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NAVAL POSTGRADUATE SCHOOL

Monterey, California



REAL-TIME APPLICATIONS IN MULTIPROCESSOR SYSTEMS

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M. Kadri Ozyurt

December 1983

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Thesis Advisor:

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Real-Time Applications in Multiprocessor Systems

by

M. Kadri Ozyurt Lieutenant J.G., Turkish Navy

Submitted in partial fulfillment of the requirements for the degree of

MASTER OF SCIENCE IN ENGINEERING SCIENCE

from the

NAVAL POSTGRADUATE SCHOOL December, 1983

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ABSTRACT

This thesis builds a simulation model of a tactical fire control system in a real time environment, using a tightly connected multi-processing system consisting of two single board computers. The additional hardware used in this project consists of an ADM-3A video terminal with a built-in retrographics feature, an MDS microprocessor development system, an analog-to-digital convertwo sets of triplet potentiometers. The potenter, and tiometers are used to feed analog information about ownship, targetship, and gun position to the simulation model, which then evaluates and computes projected target positions and gun control parameters, and displays the results.

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P

I. INTRODUCTION

A. BACKGROUND

To this date, many tactical control and decision systems have been designed and implemented in various places, where the nature of the job required fast response and decision making. The NTDS (Navy Tactical Decision System), for instance, is one such system implemented for U.S. Navy ships in 1962, and is still in use today, with recent hardware modifications.

The revolutionary developments in the LSI (Large Scale Integration) and VLSI (Very Large Scale Integration) industries during the period from 1972 to 1983 have made the costs of computing much less expensive than the costs of yesterday's systems. Today's products, which are based on these innovations in large scale integration, have proved to be more reliable and more versatile than the old systems, and they can also be tailored to the needs of whatever the nature of the requirements may be. The serviceability, availability and inexpensiveness of these products, in addition to the above-mentioned features, offer both the designers and the implementers an opportunity to take advantage of this technology.

B. DISCLAIMER

Many terms used in this thesis are registered trademarks of commercial products. Rather than citing each individual occurance of a trademark throughout this thesis, all registered trademarks referred to in this thesis will be listed below, following the name of the firm holding the trademark.

- Intel corporation, Santa Clara, California: Intel, Intel 8086, iSBC 86/12A, MULTIBUS, MDS
- Digital Research Corporation, Pacific Grove, California: CP/M, CP/M-86, PL/I-80, PL/I-86, TED, RASM-86, LINK-86, DDT-86
- EX-CELL-O Corporation, Irvine, California: REMEX Data Warehouse
- MicroPro International, San Rafael, California: Wordstar
- Micropolis Corporation, Chatsworth, California: Micropolis
- Lear Siegler, Inc., Anaheim, California: ADM-3A

C. PURPOSE OF THIS THESIS

The purpose of this thesis is to create a simulation model for real time tactical systems which can be used to study the following features:

- 1. Multiprocessor system real time performance;
- 2. System reliability;
- 3. Graphics Display;
- 4. Software Engineering.

In order to carry out these objectives, certain hardware changes had to be made. The real time applications necessitate the existence of an interruptdriven configuration originating from an accurate timer. The interconnections and the appropriate initializations, both in the timing (PIT) and the interrupt (PIC) circuits. were made on the iSBC 86/12A board, to give the required real time clock. After achieving the desired form of operation, the real time executive module (which synchronizes the operation of the simulation system programs) was tested in an interrupt-driven environment. After testing out both the real time executive and most of the simulation system programs, the individual execution times of the simulation system programs were measured with the aid of TIMES.AID, an XINCLUDE file (see Apppendix F).

It was intended to utilize two iSBC 26/12A single board computers in order to study the real time performance of a tightly connected multiprocessing scheme. The Intel MDS (Microprocessor Development System) allows the configuration of such an expansion, through its 20 bit MULTIBUS backplane. It was also planned to write the required software to prevent a single point failure and to gain a "graceful degradation" in the case of a malfunction in any of the single board computers.

The ADM-3A video terminal, with its built-in Retrographics unit, was utilized for graphics display purposes. The Retrographics card contains a 2-80A eight bit microprocessor. This allows the computations for high precision graphics to be done by the video terminal. That improves the efficiency by removing much of the overhead the iSBC 86/12A single board computers. from The Retrographics unit can also make drawings and erasures selectively. This improves the display time, which might be lost due to total erasures and redrawings. In other words, the selective erasing capability decreases the display time, such that the program may make partial erasures, erasures of single objects.

The programs written for the simulation model were made as modular as much as possible, to facilitate testing and maintenance, and to make room for future alterations. A procedure call was placed wherever a critical design decision was to be made. This procedure call gives the option of changing a critical design decision if one that is more efficient is designed. Structured programming and efficient data structures were meant to be utilized. Circular linked-lists are examples of such a programming technique.

D. THESIS ORGANIZATION

The thesis is organized into four chapters. The program listings developed to implement the simulation system are appended at the end of the text. The first chapter covers the background, the disclaimer for the trademarks used in this thesis, the intended purpose of the project and thesis organization. The second chapter covers the system configuration and the hardware components. The third chapter deals with software modules written in both PL/I-86 and RASM-86 assembly languages. The program segments are discussed in detail in this chapter. Some information about the data structures used in the development of the software are brought up, as well as the initialization of the programmable hardware components. In the final chapter. some conclusions are presented on the work involved in the implementation of the simulation system.

II. SYSTEM HARDWARE

A. SYSTEM CONFIGURATION

The existing system hardware (see Figure 2.1) uses the Intellec microcomputer development system (MDS), which allows the expandability required to set up a multi-processor system. Within the MDS, the boards required for the operation of the system are interconnected through the MULTIBUS backplane. These boards are the following: two iSBC 86/12A boards, the front panel control board, an A to D converter board, and two interface boards for disk drives. The MDS utilizes an Intel disk drive unit which has two disk drives. Standard 8-inch IEM floppy disks are used as the removable storage media.

Fach iSBC 86/12A board had a RAM capacity of up to OFFFF hexadecimal (65535 decimal) eight bit bytes. Since the MULTIBUS is a 20-bit address bus, the address space of the whole system can be expanded up to 1 megabyte. The single board computers can address this memory space by their 20-bit address bus. Each board is so wired that the first 64 Kbyte RAM segment resides on its board (2000 through FFFFH). The 64K RAM segments on each toard can be wired to be accessible from the MULTIBUS as dual ported memories in the 1-megabyte address space. One of the iSBC 86/12A boards 15 the master of the master-slave

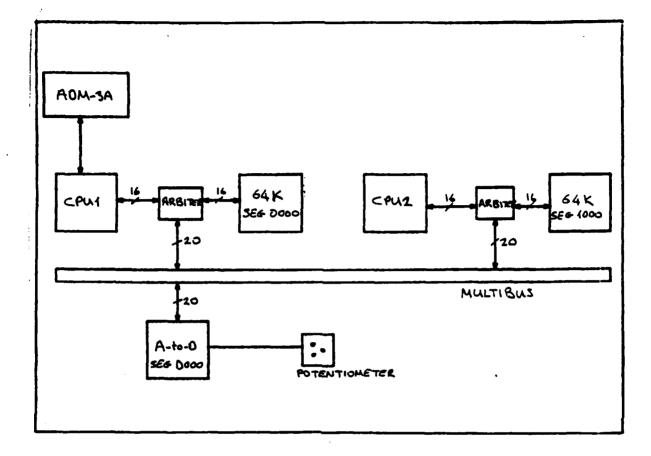


Figure 2.1 - System Interconnection

multiprocessor scheme. The master board is assigned the first S4Ksegment (00000h - 0FFFFh) where the slave board is assigned the second 64K segment (10000h-1FFFFh). In this configuration the A-to-D board occupies the segment with the addresses (D0000h-DFFFFh). The outputs of the A-tc-D converter are memory-mapped as eight bit bytes and occupy the locations DF700H through DF70FH. That is to say that they are viewed by the CPU's in the system as being

ordinary memory locations. An analogy to that is that they resemble read only memory that can only be read but not written in to.

The master iSBC 86/12A is connected to the ADM-3A video terminal through it's serial I/O connector. This is the only means of communication of the system to the outside world. The results of the simulation are fed to the video display through this interconnection. The ADM-3A video terminal has a built-in retrographics feature. The consequence of this is that the ADM-3A works not only as an alphatetic terminal but also may act as a graphics device due to the fact that the retrographics card itself has a 2-80A microprocessor built in. This microprocessor allows the high precision graphics computations to be done without the need of any other external processors. In this case, ADM-3A has four operational modes, each one of which has an impact on the simulation graphics and will be discussed later in this chapter.

B. HARDWARE

In the following subsections the individual components that comprise the system hardware are presented.

1. MDS

The Intellec Microcomputer Development System (MDS) is a complete development tool which allows the integration of both microcomputer hardware and software development. The system operates under the control of an 8086 microprocessor which supervises all system resources such as the main memory, I/O peripheral devices, and optional system facilities, such as A-to-D converters. It can support up to 7 iSBC 86/12A boards in this configuration. Some of the important boards are presented below.

a. Front Panel Control Board

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The Front Panel Control Board contains circuits for controlling the front panel options. It also provides some signals for bus control, clock generation, and the bootstrap program. A bus time-out system is included to prevent the CPU from halting operation if a nonexistent memory location or an incorrect I/O port is addressed.

This board produces two types of clock pulses:

(1) Bus Clock (10 MHz), used in Bus transactions;

(2) Common Clock (10 MHz), used by system
devices;

b. Disk Interface Boards

These two cards contain the disk controller interface for each drive in the Disk Storage Unit.

2. Single Board Computer

Intel's Single Board Computer iSBC 86/12A is used in the system. It is a member of Intel's complete line of 8- and 16-bit single board computer products and is a complete computer system on a single printed-circuit assem-The iSBC 86/12A toard includes a 16-bit central bly. processing unit (CPU), 32K tytes (32,768 tytes) of dynamic RAM, a serial communications interface (USART), three programmable parallel I/O ports, programmable timers (PIT), priority interrupt control (PIC). Multibus interface control logic, and bus expansion drivers for interfacing with other Multibus interface-compatible expansion boards. Also included is a dual port control logic to allow the iSBC 86/12A board to act as a slave RAM device to other Multibus interface masters in the system, as is the case in this project. In the current state of the hardware, the RAM capacity of both iSBC 86/12A boards is expanded up to 64K bytes by installing an iSBC 300 Multimodule RAM option. A read only memory of 16K bytes is also added to both iSBC 86/12A boards. The important components that make up the 1SBC 86/12A board are discussed in the following subsections.

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a. CPU

The iSBC 86/12A Single Board Computer is controlled by an Intel 8086 16-bit Microprocessor (CPU). The 8086 CPU includes four 16-bit general purpose registers that may also be addressed as eight 8-bit registers. In addition, the CPU contains two 16-bit pointer registers and two 16-bit index registers. Four 16-bit segment registers,

2Ø

specifically: code, data, extra, and stack segment registers; allow extended addressing to a full megabyte of memory. The CPU instruction set supports many variations of addressing modes and data transfer operations, signed and unsigned 8-bit and 16-bit arithmetic including hardware multiply and divide, and logical and string operations. The CPU architecture permits dynamic code relocation, reentrant code, and instruction lookahead.

b. Serial I/O

The serial I/O port is controlled and interfaced by an Intel 8251A USART (Universal Synchronous/Asynchronous Receiver/Transmitter) chip. The USART is individually programmable for operation in most synchronous serial data transmission formats.

In the synchronous mode, the following are programmable:

- (1) Character length
- (2) Sync character (or characters)
- (3) Parity

In the asynchronous mode the following are programmable:

- (1) Character length
- (2) Baud rate factor
- (3) Stop bits
- (4) Parity

In both the synchronous and asynchronous modes, the serial I/O port features half- or full-duplex, double buffered transmit and receive capability. The USART transmit and receive clock rates are supplied by a programmable baud rate/time generator.

c. Programmable Interval Timer

Three independent, fully programmable 16-bit interval timer/event counters are provided by in Intel 8253 Programmable Interval Timer (PIT). Each counter is capable of operating in either BCD (binary coded decimal) or binary modes; two of these counters are available to the system's programmer to generate time intervals under software control. In this thesis project the counteri is used to generate timing pulses required for the real time clock operation to the system software. These pulses are sent to the PIC via the interrupt matrix as being an IR1 input request.

d. Priority Interrupt Control

The priority interrupt control (PIC) which can be programmed to respond to edge-sensitive or level-sensitive inputs, treats each true input signal condition as an interrupt request. After resolving the interrupt priority, the PIC issues a single interrupt request to the CPU. Interrupt priorities are independently programmable under

under software control. The programmable interrupt priority modes are:

(1) Nested Priority. Each interrupt request has a fixed priority: input Ø is highest, input 7 is lowest. This mode of operation is chosen in this thesis project;

(2) Fully Nested Priority. This is essentially the same as item (1) above, with the exception that the requesting input is not locked out and pending requests are still accepted;

(3) Auto-Rotating Priority. Priorities areequal. The last received input becomes the lowest priorityinput;

(4) Specific Priority. Software assigns the priorities;

(5) Special Mask. Interrupt requests that are being serviced are masked out;

(6) Poll. The CPU's internal interrupt enable is disabled. Interrupt service is achieved by a programmer-initiated Poll command.

The iSBC 86/12A board provides two sorts of interrupts which are bus vectored (BV) and non-bus vectored (NBV). The former deals with the interrupt requests from off-board sources where the latter deals with various on-board sources. The interrupt requests are fed to the

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PIC through the jumpers of the interrupt matrix, which will be discussed in the next subsection.

e. Interrupt Matrix

Interrupt requests may originate from eighteen sources without the necessity of external hardware. The interrupt matrix connects the selected source lines to a maximum of eight selected inputs of the PIC. It is an array of pins which can be connected to each other via jumper wires. There are two types of pins. The eighteen source lines constitute the input pins, where the pins that lead to the IRO through the IR7 inputs of the PIC constitute the output pins. (See Figure 2-2 for the interconnection scheme of this thesis project.)

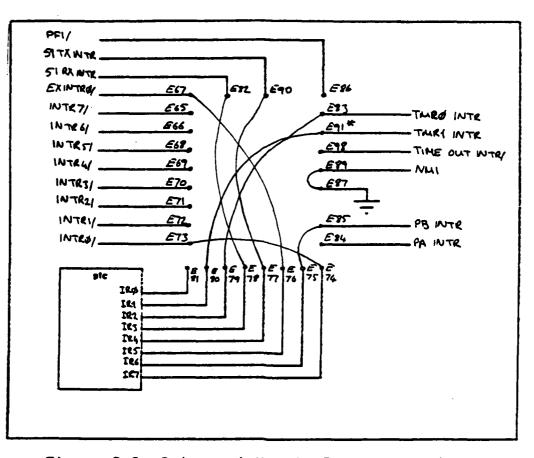
f. Dual Port RAM

The iSBC 86/12A board has an internal bus for all on-board memory and I/O operations. Hence, local (onboard) operations do not involve the MULTIBUS interface, making it available for other iSBC 86/12A boards for a multi-processor scheme. Dual port control logic is included to interface so that the iSBC 86/12A board can function as a slave RAM device (or common memory) when not in control of the Multibus interface. The CPU has priority when accessing on-board RAM. After the CPU completes its read or write operation, the controlling bus master is is allowed to access RAM and complete its operation. Where both the CPU and the controlling bus master have the need

to write or read several bytes or words to or from the onboard RAM, their operations are interleaved. For CFU access, the on-board RAM addresses are assigned from the bottom up of the 1-megabyte address space; i.e., 00000-ØEFFFh. The slave RAM address decoding logic includes jumpers and switches to allow positioning the on-board RAM into any 64-K segment of the 1-megabyte system address space. The slave RAM can be configured to allow either 16K, 32K, 48K, or 64K access by another bus master, with the installation of the iSBC 300 Multimodule RAM. In this thesis project all of the 64K-byte memory of the slave iSBC 86/12A is made accessible to the master. Furthermore. both iSBC 86/12A boards are configured to occupy the first 128K section of the 1-megabyte address space.

3. A-to-D Converter Board

This board is electrically and mechanically compatible with any iSBC 86/12A board and with MDS. Both the anolog input and output systems are contained on a single printed circuit board that is treated as ordinary memory locations by the CPU (memory mapping). This board simply gets the analog signals form the potentiometers and converts them to the digital signals compatible with TTL standards. The output of the A-to-D converter is one byte per potentiometer input, which varies from +127 to -128. This is the maximum value range a fixed binary (7) variable can assume in the PL/I language, by definition. So, the



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Figure 2.2 -Interrupt Matrix Interconnection

software system thinks of the potentiometer value as changing between +127 and -128. The A-to-D converter board is configured to occupy the segment D000: (14th 64K portion of the 1-megabyte address space).

4. ADM-3A Terminal

This terminal is the only means by which the operator communicates with the system. It is connected to the system with the master iSBC 86/12A board's serial I/O connector. It is an interactive device which is used to enter, display, and send information to a host computer,

and to receive and display information from that computer. The information exchange between the terminal and the computer is made at different baud rates, ranging up to 19200. In this scheme, a 9600 baud rate is used. The keyboard contains 59 keys. The display memory is a RAM which is capable of holding 1920 characters. Data characters are displayed on 24 or 12 equally-spaced rows, each consisting of 80 columns.

5. RG-512 Retrographics Card

The RG-512 Retrographics printed circuit board is added to the ADM-3A terminal to extend the data and graphics display capabilities with the aid of a Z-80A built-in microprocessor. The RG-512 employs the bit map method of storing graphic images. This information is in a digital memory as a rectangular array of stored bits. Each bit is mapped onto the CRT screen and can cause a bright spot to be displayed. The RG-512 displays graphs and pictures by writing the proper bits into the graphics memory. One of the important features of the RG-512 is the ability to erase portions of the screen selectively. This is desirable when the application requires the use of dynamic displays employing motion or rotation to convey information. The RG-512 has four modes of operation. These are the ADM-3A Alpha Mode, the 4010 Alpha Mode, the Point Mode, and the Vector Mode. The first one is equivalent to the operation of ADM-3A without RG-

512. The latter three modes make use of the bit map method.

6. Potentiometers

Two sets of triplet potentiometers are used as simulating analog sensor information sources. They feed anolog signals to the input of the A-to-D converter board, varying between -5V and +5V.

III. IMPLEMENTATION OF THE SOFTWARE

A. GENERAL INFORMATION

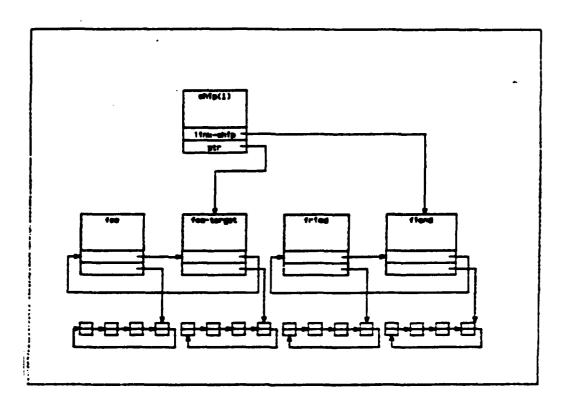
1. Modularity

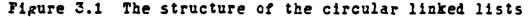
A modular and extensible simulation program is aimed at simplifying the debugging and testing phase and at facilitating possible alterations. The hierarchy of the modularity is composed of a head module and four secondlevel modules connected to the head module. These secondlevel modules are the initialization module. the simulation system module, the real time executive module. and the dynamic debugging tool module. These modules аге separately compiled PL/1-86 and RASM-86 programs. Each main module is further subdivided into third-level modules to gain a finer granularity of modularity. In the programs, two useful special features of PL/I-86 are Those are %INCLUDE and %REPLACE statements. used. Ву those statements, global declarations that are the same in the scope of the simulation program need not be declared within each and every module. Instead, they are grouped together in the GLOBALS.INP declaration file. The %REPLACE statement allows constants to be declared as in the other high-level languages like Pascal, such that the value of the constants can be changed without having to go through every program segment in which they occur.

2. Data Structures

Linear arrays and arrays of structures (records) are used in the simulation program as data structures. These data structures are then linked to each other to establish circular linked lists. Figure 3.1 (on the next page) explains the general picture of the circular linked lists in the simulation program. Fixed size data structures with fixed binary pointers are used in the program, rather than pointer data, to avoid the dynamic system overhead and to retain the benefit of random access capabilities inherent to linear arrays.

The structure SHIP has two pointers, in addition to the fields that hold specifications about the ships in the area. Those pointers, PTR and LINK-SHIP have different purposes for different ships. SHIP (1), for instance, being the ownship points to two different circular linked lists. PTR points to the enemy ship's circular linked list by pointing to the target ship which is engaged (tracked) for the sea battle. The other pointer, LINK-SHIP points to the friendly ships circular list. The PTR field of other ships, on the other hand, points to another circular linked list, WAKE, to record their past positions that will be used for tracking and display purposes, where LINK-SIP points to the other ships in their category. The reason for using circular lists in this program is the ease with





which one traverses through the lists and which does not necessitate the use of another external pointer.

3. Other Features

The PL/I exception handlers (ON body statements) are used extensively in the dynamic debugging module, and in various interactive parts of the program, to intercept the error conditions that might be raised during the testing and execution of the program. The ON body statements are contained in the ERRHAND.AID file, which is an %INCLUDE file. Upon receiving the control through a raised error condition, the statements in the file prompt

the user and give the control over to the REENTRY.PLI interactive debugging tool.

One of the exceptional features of the simulation program is the use of non-local goto statements which are unacceptable in structured style of programming. It is an inevitable requirement, by the PL/I language, to suppress the raised error conditions by a non-local goto statement. It is also used in some parts of the dynamic debugging module, in order to by-pass the flow of control over to the debugging program when the optional ERRORON boolean switch is closed. Explicit comments are offered wherever nonlocal goto statements are used, to avoid confusing the reader.

B. SOFTWARE FUNCTIONAL DESCRIPTION

In the following sections, the structure of the modules and the programs that belong to those modules are described. The program listings are presented as appendices.

1. Head Module

a. WAR.PLI

This main procedure is the head node of the hierarchical structure of the procedures used to modularize and structure the implementation of the simulation program. It contains two call statements, one of which is to the initialization module, to set up the tactical database

and to initialize various external variables that are used throughout the simulation program. The other call is for passing the control to the real time executive module where the control stays for the rest of the program execution. A listing of WAR.PLI is presented as Appendix A.

2. Initialization Module

a. INITVARS.PLI

This PL/I routine, when called from the main procedure, constructs the tactical database in an interactive manner. It first initializes the pool of available SHIP and WAKE modes for later use. It then gets interval of time which is used to update the the information about ships in the tactical area and other time dependent functions. This time interval must be equal to the period of the real-time interrupts which depend on the timing constants used during the initialization of PIT. Detailed information will be presented later under the hardware initialization section. Then the control proceeds to establish the tactical database interactively. In this session, the initial information about azimuth, range, friend or enemy are written to the proper fields. Then, a circular linked list of four nodes is composed and the PTR field is made to point to that list. Finally, dependent on whether friend or foe, that particular node is added to the appropriate circular linked list. INITVARS.PLI makes use of various internal subroutines for linked list

operations. Those subroutines simply extract a node from the pool of available nodes. After establishing the linked lists, the INITVARS.PLI initializes external variables that are used throughout the program. A listing of INITVARS.PLI is presented in Appendix B.

3. Simulation System Module

This module is composed of four PL/I programs which perform the simulation under the control of the real-time executive module. The following subsections describe the functional description of those programs. The listings of the program segments that comprise this module are presented in Appendix C.

a. TACTICAL.PLI

This routine has the highest priority among the system module programs. It first updates the position of each ship in the tactical area by calculating the relative velocity and multiplying that with the time interval, which is the period of timing interrupts that occur every 250 milliseconds. The control then proceeds to calculate the future positions of the ships, for those ships which have been in the area for more than 4 seconds, and which are included in the enemy ship circular linked list. The routine uses the polynomial least squares curve fitting method with Legendre Polynomials. The coefficients are pre-calulated for the position of the ship one second after the time of calculations, based on the past four wake

points. TACTICAL.PLI also calculates the trajectory of the travelling projectile if the gun is fired.

b. DISPLAY.PLI

This routine, which is invoked every second, is the interface of the simulation program for the ADM-3A screen. It simply traverses the circular linked lists and generates the appropriate display objects for the ships, their aim points which the own ship aims at, and the gun aim point. It also displays the travelling projectile if it is fired by the system. The routines for generating the objects are internal for the DISPLAY.PLI. Another routine, TRANSLATE, translates the cartesian coordinates to the stream of characters that represent the x and y grid coordinates of the RG-512 Retrographic Screen Memory. Finally, the internal routine DRAW puts the generated objects on the screen in vector mode. According to the key variable, D, it either sets the data level to white and puts the object on the screen, or sets the data level to black and makes selective erasures.

c. STATUS.PLI

This routine, which is invoked at every second, is the interface of the system to the ADM-3A video terminal keyboard. The commands for the system are read from the keyboard by calling the serial I/O chip (USART) interface assembly program, KEYBOARD, which will be presented under Miscellaneous Assembly Routines.

The are four boolean variables used in the simulation program. Those are ENGAGED, MAGNIFIED, FIRED. ERRORON. The STATUS.PLI sets these variables and according to the value of the parameter passed to the assembly interface routine. ENGAGED (with "E") shows if the system is engaged to any target (for TACTICAL.FLI). MAGNIFIED (with "M"), which is used in DISPLAY.PLI, shows which display scale is being used and what the reference point of the display is. Usually, the display on the screen is relative to own ship. But that can be changed so that the ship engaged is at the center to the screen by setting MAGNIFIED true. FIRED (with "F") boolean variable is used to commence the ballistic calculations and display. It is used both in TACTICAL.PLI and DISPLAY.PLI. Finally, ERRORON (with "D") is used to transfer the control to the dynamic debugging module.

d. IDLE.PLI

This program is the idle routine for the system, as it waits for a 250 millisecond timing interrupt to occur. It is the interface program for the six potentiometers which are used as sensors for the own ship velocity vector, the known ship velocity vector, and the gun elevation and bearing. The velocity vectors are composed of speed and course components. The IDLE.PLI gets this information by calling the Analog-to-Digital Converter interface assembly program, ATOD. The control then

proceeds to convert this information, which is in the fixed binary (7) form in range (-128, +127), to appropriate coordinate values, e.g. 50 knots maximum speed and the true azimuth between 0 and 360 degrees. This routine calculates own ship velocity components in cartesian coordinates for later use by TACTICAL.PLI. It also makes the initial ballistic computations for the gun.

4. Real Time Executive Model

This module works as the interrupt handler for the real time interrupts that are initiated by programmable hardware components every 250 milliseconds. It is invoked by the WAR.PLI the first time and interrupts thereafter. It responds to the timing interrupts, which tell the system that data must be collected at this point in time. The module then resolves the priorities of the simulation system module programs and arbitrates the flow of control during the execution of the system. The real time module makes use of the operating executive system primitives, which are presented in the following subsections. The labels P1 through P4 are associated with the simulation system modules TACTICAL.PLI, DISPLAY.PLI, STATUS.PLI, and IDLE.PLI, respectively. The listing of the programs included in the Real Time Module are presented in Appendix D.

a. ARBITER.A86

This assembly language program is the real "workhorse" of the entire system. It first allocates stack areas for four simulation system during the assembly time. Upon invocation by the main procedure, it initializes the programmable hardware components and transfers the control to the P4.PLI process, which in turn calls IDLE.PLI repetitively until the first timing interrupt occurs. The interrupt entry point PROCØ, where the process switching starts, is entered by the interrupt software. At this point of execution. the external variable fourthevc. which signals the system that 250 millisecond event has occured, is updated to the new value by incrementing it by one. After storing the state of the program which is interrupted during its execution, ARBITER.A86 invokes SCHEDULE.PLI to obtain the name of the ready program that has the highest priority. If there are none, the interrupted program is resumed. If there is any ready process of higher priority, then ARBITER.A86 loads the process state and gives the control over to it. During this process switching, the upper boundaries of the stack areas are checked for a rossible stack overflow, which could happen if the time interval was not sufficiently large for the system module routines to finish execution before the next timing interrupt comes. ARBITER.A86 has a second entry point, STORESTATUS, for the synchronization primitive AWAIT.PLI to

enter when the correct number of interrupts for the calling synchronization primitive have not yet occured.

b. AWAIT. PLI

This synchronization primitive is invoked as as an operating system primitive, by any process, P1 through P3. AWAIT compares the value of the external variable FOURTHEVC to the threshold value of the calling process to see if it is greater or equal to the value at which the process is to proceed. If not, then it calls the STORESTATUS entry of ARBITER.A86 to relinquish the control to the awaiting ready process, or to the P4.PLI that calls the IDLE.PLI simulation system program as the system idling routine.

c. SCHEDULE.PLI

This synchronization primitive is called by ARBITER.A86 to return the name of the highest priority ready process. It does that simply by identifying the first ready process on the list. Because the scheduler scans the list in the descending priority order, the highest priority ready process will automatically be scheduled.

d. THRESH.PLI

This routine, when invoked by processes P1 through P3, increments the corresponding thresholds in an external one dimensional array called THRESHOLD. This table is used by the AWAIT.PLI and SCHEDULE.PLI

synchronization primitives to decide whether or not a process is ready for execution.

e. P1.PLI

This process is basically an infinite loop. Within this loop, there are three subroutine calls. The routine first makes a call to AWAIT.PLI to see if it is time for it to proceed. If not, the control doesn't come back again; instead, the current state of the process is stored by ARBITER.A86 and the highest priority ready process is executed. If it is the time, the control proceeds to call the simulation system module program TACTICAL.PLI. After that, a call to THRESH.PLI is made, where the threshold value that is allocated to P1.PLI is incremented by the proper value. When the infinite do loop repeats itself. the call to the AWAIT.PLI will indicate that the process TACTICAL.PLI is not yet ready for execution and control is transfered to the highest priority ready process.

f. P2.PLI

This process is identical to P1.PLI in form except the call is to DISPLAY.PLI instead of to TACTICAL.PLI.

g. P3.PLI

This process is identical to P1.PLI in code except the call is to STATUS.PLI.

h. P4.PLI

This process is similar to P1 through P3 in the structure described above. There is only one call in the infinite do loop, which is to IDLE.PLI. This routine is always ready for execution and, basically, repeats itself until the next timing interrupt comes along.

5. Miscellaneous Assembly Routines

There are few machine dependent functions that cannot be accomplished by the high-level language PL/I-86. Assembly routines were written to interface the PLI-86 programs with the hardware of the 8086 microprocessor. These assembly routines are included within the main body ARBITER.A86. There are two parameter passing of conventions from PL/I-86 to the assembly language routines. In the first one, there is only one argument passed in the accumulator. as in a function call. In the subroutine calls. which is the case here, the address of the VECTOR that contains the pointers to the actual parameters is passed in the BX register. The following subsections give some descriptions about those assembly routines. Appendix E shows the listings of the modules.

a. KEYBOARD.A86

This routine is invoked by STATUS.PLI to read the keyboard. It is written so that the keyboard status is read to see if a key had been pressed instead of waiting indefinitely until a key was depressed, as would be the

case had the PL/I get statement been used. The program first reads in the status of the serial I/O interface chip (USART) to see if a character has been received from the keyboard. If it has, then the character is read and placed into a corresponding character variable, which is the formal parameter in the subroutive invocation. If there is no character received since the last attempt to read, the ASCII equivalent of 0 is put in key. The reason for that is that the character 0 is not teing used as a command.

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b. ATOD.486

This assembly routine is called by IDLE.PLI to read the first six Analog-to-Digital Converter Board outputs. The reason to write this assembly routine is that the Analog-to-Digital Converter ports are memory mapped to be in the segment DØØØH. The PL/I function UNSPEC works for those memory locations which are included within the first 64K bytes of memory. The assembly routine sets the proper segment and source index registers to point to those locations and makes an ordinary read operation. This value is then put in the formal parameter passed to the PLI-86 routine.

c. RINGBELL.A86

This assembly routine simply sends a bell character to the video terminal and causes a bell sound to ring. This is equivalent to sending a control G in PL/I-86.

The only difference is that it can be used in other assembly routines.

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d. WAIT.A86

This assembly routine reads in the status of the I/O interface chip and waits until the transmitter buffer is empty; i.e. the character which had been in the buffer is received by the video terminal and an acknowledgement signal is sent back to the interface chip. This routine is used by assembly routines that put out a message.

e. SUSPEND.A86

This routine simply resets the interrupt bit of the program status word (PSW) to disable the 8086 CPU from acknowledging the interrupts. It is used by the dynamic debugging system to stop the real time clock.

f. RESUME.A86

This routine first sets the interrupt flag to enable the 8086 CPU to respond to the interrupt requests. It also resets the PIT that is used to generate the timing clock pulses for the PIC.

D. INITIALIZATION OF THE PROGRAMMABLE COMPONENTS

The iSBC 86/12A board has three programmable hardware components, which were described in Chapter II. In the following subsections, the initialization sequences for

those hardware components which produce the real time synchronization are described.

1. USART Programming

The 8251A USART converts parallel output data into virtually any serial output data format. The USART also converts serial input data into the parallel data format. Prior to starting to transmit or to receive data, the USART must be loaded with a set of control words. These control words, which define the complete functional operation of the USART, must immediately follow a reset (internal or external). There are two types of control words, namely, a Mode instruction and a Command instruction. Once the Mode instruction has been sent, the Command instruction can be sent at any time prior to a read/write operation. The following assembly code is used to initialize the USART read/write mode:

MOV AL,37H OUT ØØDAH,AL

During the course of execution, the serial I/O interface routine executes the following assembly code to read the USART status and to read the receiver buffer if any character has been received:

IN	AL,ØØDAH
AND	AL,02H
JZ	KEYBOARD1
IN	AL,00DEH

2. PIT Programming

The 8253 PIT has three independent counters. They are Counter 0, Counter 1, And Counter 2. The input clock frequency is 22.1184 MHz supplied by a crystal oscillator. The input clock frequencies can be adjusted via jumpers. In this project, Counter 1, with the default factory jumper connection (E59-E60), is used. The input frequency is 153.6 KHz. It is chosen to work in mode 0 so that it gives an interrupt on terminal count. The formula:

N = TC

where

N is the count value for the counter T is the desired interrupt time interval in seconds

C is the internal frequency (Hz)

From the above formula, the count number for the counter 1 is found to be 38400 decimal (9600 her). Since it is not possible to express this number in four decimal digits, the binary count mode is selected. The counter is initialized by sending a mode control word, followed by a down-count number. Once the counter is initialized, sending the down-count number resets it to the start condition. The following sequence of code is used to program the PIT:

MOV	AL,5ØH
OUT	ØØDEH,AL
MOV	AL,ØØH
OUT	ØØD2H,AL
MOV	AL,6ØH
OUT	ØØD6H,AL
MOV	AL,96H
out	ØØD2H,AL

3. PIC Programming

The 8259A PIC is programmed in the nested mode. The master PIC with no slaves is accepted. For this particular situation, the initialization words 1,2, and 4 are The initialization word 2 is set to represent the sent. interrupt vector address. This is the address that the control is given when the interrupt 1 occurs (04H). An interrupt mask byte is used to mask out the irrelevant interrupts for the purpose of this thesis. The PIC can be reset after each interrupt simply by sending the EOI (end of interrupt) status byte to the appropriate address. The initialization sequence is as follows:

> CLI MOV AL,13H OUT ØØC2H,AL MOV AL,20H OUT ØØC2H STI MOV AL,ØFDH OUT ØØC2H,AL

To reset the PIC, MOV AL,20H OUT 00C0,AL

E. ASSEMBLY, COMPILING AND LINKING

The assembly language code was written in RASM-86 and assembled by using the RASM-86 Assembler. This assembler produces relocatable files that can then be linked with other separately compiled or assembled object files by Link This linker accepts three types of input files. Those 86. are the object file, library file, and/or an input file. Input files are very useful tools in that they include input command lines by an input file instead of writing the command line each time the programs are to be linked. The PL/I-86 compiler is used to compile PL/I programs. This compiler requires a 128 Kbyte RAM, as opposed to the PL/I-80 compiler, which requires 48K RAM.

F. TESTING

The hierarchical simulation program modules were designed and tested in a top-down manner. An extensive dynamic debugging module was used in the testing of individual modules (see Appendix F). The testing phase is first started with writing a skeletal model for the real time executive module. The PL/I-86 output statements (stubs), which printed some appropriate numbers, were inserted in the places where the simulation system module programs were invoked. The DDT86 (Dynamic Debugging Tool) was used to test and debug this skeleton program. Since it was a real-time interrupt driven program, the interrupt

enable bit (I) of the program status word was reset to zero (0) to control the flow of the program. Some error checkings are inserted in the main assembly language program. ARBITER.A86, to see whether the process switching was done correctly, or to see if any of the stacks allocated to the synchronization primitives have overflowed.

After testing the real-time executive model, the second phase was the testing of the subroutines that are used by the simulation system module programs. Appendix G shows an example for testing a subroutine.

The third phase of the testing was to test and debug the simulation system program. The dynamic debugging module was used for this purpose. This module is composed of FL/I-86 %INCLUDE files and external PL/I-86 programs. The %INCLUDE files are inserted into the various parts of the program being tested. The code of this debugging system is bordered with comment lines from the main body of the program it tests. It is not visible to the program, i.e., it brings its local and global variables with the LOCALS.AID declaration %INCLUDE file. It is possible to manipulate the system's external variables through the debugging system module, IDLE.PLI.

The PL/I-80 had been used in the early stages of the testing phases because of the existence of redundant Intel 8080 based systems in the Naval Postgraduate Micro-Lab.

But in the later stages, it was realized that this had caused some problems, due to some incompatibilities between the PL/I-80 and PL/I-86 systems. These deficiencies of the PL/I-86 have necessitated the testing and debugging of the programs to be done in the PL/I-86 based systems. The need for such a thorough debugging system module was then realized.

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IV. CONCLUSIONS

original objectives of the thesis have been The large extent. accomplished. to а The hardware interconnections intended to promote real time clock operations have been successful. A test model, which comprised the real time executive module and the PL/I put statements, instead of the simulation system modules, was developed to check the system's operation and timing. Correct results appeared on the screen. Then, the testing and debugging of the simulation system programs showed that their algorithms and operations were correct, with the exception of the simulation system program, DISPLAY.PLI. The testing of DISPLAY.PLI showed that the objects to be displayed on the video screen were not successfully put on the screen. One error, which was an automatic conversion error, was found in the routine TRANSLATE and corrected by using a step variable. However, an error still exists in the body of the procedure DRAW.

Since the testing phase of the DISPLAY.PLI has not been accomplished, the objective of synchronizing two iSBC 86/12A boards could not be accomplished. The intended purpose for this objective was to make the second iSBC 86/12A jump to a waiting loop with the initial power start up interrupt (reset), to load the assigned simulation program segment to the common RAM, and to direct it to the

beginning of that program segment, through a bus vectored interrupt.

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A hierarchical and modular program model was constructed through the use of data structures and structured programming. The modularization allows possible future changes to the programs.

The simulation system program constructed in this thesis could be a basis for further enhancements towards a complete fire control system, or a related tactical simulation system, due to its modularized nature.

The dynamic debugging module designed to test and debug the simulation system modules can be used for the purpose of testing any other programs, simply by changing the names and formats of the variables that the debugging system manipulates. The advantage of using the %REPLACE and %INCLUDE pre-processor statements, which are peculiar to the PL/I-86 version of the subset G, makes such an implementation feasible.

APPENDIX A

HEAD MODULE PROGRAM LISTINGS

A. WAR.PLI

/* Prog Name WAR.PLI : December 83 Date : Written by : M. Kadri Ozyurt For Thesis . Advisor Professor Kodres . This is the main procedure of the Purpose : It invokes the initialization modular simulation program. module to set up the target database and to initialize the external variables used throughout the simulation program. */

WAR: PROCEDURE OPTIONS (MAIN);

/*external procedures*/

DCL

(INITVARS, ARBITER) ENTRY;

/*this call to the initialization module initializes the simulation system*/

CALL INITVARS;

/*this call gives the control over to the real time
executive module*/

CALL ARBITER;

END WAR;

B. GLOBALS.INP

/*
Prog Name : GLOBALS.INP
Date : December 83
Written by : M. Kadri Ozyurt
For : Thesis
Advisor : Professor Kodres
Purpose : This %include file contains the declarations of the global variables used in the simulation program
*/

DCL

(COUNTER, SECONDS, MINUTES, HOURS, WAKE PTR, SHIP_PTR, AVAILSHIP, AVAILWAKE, P,Q, NUMBERSHIPS, NODE, TARGET, KNOWN) FIXED BIN(7) EXTERNAL, FOURTHEVC FIXED BIN(15) EXTERNAL,

D FIXED BIN(7) EXTERNAL, (DT,T PRIME,T OF,T) FLOAT EXTERNAL, (I,J,XX,YY) FIXED BIN(15), CURRENTPROC FIXED BIN(7) EXTERNAL,

(ENGAGED, MAGNIFIED, FIRED, ERRORON, DONE) BIT(1) EXTERNAL,

KEY CHARACTER (1) EXTERNAL, THRESHOLD(0:2) FIXED BIN(15) EXTERNAL, ARG(0:5) FIXED BIN(7) EXTERNAL,

(VX_OWN,VY_OWN,VX_TARGET,VY_TARGET,VX_REL, VY_REL,VX_ROUND,VY_ROUND,VR) FIXED DECIMAL EXTERNAL,

ALPHA FIXED DECIMAL EXTERNAL ,

(AX_SUM,BX_SUM,CX_SUM,AY_SUM,BY_SUM,CY_SUM, AX,BX,CX,AY,BY,CY,X_AT5,Y_AT5,R,DX_DT_AT5, DY_DT_AT5,DR_DT_AT5,X_OFFSET,Y_OFFSET,M)FIXED_DECIMAL EXTERNAL.

(01.02)(5) FIXED DECIMAL EXTERNAL .

1 SHIP(MAX_SHIPS) EXTERNAL, 2 VELOCITY, 3 COURSE FIXED DECIMAL INIT(0.0), 3 SPEED FIXED DECIMAL INIT(0.0), 2 POSITION, 3 AZIMUTH FIXED DECIMAL INIT(0.0), 3 RANGE FIXED DECIMAL INIT(0.0).

2 COORDINATES, 3 X FIXED DECIMAL INIT(0.0),3 Y FIXED DECIMAL INIT(0.0), 2 AIM, **3 X AIM FIXED DECIMAL** $INIT(\emptyset.\emptyset)$. 3 Y AIM FIXED DECIMAL INIT(Ø.Ø). 2 COUNT FIXED BIN (7) INIT(0), 2 NUMBER FIXED BIN (7) INIT(0), 2 PTR FIXED BIN (7) INIT(0). 2 LINK_SHIP FIXED BIN (7) INIT(0), 2 FRIEND BIT(1) INIT(FALSE), 1 OBJECT(MAX SHIPS) EXTERNAL. 2 LOCATIONS. 3 U (0:10) FIXED BIN(15) INIT((11) -1), 3 V (0:10) FIXED BIN(15) INIT((11) -1), 2 AIMS. 3 U AIM (0:10) FIXED BIN(15) INIT((11) -1), 3 V AIM (0:10) FIXED BIN(15) INIT((11) -1), 2 GUN, 3 U GUN (0:10) FIXED BIN(15) INIT((11) -1), 3 V_GUN (0:10) FIXED BIN(15) INIT((11) -1), 2 WAKES, 3 U_WAKE (0:10) FIXED BIN(15) INIT((11) -1) 3 UU_WAKE (0:10) FIXED BIN(15) INIT((11) -1), 3 V WAKE (0:10) FIXED BIN(15) INIT((11) -1), 3 VV WAKE (0:10) FIXED BIN(15) INIT((11) -1). 1 GUN EXTERNAL. 2 POSITION. 3 AZ FIXED DECIMAL INIT(0.0). 3 ALT FIXED DECIMAL INIT(0.0). 2 COORDINATES, 3 X GUN FIXED DECIMAL INIT(0.0), 3 Y GUN FIXED DECIMAL INIT(0.0). 1 WAKE(4) EXTERNAL. 2 COORDINATES. 3 X WAKE FIXED DECIMAL $INIT(\emptyset.\emptyset),$ 3 Y WAKE FIXED DECIMAL INIT(0.0), 2 LINK WAKE FIXED BIN(7) INIT(0). (SUSPEND, RESUME, ARBITER, INITVARS) ENTRY, KEYBOARD ENTRY (CHARACTER(1)),

ATOD ENTRY (FIXED BIN(?), FIXED BIN(?));

С. CONST.INP

/*		
Prog Name	: CONST.INP	
Date	: December 83	
For	: Thesis	
Advisor	: Professor Kodres	
Purpose	: This %include file contains the constant	
declarations */	used throughout the program.	

%REPLACE /*max number of wake nodes*/
/*max number of ship nodes*/ MAX_WAKE BY 4, MAX SHIPS BY 2, /*max number of variables*/ MAXVARS BY 82, OWN BY 1, RMAX BY 25000.0, /*ownship indicator*/ /*max gun range*/ /*max argument for SQRT */ MAXSQ BY 1.073E+09. /*max number for fixed (15)*/ TOP BY 32767.0, /*gravitational con.yd/sec2*/ G BY 10.7246, VM BY 518.0. /*muzzle velocity*/ /*x coord. for center*/ A BY 512, /*y coord. for center*/ B BY 390, K BY 1.40625, L BY 4.513, /*azimuth proportionality c.*/ /*speed proportionality c.*/ /*definition for 360 degree*/ TWO PI BY 360.0, /*definition of pi rad/180deg*/ PI BY 3.1416, 00 BY 1.0, /*legendre poly. of zero deg.*/ TRUE BY '1'B, /*boolean true*/ FALSE BY 'Ø'B, /*boolean false*/ CLEAR SCREEN BY '`\`[`L', /*char. sequence for retro.*/ CLEAR ALPHA BY '2', /* " " " " " " " '' /*in linked list terminology*/ /* Ţ, VECMOD BY **′/*** " ** POINTMOD BY /* " ... ALPHA4010 BY /* ALPHAJA BY WHITEMOD BY [ā', ['; /* 10 /* BLACKMOD BY

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*/

APPENDIX B

INITIALIZATION MODULE PROGRAM LISTINGS

A. INITVARS.PLI

```
/*
                  : INITVARS.PLI
Prog Name
                  : December 83
Date
Written by
                  : M. Kadri Ozyurt
                  : Thesis
For
                  : Professor Kodres
Advisor
                  : This routine prompts the user to give
Purpose
the time interval (dt) and constructs the tactical circular
linked list in an interactive manner. The control then
proceeds to initialize the external variables.
initvars:proc external;
/*
  dcl
*/
          %include'const.inp';
          %include'globals.inp';
  /*this iterative loop initializes the pool of available
wake nodes*/
  do i=1 to max_wake-1;
          link_wake(i)=i+1;
  end;
  link wake(max_wake)=nil;
/*this sequence initializes the pool of available ship
nodes*/
  do i=1 to max ships-1;
          link_ship(i)=i+1;
  end;
  link_ship(max_ships)=nil;
  put skip list ('Enter the time interval (dt) in seconds');
```

```
/* the following block is an interceptor for a too large
input value*/
 on overflow begin;
put list('*** too large, try again');
          goto init1;
 end;
init1:
 put skip list('>');
 get list(dt);
 revert overflow;
 put skip list ('Construction of the tactical database');
 numberships=0;
 done=false;
/* this procedure call gets a ship node from the pool of
available ships and assigns its address to ship pointer*/
  ship ptr=getship();
 do while ( done);
          put skip list('Enter the position of ', numberships,
                   'th ship in true azimuth and in yards');
/*the following
                  two on condition bodies are for the input
line at init2*/
          on error begin;
                  put list('*** bad value.try again');
                  goto init2;
          end;
          on fixedoverflow begin;
                  put list('*** too large, try again');
                  goto init2;
          end;
init2:
          put skip list('>');
          get list(azimuth(ship_ptr),range(ship_ptr));
          revert error;
          revert fixedoverflow;
          put skip list('Friend or foe (F/E)?');
          put skip list('>');
          get list(key);
/* the following sequence adds the ship node to the
appropriate circular linked list, friend ships or enemy
ships, according to the friend boolean value entered*/
          if (key='I')!(key='f') then do;
               friend(ship_ptr)=true;
              if link_ship(own)=nil then
                  link_ship(ship_ptr)=ship_ptr;
              else
                  link_ship(ship_ptr)=link_ship(own);
              /*end if#/
```

```
link_ship(own)=ship_ptr;
end /*do*/;
else do;
  friend(ship_ptr)=false;
    if ptr(own)=nil then
        link_ship(ship_ptr)=ship_ptr;
    else
        link_ship(ship_ptr)=ptr(own);
    /*end if*/
    ptr(own)=ship_ptr;
end /*if*/;
```

/* the following procedure call and the iterative loop get available nodes from the pool, construct a circular linked list of four nodes, and assign the address of the list to the ptr pcinter of the ship node which is being constructed*/

wake_ptr=getwake(); q=wake_ptr; do i=1 to 4; p=wake_ptr; wake_ptr=getwake(); link_wake(p)=wake_ptr; end /*do*/; link_wake(wake_ptr)=q; ptr(ship_ptr)=wake_ptr; number(ship_ptr)=numberships; put skip list ('Would you like to enter another 'sbir (Y/N)?'); put skip list('>'); get list(key); if (key='Y')!(key='y') then do; ship_ptr=getship(); if ship_ptr=nil then done=true; /*end if*/ end; else done=true; /*end if*/ end /*do*/; seconds=0; minutes=0; hours=0; wake ptr=0; target=0; t_prime=0.0; t_of=0.0;

t=0.0;

```
link_ship(own)=ship_ptr;
end /*do*/;
else do;
friend(ship_ptr)=false;
if ptr(own)=nil then
link_ship(ship_ptr)=ship_ptr;
else
link_ship(ship_ptr)=ptr(own);
/*end if*/
ptr(own)=ship_ptr;
end /*if*/;
```

/* the following procedure call and the iterative loop get available nodes from the pool, construct a circular linked list of four nodes, and assign the address of the list to the ptr pointer of the ship node which is being constructed*/

```
wake_ptr=getwake();
        g=wake_ptr;
        do i=1 to 4;
                p=wake_ptr;
                wake_ptr=getwake();
                link wake(p)=wake_ptr;
        end /*do*/;
        link_wake(wake_ptr)=q;
        ptr(ship_ptr)=wake_ptr;
        number(ship_ptr)=numberships;
        put skip list ('Would you like to enter another ',
                          ship (Y/N)?');
        put skip list('>');
        get list(key);
        if (key='Y')!(key='y') then do;
                ship_ptr=getship();
                if ship_ptr=nil then
                         done=true;
                 /*end if*/
                end:
        else
                done=true;
        /*end if*/
end /*do*/;
seconds=0;
minutes=0;
hours=0;
wake_ptr=0;
target=0;
t prime=0.0;
t_of=0.0;
```

t=0.0;

friend(own)=true; known=2; /* engaged ship no by the sensor */ x_offset=0.0; y_offset=0.0; m=50.0; fourthevc=0; 1=0; j=0; currentproc=4; engaged=false; magnified=false; fired=false; erroron=false; key='0'; threshold(\emptyset)=1; threshold(1)=4; threshold(2)=4; **vx** own=0.0; vy_own=0.0; vi target=0.0; vy_target=0.0; vx rel=0.0; vy_rel=0.0; **vx** round=0.0; vy_round=0.0; vr=0.0; alpha=0.0; ax_sum=0.0; bx_sum=0.0; cx_sum=0.0; ay_sum=0.0; by_sum=0.0; cy_sum=0.0; ar=0.0;tr=0.0;cr=0.0; ay=0.0;by=0.0;cy=0.0; x_at5=0.0;y_at5=0.0; r=0.0; $dx_dt_at5=0.0$; $dy_dt_at5=0.0$; $dr_dt_at5=0.0$; o1(1)=1.0; o1(2)=0.5; o1(3)=0.0; o1(4)=-0.5; o1(5)=-1.0; 02(1)=1.0; 02(2)=-0.5; 02(3)=-1.0; 02(4)=-0.5; 02(5)=1.0;/*getship. when invoked. extracts a node from the pool of available and returns a pointer value pointing to that node. It puts an error message if there is no available node*/ getship:procedure returns (fixed bin(7)); /* dcl */ %include 'globals.inp'; if availship=nil then do; put skip list('No more available ship nodes'); return(nil); end /*do*/; else do; node=availship; availship=link_ship(availship); end /*if*/;

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```
return(node);
end getship;
/*getwake does the same function as getship except for the
operations made are on wake nodes*/
getwake:procedure returns (fixed bin(7));
   /* dcl */
            %include 'globals.inp';
   if availwake=nil then do;
        put skip list('No more available wake nodes');
             return (Ø);
            end /*do*/;
   else do;
            node=availwake;
            availwake=link_wake(availwake);
   end /*if*/;
   return(node);
end getwake;
end initvars;
```

APPENDIX C

SIMULATION SYSTEM MODULE PROGRAM LISTINGS

and the second
A. TACTICAL.PLI

/* Prog Name : TACTICAL.PLI Date : December 83 Written by : M. Kadri Ozyurt For : Thesis Advisor : Professor Kodres : This external routine calculates and Purpose updates the positions of the ships in the tactical area and the future positions of the ships that belong to enemy ships circular linked list pointed to by ptr(own), and and calculates the trajectory of the travelling projectile if fired */

TACTICAL: PROCEDURE EXTERNAL;

/* DCL */ %INCLUDE 'CONST.INP'; %INCLUDE 'GLOBALS.INP';

/*following sequence of code updates the present positions
of the ships in the tactical area*/
DO I=2 TO NUMBERSHIPS;
 VI_TARGET = SPEED(I) * SIND (COURSE(I));
 VI_TARGET = SPEED(I) * COSD (COURSE(I));
 VI_REL = VX_TARGET - VX_OWN;
 VY_REL = VY_TARGET - VY_OWN;
 WAKE_PTR = PTR(I);

```
X WAKE(WAKE PTR) = X(I);
       Y WAKE(WAKE PTR) = Y(I);
X(I) = X(I) + VX_REL * DT;
       Y(I) = Y(I) + VY REL * DT;
                                          /*ptr(i) points to*/
       PTR(I) = LINK WAKE(WAKE PTR);
       IF COUNT(I) < 4 THEN
                                            /* the oldest wake*/
            COUNT(I)=COUNT(I)+1;
       /*END IF*/
  END /*DO*/;
/*calculating the future positions (aim
                                                    points)
                                                               starts
here by using the least squares method with legendre
polynomials. The coefficients are pre-calculated according
to the fifth second including the zeroth second*/
FILTERING:
  TARGET = PTR(OWN);
  DONE=TRUE;
  IF
       (TARGET=NIL) THEN
            DONE=FALSE;
  /*END IF*/
  DO WHILE ( DONE);
       IF COUNT(TARGET)=4 THEN
         BEGIN;
            AX SUM = X(TARGET) * 00;
                                            /*Leg. poly. Ø deg.*/
            BX_SUM = X(TARGET) * 01(5);/*
                                                                  */
                                                           1
           CX_SUM = X(TARGET) * 02(5);/*
                                                               ,,
                                                       ..
                                                           2
                                                                  */
            AY_SUM = Y(TARGET) * OØ;
            BY SUM = Y(TARGET) * 01(5);
           CY SUM = Y(TARGET) + 02(5);
            WA\overline{K}E_PTR = PTR(TARGET);
            J = \overline{1};
            DO WHILE (LINK WAKE(WAKE PTR) ~= PTR(TARGET));
                AI_SUM = A\overline{I}_SUM + I_WAKE(WAKE_PTR) * 00;
                BX_SUM = BX_SUM + X_WAKE(WAKE_PTR) + O1(J);
                CI_SUM = CI_SUM + I_WAKE(WAKE PTR) + O2(J);
                AY_SUM = AY_SUM + Y_WAKE(WAKE PTR) * OØ;
BY_SUM = BY_SUM + Y_WAKE(WAKE PTR) * O1(J);
CY_SUM = CY_SUM + Y_WAKE(WAKE PTR) * O2(J);
                WAKE_PTR=LINK_WAKE(WAKE PTR);
                \mathbf{J} = \mathbf{J} + \mathbf{1};
            END /*DO*/;
            AX = AX_SUM / 5.;
            BI = 2.0 * BI SUM / 5.;
CI = 2.0 * CI_SUM / 7.;
            AY = AY_SUM / 5.;
            BY = 2.0 * BY SUM / 5.;
            CI = 2.0 * CI SUM / 7.;
```

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```
X AT5 = AX - 1.5*BX + 3.5*CX;
            Y AT5 = AY - 1.5*BY + 3.5*CY;
INCONVENIANCE:
/*this
         begin block is inserted to avoid the complications
which might arise from the automatic conversions*/
          BEGIN;
               DĊL
                      (XSTEP, YSTEP) FIXED,
                      (IF.YF.RSQD) FLOAT;
                IF (ABS(X AT5)>TOP) (ABS(Y AT5)>TOP) THEN
                       RSQD = -1.0;
                 ELSE DO;
                       XSTEP=BINARY (X_AT5);
YSTEP=BINARY (Y_AT5);
                       XF=FLOAT(XSTEP);
                       YF=FLOAT(YSTEP);
                       RSOD=XF*XF+YF*YF;
                 END /*IF*/;
                IF (RSQD<0.0)!(RSQD>MAXSQ) THEN
                       R=0.0;
                ELSE
                       R = SQRT(RSQD);
                  /*END IF*/
            END INCONVENIANCE;
            IF (R=0.0) ! (R>RMAX) THEN DO;
                       \mathbf{X} AIM(TARGET)=0.0;
                   Y AIM(TARGET)=0.0;
END 7*DO*/;
            ELSE DO;
                       ALPHA = ASIN(G*R/VM**2) / 2.0;/*IN RADS*/
                       VR = VM * COS(ALPHA);
                       T PRIME = R / VR;
D\bar{X}_DT_AT5 = 3.0*CX - 0.5*BX;
                       DY DT AT5 = 3.0 \times CY - 0.5 \times BY;
                       DR DT AT5 = (X AT5 * DX DT AT5 +
                       \mathbf{Y}_{AT5} * \mathbf{D}\mathbf{Y}_{DT} \mathbf{A}\mathbf{T}5) / \mathbf{R};\mathbf{T}_{OF} = (\mathbf{R} + \mathbf{D}\mathbf{R}_{DT}_{AT5} * \mathbf{T}_{PRIME}) / \mathbf{V}\mathbf{R};
                       X AIM(TARGET) = X AT5 + DX DT AT5 * T OF;
                       Y\_AIM(TARGET) = Y\_AT5 + DY\_DT\_AT5 * T_OF;
            END /*IF*7;
       END /*IF*/;
       TARGET=LINK SHIP(TARGET);
                                                  /*next target?*/
       IF TARGET=PTR(OWN) THEN
             DONE=TRUE;
        /*END IF*/
  END /*DO*/;
```

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```
ROUNDTRACK:

/*the ballistic calculations start here*/

IF (FIRED) THEN

BEGIN;

VX_REL = VX_ROUND - VX_OWN;

VY_REL = VY_ROUND - VY_OWN;

X_GUN = X_GUN + VX_REL * DT;

Y_GUN = Y_GUN + VY_REL * DT;

T = T - DT;

IF T<=0 THEN DO;

PUT LIST('G');

FIRED = FALSE;

END /*IF*/;

END /*IF*/;
```

END TACTICAL;

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B. DISPLAY.PLI

/* Prog Name : DISPLAY.PLI Date : December 83 Written by : M. Kadri Ozyurt For : Thesis Advisor : Professor Kodres Purpose This routine first puts the time in : hours, minutes, and seconds. The control then proceeds to call DRAW subroutine in an iterative loop to erase the old objects. It then calculates the positions of the objects relative to either ownship or the ship that has been targeted according to MAGNIFIED.Then it calls DRAW to display the objects. */

DISPLAY: PROCEDURE EXTERNAL;

```
/*DCL*/
%INCLUDE 'CONST.INP';
%INCLUDE 'GLOBALS.INP';
```

```
PUT LIST('^] `I'); /*ENTER ALPHA MODE*/
IF MINUTES=Ø THEN
BEGIN;
PUT LIST ('^[= %');
PUT EDIT (HOURS)(F(2));
HOURS = HOURS+1;
IF HOURS=24 THEN
HOURS = Ø;
/*END IF*/
END /*IF*/;
IF SECONDS=Ø THEN
```

```
BEGIN;

    PUT LIST (`^[= (`);

    PUT EDIT (MINUTES)(F(2));

    MINUTES = MINUTES+1;

    IF MINUTES=60 THEN

        MINUTES = 0;

    /*END IF*/

END /*IF*/;

PUT LIST (`^[= +`);
```

```
PUT EDIT (SECONDS)(F(2));
```

```
SECONDS = SECONDS+1;
  IF SECONDS=60 THEN
      SECONDS = \emptyset;
  /*END IF*/
  PUT LIST( 'M'I'');
                                    /*HOME CURSOR*/
/*the following calls erase the objects from the screen*/
  D=0;
  DO I=1 TO NUMBERSHIPS;
          CALL DRAW(OBJECT(I).U,OBJECT(I).V,D);
          CALL DRAW(OBJECT(I).UU WAKE.OBJECT(I).VV WAKE.D);
          CALL DRAW(OBJECT(I).U AIM.OBJECT(I).V AIM.D);
  END /*DO*/;
  CALL DRAW(OBJECT(1).U_GUN,OBJECT(1).V GUN.D);
/*the
       following sequence converts the coordinates
                                                       of
                                                           the
objects to the grid coordinates of the screen and generates
the sequences of the coordinates for the objects*/
 DO I=1 TO NUMBERSHIPS;
          XX=A+BINARY((X(I)-X_OFFSET)/M);
YY=B+BINARY((Y(I)-Y_OFFSET)/M);
          IF (FRIEND(I)) THEN
              CALL GENFRIEND(IX.YY.OBJECT(I).U.OEJECT(I).V);
          ELSE
              CALL GENFOE(XX.YY.OBJECT(I).U.OBJECT(I).V);
          /*END IF*/
          CALL GENWAKE(XX,YY,OBJECT(I).U_WAKE,OBJECT(I).V_WAKE);
          XX=A+BINARY((X_WAKE(PTR(I))-X OFFSET)/M);
          YY=B+BINARY((Y_WAKE(PTR(I))-Y_OFFSET)/M);
          IF COUNT(I)=4 THEN
              CALL GENWAKE(XX.YY.OBJECT(I).UU WAKE.
                                          OBJECT(I).VV WAKE);
          XX=A+BINARY((X_AIM(I)-X OFFSET)/M);
          YY=B+BINARY((Y AIM(I)-Y OFFSET)/M);
               ((XX=A)\delta(Y\overline{Y}=B)) THEN
          IF
                IF I=1 THEN
                   CALL GENOURAIM(XX,YY,OBJECT(I).U AIM.
                                              OBJECT(I).V AIM);
                ELSE
                   CALL GENAIM(XX.YY.OBJECT(I).U AIM.
                                             OBJECT(I).V_AIM);
                /*END IF*/
           /*END IF*/
  END /*DO*/;
  XX=A+BINARY((X GUN-X OFFSET)/M);
  YY=B+BINARY((Y_GUN-Y_OFFSET)/M);
  CALL GENGUN(XX,YY,OBJECT(1).U_GUN,OBJECT(1).V_GUN);
```

/*the following sequence draws the objects by calling the routine DRAW*/ D=1; DO I=1 TO NUMBERSHIPS; CALL DRAW(OBJECT(I).U,OBJECT(I).V,D); CALL DRAW(OBJECT(I).U_WAKE,OBJECT(I).V_WAKE,D); IF (ENGAGED) THEN CALL DRAW(OBJECT(I).U_AIM, OBJECT(I).V_AIM,D); END /*DO*/; IF FIRED THEN CALL DRAW(OBJECT(1).U GUN, OBJECT(1).V GUN, D); /* the following procedures produce the sequence of screen grid coordinates for various objects*/ GENFRIEND: PROC(X,Y,U,V); DCL (X.Y) FIXED BIN(15), (U.V)(Ø:1Ø) FIXED PIN(15); $U(\emptyset) = I;$ $V(\emptyset) = Y - 8;$ U(1) = X - 6;V(1)=Y-3; U(2) = I - 6;**V**(2)=Y+3; U(3)=X; V(3)=Y+8; V(4)=Y+3; U(4) = X + 6;U(5) = X + 6;V(5)=Y-3; U(6) = X;V(6)=Y-8; U(7) = -1;**V**(7)=−1; END GENFRIEND; GENFOE: PROC(X,Y,U,V); DCL (X,Y) FIXED BIN(15), (U.V)(0:10) FIXED BIN(15); **V**(Ø)=**Y**−4; U(Q) = X + 8;U(1) = I - 8;V(1) = Y - 4;U(2)=X; **V**(2)=**Y**+8; U(3)=X+8; **V**(3)=**Y**-4; U(4) = -1;**V**(4)=-1; END GENFOE;

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ようななから、11日に、11日には11日にはないないでも時代であった。11日であるような点になったのではなれた。11日のためであった。11日のためであるのでは11日であるようでは 11日のためでは、11日に、11日のためには11日にはないないでも11日であるような点になった。11日のためのためであった。11日のためためである。11日のためのないでは、11日のためである。11日のためであった。11日の

GENWAKE:PROC(X,Y,U,∀);

DCL

(X,Y) FIXED BIN(15), (U,V)(0:10) FIXED BIN(15); ÷

U(Ø)=X;	▼(∅)=Y; ▼(1)=Y;
U(Ø)=X; U(1)=X;	
U(2) = -1;	V (2)=−1;

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END GENWAKE;

GENOURAIM: PROC(X,Y,U,V);

DCL

(X.Y) FIXED BIN(15), (U.V)(0:10) FIXED EIN(15);

U(Ø)=X+8;	∀ (Ø)=Y;
U(1) = X - 8;	V(1)=Y;
U(2)=X;	V (2)=Y;
U(3)=X;	V(3)=Y-8;
U(4)=X; U(5)=-1;	V(4)=Y+8;
U(5)=-1;	₹(5)=-1;

END GENOURAIM;

```
GENAIM:PROC(X,Y,U,V);
```

DCL

(X,Y) FIXED BIN(15), (U,V)(0:10) FIXED BIN(15);

U(Ø)=X+4;	V(Ø)=Y+4;
U(1) = X - 4;	V(1)=Y-4;
U(2)=X;	V (2)= Y ;
U(3)=X+4;	₹(3)=¥-4;
U(4) = X - 4;	V (4)=Y+4;
U(5)=-1;	V(5)=-1;

END GENAIM;

```
GENGUN:PROC(X,Y,U,V);
```

DCL

(X,Y) FIXED BIN(15), (U,V)(0:10) FIXED BIN(15);

U(2)=X+1;	V (Ø)=Y−1;
U(1)=X+1;	V (1)=Y+1;
U(2)=X-1;	V (2)=Y+1;
U(3)=X-1;	V(3)=Y-1;
U(4)=X+1;	V(4)=Y−1;
U(5)=-1;	▼(5)=-1;

END GENGUN;

/*this procedure receives two arrays and a key variable as
paramaters, and either displays the object or erases it*/
DRAW:PROC(U,V,D) EXTERNAL;

DCL

```
(U,V)(0:10) FIXED BIN(15),
(I,J,H,D) FIXED BIN(7),
RUB CHAR(1) EXTERNAL,
C(7) CHAR(1),
Z1(0:3) CHAR(1) BASED(P),
Z(0:3) BIT(8),
P POINTER;
```

```
P=ADDR(Z);
                   /*Z and Z1 share same location hereon*/
I=Ø;
DO WHILE (I<11);
          IF (D=1)
               /*enter vector set level white*/
THEN PUT LIST( ^ ] (a');
              /*enter vector set level black*/
ELSE PUT LIST( ^ ] ^ [ ',RUB);
          /*END IF*/
         DO H=1 TO 5;
                    IF (U(I) < \emptyset) THEN DQ;
                             PUT LIST ( M X [=
                                                     1);
                             RETURN;
                   END /*IF*/;
            /*this call translates the coordinates to the
               stream of bits*/
                   CALL TRANSLATE(U(I),V(I).Z);
                   I=I+1;
```

```
/*this put statement puts out the bit streams
                  as being characters*/
                     PUT \overline{E}DIT((21(J) DO J=0 TO 3)) (4A(1));
           END /*DO*/;
         /*the following two statements get the status of
         the screen. The status is not sent back until the
         screen is ready.*/
           PUT LIST(' [ E'); /*HANDSHAKE*
GET EDIT((C(J) DO J=1 TO 7))(7A(1));
                                         /*HANDSHAKE*/
  END /*DO*/;
  PUT LIST( ^^M^X');
                              /*BACK TO ALPHA*/
TRANSLATE: PROCEDURE(X,Y,Z) ;
  DECLARE (X.Y) FIXED BIN(15),
            Z(0:3) BIT(8),
            T FIXED BIN(15),
            S1 BIT (16),
            SS BIT(7).
            S BIT (8),
            I FIXED BIN(7);
  I = DIVIDE(Y.32.8);
  SS=BIT(1,7);
  S='0'B || SS;
Z(0) = '00100000'B ! S;
  T = Y - (Y/32) * 32;
  S1 = BIT (T,16);
S = SUBSTR(S1,8,8);
  Z(1) = '01100000'B ! S;
  I = DIVIDE (X.32.8);
  SS=BIT(I,7);

S='0'B || SS;

Z(2) = '00100000'B | S;

T = I - (I/32) * 32;
  S1 = BIT(T, 16);
  S = SUBSTR (S1, 8, 8);
  Z(3) = 01000000 B ! S;
END TRANSLATE;
END DRAW;
```

END DISPLAY;

Sector Andrews

C. STATUS.PLI

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/* : STATUS.PLI Prog Name Date : December 83 : M. Kadri Ozyurt Written by For : Thesis : Professor Kodres Advisor Purpose : this routine calls the assembly routine **KEYBOARD** to read the keyboard to set the boolean variables used in other routines */ STATUS: PROCEDURE EXTERNAL;

/* DCL */ %INCLUDE 'CONST.INP'; %INCLUDE 'GLOBALS.INP';

```
CALL KEYBOARD (KEY);
IF (KEY='Q')!(KEY='q') THEN
    STOP;
ELSE IF (KEY='E')!(KEY='e') THEN
    ENGAGED = TRUE;
ELSE IF ENGAGED & ((RANK(KEY)>48)&(RANK(KEY)<=57))
  THEN BEGIN;
    I=RANK(KEY)-48;
    IF NUMBER(I)=0 THEN
            LINK_SHIP(I)=I;
    ELSE IF FRIEND(I) THEN DO;
            CALL REMOVENODE(LINK SHIP(OWN),I);
            CALL ADDNODE(PTR(OWN),I);
            FRIEND(I)=FALSE;
            END;
    PTR(OWN) = I;
    END;
ELSE IF (KEY='R') ! (KEY='r') THEN
    ENGAGED=FALSE;
ELSE IF (KEY='M')!(KEY='m') THEN DO;
```

```
/*set the scale to 1/200*/
          MAGNIFIED=TRUE;
           M=200.0;
           IF ENGAGED THEN DO;
           /*set reference as the target*/
                   X OFFSET=X(PTR(OWN));
                   Y OFFSET=Y(PTR(OWN));
                   END /*DO*/:
           ELSE DO;
           /*ownship is the reference*/
                   X OFFSET=0.0;
                   Y^{OFFSET=0.0};
                   END /*DO*/;
           /*END IF*/
      END;
      ELSE IF (KEY='T')!(KEY='t') THEN DO;
      /*set the scale back to normal (1/50)*/
                   MAGNIFIED=FALSE;
                   M=50.0;
                   X_OFFSET=0.0;
Y_OFFSET=0.0;
        END /*DO*/;
      ELSE IF (KEY='F')!(KEY='f') THEN
           FIRED = TRUE;
      ELSE IF (KEY='D')!(KEY='d') THEN
           ERRORON=TRUE;
      ELSE IF (KEY='S')!(XEY='s') THEN
           SIGNAL ERROR(1);
      /*END IF*/
      KEY='0';
/*this routine removes the node pointed by QQ from
                                                             the
circular linked linked list pointed by PP*/
REMOVENODE: PROC(PP.QQ);
  DCL
           (PP,QQ) FIXED BIN(7);
           ZINCLUDE GLOBALS. INP';
  P=PP;
  P=LINK SHIP(P);
DO WHILE ( (LINK SHIP(P)=QQ));
P=LINK_SHIP(P);
  END;
  LINK SHIP(P)=LINK SHIP(QQ);
```

END;

/*this routine adds a node pointed by QQ to the circular linked list by PP*/ ADDNODE:PROC(PP,QQ); DCL (PP,QQ) FIXED BIN(7); XINCLUDE'GLOBALS.INP'; P=PP; P=LINK_SHIP(P); DO WHILE (~(LINK_SHIP(P)=PP)); P=LINK_SHIP(P); END; LINK_SHIP(P)=QQ; LINK_SHIP(QQ)=PP; END; END; •

END STATUS;

53.55

D. IDLE.PLI

/* : IDLE.PLI Prog Name : December 83 Date Written by : M. Kadri Ozyurt For : Thesis : Professor Kodres Advisor Purpose : This routine reads the A/D converter output to get the velocity vectors of ownship and а selected ship, and the gun information as azimuth and elevation. It then converts this information to real world values. It calculates ownship speed which will be used to find relative speeds later. It then computes the maximum cartesian coordinates of the splash point, and time range. of flight corresponding to the current gun position */

ことでは、この「東京などではないではない」となっていた。「東京などのでは、「東京などのない」となった。「東京などのではない」というできた。「東京などのない」というない」というない。「東京などのない」というない。

IDLE: PROCEDURE EXTERNAL;

/* DCL */ %INCLUDE 'CONST.INP'; %INCLUDE 'GLOBALS.INP';

> DO D=0 TO 5; CALL ATOD (D,ARG(D)); END /*DO*/;

/*at this point the A/D output values are fixed bin(7)
values. The following sequence converts those to fixed
decimal values*/
 COURSE(OWN)=ARG(Ø);
 SPEED(OWN)=ARG(4);
 COURSE(KNOWN)=ARG(2);
 SPEED(KNOWN)=ARG(3);
 AZ=ARG(1);
 ALT=ARG(5);

/*the following sequence converts A/D values to real time
values by using appropriate proportionality constants*/
COURSE(OWN) = COURSE(OWN) * K;

```
COURSE(KNOWN) = COURSE(KNOWN) * K;
      AZ = AZ * K;
      IF COURSE(OWN)<0.0 THEN
          COURSE(OWN) = COURSE(OWN) + TWO PI ;
      IF COURSE(KNOWN)<0.0 THEN
          COURSE(KNOWN) = COURSE(KNOWN) + TWO PI ;
      IF AZ<0.0 THEN
          AZ = AZ + TWO PI;
       IF ALT>90.0 THEN
          ALT = 90.0;
          SPEED(OWN) = SPEED(OWN)/L;
          SPEED(KNOWN) = SPEED KNOWN) / L;
/*ownship speed computations*/
      VX OWN = SPEED(OWN) * SIND(COURSE(OWN));
      VY OWN = SPEED(OWN) * COSD(COURSE(OWN));
/*when not have fired, the following makes the ballistic
computations*/
      IF
           FIRED THEN
         BEGIN;
          T OF = 2.0 * VM * SIND(ALT) / G;
          V\bar{R} = VM + COSD(ALT);
          R = VR * T OF;

X_AIM(OWN) = R * SIND(AZ);
          T = AIM(OWN) = R + COSD(AZ);
          X GUN = 0.0;
          Y GUN = 0.0;
          V\bar{X}_ROUND = VR + SIND(AZ);
          VY_ROUND = VR * COSD(AZ);
          T = T OF;
      END /*IF*7;
```

END IDLE;

APPENDIX D

REAL TIME EXECUTIVE MODULE LISTINGS

A. ARBITER.A86

; Prog Name : ARBITER.A86 ; Date : December 83 : M. Kadri Ozyurt Written by For : Thesis : Professor Kodres ;Advisor : This program contains all the assembly ; Purrose ;routines used by the simulation system. It initializes all ;programmable hardware components, responds to the timing ; interrupts, and increment the FOURTHEVC used throughout the Upon receiving ;simulation model program. interrupt it performs process switching by storing the ;requests , ;state of interrupted process in the stack area allocated for the processes and by restoring the highest ready ; process given by SCHEDULER

;GLOBALS

DGROUP GROUP FOURTHEVC, CURRENTPROC FOURTHEVC DSEG COMMON FOURTHEVC1 DW Ø CURRENTPROC DSEG COMMON CURRENTPROC1 DB 4

CSEG	
EXTRN	SCHEDULE: FAR
EXTRN	P1 : FAR
EXTRN	P2:FAR
EXTRN	P3:FAR
EXTRN	P4:FAR
PUBLIC	ARBITER
PUBLIC	STORESTATUS
PUBLIC	RINGBELL
PUBLIC	KEYBOARD
PUBLIC	ATOD
PUBLIC	SUSPEND
PUBLIC	RESIME

) DQUALDO		
ICW4 MASK1 EOI CNTR1 CNTR2 PORTC COUNT CNTRLO CNTRHI READWR RXRDY TXRDY PORTIO PORTST SEGCONV OFFCONV STACKSI LF CR BEL FS ESC FF CAN FALSE	EQU EQU EQU EQU EQU EQU EQU EQU EQU EQU	84H;INTR1 JUMP ADDR.8CH;INTR3 JUMP ADDRESS9C0H;PIC COMMAND OUTPUT PORT19C2H;PIC COMMAND OUTPUT PORT213H;PIC COMMAND WORD120H;PIC COMMAND WORD29DH;PIC COMMAND WORD49FDH;PIC COMMAND WORD49FDH;PIC COMMAND WORD49FDH;PIC CONTROL BYTE90H;PIT CONTROL PORT90H;PIT CONTROL PORT90H;PIT COUNT # OUTPUT PORT90H;PIT COUNT # LO BYTE96H;PIT COUNT # HI BYTE37H;USART STATUS MASK(READ)91H;USART STATUS MASK(WRITE)908H;USART STATUS PORT9080H;A/D CONVTR PORT SEGMENT97700H;A/D CONVTR PORT SEGMENT978;A/D CONVTR PORT OFFSET96H;A/D CONVTR PORT SEGMENT978;A/D CONVTR PORT SEGMENT978;A/D CONVTR PORT OFFSET96H;A/D CONVTR PORT OFFSET
TRUE	EQU	NOT FALSE
J RUL	しんつ	NOT LUDD
7		
ARBITER:		
	PUSH	DS
	CLI	;DISABLE INTR'S
	MOV	AX,0
	MOV	DS,AX ;SET SEGREG TO Ø
	MOV Mov	BX,INT1 WORD PTR [BX],OFFSET PROCØ ;INT1 JMP ADDRESS
	INC	BX
	INC	BX
	MOV	WORD PTR [EX],CS
	POP	DS
	MOV	BX, OFFSET STACKTBL+STACKSIZE-2
	MOV Add	CS:[BX],CS BX,STACKSIZE
	UUU	DY PILATICE

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; EQUATES

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	MOV ADD	CS:[BX],CS BX,STACKSIZE	
	MOV	CS:[BX],CS	
	ADD	BX,STACKSIZE	
		CS:[BX].CS	
	MOV	03:[DA],03	
	CLI	AT TOWA	;INIT. PIC TO
	MOV	AL,ICW1	; EDGE-TRIG., SINGLE PIC
	OUT	PIC1,AL	, EDGE-INIG., SINGDE INC
	MOV	AL,ICW2	; INT1 ADDR IS Ø4H
	OUT	PIC2, AL	
	MOV	AL,ICW4	; NOT F. NESTED, NORM. EOI
	OUT	PIC2,AL	, NOT F. WEDTED, NORME 202
	STI		
	MOV	AL, MASK1	Y THMA TO ATTOWED
	OUT		Y INT1 IS ALLOWED ;INIT. PIT
	MOM	AL, CNTR1	
	OUT	PORTC,AL	;SELECT MODE Ø ,CNTR 1
			; CLK FREQ. IS 153.6 KHZ
	MOV	AL, CNTRLO	COUNT-DOWN VALUE 9600H
	OUT	COUNT,AL	; WHICH GIVES AN INTR AT
	MOV	AL, CNTR2	
	out	PORTC,AL	
	MOV	AL, CNTRHI	; EVERY FORTH OF A SEC.
	OUT	COUNT, AL	
	MOV	AL, READWR	
	OUT	PORTST, AL	
	MOV	AX,CS	
	MOV	SS,AX	;SET STACK SEG. TO CODE
	MOV	BP,3*STACKSIZE	
	MO₹	SP, STACKTBL [BP]	
	JMP	P4	
-			
; PROCØ:			
FRUCE:	PUSH	XA	
	PUSH	EX	
		CX	
	PUSH		
	MO V Out	AL, CNTR2 Portc, Al	;RESET COUNTER
		AL, CNTRHI	RESET THE CNTR.
	MOV		
	OUT	COUNT, AL	;RESET PIC
	MOV	AL, EOI	VALODI IIO
	OUT	PIC1,AL	;SP->INTERRUPTED IP
	ADD	SP,6	; BX-> IP
	POP	BX	;AX-> " CS
	POP	AI	;SP->INTERRUPTED IP
	POP	CI AI	PUSH CS
	PUSH PUSH	BX	PUSH "IP
	PUSH	CX	,
	LOOU	VA	

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		SUB	SP,6	;SP->PUSHED CX
		POP	CX	••••••••••
		POP	BX	;RESTORE BX
		POP	AX	RESTORE AX
		INC	FOURTHEVC1	
		JMP	STORESTATUS1	
S	TORESTATU		C. ORDOINIOUI	
U	TOWEDINI	DEC	SP	
		DEC	SP	
		PUSH	AX	
		ADD	SP,4	
		POP	AX	;AX->INTERRUPTED IP
		PUSE	CS	
		PUSH	AX	
		DEC	SP	
		DEC	SP	_
		POP	AX	;RESTORE AX
		PUSHF		
		CLI		
S	TORESTATI	US1:		
			I PUSH BX I PUSH CX I	
		PUSH BP	I PUSH SI I PUSH DI I	FUSH ES
		CALL	SCHEDULE	
		MOV	NEWPROC, AL	
		MOV	AL, CURRENTPROC1	
		CMP	AL.1	
		JNZ	OUT1	
		MOV	BP,1	
		JMP	OUT4	
0	UT1:	CMP	AL,2	
•		JNZ	OUT2	
		MOV	BP, STACKSIZE +1	
		JMP	OUT4	
0	UT2:	CMP	AL.3	
•	~ ~ •	JNZ	OUT3	
		MOV	BP.2*STACKSIZE+1	
		JMP	OUT4	
0	UT3:	CMP	AL,4	
U	0101	JNZ	OUT5	
		MOV	BP.3*STACKSIZE+1	
0	UT4:	JMP	OUT4	
v	014:	ADD	BP, OFFSET STACKTBL	
		CMP	SP, BP	
		JNA	OUT6	
		DEC	BP	
		MOV	[BP],SP	
-		JMP	LOADPROC	
0	UT5:	ADD	AL,30H	
		MOV	OUT5AL,AL	
		MOV	DX, OFFSET OUT5MESS	
		JMP	ERRORMESS	

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AL, 30H **OUT6:** ADD MO₹ OUTGAL, AL MOV DX.OFFSET OUT6MESS ERRORMESS JMP LOADPROC: AL, NEWPROC MOV CMP AL,1 JNZ LOUT1 MOV BP,Ø RETURNPT JMP LOUT1: AL,2 CMP LOUTZ JNZ BP, STACKSIZE MOV JMP RETURNPT LOUT2: CMP AL,3 LOUT3 JNZ BP,2*STACKSIZE MOV RETURNPT JMP LOUT3: CMP AL,4 LOUT4 JNZ MOV BP, 3*STACKSIZE JMP RETURNPT LOUT4: ADD AL,3ØH LOUT4AL,AL DX,OFFSET LOUT4MESS MOV MOV JMP ERRORMESS **RETURNPT**: MOV CURRENTPROC1, AL MOV SP.STACKTBL [BP] POP ES I POP DI I POP SI I POP BP DX I POP CX I POP BX I POP AX POP POPF STI RETF ; -OUT ;THIS ROUTINE SYSYTEM CALL PUT ERROR MAKES A TO ; MESSAGES ERRORMESS: CL,9 MOV INT 224 MOV CL,Ø MOV DL,1 INT 224 RET

.

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;THE STACK AREAS AND VARIABLE DEFINITIONS

STACKTBL:

	Dui	
	DW	OFFSET STACKTBL+STACKSIZE-22
	RS	STACKSIZE-8
	DW	FALSE
	DW	OFFSET P1
	RS	2
	DW	OFFSET STACKTBL+2*STACKSIZE-22
	RS	STACKSIZE-8
	DW	FALSE
	DW	OFFSET P2
	RS	2
	DW	OFFSET STACKTBL+3*STACKSIZE-22
	RS	STACKSIZE-8
	DW	FALSE
	DW	OFFSET P3
	RS	2
	DW	OFFSET STACKTBL+4*STACKSIZE-22
	RS	STACKSIZE-8
	DW	FALSE
	DW	OFFSET P4
	RS	2
NEWPROC	DB	Ø
OUT 5MESS	DB	FS, ESC, FF, CAN
	DB	'RETURN FROM OUTS. AN UNKNOWN CURRENT PROCEDURE:
OUTSAL	DB	0, '\$',0
OUT6MESS	DB	FS.ESC.FF.CAN
	DB	'RETURN FROM OUTG. STACK OVERFLOW FOR THE PROC .: '
OUTGAL	DB	0, '\$',0
LOUT4MESS		FS, ESC, FF, CAN
	DB	RETURN FROM LOUT4. AN UNKNOWN NEW FROCEDURE:
LOUT4AL	DB	0, '\$',0
	- 2	~, , , , , , , , , , , , , , , , , , ,
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END

B. AWAIT.PLI

/* : AWAIT.PLI Prog Name : December 83 Date by Written : M. Kadri Ozyurt : Thesis For Advisor : Professor Kodres This synchronization primitive checks Purpose : threshold value for the calling process by the comparing value with FOURTHEVC and the corresponding threshold returns the control either to the calling process. if its threshold value is equal to and greater than FOURTHEVC, or else transfers the control to ARBITER.A86 */

C. SCHEDULE.PLI

/* : SCHEDULE.PLI Frog Name : December 83 : M. Kadri Ozyurt Date Written by For : Thesis : Professor Kodres Advisor : This synchronization Purpose primitive threshold values corresponding to compares the the processes P1 through P3, beginning from P1, to FOURTHEVC and returns the name of the first one which is equal to or greater than that value. If non of the processes meet this conditions then P4 is returned. */

وأخاره الازمار فالاركار والمستعد والمحا

schedule: procedure returns (fixed bin(7));
 dcl threshold(0:2) fixed bin(15) external,
 fourthevc fixed bin(15) external,
 i fixed bin(7);
 do i=0 to 2;
 if (fourthevc>=threshold(i)) then return (i+1);
 end;
 return (4);
end schedule;

D. THRESH.PLI

/* Prog : THRESH.PLI Name Date : December 83 Written by : M. Kadri Ozyurt For : Thesis : Professor Kodres Advisor primitive : This Purpose synchronization receives a pointer to the calling process and increments the corresponding threshold value by an assigned amount */

THRESH: PROC(1);

DCL THRESHOLD(0:2)FIXED BIN(15) EXTERNAL, i FIXED BIN(7);

IF (1=1) THEN THRESHOLD(Ø)=THRESHOLD(Ø) + 1; IF (1=2) THEN THRESHOLD(1)=THRESHOLD(1) + 4; IF (1=3) THEN THRESHOLD(2)=THRESHOLD(2)+4; RETURN; END THRESH;

WEAR AND AND AND AND

E. P1.PLI

/* Prog Name : P1.PLI : December 83 Date Written by : M. Kadri Ozyurt For : Thesis Advisor : Professor Kodres Purpose : This process is basically an infinitive loop. Once entered, it first call AWAIT.PLI to see if FOURTHEVC is equal to or greater than its threshold value. If it is, then the control proceeds to call TACTICAL.PLI. The last call in the loop is to THRESH.PLI to increment its threshold value. In the next iteration, the control will not come back since its threshold value is greater than FORTHEVC. */

P1: PROCEDURE;

DCL AWAIT ENTRY (FIXED BIN(7)), THRESH ENTRY (FIXED BIN(7)), A FIXED BIN(7), TACTICAL ENTRY; A=1; DO WHILE ('1'B); CALL AWAIT (1); CALL TACTICAL; CALL THRESH (1); END /*DO*/;

END P1;

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F. P2.PLI

/* Prog Name : P2.PLI Date : December 83 Written by : M. Kadri Ozyurt For : Thesis Advisor : Professor Kodres : The purpose of this process is idecti-Purpose cal to that of P1.PLI with the exception that the second call is to DISPLAY.PLI */

P2: PROCEDURE;

DCL AWAIT ENTRY (FIXED BIN(7)), THRESH ENTRY (FIXED BIN(7)), A FIXED BIN(7), DISPLAY ENTRY; A=2; DO WHILE ('1'B);

```
CALL AWAIT (2);
CALL DISPLAY;
CALL THRESH (2);
END /*DO*/;
END P2;
```

G. P3.PLI

/* Prog Name : P3.PLI Date : December 83 : M. Kadri Ozyurt Written by For : Thesis Advisor : Professor Kodres : The purpose of this process is the Purpose same as P1.PLI with the exception that the second call is to STATUS.PLI */

P3: PROCEDURE;

DCL AWAIT ENTRY (FIXED BIN(7)), THRESH ENTRY (FIXED BIN(7)), A FIXED BIN(7), STATUS ENTRY;

```
A=3;
DO WHILE ('1'B);
CALL AWAIT (3);
CALL STATUS;
CALL THRESH (3);
END /*DO*/;
END F3;
```

H. P4.PLI

/# Prog Name : P4.PLI Date : December 83 Written by : M. Kadri Ozyurt For : Thesis Advisor : Professor Kodres Purpose : This process is an infinitive loop in which there is only one call to IDLE.PLI repeaditively until an interrupt comes along. */

F4: PROCEDURE;

DCL AWAIT ENTRY (FIXED BIN(7)), THRESH ENTRY (FIXED BIN(7)), A FIXED BIN(7), IDLE ENTRY;

A=4; DO WHILE ('1'B); CALL IDLE; END /*DO*/; END P4;

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APPENDIX E

MISCELLANEOUS ASSEMBLY ROUTINES

A. KEYBOARD.A86

;Prog Name : KEYBOARD.A86 : December 83 ; Date ;Written by : M. Kadri Ozyurt : Thesis ;For : Professor Kodres ;Advisor ;Purpose : This program receives a formal jparameter, KEY, reads the status of the serial I/0 ; interface chip. If a character has been received from the ;keyboard. it reads this character and places it to the ;formal parameter. If there is not a character available ; it puts a ascii equivalent of zero into the parameter. ;The reason for that is that zero is not used as a keyboard ;command. The variables used here are defined in the body ; of ARBITER.A86

KEYBOARD: PUSHF CLI PUSH AX IN AL.PORTST CMP AL, RXRDY KEYBOARD1 JZ AL, PORTIO IN AL,7FH AND JMP **KEYBOARD2 KEYBOARD1:** AL.30H MOV **KEYBOARD2:** BX.[BX] MOV MOV [BX],AL POP AX POPF RET

B. ATOD.A86

· Brea Nora	ATOD. A86
;Date :	December 83
;Written by :	M. Kadri Ozyurt
;For :	Thesis
;Advisor :	Professor Kodres
;Purpose :	This program receives two parameters
;It reads the output	of the A/D converter specified by the
;second parameter an	d places it into the first parameter.
	here are defined in ARBITER.A86

ATOD:

PUSHF	
CLI	
PUSH	SI
PUSH	AX
PUSH	BX
PUSH	BX
PUSH	DS
MOV	BX,[BX]
MOV	AH,Ø
MOV	AL, [BX]
MOV	SI,AX
MOV	AX,SEGCONV
MOV	DS,AX
MOV	BX,OFFCONV
MOV	AL,[BX+SI]
POP	DS
POP	BX
MOV	BX,2[BX]
MO₹	[BX],AL
POP	BX
POP	AX
POP	SI
POPT	
ret	

; BX = . ARGUMENT(1)

;READ A/D PORT

C. RINGBELL.A86

;Prog Name	: RINGBELL.A86	
;Date	: December 83	
;Written by	: M. Kadri Ozyurt	
For	: Thesis	
;Advisor	: Professor Kodres	
;Purpose	: This program sends a bell character to	,
	. The variables used here are defined in	
;the tody of ARBITE		

	For		: Thesis
	;Advisor		: Profess : This pr
5.2	;Purpose the vide	o termi	: This pr nal. The var
2	the body	of ARB	ITER.A86.
5			
à.	RINGBELL:	DICUP	
		PUSHF Cli	
		PUSH	AX
1		CALL	WAIT
نده مد		MOV	AL,BEL
à		OUT	PORTIO,AL
	_	POP	AX
		POPF	
	•	RET	
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D. WAIT.A86

jProg Name	: WAIT.A86
;Date	: December 83
;Written by	: M. Kadri Ozyurt
;For	: Thesis
;Advisor	: Professor Kodres
;Purpose	: This program program reads the status
;of the serial I/C) chip and waits until the transmitter is
;ready to send char	acters.

;Advisor		: Thesis : Professor Ko	dres
Furpose of the	serial	: This program I/O chip and wait	program reads the s until the transmi
iready to	o send c	naracters.	
WAIT:	סווכש	A ¥	
WAIT1:	IN AND JZ	AL,ØDEH AL,1 WAIT1	;GET STATU
	RET	AX	
; ;			
		92	
	<pre>;Advisor ;Purpose ;of the ;ready to WAIT: WAIT1: ;;</pre>	;Advisor ;Purpose ;of the serial ;ready to send c WAIT: WAIT1: N AND JZ POP RET ;	<pre>;Advisor : Professor Ko ;Purpose : This program ;of the serial I/O chip and wait ;ready to send characters. WAIT:</pre>

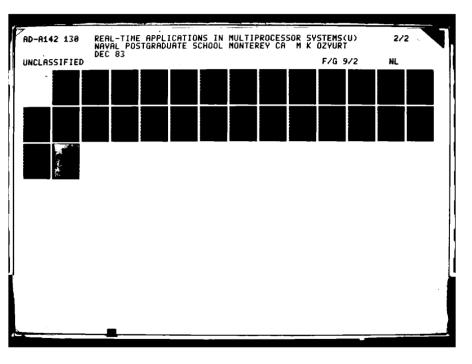
E.SUSPEND.A86

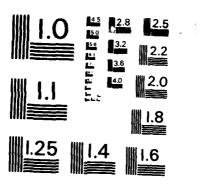
;Prog Name : SUSPEND.A86 ;Date : December 83 ;Written by : M. Kadri Ozyurt ;For : Thesis ;Advisor : Professor Kodres ;Purpose : This program stops the real time clock ;by reseting the interrupt bit of the PSW.

SUSPEND:

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F. RESUME.A86

iProg Name : RESUME.A86
jDate : December 83
iWritten by : M. Kadri Ozyurt
iFor : Thesis
iAdvisor : Professor Kodres
iPurpose : This program starts the real time clock
iby reseting the interrupt bit of the PSW. It then reset
ithe counter to zero.

:

RESUME:

PUSH	XA		
MOV	AL.CNTR2	;RESET	COUNTER
OUT	POÁTC,AL		
MOV	AL, CNTRHI		
OUT	COUNT .AL		
MOV	AL.EOI	; RESET	PIC
OUT	PIC1.AL		
POP	AX		
STI			
RET			

;

;

A. LOCALS AID

1. A. A.

/* Prog : LOCALS.AID Name Date : December 83 Written by : M. Kadri Ozyurt For : Thesis Advisor : Professor Kodres : This %include file contains the declara-Purpose tions of the variables used by the dynamic debugging module */

DCL BREAKS(0:9) LABEL, STOPS(0:9) BIT(1) EXTERNAL, (CODE1,CODE2,VALUE,H) FIXED BIN(15), BRFAKPT FIXED BIN(15) EXTERNAL, PUTVARS ENTRY, REENTRY ENTRY, BREAKPTS ENTRY, BREAKPTS ENTRY, PROMPTUSER ENTRY (FIXED BIN(7)), STORESTATUS ENTRY, TACTICAL ENTRY, DISPLAY ENTRY, IDLE ENTRY, STATUS ENTRY, CHANGEVA ENTRY (FIXED BIN(15),FIXED BIN(15));

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B. ERRHAND.AID

/*

Prog : ERRHAND.AID Name : December 83 Date : M. Kadri Ozyurt Written by : Thesis For Advisor : Professor Kodres : This %include file contains six Purpose different types of PL/I ON condition bodies. Upon intercepting any raised error condition is displayed and the control is transfered to PROMPTUSER with a number that shows which breakpoint has been past. Then the ON condition body is exited with a non-local goto statement. At the exit point the control is transfered to REENTRY which is the dynamic debugging tool. This call to REENTRY is protected during the course of normal operation with an if statement which tests the value of ERRORON. */ stops(0)=false; stops(1)=false; stops(2)=false; stops(3)=false; stops(4)=false; stops(5)=false; stops(6)=false; stops(7)=false; stops(8)=false; stops(9)=false; on error begin; put list(^^Z'); /*clear screen*/ put skip list('Error #'); /*this statement gets the code of the error condition*/ code1=oncode(); /*this call prompts the user with the # of breakpoints past and asks if the user wants to enter the dynamic debugging environment*/ call promptuser(code1); if (key='y') ! (key='Y') then goto errorexit; else if code1<=127 then do; put skip list('The program will be abandoned'); stop; end /*if*/;

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```
end /*error*/;
on fixedoverflow
  begin;
 put list(^^Z');
 put skip list('Fixedoverflow #');
  code1=oncode();
 call promptuser (code1);
 if (key='y') ! (key='Y') then goto errorexit;
  else if codel <= 127 then do;
          put skip list('The program will be abandoned');
          stop;
  end /*if*/;
end /*fixedoverflow*/;
on overflow
  tegin;
 put list(^^Z');
  put skip list('Overflow #');
  code1=oncode();
  call promptuser(code1);
  if (key='y') ! (key='Y') then goto errorexit;
  else if code1<=127 then do;
          put skip list('The program will be abandoned');
          stop;
  end /*if*/;
end /*overflow*/;
on underflow
  begin;
 put list(^^Z');
  put skip list('Underflow #');
  code1=oncode();
  call promptuser(code1);
  if (key='y') ! (key='Y') then goto errorexit;
  else if code1<=127 then do;
          put skip list('The program will be abandoned');
          stop;
  end /*if*/;
end /*underflow*/;
on zerodivide
  begin;
```

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end /*zerodivide*/;

errorexit:
 if erroron then do;
 call reentry();
 end /*if*/;

C. PROMPT.PLI

/* : PROMPT.PLI Prog Name Date : December 83 Written by : M. Kadri Ozyurt For : Thesis Advisor : Professor Kodres : This routine prompts the user in case Purpose of an error interception by ERRHAND. It puts the code of the error and which breakpoint has been past. Then the program asks if the user wants to enter the dynamic debugging environment. */ promptuser:proc(code1) external; /* dcl */ %include 'const.inp'; %include 'globals.inp'; %include 'locals.aid'; erroron=true; put list(code1, ' is detected when the module '); put skip list('The clock stopped.'); put skip list('The breakpoint #', breakpt, 'has been past.'); put skip list('Would you like to enter interactive debugging', environment (Y/N)?'); get list(key); end promptuser;

D. REENTRY.PLI

/* : REENTRY.PLI Prog Nare : December 83 Date Written by : M. Kadri Ozyurt For : Thesis Advisor : Professor Kodres : This routine is the "workhorse" of Purpose the dynamic debugging environment. It calls PUTVARS if the user wants to see the external variables. Then it calls CHANGEVA if the user wants to change any variable with in a loop until no changes are wanted. It then transfers the control to the breakpoint the user desires. */ reentry:proc external; /* dcl */ %include 'const.inp'; %include 'globals.inp'; %include 'locals.aid'; put skip list ('You have entered the interactive debugging , 'environment.'); skip list('You will be asked questions about the put 'control of program flow'); reentry1: put skip list('Do you want a listing of all variables ' (Y/N)?'); get list(key); if (key='Y') ! (key='y') then call putvars(); put skip list('Do you want to change the value of any ' 'variable(Y/N)?'); get list (key); do while (~((key = `N`) ! (key=`n`)));
 if (key = `Y`) ! (key=`y`) then do; put skip list('Enter the number and the new value ', (-32768<=value<=+32,767) of the veriable you want to'); skip list(' change in integers seperated by put a 'comma.'); on error begin; put list ('*** bad entry, try again'); goto reentry2; end /*error*/;

```
on fixedoverflow begin;
          put list('*** too large, try again');
          go to reentry2;
  end /*fixedoverflow*/;
reentry2:
  put škip list('>');
  get list (code1,value);
  revert error;
  revert fixedoverflow;
  if code1>maxvars then
      put list('invalid variable number');
  else
      call changeva(code1,value);
  /*end if*/
  put skip list('Do you want a listing again (Y/N)?');
  get list(key);
  if (key='Y')!(key='y') then call putvars();
  put skip list ('Do you want to change another
                                                      variable '
                                                         (Y/N)?');
      end /*do*/:
    else
  put list('*** bad entry. try again');
    /*end if*/
    get list(key);
end/#do#/;
put skip list ('Which breakpoint do you want to transfer the ',
                          'control (Ø thru 9, foolowed by return)?');
get list(key);
do while ((rank(key)<48) ! (rank(key)>57));
  put list('*** bad entry, try again');
  put skip list('>');
  get list(key);
end/*do*/;
code1=rank(key)-48;
breakpt=code1;
put skip list('Enter the breakpoint you want to stop (0',
'thru 9) or any non-numeral character if you do not want
                                                                   to
                                                    stop(fol. RET');
get list(key);
if ((rank(key)<48) ! (rank(key)>57)) then do;
 put skip list('The program will execute beginning from '
                                           the breakpoint',code1);
  errorop=false;
      end /*do*/;
else
    doi
  code2=rank(key)-48;
  stops(code2)=true;
  put skip list('The program will execute between the ',
                                `breakpts',code1, and ',code2);
```

```
end /*if*/;
put skip list('Is that what you want(Y/N)?');
get list(key);
if (key='N') ! (key='n') then do;
   put skip list('Do you want another run(Y/N)?');
   get list (key);
   if (key='Y') ! (key='y') then goto reentry1;
end /*if*/;
```

end reentry;

E. PUTVARS.PLI

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```
/*
                   : PUTVARS.PLI
      Name
Prog
Date
                   : December 83
Written by
                   : M. Kadri Ozyurt
For
                   : Thesis
                   : Professor Kodres
Advisor
                   : This routine puts selected external
Purpose
variables out with PL/I put edit statement.
*/
putvars:proc external;
  %replace
             mar ships by 2,
             true by 1'b.
             false by '0't;
  /*dc1*/
           %include 'globals.inp';
  i=0;
  j=0;
  put skip list('The listing of all common variables is as '
                                                        follows:');
  put skip(2) list('Fixed binary values:');
  put skip edit('(1)seconds=',seconds,'(2)minutes=',minutes,
     (3)hours=',hours,'(4)wake_ptr=',wake_ptr,'(5)I=',i)(r(format1));
  currentproc)(r(format1));
  put skip edit('(14)fourthevc=',fourthevc)(a,f(5));
  put skip(2) list('Fixed decimal values:');
  put skip(2) list('Boolean values:');
  (a,b(1),col(20),a,b(1),col(41),a,b(1),col(58),a,b(1));
  put skip(2) list('Fixed decimal values:');
  put skip edit('(19)vx_own=',vx_own,'(20)vy_own=',vy_own,
'(21)vx_target=',vx_target,'(22)vy_target=',vy_target)(r(format2));
put skip edit('(23)vx_rel=',vx_rel,'(24)vy_rel=',vy_rel,
'(25)vx_round=',vx_round,'(26)vy_round=',vy_round)(r(format2));
  put skip edit('(27)vr=',vr,'(28)alpha=',alpha,
  '(29)ax=',ax,'(30)bx=',bx)(r(format2));
put skip edit('(31)cx=',cx,'(32)ay=',ay,'(33)by=',by,'(34)cy=',cy)
                                                              (r(format2));
```

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```
put skip edit('(47)key=',key)(a(8),a(1));
put skip(2)list('Arrays:');
put skip edit('(48)threshold(0)=',threshold(0),'(49)threshold(1)=',
                 threshold(1), (50) threshold(2)=', threshold(2))
                 (a,f(5),col(26),a,f(5),col(51),a,f(5));
       put skip(2) list('Data structures:');
       put skip list('ship(1):');
       put skip edit( (51)course=', course(1), (52) speed=', speed(1
      (r(format6));
       put skip list('Ship(2):');
       put skip edit( (63)course= ', course(2), '(64) speed= ', speed(2),
       (65) azimuth=', azimuth(2), '(66) range=', range(2))(r(format4));
put skip edit('(67) x=', x(2), '(68) y=', y(2), '(69) x_aim=', x_aim(2),
       (r(format6));
      put skip edit('(79)x_vake=',x_vake(ptr(2)),'(80)y_vake='
              y_wake(ptr(2)), '(81)link_wake=',link_wake(ptr(2)))
       (a.f(7,1),col(20),a.f(7,1),col(39),a,f(1));
put skip edit ('(82)dt=',dt)(a,f(4,2));
format1:format(a,f(2),col(16),a,f(2),col(31),a,f(2),col(44),a,f(2),col(59),
                                                  a.f(5));
format2:format(a,f(4,1),col(19),a,f(4,1),col(37),a,f(4,1),col(58),a,
                                                         f(4,1));
format3:format(a,f(4,1),col(24),a,f(4,1),col(47),a,f(4,1));
format4:format(a,f(4,1),col(19),a,f(4,1),col(35),a,f(4,1),col(54),a
format5:format(a,f(7,1),col(19),a,f(7,1),col(35),a,f(7,1),col(54),a,
f(7,1));
                                                           f(7,1));
format6:format(a,f(1),col(19),a,f(2),col(35),a,f(2),col(54),a,f(2));
end putvars;
```

F. CHANGEVA.PLI

```
/*
Prog Name : CHANGEVA.PLI
Date : December 83
Written by : M. Kadri Ozyurt
For : Thesis
Advisor : Professor Kodres
Purpose : This routine changes a selected
external variable specified by the parameter passed
*/
```

changeva:proc(code1,value) external;

dcl
 (code1,value) fixed bin(15);

%include 'const.inp'; %include 'globals.inp';

```
if code1=1 then seconds=binary(value,7);
if code1=2 then minutes=binary(value,7);
if code1=3 then hours=binary(value,7);
if code1=4 then wake_ptr=binary(value,7);
if code1=5 then i=binary(value,15);
if code1=6 then t_of=binary(value,7);
if code1=7 then target=binary(value,7);
if code1=8 then own=binary(value,7);
if code1=9 then known=binary(value.7);
if code1=10 then j=binary(value,15);
if code1=11 then t=binary(value,7);
if code1=12 then t_prime=binary(value,7);
if code1=13 then currentproc=binary(value,7);
if code1=14 then fourthevc=binary(value,15);
if code1=15 then engaged=bit(value,1);
if code1=16 then magnified=bit(value,1);
if code1=17 then fired=bit(value,1);
if code1=18 then erroron=bit(value.1);
if code1=19 then vx_own=decimal(value,4,1);
if code1=20 then vy_own=decimal(value,4,1);
if code1=21 then vx_target=decimal(value,4,1);
if code1=22 then vy_target=decimal(value,4,1);
if code1=23 then vx_rel=decimal(value,4,1);
if code1=24 then vy_rel=decimal(value,4,1);
if code1=25 then vr_round=decimal(value,4,1);
if code1=26 then vy_round=decimal(value,4,1);
if code1=27 then vr=decimal(value,4,1);
```

if code1=28 then alpha=decimal(value,4,1); if code1=29 then ax=decimal(value,7,2); if code1=30 then bx=decimal(value,7,2); if code1=31 then cr=decimal(value,7,2); if code1=32 then ay=decimal(value,7,2); if code1=33 then by=decimal(value,7,2); if code1=34 then cy=decimal(value,7,2); if code1=35 then ax_sum=decimal(value,7,2); if code1=36 then br_sum=decimal(value,7,2); if code1=37 then cx_sum=decimal(value,7,2); if code1=38 then ay_sum=decimal(value,7,2); if code1=39 then by_sum=decimal(value,7,2); if code1=40 then cy_sum=decimal(value,7,2); if code1=41 then x at5=decimal(value.7,2); if code1=42 then y_at5=decimal(value,7,2); if code1=43 then r=decimal(value,7,2); if code1=44 then dx_dt_at5=decimal(value,7,2); if code1=45 then dy_dt_at5=decimal(value,7,2); if code1=46 then dr_dt_at5=decimal(value,7,2); if code1=47 then key=ascii(value); if code1=48 then threshold (\emptyset) =value; if code1=49 then threshold(1)=value; if code1=50 then threshold(2)=value; if code1=51 then course(1)=decimal(value,4,1); if code1=52 then speed(1)=decimal(value,3,1); if code1=53 then azimuth(1)=decimal(value,3,0); if code1=54 then range(1)=decimal(value,5,0); if code1=55 then x(1)=decimal(value,6,1); if code1=56 then y(1)=decimal(value.6.1); if code1=57 then x_aim(1)=decimal(value,6,1); if code1=58 then y_aim(1)=decimal(value,6,1); if code1=59 then count(1)=binary(value,7); if code1=60 then number(1)=binary(value,7); if code1=61 then ptr(1)=binary(value,7); if code1=62 then link ship(1)=binary(value,7); if code1=63 then course(2)=decimal(value,4,1); if code1=64 then speed(2)=decimal(value,3,1); if code1=65 then azimuth(2)=decimal(value,3,0); if code1=66 then range(2)=decimal(value,5,0); if code1=67 then x(2)=decimal(value,6,1); if code1=68 then y(2)=decimal(value,6,1); if code1=69 then x_aim(2)=decimal(value,6,1); if code1=70 then y_aim(2)=decimal(value,6,1); if code1=71 then count(2)=binary(value,7); if code1=72 then number(2)=binary(value,7); if code1=73 then ptr(2)=binary(value,7); if code1=74 then link_ship(2)=binary(value,7); if code1=75 then az=decimal(value,4,1); if code1=76 then alt=decimal(value,4,1); if code1=?? then x_gun=decimal(value,6,1);

if code1=78 then y_gun=decimal(value,6,1); if code1=79 then x_wake(ptr(2))=derimal(value,6,1); if code1=80 then y_wake(ptr(2))=decimal(value,6,1); if code1=81 then link_wake(ptr(2))=binary(value,7); if code1=82 then dt=float(value,7); 1111 311

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end changeva;

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G. BREAKSØ.AID

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/* Prog Name : BREAKSØ.AID : December 83 : M. Kadri Ozyurt Date Written by For : Thesis Advisor : Professor Kodres Purpose : This %include file is one of the ten %include files ,BREAKSØ through BREAKS9, that are used to insert various parts of the programs to be tested. They are protected during the normal operation of the program under test with an if statement. Within the if statement there is a call to BREAKPTS . */

BREAKS(0): BREAKPT=0; IF STOPS(BREAKPT) THEN DO; CALL BREAKPTS; GOTO BREAKS(BREAKPT); END /*IF*/;

H. BREAKPTS.PLI

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/*
                   : BREAKPTS.PLI
Prog
      Name
                   : December 83
Date
Written by
                   : M. Kadri Ozyurt
                   : Thesis
For
                   : Professor Kodres
Advisor
                   : This routine prompts the user that the
Purpose
breakpoint intended to stop has been reached. Then it asks
if the user wants to transfer the control over the dynamic
debuggig environment. If the answer is positive then it
calls REENTRY where the control stays thereafter.
*/
breakpts:proc external;
  dcl
           stops(0:9) bit(1) external,
           breakpt fixed bin(15) external.
           key char(1) external,
           erroron bit(1) external,
           reentry entry;
  stops(breakpt)='0'b;
  put skip list('***** breakpoint', breakpt, ' *****');
  put skip list('The execution halted and clock stopped.');
put skip list('Do you want to enter the interactive debugging',
                    'environment(Y/N)?');
  get list(key);
if (key='I') ! (key='y') then
           call reentry();
  else
           erroron='Ø'b;
  /*end if*/
end breakpts;
```

I. TIMES.AID

/*
Prog Name : TIMES.AID
Date : December 83
Written by : M. Kadri Ozyurt
For : Thesis
Advisor : Professor Kodres
Purpose : This %include file is inserted to
WAR.FLI to test the execution times of the individual
system routines.

put skip(2) list ('Do you want to measure the execution', , times of the modules (T/N)?'); get list (key); do while ((key='Y') ! (key='y')); put skip(2) list ('Enter the number of iterations you want ', (max 32,767).'); on error begin; put list('*** bad entry, try again.'); goto times1; end /*error*/; on fixedoverflow begin; put list('*** too large, try again.'); goto times1; end /*fixedoverflow*/; times1: put skip(2) list('>'); get list(h); revert error; revert fixedoverflow; do i=1 to 4; put skip(2) list('Ready!! Press any key to start' ,'the time check of the module'); if i=1 then put list ('IDLE.'); else if i=2 then put list(' STATUS.'); else if i=3 then put list(' TACTICAL.'); else put list(' DISPLAY.'); get list(key); do j=1 to h while (i=1); call idle; end; do j=1 to h while (i=2); call status; end; do j=1 to h while (i=3); call tactical; end; do j=1 to h while (i=4); call display; end;

```
put skip(2)list('The end of the execution .');
put skip(2)list('Enter the time measured in '
                                                  seconds.'):
         on error begin;
                  put list('*** bad entry, try again');
                  go to times2;
         end /*error*/;
          on fixedoverflow begin;
                  put list('*** bad entry,try again.');
                  goto times2;
         end /*fixedoverflow*/;
times2:
         put skip(2) list('>');
          get list(j);
          revert error;
          revert fixedoverflow;
          begin;
            dcl duration float;
            duration=float(j)/float(h);
            end;
 end /*do*/;
 put skip(2) list('Do you want another run (Y/N)?');
 get list(key);
enā /*do*/;
```

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APPENDIX G

A SAMPLE SUBROUTINE TESTING

/* Prog : P.PLI Name : December 83 Date **Vritten** : M. Kadri Ozyurt by For : Thesis Advisor : Professor Kodres : This program is written to test Purpose individual procedures in an interactive manner. At each condition The PL/I iteration new values are asked. ON Wrong entries. The endless loop can be terminated either Z from the terminal. In this particular example the procedure DRAW inside the body of DISPLAY.PLI is tested by making it external for the test purposes. */ p:proc options(main); dcl (u,v)(0:10) fixed bin(15), (1,x,y) fixed bin(15), rub char(1) external, d fixed bin(7), draw entry ((0:10) fixed bin(15), (0:10) fixed bin(15) .fixed bin(7)); on error begin; put skip list('*** bad value, try again'); goto reentry; end; on fixedoverflow begin; put skip list('*** too large, try again'); goto reentry; end;

```
reentry:
reentry:
rub=ascii(127);
do while('1'b);
    put list('^^ X');
    put skip list('enter x and y');
    put skip list('>');
    get list(x,y);
    put list('Z');
    call cen(x x u x);
       call gen(x,y,u,v);
       d=1:
       call draw(u,v,d);
call delay;
       d=Ø;
       call draw(u.v.d);
end;
delay:proc;
   dcl (i,j) fixed bin(15);
   do 1=1 to 30000;
                 do j=1 to 2;
end;
end;
end;
gen:proc(x,y,u,v);
   dcl
                 (u,v)(0:10) fixed bin(15),
(x,y) fixed bin(15);
                              v(Ø)=y;
v(1)=y+8;
u(0)=x+8;
u(1)=x;
                              v(2)=y;
u(2)=1-8;
                              v(3)=y-8;
u(3)=x;
                               v(4)=y;
u(4)=1+8;
u(5)=-1;
                               v(5)=-1;
end;
```

end p;

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APPENDIX H

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A SAMPLE PROGRAM TESTING

/* Prog : IDLE.PLI Name : December 83 Date : M. Kadri Ozyurt Written by : Thesis For : Professor Kodres Advisor : This is the testing version of the Purpose procedure IDLE.PLI under the dynamic debugging module. After the correct result from the test had been taken, the final version of the procedure was made simply removing the segment of code in between the comment lines. In order to test the program, an interactive main procedure as in Appendix G was written.

*/

Sterry P

IDLE: PROCEDURE EXTERNAL;

/* DCL */ %INCLUDE 'CONST.INP'; %INCLUDE 'GLOBALS.INP'; /******* DEBUG AID *******/ %INCLUDE 'LOCALS.AID'; %INCLUDE 'BRHAND.AID'; %INCLUDE 'BREAKSØ.AID';

/***** END AID ******/

DO D=Ø TO 5; CALL ATOD (D,ARG(D)); END /*DO*/;

```
/*at this point the A/D output values are fixed bin(7)
         The following sequence converts those to fixed
values.
decimal values*/
/****** DEBUG AID ******/
%INCLUDE 'BREAKS1.AID';
/***** END AID *******/
      COURSE(OWN) = ARG(0);
      SPEED(OWN)=ARG(4);
      COURSE(KNOWN)=ARG(2);
      SPBED(KNOWN)=ARG(3);
      AZ = ARG(1);
      ALT=ARG(5);
/***** DEBUG AID *****/
     %INCLUDE 'BREAKS2.AID';
/****** END AID *******/
/*the following sequence converts A/D values to real time
values by using appropriate proportionality constants*/
      COURSE(OWN) = COURSE(OWN) * K;
      COURSE(KNOWN) = COURSE(KNOWN) * K;
      AZ = AZ * K;
      IF COURSE(OWN)<2.0 THEN
          COURSE(OWN) = COURSE(OWN) + TWO_PI ;
      IF COURSE(KNOWN)<0.0 THEN
          COURSE(KNOWN) = COURSE(KNOWN) + TWO FI ;
      IF AZ<0.0 TEEN
          AZ = AZ + TWO PI;
       IF ALT>90.0 THEN
          ALT = 90.0;
/***** DEBUG AID ******/
     %INCLUDE 'BREAKS3.AID';
/***** END AID *******/
          SPEED(OWN) = SPEED(OWN)/L;
          SPEED(KNOWN) = SPEED(KNOWN) / L;
/*ownship speed computations*/
      VX_OWN = SPEED(OWN) * SIND(COURSE(OWN));
      VY OWN = SPEED(OWN) * COSD(COURSE(OWN));
/***** DEBUG AID *****/
     %INCLUDE 'BREAKS4.AID';
/****** END AID *******/
```

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```
/*when not have fired, the following makes the tallistic
computations*/
      IF
           FIRED THEN
         BEGIN;
          T_OF = 2.0 * VM * SIND(ALT) / G;
/***** DEBUG AID ******/
     %INCLUDE 'BREAKS5.AID';
'****** END AID *******/
          VR = VM * COSD(ALT);
          R = VR * T OF;
/***** DEBUG AID ******/
     XINCLUDE 'BREAKS6.AID';
/****** END AID *******/
          X_AIM(OWN) = R * SIND(AZ);
          T_{AIM}(OWN) = R * COSD(AZ);
          \mathbf{X}^{-} \mathbf{GUN} = \emptyset.0;
          Y^{-}GUN = 0.0;
/***** DEBUG AID *****/
     XINCLUDE 'BREAKS7.AID';
/****** END AID *******/
          VX ROUND = VR * SIND(AZ);
          VY ROUND = VR * COSD(AZ);
/***** DEBUG AID ******/
     %INCLUDE 'BREAKS8.AID';
\mathbf{T} = \mathbf{T}_{\mathbf{O}}\mathbf{F};
      END /*IF*/;
/***** DEBUG AID *****/
     %INCLUDE 'BREAKS9.AID';
/***** END AID ******/
END IDLE;
```

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