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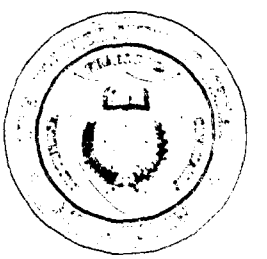

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UNCONSTRAINED MINIMIZATION BY INTERPOLATION: RATES OF CONVERGENCE

by

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ABSTRACT

We analyze the rate of convergence of a class of algorithms based on n-dimensional interpolation. In particular, we present a class of algorithms which use first order information only, while maintaining quadratic convergence.

KEY WORDS

Unconstrained minimization, Nonpolynomial interpolation, Convergence rates, Solution of equations.

1. Introduction

Most of the commonly used algorithms for the unconstrained minimization of $f: R^n \to R$ for $n > 1$, are descent methods. These are based on the iteration $x_{i+1} = x_i + \alpha_i d_i$, where $d_i \in R^n$ is a search direction, and $\alpha_i \in R$ is the step-size, usually determined by a line search as a minimizer of $f(x_i + \alpha d_i)$ over $\alpha \geq 0$.

The exception is Newton's method, which is based on interpolating f by a quadratic and minimizing this quadratic at each step of the algorithm.

For a general discussion of unconstrained minimization techniques see [1,8]. The following points are relevant to our discussion.

The classical steepest descent method uses first order (i.e., gradient) information only. Its main drawback is its linear rate of convergence.

Newton's method converges quadratically. However, it necessitates the costly computation of the Hessian at each iteration.

Other methods based on first order information are known to converge superlinearly (e.g., [6]).

Many of these methods approximate Newton's method in the sense that the search direction they generate can be shown to be a direction along which an appropriate quadratic is minimized.

A different approach is based on the unfounded assumption that algorithms having the finite termination property (i.e., solution in a finite number of steps) for a class of functions wider than the class of quadratics, are faster than those having the quadratic termination property. Thus, Jacobson and Oksman [7] generalize from quadratic termination to homogeneous functions termination. This was further generalized (see e.g., [4]).

Another step toward discarding the quadratic model has recently been taken by Davidon [5]. His motivation, however, partly coincides with ours.

In this paper, we analyze the rate of convergence of n-dimensional interpolation algorithms for unconstrained minimization. We note the following.

Most of the commonly used algorithms for one-dimensional minimization are based on polynomial interpolation. It is well known that Newton's method in this case is inefficient in the sense that quadratic convergence can be achieved using first order information only (e.g., [8, p. 142]), by using two interpolation points rather than one. This should make one doubt whether Newton's method is a suitable model for efficient algorithms.

Quadratics are inadequate for n-dimensional two-point interpolation with zero and first order information (see Davidon [5]). Therefore, non-polynomial interpolation is necessary. Our analysis shows that the rate of convergence is independent of the interpolating function. For interpolatory algorithms, therefore, the question whether the search direction coincides with Newton's direction or generalizes it, is irrelevant to the rate of convergence analysis. The same is true for termination properties.

The main difficulty in the analysis of n-dimensional interpolation algorithms is that the formulas for the error in n-dimensional interpolation are not suitable for this purpose. We overcome this difficulty by reducing the n-dimensional problem to an appropriate one-dimensional interpolation problem.

2. Minimization by Interpolation

The interpolation algorithm we study generates a sequence $\{x_i\}$ as follows. Let $s \geq 1$, $m \geq 0$ be fixed integers. Given $m+1$ approximants $x_0, \ldots, x_{m+1}$ to the solution of

$$\nabla f(x^*) = 0 \,, \tag{1}$$

we use $x_i, x_{i-1}, \ldots, x_{i-m}$ to construct a new approximant x_{i+1}. First we interpolate f by T requiring

$$T^{(k)}(x_{i-j}) = f^{(k)}(x_{i-j}) \qquad j = 0,\ldots,m;\ k = 0,\ldots,s-1 \ . \tag{2}$$

Here $f^{(1)} = \nabla f$, $f^{(2)} = \nabla^2 f$ etc., and $T: R^n \to R$ is assumed to depend on some parameters to be determined by (2). The new point x_{i+1} is determined by

$$\nabla T(x_{i+1}) = 0 \ . \tag{3}$$

In the following, we assume that equations (1)-(3) have solutions.

We define the rate (or order) of convergence of a sequence $\{x_i\}$ converging to x^* as the number p (if it exists) such that

$$\frac{\|x_{i+1}-x^*\|}{\|x_i-x^*\|^p} \to C \neq 0 \ .$$

Here $\|\cdot\|$ is a fixed arbitrary norm. Ortega and Rheinboldt [9, §9] refer to the rate p defined above as the C-order of the sequence $\{x_i\}$. When it exists, it coincides with their Q- and R-orders. We will unify our results for the C-, Q- and R-orders through the use of the C-order of convergence.

We derive the rate of convergence of the n-dimensional interpolating algorithm by establishing some difference relations for the errors $\|x_i-x^*\|$. To derive the basic difference relation we need, we pass a curve in R^n through the points x^* and $x_{i+1}, x_i, \ldots, x_{i-m}$, i.e., we determine a function $\psi: R \to R^n$ such that

$$\begin{aligned} &\psi(t_{i-j}) = x_{i-j} \qquad j = -1,0,1,\ldots,m \ , \\ &\psi(t^*) = x^* \ , \end{aligned} \tag{4}$$

where the parameter t is chosen so that

(5) $$t_{i-j} = \|x_{i-j} - x^*\| , \quad t^* = \|x^* - X^*\| = 0 .$$

We will later discuss this construction. Note, however, that the construction of Ψ is a part of the analysis of the properties of the algorithm, not a part of the algorithm itself.

Henceforth we will assume $0 \neq t_{i-j} \neq t_{i-k}$ for all i and $j,k = -1, 0, \ldots, m$. This is a natural assumption. If $t_i = 0$ for some i, the algorithm terminates, while the assumption $t_{i-j} \neq t_{i-k}$ $j,k = -1, \ldots, m$ has to be made even in the one-dimensional case (cf. Traub [14, Ch. 4]).

New define $\theta(t) = T(\Psi(t)), \phi(t) = f(\Psi(t))$. Note that Ψ, ϕ, θ depend on i. Following Traub [14] and Tamir [12,13], we will not make this dependence explicit in order to simplify notation. Equations (2) and (4) imply

(6) $$\theta^{(k)}(t_{i-j}) = \phi^{(k)}(t_{i-j}) , \quad j = 0, \ldots, m; \; k = 0, \ldots, s-1 .$$

It follows that (6) defines one-dimensional interpolation for which a convenient error formula exists (see Ostrowski [10, p. 12]). Henceforth, we assume $\theta, \phi \in C^{r+1}$ in a neighborhood of $t^* = 0$, where $r = s(m+1)$. Using the one-dimensional error formula we have

(7) $$\phi(t) - \theta(t) = \frac{\phi^{(r)}(\xi) - \theta^{(r)}(\xi)}{r!} \prod_{j=0}^{m} (t - t_{i-j})^s ,$$

where ξ is a point in the interval determined by $t, t_i, \ldots, t_{i-m}$. Note that formula (7) holds for general (not necessarily polynomial) interpolation.

We now differentiate (7) and set $t = 0$. From (1) and (2) we have $\phi'(0) = \theta'(t_{i+1}) = 0$, so that $\phi'(0) - \theta'(0) = -\theta'(0) = \theta'(t_{i+1}) - \theta'(0) = t_{i+1}\theta''(\zeta)$, where ζ is a point between t_{i+1} and 0. Differentiating the right hand side of (7) using Ralston's result [11] on the differentiation of the error term generalized for the hyperosculatory case (see [2]), we finally have

Lemma 1. Under the assumptions made above, the errors in the n-dimensional interpolation algorithm satisfy the difference relation

$$t_{i+1} = M_i \sum_{k=0}^{m} t_{i-k}^{s-1} \prod_{\substack{j=0 \\ j\neq k}}^{m} t_{i-j}^{s} + N_i \prod_{j=0}^{m} t_{i-j}^{s}, \tag{8}$$

where

$$M_i = \frac{M(\xi_i(t_{i+1}))(-1)^{r-1}\cdot s}{\theta'(\zeta(t_{i+1}))}, \qquad N_i = \frac{N(\eta_i(t_{i+1}))\cdot(-1)^{r}}{\theta'(\zeta(t_{i+1}))}$$

$$M(t) = \frac{\phi^{(r)}(t)-\theta^{(r)}(t)}{r!}, \qquad N(t) = \frac{\phi^{(r+1)}(t)-\theta^{(r+1)}(t)}{(r+1)!}$$

and where $\xi_i(t)$, $\eta_i(t)$ are in the interval determined by $t, t_{i+1}, \ldots, t_{i-m}$, and $\zeta(t_{i+1})$ is in the interval determined by $t_{i+1}, 0$. □

It follows from (8) that if the initial errors $t_0, \ldots, t_n$ are small enough, and if the coefficients $\{M_i\}$, $\{N_i\}$ are bounded, the sequence $\{t_i\}$ converges to zero, i.e., $x_i \to x^*$. Moreover, if $s \geq 2$, (8) implies

$$\frac{t_{i+1}}{t_i} \to 0, \tag{9}$$

(i.e., superlinear convergence). If $s = 1$, we assume $m \geq 2$. For $m = 2$, (8) is the basic difference relation governing the behavior of the Quadratic Fit algorithm, which is known to converge superlinearly (see Theorem 3.4.1 in Brent [3]). It is evident from (8) that the rate for $m > 2$ is not less than the rate for $m = 2$. Therefore, (9) holds for all $s \geq 1$, $m \geq 0$ if $r = s(m+1) \geq 3$. Rewriting (8) in the form

$$t_{i+1} = M_i t_i^{s-1} \prod_{j=1}^{m} t_{i-j}^{s}\left[1 + \sum_{k=1}^{n} \frac{t_i}{t_{i-k}} + \frac{N_i}{M_i} t_i\right]$$

we finally have

Lemma 2. Under the assumptions made above, and if $M_i \to M \neq 0$, the sequence t_i satisfies the difference relation

$$t_{i+1} = A_{i+1} t_i^{s-1} \prod_{j=1}^{m} t_{i-j}^{s} \tag{10}$$

with $A_{i+1} \to M$. □

Defining $y_i = \log|t_i|$ and $B_i = \log|A_i|$, (10) implies the difference equation

$$y_{i+1} - (s-1)y_i - s \sum_{j=1}^{m} y_{i-j} = B_{i+1}$$

with indicial equation

$$t^{m+1} - (s-1)t^m - s \sum_{j=0}^{m-1} t^j = 0 , \tag{11}$$

where the sum in (11) is taken as zero if $m = 0$.

Tamir [12,13] proves that under our assumptions the C-rate of convergence of the sequence $\{t_i\}$ (hence $\{x_i\}$) is given by the unique positive root of the indicial equation (11). In this case, the C-, Q-, and R-rates of convergence are exactly p, where p is the positive solution of (11).

If the limit of M_i exists and is zero, or if this limit does not exist, but the sequences $\{M_i\}$, $\{N_i\}$ are bounded, equation (8) can be rewritten in the form

$$t_{i+1} = t_i^{s-1} \prod_{j=1}^{m} t_{i-j}^{s} \{M_i[1 + \sum_{k=1}^{n} \frac{t_i}{t_{i-k}}] + N_i t_i\} ,$$

which implies that the Q- and R-rates of convergence are still at least p.

We now summarize our results:

Theorem 1. If equations (1)-(4) have solutions, the functions $T, f, \psi \in C^{(r+1)}$, if the sequences $\{M_i\}$ and $\{N_i\}$ are bounded, and if the initial errors of the interpolation algorithm are small enough, then the sequence $\{x_i\}$ converges to the solution x^* with C- (when it exists), Q- and R-rates of convergence at least p, where p is the unique positive solution of (11).

Corollary 1. The rate of convergence of the sequence generated by the interpolation algorithm is independent of the interpolating function.

The reader should note that while the curve $\psi(t)$ may be constructed in infinitely many ways, it is sufficient to establish the existence of just one such curve (for each i). We now turn our attention to this problem.

Lemma 3. If $t_0 = 0$ and $t_i \neq t_j$ for $i, j = 0, 1, \ldots, k$ the determinant

$$D = \begin{vmatrix} 1 & t_0^2 & t_0^3 & \cdots & t_0^{k+1} \\ 1 & t_1^2 & t_1^3 & \cdots & t_1^{k+1} \\ \cdots & \cdots & \cdots & \cdots & \cdots \\ 1 & t_k^2 & t_k^3 & \cdots & t_k^{k+1} \end{vmatrix}$$

does not vanish.

Proof. Since $t_0 = 0$ we have

$$D = \begin{vmatrix} t_1^2 & t_1^3 & \dots & t_1^{k+1} \\ t_2^2 & t_2^3 & \dots & t_2^{k+1} \\ \dots & \dots & \dots & \dots \\ t_k^2 & t_k^3 & \dots & t_k^{k+1} \end{vmatrix} = \prod_{i=1}^{k} t_i^2 \begin{vmatrix} 1 & t_1 & \dots & t_1^{k-1} \\ 1 & t_2 & \dots & t_2^{k-1} \\ \dots & \dots & \dots & \dots \\ 1 & t_k & \dots & t_k^{k-1} \end{vmatrix} .$$

The last determinant is a Vandermonde determinant, and since $t_i \neq t_j$ $i,j = 1,\dots,k$, it does not vanish. □

Lemma 4. Let $p(t) = \alpha t + \sum_{\substack{i=0 \\ i\neq 1}}^{k+1} a_i t^i$, and let $t_0 = 0$, $t_i \neq t_j$ for $i,j = 0,1,\dots,k$. Then the system of equations

$$p(t_j) = \beta_j \qquad j = 0,\dots,k ,$$

for the unknowns a_i $i = 0,2,3,\dots,k+1$ has a solution for all α,β_j.

Proof. This is an immediate consequence of Lemma 3. □

Since $p'(0) = \alpha$ and $p(t)$ is a polynomial, taking $k = m+2$ in Lemma 4 we have

Lemma 5. If the errors satisfy $0 \neq t_{i-j} \neq t_{i-k}$ for all i and $j,k = -1,0,\dots,m$, there exists a curve $\psi \in C^\infty$ satisfying (4). Moreover, we may require $\dot{\psi}(0) = a$ with $a \in R^n$ arbitrary. □

We can now state our main result.

Theorem 2. Assume that equations (1)-(4) have solutions, f has continuous derivatives of order $r+1$, the parameters of T depend continuously on the data through

(2), T has continuous derivatives of order $r+1$ for the appropriate values of the parameters. Assume also that $\nabla^2 f(x^*) \neq 0$ and $0 \neq t_{i-j} \neq t_{i-k}$ for all i and $j,k = -1, 0, \ldots, m$, then if the initial errors of the interpolation algorithm are small enough, the sequence $\{x_i\}$ converges to the solution x^* with C- (when it exists) Q- and R-rates of convergence at least p, where p is the unique positive solution of (11).

Proof. By the above it is sufficient to show that under the assumptions of the theorem, the sequences $\{M_i\}$, $\{N_i\}$ are bounded. This is the case if $\dot{\psi}(0)^T \nabla^2 f(x^*) \dot{\psi}(0) \neq 0$, for which it is sufficient to choose the vector a in Lemma 5 as an eigenvector of $\nabla^2 f(x^*)$ corresponding to a nonzero eigenvalue. □

Summary and Conclusion

We have shown that the rate of convergence of n-dimensional interpolation algorithms is inherited from the underlying one-dimensional interpolation, that it is independent of the interpolating functions, and is given by the unique solution of the equation

$$t^{m+1} - (s-1)t^m - s \sum_{j=0}^{m-1} t^j = 0 , \tag{12}$$

where $m+1$ interpolation points and s derivatives (of orders zero to s-1) are used.

Our work is based on the results of Traub [14] and Ostrowski [10] for the one-dimensional root-finding problem. Tamir [12,13] adapted these results for the minimization problem. In [12] he studies the rate of convergence of algorithms using function values only $(m = 0)$ with a superfluous assumption and a false conjecture. This detailed analysis is repeated in [13] for the case $m > 0$. He treats polynomial interpolation only, and shows that for fixed s and $m \to \infty$, the rate p tends to

(13) $$\frac{s}{2} + \sqrt{(\frac{s}{2})^2 + 1} \ .$$

However, he neglects to realize the effect of memory on the rate of convergence, which is implied by (11) and (12).

Indeed, for fixed s, the rate is obtained for $m = 0$ and $m = 1$ by solving the indicial equations $t - (s-1) = 0$ and $t^2 - (s-1)t - s = 0$, respectively. Therefore, $p = s - 1$ for $m = 0$, $p = s$ for $m = 1$ and $p \to \frac{s}{2} + \sqrt{(\frac{s}{2})^2 + 1}$ for $m \to \infty$. It follows that algorithms using more than two interpolation points are inefficient, and two-point algorithms are substantially faster than one-point algorithms.

In particular for $m = 1$ and $s = 2$ we have a two-point algorithm using <u>first-order</u> information with <u>second-order</u> rate of convergence (which is a well-known result in the one-dimensional case).

Note that no line search is needed in this class of algorithms, and they they may be designed to locate saddle points rather than minimum points. A line search, however, may serve as part of a globalizing procedure.

Compare also the discussion in Davidon [5] regarding the difficulty of determining the effect of memory on the performance of descent algorithms.

We have not computed the asymptotic error constant, since it depends on the norm used (see Ortega and Rheinboldt [9]). This can be computed, however, under the appropriate assumptions (cf. Tamir [12,13]). We have also made no attempt at giving the strongest results (i.e., the weakest assumptions) possible. Compare, for example, Brent [3].

Finally, note that Theorem 1 holds for infinite dimensional spaces, and that our analysis is applicable with the obvious modifications to the solution of systems of equations, for which the indicial equation analogous to (11) is

$$t^{m+1} - s \sum_{j=0}^{m} t^j = 0 \ .$$

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