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A NATURAL LANGUAGE PROCESSOR AND ITS APPLICATION TO A DATA DICTIONARY SYSTEM

THESIS

Stephen A. Wolfe Captain, USAF

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AFIT/GCS/ENG/85D-19

A NATURAL LANGUAGE PROCESSOR AND ITS APPLICATION TO A DATA DICTIONARY SYSTEM

THESIS

Presented to the Faculty of the School of Engineering of the Air Force Institute of Technology

Air University

In Partial Fulfillment of the

Requirements for the Degree of Master of Science in Computer Engineering

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Stephen A. Wolfe, A.B.

Captain, USAF

December 1985

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Acknowledgments

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Stephen A. Wolfe

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List of Acronyms

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AFIT	Air Force Institute of Technology
AFIT/ENG	AFIT School of Engineering, Department of Electrical and Computer Engineering
BNF	Backus-Naur Form
CD	Conceptual Dependency
CoIn	Constructor/Interpreter (the natural language processor)
DBMS	Data Base Management System
DFD	Data Flow Diagram
DDS	Data Dictionary System
DDT	Data Dictionary Tool (the data dictionary acces subsystem)
DEC	Digital Equipment Corporation
HIPO	Hierarchy plus Input, Process, Output
ISL	Information Sciences Laboratory
MIT	Massachusetts Institute of Technology
NIL	New Implementation of Lisp
SADT	Structured Analysis and Design Technique
SEW	Software Development Workbench
SDWE	SDW Executive

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Abstract

The development of a human-computer interface construction and interpretation tool capable of processing English-like or natural language user input is discussed. The utility of the tool is demonstrated by using it to create the natural language interface for a data dictionary system. The data dictionary's development is also documented and is used as the overall context for the presentation.

I. Introduction

Thesis Objectives

The principal objective of this thesis investigation is the development of a human-computer interface construction and interpretation tool capable of processing English-like or natural language user input. The tool is generic in nature. That is, it has the capability to create grammars and interpret input sentences for a wide variety of applications. A secondary objective of this effort is to demonstrate the utility of the tool by using it to create a natural language interface for a useful software engineering environment.

Background

This section summarizes the background material on which the remainder of this thesis effort is based. The first part of this section provides an introduction to natural language and systems designed to interpret natural language. Next, since this thesis describes the development of a software system, a general introduction to the software development process is provided. The section concludes with a description of AFIT's Software Development Workbench (SDW). As part of this thesis effort, the developed system

was interfaced with the SDW. Numerous references are provided for the interested reader.

<u>Natural Language</u>. Natural languages are the languages that we speak, read, and write. Throughout our lives we are inundated with one or more natural languages. Therefore we are quite comfortable with natural languages. They seem "natural" to us. This feature is exactly what is needed in a computer environment to make it seem natural and consequently easy to use.

Because natural languages develop in a complex, constantly-changing, uncontrolled environment, they are themselves complex and ill-defined. They are "informal" systems. These are traits which make natural languages unnatural to a computer, traits which make it difficult for a computer to interpret natural language. Computers demand completely defined, "formal" systems. Because of this difficulty (and many others), the area of natural language interpretation has been widely researched (Bobrow, <u>et al</u>, 1977) (Harris, 1977) (Hendrix, 1977) (Hendrix, 1978) (Rich, 1984) (Schank and Riesbeck, 1981) (Waltz, 1978) (Woods, 1970).

The meaning of the term natural language when used to describe a human-computer interface is different than when used to describe a human-human interface. Human-computer natural language interfaces do not contain the richness (all

of the grammar rules or vocabulary) of, say, English. The guiding principle in building a natural language human-computer interface is to include enough of a natural language so that one can command the computer, in a "natural" way, to do all of the required functions. The interface must be complete enough so as to be a help rather that a hindrance to the user; it must be "simple" enough so that it can be interpreted by a computer within any time and space constraints imposed by an application.

A general introduction to natural language and its interpretation is provided The Handbook of Artificial Intelligence, Volume I (Barr and Feigenbaum, 1981:225-232), Artificial Intelligence (Rich, 1983:295-344), and Artificial Intelligence (Winston, 1984:291-334). In addition Barr and Feigenbaum (Barr and Feigenbaum, 1981:239-321) include an introduction to grammars and their representations and a description of several parsing techniques. They also present an overview of several natural language interpreting systems including Woods' LUNAR, Winograd's SHRDLU, Schank's MARGIE, Schank and Abelson's SAM and PAM, and SRI International's LIFER. More detailed descriptions of these and other specific systems can be found in "GUS, A Frame-Driven Dialog System" (Bobrow, 1977), "User Oriented Data Base Query with the ROBOT Natural Language Query System" (Harris, 1977), "Developing a Natural Language Interface to Complex Data"

(Hendrix, <u>et al</u>, 1978), "Human Engineering for Applied Natural Language Processing" (Hendrix, 1977), <u>Inside</u> <u>Computer Understanding</u> (Schank and Riesbeck, 1981:318-372), "An English Language Question Answering System for a Large Relational Database" (Waltz, 1978), and "Transition Network Grammars for Natural Language Analysis" (Woods, 1970). The LIFER system is of particular interest since it provided many of the ideas for the natural language processor of this investigation.

The Software Development Life Cycle. The software development life cycle has been characterized many different ways (Peters, 1981:8). In this thesis effort, the cycle is broken up into five phases. They are the functional requirements analysis phase, the design phase, the implementation phase, the integration phase, and the maintenance phase. As requirements change or errors are found, this cycle is executed iteratively.

During the functional requirements analysis phase, the emphasis is on "what" the system should do. These requirements are assigned to various hardware and software components during the design phase. Also during the design phase, the defined software components are refined into interacting modules. During implementation, the modules are encoded into a computer language and are tested individually and as groups. The hardware and software components are

assembled into a system and are subjected to testing as a whole in the integration phase. Finally, in the maintenance phase, the system is used and modified as necessary.

For a detailed discussion of the requirements phase, see <u>Structured Analysis and System Specification</u> (DeMarco, 1979). The design phase is covered in <u>Reliable Software</u> <u>Through Composite Structured Design</u> (Myers, 1975) and <u>Software Design: Methods and Techniques</u> (Peters, 1981). Various aspects of the the implementation phase are covered in <u>The Design and Analysis of Computer Algorithms</u> (Aho, <u>et</u> <u>al</u>, 1974), <u>Fundamentals of Data Structures in Pascal</u> (Horowitz and Sahni, 1984), and <u>Algorithms</u> + <u>Data Structures</u> = Programs (Wirth, 1976).

The Software Development Workbench. The Software Development Workbench (SDW), which resides on the AFIT Information Sciences Laboratory (ISL) Digital Equipment Corporation (DEC) VAX-11/780 computer, was conceived and designed to help the software engineer manage the inherent complexity of developing computer software. The SDW consists of "an integrated set of automated tools to assist the software engineer in the development of quality and maintainable software" (Hadfield and Lamont, 1983:171).

The original work on the SDW was done by Steven M. Hadfield for his master's thesis (Hadfield, 1982). In his thesis, Hadfield provided motivation for the development of

an interactive and automated software development environment. He maintained that such an environment should be integrated, traceable, flexible, and user-friendly. Eventually the SDW will consist of a comprehensive set of software development tools which will help the engineer throughout the entire software development cycle. While the current SDW is usable, it does not contain a complete set of integrated tools.

Currently, the tools contained in the SDW are integrated by a menu system, the SDW Executive or SDWE, which allows one to execute any of the tools. The menu system is hierarchical. First one chooses a category of tools such as DESIGN TOOLS from the top-level menu. A menu of the tools in the chosen category is then displayed from which one designates the particular tool to be executed.

Scope of Effort

The scope of this thesis effort includes the design and implementation of the natural language processor. The processor allows a software developer to define and implement a natural language human-computer interface. It allows the developer to construct a grammar by entering and modifying the productions of the grammar. It includes an interpreter which compares user input sentences against the

defined grammar and executes any code included with the grammar productions. The processor is generic in nature and can be used to create and use a grammar for any application domain. As part of this effort the natural language processor is integrated into the SDW.

The second area that is included within the scope of this development is the design and implementation of a data dictionary system and particularly its natural language front-end. The data dictionary system is used primarily as an extended example to demonstrate the usefulness and usability of the natural language processor, but it is also meant to be a useful tool.

Standards

This section first presents standards associated with the natural language processor and then presents standards associated with the data dictionary system. Since the emphasis is on the natural language processor, the standards directly associated with it are of greatest importance.

The natural language processor should allow a software engineer to easily construct a natural language human-computer interface to a software system. If it does, then this thesis effort should be considered a success. To fulfill this goal the grammar constructor needs to provide all of the functions necessary to create and modify a

grammar, the external interfaces to these functions should be consistent, and the interpreter should be able to correctly parse grammatical sentences for any grammar created with the processor.

The sentence interpreter should be able to parse quickly enough so as to not annoy the application user. By providing positive feedback to the user to assure him/her that the program is running and performing the desired task, this time can be extended from a few seconds to perhaps a minute or more.

The data dictionary system should allow its user to manage all of the data associated with a software development effort. Not only does the data dictionary system need to provide for the storage of this data, but its human-computer interface needs to include the functionality necessary to enter and modify the data.

Approach Taken

The software development cycle as described in the Background section is followed in this development. Throughout the chapters that follow, the natural language processor is treated as a subsystem of the data dictionary system. Doing so simplifies the structure of this thesis. Also, when the natural language processor is used to

implement a human-computer interface, it becomes a subsystem of the application.

The next chapter, Chapter II, describes the functional requirements analysis phase of this effort. The design phase is the subject of Chapter III, and the implementation phase is the subject of Chapter IV. To conclude, Chapter V summarizes this development effort, presents an evaluation of the developed systems based on the standards of the previous section, and enumerates a set of recommendations for follow-on work to this effort.

Complete documentation sets, including a data dictionary, structure charts (described in Chapter III) and various other design documents, a test plan, and user manuals for both the natural language processor and the data dictionary system, are included as appendices to this thesis. Volume 2 contains complete source code listings.

II. Requirements Definition

Introduction

This chapter presents the functional requirements used in this thesis effort. First very high level, "blue sky" requirements are described. Then the overall requirements of the data dictionary system are presented. After defining the high-level requirements, the more detailed requirements of the various subsystems are discussed.

The requirements are presented in the form of data flow diagrams (DFDs). DFDs were chosen for use in this effort because of their simplicity. Structured Analysis and Design Technique (SADT) charts (Peters, 1981:63-64), another possible representation, show more information than do DFDs, but are correspondingly more difficult to create, maintain, and understand.

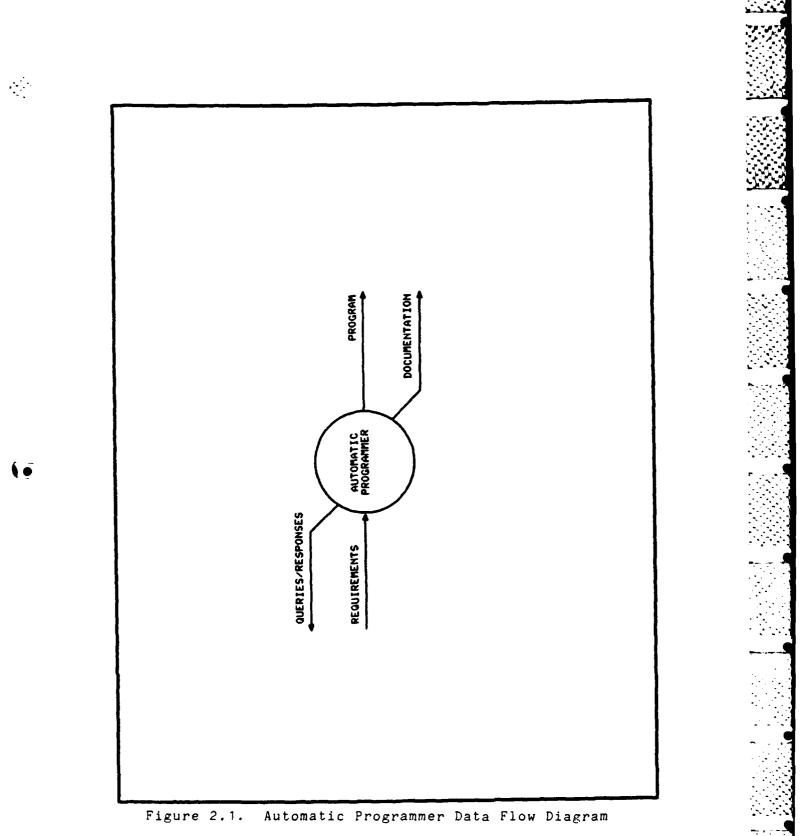
DFDs consist of four basic elements: processes, data flows, data stores, and sources/sinks. Processes, which are represented by circles, transform data. That is, they modify their inputs in some way to produce their outputs. Data flows, as their name implies, are paths along which information moves between the other three element types. Data flows are represented by arrows. Data stores are files or data bases. They are represented by a line segment or by two parallel line segments. Finally, sources and sinks,

which are represented by rectangles, are entities outside of the system which originate and receive data respectively. Each element of a DFD has associated with it a label which describes that element. For an excellent description of the mechanics and use of data flow diagrams, see <u>Structured</u> <u>Analysis and System Specification</u> (DeMarco, 1979).

In addition to the DFDs of this chapter, Appendix A contains a system data dictionary which provides more information about each of the elements contained in the DFDs.

Automatic Programming

It would be nice to have an automatic design and programming system that, given a set of inconsistent, incomplete, and ambiguous requirements, could query the user to resolve these problems and then generate, modify as requested, optimize, and fully document a program which meets the requirements (Figure 2.1). Such a system should allow input in whatever mode is most comfortable to the user, including written and spoken natural language, graphics, menu selections, examples, and mathematical formulae. Its set of output modes, for responses and queries to the user, should be similarly varied and should be user selectable. That is, if the output mode chosen by the system does not include the mode desired by the user,





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he/she should be able to specify the use of additional modes.

The documentation generated by such a system should include all documentation normally generated during the software development life cycle including, but not limited to, requirements specifications, high- and low-level design specifications (including the algorithms and data structures used), source code listings, a test plan, test specifications, and test procedures, including an extensive set of test cases. In addition, when desired by the user, the system should be able to provide correctness verification of any portion of the generated code.

All code generated by an automatic programmer should be traceable through the design to the requirements level. The same is true of the test procedures; all tests should be traceable to the requirements that they validate.

In addition, the generated system should be user-modifiable at any level--requirements, design, or implementation. The user should have ultimate control, but the system should recognize potential problems and advise the user as appropriate.

The Handbook of Artificial Intelligence, Volume II (Barr and Feigenbaum, 1982:295-379) includes an excellent introduction to automatic programming as well as an overview of the recent research in the field.

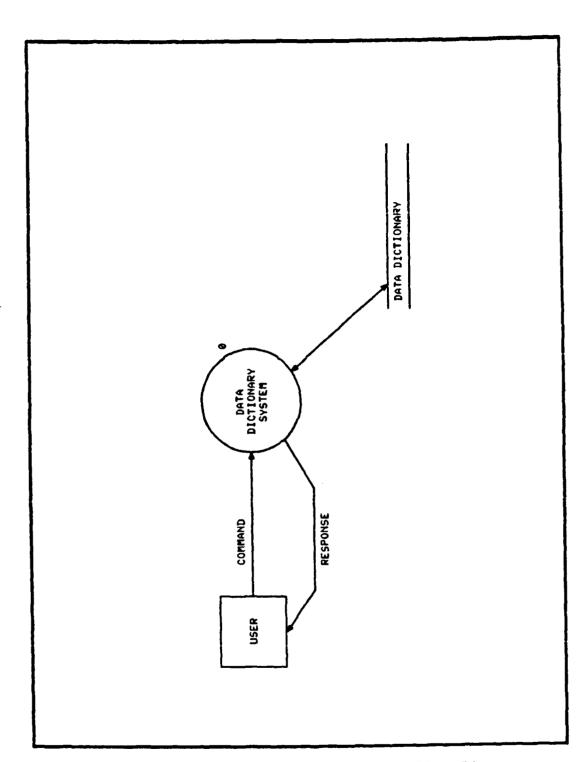
System Requirements

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The previous section described several ideas about what a computer, via an automatic programming system, should be able to do. A system which can generate quality software directly from requirements without extensive intervention by a human programmer is probably many years in the future, if it is even possible. A possibly more practical approach, given today's technology, is to concentrate on developing a set of integrated and automated tools which aid the human engineer in developing software. This is the approach taken in the SDW and is among the justifications for this thesis effort.

This section describes one such tool in terms of its functional requirements: the data dictionary system, DDS, developed for this thesis effort. Figure 2.2 shows a top-level model of the system. This model illustrates the scope of DDS: it accepts natural language input sentences from the user, interprets the sentences as commands (using the natural language processor), and retrieves information from and/or modifies the information in the data dictionary.

Figure 2.3 breaks DDS down into its major subsystems, the grammar constructor and sentence interpreter, which together comprise the natural language processor, and the data dictionary access process. The remainder of this section discusses the more detailed requirements of these

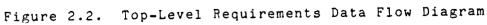


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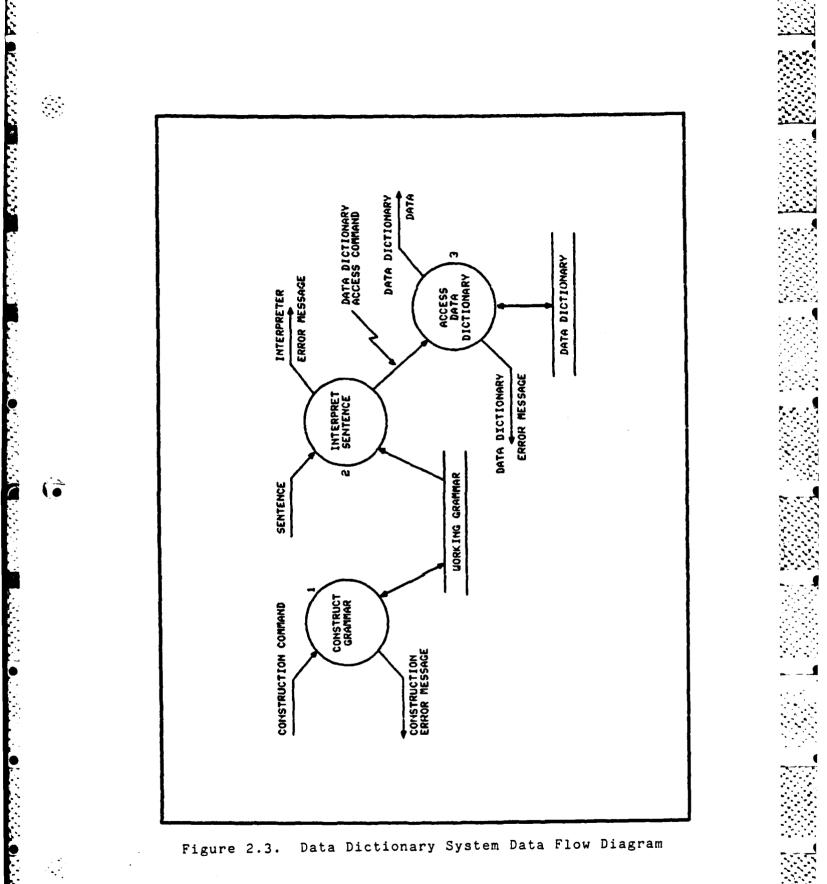
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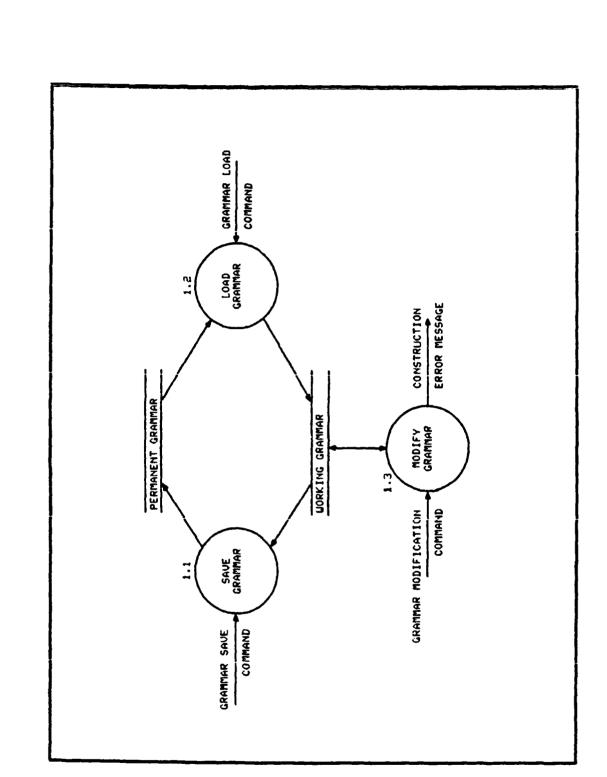
subsystems followed by a discussion of DDS's grammar requirements, its data dictionary content requirements, and, finally, its SDW interface requirements.

Grammar Constructor Requirements. The grammar constructor allows the user to define a set of production rules which collectively define a grammar. Figure 2.4 depicts the data flow for the grammar constructor. As is shown in the diagram, the grammar constructor consists of three subfunctions: SAVE GRAMMAR, LOAD GRAMMAR, and MODIFY GRAMMAR. Each of these processes is discussed in turn.

When a save command is entered by the user, the SAVE GRAMMAR process stores the current working grammar to a permanent file. This is necessary so that the grammar does not have to be created each time it is needed. The file must be in a format that allows the grammar to be retrieved into working storage.

Since the grammar is stored in a file, a method for retrieving the contents of the file is needed. This is the function of the LOAD GRAMMAR process. If a grammar is already in working storage, then it is overwritten by an incoming grammar.

The MODIFY GRAMMAR process consists of several functions which allow a user of the constructor to define the rules of a grammar. To insure that the grammar environment is in the proper condition before a grammar is



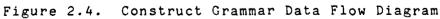
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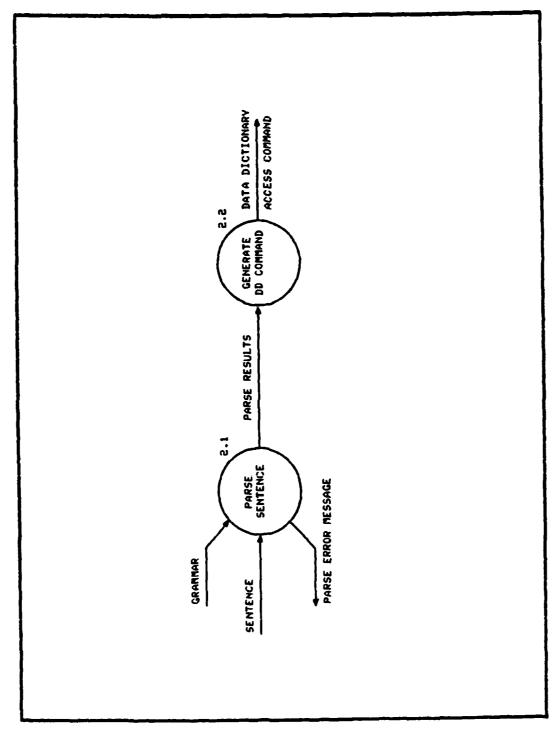
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defined, an initialization function is needed. To allow the user to initially enter the grammar's production rules, a create function is needed. To allow the user to determine what production rules exist, a list function is needed. To allow the user to view the existing production rules, a display function is needed. To allow the user to modify existing production rules, a modification function is needed. Finally, to allow the user to remove production rules which are no longer needed, a delete function is needed.

The MODIFY GRAMMAR process should be able to determine if an invalid request has been entered. If it detects an invalid request, it generates an error message.

Sentence Interpreter Requirements. The sentence interpreter has two functions: a parsing function and a command generation function. The parsing function, PARSE SENTENCE in Figure 2.5, provides the capability to determine if an input sentence is valid within the defined grammar. The parsing function also provides the capability to inform the application user that it cannot parse an input sentence. To make the system more "user-friendly", the parser should, if it is unable to parse a sentence, attempt to show the user where in an input sentence an error occurred. It should also offer suggestions as to how to correct the problem.



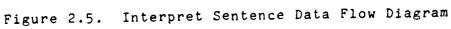
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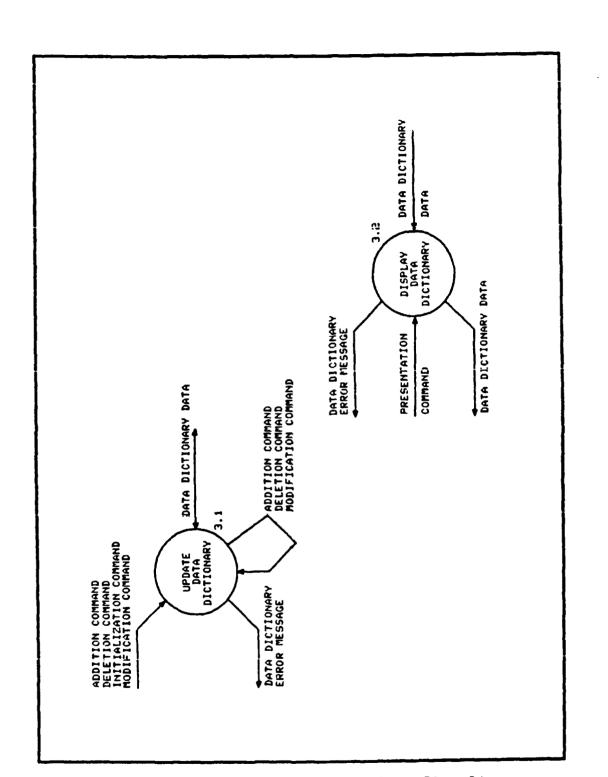
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The command generation function, GENERATE DD COMMAND in Figure 2.5, provides the capability to syntactically translate the results of a valid sentence parse into the proper execution commands of the application tool (the data dictionary access process in this case).

Data Dictionary Access Requirements. The data dictionary access process executes the commands sent from the sentence interpreter. It provides the means to update and display the data stored in the data dictionary (Figure 2.6). Therefore, it must be "complete" in the sense that it have full access to all the information stored in the data dictionary file. It must also maintain the "consistency" of the data in the data dictionary. For instance, if a module entry is deleted from the dictionary, then ALL references to that module need to be deleted. If this is not done by the access process, the information in the data dictionary can easily become inconsistent.

The UPDATE DATA DICTIONARY process shown in Figure 2.6 consists of several functions which allow the user of DDS to define the information to be stored in the data dictionary. To allow the user to define new entries in the data dictionary, an addition function is needed. To allow the user to change existing information, a modification function is needed. A reinitialization function which resets the information in an entry to its initial state can be handy,



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Figure 2.6. Access Data Dictionary Data Flow Diagram

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especially when the modifications to be made to the entry are extensive. Finally, a deletion function, which allows the user to remove an entry from the data dictionary, is needed. As is shown in Figure 2.6, the UPDATE DATA DICTIONARY process can generate its own update command. This is necessary to enforce the data consistency discussed in the previous paragraph. If the information requested to be modified is not contained in the dictionary, this process displays an error message.

The second process shown in Figure 2.6, DISPLAY DATA DICTIONARY, retrieves, formats, and presents the information stored in the data dictionary. Again, if the information requested is not contained in the dictionary, an error message is displayed.

The data dictionary access process should be independent of the natural language interpreter. That is, it should be usable without the natural language front-end. Therefore, it needs to implement all the requirements of DDS except for those that deal specifically with the natural-language human-computer interface. By making the actual data dictionary tool independent of the human-computer interface, prototypes of various kinds of interfaces and various designs of natural language interfaces can be developed and easily tested with the

access tool; then one or more can be selected and fully developed.

Grammar Requirements. The main requirements of a grammar in a natural language human-computer interface are that it be functionally "complete" for the application environment and that it be flexible. A functionally complete grammar provides at least one way for the application user to specify each function of the application. A flexible grammar provides multiple ways to specify those functions of the application that are "naturally" specified in multiple ways. Flexibility enhances ease of use.

Since DDS is to have a natural language human-computer interface, its grammar should be functionally complete and flexible. DDS's grammar should allow sentences which specify each of the operations provided by the data dictionary access mechanism. The grammar should allow access to any of the entries contained in the data dictionary. The grammar should provide for needed meta-operations, that is, operations that deal with the environment (such as exiting from the system) vice the information in data dictionary.

Data Dictionary Content Requirements. A data dictionary needs to provide storage for the information needed by an engineer during the entire software development

process. AFIT/ENG'S <u>Development Documentation Guidelines</u> and <u>Standards</u> (AFIT/ENG, 1984) was written with the intent of standardizing the composition of this information. The contents of the data dictionary should be consistent with this document since this is an AFIT thesis project.

AFIT/ENG'S documentation standard specifies several data dictionary entry types for each of the software life cycle phases. For the functional requirements analysis phase, it specifies three entry types: Activity, Data Element, and Alias. Process, Parameter, and Alias entries are specified for the design phase. The implementation phase requires Module and Variable entry types. No entries are specified for the later life cycle phases as the documentation standard is not yet complete. For each entry type, the standard describes a set of data elements which should be included in the data dictionary entry.

The data dictionary should provide for all of these entry types. As the documentation standard is updated and completed, the data dictionary contents should also be revised to remain consistent with it.

SDW Interface Requirements. Since both the data dictionary system and the natural language processor are tools which may be of help to the software developer, they should both be interfaced into the SDW. Doing so makes them easier to access and, one hopes, easier to use.

Conclusion

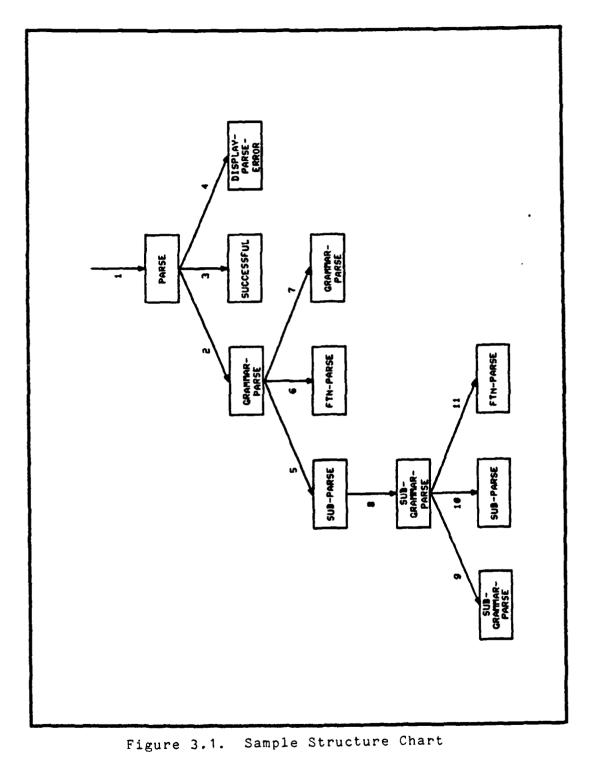
This chapter has presented and attempted to justify the results of the functional requirements analysis phase of this thesis effort. As a prelude to the discussion of the requirements, the symbology of data flow diagrams was introduced. The concept and feasibility of an automatic programming system was also briefly discussed. The next chapter describes the design phase of this effort.

III. Design

Introduction

In this chapter, the design of the natural language processing system, CoIn (<u>Constructor/Interpreter</u>), and the data dictionary system, DDS, are described and justified. The description begins with the overall functional system design of DDS. Following this, more detailed functional descriptions of CoIn's design, the data dictionary tool's (DDT) design, the grammar's design, and the grammar-to-DDT interface's design are given.

The system-level design is presented in the form of a data flow diagram (see Chapter II for a description of DFDs). The lower-level designs are in the form of structure charts which are included as Appendix B. Figure 3.1 shows an example of a structure chart. Structure charts depict the modules of a program as rectangular boxes. Each box is labeled with the name of the module it represents. The boxes are interconnected with arrows from the invoking to the invoked modules. Each arrow is labelled with a number which corresponds to a number in an accompanying table. The tables tell what information is passed between the modules. As an example, Table III-1 corresponds to the structure chart of Figure 3.1.



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parse Interfaces

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#	Passed Parameters	Туре	Returned Parameters	Туре
1	sentence grammar	data data		
2	sentence sub-grammar-list	data data	result	data
3	result	data	successful parse?	flag
4	(none)			
5	sentence sub-grammar	data data	result	data
6	sentence function	data data	result	data
7	cdr*(sentence) sub-grammar-list	data data	result	data
8	sentence sub-grammar-list	data data	result	data
9	cdr [*] (sentence) sub-grammar-list	data data	result	data
10	sentence sub-grammar	data data	result	data
11	sentence function	data data	result	data

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* The cdr function returns a list of all but the first element of a list passed to it. For instance, the cdr of the list (a b c) is the list (b c).

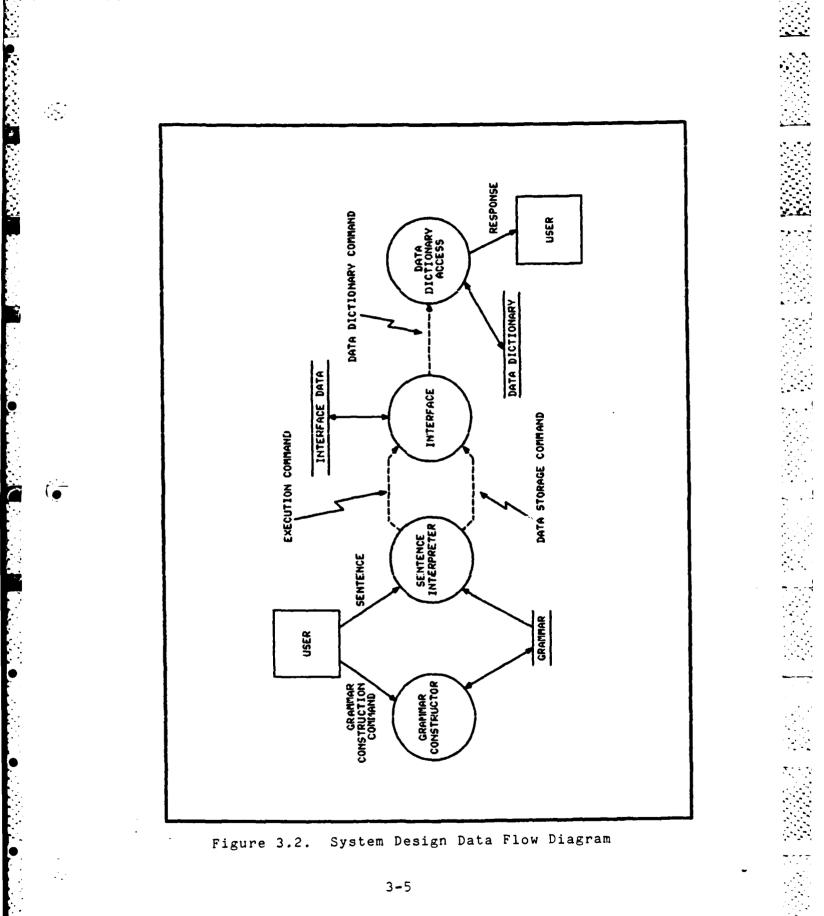
Table III-1. Sample Module Interface Table.

Structure charts were used to document the design because their notation is simple and therefore easy to use and understand, and because the SDW provides a tool (SYSFL) which partially automates their generation. Other notations which can be used to document software design include Leighton diagrams, Hierarchy plus Input, Process, Output (HIPO) charts, and structure trees. Again, for the reasons noted above, structure charts were chosen for use in this thesis effort. For more information about these design documentation methods, see <u>Software Design</u>: <u>Methods and</u> <u>Techniques</u> (Peters, 1981:44-62).

System Design

The design of DDS is divided into four functional elements: CoIn, DDT, the grammar, and the interface. As shown in Figure 3.2, CoIn consists of the GRAMMAR CONSTRUCTOR and SENTENCE INTERPRETER processes, DDT consists of the DATA DICTIONARY ACCESS process and its associated data store, the grammar consists solely of the GRAMMAR store, and finally, the interface consists of the INTERFACE process and its data store.

The EXECUTION COMMAND shown in Figure 3.2 is built into the grammar and is sent to the INTERFACE process by the SENTENCE INTERPRETER process at the end of a successful sentence parse. Similarly, the DATA STORAGE COMMANDs are



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built into the grammar and are executed by the INTERPRETER process as the input sentence is traversed. These commands tell the INTERFACE process what actions are desired and upon which data dictionary items these actions should be performed. The DATA DICTIONARY COMMANDs are generated by the INTERFACE process upon receipt of the EXECUTION COMMAND. They are based on the information in the INTERFACE DATA data store and are sent to the DATA DICTIONARY ACCESS process to be acted upon.

Natural Language Processor Design. As suggested by the requirements in Chapter II, the natural language system consists of two distinct subsystems: a grammar constructor and a sentence interpreter. The design of both of these is described in this section, but first the data structures used to represent the grammar are defined.

Data <u>Structures</u>. The first decision in the design of CoIn was the form of the data structures in which the grammar is maintained. Three structures were defined: the production record, the function record, and the meta-symbol list (Figure 3.3). These structures and their purposes are described in the following paragraphs.

<u>Production Records</u>. Production records are used to store the productions of a grammar and consist of three fields. The first field is the element which the current word of the input sentence must match for a parse to

Figure 3.3. Natural Language Processor Data Structures

proceed. This could be a literal, which the word must match exactly, the name of a sub-production record to be used as the grammar in a recursive call to the parse routine, the name of a function record to apply to the current word, an end-of-production marker which always matches but does not consume any of the input sentence, or an end-of-sentence marker which matches the end of a sentence.

The next member of a production record is a list of possible sub-productions which can be used to continue the parse should the current word parse successfully. Each of

these sub-productions are identical in structure to the parent production record. This eases the parsing task by allowing recursion.

The third and final production record element is a piece of executable source code to be invoked if the production is successfully traversed.

<u>Function Records</u>. The second structure needed by CoIn is called a function record. Function records consist of a name and a piece of executable source code which is used to determine if the word passed to a function is a member of the function's domain. If it is, then the parse of the word is considered successful.

The Meta-symbol List. The meta-symbol list structure is simply a list of the names of a grammar's production and function records. These names are called meta-symbols because they are names and not terminal symbols of the grammar. The meta-symbol list is used to keep track of the meta-symbols that have been defined so they are not accidentally redefined and so they can be easily stored by the grammar save routine.

The Constructor. The grammar constructor consists of eight subprograms which are used to build the structures described above. These include routines to initialize a grammar, create production records, create function records, add productions to an existing production record, modify an

existing function record's executable code, destroy production records, destroy function records, and save a grammar. Each of these subprograms is described below.

The grammar initialization subprogram (initialize-grammar) destroys all of the production records and function records of the current working grammar. It also empties the list of meta-symbols. This subprogram requires no input parameters.

The grammar constructor subprogram that creates production records (create-production) takes as input the name of the production record to be created and, optionally, a piece of executable code to be executed whenever the production record is successfully traversed by the interpreter. Before attempting to create a production record, this subprogram confirms that the name passed to it has not already been used.

The grammar constructor subprogram that creates function records (create-function) is similar to the one that creates production records. It too requires a name which it assigns to the function and an optional piece of executable code. It too confirms that the name passed to it has not already been used.

The subprogram that adds productions to a production record (add-production) requires three input parameters: the grammar production, a piece of executable code to be

executed whenever the production is successfully traversed by the interpreter, and the name of the production record that the production is to be added to. This subprogram first determines whether the name passed to it is a defined production record. If it is, then the subprogram adds the production and the code to the record. Otherwise, it displays an error message.

The subprogram that modifies the executable code of a function record (modify-function) requires the new code and the name of function to modify as input parameters. This subprogram determines whether the name is a defined function record and, if it is, replaces the old code with the new. Otherwise, if the name is not a defined function, it displays an error message.

The production record and function record destruction subprograms (destroy-production and destroy-function, respectively) purge the production or function record whose name matches their input parameter. If the production or function record does not exist, then an error message is displayed.

The grammar save subprogram (save-grammar) stores all of the production records and function records of the current grammar to a disk file. It also stores the list of meta-symbols. This subprogram requires no input parameters.

The Interpreter. The sentence interpreter consists of a single user-callable routine. The user provides an input sentence and a grammar to the interpreter. It attempts to match the sentence to the grammar. As a successful parse continues, the interpreter executes the associated code built into the grammar. If the interpreter is unable to successfully parse a sentence, it displays an error message on the user's terminal screen.

As was discussed in the section on data structures, the interpreter uses recursion to traverse the grammar. The objective was to model the interpreter after the data structures in order to make the code as simple and easy to understand as possible.

Data Dictionary Tool Design. DDT was designed from an object-oriented viewpoint. Goldberg defines an object as a "uniform representation of information that is an abstraction of the capabilities of a computer" (Goldberg, 1984:76). An instance of an object has associated with it a set of memory locations, called instance variables, and a set of operations, called methods, which can access the instance variables. An object's instance variables can be accessed only through its methods. This trait forces a well-defined interface to the information stored within the instance. Each "type" of object has associated with it a schema which defines the information that can be stored.

The set of objects which are instances of a particular schema is called a class. For more information about object-oriented programming, see <u>Smalltalk-80</u>, <u>The Language</u> and <u>Its Implementation</u> (Goldberg, 1984:76-80) and <u>Smalltalk-80</u>, <u>The Interactive Programming Environment</u> (Goldberg and Robson, 1983:6-9).

Object Classes. This subsection defines the object classes which make up the data dictionary. Appendix E provides a set of diagrams which show the hierarchical structure of the high-level data dictionary object classes. In these diagrams, solid lines indicate that the lower level class is explicitly part of the higher level class. Dashed lines are used when the lower level class is implicitly part of the higher level class. Figure 3.4 is an example of one of these diagrams. In this example, the PROCESS object class is only implicitly a part of the DESIGN object class. The DESIGN class actually contains a reference to the SLOT class in which is stored a set of pointers to instances of the PROCESS object class. Appendix E also includes a detailed description of each of the defined object classes including their corresponding methods. Appendix B contains structure charts for each of the defined high-level methods. Structure charts were not included for the low-level methods (i.e. methods which retrieve or set the value of a single instance variable) because of their simplicity. The

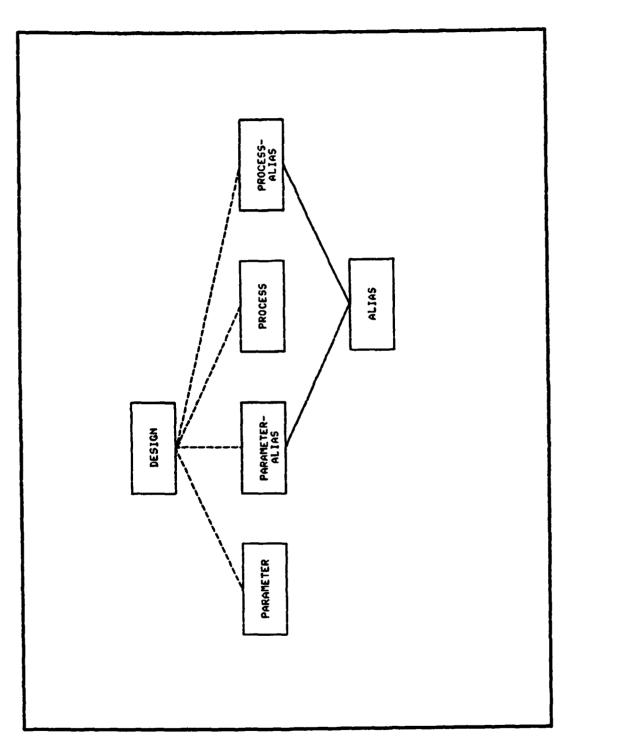


Figure 3.4. Sample Object Hierarchy Chart

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remainder of this subsection presents a textual description of the object classes.

ACTIVITY. This class corresponds to the Functional Requirements Analysis Phase 'Activity' data dictionary entry.

<u>ACTIVITY-ALIAS</u>. This class corresponds to the Functional Requirements Analysis Phase 'Alias' data dictionary entry for activities.

ALIAS. This class is a common component of the ACTIVITY-ALIAS, DATA-ELEMENT-ALIAS, PARAMETER-ALIAS, and PROCESS-ALIAS object classes. It consists of the instance variables common to all of these objects.

ASSOC. This is a primitive-level object class used to define the ASSOC data type and its operations. The ASSOC data type is similar to Lisp's association list. An association list consists of a list of pairs. The first element of the pairs can be searched by using the 'assoc' function. The 'assoc' function returns the matching pair or NIL if no match is found.

ATOM. This is a primitive-level object class used to define the ATOM data type and its operations. The ATOM data type is similar to Lisp's atom.

DATA-ELEMENT. This class corresponds to the Functional Requirements Analysis Phase 'Data Element' data dictionary entry.

DATA-ELEMENT-ALIAS. This class corresponds to the Functional Requirements Analysis Phase 'Alias' data dictionary entry for data elements.

DATE. This is a primitive-level object class used to define the DATE data type and its operations.

DESIGN. This is a high-level class which is used to keep track of all the aliases, parameters, and processes in a program design. It also is used to maintain a list of the main processes (i.e. those processes which are not invoked by another process in the design) of a program design.

ENTRY. This is a low-level object class which is used to define instance variables which contain the version number and date of entry of the object instances in which it is included. This class is included as a subclass of the ACTIVITIES, ALIAS, DATA-ELEMENT, DESIGN, IMPLEMENTATION, PARAMETER, PROCESS, and REQUIREMENTS object classes.

HEADER. This is a low-level object class which is used to define instance variables which contain the name, type, and project name of the object instances in which it is included. This class is included as a subclass of the ACTIVITIES, ALIAS, DATA-ELEMENT, DESIGN, IMPLEMENTATION, PARAMETER, PROCESS, and REQUIREMENTS object classes.

IMPLEMENTATION. This is a high-level class which is used to keep track of the modules and variables of a program. It is also used to maintain a list of main modules (i.e. those modules which are not invoked by another module in the implementation) of a program implementation.

MODULE. This class corresponds to the Implementation Phase 'Module' data dictionary entry.

LIST. This is a primitive-level object class used to define the LIST data type and its operations. The LIST data type is similar to Lisp's list.

PARAMETER. This class corresponds to the Systems Design Phase and Detailed Design Phase 'Parameter' data dictionary entries. It is also included as a subclass of the DATA ELEMENT and VARIABLE object classes.

<u>PARAMETER-ALIAS</u>. This class corresponds'to the Systems Design Phase and Detailed Design Phase 'Alias' data dictionary entries for parameters.

<u>PROCESS</u>. This class corresponds to the Systems Design Phase and Detailed Design Phase 'Process' data dictionary entries. It is also included as a subclass of the MODULE object class.

PROCESS-ALIAS. This class corresponds to the Systems Design Phase and Detailed Design Phase 'Alias' data dictionary entries for processes.

PROJECT. This is the highest-level object class. It is used as a pointer to the REQUIREMENTS, DESIGN, and IMPLEMENTATION object instances of a project.

REQUIREMENTS. This is a high-level class which is used to keep track of the activities, data elements, and aliases of a requirements analysis.

SLOT. This is a low-level class which is used to store pointers to the primitive value-storing object instances. It is also used to define which instance variables are required to be filled in and which instance variables have been filled in. Finally it is used to store a label to be printed when an instance variable's value is displayed.

TEXT. This is a primitive-level object class used to define the TEXT data type.

VARIABLE. This class corresponds to the Implementation Phase 'Variable' data dictionary entry.

Grammar Design. The design of the grammar consists of two parts. For each desired action type (e.g. add, delete, display, etc.), English-like sentences which describe the action and its object must be defined. This can be a never-ending process, since there are many sentences with the "same meaning" in the English language. The best that can be done is to try to define and implement the sentence structures that most people will use most of the time. The

minimum that must be done is to define at least one English-like way to describe each action that the data dictionary tool is capable of performing. Since the emphasis of this thesis effort was on the grammar constructor and interpreter, the second of these two approaches was taken.

The second part of the grammar definition task is to define the pieces of source code which transform the input English-like requests into a form usable by the grammar-to-DDT interface. From this perspective the grammar is really part of the interface, but, for presentation purposes, it will continue to be described separately from the interface.

Interface Design. The interface design makes use of object-oriented techniques. Here a single object class, called EVENT, is defined. This class is used to store information about what actions are to be executed by the data dictionary tool. The methods of EVENT send messages to the data dictionary to perform the desired actions. Appendix F provides a description of EVENT and its methods.

Conclusion

This chapter has described the design of DDS including the natural language processor, CoIn, and DDS's other subsystems. Structure chart and object-oriented design

methods were introduced and were then used in the design specification. Chapter IV continues the description of this effort by presenting the implementation phase.

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IV. Implementation

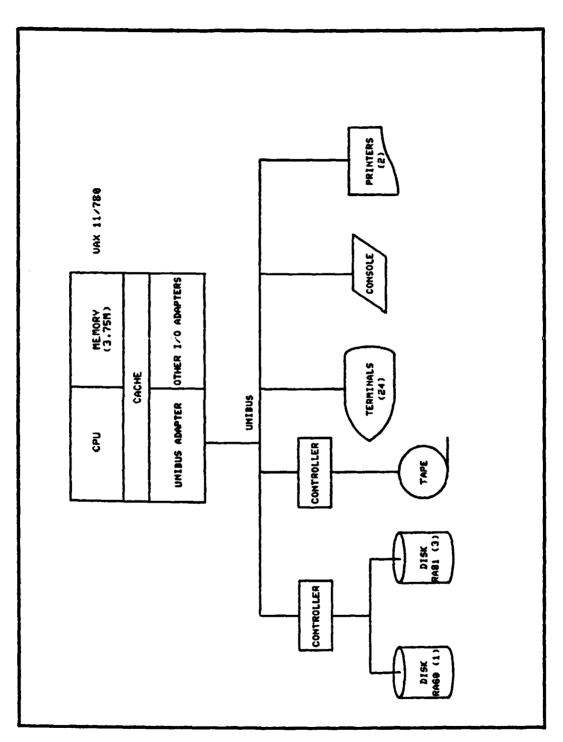
Introduction

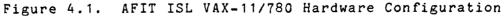
This chapter discusses the implementation of DDS including CoIn. As in previous chapters, the system level is looked at first, followed by a more detailed discussion of each of its subsystems. The final section of this chapter describes the integration of CoIn and DDS into the Software Development Workbench. Volume 2 of this thesis contains complete source code listings of the data dictionary system.

System Implementation

DDS was implemented on the AFIT Information Sciences Laboratory's DEC VAX-11/780 minicomputer under the VMS operating system Version 3.6. Figure 4.1 is a diagram of the configuration of this computer system during the time DDS was being implemented. The ISL VAX was chosen as the target machine for two reasons. First, it was available and not overloaded compared to other machines at AFIT. Second, it is the host machine of the SDW (Hadfield, 1982:17-18).

After deciding upon the target machine, the next major decision made during the implementation phase was the choice of the implementation programming language. Several languages were available on the target machine at the beginning of this phase, including C, Fortran, Lisp, Pascal,





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المربع المرب المربع and Prolog. Of these, Lisp was chosen for three reasons. These are described in the next paragraphs.

The first and most important reason for choosing Lisp as the implementation language is the development environment provided by Lisp systems. Lisp systems generally provide an interpreter, compiler, editor, and debugger combined into one nicely integrated environment. This allows the Lisp programmer to easily jump back and forth between these different tools. As each subroutine is developed, it can be tested; debugged, modified, or redesigned as necessary; compiled into object code; and integrated into the system, all without leaving the Lisp environment.

The second reason for choosing Lisp is pedagogical in nature: a student of artificial intelligence is generally expected to learn to program in Lisp. Writing code in Lisp is a necessary part of fulfilling this goal. The thesis effort provided an excellent opportunity to pursue this goal.

The existence of a large amount of Lisp code that could possibly be used in DDS was the third reason for choosing Lisp as the implementation language. Prototypes of both the grammar constructor/interpreter and the software design part of the data dictionary tool were built in other projects (Wolfe, 1985a) (Wolfe, 1985b).

The particular implementation of Lisp used is called NIL (Burke, 1984) which is an acronym for New Implementation of Lisp. NIL was developed at MIT and is based on Common Lisp (Steele, 1984). NIL is a fairly complete implementation of Lisp. Its most serious lack is that it does not include a garbage collector (a routine which reclaims discarded memory cells). The lack of a garbage collector means that one must periodically exit NIL, restart it, and then reload DDS and one's database. Fortunately, this needs to be done seldom enough that it should not be a major problem. The reason for choosing NIL was that it was the only version of Lisp installed on the ISL VAX-11 computer at the time this effort was begun.

While the entire DDS is implemented in Lisp, parts of it (the data dictionary tool and the interface) are implemented using an object-oriented language built on top of Lisp. This language is called Flavors and is included in NIL (Burke, 1984:170-178). Flavors is an environment which allows one to define and manipulate objects (Chapter III contains a discussion of objects). Since Flavors is built on top of Lisp, one can still access all of the functions of Lisp. This makes the Flavors system ideal for implementing an object-oriented system which must interface with a system written in Lisp.

<u>CoIn Implementation</u>. This subsection describes the implementation of the natural language processor. CoIn was by far the most difficult part of the data dictionary system to design and implement. Implementation of the grammar data structures was straightforward--the record structures discussed in the last chapter were simply implemented as Lisp lists. However, the grammar constructor routines were difficult to implement properly, and the development of a workable design for the sentence interpreter was an iterative process of prototype development and throwaway.

The Grammar Constructor. The grammar constructor consists of an implementation of all of the subprogram designs described in Chapter III. Each of the subprograms, except for the grammar initialization and saving routines, is implemented as a Lisp macro (along with any needed lower-level subfunctions). Macros were used to simplify the user interface. Unlike a Lisp function, a macro does not evaluate its arguments before the body of its code is executed. Since the parameters that are passed to these macros are not meant to be evaluated, using macros saves the user the trouble of quoting the parameters. Since neither the grammar initialization nor the grammar saving subprograms require any arguments, they were implemented as Lisp functions.

Implementing the grammar constructor presented a couple of problems. One difficulty was the necessity of making the constructor's routines "destructive" in nature. That is, any changes made to the grammar using the constructor must permanently alter the global data structures of the grammar. Otherwise, if the change is not global, it will be lost, and the grammar will remain unchanged. In order to facilitate this requirement, a slightly modified version of the editor presented in Chapter 7 of Artificial Intelligence Programming (Charniak, et al, 1980:84-97) was heavily relied upon in the constructor implementation. This editor is destructive in nature. The original editor was designed to edit Lisp function definitions, but since the data structures of a grammar are not functions but lists, it was necessary to modify the editor to enable it to edit any Lisp symbolic expression. A side benefit of implementing this editor is that not only is it used in the grammar constructor, but it can be used as a standalone tool to edit the grammar data structures or any other Lisp symbolic expressions one wishes to modify.

Keeping track of the positions of the expression editor global pointers within the grammar data structures was a second source of difficulty in the constructor implementation. The data structures of a complex grammar are themselves complex. To overcome these problems, many

hours of tracing code while keeping track of pointers on a hand-drawn representation of the structures was done.

The Sentence Interpreter. As mentioned previously, the design and implementation of the sentence interpreter turned out to be an iterative process. The original version of the interpreter was developed as part of a project in EE 6.23, Artificial Intelligence System Design (Wolfe, 1985a). This version provided insight into the problem but no executing code.

The second design/implementation followed the form of the data structures much more closely. It was capable of correctly interpreting a grammatical sentence (if the grammar had been carefully designed) but could not always recognize a nongrammatical sentence as being invalid. The problem was that not enough information was being returned by the low-level subroutines for the driver subroutine to recognize an error. Correcting this problem was one of the major changes in the third (and present) version of the sentence interpreter.

The present module structure of the sentence interpreter is very similar to the previous one. Most of the differences are in the detailed design and implementation of the mid-level subroutines. Both the high-level driver routines and the low-level data structure access routines are essentially unchanged from the second

version. However, the module interfaces between the high-level and the mid-level routines were modified so as to return the information necessary for the top-level module to detect nongrammatical input sentences.

The interpreter parses an input sentence left-to-right comparing the words of the sentence to the grammar. As it successfully parses the sentence, it executes the associated Lisp source code contained in the grammar. If it is unable to fully parse a sentence using a production of the grammar, it backtracks to a previous branch point in the grammar and tries again. If the interpreter is unable to fully parse a sentence using any of the grammar productions, then the sentence is considered nongrammatical and an error handling routine is invoked. Currently the error handling routine displays only a message that the sentence was invalid within the context of the supplied grammar.

Data Dictionary Implementation. The data dictionary tool (DDT) was implemented as a set of Flavors objects and their associated methods. For most of the defined object classes of the data dictionary tool (not all were implemented due to time constraints), a Flavors object was declared. The slots of the Flavors objects correspond exactly to the instance variables defined in Appendix E. Similarly, the implemented Flavors access methods correspond exactly to the defined methods listed in Appendix E.

Since the low-level objects were needed by all of the higher level objects, they were implemented first (except for the TEXT object which was not implemented in this effort). Next the objects associated with the software design phase were implemented. This choice was made because some of the code already existed from a prior project (discussed above). Finally the highest-level object (PROJECT) was implemented. Also the objects concerned with the requirements phase and implementation phase were minimally implemented at this time.

<u>Grammar Implementation</u>. Once the grammar had been defined, implementing it was quite straightforward. The grammar productions of Appendix C were entered verbatim using the grammar construction routines of CoIn. Figure 4.2 shows an example of one of these productions (the notation used in the figure is defined in Appendix C). Entry errors were corrected either by making use of the expression editor in CoIn or by deleting and reentering the erroneous productions. Implementation of the grammar brought out a need for several more grammar modification subprograms. These are further discussed in Chapter V. The Lisp code that was included within the grammar is really part of the interface, so it is discussed in the next subsection.

<grammar> ::=</grammar>
(please) <aux-grammar></aux-grammar>

Figure 4.2. Sample Grammar Production

The defined grammar was not completely implemented during this thesis effort. Effort was concentrated on the presentation part of the grammar as it was thought to be simple enough to fully implement, yet complex enough to be a valid demonstration of that CoIn can be used to implement a useful natural language human-computer interface. The presentation part of the grammar was considered relatively simple because during presentation of the data, there is no new information being added. The information that is contained in an input sentence can all be checked word-for-word against the grammar and database. The presentation part of the grammar was considered complex enough because it needs to have the capability of accessing the entire database, and it makes use of all the capabilities of the interpreter.

Other parts of the grammar that were implemented during this thesis effort include the initialization productions and the data save and quit commands. The initialization

productions were implemented because it was quite simple to do so, and the data save and quit commands because of their obvious importance.

Interface Implementation. The interface between the sentence interpreter and DDT was implemented as a single Flavors object class called EVENT. A set of methods to access this class and the internal Lisp code of the grammar, which invokes these methods, constitute the remainder of the interface implementation.

The EVENT object class was modeled after Schank's Conceptual Dependency (CD) theory (Schank and Riesbeck, 1981:10-26). According to the CD theory, every event has an actor, an action, an object, and a direction. The actor is the entity which performs the action. The action is performed upon an object and is oriented in a direction. In the case of the interface, the actor is always the computer, the action is one of the defined commands of the interface, the object is an instance of one of the data dictionary's objects, and the direction is either from the data base to the user's terminal or vice versa.

A second idea of CD theory is best described by Schank and Riesbeck: "When two sentences describe the same event in such a way that these descriptions have the same overall meaning but quite different forms, we expect out CD

representations to be identical for both descriptions." (Schank and Riesbeck, 1981:14). This idea was also used in the design and implementation of the interface. A set of command types corresponding to the operations of DDT was defined in the grammar. The various sentence forms of each of the command types is mapped into an identical representation by the interface. Upon completion of a successful parse, this representation is converted into a command that can be executed directly by the data dictionary tool. Any response by DDT is sent directly to the user's terminal.

Integration into the SDW

Using the <u>Software Development Workbench Executive</u> <u>Maintenance Guide</u> (Hadfield, 1982:355-364) as a guide, DDS and CoIn were integrated into the SDW. The first decision made for each of these systems was the choice of a two letter code to be entered by the SDW user to invoke the tool. The codes 'DD' for DDS and 'NL' for CoIn were chosen because they were available and because of their obvious mnemonic nature.

The second decision made was to determine to which of the SDW's functional groups of tools each system should be added. Since the design phase part of DDS was concentrated upon during implementation, it was decided to initially add DDS to the Design Tools functional group of the SDW. When

the requirements phase part of DDS is completed, DDS should be added to the Requirements Definition Tools functional group. Likewise, when the implementation phase part of DDS is completed, DDS should be added to one of the implementation tools functional groups, possibly the Text Editors group. The grammar constructor portion of CoIn is essentially a grammar editor, so CoIn was added to the Text Editors group.

Conclusion

This chapter discussed the implementation of DDS and its subsystems, including CoIn. The choice of the target machine was discussed, as was the choice of implementation language. This chapter concluded with a discussion concerning the integration into the SDW of DDS and CoIn. The next chapter, Chapter V, concludes this thesis by reviewing and analyzing this thesis effort and providing suggestions for further work.

V. Conclusion and Recommendations

Introduction

This thesis has described the development of a natural language processor, CoIn, and its application to a data dictionary system. DDS, the data dictionary system, consists of several distinct but interacting parts: a natural language human-computer interface implemented using CoIn, a data dictionary access mechanism, and an interface between the grammar and the data dictionary tool. The emphasis in DDS was on the first of these subsystems. Although the other subsystems make DDS a usable system, their primary role was to show that the interface constructor and interpreter is a useful and usable tool. This final chapter first presents a short summary of the system development. Following this, an analysis which relates the current system back to the standards described in Chapter I is presented. Finally, a list of recommendations for the completion and enhancement of the system is included.

Development Summary

CoIn and DDS were built using a variation of the classic software development life cycle. First an extensive literature search was done to gain a better understanding of

natural language systems, the software development process and its problems, and how these problems might be diminished by automation. The information gleaned from this search, along with prior knowledge, was used as the input to the requirements analysis phase. In this phase, sets of requirements were defined for the natural language processor, for the entire data dictionary system, and for each of its subsystems.

After generating the initial set of requirements, an iterative process of design-implement-test had to be done. This process provided necessary feedback about the completeness and consistency of the requirements. This feedback was used to modify the requirements as required. The design and implementation of the sentence interpreter in particular was a cyclic process of prototype development. Throughout the implementation, routines were tested as they were developed, both isolated from and integrated into the system. The testing done is by far inadequate (mainly due to its informality) but does suggest that DDS is reasonably error-free.

Analysis of the Current System

As has been shown in this thesis effort, CoIn does allow one to construct a usable natural language human-computer interface. The constructor does provide at

least the minimum set of routines necessary for building a grammar, and their interfaces do seem consistent with each other. The interpreter is able to correctly parse grammatical sentences within the implemented DDS grammar (Figure 5.1). The interpreter is also able to recognize and report non-grammatical sentences (Figure 5.2). Therefore, from the point of view of the natural language processor, this project should be considered a success.

The data dictionary itself is not complete and needs to be extended. This is discussed more in the next section. Enough of the system is implemented to show that it is a reasonable project. What is implemented does show the utility of the natural language processor, so the data dictionary should be considered fairly successful. Finishing DDT should be easy, albeit time-consuming.

Finally, DDS does operate in a reasonable amount of time. Several timing test were made with an moderate load on the computer system (4 users logged in). DDS was able to consistently respond in 1 to 3 seconds.

Recommendations for Future Work

DDS is incomplete. Not all of the requirements have been fulfilled by the design. Not all of the design has been implemented. More, formal testing, in accordance with the test plan (see Appendix G), needs to be done. The

Data Dictionary System (DDS) Type (help) for help. DDS-> (please show me the version and date of process is-process) Entry Version: 1 Entry Date: 29 Aug 85 DDS->

Figure 5.1. Example of a Valid Sentence Entered into DDS

DDS-> (what are the calling processes of foo alias1) *** Error - Can't parse sentence *** DDS->

Figure 5.2. Example of an Invalid Sentence Entered into DDS

following subsections describe some of the work that needs to be done for DDS to become a usable system.

<u>COIN</u>. Of the subsystems, COIN is probably the closest to being a "complete" implementation. This is natural since the emphasis of the thesis work was on this part of the system. The most severe lack in COIN is in its handling of nongrammatical input sentences. The sentence interpreter's error handling routine needs to be greatly enhanced. As it

is implemented now, it provides virtually no information to a user of the system. This is a serious limitation and needs to be investigated and corrected as soon as possible.

In Chapter IV, the statement was made that several grammar modification subprograms need to be added to the grammar constructor. There are at least two such routines. First, a routine to remove a production from a production record is needed. This can currently be done either by destroying the production record and reentering all but the production to be deleted or by using the s-expression error. Neither of these solutions is good a good one. The first can be very time consuming; the second, due to the recursive nature of production records, is difficult and error-prone.

The second subprogram that needs to be added to the grammar constructor is one that would allow the modification of the Lisp source code contained in the grammar. The need for this routine is not as critical as the need for the production deletion routine--modification of the Lisp code is fairly easy using the s-expression editor. This is more of a "nice to have" routine.

DDT. The emphasis in DDT was on the objects needed during the design phase of the software development life cycle. Therefore, the implementation of these objects is closer to being complete than the objects associated with the requirements or implementation phase. However, even the

design phase objects still need work. Files for source code have been created and instance variables have been defined for all of the object classes. Several of these object definitions have not yet been entered. While the defined access methods have all been implemented, the access methods of several object classes have not yet been defined. Before DDS can be used throughout the life cycle, these deficiencies must be eliminated.

DDT was implemented using the Flavors language of NIL. A subject worth investigating is the rehosting of DDT using a data base management system (DBMS). Unfortunately, NIL provides no means for interfacing to an external DBMS. However, DEC's VAX Lisp is now available on the ISL VAX. VAX Lisp may provide the capability needed to access external systems. Whether this is indeed the case needs to be determined.

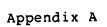
The Grammar. It is felt that the grammar as defined is adequate for interacting with DDS. This is not to imply that the grammar is all inclusive, but that it does provide at least one way to specify each of the possible operations of DDT. Extending and refining the DDS's grammar could probably be a thesis project all by itself.

One of the primary deficiencies of the grammar is that it was not completely implemented during this thesis effort. Completing the implementation should be possible for a Lisp

programmer with reasonable knowledge of the structure of DDS.

The Interface. How close the interface is to being complete is difficult to judge. As the grammar is implemented, the corresponding Lisp interface code must be included. Whether changes will be necessary to the rest of the interface depends on how this code is implemented.





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System Data Dictionary



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Introduction

This appendix contains the data dictionary for DDS and CoIn. Currently there are four entry types: ACTIVITY, ALIAS, DATA FLOW, and PROCESS. The ACTIVITY, ALIAS, and DATA FLOW types are all associated with the functional requirements analysis phase of the software development lifecycle. The PROCESS type is associated with the design phase.

DDS System Data Dictionary

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ACCESS DATA DICTIONARY NAME: TYPE: ACTIVITY **PROJECT:** DDS NUMBER: DESCRIPTION: INPUTS: DATA DICTIONARY ACCESS COMMAND DATA DICTIONARY DATA DATA DICTIONARY DATA OUTPUTS: DATA DICTIONARY ERROR MESSAGE CONTROLS: **MECHANISMS:** ALIASES: PARENT ACTIVITY: DATA DICTIONARY SYSTEM VERSION: 1 22 OCT 85 DATE: NAME: ADD-PRODUCTION (MACRO) TYPE: PROCESS **PROJECT:** DDS NUMBER: DESCRIPTION: INPUT DATA: PRODUCTION CODE TO EXECUTE PRODUCTION RECORD NAME INPUT FLAGS: OUTPUT DATA: **OUTPUT FLAGS:** GLOBAL DATA USED: *EDIT-EXP* GLOBAL DATA CHANGED: *EDIT-EXP* FILES READ: FILES WRITTEN: HARDWARE READ: HARDWARE WRITTEN: ALIASES: CALLING PROCESSES: **PROCESSES CALLED:** IS-PRODUCTION ED-RESET AUX-ADD-PRODUCTION **RELATED ACTIVITY:** VERSION: 1 29 OCT 85 DATE:

NAME: TYPE: PROJECT: PART OF: COMPOSITION: ALIASES: SOURCES: **DESTINATIONS:** VERSION: DATE:

NAME: TYPE: **PROJECT:** NUMBER: DESCRIPTION: INPUT DATA:

INPUT FLAGS: OUTPUT DATA: OUTPUT FLAGS: GLOBAL DATA USED: GLOBAL DATA CHANGED: FILES READ: FILES WRITTEN: HARDWARE READ: HARDWARE WRITTEN: ALIASES: CALLING PROCESSES: **PROCESSES CALLED:**

ED-RIGHT INSERT-PRODUCTION ED-DOWN LOOP& AUX-ADD-PRODUCTION (RECURSIVE) ED-UP

RELATED ACTIVITY: VERSION: DATE:

A-4

ADDITION COMMAND DATA FLOW DDS DATA DICTIONARY ACCESS COMMAND

UPDATE DATA DICTIONARY UPDATE DATA DICTIONARY 2 22 OCT 85

AUX-ADD-PRODUCTION PROCESS DDS

PRODUCTION CODE TO EXECUTE

EDIT-PTR

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29 OCT 85





NAME: TYPE: PROJECT: PART OF: COMPOSITION:

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ALIASES: SOURCES: DESTINATIONS: VERSION: DATE:

NAME: TYPE: PROJECT: NUMBER: DESCRIPTION: INPUTS:

OUTPUTS:

CONTROLS: MECHANISMS: ALIASES: PARENT ACTIVITY: VERSION: DATE:

NAME: TYPE: PROJECT: PART OF: COMPOSITION:

ALIASES: SOURCES: DESTINATIONS: VERSION: DATE: COMMAND DATA FLOW DDS

CONSTRUCTION COMMAND SENTENCE

USER DATA DICTIONARY SYSTEM 2 22 OCT 85

CONSTRUCT GRAMMAR ACTIVITY DDS

CONSTRUCTION COMMAND GRAMMAR CONSTRUCTION ERROR MESSAGE GRAMMAR

DATA DICTIONARY SYSTEM 2 22 OCT 85

CONSTRUCTION COMMAND DATA FLOW DDS COMMAND GRAMMAR LOAD COMMAND GRAMMAR MODIFICATION COMMAND GRAMMAR SAVE COMMAND

CONSTRUCT GRAMMAR 2 22 OCT 85 NAME: TYPE: PROJECT: PART OF: COMPOSITION: ALIASES: SOURCES: **DESTINATIONS:** VERSION: DATE: NAME: TYPE: PROJECT: NUMBER: DESCRIPTION: INPUT DATA: INPUT FLAGS: OUTPUT DATA: OUTPUT FLAGS: GLOBAL DATA USED: GLOBAL DATA CHANGED: FILES READ: FILES WRITTEN: HARDWARE READ: HARDWARE WRITTEN: ALIASES: CALLING PROCESSES: PROCESSES CALLED:

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RELATED ACTIVITY: VERSION: DATE: CONSTRUCTION ERROR MESSAGE DATA FLOW DDS RESPONSE

CONSTRUCT GRAMMAR MODIFY GRAMMAR

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CREATE-FUNCTION (MACRO) PROCESS DDS

FUNCTION RECORD NAME FUNCTION CODE

CREATE-META-SYMBOL SETF* MSG

29 OCT 85

NAME: TYPE: PROJECT: NUMBER: DESCRIPTION: INPUT DATA: INPUT FLAGS: OUTPUT DATA: OUTPUT FLAGS: GLOBAL DATA USED: GLOBAL DATA CHANGED: *META-SYMBOLS* FILES READ: FILES WRITTEN: HARDWARE READ: HARDWARE WRITTEN: ALIASES: CALLING PROCESSES: PROCESSES CALLED: RELATED ACTIVITY: VERSION: DATE: NAME: TYPE: PROJECT: NUMBER: DESCRIPTION: INPUT DATA:

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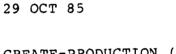
INPUT FLAGS: OUTPUT DATA: OUTPUT FLAGS: GLOBAL DATA USED: GLOBAL DATA CHANGED: FILES READ: FILES WRITTEN: HARDWARE READ: HARDWARE WRITTEN: ALIASES: CALLING PROCESSES: PROCESSES CALLED: **RELATED ACTIVITY:** VERSION:

DATE:

CREATE-META-SYMBOL SETF* MSG 1 29 OCT 85

A-7

CREATE-PRODUCTION (MACRO) PROCESS DDS PRODUCTION RECORD NAME CODE TO EXECUTE



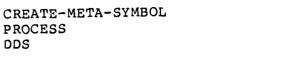
PROCESS

META-SYMBOL NAME

META-SYMBOLS

DDS

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NAME: TYPE: PROJECT: PART OF: COMPOSITION:

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ALIASES: SOURCES:

DESTINATIONS: VERSION: DATE:

NAME: TYPE: PROJECT: PART OF: COMPOSITION: ALIASES: SOURCES:

DESTINATIONS:

VERSION: DATE:

NAME: TYPE: PROJECT: PART OF: COMPOSITION: ALIASES: SOURCES:

DESTINATIONS: VERSION: DATE: DATA DICTIONARY ACCESS COMMAND DATA FLOW DDS

ADDITION COMMAND DELETION COMMAND INITIALIZATION COMMAND MODIFICATION COMMAND PRESENTATION COMMAND

GENERATE DD COMMAND INTERPRET SENTENCE ACCESS DATA DICTIONARY 1 22 OCT 85

DATA DICTIONARY DATA DATA FLOW DDS RESPONSE

ACCESS DATA DICTIONARY DISPLAY DATA DICTIONARY UPDATE DICTIONARY ACCESS DATA DICTIONARY DISPLAY DATA DICTIONARY UPDATE DATA DICTIONARY 1 22 OCT 85

DATA DICTIONARY ERROR MESSAGE DATA FLOW DDS RESPONSE

ACCESS DATA DICTIONARY DISPLAY DATA DICTIONARY UPDATE DATA DICTIONARY

1 22 OCT 85

NAME: TYPE: PROJECT: NUMBER: DESCRIPTION: INPUTS:

OUTPUTS:

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CONTROLS: MECHANISMS: ALIASES: PARENT ACTIVITY: VERSION: DATE:

NAME: TYPE: DD TYPE: PROJECT: NUMBER: DESCRIPTION: SYNONYM: VERSION: DATE:

NAME: TYPE: PROJECT: PART OF: COMPOSITION: ALIASES: SOURCES: DESTINATIONS: VERSION: DATE: DATA DICTIONARY SYSTEM ACTIVITY DDS

COMMAND DATA DICTIONARY DATA DICTIONARY RESPONSE

DDS (NONE) 2 22 OCT 85

DDS ALIAS ACTIVITY DDS

DATA DICTIONARY SYSTEM 2 22 OCT 85

DELETION COMMAND DATA FLOW DDS DATA DICTIONARY ACCESS COMMAND UPDATE DATA DICTIONARY UPDATE DATA DICTIONARY 2

22 OCT 85

NAME: TYPE: PROJECT: NUMBER: DESCRIPTION: INPUT DATA: INPUT FLAGS: OUTPUT DATA: OUTPUT FLAGS: GLOBAL DATA USED: GLOBAL DATA CHANGED: FILES READ: FILES WRITTEN: HARDWARE READ: HARDWARE WRITTEN: ALIASES: CALLING PROCESSES: PROCESSES CALLED: RELATED ACTIVITY: VERSION: DATE: NAME: TYPE: **PROJECT:** NUMBER: DESCRIPTION: INPUT DATA: INPUT FLAGS: OUTPUT DATA: OUTPUT FLAGS: GLOBAL DATA USED: GLOBAL DATA CHANGED: FILES READ: FILES WRITTEN: HARDWARE READ: HARDWARE WRITTEN: ALIASES: CALLING PROCESSES: PROCESSES CALLED: **RELATED ACTIVITY:** VERSION: DATE:

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FUNCTION RECORD NAME IS-FUNCTION DESTROY-META-SYMBOL MSG 1 29 OCT 85 DESTROY-META-SYMBOL PROCESS DDS META-SYMBOL NAME *META~SYMBOLS* *META-SYMBOLS*

DESTROY-FUNCTION

PROCESS

DDS

IS-META-SYMBOL 1 29 OCT 85

NAME: TYPE: DDS **PROJECT:** NUMBER: **DESCRIPTION:** INPUT DATA: INPUT FLAGS: OUTPUT DATA: **OUTPUT FLAGS:** GLOBAL DATA USED: GLOBAL DATA CHANGED: FILES READ: FILES WRITTEN: HARDWARE READ: HARDWARE WRITTEN: ALIASES: CALLING PROCESSES: PROCESSES CALLED: MSG **RELATED ACTIVITY:** VERSION: 1 DATE: NAME: TYPE: PROJECT: DDS NUMBER: **DESCRIPTION:** INPUTS: OUTPUTS: CONTROLS: **MECHANISMS:** ALIASES: PARENT ACTIVITY:

VERSION:

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DESTROY-PRODUCTION (MACRO) PROCESS

PRODUCTION RECORD NAME

IS-PRODUCTION DESTROY-META-SYMBOL MSG

DISPLAY DATA DICTIONARY ACTIVITY DDS

DATA DICTIONARY DATA PRESENTATION COMMAND DATA DICTIONARY DATA DATA DICTIONARY ERROR MESSAGE .

ACCESS DATA DICTIONARY 2 22 OCT 85

DISPLAY-PARSE-ERROR NAME: PROCESS TYPE: DDS **PROJECT:** NUMBER: DESCRIPTION: INPUT DATA: INPUT FLAGS: OUTPUT DATA: OUTPUT FLAGS: GLOBAL DATA USED: GLOBAL DATA CHANGED: FILES READ: FILES WRITTEN: HARDWARE READ: HARDWARE WRITTEN: ALIASES: CALLING PROCESSES: PROCESSES CALLED: RELATED ACTIVITY: 1 VERSION: DATE: NAME: TYPE: PROJECT: NUMBER: DESCRIPTION: INPUT DATA: INPUT FLAGS: OUTPUT DATA: OUTPUT FLAGS: GLOBAL DATA USED: GLOBAL DATA CHANGED: FILES READ: FILES WRITTEN: HARDWARE READ: HARDWARE WRITTEN: ALIASES: CALLING PROCESSES: **PROCESSES CALLED: RELATED ACTIVITY:**

VERSION:

DATE:

MSG 04 NOV 85 FTN-CODE PROCESS DDS FUNCTION RECORD FUNCTION CODE

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NAME: FTN-NAME TYPE: PROCESS PROJECT: DDS NUMBER: **DESCRIPTION:** INPUT DATA: FUNCTION RECORD INPUT FLAGS: OUTPUT DATA: FUNCTION NAME OUTPUT FLAGS: GLOBAL DATA USED: GLOBAL DATA CHANGED: FILES READ: FILES WRITTEN: HARDWARE READ: HARDWARE WRITTEN: ALIASES: CALLING PROCESSES: **PROCESSES CALLED: RELATED ACTIVITY:** VERSION: 1 DATE: 04 NOV 85 FTN-PARSE NAME: TYPE: PROCESS **PROJECT:** DDS NUMBER: DESCRIPTION: INPUT DATA: SENTENCE FUNCTION RECORD INPUT FLAGS: OUTPUT DATA: RESULT **OUTPUT FLAGS:** GLOBAL DATA USED: GLOBAL DATA CHANGED: FILES READ: FILES WRITTEN: HARDWARE READ: HARDWARE WRITTEN: ALIASES: CALLING PROCESSES: **PROCESSES CALLED:** FTN-NAME FTN-CODE

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RELATED ACTIVITY: VERSION: DATE:

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NAME: TYPE: PROJECT: NUMBER: DESCRIPTION: INPUTS: OUTPUTS: CONTROLS: MECHANISMS: ALIASES: PARENT ACTIVITY: VERSION: DATE:

GENERATE DD COMMAND ACTIVITY DDS

PARSE RESULTS DATA DICTIONARY ACCESS COMMAND

INTERPRET SENTENCE 1 22 OCT 85

GRAMMAR

DDS

DATA FLOW

NAME: TYPE: PROJECT: PART OF: COMPOSITION: ALIASES: SOURCES:

DESTINATIONS:

VERSION: DATE: CONSTRUCT GRAMMAR LOAD GRAMMAR MODIFY GRAMMAR PERMANENT GRAMMAR SAVE GRAMMAR WORKING GRAMMAR CONSTRUCT GRAMMAR INTERPRET SENTENCE LOAD GRAMMAR MODIFY GRAMMAR PARSE SENTENCE PERMANENT GRAMMAR SAVE GRAMMAR WORKING GRAMMAR 2 22 OCT 85





NAME: TYPE: PROJECT: PART OF: COMPOSITION: ALIASES: SOURCES: DESTINATIONS: VERSION: DATE:

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NAME: TYPE: PROJECT: PART OF: COMPOSITION: ALIASES: SOURCES: DESTINATIONS: VERSION: DATE:

NAME: TYPE: **PROJECT:** NUMBER: DESCRIPTION: INPUT DATA: INPUT FLAGS: OUTPUT DATA: OUTPUT FLAGS: GLOBAL DATA USED: GLOBAL DATA CHANGED: FILES READ: FILES WRITTEN: HARDWARE READ: HARDWARE WRITTEN: ALIASES: CALLING PROCESSES: **PROCESSES CALLED:** RELATED ACTIVITY: VERSION: DATE:

GRAMMAR LOAD COMMAND DATA FLOW DDS CONSTRUCTION COMMAND

LOAD GRAMMAR 1 22 OCT 85

GRAMMAR MODIFICATION COMMAND DATA FLOW DDS CONSTRUCTION COMMAND

MODIFY GRAMMAR 1 22 OCT 85

GRAMMAR-CODE PROCESS DDS

PRODUCTION RECORD

PRODUCTION CODE

1 04 NOV 85

NAME: TYPE: **PROJECT:** NUMBER: **DESCRIPTION:** INPUT DATA: INPUT FLAGS: OUTPUT DATA: OUTPUT FLAGS: GLOBAL DATA USED: GLOBAL DATA CHANGED: FILES READ: FILES WRITTEN: HARDWARE READ: HARDWARE WRITTEN: ALIASES: CALLING PROCESSES: PROCESSES CALLED: **RELATED ACTIVITY:** VERSION: DATE: NAME: TYPE: **PROJECT:** NUMBER: DESCRIPTION: INPUT DATA: INPUT FLAGS: OUTPUT DATA: OUTPUT FLAGS: GLOBAL DATA USED: GLOBAL DATA CHANGED: FILES READ: FILES WRITTEN: HARDWARE READ: HARDWARE WRITTEN: ALIASES: CALLING PROCESSES:

PROCESSES CALLED: RELATED ACTIVITY:

VERSION:

DATE:

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GRAMMAR-LIST PROCESS DDS

PRODUCTION RECORD

SUB-GRAMMAR LIST

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GRAMMAR-NAME PROCESS DDS

PRODUCTION RECORD

PRODUCTION NAME

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NAME: TYPE: PROJECT: NUMBER: DESCRIPTION: INPUT DATA:

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INPUT FLAGS: OUTPUT DATA: OUTPUT FLAGS: GLOBAL DATA USED: GLOBAL DATA CHANGED: FILES READ: FILES WRITTEN: HARDWARE READ: HARDWARE WRITTEN: ALIASES: CALLING PROCESSES: PROCESSES CALLED:

RELATED ACTIVITY: VERSION: DATE:

NAME: TYPE: PROJECT: PART OF: COMPOSITION: ALIASES: SOURCES: DESTINATIONS: VERSION: DATE:

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GRAMMAR-PARSE PROCESS DDS

SENTENCE SUB-GRAMMAR LIST

RESULT

GRAMMAR-CODE GRAMMAR-NAME GRAMMAR-LIST IS-PRODUCTION SUB-PARSE IS-FUNCTION FTN-PARSE RESULT-LIST GRAMMAR-PARSE (RECURSIVE) RESULT-SENTENCE

1 04 NOV 85

GRAMMAR SAVE COMMAND DATA FLOW DDS CONSTRUCTION COMMAND

SAVE GRAMMAR 1 22 OCT 85

NAME: TYPE: PROJECT: PART OF: COMPOSITION: ALIASES: SOURCES: DESTINATIONS: VERSION: DATE:

NAME:

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INITIALIZATION COMMAND DATA FLOW DDS DATA DICTIONARY ACCESS COMMAND UPDATE DATA DICTIONARY 2 22 OCT 85

INITIALIZE-GRAMMAR PROCESS DDS

TYPE: PROJECT: NUMBER: DESCRIPTION: INPUT DATA: INPUT FLAGS: OUTPUT DATA: OUTPUT FLAGS: GLOBAL DATA USED: GLOBAL DATA CHANGED: FILES READ: FILES WRITTEN: HARDWARE READ: HARDWARE WRITTEN: ALIASES: CALLING PROCESSES: PROCESSES CALLED: **RELATED ACTIVITY:** VERSION: DATE:

FOR& 1 29 OCT 85

META-SYMBOLS *META-SYMBOLS*

PROCESS TYPE: PROJECT: DDS NUMBER: DESCRIPTION: FUNCTION CODE INPUT DATA: INPUT FLAGS: OUTPUT DATA: OUTPUT FLAGS: *EDIT-PTR* GLOBAL DATA USED: GLOBAL DATA CHANGED: *EDIT-PTR* FILES READ: FILES WRITTEN: HARDWARE READ: HARDWARE WRITTEN: ALIASES: CALLING PROCESSES: SETF* PROCESSES CALLED: RELATED ACTIVITY: 1 VERSION: 29 OCT 85 DATE: INSERT-PRODUCTION NAME: PROCESS TYPE: DDS **PROJECT:** NUMBER: DESCRIPTION: PRODUCTION INPUT DATA: CODE TO EXECUTE INPUT FLAGS: OUTPUT DATA: OUTPUT FLAGS: GLOBAL DATA USED: *EDIT-PTR* *EDIT-PTR* GLOBAL DATA CHANGED: FILES READ: FILES WRITTEN: HARDWARE READ: HARDWARE WRITTEN: ALIASES: CALLING PROCESSES: ED-DOWN PROCESSES CALLED:

RELATED ACTIVITY:

VERSION:

DATE:

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INSERT-FUNCTION

NAME:

AUX-ADD-PRODUCTION 1 29 OCT 85

NAME: TYPE: PROJECT: NUMBER: DESCRIPTION: INPUTS:

OUTPUTS:

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CONTROLS: MECHANISMS: ALIASES: PARENT ACTIVITY: VERSION: DATE:

NAME: TYPE: PROJECT: PART OF: COMPOSITION: ALIASES: SOURCES: DESTINATIONS: VERSION: DATE: INTERPRET SENTENCE ACTIVITY DDS

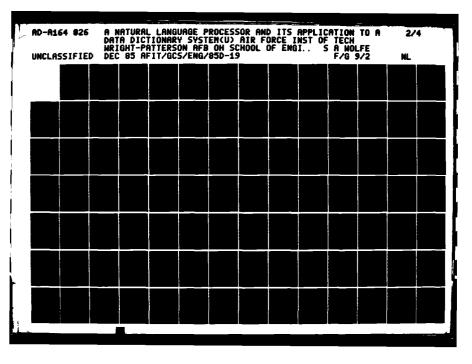
GRAMMAR SENTENCE DATA DICTIONARY ACCESS COMMAND INTERPRETER ERROR MESSAGE

DATA DICTIONARY SYSTEM 2 22 OCT 85

INTERPRETER ERROR MESSAGE DATA FLOW DDS RESPONSE PARSE ERROR MESSAGE

INTERPRET SENTENCE

2 22 OCT 85





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MICROCOPY RESOLUTION TEST CHART

NAME: TYPE: PROJECT: NUMBER: DESCRIPTION: INPUT DATA: INPUT FLAGS: OUTPUT DATA: OUTPUT FLAGS: GLOBAL DATA USED: GLOBAL DATA CHANGED: FILES READ: FILES WRITTEN: HARDWARE READ: HARDWARE WRITTEN: ALIASES: CALLING PROCESSES: **PROCESSES CALLED: RELATED ACTIVITY:** VERSION: DATE: NAME: TYPE: **PROJECT:** NUMBER: DESCRIPTION: INPUT DATA: INPUT FLAGS: OUTPUT DATA: OUTPUT FLAGS: GLOBAL DATA USED: GLOBAL DATA CHANGED: FILES READ: FILES WRITTEN:

HARDWARE READ: HARDWARE WRITTEN:

CALLING PROCESSES: PROCESSES CALLED: RELATED ACTIVITY:

ALIASES:

VERSION:

DATE:

SYMBOL TO CHECK IS-FUNCTION 1 29 OCT 85 IS-META-SYMBOL PROCESS DDS SYMBOL TO CHECK IS-META-SYMBOL *META-SYMBOLS* 1 29 OCT 85

IS-FUNCTION

PROCESS

DDS



NAME: TYPE: PROJECT: NUMBER: DESCRIPTION: INPUT DATA: INPUT FLAGS: OUTPUT DATA: OUTPUT FLAGS: GLOBAL DATA USED: GLOBAL DATA CHANGED: FILES READ: FILES WRITTEN: HARDWARE READ: HARDWARE WRITTEN: ALIASES: CALLING PROCESSES: **PROCESSES CALLED: RELATED ACTIVITY:** VERSION: DATE:

NAME:

TYPE:

PROJECT: NUMBER:

INPUTS:

OUTPUTS:

CONTROLS: MECHANISMS: ALIASES:

VERSION:

DATE:

DESCRIPTION:

PARENT ACTIVITY:

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IS-PRODUCTION PROCESS DDS

SYMBOL TO CHECK

IS-PRODUCTION

1 29 OCT 85

LOAD GRAMMAR ACTIVITY DDS

GRAMMAR GRAMMAR SAVE COMMAND GRAMMAR

CONSTRUCT GRAMMAR 2 22 OCT 85 NAME: TYPE: PROJECT: PART OF: COMPOSITION: ALIASES: SOURCES: DESTINATIONS: VERSION: DATE:

NAME: TYPE: PROJECT: NUMBER: DESCRIPTION: INPUT DATA:

INPUT FLAGS: OUTPUT DATA: OUTPUT FLAGS: GLOBAL DATA USED: GLOBAL DATA CHANGED: FILES READ: FILES WRITTEN: HARDWARE READ: HARDWARE WRITTEN: ALIASES: CALLING PROCESSES: PROCESSES CALLED:

RELATED ACTIVITY: VERSION: DATE: MODIFICATION COMMAND DATA FLOW DDS DATA DICTIONARY ACCESS COMMAND UPDATE DATA DICTIONARY

UPDATE DATA DICTIONARY 2 22 OCT 85

MODIFY-FUNCTION (MACRO) PROCESS DDS

FUNCTION CODE FUNCTION RECORD NAME

EDIT-PTR *EDIT-PTR*

IS-FUNCTION ED-RESET INSERT-FUNCTION MSG 1 29 OCT 85

NAME: TYPE: PROJECT: NUMBER: DESCRIPTION: INPUTS:

OUTPUTS:

CONTROLS: MECHANISMS: ALIASES: PARENT ACTIVITY: VERSION: DATE:

NAME: TYPE: PROJECT: NUMBER: DESCRIPTION: INPUT DATA:

INPUT FLAGS: OUTPUT DATA: OUTPUT FLAGS: GLOBAL DATA USED: GLOBAL DATA CHANGED: FILES READ: FILES WRITTEN: HARDWARE READ: HARDWARE WRITTEN: ALIASES: CALLING PROCESSES: PROCESSES CALLED: MODIFY GRAMMAR ACTIVITY DDS

GRAMMAR GRAMMAR MODIFICATION COMMAND CONSTRUCTION ERROR MESSAGE GRAMMAR

CONSTRUCT GRAMMAR 1 22 OCT 85

PARSE PROCESS DDS

SENTENCE GRAMMAR

GRAMMAR-LIST GRAMMAR-CODE GRAMMAR-PARSE SUCCESSFUL DISPLAY-PARSE-ERROR 1

RELATED ACTIVITY: VERSION: DATE:

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NAME: TYPE: PROJECT: PART OF: COMPOSITION: ALIASES: SOURCES: DESTINATIONS: VERSION: DATE:

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NAME: TYPE: PROJECT: PART OF: COMPOSITION: ALIASES: SOURCES: DESTINATIONS: VERSION: DATE:

NAME: TYPE: PROJECT: NUMBER: DESCRIPTION: INPUTS:

OUTPUTS:

CONTROLS: MECHANISMS: ALIASES: PARENT ACTIVITY: VERSION: DATE: PARSE ERROR MESSAGE DATA FLOW DDS INTERPRETER ERROR MESSAGE

PARSE SENTENCE

PARSE RESULTS DATA FLOW DDS

PARSE SENTENCE GENERATE DD COMMAND 1 22 OCT 85

PARSE SENTENCE ACTIVITY DDS

GRAMMAR SENTENCE PARSE ERROR MESSAGE PARSE RESULTS

INTERPRET SENTENCE 1 22 OCT 85

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NAME: TYPE: PROJECT: PART OF: COMPOSITION: ALIASES: SOURCES: DESTINATIONS: VERSION: DATE:

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NAME: TYPE: PROJECT: PART OF: COMPOSITION:

ALIASES: SOURCES: DESTINATIONS: VERSION: DATE:

NAME: TYPE: **PROJECT:** NUMBER: **DESCRIPTION:** INPUT DATA: INPUT FLAGS: OUTPUT DATA: OUTPUT FLAGS: GLOBAL DATA USED: GLOBAL DATA CHANGED: FILES READ: FILES WRITTEN: HARDWARE READ: HARDWARE WRITTEN: ALIASES: CALLING PROCESSES: **PROCESSES CALLED: RELATED ACTIVITY:** VERSION: DATE:

PRESENTATION COMMAND DATA FLOW DDS DATA DICTIONARY ACCESS COMMAND

DISPLAY DATA DICTIONARY 2 22 OCT 85

RESPONSE DATA FLOW DDS

CONSTRUCTION ERROR MESSAGE DATA DICTIONARY DATA DATA DICTIONARY ERROR MESSAGE INTERPRETER ERROR MESSAGE

DATA DICTIONARY SYSTEM USER 2 22 OCT 85

RESULT-LIST PROCESS DDS

RESULT

SUB-GRAMMAR LIST

04 NOV 85

1

RESULT-SENTENCE NAME: TYPE: PROCESS PROJECT: DDS NUMBER: DESCRIPTION: RESULT INPUT DATA: INPUT FLAGS: OUTPUT DATA: SENTENCE OUTPUT FLAGS: GLOBAL DATA USED: GLOBAL DATA CHANGED: FILES READ: FILES WRITTEN: HARDWARE READ: HARDWARE WRITTEN: ALIASES: CALLING PROCESSES: PROCESSES CALLED: RELATED ACTIVITY: VERSION: DATE: NAME: TYPE: PROJECT: NUMBER: **DESCRIPTION:**

OUTPUTS: CONTROLS: **MECHANISMS:** ALIASES: PARENT ACTIVITY: VERSION: DATE:

INPUTS:

SAVE GRAMMAR ACTIVITY DDS GRAMMAR

1

04 NOV 85

GRAMMAR SAVE COMMAND GRAMMAR

CONSTRUCT GRAMMAR 2 22 OCT 85





















NAME: TYPE: **PROJECT:** NUMBER: **DESCRIPTION:** INPUT DATA: INPUT FLAGS: OUTPUT DATA: OUTPUT FLAGS: GLOBAL DATA USED: GLOBAL DATA CHANGED: FILES READ: FILES WRITTEN: HARDWARE READ: HARDWARE WRITTEN: ALIASES: CALLING PROCESSES: **PROCESSES CALLED: RELATED ACTIVITY:** VERSION: DATE:

SAVE-GRAMMAR PROCESS DDS

META-SYMBOLS

GRAMMAR.LSP

TYPE: PROJECT: PART OF: COMPOSITION: ALIASES: SOURCES: DESTINATIONS: VERSION:

DATE:

NAME:

SENTENCE DATA FLOW DDS COMMAND

29 OCT 85

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INTERPRET SENTENCE PARSE SENTENCE 2 22 OCT 85







NAME: TYPE: PROJECT: NUMBER: DESCRIPTION: INPUT DATA:

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SUB-GRAMMAR-PARSE PROCESS DDS

SENTENCE SUB-GRAMMAR LIST

RESULT

INPUT FLAGS: OUTPUT DATA: OUTPUT FLAGS: GLOBAL DATA USED: GLOBAL DATA CHANGED: FILES READ: FILES WRITTEN: HARDWARE READ: HARDWARE WRITTEN: ALIASES: CALLING PROCESSES: PROCESSES CALLED:

RELATED ACTIVITY: VERSION: DATE: GRAMMAR-CODE GRAMMAR-NAME GRAMMAR-LIST IS-PRODUCTION SUB-PARSE IS-FUNCTION FTN-PARSE RESULT-LIST SUB-GRAMMAR-PARSE (RECURSIVE) RESULT-SENTENCE

1

04 NOV 85

SUB-PARSE NAME: PROCESS TYPE: PROJECT: DDS NUMBER: DESCRIPTION: INPUT DATA: SENTENCE GRAMMAR INPUT FLAGS: OUTPUT DATA: RESULT OUTPUT FLAGS: GLOBAL DATA USED: GLOBAL DATA CHANGED: FILES READ: FILES WRITTEN: HARDWARE READ: HARDWARE WRITTEN: ALIASES: CALLING PROCESSES: GRAMMAR-LIST PROCESSES CALLED: GRAMMAR-CODE SUB-GRAMMAR-PARSE RESULT-LIST **RELATED ACTIVITY:** VERSION: 1 DATE: 04 NOV 85

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NAME: TYPE: PROJECT: NUMBER: DESCRIPTION: INPUT DATA: INPUT FLAGS: OUTPUT DATA: OUTPUT FLAGS: GLOBAL DATA USED: GLOBAL DATA CHANGED: FILES READ: FILES WRITTEN: HARDWARE READ: HARDWARE WRITTEN: ALIASES: CALLING PROCESSES: PROCESSES CALLED: RELATED ACTIVITY: VERSION: DATE: NAME: TYPE: PROJECT: NUMBER: DESCRIPTION: INPUTS:

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OUTPUTS:

CONTROLS: MECHANISMS: ALIASES: PARENT ACTIVITY: VERSION: DATE: SUCCESSFUL

RESULT-SENTENCE

RESULT-LIST

04 NOV 85

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DDS

SUCCESSFUL

PROCESS

RESULT

DDS

UPDATE DATA DICTIONARY ACTIVITY

ADDITION COMMAND DATA DICTIONARY DATA DELETION COMMAND INITIALIZATION COMMAND MODIFICATION COMMAND DATA DICTIONARY DATA DATA DICTIONARY ERROR MESSAGE DELETION COMMAND MODIFICATION COMMAND 9

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ACCESS DATA DICTIONARY 2 22 OCT 85

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Appendix B

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Structure Charts

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Appendix B Table of Contents

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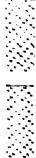


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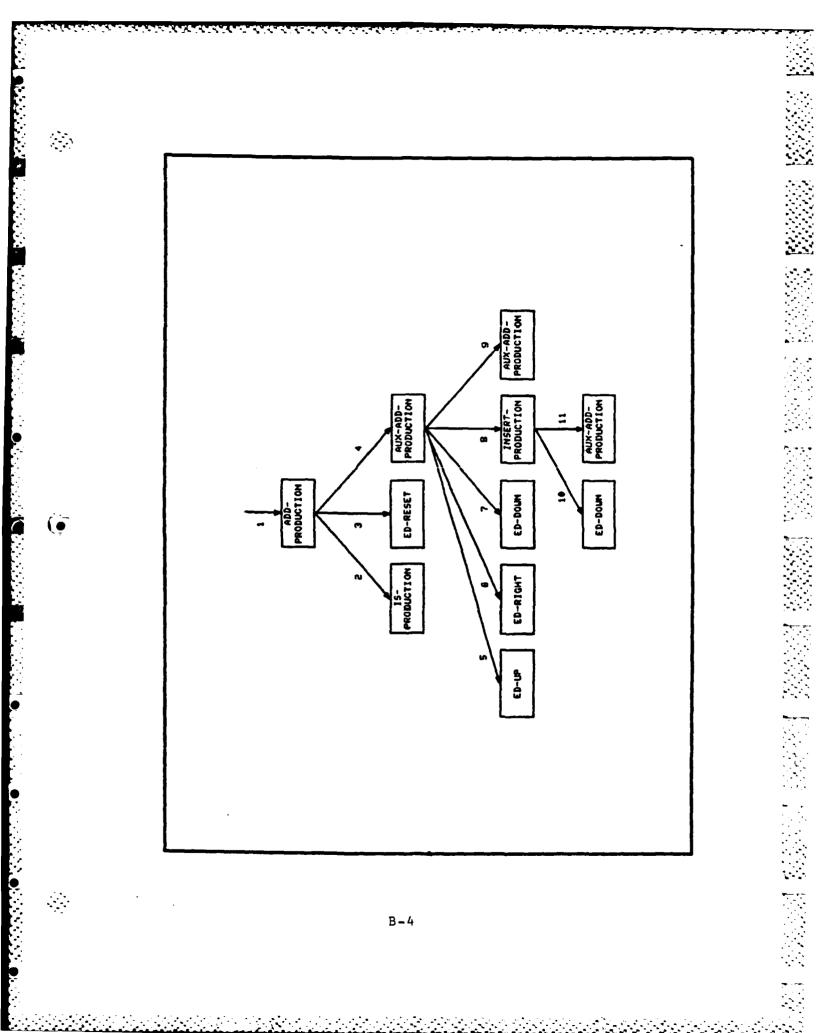
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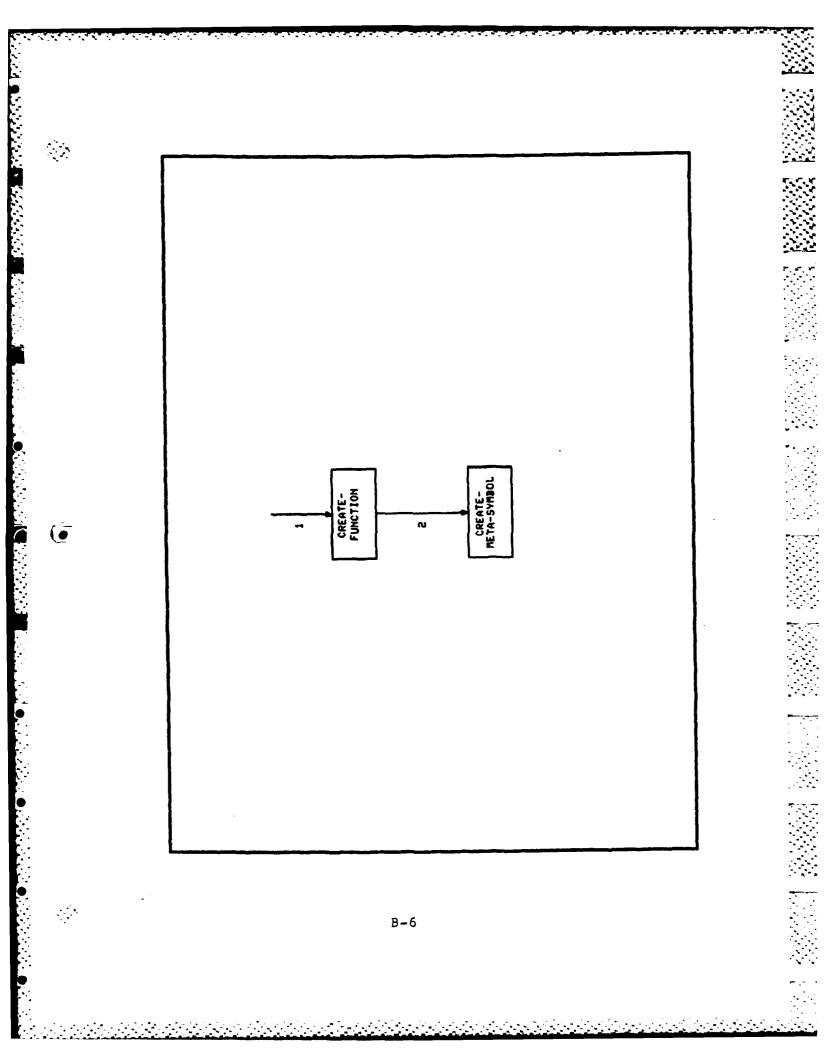
add-production Interfaces

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#	Passed Parameters	Туре	Returned Parameters	Туре
1	sentence code production name	data data data		
2	production-name	data	is production?	flag
3	*edit-exp* *edit-exp* nil	data data data		
4	sentence code	data data		
5	(none)			
6	(none)			
7	(none)			
8	sentence code	data data		
9	cdr [*] (sentence) code	data data		
10	(none)			
11	cdr [*] (sentence) code	data data		

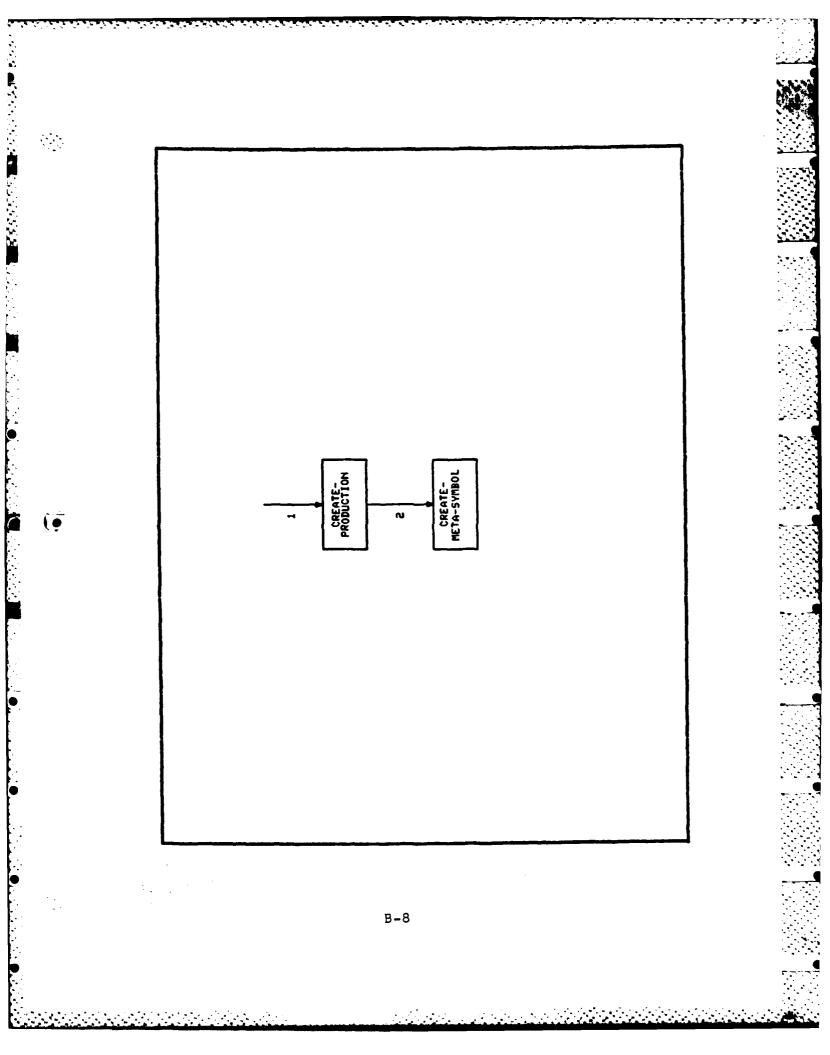
* The cdr function returns a list of all but the first element of a list passed to it. For instance, the cdr of the list (a b c) is the list (b c).



create-function Interfaces

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#	Passed Parameters	Туре	Returned Parameters	Туре
1	function-name code	data data		
2	function-name	data	function created?	flag



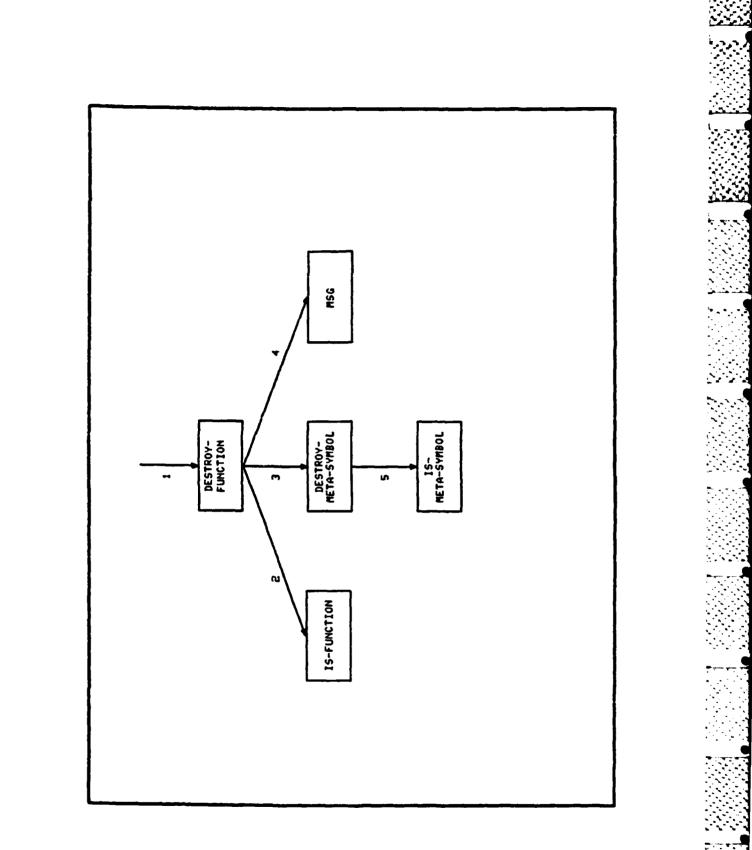
create-production Interfaces

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#	Passed Parameters	Туре	Returned Parameters	Туре
1	production-name code	data data		
2	production-name	data	production created?	flag



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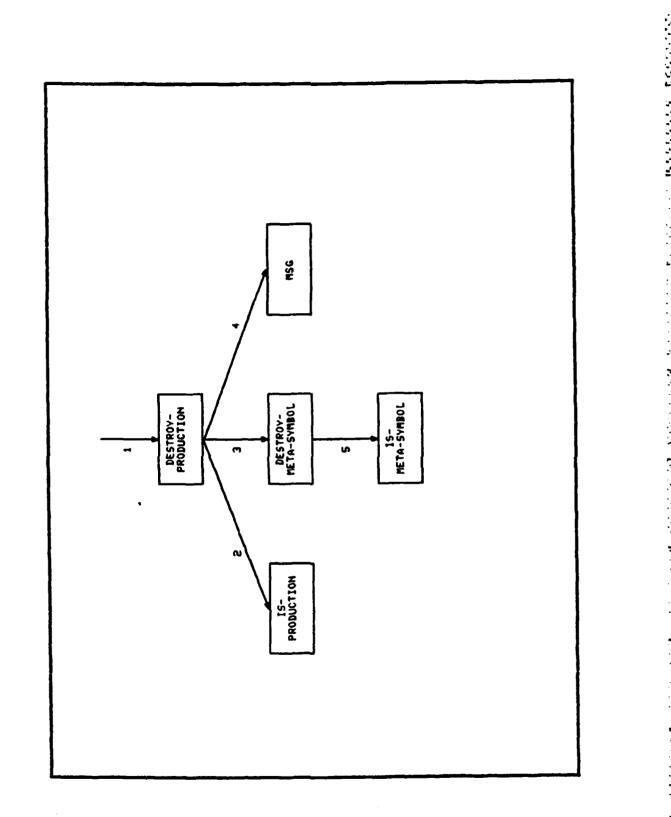


destroy-function Interfaces

#	Passed Parameters	Туре	Returned Parameters	Туре
1	function name	data		
2	function name	data	is function?	flag
3	function name	data		
4	error message text	data		
5	function name	data	is meta-symbol?	flag



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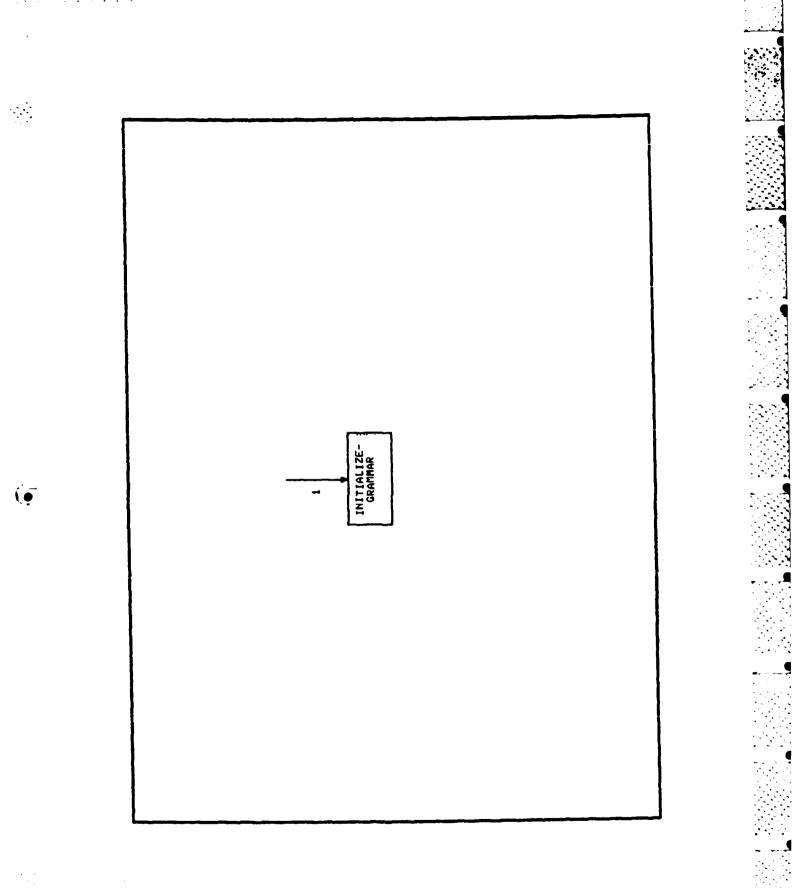
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destroy-production Interfaces

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#	Passed Parameters	Туре	Returned Parameters	Туре
1	production name	data		
2	production name	data	is production?	flag
3	production name	data		
4	error message text	data		
5	production name	data	is meta-symbol?	flag





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initialize-grammar Interfaces

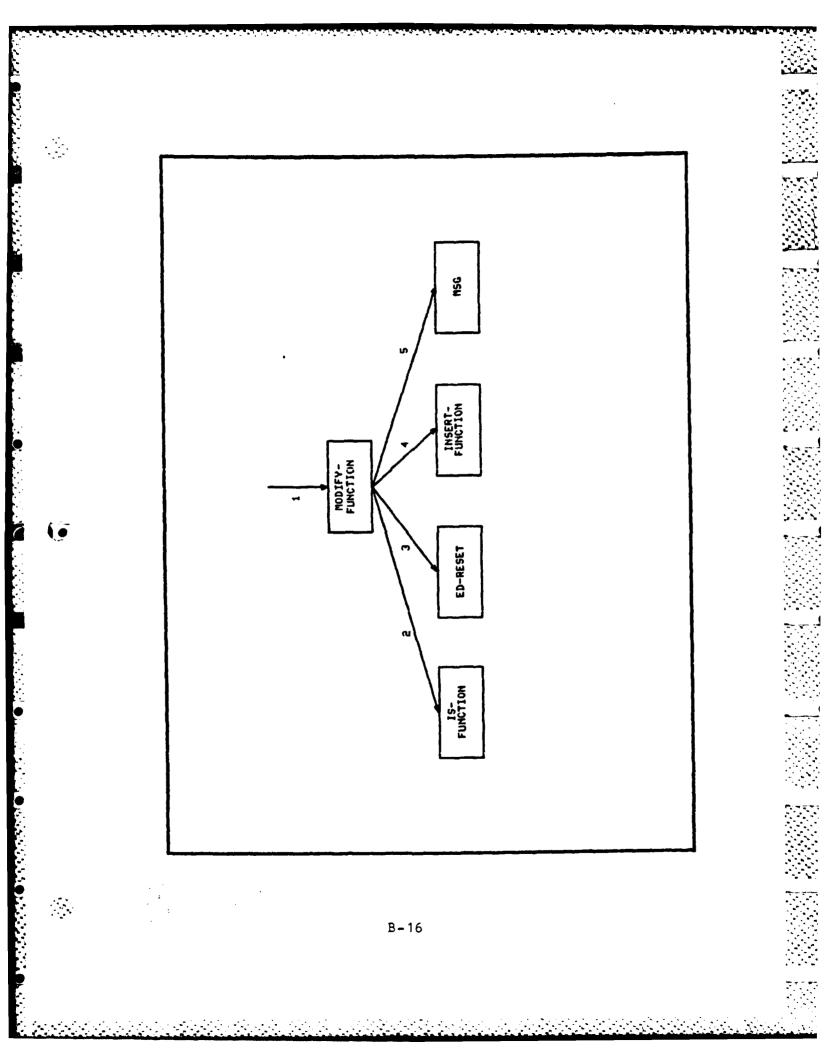
Passed Parameters Type Returned Parameters Type

l (none)

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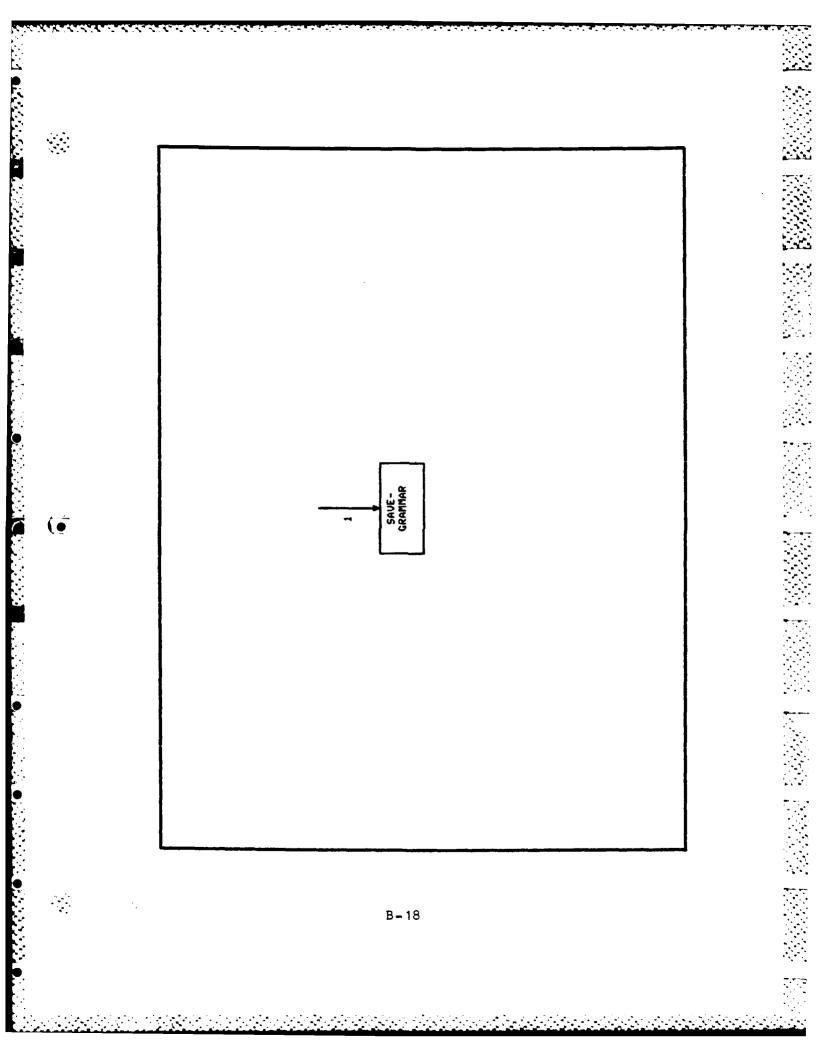


_#	Passed Parameters	Туре	Returned Parameters	Туре
1	function name	data		
2	function name	data	is function?	flag
3	*edit-ptr* *edit-ptr* nil	data data		
4	code	data		
5	error message text	data		

modify-function Interfaces

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save-grammar Interfaces

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#	Passed	Parameters	Туре	Returned	Parameters	Туре

1 (none)

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Sentence Interpreter

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SUCCESSFUL GRAMMAR-PARSE PARSE C FTN-PARSE FTN-PARSE GRAMMAR-PARSE . SUB-PARSE SUB-PARSE SUB-GRAMMAR-PARSE • .

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parse Interfaces

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_#	Passed Parameters	Туре	Returned Parameters	Туре
1	sentence grammar	data data		
2	sentence sub-grammar-list	data data	result	data
3	result	data	successful parse?	flag
4	(none)			
5	sentence sub-grammar	data data	result	data
6	sentence function	data data	result	data
7	cdr [*] (sentence) sub-grammar-list	data data	result	data
8	sentence sub-grammar-list	data data	result	data
9	cdr [*] (sentence) sub-grammar-list	data data	result	data
10	sentence sub-grammar	data data	result	data
11	sentence function	data data	result	data

* The cdr function returns a list of all but the first element of a list passed to it. For instance, the cdr of the list (a b c) is the list (b c).





































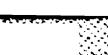


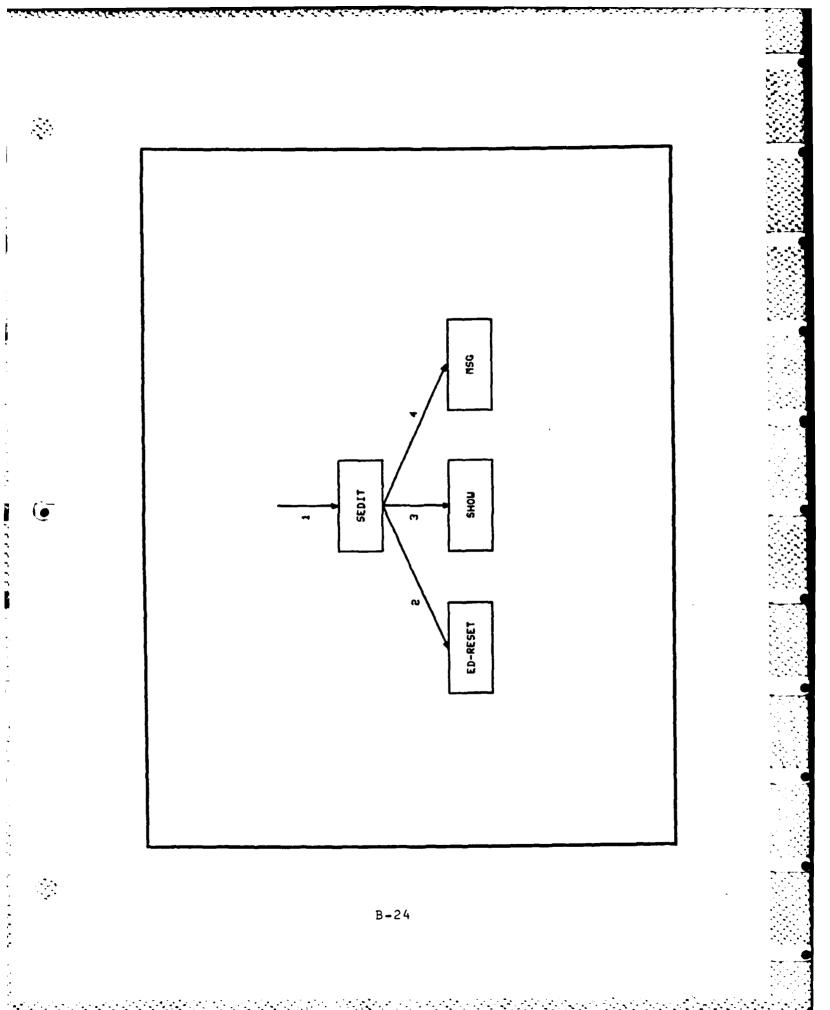
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S-expression Editor

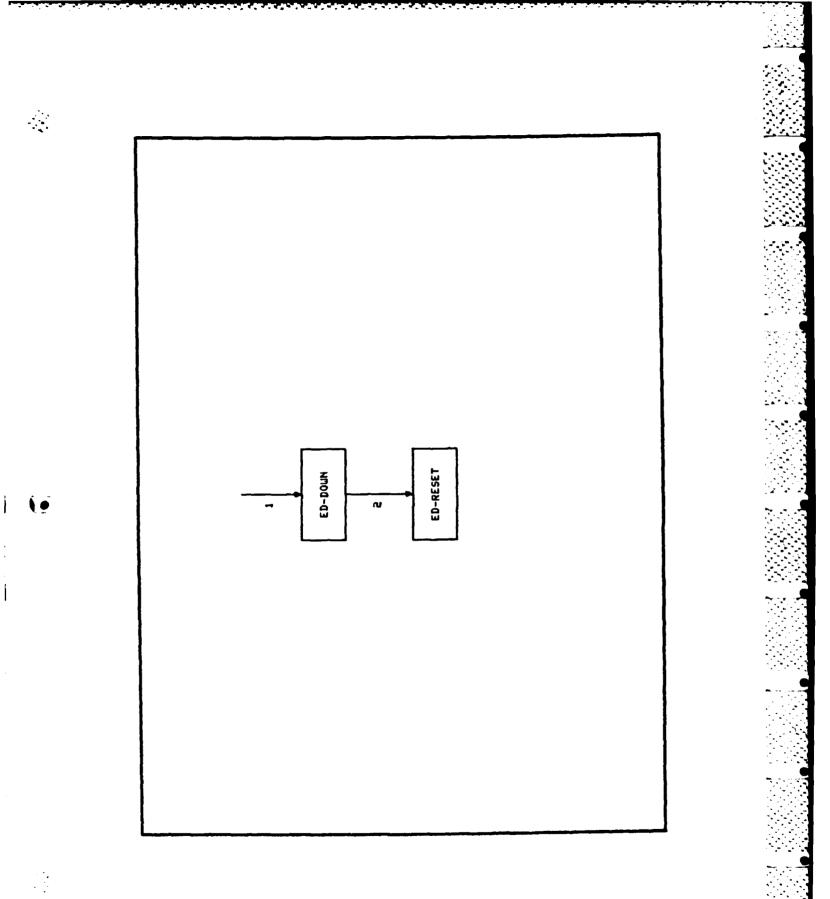






sedit Interfaces

Ŧ Passed Parameters Туре Returned Parameters Туре 1 data expression expression (edit 2 pointer) expression (edit data line) data nil (edit stack) data number of levels 3 to display data 4 error message text data



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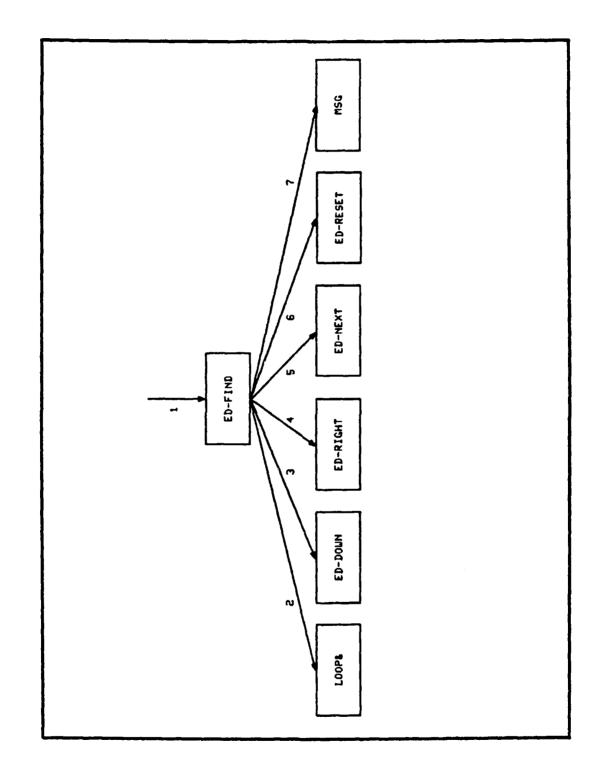
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ed-down Interfaces

#	Passed Parameters	Туре	Returned Parameters	Туре
1	(none)			
2	new edit pointer new edit line new edit stack	data data data		





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#	Passed Parameters	Туре	Returned	Parameters	Туре
1	(none)				
2	*key* edit pointer edit line edit stack	data data data data			
3	(none)				
4	(none)				
5	(none)				
6	edit pointer edit line edit stack	data data data			
7	error message text	data			

ed-find Interfaces

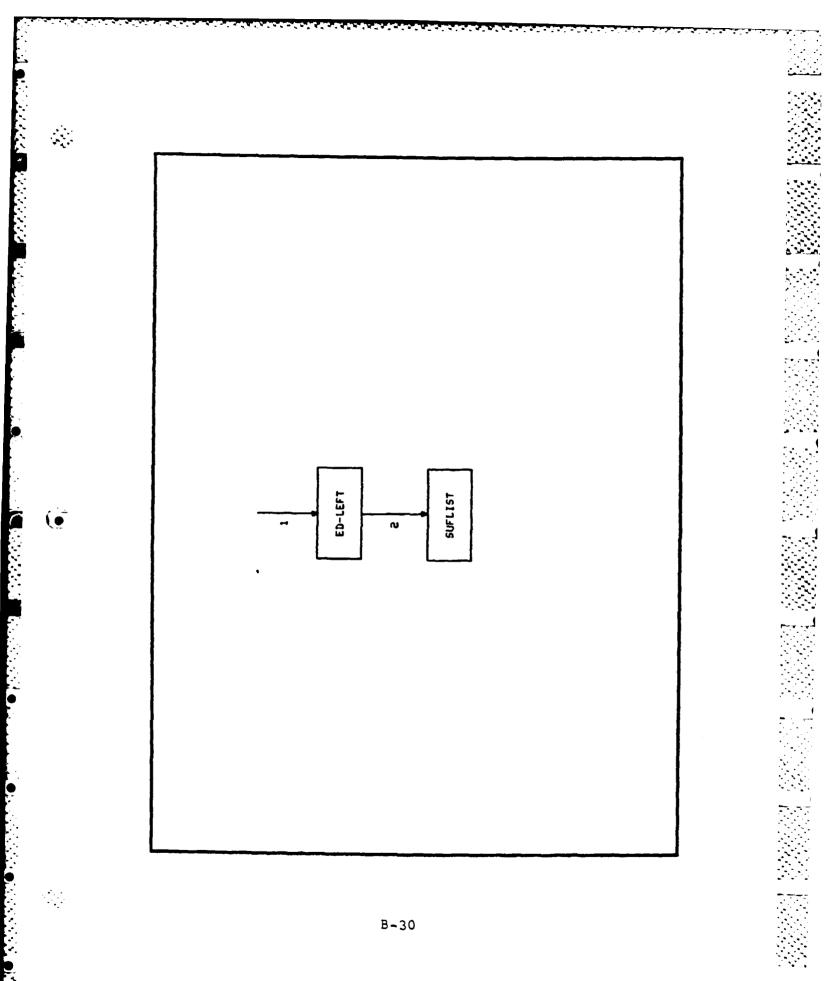
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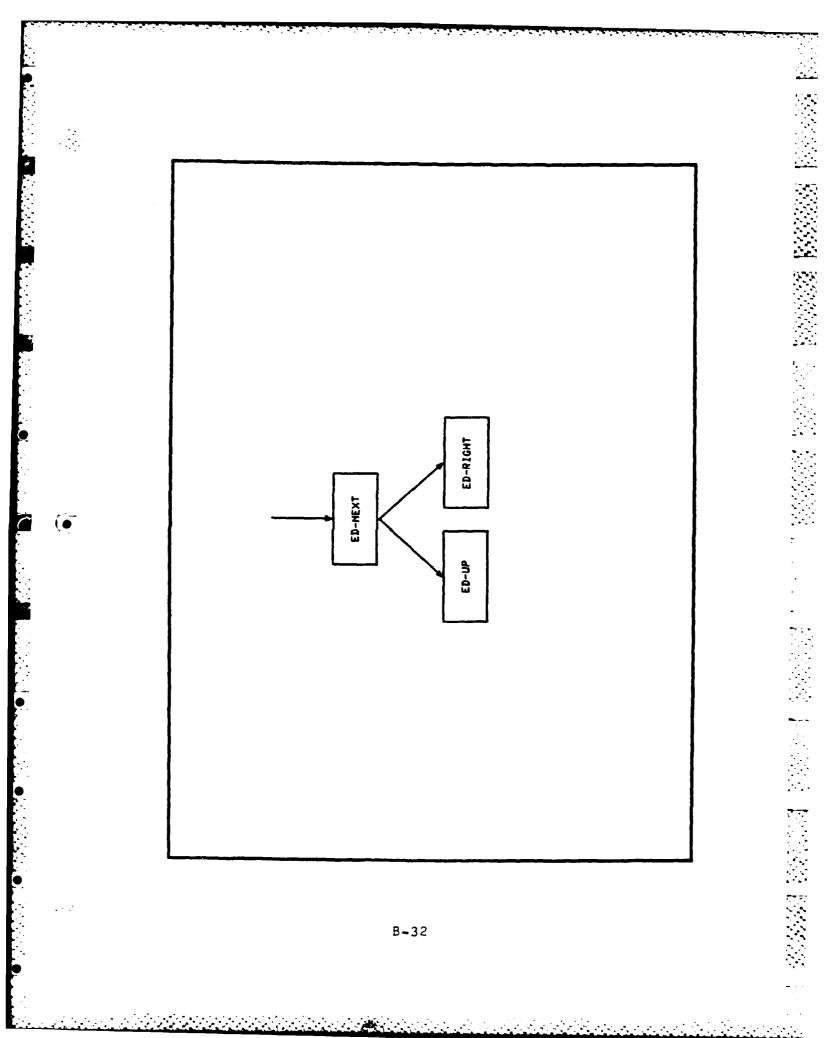
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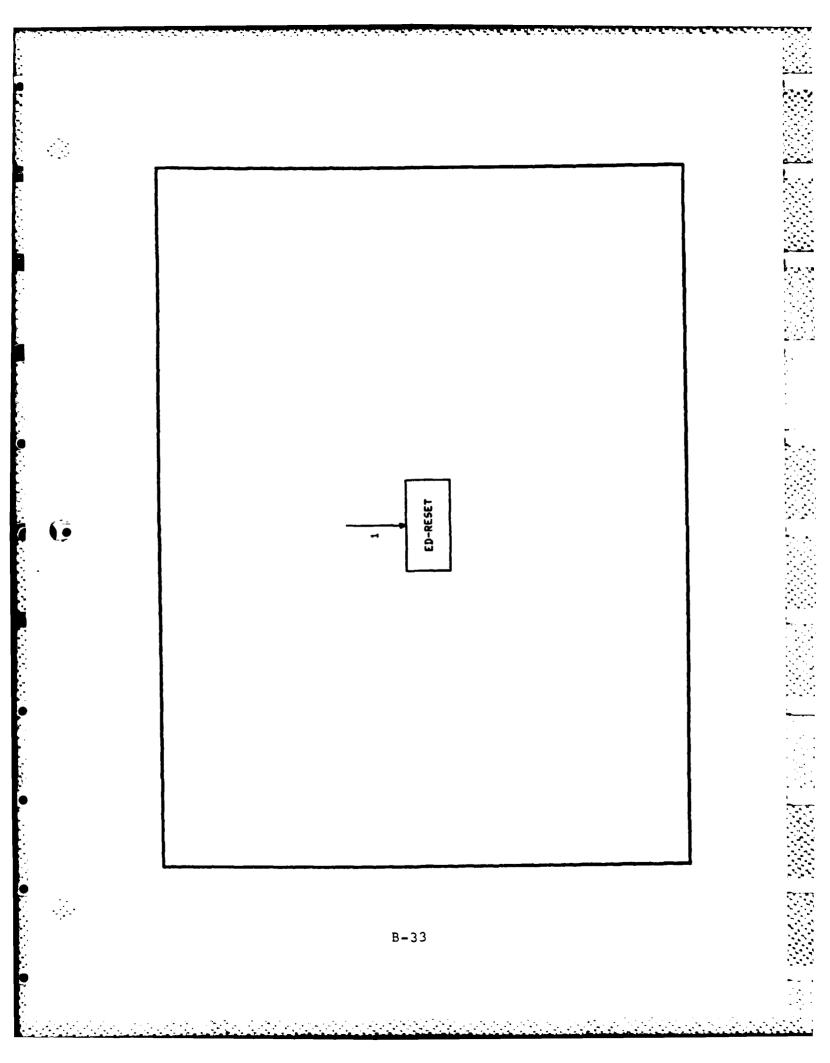
ed-left Interfaces

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#	Passed Parameters	Туре	Returned Parameters	Туре
1	(none)			
2	edit line current edit	data	new edit pointer	data
	pointer position	data		





ed-reset Interfaces

#	Passed Parameters	Туре	Returned Parameters	Туре
1	new edit pointer new edit líne new edit stack	data data data		

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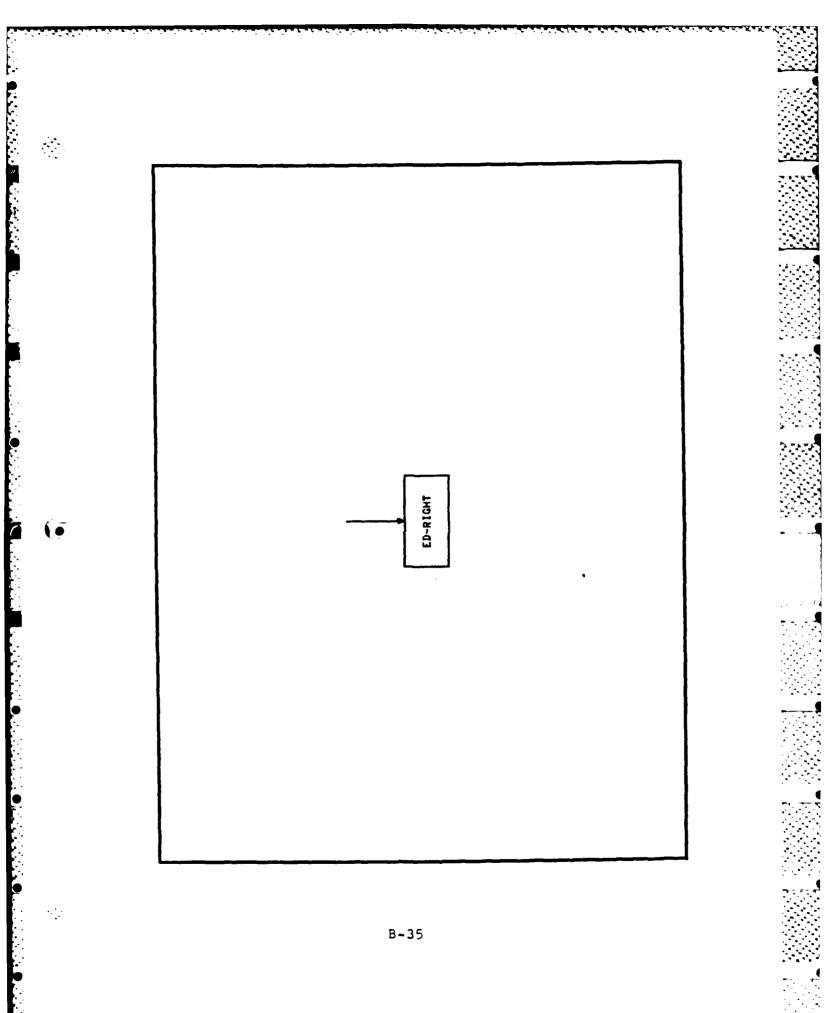
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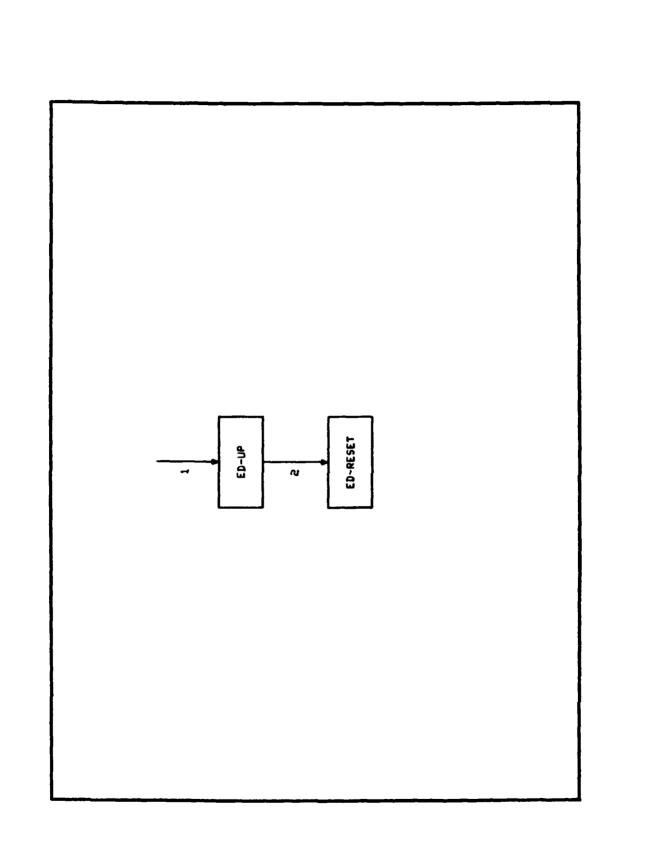
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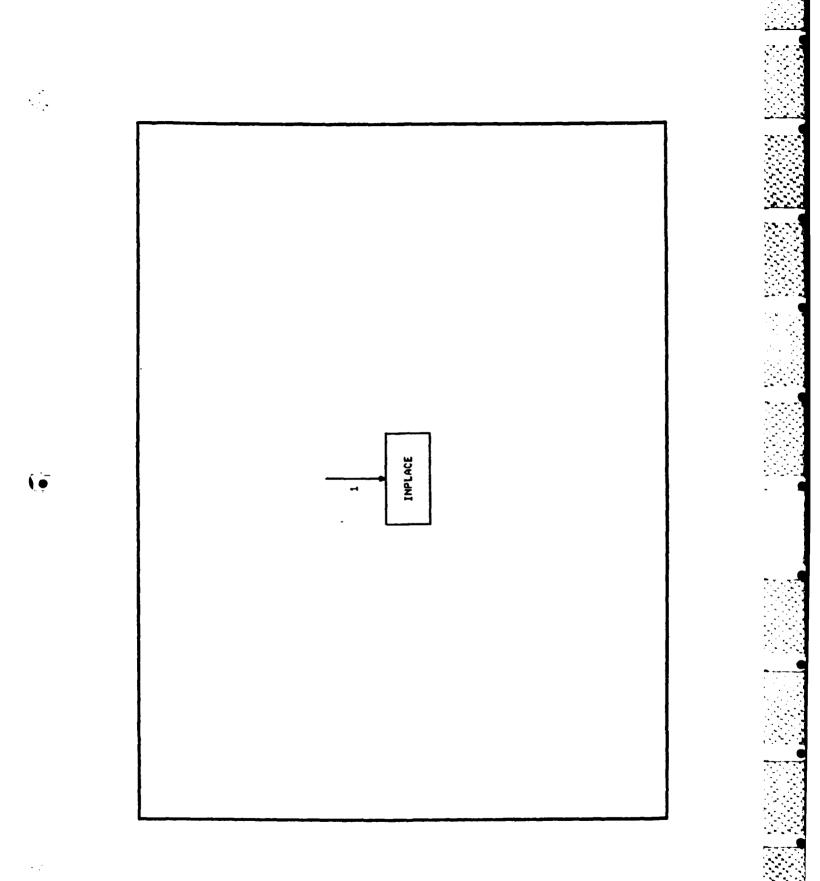


ed-up Interfaces

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#	Passed Parameters	Туре	Returned Parameters	Туре
1	(none)			
2	new edit pointer new edit line new edit stack	data data data		

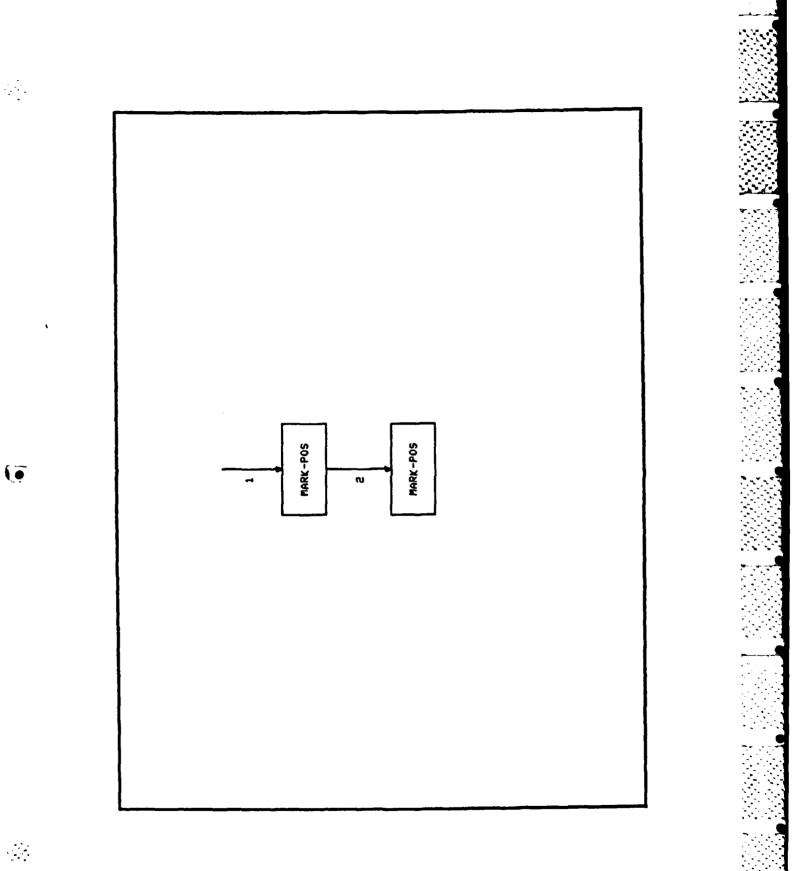


inplace Interfaces

#	Passed Parameters	Туре	Returned Parameters	Туре
1	car [*] (new edit pointer) cdr ⁺ (new edit pointer)	data data		

* The car function returns the first element of a list passed to it. For instance, the car of the list (a b c) is a.

* The cdr function returns a list of all but the first element of a list passed to it. The cdr of the above list is the list (b c).

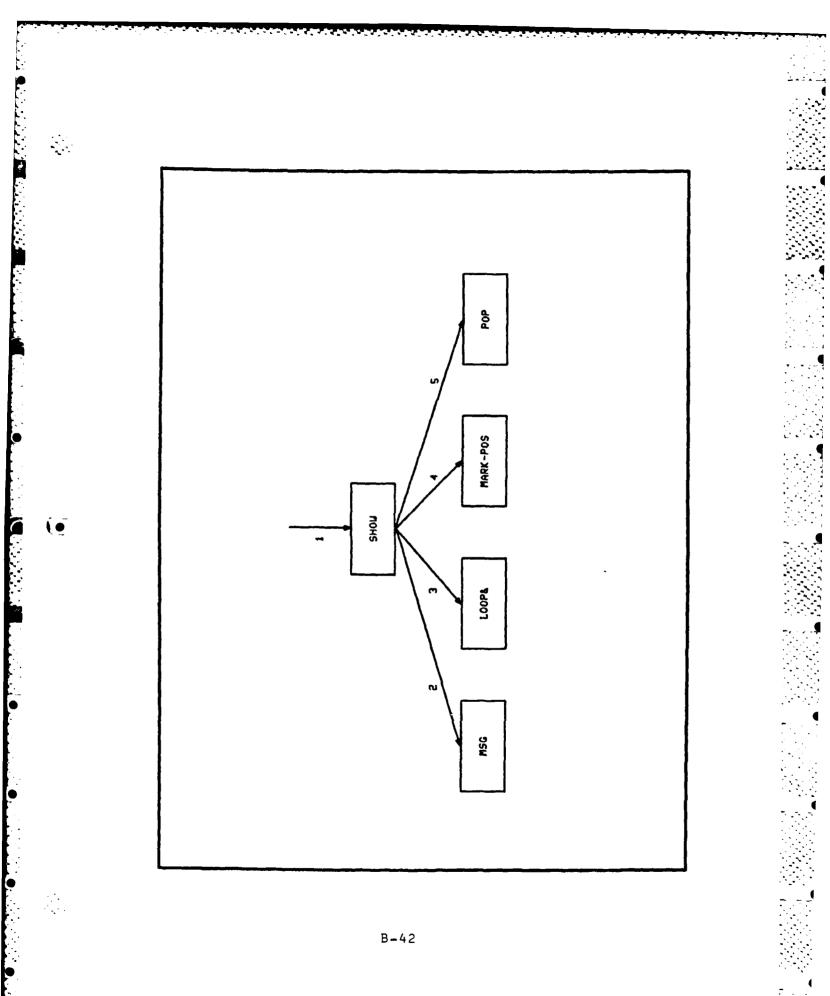


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mark-pos Interfaces

#	Passed Parameters	Туре	Returned	Parameters	Туре
1	position to mark list	data data			
2	position to mark cdr (list)	data data			

* The cdr function returns a list of all but the first element of a list passed to it. For instance, the cdr of the list (a b c) is the list (b c).



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show Interfaces

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#	Passed Parameters	Туре	Returned	Parameters	Туре
1	number of levels of the edit stack to display	data			
2	car [*] (edit pointer)	data			
3	edit pointer edit line edit stack	data data data			
4	edit pointer edit line	data data			
5	stack	data	top of st	tack	data

* The car function returns the first element of a list passed to it. For instance, the car of the list (a b c) is a.

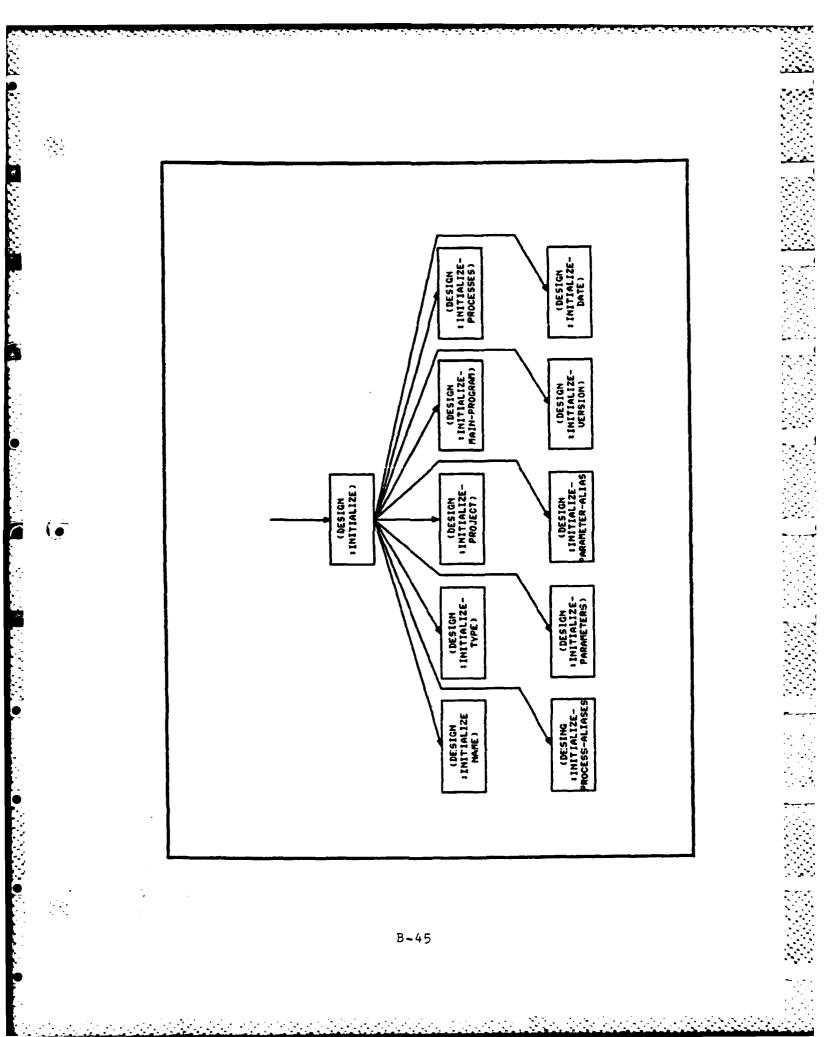
DESIGN Methods

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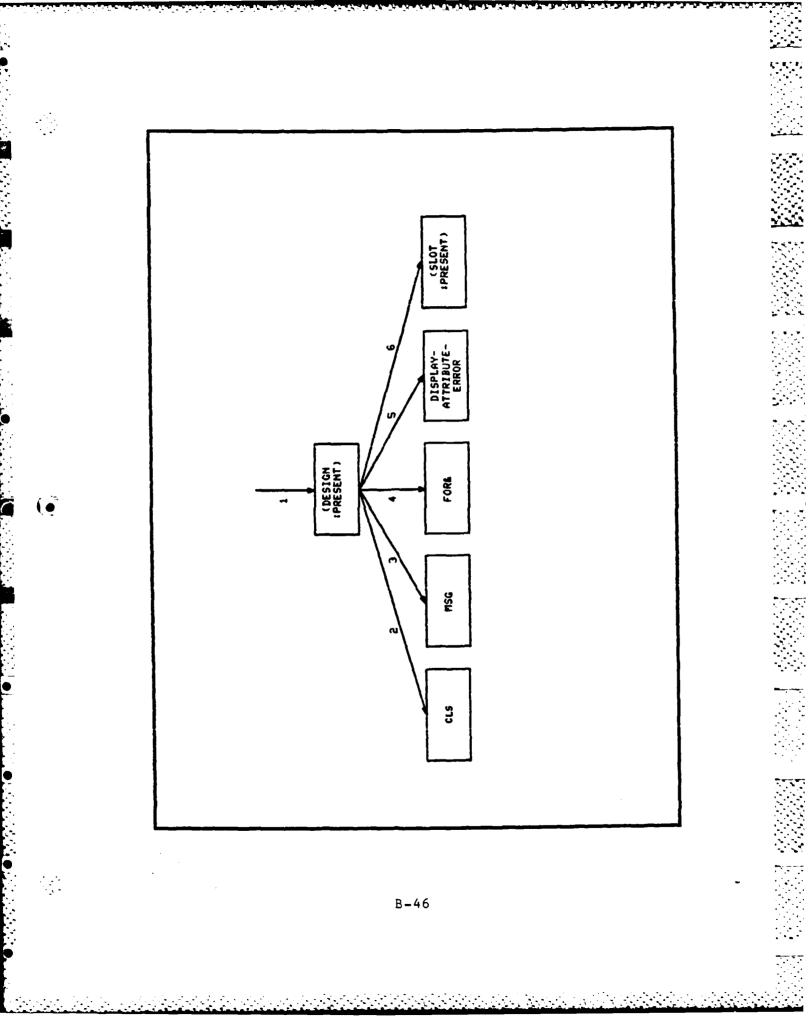
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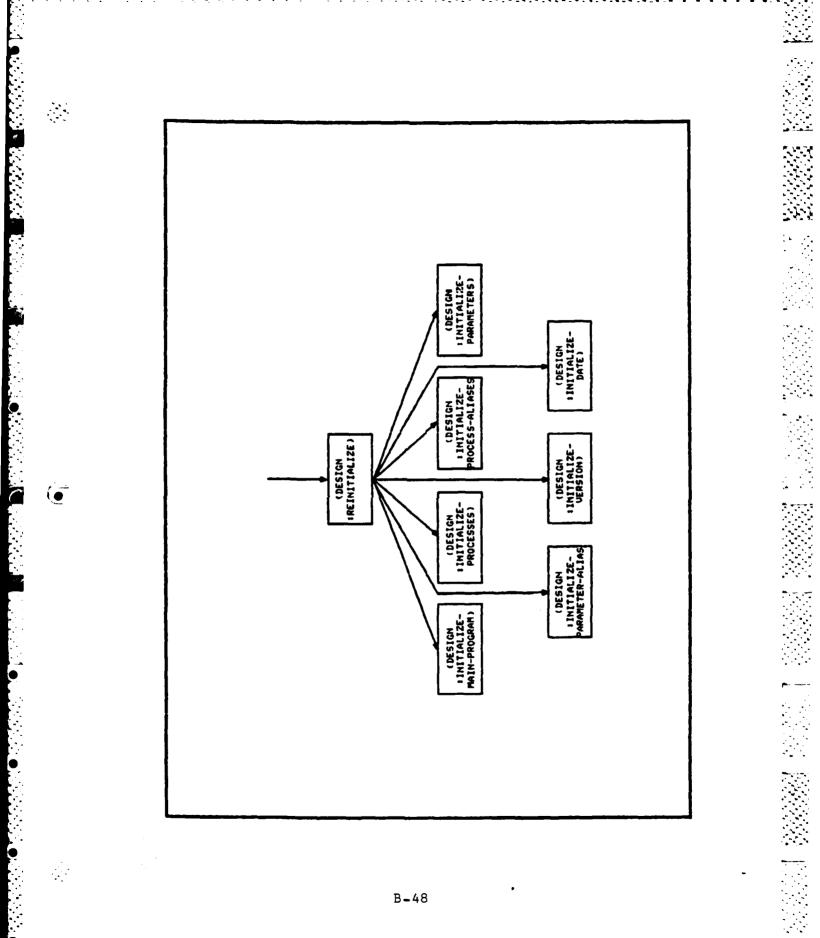


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(design :present) Interfaces

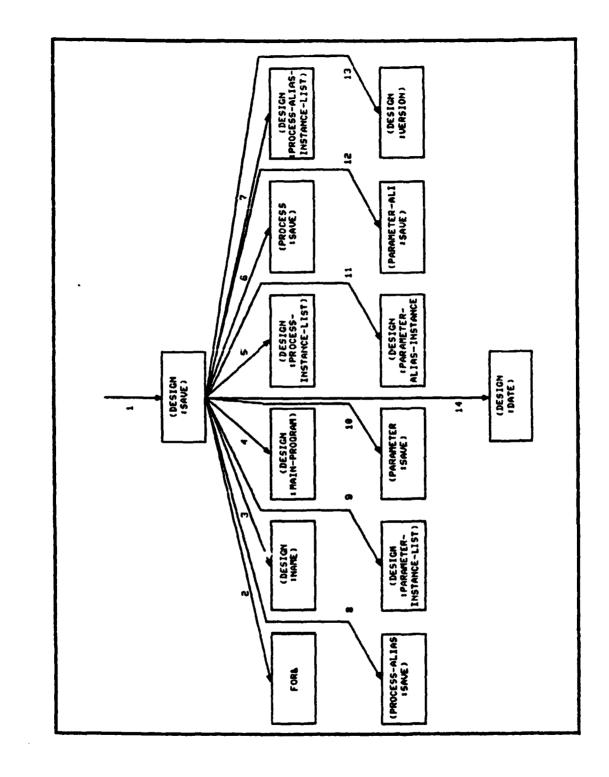
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#	Passed Parameters	Туре	Returned	Parameters	Туре
1	(none)				
2	list of instance variables	data			
3	(none)				
4	list of instance variables	data			
5	list of invalid instance variables	data			
6	(none)				





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(design :save) Interfaces

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	Passed Parameters	Туре	Returned Parameters	Туре
1	project instance pointer stream	data data		
2	list of instance pointers	data		
3	(none)		name	data
4	(none)		main-program	data
5	(none)		list of process instance pointers	data
6	process instance pointer	data		
7	(none)		list of process- alias instance pointers	data
8	process-alias instance pointer	data		
9	(none)		list of parameter instance pointers	data
10	parameter instance pointer	data		
11	(none)		list of parameter- alias instance pointers	data
12	parameter-alias instance pointer	data		
13	(none)		version	data
14	(none)		date	data















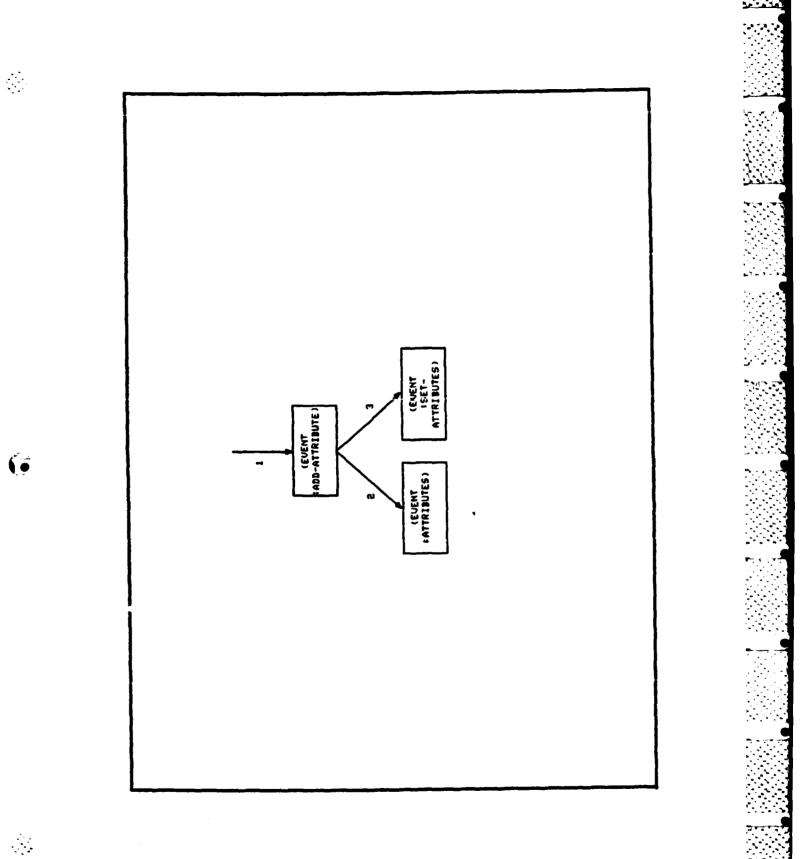






EVENT Methods





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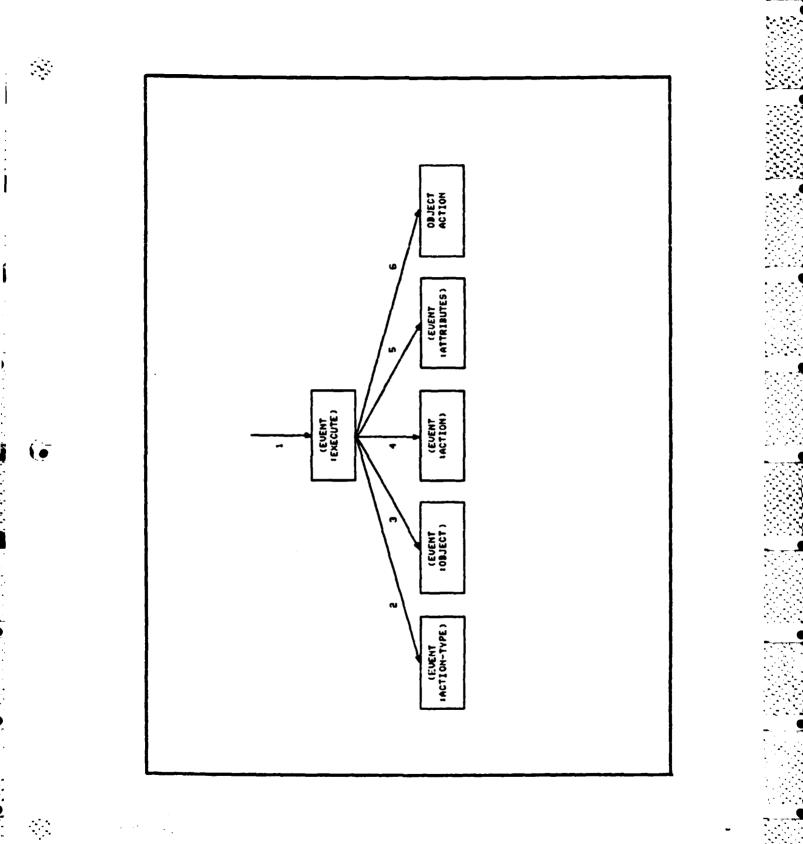


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(event :add-attribute) Interfaces

#	Passed Parameters	Туре	Returned Parameters	Туре
1	attribute	data		
2	(none)		list of attributes	data
3	list of attributes	data		

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(event	:execute)	Interfaces
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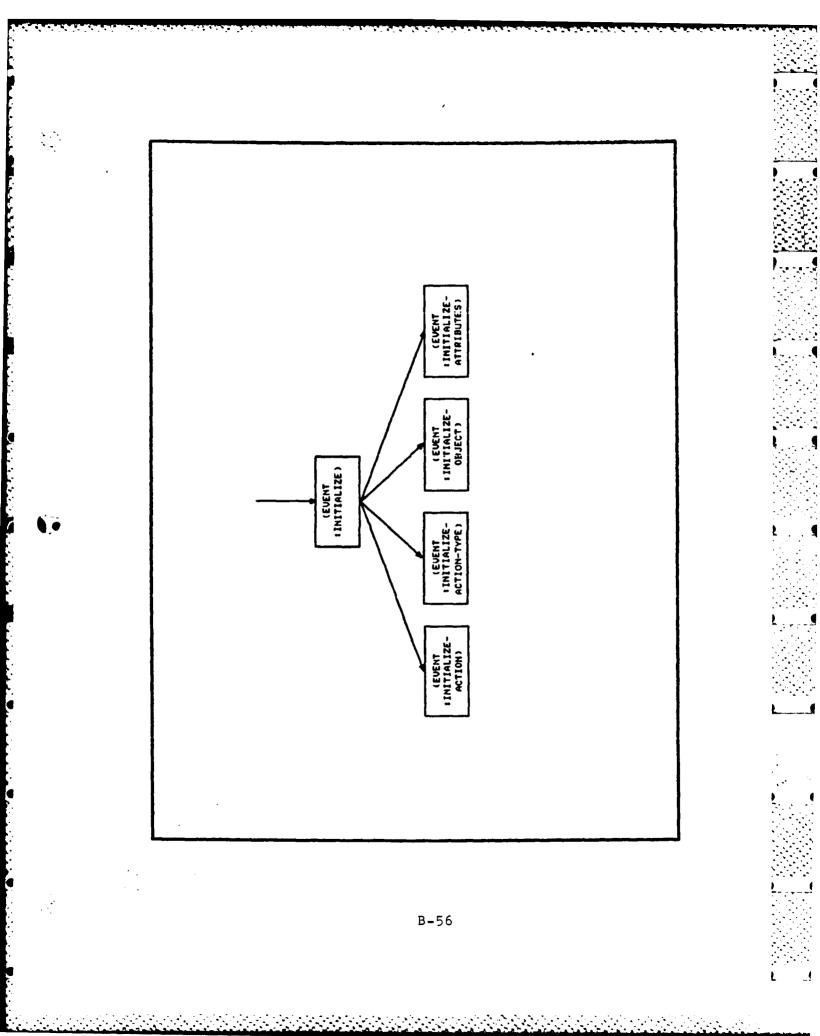
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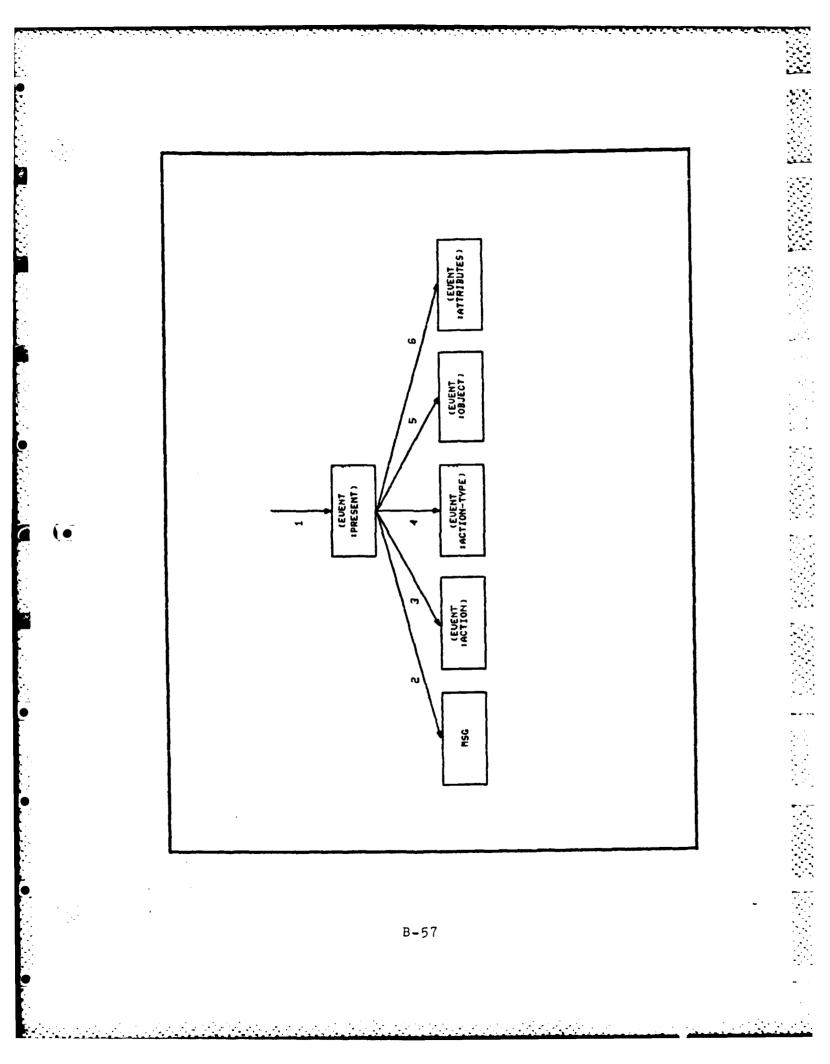
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#	Passed Parameters	Туре	Returned Parameters	Туре
1	(none)			
2	(none)		action type	data
3	(none)		object	data
4	(none)		action	data
5	(none)		attributes	data
6	object action attributes	data data data		





(event :present) Interfaces

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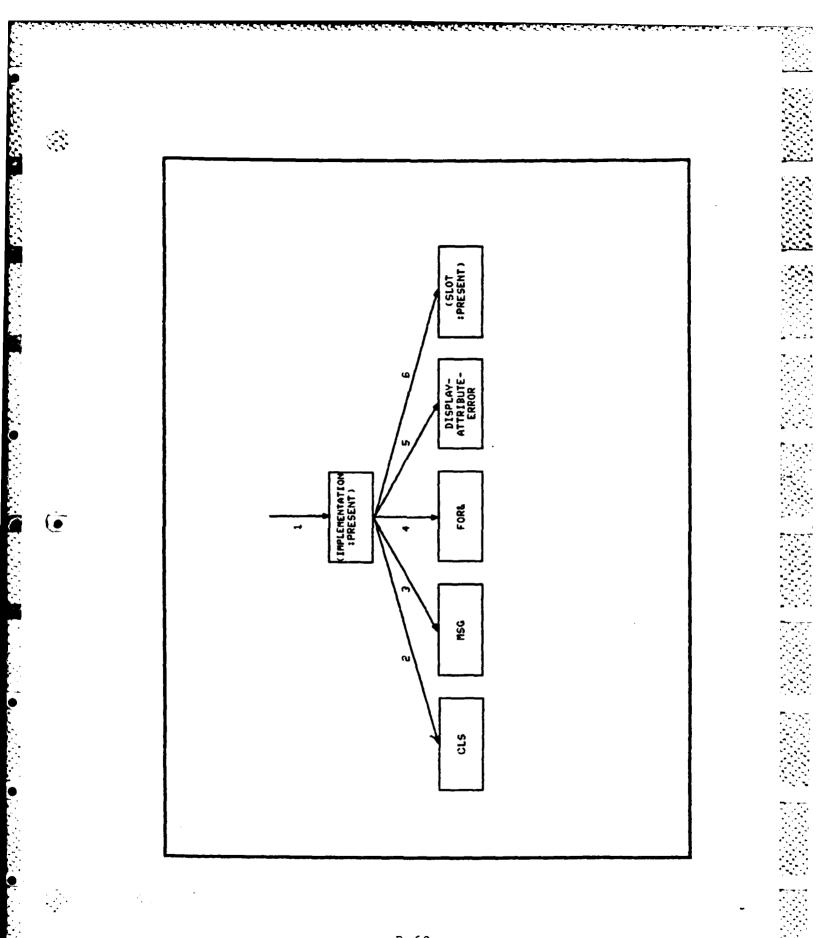
_#	Passed Parameters	Туре	Returned Parameters	Туре
1	(none)			
2	data to display	data		
3	(none)		action	data
4	(none)		action-type	data
5	(none)		object	data
6	(none)		list of attributes	data

IMPLEMENTATION Methods

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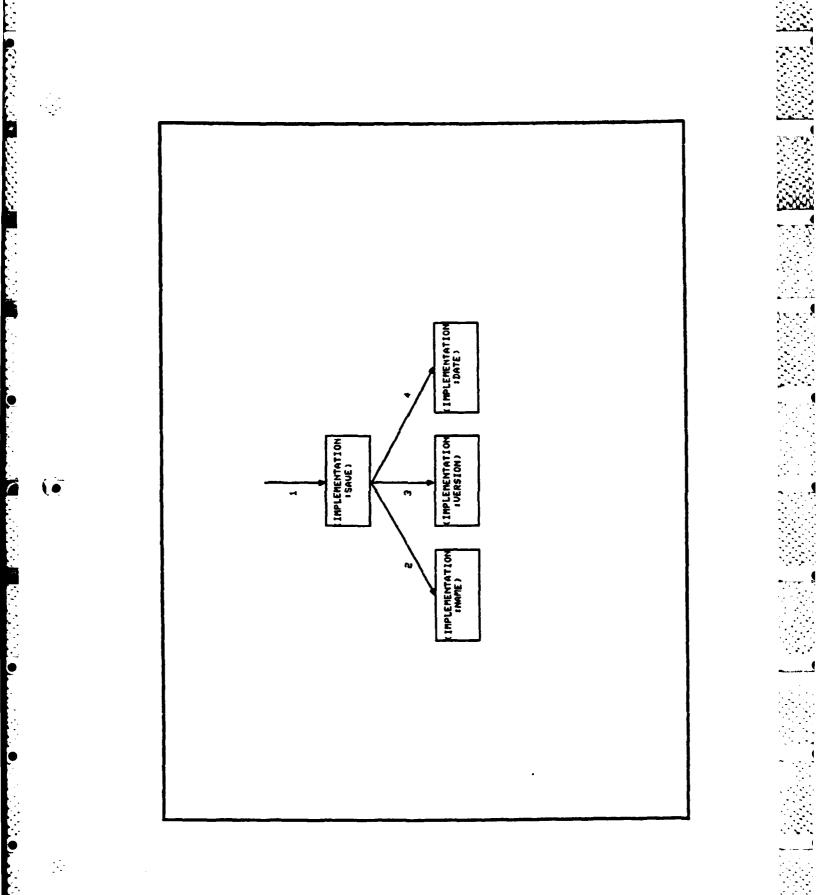
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(implementation :present) Interfaces

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#	Passed Parameters	Туре	Returned	Parameters	Туре
1	(none)		-		
2	list of instance variables	data			
3	(none)				
4	list of instance variables	data			
5	list of invalid instance variables	data			
6	(none)				



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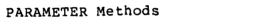
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#	Passed Parameters	Туре	Returned Parameters	Туре
1	project instance pointer stream	data data		
2	(none)		name	data
3	(none)		version	data
4	(none)		date	data







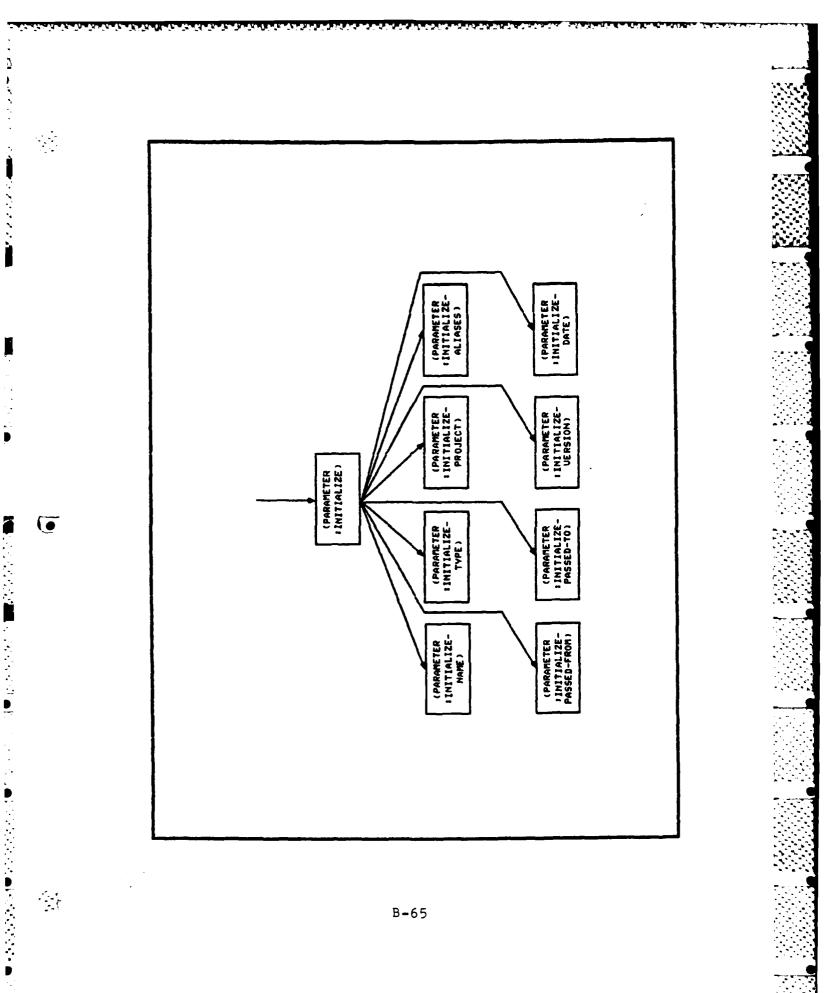


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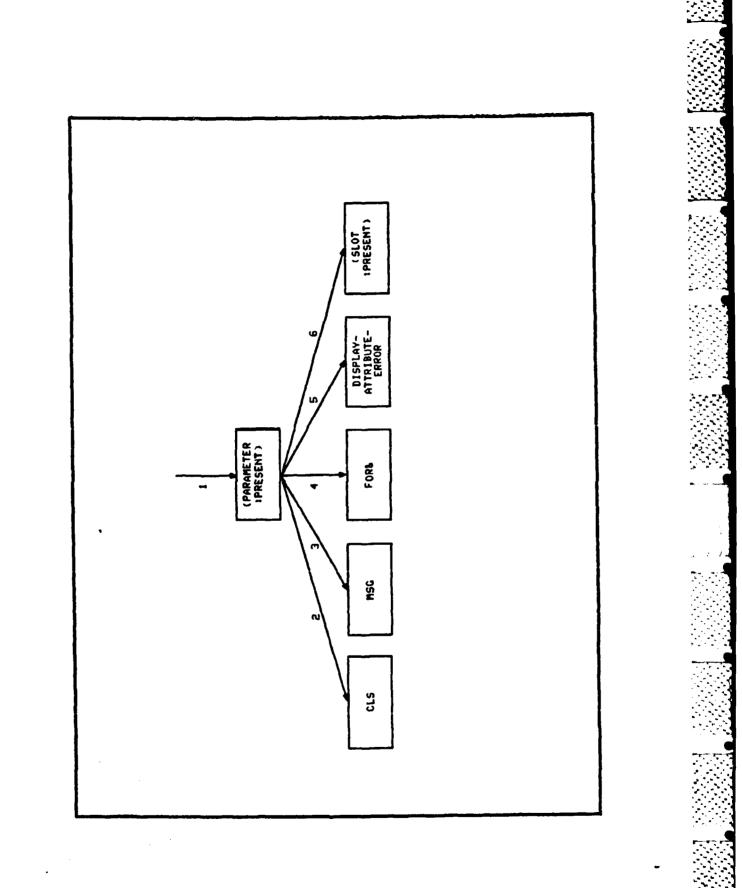
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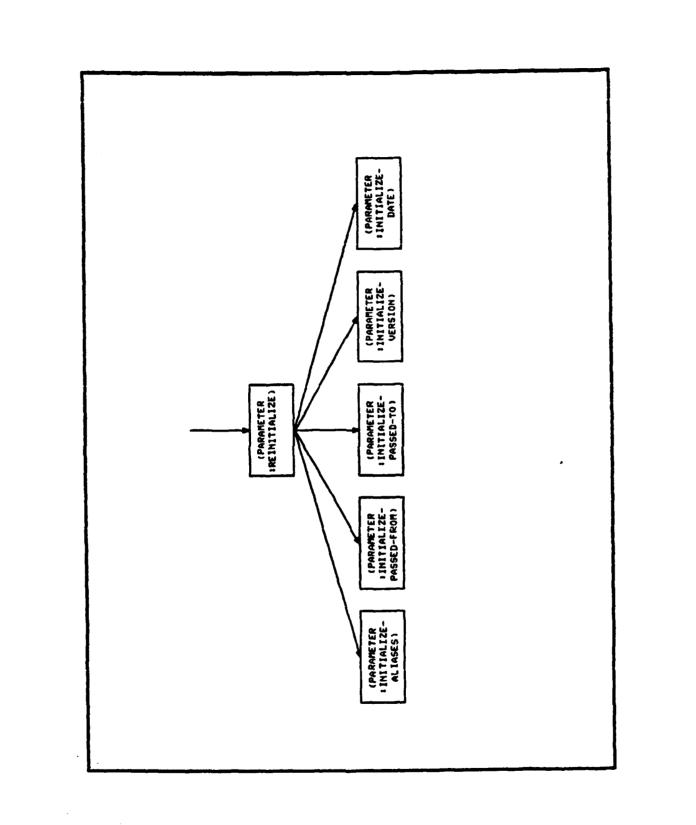
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(parameter :present) Interfaces

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#	Passed Parameters	Туре	Returned	Parameters	Туре
1	(none)				
2	list of instance variables	data			
3	(none)				
4	list of instance variables	data			
5	list of invalid instance variables	data			
6	(none)				



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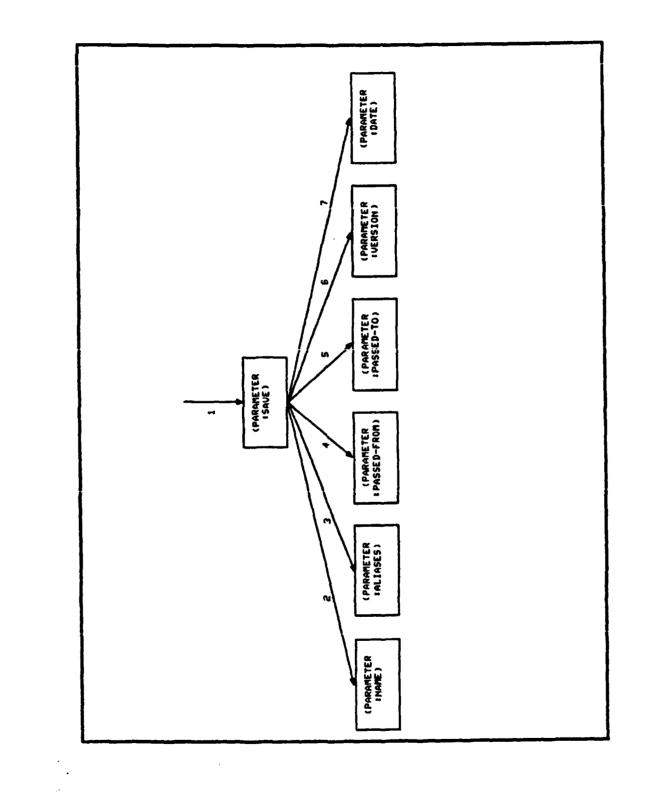
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(parameter :save) Interfaces

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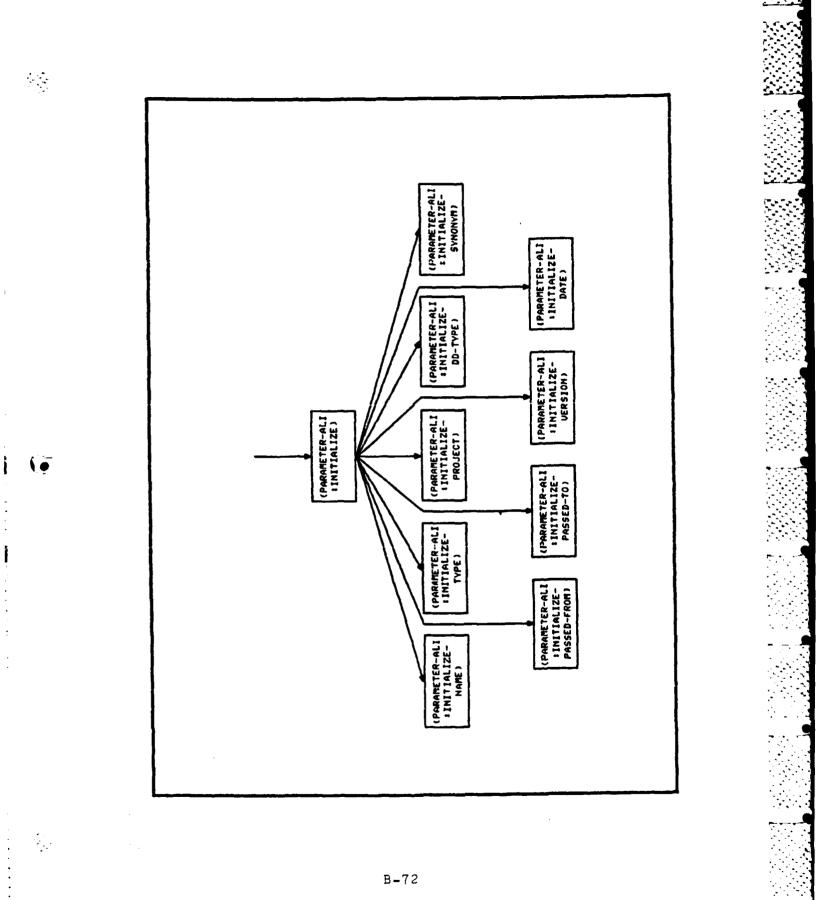
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#	Passed Parameters	Туре	Returned Parameters	Type
1	design instance pointer stream	data data		
2	(none)		name	data
3	(none)		aliases	data
4	(none)		passed-from	data
5	(none)		passed-to	data
6	(none)		version	data
7	(none)		date	data

PARAMETER-ALIAS Methods

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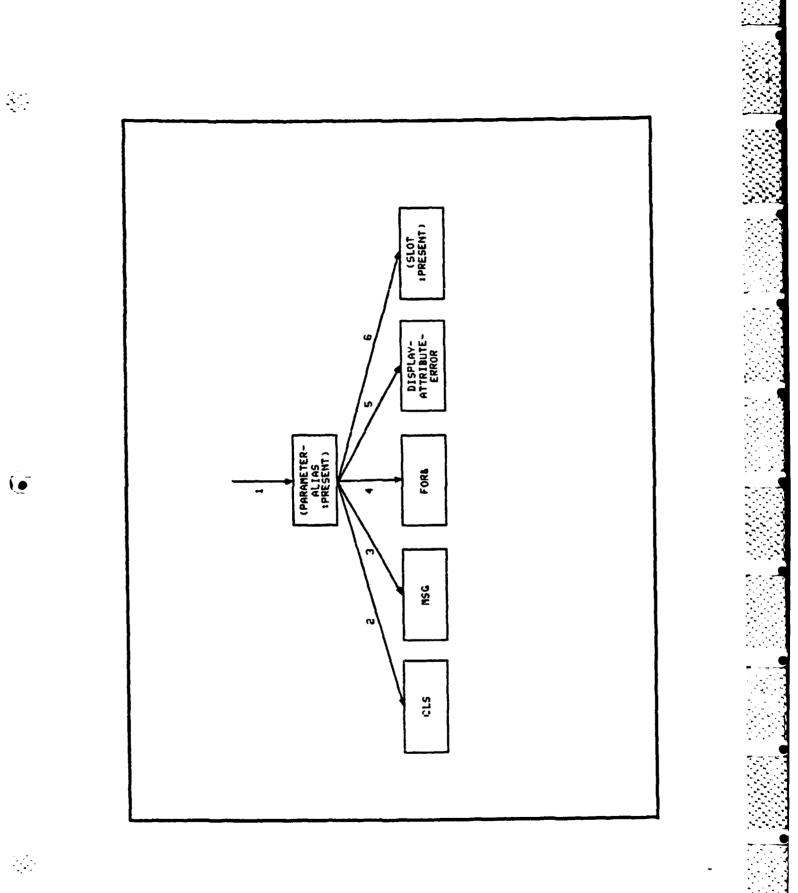


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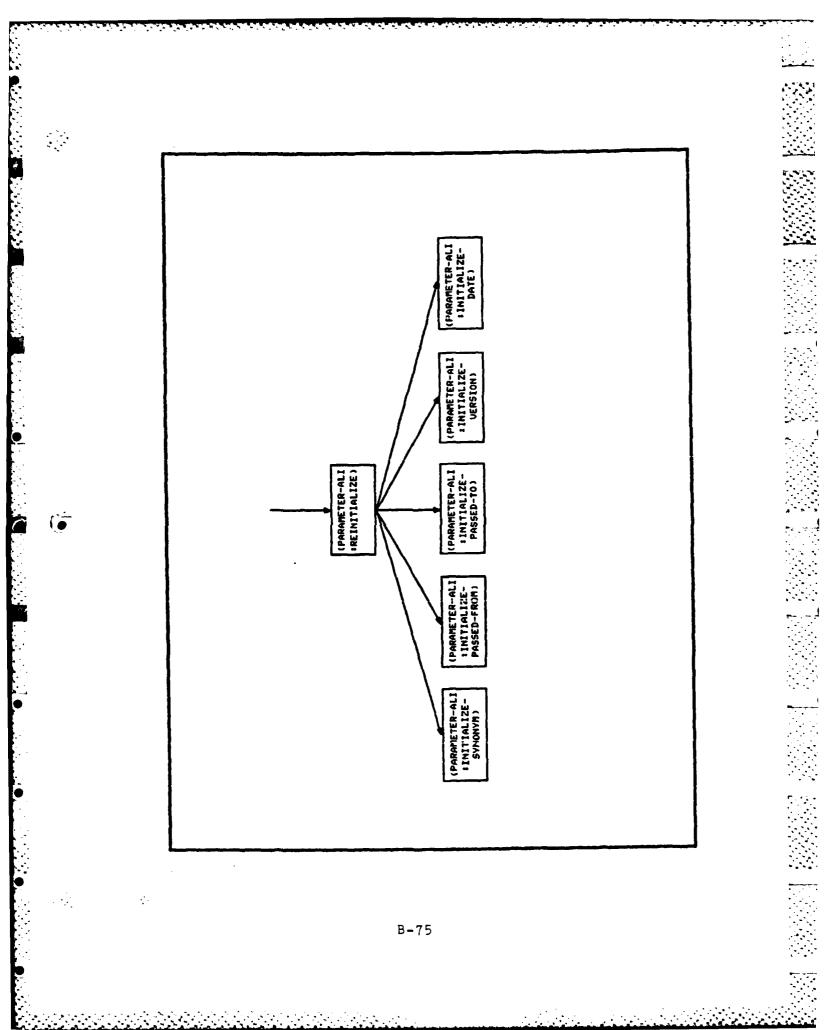
(parameter-alias :present) Interfaces

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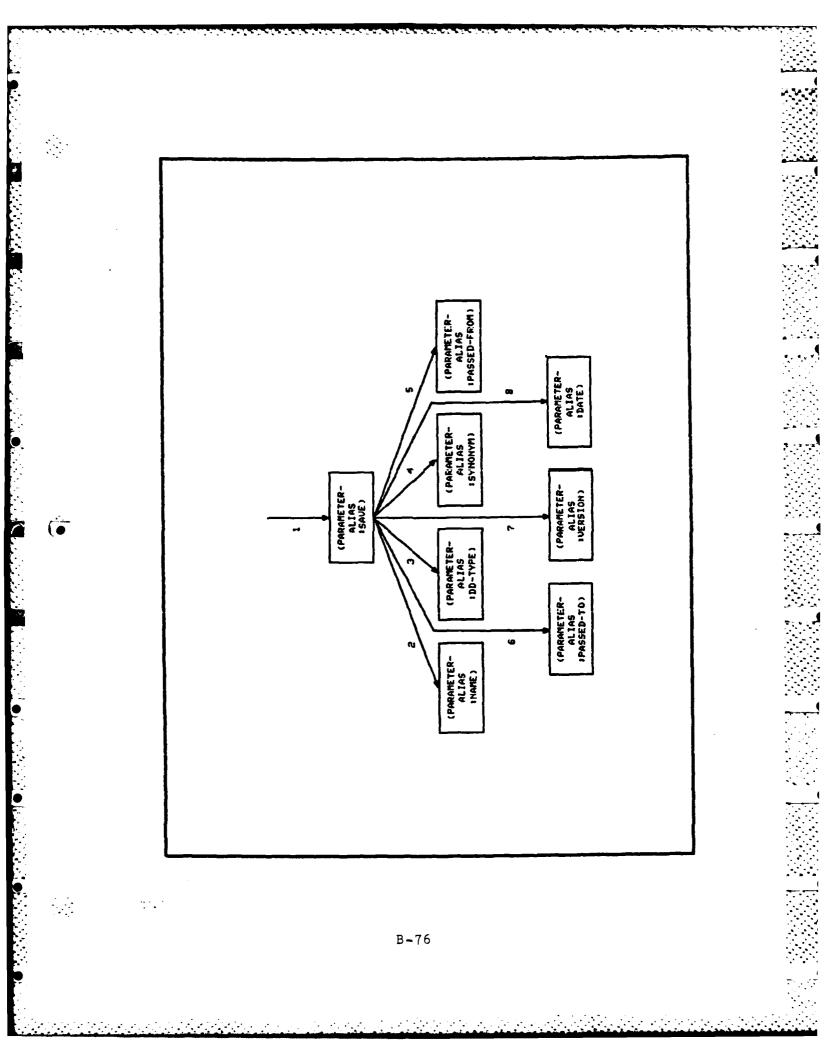
#	Passed Parameters	Туре	Returned	Parameters	Туре
l	(none)				
2	list of instance variables	data			
3	(none)				
4	list of instance variables	data			
5	list of invalid instance variables	data			
6	(none)				

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(parameter-alias :save) Interfaces

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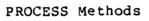
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#	Passed Parameters	Туре	Returned Parameters	Туре
1	design instance pointer stream	data data		
2	(none)		name	data
3	(none)		dd-type	data
4	(none)		synonym	data
5	(none)		passed-from	data
6	(none)		passed-to	data
7	(none)		version	data
8	(none)		date	data





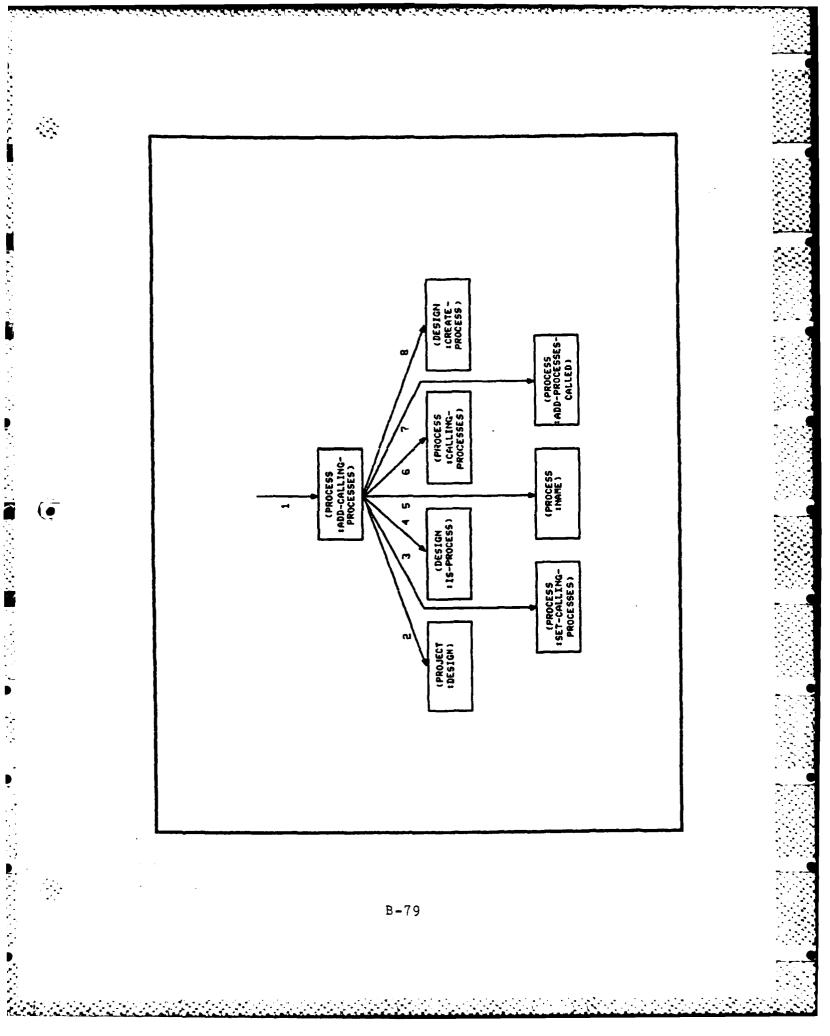
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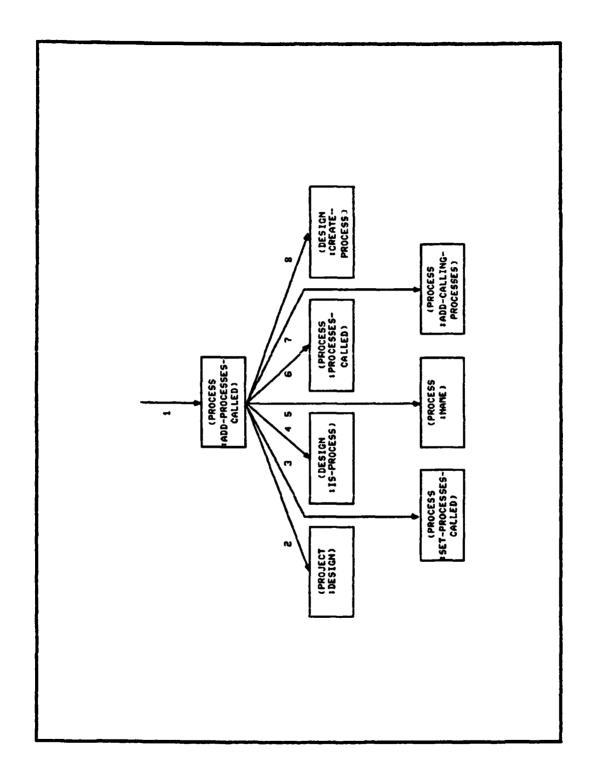


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(process :add-calling-processes) Interfaces

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#	Passed Parameters	Туре	Returned Parameters	Туре
1	list of process names	data		
2	(none)		pointer to design instance	data
3	list of process names	data		
4	process name	data	is process?	flag
5	(none)		name of current process	data
6	(none)		list of calling processes	data
7	name of current process	data		
8	process name	data	pointer to new process instance	data



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(process :add-processes-called) Interfaces

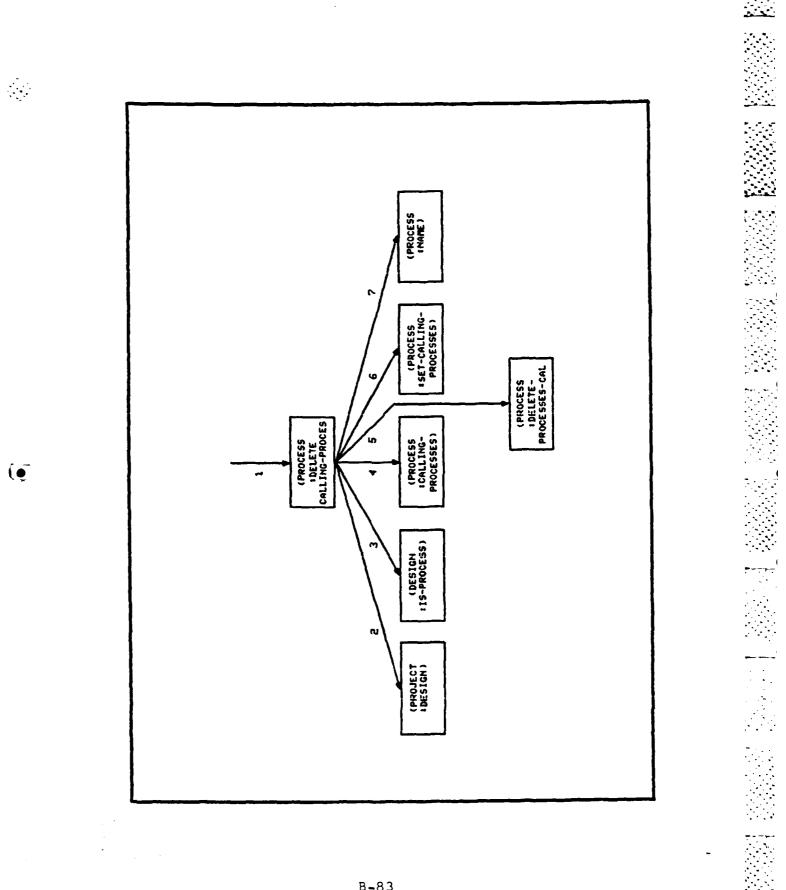
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#	Passed Parameters	Туре	Returned Parameters	Туре
1	list of process names	data		
2	(none)		pointer to design instance	data
3	list of process names	data		
4	process name	data	is process?	flag
5	(none)		name of current process	data
6	(none)		list of calling processes	data
7	name of current process	data		
8	process name	data	pointer to new process instance	data

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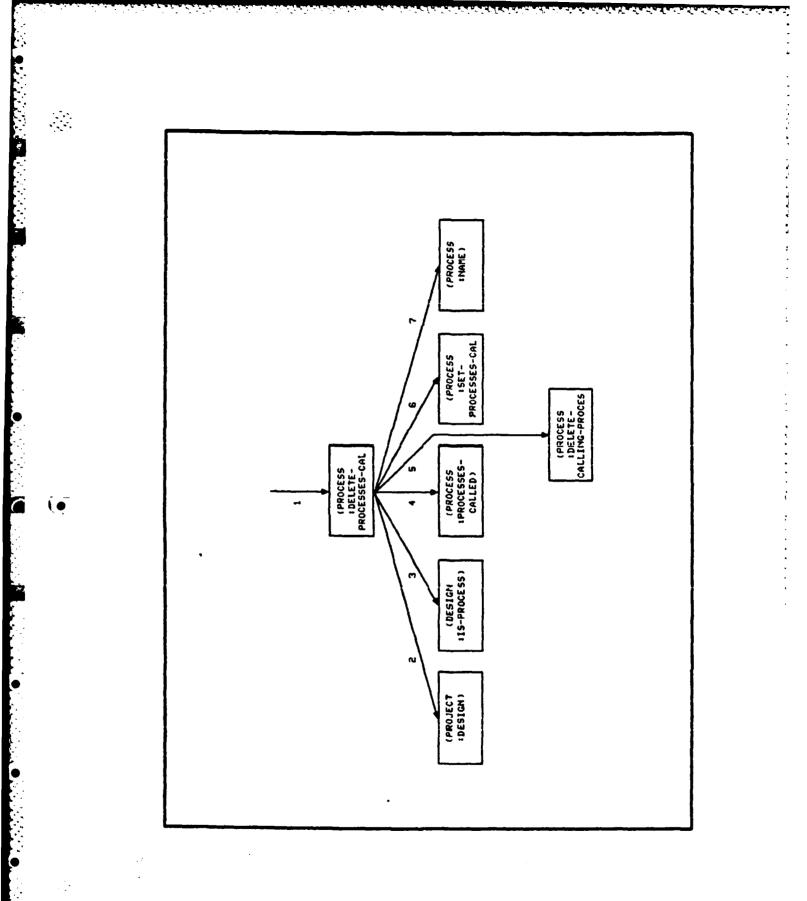
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(process :delete-calling-processes) Interfaces

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#	Passed Parameters	Туре	Returned Parameters	Туре
1	list of process names	data		
2	(none)		pointer to design instance	data
3	process name	data	is process?	flag
4	(none)		list of calling processes	data
5	name of current process	data		
6	list of process names	data		
7	(none)		name of current process	data







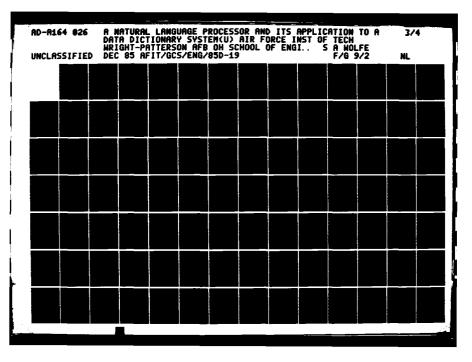
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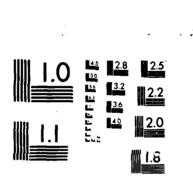
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(process :delete-processes-called) Interfaces

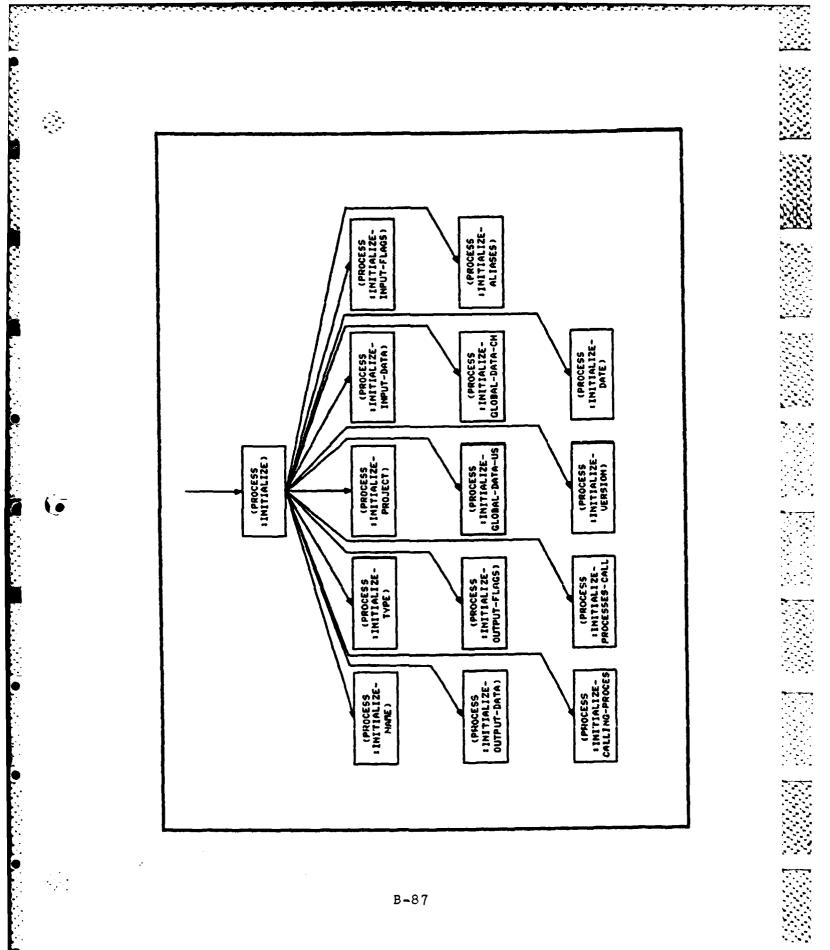
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#	Passed Parameters	Туре	Returned Parameters	Туре
1	list of process names	data		
2	(none)		pointer to design instance	data
3	process name	data	is process?	flag
4	(none)		list of calling processes	data
5	name of current process	data		
6	list of process names	data		
7	(none)		name of current process	data



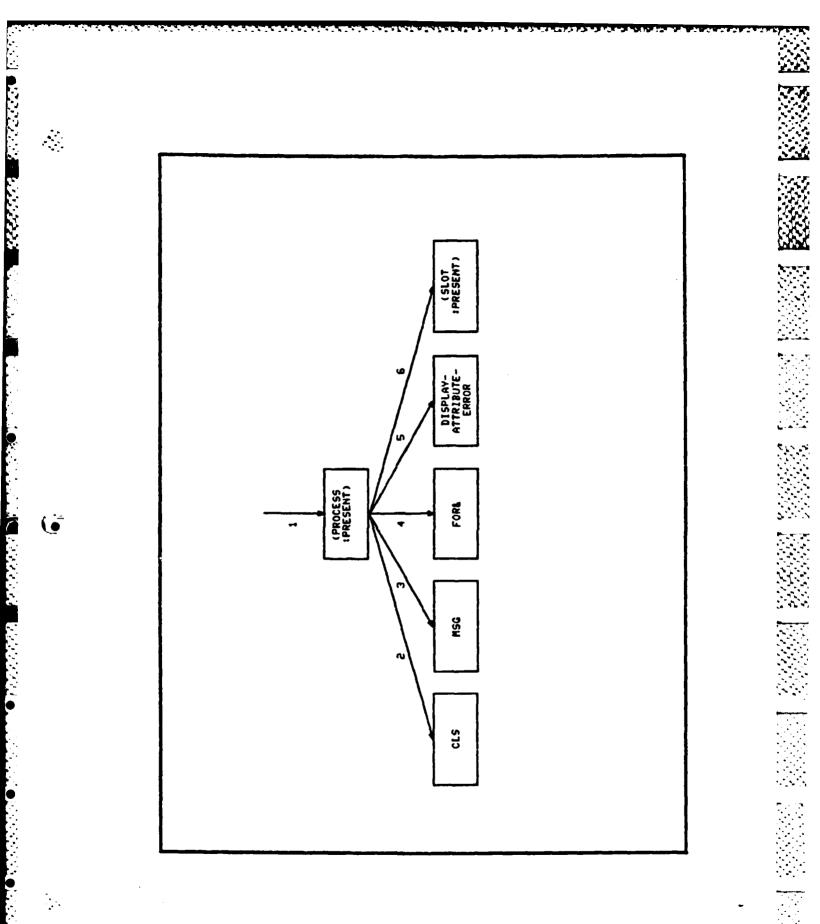




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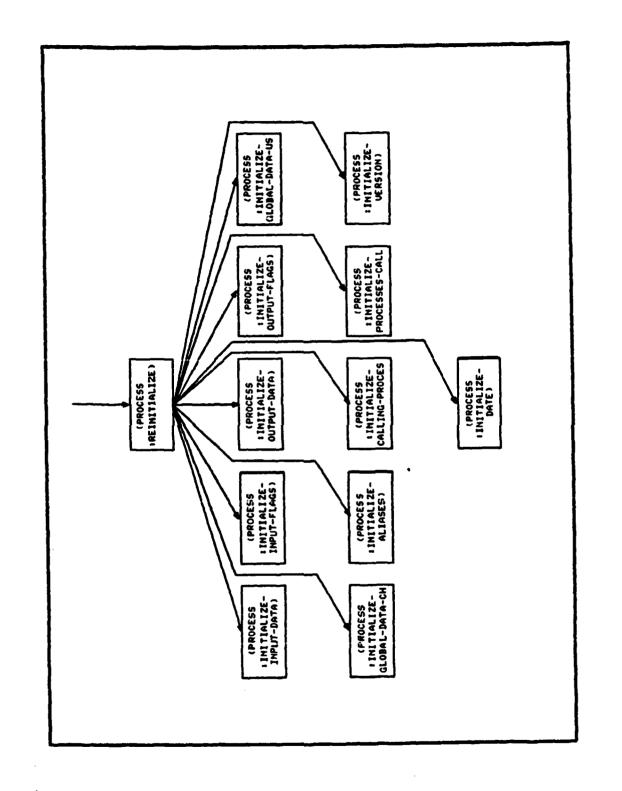




(process :present) Interfaces

#	Passed Parameters	Туре	Returned	Parameters	Туре
l	(none)				
2	list of instance variables	data			
3	(none)				
4	list of instance variables	data			
5	list of invalid instance variables	data			
6	(none)				

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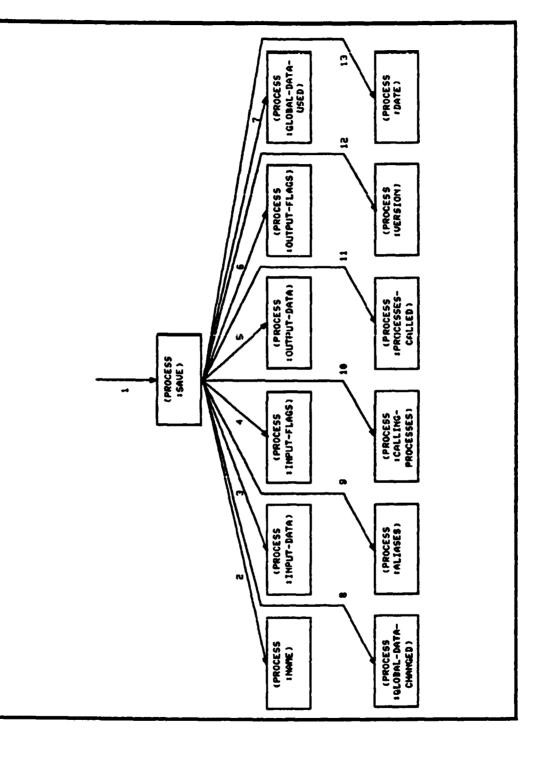








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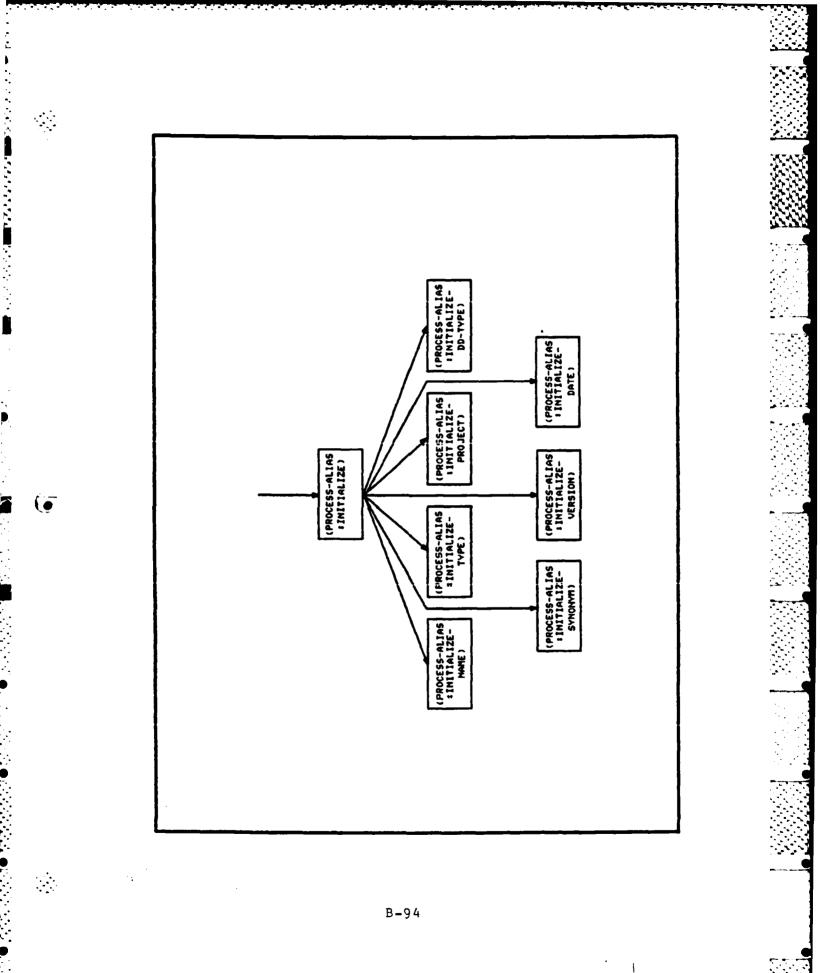
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(process :save) Interfaces

#	Passed Parameters	Туре	Returned Parameters	Туре
1	design instance pointer stream	data data		
2	(none)		name	data
3	(none)		input-data	data
4	(none)		input-flags	data
5	(none)		output-data	data
6	(none)		output-flags	data
7	(none)		global-data-used	data
8	(none)		global-data-changed	data
9	(none)		aliases	data
10	(none)		calling-processes	data
11	(none)		processes-called	data
12	(none)		version	data
13	(none)		date	data

PROCESS-ALIAS Methods



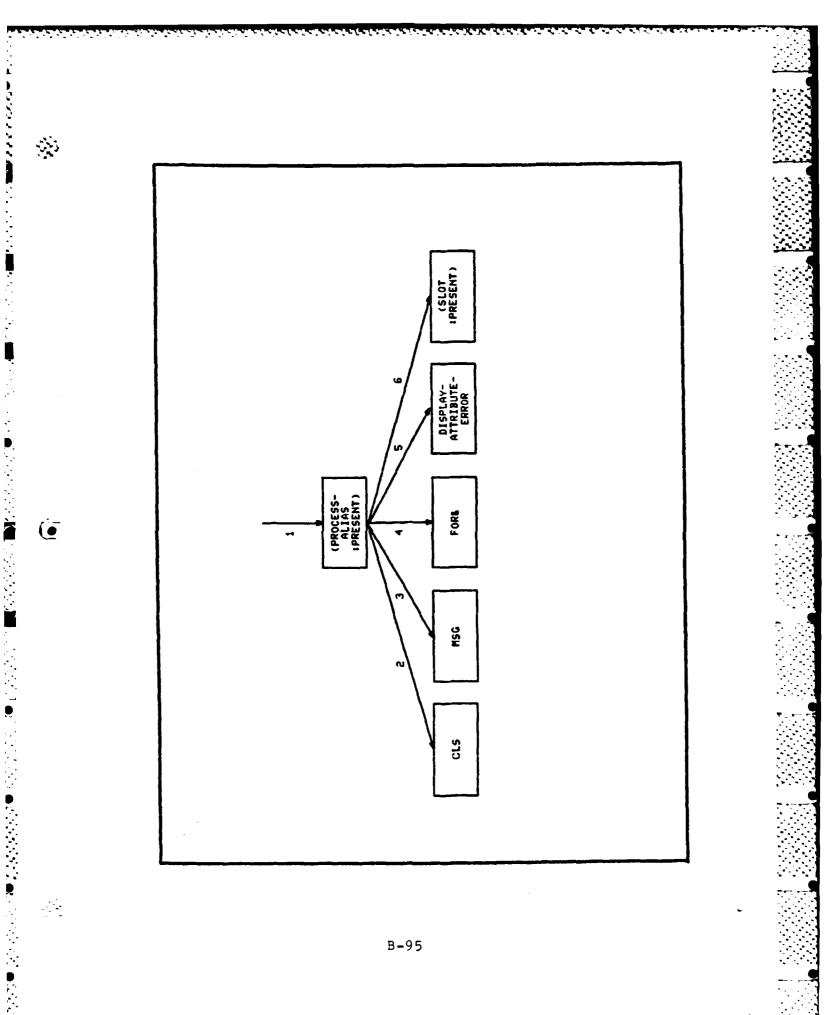
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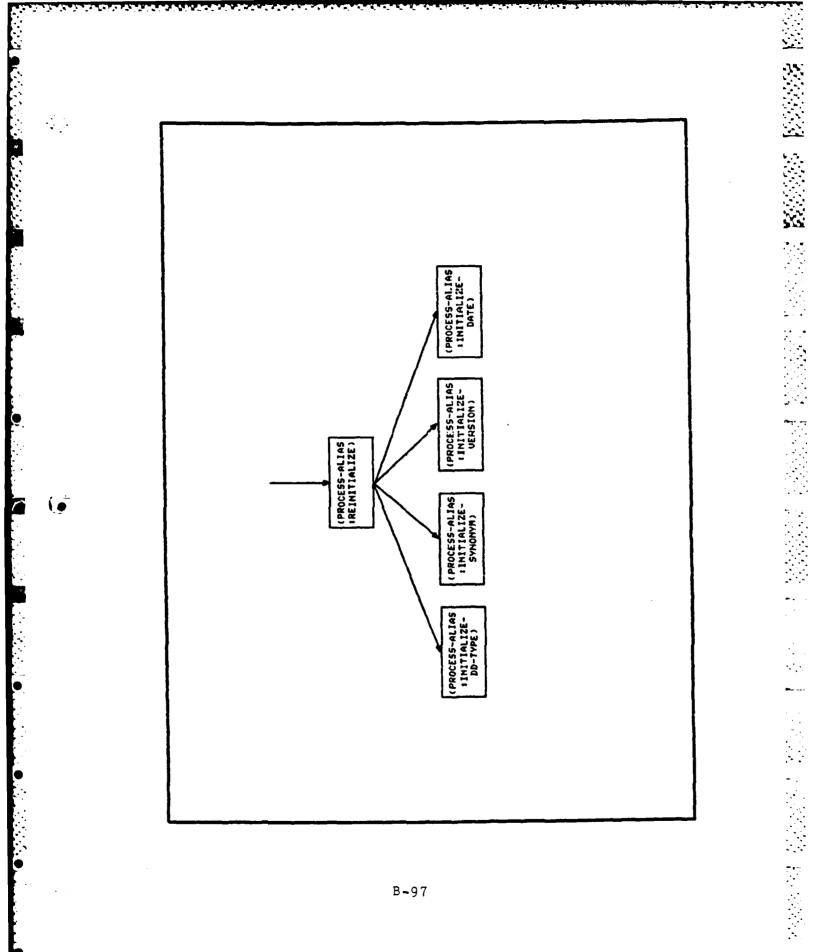


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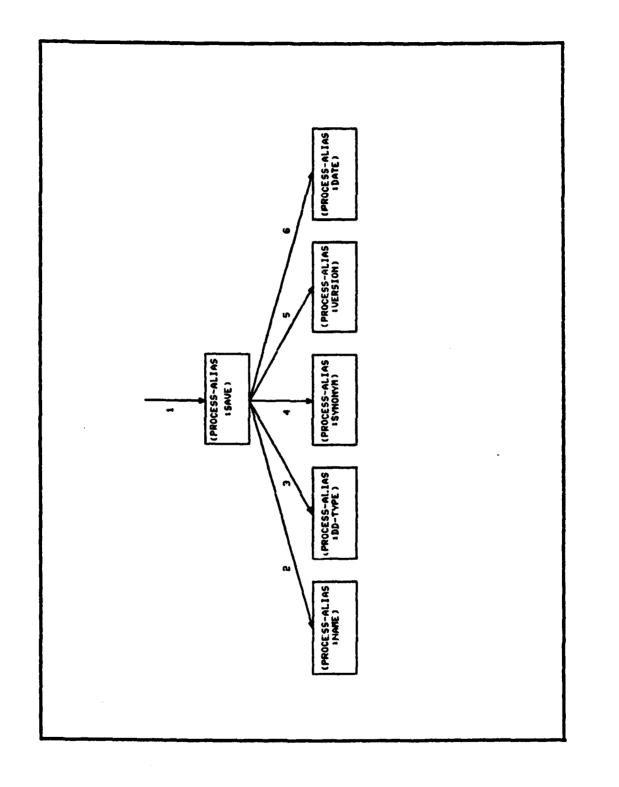
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(process-alias :present) Interfaces

_#	Passed Parameters	Туре	Returned	Parameters	Туре
1	(none)				
2	list of instance variables	data			
3	(none)				
4	list of instance variables	data			
5	list of invalid instance variables	data			
6	(none)				







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#	Passed Parameters	Туре	Returned Parameters	в Туре
1	design instance pointer stream	data data		
2	(none)		name	data
3	(none)		dd-type	data
4	(none)		synonym	data
5	(none)		version	data
6	(none)		date	data

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PROJECT Methods

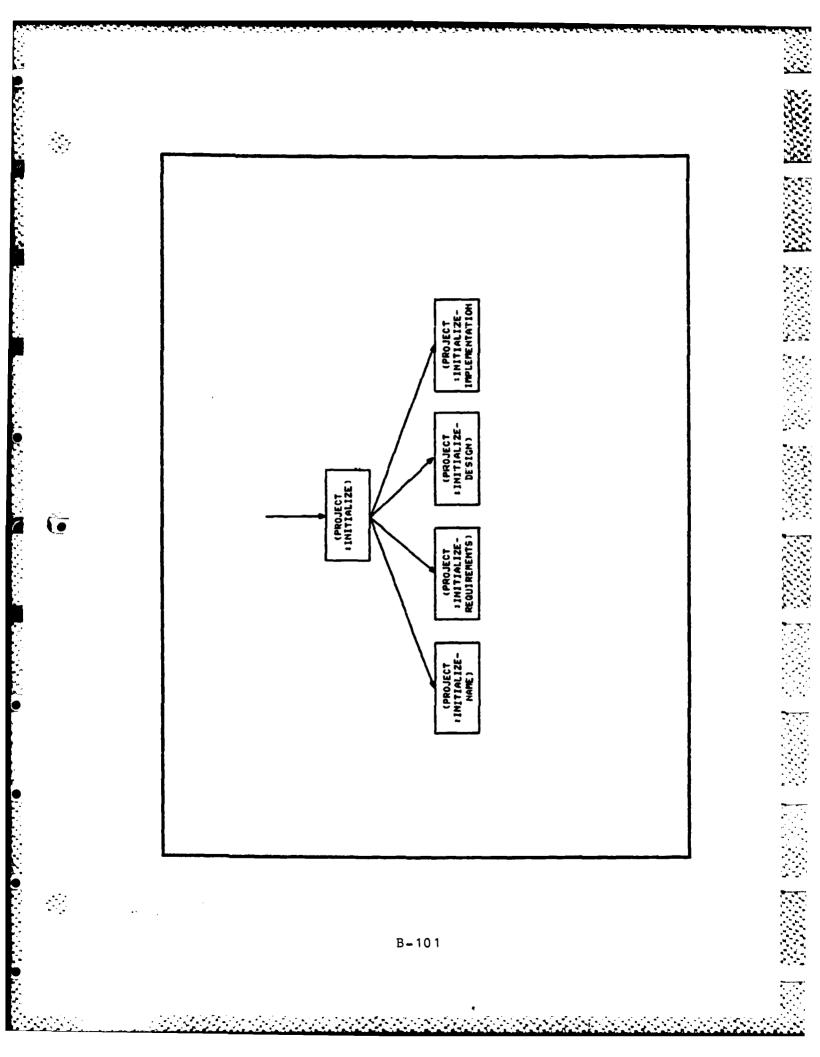
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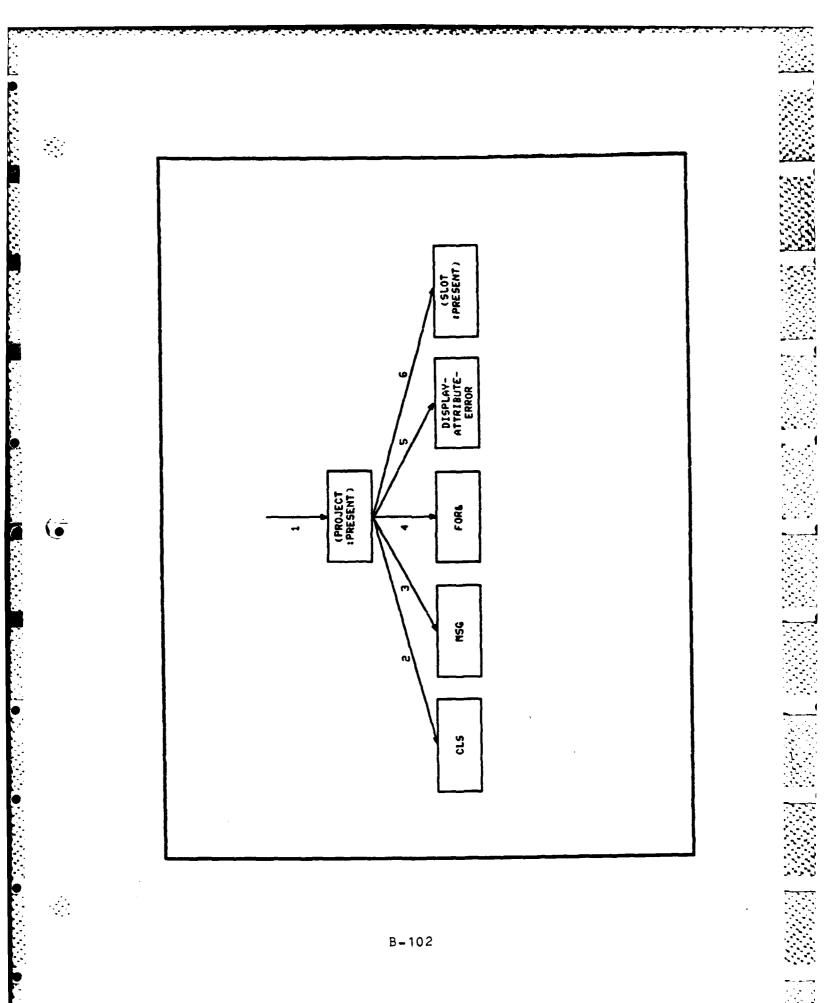
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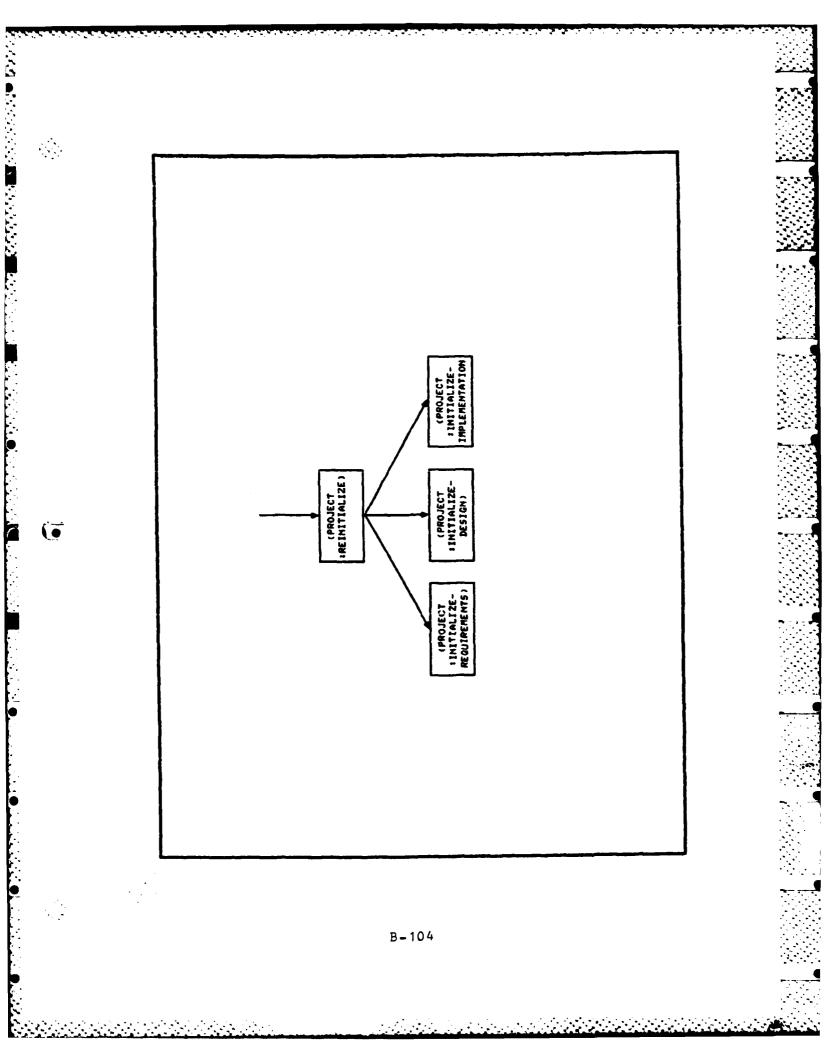
(project :present) Interfaces

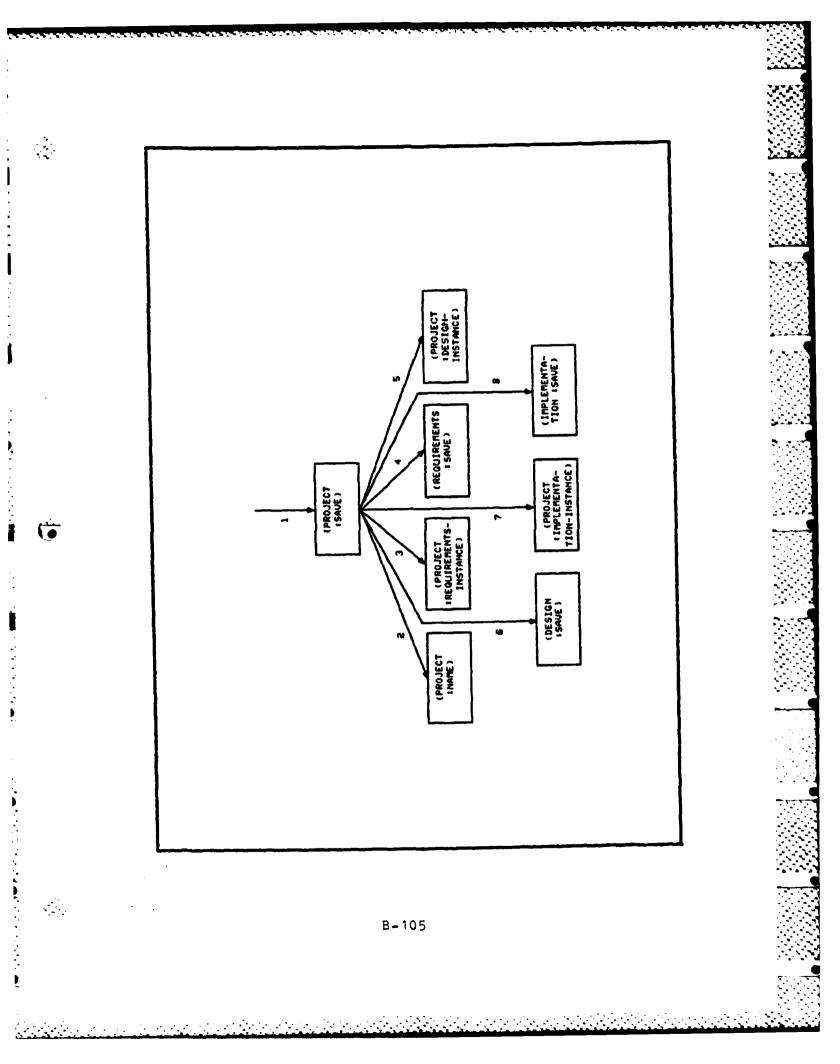
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#	Passed Parameters	Туре	Returned	Parameters	Туре
1	(none)				
2	list of instance variables	data			
3	(none)				
4	list of instance variables	data			
5	list of invalid instance variables	data			
6	(none)				

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(project :save) Interfaces

#	Passed Parameters	Туре	Returned Parameters	Туре
1	(none)			
2	(none)		project name	data
3	(none)		pointer to requirements instance	data
4	pointer to project instance output stream	data data		
5	(none)		pointer to design instance	data
6	pointer to project instance output stream	data data		
7	(none)		pointer to implementation instance	data
8	pointer to project instance output stream	data data		



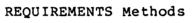




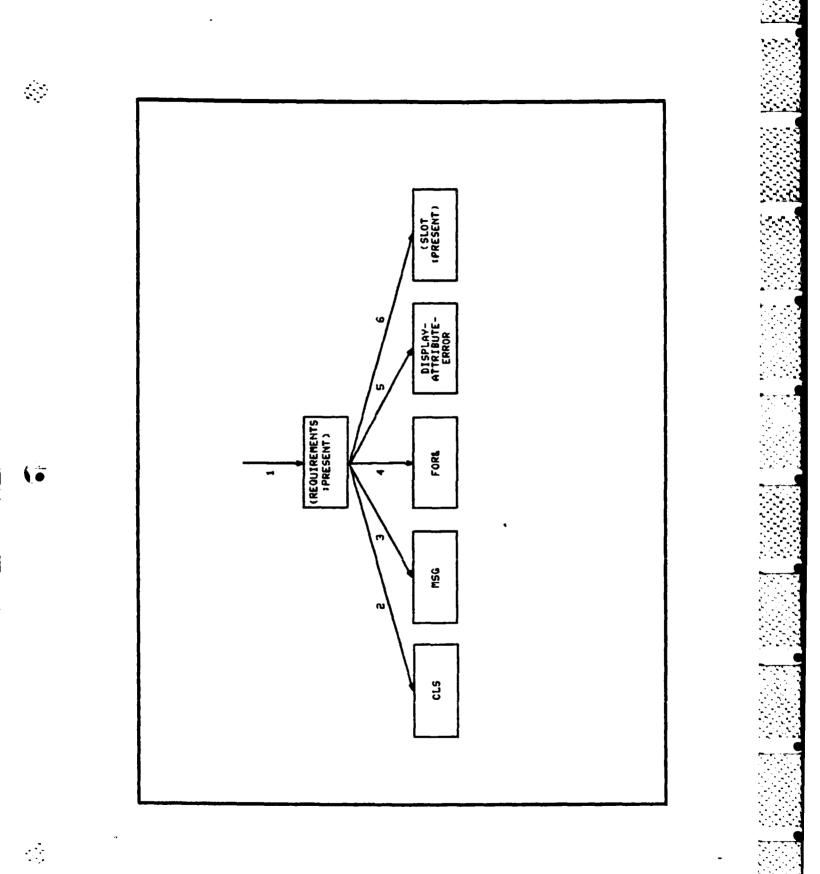














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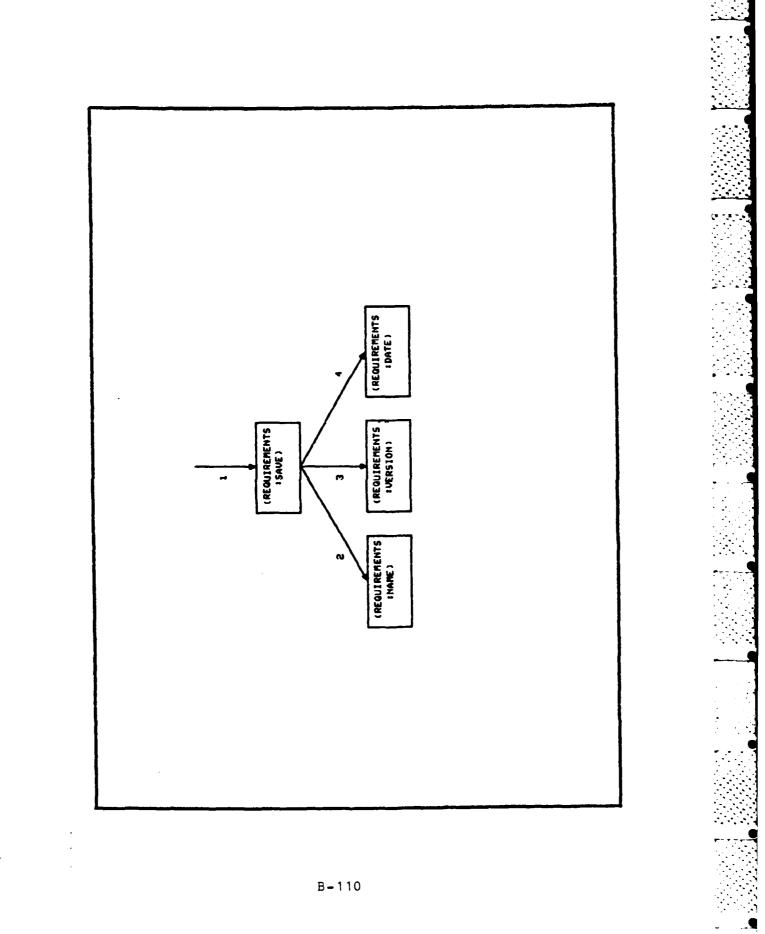
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(requirements :present) Interfaces

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#	Passed Parameters	Туре	Returned	Parameters	Туре
l	(none)				
2	list of instance variables	data			
3	(none)				
4	list of instance variables	data			
5	list of invalid instance variables	data			
6	(none)				



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(requirements :save) Interfaces

#	Passed Parameters	Туре	Returned Parameters	Туре
1	project instance pointer stream	data data		
2	(none)		name	data
3	(none)		version	data
4	(none)		date	data

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Appendix C Grammar Definition (

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Appendix C Table of Contents

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Correspondence Between Instance Variables and Attribute Values	C-10

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Introduction

This appendix defines the grammar of the Data Dictionary System (DDS). The first part of this definition specifies the productions of the grammar using a Backus-Naur Form (BNF) notation. The second part defines the correspondence between the instance variables of the various object types and literal values contained in the <attribute> production of the grammar.

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Grammar Definition

Notation is defined to be ::= [] select one optional () <> class symbol \$1 comment 9 or user input <grammar> ::= (please) <aux-grammar> <activity> ::= * Sa user-defined member of the object class ACTIVITY <activity-alias> ::= * Sa user-defined member of the object class ACTIVITY-ALIAS[†] <activity-alias-instance> ::= activity alias <activity-alias> <activity-instance> ::= activity <activity> <add> ::= [add ¶ append ¶ concatenate ¶ include ¶ insert] <add-command> ::= [<add> <literal> to (the) <attribute> [of ¶ by] <instance> ¶ <add> <literal> to <instance> <attribute> ¶ the <attribute> of <instance> is <literal> ¶ <instance> <attribute> is <literal> ¶ [<module-instance> ¶ <process-instance>] [calls ¶ invokes] <literal> ¶ [<module-instance> ¶

is a submodule of] <literal> }

<process-instance>] [is called by ¶ is invoked by ¶

<attribute> ::=

[activities ¶ activity number ¶ activity numbers where used ¶ activity # ¶ alias ¶ aliases ¶ attributes ¶ called modules ¶ called processes ¶ calling modules ¶ calling processes ¶ common data changed ¶ common data used ¶ common parameters changed ¶ common parameters used ¶ common variables changed ¶ common variables used ¶ component data elements ¶ component parameters ¶ component variables ¶ components ¶ composition ¶ control inputs ¶ controls ¶ data dictionary type ¶ data elements ¶ data type ¶ date ¶ dd type ¶ description ¶ design ¶ design data ¶ design information ¶ destination ¶ destinations ¶ driver ¶ drivers ¶ entry date ¶ entry version ¶ files input ¶ files output ¶ files read ¶ files written ¶ input files ¶ global data changed ¶ global data used ¶ global parameters changed ¶ global parameters used ¶ global variables changed ¶ global variables used ¶ globals changed ¶ globals used ¶ hardware input ¶ hardware output ¶ highest-level activity ¶ implementation ¶ implementation data ¶ implementation information ¶ input data ¶ input files ¶ input flags ¶ inputs ¶ invoked modules ¶ invoked processes ¶ invoking modules ¶ invoking processes ¶ instance variables ¶ legal range ¶ legal values ¶ list of activities ¶ list of aliases ¶ list of called modules ¶ list of called processes ¶ list of calling modules ¶ list of calling processes ¶ list of component data elements ¶ list of component parameters ¶ list of component variables ¶ list of components ¶ list of controls ¶ list of data elements ¶ list of input data ¶ list of input flags ¶ list of inputs ¶ list of invoked modules ¶ list of invoked processes ¶ list of invoking modules ¶ list of invoking processes ¶ list of legal values ¶ list of mechanisms ¶ list of modules ¶ list of modules called ¶ list of modules invoked ¶ list of output data ¶ list of output flags ¶ list of outputs ¶ list of parameter aliases ¶ list of parameters ¶ list of process aliases ¶list of processes ¶ list of processes called ¶ list of processes invoked ¶ list of related items ¶ list of related parameters ¶ list of related requirements ¶ list of submodules ¶ list of subprocesses ¶ list of values ¶ list of variables ¶ main activity ¶ main module ¶ main modules ¶ main process ¶ main processes ¶ main program ¶ main programs ¶ max legal value ¶ max value ¶ maximum legal value ¶ maximum value ¶ mechanisms ¶ min legal value ¶ min value ¶ minimum legal value ¶ minimum value ¶ module number ¶ module #

```
¶ modules ¶ name ¶ node ¶ node number ¶ node # ¶
      number ¶ output data ¶ output files ¶ output flags ¶
      outputs ¶ parameter aliases ¶ parameters ¶ parent
      activity ¶ parent data element ¶ parent parameter ¶
      parent variable ¶ part of ¶ passed from ¶ passed to ¶
      process aliases ¶ process number ¶ process # ¶
      processes ¶ processes called ¶ processes invoked ¶
      project ¶ project name ¶ range ¶ related items ¶
      related parameters ¶ related requirement ¶ related
      requirements ¶ related requirement number ¶ related
      requirements numbers ¶ requirement number ¶
      requirement # ¶ requirements ¶ requirements data ¶
      requirements information ¶ SADT activity numbers where
      used ¶ slots ¶ source ¶ sources ¶ storage type ¶
      submodules ¶ subprocesses ¶ synonym ¶ synonyms ¶ title
      ¶ top-level activity ¶ top-level program ¶ top-level
      programs ¶ type ¶ values ¶ variables ¶ version ¶ # ]
<attribute-list> ::=
    <attribute> ( [ and <attribute> ¶ <attribute-list> ] )
<aux-grammar> ::=
    [ <command> ¶ <meta-command> ]
<cls> ::=
    [ cls ¶ clear screen ¶ clear the screen ]
<command> ::=
    [ <add-command> ¶ <create-command> ¶ <delete-command> ¶
      <destroy-command> ¶ <initialize-command> ¶
      <modify-command> ¶ <present-command> ]
<create> ::=
    [ create ¶ make ]
<create-command> ::=
    [ <create> ( a ) ( new ) <object-class> ( called )
      <literal> ¶ <literal> is ( a ) ( new ) <object-class>
<data-element> ::= *
    Sa user-defined member of the object class
      DATA-ELEMENT<sup>†</sup>
<data-element-alias> ::= *
    Sa user-defined member of the object class
      DATA-ELEMENT-ALIAS†
```

```
<data-element-alias-instance> ::=
    data element alias <data-element-alias>
<data-element-instance> ::=
    data element <data-element>
<delete> ::=
    [ delete ¶ purge ¶ remove ]
<delete-command> ::=
    { <delete> <literal> ( from ) <attribute> of <instance>
      ¶ <delete> <literal> ( from ) <instance> <attribute> ]
<design> ::= *
    §a user-defined member of the object class DESIGN†
<design-instance> ::=
    design <design>
<destroy> ::=
    destroy
<destroy-command> ::=
    <destroy> <instance>
<exit> ::=
    [ all pow ¶ exit ¶ quit ]
<help> ::=
    [ h ¶ help ¶ ? ]
<implementation> ::= *
    §a user-defined member of the object class
      IMPLEMENTATION<sup>†</sup>
<implementation-instance> ::=
    implementation <implementation>
<initialize> ::=
    [ init ¶ initialize ¶ reinit ¶ reinitialize ¶ reset ]
<initialize-command> ::=
    <initialize> <instance>
```

```
<instance> ::=
    [ <activity-instance> ¶ <activity-alias-instance> ¶
      <data-element-instance> ¶
      <data-element-alias-instance> ¶ <design-instance> ¶
      <implementation-instance> ¶ <module-instance> ¶
      <parameter-instance> ¶ <parameter-alias-instance> ¶
      <process-instance> ¶ <process-alias-instance> ¶</process-alias-instance> ¶
      <project-instance> ¶ <requirements-instance> ¶
      <variable-instance> ]
teral> ::= *
    Sany string of characters delimited by blankst
<meta-command> ::=
    [ <cls> ¶ <exit> ¶ <help> ¶ <save> ]
<modify> ::=
    [ alter ¶ change ¶ mod ¶ modify ¶ replace ]
<modify-command> ::=
    [ <modify> <attribute> of <instance> to <literal> ¶
      <modify> <instance> <attribute> to <literal> ]
<module> ::= *
    §a user-defined member of the object class MODULEt
<module-instance> ::=
    module <module>
<parameter> ::= *
    §a user-defined member of the object class PARAMETERt
<parameter-alias> ::= *
    Sa user-defined member of the object class
      PARAMETER-ALIAS<sup>†</sup>
<parameter-alias-instance> ::=
    parameter alias <parameter-alias>
<parameter-instance> ::=
    parameter <parameter>
<present> ::=
    [ display ¶ present ¶ print ¶ show ( me ) ¶ tell ( me )
```

```
" type " what [ is " are ] " who is ]
```

```
<present-command> ::=
    [ <present> the <attribute-list> [ of ¶ by ] <instance>
      ¶ <present> <instance> ( <attribute-list> ) ]
<process> ::= *
    §a user-defined member of the object class PROCESS†
<process-alias> ::= *
    §a user-defined member of the object class
      PROCESS-ALIAS†
<process-alias-instance> ::=
    process alias <process-alias>
<process-instance> ::=
    process <process>
<project> ::= *
    §a user-defined member of the object class PROJECT<sup>†</sup>
<project-instance> ::=
    [ project <project> ¶ *project* ]
<requirements> ::= *
    §a user-defined member of the object class
      REQUIREMENTS†
<requirements-instance> ::=
    requirements < requirements>
<save> ::=
    [ keep ¶ save ¶ store ]
<variable> ::= *
    §a user-defined member of the object class VARIABLE†
<variable-instance> ::=
    variable <variable>
```

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Correspondence	Between	Instance	Variables
and l	Attribute	Values	

Instance Variable	Attribute Value	Implemented?
activities	activities	
	list of activities	
aliases		
	alias	
	aliases	Y
	list of aliases	
calling-processes		
	calling modules	
	calling processes	Y
	invoking modules	
	invoking processes	Y
	list of calling modules	
	list of calling processe	S
	list of invoking modules	
	list of invoking process	es
composition		
	component data elements	
	component parameters	
	component variables	
•	components	
	composition	
	list of component data	
	elements	
	list of componen	
	parameters	
	list of component variab	les
	list of components	
controls		
	control inputs	
	controls	
	list of controls	
data-elements		
autu eremented	data elements	
	list of data elements	

Instance Variable	Attribute Value	Implemented?
data-type	data type	
date	date entry date	Y Y
dd-type	data dictionary type dd type	Y
description	description	
design	design design data design information	Y
files-read	files input files read input files	
files-written	files output files written output files	
global-data-changed	common data changed common parameters change common variables changed	
	global data changed global parameters change global variables changed globals changed	
global-data-used	common data used common parameters used common variables used	
	global data used global parameters used global variables used globals used	Y Y

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Instance Variable	Attribute Value	Implemented?
hardware-input	hardware input hardware read	
hardware-output	hardware output hardware write	
implementation	implementation implementation data implementation informatio	Y nc
input-data	input data list of input data	Y
input-flags	input flags list of input flags	Y
inputs	inputs list of inputs	
main-program	driver drivers main module main modules main process main processes main program main programs top-level program top-level programs	Y Y Y Y
max-value	max legal value max value maximum legal value maximum value	
mechanisms	list of mechanisms mechanisms	

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Instance Variable	Attribute Value	Implemented?
min-value	min legal value min value minimum legal value minimum value	
modules		
	list of modules modules	
name		
	name title	Y
number		
	activity number activity # node node number node # number #	
output-data		
-	list of output data output data	Y
output-flags		
	list of output flags output flags	Y
outputs		
	list of outputs outputs	
parameter-aliases		
	list of parameter aliase parameter aliases	s Y
parameters		
	list of parameters parameters	Y
parent-activity		
	parent activity	

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Instance Variable	Attribute Value	Implemented?
part-of	parent data element parent parameter parent variable part of	
passed-from	passed from source sources	Y Y Y
passed-to .	destination destinations passed to	Y Y Y
process-aliases	list of process aliases process aliases	Y
processes	list of processes processes	Y
processes-called	called modules called processes	Y
	<pre>invoked modules invoked processes list of called modules list of called processes list of invoked modules list of invoked processes list of modules called list of modules invoked list of processes called list of processes invoked list of submodules list of subprocesses modules called modules invoked</pre>	đ
	processes called processes invoked	Y Y
	submodules subprocesses	Y

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Instance Variable	Attribute Value	Implemented?
project	project project name	Y Y
range	legal range range	
related-items	list of related items list of related paramet list of related requirements related items related parameters related requirements numbers related requirements # related requirements # related SADT activity related SADT data items requirement #	mber
requirements	requirements requirements data requirements informatio	Y
slots	attributes instance variables slots	
storage-type	storage type	
synonym	synonym synonyms	Y Y
top-level-activity	highest-level activity main activity top-level activity	

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Instance Variable	Attribute Value	Implemented?
type	type	Y
values	legal values list of legal values list of values values	
variables	list of variables variables	
version	entry version version	Y Y
where-used	activity numbers where us SADT activity numbers whe used where used	

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Appendix D

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Implemented Grammar



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Implemented Grammar	Listing	•	•	•	•	•	•	•	•	•	•	•	•	•	D- 4	

Introduction

This appendix contains a listing of the Data Dictionary System's (DDS) implemented grammar. It is presented in the same Lisp-readable format as the actual grammar disk file.

(DECLARE (SPECIAL *META-SYMBOLS*)) (SETQ *META-SYMBOLS* NIL) (DECLARE (SPECIAL <GRAMMAR>)) (CREATE-PRODUCTION <GRAMMAR>) (SETO <GRAMMAR> '(<GRAMMAR> ((PLEASE ((<AUX-GRAMMAR> ((\$ NIL NIL)) NIL)) NIL) (<AUX-GRAMMAR> ((\$ NIL NIL)) NIL)) (SEND *NEW-EVENT* :EXECUTE))) (DECLARE (SPECIAL <AUX-GRAMMAR>)) (CREATE-PRODUCTION <AUX-GRAMMAR>) (SETO <AUX-GRAMMAR> '(<AUX-GRAMMAR> ((<COMMAND> ((* NIL (SEND *NEW-EVENT* :SET-ACTION-TYPE NIL))) NIL) (<META-COMMAND> ((* NIL (SEND *NEW-EVENT* :SET-ACTION-TYPE 'META))) NIL)) NIL)) (DECLARE (SPECIAL <COMMAND>)) (CREATE-PRODUCTION <COMMAND>) (SETQ <COMMAND> '(<COMMAND> ((<PRESENT-COMMAND> ((* NIL NIL)) NIL) (<INITIALIZE-COMMAND> ((* NIL NIL)) NIL)) NIL)) (DECLARE (SPECIAL < ADD-COMMAND>)) (CREATE-PRODUCTION < ADD-COMMAND>) (SETO <ADD-COMMAND> '(<ADD-COMMAND> NIL (SEND *NEW-EVENT* :SET-ACTION ':ADD))) (DECLARE (SPECIAL < CREATE-COMMAND>)) (CREATE-PRODUCTION <CREATE-COMMAND>) (SETQ < CREATE-COMMAND> '(<CREATE-COMMAND> NIL (SEND *NEW-EVENT* :SET-ACTION ':CREATE))) (DECLARE (SPECIAL <DELETE-COMMAND>)) (CREATE-PRODUCTION <DELETE-COMMAND>) (SETQ <DELETE-COMMAND> '(<DELETE-COMMAND> NIL (SEND *NEW-EVENT* :SET-ACTION ':DELETE)))

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```
(DECLARE (SPECIAL <DESTROY-COMMAND>))
(CREATE-PRODUCTION <DESTROY-COMMAND>)
(SETQ < DESTROY-COMMAND>
      '(<DESTROY-COMMAND> NIL (SEND *NEW-EVENT* :SET-ACTION
':DESTROY)))
(DECLARE (SPECIAL <INITIALIZE-COMMAND>))
(CREATE-PRODUCTION <INITIALIZE-COMMAND>)
(SETQ <INITIALIZE-COMMAND>
      '(<INITIALIZE-COMMAND>
      ((<INITIALIZE> ((<INSTANCE> ((* NIL NIL)) NIL)) NIL))
      (SEND *NEW-EVENT* :SET-ACTION ':REINITIALIZE)))
(DECLARE (SPECIAL < MODIFY-COMMAND>))
(CREATE-PRODUCTION < MODIFY-COMMAND>)
(SETQ < MODIFY-COMMAND>
      '(<MODIFY-COMMAND> NIL (SEND *NEW-EVENT* :SET-ACTION
':MODIFY)))
(DECLARE (SPECIAL <PRESENT-COMMAND>))
(CREATE-PRODUCTION < PRESENT-COMMAND>)
(SETO
 <PRESENT-COMMAND>
 '(<PRESENT-COMMAND>
    ((<PRESENT>
       ((<INSTANCE> ((* NIL NIL) (<ATTRIBUTE-LIST> ((* NIL
NIL)) NIL)) NIL)
     (THE ((<ATTRIBUTE-LIST> ((BY ((<INSTANCE> NIL NIL))
NIL)
                      (OF ((<INSTANCE> ((* NIL NIL)) NIL))
NIL))
                    NIL))
          NIL))
       NIL))
    (SEND *NEW-EVENT* :SET-ACTION ':PRESENT)))
(DECLARE (SPECIAL < ADD>))
(CREATE-PRODUCTION < ADD>)
(SETQ
  <ADD> '(<ADD> ((ADD ((* NIL NIL)) NIL) (APPEND ((* NIL
NIL)) NIL)
           (CONCATENATE ((* NIL NIL)) NIL)
           (INCLUDE ((* NIL NIL)) NIL) (INSERT ((* NIL NIL))
NIL))
          NIL))
```

D-5

(DECLARE (SPECIAL <CREATE>)) (CREATE-PRODUCTION <CREATE>) (SETO <CREATE> '(<CREATE> ((CREATE ((* NIL NIL)) NIL) (MAKE ((* NIL NIL)) NIL)) NIL)) (DECLARE (SPECIAL <DELETE>)) (CREATE-PRODUCTION <DELETE>) (SETO <DELETE> '(<DELETE> ((DELETE ((* NIL NIL)) NIL) (PURGE ((* NIL NIL)) NIL) (REMOVE ((* NIL NIL)) NIL)) NIL)) (DECLARE (SPECIAL <DESTROY>)) (CREATE-PRODUCTION < DESTROY>) (SETQ <DESTROY> '(<DESTROY> ((DESTROY ((* NIL NIL)) NIL)) NIL)) (DECLARE (SPECIAL <INITIALIZE>)) (CREATE-PRODUCTION <INITIALIZE>) (SETQ <INITIALIZE> '(<INITIALIZE> ((REINIT ((* NIL NIL)) NIL) (REINITIALIZE ((* NIL NIL)) NIL) (RESET ((* NIL NIL)) NIL) (INIT ((* NIL NIL)) NIL) (INITIALIZE ((* NIL NIL)) NIL)) NIL)) (DECLARE (SPECIAL < MODIFY>)) (CREATE-PRODUCTION < MODIFY>) (SETO <MODIFY> '(<MODIFY> ((ALTER ((* NIL NIL)) NIL) (CHANGE ((* NIL NIL)) NIL) (MOD ((* NIL NIL)) NIL) (MODIFY ((* NIL NIL)) NIL) (REPLACE ((* NIL NIL)) NIL)) NIL))

D-6

(DECLARE (SPECIAL < PRESENT>)) (CREATE-PRODUCTION < PRESENT>) (SETQ PRESENT> '(<PRESENT> ((WHO ((IS ((* NIL NIL)) NIL)) NIL)) (WHAT ((IS ((* NIL NIL)) NIL) (ARE ((* NIL NIL)) NIL)) NIL) (TYPE ((* NIL NIL)) NIL) (TELL ((ME ((* NIL NIL)) NIL) (* NIL NIL)) NIL) (SHOW ((ME ((* NIL NIL)) NIL) (* NIL NIL)) NIL) (PRINT ((* NIL NIL)) NIL) (PRESENT ((* NIL NIL)) NIL) (DISPLAY ((* NIL NIL)) NIL)) NIL)) (DECLARE (SPECIAL <META-COMMAND>)) (CREATE-PRODUCTION <META-COMMAND>) (SETO <META-COMMAND> '(<META-COMMAND> ((<SAVE> ((* NIL NIL)) NIL) (<HELP> ((* NIL NIL)) NIL) (<EXIT> ((* NIL NIL)) NIL) (<CLS> ((* NIL NIL)) NIL)) NIL)) (DECLARE (SPECIAL <CLS>)) (CREATE-PRODUCTION <CLS>) (SETQ <CLS> '(<CLS> ((CLS ((* NIL NIL)) NIL) (CLEAR ((THE ((SCREEN ((* NIL NIL)) NIL)) NIL) (SCREEN ((* NIL NIL)) NIL)) NIL)) (CLS)))(DECLARE (SPECIAL <EXIT>)) (CREATE-PRODUCTION < EXIT>) (SETQ <EXIT> '(<EXIT> ((EXIT ((* NIL NIL)) NIL) (QUIT ((* NIL NIL)) NIL) (ALL ((POW ((* NIL NIL)) NIL)) NIL)) (SETQ *EXIT* T)))

D-7

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(DECLARE (SPECIAL <HELP>)) (CREATE-PRODUCTION <HELP>) (SETO <HELP> '(<HELP> ((HELP ((* NIL NIL)) NIL) (H ((* NIL NIL)) NIL) (? ((* NIL NIL)) NIL)) (HELP))) (DECLARE (SPECIAL < INSTANCE>)) (CREATE-PRODUCTION <INSTANCE>) (SETQ < INSTANCE> '(<INSTANCE> ((<VARIABLE-INSTANCE> ((* NIL NIL)) NIL) (<REQUIREMENTS-INSTANCE> ((* NIL NIL)) NIL) (<PROJECT-INSTANCE> ((* NIL NIL)) NIL) (<PROCESS-ALIAS-INSTANCE> ((* NIL NIL)) NIL) (<PROCESS-INSTANCE> ((* NIL NIL)) NIL) (<PARAMETER-ALIAS-INSTANCE> ((* NIL NIL)) NIL) (<PARAMETER-INSTANCE> ((* NIL NIL)) NIL) (<MODULE-INSTANCE> ((* NIL NIL)) NIL) (<IMPLEMENTATION-INSTANCE> ((* NIL NIL)) NIL) (<DESIGN-INSTANCE> ((* NIL NIL)) NIL) (<DATA-ELEMENT-ALIAS-INSTANCE> ((* NIL NIL)) NIL) (<DATA-ELEMENT-INSTANCE> ((* NIL NIL)) NIL) (<ACTIVITY-ALIAS-INSTANCE> ((* NIL NIL)) NIL) (<ACTIVITY-INSTANCE> ((* NIL NIL)) NIL)) NIL)) (DECLARE (SPECIAL <ACTIVITY-INSTANCE>)) (CREATE-PRODUCTION <ACTIVITY-INSTANCE>) (SETO <ACTIVITY-INSTANCE> '(<ACTIVITY-INSTANCE> ((ACTIVITY ((<ACTIVITY> ((* NIL NIL)) NIL)) NIL)) NIL)) (DECLARE (SPECIAL <ACTIVITY-ALIAS-INSTANCE>)) (CREATE-PRODUCTION <ACTIVITY-ALIAS-INSTANCE>) (SETO <ACTIVITY-ALIAS-INSTANCE> '(<ACTIVITY-ALIAS-INSTANCE> ((ACTIVITY ((ALIAS ((<ACTIVITY-ALIAS> ((* NIL NIL)) NIL)) NIL)) NIL)) NIL))

(DECLARE (SPECIAL <DATA-ELEMENT-INSTANCE>)) (CREATE-PRODUCTION <DATA-ELEMENT-INSTANCE>) (SETQ <DATA-ELEMENT-INSTANCE> '(<DATA-ELEMENT-INSTANCE> ((DATA ((ELEMENT ((<DATA-ELEMENT> ((* NIL NIL)) NIL)) NIL)) NIL)) NIL)) (DECLARE (SPECIAL <DATA-ELEMENT-ALIAS-INSTANCE>)) (CREATE-PRODUCTION <DATA-ELEMENT-ALIAS-INSTANCE>) (SETQ <DATA-ELEMENT-ALIAS-INSTANCE> '(<DATA-ELEMENT-ALIAS-INSTANCE> ((DATA ((ELEMENT ((ALIAS ((<DATA-ELEMENT-ALIAS> ((* NIL NIL)) NIL)) NIL)) NIL)) NIL)) NIL)) (DECLARE (SPECIAL <DESIGN-INSTANCE>)) (CREATE-PRODUCTION < DESIGN-INSTANCE>) (SETO <DESIGN-INSTANCE> '(<DESIGN-INSTANCE> ((DESIGN ((<DESIGN> ((* NIL NIL)) NIL)) NIL)) NIL)) (DECLARE (SPECIAL < IMPLEMENTATION-INSTANCE>)) (CREATE-PRODUCTION < IMPLEMENTATION-INSTANCE>) (SETO < IMPLEMENTATION-INSTANCE> '(<IMPLEMENTATION-INSTANCE> ((IMPLEMENTATION ((<IMPLEMENTATION> ((* NIL NIL)) NIL)) NIL)) NIL)) (DECLARE (SPECIAL < MODULE-INSTANCE>)) (CREATE-PRODUCTION < MODULE-INSTANCE>) (SETQ <MODULE-INSTANCE> '(<MODULE-INSTANCE> ((MODULE ((<MODULE> ((* NIL NIL)) NIL)) NIL)) NIL)) (DECLARE (SPECIAL < PARAMETER-INSTANCE>)) (CREATE-PRODUCTION < PARAMETER-INSTANCE>) (SETQ < PARAMETER-INSTANCE> '(<PARAMETER-INSTANCE> ((PARAMETER ((<PARAMETER> ((* NIL NIL)) NIL)) NIL))

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```
(DECLARE (SPECIAL < PARAMETER-ALIAS-INSTANCE>))
(CREATE-PRODUCTION < PARAMETER-ALIAS-INSTANCE>)
(SETO < PARAMETER-ALIAS-INSTANCE>
      '(<PARAMETER-ALIAS-INSTANCE>
      ( ( PARAMETER
         ((ALIAS ((<PARAMETER-ALIAS> ((* NIL NIL)) NIL))
NIL)) NIL))
      NIL))
(DECLARE (SPECIAL < PROCESS-INSTANCE>))
(CREATE-PRODUCTION < PROCESS-INSTANCE>)
(SETQ < PROCESS-INSTANCE>
      '(<PROCESS-INSTANCE>
      ((PROCESS ((<PROCESS> ((* NIL NIL)) NIL)) NIL)) NIL))
(DECLARE (SPECIAL < PROCESS-ALIAS-INSTANCE>))
(CREATE-PRODUCTION < PROCESS-ALIAS-INSTANCE>)
(SETO
  <PROCESS-ALIAS-INSTANCE>
  '(<PROCESS-ALIAS-INSTANCE>
     ((PROCESS ((ALIAS ((<PROCESS-ALIAS> ((* NIL NIL)) NIL))
NIL)) NIL))
     NIL))
(DECLARE (SPECIAL < PROJECT-INSTANCE>))
(CREATE-PRODUCTION < PROJECT-INSTANCE>)
(SETO
  <PROJECT-INSTANCE>
  '(<PROJECT-INSTANCE>
     ((*PROJECT* ((* NIL (SEND *NEW-EVENT* :SET-OBJECT
*PROJECT*))) NIL)
      (PROJECT ((<PROJECT> ((* NIL NIL)) NIL)) NIL))
     NIL))
(DECLARE (SPECIAL < REQUIREMENTS-INSTANCE>))
(CREATE-PRODUCTION < REQUIREMENTS-INSTANCE>)
(SETQ < REQUIREMENTS-INSTANCE>
      ' (< REQUIREMENTS-INSTANCE>
      ((REQUIREMENTS ((<REQUIREMENTS> ((* NIL NIL)) NIL))
NIL)) NIL))
(DECLARE (SPECIAL <VARIABLE-INSTANCE>))
(CREATE-PRODUCTION <VARIABLE-INSTANCE>)
(SETQ <VARIABLE-INSTANCE>
      '(<VARIABLE-INSTANCE>
      ((VARIABLE ((<VARIABLE> ((* NIL NIL)) NIL)) NIL))
NIL))
```

```
(DECLARE (SPECIAL <ACTIVITY>))
(CREATE-FUNCTION <ACTIVITY>)
(SETQ <ACTIVITY> '(<ACTIVITY> (LAMBDA (X) NIL)))
(DECLARE (SPECIAL <ACTIVITY-ALIAS>))
(CREATE-FUNCTION <ACTIVITY-ALIAS>)
(SETQ <ACTIVITY-ALIAS> '(<ACTIVITY-ALIAS> (LAMBDA (X) NIL)))
(DECLARE (SPECIAL <DATA-ELEMENT>))
(CREATE-FUNCTION < DATA-ELEMENT>)
(SETQ <DATA-ELEMENT> '(<DATA-ELEMENT> (LAMBDA (X) NIL)))
(DECLARE (SPECIAL <DATA-ELEMENT-ALIAS>))
(CREATE-FUNCTION <DATA-ELEMENT-ALIAS>)
(SETQ <DATA-ELEMENT-ALIAS> '(<DATA-ELEMENT-ALIAS> (LAMBDA
(X) NIL)))
(DECLARE (SPECIAL <DESIGN>))
(CREATE-FUNCTION <DESIGN>)
(SETO < DESIGN>
      '(<DESIGN>
      (LAMBDA (X)
        (LET ((INSTANCE (SEND *PROJECT* : IS-DESIGN X)))
          (COND (INSTANCE (SEND *NEW-EVENT* :SET-OBJECT
INSTANCE)))))))
(DECLARE (SPECIAL < IMPLEMENTATION>))
(CREATE-FUNCTION < IMPLEMENTATION>)
SETO < IMPLEMENTATION>
      '(<IMPLEMENTATION>
      (LAMBDA (X)
        (LET ((INSTANCE (SEND *PROJECT* : IS-IMPLEMENTATION
X)))
          (COND (INSTANCE (SEND *NEW-EVENT* :SET-OBJECT
INSTANCE)))))))
(DECLARE (SPECIAL < MODULE>))
(CREATE-FUNCTION < MODULE>)
```

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(SETQ < MODULE> '(< MODULE> (LAMBDA (X) NIL)))

```
(DECLARE (SPECIAL < PARAMETER>))
(CREATE-FUNCTION < PARAMETER>)
(SETO
 <PARAMETER>
  '(<PARAMETER>
     (LAMBDA (X)
       (LET ((INSTANCE
            (SEND (SEND *PROJECT* :DESIGN-INSTANCE)
:IS-PARAMETER X)))
      (COND (INSTANCE (SEND *NEW-EVENT* :SET-OBJECT
INSTANCE)))))))
(DECLARE (SPECIAL < PARAMETER-ALIAS>))
(CREATE-FUNCTION < PARAMETER-ALIAS>)
(SETO
 PARAMETER-ALIAS>
  '(<PARAMETER-ALIAS>
     (LAMBDA (X)
       (LET ((INSTANCE
            (SEND
           (SEND *PROJECT* :DESIGN-INSTANCE)
:IS-PARAMETER-ALIAS X)))
      (COND (INSTANCE (SEND *NEW-EVENT* :SET-OBJECT
INSTANCE)))))))
(DECLARE (SPECIAL < PROCESS>))
(CREATE-FUNCTION < PROCESS>)
(SETQ <PROCESS>
      '(<PROCESS>
      (LAMBDA (X)
        (LET ((INSTANCE
             (SEND (SEND *PROJECT* :DESIGN-INSTANCE)
:IS-PROCESS X)))
          (COND (INSTANCE (SEND *NEW-EVENT* :SET-OBJECT
INSTANCE)))))))
(DECLARE (SPECIAL < PROCESS-ALIAS>))
(CREATE-FUNCTION < PROCESS-ALIAS>)
(SETO
 <PROCESS-ALIAS>
  '(<PROCESS-ALIAS>
     (LAMBDA (X))
       (LET
      ((INSTANCE
         (SEND (SEND *PROJECT* : DESIGN-INSTANCE)
:IS-PROCESS-ALIAS X)))
      (COND (INSTANCE (SEND *NEW-EVENT* :SET-OBJECT
INSTANCE)))))))
```

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(DECLARE (SPECIAL < PROJECT>)) (CREATE-FUNCTION < PROJECT>) (SETQ <PROJECT> '(<PROJECT> (LAMBDA (X) (COND ((EQUAL X (SEND *PROJECT* :NAME)) (SEND *NEW-EVENT* :SET-OBJECT *PROJECT*)))))) (DECLARE (SPECIAL <VARIABLE>)) (CREATE-FUNCTION <VARIABLE>) (SETQ <VARIABLE> '(<VARIABLE> (LAMBDA (X) NIL))) (DECLARE (SPECIAL <ATTRIBUTE-LIST>)) (CREATE-PRODUCTION <ATTRIBUTE-LIST>) (SETQ <ATTRIBUTE-LIST> '(<ATTRIBUTE-LIST> ((<ATTRIBUTE> ((AND ((<ATTRIBUTE> ((* NIL NIL)) NIL)) NIL) (<ATTRIBUTE-LIST> ((* NIL NIL)) NIL) (* NIL NIL)) NIL))

NIL))

.

(DECLARE (SPECIAL <ATTRIBUTE>)) (CREATE-PRODUCTION <ATTRIBUTE>) (SETQ <ATTRIBUTE> '(<ATTRIBUTE> ((GLOBAL ((DATA ((CHANGED ((* NIL (SEND *NEW-EVENT* :ADD-ATTRIBUTE 'GLOBAL-DATA-CHANGED))) NIL) (USED ((* NIL (SEND *NEW-EVENT* :ADD-ATTRIBUTE 'GLOBAL-DATA-USED))) NIL)) NIL)) NIL) (DESTINATIONS ((* NIL (SEND *NEW~EVENT* :ADD-ATTRIBUTE 'PASSED-TO))) NIL) (SOURCES ((* NIL (SEND *NEW-EVENT* :ADD-ATTRIBUTE 'PASSED-FROM))) NIL) (SYNONYMS ((* NIL (SEND *NEW-EVENT* :ADD-ATTRIBUTE 'SYNONYM))) NIL) (DESTINATION ((* NIL (SEND *NEW-EVENT* :ADD-ATTRIBUTE 'PASSED-TO))) NIL) (SOURCE ((* NIL (SEND *NEW-EVENT* :ADD-ATTRIBUTE 'PASSED-FROM))) NIL) (PASSED ((TO ((* NIL (SEND *NEW~EVENT* :ADD-ATTRIBUTE 'PASSED-TO))) NIL) (FROM ((* NIL (SEND *NEW-EVENT* :ADD-ATTRIBUTE 'PASSED-FROM))) NIL)) NIL) (OUTPUT ((FLAGS ((* NIL (SEND *NEW-EVENT* :ADD-ATTRIBUTE 'OUTPUT-FLAGS))) NIL) (DATA ((* NIL (SEND *NEW-EVENT* :ADD-ATTRIBUTE 'OUTPUT-DATA))) NIL)) NIL)

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```
(INPUT
       ((FLAGS
       ((* NIL (SEND *NEW-EVENT* :ADD-ATTRIBUTE
'INPUT-FLAGS))) NIL)
     (DATA ((* NIL (SEND *NEW-EVENT* :ADD-ATTRIBUTE
'INPUT-DATA))) NIL))
       NIL)
     (ALIASES ((* NIL (SEND *NEW-EVENT* :ADD-ATTRIBUTE
'ALIASES))) NIL)
     (GLOBALS
       ((USED ((* NIL (SEND *NEW-EVENT* :ADD-ATTRIBUTE
'GLOBAL-DATA-USED)))
           NIL)
   • (CHANGED
       ((* NIL (SEND *NEW-EVENT* :ADD-ATTRIBUTE
'GLOBAL-DATA-CHANGED)))
       NIL))
       NIL)
     (INVOKING
       ((PROCESSES
       ((* NIL (SEND *NEW-EVENT* :ADD-ATTRIBUTE
'CALLING-PROCESSES)))
       NIL))
       NIL)
     (ENTRY
       ((DATE ((* NIL (SEND *NEW-EVENT* :ADD-ATTRIBUTE
'DATE))) NIL)
     (VERSION ((* NIL (SEND *NEW-EVENT* :ADD-ATTRIBUTE
'VERSION))) NIL))
       NIL)
     (SYNONYM ((* NIL (SEND *NEW-EVENT* :ADD-ATTRIBUTE
'SYNONYM))) NIL)
     (REQUIREMENTS
       ((* NIL (SEND *NEW-EVENT* :ADD-ATTRIBUTE
'REQUIREMENTS))) NIL)
     (SUBPROCESSES
       ((* NIL (SEND *NEW-EVENT* :ADD-ATTRIBUTE
'PROCESSES-CALLED))) NIL)
     (INVOKED
       ( ( PROCESSES
       ((* NIL (SEND *NEW-EVENT* :ADD-ATTRIBUTE
'PROCESSES-CALLED)))
       NIL))
       NIL)
```

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(CALLED ((PROCESSES ((* NIL (SEND *NEW-EVENT* :ADD-ATTRIBUTE 'PROCESSES-CALLED))) NIL)) NIL) (PROCESSES ((INVOKED ((* NIL (SEND *NEW-EVENT* :ADD-ATTRIBUTE 'PROCESSES-CALLED))) NIL) (CALLED ((* NIL (SEND *NEW-EVENT* :ADD-ATTRIBUTE 'PROCESSES-CALLED))) NIL) (* NIL (SEND *NEW-EVENT* :ADD-ATTRIBUTE 'PROCESSES))) NIL) (PROCESS ((ALIASES ((* NIL (SEND *NEW-EVENT* :ADD-ATTRIBUTE 'PROCESS-ALIASES))) NIL)) NIL) (PARAMETERS ((* NIL (SEND *NEW-EVENT* :ADD-ATTRIBUTE 'PARAMETERS))) NIL) (PARAMETER ((ALIASES ((* NIL (SEND *NEW-EVENT* :ADD-ATTRIBUTE 'PARAMETER-ALIASES))) NIL)) NIL) (MAIN ((PROGRAM ((* NIL (SEND *NEW-EVENT* :ADD-ATTRIBUTE 'MAIN-PROGRAM))) NIL) (PROGRAMS ((* NIL (SEND *NEW-EVENT* :ADD-ATTRIBUTE 'MAIN-PROGRAM))) NIL) (PROCESS ((* NIL (SEND *NEW-EVENT* :ADD-ATTRIBUTE 'MAIN-PROGRAM))) NIL) (PROCESSES ((* NIL (SEND *NEW-EVENT* :ADD-ATTRIBUTE 'MAIN-PROGRAM))) NIL)) NIL) (IMPLEMENTATION ((* NIL (SEND *NEW-EVENT* :ADD-ATTRIBUTE 'IMPLEMENTATION))) NIL)

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•

```
(DESIGN ((* NIL (SEND *NEW-EVENT* :ADD-ATTRIBUTE
'DESIGN))) NIL)
     (DD ((TYPE ((* NIL (SEND *NEW-EVENT* :ADD-ATTRIBUTE
'DD-TYPE))) NIL))
     NIL)
     (CALLING
       ( (PROCESSES
       ((* NIL (SEND *NEW-EVENT* :ADD-ATTRIBUTE
'CALLING-PROCESSES)))
       NIL))
       NIL)
     (VERSION ((* NIL (SEND *NEW-EVENT* :ADD-ATTRIBUTE
'VERSION))) NIL)
     (TYPE ((* NIL (SEND *NEW-EVENT* :ADD-ATTRIBUTE 'TYPE)))
NIL)
     (PROJECT
       ((NAME ((* NIL (SEND *NEW-EVENT* :ADD-ATTRIBUTE
'PROJECT))) NIL)
     (* NIL (SEND *NEW-EVENT* :ADD-ATTRIBUTE 'PROJECT)))
       NIL)
     (NAME ((* NIL (SEND *NEW-EVENT* :ADD-ATTRIBUTE 'NAME)))
NIL)
     (DATE ((* NIL (SEND *NEW-EVENT* :ADD-ATTRIBUTE 'DATE)))
NIL))
    NIL))
(DECLARE (SPECIAL <LITERAL>))
(CREATE-FUNCTION <LITERAL>)
(SETQ <LITERAL> '(<LITERAL> (LAMBDA (X) NIL)))
(DECLARE (SPECIAL <SAVE>))
(CREATE-PRODUCTION < SAVE>)
(SETQ <SAVE> '(<SAVE> ((STORE ((* NIL NIL)) NIL) (KEEP ((*
NIL NIL) > NIL)
                 (SAVE ((* NIL NIL)) NIL))
                 (SEND *PROJECT* :SAVE)))
```

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Appendix E

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DDT Object Definitions

Appendix E Table of Contents

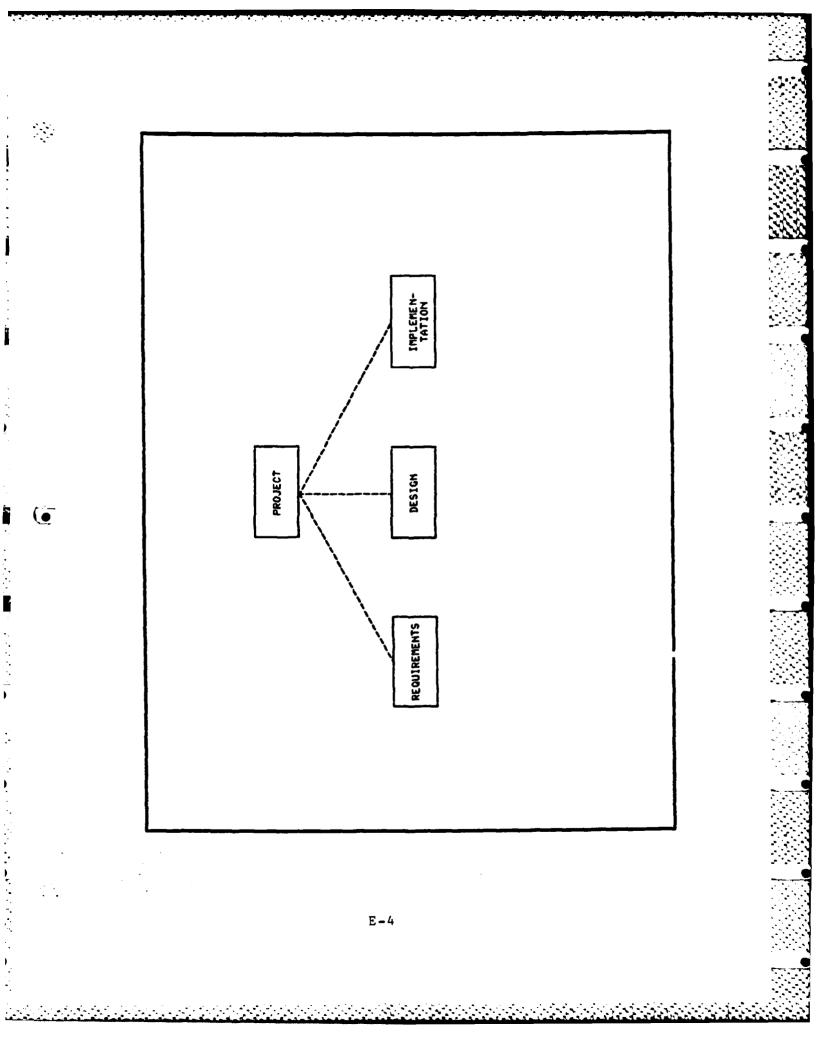
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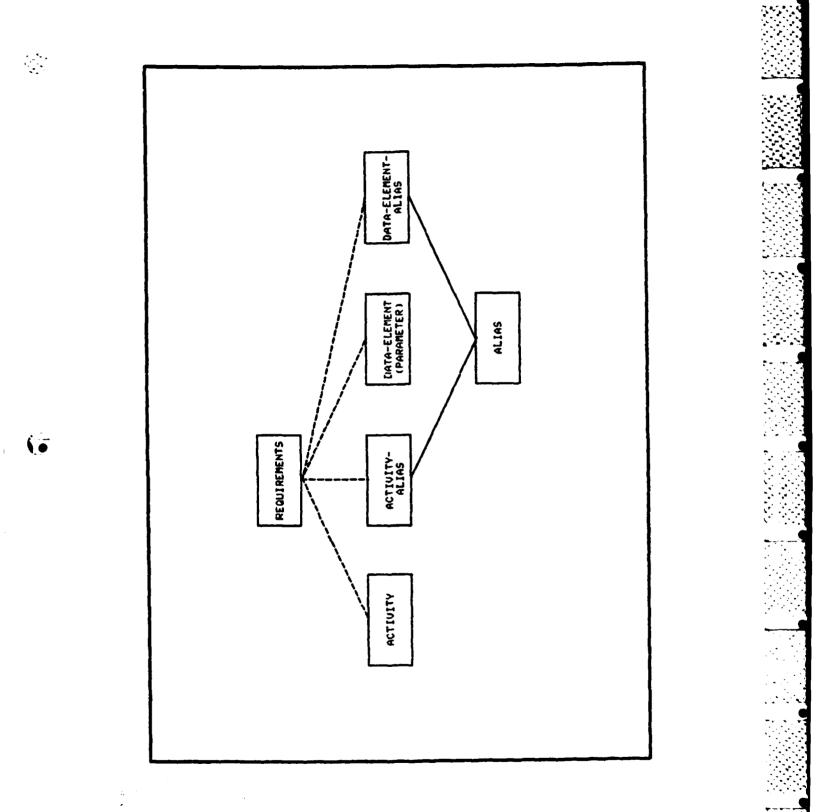
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Intr	oductio	n	•	•••	•	•	•	•	•	•	•	•	٠	•	•	•	•	•	•	•	•	E-	3
DDT	Object	Hie	ra	rch	Y	Cha	art	ts	•	•	٠	•	•	•	•	•	٠	٠	•	•	•	E-	4
DDT	Object	Def	in	iti	on	s	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•	E-	8

Introduction

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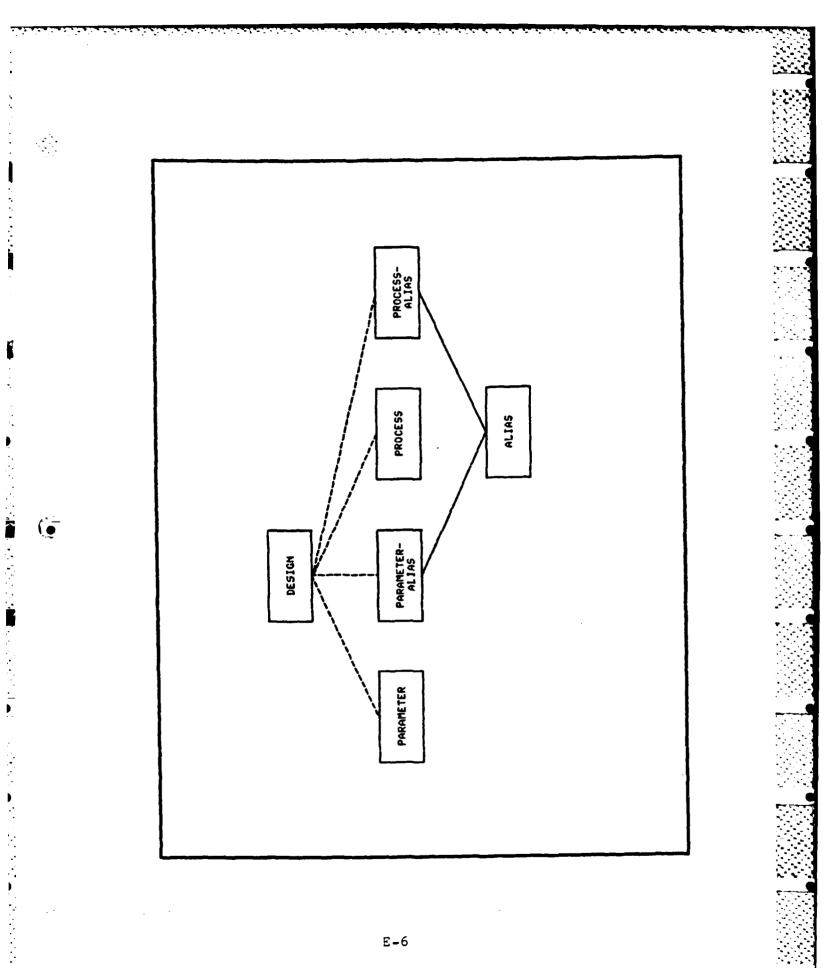
This appendix defines the object-oriented structure of the Data Dictionary Tool (DDT). The first part of this definition shows the hierarchy of DDT's object classes. The second part defines the structure of each object class in terms of its included subclasses, its instance variables, and its associated access methods.





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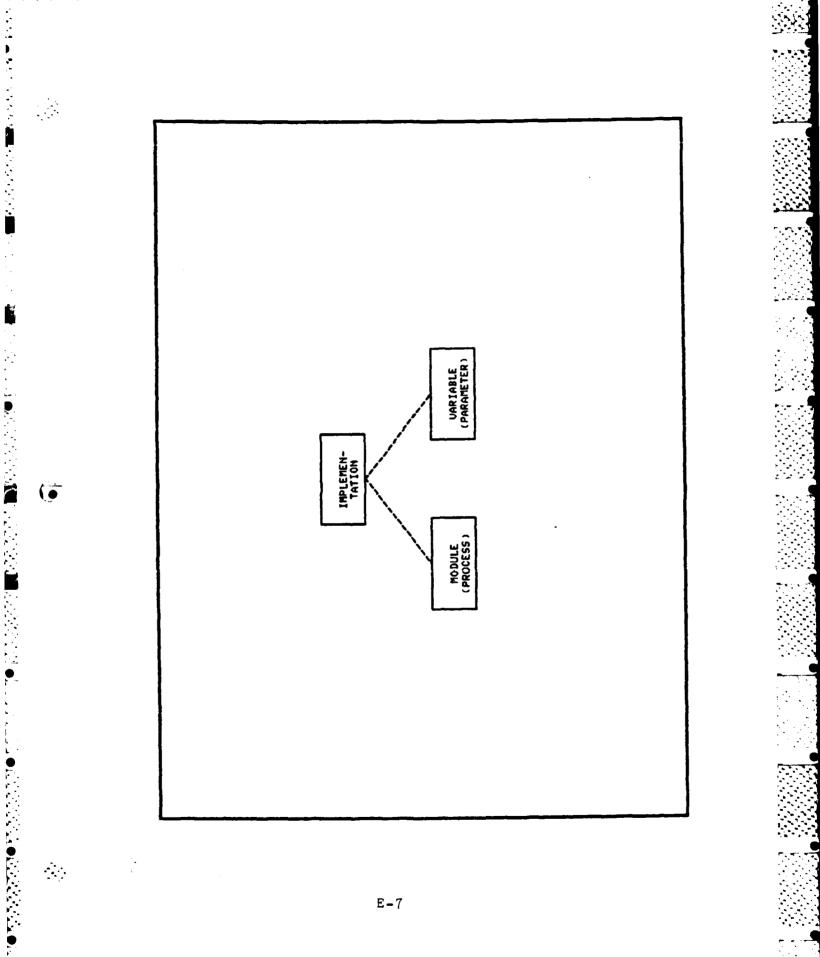
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Object Class: ACTIVITY

Subclasses HEADER ENTRY

Instance Variables NUMBER DESCRIPTION INPUTS

OUTPUTS

CONTROLS

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MECHANISMS

ALIASES PARENT-ACTIVITY RELATED-ITEMS Node number of activity Textual description of activity Assoc-list of input parameter names and numbers Assoc-list of output parameter names and numbers Assoc-list of control parameter names and numbers Assoc-list of mechanism parameter and numbers List of activity alias names Name of parent activity Paragraph number of textual requirements statement Subclasses ALIAS

Instance Variables WHERE-USED

List of SADT activity numbers

Object Class: ALIAS

Subclasses HEADER ENTRY

Instance Variables DD-TYPE

Data dictionary type ("ACTIVITY", "DATA ELEMENT", "PARAMETER", "PROCESS") Name of a defined module or parameter

Methods

SYNONYM

それましたが、この時間になるためないの時間をついていた。 読むたちの いたい

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dd-type		Retur	ns t	the val	lue	of	DD-TY	'PE	
set-dd-type	(VAL)	Sets	the	value	of	DD	-TYPE	to	VAL

synonym	Returns	the the	value	of SYNON	IYM	
set-synonym (VAL)	Sets th	e val	ue of	SYNONYM	to	VAL



Object Class: ASSOC

Subclasses none

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Instance Variable VALUE

<u>Methods</u> initialize

present

value set-value (VAL) Initializes all of ASSOC's instance variables Displays an ASSOC instance

Returns the value of VALUE Sets VALUE to VAL Object Class: ATOM

Subclasses none

Instance Variables VALUE

<u>Methods</u> initialize

present

value set-value (VAL) Initializes all of ATOM's instance variables Displays an ATOM instance

Returns the value of VALUE Sets VALUE to VAL





Subclasses PARAMETER

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Instance Variables none







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Object Class: DATA-ELEMENT-ALIAS

Subclasses ALIAS

Instance Variables WHERE-USED

List of SADT activity numbers

Object class: DATE

Subclasses none

Instance Variables VALUE

<u>Methods</u> initialize

present

Initializes all of DATE's instance variables Displays a DATE instance

value set-value (VAL)	Returns the value of VALUE Sets VALUE to VAL
day	Returns the value of the DAY subfield of VALUE
month	Returns the value of the MONTH subfield of VALUE
year	Returns the value of the YEAR subfield of VALUE

Object Class: DESIGN

Subclasses HEADER ENTRY

Instance Variables MAIN-PROGRAM PROCESSES

PROCESS-ALIASES

PARAMETERS

PARAMETER-ALIASES

<u>Methods</u> create initialize

present (ATTRIBUTE-LIST)

reinitialize (DUMMY)

save (PROJECT, STREAM)

initialize-date

initialize-mainprogram main-program set-main-program (VAL)

initialize-name

Name of main program Assoc-list of module names and pointers to module instances Assoc-list of process alias names and pointers to alias instances Assoc-list of parameter names and pointers to parameter instances Assoc-list of parameter alias names and pointers to alias instances

Creates a DESIGN instance Initializes all of DESIGN's instance variables

Displays the instance variables listed in ATTRIBUTE-LIST of a DESIGN instance; if ATTRIBUTE-LIST is empty, displays all instance variables Initializes all of DESIGN's instance variables except NAME, TYPE, and PROJECT

Saves a DESIGN instance and all of its sub-instances to STREAM

Initializes DATE

Initializes MAIN-PROGRAM Returns the value of MAIN-PROGRAM

Sets the value of MAIN-PROGRAM to VAL

Initializes NAME

create-parameteralias

initialize-parameteraliases parameter-aliases

parameter-aliasinstance-list

set-parameter-aliases
(VAL)

is-parameter-alias (VAL)

create-parameter

initialize-parameters parameters parameter-instancelist

set-parameters (VAL)

is-parameter (VAL)

Creates a PARAMETER-ALIAS instance and adds it to the list of PARAMETER-ALIASES

Initializes PARAMETER-ALIASES Returns the value of PARAMETER-ALIASES

Returns a list of pointers to the PARAMETER-ALIASes of a DESIGN

Sets the value of PARAMETER-ALIASES to VAL

Returns a pointer to the appropriate instance if VAL is a PARAMETER-ALIAS, otherwise returns NIL

Creates a PARAMETER instance and adds it to the list of PARAMETERS Initializes PARAMETERS Returns the value of PARAMETERS

Returns a list of pointers to the PARAMETERS of a DESIGN Sets the value of PARAMETERS to VAL Returns a pointer to the appropriate instance if VAL is a PARAMETER, otherwise returns NIL create-process-alias

initialize-processaliases process-aliases

process-aliasinstance-list

set-process-aliases
 (VAL)

is-process-alias (VAL) Creates a PROCESS-ALIAS instance and adds it to the list of PROCESS-ALIASES

Initializes PROCESS-ALIASES Returns the value of PROCESS-ALIASES

Returns a list of pointers to the PROCESS ALIASes of a DESIGN

Sets the value of PROCESS-ALIASES to VAL

Returns a pointer to the appropriate instance if VAL is a PROCESS-ALIAS, otherwise returns NIL

Creates a PROCESS instance and

Returns the value of PROCESSES

Sets the value of PROCESSES to

Initializes PROCESSES

PROCESSes of a DESIGN

Initializes VERSION

adds it to the list of PROCESSES

Returns a list of pointers to the

create-process

initialize-processes processes process-instance-list

set-processes (VAL)

is-process (VAL)

VAL Returns a pointer to the appropriate instance if VAL is a PROCESS, otherwise returns NIL

initialize-project Initializes PROJECT

initializ -type Initializes TYPE

initialize-version

Object Class: ENTRY

Subclasses none

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Instance Variables VERSION DATE

Version number of entry Date of entry

Methods date

version increment-version Returns the value of DATE

Returns the value of VERSION Increments the value of VERSION by 1

Object Class: HEADER

Subclasses none

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Instance Variables NAME TYPE

Name of instance Type of instance ("ACTIVITY", "ALIAS", "DATA ELEMENT", "DESIGN", "IMPLEMENTATION", "MODULE", "PARAMETER", "PROCESS", "REQUIREMENTS", "VARIABLE") Name of project

PROJECT

Methods name set-name (VAL)	Returns the value of NAME Sets the value of NAME to VAL
project	Returns the value of PROJECT
type	Returns the value of TYPE













Object Class: IMPLEMENTATION

Subclasses HEADER ENTRY

Instance Variables MAIN-PROGRAM MODULES

VARIABLES

Methods create

initialize

Creates an IMPLEMENTATION instance Initializes all of IMPLEMENTATION's instance variables

all instance variables

Name of main program

Assoc-list of module names and

pointers to variable instances

pointers to module instances Assoc-list of variable names and

present (ATTRIBUTE-LIST)

reinitialize (DUMMY)

save (PROJECT, STREAM)

Initializes all of IMPLEMENTATION'S instance variables except NAME, TYPE, and PROJECT Saves an IMPLEMENTATION instance

Displays the instance variables

ATTRIBUTE-LIST is empty, displays

listed in ATTRIBUTE-LIST of an IMPLEMENTATION instance; if

Saves an IMPLEMENTATION instance and all of its sub-instances to STREAM

Object Class: LIST

Subclasses none

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Instance Variables VALUE

<u>Methods</u> initialize

Initializes all of LIST's instance variables present Displays a LIST instance

value set-value (VAL) Returns the value of VALUE Sets VALUE to VAL

Subclasses PROCESS

Instance Variables none

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Object Class: PARAMETER

Subclasses HEADER ENTRY

Instance Variables DESCRIPTION Textual description of parameter/variable/data element DATA-TYPE Data type of parameter/variable/data element Minimum value of MIN-VALUE parameter/variable/data element (if applicable) Maximum value of MAX-VALUE parameter/variable/data element (if applicable) Range of parameter/variable/data RANGE element (if applicable) List of legal values of VALUES parameter/variable/data element (if applicable) STORAGE-TYPE Storage type of parameter/variable/data element ("FILE", "GLOBAL", "HARDWARE", "I/O", or "PASSED") Name of parent PART-OF parameter/variable/data element (if applicable) COMPOSITION List of sub-parameter/variable/data element names (if applicable) List of parameter/variable/data ALIASES element alias names (if any) PASSED-FROM List of process/module/activity names passed from List of process/module/activity PASSED-TO names passed to List of related RELATED-ITEMS SADT-data-item/parameter/requirement names/numbers

Methods create initialize

present (ATTRIBUTE-LIST)

reinitialize (DUMMY)

save (DESIGN, STREAM)

initialize-aliases aliases set-aliases (VAL)

initializecomposition composition set-composition (VAL)

initialize-data-type
data-type
set-data-type (VAL)

initialize-date

initializedescription description set-description (VAL)

initialize-max-value max-value set-max-value (VAL)

initialize-min-value min-value set-min-value (VAL)

Creates a PARAMETER instance Initializes all of PARAMETER's instance variables

Displays the instance variables listed in ATTRIBUTE-LIST of a PARAMETER instance; if ATTRIBUTE-LIST is empty, displays all instance variables Initializes all of PARAMETER's instance variables except NAME, TYPE, and PROJECT Saves a PARAMETER instance to STREAM

Initializes ALIASES Returns the value of ALIASES Sets the value of ALIASES to VAL

Initializes COMPOSITION Returns the value of COMPOSITION Sets the value of COMPOSITION to VAL

Initializes DATA-TYPE Returns the value of DATA-TYPE Sets the value of DATA-TYPE to VAL

Initializes DATE

Initializes DESCRIPTION Returns the value of DESCRIPTION Sets the value of DESCRIPTION to VAL

Initializes MAX-VALUE Returns the value of MAX-VALUE Sets the value of MAX-VALUE to VAL

Initializes MIN-VALUE Returns the value of MIN-VALUE Sets the value of MIN-VALUE to VAL

initialize-name

initialize-part-of
part-of
set-part-of (VAL)

initialize-passedfrom
passed-from
set~passed-from (VAL)

initialize-passed-to passed-to set-passed-to (VAL)

initialize-project

initialize-range
range
set-range (VAL)

initialize-related items
related-items

set-related-items
(VAL)

initialize-storagetype storage-type set-storage-type (VAL)

initialize-type

initialize-values values set-values (VAL)

initialize-version

Initializes NAME

Initializes PART-OF Returns the value of PART-OF Sets the value of PART-OF to VAL

Initializes PASSED-FROM Returns the value of PASSED-FROM Sets the value of PASSED-FROM to VAL

Initializes PASSED-TO Returns the value of PASSED-To Sets the value of PASSED-TO to VAL

Initializes PROJECT

Initializes RANGE Returns the value of RANGE Sets the value of RANGE to VAL

Initializes RELATED-ITEMS Returns the value of RELATED-ITEMS

Sets the value of RELATED-ITEMS to VAL

Initializes STORAGE-TYPE Returns the value of STORAGE-TYPE

Sets the value of STORAGE-TYPE to VAL

Initializes TYPE

Initializes VALUES Returns the value of VALUES Sets the value of VALUES to VAL

Initializes VERSION

Object Class: PARAMETER-ALIAS Subclasses ALIAS Instance Variables PASSED-FROM List of process names parameter is passed from PASSED-TO List of process names parameter is passed to Methods create Creates a PARAMETER-ALIAS instance initialize Initializes all of PARAMETER-ALIAS' instance variables present (ATTRIBUTE-LIST) Displays the instance variables listed in ATTRIBUTE-LIST of a PARAMETER-ALIAS instance; if ATTRIBUTE-LIST is empty, displays all instance variables Initializes all of reinitialize (DUMMY) PARAMETER-ALIAS' instance variables except NAME, TYPE, PROJECT, and DD-TYPE save (DESIGN, STREAM) Saves a PARAMETER-ALIAS instance to STREAM initialize-date Initializes DATE initialize-dd-type Initializes DD-TYPE Initializes NAME initialize-name initialize-passedfrom Initializes PASSED-FROM Returns the value of PASSED-FROM passed-from Sets the value of PASSED-FROM to set-passed-from (VAL) VAL Initializes PASSED-TO initialize-passed-to Returns the value of PASSED-TO passed-to Sets the values of PASSED-TO to set-passed-to (VAL) VAL initialize-project Initializes PROJECT

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initialize-synonym Initializes SYNONYM initialize-type Initializes TYPE initialize-version Initializes VERSION Object Class: PROCESS

Subclasses HEADER ENTRY

Instance Variables NUMBER DESCRIPTION

INPUT-DATA

INPUT-FLAG

OUTPUT-DATA

OUTPUT-FLAGS

GLOBAL-DATA-USED GLOBAL-DATA-CHANGED FILES-READ FILES-WRITTEN

HARDWARE-INPUT HARDWARE-OUTPUT ALIASES

CALLING-PROCESSES

PROCESSES-CALLED

RELATED-ITEMS

Number of process/module Textual description of process/module List of input data parameter/variable names List of input flag parameter/variables names List of output data parameter/variable names List of output flag parameter/variables names List of globals used List of globals used List of filenames of files read List of filenames of files written

List of process/module alias names List of calling process/module names List of processes/modules called names List of related SADT-activity-names/process-names Methods create initialize

present (ATTRIBUTE-LIST)

reinitialize (DUMMY)

save (DESIGN, STREAM)

initialize-aliases aliases set-aliases (VAL)

initialize-callingprocesses calling-processes

set-calling-processes (VAL)

add-calling-processes (LIST)

delete-callingprocesses (LIST)

initialize-date

Initializes DATE

Deletes NAME from PROCESSES-CALLED slot

initialize-description description set-description (VAL)

Initializes DESCRIPTION Returns the value of DESCRIPTION Sets the value of DESCRIPTION to VAL

Creates a PROCESS instance

instance variables

PROCESS instance; if

TYPE, and PROJECT

Initializes ALIASES

Returns the value of

CALLING-PROCESSES to VAL

necessary, Creates PROCESS instance; Adds NAME to PROCESSES-CALLED slot

CALLING-PROCESSES

Sets the value of

STREAM

all instance variables

Initializes all of PROCESS'

Initializes all of PROCESS'

Saves a PROCESS instance to

Returns the value of ALIASES Sets the value of ALIASES to VAL

Initializes CALLING-PROCESSES

For each process name in LIST: Adds to CALLING-PROCESSES;

For each process name in LIST:

Deletes from CALLING-PROCESSES;

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Displays the instance variables listed in ATTRIBUTE-LIST of a

instance variables except NAME,

ATTRIBUTE-LIST is empty, displays

initialize-files-read files-read set-files-read (VAL)

initialize-fileswritten files-written

set-files-written
 (VAL)

initialize-globalschanged globals-changed

set-globals-changed
 (VAL)

initialize-globalsused globals-used

set-globals-used (VAL)

initialize-hardwareinput hardware-input

set-hardware-input
(VAL)

initialize-hardwareoutput hardware-output

set-hardware-output
 (VAL)

initialize-input-data input-data set-input-data (VAL) Initializes FILES-READ Returns the value of FILES-READ Sets the value of FILES-READ to VAL

Initializes FILES-WRITTEN Returns the value of FILES-WRITTEN

Sets the value of FILES-WRITTEN to VAL

Initializes GLOBAL-DATA-CHANGED Returns the value of GLOBAL-DATA-CHANGED

Sets the value of GLOBAL-DATA-CHANGED to VAL

Initializes GLOBAL-DATA-USED Returns the value of GLOBAL-DATA-USED Sets the value of GLOBAL-DATA-USED to VAL

Initializes HARDWARE-INPUT Returns the value of HARDWARE-INPUT

Sets the value of HARDWARE-INPUT to VAL

Initializes HARDWARE-OUTPUT Returns the value of HARDWARE-OUTPUT

Sets the value of HARDWARE-OUTPUT to VAL

Initializes INPUT-DATA Returns the value of INPUT-DATA Sets the value of INPUT-DATA to VAL

initialize-input-flags Initializes INPUT-FLAGS Returns the value of INPUT-FLAGS input-flags Sets the value of INPUT-FLAGS to set-input-flags (VAL) VAL Initializes NAME initialize-name initialize-number Initializes NUMBER number Returns the value of NUMBER Sets the value of NUMBER to VAL set-number (VAL) initialize-output-Initializes OUTPUT-DATA data Returns the value of OUTPUT-DATA output-data Sets the value of OUTPUT-DATA to set-output-data (VAL) VAL initialize-outputflags Initializes OUTPUT-FLAGS Returns the value of OUTPUT-FLAGS output-flags set-output-flags (VAL) Sets the value of OUTPUT-FLAGS to VAL initialize-processescalled Initializes PROCESSES-CALLED processes-called Returns the value of PROCESSES-CALLED set-processes-called (VAL) Sets the value of PROCESSES-CALLED to VAL add-processes-called (LIST) For each process name in LIST: Adds to PROCESSES-CALLED; If necessary, Creates PROCESS instance; Adds NAME to CALLING-PROCESSES slot delete-processescalled (LIST) For each process name in LIST: Deletes from PROCESSES-CALLED; Deletes NAME from CALLING-PROCESSES slot initialize-project Initializes PROJECT

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initialize-related-	Initializes RELATED-ITEMS
items	Returns the value of
related-items	RELATED-ITEMS
set-related-items	Sets the value of RELATED-ITEMS
(VAL)	to VAL
initialize-type	Initializes TYPE
initialize-version	Initializes VERSION

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Object Class: PROCESS-ALIAS

Subclasses ALIAS

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Instance Variables none

Methods create initialize

present (ATTRIBUTE-LIST)

reinitialize (DUMMY)

save (DESIGN, STREAM)

initialize-date

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initialize-dd-type initialize-name

initialize-project

initialize-synonym

initialize-type

initialize-version

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Creates a PROCESS-ALIAS instance

Displays the instance variables listed in ATTRIBUTE-LIST of a PROCESS-ALIAS instance; if

ATTRIBUTE-LIST is empty, displays

Initializes all of PROCESS-ALIAS' instance variables except NAME, TYPE, PROJECT, and DD-TYPE

Saves a PROCESS-ALIAS instance to

instance variables

all instance variables

Initializes DATE

Initializes NAME

Initializes TYPE

Initializes DD-TYPE

Initializes PROJECT

Initializes SYNONYM

Initializes VERSION

STREAM

Initializes all of PROCESS-ALIAS'

Object Class: PROJECT

Subclasses ENTRY

Instance Variables NAME REQUIREMENTS

DESIGN

IMPLEMENTATION

Methods

create initialize

reinitialize (DUMMY)

present (ATTRIBUTE-LIST)

save

create-design (NAME)

initialize-design design design-instance

set-design (VAL)
is-design (VAL)

Name of project Assoc-list of requirements names and pointers to requirements instance(s) Assoc-list of design names and pointers to design instance(s) Assoc-list of implementation names and pointers to implementation instance(s)

Creates a PROJECT instance Initializes all of PROJECT's instance variables Initializes all of PROJECT's instance variables except NAME

Displays the instance variables listed in ATTRIBUTE-LIST of a PROJECT instance; if ATTRIBUTE-LIST is empty, displays all instance variables Saves a PROJECT instance and all of its sub-instances to STREAM

 Creates a DESIGN instance whose name is NAME and adds it to DESIGN
 Initializes DESIGN
 Returns the value of DESIGN
 Returns a pointer to the DESIGN
 instance
 Sets the value of DESIGN to VAL
 Returns a pointer to the
 appropriate instance if VAL is a
 DESIGN, otherwise returns NIL

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create-implementation (NAME)

initializeimplementation implementation

implementationinstance

set-implementation
 (VAL)

is-implementation
 (VAL)

initialize-name name set-name (VAL)

create-requirements (NAME)

initializerequirements requirements requirements-instance

set-requirements
(VAL)

is-requirements (VAL) Creates an IMPLEMENTATION instance whose name is NAME and adds it to IMPLEMENTATION

Initializes IMPLEMENTATION Returns the value of IMPLEMENTATION

Returns a pointer to the IMPLEMENTATION instance

Sets the value of IMPLEMENTATION to VAL

Returns a pointer to the appropriate instance if VAL is an IMPLEMENTATION, otherwise returns NIL

Initializes NAME Returns the value of NAME Sets the value of NAME to VAL

Creates a REQUIREMENTS instance whose name is NAME and adds it to REQUIREMENTS

Initializes REQUIREMENTS Returns the value of REQUIREMENTS Returns a pointer to the REQUIREMENTS instance

Sets the value of REQUIREMENTS to VAL

Returns a pointer to the appropriate instance if VAL is an REQUIREMENTS, otherwise returns NIL

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Object Class: REQUIREMENTS

Subclasses HEADER ENTRY

Instance Variables TOP-LEVEL-ACTIVITY

ACTIVITIES

DATA-ELEMENTS

ALIASES

Methods create initialize

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present (ATTRIBUTE-LIST)

reinitialize (DUMMY)

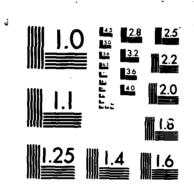
save (PROJECT, STREAM) Name of the highest level activity Assoc-list of activity names and pointers to activity instances Assoc-list of data element names and pointers to data element instances Assoc-list of alias names and pointers to alias instances

Creates a REQUIREMENTS instance Initializes all of REQUIREMENTS' instance variables

Displays the instance variables listed in ATTRIBUTE-LIST of a REQUIREMENTS instance; if ATTRIBUTE-LIST is empty, displays all instance variables Initializes all of REQUIREMENTS' instance variables except NAME, TYPE, and PROJECT

Saves a REQUIREMENTS instance and all of its sub-instances to STREAM

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Object Class: SLOT

Subclasses none

Instance Variables VALUE

REQUIRED

SET

PRESENT-LABEL

Methods create

initialize

present

present-label

set-present-label
(VAL)

required set-required (VAL)

set set-set (VAL)

value set-value (VAL) Pointer to low-level value instance Boolean - "YES" if entry required "NO" otherwise Boolean - "YES" if value has been set "NO" otherwise Label to be displayed for present operation

Creates a SLOT instance Initializes all of SLOT's instance variables Displays a SLOT instance

Returns the value of PRESENT-LABEL

Sets the value of PRESENT-LABEL to VAL

Returns the value of REQUIRED Sets the value of REQUIRED to VAL

Returns the value of SET Sets the value of SET to VAL

(VAL) Returns the value of VALUE Sets the value of VALUE to VAL

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Object Class: TEXT

Subclasses none

Instance Variables VALUE

Methods initialize

present value set-value (VAL) Initializes all of TEXT's instance variables Displays a TEXT instance Returns the value of VALUE Sets the value of VALUE to VAL Subclasses PARAMETER

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Instance Variables none









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Appendix F

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Interface Object Definition

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Introduction

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This appendix defines the object-oriented structure of the Data Dictionary System's (DDS) natural language front-end to Data Dictionary Tool (DDT) Interface. The Interface's EVENT object class is specified in terms of its included subclasses, its instance variables, and its associated access methods.



Object Class: EVENT

Subclasses none

Instance Variables ACTION ACTION-TYPE

OBJECT

ATTRIBUTES

Methods execute initialize

present

instance variables initialize-action Initializes ACTION

not

initialize-action-type initialize-attributes

add-attribute (LIST)

initialize-object

Initializes OBJECT

ATTRIBUTES

Action to perform

an action upon a data

specified by ACTION

instance variables

Displays all of EVENT's

Initializes ACTION-TYPE

Initializes ATTRIBUTES

For each attribute in LIST: Verifies not already a member of ATTRIBUTES; If not, adds to

Type of action - either NIL if

dictionary object of META if

A list of instance variables of OBJECT on which to perform the action specified by ACTION

Interprets an EVENT instance Initializes all of EVENT's

Data dictionary object on which to perform the action













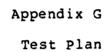












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Introduction

This appendix specifies the test plan to be used in formally testing the DDS data dictionary system including the CoIn natural language processor. This plan consists of a series of functional tests which specify what is being tested, how the test should be performed (at a high level), and what the results of the test should be. This test plan is organized to follow the System Requirements section of the requirements chapter of this thesis, Chapter II.

System-Level Tests

1. This test is to determine if DDS is able to accept a valid input sentence from the user, interpret the sentence as a command, and retrieve and/or modify the requested information from the data dictionary. The test should be performed by entering several valid requests into the system. Any information requested to be displayed should be displayed. Any information requested to be modified should be modified.

2. This test is to determine if DDS is able to reject an invalid input sentence from the user. The test should be performed by entering several invalid requests into the system. Each of the requests should cause an error message to be displayed.

Grammar Constructor Tests

1. This test is to determine if the grammar constructor is able to save a grammar from the working grammar file to a permanent grammar file. The test should be performed by entering a valid request to save a working grammar. The grammar should be saved to a permanent file.

This test is to determine if the grammar constructor is able to load a grammar from a permanent grammar file to the working grammar file. The test should be performed by entering a valid request to load a permanent grammar. The grammar should be loaded to the working grammar file.
 This test is to determine if the grammar constructor is able to initialize the working grammar. The test should be performed by first loading the working grammar from a permanent grammar and then entering a valid request to initialize the working grammar. The working grammar should be initialize the working a valid request to initialize the working grammar.

4. This test is to determine if the grammar constructor is able to create production rules. The test should be performed by entering a valid request to create a new production rule. The production rule should be created.
5. This test is to determine if the grammar constructor is able to display the existing production rules. The test should be performed by entering a valid request to display a

production rule. The requested production rule should be displayed.

6. This test is to determine if the grammar constructor is able to display a list of the existing production rules. The test should be performed by entering a valid request to display a list of the existing production rules. A list of the existing production rules should be displayed.

7. This test is to determine if the grammar constructor is able to modify an existing production rule. The test should be performed by entering a series of valid requests to modify production rules. At least one entry should be made for each defined modification command. The requested modifications should be made to the proper production rules.
8. This test is to determine if the grammar constructor is able to delete an existing production rule. The test should be performed by entering a valid request to delete a production rule. The specified production rule should be deleted.

9. This test is to determine if the grammar constructor is able to reject invalid construction requests. The test should be performed by entering a series of invalid requests. At least one entry should be made for each of the defined commands of the the constructor. Each of the requests should cause an error message to be displayed.

Sentence Interpreter Tests

1. This test is to determine if the sentence interpreter is able to accept input sentences that are valid within the defined grammar. The test should be performed by entering a series of valid sentences. At least one entry should be made which exercises each of the defined grammar productions. The interpreter should correctly parse each entry.

This test is to determine if the sentence interpreter is 2. able to reject input sentences that are invalid within the defined grammar. The test should be performed by entering a series of invalid sentences. At least one sentence should be entered for each of the following cases: a sentence whose first word is not valid within the grammar, a sentence whose last word is not valid within the grammar, a sentence which contains a word that is not valid within the grammar and in which the word is neither the first nor the last word of the sentence, a sentence which is correct within the grammar except that it contains at least one extra word at the end, and a sentence that is correct within the grammar except that it ends before the defined end-of-sentence. The interpreter should reject each entry and display an error message. The error message should attempt to show where in the input sentence the error occurred. The error message should offer suggestions as to how to correct the problem.

3. This test is to determine if the sentence interpreter is able to syntactically translate the results of a valid sentence parse into the proper execution commands of the application tool. The test should be performed by entering a series of valid sentences should generate all (if this is reasonable) or several (if it is not reasonable to generate all) of the execution commands of the application tool. For each input sentence, the proper application commands should be generated.

Data Dictionary Access Tests

 This test is to determine if the data dictionary access process is able to access all of the information stored in the data dictionary file. The test should be performed by entering a series of valid requests to access the data dictionary. As a minimum, an entry should be made which accesses an instance of each object type in the data dictionary. Also as a minimum, an entry should be made which causes each access type (add, delete, display, etc.) to be performed. For each entry, the proper access to the information stored in the data dictionary should be done.
 This test is to determine if the data dictionary access process is able to maintain the "consistency" of the data in the data dictionary. The test should be performed by entering a series of valid requests which cause each of the

defined consistency routines to be invoked. For each entry, the data dictionary access process should maintain the consistency of the information stored in the data dictionary.

3. This test is to determine if the data dictionary access process is able to add new objects to the data dictionary. The test should be performed by entering a series of valid requests to add new object instances. As a minimum, an entry should be made which adds a new instance of each defined object type. The specified instances should be added to the data dictionary.

4. This test is to determine if the data dictionary access process is able to modify existing information in the data dictionary. The test should be performed by entering a series of valid requests to modify existing object instances. An entry should be made which exercises each defined modification function. For each entry, the requested modification should be made.

5. This test is to determine if the data dictionary access process is able to reinitialize existing information in the data dictionary. The test should be performed by entering a series of valid requests to reinitialize existing object instance. As a minimum, an entry should be made which reinitializes an instance of each defined object type. The

specified instances should be reinitialized to their respective initial states.

6. This test is to determine if the data dictionary access process is able to delete existing information in the data dictionary. The test should be performed by entering a series of valid requests to delete existing object instances. As a minimum, an entry should be made which deletes an instance of each defined object type. The specified instances should be deleted from the data dictionary.

7. This test is to determine if the data dictionary access process is able to reject modification requests which reference information not contained in the data dictionary. The test should be performed by entering a series of modification requests which reference information not contained in the data dictionary but which are otherwise valid. An entry should be made which exercises each defined modification function. For each request, the data dictionary access process should respond with an error message.

8. This test is to determine if the data dictionary access process is able to display information stored in the data dictionary. The test should be performed by entering a series of valid presentation requests. As a minimum, an entry should be made which displays an instance of each

define object type. The specified instances should be displayed.

9. This test is to determine if the data dictionary access process is able to reject presentation requests which reference information not contained in the data dictionary. The test should be performed by entering a series of presentation requests which reference information not contained in the data dictionary but which are otherwise valid. For each request, the data dictionary access process should respond with an error message.

10. This test is to determine if the data dictionary access process is independent of the natural language processor. This test should be performed by inspecting the source code listings of the data dictionary to verify that it does not reference any of the modules or data local to the natural language processor. No references to the modules or data of the natural language processor should be found.

Grammar Tests

1. This test is to determine if the defined grammar of DDS is functionally complete. This test should be performed by entering a series of valid requests which invoke each of the operations provided by the data dictionary access process. Requests should be made which access instance of each of the

defined object types. The specified operation should be performed on the specified object instance.

2. This test is to determine if the defined grammar of DDS is sufficiently flexible. No test is specified to validate this requirement. The entries of users of varying expertise should be monitored to determine flexibility shortcomings.

Data Dictionary Content Tests

1. This test is to determine if the data dictionary contents of DDS is consistent with the requirements specified in <u>Development Documentation Guidelines and</u> <u>Standards</u> (AFIT/ENG, 1984). This test should performed by inspecting the source code listings to verify that they are consistent with this document. No inconsistencies should be found.

SDW Interface Tests

1. This test is to determine if the data dictionary system is interfaced into the Software Development Workbench (SDW). This test should be performed by executing the SDW and entering the proper codes to invoke DDS. The SDW should execute DDS.

2. This test is to determine if the natural language processor is interfaced into the SDW. This test should be performed by executing the SDW and entering the proper codes

to invoke the natural language processor. The SDW should execute the natural language processor.

Conclusion

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This test plan has specified a suite of tests that validate the requirements of DDS. As the requirements change, this document should be updated to remain consistent with them.

























Appendix H

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CoIn User Manual



















CoIn User Manual

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Version 1.0

Capt Stephen A. Wolfe, USAF

AFIT/GCS-85D

December 1985

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I. Introduction

CoIn is a natural language processing system. It is implemented on the AFIT Information Systems Laboratory VAX-11/780 computer. CoIn is a part of AFIT's Software Development Workbench (SDW). CoIn is written in NIL Lisp and runs within the NIL interpreter environment.

CoIn consists of two subsystems: a grammar constructor and a sentence interpreter. The grammar constructor allows you to define a grammar consisting of production records and function records (these are described in Chapter III). Using a grammar built with the constructor, the sentence interpreter is able to parse input sentences and execute Lisp code that you have built into the grammar.

II. Scope

This manual describes the use of the CoIn natural language processing system. It is assumed that the reader has a minimal working knowledge of the VMS operating system (see DEC, 1984, for more information), the SDW (Hadfield, 1982), and the NIL interpreter (Burke, et al, 1984).

III. Using the Grammar Constructor

This chapter describes how to use the grammar constructor subsystem of CoIn. This chapter is divided into four sections: Getting Started, which explains how to execute CoIn; Creating a Grammar, which explains how to create a new grammar and how to permanently store it when you're done; Modifying a Grammar, which explains how to create new production and function records and how to alter and delete existing production and function records; and Exiting which explains how to exit from CoIn to the SDW.

Getting Started

CoIn is integrated into the SDW. This makes it very easy to execute. First log into the SDW account, enter the SDW, and move to your project directory. Return to the top-level menu and enter 'ED' (the Text Editors functional group) at the prompt. The EDITOR (ED) MENU will be displayed. When prompted, enter 'NL' (the CoIn Natural Language Grammar Editor). NIL will be invoked, and CoIn will be loaded. When NIL is done loading, the message "*** CoIn Loaded ***" will be displayed. At that time, CoIn is ready for input.

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Creating a Grammar

Most of the time you will probably not want to create a grammar from scratch. You'll probably just want to add to or modify an existing grammar. However, it is possible to create a new grammar. To do so, first invoke CoIn as described in the previous section, then type "(initialize-grammar)". Be careful! If you have loaded a grammar, modified it, and not saved it, this command will destroy all of your changes!

After initializing the grammar, you can then proceed to create your grammar by using the commands as described in the next section, Modifying a Grammar. When your grammar is complete (or anytime you want to save a partial definition), you can type "(save-grammar)" to store it on a disk file. The grammar will be saved in a file called GRAMMAR.LSP in the current NIL working directory (your default directory unless you have told NIL otherwise).

Modifying a Grammar

A grammar consists of a set of production records and function records. Production records are used to store the productions of a grammar and consist of three fields: NAME, SUB-PRODUCTION LIST, and CODE (see Figure 1). The NAME field contains the element which the current word of the input sentence must match for a parse to proceed. The

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	+~=~=~+===========================+	
	<pre>¶ NAME ¶ SUB-PRODUCTION LIST ¶ CODE ¶</pre>	
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Figure 1. Production Record Fields

SUB-PRODUCTION LIST field contains a list of possible sub-productions which can be used to continue the parse should the current word parse successfully. The CODE field contains a piece of Lisp code which is executed if the production it is associated with is successfully traversed. A production must end with an end-of-sentence marker (\$) if it is at the top level of the grammar. Otherwise, it must end with an end-of-production marker (*).

Function records are used to store "functions". A function is equivalent to a Lisp predicate function. The function is applied to the current word of an input sentence and, if the word is part of the function's domain, non-NIL is returned. Otherwise, NIL is returned. A function record consists of two fields: NAME and CODE (see Figure 2). The NAME field contains the name of the function, and the CODE field contains the Lisp code to be applied.

The following subsections describe how to create a new production record, create a new function record, add a

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1	NAME V CODE V	
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Figure 2. Function Record Fields

grammar production to an existing production record, modify the CODE field of an existing function record, destroy an existing production record, destroy an existing function record, and do other miscellaneous modifications.

Creating a Production Record. To create a new production record, type "(create-production PRODUCTION-NAME (LISP-CODE))" where PRODUCTION-NAME is replaced by the name of the production record you wish to create and LISP-CODE is replaced by the Lisp code that you want executed when the production record is successfully used. For instance, if you want to create a production record called <display-message-1> that displays the message "message 1" when it is traversed, then you should enter:

(create-production <display-message-l>
 (msg "message 1"))

If the production record <display-message-l> already exists, an error message is displayed and the existing production record is left as is. The LISP-CODE field is optional.

<u>Creating a Function Record</u>. To create a new function record, type "(create-function FUNCTION-NAME (LISP-CODE))" where FUNCTION-NAME is replaced by the name of the function record you wish to create and LISP-CODE is replaced by the Lisp code that defines the function. For instance, if you want to create a function record called <number> that is defined using the Lisp "numberp" predicate, then you should enter:

The LISP-CODE field is optional but if left out, must subsequently be modified before using the grammar. If this is not done, the sentence interpreter will crash when it tries to use the function record. As is true for production records, an error message is displayed if you try to create a function record using a name that already exists.

Adding a Production. To add a grammar production to an existing production record, type "(add-production (PRODUCTION) (LISP-CODE) PRODUCTION-NAME)" where PRODUCTION is replaced by the grammar production to be added, LISP-CODE is replaced by the Lisp code you want executed when the production is successfully traversed, and PRODUCTION-NAME is

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replaced by the name of the production record to which you want the production added. For example:

(add-production (what is *) () <display>)

adds the production "(what is *)" to the existing production record <display>. As is shown, the LISP-CODE parameter may be empty. However it is not optional. The parentheses must be included. If the <display> production record does not exist, an error message is displayed.

Note the use of the end-of-production marker (*). All productions that are not at the top level (that is, they are not part of your highest level production record) must end with an end-of-production marker. All productions at the top level must end with an end-of-sentence marker (\$).

Modifying a Function. To modify the CODE field of an existing function record, type "(modify-function (LISP-CODE) FUNCTION-NAME)" where LISP-CODE is replaced by the new Lisp code that defines the function and FUNCTION-NAME is replaced by the function to be modified. For example:

(modify-function (lambda (x) (null x)) <empty-list>)

modifies the CODE field of the function record <empty-list> to be the Lisp expression (lambda (x) (null x)). If the specified function record does not exist, an error message is displayed. The LISP-CODE may be empty, but, as is true

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for the add-production operation, the parentheses must be included.

Destroying a Production Record. To delete an existing production record from a grammar, type "(destroy-production PRODUCTION-NAME)" where PRODUCTION-NAME is replaced by the name of the production record to destroy. For instance, to destroy the production record <display>, enter:

(destroy-production <display>)

As usual, if the specified production record does not exist, an error message is displayed.

Destroying a Function. To delete an existing function record from a grammar, type "(destroy-function FUNCTION-NAME)" where FUNCTION-NAME is replaced by the name of the FUNCTION to destroy. For example:

(destroy-function <number>)

deletes the <number> function record. Again, an error message is displayed if the function record does not exist.

Other Operations. In addition to the operations described above, a grammar can be modified by using a Lisp s-expression editor that is included in CoIn. The editor is a slightly modified version of a Lisp function editor described in <u>Artificial Intelligence Programming</u> by Charniak, Riesbeck, and McDermott (Charniak, et al, 1980).

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To specify a production or function record to edit, type "(sedit STRUCTURE-NAME)" where STRUCTURE-NAME is replaced by the name of the production or function record that you wish to edit. Once you have entered this command, the editor is used exactly as described in Chapter 7 of <u>Artificial</u> Intelligence Programming.

Exiting

To exit from NIL and return to the SDW, type "(quit)". This will return you to the Text Editors functional group of the SDW.

IV. Using the Sentence Interpreter

This chapter describes how to use the sentence interpreter subsystem of CoIn. This chapter has only one section, Interpreting A Sentence, which explains how to enter a request to interpret a sentence. Invoking and exiting the system is the same as is described in Chapter 3.

Interpreting a Sentence

To invoke the interpreter to parse a sentence, type "(parse '(SENTENCE) TOP-LEVEL-PRODUCTION)" where SENTENCE is replaced by the sentence to be parsed and TOP-LEVEL-PRODUCTION is replaced by the name of your top-level production record. For instance:

(parse '(show me module M1) <grammar>)

tells the interpreter to parse the sentence "show me module Ml" using the grammar specified by the top-level production record <grammar>. If your input sentence is not valid within the specified grammar, an error message is displayed.

V. <u>A Short Example</u>

This chapter provides a short, but one hopes useful, example of the use of CoIn. First a grammar is defined, then the constructor commands to implement it are shown. Finally, several sample interpreter commands are included. The defined grammar accepts any sentence consisting of a string of one or more occurrences of the word "a".

Grammar Productions

<grammar></grammar>	->	< a	•>
<a>	->	a	
<a>	->	a	<a>

Constructor Commands

(initialize-grammar)
(create-production <grammar> (msg "Parse complete" N))
(add-production (<a> \$) () <grammar>)
(create-production <a> ())
(add-production (a *) (msg "Using (a *)" N) <a>)
(add-production (a <a> *) (msg "Using (a <a>)" N) <a>)

Sample Interpreter Inputs

(parse '(a) <grammar>)
(parse '(a a) <grammar>)
(parse '(a a a a a) <grammar>)



VI. Conclusion

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This manual has described how to use the CoIn natural language processor. For an extended example of the use of CoIn, see (Wolfe, 1985), which describes the implementation of a natural language human-computer interface for a data dictionary system.

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Appendix I

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DDS User Manual

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DDS User Manual

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version 1.0

Capt Stephen A. Wolfe, USAF

AFIT/GCS-85D

December 1985

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I. Introduction

DDS is a data dictionary system which employs a natural language human-computer interface. To interact with DDS, one enters English sentences which specify to DDS the actions it is to perform. DDS is implemented on the AFIT Information Systems Laboratory VAX-11/780 computer. It is a part of AFIT's Software Development Workbench (SDW). DDS is written in NIL Lisp and runs within the NIL interpreter environment.

II. Scope

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This manual describes the use of the DDS data dictionary system. It is assumed that the reader has a minimal working knowledge of the VMS operating system (see DEC, 1984, for more information), the SDW (Hadfield, 1982), and the NIL interpreter (Burke, <u>et al</u>, 1984).

III. Using DDS

This chapter describes how to use DDS. This chapter is divided into 3 sections: Getting Started, which explains how to execute DDS; Using the System, which explains how to enter commands to the system; and Exiting which explains how to exit from DDS to the SDW.

Getting Started

DDS is integrated into the SDW. This makes it very easy to execute. First log into the SDW account, enter the SDW, and move to your project directory. Return to the top-level menu and enter 'DS' (the Design Tools functional group) at the prompt. The DESIGN TOOL (DS) MENU will be displayed. When prompted, enter 'DD' (the DDS data dictionary system). NIL will be invoked, and DDS will be loaded. When DDS is done loading, the message "*** Data Dictionary System (DDS) Loaded ***" will be displayed. To execute DDS, type "(dds)" (without the quotes). DDS will be invoked and will display the message "Data Dictionary System (DDS)" followed by a prompt "DDS->". At that time, DDS is ready for input.

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Using the System

Currently the capabilities of DDS are very limited. The only data manipulation operations that are implemented are the initialization and presentation functions. Also implemented are the clear screen, help (again limited), and exit meta-functions. The best source for determining what the current system can and cannot do is Appendix D (Implemented Grammar) of <u>A Natural Language Processor and its Application to a Data Dictionary System</u> (Wolfe, 1985). This chapter presents examples of valid input sentences for the current grammar. Please note that the input sentences are enclosed within parentheses, and that they contain no punctuation. This format must be used, or NIL will complain unmercifully.

<u>Clearing the Screen</u>. The following examples show how to clear the terminal screen in DDS:

(please clear the screen)

(cls)

Displaying the Help Message. The following example shows how to display the help message in DDS:

(help)

Exiting from DDS. The following examples show how to exit from DDS:

(quit)

(exit)

Initializing a Data Dictionary Entry. Before an entry can be initialized, it must exist. Currently, the only way to create an entry is to modify your data base (the file DB.LSP) using a text editor. In the following examples, it is assumed that P1 is a defined process and PARAM is a defined parameter.

(initialize process Pl)

(please init parameter PARAM)

Displaying a Data Dictionary Entry. Before an entry can be displayed, it must exist. See the discussion in the previous subsection concerning creation of an entry. In the following examples, it is again assumed that Pl is a defined process and PARAM is a defined parameter.

(please show me process Pl)

(what are the processes called by process P1)

(display the name type and version of parameter PARAM)

Exiting

To exit from DDS and return to NIL, type "(quit)". To exit from NIL and return to the SDW, type "(quit)" again. This will return you to the Design Tools functional group of the SDW.

IV. Conclusion

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This manual has described how to use the DDS data dictionary system. For a more complete description of DDS, see (Wolfe, 1985).

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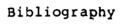












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Captain Stephen A. Wolfe was born on 23 February 1956 in Portland, Oregon. He graduated from high school in Spring Valley, California, and attended San Diego State University from which he received the degree of Bachelor of Arts in Computer Science in May 1979. He graduated with high honors and with distinction in the major. After graduating, he received a commission in the United States Air Force through the Officer Training School program. His first assignment was at Space Division in Los Angeles, California, where he was a systems analyst and applications programmer. In April 1982, he was transferred to Washington, DC, where he became the lead software engineer for the procurement of the Ground Launched Cruise Missile Weapons Control System. He entered the School of Engineering, Air Force Institute of Technology, in May 1984.

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The development of a human-computer interface construction and repretation tool capable of processing English-like of natural fanguage user input is discussed. The utility of the tool is demonstrated by using it to create the natural language interface for a data dictionary system. The data dictionary's development is also documented and is used as the overall context for the presentation.

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