

AD-R164 487

DEVELOPMENT OF A TESTBED FOR MULTISENSOR DISTRIBUTED
DECISION ALGORITHMS(U) NAVAL POSTGRADUATE SCHOOL
MONTEREY CA M A SCHON DEC 85

1/1

UNCLASSIFIED

F/G 9/2

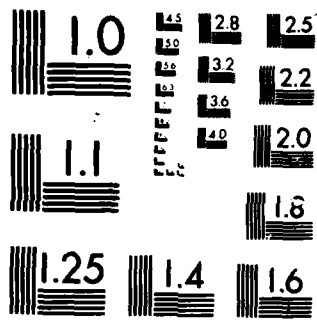
NL

A large grid of blacked-out cells covering most of the page, likely representing redacted content. The grid consists of approximately 13 columns and 8 rows of solid black rectangles.

END

F/G 9/2

NL



MICROCOPY RESOLUTION TEST CHART
NATIONAL BUREAU OF STANDARDS-1963-A

2

NAVAL POSTGRADUATE SCHOOL Monterey, California

AD-A164 487



DTIC
ELECTE
FEB 26 1986
S D

THESIS

DEVELOPMENT OF A TESTBED
FOR MULTISENSOR
DISTRIBUTED DECISION ALGORITHMS

By
Mark A. Schon

December 1985

Thesis Advisor:

Charles W. Therrien

Approved for public release; distribution is unlimited.

DTIC FILE COPY

86 2 25 030

REPORT DOCUMENTATION PAGE

1a. REPORT SECURITY CLASSIFICATION UNCLASSIFIED			1b. RESTRICTIVE MARKINGS			
2a. SECURITY CLASSIFICATION AUTHORITY			3. DISTRIBUTION/AVAILABILITY OF REPORT Approved for public release; distribution is unlimited			
2b. DECLASSIFICATION/DOWNGRADING SCHEDULE						
4. PERFORMING ORGANIZATION REPORT NUMBER(S)			5. MONITORING ORGANIZATION REPORT NUMBER(S)			
6a. NAME OF PERFORMING ORGANIZATION Naval Postgraduate School		6b. OFFICE SYMBOL (if applicable) 62	7a. NAME OF MONITORING ORGANIZATION Naval Postgraduate School			
6c. ADDRESS (City, State, and ZIP Code) Monterey, CA 93943-5100			7b. ADDRESS (City, State, and ZIP Code) Monterey, CA 93943-5100			
8a. NAME OF FUNDING/SPONSORING ORGANIZATION		8b. OFFICE SYMBOL (if applicable)	9. PROCUREMENT INSTRUMENT IDENTIFICATION NUMBER			
8c. ADDRESS (City, State, and ZIP Code)			10. SOURCE OF FUNDING NUMBERS			
			PROGRAM ELEMENT NO.	PROJECT NO.	TASK NO.	WORK UNIT ACCESSION NO.
11. TITLE (Include Security Classification) DEVELOPMENT OF A TESTBED FOR MULTISENSOR DISTRIBUTED DECISION ALGORITHMS						
12. PERSONAL AUTHOR(S) Schon, Mark A.						
13a. TYPE OF REPORT Master's Thesis		13b. TIME COVERED FROM _____ TO _____		14. DATE OF REPORT (Year, Month, Day) 1985 December		15. PAGE COUNT 85
16. SUPPLEMENTARY NOTATION						
17. COSATI CODES			18. SUBJECT TERMS (Continue on reverse if necessary and identify by block number)			
FIELD	GROUP	SUB-GROUP	Distributed Decision Processing; Computer Network; Microcomputer Clusters; Process Synchronization; Network Communication			
19. ABSTRACT (Continue on reverse if necessary and identify by block number)						
<p>Distributed decision problems arise when two or more sensors viewing the same phenomenon must work cooperatively to draw inferences about the observed situation. Typical examples are in target detection and target classification. Such problems are characterized by distributed processing of information and communication between processors over a limited bandwidth data link. This thesis presents some statistical distributed decision algorithms and describes the implementation of one of them on a set of loosely coupled multiprocessor clusters which simulate the distributed environment characterizing multisensor decision problems.</p>						
20. DISTRIBUTION/AVAILABILITY OF ABSTRACT <input checked="" type="checkbox"/> UNCLASSIFIED/UNLIMITED <input type="checkbox"/> SAME AS RPT <input type="checkbox"/> DTIC USERS			21. ABSTRACT SECURITY CLASSIFICATION UNCLASSIFIED			
22a. NAME OF RESPONSIBLE INDIVIDUAL C. W. Therrien			22b. TELEPHONE (Include Area Code) (408)646-2160		22c. OFFICE SYMBOL 62T1	

The purpose of the implementation was to investigate problems of communication and process synchronization in a pair of processor clusters performing a statistical distributed decision algorithm. This thesis describes how these communication and synchronization problems were addressed and solved.

Approved for public release; distribution is unlimited.

Development of a Testbed for Multisensor Distributed Decision Algorithms

by

Mark Alan Schon
Captain, United States Marine Corps
B.S., University of Utah, 1978

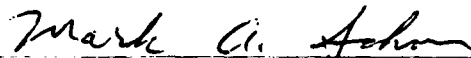
Submitted in partial fulfillment of the
requirements for the degree of

MASTER OF SCIENCE IN ELECTRICAL ENGINEERING

from the

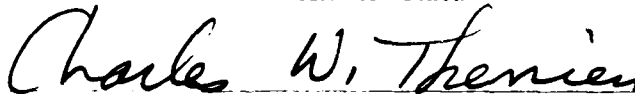
NAVAL POSTGRADUATE SCHOOL
December 1985

Author

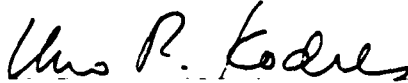


Mark A. Schon

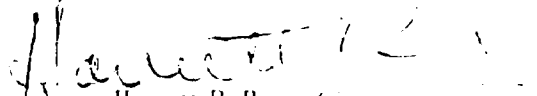
Approved by



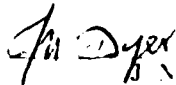
Charles W. Therrien, Thesis Advisor



Uno R. Kodres, Second Reader



Harriett B. Rigas, Chairman,
Department of Electrical and Computer Engineering



B. S. Dyer
Dean of Science and Engineering

ABSTRACT

Distributed decision problems arise when two or more sensors viewing the same phenomenon must work cooperatively to draw inferences about the observed situation. Typical examples are in target detection and target classification. Such problems are characterized by distributed processing of information and communication between processors over a limited bandwidth data link. This thesis presents some statistical distributed decision algorithms and describes the implementation of one of them on a set of loosely coupled multiprocessor clusters which simulate the distributed environment characterizing multisensor decision problems. The purpose of the implementation was to investigate problems of communication and process synchronization in a pair of processor clusters performing a statistical distributed decision algorithm. This thesis describes how these communication and synchronization problems were addressed and solved.

DISCLAIMER

Some terms used in this thesis are registered trademarks of commercial products. Rather than attempt to cite each occurrence of a trademark, all trademarks appearing in this thesis are listed below following the name of the firm holding the trademark:

1. INTEL Corporation, Santa Clara, California

8086 MULTIBUS

2. Digital Research, Pacific Grove, California

PL/I-86 LINK86

3. XEROX Corporation, Stamford, Connecticut

Ethernet Local Area Network

4. InterLAN Corporation, Westford, Massachusetts

NI3010 Ethernet Communication Controller Board

Accession For	
NTIS CRA&I	<input checked="" type="checkbox"/>
DTIC TAB	<input type="checkbox"/>
Unannounced	<input type="checkbox"/>
Justification	
By	
Distribution	
Availability Codes	
Dist	Availability or special
A-1	

TABLE OF CONTENTS

I.	INTRODUCTION	9
	A. PROBLEM DESCRIPTION	9
	B. HARDWARE/SOFTWARE CONFIGURATION	10
	C. STRUCTURE OF THE THESIS	11
II.	DISTRIBUTED DECISION ALGORITHMS	13
	A. SUMMARY OF ALGORITHMS	13
	1. Tenney - Sandell Algorithm	13
	2. Relaxation Algorithms	13
	3. The Generalized Likelihood Ratio Test	14
	4. Decision Based on the Nearest Neighbor Rule	16
	B. GENERALIZED LIKELIHOOD RATIO TEST	17
III.	THE TEST ENVIRONMENT	21
	A. HARDWARE DESCRIPTION	21
	1. The Cluster	21
	2. Real-Time Cluster Star (RTC *)	21
	B. THE OPERATING SYSTEM ENVIRONMENT	22
	1. The Synchronization Model	23
	2. Eventcount Distribution	24
	3. Data Distribution	24
	C. ALGORITHM IMPLEMENTATION	25
	1. Process Distributivity/Parallel Processing	25
	2. Process Synchronization	26
	D. RESULTS OF THE SIMULATION	28

IV. CONCLUSIONS	29
APPENDIX A: Quadratic Classifiers	30
APPENDIX B: LINK86 Input Option Files	33
APPENDIX C: Device Driver and Packet Processor Source Code	36
APPENDIX D: Distributed Decision Algorithm Source Code	59
LIST OF REFERENCES	80
INITIAL DISTRIBUTION LIST	82

LIST OF FIGURES

1. Distributed Decision Scenario	9
2. Cluster Architecture	21
3. Real-Time Cluster Star (RTC*) Architecture	22
4. Computations of Reduced Statistics	26
5. Synchronization Diagram	27

I. INTRODUCTION

A. PROBLEM DESCRIPTION

Modern military battle systems increasingly rely on the coordinated use of information from multiple sources to assess the battlefield situation. Two or more remotely located sensors may observe the same object with the purpose of drawing inferences about the observation. A common example is in the use of radars to detect and eventually classify objects for purposes of an appropriate response. In this type of scenario it is important to process the acquired information jointly to arrive at the optimum or near optimum decision.

A simple example to demonstrate the distributed decision scenario is illustrated in Fig. 1 and explained below.

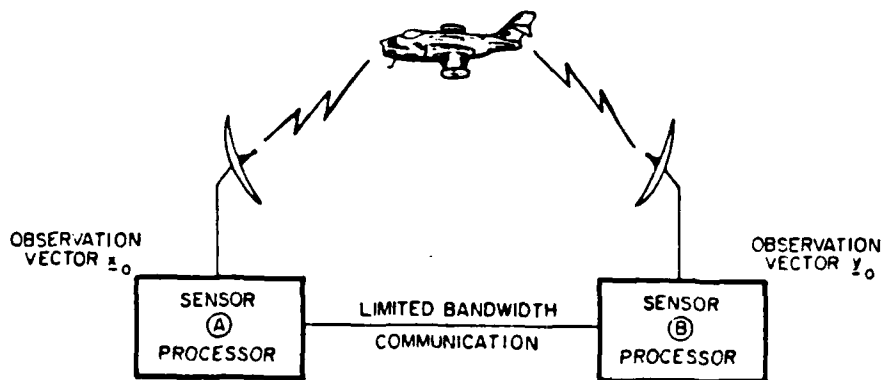


Figure 1 - Distributed Decision Scenario

Two sensors, labeled A and B, observe the same area in space to jointly make a binary decision based on the statistical properties of the observations: either a target is detected or there is no target detected. In certain situations the optimal decision made by each sensor acting individually would result in each deciding that a target exists when an optimal joint decision would decide that a target does not exist. This dichotomy points out that in order for a higher level process

to make a correct decision about the object, the distributive nature of the problem must be built into the front end statistical decision procedure.

The problem of distributed processing of the observation data to achieve optimal or near optimal decisions is discussed herein. The sensors are configured to perform computations to reduce the observation data and to communicate among themselves over a limited bandwidth channel. Algorithms which operate in this environment are called distributed decision algorithms. Algorithms which perform computations on all the observation data collected and gathered at one central location are called centralized algorithms.

This thesis deals specifically with modeling a particular class of distributed decision algorithms in a multiprocessor environment, and with related issues of process synchronization. Although centralized algorithms are not of concern here, a companion thesis [1] deals with non real-time simulation and evaluation of distributed decision algorithms and comparison with centralized algorithms.

Although the thesis deals with a particular class of distributed decision algorithms, the implementation problems presented by the algorithm would be typical of most distributed algorithms. Thus the work can be regarded as developing a test facility in which distributed decision algorithms can be tested in a realistic computational environment.

The problem is to model the processing environment of two sensors which collect data on a common object. The sensors and their associated processors then perform parallel processing to partially reduce the data and the partial results are exchanged via a local area network. A final decision about the observed object is then made at each sensor, based on the locally processed data and the exchanged information.

B. HARDWARE/SOFTWARE CONFIGURATION

The hardware/software configuration used for the modeling of the distributed decision network was the REAL-TIME CLUSTER STAR (RTC *) system. This system was developed by thesis students under the AEGIS Project Group at the Naval Postgraduate School. RTC * was designed to handle algorithms of the type

incorporated here with the appropriate synchronization and control primitives. The hardware consists of two clusters of single board computers (SBC's) sharing a common backplane with an Ethernet local area network(LAN) serving as the communications link. The operating system is a distributed multicomputer real-time executive that permits asynchronous parallel operation of processes resident on SBC's of the same cluster and in separate clusters linked by the LAN. User processes, such as the distributed decision algorithms, are resident in the local memory of each SBC. They can share data and control variables using the common memory in each cluster, as well as the backplane and the LAN data paths.

A detailed description of the hardware system and the software operating system is provided in [2]. The distributed decision algorithms are organized as a number of separate processes on various single board microcomputers in the two cluster arrangement. Process synchronization is achieved through certain **await**, **advance**, and **read** primitives to control the orderly multiple/parallel process execution as well as a sequencer to control the allocation of the LAN shared resource. Each cluster simulates the operations that would be performed by the sensor processors. Data read from disk storage simulates the input sensor observations. Each processor then performs the necessary computations to reduce the data to the statistics required for the joint decision. One set of statistics are then exchanged between clusters while another set is retained locally and computation is continued to produce a combined statistic based on the joint data. This combined statistic is then compared with a predetermined threshold to make the detection decision. Computations continue while data is available for input and the decision results are displayed on a local console of each cluster.

C. STRUCTURE OF THE THESIS

In the remainder of this thesis the distributed decision problem is defined and various distributed decision algorithms and their characteristics are discussed. The implementation of one algorithm in a distributed multiprocessor test environment is introduced and discussed in detail. Emphasis is placed on

obtaining solutions to the problems of communication and synchronization for processes operating in two remote computer systems. The specific contents of each chapter is as follows.

Chapter II presents the distributed decision problem with a discussion of a specific distributed decision algorithm. Simple examples illustrate the detection problem with a binary decision rule.

Chapter III presents the implementation of a specific detection distributed decision algorithm in the RTC* multicomputer system and discusses important issues relevant to the implementation of this type of algorithm.

Chapter IV is a summary of the findings and summarizes the results of the implementation in the RTC* multicomputer system.

II. DISTRIBUTED DECISION ALGORITHMS

A. SUMMARY OF ALGORITHMS

Alternative approaches to a simple binary (two hypothesis) decision problem are presented in this chapter. The various algorithms have overall similar characteristics in that local computations are performed by each sensor, reduced data is exchanged over a limited capacity communications channel, and final decisions are made based on the joint observations of the sensors.

The discussion here assumes that there are only two sensors involved (A and B) and that the task is to make a binary decision (H_1 : target is present, or H_2 : no target is present). Generalization of most of these methods to multiple sensors and/or multiple hypotheses is possible.

1. Tenney - Sandell Algorithm

Tenney and Sandell [3] seem to have been the first to look at distributed decision algorithms of the type described here. In their work, the observations of the two sensors are assumed to be independent when conditioned on the decision hypotheses. Such independence of observations could arise if the sensors measured different physical properties of the target (e.g. radar cross section and infrared radiation). The sensors each make a binary decision based on their own observations and send the result (a single bit) to a fusion center for arbitration. A cost criterion was devised that depends on the decisions made by each sensor and on the two hypotheses. Tenney and Sandell showed that the procedure that minimized the expected value of the cost is a likelihood ratio test at each sensor. However the thresholds used by the two sensors are coupled through some integral equations.

2. Relaxation Algorithms

Relaxation algorithms [4,5] are another way to execute distributed decisions. These algorithms are less well-founded in a theoretical sense, but seem

to work well in practice. In the relaxation algorithm each sensor makes an initial decision based on its own observations. The decisions are exchanged and each sensor may then revise its decision based on the new information. The procedure works best when there are multiple decision makers involved and may require more than a single iteration to converge.

3. The Generalized Likelihood Ratio Test

If the information exchanged between sensors is more than a single bit, but limited to, say, a single floating point number, then a whole new class of procedures can be suggested. In particular, if the observations are independent as in the Tenney-Sandell analysis, then the likelihood ratio for the joint observations factors into two parts, each depending only on the observations of a single sensor. Thus each sensor can compute the likelihood ratio (or log likelihood ratio) statistic for its own observations and send it to the other sensor. Each sensor then has the complete information required for making a decision to minimize probability of error based on the joint observations.

A more interesting problem occurs if the observations are correlated. In this case the joint likelihood ratio does not factor in such a convenient way. However, a procedure can be suggested that leads to a relatively simple decision algorithm. Let the observations acquired by sensors A and B be represented by \mathbf{x}_0 and \mathbf{y}_0 respectively. The optimal centralized test to minimize the probability of error has the form

$$\ln \frac{p_1(\mathbf{x}_0, \mathbf{y}_0)}{p_2(\mathbf{x}_0, \mathbf{y}_0)} = \ln \frac{p_1(\mathbf{x}_0)}{p_2(\mathbf{x}_0)} + \ln \frac{p_1(\mathbf{y}_0 | \mathbf{x}_0)}{p_2(\mathbf{y}_0 | \mathbf{x}_0)} \underset{H_2}{\overset{H_1}{>}} \ln T \quad (1)$$

where the subscript i on each probability density function p indicates that the density function is for hypothesis H_i . A distributed form of this test can be developed by allowing sensor A to compute the first term in (1) and allowing sensor B to compute an approximation to the second term (the conditional log likelihood ratio) by using some estimate for the observations \mathbf{x}_0 . This procedure is

known as a *generalized likelihood ratio test* [6]. In essence, when the density function involves an unknown parameter (in our case x_0 in the second term in (1)) estimates are made based on each hypothesis (x_1 for H_1 and x_2 for H_2) and used in the corresponding density function. The form of the second term then becomes

$$\ln \frac{p_1(y_0, \hat{x}_1)}{p_2(y_0, \hat{x}_2)} \quad (2)$$

If sensor B sends the result of this computation to sensor A, then the test (1), can be evaluated to make a decision. A symmetric computation can be made with the roles of A and B reversed, where the the estimates for y_0 are \hat{y}_1 and \hat{y}_2 at sensor A.

The decision rule just described has a number of essential differences from the corresponding centralized algorithm. First, since the likelihood ratio evaluated by one sensor uses an *estimate* for the other sensor's observations, the performance of the algorithm will in general be different and suboptimal when compared to the centralized test. Second, since the two sensors perform symmetric computations with the roles of x_0 and y_0 reversed, there will, in general be a region of the combined observation space where the decisions of the two sensors do not agree. The properties of this class of distributed decision algorithms is dependent on the various methods of estimating the unknown observations x_0 . If the sensors are allowed to exchange only a single statistic then the estimate for x_0 must be derived entirely from y_0 (e.g. using MAP estimation) and the resulting decision rule is of the form

$$\lambda_A(x_0) - \lambda_B(y_0) \underset{H_2}{\overset{H_1}{>}} \ln T \quad (3)$$

This limits the degree to which the distributed test can approximate the centralized test since in many cases the centralized test will not be separable.

The log likelihood ratios in (3) are computed at their respective sensors and once the primed statistic is received, it is added to the unprimed locally

computed statistic and the result is compared to the known threshold, $\ln T$. For the case of Gaussian observations, the densities, p_1 and p_2 of (1), are of the form

$$p_i = \frac{1}{(2\pi)^{\frac{N}{2}} |\mathbf{K}^{(i)}|^{\frac{1}{2}}} \exp \left[-\frac{1}{2} [\mathbf{x} - \mathbf{m}^{(i)}]^T [\mathbf{K}^{(i)}]^{-1} [\mathbf{x} - \mathbf{m}^{(i)}] \right] \quad i=1,2 \quad (4)$$

and the points where the sum of the statistics in (3) is equal to $\ln T$ establish a decision boundary for this particular decision rule. If observations \mathbf{x}_0 and \mathbf{y}_0 result in a point with value greater than the boundary value, the decision is H_1 and if the value is less than the boundary value the decision is H_2 . Decision boundaries for Gaussian density functions are generally elliptical, parabolic, or hyperbolic and define two (not necessarily connected) regions, one for each hypothesis.

4. Decision Based on the Nearest Neighbor Rule

A final form of distributed decision algorithm is based on the k-nearest neighbor rule of pattern recognition [7]. In this nonparametric decision rule, a set of observations to be tested is represented as a point in a multidimensional observation space. Also existing in this space are previously given sets of points (training data) corresponding to each of the two hypotheses. The distance of the measured observations to each of the other points is computed to determine its k nearest neighbors. If most of the neighbors correspond to H_1 , then the given observations are also associated with H_1 , otherwise the given observations are classified according to H_2 .

A distributed form of this decision rule can be developed by letting each sensor determine a small number of nearest neighbors in the \mathbf{x} or \mathbf{y} subspace. If the labels of these points and their distances from the observation data are interchanged, one can compute the distances in the \mathbf{xy} observation space and classify the observation data. This policy does not guarantee that the true nearest neighbors will always be found but allows a decision to be made without further iterations and exchange of information.

B. GENERALIZED LIKELIHOOD RATIO TEST

The algorithm based on the generalized likelihood ratio test was chosen for implementation on the distributed system. It has requirements for communication and process synchronization that are representative of distributed decision algorithms in general. The performance characteristics of the generalized likelihood ratio test are investigated in [1]. If the joint density function for vector observations \mathbf{x} and \mathbf{y} is Gaussian, then a quadratic decision boundary results. This is known as a quadratic classifier [8]. The joint density function for observations \mathbf{x} and \mathbf{y} has the form

$$p_i(\mathbf{z}) = \frac{1}{(2\pi)^{\frac{N}{2}} |\mathbf{K}^{(i)}|^{\frac{1}{2}}} \exp \left[-\frac{1}{2} \{ \mathbf{z} - \mathbf{m}^{(i)} \}^T [\mathbf{K}^{(i)}]^{-1} \{ \mathbf{z} - \mathbf{m}^{(i)} \} \right] \quad i=1,2 \quad (5)$$

where \mathbf{z} is the observation vector with elements \mathbf{x} and \mathbf{y} and $\mathbf{m}^{(i)}$ is the mean vector

$$\mathbf{m}^{(i)} = \begin{bmatrix} \mathbf{m}_x^{(i)} \\ \mathbf{m}_y^{(i)} \end{bmatrix} \quad i=1,2 \quad (6)$$

and \mathbf{K} is the covariance matrix partitioned as follows

$$\mathbf{K} = \begin{bmatrix} \mathbf{K}_x^{(i)} & \mathbf{B}_{xy}^{(i)} \\ \mathbf{B}_{xy}^{(i)T} & \mathbf{K}_y^{(i)} \end{bmatrix} \quad i=1,2 \quad (7)$$

Note that $\mathbf{K}_x^{(i)}$ is the covariance matrix for \mathbf{x} , $\mathbf{K}_y^{(i)}$ is the covariance matrix for \mathbf{y} , and $\mathbf{B}_{xy}^{(i)}$ is the cross covariance matrix between \mathbf{x} and \mathbf{y} . The marginal and conditional densities are Gaussian [9] and are given by

$$p_i(\mathbf{x}) = \frac{1}{(2\pi)^{\frac{N}{2}} |\mathbf{K}_x^{(i)}|^{\frac{1}{2}}} \exp \left[-\frac{1}{2} \{ \mathbf{x} - \mathbf{m}_x^{(i)} \}^T [\mathbf{K}_x^{(i)}]^{-1} \{ \mathbf{x} - \mathbf{m}_x^{(i)} \} \right] \quad i=1,2 \quad (8)$$

$$p_i(\mathbf{y} | \mathbf{x}) = \frac{1}{(2\pi)^{\frac{N}{2}} |\mathbf{K}_y^{(i)}|^{-\frac{1}{2}}} \exp \left[-\frac{1}{2} (\mathbf{y} - \mathbf{m}_y^{(i)})^T [\mathbf{K}_y^{(i)}]^{-1} (\mathbf{y} - \mathbf{m}_y^{(i)}) \right] \quad i=1,2 \quad (9)$$

$$p_i(\mathbf{y}) = \frac{1}{(2\pi)^{\frac{N}{2}} |\mathbf{K}_y^{(i)}|^{-\frac{1}{2}}} \exp \left[-\frac{1}{2} (\mathbf{y} - \mathbf{m}_y^{(i)})^T [\mathbf{K}_y^{(i)}]^{-1} (\mathbf{y} - \mathbf{m}_y^{(i)}) \right] \quad i=1,2 \quad (10)$$

$$p_i(\mathbf{x} | \mathbf{y}) = \frac{1}{(2\pi)^{\frac{N}{2}} |\mathbf{K}_x^{(i)}|^{-\frac{1}{2}}} \exp \left[-\frac{1}{2} (\mathbf{x} - \mathbf{m}_x^{(i)})^T [\mathbf{K}_x^{(i)}]^{-1} (\mathbf{x} - \mathbf{m}_x^{(i)}) \right] \quad i=1,2 \quad (11)$$

where the conditional covariances and means have the form

$$\mathbf{K}_{y|z}^{(i)} = \mathbf{K}_y^{(i)} - \mathbf{B}_{zy}^{(i)T} [\mathbf{K}_z^{(i)}]^{-1} \mathbf{B}_{zy}^{(i)}, \quad i=1,2 \quad (12)$$

$$\mathbf{m}_{y|z}^{(i)} = \mathbf{m}_y^{(i)} + \mathbf{B}_{zy}^{(i)T} [\mathbf{K}_z^{(i)}]^{-1} (\mathbf{x} - \mathbf{m}_z^{(i)}), \quad i=1,2 \quad (13)$$

Since the \mathbf{x} term of (13) is not available at the given sensor, an estimate of the form

$$\hat{\mathbf{x}}_i = \mathbf{m}_z^{(i)} - \mathbf{B}_{zy}^{(i)} [\mathbf{K}_y^{(i)}]^{-1} (\mathbf{y} - \mathbf{m}_y^{(i)}), \quad i=1,2 \quad (14)$$

is used. The estimate is the value of \mathbf{x} that maximizes the density $p_i(\mathbf{x} | \mathbf{y})$. Symmetric forms of (12), (13), and (14) are used for $\mathbf{K}_{z|y}^{(i)}$, $\mathbf{m}_{z|y}^{(i)}$, and $\hat{\mathbf{y}}_i$ at the other sensor.

The natural logarithm of (8) is given by

$$\ln p_i(\mathbf{x}) = -\frac{1}{2} \left[\ln(2\pi)^N + \ln |\mathbf{K}_i^{(i)}| + \{\mathbf{x} - \mathbf{m}_i^{(i)}\}^T [\mathbf{K}_i^{(i)}]^{-1} \{\mathbf{x} - \mathbf{m}_i^{(i)}\} \right] \quad i=1,2 \quad (15)$$

and the natural logarithms of (9), (10), and (11) are similarly obtained. The logarithms of the conditional likelihood ratios are then used to obtain the terms on the right side of (1). The term given by

$$\lambda_A(\mathbf{x}_0) = \ln \frac{p_1(\mathbf{x}_0)}{p_2(\mathbf{x}_0)} = \ln p_1(\mathbf{x}_0) - \ln p_2(\mathbf{x}_0) \quad (16)$$

then becomes

$$\begin{aligned} \lambda_A(\mathbf{x}_0) = & -\frac{1}{2} \left[\ln |\mathbf{K}_1^{(1)}| + \{\mathbf{x} - \mathbf{m}_1^{(1)}\}^T [\mathbf{K}_1^{(1)}]^{-1} \{\mathbf{x} - \mathbf{m}_1^{(1)}\} \right] \\ & + \frac{1}{2} \left[\ln |\mathbf{K}_2^{(2)}| + \{\mathbf{x} - \mathbf{m}_2^{(2)}\}^T [\mathbf{K}_2^{(2)}]^{-1} \{\mathbf{x} - \mathbf{m}_2^{(2)}\} \right] \end{aligned} \quad (17)$$

Expanding (17) and collecting terms leads to the form

$$\lambda_A(\mathbf{x}_0) = \mathbf{x}^T \mathbf{A} \mathbf{x} + \mathbf{b}^T \mathbf{x} + c \quad (18)$$

where

$$\mathbf{A} = \frac{1}{2} \left[[\mathbf{K}_2^{(2)}]^{-1} - [\mathbf{K}_1^{(1)}]^{-1} \right] \quad (19)$$

is an $N \times N$ matrix,

$$\mathbf{b}^T = \left[\mathbf{m}_1^{(1)T} [\mathbf{K}_1^{(1)}]^{-1} - \mathbf{m}_2^{(2)T} [\mathbf{K}_2^{(2)}]^{-1} \right] \quad (20)$$

is a $1 \times N$ vector, and

$$c = \frac{1}{2} \left[\mathbf{m}_2^{(2)T} [\mathbf{K}_2^{(2)}]^{-1} \mathbf{m}_2^{(2)} - \mathbf{m}_1^{(1)T} [\mathbf{K}_1^{(1)}]^{-1} \mathbf{m}_1^{(1)} + \ln \frac{|\mathbf{K}_2^{(2)}|}{|\mathbf{K}_1^{(1)}|} \right] \quad (21)$$

is a scalar.

The coefficients of the conditional log likelihood ratio, $\lambda_A(\mathbf{x}_0)$ are called \mathbf{A} , \mathbf{b} , and c and are derived in the same way with (12), (13), and (14) substituted for the corresponding variables. Similar coefficients are calculated for $\lambda_B(\mathbf{y}_0)$ and

$\lambda_B(y_0)$ and are listed in Appendix A along with the coefficients for $\lambda_A(x_0)$ and $\lambda'_A(x_0)$. The computations of A , b , c , A' , b' , and c' are performed prior to their use in a real-time application and are input at the start of each process as the parameters for each of the quadratic classifiers.

III. THE TEST ENVIRONMENT

A. HARDWARE DESCRIPTION

The test environment for the distributed decision algorithms, designated Real-Time Cluster Star (RTC *), consists of a highly modular hardware base and a highly flexible operating system. The hardware consists of two clusters of single board computers (SBC's), each sharing a common backplane and an Ethernet local area network (LAN) serving as the communication link. Thus each cluster can be thought of as a node of a network and each node has multiple processors on a common bus.

1. The Cluster

The cluster configuration is diagramed in Figure 2. Each cluster consists of three SBC's physically connected by the MULTIBUS. Each SBC has 64K RAM of local memory and can access an additional 64K RAM board of shared memory and a 32K RAM board of common memory on the MULTIBUS. Also connected to the MULTIBUS are hard and floppy disk drives used for bootup and input/output operations.

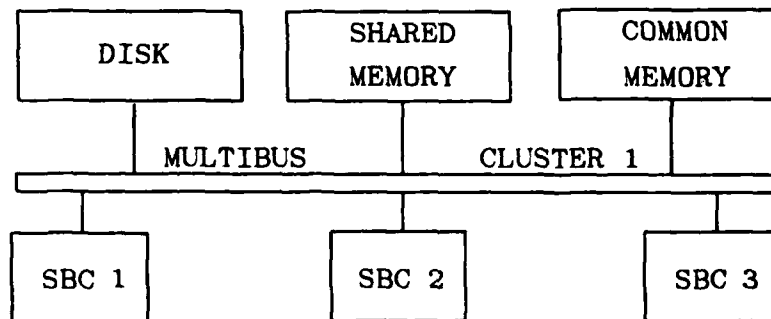


Figure 2 - Cluster Architecture

2. Real-Time Cluster Star (RTC *)

Figure 3 illustrates the RTC* architecture. It consists of two clusters connected by the Ethernet LAN. The Ethernet LAN/MULTIBUS interface is

the InterLAN NI3010 Ethernet Communications Controller Board (ECCB). This provides each cluster with its connection to the network. Further information on operating characteristics of the Ethernet LAN and RTC * use of the Ethernet LAN is available in [2,10].

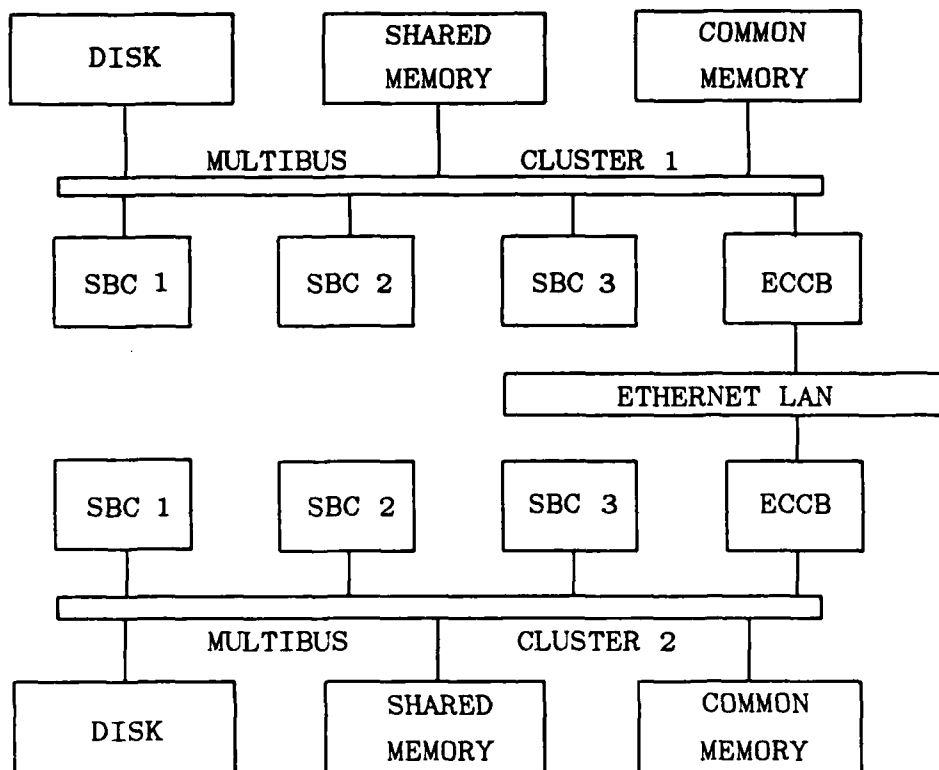


Figure 3 - Real-Time Cluster Star (RTC*) Architecture

B. THE OPERATING SYSTEM ENVIRONMENT

MCORTEX, the operating system, is a distributed multicomputer real-time executive. It allows for asynchronous operation of processes resident on SBC's in the same cluster and in separate clusters which are linked via the Ethernet LAN. System synchronization of computations in various distributed processes is accomplished using the synchronization model of Reed and Kanodia [11]. This section describes the MCORTEX system distribution of control variables, known

as **eventcounts** and **sequencers**. The modifications to the operating system, necessary to distribute user data throughout the system, are also discussed in this section.

1. The Synchronization Model

The MCORTEX operating system is based upon a synchronization model which is **event** oriented. Processes coordinate various activities by signaling and observing events using synchronization variables known as **eventcounts** and **sequencers**. An **eventcount** is a variable created by the user to signal the occurrence of an associated **event**. **Eventcounts** are initialized with the value zero and incremented by one each time the associated **event** occurs. The mechanism used to signal this occurrence is a call to a system primitive, the **advance**, which causes the **eventcount** to be incremented by one. A call to another system primitive, the **await**, causes a process to wait until the designated **eventcount** has reached a designated threshold. Once the **eventcount** value is equal to or greater than the threshold value the process may continue its execution. Therefore, processing at distributed locations may be controlled using **eventcounts** which are signaled and observed with the **advance** and the **await** primitives.

A **sequencer** is a variable provided by the system to control the allocation of a system shared resource. The **sequencer** is a positive integer number generator which starts with zero. It increments by one after providing its current value to any process which requests its associated shared resource. The **ticket** operation is the mechanism used to obtain a number from the **sequencer**. The number obtained is used as a threshold value in the **await** call to a system **eventcount** which is also associated with the shared resource. As users of the shared resource relinquish it, they increment the associated **eventcount** with the **advance**. This allows the user with the **ticket** value which matches the **eventcount** to gain access. An example of a shared resource controlled by a **sequencer** is the Ethernet LAN.

2. Eventcount Distribution

The kernel of MCORTEX is resident on each SBC and schedules processes for execution. A process runs until it invokes one of the system primitives, the **advance** or the **await**, which results in the actions described in Section B.1. The **advance** of an **eventcount**, which is used only within one cluster, causes an update of that clusters eventcount value. Processes in the same cluster, which are awaiting the eventcount, may then continue to execute. Update of eventcounts required for intercluster synchronization are **packetized** for transfer, via the Ethernet LAN, to the other cluster. The operating system procedure which accomplishes the transfer is located on SBC 1 of each cluster and is referred to in this thesis as the **driver**. The **driver** is the system software modified to allow for user data transfer between clusters.

3. Data Distribution

Data which must be shared between processes of the same cluster is made accessible through the use of pointers to access the local cluster **shared** memory locations. In the RTC* system, buffering of data must be done explicitly by user processes since no means of dynamic allocation presently exists. In this thesis, the real-time application requires the immediate use of the data generated, which precludes the need for buffering. Static storage locations, which are overwritten, are used for transfer of data throughout the system.

Data transfer from one cluster to another is accomplished by first establishing an absolute address in the local cluster **shared** memory to receive the data to be transferred. A pointer is used to access the absolute address in **shared** memory and the data value based at the pointer is updated. The system **driver** is then notified that a data value is ready for transfer. The Ethernet LAN sequencer provides the ticket to the user process for this data transfer. Once the ticket for this data value matches the eventcount associated with the Ethernet LAN, the data value is transferred to the driver's transmit data block in the appropriate data field in local cluster shared memory. The driver then causes the necessary calls to system subroutines to allow packetization and transfer over the Ethernet.

At the receiving end the message is processed by the local ECCB and the data is placed in the receive data block. The driver then stores the data value at the absolute address designated in the receiving clusters shared memory. Another pointer is then used in the receiving process to access the absolute address in *shared* memory. The data value based at the pointer is then available for further computations in this cluster. When the eventcount associated with this data transfer is updated via a similar procedure, the remaining computations are performed. User process eventcounts prevent the generation of additional data until the remaining computations in the present iteration are complete.

Appendix B provides an explanation of the steps necessary to create the system driver and user command files. The driver modifications required to transmit and receive data values for the distributed decision algorithms are shown in upper case lettering in the system procedure SYSDEV.PLI in Appendix C. User defined pointers and variable basing are shown and described further in the user procedures PA2, PA3, PB2, and PB3 in Appendix D.

C. ALGORITHM IMPLEMENTATION

Each cluster can be viewed as representing the set of local processors of a particular sensor which obtains large volumes of raw observation data from a target for initial processing. Decision rule parameters and raw observation data are read from local disk storage to the processes of two SBC's in a cluster. Two identical data sets are processed in parallel to generate a different reduced statistic in each processor. One statistic is to be used locally (at the same sensor) in further computation while the other is to be sent to the remote sensor for use in further computations. The local sensor then receives a reduced statistic from the remote sensor to combine with its locally retained statistic. The final result of the combined statistics is then compared to a decision threshold and the decision is displayed at a local sensor terminal.

1. Process Distributivity/Parallel Processing

The implementation of the decision rule described by (1) is accomplished with the following organization. The sensors associated with the two system

clusters, as well as the clusters themselves, are referred to as **SENSOR A** and **SENSOR B**. As illustrated in Figure 4, each sensor uses two processes labeled PA2(PB2) and PA3(PB3). Process computations take place in time order from left to right and computations shown above/below one another are performed in parallel.

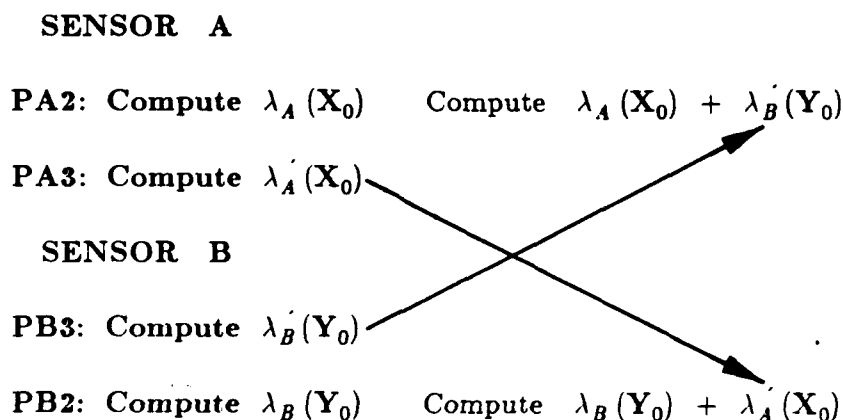


Figure 4 - Computations of Reduced Statistics

Processes PA2(PB2) and PA3(PB3) are resident on SBC 2 and SBC 3, respectively, at each sensor. Computations are performed as shown, with the primed statistics exchanged between sensors to allow further computations in processes PA2 and PB2. The detailed computations discussed in Chapter II are shown in user processes PA2, PA3, PB2, and PB3 of Appendix D.

2. Process Synchronization

Synchronization of events during the decision rule computations is crucial for accurate and meaningful results. As illustrated in Figure 5, the careful synchronization of time critical events is coordinated with the use of two distributed eventcounts at each sensor. The A1EVC eventcount of Sensor A is advanced to signal the availability of the statistic $\lambda'_A(x_0)$ for use in PB2 of Sensor B and the B1EVC eventcount of Sensor B signals PA2 of Sensor A that $\lambda'_B(y_0)$ is available. The A2EVC and B2EVC eventcounts control the timing of the next input operation at both sensors to ensure correct correspondence of the

observation data. In distributed processing multicomputer systems, it is essential that all threshold values used in the calls to the await primitives for comparison to the eventcounts, be initialized properly to ensure continued operation of the real-time system.

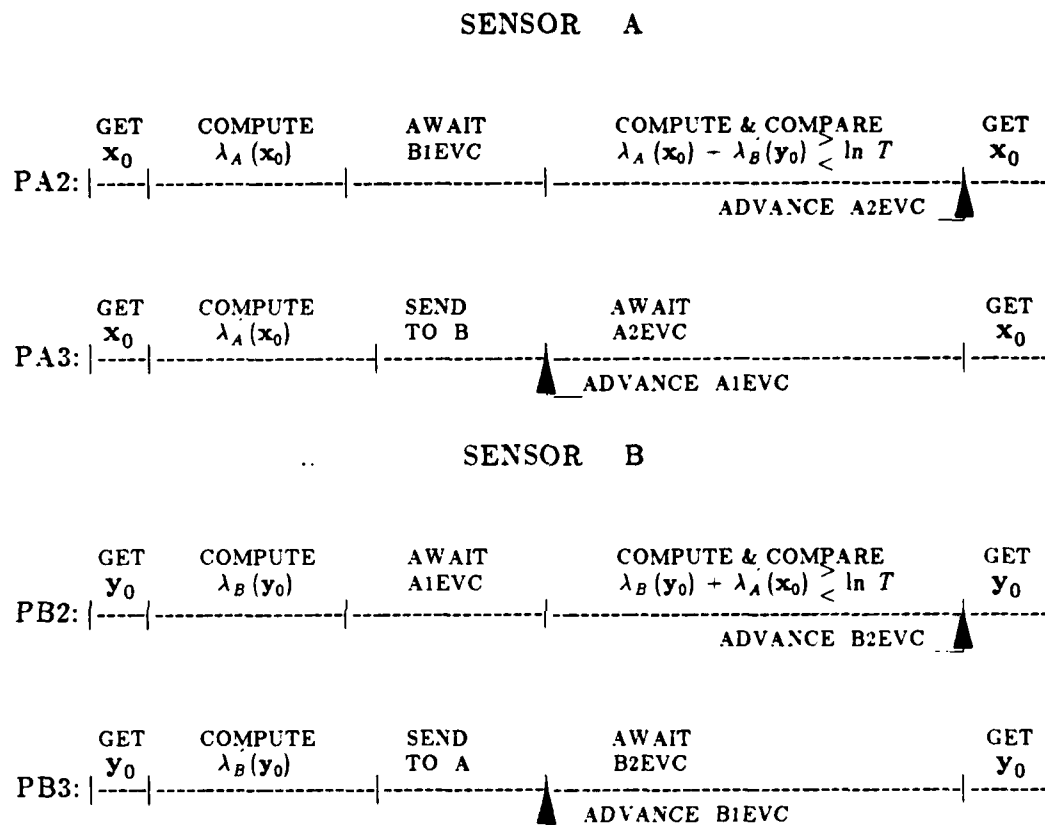


Figure 5 - Synchronization Diagram

As one might expect, there is a need to ensure that the required statistic, $\lambda_A(x_0)$ or $\lambda_B(y_0)$, is available for use prior to advancing the A1EVC or B1EVC eventcounts. This is insured by the forced synchronization of events inherent in the sensor to sensor transfer of user data and eventcount updates. The statistic to be transferred is stored in shared memory and transferred as described earlier. Once the Ethernet LAN sequencer ticket value is obtained for the data transfer and the request is placed in the ERB queue, the appropriate A1EVC or B1EVC

eventcount is advanced causing a system request for a ticket value from the same sequencer. This places the eventcount transfer request, which will signal the availability of data, behind the data transfer request in the same ERB queue. Therefore, when the eventcount is finally updated at the remote sensor the statistic required will be in place and available.

In the final stage of computation the reduced statistic retained locally and the statistic received from the remote sensor are added in processes PA2 and PB2 of each sensor and compared to a threshold (see Figure 2). The reduced statistics $\lambda_A(x_0)$ and $\lambda_B(y_0)$ are added and compared to the threshold at sensor A. Similarly, $\lambda_B(y_0)$ and $\lambda_A(x_0)$ are added and compared to the threshold at Sensor B. Results of the threshold decision are tabulated on the local consoles of each sensor and the loop begins again with the next observation vector read from disk. The processing of input observation vectors continues, simulating real-time operation until the vector files are depleted.

D. RESULTS OF THE SIMULATION

In the development and use of the test environment it was verified that it is important to distribute computation among processors to better utilize the available computational ability and minimize interprocess communication. Processes at each sensor were broken up and distributed among the available processors to gain increased computational advantages. Since processes at remote sensors had to be carefully synchronized, specific semaphore-like mechanisms were made available to provide this synchronization over the network. The specific mechanisms used in this implementation are the **await** and the **advance**. Correct operation of these synchronization mechanisms over the network depends on the prompt and orderly communication of protected variables used by the synchronization mechanisms. This orderly communication is achieved by the ticket operation. Successful implementation of a distributed decision algorithm requires the availability of all of these control mechanisms.

IV. CONCLUSIONS

The process of distributed decision making by two cooperating sensors observing a common phenomenon was introduced in this thesis. Decisions reached in this cooperative way produce more reliable results than those of sensors acting alone. Such decision procedures are characterized by the need to perform local computations at each sensor and to communicate partial results to the other sensor. Although several types of algorithms were cited to accomplish the desired distributed decision procedures, all have similar computation, communication, and process synchronization requirements.

A particular distributed decision algorithm based on the *generalized likelihood ratio test* was implemented to explore the computation, communication, and synchronization problems. The implementation was accomplished on a two node network connected via an Ethernet local area network. Each node of the network contained the required number of identical microprocessors sharing a common bus, shared memory, and network interfacing.

Problems of intercluster as well as intracluster synchronization of events between processes to ensure the timely input of observation data and the coordinated computation using the shared data from the opposite cluster were tested and resolved. Initial results using the *generalized likelihood ratio test* algorithm demonstrated the feasibility of performing the computations involved in the distributed decision algorithms in a realistic environment. The requirement for carefully designed, network-wide process control mechanisms was also found to be essential. The specific procedures used were discussed in the body of the thesis.

APPENDIX A

Quadratic Classifiers

Specific formulas for the quadratic classifiers, $\lambda_A(\mathbf{x}_0)$, $\lambda_A'(\mathbf{x}_0)$, $\lambda_B(\mathbf{y}_0)$, and $\lambda_B'(\mathbf{y}_0)$ described in Chapter II are provided in this appendix. Each quadratic classifier was derived similar to $\lambda_A(\mathbf{x}_0)$, in Chapter II, Section B. The coefficients, \mathbf{A} , \mathbf{b}^T , c , \mathbf{A}' , $\mathbf{b}^{T'}$, and c' , the necessary expanding equations for variables $\mathbf{K}_z^{(i)}|_y$, $\mathbf{m}_z^{(i)}|_y$, $\mathbf{K}_y^{(i)}|_z$, and $\mathbf{m}_y^{(i)}|_z$, and the estimates, $\hat{\mathbf{y}}_i$ and $\hat{\mathbf{x}}_i$, are given as functions of the known terms, $\mathbf{K}_z^{(i)}$, $\mathbf{m}_z^{(i)}$, $\mathbf{K}_y^{(i)}$, $\mathbf{m}_y^{(i)}$, and $\mathbf{B}_{zy}^{(i)}$.

The coefficients computed for

$$\lambda_A(\mathbf{x}_0) = \mathbf{x}^T \mathbf{A} \mathbf{x} + \mathbf{b}^T \mathbf{x} + c$$

are

$$\mathbf{A} = \frac{1}{2} \left[[\mathbf{K}_z^{(2)}]^{-1} - [\mathbf{K}_z^{(1)}]^{-1} \right]$$

$$\mathbf{b}^T = \left[[\mathbf{m}_z^{(1)}]^T [\mathbf{K}_z^{(1)}]^{-1} - [\mathbf{m}_z^{(2)}]^T [\mathbf{K}_z^{(2)}]^{-1} \right]$$

$$c = \frac{1}{2} \left[[\mathbf{m}_z^{(2)}]^T [\mathbf{K}_z^{(2)}]^{-1} \mathbf{m}_z^{(2)} - [\mathbf{m}_z^{(1)}]^T [\mathbf{K}_z^{(1)}]^{-1} \mathbf{m}_z^{(1)} + \ln \frac{|\mathbf{K}_z^{(2)}|}{|\mathbf{K}_z^{(1)}|} \right]$$

The coefficients computed for

$$\lambda_A'(\mathbf{x}_0) = \mathbf{x}^T \mathbf{A}' \mathbf{x} + \mathbf{b}^{T'} \mathbf{x} + c'$$

are

$$\mathbf{A}' = \frac{1}{2} \left[[\mathbf{K}_z^{(2)}|_y]^{-1} - [\mathbf{K}_z^{(1)}|_y]^{-1} \right]$$

$$\mathbf{b}^{T'} = \left[[\mathbf{m}_z^{(1)}|_y]^T [\mathbf{K}_z^{(1)}|_y]^{-1} - [\mathbf{m}_z^{(2)}|_y]^T [\mathbf{K}_z^{(2)}|_y]^{-1} \right]$$

$$c' = \frac{1}{2} \left[\mathbf{m}_z^{(2)|y}{}^T [\mathbf{K}_z^{(2)|y}]^{-1} \mathbf{m}_z^{(2)|y} - \mathbf{m}_z^{(1)|y}{}^T [\mathbf{K}_z^{(1)|y}]^{-1} \mathbf{m}_z^{(1)|y} + \ln \frac{|\mathbf{K}_z^{(2)|y}|}{|\mathbf{K}_z^{(1)|y}|} \right]$$

$$\mathbf{K}_z^{(1)|y} = \mathbf{K}_z^{(1)} - \mathbf{B}_{zy}^{(1)} [\mathbf{K}_y^{(1)}]^{-1} \mathbf{B}_{zy}^{(1)T}$$

$$\mathbf{K}_z^{(2)|y} = \mathbf{K}_z^{(2)} - \mathbf{B}_{zy}^{(2)} [\mathbf{K}_y^{(2)}]^{-1} \mathbf{B}_{zy}^{(2)T}$$

$$\mathbf{m}_z^{(1)|y} = \mathbf{m}_z^{(1)} + \mathbf{B}_{zy}^{(1)} [\mathbf{K}_y^{(1)}]^{-1} [\mathbf{y} - \mathbf{m}_y^{(1)}]$$

$$\mathbf{m}_z^{(2)|y} = \mathbf{m}_z^{(2)} + \mathbf{B}_{zy}^{(2)} [\mathbf{K}_y^{(2)}]^{-1} [\mathbf{y} - \mathbf{m}_y^{(2)}]$$

$$\hat{\mathbf{y}}_1 = \mathbf{m}_y^{(1)} + \mathbf{B}_{zy}^{(1)T} [\mathbf{K}_z^{(1)}]^{-1} [\mathbf{x} - \mathbf{m}_z^{(1)}]$$

$$\hat{\mathbf{y}}_2 = \mathbf{m}_y^{(2)} + \mathbf{B}_{zy}^{(2)T} [\mathbf{K}_z^{(2)}]^{-1} [\mathbf{x} - \mathbf{m}_z^{(2)}]$$

The coefficients computed for

$$\lambda_B(\mathbf{y}_0) = \mathbf{y}^T \mathbf{A} \mathbf{y} + \mathbf{b}^T \mathbf{y} + c$$

are

$$\mathbf{A} = \frac{1}{2} \left[[\mathbf{K}_y^{(2)}]^{-1} - [\mathbf{K}_y^{(1)}]^{-1} \right]$$

$$\mathbf{b}^T = \left[\mathbf{m}_y^{(1)T} [\mathbf{K}_y^{(1)}]^{-1} - \mathbf{m}_y^{(2)T} [\mathbf{K}_y^{(2)}]^{-1} \right]$$

$$c = \frac{1}{2} \left[\mathbf{m}_y^{(2)T} [\mathbf{K}_y^{(2)}]^{-1} \mathbf{m}_y^{(2)} - \mathbf{m}_y^{(1)T} [\mathbf{K}_y^{(1)}]^{-1} \mathbf{m}_y^{(1)} + \ln \frac{|\mathbf{K}_y^{(2)}|}{|\mathbf{K}_y^{(1)}|} \right]$$

The coefficients computed for

$$\lambda_B'(\mathbf{y}_0) = \mathbf{y}^T \mathbf{A}' \mathbf{y} + \mathbf{b}'^T \mathbf{y} + c'$$

are

$$\mathbf{A}' = \frac{1}{2} \left[[\mathbf{K}_y^{(2)}]^{-1} - [\mathbf{K}_y^{(1)}]^{-1} \right]$$

$$\mathbf{b}^T = \left[[\mathbf{m}_y^{(1)}|_z]^T [\mathbf{K}_y^{(1)}|_z]^{-1} - [\mathbf{m}_y^{(2)}|_z]^T [\mathbf{K}_y^{(2)}|_z]^{-1} \right]$$

$$c = \frac{1}{2} \left[[\mathbf{m}_y^{(2)}|_z]^T [\mathbf{K}_y^{(2)}|_z]^{-1} \mathbf{m}_y^{(2)}|_z - [\mathbf{m}_y^{(1)}|_z]^T [\mathbf{K}_y^{(1)}|_z]^{-1} \mathbf{m}_y^{(1)}|_z + \ln \frac{|\mathbf{K}_y^{(2)}|_z|}{|\mathbf{K}_y^{(1)}|_z|} \right]$$

$$\mathbf{K}_y^{(1)}|_z = \mathbf{K}_y^{(1)} - \mathbf{B}_{xy}^{(1)T} [\mathbf{K}_z^{(1)}]^{-1} \mathbf{B}_{xy}^{(1)}$$

$$\mathbf{K}_y^{(2)}|_z = \mathbf{K}_y^{(2)} - \mathbf{B}_{xy}^{(2)T} [\mathbf{K}_z^{(2)}]^{-1} \mathbf{B}_{xy}^{(2)}$$

$$\mathbf{m}_y^{(1)}|_z = \mathbf{m}_y^{(1)} + \mathbf{B}_{xy}^{(1)T} [\mathbf{K}_z^{(1)}]^{-1} [\mathbf{x} - \mathbf{m}_z^{(1)}]$$

$$\mathbf{m}_y^{(2)}|_z = \mathbf{m}_y^{(2)} + \mathbf{B}_{xy}^{(2)T} [\mathbf{K}_z^{(2)}]^{-1} [\mathbf{x} - \mathbf{m}_z^{(2)}]$$

$$\hat{\mathbf{x}}_1 = \mathbf{m}_z^{(1)} + \mathbf{B}_{xy}^{(1)} [\mathbf{K}_y^{(1)}]^{-1} [\mathbf{y} - \mathbf{m}_y^{(1)}]$$

$$\hat{\mathbf{x}}_2 = \mathbf{m}_z^{(2)} + \mathbf{B}_{xy}^{(2)} [\mathbf{K}_y^{(2)}]^{-1} [\mathbf{y} - \mathbf{m}_y^{(2)}]$$

APPENDIX B

LINK86 Input Option Files

When linking files to create a command file for use on each SBC, the following command is invoked with the appropriate user filename: LINK86 filename [I]. The "I" in square brackets invokes the input file option which directs LINK86 to obtain further command line input from the designated input file. As an example, the modules listed in CA.INP are linked with the command: LINK86 CA[I], where the "I" indicates that CA.INP contains the names of the files to be linked. The name preceding the equal sign is the filename assigned to the command file. LINK86 CA[I] produces the command file CA.CMD, which is the system driver for Sensor A (cluster A). All files listed in the input file must be on the logon disk and must be of type object (.obj). Object files are generated by compiling files of type PLI (.pli) or A86 (.a86). The above steps also apply for linking the system driver files for Sensor B as well as the user files, processes PA2, PA3, PB2, and PB3, to create the respective command files CB.CMD, NUM12.CMD, NUM13.CMD, NUM22.CMD, and NUM23.CMD.

```
*****
*****
**
** CA.INP input option file **
**
*****
```

```
CA =
SYSINIT [CODE[AB[439]],DATA[AB[800],M[0],AD[82]],MAP[ALL]],
SYSDEV,
ASMROUT,
GATEMOD
```

```
*****
*****
**
** NUM12.INP input option file **
**
*****
```

```
NUM12=
SBC2INIT [CODE[AB[439]],DATA[AB[300],M[0],AD[92]],MAP[ALL]],
PA2,
GATEMOD
```

```
*****
*****
**
** NUM13.INP input option file **
**
*****
```

```
NUM13=
SBC3INIT [CODE[AB[439]],DATA[AB[220],M[0],AD[92]],MAP[ALL]],
PA3,
GATEMOD
```

```
*****
*****
**                               CB.INP input option file                               **
*****
*****
```

```
CB =
SYSINITB [CODE[AB[439]],DATA[AB[800],M[0],AD[82]],MAP[ALL]].
SYSDEV,
ASMROUT,
GATEMOD
```

```
*****
*****
**                               NUM22.INP input option file                               **
*****
*****
```

```
NUM22=
SEC2INIT [CODE[AB[439]],DATA[AB[800],M[0],AD[82]],MAP[ALL]].
PR2,
GATEMOD
```

```
*****
*****
**                               NUM23.INP input option file                               **
*****
*****
```

```
NUM23=
SEC3INIT [CODE[AB[439]],DATA[AB[800],M[0],AD[82]],MAP[ALL]].
PR3,
GATEMOD
```

APPENDIX C

Device Driver and Packet Processor

Source Code

This code consists of PL/I-86 and 8086 assembly language modules. When linked as described in Appendix A and loaded in local memory of SBC #1 of each cluster, the driver handles the systemwide distribution of user data and eventcounts via the local area network.

Initialization modules (SYSINITA & SYSINITB), each for their own cluster, define cluster addresses, create user eventcounts, establish eventcount distribution, and create the procedure space, under operating system control, for the driver. SYSDEV. The system definitions file, SYSDEF and the file NI3010.DCL are required when compiling SYSDEV. Any user eventcounts, sequencers, or shared variable pointers which are defined in SYSDEF must be updated when these items change with new synchronization and control schemes.

SYSINITA and SYSINITB must also be updated whenever changes are made to user eventcounts or their distribution. Recompilation and relinking are also necessary to produce the updated command files CA.CMD and CB.CMD.

```

*****
*****
**
** CLUSTER A INITIALIZATION MODULE SYSINITA.PLI **
**
*****
*****

```

```

SYSINITA: proc options (main);
    %include 'sysdef.pli';

    %replace

    EVC_TYPE          by '00'b4;

    /* main */

    call define_cluster ('0001'b4); /* must be called
                                     prior to creating
                                     evc's */

    /**** USER ****/

    CALL CREATE_EVC (A1EVC);
    CALL CREATE_EVC (A2EVC);
        CALL CREATE_EVC (B1EVC);

    /*** SYSTEM ***/

    call create_evc ('FBF_READ');
    call create_evc ('ERF_WRITE');
    call create_seq ('ERF_WRITE_REQUEST');

    /* distrib. map called after eventcounts have
       been created */

    /* local and remote copy of A1EVC needed */

    call distribution_map (EVC_TYPE, A1EVC, '0003'b4);

    call create_proc ('fc'b4, '90'b4,
                    '2941'b4, '0270'b4, '0053'b4,
                    '0439'b4, '0800'b4, '0200'b4);

    call await ('fe'b4, 'c1'b4);

END SYSINITA;

```



```

*****
*****
**
** CLUSTER B INITIALIZATION MODULE SYSINITB.PLI **
**
*****
*****

```

```

SYSINITB: proc options (main);
    %include 'sysdef.pli';
    %replace
EVC_TYPE          by '00'b4;
    /* main */
call define_cluster ('0002'b4); /* must be called
                                prior to creating
                                evc's */
    /**** USER *****/
CALL CREATE_EVC (A1EVC);
CALL CREATE_EVC (B1EVC);
    CALL CREATE_EVC (B2EVC);
    /**** SYSTEM *****/
call create_evc (PRB_READ);
call create_evc (PRB_WRITE);
call create_seq (PRB_WRITE_REQUEST);
    /* distrib. map called after eventcounts have
    been created */
    /* local and remote copy of B1EVC needed */
call distribution_map (EVC_TYPE, B1EVC, '0003'b4);
call create_proc ('fc'b4, '80'b4,
                 '0941'b4, '0800'b4, '0053'b4,
                 '0439'b4, '0800'b4, '0600'b4);
call await ('fe'b4, '01'b4);
END SYSINITB;

```

```

/*****
/*****
/** FILE SYSDEF.PLI MARK A. SCHON 24 JUL 85 **/
/**=====**/
/** This section of code is given as a PLI file to be **/
/** %INCLUDE'd with SYSDEV.PLI. ENTRY declarations are **/
/** made for all available MCORTEX functions. **/
/*****
/*****

```

DECLARE

```

advance ENTRY (BIT (8)),
/* advance (event_count_id) */
await ENTRY (BIT (8), BIT (16)),
/* await (event_count_id, awaited_value) */
create_evt ENTRY (BIT (8)),
/* create_evt (event_count_id) */

create_proc ENTRY (BIT (8), BIT (8),
                  BIT (16), BIT (16), BIT (16),
                  BIT (16), BIT (16), BIT (16)),
/* create_proc (processor_id, processor_priority, */
/* stack_pointer_highest, stack_seg, ip */
/* code_seg, data_seg, extra_seg) */

create_seq ENTRY (BIT (8)),
/* create_seq (sequence_id) */

preempt ENTRY (BIT (8)),
/* preempt (processor_id) */

read ENTRY (BIT (8)) RETURNS (BIT (16)),
/* read (event_count_id) */
/* RETURNS current_event_count */

ticket ENTRY (BIT (8)) RETURNS (BIT (16)),
/* ticket (sequence_id) */
/* RETURNS unique_ticket_value */

define_cluster ENTRY (bit (16)),
/* define_cluster (local_cluster_address) */

distribution_map ENTRY (bit (2), bit (2), bit (16)),
/* distribution_map (distribution_type, id, cluster_addr) */

add2bit16 ENTRY (BIT(16), BIT(16)) RETURNS (BIT (16));
/* add2bit16 ( a_16bit_#, another_16bit_#) */
/* RETURNS a_16bit_# + another_16bit_# */

```

%replace

/*-----

*** EVCSID's ***

(1) USER */

A1EVC BY '01'B4,

A2EVC BY '02'B4,

B1EVC BY '03'B4,

B2EVC BY '04'B4,

/* (2) SYSTEM */

ERB_READ by 'fc'b4,
ERB_WRITE by 'fd'b4.

/*-----

*** SEQUENCER NAMES ***

(1) USER

USER PROCESSES USE ERB_WRITE_REQUEST ONLY.

(2) SYSTEM */

ERB_WRITE_REQUEST by 'ff'b4,

/*-----

*** SHARED VARIABLE POINTERS ***

(1) USER */

PB BY '2000'B4,
PC BY '8DD0'B4,

/* (2) SYSTEM */

block_ptr_value by '8000'b4,

xmit_ptr_value by '8008'b4,

rcv_ptr_value by '8666'b4,

END_RESERVE by 'FFFF'b4;

```

*****
*****
**
**
**
**
*****
*****

```

NI3010.DCL FILE

```
%replace
```

```
/* I/O port addresses
```

These values are specific to the use of the INTERLAN NI3010 MULTIBUS to ETHERNET interface board. Any change to the I/O port address of '00b0' hex (done so with a DIP switch) will require a change to these addresses to reflect that change. */

```

command_register      by 'b0'b4,
command_status_register  by 'b1'b4,
transmit_data_register  by 'b2'b4,
interrupt_status_reg   by 'b5'b4,
interrupt_enable_register  by 'b9'b4,
high_byte_count_reg    by 'bc'b4,
low_byte_count_reg     by 'bd'b4,

```

```
/* end of I/O port addresses */
```

```

/* Interrupt enable status register values */
disable_ni3010_interrupts  by '00'b4,
ni3010_intrpts_disabled    by '00'b4,
receive_block_available     by '04'b4,
transmit_dma_done           by '06'b4,
receive_dma_done            by '07'b4,

```

```
/* end register values */
```

```
/* Command Function Codes */
```

```

module_interface_loopback  by '01'b4,
internal_loopback          by '02'b4,
clear_loopback             by '03'b4,
go_offline                 by '08'b4,
go_online                  by '09'b4,
    onboard_diagnostic     by '0a'b4,
clr_insert_source         by '0e'b4,
load_transmit_data        by '2e'b4,
load_and_send             by '29'b4,
load_group_addresses      by '2a'b4,
reset                    by '3f'b4,

```

```
/* end Command Function Codes */
```

```

*****
*****
**
** CLUSTER A - ADDRESS.DAT FILE - USED BY SYSDEV.PLI **
** - 1ST THREE VALUES USED IN SUBROUTINE **
** program_group_addresses **
** - LAST TWO USED IN MAIN PROGRAM SYSDEV **
** TO IDENTIFY THE LOCAL CLUSTER ADDRESS. **
**
*****
*****

```

```

1.
'00000000'b, '00000001'b,
'00000000'b, '00000001'b

```

```

*****
*****
**
** CLUSTER B - ADDRESS.DAT FILE - USED BY SYSDEV.PLI **
** - 1ST THREE VALUES USED IN SUBROUTINE **
** program_group_addresses **
** - LAST TWO USED IN MAIN PROGRAM SYSDEV **
** TO IDENTIFY THE LOCAL CLUSTER ADDRESS. **
**
*****
*****

```

```

1.
'00000000'b, '00000010'b,
'00000000'b, '00000010'b

```

SYSDEV: PROCEDURE;

```

/* Date: 24 JULY 1985
Programmer: MARK A. SCHON (MODIFIED CODE FROM
PREVIOUS THESIS[2])

```

Module Function: To serve as the Ethernet Communication Controller Board, ECCB (NI3010) device handler. This process is scheduled under MCORTEX and consumes Ethernet Requests Packets (ERP) generated by the SYSTEMSIC located in LEVEL2.SPC & by USER PROGRAMS. It also processes any inbound packets by analyzing the packet contents and making the appropriate MCORTEX calls. */

%replace

erc_type by '00'b4,
erb_block_len by 20,
erb_block_len_m1 by 19,
infinity by 32767;

%include 'sysdef.pli';

DECLARE

```
1   erb(0:erb_block_len_m1) based (block_ptr),
    2   command            bit (8),
    2   type_name         bit (8),
    2   name_value         bit (16),
    2   remote_addr        bit (16),

    1   transmit_data_block based (xmit_ptr),
    2   destination_address_a
         bit (8) ,
    2   destination_address_b
         bit (8),
    2   destination_address_c
         bit (8),
         2   destination_address_d
         bit (8),
    2   destination_address_e
         bit (8) ,
    2   destination_address_f
         bit (8) ,
    2   source_address_a
         bit (8) ,
    2   source_address_b
         bit (8),
    2   source_address_c
         bit (8),
         2   source_address_d
         bit (8),
    2   source_address_e
         bit (8) ,
    2   source_address_f
         bit (8) ,
         2   type_field_a
         bit (8) ,
         2   type_field_b
         bit (8),
    2   data (4) bit (8),
```

```

                2 USER_DATA (12) FLOAT,
                (TX_DATA_PTR,XMIT_PTR) POINTER,
/* HIGH MEMORY ADDRESSES OF TX_DATA_PTR AND XMIT_PTR      */
/*                               ASSIGNED IN SYSDEV        */

```

```
DATA_TO_SEND FLOAT BASED(TX_DATA_PTR),
```

```
1 receive_data_block based (rcv_ptr),
```

```

                2 frame_status          bit (8),
                2 null_byte             bit (8) ,
                2 frame_length_lsb     bit (8) ,
                2 frame_length_msb     bit (8) ,
                2 destination_address_a bit (8) ,
                2 destination_address_b bit (8) ,
                2 destination_address_c bit (8) ,
                2 destination_address_d bit (8) ,
                2 destination_address_e bit (8) ,
                2 destination_address_f bit (8) ,
                2 source_address_a      bit (8) ,
                2 source_address_b      bit (8) ,
                2 source_address_c      bit (8) ,
                2 source_address_d      bit (8) ,
                2 source_address_e      bit (8) ,
                2 source_address_f      bit (8) ,
                2 type_field_a          bit (8) ,
                2 type_field_b          bit (8) ,
                2 data(4)               bit (8) ,
                2 USER_DATA (12)       FLOAT,
                2 crc_msb               bit (8) ,
                2 crc_upper_middle_byte bit (8) ,
                2 crc_lower_middle_byte bit (8) ,
                2 crc_lsb               bit (8) ,

```

```
(RX_DATA_PTR,RCV_PTR,BLOCK_PTR) POINTER,
```

```

/* HIGH MEMORY ADDRESSES OF RX_DATA_PTR,RCV_PTR,&        */
/*                               BLOCK_PTR ARE ASSIGNED IN FILE SYSTEM */

```

```
DATA_ARRIVED FLOAT BASED(RX_DATA_PTR),
```

```

                index fixed bin (15),
                (addr_e, addr_f) bit (8),
                address file,
                copy_ie_register bit (8),
                (cluster_addr,erb_write_value,i) bit (16),
                (j,k) fixed bin (15),
                reg_value bit (8) ,
                write_io_port entry (bit (8), bit (8)).

```

```

        read_io_port  entry (bit (8), bit (8)),

        initialize_cpu_interrupts  entry,
        enable_cpu_interrupts     entry,
        disable_cpu_interrupts    entry,
        write_bar entry (bit(16));

        /* end module listing */

%replace

/* codes specific to the Intel 8259a Programmable
   Interrupt Controller (PIC) */

        icw1_port_address  by 'c2'b4,
        icw2_port_address  by 'c2'b4,
        icw4_port_address  by 'c2'b4,
        ocw_port_address   by 'c2'b4,

        /* note: icw ==> initialization control word
           ocw ==> operational command word */

        icw1  by '13'b4,

/* single PIC configuration, edge triggered input */

        icw2  by '40'b4,

/* most significant bits of vectoring byte; for an
   interrupt 5, the effective address will be
   (icw2 + interrupt #) * 4 which will be
   (40 hex + 5) * 4 =114 hex */

        icw4  by '0f'b4,

/* automatic end of interrupt and buffered mode/master */

        ocw1  by '3f'b4;

        /* unmask interrupt 4 (bit 4), */
        /* interrupt 5 (bit 5), and */
*/

        /* interrupt 6 (bit 6), mask all others */

        /* end 8259a codes */

/* include constants specific to the NI3010 board */
#include 'ni3010.dcl';

```



```

/*****
                                     */
                                     /* Main Body */

call write_io_port(interrupt_enable_register,
                   disable_ni3010_interrupts);
call initialize_pic;
call initialize_cpu_interrupts;
call read_io_port (command_status_register, reg_value);
call perform_command (reset);

call program_group_addresses;
/* assignments to the source and destination address
   fields that will not change */

call perform_command (clr_insert_source);
/* NI3010 performance is enhanced in this mode */

/*****
                                     */
/* ASSIGN POINTER VALUES, PREVIOUSLY DEFINED -FILE SYSDEF */
/* TX_DATA_PTR <-- PP = 8CC0 A FLOAT BLOCK OF 4-BYTES */
/* RX_DATA_PTR <-- PC = 8DD0 A FLOAT BLOCK OF 4-BYTES */
/*
                                     */
/* BLOCK_PTR <-- BLOCK_PTR_VALUE = 8000 THE ECB 120-BYTES */
/* RCV_PTR <-- RCV_PTR_VALUE = 8666 THE RDB 66-BYTES */
/* XMIT_PTR <-- XMIT_PTR_VALUE = 82C8 THE TDB 74-BYTES */
/*****
                                     */

UNSPEC(TX_DATA_PTR) = PP;
UNSPEC(RX_DATA_PTR) = PC;
unspec(block_ptr) = block_ptr_value;
unspec(rcv_ptr) = rcv_ptr_value;
unspec(xmit_ptr) = xmit_ptr_value;

/* make one time assignments to transmit data block */

transmit_data_block.destination_address_a = '03'b4;
transmit_data_block.destination_address_b = '22'b4;
transmit_data_block.destination_address_c = '07'b4;
transmit_data_block.destination_address_d = '0C'b4;
transmit_data_block.source_address_a = '03'b4;
transmit_data_block.source_address_b = '32'b4;
transmit_data_block.source_address_c = '27'b4;
transmit_data_block.source_address_d = '32'b4;

/* get the local cluster address - file was
   opened in proc program_group_addresses */

get file (address) list (addr_e, addr_f);
transmit_data_block.source_address_e = addr_e;
transmit_data_block.source_address_f = addr_f;

```

```

cluster_addr = addr e || addr f;
put skip (2) edit ('*** CLUSTER ',cluster_addr,
                  ' Initialization Complete ***')
                  (col(15),a,b4(4),a);
i = '0001'b4;
call perform_command (go_online);

/* at this point copy_ie_reg = RBA , but
   ie_reg on NI3010 is actually disabled */
call disable_cpu_interrupts;

do k = 1 to infinity;
  /* note: interrupt not allowed during a
     call to MCCORTEX primitive */

  erb_write_value = read(ERB_WRITE);
  /* In the MXTRACE version of the RTOS
     all primitive calls clear and
     set interrupts (diagnostic message
     routines), so the NI3010 interrupts
     must be disabled on entry to MXTRACE */
  do while (erb_write_value < 1);
  /* busy waiting */
    erb_write_value = read(ERB_WRITE);
  copy_ie_register=receive_block_available;
  call write_io_port(interrupt_enable_register,
                    receive_block_available);
    call enable_cpu_interrupts;
  /* if a packet has been received, this
     is when an interrupt may occur - can
     see that outbound packets are always
     favored. */
  do j = 1 to 1000;
    /* interrupt window for packets received */
  end; /* do j */
  call disable_cpu_interrupts;
  if (copy_ie_register = receive_dma_done) then
  do;
    /* receive DMA operation started, so let
       finish. */
    call enable_cpu_interrupts;
    do while (copy_ie_register = receive_dma_done);
    end;
    call disable_cpu_interrupts;
  end; /* ift */

  copy_ie_register = disable_ni3010_interrupts;
  call write_io_port(interrupt_enable_register,
                    disable_ni3010_interrupts);

end; /* busy */

```

```

/* ERB has an ERP in it, so process it */
/* no external interrupts (RBA) until
   the ERP is consumed and the packet
   gets sent */
index = mod((fixed(i) - 1), erb_block_len);
        /* 32k limit on parameter to fixed fcn. */

transmit_data_block.data(1) = erb(index).command;
transmit_data_block.data(2) = erb(index).type_name;

transmit_data_block.data(3) =
        substr(erb(index).name_value,9,8);

transmit_data_block.data(4) =
        substr(erb(index).name_value,1,8);

IF (ERB(INDEX).COMMAND = 1) THEN DO;
        TRANSMIT_DATA_BLOCK.USER_DATA(1) = DATA_TO_SEND;
END;

transmit_data_block.destination_address_e=
        substr(erb(index).remote_addr, 1,8);

transmit_data_block.destination_address_f=
        substr(erb(index).remote_addr, 9,8);

call advance (ERP_PPAD); /* caution here !!!!
   an ADVANCE will result in a
   call to VP$SCHEDULER, which
   will set CPU interrupts on exit.
   It's the reason NI3010 interrupts
   are disabled first in the
   Do While loop above. */

/* packet ready to go, so send it */
call transmit_packet;

/* copy_ie_register = RBA , but not actual register */
call disable_cpu_interrupts;

/* setting up for next ERP consumption */
i = add2bit16(i, '0001'b4);
end; /* do forever */

/* end main body */

```

```
/******
```

```
initialize_pic:      procedure;
```

```
    DECLARE
```

```
        write_io_port entry (bit (8) , bit(8));
```

```
    call write_io_port (icw1_port_address,icw1);  
    call write_io_port (icw2_port_address,icw2);  
    call write_io_port (icw4_port_address,icw4);  
    call write_io_port (ocw_port_address,ocw1);
```

```
end initialize_pic;
```

```
/******
```

```
perform_command:    procedure (command);
```

```
    DECLARE
```

```
        command bit (8) ,  
        reg_value bit (8) ,  
        srf bit (8) ,  
        write_io_port entry (bit (8) ,bit (8) ) ,  
        read_io_port entry (bit (8) ,bit (8) );
```

```
    /* end declarations */
```

```
    srf = '0'b4;  
    call write_io_port (command_register,command);  
    do while ((srf & '01'b4) = '00'b4);  
        call read_io_port (interrupt_status_reg,srf);  
    end; /* do while */  
        call read_io_port  
        (command_status_register,reg_value);  
    if (reg_value > '01'b4) then  
    do;  
        /* not (SUCCESS or SUCCESS with Retries) */  
  
        put skip edit ('*** ETHERNET Board Failure ***'  
            (col(20),a);  
            /* when this occurs, run the diagnostic  
            routine T3010/Cx, where x is the  
            current cluster number */
```

```
        stop;  
    end; /* itd */
```

```
end perform_command;
```

```
/***/
```

```
transmit_packet: procedure external;
```

```
DECLARE
```

```
  srf bit (8) ,  
  reg_value bit (8) ,  
  write_io_port entry (bit (8) ,bit (8) ),  
  read_io_port entry (bit (8) ,bit (8) ),  
  enable_cpu_interrupts entry,  
  disable_cpu_interrupts entry,  
  write_bar entry (bit(16));
```

```
  /* begin */
```

```
  srf = '0'b4;
```

```
  call write_bar (xmit_ptr value);  
  call write_io_port(high_byte_count_reg,'20'b4);  
  call write_io_port(low_byte_count_reg,'3c'b4);  
  copy_ie_register = transmit_dma_done;  
  call write_io_port(interrupt_enable_register,  
                    transmit_ima_done);  
  
  call enable_cpu_interrupts;  
  do while (copy_ie_register = transmit_dma_done);  
  end; /* loop until the interrupt handler  
        takes care of the TDD interrupt -  
        it sets copy_ie_register = RBA */  
  call perform_command (load_and_send);
```

```
end transmit_packet;
```

```
/***/  
*****/
```

```
HL_interrupt_handler: procedure external;
```

```
/* This routine is called from the low level  
S086 assembly language interrupt routine */
```

```
DECLARE
```

```
  write_io_port entry (bit (8) ,bit (8) ),  
  read_io_port entry (bit (8) ,bit (8) ),  
  enable_cpu_interrupts entry,  
  disable_cpu_interrupts entry,  
  write_bar entry (bit(16));
```

```

/* begin */

call write_io_port(interrupt_enable_register,
                   disable_ni3010_interrupts);

if (copy_ie_register = receive_block_available)
then do;

    call write_bar(rcv_ptr_value);
    call write_io_port(high_byte_count_reg, '05'b4);
    call write_io_port(low_byte_count_reg, 'f2'b4);

    /* initiate receive DMA */

    copy_ie_register = receive_dma_done;
    call write_io_port(interrupt_enable_register,
                       receive_dma_done);

end; /* do */
else
if (copy_ie_register = receive_dma_done) then
do;
    call process_packet;
    copy_ie_register = receive_block_available;
    call write_io_port(interrupt_enable_register,
                       receive_block_available);
end; /* if then do */
else
if (copy_ie_register = transmit_dma_done)
then do;

    copy_ie_register = receive_block_available;
    /* NI3010 interrupts disabled on entry */
end; /* if then do */

end HL_interrupt_handler;

/*****
process_packet: procedure;

DECLARE

    DATA_ARRIVED FLOAT BASED(RX_DATA_PTR),

    RX_DATA_PTR POINTER,

    local_evc_value bit (16),
    data_ptr pointer,
    remote_evc_value bit (16) based (data_ptr);

```

```

if (receive_data_block.data(1) = evc_type) then
do;
    data_ptr = addr(receive_data_block.data(3));
    /* remote_evc_value now has a value */
    local_evc_value = read(receive_data_block.data(2));
    do while (local_evc_value < remote_evc_value);
        call advance (receive_data_block.data(2));
        local_evc_value = add2bit16(local_evc_value,
                                     '0001'b4);
    end;
    call disable_cpu_interrupts;
    /* this must be done due to setting of
       cpu interrupts by calls to MCO RTEX's
       VPSCHEDULER via ADVANCE */
    end; /* itd */
/* IF DATA IS IN THIS RDB THEN TRANSFER IT TO USER HIGH */
/* MEMORY */
    ELSE DO;
        UNSPEC(RX_DATA_PTR) = PC;
        DATA_ARRIVED = RECEIVE_DATA_BLOCK.USER_DATA(1);
    END;
end process_packet;
/*****/
program_group_addresses: procedure;
    DECLARE
        1 group_addr(40) based (group_ptr),
        2 mc_group_field_a
            bit (8),
        2 mc_group_field_b
            bit (8),
        2 mc_group_field_c
            bit (8),
        2 mc_group_field_d
            bit (8),

```

```

2 mc_group_field_e
  bit (8),
2 mc_group_field_f
  bit (8);

```

```

DECLARE

```

```

(group_ptr,p) pointer,
(field_e, field_f) bit (8),
bit_8_groups bit (8) based (p),
(i,num_groups,groups_times_6) fixed bin (7);

```

```

unspec(group_ptr) = xmit_ptr_value;
open file (address) stream input;
get file (address) list (num_groups);
do i = 1 to num_groups;

```

```

  group_addr(i).mc_group_field_a = '03'b4;
  group_addr(i).mc_group_field_b = '00'b4;
  group_addr(i).mc_group_field_c = '00'b4;
  group_addr(i).mc_group_field_d = '00'b4;
  get file (address) list (field_e,field_f);
  group_addr(i).mc_group_field_e = field_e;
  group_addr(i).mc_group_field_f = field_f;

```

```

end; /* do i */

```

```

call disable_cpu_interrupts;
call write_bar (xmit_ptr_value);
call write_io_port(high_byte_count_reg, '00'b4);
groups_times_6 = 6 * num_groups;
p = addr (groups_times_6);
call write_io_port(low_byte_count_reg, bit_8_groups);
copy_ie_register = transmit_dma_done;
call write_io_port(interrupt_enable_register,
  transmit_dma_done);
call enable_cpu_interrupts;

```

```

do while (copy_ie_register = transmit_dma_done);
end; /* loop until the interrupt handler
  takes care of the TTD interrupt -
  it sets COPY_IF_REG = FBA */

```

```

call perform_command(load_group_addresses);

```

```

end program_group_addresses;

```

```

/*****/

```

```

end; /* system device handler and packet processor */

```



```

*****
*****
**                               ASMROUT.A86 FILE                               **
*****
*****

```

```
extrn hl_interrupt_handler : far
```

```

public write_io_port
public read_io_port
public write_bar
public initialize_cpu_interrupts
public enable_cpu_interrupts
public disable_cpu_interrupts
;*****

```

```
write_io_port:
```

```
; Parameter Passing Specification:
```

```

;                               entry                               exit
;
; parameter 1    <port address>          <unchanged>
;
; parameter 2    <value to be outputted> <unchanged>
;
;

```

```
dseg
```

```
port_address  rb  1
```

```
cseg
```

```

push bx! push si! push dx! push ax
mov  si, [bx]
mov  al, [si]
mov  port_address, al
mov  si, 2[bx]
mov  al, [si]
mov  dl, port_address
mov  dh, 00h
out  dx, al
pop  ax! pop dx! pop si! pop bx
ret

```

```
;*****
```

read_io_port:

```
; Parameter Passing Specification
;
;                               entry                exit
; parameter 1    <port address>    <unchanged>
; parameter 2    <meaningless>     <register value>
```

```
    cseg
push  bx! push si! push dx! push ax
    mov  si, [bx]
    mov  al, [si]
    mov  port_address, al
mov   si, 2[bx]
mov   dl, port_address
    mov  dh, 00h
    in   al, dx
    mov  [si], al
pop   ax! pop dx! pop si! pop bx!
    ret
```

write_bar:

```
; Parameter Passing Specification
;
; parameter 1 (and only): the address of the data block
;                               to be transmitted or received.
```

```
    dseg

    e_bar_port    equ 0b9h
    h_bar_port    equ 0bah
    l_bar_port    equ 0bth
    temp_e_byte   rb 1
    temp_es       rw 1
```

cseg

```
; This module computes a 24 bit address from a 32 bit
; address - actually a combination of the ES register
; and the IP passed via a parameter list.
```

```
    push bx! push ax! push cx! push es! push dx! push si
```

```
    mov  dx, 0800h    ; shared memory segment
```

```

mov     es, dx
mov     temp_es, es
mov     dx, es
        mov     si, [bx]
        mov     ax, [si]
        mov     cl, 12
        shr     dx, cl
        mov     temp_e_byte, dl
mov     dx, temp_es
        mov     cl, 4
        shl     dx, cl
        add     ax, dx
        jnc     no_add
add_1:  inc     temp_e_byte
no_add: out     l_bar_port, al
        mov     al, ah
        out     h_bar_port, al
        mov     al, temp_e_byte
        out     e_bar_port, al
        pop     si! pop dx! pop es! pop cx! pop ax! pop bx
ret

```

```

;-----
initialize_cpu_interrupts:

```

```

; Module Interface Specification:

```

```

;   Caller:           Ethertest(PL/I) Procedure

```

```

;   Parameters:       NONE

```

```

initmodule cseg common
        org 114h
        int5_offset  rw 1
        int5_segment rw 1

```

```

        cseg
        push bx
push     ax
        mov     bx, offset interrupt_handler
        mov     ax, 3
        push  ds
        mov     ds, ax

```

```
        mov ds:int5_offset, bx
        mov bx, cs
        mov ds:int5_segment, bx
        pop ds
    pop ax
        pop bx
        sti
        ret
```

enable_cpu_interrupts:

; Module Interface Specification:

; Caller: Ethertest(PL/I) Procedure

; Parameters: NONE

```
        sti
        ret
```

disable_cpu_interrupts:

; Module Interface Specification:

; Caller: Ethertest(PL/I) Procedure

; Parameters: none

```
        cli
        ret
```

interrupt_handler:

```
; IP, CS, and flags are already on stack
; save all other registers
```

```
push ax
        push bx
push cx
push dx
```

```
push si
push di
push bp
push ds
push es
    call hl_interrupt_handler ; high level source
                                ; routine
; restore registers

pop es
pop ds
pop bp
pop di
pop si
pop dx
pop cx
pop bx
pop ax
    sti
    iret
```

end

APPENDIX D

Distributed Decision Algorithm

Source Code

PA2, PA3, PB2, and PB3, the distributed user processes which implement the distributed decision algorithm described in Chapter III, are documented herein. Note that the systems file SYSDEF, described in Appendix B, must also be available for compilation of each user process.

Processes PA2 and PA3 are linked as described in Appendix A. Their associated command files NUM12.CMD and NUM13.CMD are loaded into local memory of SBC #2 and SBC #3 respectively in cluster A at runtime. PB2 and PB3 produce NUM22.CMD and NUM23.CMD which are loaded into the memories of cluster B in the same way.

Processes are loaded when requested under MCORTEX control and execution begins and continues until an await state is encountered. Once all processes have been loaded, the various await states will be satisfied by advances of eventcounts in other processes and operation will continue until all input data vectors are processed.

```

/*****
/*****
/*
/* PA2 is resident in local memory of SBC 2, CLUSTER A. */
/*
/* This procedure performs the following operations: */
/*
/* 1. Loads quadratic equation parameters A,B,C,D. */
/* 2. Reads sensor A observation vectors from disk. */
/* 3. Computes LLF ( LAMBDA_A_X ) for local use. */
/* 4. Computes ( LAMBDA_A_X + LAMBDA_PP_Y ) the */
/* sum of the local and remote sensor LLR's. */
/* 5. Compares the result to the decision threshold */
/* and displays the final result and decision. */
/* 6. Performs steps 2-5 for each input vector. */
/*
/*****
/*****

```

PA2: PROCEDURE;

```

XINCLUDE 'SYSDEF.PLI';
%REPLACE
PC BY '0DD0'B4, /* P3 IS SET TO THIS VALUE */
TRUE BY '1'B,
FALSE BY '0'B,
ONE BY '0001'B4;

```

```

DECLARE
/*****
/* PARFILE CONTAINS THE FOLLOWING PARAMETERS */
/*
/* 1. MATRIX/VECTOR DIMENSION. */
/* 2. D DIAGONAL ELEMENTS OF THE MATRIX-A. */
/* 3. COL BY COL ELEMENTS BELOW DIAGONAL OF */
/* MATRIX-A. */
/* 4. D ELEMENTS OF VECTOR-B. */
/* 5. SCALAR-C. */
/* 6. THRESHOLD. */
/*
/* DATFILE CONTAINS THE FOLLOWING VALUES */
/*
/* 1. D-ELEMENT X-VECTORS. */
/*
/*****

```

```

(PARFILE,DATFILE) FILE,
EOF BIT(1) STATIC INIT(FALSE),
(I,J,D,N) FIXED,
(A(528),B(32),C,T2,X(32),THRESH,LAMBDA_A_X) FLOAT,
K BIT(16) STATIC INIT('0000'B4),

```

```

/*****
/*
/* P3 SET TO PC TO BE ADDED TO SEGMENT ADDF 0800
/*
/*****

```

P3 POINTER,

```

/*****
/*
/* BASE LAMBDA_BP_Y AT P3 = PC (OFFSET ADD TO DATA
/* SEGMENT = 0800 )
/*
/*****

```

LAMBDA_BP_Y FLOAT BASED(P3);

```

/* SET POINTERS TO VALUES INDICATED IN REPLACE ABOVE */

```

UNSPEC(P3) = PC;

```

/*****
/*
/* INPUT PARAMETERS FROM DISK FILE
/*
/* MATRIX & VECTOR DIMENSION (D = INTEGER)
/*
/* CALCULATE N = # OF MATRIX ELEMENTS TO INPUT
/*
/* MATRIX-A (SYMMETRIC)
/*
/* DIAGONAL ELEMENTS FIRST (# = D)
/* COLUMNS BELOW DIAGONAL NEXT (# = N-D)
/*
/* VECTOR-B (D ELEMENTS)
/*
/* SCALAR-C (1 NUMBER)
/*
/* THRESHOLD (1 NUMBER)
/*
/*****

```

OPEN FILE(PARFILE) STREAM INPUT;

GET FILE(PARFILE) LIST (D);

N = ((D * D) + D) / 2;

DO I=1 TO N;

GET FILE(PARFILE) LIST A(I);

END;


```

DO I=1 TO D;
  GET FILE(PARFILE) LIST (B(I));
END;

GET FILE(PARFILE) LIST (C,THRESH);

PUT SKIP LIST ('DIMENSION =',D,'THRESHOLD =',THRESH);

/*****
/*
/*      INPUT AND PROCESS X-VECTORS
/*
/*
*****/

ON ENDFILE(DATFILE) EOF = TRUE;
OPEN FILE(DATFILE) STREAM INPUT;

DO WHILE(EOF = FALSE);

  K = ADD2BIT16(K,ONE);

  PUT SKIP(2);
  DO I=1 TO D;
    GET FILE(DATFILE) LIST (X(I));
    PUT SKIP LIST('X      ('I,') =',X(I));
  END;

/* CALC  LAMBDA_A_X = (X-TRANS)*(A _MATRIX)*(X)      */

  LAMBDA_A_X = 0;
  DO J=1 TO D-1;
    DO I=J+1 TO D;
      LAMBDA_A_X = LAMBDA_A_X + A(I+J+1)*X(I)*X(J);
    END;
  END;

  T2 = 0;
  DO I=1 TO D;
    T2 = T2 + ( A(I)*X(I)*X(I));
  END;
  LAMBDA_A_X = (2*LAMBDA_A_X ) + T2;

/* ADD LAMBDA_A_X TO ( B-VECTOR)*(X) & STORE      */

  DO I=1 TO D;
    LAMBDA_A_X = LAMBDA_A_X + ( B(I) * X(I));
  END;

/* ADD LAMBDA_A_X TO C & STORE IN LAMBDA_A_X      */

  LAMBDA_A_X = LAMBDA_A_X + C;

```

```

/*****
/*
/* AWAIT LAMBDA_BP_Y CALCULATED IN THE OTHER CLUSTER */
/*
/*****

      CALL AWAIT(B1EVC,K);

      PUT SKIP(2) LIST('LAMBDA_A_X =',LAMBDA_A_X );
      PUT SKIP LIST ('LAMBDA_BP_Y =',LAMBDA_BP_Y);

/*****
/*
/* ADD THE LAMBDA_BP_Y VALUE RECEIVED FROM */
/* THE OTHER CLUSTER VIA THE ETHERNET TO */
/* THE LAMBDA_A_X VALUE CALCULATED IN THIS */
/* CLUSTER, AND COMPARE TO THE THRESHOLD. */
/*
/*****

      T2 = LAMBDA_A_X + LAMBDA_BP_Y;

      IF (T2 > THRESH) THEN DO;

          PUT SKIP LIST('RESULT      =',T2,' IS > THRESHOLD ');

          END;
          ELSE DO:

              PUT SKIP LIST('RESULT      =',T2,' IS < THRESHOLD ');

          END;

      DO I=0 TO 1000;
          DO J=0 TO 500; /* DELAY LOOP */
              END;
          END;

/*****
/*
/* NOTIFY BOARD 3 TO CONTINUE WITH NEXT INPUT VECTOR */
/*
/*****

      CALL ADVANCE(A2EVC);

END;          /* END OF DO WHILE (VCF = FALSE) LOOP */

      PUT SKIP(3) LIST('END OF INPUT DATA');

END PA2;

```

```

/*****
/*****
/*
/* PA3 is resident in local memory of SBC 3, CLUSTER A. */
/*
/* This procedure performs the following operations: */
/*
/* 1. Loads quadratic equation parameters A,B,C,D. */
/* 2. Reads sensor A observation vectors from disk. */
/* 3. Computes the Conditional LLR ( LAMBDA_AP_X ) */
/* to send to sensor B for further computation. */
/* 4. Submits a request into the ERB queue to send */
/* the CLLR statistic to sensor B. */
/* 5. Advances eventcount A1EVC to signal sensor B */
/* that its awaited statistic is available. */
/*
/*****
/*****

```

PA3: PROCEDURE;

```

%INCLUDE 'SYSDEF.PLI';
%REPLACE

```

```

PA BY '8000'B4, /* P1 IS SET TO THIS VALUE */
PB BY '8007'B4, /* P2 IS SET TO THIS VALUE */
ERB_BLOCK_LENGTH BY 24, /* USED TO CONTROL */
ERB_BLOCK_LENGTH_M1 BY 19, /* ERB SIZE */
TRUE BY '1'B,
FALSE BY '0'B,
ONE BY '0001'B4;

```

DECLARE

```

/*****
/* PARAFIELD CONTAINS THE FOLLOWING PARAMETERS */
/*
/* 1. MATRIX/VECTOR DIMENSION. */
/* 2. D DIAGONAL ELEMENTS OF THE MATRIX-AP. */
/* 3. COL BY COL ELEMENTS BELOW DIAGONAL OF */
/* MATRIX-AP. */
/* 4. D ELEMENTS OF VECTOR-BP. */
/* 5. SCALAR-CP. */
/*
/* DATAFILE CONTAINS THE FOLLOWING VALUES */
/*
/* 1. D-ELEMENT X-VECTORS. */
/*
/*****

```

```

(PARAFIELD,DATAFILE) FILE,
EOF BIT(1) STATIC INIT(FALSE).

```

(I,J,D,N) FIXED,
(AP(528),BP(32),CP,T1,X(32)) FLOAT,

```
/**/
/*
/* INDEX VARIABLES AND CONSTANTS USED FOR
/*
/* INDEXING IN THE ERB (ERB_INDEX)
/* SEQUENCING & CONTROL( II, JJ, K )
/* IDENTIFYING DATA TRANSFER(DATA_TYPE)
/* IDENTIFYING OPPOSITE CLUSTER_ADDRESS
/*
/**/
```

ERB_INDEX FIXED,
 (II, JJ) BIT(16),
K BIT(16) STATIC INIT('0000'34),
DATA_TYPE BIT(8) STATIC INIT('01'84),
CLUSTER_ADDRESS BIT(16) STATIC INIT('0002'84),

```
/**/
/*
/* PCINTERS ARE USED IN THE FOLLOWING MANNER
/*
/* P1 SET TO PA TO BE ADDED TO SEGMENT ADDR 0800
/* P2 SET TO PB TO BE ADDED TO SEGMENT ADDR 0800
/*
/**/
```

(P1, P2) POINTER,

```
/**/
/*
/* THE ETHERNET REQUEST BLOCK (ERB)
/*
/* ETHERNET REQUEST PACKET (ERP) STRUCTURE
/*
/* IS USED IN THE FOLLOWING MANNER
/*
/* COMMAND = 1   FOR DATA TRANSFER OVER E-NET
/* TYPE        (NOT USED BY THIS PROCEDURE)
/* VALUE       (NOT USED BY THIS PROCEDURE)
/* REMOTE_ADDR = CLUSTER_ADDRESS OF DESTINATION
/*
/**/
```

1 ERP(7:ERB_BLOCK_LENGTH_M1) BASED (P1),
2 COMMAND BIT(8),
2 TYPE BIT(8),
2 VALUE BIT(16),
2 REMOTE_ADDR BIT(16),

```

/*****
/*
/*  BASE LAMBDA_AP_X  AT P2 = PB (OFFSET ADD TO DATA  */
/*                               SEGMENT = 0800 )      */
/*
/*
/*****

```

```

      LAMBDA_AP_X FLOAT BASED(P2);

```

```

/* SET PCINTERS TO VALUES INDICATED IN REPLACE ABOVE */

```

```

      UNSPEC(P1) = PA;

```

```

      UNSPEC(P2) = PB;

```

```

/*****
/*
/*  INPUT PARAMETERS FROM DISK FILE
/*
/*  MATRIX & VECTOR DIMENSION (D = INTEGER)
/*
/*  CALCULATE N = # OF MATRIX ELEMENTS TO INPUT
/*
/*  MATRIX-AP (SYMMETRIC)
/*
/*  DIAGONAL ELEMENTS FIRST (# = D)
/*  COLUMNS BELOW DIAGONAL NEXT (# = N-1)
/*
/*  VECTOR-BP (D ELEMENTS)
/*
/*  SCALAR-CP (1 NUMBER)
/*
/*****

```

```

      OPEN FILE(PARAFILE) STREAM INPUT;

```

```

      GET FILE(PARAFILE) LIST (D);

```

```

      N = ((D * D) + D) / 2;

```

```

      DO I=1 TO N;
        GET FILE(PARAFILE) LIST (AP(I));
      END;

```

```

      DO I=1 TO D;
        GET FILE(PARAFILE) LIST (BP(I));
      END;

```

```

      GET FILE(PARAFILE) LIST (CP);

```

```

      PUT SKIP LIST ('DIMENSION =', D);

```

```

/*****
/*
/*      INPUT AND PROCESS X-VECTORS
/*
/*
/*****

      ON ENDFILE(DATAFILE) EOF = TRUE;

      OPEN FILE(DATAFILE) STREAM INPUT;

      DO WHILE(EOF = FALSE);

          CALL AWAIT(A2EVC,K);

          PUT SKIP(2);

          DO I=1 TO D;

              GET FILE(DATAFILE) LIST (X(I));

              PUT SKIP LIST('X      ('.I.') ='.X(I));

          END;

/* STORE (X-TRANS)*(AP-MATRIX)*(X) IN LAMBDA_AP_X      */

      LAMBDA_AP_X = 0;

      DO J=1 TO D-1;
        DO I=J+1 TO D;
          LAMBDA_AP_X = LAMBDA_AP_X+(AP(I+J+1)*X(I)*X(J));
        END;
      END;

      T1 = 0;

      DO I=1 TO D;
        T1 = T1 + (AP(I)*X(I)*X(I));
      END;

      LAMBDA_AP_X = (2*LAMBDA_AP_X) + T1;

/* ADD LAMBDA_AP_X TO (BP-VECTOR)*(X) & STORE      */

      DO I=1 TO D;
        LAMBDA_AP_X = LAMBDA_AP_X + (BP(I) * X I));
      END;

/* ADD LAMBDA_AP_X TO CP & STORE IN LAMBDA_AP_X      */

      LAMBDA_AP_X = LAMBDA_AP_X + CP;

```

```

/*****
/*
/*      GET A TICKET TO ENABLE A WRITE TO THE ERB      */
/*
/*****

      II = TICKET(ERB_WRITE_PREQUEST);

/* II NOW HAS THE VALUE OF THE TICKET RETURNED      */

      JJ = READ(ERB_WRITE);

/* JJ NOW HAS THE VALUE OF ERB_WRITE      */

      DO WHILE(JJ < II);
          JJ = READ(ERB_WRITE);
      END;

/* IF ETHERNET REQUEST BLOCK (ERB) IS FULL-BUSY WAIT */
      JJ = READ(ERB_READ);
      DO WHILE((II - JJ) >= ERB_BLOCK_LENGTH);
          JJ = READ(ERB_READ);
      END;

/*****
/*
/* WRITE TO ERB WHEN A SLOT IS OPEN      */
/* COMMAND = 1 FOR DATA TO BE TRANSFERRED      */
/* REMOTE_ADDR = DESTINATION CLUSTER ADDRESS      */
/*
/*****

      ERB_INDEX = MOD(II, ERB_BLOCK_LENGTH);
      ERB(ERB_INDEX).COMMAND = DATA_TYPE;
      ERB(ERB_INDEX).REMOTE_ADDR = CLUSTER_ADDRESS;

/*****
/*
/*      NOTIFY MCORTEX THAT ERB WRITE IS COMPLETE      */
/*
/*****

      CALL ADVANCE(ERB_WRITE);

```

```

/*****
/*
/*      AN ETHERNET REQUEST PACKET (ERP) IS NOW SETUP /*
/* IN THE ETHERNET REQUEST BLOCK (ERB). THIS WILL /*
/* SIGNAL THE DRIVER PROCEDURE ON BOARD 1 TO FETCH /*
/* THE DATA STORED IN COMMON MEMORY (LAMBDA_AP_X) AT /*
/* ADDRESS 0800:8CC0-0800:8CC3 & MOVE IT TO /*
/* ADDRESS 0800:80DA-0800:80DD (TRANSMIT_DATA_BLOCK) /*
/* ALSO IN COMMON MEMORY TO BE PACKETIZED AND SENT TO /*
/* THE RECEIVE_DATA_BLOCK (RDB) OF THE OTHER CLUSTER /*
/* ADDRESS 0800:867C-0800:867F WHERE IT IS MOVED TO /*
/* ADDRESS 0800:8DD0-0800:8DD3 IN THE OTHER CLUSTERS /*
/* COMMON MEMORY (LAMBDA_AP_X). /*
/*
/*****

/*****
/*
/*      NOTIFY OTHER CLUSTER THAT DATA IS READY /*
/*
/*****

      CALL ADVANCE(A1EVC);

      K = ADD2BIT16(K,ONE);

END;      /* END OF DO WHILE (ECF = FALSE) LOOP */

      PUT SKIP(3) LIST('END OF INPUT DATA');

END PA3;

/*****
/*****
/*
/*      PB2 is resident in local memory of SBC 2, CLUSTER B. /*
/*
/*      This procedure performs the following operations: /*
/*
/*      1. Loads quadratic equation parameters A,B,C,D. /*
/*      2. Reads sensor P observation vectors from disk. /*
/*      3. Computes LLR ( LAMBDA_P_Y ) for local use. /*
/*      4. Computes ( LAMBDA_B_Y + LAMBDA_AP_X ) the /*
/*          sum of the local and remote sensor LLR's. /*
/*      5. Compares the result to the decision threshold /*
/*          and displays the final result and decision. /*
/*      6. Performs steps 2-5 for each input vector. /*
/*
/*****
/*****

```


PB2: PROCEDURE;

%INCLUDE 'SYSDEF.PLI';
%REPLACE

PC BY '8DD2'B4. /* P3 IS SET TO THIS VALUE */
TRUE BY '1'B,
FALSE BY '0'B,
ONE BY '0001'B4;

DECLARE

```
/******  
/* PARFILE CONTAINS THE FOLLOWING PARAMETERS */  
/*  
/* 1. MATRIX/VECTOR DIMENSION. */  
/* 2. D DIAGONAL ELEMENTS OF THE MATRIX-A. */  
/* 3. COL BY COL ELEMENTS BELOW DIAGONAL OF */  
/* MATRIX-A. */  
/* 4. D ELEMENTS OF VECTOR-B. */  
/* 5. SCALAR-C. */  
/* 6. THRESHOLD. */  
/*  
/* DATFILE CONTAINS THE FOLLOWING VALUES */  
/*  
/* 1. D-ELEMENT Y-VECTORS. */  
/*  
/******
```

(PARFILE,DATFILE) FILE,
EOF BIT(1) STATIC INIT(FALSE),
(I,J,D,N) FIXED,
(A(528),B(32),C,T2,Y(32),THRESH,LAMBDA_B_Y) FLOAT,
K BIT(16) STATIC INIT('0000'B4),

```
/******  
/*  
/* P3 SET TO PC TO BE ADDED TO SEGMENT ADDR 0P00 */  
/*  
/******
```

P3 POINTER,

```
/******  
/*  
/* EASE LAMBDA_AP_X AT P3 = PC (OFFSET ADD TO DATA */  
/* SEGMENT = 0800 ) */  
/*  
/******
```

LAMBDA_AP_X FLOAT BASED(P3);

/* SET POINTERS TO VALUES INDICATED IN REPLACE ABOVE */

UNSPEC(P3) = PC;

```

/*****
/*
/*      INPUT PARAMETERS FROM DISK FILE
/*
/*      MATRIX & VECTOR DIMENSION (D = INTEGER)
/*
/*      CALCULATE N = # OF MATRIX ELEMENTS TO INPUT
/*
/*      MATRIX-A (SYMMETRIC)
/*
/*      DIAGONAL ELEMENTS FIRST (# = D)
/*      COLUMNS BELOW DIAGONAL NEXT (# = N-D)
/*
/*      VECTOR-B (D ELEMENTS)
/*
/*      SCALAR-C (1 NUMBER)
/*
/*      THRESHOLD (1 NUMBER)
/*
*****/

```

OPEN FILE(PARFILE) STREAM INPUT;

GET FILE(PARFILE) LIST (D);
N = ((D * D) + D) / 2;

DO I=1 TO N;
GET FILE(PARFILE) LIST (A(I));
END;

DO I=1 TO D;
GET FILE(PARFILE) LIST (B(I));
END;

GET FILE(PARFILE) LIST (C, THRESH);

PUT SKIP LIST ('DIMENSION =', D, 'THRESHOLD =', THRESH);

```

/*****
/*
/*      INPUT AND PROCESS Y-VECTORS
/*
*****/

```

ON ENDFILE(DATAFILE) EOF = TRUE;

OPEN FILE(DATAFILE) STREAM INPUT;

DO WHILE(EOF = FALSE);

```

        K = ADD2BIT16(K,ONE);

        PUT SKIP(2);
        DO I=1 TO D;
            GET FILE(DATFILE) LIST (Y(I));
            PUT SKIP LIST('Y      ('.I.') =',Y(I));
        END;

/* CALC LAMBDA_B_Y = (Y-TRANS)*(A_MATRIX)*(Y)      */
        LAMBDA_B_Y = 0;
        DO J=1 TO D-1;
            DO I=J+1 TO D;
                LAMBDA_B_Y = LAMBDA_B_Y +( A(I+J+1)*Y(I)*Y(J));
            END;
        END;

        T2 = 0;
        DO I=1 TO D;
            T2 = T2 + ( A(I)*Y(I)*Y(I));
        END;

        LAMBDA_P_Y = (2*LAMBDA_P_Y ) + T2;

/* ADD LAMBDA_B_Y TO ( B-VECTOR)*(Y) & STORE      */
        DO I=1 TO D;
            LAMBDA_B_Y = LAMBDA_B_Y + ( B(I) * Y(I));
        END;

/* ADD LAMBDA_B_Y TO C & STORE IN LAMBDA_B_Y      */
        LAMBDA_P_Y = LAMBDA_B_Y + C;

/*****
/*
/* AWAIT LAMBDA_AP_X CALCULATED IN THE OTHER CLUSTER
/*
*****/

        CALL AWAIT(A1EVC,K);
        PUT SKIP(2) LIST('LAMBDA_B_Y =',LAMBDA_P_Y );
        PUT SKIP LIST ('LAMBDA_AP_X =',LAMBDA_AP_X);

/*****
/* ADD THE LAMBDA_AP_X VALUE RECEIVED FROM
/* THE OTHER CLUSTER VIA THE ETHERNET TO
/* THE LAMBDA_P_Y VALUE CALCULATED IN THIS
/* CLUSTER, AND COMPARE THE RESULT TO THE
/* THRESHOLD VALUE.
*****/

```

```

T2 = LAMBDA_B_Y + LAMBDA_AP_X;
IF (T2 > THRESH) THEN DO;
  PUT SKIP LIST('RESULT      =',T2,'IS > THRESHOLD ');
END;
ELSE DO;
  PUT SKIP LIST('RESULT      =',T2,'IS < THRESHOLD ');
END;

DO I=0 TO 1000;
  DO J=0 TO 500; /* DELAY LOOP */
  END;
END;

/*****
/*
/* NOTIFY BOARD 3 TO CONTINUE WITH NEXT INPUT VECTOR */
/*
*****/

      CALL ADVANCE(B2EVC);

END;      /* END OF DO WHILE (EOF = FALSE) LOOP */

      PUT SKIP(3) LIST('END OF INPUT DATA');

END PB2;

/*****
*****/
/*
/* PB3 is resident in local memory of SRC 3, CLUSTER B. */
/*
/* This procedure performs the following operations: */
/*
/* 1. Loads quadratic equation parameters A,B,C,D. */
/* 2. Reads sensor B observation vectors from disk. */
/* 3. Computes the Conditional LLR ( LAMBDA_PP_Y ) */
/*    to send to sensor A for further computation. */
/* 4. Submits a request into the FRB queue to send */
/*    the CLLR statistic to sensor A. */
/* 5. Advances eventcount B1EVC to signal sensor A */
/*    that its awaited statistic is available. */
/*
*****/
/*****
*****/

```

PP3: PROCEDURE:

```

%INCLUDE 'SYSDEF.PLI';
%REPLACE

```

```

PA   BY '0000'B4, /* P1 IS SET TO THIS VALUE */
PB   BY '8CC0'B4, /* P2 IS SET TO THIS VALUE */

ERB_BLOCK_LENGTH BY 20, /* USED TO CONTROL */
ERE_BLOCK_LENGTH_M1 BY 19, /* ERE SIZE */

TRUE BY '1'B,
FALSE BY '0'B,

ONE BY '0001'B4;

```

DECLARE

```

/*****
/*   PARAFIELD CONTAINS THE FOLLOWING PARAMETERS   */
/*
/*       1. MATRIX/VECTOR DIMENSION.             */
/*       2. D DIAGONAL ELEMENTS OF THE MATRIX-AP. */
/*       3. COL BY COL ELEMENTS BELOW DIAGONAL OF */
/*                                     MATRIX-AP.  */
/*       4. D ELEMENTS OF VECTOR-BP.             */
/*       5. SCALAR-CP.                           */
/*
/*   DATAFILE CONTAINS THE FOLLOWING VALUES   */
/*
/*       1. D-ELEMENT Y-VECTORS.                 */
/*
*****/

```

```

(PARAFIELD,DATAFILE) FILE,
ECF BIT(1) STATIC INIT(FALSE),
(I,J,D,N) FIXED,
(AP(528),BP(32),CP,T1.Y(32)) FLOAT,

```

```

/*****
/*
/*   INDEX VARIABLES AND CONSTANTS USED FOR   */
/*
/*       INDEXING IN THE ERR (ERB_INDEX)        */
/*       SEQUENCING & CONTROL( II,JJ,K )      */
/*       IDENTIFYING DATA TRANSFER(DATA_TYPE) */
/*       IDENTIFYING OPPOSITE CLUSTER_ADDRESS */
/*
*****/

```

```

ERB_INDEX      FIXED,
(II,JJ)        BIT(16),
K              BIT(16) STATIC INIT('0000'B4),
DATA_TYPE      BIT(8)  STATIC INIT('21'B4),
CLUSTER_ADDRESS BIT(16) STATIC INIT('0001'B4),

```

```

/*****
/*
/*   POINTERS ARE USED IN THE FOLLOWING MANNER   */
/*
/*   P1 SET TO PA TO BE ADDED TO SEGMENT ADDR 0800 */
/*   P2 SET TO PB TO BE ADDED TO SEGMENT ADDR 0800 */
/*
*****/

```

(P1,P2) POINTER,

```

/*****
/*
/*   THE ETHERNET REQUEST BLOCK (ERB)           */
/*
/*   ETHERNET REQUEST PACKET (ERP) STRUCTURE   */
/*
/*   IS USED IN THE FOLLOWING MANNER           */
/*
/*   COMMAND = 1   FOR DATA TRANSFER OVER E-NET */
/*   TYPE          (NOT USED BY THIS PROCEDURE) */
/*   VALUE         (NOT USED BY THIS PROCEDURE) */
/*   REMOTE_ADDR = CLUSTER_ADDRESS OF DESTINATION */
/*
*****/

```

```

1 ERB(0:ERB_BLOCK_LENGTH_M1) BASED (P1),
2 COMMAND      BIT(8),
2 TYPE        BIT(8),
2 VALUE       BIT(16),
2 REMOTE_ADDR BIT(16).

```

```

/*****
/*
/*   BASE LAMBDA_PP_Y AT P2 = PB 'OFFSET ADD TO DATA */
/*                                     SEGMENT = 0800 ) */
/*
*****/

```

LAMBDA_BP_Y FLOAT BASED(P2);

/* SET POINTERS TO VALUES INDICATED IN REPLACE ABOVE */

UNSPEC(P1) = PA;

UNSPEC(P2) = PB;

```

/*****
/*
/*      INPUT PARAMETERS FROM DISK FILE
/*
/*      MATRIX S VECTOR DIMENSION (D = INTEGER)
/*
/*      CALCULATE N = # OF MATRIX ELEMENTS TO INPUT
/*
/*      MATRIX-AP (SYMMETRIC)
/*
/*      DIAGONAL ELEMENTS FIRST (# = D)
/*      COLUMNS BELOW DIAGONAL NEXT (# = N-D)
/*
/*      VECTOR-EP (D ELEMENTS)
/*
/*      SCALAR-CP (1 NUMBER)
/*
*****/

```

```

OPEN FILE(PARAFILE) STREAM INPUT;

```

```

GET FILE(PARAFILE) LIST (D);
N = ((D * D)+D)/2;

```

```

DO I=1 TO N;
  GET FILE(PARAFILE) LIST (AP(I));
END;

```

```

DO I=1 TO D;
  GET FILE(PARAFILE) LIST (BP(I));
END;

```

```

GET FILE(PARAFILE) LIST (CP);

```

```

PUT SKIP LIST ('DIMENSION =',D);

```

```

/*****
/*
/*      INPUT AND PROCESS Y-VECTORS
/*
*****/

```

```

ON ENDFILE(DATAFILE) EOF = TRUE;

```

```

OPEN FILE(DATAFILE) STREAM INPUT;

```

```

DO WHILE(EOF = FALSE);

```

```

  CALL AWAIT(E2EVC,K);

```

```

  PUT SKIP(2);

```

```

DO I=1 TO D;
  GET FILE(DATFILE) LIST (Y(I));
  PUT SKIP LIST('Y      (',I,') =',Y(I));
END;

/* STORE (Y-TRANS)*(AP-MATRIX)*(Y) IN LAMBDA_BP_Y      */
LAMBDA_BP_Y = 0;
DC J=1 TO D-1;
DO I=J+1 TO D;
  LAMBDA_BP_Y = LAMBDA_BP_Y+(AP(I+J+1)*Y(I)*Y(J));
END;
END;

T1 = 0;

DO I=1 TO D;
  T1 = T1 + (AP(I)*Y(I)*Y(I));
END;

LAMBDA_BP_Y = (2*LAMBDA_BP_Y) + T1;

/* ADD LAMBDA_BP_Y TO (PP-VECTOR)*(Y) & STORE      */
DO I=1 TO D;
  LAMBDA_PP_Y = LAMBDA_BP_Y + (PP(I) * Y(I));
END;

/* ADD LAMBDA_PP_Y TO CP & STORE IN LAMBDA_PP_Y      */
LAMBDA_PP_Y = LAMBDA_PP_Y + CP;

/*****
/*
/* GET A TICKET TO ENABLE A WRITE TO THE ERB      */
/*
*****/

II = TICKET(ERB_WRITE_REQUEST);

/* II NOW HAS THE VALUE OF THE TICKET RETURNED      */
JJ = READ(ERB_WRITE);

/* JJ NOW HAS THE VALUE OF ERB_WRITE      */
DO WHILE(JJ < II);
  JJ = READ(ERB_WRITE);
END;

```



```

/* IF ETHERNET REQUEST BLOCK (ERB) IS FULL-BUSY WAIT */
    JJ = READ(ERB_READ);
    DO WHILE((II - JJ) >= ERB_BLOCK_LENGTH);
        JJ = READ(ERB_READ);
    END;

/*****
/*
/* WRITE TO ERB WHEN A SLOT IS OPEN
/* COMMAND = 1 FOR DATA TO BE TRANSFERRED
/* REMOTE_ADDR = DESTINATION CLUSTER ADDRESS
/*
*****/

    ERB_INDEX = MOD(II,ERB_BLOCK_LENGTH);
    ERB(ERB_INDEX).COMMAND = DATA_TYPE;
    ERB(ERB_INDEX).REMOTE_ADDR = CLUSTER_ADDRESS;

/*****
/*
/* NOTIFY MCORTEX THAT ERP WRITE IS COMPLETE
/*
*****/

    CALL ADVANCE(ERB_WRITE);

/*****
/*
/* AN ETHERNET REQUEST PACKET (ERP) IS NOW SETUP
/* IN THE ETHERNET REQUEST BLOCK (ERB). THIS WILL
/* SIGNAL THE DRIVER PROCEDURE ON BOARD 1 TO FETCH
/* THE DATA STORED IN COMMON MEMORY (LAMBDA_PP_Y) AT
/* ADDRESS 0800:4CC0-0800:80C3 & MOVE IT TO
/* ADDRESS 2800:80DA-0800:80DD (TRANSMIT_DATA_BLOCK)
/* ALSO IN COMMON MEMORY TO BE PACKETIZED AND SENT TO
/* THE RECEIVE_DATA_BLOCK (RDB) OF THE OTHER CLUSTER
/* ADDRESS 0800:867C-0800:867F WHERE IT IS MOVED TO
/* ADDRESS 2800:8DD0-0800-8DD3 IN THE OTHER CLUSTERS
/* COMMON MEMORY (LAMBDA_PP_Y).
/*
*****/

/*****
/*
/* NOTIFY OTHER CLUSTER THAT DATA IS READY
/*
*****/

```

```
CALL ADVANCE(B1EVC);  
K = ADD2BIT16(K,ONE);  
END;      /* END OF DO WHILE (EOF = FALSE) LOOP */  
PUT SKIP(3) LIST('END OF INPUT DATA');  
END PB3;
```

LIST OF REFERENCES

1. Hahn, S.C., Analysis of a Distributed Decision Algorithm, M.S. Thesis, Naval Postgraduate School, Monterey, California, December 1985.
2. Brewer, D.J., A Real-Time Executive for Multiple-Computer Clusters, M.S. Thesis, Naval Postgraduate School, Monterey, California, December 1984.
3. Tenney, R.R. and Sandell, N.R., "Detection with Distributed Sensors," IEEE Trans. Aerospace and Electronic Systems, Vol. AES-17, No.4, July 1981.
4. Rosenfeld, A., Hummel, R.A. and Zucker, S.W., "Scene Labeling by Relaxation Operations", IEEE Trans. Systems, Man, and Cybernetics, Vol. 6, 1976 pp 420-433.
5. Haralick, R.M., "Decision Making in Context", IEEE Trans. Pattern Anal. Machine Intelligence, Vol. PAMI-5, No.4, July 1983.
6. Van Trees, H.L., Detection Estimation and Modulation Theory, Part I, John Wiley & Sons, New York, 1968.
7. Duda, R.O. and Hart, P.E., Pattern Classification and Scene Analysis, John Wiley & Sons, New York, 1973.
8. Fukunaga, K., Introduction to Statistical Pattern Recognition, Academic Press, New York, 1972.
9. Helstrom, C.W., Probability and Stochastic Processes for Engineers, Macmillan Publishing Company, New York, 1984.

10. Xerox Corporation, The Ethernet – A Local Area Network: Data Link Layer and Physical Layer Specifications, Version 1.0, September 1980.
11. Reed, D.P. and Kanodia, R.J., "Synchronization with Eventcounts and Sequencers", Communication of the ACM, Volume 22, pp. 115-123, February 1979.

INITIAL DISTRIBUTION LIST

		No. Copies
1.	Defense Technical Information Center Cameron Station Alexandria, Virginia 22304-6145	2
2.	Library, Code 0142 Naval Postgraduate School Monterey, California 93943-5100	2
3.	Department Chairman, Code 62 Electrical and Computer Engineering Dept. Naval Postgraduate School Monterey, California 93943-5100	1
4.	Professor Charles W. Therrien, Code 62Ti Electrical and Computer Engineering Dept. Naval Postgraduate School Monterey, California 93943-5100	3
5.	Professor Uno R. Kodres, Code 52Kr Computer Science Department Naval Postgraduate School Monterey, California 93943-5100	2
6.	Professor M. L. Cotton, Code 62Cc Electrical and Computer Engineering Dept. Naval Postgraduate School Monterey, California 93943-5100	1
7.	Professor R. Panholzer, Code 62Pz Electrical and Computer Engineering Dept. Naval Postgraduate School Monterey, California 93943-5100	1
8.	Capt. Mark A. Schon 1801 Artillery Ridge Road Fredericksburg, Virginia 22401	3

9. Daniel Green, Code 20E 1
Naval Surface Weapons Center
Dahlgren, Virginia 22449
10. Capt. J. Donegan, USN 1
PMS 400B5
Naval Sea System Command
Washington, D.C. 20362
11. RCA AEGIS Depository 1
RCA Corporation
Government Systems Division
Mail Stop 127-327
Moorestown, New Jersey 08057
12. Library, Code E33-05 1
Naval Surface Weapons Center
Dahlgren, Virginia 22449
13. Dr. M.J. Gralia 1
Applied Physics Laboratory
John Hopkins Road
Laurel, Maryland 20707
14. Dana Small 1
Code 8242, NOSC
San Diego, California 92152
15. Lt. Robin K. Weinhold-Schon 1
4582 Marlwood Way
Virginia Beach, Virginia 23462
16. Susan M. Schon 1
Route 1, Box 531
Beaufort, North Carolina 28516
17. Otto W. Schon 1
509 Stafford Avenue
Bristol, Connecticut 06010

- | | | |
|-----|--|---|
| 18. | Kevin and Michael Schon
1914 Nelson Avenue
Memphis, Tennessee 38104 | 2 |
| 19. | Capt. "Wild Bill" Johnson
2708 Pine Manor Lane
Albany, Georgia 31707 | 1 |
| 20. | Professor P. Moose, Code 62Me
Electrical and Computer Engineering Dept.
Naval Postgraduate School
Monterey, California 93943-5100 | 1 |
| 21. | Mr. D. Cowan
NWC China Lake, Code 31507
China Lake, California 93555 | 1 |
| 22. | Dr. P. Krueger
AIR 320H
Naval Air Systems Command
Washington, D.C. 20361 | 1 |

END

FILMED

386

DTIC