraduate School (NPS) and scheduled to be launched from the space shuttle in 1996. The satellite will use half-duplex, spread-spectrum communications and will support a store and forward bulletin board system for use by the amateur "HAM" radio community. PANSAT will operate autonomously in the performance of many of its functions, while also carrying out commands issued by the ground control station located in Monterey, CA at NPS.

#### **B.** GOALS OF THE FLIGHT SOFTWARE

The flight software will control the autonomous operation of the satellite, allow amateur radio operators access to the onboard bulletin board or "mailbox" system, and provide the means for the satellite to respond to commands from the ground control station. The major functional areas of the flight software are listed in Table 1.1. There are other software functions which are equally important to the PANSAT project, but outside the scope of the flight software. These include the "bootstrap" software which will control initial configuration of the satellite systems upon launch or reboot, "client"

ground station software for use by the amateur radio community, and "commanding" software for the ground control station.

#### TABLE 1.1 FUNCTIONAL AREAS OF THE FLIGHT SOFTWARE

- 1. Communications Protocols for Amateur Radio User Access and File Transfer
- 2. Telemetry Gathering and Storage
- 3. Control of Satellite Hardware Systems
- 4. Command Interpretation and Response to NPS Ground Station Control

#### C. SCOPE OF THIS THESIS

This thesis is the result of the initial attempt at specification of the flight software for PANSAT. The hardware systems are still evolving, which means that hardware control, command interpretation, and telemetry gathering requirements are still being defined. The only requirements which can be completely defined this early and remain essentially unchanged, regardless of the final hardware configuration of the satellite, are the communications protocol and mail handling functions. For this reason, the actual preliminary software specification, which can be found in Appendix A, concentrates on the mailbox system and the transfer protocols for uploading and downloading files. The preliminary architecture for the remaining software modules is sketched out, and interfaces have been defined between all modules.

Some existing third party software, designed specifically for communications satellites, will be used aboard PANSAT to support the applications software being developed at NPS. The commercial software used will be introduced in Chapter II. A PANSAT file header has been developed to assist in proper maintenance of the mail storage system. The elements of this file header will be explained in Chapter III. Chapter IV will discuss the specification language, Estelle, in which Appendix A is written. Chapters V through VIII will explain the functionality of the software modules which have received the most attention in Appendix A, those dealing with file transfer and mailbox control. Chapter IX will introduce the preliminary goals for the remainder of the flight software, which has yet to be specified in detail. Chapter X will contain conclusions and recommendations for further work. Appendix B contains a graphical interpretation of the specification in Appendix A. This interpretation uses data flow diagrams which provide a visual means of identifying the software architecture and module interfaces. Appendix C will provide some details of the syntax of Estelle.

#### II. THIRD PARTY SOFTWARE

#### A. SPACE CRAFT OPERATING SYSTEM

The role of DOS (Disc Operating System) on a personal computer is filled by SCOS (Space Craft Operating System) on the computer aboard PANSAT. SCOS is a real time, multi-tasking, operational environment designed specifically for the needs of a small satellite. [Ref. 1] It supports applications written in the "C" programming language, and many of the primitive functions familiar to "C" programmers are available through the Space Craft Operating System. These include file management capabilities, dynamic memory allocation, bit and byte manipulations, and logical and mathematical operations. Within the SCOS operating environment, the various modules which make up the flight software are running as concurrent processes, able to pass data among themselves. They are put to "sleep" when they are not needed, saving CPU time, and are "woken up" again whenever their services are required.

#### B. BAX

#### 1. AX.25

AX.25 is a variant of the CCITT X.25 link-layer protocol, and is designed to provide reliable data transport between two signaling terminals. [Refs. 2, 3] This asynchronous data transfer protocol is currently in wide use by the amateur radio community for packet communications, and has thus been selected for use in

communications with PANSAT. To incorporate AX.25 functionality into the communications software aboard PANSAT, a program called BAX is employed.

BAX is the BekTek corporation's implementation of the AX.25 protocol. It is designed to work with the Space Craft Operating System. Because of the availability of the BAX software, the actual functionality of the AX.25 protocol is transparent to the applications programs developed for the satellite. The software designer and programmer need only consult the BAX manual [Ref. 4] to discover how to access the capabilities required. For the amateur radio operator, the AX.25 protocol is normally implemented by a piece of dedicated hardware known as a TNC (terminal node controller), or by software contained in a PC (personal computer) based system.

### 2. BAX Application Programs

In order to utilize the capabilities of BAX, a "BAX application" program must be written. In the PANSAT flight software specification, the primary BAX application is known as the DATA\_TRANSFER module (see Chapter V). Other modules, such as GROUND\_CONTROL, may also access the services of BAX.

BAX has the capability of receiving frames addressed to various applications which are distinguished from each other by having different ssid's (subsystem identification numbers). PANSAT will be addressed by a multi-character Amateur Radio Callsign, "PANSAT", for example. This callsign may be modified by use of ssid's to access different functions aboard the satellite. For instance, the data transfer module will have ssid '1', and the HAMs will send their mail to "PANSAT-1". The command interpreter, GROUND\_CONTROL, may have ssid '2', so that commands from the NPS

ground station could be sent there directly by BAX without having to be in the same format as that recognized by the bulletin board system.

BAX communicates directly with the hardware drivers which operate the radio equipment on the satellite. As incoming frames are received, BAX handles all of the AX.25 protocol requirements, and notifies the appropriate application module of the receipt and the source of each transmission. In the case of the data transfer module, the frames passed to it by BAX must be assembled into the packets required at the next higher protocol level. This is the level of the PACKET\_TRANSFER modules described in Chapters VI and VII. When the data transfer module receives a packet from a packet transfer module for transmission to a ground user, the data transfer module breaks the data into frames which it passes "down" to BAX for transmission in accordance with the AX.25 protocol.

#### 3. BAX Functions

The accessing of most BAX functions by PANSAT application modules is represented in the software specification as messages being passed through an Abstract\_Bax\_Channel. In fact in the specification model, all communication between modules is accomplished via "channels", each channel having certain types of messages defined which can be passed through it. These channel and message definitions form the interface specification between software modules (see Chapter III and Appendix A). The BAX functions accessed are listed and briefly explained in Table 2.1.

TABLE 2.1 BAX FUNCTIONS		
Function	Purpose	
qax_input	BAX informs PANSAT application of user connection request, user disconnect, or incoming data frame.	
qax_claim	PANSAT application tells BAX what callsign and said it will be using.	
qax_data	PANSAT application passes data to BAX for transmission.	
qax_busy	PANSAT application informs BAX that it is "busy" and will not be accepting incoming frames.	
qax_unbusy	PANSAT application informs BAX that it is no longer "busy" and will once again accept incoming frames.	
qax_con_acpt	PANSAT application accepts a user connection request	
qax_con_rej	PANSAT application rejects a user connection request	
qax_connect	PANSAT application sends a connection request to ground station	
qax_ui	PANSAT application sends an unnumbered information frame to a ground station.	
qax_ disconnect	PANSAT application disconnects from a ground station.	

#### C. FILE TRANSFER LEVEL 0

File Transfer Level 0 (FTL0) is an asynchronous connected mode file transfer protocol developed by Jeff Ward and Harold E. Price for use with the PACSATs (packet satellites). In a connected mode protocol, there is a virtual link between each user and the satellite, with each transmission having a specific destination. This is in contrast to a broadcast, or unconnected mode protocol, in which communications are intended to be picked up by anyone listening.

An implementation of FTLO is currently available to amateur radio operators in the form of the program "PG" along with several utility programs that work along with it, such as "PHS" and "PFHADD". Although the specification for FTLO contains provisions for both uploading and downloading files from a satellite, only the uploading capabilities are implemented by the current version of "PG". For downloading from the satellites which currently employ FTLO, the non-connected mode, "PACSAT Broadcast Protocol" [Ref. 5], is used. This is implemented by the program "PB" and it's utilities.

The specification for FTLO [Ref. 6] is used as the motivation for the specification of the PACKET\_TRANSFER module aboard PaANSAT (Chapters VI and VII). The specification of the packet transfer module in Appendix A is much more detailed than [Ref. 6], in an attempt to show how the protocol will actually be implemented by the software aboard the satellite at the lowest possible level. The PANSAT implementation will employ an FTLO-like connected-mode protocol for both upload and download.

An effort has been made to remain as compatible as possible with any other FTLO implementation. Some of the specific requirements of PANSAT and certain design decisions have led to some variation from [Ref. 6]. As source code for "PG" was not available, it is unknown at this point whether that software will actually be able to communicate with PANSAT. PANSAT-specific ground station software, capable of communicating with the packet transfer module specified here, will be developed by NPS and made available to the amateur radio community.

#### III. THE PANSAT FILE HEADER

#### A. FUNCTION

Each file maintained in the mail box memory of the satellite must begin with a PANSAT file header. The header includes information such as the file number, file name, file length, source callsign, destination callsigns, upload time, and expiration time. The information in the header is necessary for the proper maintenance and administration of the mail box. It can also be used by a client to determine which files onboard the satellite may be of interest. The select\_cmd makes use of the various fields of the PANSAT file header in its selection criteria (see Chapter VI).

#### B. STRUCTURE

The PANSAT file header is inspired by, but is not the same as, the PACSAT File Header developed by Jeff Ward and Harold Price [Ref. 7]. It is arranged as a variable length array of unsigned characters (bytes). The fields nearest to the beginning of the header have fixed positions and fixed lengths. The lengths of other fields are specified within the header itself, causing the positions of the later fields to be variable, and dependant on the fields ahead of them.

The byte positions, field names and formats are listed in Table 3.1. Note that positions and field lengths through byte 41 are fixed. Each of the fields "Destination 1" through "Destination 7" is either present, with a fixed length of 6 bytes, or absent

"Title" and "Keywords" have variable length, based upon the contents of the fields "Title Length" and "Keyword Length", respectively. The column Const refers to the constant name given to the associated field in the specification of Appendix A.

TABLE 3.1 PANSAT FILE HEADER FIELDS			
Byte(s)	Const	Name	Format
[01] fixed	fl	Flag	<0xBB><0x55>
[25] fixed	mn	Mail Number	ulong
[69] fixed	ml	File Length	ulong
[10] fixed	ft	File Type	uchar
[11] fixed	ct	Compression Type	uchar
[1213] fixed	bo	Body Offset	uint
[14] fixed	dc	Download Count	uchar
[1520] fixed	sc	Source	array[6] of uchar
[21] fixed	pr	Priority	uchar
[2225] fixed	ut	Upload Time	ulong
[2629] fixed	et	Expire Time	ulong
[3037] fixed	na	PANSAT File Name	array[8] of uchar
[3840] fixed	ex	PANSAT File Extension	array[3] of uchar
[41] fixed	nd	Number of Destinations	uchar
[4247] approx.	ds	Destination 1	array[6] of uchar
[4853] approx.		Destination 2	array[6] of uchar
[5459] approx.		Destination 3	array[6] of uchar
[6065] approx.		Destination 4	array[6] of uchar
[6671] approx.		Destination 5	array[6] of uchar

TABLE 3.1 PANSAT FILE HEADER FIELDS			
Byte(s)	Const	Name	Format
[7277] approx.		Destination 6	array[6] of uchar
[7883] approx.	,	Destination 7	array[6] of uchar
[84] approx.		Title Length	uchar
[85114] approx.	ti	Title	array[30] of uchar
[115] approx.		Keyword Length	uchar
[116195] approx.	kw	Keywords	array[80] of uchar
[196197] approx.		Header Checksum	uint
[198199] approx.		Body Checksum	uint

#### C. FIELDS TO BE FILLED IN BY SOURCE

A PANSAT file header must be prepended to any file before it is uploaded to the satellite. Certain fields within the header must be completed by the user station where the file originates, while other fields are filled in by the satellite once the file has been completely uploaded. The user must place all zeros in those fields which the satellite will complete. These satellite responsible fields will all be of fixed length. The fields for which the user is responsible are as follows:

#### 1. Flag

The flag indicates that this is the beginning of a file with a PANSAT file header. The flag must always consist of the same two bytes: '0xBB' followed by '0x55'. Example of the flag field: 10111011 01010101.

## 2. Mail Number

A mail number, or file number, is assigned by PANSAT to each file. This number is not known to the user until it is provided by the satellite in an upload\_go\_resp following an upload\_cmd from the ground station. (See Chapter VI, section E.). The user software has two options here. The simplest is to leave 4 bytes of 0's in this field, and let the satellite update it after the upload. If the upload is interrupted, however, it will be the responsibility of the user software to "remember" the number associated with the partially uploaded file, and to provide it to the satellite in the next upload\_cmd which will continue the process. The obvious place to store the number is in the file header. For this reason, it may make more sense to choose the second option, which is to place the proper number in the header before transmission of the file begins. Of course, this will also necessitate adjusting the header checksum before transmitting. (See subsection 12).

#### 3. File Length

The file length is a four byte unsigned integer (ulong). The source software must place in this field the number of bytes contained in the file, including the PANSAT file header.

## 4. File Type

The file type is a one byte field which indicates the format of the file body. The satellite does not care about the file format, as it treats all files simply as arrays of bytes. The information in this field is for the use of anyone who downloads the file, so that they will know how it must be read. The contents of this field will be interpreted as in Table 3.2. The first eleven of these types are the same as those defined by Price and Ward in [Ref. 7], and some of them might never be used aboard PANSAT. They are included for completeness, and to provide as much parallelism as possible between this specification and FTLO.

TABLE 3.2 FILE TYPES		
Field Contents	File Type	
00000000	ASCII file	
0000001	RLI/MBL message body. Single message.	
00000010	RLI/MBL import/export file. Multiple message.	
00000011	UoSAT Whole Orbit Data	
00000100	Microsat Whole Orbit Data	
00000101	UoSAT CPE Data	
00000110	MS/PC-DOS .exe file	
00000111	MS/PC-DOS .com file	
00001000	Keplerian elements NASA 2-line format	
00001001	Keplerian elements "AMSAT" format	
00001010	Simple ASCII text file, but compressed	

TABLE 3.2 FILE TYPES			
Field Contents	File Type		
10100000	PANSAT short format telemetry file		
10100001	PANSAT long format telemetry file		
10100010	PANSAT bax telemetry file		
11111110	User defined type.		

### 5. Compression Type

If the body of the file is compressed, the source must indicate the type of compression used in this one byte field. Again, the satellite does not care whether or not a file is compressed, or if so, how. This information is for the use of the downloading user only. Note that no matter what file format or compression type is used in the file body, the PANSAT file header will always be uncompressed ASCII text. Compression types are indicated by Table 3.3. "User defined type" in Table 3.2 and "Other" in Table 3.3 indicate that a file format or compression type not listed is being used. The user must know the type, perhaps based on the source or title.

TABLE 3.3 COMPRESSION TYPES		
Field Contents Compression Type		
00000000	No compression	
00000001	PKARC	
00000010	PKZIP	
00000011	Other	

#### 6. Body Offset

The body offset is a two byte unsigned integer (uint). The source must enter in this field the byte number at which the file body begins; that is, the number of the byte following the last byte in the PANSAT file header. This is where the file format and compression type will take effect, as far as the ground user is concerned. Note that the first byte in the file header is number 0. If there are 200 bytes in the file header, then the body offset will be '200' (0x00C8).

#### 7. Source

The source field identifies the origin of the file, or the ground station from which it was uploaded. The uploading user's HAM callsign, consisting of six ASCII characters, must be entered in this field by the client software.

#### 8. Priority

No particular use for the one byte priority field is currently specified for the satellite software. The user is free to use this field for his own purposes, such as to indicate the relative urgency of messages to addressees who share the same interpretation for this field. Any one byte bit pattern may be entered in the field, as long as the header checksum takes the contents into account.

#### 9. Number of Destinations and Destination 1 through Destination 7

If the message to be uploaded is intended for receipt by between 1 and 7 individual destination stations, then this is the number which is placed in the one byte unsigned integer of the "Number of Destinations" field. The appropriate number of

"Destination" fields are then used to indicate the HAM callsigns of the addressees. Any unused destination fields are left out of the header. If the source wishes to indicate that a message is intended for all users, then the number '0x00' is placed in the "Number of Destinations" field, and no destination fields are used.

To modify the "all users" destination, the uploading station may choose to include a "source path" or a "destination path" to further define the intended audience for the file. If a source path is to be included, then the number '0x08' is placed in the "Number of Destinations" field, and all 7 of the destination fields are included as a single 42-byte path field. Any ASCII string may be placed in this field to indicate a source path or other source identification. Similarly, if a destination path is to be included, the number '0x09' is placed in the "Number of Destinations" field, and the 42-byte path field is used to indicate a destination path or to identify the intended audience.

The satellite will not attempt to interpret destination paths or identifications. It is up to the potential downloaders to use this information, either by reading it after downloading file directories, or by providing strings to compare with the path field in select\_cmds (see Chapter VI).

#### 10. Title Length and Title

The "Title" field is a variable length array of from 0 to 30 bytes. The "Title Length" field must be entered by the source to indicate the actual length of the title. The title should be an ASCII string which will indicate to potential downloaders the contents of the file body. If there is an original file name, which it is important to keep with the file, it may be entered here. PANSAT does not otherwise retain original file names,

assigning its own after upload. The title information is for the use of potential downloaders only, and the satellite does not attempt to interpret this field.

#### 11. Keyword Length and Keywords

Like the "Title" field, the "Keywords" field is a variable length array of ASCII characters. The "Keyword Length" field must be used to specify the actual length, of between 0 and 80 bytes. Keywords should be separated from each other by one or more spaces. The satellite does interpret keyword information, but will attempt to find keywords of interest within this field if so directed by a *select cmd*.

## 12. Header Checksum and Body Checksum

The header and body checksums are used to verify the integrity of a file after uploading or downloading is complete. The body checksum must be calculated first, since it is included in the header and is thus a factor in the header checksum. All bytes in the body of the file are added together as unsigned 8-bit integers. The least significant two bytes of the resulting sum are placed in the body checksum field. The header checksum is the result of adding all bytes in the header together, except for the header checksum itself, and taking the least significant two bytes of the sum. The source must take care to update the checksums if any part of the file or header is changed before the actual upload begins. When the file number to be used has been identified by the satellite, for instance, if the source then replaces the zeroes in the "File Number" field with the proper number, those four bytes must also be added to the header checksum.

either or both of them may be left out. In this case, the fields must be filled with all zeroes. The satellite will not fill in or update "all zero" checksum fields. When the satellite performs file integrity checks, any all zero checksum will be ignored and that check will be skipped. Because of this, corrupted files may remain aboard the satellite undetected. It is up to the file's source to determine whether checksums are required. Currently, there is no way to distinguish between "no checksum" and a checksum which is actually '0'. Consequently, any checksum which is calculated to be exactly zero will be ignored. This can be remedied by adding a character to the Keywords or Title field, adjusting the appropriate length field and recalculating the checksum.

#### D. FIELDS TO BE FILLED IN BY SATELLITE

As previously stated, the uploading source of a file may choose to leave the "File Number" field of the PANSAT file header filled with zeroes. If the satellite finds that this has been done, it will fill in this field and update the header checksum appropriately. When the satellite "updates" a checksum, it does so simply by adding to it any bytes with which it has replaced zeroes. When nonzero field contents must be changed, the existing bytes are subtracted from the checksum, and the new bytes added to it. This happens, for instance, when the "Download Count" is updated (see below). The least significant two bytes of the sum are placed into the checksum field. The checksum is not completely recalculated, as this would invalidate the purpose of checking the integrity of the bytes uploaded.

There are several additional fields within the PANSAT file header which must be filled in by the satellite. The satellite software updates the header checksum whenever it places information in any of these fields. The satellite will never change a body checksum. The satellite responsible fields are described in the following subsections.

#### 1. Download Count

In this one byte field, the satellite software keeps track of how many times a particular file has been successfully downloaded. For a file addressed to "all", the download count is incremented each time a download is completed. For a file addressed to between one and seven individual callsigns, the count is incremented only when a download is completed to one of the intended addressees. The information in this field is used to determine whether a file has been previously downloaded when the default selection list is being formed (see Chapter VI, section K). The satellite software looks at the header of each file to see if the current client is one of the addressees. If there are five addressees listed, for instance, and the client is one of them, but the download count is already at '5', it is assumed that the client has already downloaded this file.

#### 2. Upload Time and Expire Time

The "Upload Time" field is filled in after a complete file has been successfully uploaded. When the final bytes of a file have been received, and the file has passed the integrity checks (such as checksums), the satellite "stamps" it with its current onboard time. This time is in the form of a four byte unsigned integer which is a count of the number of seconds since January 1, 1970 (the UTC, or Universal Time

Constant). Then the current amount of time allotted to each file to stay aboard the satellite is added to the upload time to form the expiration time. This number is placed in the Expire Time field. The amount of time allowed for each file may change based upon satellite loading. When the expiration time for a file is exceeded by the clock onboard the satellite, that file is discarded.

#### 3. PANSAT File Name and PANSAT File Extension

As each file is uploaded to the satellite, the satellite software assigns a DOS file name and extension to it. This is the file name which will be used by the onboard file management system to access the file. It is also used to easily associate each file with it's source without having to read any header fields. The 8-byte file name assigned consists of the 6 character source callsign preceded by two ASCII spaces. The file extension consists of 3 ASCII numerals (0 through 9), which indicate the sequence of files uploaded from this particular source. For instance, the first file uploaded by callsign ABCDEF would be named "ABCDEF.001", the second would be

" ABCDEF.002", etc. Extensions are repeated after number "999". It is unlikely that any file would remain with an extension that is up for re-use. But if that happens, the next unused extension will be assigned instead.

Certain file names are used by the satellite to indicate particular kinds of files generated aboard the satellite, rather than uploaded by users. These include "BULLETIN.xxx" and "USRTELEM.xxx". Files with the name "BULLETIN" contain information of general interest posted by the satellite or ground control operators and addressed to all users. Files with the name "USRTELEM" contain satellite telemetry

which may be of interest to users. USRTELEM files will be in the format "PANSAT short telemetry file". This format has not been completely specified as yet, and will be published at a later date.

#### IV. THE SPECIFICATION LANGUAGE - ESTELLE

#### A. FORMAL DESCRIPTION TECHNIQUE

A formal description technique (FDT) is a method of precisely defining the behavior of a system. It is generally advantageous to employ an FDT in the design and specification of software because descriptions produced in this way tend to be more complete, consistent, precise, concise, and unambiguous than descriptions produced in a natural language, such as English. For the specification of the PANSAT flight software, the language Estelle has been chosen. Estelle is a formal description technique which is based on an extended state transition model and uses much of the familiar syntax of the programming language Pascal. [Ref. 8]

As stated in Chapter II, the operating system chosen for the computer aboard PANSAT supports software written in the "C" programming language. For this as well as other reasons, such as development and debugging tools currently available to the Space Systems Academic Group at NPS, the implementation languages for the flight software will be "C", "C++", and assembly code as required. In spite of this, there are many reasons for developing the software specification in a description language like Estelle, prior to implementing it in a compilable language such as "C". Some of these reasons are addressed in the following sections.

#### B. CLARITY

One of the most important aspects of a software specification is clarity. The purpose of the specification is to clearly communicate the intended behavior of the program to those who must actually write the software (both the original version and later revisions) as well as to those who must use it. The behavior described by the specification must be verifiable to be the correct behavior by the systems designers who define the requirements of the system. A programming language like "C" is certainly very precise, but is often lacking in the required clarity, at least as far as humans other than the original programmer are concerned.

A particular "C" statement is written in a particular way and will cause a particular event to occur. What is not obvious is whether the particular event that occurs is exactly the event intended. When some software requirement is translated directly from an English description into a programming language implementation, there are several dangers. First of all, it is difficult to guarantee that the English description is sufficiently unambiguous that it will be understood and translated in exactly the same way by everyone. Second, if the translation is off somewhat and the software written implements a slightly different requirement than that intended, it can be difficult to catch the error by examining the code. Third, since the code is more precise than the original English description, it may be tempting to use it as the description of required behavior as the program is debugged and modified. Some programming languages, "C" in particular, are sufficiently terse that it can be difficult to extract a complete understanding of the intended behavior directly from the code without intense examination. Comments are

used to alleviate this problem - and we are back to the ambiguities of the English language. Of course, even if a precise description of behavior is extracted from the code and comments, it may no longer be the intended one.

The Pascal syntax used in Estelle, though more precise and unambiguous than English, is more obvious and easily readable than "C". Simple, well-understood, and extremely precise programming language constructs are used. These include while statements, if-then-else constructs, and for loops, as well as function and procedure calls [Ref. 8]. The specific Pascal syntax used in Appendix A is summarized in Appendix C, Table C.1. Pascal was developed as an educational language and is designed specifically to be clear and easily understood; the syntax is very straight forward. The intricate and often inscrutable statement construction of a high-powered language such as "C" is avoided.

Using Estelle, an English description of required behavior can be translated into a precise series of program-like statements. These statements are sufficiently readable that the resulting behavior can be easily analyzed and compared with the intended requirements. An ambiguous requirement statement is made crystal clear, once it has been set down in the proper series of precise Estelle statements. Once the formal specification is in place, there should be only one way to translate it, the correct way. Any software implementation must then be checked against the required behavior imparted by the Estelle description. When the program does not act in a useful way, it can be easily determined whether the original requirements statement was at fault, or whether the program code is flawed. When the software must be modified throughout

the life cycle of the host system, the originally intended behavior of the existing code will be more easily understood from the specification than from the code itself.

Of course, the specification must be maintained up-to-date along with the code. If the system requirements change, this must be reflected in the specification. The specification should always be the most accurate description of the currently intended behavior of the software system.

#### C. STATE MACHINE MODEL

Many software systems, including communications protocols, can be modeled as state machines. A major function of the PANSAT flight software is to implement communications and file transfer protocols between the ground users and the satellite. State machines provide a convenient way of modeling the software and describing its required behavior. Estelle extends the syntax of Pascal to include constructs specifically designed to clearly convey a state machine architecture. The behavior of each module is defined by its reactions to each legal stimulus it may receive while in each specific state. Even where several different states are not required for the proper functioning of a module, the state machine architecture still provides a convenient way to show the module's reactions to different inputs, and provides a means of identifying what inputs are anticipated and legal and what inputs are illegal or unexpected.

While individual statements primarily use common Pascal syntax, the hierarchy of the program modules and the module interfaces are defined by the Estelle state machine model. There are Estelle-specific reserved words which are used to establish the state machine architecture and to define other aspects of the specification which are beyond the scope of the Pascal syntax. These reserved words, the specification segments with which they are associated, and their functions are listed in Appendix C, Table C.2.

#### D. MODULE COMMUNICATIONS

Communications between various program modules are very clearly defined in Estelle. What types of information are passed between precisely which program modules is thoroughly spelled out. The set of channel definitions, which controls the flow of information between modules, is also the module interface definition.

In the software specification, several channels are defined. Each channel definition includes a list of the message types which can be "passed" through that channel. Each end of the channel is named and the message types are direction-specific. For instance, module 'A', attached to the 'User' end of a particular channel, may request information from module 'B', at the 'Provider' end, using one of several different 'request' messages. Module 'B' will reply using one of a completely different set of 'response' messages.

The name of a message type and the channel it is passed through may in itself provide all the information that is needed. In other situations, specific parameters must be passed. The parameters to be passed with each message are listed in parenthesis next to the message name in the channel definition. Estelle is a strongly typed language, and this requirement extends to the parameters passed between modules. The type of each parameter is indicated in the message definition.

Each program module has a module header definition and a module body definition. The module header definition includes a list of all the "interaction points" available to the module. These interaction points are channels, and the end of the channel to which the module is attached is indicated for each. The interaction points are the only means by which information can be passed from one module to another. The nature of each information exchange is thus precisely defined. In the modvar section at the end of the software specification, channels are attached explicitly between the various modules. The channel definitions, module header definitions, and the modvar section, taken together, completely define the architecture of the software system.

#### E. DETAIL AND ABSTRACTION

One final advantage of using a formal description technique such as Estelle, is that various levels of abstraction can be used to clarify the specification. Abstraction can be used to ignore details irrelevant to the context at any point, so that the local complexity of the description can be decreased and the overall understanding increased. Abstraction can also be used to continue with a description even though some essential details are not yet known. Commonly used functions, such as those assumed to be readily available from the operating system, can be defined as "primitives", the actual details of their internal implementations unimportant. Hardware specific details which are not known when the specification is being developed can be defined abstractly, with the specifics to be filled in later.

In contrast, any level of detail desired can be included. Thus, if minute details of the specific hardware implementation to be used are known, they can be indicated in the specification to avoid mistakes. Detailed algorithms which demonstrate a method for obtaining the specific results desired can be drawn out. In the specification of Appendix A, the communications protocols and mailbox control are described in somewhat minute detail, at a level where specific hardware requirements are unimportant. The portions of the specification dependant upon hardware configurations are merely indicated in a high-level architecture, with all details to be worked out as more information becomes available.

#### V. DATA TRANSFER MODULE

#### A. FUNCTION

The DATA\_TRANSFER module provides the interface between the high level file transfer protocol used by PANSAT and the BAX link-level AX.25 protocol software. It is the primary "BAX application program." The PACKET\_TRANSFER module, described in Chapters VI and VII, relies on the data transfer module to reassemble the AX.25 level frames passed from BAX into the complete packets uplinked from the ground station. The data transfer module also receives packets from the packet transfer module, breaks them down into frames, and passes them on to BAX to be transmitted to the intended user.

#### B. THE BAX CONTROL BLOCK

Communication with the BAX program is accomplished via the BAX functions listed in Chapter II. These are represented in the Estelle specification of Appendix A by the message types within the Abstract\_Bax\_Channel. Many of these messages have a parameter of the type Control\_Block. The control block is a data structure defined in [Ref. 4] which carries much of the actual information passed between BAX and the application program. QAX\_CLEAN\_CB is a BAX function which provides a control block structure with all fields initialized to zero. This is the only BAX function referenced in the specification by a procedure call rather than by a message type.

The definition of the Control\_Block type appears somewhat differently in Appendix A than in [Ref. 4]. It has been altered to match the syntax and avoid the reserved words of the remainder of the specification and includes only those fields which are actually used by the data transfer module. The control block fields used are listed in Table 5.1 along with the information each conveys.

	TABLE 5.1 BAX CONTROL BLOCK FIELDS				
App. A Name	[Ref. 4] Name	Туре	Information		
link	channel	uint	Indicates which of the 30 possible BAX links a frame has come in on, or which it should be sent out over. In effect, designates the ground user at the other end.		
kind	type	enumerated Frame_Type	Indicates the type of information carried by the control block. If qat_data, then the data from a data frame has been placed in an Fdata buffer, included as another message parameter. If qat_state, then the state of the link has changed, and the new state is indicated by the 'l_state' field. If qat_ui, then an unnumbered information frame has been received.		

	TABLE 5.1 BAX CONTROL BLOCK FIELDS				
App. A Name	[Ref. 4] Name	Туре	Information		
1_state	state	enumerated Link_State	Indicates the new link state in a qat_state kind of Control_Block. Only two of these states concern the data transfer module: qas_connect_pend indicates that a user has requested to be connected with the satellite. qas_disconnected indicates that a user link has been terminated The reason for termination can be determined from the 'why' field.		
why	cause	enumerated Cause	Indicates the reason for a link state change. Causes include qac_local (action of the satellite), qac_remote (action of the ground station), qac_remotefrmr (AX.25 protocol error) and qac_timeout (maximum number of frame retries exceeded).		
my_call	struct AX25_ADDR my_call	Callsign_Type = array[6] of uchar	Indicates the application program's call sign, which is always the satellite's call sign.		
my_ssid		uchar	Indicates the application's subsystem identification number (to distinguish the data transfer module from the ground control module, for instance).		

	TABLE 5.1 BAX CONTROL BLOCK FIELDS			
App. A Name	[Ref. 4] Name	Туре	Information	
his_call	struct AX25_ADDR	Callsign_Type	Indicates the ground station's call sign.	
his_ssid	his_call	uchar	Indicates the ground station's subsystem identification number	
t1	t1	uchar	The number of seconds to use for the link-level frame timeout-timer. If a frame acknowledgement is not received within t1 seconds of transmission, the frame must be retransmitted.	
maxframe	maxframe	uchar	The maximum number of frames "in flight" at one time - the link-level sliding window size. Must be 1 - 7.	
retry	retry	uchar	The maximum number of times to retry a frame transmission before terminating the link.	
paclen	paclen	uint	The maximum size of the data field on an outgoing frame.  Must be <= 256 bytes.	

### C. STATES OF THE DATA TRANSFER MODULE

The data transfer module has only two states, NORMAL and BUSY. The module is initialized in the NORMAL state, and is expected to remain in that state for the majority of the time. A transition to the BUSY state occurs only as the result of a

message from the either the GROUND\_CONTROL module or the AUTO\_CONTROL module.

If the performance of the satellite, as judged by the onboard decision-making software or by the ground control station, deteriorates to the point where it seems beneficial to allow fewer users to access the satellite for a period of time, a lockout message can be sent to the data transfer module. The type of lockout, ('l\_kind') may be new-user or all-user. When a new-user lockout message is received, the data transfer module remains in the NORMAL state, but rejects all new user connection requests. The satellite will continue communications with all users already logged on when the message is received. When the data transfer module receives an all-user lockout message, the state will change to BUSY and incoming communications from everyone except the NPS ground control station will be rejected. The data transfer module will send a 'busy' message to every BAX link, and BAX will respond to any frame (except those from NPS) with a "receive-not-ready" frame.

The control software may also find it necessary to turn the transmitter off for an extended period of time, such as during a battery recharge. When this occurs, and the transmitter will not be ready at a moment's notice, a 'transmitter.off' message will be sent to the data transfer module. This message will not change the state of the module, which will still be able to receive any incoming frames, but it will change the state variable 'transmit\_ok' to false. When this occurs, the data transfer module will not attempt to transmit any frames, and any logged-in users will most likely disconnect due to frame time-outs.

#### VI. PACKET TRANSFER MODULE - PACKET TYPES

#### A. FILE TRANSFER LEVEL 0

The packet transfer module specified in Appendix A has its origins in the File Transfer Level 0 (FTL0) Pacsat Protocol developed by Jeff Ward and Harold E. Price [Ref. 7]. The basic data structures and state transitions are functionally equivalent to FTL0, with some modifications. In order to make use of the satellite software specified in Appendix A, the corresponding ground station software must be developed which will produce packets in the proper format to be interpreted by PANSAT. Therefore, the bit-level structure of the ground station packets is described below, as well as the proper ground interpretation of the packets originating on the satellite.

#### B. PACKET FORMAT

Each packet to be transmitted consists of an information field of 0 to 2047 bytes, preceded by a two byte header. The header identifies the type of packet and indicates the number of bytes in the information field. The packet structure is defined as follows in the software specification:

Packet Type = record

length lsb:

uchar:

hl:

uchar;

info:

Pdata:

end;

This structure indicates that the header portion of the packet consists of the two unsigned characters (octets) 'length\_lsb' and 'hl'. The information field is given the type 'Pdata', which is defined in the specification as an array of 0 to 2047 unsigned characters. Since this is a variable length array, its length must be indicated in the header.

The octet 'length\_lsb' contains the least significant 8 bits of the data length. The octet 'hl' contains the 3 most significant bits of the data length, as well as an indication of the type of packet. The bits of 'hl' are labeled '76543210'. Bits 7-5 are the 3 most significant bits of the data length, and must be prepended to the 'length\_lsb' to give the full length of the information field. Bits 4-0 of 'hl' provide a number from 0 to 31. This number is decoded into packet type as indicated in Table 6.1.

TABLE 6.1 PACKET TYPES				
Packet Number   Specification Constant   Pa		Packet Name		
0	data	Data		
1	data_end	Data End		
2	login_resp	Login Response		
3	upload_cmd	Upload Command		
4	ul_go_resp	Upload Go Response		
5	ul_error_resp	Upload Error Response		
6	ul_ack_resp	Upload Acknowledged Response		
7	ul_nak_resp	Upload Not Acknowledged Response		
8	download_cmd	Download Command		
9	dl_error_resp	Download Error Response		
10	dl_aborted_resp	Download Aborted Response		
11	dl_completed_resp	Download Completed Response		

TABLE 6.1 PACKET TYPES				
Packet Number Specification Constant		Packet Name		
12	dl_ack_cmd	Download Acknowledged Command		
13	dl_nak_cmd	Download Not Acknowledged Command		
14	dir_short_cmd	Directory Short Command		
15	dir_long_cmd	Directory Long Command		
16	select_cmd	Select Command		
17	select_resp	Select Response		
18 - 29		reserved		
30	del_cmd	Delete Command		
31	del_resp	Delete Response		

The ground software and satellite software are peer entities at this level, rather than master and slave. However, since the ground must initiate all data exchanges, with the satellite acting as a server responding to requests made from the ground, the identifier 'cmd' is used to indicate packets originating from the ground, while 'resp' indicates packets sent from the satellite. The data and data\_end packets can originate from either the ground station or the satellite. Explanations of each of the packet types and the contents of their information fields are given in the following sections.

#### C. THE DATA AND DATA END PACKETS

Any file to be transmitted, either from the ground or from the satellite, will be broken up into an appropriate number of *data* packets, depending on its length. The

information field of each data packet will be the bytes from the file to be transmitted. Bits 4-0 of the 'hl' field of the packet header will be '00000', identifying the packet as containing file data in its information field. Bits 7-5 of 'hl' and the octet 'length\_lsb' will together indicate the number of bytes of file data being transmitted in this packet. The transmission of the end of the file will be indicated by sending a data\_end packet immediately after the transmission of the last data packet. The data\_end packet has no 'info' field.

Bit-level examples of the various packet types will be given in Tables 6.2A through 6.2S. In these examples, 0's and 1's will be shown where particular bit patterns must be used. Where arbitrary bit patterns may be present, other symbols will be used. The intended meanings of these symbols will be made clear in the "interpretation" section of each table.

TABLE 6.2A EXAMPLE OF A data PACKET				
length_lsb	hl	info		
LLLLLLL	ннн 00000	dddddddd ddddddd		
Interpretation				
low order bits of the data length. HHH = high order bits of the data length. file data in bytes				

TABLE 6.2B	TABLE 6.2B EXAMPLE OF A data_end PACKET		
length_isb hi			
00000000	00 000 00001		
Interpretation			
no info field '00001' = data_end packet.			

#### D. THE LOGIN RESPONSE PACKET

Neither the FTLO protocol of Ward and Price, nor the packet transfer protocol specified here, has an explicit Login Command packet. A login request from the user is made implicitly whenever a data link is established between the ground station and the satellite on the lower, AX.25, data transfer level. When the AX.25 protocol software, BAX, recognizes a "connection request" frame from a new user, it informs the satellite data transfer module. The data transfer module tells BAX whether to accept the connection or not. If the connection is accepted, BAX sends the "accept connection" frame to the ground user, and the data transfer module informs the packet transfer module that a data link has been established. At this point, the satellite packet transfer protocol calls for the transmission of a Login Response packet.

The purpose of the *login\_resp* packet is to inform the user of the time onboard the satellite when the data link is established. The *login\_resp* packet also has a one byte login flag. This flag indicates whether the user currently has an active selection list (explained below) and whether the satellite requires Pacsat file headers. The Pacsat file header was developed by Jeff Ward and Harold Price for use with their Pacsat Protocol

Suite, which includes FTLO [Ref. 8]. A PANSAT file header has been developed which does not match the Pacsat file header of Ward and Price, and so the login flag in the login\_resp packet will always indicate that Pacsat file headers are not required. PANSAT file headers will always be required for files uploaded to PANSAT. The information field of the login\_resp packet includes a 4-byte unsigned integer indicating the login time (the number of seconds since January 1, 1970), followed by a 1-byte login flag. Thus, the information length indicated by the header must be 5. Bits 7-4 of the login flag will be '0000'. Bit 3, the 's' bit, will be '1' if the client already has an active selection list, and will be '0' if not. Bit 2 will be '0', indicating that Pacsat file headers are not used. Bits 1 and 0 indicate the protocol version number. They will be '00' in the case of the protocol specified in Appendix A.

TABLE 6.2C EXAMPLE OF A login_resp PACKET			
length_lsb	hl	info	
00000101	000 00010	11111111 11111111 111111111 111111111 1111	
Interpretation			
5 bytes in info field login_resp packet   4 byte login time login flag: 's' = '1' indicates active selection list			

## E. THE UPLOAD COMMAND, UPLOAD GO RESPONSE, AND UPLOAD ERROR RESPONSE PACKETS

Before uploading a file to the satellite, the client software must first determine whether there is room in the mail box and if the satellite will accept the upload. To

make this determination, the ground station transmits the upload\_cmd. The satellite will respond with the ul\_go\_resp if the upload will be allowed and with the ul\_error\_resp if not.

The information field of the *upload\_cmd* contains a 4-byte file number followed by a 4-byte file length. If this is the first request to upload a particular file, the file number must be '0x00000000'. The file length must be the actual length of the file which is intended for upload, including the PANSAT file header, which must be prepended to each file. When the satellite receives the *upload\_cmd*, it will determine if there is room for a file of the indicated length. If there is, the *ul\_go\_resp* will include a file number to be assigned to the file in the mail box aboard the satellite. When the client receives this file number, it may be placed in the PANSAT file header before upload. If the client does not place the proper file number in the PANSAT file header before upload, then that field must contain all 0's and the satellite will make the correction once the file has been successfully uploaded.

If the upload request is for the continuation of a previously interrupted upload, the upload\_cmd must contain the actual file number previously assigned by the satellite. The file length must still indicate the full length of the file, regardless of how much of the file was previously uploaded. If the satellite can accept this continued upload, the ul\_go\_resp will include the offset at which the client should begin the transmission of the file. To determine this offset, the satellite simply inspects its partial copy of the file to see how many bytes it has previously received. The complete information field of the ul\_go\_resp includes the 4-byte file number either newly or previously assigned to the

file, followed by the 4-byte file offset. If no part of the file has been previously uploaded, then the indicated offset will be '0'.

If there is no room for the file, if the upload\_cmd includes a non-zero file number which does not correspond to any file onboard the satellite, or if the satellite determines that upload of the indicated file has already been completed, then an ul\_error\_resp is transmitted rather than an ul\_go\_resp. The ul\_error\_resp has a 1-byte information field which simply indicates one of several possible error conditions. The possible errors associated with an upload command and their corresponding bit patterns are indicated in Table 6.3A.

TABLE 6.3A ERROR CODES		
Bit Pattern	Error Code	
00000001	er_ill_formed_cmd	
00000010	)10 er_bad_continue	
00000100 er_no_such_file_number		
00001100	er_file_complete	
00001101	er_no_room	

The meanings of most of these error codes are obvious from their names. The code er\_bad\_continue is issued when the file number in the upload\_cmd is non-zero but the file length in the upload\_cmd does not agree with the file length stored in the PANSAT header of the partially uploaded file.

TABLE 6.2D EXAMPLES OF upload_cmd PACKETS			
length_lsb	bl	info	
00001000	000 00011	00000000 00000000 00000000 00000000 LLLLLLLL	
00001000	000 00011	nnnnnnn nnnnnnn nnnnnnn nnnnnnn LLLLLLLL	
		Interpretations	
8 bytes in info field	upload_ cmd packet	New upload - unknown file number 4 byte file length	
8 bytes in info field	upload_ cmd packet	4 byte file number 4 byte file length	

TABLE 6.2E EXAMPLE OF AN ul_go_resp PACKET				
length_ksb	hl	info		
00001000	000 00100	nnnnnnn nnnnnnn nnnnnnn nnnnnnnn 0000000		
	Interpretation			
8 bytes in ul_go_resp 4 byte file number info field packet 4 byte file offset at which to begin upload				

TABLE 6.2F EXAMPLE OF AN ul_error_resp PACKET				
length_isb hl info				
00000001	000 00101	cececcec		
Interpretation				
1 byte in info field ul_error_resp 1 byte error code				

## F. THE UPLOAD ACKNOWLEDGED RESPONSE AND UPLOAD NOT ACKNOWLEDGED RESPONSE PACKETS

When the client software on the ground receives an ul\_go\_resp from the satellite, it will commence to upload the file in a series of data packets, starting with the byte of the file indicated by the offset in the ul\_go\_resp. Once the data packet containing the last byte of the file has been transmitted, the data\_end packet must be sent. After receiving the data\_end packet, the satellite will check the integrity of the file, as will be explained in Chapter VIII. If the file passes all checks and is successfully stored aboard the satellite, an ul\_ack\_resp will be transmitted to the user. If the file is found to be defective, an ul\_nak\_resp is transmitted instead and the file is discarded. The satellite will remember the file number, however, so that the user can later attempt another upload of the same file.

The ul\_ack\_resp has no information field. The ul\_nak\_resp has a 1-byte information field which consists of one of the error codes of Table 6.3B.

TABLE 6.3B ERROR CODES	
Bit Pattern Error Code	
00001101	er_no_room
00001110	er_bad_header
00001111	er_header_check
00010000	er_body_check

The er\_bad\_header code is sent when the PANSAT file header is missing, incomplete, or incorrect. The code for er\_header\_check is sent when the checksum on the header fails and er\_body\_check is sent when the checksum on the file body fails.

An ul\_nak\_resp may also be sent by the satellite before a data\_end packet is received if the satellite needs to terminate the upload for any reason. If the ground station receives an ul\_nak\_resp, it should immediately stop sending data packets, and transmit a data\_end packet if it has not already done so.

TABLE 6.2G EXAMPLE OF AN ul_ack_resp PACKET		
length_lsb hl		
00000000	000 00110	
Interpretation		
no info field ul_ack_resp		

TABLE 6.2H EXAMPLE OF AN ul_nak_resp PACKET		
length_lab	PI	info
00000001	000 00111	ccccccc
Interpretation		
1 byte in info field	ul_nak_resp	1 byte error code

#### G. THE DOWNLOAD COMMAND PACKET

Prior to using download\_cmd packets to request files to be downloaded from PANSAT, the client must establish an active selection list onboard the satellite. This is achieved by using the select\_cmd which is explained below. Once the client has used the selection list to begin downloading a file or to obtain file directories (see Directory Commands below), the selection list need not remain active to continue downloading, as long as the client knows the file number for each file requested.

Each download\_cmd packet is used to request a single file. The information field of this packet contains the 4-byte file number for the file requested followed by the 4-byte file offset from which transmission of the file should begin. In FTLO, a file number of '0x00000000' is used to indicate the next file in the active selection list, proceeding from newer files toward older files, while '0xFFFFFFFF' requests the next file in the list proceeding from older files toward newer files. In the current packet transfer specification for PANSAT, both '0x00000000' and '0xFFFFFFFF' will result in requesting the next file in the active selection list proceeding from older files toward newer files. I' he actual file number of the file requested is known, then this number

is placed in the download\_cmd packet. The file offset should be '0' if this is a new download request, and should indicate the byte number from which to proceed if this is a download continuation.

TABLE 6.21 EXAMPLES OF download_cmd PACKETS		
length_lsb	bì	info .
00001000	000 01000	00000000 00000000 00000000 00000000
00001000	000 01000	11111111 11111111 11111111 11111111 00000000
00001000	000 01000	nnnnnnn nnnnnnn nnnnnnn nnnnnnn
		Interpretations
8 bytes in info field	download_ cmd	requesting next file in select list begin download at beginning of file
8 bytes in info field	download_ cmd	requesting next file in select list begin download at beginning of file
8 bytes in info field	download_ cmd	file number of requested file file offset at which to begin download

When requesting the "next" file in the selection list, the offset should always be '0', since this should only be used to request a new file. Once a download has been interrupted and subsequently continued, the file number should already be known, and this information as well as the offset should be used in the download\_cmd. If the file offset indicated by the ground station is equal to or greater than the length of the file stored on the satellite, no error is generated. Instead, the satellite transmits a data\_end packet immediately, with no preceding data packets.

#### H. THE DOWNLOAD ERROR RESPONSE PACKET

When the satellite receives a properly formatted download\_cmd which it is able to respond to, it immediately begins downloading the file in a series of data packets. Once the last byte of the file has been transmitted, the satellite sends a data\_end packet. If, however, the satellite cannot service the download\_cmd for any reason, it will transmit a dl\_error\_resp packet.

The dl\_error\_resp information field consists of a 1-byte error code. The possible errors are shown in Table 6.3C.

TABLE 6.3C ERROR CODES	
Bit Pattern Error Code	
00000100	er_no_such_file_number
00000101 er_selection_empty	

The code er\_no\_such\_file\_number is used if a specific file has been requested, the file number of which is not found in the mail box. The code er\_selection\_empty is used when the "next" file is requested, but the user currently has no active selection list. The specification for FTLO also includes error codes dealing with file forwarding capabilities which are not implemented on PANSAT. These codes are included in the specification of Appendix A for the sake of completion, to ensure they will not be used for any PANSAT specific definitions. They will not be included in any dl\_error\_resp packets from PANSAT, however. These FTLO error codes, unused by PANSAT, include er\_already\_locked and er\_no\_such\_destination.

TABLE 6.21 EXAMPLE OF A dl_error_resp PACKET		
length_lsb	Ы	info
00000001	000 01001	ccccccc
Interpretation		
1 byte info field	dl_error_resp	1 byte error code

# I. THE DOWNLOAD ACKNOWLEDGED COMMAND, DOWNLOAD COMPLETED RESPONSE, DOWNLOAD NOT ACKNOWLEDGED COMMAND, AND DOWNLOAD ABORTED RESPONSE PACKETS

When the client software receives a data\_end packet from the satellite, it knows the downloaded file is complete. It performs any desired integrity checks on the file (such as checking header and body check sums) to determine whether the download was completed successfully. If the file has been received satisfactorily, the ground station must transmit a dl\_ack\_cmd. The satellite responds to a dl\_ack\_cmd with a dl\_completed\_resp to end the download process. If the ground software does not find the downloaded file to be satisfactory, it transmits a dl\_nak\_cmd, to which the satellite responds with a dl\_aborted\_resp.

In the FTLO specification, the information field of the dl\_ack\_cmd consists of a one byte 'register\_destination'. This information is used by a Pacsat to complete some of the file forwarding operations which are not implemented by the current PANSAT software specification. Therefore, the information field is not necessary in a dl\_ack\_cmd sent to PANSAT, and it will be ignored if it is included. This single byte of information may

be adapted for use by PANSAT at a later date. The dl\_completed\_resp transmitted by the satellite has no information field. The dl\_nak\_cmd and the dl\_aborted\_resp likewise have no information fields.

TABLE 6.2K EXAMPLE OF A dl_ack_cmd PACKET		
length_lsb	<b>bl</b>	info
00000001	000 01100	XXXXXXX
	Interpretatio	n
1-byte info field	dl_ack_cmd	register_destination

TABLE 6.2L EXAMPLE OF A dl_completed_resp PACKET		
length_lsb hl :		
00000000	000 01011	
Interpretation		
no info field dl_completed_resp		

TABLE 6.2M EXA	MPLE OF A dl_nak_cmd PACKET	
length_lsb hl		
00000000	000 01101	
	Interpretation	
no info field dl_nak_cmd		

TABLE 6.2N EXAMPLE OF A dl_aborted_resp PACKET		
length_lsb hl		
00000000	000 01010	
Interpretation		
no info field dl_aborted_resp		

#### J. THE DIRECTORY COMMAND PACKETS

FTLO specifies two directory commands, the dir\_short\_cmd and the dir\_long\_cmd.

For PANSAT, there is only one directory command, and either of these two packet types will invoke it. The results of each command will be exactly the same. The information field of a dir\_cmd is a 4-byte file number. This number indicates the file for which a directory entry is requested. A directory entry consists of the PANSAT file header from the file of interest. From this file header, the ground station software can determine any necessary information about the file. The user can decide from this information whether to request the file for download.

In FTLO, a file number of '0x00000000' is used to request the directory entries for the next 10 files in the active selection list, proceeding from newer files toward older files, while '0xFFFFFFF' requests the next 10 file directories proceeding from older files toward newer files. In the current packet transfer specification for PANSAT, both '0x0000000' and '0xFFFFFFFF' will result in requesting entries for the next 10 files in the active selection list proceeding from older files toward newer files. If the client

has no currently active selection list, then a directory entry can only be requested for a file for which the file number is already known.

When the satellite receives a correctly formatted dir\_cmd which it can respond to, it sends the requested information down in a data packet. Since PANSAT file headers are at most 200 bytes long, 10 of them will fit in a single packet. Thus, after one data packet is transmitted, a data\_end packet will immediately be sent. If the satellite is unable to respond to the dir\_cmd, it will send a dl\_error\_resp indicating the reason. The error character contained in this packet will be either er\_selection\_empty or er no such file number.

	TABLE 6.20 EXAMPLES OF dir_cmd PACKETS		
length_lsb	hl	info	
00000100	000 01110	00000000 00000000 00000000 00000000	
00000100	000 01111	11111111 11111111 11111111 11111111	
00000100	000 01111	nnnnnnn nnnnnnn nnnnnnn	
		Interpretations	
4 bytes in info field	dir_short_ cmd	requesting directory entries for next 10 files in select list	
4 bytes in info field	dir_long_ cmd	requesting directory entries for next 10 files in select list	
4 bytes in info field	dir_long_ cmd	file number of file for which directory entry is requested	

#### K. THE SELECT COMMAND AND SELECT RESPONSE PACKETS

The select\_cmd is the means by which the user designates files to be placed in an active selection list onboard the satellite. Once this list has been established, it can be used to request file directories or files for download. The select\_cmd specified by FTLO assumes use of Pacsat file headers in its structure. A different select\_cmd structure is specified here, which is based upon the PANSAT file header. This structure is related to, but not exactly the same as, the select\_cmd specified by Price and Ward [Ref 1.].

Once a select\_cmd is recognized, if it is not in the PANSAT structure specified here, then a default selection list will be compiled. This list will be comprised of all mail addressed to the requesting user which has not been previously downloaded, all bulletins and user-accessible telemetry and messages addressed to "all". The satellite will send a select\_resp packet to the user. The information field of this packet consists of a two byte integer which indicates the number of files in the selection list.

If the PANSAT select structure is used, the selection list will be assembled according to the selection criteria contained in the select\_cmd. If the satellite can interpret the select\_cmd and compile the corresponding selection list it will transmit the select\_resp packet. In this case, the two byte integer in the information field indicates the number of files matching the selection criteria. An active list consisting of those file numbers will be maintained aboard the satellite. If the select\_cmd appears to be in the PANSAT format but cannot be successfully parsed by the satellite software, then a dl\_error\_resp is transmitted with the error code: 00001000 er poorly formed sel.

The select\_cmd has a variable length information field. The information contained in the field consists of a PANSAT specific flag, the number of selection criteria present, and the criteria themselves. The selection criteria are combined with each other using the operators and and or, forming a restricted type of postfix logical equation. In this postfix equation, each logical operator is preceded by its two operands. The first two selection criteria are combined logically, according to the first operator, to form the single operand true or false. This operand is then followed by another selection and another operator. The second operator combines the two operands preceding it to form a single operand. The process is continued until the last logical operator present, which will be the last component of the equation, has been used to combine its two operands. The result will be a single value of true or false. Each file for which the selection equation yields a value of true will have its file number added to the active selection list aboard the satellite.

The first byte of the select\_cmd information field should be '0xFF', a flag indicating that the select\_cmd is in the PANSAT format. pon recognizing this flag, the satellite will attempt to translate the select\_cmd into the appropriate logical equation. If this flag is not present, the default selection list described above will be compiled and the remainder of the select structure will be discarded.

The second byte in the information field is an unsigned integer indicating the number of selection criteria contained in the remainder of the structure. The selection criteria can be defined as follows:

#### selection = record

relop:

uchar;

header item:

uchar;

item len:

uchar:

compare\_item: array[item\_len] of uchar;

#### end:

Bit '7' of the one byte 'relop' (relational operator) must always = '0'. Bits '654' have the interpretations shown in Table 6.4, and bits '3210' are translated as indicated in Table 6.5. The 'header item' identifies which item in the PANSAT file header to compare the 'compare\_item' with. The one byte 'header\_item' is decoded as Table 6.6 indicates.

TABLE 6.4 BITS '654' OF THE RELATIONAL OPERATOR	
Bits 654	Relation
000	equal to
001	greater than
010	less than
011	not equal to
100	greater than or equal to
101	less than or equal to

TABLE 6.5 BITS '3210' OF THE RELATIONAL OPERATOR			
Bits 3210	Interpretation of 'compare_item'	Notes	
0000	Multi-byte unsigned integer	'item_len' must equal 1, 2 or 4.	
0011	Array of characters, convert to lower case before comparison	Valid only with Bits '654' = '000' or '011'	

TABLE 6.6 HEADER ITEMS			
'header_item' Bit Pattern	Const	Name of Header Field	
00000000	fl	Flag	
00000010	mn	Mail Number (File Number)	
00000110	ml	Mail Length	
00001010	ft	File Type	
00001011	ct	Compression Type	
00001100	bo	Body Offset	
00001110	dc	Download Count	
00001111	sc	Source Call Sign	
00010101	pr	Priority	
00010110	ut	Upload Time	
00011010	et	Expire Time	
00011110	na	PANSAT File Name	
00100110	ex	PANSAT File Extension	
00101001	nd	Number of Destinations	
00101010	ds	Destination Call Signs or Path	
01010100	ti	Title	
01110100	kw	Keywords	

The short integers formed by the 'header\_item' bit patterns correspond to the normal byte offsets within the PANSAT file header of the beginning of each of the listed header fields. This is useful in other areas of the software specification.

The one byte integer 'item\_len' gives the byte length of the last item in the 'selection', the 'compare\_item.' The compare item is interpreted, as indicated by the

'relop', as either an unsigned one, two or four byte integer, or an array of characters.

The relational operation specified by the 'relop' is performed between the designated header item and the compare item. If the relation 'header\_item' 'relop' 'compare\_item' is satisfied, then this selection equation is evaluated as true.

If the user has only one criteria for selection, the Select\_Structure can end after just one 'selection'. The user may, however, specify multiple selection criteria. For this purpose, the bit patterns for logical operators are defined as in Table 6.7.

TABLE 6.7 LOGICAL OPERATORS			
'logop' Bit Pattern	Logical Operation		
10000000	and		
11100000	or		

A completed Select\_Structure may appear as follows:

OxFF num selection selection logop selection logop selection logop, etc.

When 'equal to' string comparisons are made between 'compare\_items' and certain header fields, the comparison is defined as successful if the 'compare\_item' string is found anywhere within the header item string. In these same fields, a 'not equal to' comparison is successful if the 'compare\_item' string is not found anywhere within the header item string. Header fields for which this applies are the "Title" and "Keywords" fields. This can also apply to the destination fields under certain circumstances, which will be elaborated on below.

To indicate that a file is addressed to "all", the source places the number '0x00' in the "Number of Destinations" field, and the following 7 "Destination" fields are left out of the header completely. A one-byte integer comparison between the "Number of Destinations" field and the number '0x00' can determine that a file is addressed to all users. When there are between 1 and 7 individual destination call signs, this number is placed in the "Number of Destinations" field, and the appropriate number of "Destination" fields are included. To find files which are addressed to a specific user, a string comparison can be made between a 6-byte call sign as the 'compare\_item' and the header item "Destination Call Signs or Path". This header item refers to all "Destination" fields present. The satellite will compare the designated call sign to each destination listed, and the comparison will be successful if a match is found with any of them.

It may be that the user wishes to designate an audience for the file which is broader than 7 individual call signs but narrower than "all" users in the world. In order to achieve this, the user places the number '0x08' or the number '0x09' in the "Number of Destinations" field, and all 7 "Destination" fields are then included as a single 42-byte array. In this combined field can be placed information to further define the audience for which the file is intended. If the "Number of Destinations" is '0x08', then the file is addressed to "all", and the source has included path, location or other information about himself in the following "Path" field. If the "Number of Destinations" is '0x09', then the source has included further information about the intended audience in the following "Path" field. Users may use this information in select\_cmds, in which case

any 'compare\_item' will be searched for anywhere within the "Path" field. The header item to use within the 'selection' will again be "Destination Call Signs or Path", but this time the satellite will look for any matching string, not simply matching 6-byte call signs. Source or destination path information can probably better be used by ground software as the result of dir\_cmds, in which case the software can present the user with any information which may help the user in choosing individual files to download.

1	TABLE 6.2P EXAMPLES OF select_cmd PACKETS		
length_lsb	hl	info	
00001001	000 10000	1111111 00000001 0 001 0000 00010110 00000100 ttttttt ttttttt ttttttt	
00011001	000 10000	11111111 00000011 0 000 0011 01110100 00000100 01000001 01010010 01001101 01011001 0 000 00	
	Interpretations		

T	TABLE 6.2P EXAMPLES OF select_cmd PACKETS			
9 bytes in info field	select_cmd packet	PANSAT select_cmd flag, 1 'selection' 'relop' - greater than a multibyte unsigned integer 'header_item' - Upload Time 4 byte 'compare_item' - last login time  ( This command requests all files which were uploaded after the user's last login time.)		
25 bytes in info field	select_cmd packet	PANSAT select_cmd flag, 3 'selections' 'relop' - equal to an array of characters 'header_item' - Keyword 4 characters in the 'compare_item' 'compare_item' - "ARMY" 'relop' - equal to an array of characters 'header_item' - Keyword 4 characters in the 'compare_item' 'compare_item' - "NAVY" 'logop' - or 'relop' - greater than a multibyte unsigned integer 'header_item' - Upload Time 4 byte 'compare_item' - last login time 'logop' - and  (This command requests all files which were uploaded after the user's last login time and which also have either "ARMY" or "NAVY" as a keyword.)		

TABLE 6.2Q EXAMPLE OF A select_resp PACKET			
length_lsb	hl	info	
00000010	000 10001	nnnnnn nnnnnnn	
Interpretation			
2 byte info field select_resp Number of files placed in select list			

#### L. THE DELETE COMMAND AND DELETE RESPONSE PACKETS

The specification for FTLO does not allow for a user-requested file deletion. It assumes that files will simply remain aboard the satellite until their expiration dates are exceeded, or until the satellite itself or the ground controllers cause the files to be deleted. This del\_cmd, therefore, has no equivalent in FTLO.

The information field of the <u>del\_cmd</u> contains only the 4-byte file number of the file to be deleted. The satellite responds to a <u>del\_cmd</u> with a <u>del\_resp</u> packet. The information field of this packet merely contains one of the one-byte error codes of Table 6.3D.

TABLE 6.3D ERROR CODES			
Bit Pattern Error Code			
00000000	no_error		
00000100	er_no_such_file_number		
10010000	er_permission_denied		

If the satellite indicates no\_error in the del\_resp, then the file has been successfully deleted. A user may only delete files uploaded by him or addressed to him as the sole destination. The satellite ensures these criteria are met by inspecting the appropriate fields in the file's header before any deletion is carried out. An attempt to delete any other file will result in er permission denied.

TABLE 6.2R EXAMPLE OF A del_cmd PACKET				
length_lsb hl info				
00000100 000 11110 nnnnnnn nnnnnnn nnnnnnn nnnnnnn				
Interpretation				
4-byte info field del_cmd file number of file to be deleted				

TABLE 6.2S EXAMPLE OF A del_resp PACKET				
length_lsb hl info				
0000001	000 11111	ccccccc		
Interpretation				
1-byte info field del_resp 1-byte error code				

#### VII. PACKET TRANSFER MODULE - STATE TRANSITIONS

#### A. TRIGGERS

The operation of the packet transfer module is based upon state transitions triggered by the receipt of packets from the user and messages from other flight software modules. When the user on the ground sends a *cmd* packet to the satellite, the response will depend, in part, on what state the packet transfer module is in.

The architecture of the packet transfer state machine can easily be seen in the Estelle specification of Appendix A. The trans section of the module definition clearly shows all possible transitions from one state to another, along with what packet or message triggers each transition, and what action is taken as a result. This textual description can be translated into a more visual format by means of state transition graphs, such as those included with the data flow diagrams of Appendix B.

#### **B.** INSTANTIATIONS

It is intended that multiple users should have access to the mail box onboard the satellite "simultaneously". This is achievable because all user transmissions to the satellite are packetized. BAX, the AX.25 data transfer level software, can administer up to 30 user links at once. As each AX.25 frame is received, BAX determines which user it is from, and deals with it according to the AX.25 protocol and the state of the link with that particular user.

"simultaneous" users, there must be a copy of the packet transfer module associated with each virtual link. The fact that multiple copies of the packet transfer module are initialized can be seen in the modvar section at the end of the Estelle software specification. This says that one packet transfer module is created for each link. The definition of "Link\_Type", near the beginning of the specification, indicates that there are between 0 and 30 links. When a packet transfer module receives messages from, or sends messages to, another software module, the specific instantiation involved is indicated by the initialization parameter 'link'. The packet transfer modules, as well as the channels associated with them, are referenced as array elements; the 'link' each is associated with acts as the array index. For instance, in the module header definition of the MAILBOX\_CONTROL\_TYPE, it is stated that this module has an array of 30 (maxlinks) Mailbox\_Access\_Channels. In the modvar section, each of these channels is connected to a different copy of the packet transfer module.

#### C. TRANSITIONS

The state transitions for a particular instantiation of the packet transfer module are affected only by packets from the user associated with the link being administered by that module. The data transfer module, as explained in Chapter V, must assemble complete packets from the frame data sent to it by BAX. Each packet is sent to the appropriate packet transfer module, depending upon which BAX link it was received on. Likewise, as the packet transfer modules send resp packets to the data transfer module for

transmission via BAX, the data transfer module must break each packet up into frames and send them via the appropriate BAX link to reach the intended recipient.

While a packet transfer module is in any particular state, only certain packets from the user will have meaning. An unexpected packet will cause any actions in progress (such as uploading or downloading a file) to be aborted. FTLO defines unexpected or incorrect packets as sufficient cause to terminate the link with a user. The specification in Appendix A, however, only calls for the packet transfer module to return to a waiting state after abandoning any action in progress. At this point, the module is ready to accept any valid command from the user. The user will be informed of the problem via an appropriate error\_resp packet. After receipt of any error message, the user should assume that the satellite is waiting for the ground station to initiate a new action.

#### D. STATES

FTLO was designed primarily for use with satellites with full duplex capabilities. For this reason, it maintains two separate state machines, one for the uplink process and the second for the downlink process. PANSAT is a half-duplex communications satellite. The two state machines of FTLO have been combined into a single machine in the specification of the PANSAT packet transfer module. The states are listed in Table 7.1. Explanations are included in the following subsections.

TABLE 7.1 PACKET TRANSFER STATES			
State Identifier Explanation of State			
UL/DL_UNINIT	Upload/Download Uninitiated		
UL/DL_CMD_WAIT	UL/DL_CMD_WAIT Waiting for an Upload or Download Command		
WAIT_MAILBOX Waiting for a Message from the Mailbox Control Module			
UL_DATA_RX	Ready to Uplink Data		
UL_ABORT	Upload Aborted		
DL_FILE_DATA	Downloading a File		

### 1. UL/DL UNINIT

UL/DL\_UNINIT is the state into which the packet transfer module is first initialized, before a user link has been established with it. In this state, the module does nothing but wait to be assigned a user. Upon receipt of the 'connection' message from the data transfer module, the packet transfer module asks the mailbox control module whether or not the new user has an active selection list, and moves into the WAIT\_MAILBOX state to await a reply. When the reply message is received, the module will enter the UL/DL\_CMD\_WAIT state. The packet transfer module returns to the UL/DL\_UNINIT state when it is sent a 'disconnect' message by the data transfer module, regardless of what state it is in when this message is received.

TABLE 7.2 STATE TRANSITIONS FROM UL/DL_UNINIT				
Received Message	Action	Next State		
connection	active_sl_req message (Asks the mailbox control module if there is an active selection list for this user.)	WAIT_MAILBOX		

# 2. UL/DL\_CMD\_WAIT

In the UL/DL\_CMD\_WAIT state, the packet transfer module is waiting for a packet from the user which will initiate either an upload process or a download process. Packets which can be legally received while in this state are listed in Table 7.3, along with the resultant actions and transitions. Any other, unexpected, packets will result in an ul\_error\_resp packet with the error code er\_ill\_formed\_cmd, and the module will remain in the state UL/DL\_CMD\_WAIT.

TABLE 7.3 STATE TRANSITIONS FROM UL/DL_CMD_WAIT			
Received Packet or Message	Action	Next State	
upload_cmd	mail_num_req message (Request a new file number or a current file offset from the mailbox control module.)	WAIT_MAILBOX	
del_cmd	mail_del_req message (Request that the mailbox control module delete a file.)	WAIT_MAILBOX	
select_cmd	mselect_req message (Request the mailbox control module form a selection list.)	WAIT_MAILBOX	

TABLE 7.3 STATE TRANSITIONS FROM UL/DL_CMD_WAIT				
Received Packet or Message	Action	Next State		
dir_short_cmd dir_long_cmd	dir_req message (Request directory information from the mailbox control module.)	AIT_MAILBOX		
download_cmd	mail_req message (Request file data from mailbox control module.)	DL_FILE_DATA		
dl_nak_cmd	none	UL/DL_CMD_WAIT		
disconnect	none	UL/DL_UNINIT		
other packets	ul_error_resp packet	UL/DL_CMD_WAIT		

### 3. WAIT MAILBOX

As can be seen in Table 7.3, most packets received while in the UIL/DL\_CMD\_WAIT state result in a transition to the WAIT\_MAILBOX state, with no immediate response packet to the user. This is because the packet transfer module requires information from the mailbox control module before it can make a proper reply to the user. The mailbox control module analyzes each information request message and replies with an appropriate response message. The response of the mailbox control module will determine which state the packet transfer module will enter when it leaves the WAIT\_MAILBOX state, as well as what packet it sends to the user. The WAIT\_MAILBOX state may also be entered from the UL/DL\_UNINIT state as shown above, or the UL\_DATA\_RX state, as will be explained below. Table 7.4 summarizes the mailbox access channel messages just prior to a transition to the WAIT\_MAILBOX

state, the possible reply messages from the mailbox control module, and the resulting further actions and state transitions of the packet transfer module. No user command packets are expected while in the WAIT\_MAILBOX state, as the user should still be waiting for a reply from the last packet sent to the satellite.

TABLE 7.4 STATE TRANSITIONS FROM WAIT_MAILBOX				
Message	Reply Message	Action	Next State	
active_sl_req	active_sl_resp	login_resp packet	UL/DL_CMD_ WAIT	
mail_num_req	mail_num_resp, no errors	ul_go_resp packet	UL_DATA_RX	
	mail_num_resp, error	ul_error_resp packet	UL/DL_CMD_ WAIT	
mail_recv	mail_recv_resp, no errors	Change current upload offset	UL_DATA_RX	
	mail_recv_resp, error	ul_nak_resp packet	UL_ABORT	
mail_close_req	mail_close_resp, no errors	ul_ack_resp packet	UL/DL_CMD_ WAIT	
	mail_close_resp, error	ul_nak_resp packet	UL_ABORT	
mail_del_req	mail_del_resp	del_resp packet	UL/DL_CMD_ WAIT	
mselect_req	mselect_resp, no errors	select_resp packet	UL/DL_CMD_ WAIT	
	mselect_resp, error	dl_error_resp packet	UL/DL_CMD_ WAIT	

TABLE 7.4 STATE TRANSITIONS FROM WAIT_MAILBOX				
Message	Reply Message	Action	Next State	
dir_req	directory, no errors	data packet, data_end packet	UL/DL_CMD_ WAIT	
	directory, error	dl_error_resp packet	UL/DL_CMD_ WAIT	
disconnect			UL/DL_ UNINIT	

# 4. UL\_DATA\_RX

The packet transfer module enters the UL\_DATA\_RX state after the user has sent an *upload\_cmd* and the mailbox control module has replied to the resulting inquiry with the appropriate file number or offset. That is, this state is first entered from the WAIT\_MAILBOX state. When the module is in the UL\_DATA\_RX state, it is ready to receive *data* packets from the user. As each packet it received, the packet transfer module passes the file data on to the mailbox control module for storage, entering the WAIT\_MAILBOX state each time to await acknowledgement. When the *data\_end* packet is received, the packet transfer module requests that the mailbox control module close the file and conduct integrity checks on it. The result of these checks will determine whether the packet transfer module returns directly to the UL\_DL\_CMD\_WAIT state, or goes into the UL\_ABORT state, as indicated in Table 7.4. The state transitions out of UL\_DATA\_RX are summarized in Table 7.5. The only legal user packets which can be received while in this state are *data* and *data\_end*. If the packet transfer module receives an unexpected packet while in this state, it will send a

'mail\_close\_req' to the mailbox control module, an ul\_error\_resp packet to the user, and then return to the UL/DL\_CMD\_WAIT state. If the user becomes disconnected while the module is in the UL\_DATA\_RX state, it will send a 'mail\_close\_req' message to the mailbox control module and then return to the UL/DL\_UNINIT state.

TABLE 7.5 STATE TRANSITIONS FROM UL_DATA_RX				
Received Packet or Message	Message Sent	Next State		
data	mail_recv	WAIT_MAILBOX		
data_end	mail_close_req	WAIT_MAILBOX		
other packets	mail_close_req, ul_error_resp Packet	UL/DL_CMD_WAIT		
disconnect	mail_close_req	UL/DL_UNINIT		

### 5. UL\_ABORT

The UL\_ABORT state is entered whenever a problem is found with an ongoing upload prior to receipt of the data\_end packet. While the packet transfer module is in the UL\_ABORT state, all data packets are discarded. It will remain in this state until a data\_end packet is received, an unexpected packet is received, or the user is disconnected. The state transitions out of UL\_ABORT are summarized by Table 7.6.

TABLE 7.6 STATE TRANSITIONS FROM UL_ABORT			
Received Packet or Message	Packet Sent	Next State	
data	none	UL_ABORT	
data_end	none	UL/DL_CMD_WAIT	
other packets	ul_error_resp	UL/DL_CMD_WAIT	
disconnect		UL/DL_UNINIT	

### 6. DL\_FILE\_DATA

The state DL\_FILE\_DATA is entered from the UL/DL\_CMD\_WAIT state whenever a properly formatted download\_cmd is received from the user. If a badly formatted download\_cmd is received, the user will be sent a dl\_error\_resp and the packet transfer module will remain in the UL/DL CMD\_WAIT state.

Just prior to entering the DL\_FILE\_DATA state, the packet transfer module sends a 'mail\_req' message to the mailbox control module. While in the DL\_FILE\_DATA state, the packet transfer module simply waits for 'mail\_resp' messages from the mailbox module containing file data to be transmitted to the user. As each piece of the file arrives, it is sent on to the user in a data packet. When the mailbox control module indicates that the last byte of the file has been provided, a data\_end packet is sent to the user. The packet transfer module remains in the DL\_FILE\_DATA state until either a dl\_ack\_cmd or a dl\_nak\_cmd is received from the ground. Then it returns to the UL/DL\_CMD\_WAIT state. The satellite takes no particular action upon receipt of a dl\_nak\_cmd. It will be the responsibility of the ground station to request a

new download of the same file at a future time if the user so desires. If the file number requested for download does not exist, a <u>dl\_error\_resp</u> packet will be transmitted and the module will return immediately to the UL/DL\_CMD\_WAIT state. The state transitions from DL\_FILE\_DATA are shown in Table 7.7.

TABLE 7.7 STATE TRANSITIONS FROM DL_FILE_DATA				
Received Message or Packet	Packet Sent	Next State		
mail_resp, error	dl_error_resp	UL/DL_CMD_WAIT		
mail_resp, no error	data	DL_FILE_DATA		
mail_resp, end of file	data_end	DL_FILE_DATA		
dl_ack_cmd	dl_completed_resp dl_ack message	UL/DL_CMD_WAIT		
dl_nak_cmd	dl_aborted_resp	UL/DL_CMD_WAIT		
other Packets	dl_error_resp	UL/DL_CMD_WAIT		
disconnect		UL/DL_UNINIT		

#### VIII. MAILBOX CONTROL MODULE

#### A. FUNCTION

The primary role of the MAILBOX\_CONTROL module is to keep track of the mail files which have been uploaded to PANSAT from users on the ground. It also keeps track of user-accessible telemetry files which have been prepared by the TELEMETRY\_GATHER module for downloading to interested users, and "bulletins" which have been posted by the ground control station for the information of all PANSAT clients. The users' active selection lists are also maintained by the mailbox control module.

The mailbox control module has only one state, WAIT. This state name exemplifies the method employed by the module to carry out its duties. It "waits" until it receives a request for information or a packet of file data from the packet transfer module, or is notified of a file posted by the telemetry module or by the ground control module. Most housekeeping functions within the "mailbox" are triggered by receipt of these messages. The mailbox control module responds to the received message, carries out any necessary activity, and then continues waiting until the next message arrives. A few administrative functions, such as purging all mail, must be directed by special command messages from the ground control or auto control modules.

#### B. SOURCE RECORDS

The method employed by the mailbox control module to keep track of all uploaded files and all users' active selection lists is a linked list of Source\_Records. The Source\_Record type is a data structure which contains information which links every stored file with the source from which it was originally uploaded, as well as an active selection list for any source (user) that has requested one. The fields of the Source\_Record are listed in Table 8.1, along with the function of each field.

	TABLE 8.1 FIELDS OF THE SOURCE RECORD		
Field	Туре	Function	
source_num	uint	Contains a unique integer assigned to each ground user who has uploaded any files currently stored onboard the satellite or has an active selection list. Used as the first 2 bytes in the file numbers assigned to each file uploaded by this user.	
call	Callsign_Type	The call sign belonging to the client assigned the above 'source_num'. The call sign will be used as the DOS file name for all files uploaded by this client. Each file will be assigned an extension from "001" to "999".	
selected	Select_List	The Select_List structure includes the fields 'num_sel', a uint indicating the number of files in the client's selection list, and 'sel', a variable length array of the mail numbers of those files. 'num_sel' must be <= max_mail, the maximum number of files allowed in one selection list. A 'num_sel' equal to '0' indicates "no active selection list".	

	TABLE 8.1 FIELDS OF THE SOURCE RECORD			
Field	Туре	Function		
next_mail	uint	The index into the 'sel' array which marks the "next" mail file in the selection list not yet downloaded by the client. When 'next_mail' becomes > = 'num_sel', the selection list is "empty" if another request to download the "next" file arrives.		
next_dir	uint	The index into the 'sel' array which marks the "next" mail file in the selection list for which a directory entry has not yet been downloaded by the client. When 'next_dir' becomes >= 'num_sel', the selection list is "empty" if another request to download directories for the "next" 10 files arrives. When both 'next_mail' and 'next_dir' are >= 'num_sel', 'num_sel' reverts to '0' and the client no longer has an active selection list.		
next_ext	File_Ext = 000999	The next file extension to be used on a file uploaded by this client. In binary form, the 'ext' is used as the last 2 bytes in the file number assigned to the file. In ascii form, it forms the 3 character DOS file extension.		
num_act	uchar	The number of files uploaded by this client which are still being stored aboard the satellite.		
next_num	^Source_Record (pointer to Source_Record)	Pointer to the next Source_Record, numerically by 'source_num'.		
next_call	^Source_Record	Pointer to the next Source_Record, alphabetically by 'call'.		

#### C. RESPONSE TO MESSAGES FROM THE PACKET TRANSFER MODULE

By far the greatest number of messages received by the mailbox control module originate from the packet transfer module, via the Mailbox\_Access Channel. Chapter VII lists many transitions of the packet transfer module to the WAIT\_MAILBOX state. These transitions indicate that the packet transfer module has requested information from the mailbox control module and is awaiting a reply. The packet transfer module also sends messages to the mailbox control module while remaining in the UL\_DATA\_RX state. The activities of the mailbox control module triggered by each message type from the packet transfer module, and the required reply messages, are addressed in the following subsections.

## 1. The 'active\_sl\_req' and 'active\_sl\_resp' Messages

When the packet transfer module sends an 'active\_sl\_req' message, it is inquiring whether there is an active selection list for a particular user. The mailbox control module must check the source records to see if the user has an active selection list or not. An 'active\_sl\_resp' message is returned to the packet module, indicating true if the user does have an active list, and false otherwise.

# 2. The 'mail\_num\_req' and 'mail\_num\_resp' Messages

A 'mail\_num\_req' message indicates that a user wants to upload a file. If this is a new file, a file number is required for it. If it is an upload continuation, the current file offset is needed. The mailbox module must ensure that there is enough room in memory to store a file of the length indicated in the message. If the indicated file number is '0x00000000', the mailbox will get the user's source number and next extension from the source records (or assign a new source number if necessary) and form a new file number. If the file number in the 'mail\_num\_req' message is not '0x00000000', the mailbox module will find the current length of it's partial copy of the file. In the 'mail\_num\_resp' message, the mailbox module will supply the packet module with the required file number or offset for the upload, or indicate that an error has occurred (such as insufficient space or incorrect file number).

### 3. The 'mail recv' and 'mail recv resp' Messages

The 'mail\_recv' message passes file data which has been received from a user to the mailbox module. The data must be appended to the appropriate file. The mailbox module attempts to find and open the file to which the data belongs and append it. The 'mail\_recv\_resp' will indicate whether the data has been stored successfully or whether an error has occurred.

# 4. The 'mail\_close\_req' and 'mail\_close\_resp' Messages

The 'mail\_close\_req' message can indicate one of two situations. Either a data\_end packet has arrived, indicating that an upload has been completed, or an upload has been interrupted due to user disconnect or an unexpected packet. If an upload has

been completed, the packet module will indicate this by setting the 'req\_resp' parameter of the message to true, requesting a response. In this case, the mailbox module will check the integrity of the uploaded file and report the results in the 'mail\_close\_resp' message. If 'req\_resp' is set to false, the upload has been interrupted and the mailbox module will simply close the file and wait for the upload to be continued at a later time.

### 5. The 'mselect req' and 'mselect resp' Messages

An 'mselect\_req' message forwards to the mailbox module the Select\_Structure of a client requesting to form a new active selection list. The mailbox module must parse the Select\_Structure and either prepare the default selection list or evaluate the selection equation with respect to each file in the mail box. The file number of each matching (or default) file will be placed in the client's selection list. There is a maximum number of file numbers which can be placed in any selection list. When this number is reached, further selection will be discontinued. The 'mselect\_resp' message indicates how many files have been placed in the selection list, or if an error has occurred. Any prior existing list will be discarded.

# 6. The 'mail\_req' and 'mail\_resp' Messages

The packet transfer module sends a 'mail\_req' message in order to obtain file data for download to a client. A file number and offset will be included in the message.

If the "next" file in the selection list is requested, the indicated file number will be

'0x0000000', and the mailbox module must consult the client's source record to determine the actual file number of the next file in the list. The offset for the "next" file will always be zero. When the next data set from that file is requested, the file number and appropriate non-zero offset will be known, and included in the 'mail\_req' message. The mailbox module will begin at the appropriate file offset and begin copying bytes into the data buffer. It will copy either the number of bytes which will fit into one packet, or the remaining bytes in the file, whichever is less. Either the data buffer or an error indication will be sent back to the packet module in the 'mail\_resp' message. When the end of a file has been sent to the packet module, the mailbox module responds to the next 'mail\_req' with an empty data buffer and no error code. This indicates to the packet buffer that it is time to send the data end packet.

### 7. The 'dl\_ack' Message

The packet transfer module will send a 'dl\_ack' message to the mailbox control module after receiving a dl\_ack\_cmd packet from the user. Only if a 'dl\_ack' message is received will the mailbox module change the 'next\_mail' field in the user's source record. The 'next\_mail' pointer is only advanced after the file indicates has been successfully downloaded to the client. The number of the file acknowledged will be included in the 'dl\_ack' message along with the client's call sign. The file number must match that indicated by the 'next\_mail' field of the client's source record for the field to be updated.

### 8. The 'dir\_req' and 'directory' Messages

The 'dir\_req' message requests directory information for either the file number indicated, or the "next" ten files in the client's active selection list. Directory information for a file is simply a copy of the PANSAT file header. If a file number is indicated, the mailbox module places a copy of the appropriate header in the data buffer which is send back with the 'directory' message. If the "next" 10 entries are requested, the mailbox module consults the source record to determine whether there is an active list, and if so, which is the "next" file for which a directory entry has not yet been sent. The headers are copied for the next 10 files on the list, beginning with the one marked by 'next\_dir'. If there are less than 10 remaining on the list, they are all sent. There is no downland acknowledge associated with directories, and the 'next\_dir' counter is automatically advanced when the 'directory' message is sent back to the packet transfer module.

### 9. The 'mail del req' and 'mail del resp' Messages

The packet transfer module sends a 'mail\_del\_req' when a user wishes to delete a file from the satellite's mailbox. The mailbox module must first ensure that the user in question is authorized to delete the indicated file. A user may only delete a file which they have uploaded, or one which is addressed to them as the sole recipient. The mailbox module knows who uploaded the file, since the file name is the same as the source call sign. It can consult the destination fields of the file header to determine

whether the requesting user is the sole recipient. If the deletion is authorized, it will be carried out, and a no\_error indication returned to the packet module in the 'mail del resp' message. Otherwise, the er permission\_denied code will be returned.

#### D. RESPONSE TO MESSAGES FROM OTHER MODULES

The mailbox control module may also be tasked to respond to messages from modules other than the packet transfer module. These messages may come via one of the Mailbox\_Admin\_Channels or via the Telemetry\_Storage\_Channel. The latter channel is connected to the TELEMETRY\_GATHER module, while one copy of the former is connected to the AUTO\_CONTROL module and another is connected to the GROUND\_CONTROL module. None of these three modules has been completely specified, and the requirements for them are still evolving. Some possible functions for them have been suggested, and those which impact upon the mailbox control module will be discussed in the following subsections.

## 1. The 'list\_mail' and 'mail\_list' Messages

The NPS ground control station personnel retain the right to inspect all messages in the mailbox, regardless of the upload sources or the addressees. The ground control station, when it is so desired, can request a list of all files currently maintained in the memory, or a partial list of only those files "from" or "to" a particular call sign.

This command is received by the GROUND\_CONTROL module, which responds

requesting the appropriate file list from the mailbox control module using a 'list\_mail' message. The mailbox module responds with a 'mail\_list' message which indicates the number of files matching the criteria of the 'list\_mail' message and provides a list of all of the appropriate file numbers. From this list, the ground control module or the ground control station personnel can then choose files to download.

### 2. The 'post\_bulletin' and 'delete\_bulletin' Messages

The ground control module has the same access to the file handling facilities of the Space Craft Operating System as does the mailbox control module. For this reason, it does not need to go through the mailbox module in order to "post" a bulletin, which really consists only of storing a file with the name "BULLETIN.xxx" in the mail storage area. (File lists such as those discussed in the previous subsection are requested from the mailbox module merely to take advantage of its enhanced association capabilities using the source records it maintains.) The mailbox module should, however, maintain a complete set of source records, including one for the ground control station. When a bulletin is posted, the ground control module informs the mailbox module using a 'post\_bulletin' message, so that an appropriate file number can be assigned and the source record can be updated. Similarly, when a bulletin is deleted, the mailbox module is informed by a 'delete bulletin' message.

# 3. The 'full\_mailbox' and 'purge\_mail' Messages

Whenever a user requests to upload a file, the mailbox module must first determine whether there is room for it in the memory. If it finds that there is not enough room, it does some "house-cleaning", deleting all files which have passed their expiration dates. This is the only time the mailbox module deletes files on its own, so that many files may actually remain onboard the satellite for longer than the nominal time allowed. After the mailbox module has deleted all files which have expired, it once again checks to see if there is enough room to upload the new file as requested by the user. If there is still not enough room, the mailbox module must deny the request to upload. At the same time, it informs the AUTO\_CONTROL module of the problem with a 'full\_mailbox' message.

Perhaps in response to a 'full\_mailbox' message, or perhaps in obedience to a ground station command, or for some other pressing reason, the ground control or auto control module can direct the mailbox module to "purge" the mail box. The 'purge\_mail' message will indicate whether all mail files should be deleted, or all files posted prior to some designated upload time, or all files "from" or "to" a particular call sign. This purge is done via the mailbox module, so that it will have the chance to update all affected source records.

### 4. The 'store\_user\_telem' and 'delete\_user\_telem' Messages

Like the ground control and auto control modules, the TELEMETRY\_GATHER module also has complete access to the SCOS file management capabilities. When user-accessible telemetry data is to be posted, it merely saves a file called "USRTELEM.xxx" in the mail storage area. These telemetry files can also be deleted by the telemetry module when they become outdated. In the interest of maintaining a complete set of source records, the mailbox control module is informed of these actions via the 'store\_user\_telem' and 'delete\_user\_telem' messages.

#### IX. REMAINING MODULES

#### A. TELEMETRY GATHERING MODULE

In the current PANSAT flight software specification, 14 separate software modules have been defined at the module header definition level. Of these, detailed module body definitions have been developed The DATA TRANSFER, PACKET TRANSFER, and MAILBOX CONTROL modules are described in Chapters V through VIII of this thesis. A preliminary module body definition for the PASSWORD CONTROL module has been written, but will not be released to the general public. Two modules, PRIMITIVE AX25 and PRIMITIVE SW LOADER, are actually commercial software products, BAX and PHTX. The capabilities of these programs will be accessed by various PANSAT modules, but no body definitions will be written for them, because there is no need to specify existing software, only the interfaces to it. The body definitions of the remaining 8 modules will be highly dependent upon the actual hardware configuration of the satellite, which is still undergoing daily design changes. Central to the operation of these remaining modules will be the operation of the TELEMETRY GATHER module.

The function of the telemetry gathering module is to collect data on the operation of the satellite from which control decisions can be made, both by the automatic control module (AUTO\_CONTROL) and the ground control station personnel at NPS. In order to obtain much of this data, the telemetry gathering module has direct control over the

A/D\_DRIVER module which operates the analog-to-digital converters and associated multiplexors in order to obtain relevant sensor data, such as battery voltages or solar array temperatures. Other telemetry information will come from the BAX and SCOS software, which maintain various statistics about the communications and operating environments.

The hardware telemetry points which have been defined thus far are listed in Table 9.1. The best situation is for each point to stay within the expected or "nominal" range. When a reading goes outside the nominal range, there is still no serious system degradation unless it also goes outside the "operating range". At this point, there may be no immediate danger to the system, but a trend may have started which will soon lead to operational difficulties. When a reading goes outside the "red alert" range, immediate correctional actions must be initiated, if they have not been already. System failure could be imminent. Many of the exact values for these ranges have not yet been determined. The proper preventive and/or correctional steps to be taken in each situation are also still under study. The values contained in Table 9.1 are the best estimates available at this time, but are subject to change. Those readings for which no estimated values have yet been determined are marked with "tbd". The "totals" listed are for the sensors controlled by one Digital Control System (DCS) board, on which will be running one copy of the flight software. The current design calls for the entire DCS to be duplicated, and for each board to be attached to its own complete and separate set of sensors.

TABLE 9.1 HARDWARE TELEMETRY POINTS							
Point	Nom	Nominal		Operating		Red Alert	
	Min.	Max.	Min.	Max.	Min.	Max.	
Solar Array Temperatures (17 total)	O° C	50° C	-20° C	120° C	-30° C	140° C	
Battery Voltages (2 total)	12 V	13 V	11.5 V	13.5 V	10 V	15 V	
Battery Temperatures (4 total)	-1.1° C	10° C	-6.7° C	26.7° C	-15° C	50° C	
Battery Discharge Currents (2 total)	tbd	tbd	tbd	tbd	tbd	tbd	
Electrical Power System (EPS) Bus Voltage (1 total)	tbd	tbd	tbd	tbd	tbd	tbd	
EPS Board Temperature (2 total)	°С	40° C	-10° C	50° C	tbd	tbd	
Transmitter Current (1 total)	tbd	tbd	tbd	tbd	tbd	tbd	
Transmitter RF Power (2 total)	tbd	tbd	tbd	tbd	tbd	tbd	
Transmitter Temperature (2 total)	0° С	40° C	-10° C	50° C	tbd	tbd	
Received Signal Strength (2 total)	tbd	tbd	tbd	tbd	tbd	tbd	
Receiver Temperature (2 total)	0°C	40° C	-10° C	50° C	tbd	tbd	
Sense Relays for State of Communications Hardware (total tbd)	tbd	tbd	tbd	tbd	tbd	tbd	
DCS Board Temperature (2 total)	0°C	40° C	-10° C	50° C	tbd	tbd	

The telemetry gathering module maintains a list of sensor points with timing intervals and expected operating ranges for each. This list can be updated by commands from the ground control station, which can cause points to be added or deleted, or can change the timing intervals for obtaining readings from various points. Some timing intervals may be changed dynamically by the automatic control module or the telemetry gathering module itself, based upon trends in the readings or upon reading which are out of the expected ranges.

Table 9.2 lists some "operating environment telemetry points" which can be gathered by SCOS, and passed to the telemetry gathering module for inclusion in the telemetry files. Table 9.3 lists some "communications environment telemetry points" which can be gathered by BAX, and may be of interest to the ground station controllers. BAX has the capability to maintain a file of this data itself and to download it directly to the ground control station. Whether the information will be passed to the telemetry gathering module to be included with the rest of the telemetry, or whether this separate "BAX telemetry" file will be maintained and passed to the ground control station as the result of a separate ground station command, has not yet been determined. Table 9.4 contains a list of other general system data which may be collected by the telemetry gathering module directly from the satellite hardware or from the other software modules. In some cases, such as the data points listed under "LOGIN" and "MAILBOX", existing module specifications will have to be modified in order to require the software to gather the data required by the telemetry module. Such modifications

will be postponed until it has been decided which of these data points will be of most interest to the ground station controllers, and what sampling intervals will be required.

TABLE 9.2 SCOS TELEMETRY POINTS			
Data Point	Data Type	Description	
Scheduler Events	List of numbers	Operating System multi- tasking events (tasks running & scheduled).	
Timer Events	List of numbers	Operating System tasks in queue.	
SCOS Service Calls	List of numbers	General Operating System information.	

TABLE 9.3 BAX TELEMETRY POINTS			
BAX Data Point	Data Type	Description	
smallct	uint	Count of received frames containing < 32 bits.	
nonint	uint	Count of received frames with a bit length not evenly divisible by 8	
bigent	uint	Count of received frames that exceed maximum size.	
abortent	uint	Count of received frames that are aborted.	
overcnt	uint	Count of receiver overruns.	
crc	uint	Count of receiver crc errors.	
tx_aborted	uint	Count of transmitted frames aborted or flushed.	
tx_under	uint	Count of transmitted frame underruns.	
tx_abort_call	uint	Count of calls to qio_abort/flush.	
qiocurrx	uint	Current number of frames in receiver queue.	
qiomaxrx	uint	Maximum number of frames in receiver queue.	

TABLE 9.3 BAX TELEMETRY POINTS			
BAX Data Point	Data Type	Description	
qiocurtx	uint	Current number of frames in transmitter queue.	
qiomaxtx	uint	Maximum number of frames in transmitter queue.	
poolfail	uint	Number of "pool gets" that failed.	
retry_exceeded	uint	Count of times the maximum number of frame retires has been exceeded.	
quitottx	ulong	Total number of transmitted frames.	
quitotrx	ulong	Total number of frames received with no errors.	
tdatain	ulong	Total number of data bytes received.	
tdataout	ulong	Total number of data bytes transmitted.	
tdigi	ulong	Total number of digipeated frames.	
daytime	ulong	Total number of 50msec intervals that have expired since system startup.	
start_time	ulong	Startup time in seconds. (UTC)	
Following are	counts of	f frames types defined in the AX.25 protocol [Ref. 2].	
<i> in</i>	ulong	Number of "information" frames received.	
<rr> in</rr>	ulong	Number of "receive ready" frames received.	
<rnr> in</rnr>	ulong	Number of "receive not ready" frames received.	
<rej> in</rej>	ulong	Number of "reject" frames received.	
<dm> in</dm>	ulong	Number of "disconnect mode" frames received.	
<sabm> in</sabm>	ulong	Number of "set asynchronous balanced mode" (connect request) frames received.	
<disc> in</disc>	ulong	Number of "disconnect request" frames received.	
<ua> in</ua>	ulong	Number of "unnumbered acknowledge" frames received.	
<frmr> in</frmr>	ulong	Number of "frame reject" frames received.	

TABLE 9.3 BAX TELEMETRY POINTS			
BAX Data Point	Data Type	Description	
<inv> in</inv>	ulong	Number of "invalid" frames received.	
<ui> in</ui>	ulong	Number of "unnumbered information" frames received.	
<1> out	ulong	Number of "information" frames transmitted.	
<rr> out</rr>	ulong	Number of "receive ready" frames transmitted.	
<rnr> out</rnr>	ulong	Number of "receive not ready" frames transmitted.	
<rej> out</rej>	ulong	Number of "reject" frames transmitted.	
<dm> out</dm>	ulong	Number of "disconnect mode" frames transmitted.	
<sabm>out</sabm>	ulong	Number of "set asynchronous balanced mode" frames transmitted.	
<disc> out</disc>	ulong	Number of "disconnect request" frames transmitted.	
<ua> out</ua>	ulong	Number of "unnumbered acknowledge" frames transmitted.	
<frmr>out</frmr>	ulong	Number of "frame reject" frames transmitted.	
<inv> out</inv>	ulong	Number of "invalid" frames transmitted.	
<ui> out</ui>	ulong	Number of "unnumbered information" frames transmitted.	

TAB	LE 9.4 GENERA	L SYSTEM TELEMETRY POINTS	
Data Point	Data Type	Data Type Description	
		LOGIN Data	
Logins	uint	Number of user logins.	
Logouts	uint	Number of user logouts (requested disconnects).	
ALogins	uint	Number of authorized logins.	

TABLE 9.4 GENERAL SYSTEM TELEMETRY POINTS			
Data Point	Data Type	Description	
UALogin	uint	Number of unauthorized login attempts.	
UALtime	UTC	Unauthorized login attempt time stamp.	
Uuser	array of Callsign_ Type	List of Undesirable Users.	
MAILBOX Data			
RMail	uint	Count of received mail.	
SMail	uint	Count of sent mail.	
StMail	uint	Count of stored mail.	
Stor	ulong	Amount of storage used.	
	Commu	nication System Data	
Receiver	boolean	Receiver A or B selected.	
Transmitter	boolean	Transmitter A or B selected.	
Mode	boolean	Spread Spectrum turned On or Off.	
Atten	uint	Attenuation Level 1 through 8 selected.	
Digital Control System Data			
DCS			
DCS SWver	Digital	Control System Data	
	Digital boolean	Control System Data  DCS A or B selected.	
SWver	Digital boolean uint	Control System Data  DCS A or B selected.  Software version in use.	
SWver date	Digital boolean uint UTC	Control System Data  DCS A or B selected.  Software version in use.  Current satellite date and time.  EDAC (error detection and correction) SEU	
SWver date SEUc	Digital boolean uint UTC uint	Control System Data  DCS A or B selected.  Software version in use.  Current satellite date and time.  EDAC (error detection and correction) SEU (single event upset) count.	

As the telemetry gathering module completes each round of readings, it updates a file of the "current telemetry" which is accessible to the automatic control module. The automatic control module makes use of this data in its autonomous control of the satellite hardware systems. The telemetry gathering module also stores telemetry data in a telemetry history file, which will continue to grow and store past data until it is purged by a command from the ground control station, or it reaches a pre-determined maximum size. If the maximum file size is exceeded before the file can be downloaded and then purged by the ground control station, the most recent entries for each data point will be maintained, and older entries deleted, in order to control the size of the file. The ground control station will use this larger telemetry file to analyze trends in satellite performance, and to make control decisions beyond the scope of those made by the automatic control module. The telemetry gathering module will also maintain shorter telemetry files containing data which may be of interest to the amateur radio users who access the satellite mail box system. These files are stored in the mail area with the file name "USRTELEM.xxx".

#### **B.** AUTOMATIC CONTROL MODULE

The AUTO\_CONTROL module carries out periodic functions, such as battery conditioning, on a time scheduled basis. It also carries out aperiodic functions. As indicated in the previous section, the AUTO\_CONTROL module makes use of the data

collected by the telemetry gathering module to make decisions about the control of the satellite hardware. It also maintains a "time\_tagged" command buffer which lists activities which should take place at a particular time in the future. This command buffer is updated by the GROUND\_CONTROL module as a result of ground control station commands. The Digital Control System design includes hardware timers which can be programmed by the automatic control module to interrupt the microprocessor at designated time intervals to initiate periodic events or to produce a set of one-time-only interrupts to initiate events controlled by the command buffer. Software timers may also be used for some of the automatic control module functions.

Most of the control functions carried out by the automatic control module will likely be based on a "table look-up" system. When a timer interrupt occurs, an interrupt vector table will contain the address of the appropriate subroutine needed to carry out the scheduled activity. Another table of subroutine addresses will be indexed based on combinations of telemetry readings which call for some action to be taken. These subroutines and tables will be developed as more is learned about the specific requirements of the hardware as it is designed.

Table 9.5 lists some possible functions of the automatic control module which have been identified thus far. The EPS\_DRIVER, COMM\_DRIVER and DCS\_DRIVER modules contain the software drivers required for direct digital control of the electric power system, communications, and digital control system hardware. The services of

these modules will be accessed as necessary by the automatic control module in order to carry out functions listed in Table 9.5.

TABLE 9.5 AUTOMATIC CONTROL MODULE FUNCTIONS		
Function	Description	
Electric Power Supply Control	Turn hardware components off and on as necessary to conserve power, allow battery conditioning, etc.	
Condition Batteries	Periodically discharge and recharge batteries in order to prevent battery "memory".	
Systems Test Management	Carry out periodic systems checks in addition to normal telemetry gathering. Save test data for download to ground control station.	
Communications Control	Transmitter/Receiver component select.	
Transmitter Output Power Control	Set level of transmitter power.	
Automatic Subsystem Select	Select alternate subsystem upon time-out waiting for response of a primary subsystem.	
Real Time Clock Control	Set and remove times for periodic interrupts.	
Send Messages	Send periodic messages to the ground control station via BAX.	
Copy Vital Statistics	Transfer vital operating system information to alternate processor.	
RAM Wash	Periodic reading/writing of system RAM to enable Error Detection and Correction functions.	
Digital Control System Health Check	Periodic signal to EPS to ensure proper operation of active DCS. EPS will disable a malfunctioning DCS board and "boot" the alternate when the proper signal is not received on time.	

TABLE 9.5 AUTOMATIC CONTROL MODULE FUNCTIONS	
Function	Description
User Lockout.	Message to data transfer module locking out all or new users.

#### C. GROUND CONTROL MODULE

The GROUND\_CONTROL module contains the command interpreter and the functionality required to carry out commands transmitted by the ground control station at NPS. Ground control packets will be passed directly to the ground control module by BAX, since they will be addressed specifically to the ssid (subsystem identification number) for this module. All ground station commands will be subject to verification by including a time varying password. The PASSWORD\_CONTROL module will keep track of the current password aboard the satellite, and will provide this information as necessary to the ground control module. Similar software will track the current password for the ground control station. There will be facilities for determining the current password aboard the satellite, in case the two systems lose synchronization for any reason. The specification of the password control module contains proprietary information, and will not be published for general release.

Once a command has been received from the ground control station, the password has been verified, and the command has been interpreted, the ground control module either carries out the command directly, or communicates with other software modules as necessary to utilize their capabilities. A ground command may involve updating the

time-tagged command list of the automatic control module, or varying the time intervals for periodic events carried out by the automatic control or telemetry gathering modules. It may initiate a one-time-only corrective action, or change a basic system parameter. Some ground station commands simply involve the acquisition of information for use by the ground control station software or personnel.

Some possible functions of the ground control module which have been identified thus far are listed in Table 9.6. The PRIMITIVE\_SW\_LOADER module, which is actually the commercial program "PHTX", is designed to work directly with BAX to upload software. This module will be utilized by the ground control module when a command is received to upload new software. In this way, the flight software can be updated as necessary to correct errors or increase functionality.

TABLE 9.6 FUNCTIONS OF THE GROUND CONTRC. MODULE	
Function	Description
Command Interpretation/Validation.	Process a received ground station command.
Update Time-Tagged Command Buffer.	Schedule events to be carried out at a future time by the automatic control module, or delete events from the command buffer.
Set Control Rates.	Update time intervals or list of periodic functions of the automatic control module.
Set Telemetry Polling Rates.	Update time intervals used by the telemetry gathering module for particular telemetry points.  Add or delete telemetry points.

TABLE 9.6 FUNCTIONS OF THE GROUND CONTROL MODULE		
Function	Description	
Software Upload.	Upload, and store new or updated software modules.	
Run Software.	Begin using newly uploaded or alternate software module.	
Delete Software.	Delete specified software module.	
Copy Software.	Copy verified software to alternate processor.	
Boot ROM.	Reboot PANSAT from ROM (read only memory).	
Boot OS.	Load a new operating system and transfer con	
Read OS Information.	Download the current operating system pointers and parameters.	
List Mail.	Download a list of all mail messages and bulletins currently stored.	
Dump Mail.	Download system bulletins and mail in bulk.	
Post Bulletin.	Post a system bulletin in the mailbox area for all users.	
Remove Bulletin.	Remove a system bulletin.	
Purge Mail.	Purge all or selected mail from the mailbox storage.	
Read Current Telemetry.	Download the current telemetry file maintained by the telemetry gathering module.	
Read Stored Telemetry.	Download the telemetry history file maintained by the telemetry gathering module.	
Purge Stored Telemetry.	Delete all or portions of the telemetry history file.	
Read Data.	Download an arbitrary block of data, specified by address pointer, from the file storage area or system RAM.	
Set Real Time Clock.	Set satellite's real time clock to a specified time.	
Read Real Time Clock.	Download current time on satellite's real time clock.	

TABLE 9.6 FUNCTIONS OF THE GROUND CONTROL MODULE		
Function	Description	
Subsystem Power Control.	Turn power on/off to a particular subsystem.	
Condition Battery.	Discharge/Recharge specified battery.	
Trickle Charge Battery.	Trickle charge specified battery.	
Charge Battery.	Quick charge of specified battery.	
Select Battery.	Select specified redundant battery.	
Select Receiver.	Select specified redundant receiver.	
Select Transmitter.	Select specified redundant transmitter.	
Select Processor.	Select specified redundant digital control system board.	
Set Mode.	Select communications mode: spread spectrum or BPSK.	
Set Maximum Transmitter Power.	Set maximum allowable amplitude of transmitter power.	
Set Attenuation.	Set attenuation level of the active transmitter.	
Switch to Super User Mode	Functions requiring super user mode are tbd.	
Exit Super User Mode.		
Read Event Log.	Download Event Log maintained by the event logging module.	
Purge Event Log.	Delete all or portions of the event log,.	
Read Time-Tagged Command Buffer.	Download the time-tagged command buffer.	
Purge Time-Tagged Command Buffer.	Delete the entire time-tagged command buffer.	
User Lockout.	Message to data transfer module locking out all or new users.	

### D. EVENT LOGGING MODULE.

The purpose of the EVENT\_LOGGING module is to maintain a history of all the significant events which happen and commands which are carried out aboard the satellite. It is hoped that this event log will be helpful in trouble shooting problems aboard the satellite, or merely in studying its operation. The event logging module differs from the telemetry gathering module in one major respect. The telemetry gathering module periodically polls the hardware and other software modules, gathering a predetermined list of specified data. The event logging module waits to receive event messages from other modules, informing it of aperiodic events which are deemed significant in some way.

A list of "significant" events will need to be determined, so that the exact nature of the event messages can be defined in the software specification. Some possible events include the occasion of a full mailbox, a telemetry reading beyond the "operating range" (this will also be listed in the telemetry files, of course, but may stand out more here, or be associated with some other event which will make trouble shooting and correction easier), user connections lost because the transmitter has been shut down for power reasons, etc. An event log entry will also be made each time an automatic command function or a ground control command is carried out. The exact format of the event log entries will be developed as the list of significant events and useful information is further defined.

### X. CONCLUSIONS AND RECOMMENDATIONS

#### A. THE USE OF ESTELLE

The formal description technique, Estelle, has proven to be a valuable tool in creating a software specification. Its methods of defining state machine behavior and its channel and message definitions have provided a unique way of visualizing a system, and seeing how all of the pieces fit together. The various levels of abstraction greatly facilitate the advancement of a project, even when all details are not yet known. When details are known, Estelle provides ample means of specification at the lowest possible levels, and the flexibility to define algorithms both simple and complex.

In order to make Estelle even more useful in this project, a few modifications have been made to it. For instance, since "C" has already been chosen as the implementation language, a few data types have been defined to more closely match familiar structures in "C". Array indices start at 0 in this specification, as they do in "C". Multiple dimension arrays are indexed by multiple sets of brackets, "var[i][j]", rather than by multiple indices within one set of brackets, "var[i, j]". The names of the primitive data types are borrowed from "C": "uchar", "uint "and "ulong". Many of the primitive functions and procedures are functions familiar to "C" programmers. In addition, various font modifications have been used to make elements of the Estelle and Pascal syntax stand out, so that their meanings are more obvious in the context. Bold is used

to indicate reserved words, user-defined data types begin with Capital\_Letters, constants are written in *italics*, etc.

Many of the more complex capabilities of Estelle are not utilized, since they are somewhat confusing and are not needed to make clear the intended behavior of the software being defined. The greatest drawback of Estelle is the specification of Estelle itself, [Ref. 8]. [Ref. 8] is very difficult to read and sometimes impossible to understand. For those interested in using Estelle in future software specification projects, it is recommended that only the Annexes be read. These contain all the information needed, as well as adequate examples to provide understanding of how this language can actually be used.

#### **B.** RECOMMENDATIONS FOR FURTHER WORK

This thesis provides a preliminary specification for the flight software of the Petite Amateur Navy Satellite. As much information as is currently available concerning the high-level operational requirements of the satellite has been included. A software architecture has been provided which defines the individual software modules and their interfaces. Detailed definitions for the bodies of the communications and file transfer protocol modules have been developed.

There is obviously much work remaining to be done. The module body definitions for the telemetry gathering, analog to digital conversion, automatic control, ground control, electronic power system driver, communication driver, digital control system driver, and event logging modules must be developed. The channel types and message

interfaces between these remaining modules and between them and the existing modules must be defined in greater detail. Once the complete, detailed specification is available for the entire flight software system, the actual code must be written and tested. The ground software and the bootstrap software must be specified and coded, and the interfaces between these programs and the flight software must be tested. Hardware designs must be completed and tested before any software specifications can actually be finalized. A start has been made, and the beginnings of a road map have been drawn. Much more effort will be required before this project is completed.

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## APPENDIX A - ESTELLE SOFTWARE SPECIFICATION

## specification Flight\_Software;

```
type
                                      { Primitive Types
                                       Note: the notation '0xhh' is used to refer to
                                         hexadecimal numbers, with the h's
                                        representing the hex digits 0..F. The number
                                         of bytes in each hex number is the number of
                                         digits divided by 2.
                                       8 bits of binary data or 1 byte unsigned integer or
   uchar = 0x00..0xFF;
                                         1 ASCII character.
   uint
          = 0x0000..0xFFFF;
                                      { 2 byte unsigned integer
   ulong = 0x00000000..0xFFFFFFFF; { 4 byte unsigned integer
                                     { Positive and Negative integers as defined on
   int
                                         implementation hardware.
                                      { 'Global' Constant Declarations
const
                                      { Maximum length of a file onboard the satellite
   max file length
                    = any ulong;
                    = any uint;
                                      { Maximum pieces of mail in Select list
   max mail
   password length = any uchar;
                                         { Number of characters in the password.
                                      { Max number of bytes in the data field of a packet. }
   max pdat
                    = 2047:
                                      { Maximum length of data field in a BAX frame.
   max fdat
                    = 256:
                    = 30:
                                     { Max channels allowed by BAX.
   maxlinks
                    = any Callsign Type; { Pansat's call sign
   pansat call
   nps call
                    = any Callsign Type;
                                            { NPS' call sign.
                               = 0x00: { Error Codes
   no error
   er ill formed cmd
                               = 0x01:
                                         { Incorrect or unexpected command
                               = 0x02:
   er bad continue
   er_server_fsys
                               = 0x03:
   er no such file number
                               = 0x04;
                               = 0x05:
   er selection empty
   er mandatory field missing = 0x06;
   er no pfh
                               = 0x07:
   er poorly formed sel
                               = 0x08:
   er already locked
                               = 0x09:
   er no such destination
                               = 0x0A:
   er_file_complete
                               = 0x0C; { 0x0B was a repeat of 0x05
                                                                                        }
   er no room
                               = 0x0D:
```

```
er bad header
                           = 0x0E;
er header check
                           = 0x0F;
er body check
                           = 0x10;
                                   { PANSAT specific; not in FTL0.
                                                                                   }
er permission denied
                           = 0x90:
                                 { Global Type Declarations
                                 { . . = 'To be determined'
Tooint Record
Telem Data Type
EPS Cmd Type
EPS_Resp_Type
Comm Cmd Type
Comm Resp Type
DCS_Cmd_Type
DCS Resp Type
Sat Cmd_Type
Sat Resp Type
Event_Report_Type
Callsign Type
                 = array[6] of uchar;
                                        { Any even number of hex digits surrounded }
Byte String = "
                 array[] of uchar ";
                                           by double quotes. Refers to raw binary
                                           data matching the pattern of the hex digits}
                                           and of the same length.
Pdat_Len
                    = 0..max_pdat; { Number of bytes of info in a packet.
                    = 0..max fdat; { Number of bytes of info in a frame.
Fdat Len
                                     { Each member of an enumerated type is assumed}
Lockout_Type
                    = (all, new);
                                        to be associated with a distinct uchar.
Frame Type
                    = (qat data, qat state, qat_ui);
Link_State
                    = (qas connect pend, qas connected,
                       qas_connecting, qas_disconnected,
                       qas_disconnecting, qas_framereject);
                    = (qac_local, qac_remote, qac_remotefrmr,
Cause
                       qac timeout);
                    = array[password length] of uchar;
Password_Type
                    = array[max file length] of uchar;
File Type
Pdata
                    = array[max pdat] of uchar; { Info field of a packet
                    = array[max fdat] of uchar; { Info field of a frame
Fdata
                    = 0..max_mail;
Num Mail
Direction
                    = (left, right);
Packet Type
                    = record
   length_isb:
                 uchar:
   hl:
                 uchar:
   info:
                 Pdata:
end:
Link_Type
                    = 0 \dots maxlinks - 1;
```

Name_Type	= array[12]	of uchar; { DOS file name. Character in 9th
		{ position must be '.' }
File_List	= array[](	of Name_Type; { Variable len array of file names. }
Bit_Type	= 01;	
Control_Block	= record	{ Includes only BAX fields which are used }
link:	uint;	{ 'channel' in BAX manual }
kind:	Frame_Type;	{ 'type' in BAX manual }
l_state:	Link_State;	{ 'state' in BAX manual }
why:	Cause;	{ 'cause' in BAX manual }
my_call:	Callsign_Type;	{ AX25_ADDR or AX25_CALL }
my_ssid:	uchar;	{ in BAX manual }
his_call:	Callsign_Type;	{ Client's Call sign }
his_ssid:	uchar;	{ Client's SSID - may not be needed }
t1:	uchar;	{ t1 frame ACK/NAK timeout timer value }
maxframe:	uchar;	{ frame sliding window size }
retry:	uchar;	{ maximum number of retries for an out frame }
paclen:	uint;	{ maximum size of info field in outgoing packet }
end;		,

```
{ Channel Definitions
                                                                                         }
channel Abstract Bax Channel (Bax End. Data Transfer End):
   by Bax End:
       gax input(in cb: Control Block; idata: Fdata; datal: Fdat Len);
   by Data Transfer End:
       gax claim( out cb: Control Block; grab: uchar);
       gax data( link: Link Type; out cb: Control Block; odata: Fdata; datal; Fdat Len);
       gax busy( link: Link Type);
       gax con acpt;
       gax con rej:
       qax unbusy( link: Link Type);
       gax connect( out cb: Control Block);
       gax ui(link: Link Type; out cb: Control Block; odata: Fdata; datal; Fdat Len);
       gax disconnect( link: Link Type; wait: boolean);
channel Abstract Packet Channel (Data Transfer End, Packet Transfer End);
   by Data Transfer End:
      connection( callsign: Callsign Type);
       disconnect:
       command_packet( command: Packet_Type; datal: Pdat_Len);
   by Packet Transfer End:
       response packet (response: Packet Type);
channel Mailbox Access Channel (Packet Transfer End, Mailbox End);
   by Packet Transfer End:
      active sl req( client call: Callsign_Type); { Does client have an active select list?
       mail_num_req( client_call: Callsign_Type; mail_number, length: ulong);
       mail_recv( mail_number, offset, length: ulong; mail: Pdata);
       mail close req (mail number, offset: ulong; req resp; boolean);
       mselect req( client call: Callsign Type; select struct: Pdata);
       mail req( client call: Callsign Type; mail number, offset: ulong);
      dl_ack( client_call: Callsign_Type; mail_number: ulong);
      dir req( client call: Callsign Type; mail number: ulong);
       mail_del_req( client_call: Callsign Type; mail_number, length: ulong);
   by Mailbox End:
      active sl resp(sl: boolean);
                                      { true if client has an active select list.
                                                                                          }
       mail num resp( mail number, offset: ulong; error code: uchar);
       mail recv resp( error code: uchar);
       mail_close resp( error code: uchar);
      mail_send( mail_number, length: ulong; mail: Pdata);
      mselect resp( num sel: Num Mail; error code: uchar):
      mail_resp( mail: Pdata; mail number, length: ulong; no al: boolean);
      directory( len: Pdat Len; dir: Pdata; no al: boolean);
      mail_del_resp( error_code: uchar);
```

```
channel Mailbox Admin Channel (Control End, Mailbox End);
   by Control End:
      list mail (bulletins, messages, from, to: boolean; callsign: Callsign Type);
      post bulletin( bulletin: Name Type);
      delete bulletin( bulletin name: Name Type);
      purge mail(all, from, to: boolean; callsign: Callsign Type; post time: ulong);
   by Mailbox End:
      mail list( num files: uint; mail: File List);
      full mailbox:
channel Telemetry Storage Channel (Telemetry End, Mailbox End);
   by Telemetry End:
      store user telem( telem: Name Type);
      delete user telem( telem file: Name Type);
channel Password Control Channel (Control End, Password End);
   by Control End:
      password change request;
      request current password;
   by Password End:
      password( pswd:Password Type);
channel Data Transfer Control Channel (Control End, Data Transfer End);
   by Control End:
      change params( out cb: Control Block):
      lockout( l kind: Lockout Type);
      unlock(1 kind: Lockout Type);
      transmitter( off: boolean);
   by Controlled End:
      acknowledge:
channel Telemetry Control Channel (Control End, Telemetry_Gather_End);
   by Control End:
      add point( point: Tpoint Record ):
      delete point( point: Tpoint Record );
      change_timing( point: Tpoint_Record);
      read current telem;
      read stored telem;
      purge stored telem;
   by Telemetry Gather End:
      ack point change( error: uchar);
      current telem( telem: Cur Telem Type);
      stored telem( telem: Full Telem Type);
```

```
channel A/D Control Channel (Command End, A/D Converter End);
   by Command End:
      warmup(device num: uchar);
      start_conversion( telem_point: uchar);
      report data( telem point: uchar);
   by A/D Converter End:
      device ready( device num: uchar);
      data ready( telem_point: uchar);
      telem data( t data: Telem Data Type);
channel SW Load Control Channel (Control End, Loader End);
   by Control End:
      upload( new software: Name Type; sw address: ulong);
   by Loader End:
      upload begin( new software: Name Type);
      upload_complete( new_software: Name_Type);
channel EPS Control Channel (Control End, EPS Driver End);
   by Control End:
      eps cmd( cmd: EPS Cmd Type);
   by EPS Driver End:
      eps_resp( resp: EPS_Resp_Type);
channel Comm_Control_Channel( Control_ind, Comm_Driver_End);
   by Control End:
      comm_cmd( cmd: Comm_Cmd_Type:),
   by Comm Driver End:
      comm resp( resp: Comm Resp Type);
channel DCS Control Channel (Control End, DCS Driver End);
   by Control End:
      dcs_cmd( cmd: DCS_Cmd Type);
   by DSC Driver End:
      dcs_resp( resp: DCS_Resp_Type);
channel Satellite Control Channel (Ground Control End, Auto Control End);
   by Ground Control End:
      sat cmd( cmd: Sat_Cmd_Type);
   by Auto Control End:
      sat_resp( resp: Sat_Resp_Type);
channel Event Log Channel (Event End, Log End);
   by Event End:
      event report( report: Event Report Type);
```

function GET_TIME: ulong; primitive;	{ Global function declarations } { Returns a 32-bit unsigned integer indicating the number of seconds since January 1, 1970. }
function C_BIT_SHIFT( d: Direction; primitive	<pre>f num: uchar; b: uchar): uchar; { Bit-wise shift of the byte specified by 'b' in the } { direction specified by 'd'. 'num' specifies the } { number of bit positions to shift. Returns a 1 } } byte answer. } </pre>
function I_BIT_SHIFT( d: Direction; primitive;	num: uint; b: uint): uint; { Bit-wise shift of the uint specified by 'b' in the } { direction specified by 'd'. 'num' specifies the } { number of bit positions to shift. Returns a 2 } { byte answer. }
function C_BIT_AND( a, b: uchar): primitive;	uchar; { Returns Bit-wise AND of the bytes 'a' and 'b'. }
function I_BIT_AND( a, b: uint): uiprimitive;	nt; { Returns the bit-wise AND of the uints 'a' and 'b'. }
function GET_LSB( number: uint): primitive;	uchar; { Receives a 16 bit number and returns the } least significant 8 bits. }
<pre>function GET_MSB( number: uint): primitive;</pre>	uchar; { Receives a 16 bit number and returns the most significant 8 bits }
function INT( short: uchar): uint; primitive;	{ Receives an 8 bit unsigned number and extends it to { 16 bit unsigned number by prepending 8 0's. }
<pre>procedure QAX_CLEAN_CB( cb: Cc primitive;</pre>	ontrol_Block); { Initializes all fields of the control block structure } { to 0. This is a procedure provided by BAX. }
function FORMAT_EVENT_REPOR'external;	T( event: uint; time: ulong): Event_Report_Type; { This function prepares an Event_Report to be sent } { to the EVENT_LOG module. This function is } { external since the structure of the { Event_Report_Type has not yet been determined.} { The parameter list may have to be modified when} { this function is further defined. }

	{ Module Header Definitions	}
module PRI	MITTIVE_AX25_TYPE systemprocess; { BAX.	}
bax: end;	array[4] of Abstract_Bax_Channel(Bax_End) individual queue;	
module DA7	TA_TRANSFER_TYPE systemprocess; {Between BAX and FTLO	}
pc:	<pre>array[maxlinks] of Abstract_Packet_Channel( Data_Transfer_End) individual queue;</pre>	
bax: cc:	Abstract_Bax_Channel( Data_Transfer_End) individual queue; array[2] of Data_Transfer_Control_Channel( Data_Transfer_End) common queue;	
el: end;	Event_Log_Channel( Event_End) common queue;	
module PAC	CKET_TRANSFER_TYPE systemprocess( link: Link_Type); { FTL0	}
pc: mc: el: end;	Abstract_Packet_Channel( Packet_Transfer_End) individual queue; Mailbox_Access_Channel( Packet_Transfer_End) individual queue; Event_Log_Channel( Event_End) common queue;	
module MA	ILBOX_CONTROL_TYPE systemprocess;	
mc:	array[maxlinks] of Mailbox_Access_Channel( Mailbox_End) individual queue;	
cc: ts:	array[2] of Mailbox_Admin_Channel( Mailbox_End) common queue; Telemetry_Stroage_Channel( Mailbox_End) individual queue;	
ci: end;	Event_Log_Channel( Event_End) common queue;	
module PAS	SWORD_CONTROL_TYPE( first_password: Password_Type; shuffle:Shuffle Type); systemprocess;	
ip		
cc: el:	array[2] of Password_Control_Channel( Password_End) common queue; Event Log Channel( Event End) common queue;	
end:		

```
AUTO CONTROL TYPE systemprocess;
                                                 { Automatic Housekeeping Functions }
      bex:
             Abstract Bax Channel (Data Transfer End) individual queue;
            Data Transfer Control Channel (Control End) individual queue;
      acd:
             Telemetry Control Channel (Control End) individual queue:
      act:
            Password Control Channel (Control End) individual queue:
      acp:
      acm: Mailbox Admin Channel (Control End) individual queue;
            EPS Control Channel (Control End) individual queue:
      acom: Comm Control Channel (Control End) individual queue;
      acdc: DCS Control Channel (Control End) individual queue:
             Satellite Control Channel (Auto Control End) individual queue;
      SC:
      el:
             Event Log Channel (Event End) common queue:
   end:
module GROUND CONTROL TYPE systemprocess;
                                                    { Command Functions
   ip
             Abstract_Bax Channel( Data_Transfer_End) individual queue;
      bax:
             Data Transfer Control Channel (Control End) individual queue;
      ccd:
             Telemetry Control Channel (Control End) individual queue:
      cct:
            Password Control Channel (Control End) individual queue;
      ccp:
      ccl:
             SW Load Control Channel (Control End) individual queue;
      ccm: Mailbox Admin Channel (Control End) individual queue:
             EPS Control Channel (Control End) individual queue;
      ccom: Comm_Control_Channel( Control_End) individual queue;
      ccdc: DCS Control Channel (Control End) individual queue;
      sc:
             Satellite_Control_Channel( Ground_Control_End) individual queue;
      el:
            Event_Log_Channel( Event_End) common queue;
   end:
module PRIMITIVE SW_LOADER TYPE systemprocess; { PHTX
                                                                                   }
   ip
             SW Load Control Channel (Loader End) individual queue;
      cc:
             Abstract Bax Channel (Data Transfer End) individual queue;
   end:
module TELEMETRY GATHER TYPE systemprocess; { Automatic Telemetry Gathering }
   ip
             array[2] of Telemetry Control Channel (Telemetry Gather End)
      cc:
             common queue:
             A/D Control Channel (Command End) individual queue;
      ad:
             Event Log_Channel( Event_End) common queue;
      el:
      ts:
             Telemetry Storage Channel (Telemetry End) individual queue:
   end;
```

```
module A/D_DRIVER_TYPE systemprocess; { Driver for Analog-Digital Conv HW
            A/D_Control_Channel( A/D_Converter_End) individual queue;
   end:
module EVENT_LOGGER_TYPE systemprocess;
     el: array[maxlinks + 6] of Event_Log_Channel( Log_End) common queue;
module EPS_DRIVER_TYPE systemprocess;
            array[2] of EPS_Control_Channel( EPS_Driver_End) common queue;
   end;
module COMM_DRIVER_TYPE system process;
            array[2] of Comm_Control_Channel( Comm_Driver_End) common queue;
      cc:
   end;
module DCS_DRIVER_TYPE systemprocess;
           array[2] of DCS_Control_Channel( DCS_Driver_End) common queue;
   end:
                                 { Module Body Definitions
                                                                             }
body PRIMITIVE AX25_BODY for PRIMITIVE AX25_TYPE;
                                                                  external;
```

```
body DATA_TRANSFER_BODY for DATA_TRANSFER_TYPE;
                                     { AX.25 handler - uses resources of BAX
                    = 0x01:
                                     { SSID of this module
      mail ssid
      maxclients
                    = any uchar;
                                      the max number of 'active users at any time
      t timeout
                                     { number of seconds for frame time-out timer.
                    = any uchar:
      max_frames = 0x07;
                                     { frame sliding-window size.
                                     { max # of retries for outgoing frame
                    = any uchar:
      max tries
      packet_length = max_pdat + 2;
                                        { max size of FTL0 level packet
   type
      Client Num
                    = 0..maxclients;
      Client
                    = record
          callsign:
                              Callsign_Type;
          last comm time:
                              ulong:
          data in progress:
                              boolean:
      end:
      Pac Data = array[packet length] of uchar;
      Client Array = array[maxlinks] of Client;
      Data Record = record
          running length: uint;
          final length:
                          uint;
          data:
                          Pac Data;
      end:
      Data Array = array[maxlinks] of Data Record;
   var
      data:
                           Pac Data;
      length:
                           uint:
      in dat:
                           Data Array: { Array of incoming data on each link
                                                                                      }
      clients:
                           Client_Array;
      cb:
                           Control Block:
                           Client Num;
      num clients:
      new_user_lockout:
                           boolean:
      all user lockout:
                           boolean;
      packet:
                           Packet Type:
      transmit ok:
                           boolean:
      i:
                           uchar;
                                     { general purpose loop counter/ index
                                                                                      }
      b:
                           Bit_Type;
      1:
                          Link_Type;
             NORMAL, BUSY;
   state
                                        { States of DATA_TRANSFER_BODY
                                                                                      }
             EITHER = [NORMAL, BUSY];
   stateset
```

```
function CONCAT( a, b: Pac_Data): Pac_Data;
                                 { Concatenates the array 'b' to the end of array
primitive:
                                     'a', and returns the combined array of uchars.
function PACKET LEN( d: Pdata):
begin
   PACKET_LEN := I_BIT_SHIFT( left, 3, INT( C_BIT_AND( d[1], 0xE0)));
   PACKET LEN := PACKET LEN + INT( d[0]) + 2;
end;
procedure FILL_PACKET( data: Pac_Data; var packet: Packet_Type);
                                 { Takes the uchars from array 'data' and places
primitive;
                                     them, in order, into the record structure of
                                     'packet'.
initialize
                                 { DATA_TRANSFER_BODY
to NORMAL
begin
   for i := 0 to maxlinks do clients[i].callsign := 'none';
   QAX_CLEAN_CB( cb);
   cb.my call := pansat call;
   cb.my_ssid := mail_ssid;
   cb.tl := t timeout;
   cb.maxframe := max_frames;
   cb.retry := max tries;
   cb.paclen := max fdat;
   num clients := 0;
   new user lockout := false;
   all user lockout := false;
   transmit ok := true;
   output bax.qax_claim( cb);
end:
trans
from EITHER to same
when bax.qax input
provided in cp.kind = qat_state and in_cb.l_state = qas_disconnected
begin
   if in cb.callsign < > nps call then num clients := num clients - 1;
   clients[in cb.link].callsign := 'none';
   output pc[in_cb.link].disconnect;
end:
```

```
e cc[b].transmitter
   vided off
   transmit_ok := false;
from EITHER to same
when cc[b].transmittter
provided not off
begin
   transmit_ok := true;
from EITHER to same
when cc[b].change_params
begin
   cb := out_cb;
end:
from EITHER to same
when bax.qax_input
provided in cp.kind = qat_ui
                                { No action required - discard frame
begin
end:
from NORMAL to same
when cc[b].lockout
provided 1_kind = new
begin
   new user_lockout := true;
end;
from NORMAL to same
when cc[b].unlock
provided | kind = new
begin
   new_user_lockout := false;
```

end:

```
from EITHER to same
when pc[l].response packet
provided transmit ok
   length := PACKET_LEN( response) + 2;
   i := 0;
   while i < length do begin
       while i < out cb.paclen and i < length do begin
          data[i] := response[i]:
          i := i + 1:
       end:
       output bax.qax data(1, out cb, data, i);
       length := length - i;
       i := 0:
   end;
end:
from NORMAL to same
when bax.gax input
provided in cb.kind = gat state and in cb.1 state = gas connect pend
begin
   if in cb.his call = nps call then begin
       clients[in cb.link].callsign := nps call:
       clients[in cb.link].last comm time := GET TIME();
       output bax.qax con acpt;
       output pc[in_cb.link].connection( nps call);
   end;
   else
       if num_clients < maxclients and not new user lockout then begin
          num clients := num clients + 1;
          clients[in cb.link].callsign := in cb.his call;
          clients[in cb.link].last comm time := GET TIME();
          output bax.qax con acpt;
          output pc[in_cb.link].connection(in cb.his call);
       else output bax.qax con rej;
end:
```

```
from NORMAL to BUSY
when cc[b].lockout
provided 1_kind = all
begin
   all user lockout := true;
   new user lockout := true;
   for i = 1 to maxlinks do
      if clients[i].callsign < > 'none' and clients[i].callsign < > nps call then
          output qax_busy(i);
end:
from BUSY to NORMAL
when cc[b].unlock
provided | kind = all
begin
   all user lockout := false;
   for i = 1 to max links do
      if clients[i].callsign < > 'none' and clients[i].callsign < > nps call then
          output qax_unbusy(i);
end;
from BUSY to same
when bax.gax input
provided in cb.kind = qat state and
          in cb.1 state = qas connect pend and
          in cb.his call = nps call
begin
   clients[in cb.link].callsign := nps call;
   clients[in cb.link].last comm time := GET TIME();
   output bax.qax con acpt;
   output pc[in_cb.link].connection( nps call);
end;
```

```
from EITHER to same
   when bax.gax.input
   provided in cb.kind = qat_data
   begin
      i := in cb.link;
      clients[i].last comm time := GET TIME();
      if clients[i].data in progress then begin
          data := in dat[i].data;
          length := in dat[i].running_length;
          in dat[i].data := CONCAT( data[0..length], idata);
          length := length + datal;
          if length < in_dat[i].final_length then
             in dat[i].running_length := length;
          else begin
              clients[i].data in_prograss := false;
              FILL PACKET( in dat[i].data, packet);
              output pc[i].command packet( packet, in_dat[i].final_length-2);
          end:
      else begin
          length := PACKET LEN( idata);
          if datal < length then begin
              clients[i].data in progress := true;
              in_dat[i].data := idata;
              in dat[i].running length := datal;
              in dat[i].final length := length;
          else begin
              FILL PACKET( idata, packet);
              output pc[in cb.link].command_packet( packet, length - 2);
          end:
       end:
   end;
end; { of Data_Transfer_Body }
```

# body PACKET\_TRANSFER\_BODY for PACKET\_TRANSFER\_TYPE;

```
{ Constants for PACKET TRANSFER BODY
                                     { Packet Types
      data
                            = 0x00;
                           = 0x01:
      data end
                           = 0x02;
      login resp
                           = 0x03;
      upload cmd
                           = 0x04:
      ul go resp
      ul error resp
                           = 0x05;
                           = 0x06:
      ul ack resp
                           = 0x07:
      ul nak resp
      download cmd
                            = 0x08:
      dl error resp
                           = 0x09;
                           = 0x0A:
      dl aborted resp
      dl completed resp
                           = 0x0B;
      dl ack cmd
                           = 0x0C:
      dl nak cmd
                           = 0x0D:
                                     { There is no difference between the short and long }
      dir short cmd
                            = 0x0E:
                                          dir formats- both send complete headers.
      dir long cmd
                            = 0x0F: {
      select cmd
                           = 0x10:
                            = 0x11;
      select_resp
                                     { Delete a file... not provided for in FTL0.
                            = 0x1E:
      del cmd
                                         For del cmd, the following packet fields apply:
                                         length lsb := 0x04; hl := 0x1E;
                                         info[0..3] := mail number: ulong.
                            = 0x1F; { Not provided for in FTL0.
      del_resp
                                         Thefollowing packet fields apply:
                                         lenght lsb := 0x01; hl := 0x1F;
                                         info[0] := error_code: uchar.
                                      { These are the FTLO login flags, assuming
                            = 0x00;
      no active list
                            = 0x08;
                                         PACSET File Headers are Not used
      active_list
                           Packet_Type;
var
      p:
      client callsign:
                           Callsign_Type;
       selection active:
                           uchar;
      err code:
                           uchar;
                                       mail number of file currently being uploaded.
      current ul mail:
                           ulong;
      current ul offset:
                                       Number of next byte to be uploaded in current file.}
                           ulong:
                                       mail number of file currently being downloaded.
      current dl mail:
                           ulong;
      current_dl_offset:
                                      { Num of next byte to be downloaded in current file.}
                           ulong;
      data length:
                           Pdat Len;
                                      { Raw select instruction.
                                                                                         }
       select:
                           Pdata:
```

```
{ States of PACKET TRANSFER BODY
         UL/DL UNINIT, WAIT MAILBOX, UL/DL CMD WAIT, UL DATA RX,
state
         UL ABORT, DL FILE DATA;
         ANY = [ UL/DL_UNINIT, WAIT_MAILBOX, UL/DL_CMD_WAIT,
stateset
                  UL_ABORT, DL_FILE_DATA];
                              { Function Declarations for
                              { PACKET TRANSFER BODY
function CURRENT COMMAND( packet: Packet Type): uint;
begin
   CURRENT COMMAND := C_BIT_AND( packet.hl, 0x1f);
end;
                              { Procedure declarations for
                              { PACKET_TRANSFER_BODY
procedure FORMAT LOGIN RESP(login_flag: uchar; var packet: Packet_Type);
   var
         login_time: ulong;
begin
   packet.length lsb := 0x05;
                              { 5 byte information field
                              { login_resp Packet Type
   packet.hl := login_resp;
   packet.info[0..3] := GET TIME();
   packet.info[4] := login_flag;
end;
procedure FORMAT_UPLOAD_GO_RESP( file_no, offset: ulong; var packet:
                                                          Packet Type);
begin
                              { 8 byte information field
                                                                            }
   packet.length lsb := 0x08;
   packet.hl := UPLOAD GO RESP;
   packet.info[0..4]:= file no;
   packet.info[5..7] := offset;
end:
procedure FORMAT NI RESP( tag: uchar; var packet: Packet_Type);
                              { No information
                                                                            }
   packet.length lsb := 0x00;
   packet.hl := tag;
end:
```

```
precedure FORMAT UL ERROR RESP( error: uchar; var packet: Packet_Type);
   packet.length lsb := 0x01;
   packet.hl := ul error resp;
   pecket.info[0] := error;
end:
procedure FORMAT_UL_NAK_RESP( error: uchar; var packet: Packet_Type);
   packet.length_lsb := 0x01;
   packet.hl := ul_nak resp;
   pecket.info[0] := error;
end:
procedure FORMAT SELECT RESP( num: uint; var packet: Packet Type);
begin
   packet.length lsb := 0x02;
   packet.hl := select resp;
   packet.info[0..1] := num;
end;
procedure FORMAT DL ERROR RESP( error: uchar; var packet: Packet Type);
   packet.length lsb := 0x01;
   packet.hl := dl error resp;
   packet.info[0] := error;
end:
procedure FORMAT DEL RESP( error: uchar; var packet: Packet Type);
begin
   packet.length lsb := 0x01;
   packet.hl := del resp;
   packet.info[0] := error;
end:
```

```
procedure FORMAT_DATA( len: Pdat_Len; dat: pdata; var packet: Packet_Type);
                      uchar:
         high byte:
                      uint:
         msb:
begin
   packet.length lsb := GET LSB(len);
   high byte := GET MSB( len);
   msb := I_BIT_SHIFT( left, 5, INT( high_byte));
   packet.hl := GET_LSB(msb);
   packet.info[0..len-1] := dat[0..len-1];
end:
                                { PACKET_TRANSFER_BODY
initialize
to UL/DL_UNINIT
begin
end:
                                {Transition Part of PACKET_TRANSFER_BODY}
from ANY to UL/DL_UNINIT
when pc.disconnect
                                { Link has been terminated by client or satellite.
begin
end:
                                { No action required.
from UL/DL UNINIT to WAI. MAILBOX
when pc.connection
begin
   client callsign := callsign;
   output mc.active sl req(callsign);
end:
from WAIT_MAILBOX to UL/DL_CMD_WAIT
when mc.active_sl_resp
begin
   if sl then selection active := active list;
   else selection_active := no_active_list;
   FORMAT LOGIN_RESP( selection_active, p);
   output pc.response_packet( p);
end:
                                    { Default condition for unexpected packet or
from UL/DL CMD WAIT to same
when others
                                       format.
begin
   FORMAT_UL_ERROR_RESP( er_ill_formed_cmd, p);
   output pc.response_packet( p);
end;
```

```
IN ULADL_CMD_WAIT to WAIT MAILBOX
   un pc.command packet
provided CURRENT_COMMAND(command) = upload cmd
   cutput mc.mail_num_req( client_callsign, command[2..5], command[6..9]);
from WAIT MAILBOX to UL/DL CMD WAIT
when mc.mail_num_resp
provided error code < > no error
begin
   FORMAT UL ERROR RESP(error code, p);
   output pc.response_packet(p);
end:
from WAIT MAILBOX to UL_DATA RX
when mc.mail num resp
provided error code = no error
begin
   current ul mail := mail number;
   current ul offset := offset;
   FORMAT UPLOAD GO RESP( mail number, offset, p);
   output pc.response_packet(p);
end:
from UL DATA RX to UL/DL_UNINIT
when pc.disconnect
                               { data link terminated by client or satellite.
begin
   output mc.mail_close_req( current ul mail, current ul offset, false);
from UL_DATA RX to UL/DL CMD WAIT
when others
                               { Default condition for unexpected packet or
begin
                               { format.
   output mc.mail_close_req( current_ul_mail, current_ul_offset, false);
   FORMAT_UL_ERROR_RESP( er_ill_formed_cmd, p);
   output pc.response_packet( p);
end:
```

```
from UL DATA RX to WAIT MAILBOX
when fc.command packet
provided CURRENT COMMAND(command) = data end
   output mc.mail_close_req( current_ul_mail, current_ul_offset, true);
end:
from WAIT_MAILBOX to UL/DL_CMD_WAIT
when mc.mail close_resp
begin
   if error code = no error then FORMAT NI RESP(ul_ack_resp, p);
   else FORMAT_UL_NAK_RESP( error_code, p);
   output pc.response_packet( p);
end:
from UL DATA RX to WAIT MAILBOX
when pc.command packet
provided CURRENT COMMAND(command) = data
   data_length := datal;
   output mc.mail recv( current_ul_mail, current_ul_offset, lata_length, command.info);
end:
from WAIT_MAILBOX to UL_DATA_RX
when mc.mail recv resp
provided error code = no error
begin
   current_ul_offset := current_ul_offset + data_length;
end:
from WAIT MAILBOX to UL ABORT
when mc.mail recv resp
provided error code < > no error
begin
   FORMAT UL NAK RESP( error code, p);
   output pc.response_packet(p);
end:
from UL_ABORT to UL/DL_CMD_WAIT
                               { Default condition for unexpected packet or
when others
begin
                                 format.
   FORMAT_UL_ERROR_RESP( er_ill_formed_cmd, p);
   output pc.response_packet( p);
end;
```

```
UL ABOUT to UL/DL_CMD_WAIT
     ind CURRENT_COMMAND(command) = data end
                              { No action required
                                                                          }
from UL_ABORT to some
when pc.command packet
provided CURRENT COMMAND(command) = data
bezin
                              { No action required
                                                                          }
end:
from UL/DL_CMD_WAIT to WAIT_MAILBOX
when pc.command packet
provided CURRENT_COMMAND( command) = del cmd
begin
   output mc.mail_del_req( client_callsign, command.info[0..3]);
end:
from WAIT_MAILBOX to UL/DL_CMD_WAIT
when mc.mail_del_resp
begin
  FORMAT_DEL_RESP( error_code, p);
   output pc.response packet(p);
end:
from UL/DL_CMD_WAIT to WAIT_MAILBOX
when pc.command packet
provided CURRENT COMMAND(command) = select cmd
begin
   select := command.info[0..datal-1];
   output mc.mselect req( client callsign, select);
end;
```

```
from WAIT MAILBOX to UL/DL CMD WAIT
when mc.mselect resp
   if error code = no error then begin
      selection active := active list;
      FORMAT SELECT RESP( num sel, p);
   else begin
      selection active := no active list;
      FORMAT_DL_ERROR_RESP( error code, p);
   output pc.response packet(p);
from UL/DL_CMD WAIT to WAIT MAILBOX
when pc.command packet
provided (CURRENT COMMAND(command) = dir short cmd
         or CURRENT COMMAND( command) = dir long cmd)
         and (
         (command.info[0..3] < > 0x00000000 and command.info[0..3] < >
                                                             OxFFFFFFFF)
         or selection active = active list)
begin
  if command.info[0..3] = 0x00000000 or command.info[0..3] = 0xFFFFFFF then
      output mc.dir req(client callsign, 0x00000000);
      output mc.dir_req( client callsign, command.info[0..3]);
end:
from UL/DL CMD WAIT to same
when pc.command packet
provided (CURRENT_COMMAND(command) = dir short cmd
         or CURRENT_COMMAND( command) = dir long_cmd)
         and not (
         (command.info[0..3] < > 0x00000000 and command.info[0..3] < >
                                                             0xFFFFFFFF)
         or selection active = active list)
   FORMAT DL_ERROR_RESP( er_selection_empty, p);
   output pc.response packet(p);
end:
```

```
UE/DL_CMD_WAIT to some
     BC-0000
     ded CURRENT_COMMAND(command) = dl_nak cmd
                                  I no action required
                                                                             }
from WAIT MAILBOX to UL/DL_CMD_WAIT
when mc.directory
begin
                               { Each File Header must = < 200 bytes
   If no al then selection active := no active list:
   if len < > 0 then begin
      FORMAT DATA(len, dir, p) { Assumes 10 file headers/DataPacket
      output pc.response packet(p);
      FORMAT NI RESP( data end, p);
      output pc.response_packet(p);
   end:
   else begin
      FORMAT DL ERROR RESP(dir[0], p)
      output pc.response packet(p);
   end:
end:
from UL/DL CMD WAIT to same
when pc.command
provided CURRENT COMMAND(command) = download cmd and (
         (command.info[0..3] = 0x00000000 or command.info[0..3] = 0xFFFFFFFF)
         and selection active = no active list)
begin
   FORMAT DL ERROR_RESP( er_selection_empty, p);
   output pc.response packet(p);
end:
from UL/DL CMD WAIT to DL FILE DATA
when pc.command
provided CURRENT_COMMAND( command) = download_cmd and (
         (command.info[0..3] < > 0x00000000 and command.info[0..3] < >
                                                              0xFFFFFFFF)
         or selection active = active list)
begin
   current dl offset := command.info[4..7]:
   current dl mail := command.info[0..3]:
   if current dl_mail := 0xFFFFFFFF then current dl_mail := 0x00000000;
   output mc.mail req( client callsign, current dl mail, current dl offset):
end:
```

```
from DL_FILE_DATA to UL/DL_CMD_WAIT
when mc, mail resp
provided length = 0 and mail[0] < > 0 { Error flag from mailbox
heela
   If no al then selection active := no active list:
   FORMAT DL ERROR RESP( mail[0], p);
   cutput pc.response_packet( p);
end:
from DL FILE DATA to same
when mc.mail resp
provided length < > 0 or (mail[0] = 0 and length = 0)
begin
   if no al then selection_active := no active list;
   if length = 0 then begin
      FORMAT NI RESP( data end, p);
      output pc.response packet(p);
   end:
   else begin
      FORMAT DATA( length, mail, p);
      output pc.response_packet(p);
      current dl offset := current dl offset + length;
      if current dl mail = 0x00000000 then current dl mail := mail number:
      output mc.mail req( client callsign, current dl mail, current dl offset);
   end:
end:
from DL_FILE DATA to UL/DL CMD WAIT
when pc.command packet
provided CURRENT COMMAND(command) = dl ack cmd
begin
   FORMAT_NI_PACKET( dl_completed_resp, p);
   output pc.response packet(p);
   output mc.dl ack( client callsign, current dl mail);
end:
from DL FILE DATA to UL/DL CMD WAIT
when pc.command packet
provided CURRENT COMMAND(command) = DL NAK DMD
begin
   FORMAT_NI_PACKET( dl_aborted_resp, p);
   output pc.response_packet( p);
end:
```

## body MAILBOX\_CONTROL\_BODY for MAILBOX\_CONTROL\_TYPE;

```
const
                                   { End Of File marker used by operating system.
   eof
   null
                                    The null pointer. A pointer which is null points
                                      to nothing, and marks the end of a linked list.
   empty_string
                                   { An empty string as defined by operating system.
   mailbox full
                         = any uint; { Parameter for FORMAT EVENT REPORT
   grab
                         = 0x01; { A parameter needed by 'qax_claim'.
                         = 0x01:
   mailbox ssid
   mail_flag
                         = 0xbb55:
   min file length
                         = 0x000000029:
                                          { Min of 41 bytes in the initialized mail file. }
                         = 0x03E7;
                                          { Higest mail name extension = 999 dec.
   max ext
                                          { Default mail life, in seconds from upload.
   default stay time
                         = any ulong:
                                          { Number of different file types allowed.
   numtypes
                         = any uchar;
                                             { Number of different file compression
   numcomps
                        = any uchar;
                                                 methods allowed.
                                    Item numbers for fields in the PANSAT file
                                      header, for use in 'select' statements.
   fl
                         = 0x00:
                                    flag
                         = 0x02:
                                   { mail number
   mn
                         = 0x06:
                                   { length
   ml
   ft
                         = 0x0A;
                                   { file type
                         = 0x0B:
                                   { compression type
   Ct
                                   { body offset
   bo
                         = 0x0C;
                         = 0x0E:
                                   { download count
   dc
                         = 0x0F:
                                   { source
   SC
                                   { priority
   pr
                         = 0x15:
                         = 0x16;
                                   { upload time
   ut
                         = 0x1A:
                                  { expire time
   et
                         = 0x1E;
                                   { mail name
   na
   ex
                         = 0x26:
                                   { mail extension
   nd
                         = 0x29:
                                   { num destinations
   ds
                         = 0x2A:
                                    destination callsigns or paths
   ti
                         = 0x54:
                                   { title
                         = 0x74;
                                   { keywords
   kw
                                   {Relational operators in 'select struct'
   pan sel
                         = 0xFF:
                                    equal to an unsigned 1,2 or 4 byte integer
   equal int
                         = 0x00:
                         = 0x03;
                                   { equal to a string
   equal str
   great int
                         = 0x10:
                                   { greater than an unsigned 1, 2 or 4 byte integer
                                   { less than an unsigned 1, 2 or 4 byte integer
   less int
                         = 0x20;
                                   { not equal to an unsigned 1, 2 or 4 byte integer
   not equ int
                         = 0x30:
   not_equ_str
                         = 0x33;
                                   { not equal to a string
```

```
= 0x40; { greater than or equal to an unsigned integer
   gr_equ_int
   ie_equ_int
                        = 0x50; { less than or equal to an unsigned integer
                        = 0x80; { logical 'and'
   l and
   l_or
                        = 0xE0; { logical 'or'
                                  { types for MAILBOX_CONTROL_BODY.
type
   File Ext
                           = 0..max ext;
   Ext Type
                           = array[3] of uchar;
   Mail_Array
                           = array[max_mail] of ulong; { Array of mail_numbers }
   Select_List
                           = record
      num sel: Num Mail;
                Mail_Array;
      sel:
   end;
   Source Record
                           = record
      source_num:
                           uint:
      call:
                           Callsign Type;
      selected:
                           Select List;
      next mail:
                           Num Mail;
      next dir:
                           Num Mail;
      next ext:
                           File Ext;
                           uchar;
      num act:
                           ^Source_Record; { Pointer to a source record.
      next num:
      next call:
                           *Source Record;
   end;
   Letter Array
                           = array[26] of \(^Source_Record;\)
   File_Desig_Array
                           = array[numtypes] of uchar;
   Compression Array
                           = array[numcomps] of uchar;
   S Name
                           = array[8] of uchar;
   File_Order
                           = (date, name);
var
   done:
                       boolean:
   filetype:
                       File Desig Array;
   comptype:
                       Compression_Array;
   next_source_num:
                       uint:
   nsel:
                       Num_Mail;
   file name:
                       Name_Type;
   rfile, tfile:
                       ^File_Type;
   first let:
                       Letter_Array;
   mail head:
                       *Source Record;
   temp1, temp2:
                       ^Source_Record;
                       ulong;
   mail num:
                       Ext_Type;
   ext:
```

```
time:
                        ulong;
   link:
                        Link Type;
   file length:
                        ulong:
   file offset:
                        ulong:
   body_offset:
                        uint;
   file error:
                        uchar;
   selectl:
                        Select List;
                        uchar;
   num dest:
                        Callsign Type;
   CS:
   i:
                                   { loop counter}
                        uchar:
   j, k:
   dir dat:
                        Pdata:
                        Event_Report_Type;
   report:
                        File Order;
   order:
                        S Name;
   sname:
                        File_Type;
   d file:
state WAIT;
function STRING COMPARE( str1, str2: ^Byte Array):
                                                           boolean;
                                   { Compares 'str1' to 'str2' and returns true if they
primitive;
                                   { are the same, otherwise returns false.
function STRING FIND( str1, str2: 'Byte Array): boolean;
                                   { Looks to see if 'strl' is contained anywhere within }
primitive;
                                      'str2'. Returns true if it is, and false if its not.
function GET_LENGTH( file name: Name Type): ulong;
primitive;
                                   { Returns the length of the stored file 'file name'
                                                                                       }
function GET_LSI( number: ulong): uint;
primitive;
                                   { Takes a 4 byte number and returns the least
                                   { significant 2 bytes.
function GET_MSI( number: ulong): uint;
primitive:
                                   { Takes a 4 byte number and returns the most
                                      significant 2 bytes.
function MEM_SPACE(): ulong;
primitive:
                                      { Returns the number of bytes of available
                                          space in mail box memory.
```

runction SIZE_OF( type_indicator)	: uint;	
primitive;	Takes as an argument any type and returns the number of bytes needed to store a variable of that type.	}
function ALLOCATE( size: uint ): primitive;	pointer_type; { Allocates a block of dynamic memory. The number of bytes in the block is indicated by	<b>}</b>
• •	'size'. The function returns a pointer to the newly allocated block.	}
function EXISTS( file_name: Name	_Type): boolean;	
primitive;	{ Takes a complete DOS file name as an argument { and returns	} }
·	{ true if an active file by that name currently exists { in the mass storage memory, otherwise returns	}
	false.	}
function GET_FIRST_FILE( order:	File_Order; fn: S_Name): Name_Type;	
primitive;	{ Returns the name of the first active (not deleted)	}
	{ mail file in the mass storage memory. "First" is { defined as the oldest file ( the one with the	} }
	{ earliest creation date) if 'order' is date. If 'order	}
	{ is name then 'fn' is a DOS file name minus the extension, and the file name returned is the one	} }
	with the "first" alpha-numeric extension	}
	{ associated with the 'fn' given. If no file matches the critria given, then empty_string is returned.	} }
function GET_NEXT_FILE( order:	File_Order; prev_file: Name_Type ): Name_Type;	
primitive;	{ Starting at 'prev_file', searches the mail area of	}
	{ mass storage memory for the next active file. If	}
	{ 'order' is <i>date</i> , the next file is the one with the	}
	{ next later creation date. If 'order' is name then { the next file is the one with the same leading 8	} }
	{ characters and the next higher alpha-numeric	} }
	extension. This function returns the complete file	}
	ame, if found and returns empty_string if no file	-
	{ matching the criteria exists.	}

function OPEN_FILE( file_name: ) primitive;	Name_Type): 'File_Type;  { Opens the file designated by 'file_name' for }  { reading or writing. Returns a pointer to the }  beginning of the file. If the file does not already}  exist, it will be created, and will be empty except}  for an eof mark.
function READ_FILE( qty, size: uprimitive;	int; var file_ptr: ^File_Type): Byte_Array; { Reads blocks of bytes from memory, starting at the} { location indicated by 'file_ptr'. The number of } { blocks is determined by 'qty' and the number of } { bytes in each block is determined by 'size'. The } { bytes are placed in the (pre_allocated) variable or} { buffer space designated on the left side of an } { assignment statement of which this function call } { is the right side. After the read, 'file_ptr' will } { point to the byte following the last byte read. }
<pre>procedure WRITE_FILE( v: Byte_ primitive;</pre>	Array; num_bytes: uint; var file_ptr: ^File_Type);  { Writes the number of bytes indicated by  { 'num_bytes' to memory starting at the location }  { indicated by 'file_ptr'. The bytes are copied }  { beginning from the first byte of 'v'. 'v' can be a}  { variable, buffer name or file pointer. After the }  { write, 'file_ptr' will point to the byte following }  { the last byte written. If there is already data in }  { the file at the position indicated by 'file_ptr', that}  { data will be overwritten. If writing to the end of }  a file, the eof marker will be moved to indicate }  the new end of the file.
<pre>procedure FILE_SEEK( num_bytes primitive;</pre>	s: int; var file_ptr: ^File_Type); { Moves the file pointer 'file_ptr' the number of } { bytes designated by 'num_bytes', without reading}
procedure FILE_SEEK_SET( numprimitive;	_bytes: int; var file_ptr: ^File_Type); { Same as 'FILE_SEEK', except that 'file_ptr' is firs} { moved to the beginning of the file, and then } { advanced the number of bytes indicated by } { 'num_bytes'. }
<pre>procedure DELETE_FILE( file_na primitive;</pre>	me: Name_Type); { Deletes the file designated by 'file_name' }

precedure CLOSE_FILE( primitive;	file_name: Name_Type); { Closes the file designated by 'file_name'	
procedure FREE(var no		
primitive;	{ Deallocates a dynamic memory node, and makes	}
	{ the pointer null.	}

```
procedure DECREMENT_MSG( tcall: Callsign_Type; var m_head: ^Source Record:
                                                        letter: Letter Array);
         temp1, temp2, head, del_node: ^Source_Record;
                                       uchar:
         index:
begin
   del node := null:
   index := tcall[0] - 0x41;
   head := letter[index];
   if head < > null then begin
      if head -> call = tcall then begin
          head -> num act := head -> num act - 1:
          if head -> num act <= 0 and head-> selected.num sel <= 0 then begin
             del node := head:
             letter[index] := head -> next call;
          end:
      end:
      else begin
          temp2 := head;
          temp1 := head -> next call;
          while temp1 < > null do begin
             if temp1 -> call = tcall then begin
                temp1 -> num act := temp1 -> num act - 1;
                if temp1 -> num act \leq 0 and temp1-> selected.num sel \leq 0
                then begin
                    del node := temp1;
                    temp2 -> next call := temp1 -> next_call;
                end;
             end:
             else begin
                temp2 := temp1;
                temp1 := temp1 -> next call;
             end:
          end:
      end:
      if del node < > null then begin
          head := m head:
          if head = del node then begin
             m head := head -> next num;
             FREE( del node);
          end:
          else begin
             temp2 := head;
             temp1 := head -> next num;
             while temp1 < > null do begin
```

```
If temp! = del_node then begin
                   temp2 -> next_num := temp1 -> next_num;
                   FREE( del_node);
                end;
                else begin
                   temp2 := temp1;
                   temp1 := temp1 -> next_num;
                end;
             end;
      end:
   end;
end;
function GET_EXT( num_ext: uint): Ext_Type;
         digit: uint;
   var
begin
   digit := num ext/100;
   GET_EXT[0] := digit + 0x0030;
   num_ext := num_ext - (digit * 100);
   digit := num ext/10;
   GET_EXT[1] := digit + 0x0030;
   num_ext := num_ext - (digit * 10);
   GET_EXT[2] := num_ext + 0x0030;
end;
```

```
procedure INCREMENT MSG( tcall: Callsign Type; var next sn: uint;
   m head: Source Record; letter: Letter_Array; mail_num: ulong; ext: Ext_Type);
         VAL
         index:
                                 uchar;
bezin
  index := tcall[0] - 0x41;
  temp1 := letter[index];
   while temp1 < > null and temp1->call < > tcall do begin
      temp2 := temp1;
      temp1 := temp1->next call;
  end:
  if temp1 = null then begin
      new node := ALLOCATE( 1, SIZE OF(Source Record));
      new node-> source num := next sn;
      GET NEXT NUM( m head, next sn);
      new node->call := tcall;
      new node-> selected.num sel := 0x0000;
      new node->next_ext := 0x0002;
      new node-> num act := 0x01;
      new node->next num := null;
      new node->next call := null;
      temp2->next call := new node;
      mail num := new node->source num *0x00010000 + 0x0001;
      ext := "001":
      temp1 := m head;
      if temp1->source num > new node->source num then begin
         new_node->next_num := m_head;
         m head := new node;
      end:
      else begin
         while temp1->source_num < new_node->source_num and
         templ->next num < > null do begin
            temp2 := temp1;
            temp1 := temp1->next num;
         if temp1->next num = null then
            temp1->next_num := new_node;
         else begin
            new node->next num := temp1;
            temp2->next_num := new_node;
         end;
      end:
   end:
```

```
temp1->num_act := temp1->num_act + 1;
      mail num := temp1->source num * 0x00010000 + temp1->next ext;
      ext := GET_EXT( temp1->next_ext);
      If temp1->next ext < max ext then
          temp1->next_ext := temp1->next ext +1;
      else templ->next_ext := 0x0001;
   end:
end:
procedure COMPACT_MAIL( var mlist: ^Source_Array; llist: Letter Array);
begin
                                 { This function deletes mail files which are past their}
                                    expiration dates.
             this file:
                          Name_Type;
      ASL
                          ^File_Type;
             rfile:
             now:
                          ulong;
             expire:
                          ulong;
             call:
                          Callsign_Type;
   now := GET TIME();
   this file := GET_FIRST_FILE( date, empty_string);
   while this_file < > empty_string do begin
      rfile := OPEN FILE( this file);
      FILE SEEK( 26, rfile);
      expire := READ FILE( 1, 4, rfile);
      if expire \leq now then begin
          FILE_SEEK( 2, rfile);
          call := READ FILE( 1, 6, rfile);
          DELETE FILE( this file):
          DECREMENT_MSG( call, mlist, llist)
      end:
      else CLOSE FILE( this file);
      this file := GET_NEXT_FILE( date, this_file);
   end:
end:
```

```
precedure GET_NEXT_NUM( mlist: ^Source_Record, var next sn);
                               { Finds the next unused source number.
      temp: "Source_Record;
      done: boolean;
   done := false:
   while not done do begin
      if next an < 0xFFFF then
         next sn := next sn + 1;
      eine next sn := 0x0001;
      temp := mlist:
      while temp < > null and temp-> source_num < next_sn do
         temp := temp->next_num;
      if temp->source num < > next sn then
         done := true;
   end:
end:
function MAKE_FILE_NAME( source: Callsign_Type; ext: Ext_Type): Name_Type;
begin
  MAKE FILE NAME[0..1] := " ";
   MAKE FILE NAME[2..7] := source;
   MAKE FILE NAME[8] := '.';
   MAKE FILE NAME[9..11] := ext;
end:
```

```
CET_NAME (mist: Source_Record; mnum: ulong): Name_Type;
                  uint:
                  *Source Record:
         node:
                  Ext Type;
         SOURCE:
                  Callsign Type;
   s := GET MSI( mnum);
                               { source
   e := GBT LSI( mnum);
                               { extension
   ext := GET_EXT(e);
   node := mlist:
   while node < > mull and node-> source num < > s do
      node := node->next_num;
   if node := null then
      GET NAME := empty string;
   else begin
      source := node->call;
      GET NAME := MAKE FILE NAME( source, ext);
   end:
procedure INITIALIZE MAIL FILE( source: Callsign Type; mnum: ulong);
                                  ext: Ext Type; length: ulong);
         new_file: Name_Type;
                  ^File_Type;
begin
   new_file := MAKE_FILE_NAME( source, ext);
   f := OPEN FILE( new_file);
   WRITE_FILE( mail_flag, 2, f);
   WRITE, FILE( mnum, 4, f);
   WRITE FILE(length, 4, f);
   CLOSE_FILE( new_file);
                               { Of Procedure INITIALIZE MAIL FILE.
end:
```

```
precedure RE_INIT_FILE( file_name: Name_Type; mail_num: ulong);
         r file:
                    ^File_Type;
         f length: ulong:
   r_file := OPEN_FILE( file_name);
   FILE SEEK( fl. r file):
   f length := READ_FILE( 1, SIZE_OF( ulong), r_file);
   DELETE FILE( file name):
   INITIALIZE MAIL FILE (file name[2..7], mail num, file name[9..11], f length);
end:
function CRC CHECKS OUT( file name: Name Type; start: uint; stop: ulong;
                                                            crc: uint): boolean;
                                 { This algorithm assumes the crc is a simple check
                                    sum.
                          ulong;
         num bytes, i:
   var
                          ^File_Type;
         r_file:
                          uint;
          sum:
         next char:
                          uchar:
   if crc = 0x00 then CRC_CHECKS OUT := true;
   else begin
      sum := 0x00:
      r file := OPEN FILE( file name);
      FILE SEEK( start, r file);
      for i := 1 to (stop - start) do begin
         next char := READ FILE(1, 1, r file);
         sum := sum + next char;
      end:
      CLOSE FILE( file name);
      CRC CHECKS OUT := ( sum = crc);
   end;
end:
function CHANGE CASE( str: Byte Array; len: uchar):
                                                        Byte Array:
         i: uchar;
   VAL
                                 { Changes any ASCII upper case letters found within}
begin
                                    'str' to lower case. Returns the modified string. }
   for i := 0 to len-1 do
      if str[i] > 0x40 and str[i] < 0x5B then str[i] := str[i] + 0x20;
   CHANGE CASE := str:
end:
```

```
ADER CHECK file name: Name Type; mail_num: ulong;
      flietype: File Desig Array; comptype: Compression Array):
                       booleen:
      good, ok:
      r file:
                       'File Type;
      c, d:
                       uchar:
      i, body_offset:
                       uint:
                       ulong;
      lo:
                      Callsign_Type;
      call:
      ext:
                       Ext_Type;
good := true;
r file := OPEN FILE( file name);
i := READ_FILE(1, 2, r_file);
                               { Read flag
If i < > mail flag then good := false;
lo := READ FILE(1, 4, r file); { Read mail number
if lo < > mail_num then good := false;
lo := READ FILE(1, 4, r file); { Read length
if lo < > GET LENGTH( file name) then good := false;
c := READ FILE(1, 1, r file); { Read file type
ok := false:
for i := 0 to numtypes - 1 do { Check file type against all valid file types
   if c = filetype[i] then ok := true;
if not ok then good := false; { HEADER_CHECK, continued.
c := READ FILE(1, 1, r file); { Read compression type
ok := false:
for i := 0 to nuncomps - 1 do { Check compression_type against all valid
   if c = comptype[i] then ok := true; { compression_types
if not ok then good := false;
body_offset := READ_FILE(1, 2, r_file); { Read body_offset
c := READ_FILE(1, 1, r file); { Read download_count
call := READ_FILE(1, 6, r file); { Read source
call := CHANGE CASE( call, 6);
if call < > CHANGE CASE(this file[2..7], 6) then
                             { Check source vs file name
   good := fake:
c := READ_FILE(1, 1, r_file); { Read priority. What to do with it is undefined.}
lo := READ FILE( 1, 4, r file); { Read upload time
if lo > = GET_TIME() then good := false;
lo := READ FILE( 1, 4, r file); { Read expiration time
if lo < GET_TIME() then good := false;
i := READ_FILE(1, 2, r file); { Read leading spaces in file name.
if i < > this file[0..1] then good := false;
call := READ_FILE( 1, 6, r_file);
                                   { Read file name
if CHANGE CASE(call, 6) < > CHANGE CASE(this file[2..7], 6) (ibe.)
```

```
good := false:
                                 { Check vs known file name
   ext := READ_FILE(1, 3, r_file); { Read file extension
   # CHANGE_CASE(ext, 6) < > CHANGE_CASE(file_name[9..11], 6) then
                                 { Check vs known file ext
      good := false;
   c := READ_FILE(1, 1, r_file); { Read num_destinations
   If c > 0x09 then good := false;
   if good then begin
      if c > 0x07 then c := 0x07;
      for i := 1 to c do
                                 { Read past all destinations
          call := READ_FILE( 1, 6, r_file);
      body offset := body offset - min file length - c*6;
      c := READ_FILE(1, 1, r_file); { Read title_length
                                { Read title
      for i := 1 to c do
          d := READ FILE(1, 1, r_file);
      body_offset := body_offset - 1 - c;
      c := READ_FILE(1, 1, r_file); { Read keyword_length
                                 { Read keywords
      for i := 1 to c do
          d := READ_FILE(1, 1, r_file);
      body_offset := body_offset - 1 - c;
      body_offset := body_offset - 4; { Account for 2 checksum uints
      if body offset < > 0 then good := false;
   end:
   HEADER_CHECK := good;
                                                                                 }
                                 { of HEADER_CHECK( )
end;
function MSG TO( file name: Name Type; call: Callsign Type): boolean;
                       ^File Type;
   Var
          num dest, i: uchar;
                       Callsign Type;
          dest:
begin
   MSG TO := false;
   f := OPEN_FILE( file_name);
   FILE SEEK( nd, f);
   num dest := READ_FILE(1, 1, f);
   if num dest > 0 and num dest < 0x08 then do
      for i := 1 to num dest do begin
          dest := READ FILE( 1, SIZE OF(Callsign_Type), f);
          if dest := call then MSG_TO := true;
      end:
   CLOSE FILE( file name);
end:
```

```
DEFAULT SELECT( call: Calleign_Type; var m_head: ^Source_Record;
                          cs array: Letter Array; next sn: uint; nsel: Num_Mail);
                   Name Type;
   file name:
   f:
                   ^File_Type;
                   *Source Record;
   temp1, temp2:
                   *Source Record;
   new node:
                   uchar;
   num dest, i:
                             { download count
   dl count, j:
                   uchar:
   index:
                   uint:
                   Num Mail;
   num found:
                   Mail Array:
   s list:
                             { mail number, upload time
   m_num, ul_time: ulong;
num found := 0:
index := 0;
file name := GET FIRST FILE( date, empty string);
while file name < > empty string do begin
   f := OPEN FILE( file name)
   FILE SEEK( mn, f);
   m num := READ FILE(1, 4, f);
   FILE SEEK( dc - ml, f);
   dl_count := READ_FILE( 1, 1, f);
   FILE SEEK( w - sc, f);
   ul time := READ FILE( 1, 4, f);
   FILE SEEK( nd - er, f);
   num dest := READ FILE( 1, 1, f);
   CLOSE FILE( file name);
   if ul time > 0 then
                             { Only consider completely uploaded files.
      if num_dest = 0x00 or num_dest > 0x07 then begin { to 'ALL'
          num found := num found + 1;
          s list[index..index + 3] := m num;
         index := index + 4;
      end:
      else if dl count < num dest then
          if MSG TO( file name, call) then begin
             num found := num found + 1;
             s list[index..index + 3] := m_num;
             index := index + 4:
   file name := GET NEXT_FILE( date, file_name);
end:
i := call[0] - 0x41;
temp1 := cs_array[i];
```

```
while temp1 < > null and temp1->call < > call do begin
      temp2 := temp1;
      temp1 := temp1->next call:
   If temp1 = null then begin
      new node := ALLOCATE( 1, SIZE OF(Source Record));
      new node := source num := next sn:
      GET_NEXT_NUM( m_head, next_sn);
      new node->call := call;
      new_node-> selected.num_sel := num_found;
      new node->selected.sel := s list;
      new node->next mail := 0x0000;
      new node->next dir := 0x0000;
      new_node > next_ext := 0x0001;
      new node->num act := 0x00;
      new node->next num := null;
      new node->next call := null;
      templ := m head:
      if temp1->source num > new node->source num then begin
         new node->next num := m head;
         m head := new node;
      end:
      else begin
         while temp1->source num < new node->source num and
         temp1->next num < > null do begin
            temp2 := temp1;
            temp1 := temp1-> next num;
         end:
         if temp1 -> next num = null then
            temp1->next_num := new_node;
         else begin
            new node->next num := temp1;
            temp2->next num := new node;
         end:
      end:
   end:
   else begin
      temp1->selected.num sel := num found;
      temp1->selected.sel := s list;
      temp1->next mail := 0x0000;
      temp1->next dir := 0x0000;
   end:
   nsel := num found;
end:
```

```
PANSAT SELECTION( a struct: Pdata; call: Callsign Type;
ver m head: Source Record; cs array: Letter Array; next sn: uint; nsel: Num Mail;
error: ucher);
          file name:
                           Name Type;
                           'File_Type;
          f:
          s list:
                           Select List:
          temp1, temp2:
                           *Source Record:
          new node:
                           *Soruce Record:
          dest_call:
                           Callsign Type;
          list i:
                           uint:
          m_num:
                           ulong:
          abort, ok:
                           boolean:
          num s, i:
                           uchar;
                           uint:
          struct i:
          header item:
                           uchar:
                           boolean:
          first, second:
          relop, logop, j:
                           uchar;
          item_len, h_len: uchar:
          s int, hs int:
                           uchar;
          m_int, hm_int:
                           uint;
          l int, hl int:
                           ulong:
          compare item:
                           Byte_Array;
                           Byte Array;
          h string:
          num_dest, t_len: uchar;
          ul time:
                           ulong:
begin
   abort := false:
   num_s := s_struct[1];
   s list.num sel := 0;
   list i := 0:
   file name := GET FIRST FILE( date, empty string);
   while file_name < > empty_string and not about do begin
      f := OPEN FILE( file name);
      FILE SEEK( mn, f);
      m num := READ FILE(1, 4, f);
      FILE SEEK( ut - ml, f);
      ul time := READ_FILE( 1, 4, f);
      if ul time > 0 then begin
          ok := false:
          FILE SEEK( nd - et, f);
          num dest := READ_FILE( 1, 1, f);
          if num dest = 0x00 or num dest > 0x07 then ok := true;
          else
             for j := 1 to num dest do begin
```

```
dest call := READ_FILE(1, 6, f);
      if CHANGE CASE( dest call, 6) = CHANGE CASE( call, 6) then
          ok := true:
   end:
if ok then begin
   struct i := 2;
   first := true:
   second := false:
   i := 0
   while i < num s - 1 and not abort do begin
      relop := s struct[struct i];
      struct i := struct i + 1;
      h item := s struct[struct i]:
      struct i := struct i + 1;
      item_len := s_struct[struct i]:
      struct_i := struct_i + 1;
      if relop = equal str or relop = not equ str then begin
          compare item := s struct[struct i..(struct i + item len - 1)];
          compare item := CHANGE CASE( compare item, item len);
          struct i := struct_i + item_len;
          if h item < ds then begin
             FILE SEEK SET( h item, f);
             h string := READ FILE( item len, 1, f);
             h string := CHANGE CASE(h string, item len);
             second := STRING COMPARE( h string, compare item);
             if relop = not equ str then second := not second;
          end:
          else begin
             FILE SEEK SET( nd, f);
             num dest := READ FILE( 1, 1, f);
             if h item = ds then begin
                if num dest > 0 and num dest < 0x08 then
                    if item len = 6 then begin
                       for i := 1 to num dest do begin
                          h string := READ FILE( 1, 6, f);
                          h_string := CHANGE_CASE( h_string, 6);
                          if not second then
                              second := STRING COMPARE( h string,
                              compare item);
                       end;
                    end:
                    else second := false;
                else if num dest > 0x07 then begin
                    h string := READ FILE( 42, 1, 1);
```

```
h_string := CHANGE CASE( h_string, 42);
             second := STRING FIND( compare item, h string);
          else second := false;
          if relop = not equ str then second := not second;
      end:
      else
          if h item = ti then begin
             FILE SEEK( ds + num dest*6, f):
             t len := READ FILE(1, 1, 1);
             h string := READ FILE( t len, 1, f);
             h_string := CHANGE CASE( h string, t len);
             second := STRING FIND( compare item, h string);
             if relop = not_equ_str then second := not second;
          end:
          else
             if h item = kw then begin
                FILE SEEK( ds + num dest*6, f);
                t len := READ FILE(1, 1, f);
                FILE SEEK( t len, 1, f);
                t_{len} := FILE_READ(1, 1, f);
                h string := READ FILE( t len, 1, f):
                h string := CHANGE CASE( h_string, t_len);
                 second := STRING_FIND( compare_item, h_string);
                if relop = not equ str then second := not second;
             else abort := true;
   end:
end:
else
   if h item < na then begin
      FILE SEEK SET( h item, f);
      if item len = 1 then begin
          s int := s struct[struct i];
          struct i := struct i + 1;
          hs int := READ FILE( 1, 1, f);
          hl int := 0x00000000 + hs int;
          1 int := 0x00000000 + s int;
      end:
      else if item len = 2 then begin
          m_int := s_struct[struct_i..struct_i + 1];
          struct i := struct i + 2;
          hm int := READ FILE( 1, 2, f):
          hl int := 0x00000000 + hm int;
```

```
1_{int} := 0x00000000 + m_{int};
          end:
          else if item len = 4 then begin
             1_int := s_struct[struct_i..struct_i + 3];
              struct i := struct_i + 4;
              hl int := READ FILE(1, 4, f);
          end;
          else abort := true;
          if not abort then
              if relop = equal int then
                 second := (hl int = l int);
              else if relop = great int then
                 second := (hl_int > l_int);
              else if relop = less in then
                 second := ( hl_int < l_int);</pre>
              else if relop = not equ int then
                 second := (hl_int < > l_int);
              else if relop = gr equ int then
                 second := (hl_int > = l_int);
              else if relop = le equ int then
                 second := (hl_int <= l_int);
              else abort := true;
       end:
       else abort := true;
   if not abort then begin
      i := i + 1;
       logop := s struct[struct_i];
       if logop = l and then begin
          struct i := struct i + 1;
          first := (first and second);
       end;
       else if logop = l or then begin
          struct i := struct i + 1;
          first := (first or second);
       else first := (first and second);
   end:
end:
if first then begin
   s list.num sel := s list.num sel + 1;
   s list.sel[list_i..list_i + 3] := m_num;
   list i := list i + 4;
end:
```

end:

```
CLOSE FILE( file name):
  file_name := GET_NEXT_FILE( date, file_name);
if abort them
   error := er poorty formed_sel;
else begin
   error := no error;
   i := call[0] - 0x41;
   temp1 := cs_array[j];
   while temp1 < > null and temp1->call < > call do begin
      temp2 := temp1;
      temp1 := temp1->next_call;
   end:
   if temp1 = null then begin
      new_node := ALLOCATE( 1, SIZE_OF(Source_Record));
      new node := source num := next sn;
      GET NEXT NUM( m_head, next_sn);
      new node->call := call;
      new node->selected.num sel := s list.num sel;
      new node-> selected.sel := s list.sel;
      new node->next mail := 0x0000;
      new node->next dir := 0x0000;
      new node->next ext := 0x0001;
      new node->num act := 0x00;
      new_node->next_num := null;
      new node->next call := null;
      temp1 := m head;
      if temp1->source num > new_node->source_num then begin
          new node->next num := m_head;
          m head := new node;
      end;
       else begin
          while temp1->source_num < new_node->source_num and
          temp1->next num < > null do begin
             temp2 := temp1:
             temp1 := temp1-> next_num;
          if temp1 -> next num = null then
             temp1->next num := new node;
          else begin
             new node->next num := temp1;
             temp2->next num := new_node;
          end:
```

```
end;
      end;
      else begin
          temp1-> selected.num sel := s list.num sel;
          temp1->selected.sel := s list.sel;
          temp1->next mail := 0x0000;
          temp1->next dir := 0x0000;
      end:
      nsel := s_list.num_sel;
end:
procedure COPY_HEADER( file_name: Name_Type; var buffer: Byte_Array;
                                                             index: ulong);
                       ^File_Type;
   YEL
         r file:
         body_offset: uint;
begin
   r_file := OPEN_FILE( file_name);
   FILE SEEK( bo, r file);
   body_offset := READ_FILE(1, 2, r_file);
   FILE SEEK SET( 0, r file);
   buffer[index..index + body_offset - 1] := READ_FILE( body_offset, 1, r_file);
   CLOSE FILE( file_name);
   index := index + body_offset;
end:
```

```
to WAIT
                                   PACSAT File Types from H. Price
begin
                                   ascii
   filetype[0] := 0x00;
                                   RLI/MBL message body. Single message.
   filetype[1] := 0x01;
                                  { RLI/MBL import/export file. Multiple message.
   filetype[2] := 0x02;
                                  { UoSAT Whole Orbit Data.
   filetype[3] := 0x03;
                                   Microsat Whole Orbit Data.
   filetype[4] := 0x04;
                                  { UoSAT CPE Data.
   filetype[5] := 0x05;
                                  { MS/PC-DOS .exe file.
   filetype[6] := 0x06;
                                   MS/PC-DOS .com file.
   filetype[7] := 0x07;
                                  { Keplerian elements NASA 2-line format.
   filetype[8] := 0x08;
                                  { Keplerian elements "AMSAT" format.
   filetype[9] := 0x09;
                                   Simple ASCII text file, but compressed.
   filetype[10] := 0x0A;
                                   PANSAT File Types.
                                  { PANSAT short telemetry file.
   filetype[11] := 0xA0;
                                   PANSAT long telemetry file.
   filetype[12] := 0xA1;
                                   PANSAT bax telemetry file.
   filetype[13] := 0xA2;
                                   User defined type. User must know. 0xFF
   filetype[14] := 0xFE;
                                     'ESCAPE' not implemented in PANSAT file
                                      headers.
   for i := 15 to numtypes - 1 do
                                  { Extra space for types defined later.
       filetype[i] := 0x00;
                                  { PACSAT file compression types - H. Price.
                                   No compression
    comptype[0] := 0x00;
                                   Body compressed using PKARC.
    comptype[1] := 0x01;
                                   Body compressed using PKZIP.
    comptype[2] := 0x02;
                                  Other, user-known compression type.
    comptype[3] := 0xFE;
    for i := 4 to numcomps - 1 do
                                  { Extra space for compression types defined later.
       comptype[i] := 0x00;
    next source num := 0x0001;
    for i := 0 to 25 do
       first_let[i] := null;
    mail head := null;
end:
```

```
trans
from WAIT to same
when mc[link].mail num rea
bezin
   if mail number[link] = 0x00000000 then begin
      If length[link] > MEM SPACE() then
          COMPACT MAIL( mail head, first let):
      if length[link] <= MEM SPACE() then begin
         INCREMENT_MSG( client_call[link], next_source_num, mail_head, first_let,
          mail_num, ext);
         INITIALIZE MAIL FILE( client call, mail num, ext, length(link));
          output mc[link].mail num resp( mail num, 0x00000000, no error);
      end:
      else begin
          output mc[link].mail num resp( 0x00000000, 0x00000000, er no room);
          time := GET TIME();
         report := FORMAT_EVENT_REPORT( mailbox full, time);
          output el.event report( report);
          output cc.full mailbox;
      end:
   end:
   else begin
      file name := GET NAME( mail head, mail number[link]);
      if file name = empty string then
          output mc[link].mail num resp( mail number[link], 0x00000000.
          er no such file number);
      else begin
         rfile := OPEN FILE( file name);
         FILE SEEK( 6, read file);
         file length := READ FILE( 1, SIZEOF(ulong), rfile);
         if length[link] < > file length then
             output mc[link].mail_num_resp( mail_number, 0x00000000,
              er bad continue);
         else begin
             file of GET LENGTH (file name) + 1;
             if file offset > = length[link] then
                output mc[link].mail num resp( mail number[link], 0x00000000,
                 er file complete);
             else begin
                if file offset < 42 then file offset := 0;
                output mc[link].mail num resp( mail number[link], file offset,
                 no error);
             end;
         end:
```

```
from WAIT to same
when mc[link], mail close req
  file error := no error;
  header crc := 0x0000;
   file name := GET NAME( mail_head, mail_number[link]);
   If file name = empty string then file error := er_no_room;
   else begin
      rfile := OPEN FILE( file_name);
      FILE SEEK( mn, rfile);
      mail num := READ FILE( 1, 4, rfile);
      file length := READ FILE( 1, 4, rfile);
      if GET LENGTH( file name) >= file length then begin
         if mail num = 0x00000000 then begin
             FILE SEEK SET( mn, rfile);
             WRITE FILE( mail number[link], 4, rfile);
             for i := 0 to 3 do
                header crc := header crc + mail number[link][i];
          end;
          time := GET TIME():
          FILE SEEK SET( ut, rfile);
          WRITE_FILE( time, 4, rfile); { Write upload time.
         for i := 0 to 3 do
             header crc := header crc + time[i];
          time := time + default_stay_time;
          WRITE FILE( time, 4, rfile); { Write expire_time.
          for i := 0 to 3 do
             header crc := header crc + time[i];
          WRITE_FILE( file_name[0..7], 8, rfile);
          for i := 0 to 7 do
```

}

}

header crc := header crc + file name[i];

header crc := header crc + file name[i];

j := READ FILE(1, 1, rfile); { Read title length.

WRITE FILE( file name[9..11], 3, rfile);

num\_dest := READ\_FILE( 1, 1, rfile);
if num\_dest > 0 and num\_dest < 8 then
FILE\_SEEK( num\_dest\*6, rfile);</pre>

for i := 9 to 11 do

else if num\_dest > 7 then FILE SEEK( 42, rfile);

```
FILE SEEK( j, rfile);
         j := READ_FILE( 1, 1, rfile); { Read keyword_lenght.
         FILE SEEK( i. rfile);
         i := READ FILE( 1, 2, rfile); { Read header check sum
         If i > 0 then begin
             header crc := header crc + i;
             FILE SEEK( -2, rfile);
             WRITE FILE( header crc, 2, rfile);
         end:
      end:
      CLOSE FILE( file name);
      if req resp[link] then begin
          if not HEADER_CHECK( file_name, mail_number[link], filetype, comptype)
          then file error := er bad header;
          ekse begin
             rfile := OPEN FILE( file name);
             FILE SEEK( bo, rfile);
             body offset := READ FILE( 1, 2, rfile);
             FILE SEEK SET(body offset - 4, rfile);
             header crc := READ FILE(1, 2, rfile);
             body crc := READ FILE(1, 2, rfile);
             CLOSE FILE( file name);
             if not CRC CHECKS OUT( file name, 0, body offset, header crc) then
                file error := er header check;
             else if not CRC CHECKS OUT( file name, body offset, file length,
             body crc)
             then file_error := er_body_check;
          end;
          output mc[link].mail close resp( file error);
          if file error < > no error then
             RE INIT FILE( file name, mail number[link]);
      end:
end:
```

```
from WAIT to some
   en mcflinki.mail recv
   file name := GET NAME( mail head, mail number[link]);
   If length[link] > MEM SPACE() then
      COMPACT MAIL( mail head, first let);
   if length[link] <= MEM SPACE() and file name < > empty string then begin
      rfile := OPEN FILE( file name);
      FILE SEEK( offset[link], rfile);
      WRITE FILE( mail[link], length[link], rfile);
      CLOSE FILE( file name);
      output mc[link].mail recv resp( no error);
      if file name = empty string then
          output mc[link].mail recv resp( er bad continue);
      else begin
          output mc[link].mail_recv_resp( er_no_room);
          time := GET TIME();
          report := FORMAT EVENT REPORT( mailbox full, time);
          output el.event_report( report);
          output cc.full mailbox;
      end:
end:
from WAIT to same
when mc[link].mselect req
begin
   file error := no error;
   if select struct[link][0] = pan sel then
      PANSAT SELECTION( select struct[link], client call[link], mail head, first let.
      next source num, nsel, file error);
   else DEFAULT SELECT( client_call[link], mail_head, first let, next source num,
   nsel):
   output mc[link].mselect_resp( nsel, file error);
end:
```

```
from WAIT to same
when mc[link].mail_del_req
begin
   file error: = er permission denied:
   file name = GET_NAME( mail head, mail number[link]);
   if file_name = empty_string then
      file_error := er_no_such_file_number;
   else if client call[link] = nps call then
      file_error := no_error;
   else begin
      if file name[2..7] < > client call[link] then begin { A client may only delete a}
          rfile := OPEN FILE( file name); { file which he has uploaded or which is}
          FILE SEEK( nd. rfile): {
                                                   addressed to him.
          num dest := READ FILE( 1, 1, rfile);
          if num dest = 0x01 then begin
             cs := READ FILE( 1, SIZE OF(Callsign Type), rfile);
             if cs = client_call[link] then file_error := no error;
          CLOSE FILE( file name);
      end;
      else file_error := no_error;
   if file error = no error then begin
      DELETE FILE( this file);
      DECREMENT_MSG(file_name[2..7], mail head, first let);
   end:
   output mc[link].mail_del_resp( file_error);
end:
```

```
from WAIT to some
when mc[link].dir req
   If mail_number[link] < > 0x00000000 then begin
      file name := GET NAME( mail head, mail number[link]);
      If file name = empty string then
          output mc[link].directory(0, er no such file number, false);
      else begin
          body offset := 0;
          COPY HEADER( file name, dir dat, body offset);
          output mc[link].directory( body offset, dir dat, false);
      end:
   end:
   else begin
      i := client call[link][0] - 0x41:
      temp1 := first let[i];
      while temp1 < > null and temp->call < > client call[link] do begin
          temp2 := temp1:
                              temp1 := temp1->next call;
      end:
      if temp1 := null or temp1-> selected.num sel = 0x0000 then
          output mc[link].directory(0, er selection empty, true);
      else begin
                    file length := 0:
          i := 0:
          while temp1->next dir \leq temp1->selected.num sel and j \leq 10 do begin
             mail num := temp1-> selected.sel[temp1-> next dir]:
             file name := GET NAME( mail head, mail num);
             if file name < > empty string then begin
                 COPY HEADER( file name, dir dat, file length);
                i := i + 1;
             end:
             temp1 - next dir := temp1 - next dir + 1:
          if temp1->next dir > temp1->selected.num sel then begin
             done := true;
             if temp->next mail > temp1->selected.num sel then
                temp1-> selected.num sel := 0x0000;
          end:
          else done := false:
          if file length = 0 then
             output mc[link].directory(0, er selection empty, true);
          else output mc[link].directory( file length, dir dat, done);
       end;
   end:
end:
```

```
from WAIT to same
when mc[link].mail_req
bezin
   }
      file name := GET NAME( mail head, mail number[link]);
      if file name = empty string then
         output mc[link].mail resp( er_no_such_file_number, mail_number[link], 0,
         false):
   end:
                                { "Next" file in selection list requested.
                                                                                }
   else begin
      j := client call[link][0] - 0x41;
      temp1 := first_let[j];
      while temp1 < > null and temp->call < > client_call[link] do
         temp1 := temp1->next call;
      if (temp1 = null or temp1-> selected.num sel = 0x0000 or
      temp1->next mail > temp1->selected.num_sel) then begin
         if temp1 = null or temp1-> selected.num sel = 0x00 then
             done := true:
         else done := false:
         output mc[link].mail resp( er selection empty, 0, 0, done);
      end:
                                { Active selection list found
                                                                                }
      else begin
         mail num := temp1-> selected.sel[temp1-> next mail]
         file_name := GET_NAME( mail_head, mail_num);
          if file name := empty string then begin
             templ->next mail := templ->next mail + 1;
             if temp1->next mail > temp1->selected.num sel then
                if temp1->next dir > temp1->selected.num sel then begin
                   temp1-> selected.num sel := 0x0000;
                   output mc[link].mail resp( er selection empty, 0, 0, true);
                end:
                else output mc[link].mail resp( er selection empty, 0, 0, false);
                output mc[link].mail_resp( er_no_such_file number, mail num, 0,
                false);
          end;
          else begin
                                { Good file number found in select list.
                                                                                }
             rfile := OPEN FILE( file name);
             FILE SEEK( ml, rfile);
             file length := READ FILE( 1, 4, rfile);
             if file length <= offset[link] then begin
                output mc[link].mail resp(0, mail num, 0, false);
                CLOSE FILE( file name);
             end:
```

```
ele: begin
                 FILE SEEK SET( offset[link], rfile);
                 if file_length - offset[link] > max_pdat then begin
                    dir_dat[0..max pdat - 1] := READ FILE( max pdat, 1, rfile);
                     output mc[link].mail_resp( dir_dat, mail_num, max_pdat, false);
                 end;
                 else begin
                    file_offset := file_length - offset[link];
                     dir_dat[0.. file_offset - 1] := READ FILE( file_offset, 1, rfile);
                     output mc[link].mail resp(dir dat, mail num, file offset, false);
                 CLOSE_FILE( file_name);
              end;
          end;
      end;
   end;
end;
```

```
from WAIT to same
when mc[link].dl ack
begin
   i := client call[link][0] - 0x41;
   temp1 := first let[j];
   while temp1 < > null and temp->call < > client call[link] do
      temp1 := temp1->next call;
   if (temp1 < > null and temp1-> selected.num sel < > 0x0000 and
   templ->next mail <= templ->selected.num sel) then
      if temp1-> selected.selftemp1-> next mail = mail number[link] then begin
          temp1->next mail := temp1->next mail + 1;
          if temp1->next mail > temp1->selected.num sel then
             if temp1->next dir > temp1->selected.num sel then
                temp1-> selected.num sel := 0x0000;
      end;
   file name := GET NAME( mail head, mail number[link]);
   if file_name < > empty_string then begin
      rfile := OPEN FILE( file name);
      FILE SEEK SET( dc, rfile);
      k := READ FILE(1, 1, rfile);
      done := false:
      FILE SEEK SET( nd, rfile);
      num dest := READ FILE( 1, 1, rfile);
      CLOSE FILE( file name):
      if num_dest = 0x00 or num_dest > 0x07 then done := true;
      else if MSG TO( file name, client call[link]) then done := true;
      if done and k < 255 then begin
          rfile := OPEN_FILE( file_name);
         k := k + 1;
         FILE_SEEK( dc, rfile);
          WRITE FILE( k, 1, rfile);
         FILE SEEK( ds - sc, rfile);
         if num dest < 0x08 then FILE SEEK( num_dest*6, rfile);
          else FILE_SEEK( 42, rfile);
         j := READ FILE(1, 1, rfile); { Read title length.
                                                                                  }
         FILE SEEK( j, rfile);
         j := READ FILE( i. 1, rfile); { Read keyword length.
         FILE SEEK( j, rfile
         header crc := READ FILE(1, 2, rfile); { Read header check sum.
                                                                                  }
          if header crc > 0 then begin
             header crc := header crc + 0x0001;
             FILE SEEK( -2, rfile);
             WRITE FILE( header crc, 2, rfile);
         end:
```

```
CLOSE_FILE( file_name);
    end;
end;
```

```
from WAIT to same
when cc.list mail
begin
   num files := 0;
                                                                                   }
                                 { List all bulletins
   if bulletins then begin
      file name := GET FIRST FILE( name, "BULLETIN");
       while file_name < > empty_string do begin
          mail[num_files] := file_name;
          num files := num files + 1;
          file name := GET NEXT FILE( name, file name);
      end:
   end:
   if messages then
                                 { List all messages 'to' a certain callsign.
                                                                                   }
      if to then begin
          file name := GET FIRST FILE( date, empty string);
          while file name < > empty_string do begin
             if MSG TO( callsign) then begin
                 mail[num files] := file name;
                 num files := num files + 1;
             file_name := GET_NEXT_FILE( date, file_name);
          end;
       end:
                                 { List all messages 'from' a certain callsign.
                                                                                   }
       else if from then begin
          sname[0..1] := " ";
                                 sname[2..7] := callsign;
          file_name := GET_FIRST_FILE( name, sname);
          while file_name < > empty_string do begin
             mail[num files] := file name;
              num files := num files + 1;
             file name := GET NEXT FILE( name, file_name);
          end:
       end;
                                                                                   }
                                  { List all messages
       else begin
          file_name := GET_FIRST_FILE( date, empty_string );
          while file name < > emptry_string do begin
              if file_name[0..7] < > "BULLETIN" then begin
                 mail[num files] := file name;
                 num_files := num_files + 1;
              file name := GET NEXT FILE( date, file_name);
          end:
       end:
   output cc.mail list(num_files - 1, mail);
end;
```

```
from WAIT to same
when cc.post_bulletin
begin
   INCREMENT MSG( "BULLETIN", next source num, mail head, first let,
   mail_num, ext);
   rfile := OPEN FILE( bulletin);
                                     { File has already been created and written by the}
   FILE_SEEK( mn, rfile);
                                     command module. Here, we just keep track
   WRITE_FILE( mail_num, 4, rfile);
                                        { of how many bulletins are active, and
   FILE_SEEK( nd - ml, rfile);
                                        write appropriate mail number into the file
   num_dest := READ_FILE(1, 1, rfile);
                                            { header. In this case, the mail_number
   if num dest > 0x07 then
                                        will not accurately reflect the extension of the}
      FILE_SEEK( 42, rfile);
                                        file name, since this will be assigned by the
   j := READ_FILE(1, 1, rfile); {
                                        writing module.
   FILE_SEEK( j, rfile);
   j := READ FILE(1, 1, rfile);
   FILE SEEK( j, rfile);
   header crc := READ FILE(1, 2, rfile);
   if header_crc > 0 then begin
      for i := 0 to 3 do
          header_crc := header_crc + mail_num[i];
      FILE SEEK( -2, rfile);
      WRITE_FILE( header_crc, 2, rfile);
   CLOSE_FILE( file_name);
end:
from WAIT to same
when cc.delete bulletin
begin
   DECREMENT_MSG( "BULLETIN", mail_head, first_let);
end;
```

```
from WAIT to same
when cc.purge mail
begin
   if not all then begin
      if to then begin
                                 { Purge all messages 'to' a certain callsign.
                                                                                  }
         file name := GET FIRST FILE( date, empty_string);
         while file_name < > empty_string do begin
             if MSG TO( callsign) then begin
                if post time > 0 then begin
                   rfile := OPEN FILE( file name);
                   FILE SEEK( ut. rfile):
                   ul time := READ FILE( 1, 4, rfile);
                   CLOSE FILE( file name):
                   if ul time < post time then begin
                       DELETE FILE( file name);
                       DECREMENT MSG(file name[2..7], mail head, first let);
                   end:
                end;
                else begin
                   DELETE FILE( file name);
                   DECREMENT MSG(file name[2..7], mail head, first let);
                end;
             end:
             file_name := GET_NEXT_FILE( date, file_name);
         end:
      else if from then begin { Purge all messages 'from' a certain callsign.
         sname[0..1] := " ";
         sname[2..7] := callsign;
         file_name := GET_FIRST_FILE( name, sname);
         while file name < > empty string do begin
             if post time > 0 then begin
                rfile := OPEN FILE( file name);
                FILE SEEK( w. rfile):
                ul time := READ FILE( 1, 4, rfile);
                CLOSE FILE( file_name);
                if ul_time < post time then begin
                   DELETE FILE( file name);
                   DECREMENT MSG( callsign, mail head, first let);
                end;
             end:
             else begin
                DELETE FILE( file name);
                DECREMENT MSG( callsign, mail head, first let);
```

```
file_name := GET_NEXT_FILE( name, file_name);
      else begin
                                   { Purge all messages.
         file name := GET FIRST FILE( date, empty string);
         while file name < > empty string do begin
            if file_name[0..7] < > "BULLETIN"
            and file_name[0..7] < > "USRTELEM" then begin
                if post time > 0 then begin
                   rfile := OPEN FILE( file name);
                   FILE SEEK( w, rfile);
                   ul time := READ FILE( 1, 4, rfile);
                   CLOSE FILE( file name);
                   if ul_time < post_time then begin
                      DELETE FILE( file name);
                      DECREMENT MSG(file name[2..7], mail head, first let);
                   end;
                end;
                else begin
                   DELETE_FILE( file_name);
                   DECREMENT MSG(file name[2..7], mail head, first let);
                end:
            end;
             file name := GET NEXT FILE( date, file name);
         end:
      end;
   end:
end:
```

```
from WAIT to same
   when ts.store usr telem
   begin
      INCREMENT MSG( "USRTELEM", next source num, mail head, first let,
      mail num, ext);
      rfile := OPEN FILE( telem); { File has already been created and written by the
      FILE SEEK( mn, rfile);
                                        telemetry module. Here, we just keep track
      WRITE FILE( mail num, 4, rfile);
                                           { of how many usr telem files are active, and }
      FILE SEEK( nd - ml, rfile);
                                           write appropriate mail number into the file
      num_dest := READ_FILE( 1, 1, rfile);
                                              { header. In this case, the mail number
      if num_dest > 0x07 then
                                           will not accurately reflect the extension of the
          FILE SEEK( 42, rfile);
                                           file name, since this will be assigned by the
      j := READ FILE( 1, 1, rfile); {
                                           writing module.
      FILE SEEK( i, rfile);
      j := READ FILE(1, 1, rfile);
      FILE SEEK( j, rfile);
      header crc := READ FILE(1, 2, rfile);
      if header crc > 0 then begin
          for i := 0 to 3 do
             header crc := header crc + mail num[i];
          FILE SEEK( -2, rfile);
          WRITE FILE( header crc, 2, rfile);
      end:
      CLOSE FILE( telem);
   end:
   from WAIT to same
   when ts.delete user telem
                                     { Telemetry module has already deleted the file and }
                                        is only notifying the mail box module.
      DECREMENT MSG( "USRTELEM", mail head, first let);
   end;
                                     { of MAILBOX CONTROL BODY
end;
```

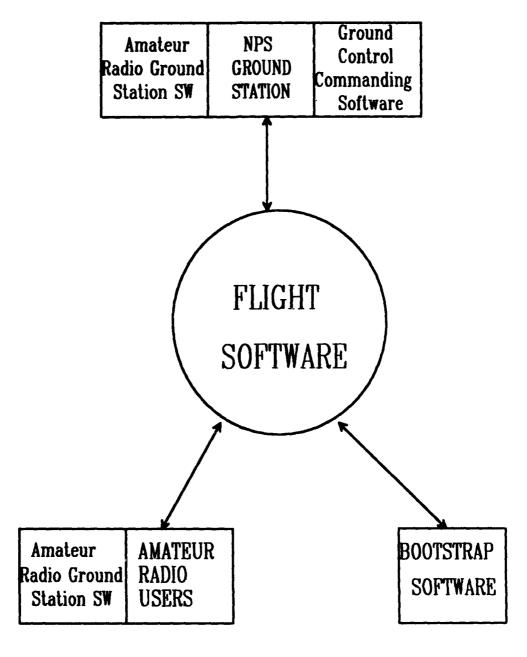
body	PASSWORD_CONTROL_BODY for PASSWORD_CONTROL_TYPE;	external;
body	AUTO_CONTROL_BODY for AUTO_CONTROL_TYPE;	external;
body	GROUND_CONTROL_BODY for GROUND_CONTROL_TYPE;	external;
body	PRIMITIVE_SW_LOADER for PRIMITIVE_SW_LOADER_TYPE;	external;
body	TELEMETRY_GATHER_BODY for TELEMETRY_GATHER_TYPE;	external;
body	A/D_DRIVER_BODY for A/D_DRIVER_TYPE;	external;
body	EVENT_LOGGER_BODY for EVENT_LOGGER_TYPE;	external;
body	EPS_DRIVER_BODY for EPS_DRIVER_TYPE;	external;
body	COMM_DRIVER_BODY for COMM_DRIVER_TYPE;	external;
body	DCS DRIVER BODY for DCS DRIVER TYPE:	external:

```
(Module instantiation and channel connection
modvar
                                 { section for Flight Software Specification.
   Primitive AX25:
                        PRIMITIVE AX25 TYPE;
   Primitive SW Loader:
                        PRIMITIVE SW LOADER TYPE:
   Data Transfer:
                        DATA TRANSFER TYPE:
   Packet Transfer:
                        array[maxlinks] of PACKET TRANSFER TYPE;
                        MAILBOX CONTROL TYPE;
   Mailbox Control:
   Ground Control:
                        GROUND CONTROL TYPE;
   Auto Control:
                        AUTO CONTROL TYPE:
                        EVENT_LOGGER_TYPE;
   Event Logger:
   Password Control:
                        PASSWORD CONTROL TYPE;
                        TELEMETRY_GATHER_TYPE;
   Telemetry_Gather:
   A/D Driver:
                        A/D_DRIVER_TYPE;
   EPS Driver:
                        EPS DRIVER TYPE;
   Comm Driver:
                        COMM DRIVER TYPE;
   DCS Driver:
                        DCS DRIVER TYPE;
initialize
                                 { Initialization Part of the Specification
                                                                              }
begin
   init Primitive AX25 with PRIMITIVE AX25 BODY;
   init Primitive SW Loader with PRIMITIVE SW LOADER BODY;
   init Data Transfer with DATA TRANSFER BODY;
   init Mailbox Control with MAILBOX CONTROL BODY:
   init Ground Control with GROUND CONTROL BODY;
   init Auto Control with AUTO_CONTROL_BODY;
   init Event Logger with EVENT LOGGER BODY:
   init Password Control with PASSWORD_CONTROL_BODY;
   init Telemetry_Gather with TELEMETRY_GATHER BODY;
   init A/D Driver with A/D DRIVER BODY;
   init EPS Driver with EPS DRIVER BODY:
   init Comm Driver with COMM DRIVER BODY;
   init DCS Driver with DCS_DRIVER_BODY;
   all link: Link Type do begin
      init Packet Transfer[link] with PACKET TRANSFER BODY:
      connect Data_Transfer.pc[link] to Packet_Transfer[link].pc;
      connect Mailbox Control.mc[link] to Packet_Transfer[link].mc;
      connect Event Logger.el[link] to Packet Transfer[link].el;
   connect Primitive_AX25.bax[0] to Data_Transfer.bax;
   connect Primitive_AX25.bax[1] to Ground_Control.bax;
   connect Primitive_AX25.bax[2] to Auto_Control.bax;
   connect Primitive_AX25.bax[3] to Primitive_SW_Loader.bax;
   connect Ground_Control.ccd to Data_Transfer.cc[0];
   connect Ground Control.cct to Telemetry_Gather.cc[0];
```

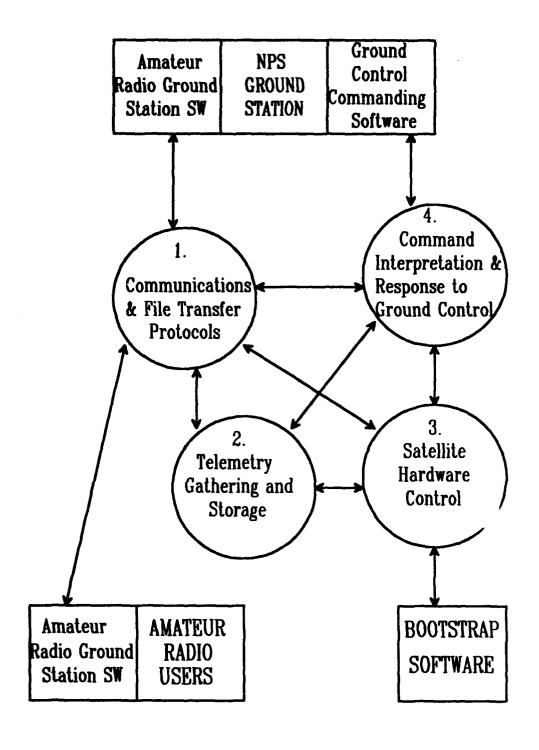
```
t Ground Control.ccl to Primitive SW Londer.cc[0]:
   connect Ground Control.ccp to Password Control.cc[0];
   connect Ground Control.ccm to Mailbox Control.cc[0]:
   connect Ground Control.cce to EPS Driver.cc[0];
   connect Ground Control.ccom to Comm Driver.cc[0];
   connect Ground Control.ccdc to DCS Driver.cc[0];
   connect Ground Control el to Event Logger el [maxlinks]:
   connect Ground Control.sc to Auto Control.sc;
   connect Auto Control.acd to Data Transfer.cc[1];
   connect Auto Control.act to Telemetry Gather.cc[1]:
   connect Auto Control.acp to Password Control.cc[1];
   connect Auto_Control.acm to Mailbox_Control.cc[1];
   connect Auto Control.ace to EPS Driver.cc[1];
   connect Auto Control.acom to Comm Driver.cc[1];
   connect Auto Control.acdc to DCS Driver.cc[1]:
   connect Auto Control.el to Event Logger. el[maxlinks + 1];
   connect Telemetry Gather.ad to A/D Driver.ad;
   connect Telemetry Gather.el to Event Log.el[maxlinks + 2];
   connect Telemetry Gather.ts to Mailbox Control.ts;
   connect Mailbox Control.el to Event Logger.el[maxlinks + 3];
   connect Data Transfer.el to Event Logger.el[maxlinks + 4];
   connect Password Control.el to Event Logger.el[maxlinks + 5];
end:
                                     { of Flight Software Specification
                                                                                       }
end.
```

### APPENDIX B - DATA FLOW DIAGRAMS

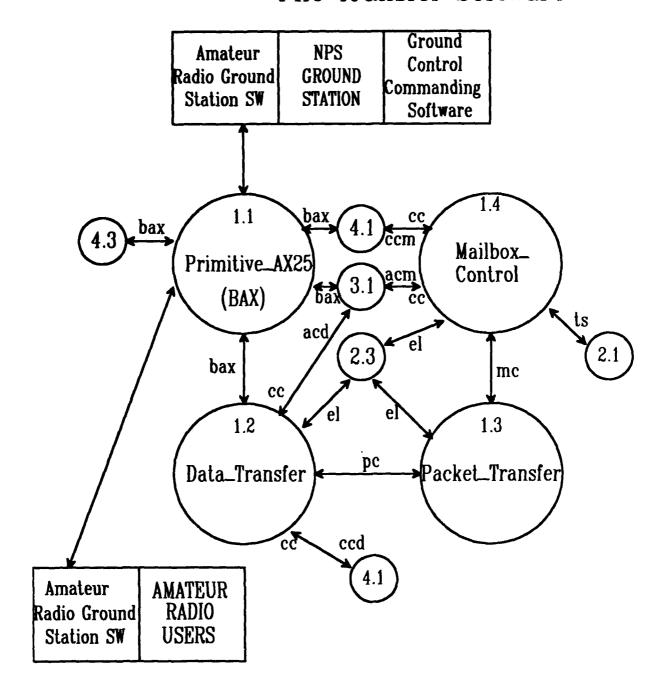
# DFD: CONTEXT DIAGRAM



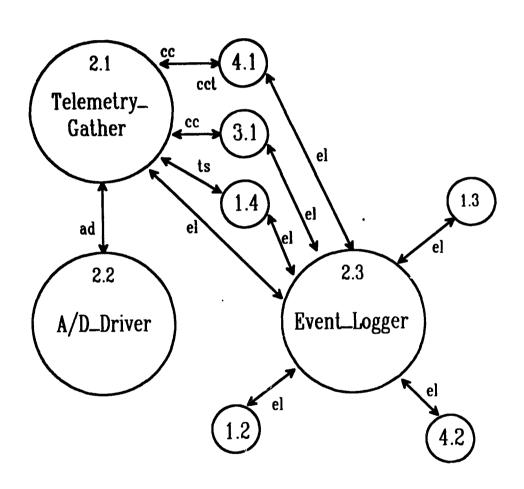
# DFD 0 - FLIGHT SOFTWARE



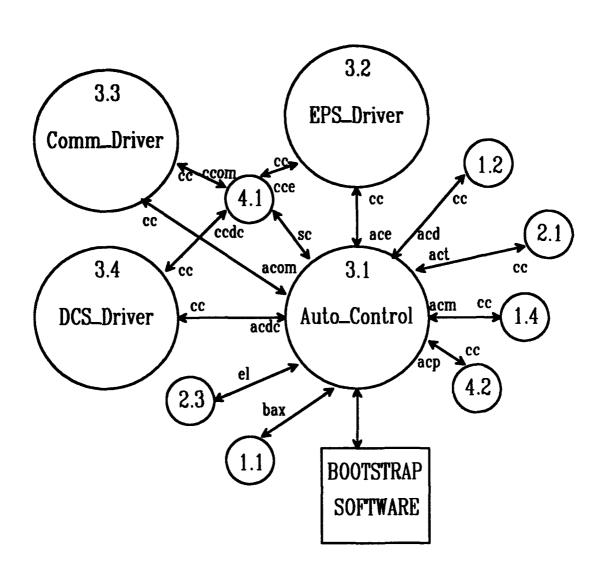
DFD 1 - Communications & File Transfer Software



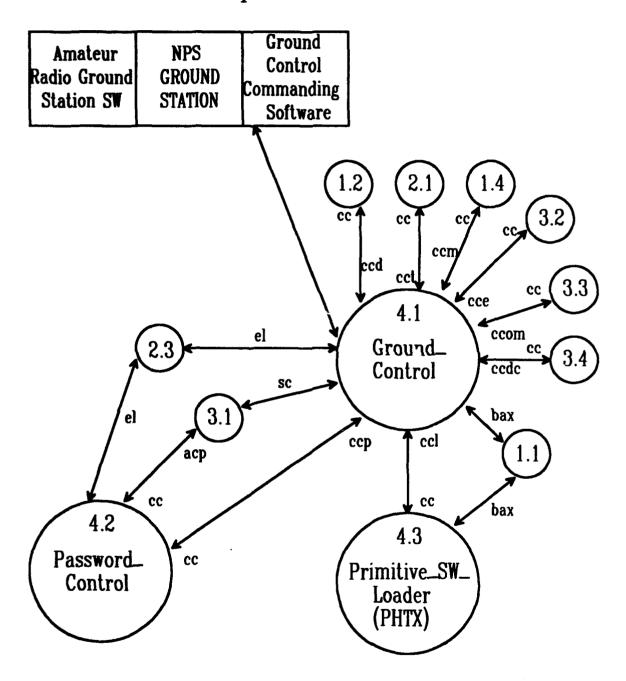
# DFD 2 - Telemetry Gathering & Storage



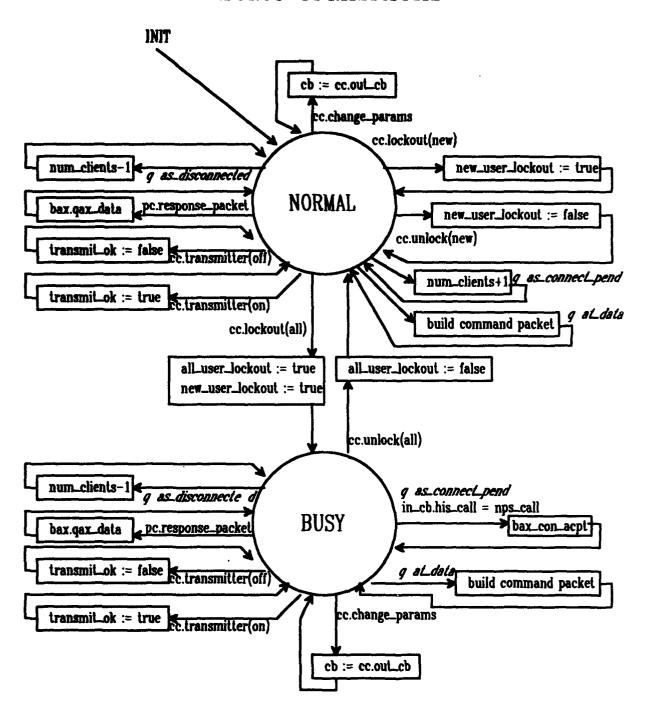
# DFD 3 - Satellite Hardware Control



DFD 4 - Command Interpretation & Response to Ground Control

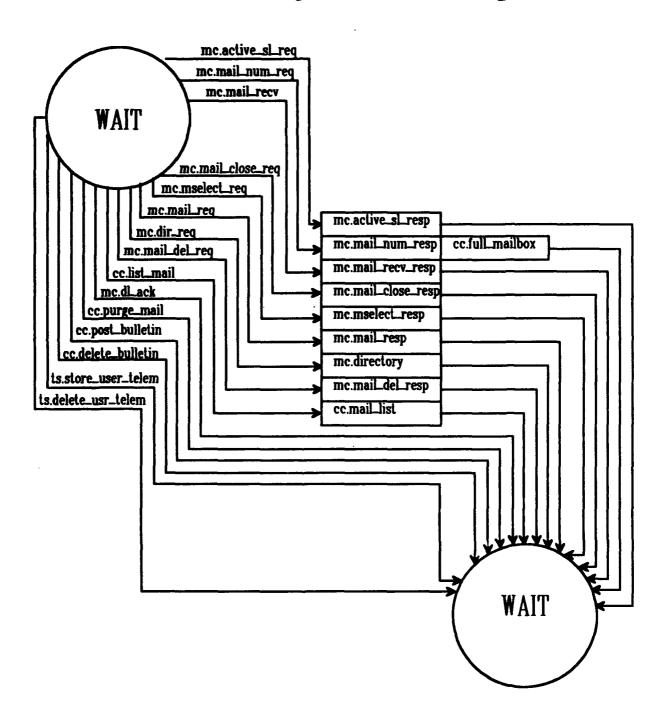


# DFD 1.2 - Data Transfer Module State Transitions



DFD 1.3 - Packet\_Transfer Module **State Transitions** d Lnek.cmd others d ownload\_cmd pc.disconnect (from any state) pc.disconnect UL/DL\_CMD\_WAI UL/DL\_UNINIT u Lack\_resp / ogin\_resp pc.connection mc.active\_sl\_resp mc.active\_sl\_req u pload\_cmd mc.mail\_close\_resp, mc.mail\_num\_rec mc.mail\_del\_r-c mc.mselect\_req mc.dir\_req WAIT\_MAILBOX , mc.directory d ala/dLerror\_resp mc.mail\_del\_resp mc.mselect\_resp me.mail\_recv\_resp. mc.mail\_num\_resp,error I\_error\_res mc.mail\_num\_resp, no error mdmail\_close\_resp. mc.mail\_recv\_resp, no error u Lnak\_resp d Lerror\_resp mc.mail\_req mc.mail\_resp. mc.mail\_recv mc.mail\_close\_req **UL\_ABORT** DL\_FILE\_DATA UL\_DATA\_RX d ata\_end others others u Lerror\_resp ack\_cmd & m c.mail\_close\_reg mc.mail\_resp. d Lcompleted\_resp no error

# DFD 1.4 - Mailbox\_Control Module Response to Messages



#### APPENDIX C - ESTELLE SYNTAX

Table C.1 contains the subset of Pascal syntax which was utilized in Appendix A.

Note that cells with double lines on top and bottom contain definitions somewhat modified from those found in [Ref. 8]. A more complete lexicon and construction rules can be found in Annex C of [Ref. 8].

### Key:

- 1) Each definition ends with a period, ".".
- 2) The or symbol, "|", denotes a choice among options.
- 3) Components enclosed by square brackets, "[]", are optional.
- 4) Parentheses, "()", are used for grouping components in order to clarify definitions.
- 5) Components enclosed by curly brackets, "{ }", may be included zero or more times.
- 6) Symbols shown within quotes, " ", must be typed exactly as they appear. (They will be found in **bold**-face type within the specification.)

# TABLE C.1 PASCAL SYNTAX USED IN APPENDIX A letter = "a" | "b" | ... | "z" | "A" | "B" | ... | "Z". digit = "0" | "1" | "2" | "3" | "4" | "5" | "6" | "7" | "8" | "9" | "A" | "B" | "C" | "D" | "E" | "F". special-symbol = "+" | "-" | "\*" | "=" | "<" | ">" | "[" | "]" | "(" | ")" | ";" | "^" | "< >" | "<=" | ">=" | ":=" | ".." | "word-symbol.

## TABLE C.1 PASCAL SYNTAX USED IN APPENDIX A word-symbol = "and" | "array" | "begin" | "case" | "const" | "do" | "downto" | "else" | "end" | "false" | "for" | "function" | "if" | "in" | "not" | "null" | "of" | "or" | "procedure" | "record" | "repeat" | "then" | "to" | "true" | "type" | " until" | "var" | "while". identifier = letter { letter | digit }. unsigned-integer = digit {digit} | "0x" digit {digit}. character-string = "' string-character { string-character } "'. comment = "{" any-sequence-of-characters-and-separations-of-lines-notcontaining-right-brace "}". block = constant-definition-part type-definition-part variable-declaration-part procedure-and-function-declaration-part statement-part. constant-definition-part = [ "const" constant-definition ";" { constant-definition type-definition-part = [ "type" type-definition ";" { type-definition ";" } ]. variable-declaration-part = [ "var" variable-declaration ";" { variable-declaration procedure-and-function-declaration-part = { (procedure-declaration | functiondeclaration ) ";" }. statement-part = compound-statement. constant-definition = identifier "=" constant. constant = [sign] (unsigned-integer | constant-identifier) | character-string. constant-identifier = identifier. type-definition = identifier "=" type-denoter. type-denoter = ordinal-type | new-type. **new-type** = **new-ordinal-type** | **new-structured-type** | **new-pointer-type**.

### TABLE C.1 PASCAL SYNTAX USED IN APPENDIX A

structured-type-identifier = type-identifier.

pointer-type-identifier = type-identifier.

type-identifier = identifier.

ordinal-type = new-ordinal-type | ordinal-type-identifier.

new-ordinal-type = enumerated-type | subrange-type.

ordinal-type-identifier = uchar | uint | ulong | int | boolean.

uchar = 8-bits-binary-data-or-1-byte-unsigned-integer-or-1-ascii-character.

uint = 2-byte-unsigned-integer.

ulong = 4-byte-unsigned-integer.

int = [sign] unsigned-integer.

sign = "+" | "-".

boolean = "true" | "false".

enumerated-type = "(" identifier-list ")".

identifier-list = identifier ( "," identifier ).

subrange-type = constant ".." constant.

structured-type = new-structured-type | structured-type-identifier.

new-structured-type = array-type | record-type.

array-type = "array" "[" number-components "]" { "[" number-components "]" }

"of" component-type.

number-components = unsigned-integer.

component-type = type-denoter.

record-type = "record" field-list "end".

field-list = record-section { ";" record-section }.

### TABLE C.1 PASCAL SYNTAX USED IN APPENDIX A

record-section = identifier-list ":" type-denoter.

pointer-type = new-pointer-type | pointer-type-identifier.

new-pointer-type = "^" type-identifier.

variable-declaration = identifier-list ":" type-denoter.

variable-access = entire-variable | component-variable | identified-variable.

entire-variable = variable-identifier.

variable-identifier = identifier.

component-variable = indexed-variable | field-designator.

array-variable = variable-access.

index-expression = expression.

field-designator = record-variable "." field-specifier | field-identifier.

record-variable = variable-access.

field-specifier = field identifier.

field-identifier = identifier.

identified-variable = "^" pointer-variable.

pointer-variable = variable-access.

procedure-declaration = procedure-heading ";" directive | procedure-identification ";" procedure-block | procedure-heading ";" procedure-block.

procedure-heading = "procedure" identifier [ formal-parameter-list ].

procedure-identification = "procedure" procedure-identifier.

procedure-identifier = identifier.

procedure-block = block.

```
TABLE C.1 PASCAL SYNTAX USED IN APPENDIX A
function-declaration = function-heading ";" directive
                       | function-identification ";" function-block
                       function-heading ";" function-block.
function-heading = "function" identifier [ formal-parameter-list ] ":" result-type.
function-identification = "function" function-identifier.
function-identifier = identifier.
function-block = block.
result-type = type-denoter.
formal-parameter-list = "(" formal-parameter-section { ";" formal-parameter-
                  section \ ")".
formal-parameter-section = value-parameter-specification
                             variable-parameter-specification.
value-parameter-specification = identifier-list ":" type-identifier.
variable-parameter-specification = "var" identifier-list ":" type-identifier.
expression = simple-expression [ relational-operator simple-expression ].
simple-expression = [sign] term { adding-operator term }.
term = factor { multiplying-operator factor }.
factor = variable-access | unsigned-constant | function-designator |
                  set-constructor | "(" expression ")" | "not" factor.
unsigned-constant = unsigned-number | character-string | constant-identifier |
                  "null".
set-constructor = "[" [ member-designator { "," member-designator } ] "]".
member-designator = expression [ ".." expression ].
multiplying-operator = "*" | "and".
adding-operator = "+" | "-" | "or".
relational-operator = "=" | "< >" | "<" | ">" | "<=" | ">=" | "in".
```

boolean-expression = expression.

# TABLE C.1 PASCAL SYNTAX USED IN APPENDIX A function-designator = function-identifier [ actual-parameter-list ]. actual-parameter-list = "(" actual-parameter { "," actual-parameter } ")". actual-parameter = expression | variable-access | procedure-identifier | function-identifier. statement = (simple-statement | structured-statement). simple-statement = empty-statement | assignment-statement | procedure-statement. empty-statement = .assignment-statement = (variable-access | function identifier) ":=" expression. procedure-statement = procedure-identifier ([actual-parameter-list]). structured-statement = compound-statement | conditional-statement | repetitive-statement. statement-sequence = statement { ";" statement }. compound-statement = "begin" statement-sequence "end". conditional-statement = if-statement | case-statement. if-statement = "if" boolean-expression "then" statement [ else-part ]. else-part = "else" statement. case-statement = "case" case-index "of" case-list-element { ";" case-list-element } [ ";" ] "end". case-list-element = case-constant-list ":" statement. case-constant-list = case-constant $\{ ", " \text{ case-constant } \}$ . case-constant = constant.case-index = expression.repetitive-statement = repeat-statement | while-statement | for-statement. repeat-statement = "repeat" statement-sequence "until" boolean-expression. while-statement = "while" boolean-expression "do" statement.

### TABLE C.1 PASCAL SYNTAX USED IN APPENDIX A

for-statement = "for" control-variable ":=" initial-value ( "to" | "downto" )
final-value "do" statement.

control-variable = entire-variable.

initial-value = expression.

final-value = expression.

The following table lists Estelle-specific reserved words which have been used in Appendix A. The expected location and function associated with each is also indicated.

TABLE C.2 ESTELLE-SPECIFIC RESERVED WORDS			
Reserved Word	Location	Function	
specification	Beginning of entire specification block.	Identifies the name of the specification.	
any	In a constant declaration.	Declares that a value of the indicated type must be chosen during implementation.	
••	By itself, on the right- hand side of a type definition.	Indicates that the actual internal details of the type have not yet been determined. The final definition may be implementation-dependant.	
channel	In the channel definition section, which immediately follows the global constant, type and and variable declaration sections.	Indicates the beginning of a channel definition. Followed by the channel name, and then, within parentheses, the end-point names.	

TABLE C.2 ESTELLE-SPECIFIC RESERVED WORDS		
Reserved Word	Location	Function
by	Within a channel definition.	'by' is followed by one of the channel end point names and then by a list of the messages which can be sent from that end point. Following each message name, in parentheses, is a list of the parameters for that message type.
module	In the module header definition section, which follows the global function and procedure declarations which follow the channel definition section.	Indicates the beginning of a module header definition. "module" is followed by the name of the module type. The module header definition defines the interfaces with the module.
system process	In the module header definition following the name of the module type.	Indicates that the module is an autonomous process, not a subprocess enclosed within another.
definition.  interface points for the mode followed by the channel nation of which is given a channel among those defined in the definition section. In parer indicated which end point to plays the role of. That in the defines the type of messages		Indicates the beginning of the list of interface points for the module. It is followed by the channel names, each of which is given a channel type from among those defined in the channel definition section. In parentheses is indicated which end point this module plays the role of. That in turn, defines the type of messages which can be sent by this module.
individual queue	In an interface point declaration within a module header definition.	Indicates that all messages to or from this module via the indicated channel will be maintained in an individual queue for that module alone.

TABLE C.2 ESTELLE-SPECIFIC RESERVED WORDS		
Reserved Location Function Word		Function
queue queue	In an interface point declaration within a module header definition.	Indicates that the module will share the queue for this channel with all other "common queue" modules playing the role of the same kind of end point for the same kind of channel. Or, if a module has an array of channels of the same type, all of the channels may use a common queue.
definition section, which body definiti		Indicates the beginning of a module body definition. This is where the actual behavior of the module is defined.
external	In the module body definition section, following the body name and the module type.	Indicates that the body definition for this module is external to the current specification. It may not yet have been developed, may be under development by another team, or it may be completely external to the system at hand, with only the interface defined by the module header definition being of any importance.
state	Within a module body definition, following the local const, type and var declaration sections.	Marks the beginning of the list of state names for this module.

TABLE C.2 ESTELLE-SPECIFIC RESERVED WORDS			
Reserved Word	Location	Function	
stateset	Within a module body definition, following the list of state names.	Defines a common name to be used for several states, when they have similar transitions. If the module state machine reacts the same way to a particular stimulus when in any one of several different states, these states may be grouped together by a stateset so that the behavior need only be written out once for all the affected states.	
primitive	In a function or procedure declaration.	Indicates that the algorithmic details of the function or procedure are not included in the present specification. The function or procedure may be deemed to be commonly understood or readily available from the operating system, or the details may simply not be relevant to the aspect of the system currently under consideration.	
initialize	Following the last local function or procedure declaration within a module body definition.	Indicates the initial state of the module state machine when it is instantiated. Statements between the "begin" and "end" keywords may be used to set up initial variable values, and to take any other automatic "start-up" actions.	
initialization section in transition section of the module body  the module body  transition section of the module body		Indicates the beginning of the state transition section of the module body. All possible state transitions will be listed within this section.	
from	In the transition section of the module body definition.	Indicates the state from which the transition takes place.	

TABLE C.2 ESTELLE-SPECIFIC RESERVED WORDS			
Reserved Word	Location	Function	
to	In the transition section of the module body definition.	Indicates the state in which the module will be, following the transition.	
when	Following the "from - to" clause of the transition statement.	Identifies the stimulus which may trigger the transition. "when" is usually followed by the name of an interface point, with a period, ".", and the kind of message from that channel which could trigger the transition. Any parameters associated with the incoming message may be referenced directly by their value-parameter names within the "begin - end" block of the transition statement.	
provided	Following the "when" clause of the transition statement.	Indicates any further conditions for the transition to occur. 'provided' is usually followed by an expression in which one of the parameters of the "when message" is compared with a necessary condition. The transition only occurs when the module is in the "from state", the "when message" arrives, and the message parameter meets the necessary condition. When the transition occurs, all of the statements between "begin" and "end" are executed before the module enters the "to state" and waits for the next stimulus.	

TABLE C.2 ESTELLE-SPECIFIC RESERVED WORDS			
Reserved Word	Location	Function	
output	Within the "begin - end" block of a transition statement.	Indicates that a message is to be sent out at the interface point indicated. The interface point name is shown on the left side of a period, with the type of message on the right side. If there are any message parameters, variables of the appropriate types must be prepared with the proper values, and must be included in parentheses following the message name.	
modvar	Following all module body definitions.	Indicates the beginning of the module instantiation and channel connection section of the specification.	
instantiation section exactly how many module type will b		Indicates the beginning of the module instantiation section, which will define exactly how many copies of each module type will be created, using which module bodies.	
init	In the module instantiation section.  "init" indicates initialize, but means instantiate. Use the body indicated by the "with (there could be more than or body for a particular module)		
connect	In the modvar section, following the module instantiations.	Indicates that the interface points of two modules will be connected together. Further defines which specific channels go between which specific module instantiations.	

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