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In-House Report September 1979

NONLINEAR CIRCUIT ANALYSIS PROGRAM (NCAP) DOCUMENTATION Programmer's Manual

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UNCLASSIFIED SECURITY CLASSIFICATION OF THIS PAGE (When Date Entered) READ INSTRUCTIONS BEFORE COMPLETING FORM REPORT DOCUMENTATION PAGE REPORT NUMBER 2. GOVT ACCESSION NO. 3. RECIPIENT'S CATALOG NUMBER RADC-TR-79-245 Vol III (of three) PERIOD COVERED NONLINEAR CIRCUIT ANALYSIS PROGRAM (NCAP) end Phase Report DOCUMENTATION Volume till May 178 - Jul Programmer's Manual CONTRACT OR GRANT NUMBER(s) . AUTHOR(A) Jon B. Valente In-House and Contract Sharon Stratakos F30602-79-C-0011 JUN PERFORMING ORGANIZATION NAME AND ADDRESS Rome Air Development Center (RBCT) Griffiss AFB NY 13441 62702F Syracuse University, Syracuse NY 13210 23380314/23380317 11. CONTROLLING OFFICE NAME AND ADDRESS TZ. REPORT DATE Rome Air Development Center (RBCT) 13. NUMBER OF PAG Griffiss AFB NY 13441 15. SECURITY CLASS. (of this report 14. MONITORING AGENCY NAME & ADDRESS(if different from Controlling Office) Same UNCLASSIFIED 154. DECLASSIFICATION DOWNGRADING 16. DISTRIBUTION STATEMENT (of this Report) Approved for public release; distribution unlimited. 17. DISTRIBUTION STATEMENT (of the abstract entered in Block 20, if different from Report) Same 18. SUPPLEMENTARY NOTES RADC Project Engineer: Jon B. Valente (RBCT) 19. KEY WORDS (Continue on reverse side if necessary and identify by block number) Electromagnetic Compatibility Computer Program Nonlinear Analysis Volterra Analysis Nonlinear Circuit Analysis Computer-aided Circuit Analysis 20. ABSTRACT (Continue on reverse side if necessary and identify by block number) The Nonlinear Circuit Analysis Program (NCAP) is a circuit analysis code which uses the Volterra approach to solve for the transfer functions and node voltages of nonlinear circuits. To increase the transportability, the code was written in ANSI FORTRAN. The code was revised and documented in a joint effort using both in-house and contracted manpower. This documentation is a result of

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hual, the User's Manual, and the Programmer's Manual.

hat effort. The documentation is made up of three volumes: the Engineering

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The Engineering Manual, Volume I, contains the introduction to the documentation, a description of the theory upon which NCAP is based, and the procedure for obtaining the input parameter data.

The User's Manual, Volume II, contains a detailed description of the NCAP input language. The description includes the input commands and data requirements. Examples are used throughout the manual to aid the user in understanding the input language.

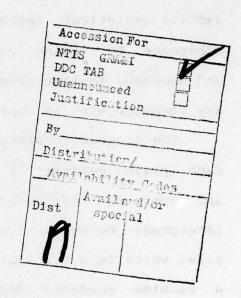
The Programmer's Manual, Volume III, contains a subroutine by subroutine description of the code. Each subroutine has a trace map, functional flow diagram and narrative which will help the programmer to understnad the structure of NCAP.

Table of Contents

INTRODUCTION		
MAIN		1
ALPHA		2
DATOUT		3
DECIF		4
DECIM		5
DRVOUT		3 4 5 . 6 7
ECIN		7
ENDIN		8
GENIN		9
GENOUT		10
INTEG		11
JFETIN		12
LCIN		13
LDSIN		14
MODGIN		15
MODIN		16
NCIN		17
NDSIN		18
NODEIN		19
PACKL		20
PACKR		21
PLOTIN		22
PRCSS		23
PSIN		24
PTIN		25
SCAN		26
SCIN		27
SDIN		28
SHIFT		29
TRNIN		30
VDIN	Accession For	31
VTIN		32
	NTIS GRA&I DDC TAB	
PHASE1		33
CPMTJF	Justification	34
CPMTSD	JUSCIII	35
CPMTVD		36
CPPT	By	37
CPT		38
CPVT	Aveilability Codes	39
DERIDS	Linail Bud Or	40
FIDS	annoin!	41
IMLOC	Dist	42
MTPT	NOSA	43
MTT MTVT	N RU	44
MTVT	14	45

PHASE 2		46
MTLDS		47
MTNDS		48
ZIADD		49
PHASE3		50
EETA CONTRL		51
CROSS		52
CURGEN		53
CURJF		54 55
CURNC		
CURNDS		56 57
CURPT		58
CURSD		59
CURT		60
CURVD		61
CURVT		62
FNCTN		63
FRPRM		64
LOCTE		65
QFN		66
SSCODE		67
TRIANG		68
UPFRQ		69
UPSEQ		70
PHASE4		71
HHMMSS		72
PRNTIM		73
PRSET		74
TFTONV		75
PHASE5		76
CHNGE		77
OWACRE		
PHASE6		78
PHASE7		79
THROLT		19
UTILITY		
CXADD		80
CXDIV		81
CXMPY		82
CXPOL		83
CXSUB		84
DATARD		85
DATAWR		86
DRIVRD		87
DRIVWR		88
ERROR	1	89
FUN		90
LSHIFT		91
PSHIFT		92
	iv	

SSCOD2 TFRD TFWR TIMOUT



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INTRODUCTION

RADC NCAP is a Fortran program which computes the transfer functions of electronic circuits. Although the program is large and its analytical technique complex, the modular structure, adherence to naming conventions for subprograms and variables, and numerous in-line comments allow NCAP to be readily adapted to any computer with an appropriate Fortran compiler.

The program consists of eight phases, numbered 0 through 7. Each phase performs a distinct portion of the circuit analysis and operates independently of the other phases. The only interphase communication is by shared disk files: the driver file, which is a translation of the NCAP asterisk input cards to a machine readable description of the circuit analyses to be performed, and the data file, which contains all circuit element input data, calculated device parameters, admittance matrices, and transfer function vectors. Although several other disk files are used by NCAP, their function is to conserve core storage and to ease the transmission of internally generated data between the which comprise individual subprograms phases. Detailed descriptions of the NCAP disk files are contained in the narrative description of Phase 1.

Phase Ø is the input processor for NCAP. It reads and interprets the input deck, mapping the input cards to appropriate driver and data file records. Phase 1 calculates the device parameters for each circuit element, collects and tabulates the circuit's frequencies, and determines the size of the admittance

matrices. Phase 2 constructs the admittance matrices, one for each possible combination of the circuit's frequencies. Phase 3 constructs the current vectors and calculates the transfer functions for each frequency combination. Phase 4 prints the results from the circuit analysis performed in Phases 1-3 and controls frequency sweeping. Phase 5 controls linear component sweeping, Phase 6 controls device modification, and Phase 7 controls generator modification.

Since numerous circuit analyses may be specified by a single NCAP input deck, the path of execution through the program phases is not necessarily sequential. Execution always begins at Phase Ø and proceeds sequentially through Phases 1-4 to perform the first circuit analysis. From Phase 4, program execution either reverts back to Phase 1 to initiate a new analysis if frequency sweeping is specified, or proceeds to Phase 5 if frequency sweeping is not specified or after all such sweeps have been satisfied. In a similar fashion, Phases 5, 6, and 7 may either cycle back to Phase 1 or proceed on the next phase depending on the linear component sweeping, device modification, and generator modifications specified in the input deck. Program execution ends with Phase 7 after the last (if any) generator modification has been effected.

Each phase is composed of a principle subprogram which controls its general operation, a group of secondary subprograms which perform specific operations for individual circuit elements or NCAP functions, and in some cases, additional support subprograms which perform operations unique to that phase. The

program is organized sequentially by phases. Within each phase, the principle subprogram appears first, followed by the secondary and support subprograms in alphabetical order. A group of shared support subprograms, such as those which perform disk input/output or complex arithmetic, follow Phase 7 and appear in alphabetical order.

The principle subprograms of each phase are subroutines, with the exception of Phase Ø whose principle, in order to satisfy the requirements of Fortran, is NCAP's main program. These principle subroutines are named PHASEØ, PHASE1,...,PHASE7. With the exception of but two function subprograms, the remainder of the NCAP subprograms are subroutines.

Wherever possible, the subprograms are named according to specific conventions. Subprograms which perform specific functions related to circuit elements are prefixed or suffixied with a device identifier:

GEN = Generator

JFET (or JF) = Junction Field Effect Transistor

LC = Linear Components

LDS = Nonlinear Dependent Source

NC = Nonlinear Components

NDS = Nonlinear Dependent Source

PT = Vacuum Pentode

SD = Semiconductor Diode

T (or TRN) = Bipolar Junction Transistor

VD = Vacuum Diode

VT = Vacuum Triode

Furthermore, within each phase the secondary subprogram names contain functional identifiers:

IN = Read and interpret input cards

CP = Calculate parameters

MT = Create matrix elements

CUR = Calculate current elements

Together the device and functional identifier describe the purpose of the subprogram: GENIN = input generator card sequence, CPPT = calculate pentode parameters, MTT = create transistor matrix elements, CPMTVD = calculate parameters and create matrix elements for vacuum diode, CURNDS = calculate current elements for nonlinear dependent source.

At the support level, subprograms which perform complex arithmetic are prefixed by CX (CXADD, CXDIV, etc), while disk I/O routines are suffixed by RD and WR (DATARD, DRIVWR, etc.). Support subprograms whose functions are too specific to be categorized are named as descriptively as possible: LOCTF = locate transfer function, FRPRM = create frequency permutation.

The program code for subroutines PHASE1 through PHASE7 are all organized alike. Execution through these routines is controlled by reading and processing the driver file records sequentially. Each driver record contains a functional identifier or mode, which serves as the index of a computed GO TO, selecting the proper code segment to process that record. The coding for each driver function is arranged numerically by mode within the subroutine and begins with the statement number equal to the value of the mode. Additional statement numbers

within a code segment are assigned in increments of 100. For example, a section of transistor code would begin with statement 9 (the transistor driver mode), and proceed through 109, 209, 309, and so on.

In a similar fashion, the IN family of subroutines (input card processors) share a common organization. Execution through these subroutines is based on a computed GO TO using the card type identifier as an index. The coding for each card type is arranged numerically within the subroutines and statement numbers are allocated in increments of 100 within code segments.

The narrative descriptions of the NCAP subprograms which follow are arranged in the order in which they appear in the program by phases, and within phases, by alphabetical order. Each subprogram description contains a brief statement of purpose, followed by a variable list, subroutines called, calling programs, and a detailed narrative of the program code. Wherever possible mathematical algorithms are summarized and tables of all possible computed results are presented.

To avoid repetition, variables which are used globally are listed only in the Phase Ø description or in the first principle subprogram in which they are used. In the secondary and support subprogram descriptions, only local variables (or in some cases less frequently used global variables) are listed.

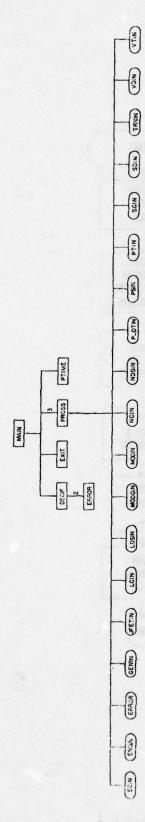
Machine dependent code has been clearly identified in both the program listing and narrative descriptions in order to ease the adaptation of NCAP to various computer systems.

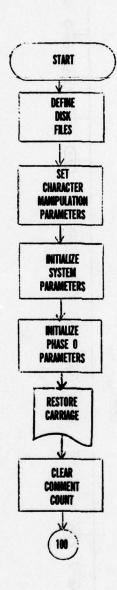
Functional flow diagrams for each subprogram have been

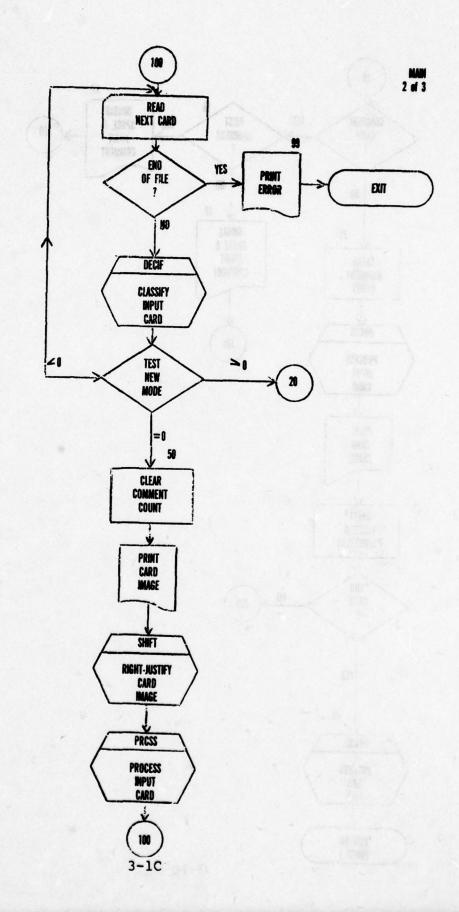
included to provide an insight into the general flow of logic through the program. They are not intended to describe the analytical techniques employed in the program or to depict the program code in detail. Individual symbols of the flow diagrams generally represent several lines of program code or complete code segments. Wherever possible, diagram symbols have been labelled with appropriate statement numbers from the program.

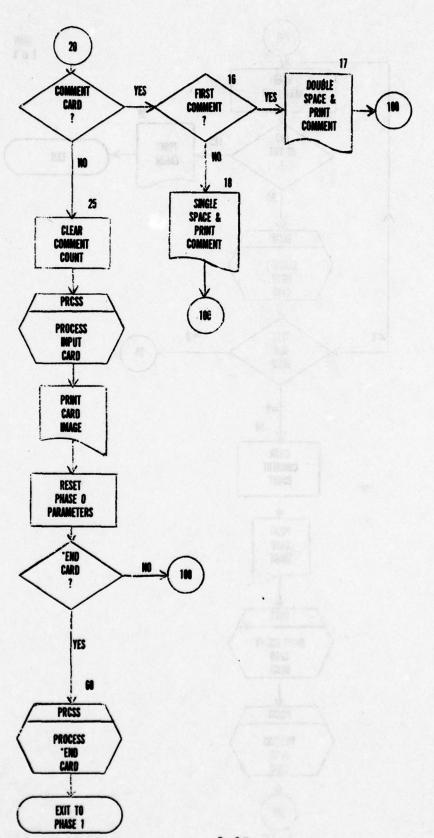
MAIN

ALPHA DATOUT DECIF DECIM DRVOUT ECIN ENDIN GENIN GENOUT INTEG **JFETIN** LCIN LDSIN MODGIN MODIN NCIN NDSIN NODEIN PACKL PACKR PLOTIN PRCSS PSIN PTIN SCAN SCIN SHIFT TRNIN VDIN VTIN









MAIN 3 of 3 NAME: MAIN

TYPE: MAIN PROGRAM

GENERAL PURPOSE:

Disk file definition, system initialization, control of input phase

VARIABLES:

BUFF = I/O buffer for disk records

CARD = Input buffer for card image. One character
 per word, alpha representation, left-justified

DATREC = Next available location on data file (20)

DRVREC = Driver file (21) address pointer

IMAGE = Print buffer for input card images

INP = Logical unit number of card reader

IOUT = Logical unit number of line printer

JAPT = Input error switch: off = 0, on = 1

LNGTH = Length of data record in words

MAXIT = Total number of frequency sweep iterations

MDESV = Mode of previous *-card

MISC = Driver record parameters unique to each mode

MODE = Identifies circuit element, device model, or NCAP function

l = Drive file header

2 = Comment or no-op

3 = Generator

- 4 = Linear Components
- 5 = Nonlinear Components
- 6 = Vacuum Diode
- 7 = Vacuum Triode
- 8 = Vacuum Pentode
- 9 = Pipolar Junction Transistor
- 10 = Semiconductor Diode
- 11 = Not used
- 12 = Frequency Sweep
- 13 = Start Circuit
- 14 = End Circuit
- 15 = Not used
- 16 = End
- 17 = Print Select
- 18 = Modify
- 19 = Not used
- 20 = Junction Field Effect Transistor
- 21 = Generator Modification
- 22 = Plot
- 23 = Not used
- 24 = Linear Component Sweep
- 25 = Linear Dependent Source
- 26 = Nonlinear Dependent Source
- MODFY = Modify switch : 0 = Modify not in effect
 - 1 = Modify in effect
- NAPT = Error switch for node card.
 - 0 = Node card present

1 = Node card missing

NBCHAR = Number of bits per character

NBWORD = Number of bits per word

NCHAR = Number of characters per word

NCOM = Number of successive comment cards

NCURV = Number of curves defined by circuit analysis

NEW = Identifies input card being processed:

-1 = Unidentified card, card out of order

0 = Parameter card

1 = Not used

2 = Comment card

>2 = *-card

NFR = Number of frequency values for abcissa of plot

NIT = Frequency sweep iteration counter

NLIST = Print select on/off switch: 0 = on, 1 = off

NPLOT = Number of plot specifications in input deck

NVEC = Number of t.f. values for ordinate of plot

NXT = Counter for data values being input

ORDER = Total number of nodes in circuit;

Order of admittance matrix

SCREC = Location of *START CIRCUIT driver record

STADD = Data file (20) address pointer

TIME = Time elapsed in single program segment

TSTART = Start time of single program segment

TSTOP = Stop time of single program segment

TTOTAL = Total time elapsed all executions single

program segment

TYPE = Identifies type of parameter card:

- 1 = *-card
- 2 = Data values only
- 3 = Impedance
- 4 = Frequency
 - 5 = Order
 - 6 = Node
 - 7 = Amplitude
 - 8 = Resistor
 - 9 = Capacitor
 - 10 = Inductor
 - 11 = Constants
 - 12 = PP

 - 14 = AC
- 15 = Label
 - 16 = FP
 - 17 = RB
 - 18 = VC
 - 19 = CC
 - 20 = VV
 - 21 = CV
 - 22 = Off
 - 23 = Øn

SUBROUTINES CALLED:

DECIF PRCSS PTIME RANSIZE RSHIFT SHIFT

CALLING PROGRAMS.

NONE

DESCRIPTION .

NCAP's main program is the primary routine of Phase 0. It performs initialization of system parameters and reads the input deck. System initialization begins with the definition of NCAP's disk files. Because different Fortran compilers employ various methods for specifying random access files, this function is considered machine dependent. Therefore the structure and method of accessing NCAP's disk files are presented here to enable the reader to translate the particular random access I/O requirements of NCAP into appropriate file definition statements.

The data file, logical unit 20, contains the NCAP input data and all internally generated data and results such as device parameters, admittance matrices, and transfer function vectors. The data file is structured as an unlimited number of contiguous one - word physical records, each of which is addressable according to its relative location in the file. A logical record, such as a device data record or a transfer function vector, consists of an arbitrary number of physical records and is accessed by the first word of the record and the record length. The integer variables STADD and DATREC are used as data file record pointers. Subroutine DATAWR causes data to be written from core storage to file 20, while subroutine DATARD

causes data to be read from file 20 to core storage.

The driver file, logical unit 21, is a mapping of the * input cards to disk storage. It contains the information required to define the circuit and the circuit analyses to the system, and is used to control the operation of the various NCAP phases. The driver file is initially created during Phase Ø as input cards are processed, and is modified or updated as necessary during subsequent phases. Each * - card in the input deck maps to one or more driver file records, which originate from the first ten words of global common. The first three words of each driver record, MODE, STADD, and LNGTH, identify the function and define the boundaries of its associated data record on file 20. The additional words MISC(1) through MISC(7) contain parameters and pointers unique to each NCAP function. The driver file is accessed by subroutines DRIVWR and DRIVRD. Individual driver records are addressed by the record number stored in DRVREC.

File 22 contains frequency data used as the abscissa for plotting NCAP output. Each record contains ten frequency values, one for each of the ten possible plot specifications in the NCAP input deck. The number of records in the file is maintained in the integer word NFR and depends upon the number of circuit analyses performed in the NCAP run. File 22 is accessed directly by random file read/write statements.

File 23 contains the transfer function values used as the ordinate for plotting NCAP output. Each record contains twenty words and stores one complex (two-word) transfer function value

for each of the ten possible plot specifications in the NCAP input deck. The number of records in the file is maintained in the integer word NVEC and depends upon the number of circuit analysis performed in the NCAP run. File 23 is accessed directly by random file read/write statements.

File 24 contains the transfer function table created in Phase 3 by subroutine CONTRL. It is made up of the lower-order transfer function vectors required by the current generator subroutines, and consists of a large number of contiguous one-word physical records, each of which is addressable. Each logical record stores one complex transfer function vector whose length depends upon the number of nodes in the circuit. Transfer function vectors are accessed according to LOCT, the address of the first word of the record, and LT, the length of the record. Transfer function values at a particular node are accessed by the first word of the vector plus a displacement factor derived from the node number. Subroutine TFWR and TFRD perform input/output for file 24.

File 25 is used to store the current vector calculated for one frequency combination in Phase 3. The file contains one logical record, the length of which is determined by the number of nodes in the circuit. Current vector values for each node of the circuit are individually addressable. Subroutines TFWR and TFRD perform input/output for file 25.

The system initialization portion of Phase 0 begins with the definition of NCAP'S machine dependent character manipulation parameters which are transmitted to the subroutines which use

them through the labelled common area WORDSZ. The integer variables NCHAR (number of characters per word), NBCHAR (number of bits per character), and NBWORD (number of bits per word) describe precisely the word length of the host computer and its method of representing alphanumeric characters internally. The alphanumeric left-justified blank, LBLANK, is defined in a DATA statement. The right-justified blank character, RBLANK, is calculated in the main program by a call to subroutine SHIFT.

Next the logical unit numbers of the primary input and output devices, usually the card reader and line printer, are stored in INP and IOUT respectively. The frequency sweep parameters NIT and MAXIT, the plot parameters NCURV, NFR, and NVEC, and the print select and plot switches NLIST and NPLOT are initialized in the main program and updated and tested during the execution of subsequent phases.

The input phase begins with the initialization of Phase 0 pointers, switches, and data areas. The data buffer BUFF(I), I=1, 250 is cleared and the driver file record number DRVREC is set to 2, leaving an open record for the driver file header which is generated internally at the conclusion of Phase 0. DATREC, the next available storage location on data file 20, is initialized to 1, as is the data file address pointer STADD. The input error switch is turned off by JART = 0, while the node card error switch is initially turned on by NABT = 1. The data value and comment card counts NXT and NCOM are cleared and the Start Circuit driver record number SCREC is set to zero.

Since the exact content of the NCAP input deck is not known 3-1L

in advance, input processing is not merely a matter of reading and storing data in a prearranged fashion. Instead the subroutines of Phase Ø act in concert to interpret the input cards, translating them into appropriate driver and data records for use by subsequent NCAP phases.

The input deck consists of a series of *-control card sequences which define circuit topology, circuit excitation, linear and nonlinear elements and models, solution modification, and output. Each such card sequence begins with an *-control card defining its function, followed by parameter cards that supply additional information and/or precise values for that function.

Phase Ø is driven primarily by the *-control cards. Whenever an *-control card is encountered in the input stream, the previous function is closed by writing its driver and data records to disk, and a new function is opened. Parameter (nonasterisk) cards are passed along to a subroutine unique to the function defined by the previous *-card. There the parameter card is processed according to an anticipated format and its data merged with any previous data for that function. Input continues in this fashion until the *END card is processed, at which time program control transfers to PHASE1.

Input processing is controlled by NCAP'S main program and begins by initializing the function identifier MODE to 1, indicating that no input card sequences have yet been processed. A card is read in left-justified alphanumeric format, one character per word, into the integer array CARD. If an

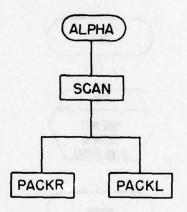
unanticipated end-of-file is encountered, execution terminates with a CALL EXIT. Otherwise the card image is transmitted to subroutine DECIF where it is examined and assigned a functional integer code which is returned to the main program in NEW. Upon return from DECIF, NEW<0 indicates an unidentified card which is not processed further, and after printing an appropriate error message program control transfers to read the next card. Otherwise the card is either a comment (**), an *-control card, or a parameter (non-asterisk) card. Comment cards, identified by NEW=2, are printed on the line printer but otherwise ignored by Phase 0.

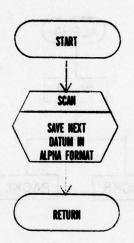
When an *-control card is encountered in the input deck, subroutine DECIF returns a value of NEW>2 to the main program. This causes the old function identified by MODE to be closed and the new function identified by NEW to be opened. Control passes from the main program, through subroutine PRCSS, to an appropriate input processing subroutine which closes the previous *-function by writing its driver and data records to disk, and printing any error messages associated with it. Upon return to the main program the new *-card is printed, the mode of the previous function is saved in MDESV, and the modify switch MODFY is turned off. The new function is opened by setting MODE=NEW and the data value counter, driver record parameters, and data buffer are reinitialized.

Parameter card images (NEW=0) are output on the line printer, and then transmitted through subroutine PRCSS to an appropriate input processing subroutine according to the value of

MODE. There the card is reformatted and interpreted according to the value of TYPE. Upon return to the main program, control transfers to read the next card.

Input continues until the *END (MODE=16) card is encountered. At this point the input processor ENDIN causes program control to transfer to PHASEl if no errors were detected in the input. Otherwise NCAP execution terminates.





NAME · ALPHA

TYPE: SUBROUTINE

GENERAL PURPOSE.

Decodes and formats an alphabetic input value by converting the characters stored in CARD(IPOS) through CARD(J-1) to an alphanumeric value returned in STRING(1)

VARIABLES:

CARD = Right-justified alphanumeric card image

IPOS = Position of first character to be formatted

ISW = End of data switch; Off = 0, On = 1

J = Position of last character to be formatted+1

STRING = Alphanumeric string representation of input

value

SUBROUTINES CALLED.

SCAN

CALLING PROGRAMS:

GENIN JEETIN LCIN LDSIN MODGIN

NCIN NDSIN NODEIN PLOTIN PSIN

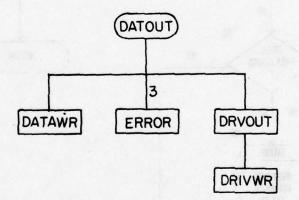
SDIN

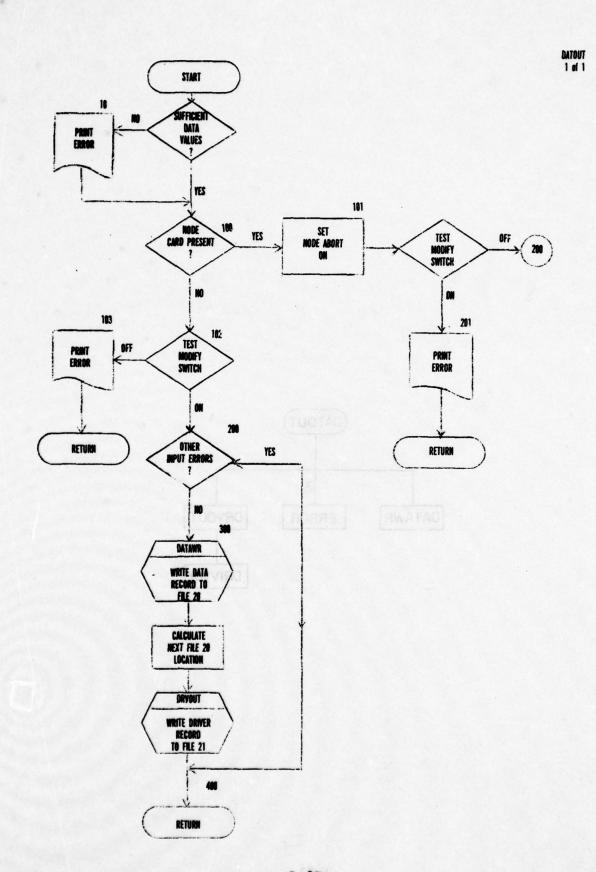
DESCRIPTION .

Subroutine ALPHA decodes and formats a free-form alphabetic input value. Input to the routine is the right-justified alphanumeric card image stored in the integer array CARD and the character position IPOS at which the datum begins.

The contiguous characters CARD(IPOS) through CARD(J-1) which comprise the alphabetic input value are extracted from the card image and converted to a left-justified alphanumeric string by subroutine SCAN. The resulting alphabetic datum stored in the integer STRING(1) is returned to the calling program without further processing.

Additional arguments returned by subroutine ALPHA are the updated character position IPOS, which indicates the first character position of the next value on the card, and the end-of-data switch ISW.





NAME: DATOUT

TYPE · SUBROUTINE

GENERAL PURPOSE .

Writes input data to file 20 and driver records to file 21. Output is by-passed if input errors exist. Clears I/O buffer.

VARIABLES:

J = Number of parameters in the data record

N = Length of data record to be output (in words)

SUBPOUTINES CALLED.

DATAWR DRVOUT ERROR

CALLING PROGRAMS:

PTIN SCIN SCIN TRNIN

VDIN VTIN JEETIN

DESCRIPTION .

Subroutine DATOUT handles the output of data and driver records for many of the Phase Ø input processing routines, closing the function defined by an *-card sequence. Arguments transmitted to the subroutine are J, the number of parameters

anticipated and N, the length of the data record in words.

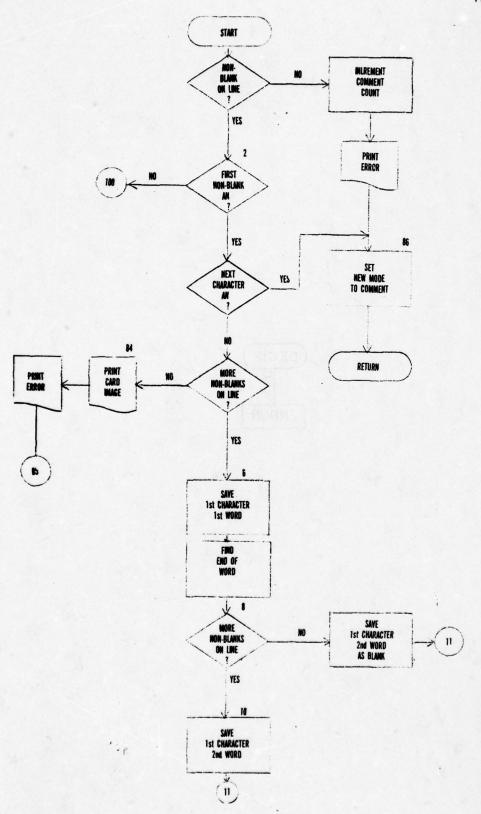
First the input data is tested for errors. Subroutine DATOUT exits with an appropriate error message if any of the following input errors are detected:

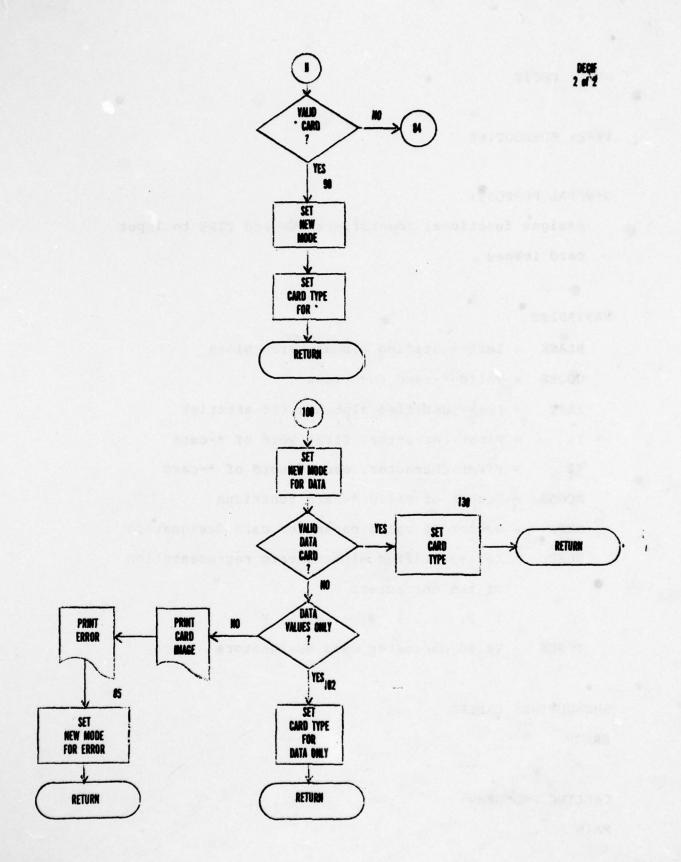
- The parameter count NXT is less than the anticipated number of data values
- 2) A node card was encountered in a modify card sequence (NABT=0 and MODFY=1)
- 3) The node card was missing from a circuit element definition (NABT=1 and MODFY=0)
- 4) The input error switch is on (JABT=1)

If no input errors are detected, the data record is written from the common data buffer BUFF to file 20 by subroutine DATAWR. The data file record number DATREC is updated to point beyond the new data record by DATREC = DATREC + N.

Subroutine DRVOUT is called to complete the driver record definition and to write it to file 21, closing the function. Subroutine DATOUT then returns to the calling program.







NAME: DECIF

TYPE: SUBROUTINE

GENERAL PURPOSE:

Assigns functional identifiers NEW and TYPE to input card images

VARIABLES:

BLANK = Left-justified alphanumeric blank

CODES = Valid *-card functions

IAST = Left-justified alphanumeric asterisk

Il = First character, first word of *-card

12 = First character, second word of *-card

NCODS = Number of valid *-card functions

NTPS = Number of valid parameter card designators

NUMS = Left-justified alphanumeric representation

of the characters

1, 2, ... 9, 0, +, -, ., E

TYPES = Valid parameter card designators

SUBROUTINES CALLED:

ERROR

CALLING PROGRAMS:

MAIN

DESCRIPTION:

Subroutine DECIF examines the card image transmitted to it in the integer array CARD and assigns NEW and TYPE codes according to its contents. The card image is stored in left-justified alphanumeric format, one character per word in the integer array CARD.

Although the system cannot anticipate the exact make-up of an input deck, it does recognize five basic card types:

- An *-control card containing an * followed by one or more keywords. Such cards are classified by TYPE=1 and NEW=3,4,...26, depending on the keywords.
- 2) A parameter card containing an identifying keyword. These cards are classified by NEW=0 and TYPE= 3,4,...23 depending on the keyword.
- 3) A parameter card containing only numeric values and classified by NEW=0 and TYPE=2.
- 4) A comment card containing ** in the first two nonblank character positions. Comments are classified by NEW=2.
- 5) An unidentifiable card constituting an error and classified by NEW=-1.

The function of subroutine DECIF is to determine which of these card types applies to the given card image and to assign corresponding NEW and TYPE codes which serve as the basis for further processing by Phase 0. A complete tabulation of NEW and TYPE codes is presented below:

NEW	CODES	DESCRIPTION
(MODE)	(Character Pair)	neilos DECIF examina
-1	None	Error Condition
Ø	None	Parameter Card
1	None	Driver File Header
		(Internally Created)
2	None	Comment
3	ewellor G_na painis	Generator
40 300	LC	Linear Components
5	NC	Nonlinear Components
6	vD as primi	Vacuum Diode
7	vr vi 621	Vacuum Triode
8	VP	Vacuum Pentode
9 9 1 9	olasaca _T ico patai	Transistor
10	SD	Semiconductor Diode
- 011 OW3	XX XX	Not Used
12	916 84 XX	Frequency Sweep
		(Internally Created)
13	sc	Start Circuit
14	EC	End Circuit
o doidw15	xx	Not Used
16		
17	PS dw act	Print Select
18	A complete Manager	Modify
19	xx	Not Used

20	THE TOT UNITED	Junction Field Effect
	8-Paremeters for J	Transistor
. 21	C-Persaxxters for J	Generator Modification
	Analytic Parameter	(Internally Created)
22	Tobel qt Plot	Plot
23	XX	Not Used
24	2 - 1 - XX	Linear Component Sweep
	Voltaged Controlled	(Internally Created)
25	LD	Linear Dependent Source
26	ND -	Nonlinear Dependent Source
	gi erryale	
TYPE	TYPES	DESCRIPTION
	(Character Pair)	
1	None	*-Control Card
2	None	Parameter Card With Only
		Numeric Values
3	IM	Impedance
, 400 4 99 x P	at the godFR can layers	Frequency
10 1 5 100 0	OR OR	Orders for Print Select or
	A communication bins retu	Plot di reagge recike etale
of te 6 made	NO NO	Nodes for Print Select or
	roer in which dake at	Plot
71036	at seals and the sees	Amplitude
clob 8 evolv	that and Padalog d	Pesistor
9	6 C_ 1 mis 12	Capacitor
10	I.	Inductor

11	out 13 to 19 CO many small	Constants for JFET
12	BP TELEBRAT	B-Parameters for JFET
13 ms	decidifie CP december	C-Parameters for JFET
14	(betaes) (AC onse fall)	Analytic Parameters for JFET
15	LA	Label for Plot
16	FB seed sow	Forward Bias
17	ed Jagaog RB	Reverse Bias
18	ve de la company	Voltage Controlled Current
		Source
19	Jaubauger CC unity	Current Controlled Current
		Source
20	vv = 1 = 1	Voltage Controlled Voltage
		Source
21	cv	Current Controlled Voltage
		Source
22	ta pred 1 Objected 3	Print Select Off
23	Paule 7 ON Terrus	Print Select On

Pecause the NCAP user expresses his input in free-form, there is no requirement for card input to begin in column 1 or that data values appear in particular card columns. A data value is defined as a string of contiguous non-blank alphanumeric characters, and although the order in which data are placed on cards is critical, the spacing between data values is arbitrary. Therefore subroutine DECIF must isolate the individual data values on the card and then classify the card according to those values.

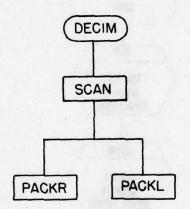
Classification begins by determining the presence or absence of an * in the first non-blank character position. Starting at CARD(1) the image is scanned column-by-column to locate the first non-blank character. If no non-blanks are found, the card contains no data and subroutine DECIF exits to the main program treating the blank card as a comment. Otherwise the first non-blank located in CARD(I) is tested for the presence of an *. If an * is found, the card is either an *-control card or a comment card. If CARD(I+1) is also an asterisk, a comment is assumed and classified by NEW=2. After incrementing the comment count NCOM, subroutine DECIF exits to the main program.

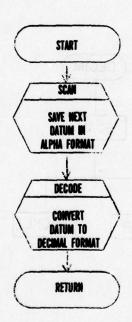
If CARD(I+1) is not an asterisk, an *-control card is assumed. The first character of each of the next two words is extracted and compared against the table of valid character pairs stored in CODES(J), J=1, NCODS. If a match is found, the card is classified by TYPE=1 and NEW=J+2. If no match is found, the unidentifiable card is classified by NEW=-1 and exit occurs after an appropriate error message is printed.

In the event that the first non-blank character CARD(I) is not an asterisk, a parameter card is assumed and classified by NEW=0. Such cards are further examined by comparing the first two characters of the first word on the card against a table of valid character pairs stored in TYPES(J), J=1, NTPS. If a match is found, the card is classified by TYPE=J+2 and exit occurs. Otherwise the first non-blank character is compared against the table of valid numeric characters stored in NUMS(J), J=1, 14. No

match indicates an unidentifiable card which is assigned NEW=-1. If the character is a valid numeric, the card is assumed to contain only numeric data values and is further classified by TYPE=2. In either event return is made to the main program where further processing of the card image takes place.

New B. . Even cords are further examined by consising the first





NAME: DECIM

TYPE: SUBROUTINE

GENERAL PURPOSE: A TEMPORAL PURPOSE: A TEMPORAL PURPOSE STATE OF THE PUR

Decodes and formats a decimal input value by converting the characters stored in CARD(IPOS) through CARD(J-1) to a floating point number returned in X

VARIABLES:

CARD = Right-justified alphanumeric card image

IPOS = Position of first character to be formatted

ISW = End of data switch: Off = \emptyset , On = 1

J = Position of last character to be formatted+l

STRING = Alphanumeric string representation of input

value value

x = Floating point value returned add as simble bear add

SUPPOUTINES CALLED.

SCAN

CALLING PROCPAMS .

CENIN JEETIN LCIN LDSIN MODCIN

MCIN MOSIN PTIN SDIN TRNIN

VIIN VIIN

DESCRIPTION .

Subroutine DECIM decodes and formats a free-form decimal input value. Input to the routine is the right-justified alphanumeric card image stored in the integer array CARD and the character position IPOS at which the datum begins.

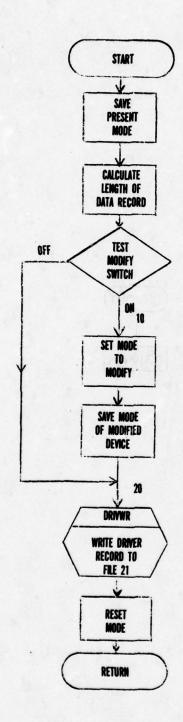
The contiguous characters CARD(IPOS) through CARD(J-1) which comprise the decimal input value are extracted from the card image and converted to a right-justified alphanumeric string by subroutine SCAN. The character string returned in the integer array STRING is then converted from alphanumeric representation to a floating point value X by:

DECODE (STRING, 1) X

1 FORMAT (E12.0)

Arguments returned by subroutine DECIM are the decimal input value X, the updated character position IPOS which indicates the first character position of the next data value on the card, and the end-of-data switch ISW.





NAME · DRVOUT

TYPE SUPROUTINE

GENERAL PURPOSE: WARREN WARREN OF BELLEVIOLE BARREN WOLLDEN

Completes definition of driver parameters and writes
driver record to file 21

is a modify function (maDFY=1), MODE is set to is and the mode of

VARIABLES.

JJ = Intermediate storage of MODE

SUPROUTINES CALLED.

DRIVWR

CALLING PROGRAMS:

DATCUT ECIN ENDIN LCIN LDSIN

NCIN PSIN SCIN

DESCRIPTION .

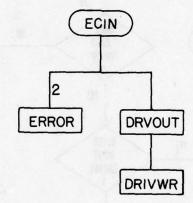
Subroutine DPVCUT completes the definition of driver record parameters for some of the Phase 0 input processing subroutines and writes the driver record to file 21.

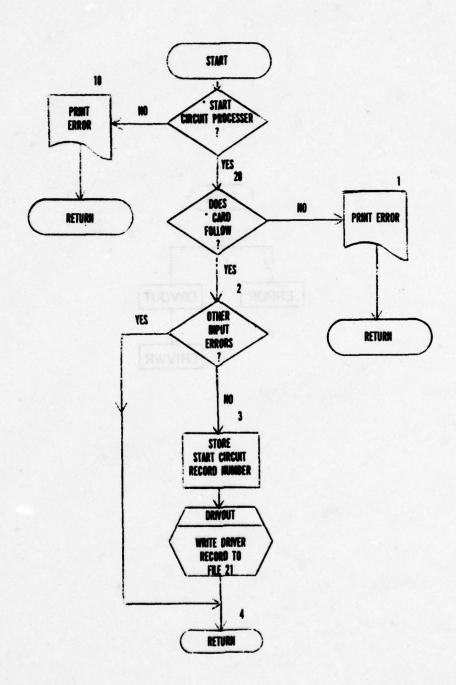
The MCDE of the present card sequence is saved in JJ and the data record length is calculated by LNGTH=DATREC-STADD. If the modify switch is off (MCDFY=0), the driver record is written to file 21 from the first ten words of global common by subroutine DQIVWR.

A MODIFY card sequence carries the MODE of the device it modifies during input processing, but is written to the driver file as MODE=18. Therefore, if the card sequence being processed is a modify function (MODFY=1), MODE is set to 18 and the mode of the device being modified is taken from MDESV and saved in JJ before the driver record is written by subroutine DRIVWR.

Upon return from subroutine DRIVWR the record number for the next data record STADD is taken from DATREC, the mode of the present card sequence is reset from JJ, and subroutine DRVOUT returns to the calling program.

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NAME: ECIN al radmont brosser revisit divority brade and il

TYPE: SUBROUTINE . nols/back sosts estal squage de . delw estat . Miss

GENERAL PURPOSE: THE WAS MAINTANAMED AND A SECOND OF THE SECOND SECO

Processes *END CIRCUIT card

VARIABLES: Sept 1408 200 000 PRIME No. 4800 Day 300 260 360 360 360 360

Start clicula ariver record number is stored in MIRCILL SHOW

ERROR

CALLING PROGRAMS:

PRCSS

DESCRIPTION:

Subroutine ECIN is the input processor for the END CIRCUIT card. Its function is to translate the end circuit definition into an appropriate driver record for use by subsequent NCAP phases.

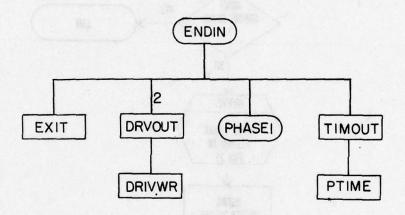
Because there are no data cards associated with the end circuit input, TYPE=1 is the only code recognized by subroutine ECIN. Any other value of TYPE constitutes an input error which is handled by subroutine ERROR before exit to the calling program.

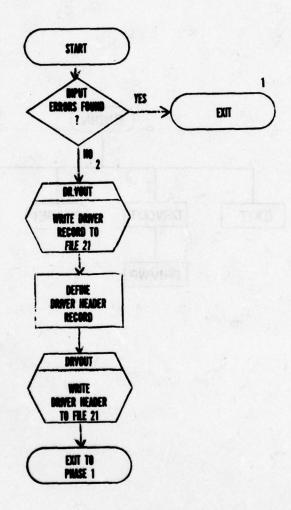
If the start circuit driver record number is undefined (SCREC=0) or the input error switch is on (JABT=1), subroutine ECIN exits with an appropriate error condition. Otherwise the END CIRCUIT driver record is written to file 21 from the first ten words of global common by subroutine DRVOUT.

The end circuit driver record is identified by MODE=14. Because there is no data associate with an end circuit function, the driver parameters STADD and LNGTH are not applicable. The start circuit driver record number is stored in MISC(1). The other MISC driver parameters are not used.

After the driver record has been written to disk, subroutine ECIN returns to the calling program.

ov subroading pance refore exist to the calling





NAME: ENDIN

sobroutine DRVOUT. The END driver record is identified by

TYPE: SUBROUTINE

GENERAL PURPOSE:

Processes *END card, creates driver file header, closes

Phase 0 processing

VARIABLES:

NONE of marriage at 510001 repend off . 6480 for our successful revision

SUBROUTINES CALLED:

DRVOUT PHASE1 TIMOUT

CALLING PROGRAMS:

PRCSS

DESCRIPTION:

Subroutine ENDIN is the input processor for the END card. Its function is to translate the end definition into an appropriate driver record for use by subsequent NCAP phases and to close the Phase Ø processing.

If errors were encountered in the input deck during Phase Ø processing (JABT = 1), NCAP execution terminates with a CALL EXIT.

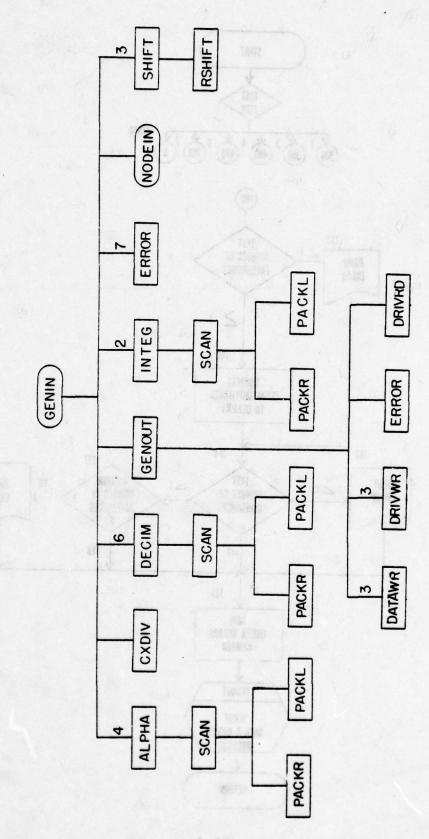
Otherwise the END driver record is written to file 21 by

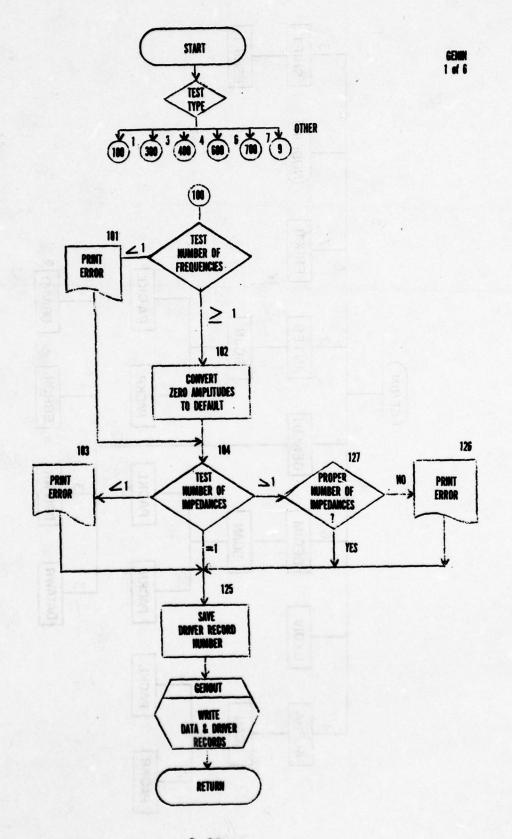
subroutine DRVOUT. The END driver record is identified by MODE=16. Because there is no data associated with the end function (it simply serves as a delimiter), the driver parameters STADD and LNGTH are not applicable. The MISC driver parameters are not used.

The Phase Ø processing is closed by creating the driver file header and writing it to the first record of file 21. The header record is identified by MODE=1 and contains the next available file 20 storage location in MISC(1) and MISC(4). The other driver parameters are not used. The header record is written to the driver file by subroutine DRVOUT and program control transfers to subroutine PHASE1.

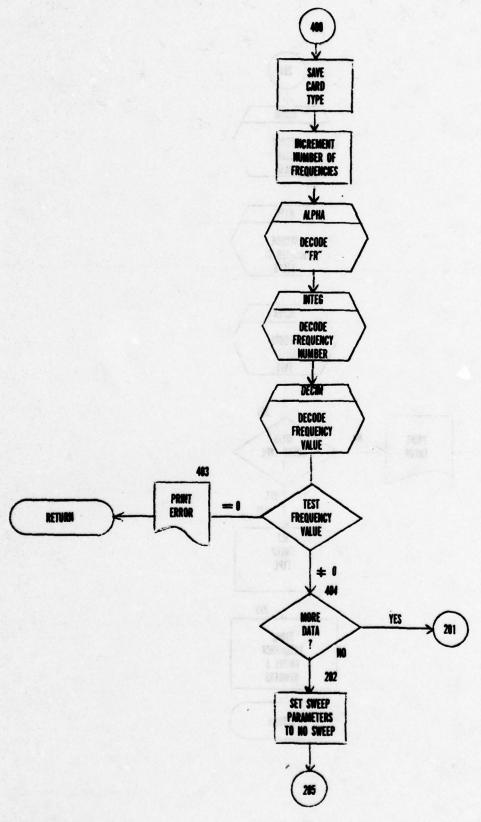
Subroating EMDIN is the isput processor for the END card.

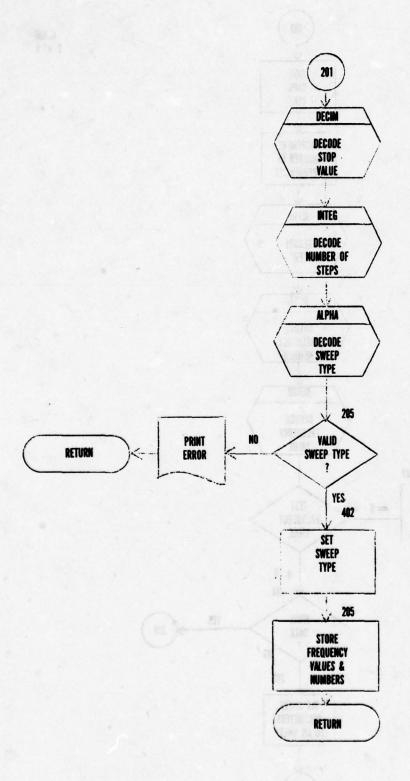
appropriate driver record for use by subsequent ACAP phases and

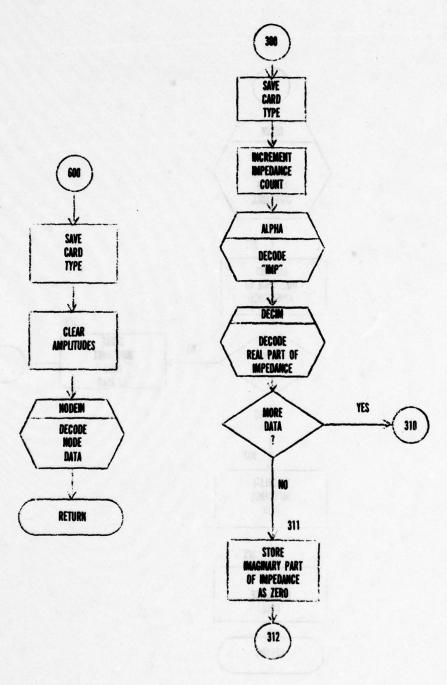


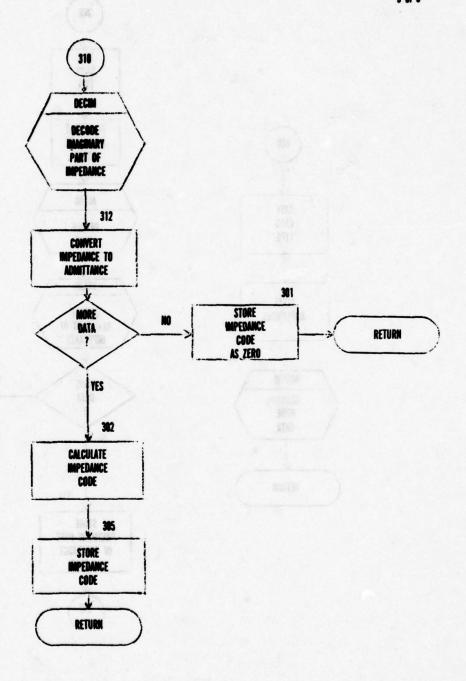


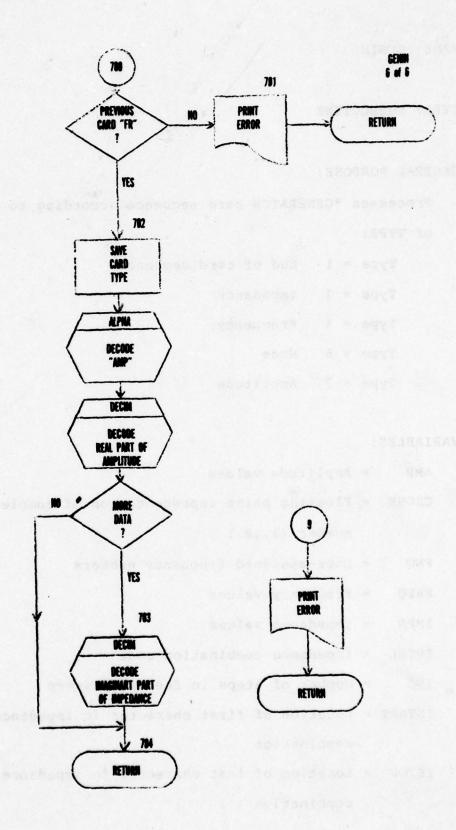












NAME: GENIN

TYPE: SUBROUTINE

GENERAL PURPOSE:

Processes *GENERATOR card sequence according to the value of TYPE:

Type = 1 End of card sequence

Type = 3 Impedance

Type = 4 Frequency

Type = 6 Node

Type = 7 Amplitude

VARIABLES:

AMP = Amplitude values

CXONE = Floating point representation of complex number (1.,0.)

FNO = User-assigned frequency numbers

FREQ = Frequency values

IMPS = Impedance values

IMTBL = Impedance combination code

INC = Number of steps in frequency sweep

ISTART = Location of first character in impedance combination

ISW = End of data switch: Off = 0, On = 1

ITSV = Type code of last parameter card processed

ITYPE = Sweep type code: 0 = No sweep

l = Linear

2 = Logarithmic

JRECO = Address of driver record

NFREQ = Number of frequencies this generator

NIMP = Number of impedances this generator

RBLANK = Right-justified alphanumeric blank

STRING = Alphanumeric string representation of input

value value

STOP = Stop value for sweep

TYPES = Valid frequency sweep types

SUBROUTINES CALLED:

ALPHA CXDIV DECIM ERROR

GENOUT INTEG NODEIN

CALLING PROGRAMS:

PRCSS

DESCRIPTION:

Subroutine GENIN is the input processor for GENERATOR card sequences. Its function is to decode the generator input data according to an anticipated format, translating the input into appropriate data and driver records for use by subsequent NCAP phases.

Each call to subroutine GENIN causes a single NCAP input card to be processed. The card image and its type code are transmitted through common storage to subroutine GENIN which processes the card according to its type as follows:

brooms toward to amount a court

TYPE=1 End of Card Sequence

The closing of a generator definition begins by testing the frequency and impedance specifications for errors. Each generator definition must contain at least one frequency card (NFREQ.GT.0) and either a constant impedance (NIMP=1) or one impedance card for each possible frequency combination (NIMP=2 of an improper number of frequencies and/or impedances have been specified, subroutine GENIN exits with an appropriate error condition.

Any undefined amplitude values are assigned the default value:

$$AMP(1,I) = 1.0$$

AMP
$$(2,I) = 0.0$$

The amplitude data is then stored in the common data buffer at BUFF(11) - BUFF(30). Subroutine GENOUT is called to complete the definition of the generator driver and data records and to write them to disk, after which subroutine GENIN returns to the calling program.

TYPE=3 Impedance Card

The card type card is saved in ITSV, the number of impedances NIMP is incremented, and the IMP identifier is decoded by subroutine ALPHA. The real part of the impedance value is decoded and formatted by subroutine DECIM and stored in IMPS(1, NIMP). If the end-of-data is encountered (ISW=1), the imaginary part of the impedance value is assigned the default value zero. Otherwise the imaginary part of the impedance value is decoded and formatted by subroutine DECIM and stored in IMPS(2, NIMP). The impedance value is then converted to an admittance by subroutine CXDIV.

At this point the end-of-data switch ISW is used to determine the presence of an impedance combination on the card. If the end-of-data switch is on (ISW=1), subroutine GENIN exits after setting the impedance code BUFF(NIMP + 40) = 0 to indicate a constant impedance.

If the end-of-data is off (ISW=0), the impedance combination is extracted from the card image, converted to an impedance code, and appended to the impedance code table at BUFF(NIMP + 40). Each digit of the impedance combination is represented in the impedance code IMTBL by the bit position of like number. For example, if the digit 3 is included in the combination, the 3rd bit of IMTBL is turned on by IMTBL = 2^2 . Beginning with the first digit of the combination stored at CARD (IPOS), each digit is mapped to the proper bit position and accumulated in IMTBL as follows:

IMTPL = IMTPL + (2** (CARD(J-1) -1))

This process continues under the control of the index J until either CARD(J) is a blank or six digits have been used. The resulting code is placed in the impedance code table at BUFF(NIMP + 40) and subroutine GENIN returns to the calling program.

TYPE=4 Frequency Card

The card type code is saved in ITSV, the number of frequencies NFREC is incremented, and the FR identifier is decoded by subroutine ALPHA. The frequency number is decoded and formatted by subroutine INTEG and stored in FNO(NFREQ). The frequency value is decoded and formatted by subroutine DECIM and stored in FREQ(NFREC). If the frequency number is zero, subroutine GENIN exits with an appropriate error condition.

At this point the end-of-data switch ISW is used to determine the presence of frequency sweep parameters on the card. If the end-of-data switch is on (ISW=1), the parameters ITYPE and INC are set to indicate no sweep. Otherwise the STOP, INC, and ITYPE parameters are decoded and formatted by subroutine DECIM, INTEG, and ALPHA respectively.

The sweep type is tested against a table of valid alphanumeric types. If no match is found, subroutine GENIN exits with an appropriate error condition. If a match is found, the sweep type is encoded in ITYPE(NFREC), the frequency number and value are placed in the common data buffer BUFF, and subroutine

GENIN returns to the calling program.

TYPE=6 Node Card

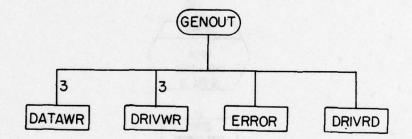
The card type code is saved in ITSV and the amplitude data area AMP is cleared. The nodes of connection are extracted from the card image, formatted, and stored in MISC(3) and MISC(4) of the driver record by subroutine NODEIN. Subroutine GENIN returns to the calling program.

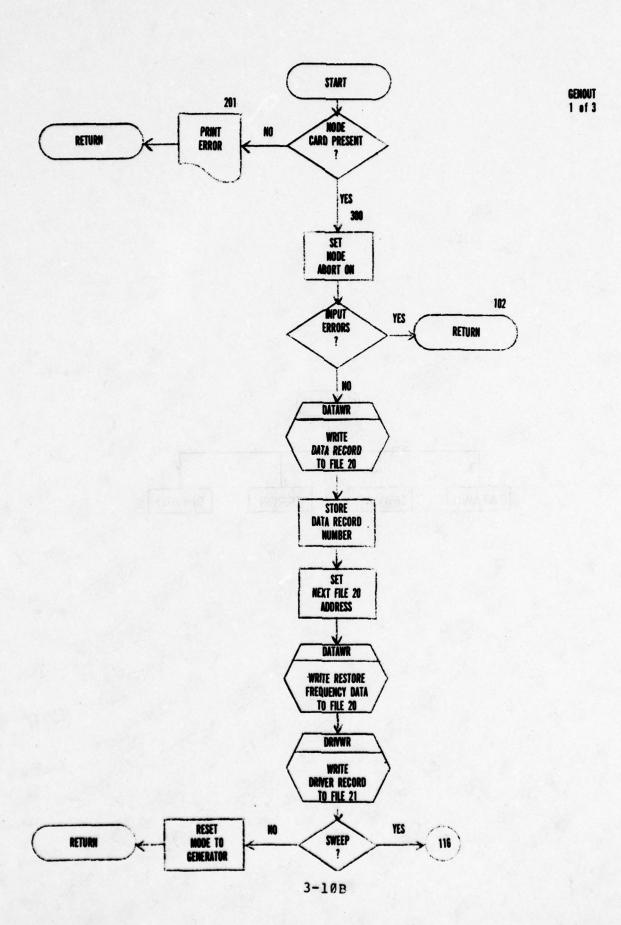
TYPE=7 Amplitude Card

To assure that amplitude data are properly associated with frequency data, each amplitude card must physically follow a frequency card in the input deck. Therefore if the last card processed was not a frequency card (ITSV \neq 4), subroutine GENIN exits with an appropriate error condition.

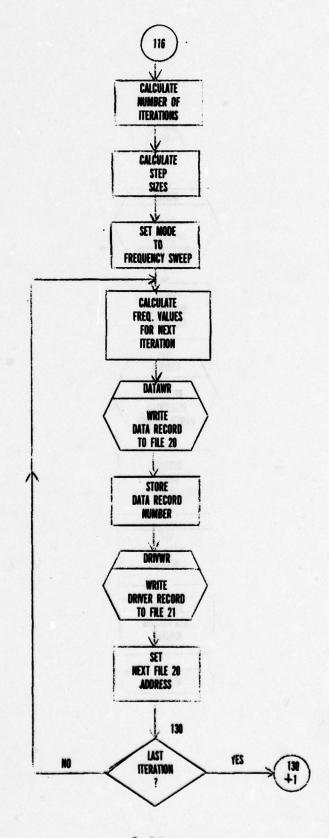
Otherwise the card type code is saved in ITSV and the AMP identifier is decoded by subroutine ALPHA. The real and imaginary parts of the complex amplitude are decoded and formatted by subroutine DECIM and stored in AMP(1, NFREQ) and AMP(2, NFREQ).

Any other value of TYPE constitutes an input error which is handled by subroutine ERROR before exit to the calling program.











NAME: GENOUT

TYPE: SUBROUTINE

GENERAL PURPOSE:

Tests for errors in generator and modify generator input, writes input data to file 20 and driver record to file 21, creates and writes frequency sweep data and driver records

VARIABLES:

FNO = User-assigned frequency numbers

FREQ = Frequency values

INC = Number of steps in frequency sweep

ITYPE = Sweep type code; 0 = No sweep

l = Linear

2 = Logarithmic

JRECO = Address of generator or modify driver record

NFREQ = Number of frequencies this generator

NIT = Maximum number of sweep iterations this generator

STRING = Alphanumeric string representation of input value

STEP = Sweep increment value

STOP = Stop value for sweep

SUBROUTINES CALLED:

DATAWR DRIVRD DRIVWR ERROR

CALLING PROGRAMS:

GENIN MODGIN

DESCRIPTION .

Subroutine GENOUT handles the output of data and driver records for the generator input processing subroutines GENIN and MODGIN, closing the generator and modify generator definitions.

First the input processed by GENIN or MODGIN is tested for errors. If the node card is missing (NABT #0) or the input error switch is on (JABT=1), return is made to the calling program with an appropriate error condition. Otherwise the data record is written to file 20 from the common data buffer BUFF by subroutine DATAWR.

The generator data record contains 230 words allocated as follows.

BUFF(1)-(10) Frequency Values

BUFF(11)-(30) Complex Amplitudes

BUFF(31)-(40) User-assigned Frequency Numbers

BUFF(41)-(103) Impedance Combination Codes

BUFF(104)-(230) Complex Impedance Values

After the data record has been written to disk, its record number is stored at STADD and the next available file 20 storage location is updated to point beyond the generator data by DATREC=DATREC+230. The frequency values, amplitudes, and

frequency numbers stored in BUFF(1)-(40) are then rewritten to the data file at record number DATREC and the next available file 20 storage location is updated to point beyond the re-written data by DATREC=DATREC+40. (This re-write file is later used by Phase 4 to restore the original frequency data after execution of a frequency sweep.)

The driver record parameters defined in the input processor, GENIN or MODGIN, are transmitted through common storage to subroutine GENOUT. Here the remaining driver parameters are defined and the record is written by subroutine DRIVWR to file 21 from the first ten words of global common. The complete driver record is comprised of:

MODE=3 or 21	Generator or Modify	Generator

Topics of the tricing Identifier

STADD	Data File	Record Number	
DINDU	Data Lite	Mecoro Mumber	

LNGTH=230 Length of Data Record in Words

MISC(1)=NFREQ Number of Frequencies

MISC(2)=NIMP Number of Impedances

MISC(3) Positive Node of Connection

MISC(4) Negative Node of Connection

MISC(6)=NIT Number of Sweep Iterations

Other MISC driver parameters are not used.

Unlike most NCAP input card sequences which map to a single driver and data record, generator and modify generator card sequences result in one or more driver records, each with an associated data record. The number of driver records is related to the number of sweep iterations defined in the card sequence.

If frequency sweeping is not specified in the card sequence (i.e., ITYPF(I)=0 for all I=1, NFREQ), the card sequence has only a single MODE=3 or 21 driver record associated with it and subroutine GENOUT returns to the calling program without further processing.

Otherwise the sweep specifications must be translated into appropriate driver and data records. First the number of sweep iterations, NIT, is derived from the largest INC parameter specified in the generator or modify generator card sequence. (Each of the NIT sweep iterations causes one driver and data record to be created internally and appended to the disk files.)

Then a sweep increment STEP(I) is calculated for each frequency value in the generator definition. The increment is a function of the original frequency value FREQ(I) and the STOP, INC, and ITYPE parameters associated with that frequency as follows:

1) For ITYPE(I)=1, linear sweep is indicated and:

STEP(I) =
$$\frac{\text{STOP}(I) - \text{FREQ}(I)}{\text{INC}(I) - 1}$$
.

2) For ITYPE(I)=2, logarithmic sweep is indicated and.

STEP(I) =
$$\frac{\text{STOP}(I)}{\text{FREQ}(I)}$$
 INC(I)-1.

After the increments have been calculated, the sweep driver and data records are created and appended to the disk files. Like standard driver records, sweep driver records are developed in the first ten words of global common. A frequency sweep driver record consists of:

MODE=12 Frequency Sweep Identifier

STADD Data File Record Number

LNGTH=40 Length of Data Record in Words

MISC(1)=NFREQ Number of Frequencies

MISC(2)=NIMP Number of Impedances

MISC(3) Positive Node of Connection

MISC(4) Negative Node of Connection

MISC(6)=NIT Number of Sweep Iterations

Other MISC driver parameters are not used.

Frequency sweep data records have the same structure as the first 40 words of the original generator data record and are developed in the common data buffer BUFF by an iterative process. The first sweep data record is derived from the original generator data by performing an incrementation operation on each frequency value BUFF(J), J=1,NFREQ according to its corresponding STEP(J), STOP(J), and ITYPE(J) parameters:

 For ITYPE(J)=0, no sweep is indicated. No incrementation is performed and the frequency value remains constant at BUFF(J).

2) For ITYPE(J)=1, linear sweep is indicated and the incrementation is additive:

BUFF(J) =BUFF(J) + STEP(J)

unless STOP(J) is exceeded, in which case the frequency remains constant at its last value.

3) For ITYPE(J)=2, logarithmic sweep is indicated and the incrementation is multiplicative:

BUFF(J) = BUFF(J) * STEP(J)

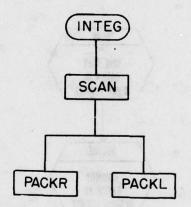
If STOP(J) is exceeded, the frequency value remains constant at STOP(J).

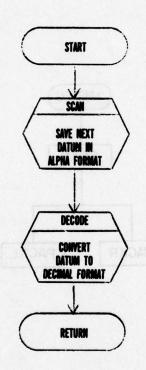
After every frequency value in the buffer has been updated, the new data record is written to file 20 at record number DATREC by subroutine DATAWR. The next available file 20 storage location is updated to DATREC=DATREC + 40 and subroutine DRIVWR is called to write the sweep driver record to file 21. The data record number STADD is updated and program control transfers back to process the next sweep iteration. The creation of sweep records continues under the control of the index I until all NIT sweep iterations have been translated into data and driver records.

The original generator or modify generator driver record is read from the driver file by subroutine DRIVRD according to the record number JRECO transmitted to subroutine GENOUT as an

argument. After storing the sweep iteration count NIT in MISC(6), the updated driver record is re-written by subroutine DRIVWR.

After the MODE is restored to 3 and the data file record number STADD is updated, subroutine GENOUT returns to the calling program.





NAME: INTEG

TYPE · SUBROUTINE

GENERAL PURPOSE:

Decodes and formats an integer input value by converting the characters stored in CARD(IPOS) through CARD(J-1) to a fixed point number returned in K

VARIABLES:

CARD = Right-justified alphanumeric card image

IPOS = Position of first character to be formatted

ISW = End of data switch: off = 0, on = 1

J = Position of last character to be formatted+1

K = Integer value returned

STRING = Alphanumeric string representation of input value

SUBROUTINES CALLED:

SCAN

CALLING PROGRAMS:

GENIN LCIN LDSIN MODGIN NCIN

NDSIN NODEIN PLOTIN PSIN

DESCRIPTION:

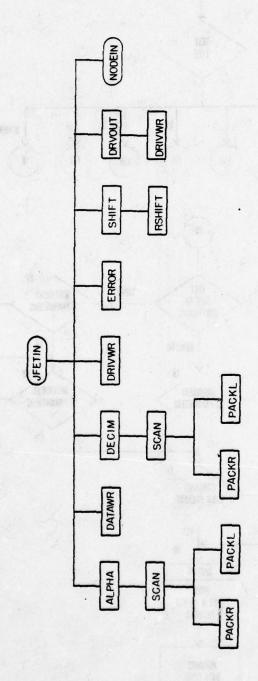
Subroutine INTEG decodes and formats a free-form integer input value. Input to the routine is the right-justified alphanumeric card image stored in the integer array CARD and the character position IPOS at which the datum begins.

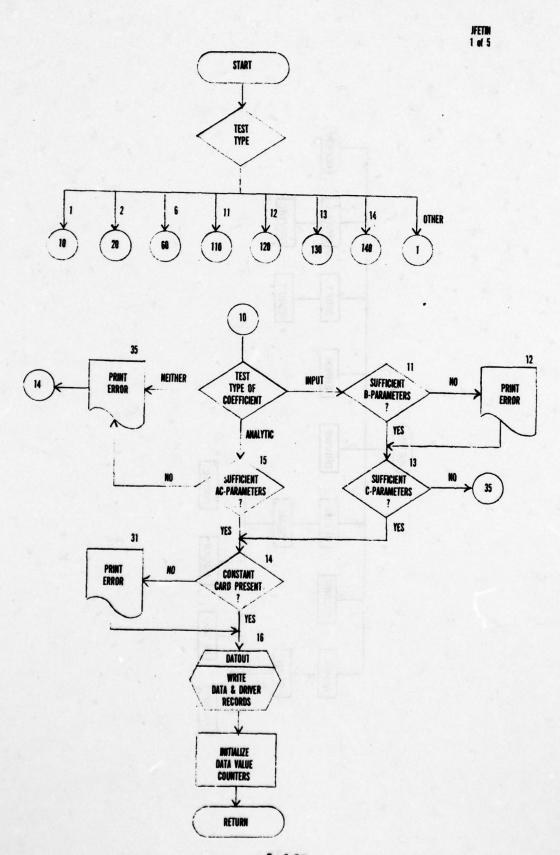
The contiguous characters CARD(IPOS) through CARD(J-1) which comprise the integer input value are extracted from the card image and converted to a right-justified alphanumeric string by subroutine SCAN. The character string returned in the integer array STRING is then converted from alphanumeric representation to an integer value K by:

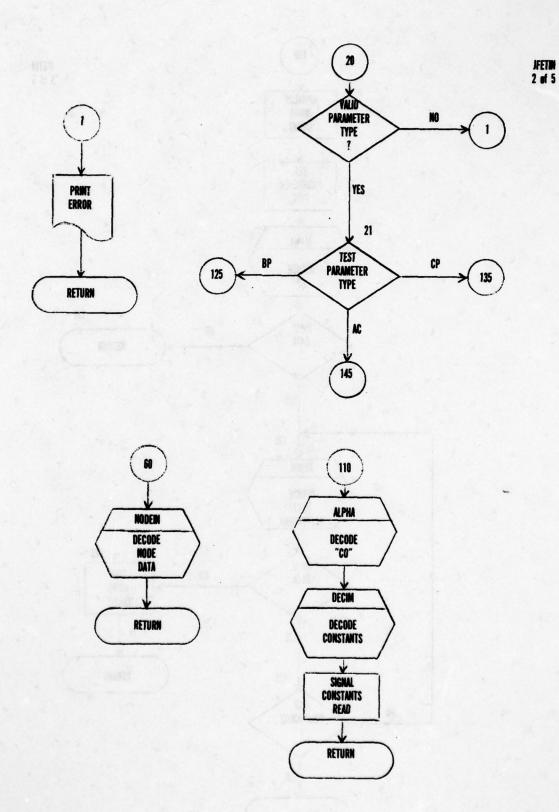
DECODE (STRING, 2) K

2 FORMAT (I12)

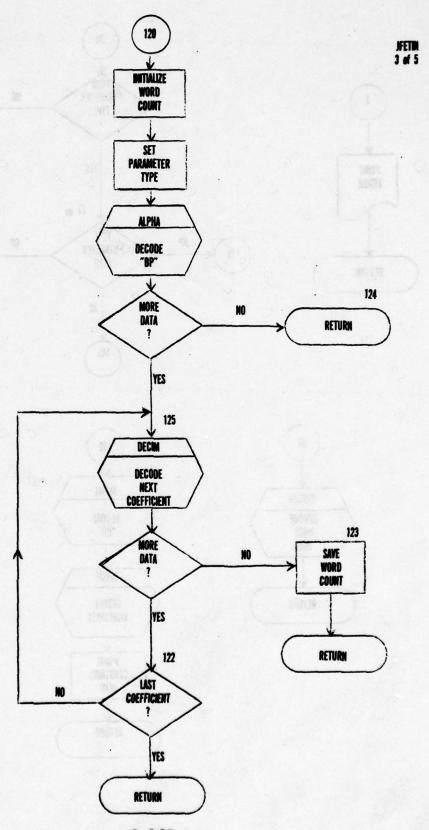
Arguments returned by subroutine INTEG are the integer input value K, the updated character position IPOS which indicates the first character position of the next data value on the card, and the end-of-data switch ISW.



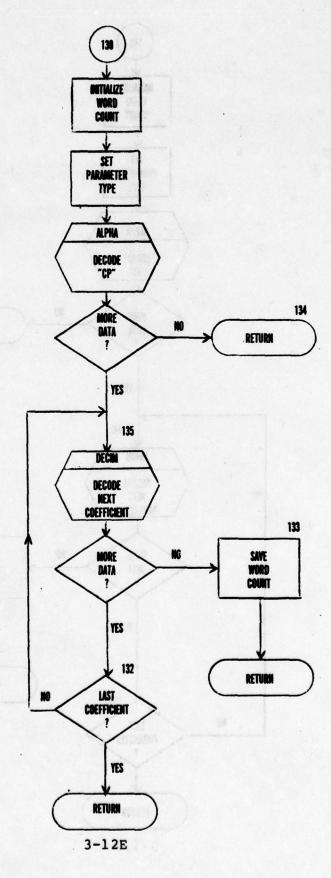


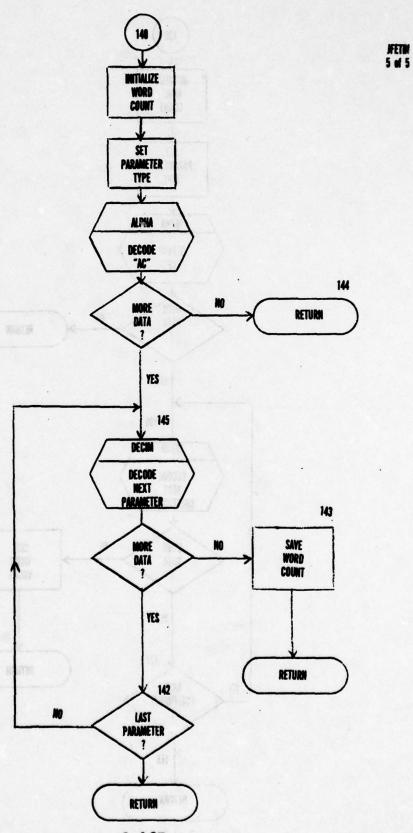


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NONLINEAR CIRCUIT ANALYSIS PROGRAM (NCAP) DOCUMENTATION. VOLUME--ETC(U)
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3-12D





3-12F

NAME: JETIN

TYPE: SUBROUTINE AND INCOME OF THE PROPERTY OF

GENERAL PURPOSE:

Processes *JFET card sequence according to the value of TYPE:

Type = 1 End of card sequence

Type = 2 Parameter card

Type = 6 Node

Type = 11 Constants

Type = 12 BP

Type = 13 CP

Type = 14 AC

VARIABLES: No. of the same was bree (reference) together bleek

B = B parameters

C = C parameters

CGD = Input constant

ICOEF = Parameter type code: 1 = BP

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doldw William enlawardes of an 3 = AC amos appares hardtmens t

ISW = End of data switch: Off = 0, On = 1

NW = Word counter for parameter input

NXTAC = Word counter for AC input

NXTB = Word counter for BP input

NXTC = Word counter for CP input

RS = Input constant

STRING = Alphanumeric string representation of input value

SUBROUTINES CALLED:

ALPHA DATOUT DECIM DRIVWR

ERROR NODEIN

CALLING PROGRAMS:

PRCSS

DESCRIPTION:

Subroutine JFETIN is the input processor for JFET (Junction Field Effect Transistor) card sequences. Its function is to decode the JFET input data according to an anticipated format, translating the input into appropriate data and driver records for use by subsequent NCAP phases.

Type = 1 8nd of card secornes

Type w 2 Parameter ceen

Each call to subroutine JFETIN causes a single NCAP input card to be processed. The card image and its TYPE code are transmitted through common storage to subroutine JFETIN which processes the card according to its TYPE as follows:

TYPE=1 End of Card Sequence

The closing of a JFET definition begins by testing the input $3-12\mathrm{H}$

- for errors. Subroutine JFETIN exits with an appropriate error condition if any of the following input errors are detected.
 - 1) No BP or CP coefficients were processed (NXTE=0)
 or (NXTC=0) when coefficient input was specified
 (MISC(2)=1)
 - 2) Fewer than 8 parameters were processed (NXTAC.LT.8) when analytic generation of coefficients was specified (MISC(2)=2)
- 3) CONSTANTS card missing (MISC(1)=0)
 Otherwise subroutine DATOUT is called to complete the driver record and to close the JFET definition by writing the driver and data records to disk.

The driver record is written to file 21 from the first ten words of global common and consists of:

MODE=20 JFET Identifier

STADD Data File Record Number

LNGTH=64 Length of Data File in Words

MISC(2)=1 or 2 Method of Deriving Coefficients.

l=Input, 2=Analytic generation

MISC(3) Base Node of Connection

Other MISC driver parameters are not used.

The JFET data record, written to file 20 from the common data buffer BUFF is 64 words in length. This record allocates sufficient disk storage for the JFET input values as well as all other data generated internally by its model in subsequent NCAP phases:

PUFF(1)-(10)

B-Parameters Input

BUFF(11)-(20) C-Parameters Input
BUFF(21)-(22) Constants Input
BUFF(23)-(32) Nonlinear Coefficients

BUFF (33) - (64) Admittance Submatrix

After returning from subroutine DATOUT, the parameter counts NXTB, NXTC, and NXTAC and the parameter type code ICOEF are cleared, and return is made to the calling program.

TYPE=2 Parameter Card

If the parameter type code is undefined (ICOEF.LT.0), subroutine JFETIN exits with an appropriate error condition. Otherwise the parameters are decoded and formatted according to the parameter code ICOEF which identifies the type of parameters being processed (1=BP, 2=CP, 3=AC) and causes them to be input accordingly.

TYPE=6 Node Card

The base node of connection is extracted from the card image, formatted, and stored in MISC(3) by subroutine NODEIN.

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TYPE=11 Constants Card

The CONSTANTS identifier is decoded by subroutine ALPHA, and the constant values CGD and RS are decoded and formatted by subroutine DECIM. MISC(1) is set to 1 to signal that the constants have been read and subroutine JFETIN returns to the calling program.

TYPE=12 PP Card

The B-parameter count is initialized (NXTR=1), MISC(2) is set to indicate that coefficients are to be read rather than analytically generated, and the parameter code is set to indicate B-parameters (ICOEF=1). The BP identifier is decoded by subroutine ALPHA and if the end-of-data is encountered (ISW=1), subroutine JFETIN returns to the calling program.

Otherwise from one to ten decimal values are extracted from the card image, formatted, and stored in the common data buffer. B-coefficients appearing on the BP card as well as those input on unlabelled parameter cards following the BP card (TYPE=2 and ICOEF=1) are processed at this point. The number of values decoded and their placement in the buffer depend on the P-parameter count, NXTB, and the end-of-data switch, ISW. NXTB contains the next buffer location available for parameter storage, which begins at P(NXTB) and proceeds through successive buffer locations as data values are decoded by subroutine DECIM.

If an end-of-data is encountered (ISW=1) before P(10) is filled, the parameter count is reset to the next available buffer

location (NXTP=NW) and control returns to the calling program in anticipation of additional parameter cards.

TYPF=13 CP Card

The C-parameter count is initialized (MXTC=1). MISC(2) is set to indicate that coefficients are to be read rather than analytically generated, and the parameter code is set to indicate C-parameters (ICOFF=2). The CP identifier is decoded by subroutine ALPHA and if the end-of-data is encountered (ISW=1), subroutine JFETIN returns to the calling program.

Otherwise the decoding, formatting, and storage technique for C-parameters is identical to that of B-parameters except that NXTC is used for the parameter count and storage is from C(1) through C(10).

TYPE=14 AC Card

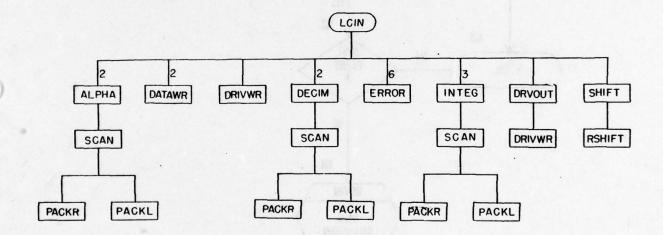
MISC(2) is set to indicate that coefficients are to be analytically generated rather than read, the AC-parameter count is initialized (NXTAC=1), and the parameter code is set to indicate AC-parameters (ICOEF=3). The AC identifier is decoded by subroutine ALPHA and if the end-of-data is encountered (ISW=1), subroutine JFETIN returns to the calling program.

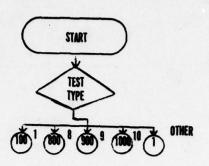
tours to reduce ent tries side is because or little to

Otherwise the decoding, formatting, and storage technique

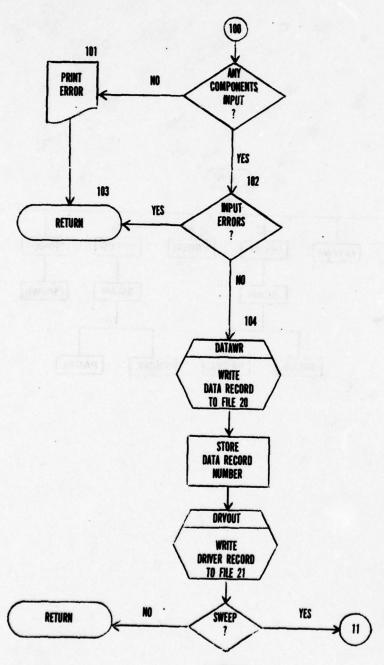
for AC-parameters is identical to that of B-parameters except that NXTAC is used for the parameter count.

Any other value of TYPE constitutes an input error which is handled by subroutine ERROR before exit to the calling program.

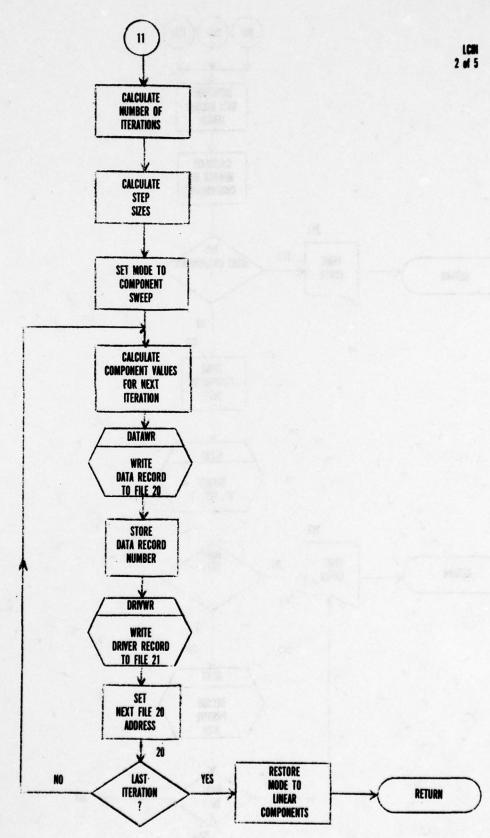


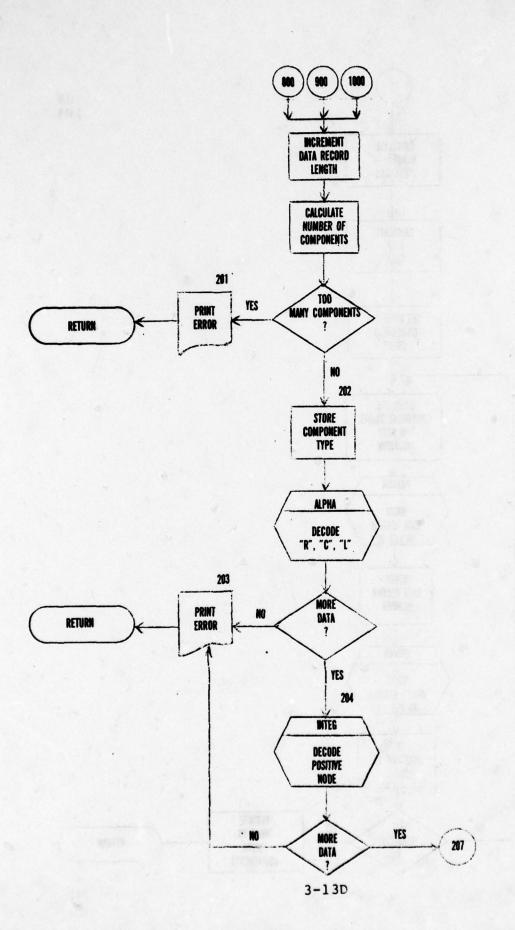




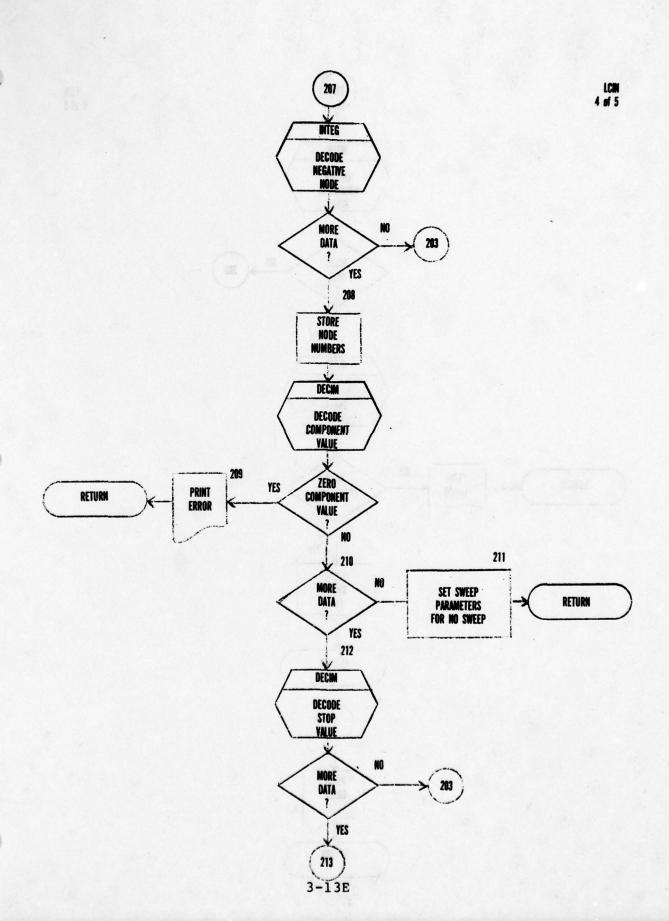


3-13B

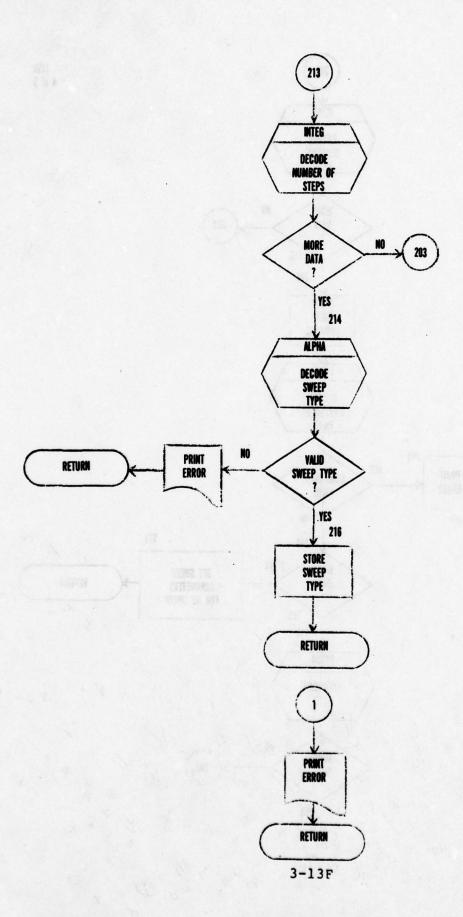




LCM 3 of 5







NAME: LCIN

TYPE: SUBROUTINE

GENERAL PURPOSE:

Processes *LINEAR COMPONENTS card sequence according

DATATR RECENT DRIVER

to the value of TYPE:

Type = 1 End of card sequence

Type = 8 Resistor

Type = 9 Capacitor

Type = 10 Inductor

VARIABLES:

INC = Number of increments for sweep

ISW = End of data switch; Off = 0, On = 1

ITYPE = Sweep code; 0 = No sweep

l = Linear

2 = Logarithmic

M = Intermediate storage of sweep type

NCOMP = Length of linear component data record in words

NIT = Maximum number of sweep iterations

NODEA = Positive node of connection

NODEB = Negative node of connection

NUM = Number of linear components this card

sequence

STRING = Alphanumeric string representation of input

value

STEP = Sweep increment value

STOP = Sweep stop value

TYPES = Valid sweep types

SUBROUTINES CALLED:

ALPHA DATAWR DECIM DRIVWR DRVOUT

ERROR INTEG

CALLING PROGRAMS:

PRCSS

DESCRIPTION:

Subroutine LCIN is the input processor for LINEAR COMPONENTS card sequences. Its function is to decode the linear components input data according to an anticipated format, translating the input into appropriate data and driver records for use by subsequent NCAP phases.

Unlike most NCAP input card sequences which map to a single driver record and a data record of fixed length, linear components card sequences result in one or more driver records, each with an associated data record of variable length. The number of driver records generated is related to the number of linear component sweep iterations defined in the card sequence, while the length of each data record depends on the number of components defined in the card sequence.

Each call to subroutine LCIN causes a single NCAP input card

to be processed. The card image and its TYPE code are transmitted through common storage to subroutine LCIN which processes the card according to its TYPE as follows:

TYPE=1 End of Card Sequence

The closing of a linear components definition begins by testing the input for errors. If no linear component cards were processed (NCOMP.LE.0) or if the input error switch is on (JAET=1), return is made to the calling program with an appropriate error condition. Otherwise the data record is written to file 20 from the first NCOMP words of the common data buffer BUFF by subroutine DATAWR.

The linear components data record is structured as a series of NUM contiguous 4-word entries, one for each linear component defined in the card sequence:

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from the largest IMC canadater specified in the linear component

card requested the notice of the greek transmission bases

PUFF(1)-(4) First Component Definition

BUFF(5)-(8) Second Component Definition

PUFF(NCOMP-3)-(NCOMP) Last Component Definition

The first word of each entry contains an integer code representing the component type: 1 = Resistor, 2 = Capacitor, 3 = Inductor. The second and third words represent the positive and negative nodes of connection. The fourth word contains the

value of the component:

BUFF(1),(5),...(NCOMP-3) Component Codes

PUFF(2),(6),...(NCOMP-2) Positive Nodes of

Connection

BUFF(3),(7),...(NCOMP-1) Negative Nodes of

Connection

BUFF(4),(8),...(NCOMP) Component Values

The linear components driver record is written by subroutine DRVOUT to file 21 from the first ten words of global common, where:

MODE=4 Linear Component Identifier

STADD Data File Record Number

LNGTH=NCOMP Length of Data Record in Words

MISC(1)=NCOMP Number of Linear Components x 4

Other MISC driver parameters are not used.

If component sweeping is not specified in the card sequence (i.e. INC(I)=0 for all I=1,NUM), subroutine LCIN returns to the calling program without further processing. Otherwise the sweep specifications are translated into appropriate driver and data records. First the number of sweep iterations, NIT, is derived from the largest INC parameter specified in the linear component card sequence. (Each of the NIT sweep iterations causes one driver and data record to be generated internally and appended to the disk files.)

Then a sweep increment STEP(I) is calculated for each of the NUM components in the card sequence. The increment is a function of the original component value stored at BUFF (I*4) and the

STOP, INC, and ITYPE parameters associated with the component as follows:

- 1) For ITYPE(I)=0, no sweep is assumed and STEP(I)=0.
- 2) For ITYPE(I)=1,linear sweep is indicated and:

$$STEP(I) = \underbrace{STOP(I) - BUFF(J)}_{INC(I)-1}$$

3) For ITYPE(I)=2, logarithmic sweep is indicated:

STEP(I) =
$$\begin{bmatrix} \frac{1}{\text{STOP}(I)} \\ \frac{1}{\text{EUFF}(J)} \end{bmatrix}$$

After the increments have been calculated the sweep driver and data records are created and appended to the disk files. Like standard driver records, sweep driver records are developed in the first ten words of global common. A component sweep data record consists of:

MODE=24 Component Sweep Identifier

STADD Data File Record Number

LNGTH=NCOMP Length of Data Record in Words

Number of Components x 4

MISC(1)=NCOMP

Sweep data records have the save structure as original linear component data records and are developed in the common data buffer BUFF by an iterative process. The first sweep data record is derived from the original linear component data record in BUFF by performing an incrementation operation on each component value according to its corresponding STEP, STOP, and

ITYPE parameters. The NUM component values in the buffer are incremented one-by-one under the control of the index J. Each component has a corresponding STEP(J), STOP(J), and ITYPE(J) which define the incrementation performed on its value stored at BUFF(J*4):

- For ITYPE(J)=0, no sweep is indicated. No incrementation is performed and BUFF(J*4) remains constant at its original value.
- 2) For ITYPE(J)=1, linear sweep is indicated and the incrementation is additive. For K=J*4:

BUFF(K) = BUFF(K) + STEP(J)

unless STOP(J) is exceeded, in which case BUFF(K) remains constant at its last value.

3) For ITYPE(J)=2, logarithmic sweep is indicated and the incrementation is multiplicative. For K=J*4:

BUFF (K) = BUFF (K) *STEP (J)

If STOP(J) is exceeded, BUFF(K) remains constant
at STOP(J).

After every component value in the buffer has been updated, the new data record is written to file 20 at record number DATREC by subroutine DATAWR. Then subroutine DRIVWR is called to write the sweep driver record to file 21. The data record number STADD is updated and program control transfers back to process the next sweep iteration. The creation of sweep records continues under the control of the index I until all NIT sweep iterations have been translated to data and driver records. Subroutine LCIN returns to the calling program after restoring MODE to 4.

TYPE=8, 9, or 10 Resistor, Capacitor, or Inductor Card

The data record length NCOMP is incremented to accommodate a new four-word component definition and the component count NUM is calculated. If more then 50 linear components have ben processed in the card sequence (NUM.GT.50), subroutine LCIN exits with an appropriate error condition. Otherwise the new component code is set to TYPE-7 and stored in BUFF(NCOMP-3).

Subroutine LCIN anticipates the input of a component identifier, two nodes of connection, and a component value for each component in the card sequence. Furthermore, for element sweeping, three sweep parameters are also expected. During decoding of the input card, if an end-of-data is encountered (ISW=1) before all the required parameters have been processed, subroutine LCIN exits with an appropriate error condition.

The R, C, or L identifier is decoded by subroutine ALPHA. The nodes of connection are decoded and formatted by subroutine INTEG and stored in BUFF(NCOMP-2) and BUFF(NCOMP-1). The component value is decoded and formatted by subroutine DECIM and stored in PUFF(NCOMP). If the component value is zero, subroutine LCIN exits with an appropriate error condition.

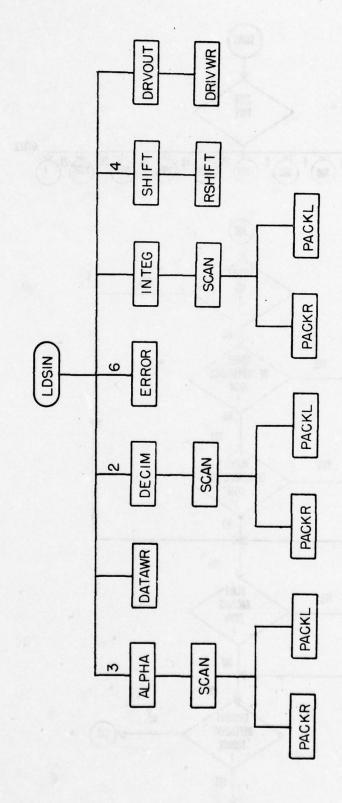
At this point ISW is used to determine the presence of component sweep parameters on the card. If the end-of-data switch is on (ISW=1), the subroutine assumes no sweeping and exits after disabling the sweep parameters ITYPE and INC.

Otherwise the STOP, INC, and ITYPE parameters are decoded and formatted by subroutine DECIM, INTEG, and ALPHA respectively.

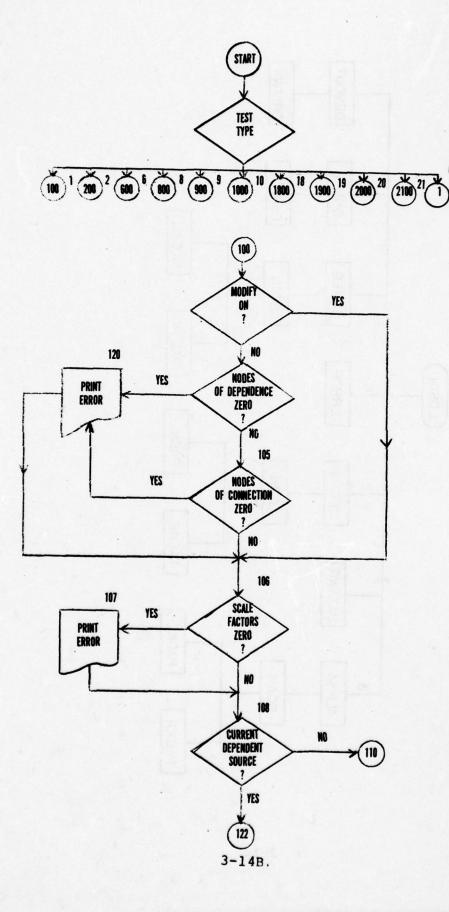
The sweep type is tested against a table of valid alphanumeric types. If a match is found the sweep type is encoded in ITYPE(NUM) and subroutine LCIN returns to the calling program. If no match is found, subroutine LCIN exits with an appropriate error condition.

Any other value of TYPE constitutes an input error which is handled by subroutine ERROR before exit to the calling program.

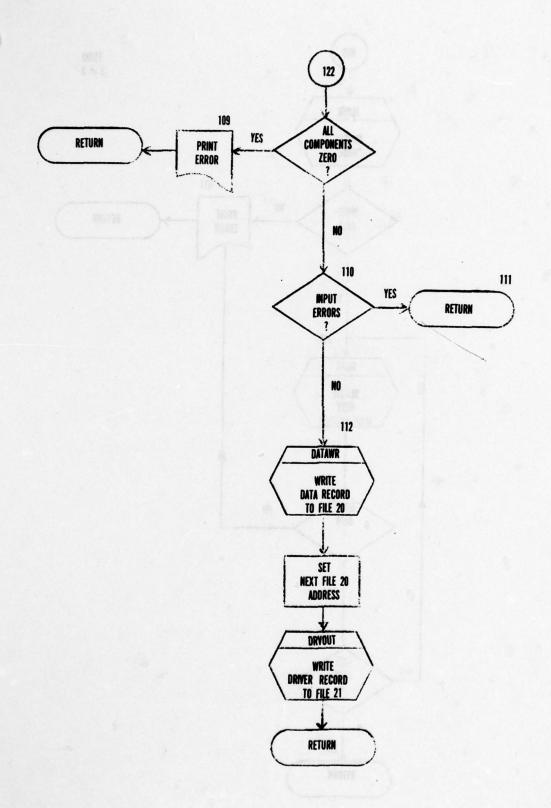
set to TFFE-1 and stored in EDFF (MCORF-3),

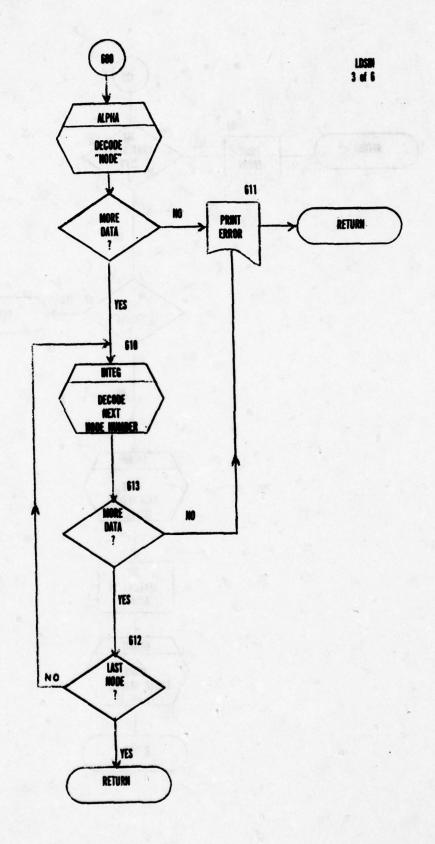


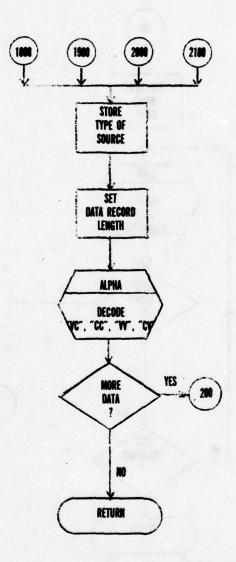


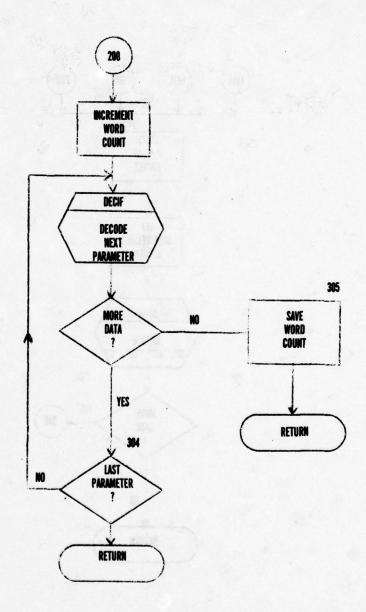


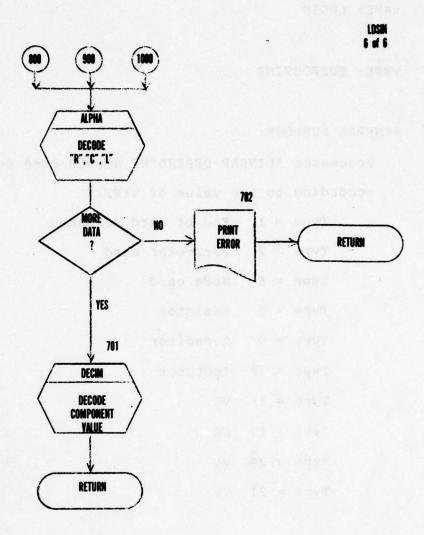


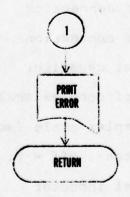












NAME: LDSIN

TYPE: SUPROUTINE

GENERAL PURPOSE:

Processes *LINEAR DEPENDENT SOURCE card sequence according to the value of TYPE:

Type = 1 End of card sequence

Type = 2 Parameter card

Type = 6 Node card

Type = 8 Resistor

Type = 9 Capacitor

Type = 10 Inductor

Type = 18 VC

Type = 19 CC

Type = 20 VV

Type = 21 CV

VARIABLES:

A = Positive node of connection

B = Negative node of connection

C = Value of parallel capacitor

CI = Imaginary part of complex scale factor

CR = Real part of complex scale factor

ISW = End of data switch: Off = 0, On = 1

L = Value of parallel inductor

NW = Word counter for parameter input

R = Value of parallel resistor

STRING = Alphanumeric string representation of input value

SOURCE = Type of source · 1 = VC

2 = CC

3 = VV

XX = Positive node of dependence

YY = Negative node of dependence

SUBROUTINES CALLED:

ALPHA DATAWR DECIM DRIVWR

ERROR INTEC

CALLING PROGRAMS:

PRCSS and and and an all the same of the s

DESCRIPTION:

Subroutine LDSIN is the input processor for LINEAR DEPENDENT SOURCE card sequences. Its function is to decode the linear dependent source input data according to an anticipated format, translating the input into appropriate data and driver records for use by subsequent NCAP phases.

Each call to subroutine LDSIN causes a single NCAP input card to be processed. The card image and its TYPE code are transmitted through common storage to subroutine LDSIN which processes the card according to its TYPE as follows:

TYPE=1 End of Card Sequence

The closing of a linear dependent source definition begins by testing the input for errors. Subroutine LDSIN exits with an appropriate error message if any of the following input errors are detected:

- 1) All four nodes (A, B, XX, and YY) are zero and the card sequence is not a modify function (MODFY=0)
- 2) The complex scale factor (CR, CI) is zero
- 3) For current controlled sources, all three parallel components (R, C, and L) are zero
- 4) The input error switch is on (JABT=1)

 Otherwise the linear dependent source data record is written to file 20 from the common data buffer BUFF and the file 20 record number DATREC is updated to point beyond the new data record.

Because the linear dependent sources do not transmit any internally generated data between NCAP phases, their data records contain only input values. The length of a linear dependent source data record depends on the type of source specified: voltage controlled sources require 2 words of data storage, while current controlled sources require 5 words.

After the data record has been written to disk, subroutine DPVOUT is called to complete the driver record definition and write it to file 21. The linear dependent source driver record is developed in the first ten words of global common and consists

of .

MODE=25	Linear Dependent Source Identifier
STADD	Data File Record Number
LNGTH=2 or 5	Length of Data Record in Words
MISC(1) = SOURCE	Source Code: 1=VC, 2=CC, 3=VV,
	4=CV.
MISC(3)=XX	Positive Node of Dependence
MISC(4)=YY	Negative Node of Dependence
MISC (5) = A	Positive Node of Connection
MISC(6)=P	Negative Node of Connection

After returning from subroutine DRVOUT, subroutine LDSIN returns to the calling program.

TYPE=2 Parameter Card

One or two decimal input values are extracted from the card image, formatted and stored in the data buffer. The number of values decoded and their placement in the buffer depend on the parameter count, NXT, and the end-of-data switch, ISW. Upon entering the subroutine, NXT contains the number of values input on previous data cards. Parameter storage begins at BUFF(NW) where NW=NXT+1 and proceeds through successive locations as data values are decoded by subroutine DECIM. If an end-of-data is encountered (ISW=1) before PUFF(2) is filled, the parameter count is reset to the last buffer location used (NXT=NW) and control returns to the calling program in anticipation of additional

parameter cards.

TYPE=6 Node Card

The NODE identifier is decoded by subroutine ALPHA. Then the nodes of dependence and nodes of connection are extracted from the card image and formatted by subroutine INTEG and stored in MISC(3) through MISC(6). If an end-of-data is encountered (ISW=1) before four node numbers have been decoded, subroutine LDSIN exits with an appropriate error condition.

TYPE= 8, 9, or 10 Resistor, Capacitor, or Inductor Card

The R, C, or L identifier is decoded by subroutine ALPHA.

If an end-of-data is encountered (ISW=1), subroutine LDSIN exits with an appropriate error condition. Otherwise the component value is extracted from the card image, formatted and stored at BUFF(TYPE-5) by subroutine DECIM.

TYPE= 18, 19, 20, or 21 VC; CC, VV, or CV Card

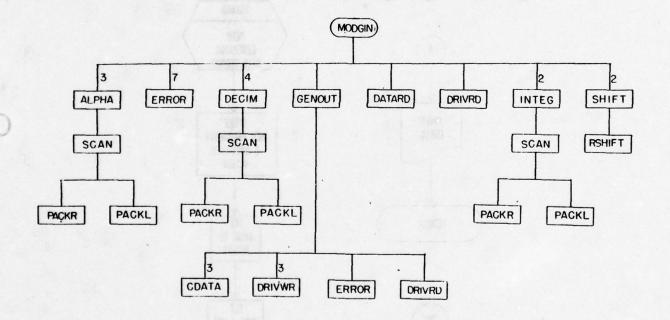
The SOURCE type is calculated according to TYPE=17 and the data record length is set: LNGTH=2 for voltage controlled sources, LNGTH=5 for current controlled sources. The VC, CC, VV,

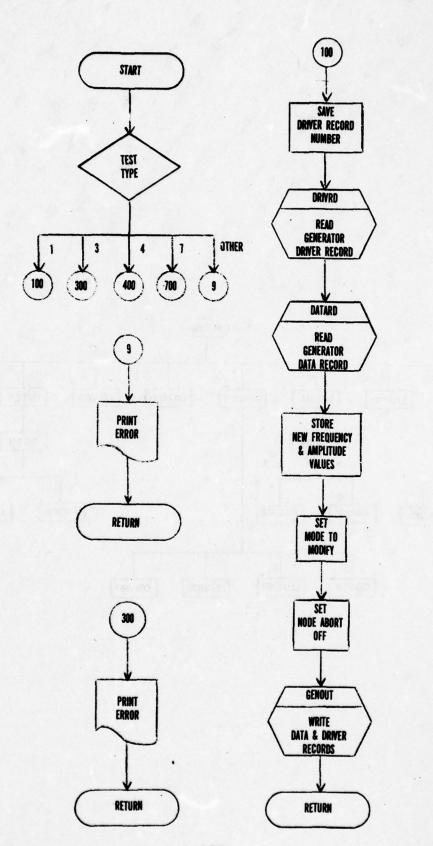
or CV identifier is decoded by subroutine ALPHA. If an end-of-data is encountered (ISW=1), subroutine LDSIN returns to the calling program. Otherwise control transfers to TYPE=2 processing for the remaining data on the card.

Any other value of TYPE constitutes an input error which is handled by subroutine ERROR before exit to the calling program.

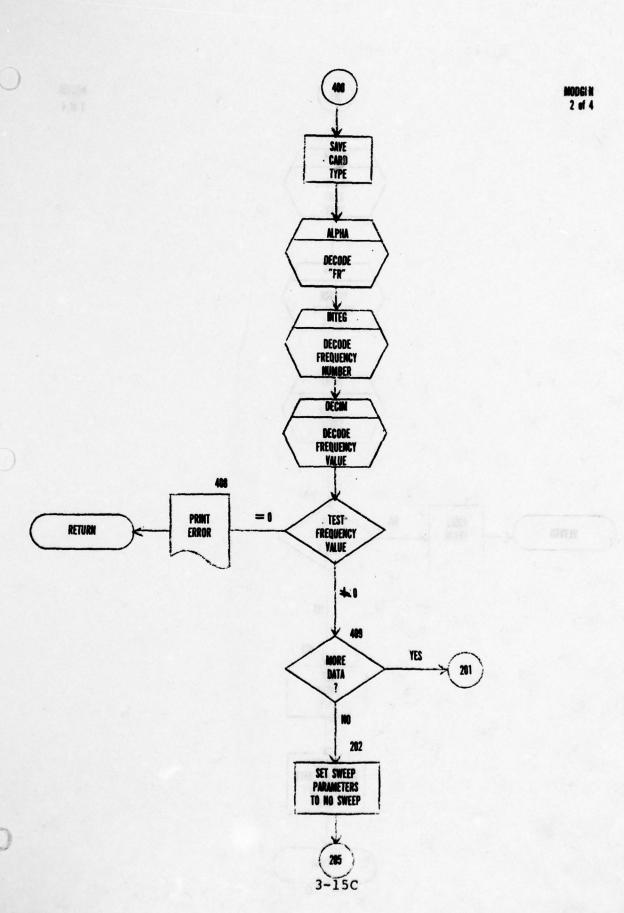
or CV identifier is decoded by subroutine Alpha. If an end-ef-date is encountered (ISW=1), subroutine tosin returns to the coling program. Coherwise control transfers to TYPE=2 processing for elected the card.

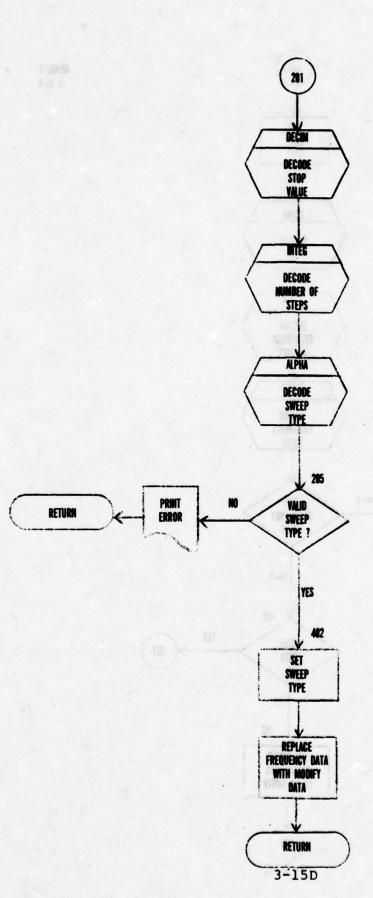
Any other value of TYPE conditiones an imput error which is handled by subvoiding ERROR before exit to the calling program.





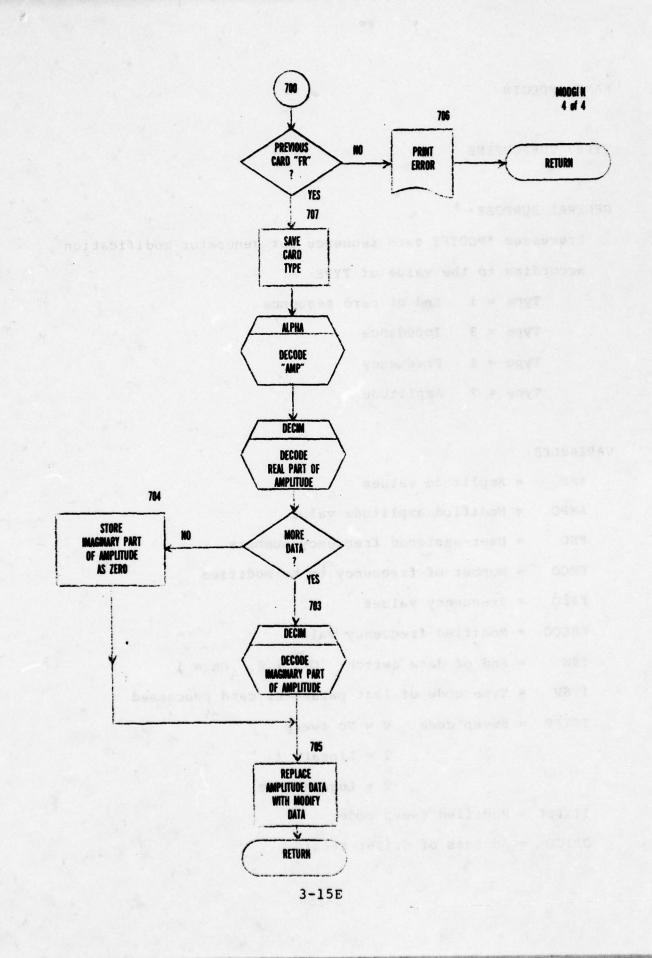
MODGM 1 of 4





MODEL 3 of 4

1



NAME · MODGIN

TYPE · SUBROUTINE

GENERAL PURPOSE .

Processes *MODIFY card sequence for generator modification according to the value of TYPE:

Type = 1 End of card sequence

Type = 3 Impedance

Type = 4 Frequency

Type = 7 Amplitude

VARIABLES:

AMP = Amplitude values

AMPO = Modified amplitude value

FNO = User-assigned frequency numbers

FNOO = Number of frequency to be modified

FREC = Frequency values

FRECO = Modified frequency value

ISW = End of data switch. Off = 0, On = 1

ITSV = Type code of last parameter card processed

ITYPE = Sweep code · 0 = No sweep

1 = Linear

2 = Logarithmic

ITYPEC = Modified sweep code

JRECO = Address of driver record

NFREQ = Number of frequencies this generator

NIT = Number of sweep iterations in generator being modified

STRING = Alphanumeric string representation of input value

STOP = Stop value for sweep

STOPC = Modified stop value

TYPES = Valid frequency sweep types

SUBROUTINES CALLED.

ALPHA DATARD DECIM DRIVRD FRROR
GENOUT INTEG

CALLING PROCRAMS.

PRCSS

DESCRIPTION .

Subroutine MODGIN is the input processor for MODIFY GENERATOR card sequences. Its function is to decode the generator modification input data according to an anticipated format, translating the input into appropriate data and driver records for use by subsequent NCAP phases.

Each call to subroutine MODGIN causes a single NCAP input card to be processed. The card image and its TYPE code are transmitted through common storage to subroutine MODGIN which processes the card according to its TYPE as follows:

TYPE=1 End of Card Sequence

The modify generator data record is created by replacing the frequency and amplitude values of the previous generator record with the updated FREQ and AMP values processed by subroutine MODGIN. After saving the present driver record number in JRECO, the previous generator driver record is read by subroutine DRIVRD (its record number is obtained by backspacing over the NIT frequency sweep driver records which are stored between it and the present driver record at DRVREC). Using the record number STADD, the previous generator data is read from file 20 by subroutine DATARD and stored in the common data buffer BUFF.

The updated frequency values are stored in BUFF(1)-(10) and the updated amplitudes in BUFF(11)-(30). MODE is set to 21 to identify the generator modification function, the node card switch is disabled, and subroutine GENOUT is called to write the data and driver records to disk.

TYPE=3 Impedance Card

Since generator impedances cannot be modified, any impedance card in a modify generator card sequence constitutes an input error which is handled by subroutine ERROR before exit to the calling program.

TYPE=4 Frequency Card

The card type code is saved in ITSV and the FR identifier is decoded by subroutine ALPHA. The frequency number is decoded and formatted by subroutine INTEG and stored in FNOO. The frequency value is decoded and formatted by subroutine DECIM and stored in FREÇO. If the frequency value is zero, subroutine MODGIN exits with an appropriate error condition.

At this point the end-of-data switch ISW is used to determine the presence of frequency sweep parameters on the card. If the end-of-data in on (ISW=1), the parameters ITYPEO and INCO are set to indicate no sweep. Otherwise the sweep parameters STOPO, INCO, and ITYPEO are decoded and formatted by subroutines DECIM, INTEG, and ALPHA respectively.

The sweep type is tested against a table of valid alphanumeric types. If no match is found, subroutine MODGIN exits with an appropriate error condition. If a match is found, the sweep type is encoded in ITYPEO.

The number of the frequency being modified FNOO is compared against the frequency numbers defined for the previous generator FNO(I), I=1,10. If no match is found, subroutine MODGIN exits with an appropriate error condition. If a match is found, the table index I serves as the index for replacing FREQ(I), STOP(I), INC(I), and ITYPE(I) with the modified values, after which subroutine MODGIN returns to the calling program.

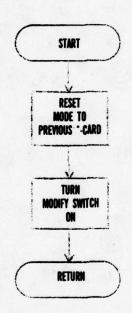
TYPE=7 Amplitude Card

In order for amplitude data to be associated with the proper frequency, each amplitude card must physically follow a frequency card in the input deck. Therefore if the last card processed was not a frequency card (ITSV \neq 4), subroutine MODGIN exists with an appropriate error condition.

Otherwise the card type code is saved in ITSV and the AMP identifier is decoded by subroutine ALPHA. The real and imaginary parts of the complex amplitude are decoded and formatted by subroutine DECIM and stored in AMPO(1) and AMPO(2). The number of the frequency being modified FNOO is compared against the frequency numbers defined for the previous generator FNO(I), I=1,10. If no match is found, subroutine MODGIN exits with an appropriate error condition. If a match is found, the table index I is used to replace the original amplitude data AMP(1,I) and AMP(2,I) with the modified amplitude, after which subroutine MODGIN returns to the calling program.

Any other value of TYPE constitutes an input error which is handled by subroutine ERROR before exit to the calling program.

MODIN 1 of 1



resets bord to identify the develop helps modified, turned the

NAME: MODIN

TYPE: SUBROUTINE

CENERAL PURPOSE:

Process *MODIFY card, resets MODE to identify device being modified, turns modify switch 'ON' so remainder of card sequence is processed according to MODE of device being modified

VAPIABLES:

NONE

SUBROUTINES CALLED:

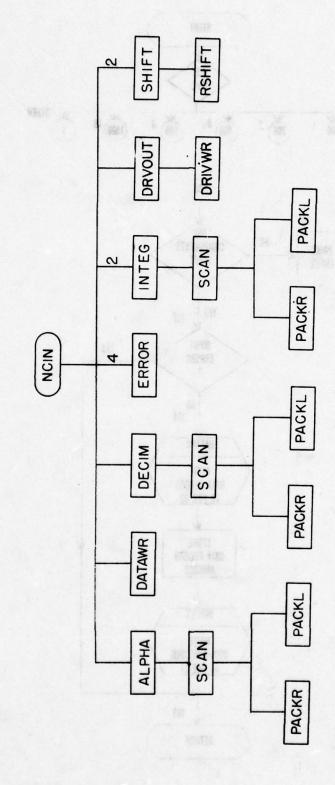
NONE

CALLING PROCRAMS:

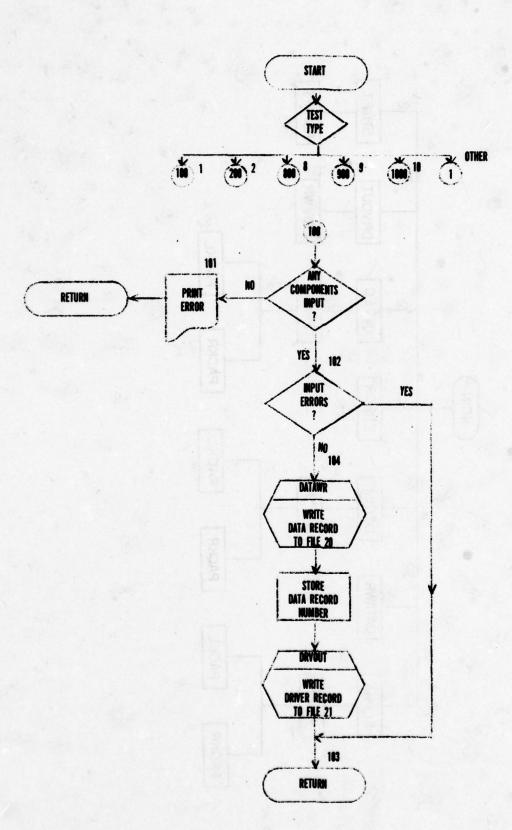
PRCSS

DESCRIPTION .

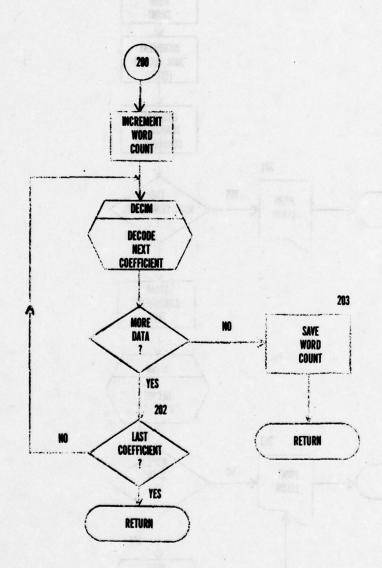
Subroutine MODIN initiates the processing of a MODIFY card sequence. Rather than processing the input cards itself, it resets MODE to identify the device being modified, turns the modify switch on (MODFY=1), and returns to the calling program where the remaining cards of the sequence are processed according to the MODE of the device being modified.



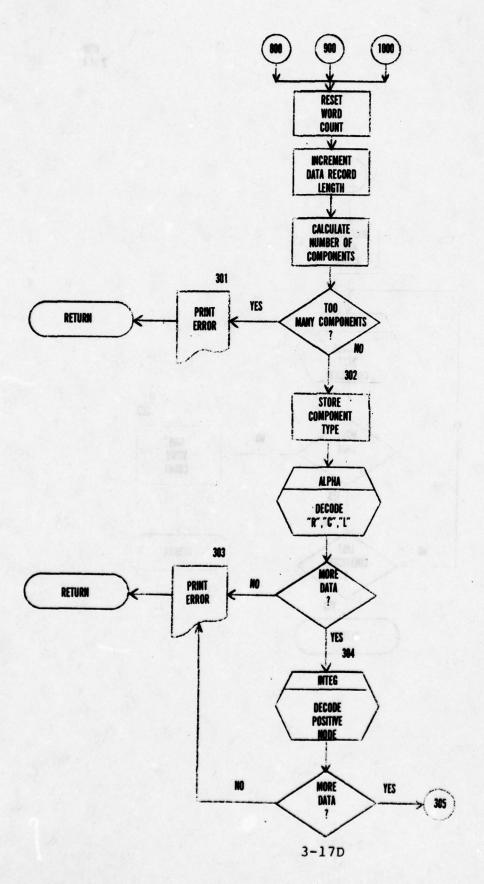
0

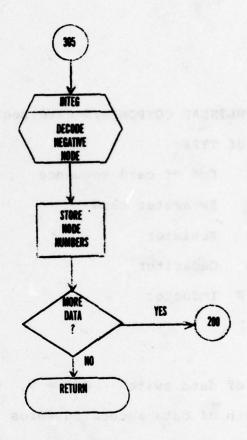


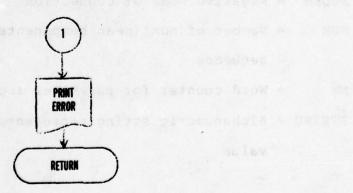
NCM 1 of 4











NAME: NCIN

TYPE · SUBROUTINE

GENERAL PURPOSE:

Processes *NONLINEAR COMPONENTS card sequence according to the value of TYPE:

Type = 1 End of card sequence

Type = 2 Parameter card

Type = 8 Resistor

Type = 9 Capacitor

Type = 10 Inductor

VARIABLES.

ISW = End of data switch \cdot Off = 0, On = 1

NCOMP = Length of data record in words

NCDEA = Positive node of connection

NCDER = Negative node of connection

NUM = Number of nonlinear components this card sequence

NW = Word counter for parameter input

STRING = Alphanumeric string representation of input value

SUBROUTINES CALLED.

ALPHA DATAWR DECIM DRVOUT ERROR INTEG

CALLING PROGRAMS:

PRCSS

DESCRIPTION .

Subroutine NCIN is the input processor for NONLINEAR COMPONENTS card sequences. Its function is to decode the nonlinear components input data according to an anticipated format, translating the input into appropriate data and driver records for use by subsequent NCAP phases.

Each call to subroutine NCIN causes a single NCAP input card to be processed. The card image and its TYPE code are transmitted through common storage to subroutine NCIN which processes the card according to its TYPE as follows:

TYPE = 1 End of Card Sequence

The closing of a nonlinear components definition begins by testing the input for errors. If no linear component cards were processed (NCOMP.LE.0) or if the input error switch is on (JAPT=1), return is made to the calling program with an appropriate error condition. Otherwise the data record is written to file 20 from the first NCOMP words of the common data buffer BUFF by subroutine DATAWR.

The nonlinear components data record is structured as a series of NUM contiguous 13-word entries, one for each nonlinear

component defined in the card sequence:

BUFF(1)-(13)

First Component Definition

BUFF(14)-(26)

Second Component Definition

PUFF(NCCMP-12)-(NCCMP)Last Component Definition

The first word of each entry contains an integer code representing the component type: l=resistor, 2=capacitor, 3=inductor. The second and third words represent the positive and negative nodes of connection. The last ten words contain from one to ten nonlinear coefficients which define the element:

PUFF(1), (14)...(NCOMP-12) Component Codes

PUFF(2),(15)...(NCOMP-11)
Positive Nodes of Connection

BUFF(3),(16)...(NCCMP-10) Negative Nodes of Connection

PUFF(4),(17)...(NCOMP-9) First Order Coefficients

BUFF(5),(18)...(NCOMP-8) Second Order Coefficients

PUFF(13), (26)...(NCCMP) Tenth Order Coefficients

The nonlinear components driver record is completed and written by subroutine DRVOUT to file 21 from the first ten words of global common, where.

MODE=5

Nonlinear Component Identifier

STADD

Data File Record Number

LNGTH=NCOMP Length of Data Record in Words

MISC(1)=NCOMP Number of Nonlinear Components x 4

Other MISC driver parameters are not used.

TYPE=2 Parameter Card

From one to ten decimal parameters are extracted from the card image, formatted, and stored in the data buffer by subroutine DECIM. The number of values decoded and their placement in the buffer depend on the parameter count, NXT, and the end-of-data switch, ISW. Upon entering the subroutine, NXT contains the number of coefficient values input on previous data cards. Parameter storage starts at BUFF(NCOMP-10+NW) where NW=NXT+1 and proceeds through successive buffer locations as coefficient values are decoded by subroutine DECIM. If an end-of-data is encountered (ISW=1) before PUFF(NCOMP) is filled, the parameter count is reset to the last buffer location used (NXT=NW) and control returns to the calling program in anticipation of additional parameter cards.

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TYPE=8, 9, OR 10 Resistor, Capacitor, or Inductor Card

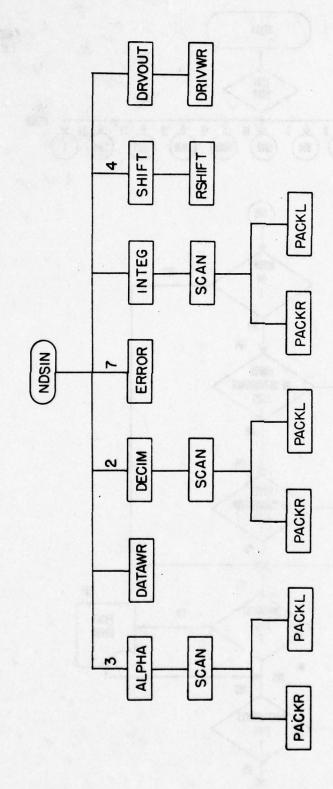
The parameter count is reset to NXT=0, the data record length is incremented to accommodate a new 13-word component definition, and the component count NUM is calculated. If more

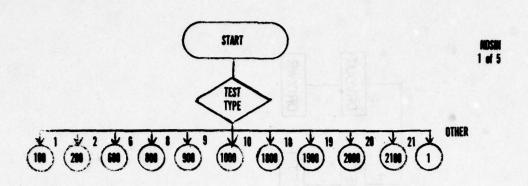
than 10 nonlinear components have been processed in the card sequence (NUM.GT.10), subroutine NCIN exits with an appropriate error condition. Otherwise the new component code is calculated by TYPE-7 and stored in BUFF(NCOMP-12).

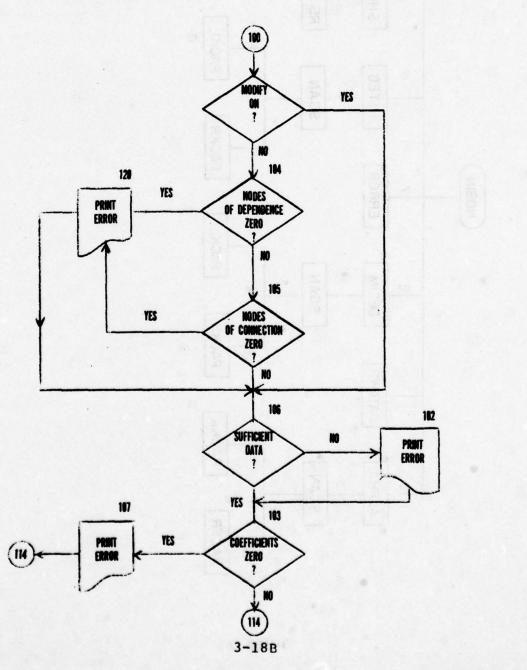
For each nonlinear component card, subroutine NCIN anticipates the input of a component identifier and two nodes of connection on one card. During decoding of the input card, if an end-of-data is encountered (ISW=1) before these parameters have been processed, subroutine NCIN exits with an appropriate error condition.

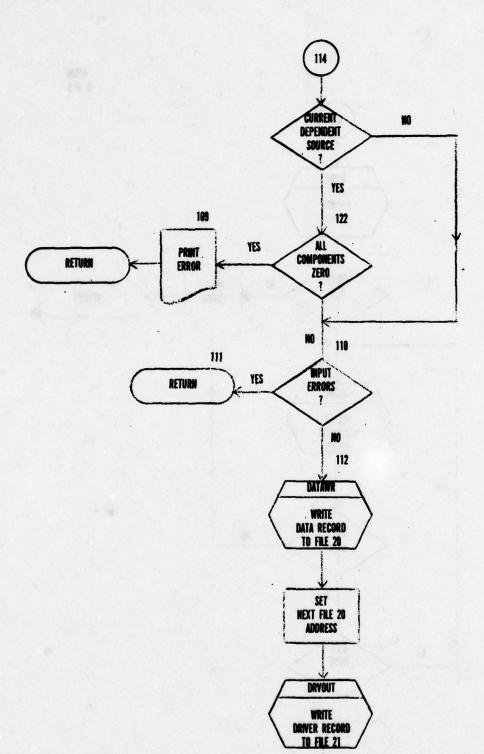
The R, C, or L identifier is decoded by subroutine ALPHA. The nodes of connection are decoded and formatted by subroutine DECIM and stored in PUFF(NCOMP-11) and BUFF(NCOMP-10). If the end-of-data switch is off (ISW=0), control transfers to decode the coefficients which follow on the card according to TYPE=2 processing. If the end-of-data switch is on (ISW=1), subroutine NCIN returns to the calling program.

Any other value of TYPE constitutes an input error which is handled by subroutine ERROR before exit to the calling program.



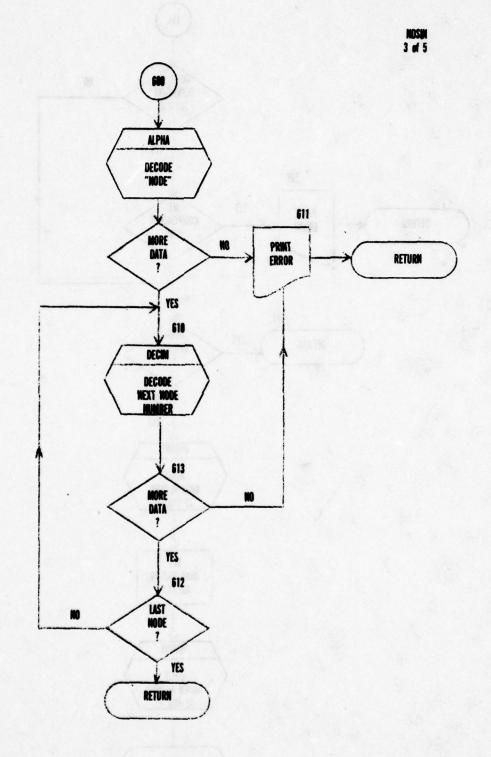


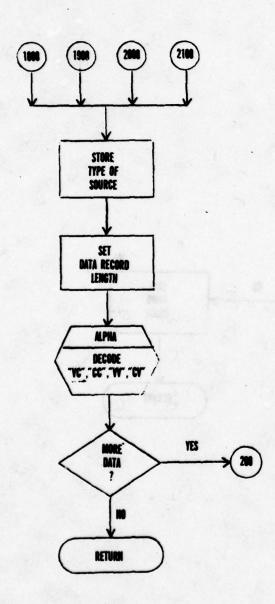


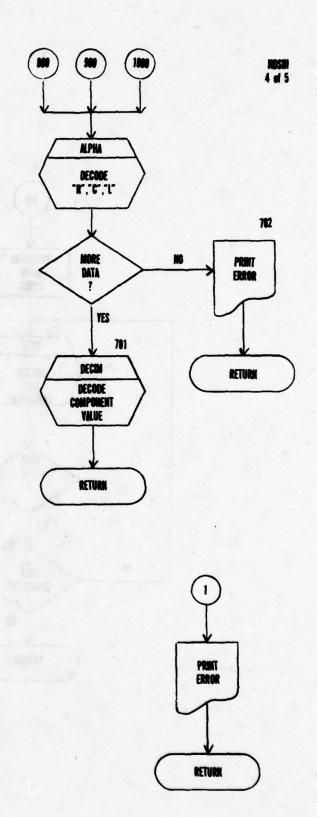


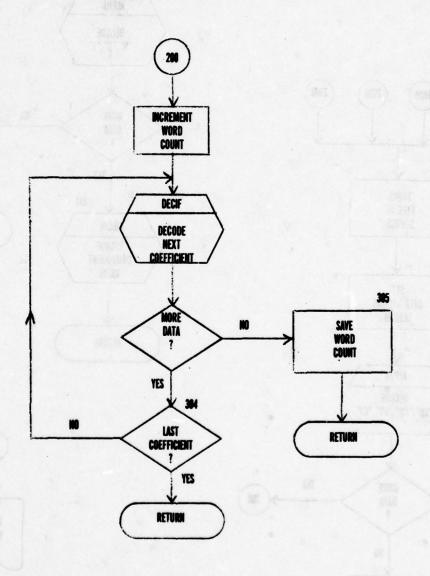
NDSIN 2 of 5

RETURN









NAME: NDSIN

TYPE: SUBROUTINE

GENERAL PURPOSE:

Processes *LINEAR DEPENDENT SOURCE card sequence according to the value of TYPE:

Type = 1 End of card sequence

Type = 2 Parameter card

Type = 6 Node card

Type = 8 Resistor

Type = 9 Capacitor

Type = 10 Inductor

Type = 18 VC

Type = 19 CC

Type = 20 VV

Type = 21 CV

VARIABLES:

A = Positive node of connection

B = Negative node of connection

C = Value of parallel capacitor

ISW = End of data switch: off = 0, on = 1

L = Value of parallel inductor

NW = Word counter for parameter input

R = Value of parallel resistor

STRING = Alphanumeric string representation of input

value

SOURCE = Type of source; 1 = VC

2 = CC

3 = VV

4 = CV

XX = Positive node of dependence

YY = Negative node of dependence

SUPROUTINES CALLED:

ALPHA DATAWR DECIM DRIVWR

ERROR INTEG

CALLING PROGRAMS:

PRCSS

DESCRIPTION .

Subroutine NDSIN is the input processor for NONLINEAR DEPENDENT SOURCE card sequences. Its function is to decode the nonlinear source input data according to an anticipated format, translating the input into appropriate data and driver records for use by subsequent NCAP phases.

Each call to subroutine NDSIN causes a single NCAP input card to be processed. The card image and its TYPE code are transmitted through common storage to subroutine NDSIN which processes the card according to its TYPE as follows:

TYPE=1 End of Card Sequence

The closing of a nonlinear dependent source definition begins by testing the input for errors. Subroutine NDSIN exits with an appropriate error message if any of the following input errors are detected:

- All four nodes (A, B, XX, and YY) are zero and the card
 sequence is not a modify function (MODFY=0)
- No coefficients were processed (NXT=0)
- 3) The coefficients are all zero
- 4) For current controlled sources, all three parallel components (R, C, and L) are zero
- 5) The input error switch is on (JABT=1)
 Otherwise the nonlinear dependent source data record is written to file 20 from the common data buffer BUFF and the file 20

record number DATREC is updated to point beyond the new data

record.

The data record allocates sufficient disk storage for the nonlinear dependent source input values as well as all other data generated internally by its model in subsequent NCAP phases. The length of a nonlinear dependent source data record depends on the type of source specified: voltage controlled sources require 20 words of storage for input values only, while current controlled sources require 86 words of storage as follows:

PUFF(1)-(20) Input Values

PUFF(21)-(23) Component Values

PUFF(24)-(86) Impedance Table

After the data record has been written to disk, subroutine
DRVOUT is called to complete the driver record and write it to
file 21. The nonlinear dependent source driver record, taken
from the first ten words of global common, consists of:

MODE=26	Nonlinear Dependent Source
	Identifier 59109195 91
STADD	Data File Record Number
LNGTH=20 or 86	Length of Data Record in Words
MISC(1) = SOURCE	Source Code: 1=VC, 2=CC, 3=VV
	4=CV
MISC(3)=XX	Positive Node of Dependence
MISC(4)=YY	Negative Node of Dependence
MISC(5)=A	Positive Node of Connection
MISC(6)=P	Negative Node of Connection

After returning from subroutine DRVOUT, subroutine NDSIN exits to the calling program.

TYPE=2 Parameter Card add to the second and the sec

From one to ten nonlinear coefficients are extracted from the card image, formatted, and stored in the data buffer. The number of values decoded and their placement in the buffer depend on the parameter count, NXT, and the end-of-data switch, ISW. Upon entering the subroutine, NXT contains the number of

coefficients input on previous data cards. Coefficient storage begins at BUFF(NW*2-1) where NW=NXT+1, and proceeds through successive odd buffer locations as the coefficients are decoded by subroutine DECIM. The corresponding even buffer locations BUFF(NW*2), which are allocated for future implementation of imaginary parts of complex coefficients, are zeroed. If an end-of-data is encountered (ISW=1) before BUFF(19) is filled, the parameter count is reset to the last buffer location used (NXT=NW) and control returns to the calling program in anticipation of additional parameter cards.

TYPE=6 Node Card

The NODE identifier is decoded by subroutine ALPHA. Then the nodes of dependence and nodes of connection are extracted from the card image and formatted by subroutine INTEC and stored in MISC(3) through MISC(6). If an end-of-data is encountered (ISW=1) before four node numbers have been decoded, subroutine NDSIN exits with an appropriate error condition.

TYPE=8, 9, or 10 Resistor, Capacitor, or Inductor

The R, C, or L identifier is decoded by subroutine ALPHA.

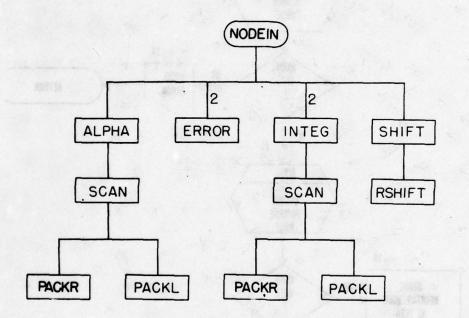
If an end-of-data is encountered (ISW=1), subroutine NDSIN exits with an appropriate error condition. Otherwise the component value is extracted from the card image, formatted, and stored at

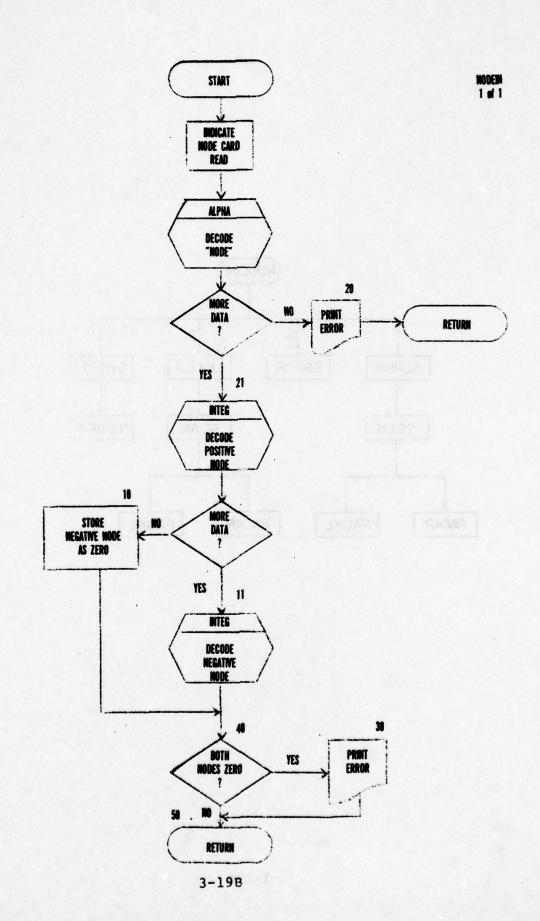
PUFF (TYPE+13) by subroutine DECIM.

TYPE= 18, 19, 20, or 21 VC, CC, VV, or CV Card

The SOURCE type is calculated according to TYPE-17 and the data record length is set: LNGTH=20 for voltage controlled sources, LNGTH=86 for current controlled sources. The VC, CC, VV, or CV identifier is decoded by subroutine ALPHA. If the end-of-data is encountered (ISW=1), subroutine NDSIN returns to the calling program. Otherwise control transfers to TYPE=2 processing for the remaining data on the card.

Any other value of TYPE constitutes an input error which is handled by subroutine ERROR before exit to the calling program.





NAME: NODEIN

TYPE: SUBROUTINE

GENERAL PURPOSE:

Processes NODE cards

VARIABLES.

ISW = End of data switch: Off = \emptyset , On = 1

NODEA = Positive node of connection

NODEB = Negative node of connection

STRING = Alphanumeric string representation of input

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SUBROUTINES CALLED:

ALPHA ERROR INTEG

CALLING PROGRAMS:

GENIN JFETIN PTIN SDIN TRNIN

VDIN VTIN

DESCRIPTION .

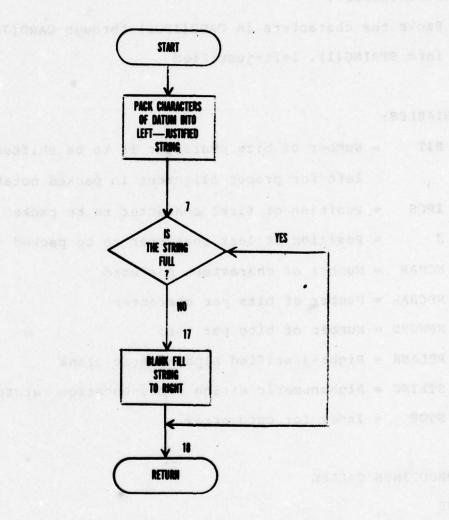
Subroutine NODEIN decodes and formats the standard NCAP node card. The card image is transmitted through common storage to subroutine NODEIN where it is processed according to an anticipated format. The node specifications are returned to the

calling program through the arguments NODEA and NODEB.

First the node error switch is turned off (NABT=0) and the NODE identifier is decoded by subroutine ALPHA. If an end-of-data is encountered (ISW=1), subroutine NODEIN exits with an appropriate error condition.

If the end-of-data switch is off (ISW=0), the positive node of connection is extracted from the card image, formatted, and stored in NODEA by subroutine INTEG. If the end-of-data is encountered (ISW=1), the negative node of connection is assumed to be ground (NODEB=0). Otherwise the negative node is extracted from the card image, formatted and stored in NODEB by subroutine INTEG.

In the event that both nodes are zero, subroutine NODEIN exits with an appropriate error condition.



NAME · PACKL

TYPE: SUBROUTINE

GENERAL PURPOSE:

Packs the characters in CARD(IPOS) through CARD(J-1) into STRING(1), left-justified

VARIABLES:

PIT = Number of bits character is to be shifted

left for proper alignment in packed notation

IPCS = Position of first character to be packed

J = Position of last character to be packed + 1

NCHAR = Number of characters per word

NPCHAR = Number of bits per character

NBWORD = Number of bits per word

RPLANK = Right-justified alphanumeric blank

STRING = Alphanumeric string representation returned

STOP = Index for card array

SUBROUTINES CALLED:

NONE

CALLING PROGRAMS:

SCAN

DESCRIPTION :

Subroutine PACKL belongs to a group of Phase 0 character manipulation routines which format NCAP input data. Its function is to pack the characters comprising an alphabetic datum extracted by subroutine SCAN into a left-justified alphanumeric string in preparation for formatting by subroutine ALPHA.

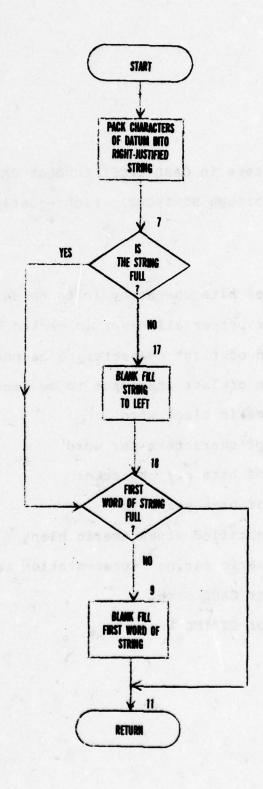
Input to the subroutine is the right-justified card image stored in the integer array CARD, the character position IPOS at which the datum begins, and the character position J at which the datum ends. Output from the routine is the left-justified alphanumeric string representation returned in STRING(1). The required machine dependent character manipulation parameters NCHAR (number of characters per word), NBCHAR (number of bits per character), and NPWORD (number of bits per word) are transmitted to subroutine PACKL through the labelled common area WORDSZ.

The NEWORD bits of STRING(1) contain NCHAR characters, each occupying NBCHAR bits. The characters are stored in the string from left to right, where the leftmost character of the datum taken from CARD (IPOS) is placed in the most significant bits of STRING(1) and so on.

The packing technique is arithmetic: each right-justified character to be packed, CARD(STOP), is shifted left into the correct position by multiplication by an appropriate power of 2, (2 ** BIT), and then simply added to the string, STRING(1). The variable STOP indexes the CARD array, while BIT represents the number of bits a character must be shifted left for proper positioning within the string.

If fewer than NCHAR characters are input to the routine, the string is filled to the right with alphanumeric blanks. If more than NCHAR characters are input, the right-most characters are truncated.





NAME: PACKR

TYPE: SUBROUTINE

GENERAL PURPOSE:

Packs the characters in CARD(IPOS) through CARD(J-I) into STRING(1) through STRING(2), right-justified

VARIABLES:

BIT = Number of bits character is to be shifted

left for proper alignment in packed notation

IPOS = Position of first character to be packed

J = Position of last character to be packed + 1

LBLANK = Alphanumeric blank word

NCHAR = Number of characters per word

NBCHAR = Number of bits per character

NEWORD = Number of bits per word

RBLANK = Right-justified alphanumeric blank

STRING = Alphanumeric string representation returned

STOP = Index for CARD array

WORD = Index for STRING array

SUBROUTINES CALLED:

NONE

CALLING PROGRAMS:

SCAN

DESCRIPTION:

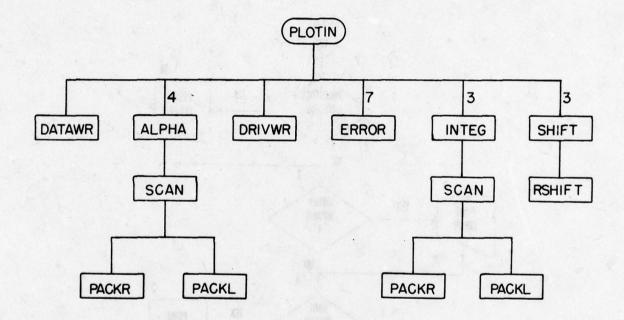
Subroutine PACKR belongs to a group of Phase Ø character manipulation routines which format NCAP input data. Its function is to pack the characters comprising a numeric data value extracted by subroutine SCAN into a right-justified alphanumeric string in preparation for formatting by subroutines DECIM and INTEG.

Input to the routine is the right-justified card image stored in the integer array CARD, the character position IPOS at which the datum begins, and the character position J at which the datum ends. Output from the routine is the right-justified alphanumeric string representation of the datum returned in STRING(1) and STRING(2). The required machine dependent character manipulation parameters NCHAR (number of characters per word), NBCHAR (number of bits per character), and NBWORD (number of bits per word) are transmitted to subroutine PACKR through the labelled common area WORDSZ.

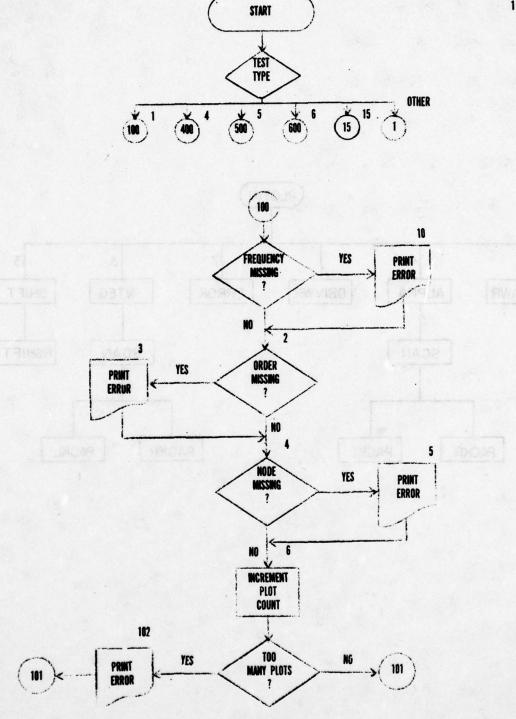
The NBWORD bits of each word of the alphanumeric string representation contain NCHAR characters, each occupying NBCHAR bits. The characters are stored from right to left (from the least significant to the most significant bit), where the least significant character of the string occupies bit positions 1 through NBCHAR and so on. STRING(2) is packed with the NCHAR characters from CARD(J-1), CARD(J-2)..., CARD(J-NCHAR). STRING(1) is packed with the remaining characters through CARD(IPOS). The elements of STRING are blank filled to the left if necessary.

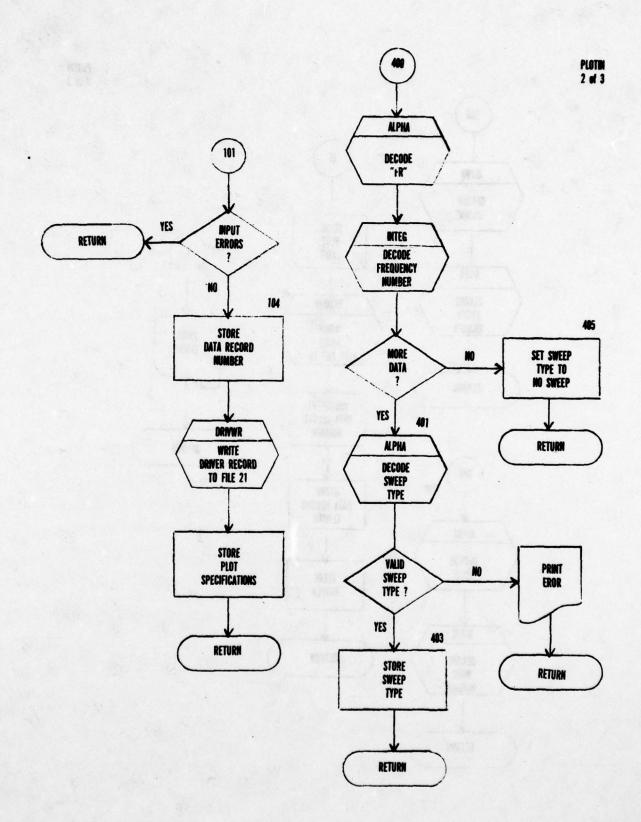
The packing technique is arithmetic: each right-justified character to be packed, CARD(STOP), is shifted left into the correct position by multiplication by an appropriate power of 2, (2 * PIT), and then simply added to the string at STRING(WORD). The variables STOP and WORD index the CARD and STRING arrays, while BIT represents the number of bits a character must be shifted left for proper positioning within the string.

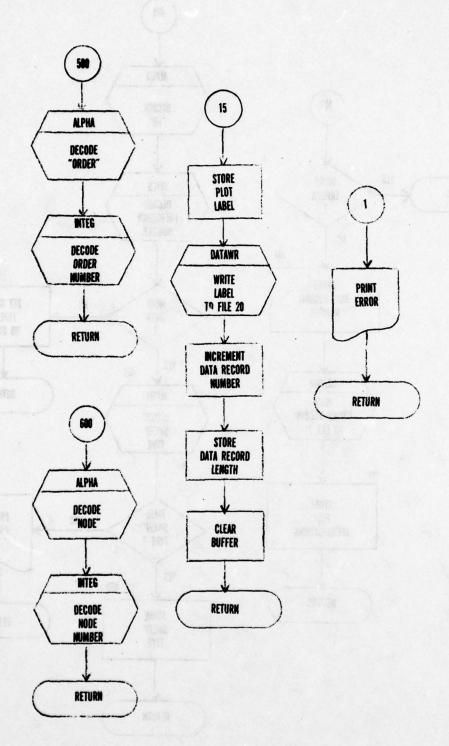
If fewer than 2 x NCHAR characters are input to the routine, the string is filled to the left with alphanumeric blanks. If more than 2 x NCHAR characters are input, the more significant characters are truncated.











NAME: PLOTIN

TYPE: SUBROUTINE

GENERAL PURPOSE:

Processes *PLOT card sequence according to the value of TYPE

Type = 1 End of card sequence

Type = 4 Frequency

Type = 5 Order

Type = 6 Node

Type = 15 Label

VARIABLES:

IFR = User-assigned number of frequency to be
 plotted

INOD = Modes to be plotted

IORD = Orders to be plotted

ISW = End of data switch · Off = 0, On = 1

ITYP = Type of plot; 1 = Linear

2 = Logarithmic

STRING = Alphanumeric string representation of input
value

TYPES = Valid frequency sweep types

SUPROUTINES CALLED:

ALPHA DATAWR DRIVWR ERROR INTEG

CALLING PROGRAMS:

PRCSS

DESCRIPTION .

Subroutine PLOTIN is the input processor for PLOT card sequences. Its function is to decode the plot input data according to an anticipated format, translating the input into appropriate data and driver records for use by subsequent NCAP phases.

Each call to subroutine PLOTIN causes a single NCAP input card to be processed. The card image and its TYPE are transmitted through common storage to subroutine PLOTIN which processes the card according to its TYPE as follows:

IFB = Unri-sesigned number of frequency to de

TYPE=1 End of Card Sequence

The closing of a plot definition begins by testing the input for errors. Subroutine PLOTIN exits with an appropriate error message if any of the following input errors are detected:

- No frequency data was processed (MISC(1)=0)
- 2) No order data was processed (MISC(3)=0)
- No node data was processed (MISC(4)=0)
- 4) More than 10 plot specifications were processed (NPLOT.GT.10)

5) The input error switch is on (JABT=1)

If no input errors are detected, the file 20 record number DATREC is updated to point beyond the plot data record and subroutine DRIVWR is called to write the driver record to file 21.

The PLOT driver record, taken from the first ten words of global common, consists of:

MODE=22 Plot Identifier

STADD Data File Record Number

LNGTH=0,50, or 100 Length of Data Record in Words

MISC(1) Frequency Number

MISC(2) Type of Plot: l=Linear

2-Logarithmic

MISC(3) Order Number

MISC(4) Node Number

Other MISC driver parameters are not used.

After the driver record has been written to disk, the plot specifications are appended to the plot table in the labelled common area PLOTFL and subroutine PLOTIN returns to the calling program.

TYPE=4 Frequency Card

The FR identifier is decoded by subroutine ALPHA and the frequency number is extracted from the card image, formatted, and

stored in MISC(1) by subroutine INTEG. If an end-of-data is encountered (ISW=1), the plot type is set to "no sweep" by MISC(2)=1 and subroutine PLOTIN returns to the calling program.

Otherwise the plot type is extracted from the card image by subroutine ALPHA and tested against a table of valid alphanumeric types. If no match is found, subroutine PLOTIN exits with an appropriate error condition. If a match is found, the plot type is encoded in MISC(2) and subroutine PLOTIN returns to the calling program.

TYPE=5 Order Card Andrews Andrews Andrews

The ORDER identifier is decoded by subroutine ALPHA and the order number is extracted from the card image, formatted, and stored in MISC(3) by subroutine INTEG. Subroutine PLOTIN returns to the calling program.

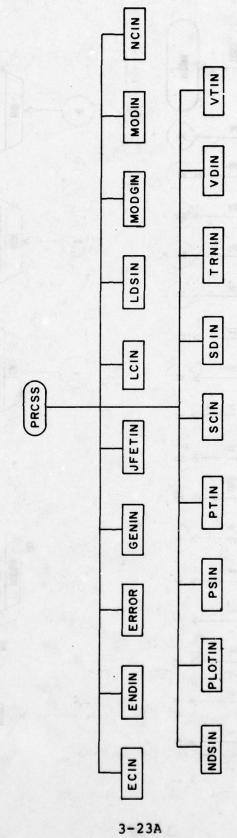
TYPE=6 Node Card

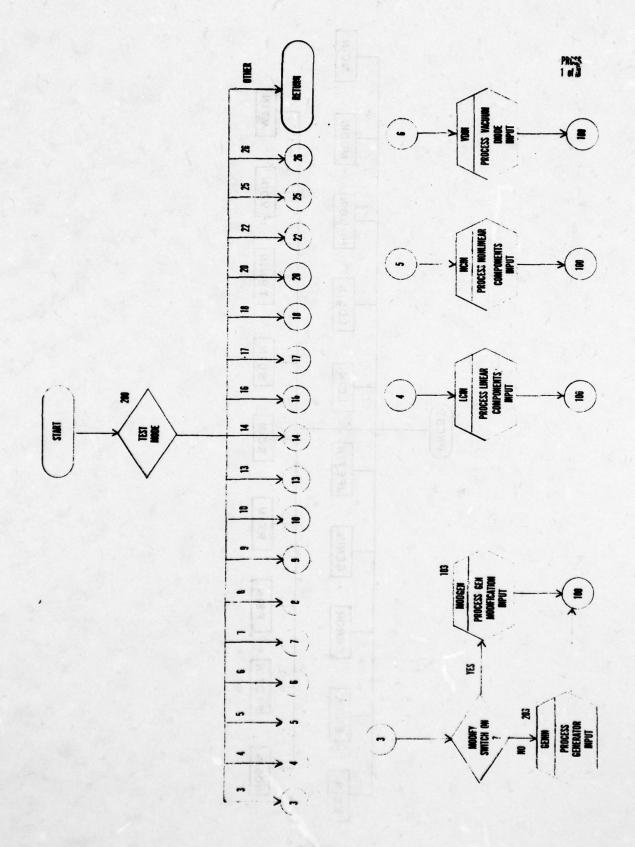
The NODE identifier is decoded by subroutine ALPHA and if an end-of-data is encountered (ISW=1), subroutine PLOTIN exits with an appropriate error condition. Otherwise the node number is extracted from the card image, formatted, and stored in MISC(4) by subroutine INTEG. Subroutine PLOTIN returns to the calling program.

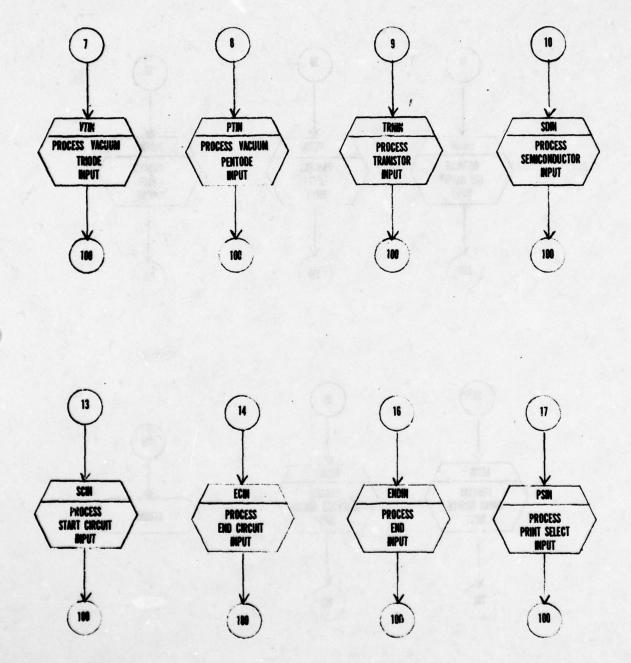
TYPE=15 Label Card

The 50-character label specification is extracted from the card image, stored in the common data buffer PUFF(1)-(50) and written to file 20 by subroutine DATAWR. The data file record number DATREC is updated and the data record length is incremented. After clearing the data buffer, subroutine PLOTIN returns to the calling program.

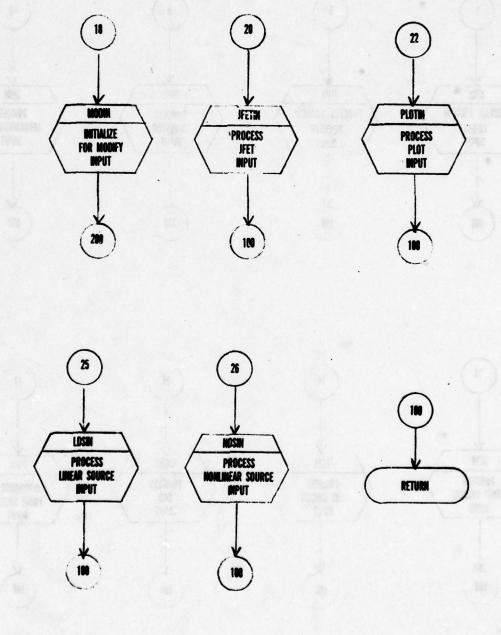
Any other value of TYPE constitutes an input error which is handled by subroutine ERROR before exit to the calling program.







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NONLINEAR CIRCUIT ANALYSIS PROGRAM (NCAP) DOCUMENTATION. VOLUME--ETC(U)
SEP 79 J B VALENTE , S STRATAKOS F30602-79-C-0011 AD-A076 317 UNCLASSIFIED RADC-TR-79-245-VOL-3 NL 3 OF 8 ADA 076317 1



NAME: PRCSS

TYPE: SUBROUTINE

GENERAL PURPOSE:

Controls input processing of *-card sequences according to the value of MODE

Modes 12 2, 15 12 12, 15, 19 21 23 a

nemer lis or bost lon to me about

VARIABLES:

NONE

SUBROUTINES CALLED.

ECIN	ENDIN	ERPOR	GENIN	JFETIN
LCIN	LDSIN	MODGIN	MODIN	NCIN
NDSIN	PLOTIN	PSIN	PTIN.	SCIN
SDIN	TRNIN	VDIN	VTIN	

CALLING PROGRAMS:

MAIN

DESCRIPTION:

Subroutine PRCSS controls the processing of NCAP input card sequences. It uses the MODE of the previous *-card transmitted from the main program through common storage as the index of a computed GO TO which transfers control to an appropriate input processing subroutine. There the input card is processed

according to an anticipated format and translated into appropriate data and driver records.

Modes 1, 2, 11, 12, 15, 19, 21, 23 and 24 imply error conditions since they are either created internally by the program or not used at all. When any of these modes are encountered by subroutine PRCSS, an error message is printed and control returns to the main program. The valid MODE values and their corresponding input processors are tabulated below:

MODE	Input Processing	Comments
	Routine	
		GNATAO SAMETURES SUS
3	GENIN/MODGIN	Generator/Generator Modify
4	LCIN	Linear Components
5	NCIN	Nonlinear Components
6	VDIN	Vacuum Diode
7	VTIN	Vacuum Triode
8	PTIN	Vacuum Pentode
9	TRNIN	Transistor
10	SDIN	Semiconductor Diode
13	SCIN	Start Circuit
14	ECIN	End Circuit
16	ENDIN	End
17	PSIN	Print Select
18	MODIN	Modify
20	JFETIN	Junction Field Effect
		Transistor

22	PLOTIN	Plot
25	LDSIN	Linear Dependent Source
26	NDSIN	Nonlinear Dependent Source

With the exception of MODE=3, 16, and 18 program control passes directly from the computed GO TO to the appropriate input processor and from there back to the main program.

* MODIFY card sencences result in a slightly different path through subroutine PRCSS. For MODE=18, program control passes from the computed GO TO to subroutine MODIN. There, rather than processing the cards of the modify sequence directly, the value of MODE is reset to identify the device being modified and the modify switch is turned on (MODFY=1). Upon return from MODIN control passes back to the computed GO TO of subroutine PRCSS where the updated value of MODE causes the cards of the modify sequence to be processed by the input subroutine corresponding to the device being modified.

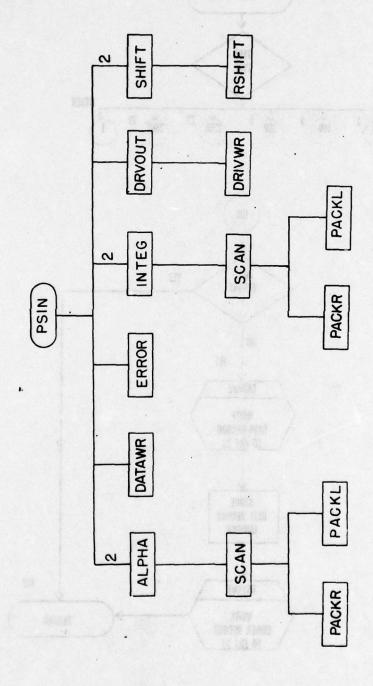
Because of the complexity of generator input specifications and the restrictions on generator modification, there is a special input processor MODGIN for generator modification cards. It is invoked by subroutine PRCSS when MODE=3 and MODFY=1, indicating that generator modification is in effect. Otherwise the standard generator input routine CENIN is called to process GENERATOR card sequences.

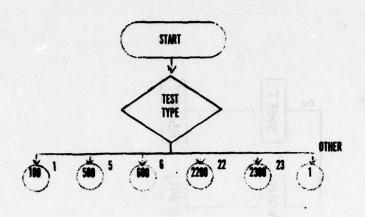
An * END card (MODE=16) signals the end of Phase 0 processing. For MODE=16, program control passes from the computed GO TO in PRCSS to subroutine ENDIN. There the input

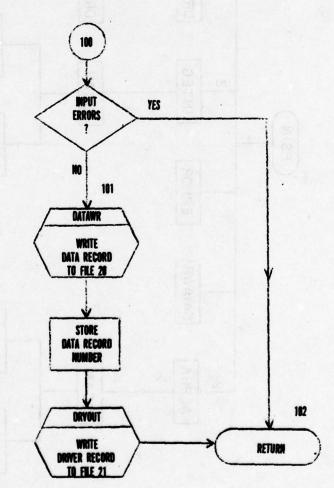
phase is terminated and the program either proceeds to subroutine PHASEl or aborts depending on the presence of errors in the input deck.

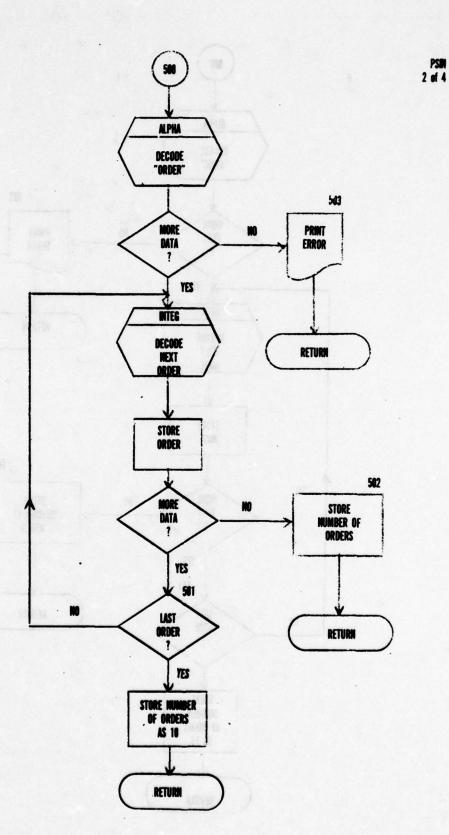
processor and from there back to the roam program.

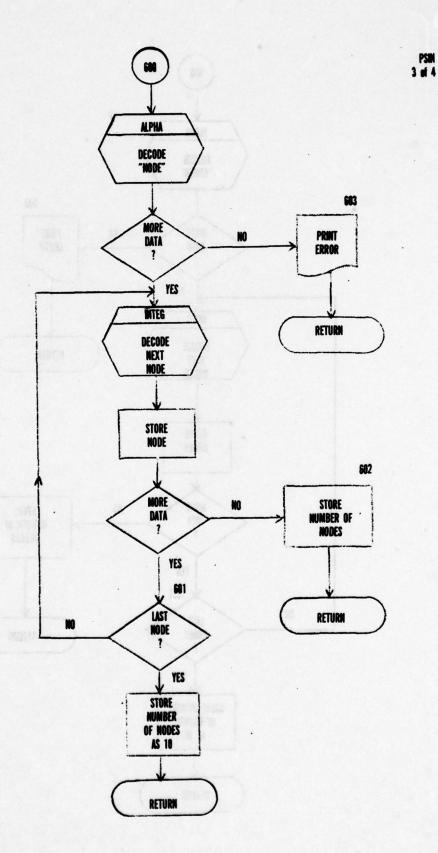
winds and we dissolve a 2009 of or or or or

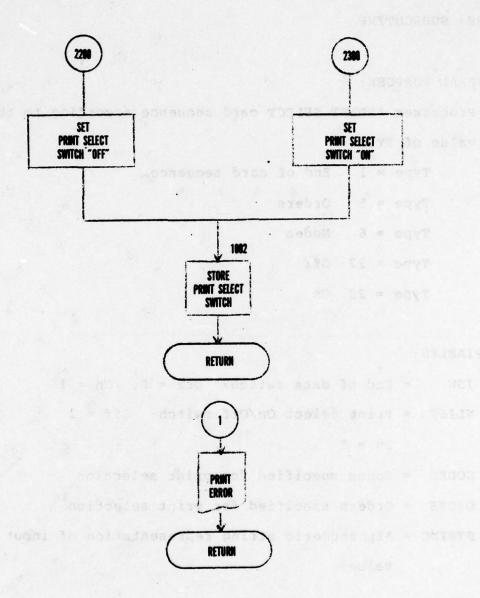












NAME: PSIN

TYPE: SUBROUTINE

GENERAL PURPOSE:

Processes *PRINT SELECT card sequence according to the value of TYPE:

Type = 1 End of card secuence

Type = 5 Orders

Type = 6. Nodes

Type = 22 Off

Type = 23 On

VAPIABLES:

ISW = End of data switch: Off = 0, On = 1

NLIST = Print select On/Off switch · Off = 1

On = 0

NODES = Nodes specified for print selection

ORDRS = Orders specified for print selection

STRING = Alphanumeric string representation of input value

SUPROUTINES CALLED:

ALPHA DATAWR DRVOUT ERROR INTEG

CALLING PROGRAMS:

PRCSS

DESCRIPTION:

Subroutine PSIN is the input processor for PRINT SELECT card sequences. Its function is to decode the print select input data according to an anticipated format, translating the input into appropriate data and driver records for use by subsequent NCAP phases.

Fach call to subroutine PSIN causes a single NCAP input card to be processed. The card image and its TYPE code are transmitted through common storage to subroutine PSIN which processes the card according to its TYPE as follows:

TYPE=1 End of Card Sequence

If the input error switch is on (JAPT=1), subroutine PSIN exits with an appropriate error condition. Otherwise the print select data record is written to file 20 from the common data buffer PUFF, and the data record number DATREC is updated to point beyond the new data record.

Since the print select feature does not cause any internally generated data to be transmitted between NCAP phases, the 20-word data record contains only input values:

PUFF(1)-(10) Orders to be Printed

PUFF(11)-(20) Nodes to be Printed

After the data record has been written to disk, subroutine DRVOUT is called to complete the driver record and write it to file 21. The print select driver record, taken from the first ten words of global common, consists of:

MODE=17 Print Select Identifier

STADD Data File Record Number

LNGTH=20 Length of Data Record in Words

MISC(1) Number of Orders Specified

MISC(2) Number of Nodes Specified

MISC(3) = \emptyset or 1 ON/OFF switch: \emptyset =on, 1=off

Other MISC driver parameters are not used.

After returning from subroutine DRVOUT, subroutine PSIN returns to the calling program.

TYPE=5 Order Card

The ORDER identifier is decoded by subroutine ALPHA, and if an end-of-data is encountered (ISW=1), subroutine PSIN exits with an appropriate error condition. Otherwise from one to ten integers specifying the orders of analysis to be printed during Phase 4 are decoded and formatted by subroutine INTEG and stored in ORDRS(1)-(10) under the control of the index KZ. After decoding, each ORDRS(KZ) is placed in the print select data record at BUFF(KZ). If an end-of-data is encountered before ORDPS(10) is filled, the number of orders is stored in MISC(1). Otherwise MISC(1)=10 and subroutine PSIN returns to the calling

program.

TYPE=6 Node Card

The NODE identifier is decoded by subroutine ALPHA. Processing of the print select node data is identical to that of the order data except that the input values are decoded to NOPES(1)-(10) and stored in the print select data record beginning at PUFF(KZ+10). The node count is saved in MISC(2).

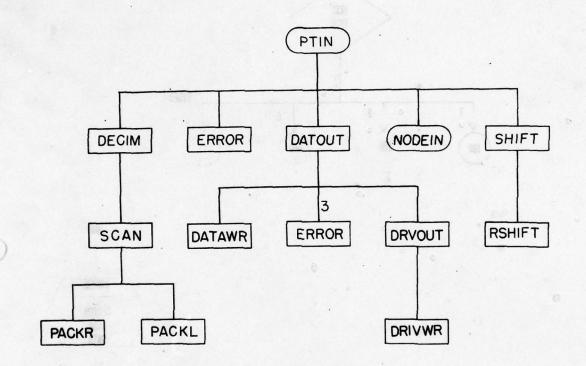
TYPE=22 Off Card

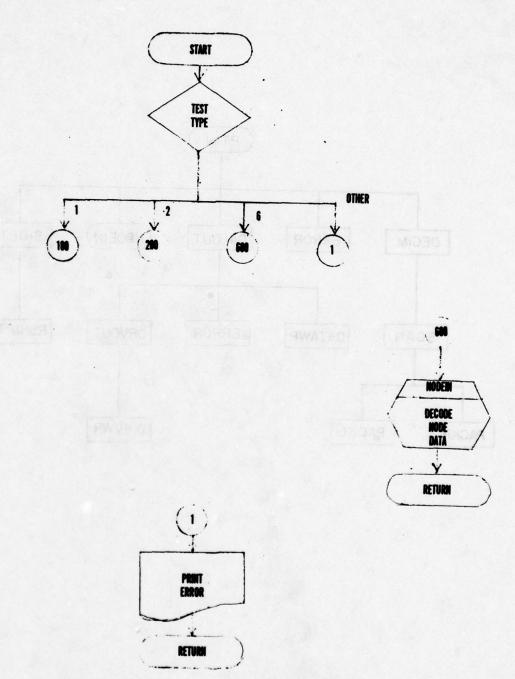
The print select switch is set to OFF by NLIST=1, saved in MISC(3), and subroutine PSIN returns to the calling program.

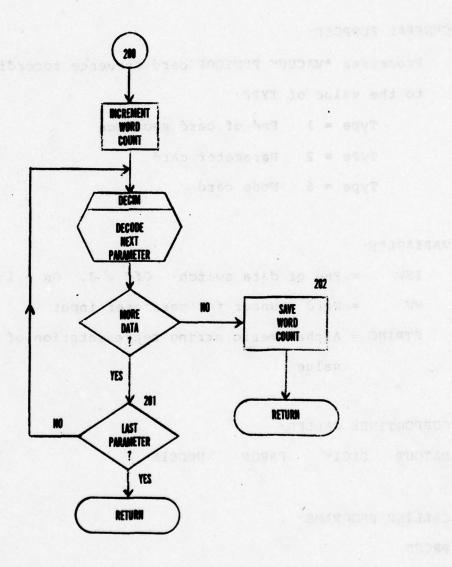
TYPE=23 On Card

The print select switch is set to ON by NLIST=0, saved in MISC(3), and subroutine PSIN returns to the calling program.

Any other value of TYPE constitutes an input error which is handled by subroutine FPFCF before exit to the calling program.







care merconces. The function is to decode the vacuum i entode

MAME: PTIN

TYPF · SUPPOUTINE

CENERAL PURPOSE:

Processes *VACUUM PENTODE card sequence according to the value of TYPE:

· Type = 1 End of card sequence

TyPe = 2 Parameter card

Type = 6 Node card

VARIABLES:

ISW = End of data switch: Off = 0. On = 1

NW = Word counter for parameter input

STRING = Alphanumeric string representation of input

value

SUBROUTINES CALLED.

DATOUT DECIM ERROR NODEIN

CALLING PROGRAMS.

PRCSS

DESCRIPTION :

Subroutine PTIN is the input processor for VACUUM PENTODE card sequences. Its function is to decode the vacuum pentode

input data according to an anticipated format, translating the input into appropriate data and driver records for use by subsequent NCAP phases.

Each call to subroutine PTIN causes a single NCAP input card to be processed. The card image and its TYPE code are transmitted through common storage to subroutine PTIN which processes the card according to its TYPE as follows:

terd image, formatted, and stored in the data butler by

TYPE=1 End of Card Sequence

Subroutine DATCUT is called to complete the driver record and to close the vacuum pentode definition. The input is tested for errors and the data and driver records are written to disk. The driver record is written to file 21 from the first ten words of global common, where:

MODE=8 Vacuum Pentode Identifier

STADD Data File Record Number

LNCTH=107 Length of Data Record in Words

MISC(3) Base Node of Connection

Other MISC driver parameters are not used.

The vacuum pentode data record, written to file 20 from the common data buffer BUFF, is 107 words in length. This record allocates sufficient disk storage for the vacuum pentode input values as well as all other data generated internally by its model in subsequent NCAP phases.

PUFF(1)-(12) Input Values

BUFF(13)-(75)

Nonlinear Parameters and Coefficients BUFF(76)-(107) Admittance Submatrix

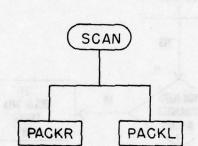
TYPE=2 Parameter Card

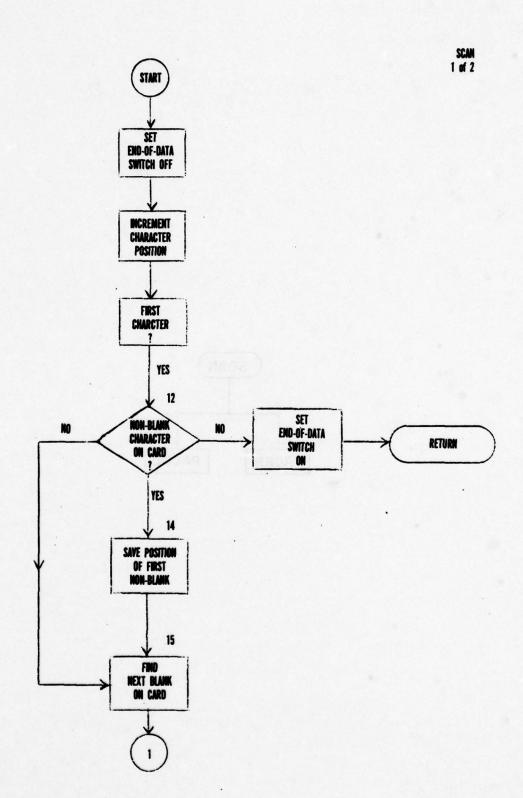
From one to twelve decimal parameters are extracted from the card image, formatted, and stored in the data buffer by subroutine DECIM. The number of values decoded and their placement in the buffer depend on the parameter count, NXT, and the end-of-data switch, ISW. Upon entering the subroutine, NXT contains the number of values input on previous data cards. Parameter storage begins at BUFF (NW) where NW=NXT+1 and proceeds through successive buffer locations as data values are decoded by subroutine DECIM. If an end-of-data is encountered (ISW=1) before BUFF(12) is filled, the parameter count is reset to the last buffer location used (NXT=NW) and control returns to the calling program in anticipation of additional parameter cards.

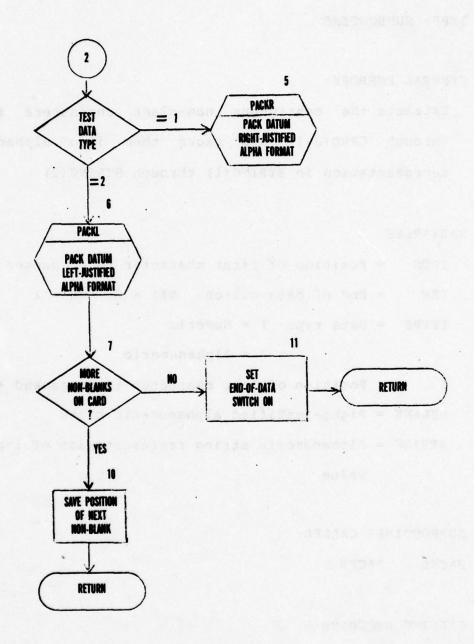
TYPE=6 Node Card

The base node of connection is extracted from the card image, formatted, and stored in MISC(3) of the driver record by subroutine NODEIN.

Any other value of TYPE constitutes an input error which is handled by subroutine ERROR before exit to the calling program.







NAME: SCAN

TYPE: SUPROUTINE

CENERAL PURPOSE:

Extracts the contiguous non-blank characters in CARD(IPOS) through CARD(J-1) and packs them into alphanumeric string representation in STRING(1) through STRING(2)

VARIABLES:

IPOS = Position of first character to be packed

ISW = End of data switch: Off = \emptyset , On = 1

ITYPE = Data type · 1 = Numeric

2 = Alphanumeric

J = Position of last character to be packed + 1

PPLANK = Pight-justified alphanumeric blank

STRINC = Alphanumeric string representation of input value

SUPROUTINFS CALLED:

PACKL PACKR

CALLING PROCRAMS:

ALPHA DECIM INTEC

DESCRIPTION .

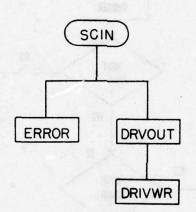
Subroutine SCAN belongs to a group of Phase Ø character manipulation routines which format NCAP input data. Its function is to extract the contiguous alphanumeric characters constituting a data value from an input card image and to pack those characters into the integer array STRING for formatting by subroutines ALPHA, DECIM, and INTEC.

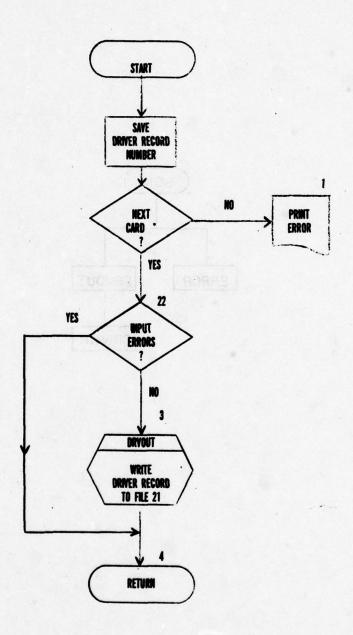
Input to the subroutine is the right-justified alphanumeric card image stored in the integer array card, the character position IPOS at which the scan begins, and the anticipated data type represented by ITYPE=1 for numeric data or ITYPE=2 for alphabetic data. Output from the routine is the alphanumeric string representation STRING, the character position IPOS at which the scan terminated, and the end-of-data switch ISW.

First the end-of-data switch is initialized off by ISW=Ø. Then starting at CARD(I) the image is scanned column-by-column to locate the first non-blank character. If no non-blank is encountered, the end-of-data switch is turned "on" and exit occurs. Otherwise the position of the first non-blank is saved in IPOS and the scan proceeds across the non-blank characters until a blank is encountered at CARD(J).

At this point the datum in CARD(IPOS) through CARD(J-1) is packed into STRING either right or left-justified according to the anticipated data type. Numeric data (ITYPE=1) are packed right-justified by subroutine PACKR, while alphabetic data (ITYPE=2) are packed left-justified by subroutine PACKL.

After the packing is complete, the scan resumes at CARD(J) to determine if there are additional data on the card and to position the character pointer accordingly. In the event that the remainder of the card is blank, the end-of-data switch is turned on (ISW=0) and exit occurs. Otherwise the character pointer IPOS is set to the starting location of the next datum and control returns to the calling program.





NAME: SCIN OF STREET STREET STREET STORY OF STREET STREET STREET STREET STREET STREET

TYPE: SUBROUTINE

GENERAL PURPOSE:

Processes *START CIRCUIT card

VARIABLES: 1 5700001 Sand the TIME TO TRATE SAN DELWISHED . Role 5700

NONE TO LEGGE to Street test series of more in elli of

SUBROUTINES CALLED:

DRVOUT ERROR

CALLING PROGRAMS:

PRCSS adoing without disconnection of bonu ins benillab are dud in assign

DESCRIPTION:

Subroutine SCIN is the input processor for the START CIRCUIT card. Its function is to translate the start circuit input card into an appropriate driver record for use by subsequent NCAP phases. The start circuit driver record and its companion end circuit driver record serve as delimiters for the NCAP driver file. The address of the start circuit driver record SCREC is saved by subroutine SCIN and transmitted through common storage to the end circuit input processor ECIN.

Processing of the start circuit function begins by testing

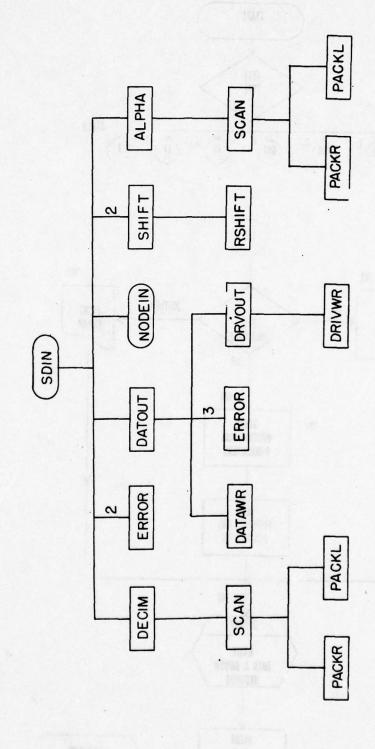
the TYPE code of the input card. Pecause there are no data cards associated with the start circuit input, TYPE=1 is the only code recognized by subroutine SCIN. Any other value of TYPE constitutes an input error which is handled by subroutine ERROR before exit to the calling program.

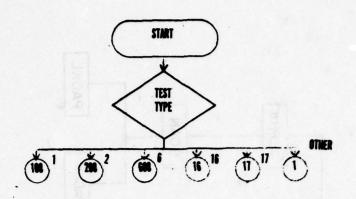
If the input error switch is on (JABT=1), subroutine SCIN returns to the calling program with an appropriate error condition. Otherwise the START CIRCUIT driver record is written to file 21 from the first ten words of global common by subroutine DRVOUT.

The start circuit driver record is identified by MODE=13. Because there is no data associated with a start circuit function, the driver parameters STADD and LNGTH are not applicable. The MISC driver parameters are not defined during Phase Ø, but are defined and used in subsequent NCAP phases to store the disk addresses of admittance matrices and transfer functions.

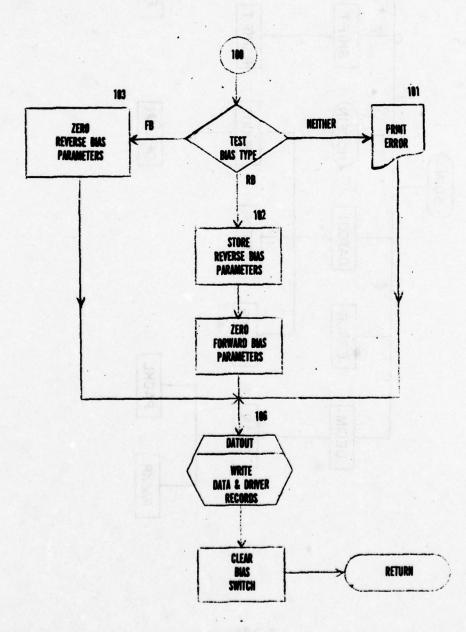
After the driver record has been written to disk, subroutine SCIN returns to the calling program.

by consequence will and transmitted through corner storage

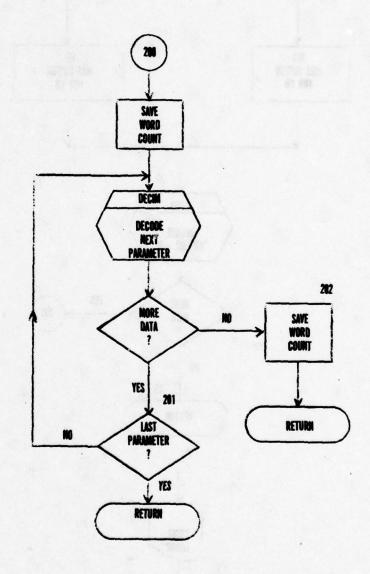


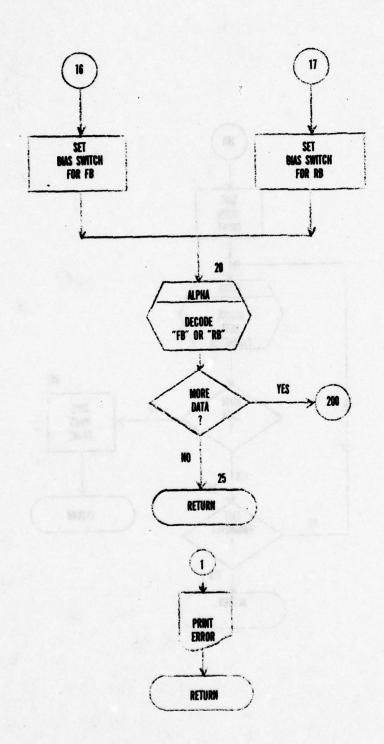






3-28B





NAME: SDIN

TYPE: SUBROUTINE

GENERAL PURPOSE:

Processes *SEMICONDUCTOR DIODE card sequence according to the value of TYPE:

Type = 1 End of card sequence

Type = 2 Parameter card

Type = 6 Node card

Type = 16 Forward Bias

Type = 17 Reverse Bias

VARIABLES:

IBSW = Bias switch; 1 = Forward, 2 = Reverse

ISW = End of data switch · Off = 0, On = 1

NW = Word counter for parameter input

STRING = Alphanumeric string representation of input

value

SUBROUTINES CALLED:

ALPHA DATOUT DECIM ERROR NODEIN

CALLING PROGRAMS:

PRCSS

DESCRIPTION .

Subroutine SDIN is the input processor for SEMICONDUCTOR DIODE card sequences. Its function is to decode the semiconductor input data according to an anticipated format, translating the input into appropriate data and driver records for use by subsequent NCAP phases.

Each call to subroutine SDIN causes a single NCAP input card to be processed. The card image and its TYPE code are transmitted through common storage to subroutine SDIN which processes the card according to its TYPE as follows:

TYPE=1 End of Card Sequence

If no forward or reverse-bias specification was processed (IBSW=0), subroutine SDIN exits with an appropriate error condition. If the diode is reverse-biased (IBSW=2), the four input values are moved from BUFF(1)-(4), which are reserved for forward-bias data, to BUFF(5)-(8) and the forward-bias data area is cleared. If the diode is forward-biased, the reverse-bias data area BUFF(5)-(8) is cleared.

Subroutine DATOUT is called to complete the driver record and to close the semiconductor diode definition. The driver record is written to file 21 from the first ten words of global common, where:

MODE=10

Semiconductor Diode Identifier

STADD

Data File Record Number

LNGTH=32

Length of Data File in Words

MISC(3)

Positive Node of Connection

MISC(4)

Negative Node of Connection

The other MISC parameters are not used.

The semiconductor diode data record, written to file 20 from the common data buffer BUFF, is 32 words in length. This record allocates sufficient disk storage for the semiconductor diode input values as well as all other data generated internally by its model in subsequent NCAP phases:

BUFF(1)-(8)

Input Values

PUFF(9)-(30)

Monlinear Parameters and Coefficients

PUFF (31) - (32)

Admittance Submatrix

After return from subroutine DATOUT, the bias switch IBSW is cleared and subroutine SDIN returns to the calling program.

TYPE=2 Parameter Card

From one to four decimal parameters are extracted from the card image, formatted, and stored in the data buffer by subroutine DECIM. The number of values decoded and their placement in the buffer depend on the parameter count, NXT, and the end-of-data switch, ISW. Upon entering the subroutine, NXT contains the number of values input on previous data cards. Parameter storage begins at PUFF(NW) where NW=NXT+1 and proceeds through successive buffer locations as data values are decoded by subroutine DECIM. If an end-of-data is encountered(ISW=1) before

BUFF(4) is filled, the parameter count is reset to the last buffer location used (NXT=NW) and control returns to the calling program in anticipation of additional parameter cards.

TYPE=6 Node Card

The nodes of connection are extracted from the card image, formatted, and stored in MISC(3) and MISC(4) by subroutine NODEIN.

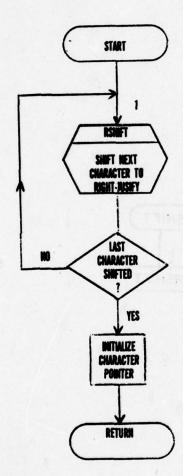
TYPE=16 or 17 FB or RP Card

The bias switch is set: l=forward-bias, 2=reverse-bias. The FP or RP identifier is decoded by subroutine ALPHA. If an end-of-data is encountered (ISW=1), subroutine SDIN returns to the calling program. Otherwise control transfers to TYPE=2 processing for the remaining data on the card.

Any other value of TYPE constitutes an input error which is handled by subroutine ERROR before exit to the calling program.



12



TYPE. SUBROUTINE

GENERAL PURPOSE:

Right-justifies the card image stored in CARD(1) through CARD(80)

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VARIABLES:

CARD = Card image, one character per word, left-justified

I = Index for CARD array

SUBROUTINES CALLED:

RSHIFT

CALLING PROGRAMS:

MAIN

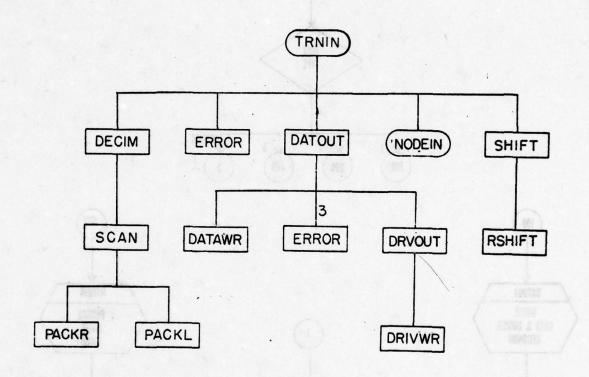
DESCRIPTION:

Subroutine SHIFT belongs to a group of Phase Ø character manipulation routines which format NCAP input data. Input cards are read by NCAP's main program into the integer array CARD under FORMAT (80A1) which places one alphanumeric character in the most significant bits each word of the card image (left-justification). The binary equivalents of such characters are frequently negative numbers or very large positive numbers

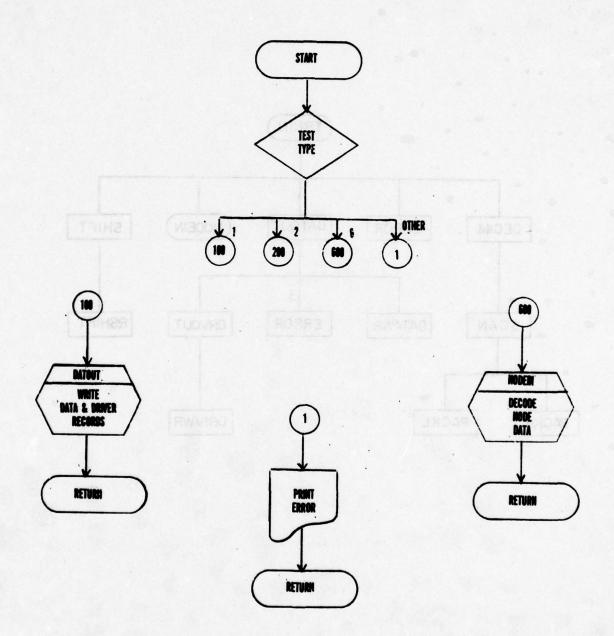
which are difficult to manipulate arithmetically because of overflow, propagation of the sign bit, and end-around carry.

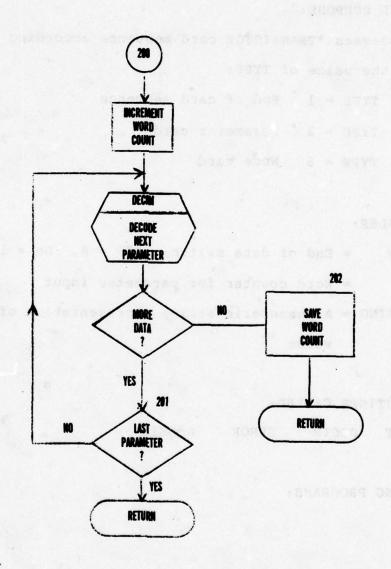
By shifting each character of the card image from the most significant to the least significant bits and zero-filling to the left, every character can be represented as a small positive integer which can be easily manipulated by arithmetic means.

Subroutine SHIFT performs this conversion on the input card image in preparation for character manipulation by the formatting routines of Phase 0. The 80 left-justified card characters are transmitted to the routine in the integer array CARD, right-justified by subroutine RSHIFT, and returned to the calling program in CARD.



0





NAME: TRNIN

TYPE: SUBROUTINE

GENERAL PURPOSE:

Processes *TRANSISTOR card sequence according to the value of TYPE:

TYPE = 1 End of card sequence

TYPE = 2 Parameter card

TYPE = 6 Node card

VARIABLES:

ISW = End of data switch; Off = \emptyset , On = 1

NW = Word counter for parameter input

STRING = Alphanumeric string representation of input

value

SUBROUTINES CALLED:

DATOUT DECIM ERROR NODEIN

CALLING PROGRAMS:

PRCSS

DESCRIPTION:

Subroutine TRNIN is the input processor for TRANSISTOR card sequences. Its function is to decode the bipolar junction

transistor input data according to an anticipated format, translating the input into appropriate data and driver records for use by subsequent NCAP phases.

Each call to subroutine TRNIN causes a single input card to be processed. The card image and its TYPE code are transmitted through common storage to subroutine TRNIN which processes the card according to its TYPE as follows:

subsculing balls. The rader of values decoded and castr

TYPE=1 End of Card Sequence

Subroutine DATOUT is called to complete the driver record and close the transistor definition. The input is tested for errors and the data and driver records are written to disk. The driver record is written to file 21 from the first ten words of global common, where:

MODE=9 Transistor Identifier

STADD Data File Record Number

LNGTH=102 Length of Data Record in Words

MISC(3) Pase Node of Connection

The other MISC driver parameters are not used.

The transistor data record, written to file 20 from the common data buffer BUFF, is 102 words in length. This record allocated sufficient disk storage for the transistor input values as well as all other data generated internally by its model for use by subsequent NCAP phases:

BUFF(1)-(20) Input Values

BUFF(21)-(70) Nonlinear Parameters and Coefficients

through Cobren propose to some potent

BUFF(71)-(102) Admittance Submatrix

TYPE=2 Parameter Card

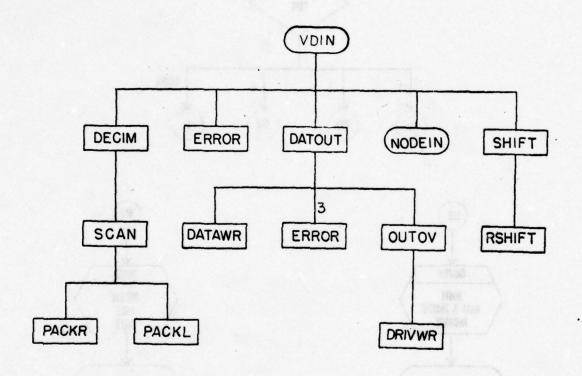
From one to sixteen decimal parameters are extracted from the card image, formatted, and stored in the data buffer by subroutine DECIM. The number of values decoded and their placement in the buffer depend on the parameter count, NXT, and the end-of-data switch, ISW. Upon entering the subroutine, NXT contains the number of values input on previous data cards. Parameter storage begins at BUFF(NW) where NW=NXT+1 and proceeds through successive buffer locations as data values are decoded by subroutine DECIM. If an end-of-data is encountered (ISW=1) before BUFF(16) is filled, the parameter count is reset to the last buffer location used (NXT=NW) and control returns to the calling program in anticipation on additional parameter cards.

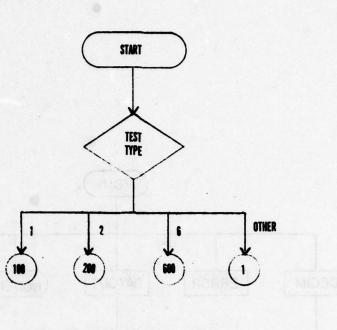
Each call to subtoution IRNIN causes a simila limit dardito

TYPE=6 Node Card

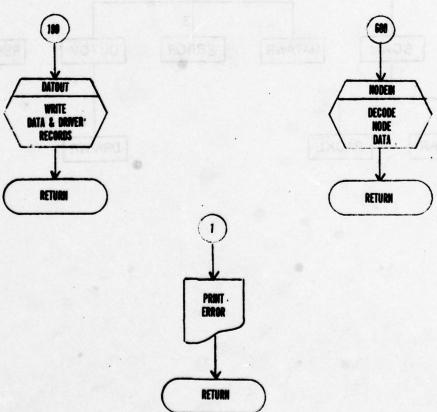
The base node of connection is extracted from the card image, formatted, and stored in MISC(3) by subroutine NODEIN.

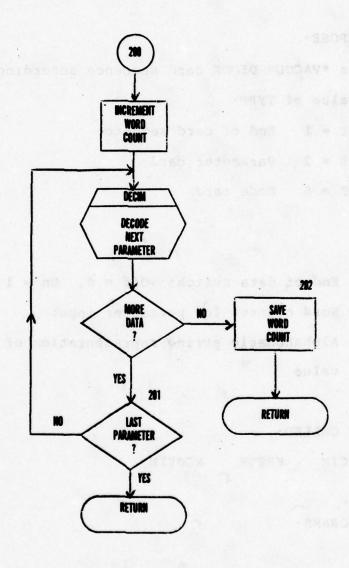
Any other value of TYPE constitutes an input error which is handled by subroutine EPROR before exit to the calling program.





VDM 1 of 2





0

NAME: VDIN

TYPE: SUBROUTINE

GENERAL PURPOSE:

Processes *VACUUM DIODE card sequence according to the value of TYPE.

TYPE = 1 End of card sequence

TYPE = 2 Parameter card

TYPE = 6 Node card

VARIABLES:

ISW = End of data switch; Off = 0, On = 1

NW = Word counter for parameter input

STRING = Alphanumeric string representation of input value

SUPROUTINES CALLED:

DATOUT DECIM ERROR NODEIN

CALLING PROGRAMS:

PPCSS

DESCRIPTION :

Subroutine VDIN is the input processor for VACUUM DIODE card secuences. Its function is to decode the vacuum diode input data

according to an anticipated format, translating the input into appropriate data and driver records for use by subsequent NCAP phases.

Each call to subroutine VDIN causes a single NCAP input card to be processed. The card image and its TYPE code are transmitted through common storage to subroutine VDIN which processes the card according to its TYPE as follows:

foresterning wie stored the the fats buffereiby

TYPE=1 End of Card Sequence

Subroutine DATOUT is called to complete the driver record and to close the vacuum Diode definition. The input is tested for errors and the data and driver records are written to disk. The driver record is written to file 21 from the first ten words of global common, where:

MODE=6	Vacuum Diode Identifier
STADD	Data File Record Number
LNGTH=15	Length of Data Record in Words
MISC(3)	Positive Node of Connection
MISC(4)	Negative Node of Connection

Other MISC driver parameters are not used.

The vacuum diode data record, written to file 20 from the common data buffer PUFF, is 15 words in length. This record allocates sufficient disk storage for the vacuum diode input values as well as all other data generated internally by its model in subsequent NCAP phases:

BUFF(1)-(3) Input Values

Sect eall to subresting verm casees a single work imple ears

BUFF(4)-(13) Nonlinear Coefficients

PUFF(14)-(15)

Admittance Submatrix

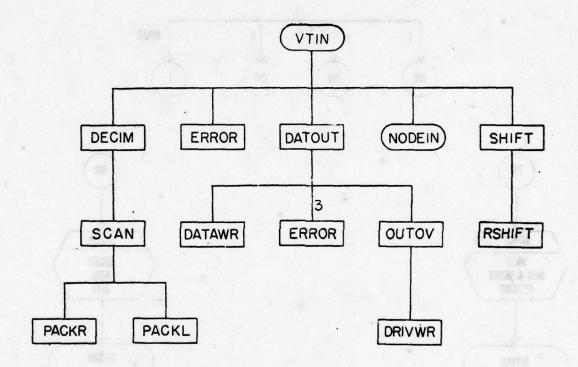
TYPE=2 Parameter Card

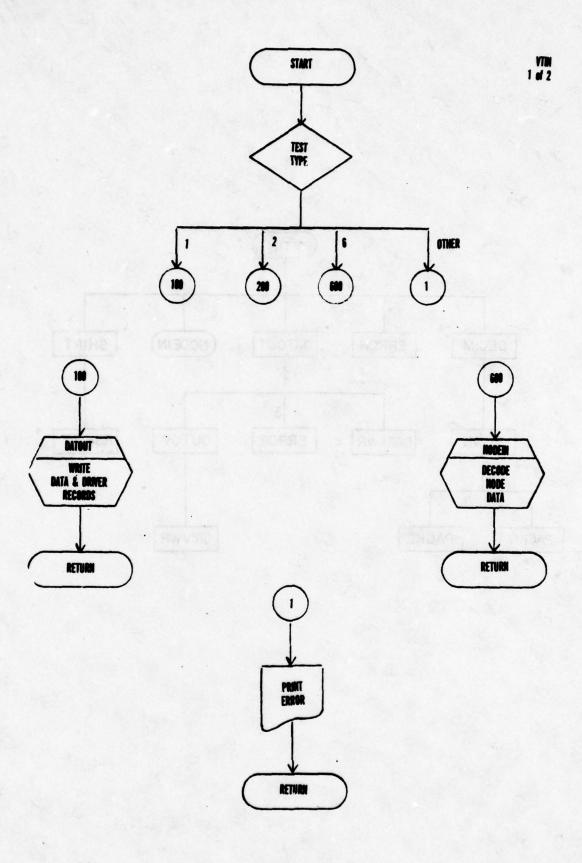
From one to three decimal parameters are extracted from the card image, formatted, and stored in the data buffer by subroutine DECIM. The number of values decoded and their placement in the buffer depend on the parameter count, NXT, and the end-of-data switch, ISW. Upon entering the subroutine, NXT contains the number of values input on previous data cards. Parameter storage begins at BUFF (NW) where NW=NXT+1 and proceeds through successive buffer locations as data values are decoded by subroutine DECIM. If an end-of-data is encountered(ISW=1) before PUFF(3) is filled, the parameter count is reset to the last buffer location used (NXT=NW) and control returns to the calling program in anticipation of additional parameter cards.

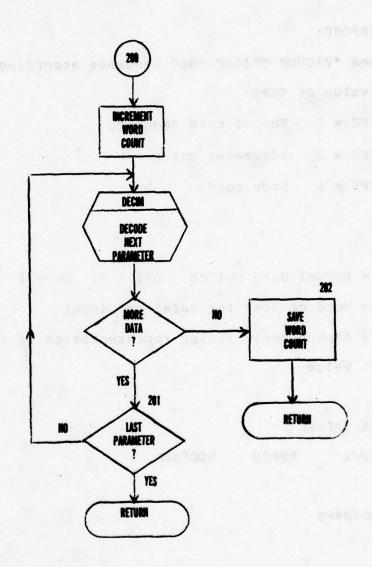
TYPE=6 Node Card

The nodes of connection are extracted from the card image, formatted, and stored in MISC(3) and MISC(4) by subroutine NODEIN.

Any other value of TYPE constitutes an input error which is handled by subroutine ERROR before exit to the calling program.







NAME: VTIN

TYPE: SUPROUTINE

CENERAL PURPOSE:

Processes *VACUUM TRIODE card sequence according to the value of TYPE:

. TYPE = 1 End of card sequence

TYPE = 2 Parameter card

TYPE = 6 Node card

VARIABLES:

ISW = End of data switch; Off = 0, On = 1

NW = Word counter for parameter input

STRING = Alphanumeric string representation of input value

SUBROUTINES CALLED:

PATOUT DECIM ERROR NODEIN

CALLING PROGRAMS .

PRCSS

DESCRIPTION .

Subroutine VTIN is the input processor for VACUUM TRIODE card sequences. Its function is to decode the vacuum triode

input data according to an anticipated format, translating the input into appropriate data and driver records for use by subsequent NCAP phases.

Each call to subroutine VTIN causes a single NCAP input card to be processed. The card image and its TYPE code are transmitted through common storage to subroutine VTIN which processes the card according to its TYPE as follows:

subrouting DECIF. The number of Values decoded and their

TYPE=1 End of Card Sequence

Subroutine DATOUT is called to complete the driver record and to close the Vacuum Triode definition. The input is tested for errors and the data and driver records are written to disk.

The driver record is written to file 21 from the first ten words of global common, where

MODE=7 Vacuum Triode Identifier

STADD Data File Record Number

LNGTH=40 Length of Data Record in Words

MISC(3) Base Node of Connection

Other MISC driver parameters are not used.

The vacuum triode data record, written to file 20 from the common data buffer, PUFF, is 40 words in length. This record allocates sufficient disk storage for the vacuum triode input values as well as all other data generated internally by its model in subsequent NCAP phases:

PUFF(1)-(9) Input Values

BUFF(10)-(22)

Nonlinear Parameters and Coefficients

BUFF(23)-(40)

Admittance Submatrix

TYPE=2 Parameter Card

From one to nine decimal parameters are extracted from the card image, formatted, and stored in the data buffer by subroutine DECIM. The number of values decoded and their placement in the buffer depend on the parameter count, NXT, and the end-of-data switch, ISW. Upon entering the subroutine, NXT contains the number of values input on previous data cards. Parameter storage begins at BUFF(NW) where NW=NXT+1 and proceeds through successive buffer locations as data values are decoded by subroutine DECIM. If an end-of-data is encountered (ISW=1) before BUFF(9) is filled, the parameter count is reset to the last buffer location used (NXT=NW) and control returns to the calling program in anticipation of additional parameter cards.

TYPE=6 Node Card

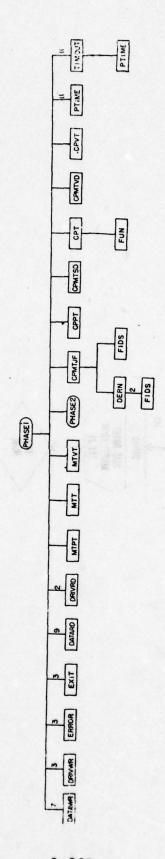
The base node of connection is extracted from the card image, formatted, and stored in MISC(3) by subroutine NODEIN.

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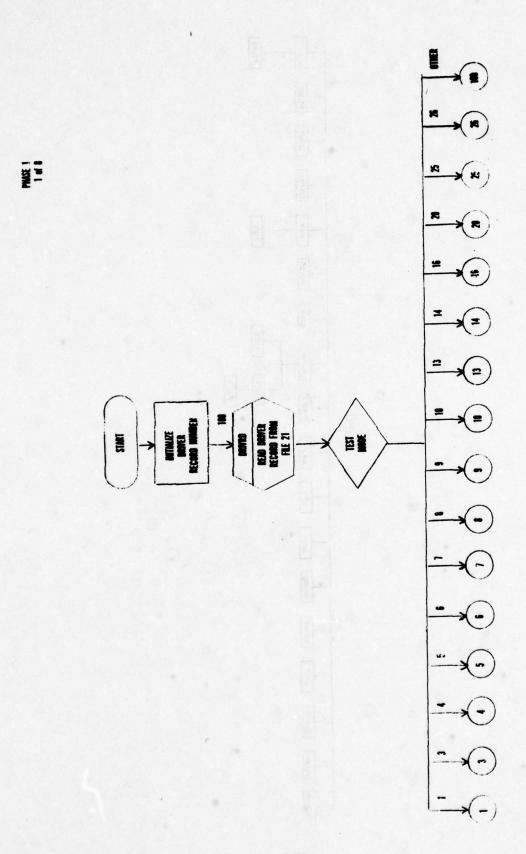
Any other value of TYPE constitutes an input error which is handled by subroutine ERROR before exit to the calling program.

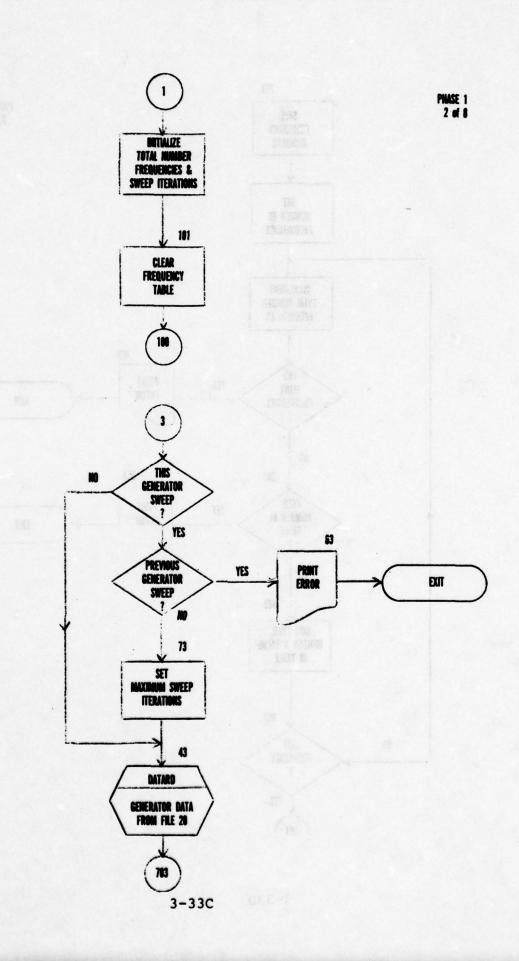
PHASE1

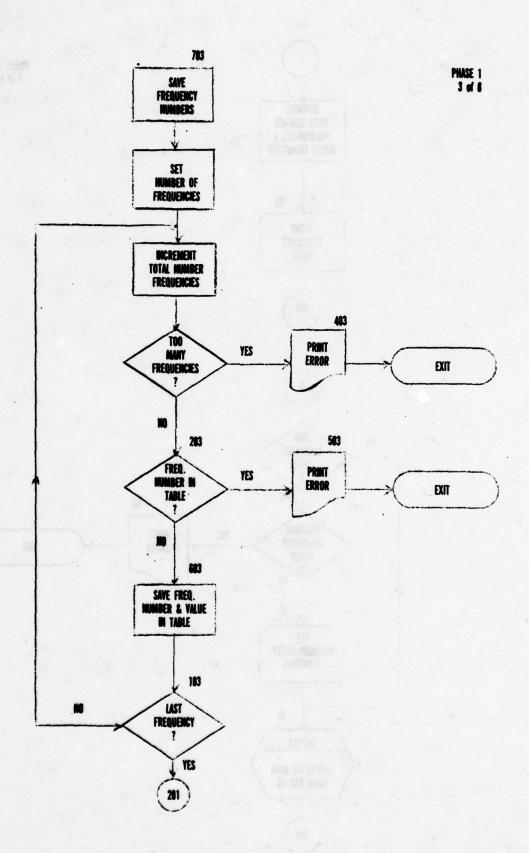
CPMTJF CPMTSD CPMTVD CPPT CPT CPVT DERIDS FIDS IMLOC MTPT MTT MTVT



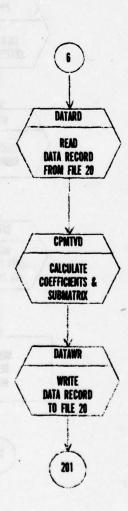
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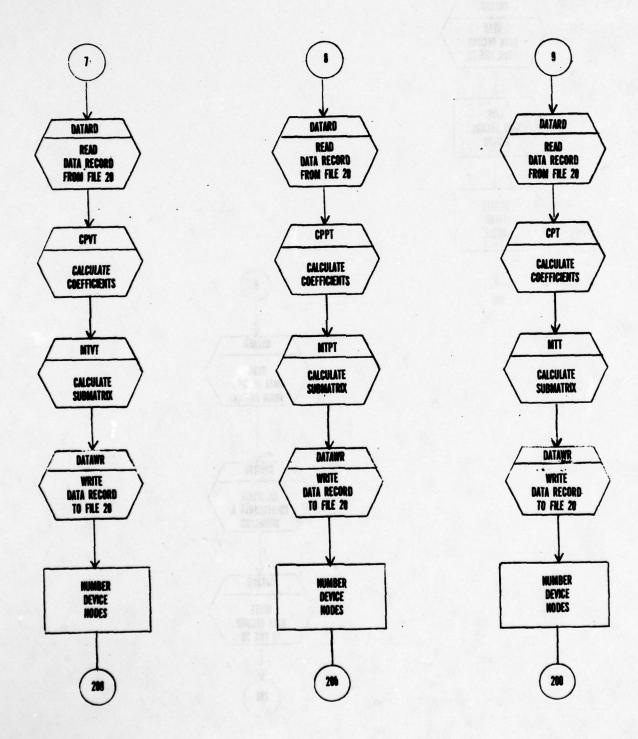




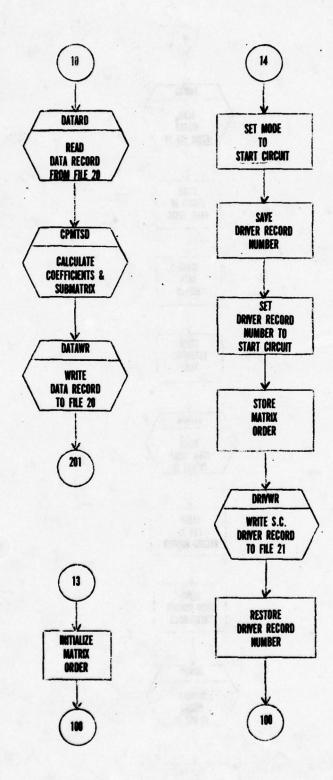


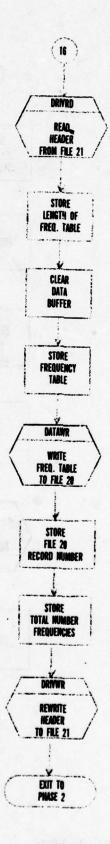


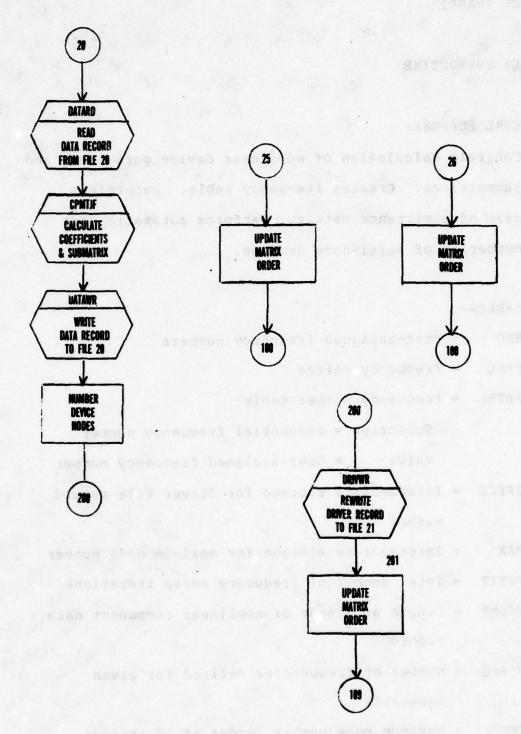
PHASE 1 4 of 8











NAME: PHASE1

TYPE: SUBROUTINE

GENERAL PURPOSE:

Controls calculation of nonlinear device parameters and submatrices. Creates frequency table. Determines size of admittance matrix. Performs automatic node numbering of multi-node devices.

VARIABLES:

FNO = User-assigned frequency numbers

FREQ = Frequency values

FRTBL = Frequency number table:

Subscript = Sequential frequency number

Value = User-assigned frequency number

JRECS = Intermediate storage for driver file record number

MAX = Intermediate storage for maximum node number

MAXIT = Total number of frequency sweep iterations

NCOMP = Length of linear or nonlinear component data record

NFREO = Number of frequencies defined for given generator

NODE = Maximum node number (order of admittance matrix)

NORDR = Total number defined frequencies (order of analysis)

SUBROUTINES CALLED:

CPMTJF CPMTSD CPMTVD CPPT CPT

CPVT DATARD DATAWR DRIVRD DRIVWR

FRROR MTPT MTT MTVT

PHASE2 PTIME TIMOUT

CALLING PROGRAMS:

ENDIN PHASE4 PHASE5 PHASE6 PHASE7

DESCRIPTION:

Subroutine PHASEl is the principal routine on NCAP's first computational phase. It creates the frequency table, controls the calculation of nonlinear device parameters and admittance submatrices, performs automatic node numbering of multi-node circuit elements, and determines the order of the admittance matrix.

The path of execution through subroutine PHASEl is governed by a sequential scan of the driver file. The scan is initialized by setting the driver record number to the start of the file, DRVEC=1. Subroutine DRIVRD reads a driver record from file 21 and places it in the first ten words of common storage. PHASEl responds to the MODE of the driver record, causing it to be processed as follows:

MODE=1 Driver File Header

The frequency count is initialized to NORDR=0, the frequency number table is cleared by FRTBL(I)=0 for I=1, 2,...10, and the number of frequency sweep iterations is set to MAXIT=0. The next file 20 location available for data storage is defined by DATREC=MISC(1) and control returns to the driver file scan.

MODE=13 Start Circuit

The order of the admittance matrix, defined as the largest node number in the circuit, is initialized to NODE=1 and control returns to the driver file scan.

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MODE=3 Generator

Processing of the generator driver record begins by determining if more than one generator employs frequency sweeping. If the present generator does not specify sweeping (MISC(6) =0), the test is by-passed. However, if the present generator employs sweeping (MAXIT>0) and a previous generator also specified sweeping (MAXIT≠0), an appropriate error message is printed and program execution terminates. Ctherwise the number of frequency sweep iterations, MAXIT, is set from MISC(6)

of the generator driver record.

The generator data record is read from file 20 by subroutine DATARD and stored in the common data buffer PUFF. The user-assigned frequency numbers are transferred from PUFF(30)-(40) to the integer vector FNO and the number of frequencies defined for the generator, NFREQ, is set from MISC(1).

Next the frequency numbers and values are appended to the frequency table as follows: for each defined frequency I=1, NFREC, the sequential frequency number is incremented by POPPEPOPP+1. If the frequency table is full (NORDE.CT.6), an appropriate error message is printed and program execution terminates. The user-assigned frequency number is tested for uniqueness against those already in the table. If FNO(I) = FRTEL(J) for any J=1, 2,...(NORDP-1), then an error message is printed and execution stops. Otherwise the frequency number and value are appended to the table by:

FRTPL (NORDR) = FNØ (I)

FFEQ (NORDE) = BUFF (I)

After all the generator frequencies have been placed in the table, the order of the admittance matrix is updated to the maximum of (NODE, MISC(3), MISC(4)) and control returns to the driver file scan.

MODE=4 Linear Components

The linear components data record is read from file 20 by subroutine DATARD and stored in the common data buffer BUFF. The length of the data record, NCOMP, is set from MISC(1). For each four-word component definition the nodes of connection are extracted from the data buffer and stored in K and L and the order of the admittance matrix is updated to the maximum of (NODE, K, L). Control then returns to the driver file scan.

MODE=5 Nonlinear Components

The nonlinear components data record is read from file 20 by subroutine DATAPD and stored in the common data buffer PUFF. The length of the data record, NCOMP, is set from MISC (1). For each 13-word component definition the nodes of connection are extracted from the data buffer and stored in K and L and the order of the admittance matrix is updated to the maximum of (NODE, K, L). Control then returns to the driver file scan.

MODE=6 Vacuum Diode

The vacuum diode data record is read from file 20 by subroutine DATARD and stored in the common data buffer PUFF. Subroutine CPMTVD is called to calculate the nonlinear parameters and admittance submatrix which are appended to the data record. The data record is rewritten to file 20 by subroutine DATAWR, the

order of the admittance matrix is updated to the maximum of (NODE, MISC(3), MISC(4)), and control returns to the driver file scan.

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the driver record is rewritten to the art by

MODE=7 Vacuum Triode

The vacuum triode data record is read from file 20 by subroutine DATAPD and stored in the common data buffer BUFF. Subroutine CPVT and MTVT are called to calculate the vacuum triode nonlinear parameters and admittance submatrix which are appended to the data record. The data record is rewritten to file 20 by subroutine DATAWR and the nodes of the device are numbered by:

MISC(4) = MISC(3) + 2

The driver record is rewritten to file 21 by subroutine DRIVWR, the order of the admittance matrix is updated to the maximum of (NODE, MISC(3), MISC(4)), and control returns to the driver file scan.

MODE=8 Vacuum Pentode

The vacuum pentode data record is read from file 20 by subroutine DATARD and stored in the common data buffer BUFF. Subroutine CPPT and MTPT are called to calculate the vacuum pentode nonlinear parameters and admittance submatrix which are

appended to the data record. The data record is rewritten to file 20 by subroutine DATAWR and the nodes of the device are numbered by:

MISC(4) = MISC(3) + 3

The driver record is rewritten to file 21 by subroutine DRIVWR, the order of the admittance matrix is updated to the maximum of (NODE, MISC(3), MISC(4)), and control returns to the driver file scan.

MODE=9 Fipolar Junction Transistor

The transistor data record is read from file 20 by subroutine DATARD and stored in the common data buffer BUFF. Subroutines CPT and MTT are called to calculate the transistor nonlinear parameters and admittance submatrix which are appended to the data record. The data record is rewritten to file 20 by subroutine DATAWR and the nodes of the device are numbered by:

MISC(4) = MISC(3) + 3

The driver record is rewritten to file 21 by subroutine DRIVWR, the order of the admittance matrix is updated to the maximum of (NODE, MISC(3), MISC(4)), and control returns to the driver file scan.

MODE=10 Semiconductor Diode

The semiconductor diode data record is read from file 20 by subroutine DATARD and stored in the common data buffer BUFF. Subroutine CPMTSD is called to calculate the nonlinear parameters and admittance submatrix which are appended to the data record. The data record is rewritten to file 20 by subroutine DATAWR, the order of the admittance matrix is updated to the maximum of (NODE, MISC(3), MISC(4)), and control returns to the driver file scan.

MODE=20 Junction Field Effect Transistor

The junction field effect transistor data record is read from file 20 by subroutine DATARD and stored in the common data buffer EUFF. Subroutine CPMTJF is called to calculate the nonlinear parameters and admittance submatrix which are appended to the data record. The data record is rewritten to file 20 by subroutine DATAWR and the nodes of the device are numbered by:

MISC(4) = MISC(3) + 3

The driver record is rewritten to file 21 by subroutine DRIVWR, the order of the admittance matrix is updated to the maximum of (NODE, MISC(3), MISC(4)), and control returns to the driver file scan.

MODE=25 Linear Dependent Source

The order of the admittance matrix is updated to the maximum of (NODE, MISC(2), MISC(3), MISC(4), MISC(5)) and control returns to the driver file scan.

ne data reput da rewitten to file 26 by subrouting parage, tou

MODE=26 Nonlinear Dependent Source

The order of the admittance matrix is updated to the maximum of (NODE, MISC(2), MISC(3), MISC(4), MISC(5)) and control returns to the driver file scan.

MODE=14 End Circuit

The driver file record number is saved by JRECS=DRVREC and reset to that of the start circuit driver record by DRVREC =MISC(1). An updated start circuit (MCDE=13) driver record is created with the order of the admittance matrix, NCDE, stored in MISC(4). The start circuit driver record is rewritten to file 21 by subroutine DRIVWR, the driver file record number is restored to DRVREC=JPECS, and control returns to the driver file scan.

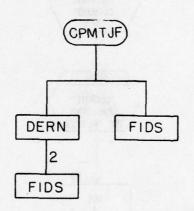
MODE=16 End

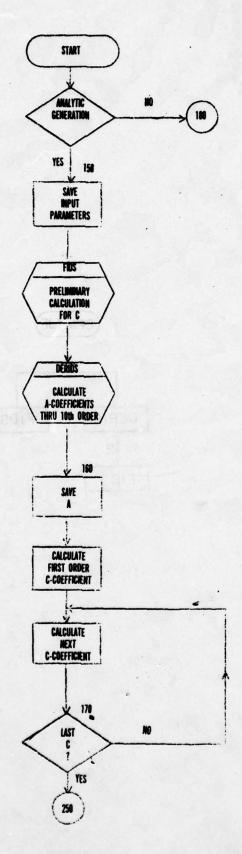
The driver file header is read from file 21 by subroutine DRIVRD. The frequency table is transferred from the FREC and 3-33R

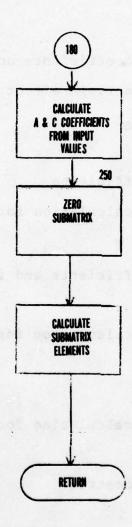
by subroutine DATAWR. The record number and length of the frequency table are stored in the driver header record by STADD=DATREC and LNCTH=20. The next available file 20 storage location is updated by DATREC=DATREC+LNGTH and stored in MISC(1). The frequency count NOPDR is stored in MISC(2) and the driver header is rewritten to file 21 by subroutine DRIVWR.

The FND driver record signals the end of the driver file scan and the completion of Phase 1 processing. Program control transfers to subroutine PHASE2.

For other values of MODE encountered in the driver file scan, control transfers back to the scan without further processing by subroutine PHASE1.







NAME: CPMTJF

TYPE: SUBROUTINE

GENERAL PURPOSE:

Calculates nonlinear coefficients and creates admittance submatrix for junction field effect transistor

VARIABLES:

A = Nonlinear coefficients

AX = Intermediate calculation for A

B = Input values

C = Nonlinear coefficients and input values

CGD = Input value

D = Intermediate calculation for A

FK = Input value

FM = Input value

GIDS = Intermediate calculation for C

IDMAX = Input value

PJ = Admittance submatrix

RHO = Input Value

Q = Intermediate storage for input values

RS = Input value

VBI = Input value

VGS = Intermediate storage for input value

VGSI = Input value

VO = Input value

VP = Input value

VSUM = Intermediate calculation for C

W = Coefficient used in calculation for C

WM = Exponent used in calculation for C

ZFACT = Multiplicative accumulator for factorial

SUBROUTINES CALLED.

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CALLING PROGRAMS:

PHASE1

DESCRIPTION:

Subroutine CPMTJF calculates the device parameters and admittance submatrix for the junction field effect transistor. The parameters are derived from the input values stored in the first 22 words of the JFET data record which is transmitted to subroutine CPMTJF in the common data buffer. The calculated device parameters and admittance submatrix are stored in the data record which is returned to subroutine PHASE1 and rewritten to file 20 for use by subsequent NCAP phases.

The FET model includes analytic generation of the nonlinear coefficients A and C from input values or the direct calculation of A and C coefficients from input parameters. The method of calculating coefficients is determined by the value of MISC(2) in the JFET driver record. For MISC(2)=2 the coefficients are

generated analytically; otherwise the coefficients are calculated directly.

The analytic generation of coefficients begins by storing the input B-parameters in the vector Q and saving B(5) in VGS. Then the function VGSI is calculated iteratively under control of the index I. VGSI is recursive in IDS, a function evaluated by subroutine FIDS. (Since the calculation of IDS is presented in detail in the narrative description of subroutine FIDS, it is not repeated here.) The algorithm for computing VGSI begins by evaluating IDS at the operating point X0=VGSI, an input parameter. Upon return from subroutine FIDS, the calculated value for IDS(X0) is stored in GIDS.

For successive iterations, the operating point is recomputed as $X\emptyset = VGSI - GIDS * RS$ and the function $IDS(X\emptyset)$ is evaluated again. After the tenth iteration, the final result is added to $V\emptyset$ to form the expression:

The parameters are delived from the imput value accided in the

VSUM = VØ + VGSI

The first order coefficient C(1) is then defined as.

 $C(1) = FK * |VSUM|^{-FM}$

$$A(3) = A(3) + 2. * B(2)^{2} * RS$$

$$A(4) = A(4) + 5. * B(2) * B(3) * RS + 5. * B(2)^2 * RS^2$$

$$D^6$$

$$AX = (6. * B(2) * B(4) * RS) + (3. * B(3)^{2} * RS)$$

$$D^{7}$$

$$A(5) = A(5) + AX$$

$$AX = \frac{21. * B(2)^2 * B(3) * RS^2}{D^8}$$

$$A(5) = A(5) + AX + 14. * B(2) *RS^{3}$$

The C-coefficients are calculated directly from the input C-vector according to:

$$C(I) = C(I) / I!$$
 for $I = 2,3,...10$

After the nonlinear coefficients have been calculated, subroutine CPMTJF calculates the complex admittances for the 4-node junction field effect transistor and stores them in the submatrix PJ. The submatrix storage area is first zeroed, and then the admittances are placed in the submatrix as follows:

Higher order coefficients C(I), I=2,3,...10 are derived by differentiating C(I-1) and dividing by (I-1)! Since each C(I-1) is a function of the form mx, its derivative is computed directly as nmx. The coefficient of the derivative is developed in W, the exponent in WM, and the factorial expression is accumulated in ZFACT/I.

The analytic generation of A-coefficients is performed by subroutine DERIDS, which calculates the coefficients up to sixth order by differentiating the function IDS.

For the direct calculation of the FET nonlinear coefficients, the A-coefficients are derived from the input values RS and B. First the intermediate calculations to simplify expressions.

$$D = 1. - B(1) * RS$$

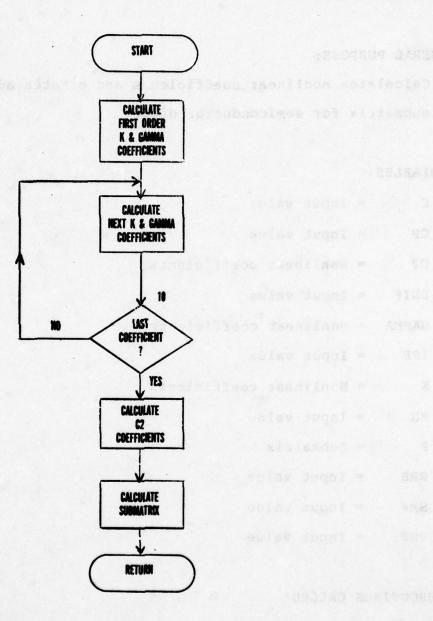
$$A(I) = B(I)/2^{I+1}$$
 for $I = 2,3,...10$

Then more specifically:

$$A(1) = B(1) / D$$

On a decree that it is the service of order or the parties of the

IMAG	PJ(2,1,4)=	0.	PJ(2,2,4)=	0.	PJ(2,3,4)=		0.	PJ(2,4,4)=	0.
REAL	PJ(1,1,4)=	0.	PJ(1,2,4)=	0.	PJ(1,3,4)=		-1./RS	PJ(1,4,4)=	1./RS
IMAG	PJ(2,1,3)=	-c(1)	PJ(2,2,3)=	0.	PJ(2,3,3)=		c(1)	PJ(2,4,3)=	.0.
REAL	PJ(1,1,3)=	0.	PJ(1,2,3)=	-A(1)	PJ(1,3,3)=	۸(1)+	1./RS	PJ(1,4,3)=	-1./RS
IMAG	PJ(2,1,2)=	-CGD	PJ(2,2,2)=	ഡാ	PJ(2,3,2)=		0.	PJ(2,4,2)=	.0
REAL	PJ(1,1,2)=	0.	PJ(1,2,2)=	.0	PJ(1,3,2)=		0.	PJ(1,4,2)=	0.
IMAG	PJ(2,1,1)= . C(1)	dDO+	PJ(2,2,1)=	-can	PJ(2,3,1)=		-c(1)	PJ(2,4,1)=	.0
REAL	PJ(1,1,1)=	0.	PJ(1,2,1)=	A(1)	PJ(1,3,1)=		-4(1)	PJ(1,4,1)=	0.



NAME: CPMTSD

TYPE: SUBROUTINE

GENERAL PURPOSE:

Calculates nonlinear coefficients and creates admittance submatrix for semiconductor diode

VARIABLES:

C = Input value

CP = Input value

C2 = Nonlinear coefficients

DNIF = Input value

GAMMA = Nonlinear coefficients

IFB = Input value

K = Nonlinear coefficients

MU = Input value

P = Submatrix

RRB = Input value

SMK = Input value

VRB = Input value

SUBROUTINES CALLED:

NONE

CALLING PROGRAMS .

PHASE1

DESCRIPTION:

Subroutine CPMTSD calculates the device parameters and admittance submatrix for the semiconductor diode. The parameters are derived from the input values stored in the first eight words of the semiconductor diode data record which is transmitted to subroutine CPMTSD in the common data buffer. The calculated device parameters and admittance submatrix are stored in the data record which is returned to subroutine PHASE1 and rewritten to file 20 for use by subsequent NCAP phases.

The semiconductor diode parameters include the coefficients K and C2 for the forward-biased model and the coefficients GAMMA for the reverse-biased model.

The calculation of coefficients K for the forward-bias model is recursive. The linear coefficient is given by:

$$K(1) = IFB/(.026 * DNIF)$$

Higher order coefficients are calculated according to:

$$K(I) = K(1) * K(I-1) / (I * IFB)$$

The C2 parameters are defined as follows:

$$C2(1) = C + IFB * CP$$

$$C2(2) = CP$$

The calculation of coefficients GAMMA for the reverse-bias model is recursive. The linear coefficient is given by:

Higher order coefficients are calculated according to:

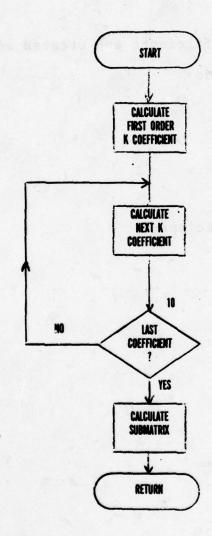
$$GAMMA(I) = (I-1-MU) * CAMMA(I-1)/(I*VRP)$$

The admittance submatrix for the semiconductor diode is the single complex number represented by:

Subsouring CPSI D in the common data buffer The calculated

$$P(1) = K(1) + (1./RRB)$$

$$P(2) = C2(1) + GAMMA(1)$$



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NONLINEAR CIRCUIT ANALYSIS PROGRAM (NCAP) DOCUMENTATION. VOLUME--ETC(U)
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NAME: CPMTVD

TYPE · SUBROUTINE

GENERAL PURPOSE:

Calculates nonlinear coefficients and creates admittance submatrix for vacuum diode

VARIABLES:

C = Input value

G = Input value

K = Nonlinear coefficients

P = Submatrix

VPO = Input value

SUBROUTINES CALLED:

NONE

CALLING PROGRAMS:

PHASE1

DESCRIPTION .

Subroutine CPMTVD calculates the device parameters and admittance submatrix for the vacuum diode. The parameters are derived from the input values stored in the first three words of the vacuum diode data record which is transmitted to subroutine

CPMTVD in the common data buffer. The calculated device parameters and admittance submatrix are stored in the data record which is returned to subroutine PHASEI and rewritten to file 20 for use by subsequent NCAP phases.

The calculation of coefficients K for the vacuum diode's single nonlinearity is recursive. By definition the linear coefficient is given by:

$$K(1) = 1.5 * G * SQRT(VPO)$$

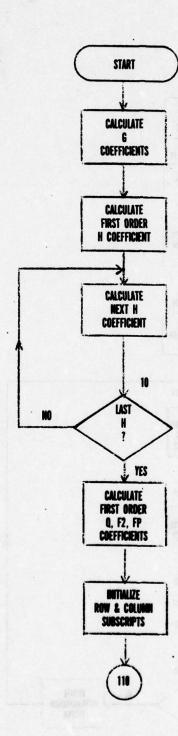
Figher order coefficients are calculated according to:

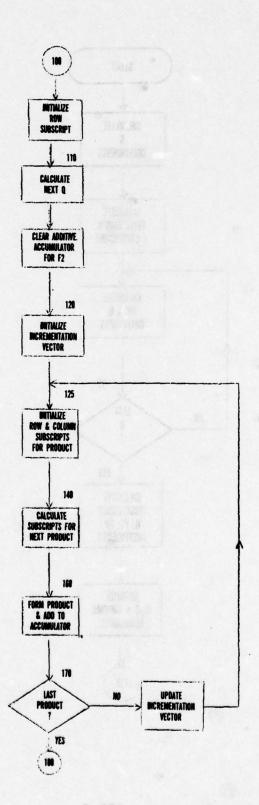
$$K(I) = (2.5-I) * K(I-1)/I*VPO$$

The admittance submatrix for the vacuum diode is the single complex number represented by:

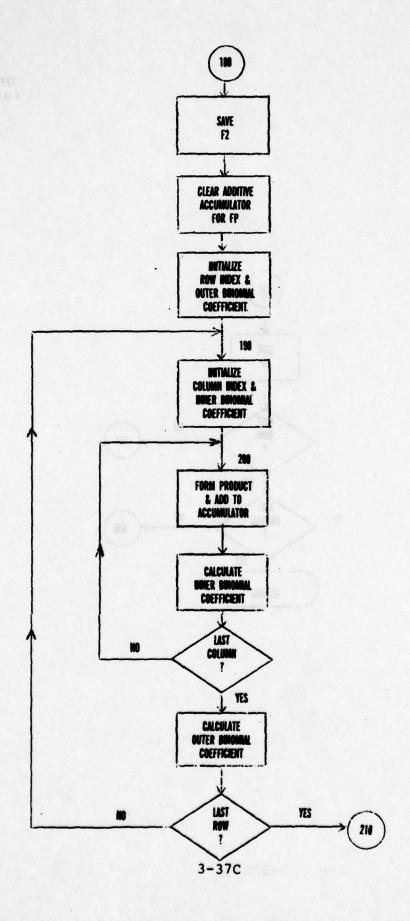
$$P(1) = K(1)$$

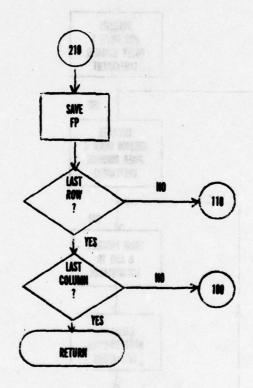
$$P(2) = C$$











NAME: CPPT

TYPE: SUPROUTINE

GENERAL PURPOSE .

Calculates nonlinear coefficients for vacuum pentode

- Presimilarly collected for PF and PF

anisy ment -

VARIABLES:

A = Intermediate calculation

D = Input value

EM = Input value

FP = Nonlinear coefficients

F2 = Nonlinear coefficients

G = Nonlinear coefficients

GO = Input value

H = Nonlinear coefficients

IC = Binomial coefficient for JP

INDX = Incrementation vector for differentiation algorithm

JC = Pinomial coefficient for J2

JP = Row subscript posterior and the street works and

J2 = Column subscript

KP = Row subscripts for differentiation algorithm

K2 = Column subscripts for differentiation
algorithm

MU = Input value

PHI = Input value

Preliminary calculation for FP and F2
nonlinearities

S = Additive accumulator

VP = Input value

V1 = Input value

VlT = Input value

V2 = Input value

SUBROUTINES CALLED:

NONE

CALLING PROGRAMS:

PHASE 1

DESCRIPTION :

Subroutine CPPT calculates the device parameters for the vacuum pentode. These parameters consist of the coefficients for the four nonlinearities of the device model G, H, F2, and FP.

The parameters are derived from the input values stored in the first twelve words of the pentode data record which is transmitted to subroutine CPPT in the common data buffer. Calculated device parameters are stored in the data record which is returned to subroutine PHASE1 and rewritten to file 20 for use by subsequent NCAP phases.

The calculaton of coefficients for the H nonlinearity begins by defining.

A=V1+PH1+V2/MU

H(1) = SQRT(A**3)

Higher order coefficients are calculated recursively according to:

H(I+1) = (2.5-FLOAT(I))*H(I)/A

The G coefficients are calculated by:

G(1) = GO*(1.-V1/V1T)

G(2) = -GO/VIT

For the F2 coefficients, the computations begin by defining:

Q(1,1) = D*(VP/V2) **EM

F2(1,1)=1./(1.+0(1,1))

Higher order coefficients are generated by a differentiation algorithm and stored in the upper triangle of the array F2(5,5). Within the first row, each successive column entry F2(1,J2+1) is the derivative with respect to V2 of the preceding column entry F2(1,J2). Similarly, within a given column, each successive row entry F2(JP+1, J2+1) is the derivative with respect to VP of the preceding row entry F2(JP,J2+1).

Because F2 is a function of O, an upper triangular array O(5,5) is constructed to simplify the arithmetic. The elements of the O array are derived and stored according to the same differentiation process as the F2 array.

The FP coefficients are calculated directly from the F2 and O arrays according to

$$FP(JP+1,J2+1) = \sum_{I=0}^{JP} {JP \choose I} \sum_{J=0}^{J2} {J^2 \choose J} F_2(JP-I+1,J2-J+1) Q(I+1,J+1)$$

In subroutine CPPT the elements of the Q, F2, and FP arrays are generated by columns within a double loop. The outer loop is indexed by J2 which ranges from 0 to 4 and serves as the basis for the column subscript. The inner loop is indexed by JP which serves as the basis for the row subscript. JP ranges from 1 to (4-J2) for the first column, effectively omitting calculations for the linear coefficients already evaluated, and from 0 to (4-J2) for the remaining columns. The indexing technique is summarized as follows:

STEP 1: Initialization

- a) Set column index J2=0 and row index JP=1. Go to Step 2 for first column.
- b) Set row index JP=0 for successive columns.

STEP 2: Calculate Coefficients

- a) Calculate Q(JP+1,J2+1)
- b) Calculate F2(JP+1,J2+1)
- c) Calculate FP(JP+1,J2+1)

STEP 3: Increment and Test Row Index

- a) Increment JP=JP+1
- b) If JP.LE. (4-J2) go to Step 2 to calculate next row entry. Otherwise column is complete; go to Step 4.

STEP 4 Increment and Test Column Index

- a) Increment J2=J2+1
- b) If J2.LE.4 go to Step 1b to begin next column. Otherwise calculations complete.

The elements O(JP+1,J2+1) are calculated recursively. For the first row $(JP=\emptyset)$, each successive column element is a function of the preceding element:

Q(JP+1,J2+1) = Q(JP+1,J2) * (-EM-FLOAT(J2-1))/V2 For the remaining rows (JP>0), each successive row element is a function of the preceding row element:

$$Q(JP+1,J2+1) = Q(JP,J2+1) * (EM-FLOAT(JP-1))/VP$$

The elements F2(JP+1,J2+1) are calculated according to a differentiation algorithm similiar to that employed in the bipolar junction transistor (see subroutine CPT). The algorithm is based on the rule of elementary calculus which defines the nth derivative of a product of functions to be a sum of products of lower order derivatives of those functions. Given:

$$F_2(1,1) = [1 + Q(1,1)]^{-1}$$

then its derivative with respect to V2, the next column element, can be expressed as:

$$F_2(1,2) = -Q(1,2) * F_2(1,1) * F_2(1,1)$$

Similarly the derivative with respect to VP, the next row element can be expressed as:

$$F_2(2,1) = -Q(2,1) * F_2(1,1) * F_2(1,1)$$

These products of functions yield sums of products of lower order derivatives when differentiated:

$$F_2(1,3) = [Q(1,3) * F_2(1,1) * F_2(1,1)] + [Q(1,2) * F_2(1,2) * F_2(1,1)] + [Q(1,2) * F_2(1,1) * F_2(1,2)]$$

$$F_2(3,1) = [Q(3,1) * F_2(1,1) * F_2(1,1)] + [Q(2,1) * F_2(2,1) * F_2(1,1)] + [Q(2,1) * F_2(1,1) * F_2(2,1)]$$

Each additive term is the product of one Q and two lower-order F2's.

In subroutine CPPT the row subscripts are stored in the integer vector KP(3), where KP(1) is the row subscript of the Q element of the product, KP(2) is the row subscript of the first

F2 element, and KP(3) is the row subscript of the second F2 element. The integer vector K2(3) serves the same purpose for column subscripts. In this context the subscript vectors KP= (1, 1, 1) and K2=(3, 1, 1) represents the product Q(1,3) * F2(1,1) * F2(1,1).

The additive terms are generated under the control of the vector INDX(L), L=1, 2...(J2+JP) where INDX(l)=1 and each INDX(L) ranges from 1 to 3 for all other values of L. The elements of the INDX, KP, and K2 vectors interact to form the complete expression for each F2 element:

The subscripts for each term are derived by incrementing the initial subscript vectors KP=(1, 1, 1) and K2=(1, 1, 1) by INDX(L) according to:

$$KP(INDX(L))=KP(INDX(L))+1$$
 for L=1, 2 ... J2

$$K2(INDX(L))=K2(INDX(L))+1$$
 for $L=(J2+1)$, ... $(J2+JP)$

It is important to note that INDX(1) is always equal to 1, and that the first J2 elements of INDX define the incrementation of the column subscripts K2 while the remaining INDX elements

define the incrementation of the row subscripts KP. Furthermore for given JP and J2 (where F2(JP+1,J2+1) is the element being calculated), there are exactly (JP+J2) defined elements in INDX.

Together these constraints force the correct formulation for F2(1,2) and F2(2,1) which are simply products of the form Q*F2*F2 and do not involve summations. The formulation of F2(1,2) is summarized below where $JP=\emptyset$, J2=1, and the single element INDX=(1) increments K2:

INDX=(1)
$$\blacktriangleright$$
 KP=(1,1,1), K2=(2,1,1) \blacktriangleright Q(1,2)*F2(1,1)*F2(1,1)

The calculation for F2(2,1) is similar. JP=1, J2=0, and the single element INDX=(1) increments KP:

INDX=(1)
$$\Rightarrow$$
 KP=(2,1,1), K2=(1,1,1) \Rightarrow Q(2,1)*F2(1,1)*F2(1,1)

The calculation of F2(1,3) involves the summation for INDX(2)=1 through 3 with INDX(1)=1 and results in three additive terms. The formulation of F2(1,3) is summarized below where $JP=\emptyset$, J2=2, and both elements of the INDX vector increment K2:

INDX= (1,1)
$$\implies$$
 KP= (1,1,1), K2= (3,1,1) \implies Q(1,3) * F2(1,1) * F2(1,1)

INDX=(1,2)
$$\Rightarrow$$
 KP=(1,1,1), K2=(2,2,1) \Rightarrow Q(1,2) * F2(1,2) * F2(1,1)

INDX=(1,3)
$$\blacktriangleright$$
 KP=(1,1,1), K2=(2,1,2) \blacktriangleright Q(1,2) * F2(1,1) * F2(1,2)

The calculation for F2(3,1) is similiar. JP=2, J2=0 and both elements of the INDX vector increment KP:

INDX=(1,1)
$$\Rightarrow$$
 KP=(3,1,1), K2=(1,1,1) \Rightarrow C(3,1) * F2(1,1) * F2(1,1)

INDX=(1,2)
$$\implies$$
 KP=(2,2,1), K2=(1,1,1) \implies Q(2,1) * F2(2,1) * F2(1,1)

INDX=(1,3)
$$\implies$$
 KP=(2,1,2), K2=(1,1,1) \implies C(2,1) * F2(1,1) * F2(2,1)

Figher order derivatives involve more summations and are therefore more complicated. Consider for example F2(2,3) which contains nine additive terms. In this case JP=1, J2=2.

$$F_{2}(2,3) = \int_{\text{INDX}(2)}^{3} \sum_{\text{INDX}(3)}^{3} Q(KP(1),K2(1)) * F_{2}(KP(2),K2(2))$$

$$= 1 = 1$$

$$* F_{2}(KP(3),K2(3))$$

The first J2=2 elements of INDX increment the initial column subscripts K2=(1,1,1) and the remaining 1 element of INDX increments the initial row subscripts KP=(1,1,1). The complete formulation is tabulated below.

INDX=(1,1,1)
$$\blacktriangleright$$
 KP=(2,1,1), K2=(3,1,1) \blacktriangleright Q(2,3)*F2(1,1)*F2(1,1)

INDX=(1,1,2) \blacktriangleright KP=(1,2,1), K2=(3,1,1) \blacktriangleright Q(1,3)*F2(2,1)*F2(1,1)

INDX=(1,1,3) \blacktriangleright KP=(1,1,2), K2=(3,1,1) \blacktriangleright Q(1,3)*F2(1,1)*F2(2,1)

INDX=(1,2,1) \blacktriangleright KP=(2,1,1), K2=(2,2,1) \blacktriangleright Q(2,2)*F2(1,2)*F2(1,1)

INDX=(1,2,2) \blacktriangleright KP=(1,2,1), K2=(2,2,1) \blacktriangleright Q(1,2)*F2(2,2)*F2(1,1)

INDX=(1,2,3) \blacktriangleright KP=(1,1,2), K2=(2,2,1) \blacktriangleright Q(1,2)*F2(1,2)*F2(2,1)

INDX=(1,3,1) \blacktriangleright KP=(2,1,1), K2=(2,1,2) \blacktriangleright Q(2,2)*F2(1,1)*F2(1,2)

INDX=(1,3,3) \blacktriangleright KP=(1,2,1), K2=(2,1,2) \blacktriangleright Q(1,2)*F2(2,1)*F2(1,2)

INDX=(1,3,3) \blacktriangleright KP=(1,1,2), K2=(2,1,2) \blacktriangleright Q(1,2)*F2(2,1)*F2(1,2)

The algorithm used in subroutine CPPT to compute F2(JP+1, J2+1) can be summarized as follows:

STEP 1 Initialization

- a) Initialize additive accumulator S=0.
- b) Initialize INDX(I)=1 for I=1,2,...(JP+J2)

STEP 2. Calculate One Additive Term

- a) Initialize row subscripts KP=(1,1,1) and column subscripts K2=(1,1,1). If J2=0 go to 2C. Otherwise go to 2b.
- b) Increment column subscripts by:

 K2(INDX(L))=K2(INDX(L))+1 for L=1,2,...,J2

- c) Increment row subscripts by:
 KP(INDX(L))=KP(INDX(L))+1 for L=(J2+1),...(JP+J2)
- d) Form product:

O(KP(1),K2(1)) * F2(KP(2),K2(2)) * F2(KP(3),K2(3))and add to accumulator S.

STEP 3: Update and Test INDX

- a) Initialize subscript T=JP+J2.
- b) If T.LE.1, summation complete in S, go to Step 4.
 Otherwise go to 3c.
- c) Increment INDX(T)=INDX(T)+1. If INDX(T).LF.3, go to Step 2. Otherwise go to 3d.
- d) Set INDX(T)=1.

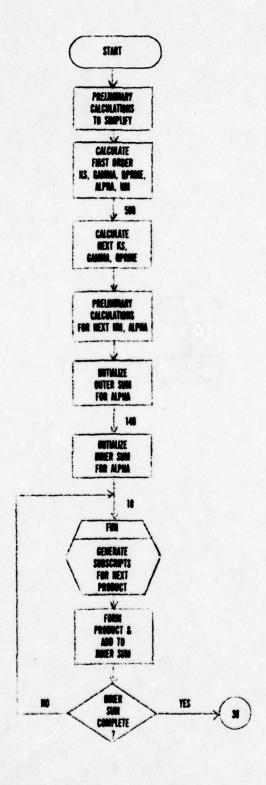
 Decrement T=T-1. Go to 3b.

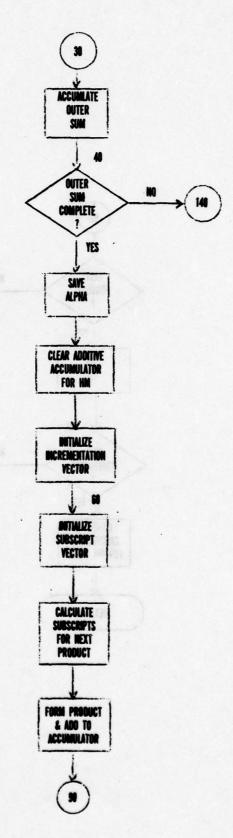
STEP 4: Set F2(JP+1,J2+1) = -S and Exit

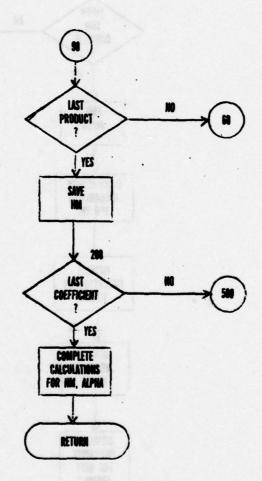
After each Q(JP+1,J2+1) and F2(JP+1,J2+1) have been computed, then FP(JP+1,J2+1) is accumulated in S according to:

S = S + FLOAT(IC*JC) * F2(JP-I+1,J2-J+1) * Q(I+1,J+1) where I ranges from Ø to JP, J ranges from Ø to J2, IC is the binomial coefficient [$_{\rm I}^{\rm J2}$] and JC is the binomial coefficient [$_{\rm I}^{\rm J2}$].









NAME: CPT

TYPE: SUBROUTINE

GENERAL PURPOSE .

Calculates nonlinear coefficients for bipolar junction transistor

Toremarko Macalabia GRADI

it d Calculated parameter

VARIABLES:

A = Input value

AA = Intermediate calculations for ALPHA
nonlinearity

ALPHA = Nonlinear coefficients

ALPHS = Intermediate calculations for ALPHA nonlinearity

AP = Multiplicative accumulator

C = Input value

CJ = Input value

CPRMF = Input value

FIJ = Subscripts for polynomial power series
expansion coefficient

CAMMA = Nonlinear coefficients

HC2 = Nonlinear coefficients

PFEMX = Input value

HM = Nonlinear coefficients

HMO = Calculated parameter

IC = Input value

ICHMO = Calculated parameter

ICMAX = Input value

IE = Calculated parameter

K = Input value

KS = Nonlinear coefficients

MINDX = Incrementation vector for differentiation

algorithm

MORD = Subscripts for differentiation algorithm

MU = Input value

N = Input value

CPRME = Intermediate calculation for HM nonlinearity

RATIC = Preliminary calculation

PP = Input value

RPI = Calculated parameter

RC = Input value

RCI = Calculated parameter

REF = Input value

RE1 = Calculated parameter

S = Additive accumulator

SCC = Input value

SCO = Calculated parameter

SIC = Preliminary calculation

SS = Additive accumulator

SWTCH = Control switch for 'FUN' algorithm

-1 = Initialize algorithm

Ø = Continue algorithm

1 = Algorithm complete

VCP = Input value

VCBO = Input value

SUBROUTINES CALLED.

FUN

CALLING PROGRAMS .

PHASE1

DESCRIPTION .

Subroutine CPT calculates the device parameters for the bipolar junction transistor. These parameters include the coefficients for the four nonlinearities of the device model (KS, CAMMA, HM, and ALPHA) and the miscellaneous values HC2, ICHMØ, PE1, IF, RBI, RCI, C(2), and SCØ which are used in the admittance submatrix and current vector subroutines.

The parameters are derived from the input values stored in the first twenty words of the transistor data record which is transmitted to subroutine CPT in the common data buffer. Calculated device parameters are stored in the data record which is returned to subroutine PFASEL and rewritten to file 20 for use by subsecuent NCAP phases.

The code consists of two main sections. The miscellaneous device parameters, first order coefficients, and preliminary calculations, which are derived from straightforward arithmetic expressions, are computed first. Then the higher order

coefficients are calculated recursively within a loop under the control of the index I which represents the order of the coefficients being calculated.

The calculation of coefficients for the KS nonlinearity uses the arithmetic statement function.

ALG18(X) = ALOG(X) * .4342944819

which computes log of a real argument. The following preliminary calculations are performed to simplify arithmetic expressions.

RATIO = ALGIØ(IC/ICMAX)

PMO = 1./(1.-(VCE/VCEO)**N)

IE = IC*(1.+HFEMX+A*PATIO*RATIC)/PMO/HFEMX

RF1 = RFF * .02567041/IE

Py definition the linear coefficient KS(1) is given by:

KS(1) = 1./PF1

Figher order coefficients are calculated recursively according to

KS(I) = KS(I-1)/FLOAT(I) * REI * IE

The coefficients for the GAMMA nonlinearity are also calculated recursively, where:

GAMMA(1) = K/VCE**MU

and for higher order coefficients:

GAMMA(I) = -GAMMA(I-1) * (MU-FLOAT(I-2))/VCB/FLOAT(I)

The calculation of coefficients for the HM nonlinearity begins with

HMC = 1./(1.-(VCP/VCPC)**N)

(PRME(1) = N * VCP**(N-1.)/VCPO**N

The linear coefficient HM(1) is a product of functions.

HM(1) = HMO * HMO * QPRME(1)

Successive higher order coefficients HM(I),I=2,3,...7 are calculated by differentiating the previous coefficient HM(I-1) with respect to VCE. The algorithm used to compute the higher order coefficients of the HM nonlinearity is derived from the rule of elementary calculus which defines the nth derivative of a product of functions to be a sum of products of lower order derivatives of those functions. In particular, the second order coefficient HM(2), obtained by differentiating HM(1), is a sum of products formulated as follows:

HM(2) = [HM(1) *HM(0) *QPRME(1)] + [HM(0) *HM(1) *QPRME(1)]+ [HM(0) *HM(0) *QPRME(2)]

where HM(1) is the derivative of $HM(\emptyset)$ and QPRME(2) is the derivative of QPRME(1). Note that each additive term is the product of two lower-order HM's and a QPRME. (The equivalency of HMO and $HM(\emptyset)$ is forced by placing HMO and the HM vector in successive common storage locations.)

The third order coefficient HM(3), obtained by differentiating HM(2), consists of nine additive terms, where each term is the product of two lower-order HM's and a CPRME. In general each coefficient HM(I) consists of 3 terms of the form HM * HM * CPRME where the highest derivative (subscript) of the HM's is (I-1).

In subroutine CPT these subscripts are stored in the integer vector MORD where MORD(1) is the subscript of the first HM in the product, MORD(2) is the subscript of the second HM, and MORD(3)

is the subscript of QPRME. In this context MORD= $(\emptyset,\emptyset,1)$ represents the product $HM(\emptyset)*HM(\emptyset)*QPRME(1)$. The 3 additive terms are generated under the control of the vector MINDX(L), L=1,2...(I-1) where each MINDX(L) ranges from 1 to 3. The elements of the MINDX and MORD vectors interact to formulate the complete expression for HM(I):

$$HM(I) = \begin{array}{c} 3 \\ \Sigma \\ MINDX(I-1) \end{array}$$

$$= 1 \\ = 1 \\ = 1 \\ 3 \\ \Sigma \\ MINDX(1) \\ MINDX(2) \\ MINDX(1) \\ M$$

The subscripts for each term are derived by incrementing the initial subscript vector $MORD=(\emptyset,\emptyset,1)$ by MINDX(L) according to:

$$MORD(MINDX(L)) = MORD(MINDX(L)) + 1$$
 for L=1,2,...(I-1)

For example the calculation of HM(2) involves the single summation for MINDX(1) ranging from 1 to 3:

$$HM(2) = \sum_{\substack{\text{MINDX}(1) \\ =1}}^{3} HM(MORD(1)) * HM(MORD(2)) * QPRME(MORD(3))$$

The subscripts for each of the three additive terms are generated by incrementing MORD= $(\emptyset,\emptyset,1)$ by MINDX as follows:

MINDX=(1)
$$\longrightarrow$$
 MORD(1,0,1) \longrightarrow HM(1) * HM(0) * QPRME(1)
MINDX=(2) \longrightarrow MORD(0,1,1) \longrightarrow HM(0) * HM(1) * QPRME(1)
MINDX=(3) \longrightarrow MORD(0,0,2) \longrightarrow HM(0) * QPRME(2)

The calculation of HM(3) involves the double summation for

MINDX(1)=1 to 3 and MINDX(2)=1 to 3 and results in nine additive terms as follows:

MINDX=
$$(1,1)$$
 \longrightarrow MORD= $(2,0,1)$ \longrightarrow HM (2) * HM (0) * QPRME (1) MINDX= $(2,1)$ \longrightarrow MORD= $(1,1,1)$ \longrightarrow HM (1) * HM (1) * QPRME (1) MINDX= $(3,1)$ \longrightarrow MORD= $(1,0,2)$ \longrightarrow HM (1) * HM (0) * QPRME (2) MINDX= $(1,2)$ \longrightarrow MORD= $(1,1,1)$ \longrightarrow HM (1) * HM (1) * QPRME (1) MINDX= $(2,2)$ \longrightarrow MORD= $(0,2,1)$ \longrightarrow HM (0) * HM (2) * QPRME (1) MINDX= $(3,2)$ \longrightarrow MORD= $(0,1,2)$ \longrightarrow HM (0) * HM (1) * QPRME (2) MINDX= $(1,3)$ \longrightarrow MORD= $(0,1,2)$ \longrightarrow HM (1) * HM (0) * OPRME (2) MINDX= $(2,3)$ \longrightarrow MORD= $(0,1,2)$ \longrightarrow HM (0) * HM (1) * QPRME (2) MINDX= $(2,3)$ \longrightarrow MORD= $(0,1,2)$ \longrightarrow HM (0) * HM (1) * QPRME (2) MINDX= $(3,3)$ \longrightarrow MORD= $(0,0,3)$ \longrightarrow HM (0) * HM (0) * QPRME (3)

The algorithm used in subroutine CPT to compute HM(I) can be summarized as follows:

STEP 1: Initialization

- a) Define II=I-l
- b) Calculate OPRME(I)
- c) Initialize additive accumulator S=0.
- d) Initialize MINDX(L)=1 for L=1,2...II

STEP 2: Create One Additive Term

- a) Initialize subscript vector MORD=(0,0,1)
- b) Calculate subscripts by:
 MORD(MINDX(L))=MORD(MINDX(L))+1
 for L=1,2,...II
- c) Form product:

HM (MORD(1)) * HM (MORD(2)) * QPRME(MORD(3))
and add to accumulator S

STEP 3: Update and Test MINDX

- a) Initialize index T=0
- b) Increment T=T+1. If T>II algorithm complete with HM(I) in S. Otherwise go to Step 3c.
- c) Increment MINDX(T)=MINDX(T)+1. If MINDX(T)≤3, go to Step 2. Otherwise go to 3d.
- d) Set MINDX(T)=1 and go to Step 3b.

The final computation for the HM nonlinearity is to divide each HM(I), I=1,2,...7 by I! which is accumulated in P:

HM(I) = HM(I)/P

The calculation of coefficients for the ALPHA nonlinearity begins with:

ALOGE = .4342944819

RATIO = ALG10 (IC/ICMAX)

AA(1) = 1. + HFEMX + A * RATIO * (RATIO+2.*ALOGE)

By definition the first order coefficient ALPHA(1) is given by: ALPHA(1) = 1./AA(1)

For higher order coefficients, each ALPHA(I), I=2,3,...7 is calculated according to the expression:

ALPHA(I) =
$$-\frac{1}{AA(1)^{I}}$$
 $\sum_{J=1}^{I-1}$ ALPHA(J) F_{I-1}^{J} (AA)

where $F_{I-1}^{J}(AA)$, a sum of products of AA's, is the polynomial power series expansion coefficient generated by subroutine FUN.

Each AA(I), I=2,3,...7 is calculated according to the expression:

where:

$$P = Q = 1.$$
 for $I=2$
 $Q = P - Q$
 $P = -(I-1) * P$ for $I>2$

Each F_{I-1}^J (AA) is computed by repeated calls to subroutine FUN under control of the variable SWTCH. For SWTCH=-1 the algorithm of subroutine FUN is initialized and each call returns a set of integers in the vector FIJ which determines the subscripts of the AA's to be multiplied together to form one

additive term of the expression F_{I-1}^J (AA). Upon return from subroutine FUN, if SWTCH=0 additional terms of the expression remain to be calculated and FUN is called again. When the expression is complete, SWTCH is returned as 1.

The algorithm used in subroutine CPT for calculating the sum

shere I. (Mil. a sum of products of MA's, is the polynomial power

$$\Sigma$$
 ALPHA(J) F_{I-1}^{J} (AA)

can be summarized as follows:

STEP 1: Initialization

- a) Define II=I-1
- b) Initialize summation S=0.
- c) Initialize index J=1

STEP 2. Calculate F_{I-1} (AA)

- a) Initialize SWTCH=-1 and summation SS=0.
- b) Call subroutine FUN to generate subscripts FIJ(L), L=1,2,...J for one additive term of the expression
- c) Form the product:

and add to summation SS.

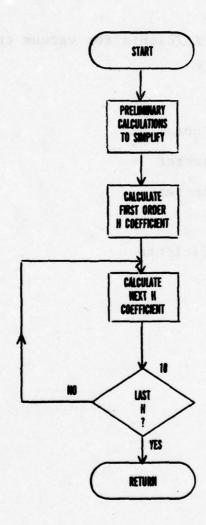
d) If SWTCH=0, go to 2b. Otherwise expression is complete and stored in SS: go to Step 3.

Step 3: Add New Term, Update and Test J

- a) Form the product ALPHA(J)*SS and add to summation S.
- b) Increment J=J+1. If J.LE.(I-1) go to Step 2. Otherwise algorithm complete with summation in S.

ALPHA(I) is then completed by multiplying the summation S by $-1/AA(1)^{I}$. The final computation for the ALPHA nonlinearity is to multiply each ALPHA(I), I=1,2,...7 by HFEMX which is accumulated in Q:

ALPHA(I) = ALPHA(I) * Q



NAME: CPVT

TYPE: SUBROUTINE

GENERAL PURPOSE:

Calculates nonlinear coefficients for vacuum triode

VARIABLES.

A = Preliminary calculation

G = Calculated parameter

GP = Calculated parameter

GO = Input value

H = Nonlinear coefficients

MU = Input value

PHI = Input value

VGO = Input value

VGT = Input value

VPO = Input value

SUBROUTINES CALLED.

NONE

CALLING PROGRAMS:

PHASE1

DESCRIPTION:

Subroutine CPVT calculates the device parameters for the vacuum triode. These parameters include the coefficients H for the vacuum triode's nonlinearity and the parameters G and GP which are used in the admittance submatrix and current vector subroutines.

The parameters are derived from the input values stored in the first nine words of the vacuum triode data record which is transmitted to subroutine CPVT in the common data buffer. The calculated device parameters are stored in the data record which is returned to subroutine PHASEL and rewritten to file 20 for use by subsequent NCAP phases.

The parameters G and GP are given by:

G = GO * (1.-VCO/VGT)

GP = -GO/VGT

The preliminary calculation of A is performed to simplify the arithmetic:

A = VGO + (VPO/MU) + PHI

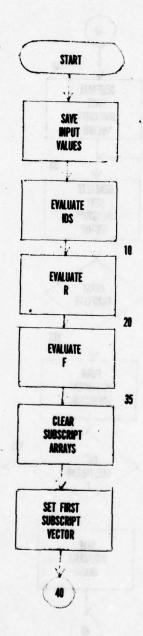
Then by definition the linear coefficient H(l) is given by:

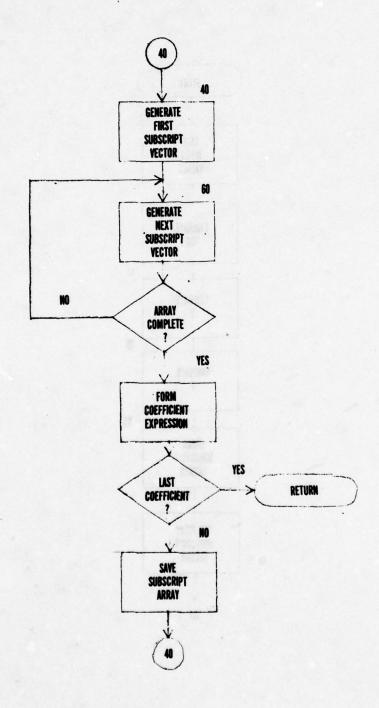
H(1) = SQRT(A**3)

Higher order coefficients are calculated recursively according to:

the vacuum titode a nonlinearity and the parameters & and de.

H(I+1) = (2.5-I) * H(I)/(I*A)





NAME: DERIDS

TYPE: SUBROUTINE

GENERAL PURPOSE:

Analytically generates nonlinear coefficients for junction field effect transistor

VARIABLES:

A = Nonlinear coefficients

COEFR = Intermediate calculation

DZ = Intermediate calculation

EXPR = Intermediate calculation

F = d(IDS)/d(X)

FACT = Accumulator for factorial

IDMAX = Input value

IDS = Value of function IDS

ISEQ = Subscript array

KS = Intermediate calculation

NO = Order of derivative

PROD = Multiplicative accumulator for derivative

Q = Intermediate storage for input values

R = d(X)/d(VGSI)

RHO = Input value

RATIO = Intermediate calculation

SEQ = Subscript array

SUM = Additive accumulator for derivative

VBI = Input value

VGSI = Input value

VP = Input value

X = Intermediate calculation

SUBROUTINES CALLED:

NONE

CALLING PROGRAMS .

CPMTJF

DESCRIPTION:

Subroutine DERIDS analytically generates the nonlinear coefficients A for the junction field effect transistor. Arguments transmitted to the subroutine are the JFET input values Q. The calculated coefficients are returned to the calling program in A.

The nonlinear coefficients A(N), N = 1,6 are defined as:

$$A(N) = \frac{1}{N!} \frac{d^{N}(IDS)}{d(VGSI)}$$

and are calculated by repeatedly differentiating the function IDS with respect to VGSI, where:

IDS =
$$3*IDMAX*RHO$$
 $\left[\frac{-RHO}{2} + \frac{1}{2} \left[RHO^2 + 4(\frac{1}{3} - RATIO + \frac{2}{3}(RATIO)^{3/2})\right]^{1/2}\right]$

and:

$$RATIO = \frac{VGSI + VBI}{VP + VBI}$$

First the function is simplified according to:

Then the first seven derivatives of the function X with respect to VGSI are calculated and stored in R. The first derivative is calculated directly according to:

Higher order derivatives R(I), I = 2,7 are calculated according to

COEFR = COEFR * EXPR * DZ

EXPR = EXPR - 1.

R(I) = COEFR * RATIO ** EXPR

Next the first seven derivatives of IDS with respect to X are evaluated and stored in F. First:

$$N = .5$$

$$F(1) = KS * X ** N$$

The derivatives F(I), I = 2.8 are calculated recursively as:

$$F(I) = (N - (I-1) + 1.) * 1./X * F(I-1)$$

Finally the derivatives of IDS with respect to VGSI are evaluated and stored in A. The first derivative is given by:

$$A(1) = F(2) * R(1)$$

Each higher order derivative A(I), I=2,6 is calculated by differentiating the previous A(I-1). The algorithm used to compute A(I) is derived from the rule of elementary calculus which defines the nth derivative of a product of functions to be a sum of products of derivatives of those functions. In particular, the second order derivative A(2) is a sum of products formulated as follows:

$$A(2) = [F(3) * R(1) * R(1)] + [F(2) * R(2)]$$

where F(3) * R(1) is the derivative of F(2) and R(2) is the derivative of R(1). Note that each additive term is the product of one F and one or more R's, and can therefore be differentiated according to the sum of products rule.

In subroutine DERIDS the subscripts of the F's and R's which constitute a derivative of IDS are developed in the integer array ISEQ, where each row of the ISEQ array contains the subscripts for one additive term of the derivative. The derivatives can then be represented by:

$$A(NO-1) = \sum_{J=1}^{K1} \left[F(ISEQ(1,J)) * \prod_{I=1}^{NO} R(ISEQ(I,J)) \right]$$

Note that the first row element of ISEQ maps to the subscript of F, while the remaining NO-1 elements map to the subscripts of R.

The elements of the ISEQ array are developed recursively. Once a derivative has been calculated, its ISEQ subscript array is saved in the array SEQ which serves as the basis for calculating the ISEQ subscript array for the next derivative.

The algorithm used by subroutine DERIDS for calculating the nonlinear coefficients A can be summarized as follows:

STEP 1: Initialization

- a) Zero subscript arrays ISEQ and SEQ
- b) Set first subscript vector: SEQ(1,1)=1

- c) Initialize FACTorial product: FACT=1.
 - d) Initialize order of derivative: NO=2
- e) Initialize SEQ row index: K=1
- f) Initialize ISEQ row index: Kl=1

STEP 2: Create First Additive Term

- a) Initialize column index: MINDX=1
- b) Calculate first two row elements:

 ISEQ(1,K1) = SEQ(1,K) + 1

 ISEO(2,K1) = 1
 - c) If SEQ(L,K) = Ø for some L=2, NO, term complete.
 Go to Step 3.
- d) Otherwise, calculate next row element:
 ISEQ(L+1,K1) = SEQ(L,K)

STEP 3: Calculate Next Additive Term

- a) Increment column indes: MINDX=MINDX+1
 - b) Increment ISEQ row index: Kl=Kl + 1
- c) If SEQ(MINDX,K) = 0, term complete.

 Go to Step 4.

 - e) Go to Step 3a.

STEP 4: Test for Subscript Array Complete

- a) Increment SEQ row index: K=K+1
- b) If SEQ(1,K) = 0, array complete.
 Go to Step 5.
- c) Otherwise go to Step 2.

STEP 5: Form Coefficient Expression

- a) Clear additive accumulator: SUM=0.
- b) Initialize row index J=1
- c) Initialize multiplicative accumulator:
 PROD = F(ISEQ(1,J))
- d) Form product:

PROD = PROD * R(ISEQ(I,J)) for I=2, NO

e) Accumulate sum of products:

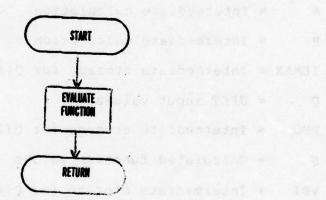
SUM = SUM + PROD

- f) Increment and test J. If J.GT. Kl. e expression complete. Go to Step 5g. Otherwise go to Step 5c.
- g) Calculate factorial: FACT=FACT * (NO-1)
- h) Store coefficient: A(NO-1) = SUM/FACT

STEP 6: Test for Algorithm Complete

- a) Increment NO = NO + 1. If NO=7, algorithm complete.
- b) Otherwise save subscript array in SEQ.
- c) Go to Step le.

FIDS of 1



NAME: FIDS

TYPE · SUBROUTINE

GENERAL PURPOSE .

Evaluates the function IDS for use in calculating coeffients for the junction field effect transistor

VARIABLES:

A = Intermediate calculation

B = Intermediate calculation

IDMAX = Intermediate storage for Q(1)

O = JFET input values

PHO = Intermediate storage for Q(2)

S = Calculated function values

VBI = Intermediate storage for C(4)

VP = Intermediate storage for Q(5)

x = Operating point x-direction

Y = Operating point y-direction

SUBPOUTINES CALLED.

NONE

CALLING PROCRAMS .

CPMTJF DERN

DESCRIPTION:

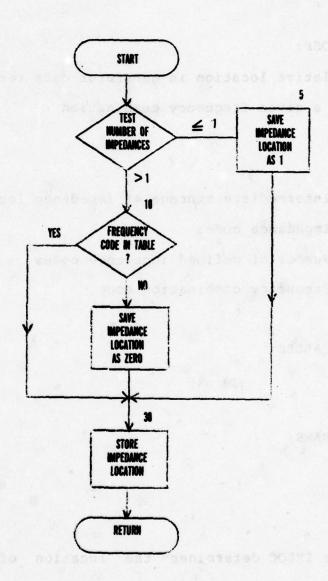
Subroutine FIDS evaluates the function IDS for use by the junction field effect transistor model. Arguments transmitted to the subroutine are: X and Y, the operating points, and Q, the transistor input values. Since IDS is a function of one variable, the second operating point is transmitted as Y=0. The calculated function IDS(X,Y) is returned to the calling program in S.

First the input values are stored as IDMAX=Q(1),RHO=Q(2), VP=Q(3), and VBI=Q(4). Next the expression for IDS is simplified by:

$$B = \frac{X + VBI}{VP + VBI}$$

Then IDS(X) is calculated according to:

IDS = 3 * A * IDMAX *
$$(-\frac{A}{2})$$
 * $\frac{1}{2}\sqrt{4(\frac{1}{3} - B + \frac{2}{3}(B)^{\frac{3}{2}} + A^2}$



NAME: IMLOC

TYPE: FUNCTION

GENERAL PURPOSE:

Finds relative location in generator data record of impedance value for a given frequency combination

VARIABLES:

I .= Intermediate storage of impedance location

IMTBL = Impedance codes

NIMPS = Number of defined impedance codes

SCODE = Frequency combination code

SUBROUTINES CALLED:

NONE

CALLING PROCRAMS .

PHASE 2

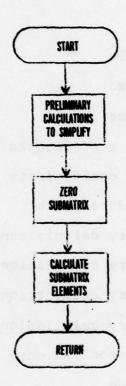
DESCRIPTION .

Function IMLOC determines the location of a generator's impedance (relative to the first impedance in the generator data record) for a particular frequency combination. Arguments transmitted to the function are: the frequency combination code, SCODE; a table of all frequency combination codes for which

impedances are stored in the generator data record, IMTBL; and the number of impedances in the data record, NIMPS. The impedance location is returned to the calling program in IMLOC.

Every NCAP generator has either a constant impedance over all frequencies or one impedance for each possible combination of its input frequencies. The impedances are converted to admittances in Phase 0 and stored with their corresponding frequency combination codes in the generator data record so that the Ith admittance in the record corresponds to the Ith frequency combination code in IMTBL.

For a generator with a constant impedance (NIMPS=1), the impedance location is simply IMLOC=1. Otherwise the location is determined by comparing SCODE to each frequency combination code IMTBL(I) for I=1, NIMPS. If a match is found, the table position I becomes the location returned in IMLOC. If no match is found, the location returned is IMLOC=0.



NAME: MTPT

TYPE: SUPROUTINE

CENERAL PURPOSE:

Create submatrix for vacuum pentode

VARIABLES:

CGC = Input value

CPC = Input value

CPG = Input value

FP = Nonlinear coefficients

F2 = Nonlinear coefficients

G = Calculated parameter

GX1 = Preliminary calculation

GX2 = Preliminary calculation

CYl = Preliminary calculation

CY2 = Preliminary calculation

P = Nonlinear coefficients

MU = Input value

P = Submatrix

SUPROUTINES CALLED:

NONE

CALLING PROGRAMS:

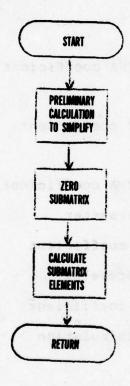
PHASE1

DESCRIPTION:

Subroutine MTPT calculates the complex admittances for the four-node vacuum pentode and stores them by columns in the submatrix PE(32). The submatrix elements are derived from the input data and device parameters stored in the pentode data record which is transmitted to subroutine MTPT through the common data buffer. The admittance submatrix is stored in the last 32 words of the data record which is returned to subroutine PHASE1 and rewritten to file 20 for use for subsequent NCAP phases.

First the intermediate calculations for GX1, GX2, GY1, and GY2 are performed. The submatrix storage area is zeroed and the admittances are defined as follows:

P(26) = -CGC	P(28)=	P (30) -	P(32)= CGC+CPC
P(25)=	P(27)= -GX2-GY2	P(29)= -GX1-GY1	P(31) = GX2+GY2+ GX1+GY1
P(18)= -CPG	P(20)=	P(22)= CPG+CPC	P(24)=-CPC
P(17)= 0.	P(19)= 0.	P(21)= Ø.	P(23)=
P(10)=	P(12)= 0.	P(14)=	P(16)=
P(9)=	P(11) = GX2	P(13) = GY1	P(7) = P(8) = P(15) = 1 $-Gx2-Gx1 - CGC - Gy2-Gy1$
P(2) =	P(4) = 0.	P(6)= -CPG	P(8)=
P(1) = 0.	P(3) = GX2	P(5) = GX1	P(7)=-Gx2-Gx
	P(2) = P(9) = P(10) = P(17) = P(18) = P(25) = CGC + CPG 0. 0. 0CPG 0.	= P(2) = P(9) = P(10) = P(17) = P(18) = P(25) = CGC+CPG	P(2) = P(9) = P(10) = P(17) = P(18) = P(25) = CGC + CPG



NAME: MTT

TYPE · SUBROUTINE

GENERAL PURPOSE:

Creates submatrix for bipolar junction transistor

VARIABLES:

ALPH1 = 1st order ALPHA coefficient

Cl = Input value

C2 = 1st order HC2 coefficient

C3 = Input value

CAM1 = 1st order GAMMA coefficient

HMO = Calculated parameter

HMl = 1st order HM coefficient

ICHMO = Calculated parameter

Kl = 1st order KS coefficient

Ml = Preliminary calculation

PE = Submatrix

PEZ = Zero submatrix

RPI = Input value

PCI = Input value

SUPPOUTINES CALLED:

NONE

CALLING PROGRAMS:

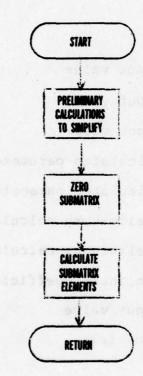
PHASE1

DESCRIPTION:

Subroutine MTT calculates the complex admittances for the four-node bipolar junction transistor and stores them in the complex submatrix PE(2,4,4). The submatrix elements are derived from the input values and device parameters stored in the transistor data record which is transmitted to subroutine MTT through the common data buffer. The admittance submatrix is stored in the last 32 words of the data record which is returned to subroutine PHASE1 and rewritten to file 20 for use by subsequent NCAP phases.

First the intermediate parameter Ml is calculated and the submatrix storage area is zeroed. The admittance submatrix elements are defined as follows:

Imag	PE(2,1,4) = -C1	PE(2,2,4)	Ş	PE(2,3,4) =	0.	PE(2,4,4) = C1+C2
Real	PE(1,1,4) = 0.	PE(1,2,4) = KI* (ALPH1	*IIMO-1.)	PE(1,3,4) = -ALPH1*	HMO*KI	PE(1,4,4) = K1
у Гинд	PE(2,1,3) = -C3	PE(2,2,3) =	-CAMI	PE(2,3,3) =	C3+GAMI	PE(2,4,3) = 0.
Real	PE(1,1,3) = 0.	PE(1,2,9) =	-RCI-MI	PE(1,3,3) =	RCI+MI	PE(1,4,3) = 0.
Ітва	PE(2,1,2) = 0.	PE(2,2,2) =	C2+G4MI	PE(2,3,2) =	-GAMI	PE(2,4,2) = -C2
Real	'PE(1,1,2) = -RBI	PE(1,2,2) = RBI+RCI+	KI* (1 Alphi*HPO)	PE(1,3,2) = -RCI+	ALPH1*HMD *KI	PE(1,4,2) = -Kl
Ітва	PE(2,1,1) = C1+C3	PE(2,2,1) =	0.	PE(2,3,1) =	ទុ	PE(2,4,1) =
Real	PE(1,1,1) = RBI	PE(1,2,1) =	-RBI+MI	PE(1,3,1) =	-M1	PE(1,4,1) = 0.



NAME: MTVT

TYPE · SUPROUTINE

GENERAL PURPOSE:

Creates submatrix for vacuum triode

VARIABLES:

CGC = Input value

CPC = Input value

CPG = Input value

G = Calculated parameter

GP = Calculated parameter

GX = Preliminary calculation

CY = Preliminary calculation

H = Nonlinear coefficients

MU = Input value

P = Submatrix

SUBROUTINES CALLED:

MONE

CALLING PROGRAMS:

PHASE1

DESCRIPTION .

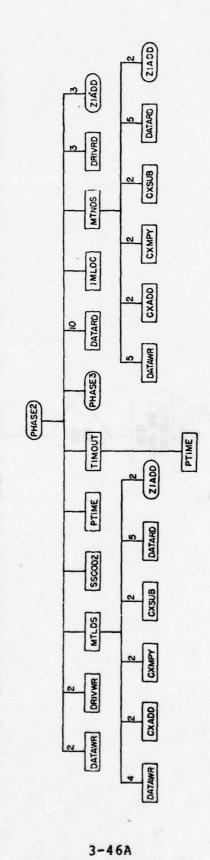
Subroutine MTVT calculates the complex admittances for the three-node vacuum triode and stores them by columns in the submatrix P(18). The submatrix elements are derived from the input values and device parameters stored in the vacuum triode data record which is transmitted to subroutine MTVT through the common data buffer. The admittance submatrix is stored in the last 18 words of the data record which is returned to subroutine PHASE1 and rewritten to file 20 for use by subsequent NCAP phases.

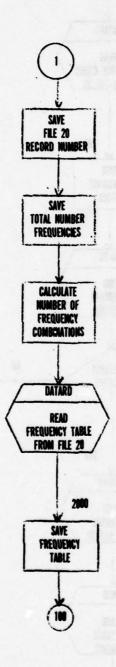
First the intermediate parameters GX and GY are calculated and the submatrix storage area is zeroed. The submatrix elements are defined as follows:

Real	Imag	Real	Imag	Real	Imag
D(1) -	D/2)	-/	- / 0 \		
P(1)=	P(2)=	P(7)=	P(8)=	P(13)=	P(14)=
Ø	CGC+CPG	Ø	-CPG	0.	-cgc
P(3)=	P(4)=	P(9)=	P(10) =	P(15)=	P(16)=
GX	-CPG	GY	CPG+CPC	-GX-GY	-CPC
P(5)=	P(6)=	P(11)=	P(12)=	P(17)=	P(18)=
-GX	-ccc	GY	-CPC	-GX-GY	CGC+CPC

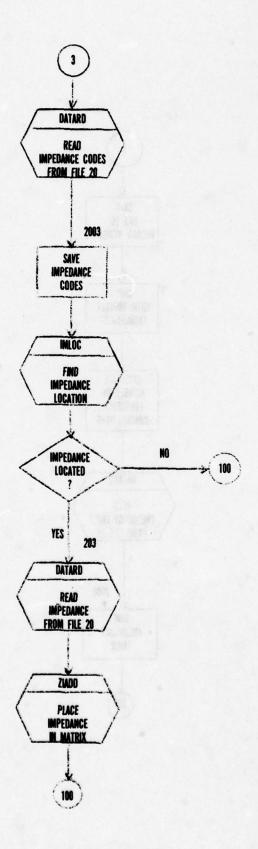
PHASE 2

MTLDS MTNDS ZIADD

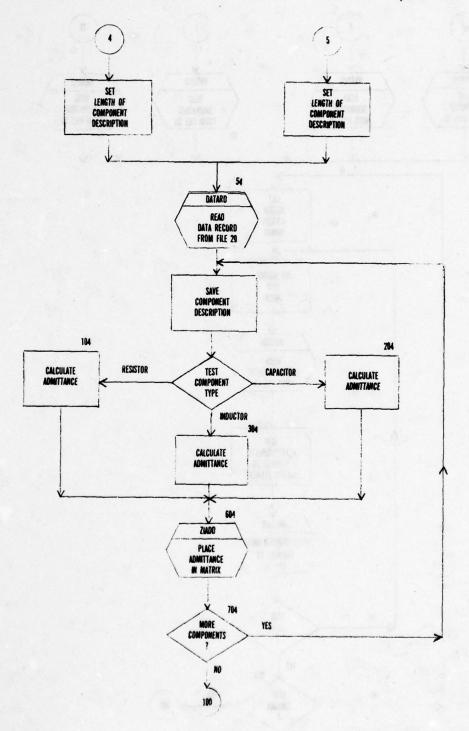






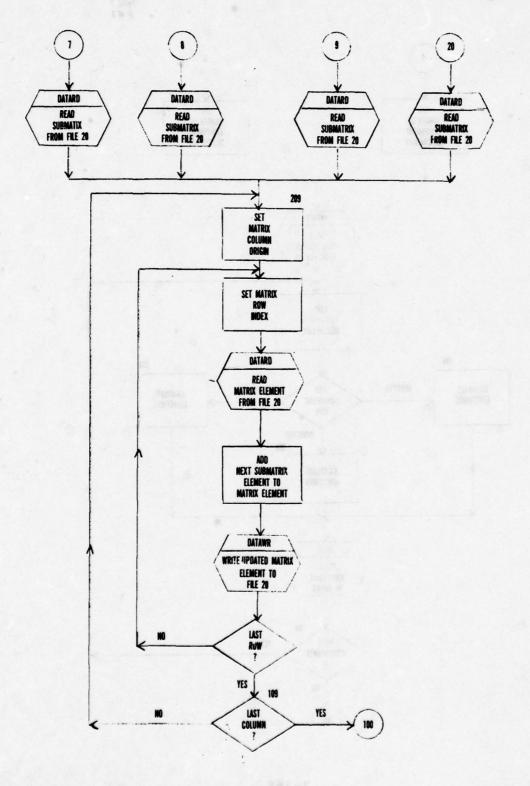




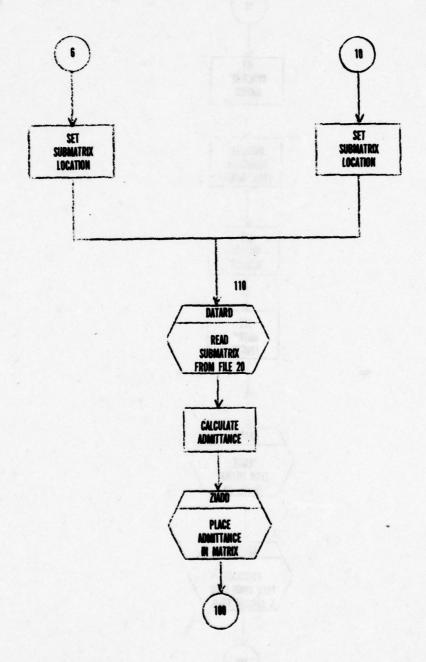




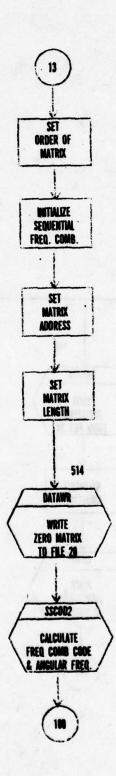
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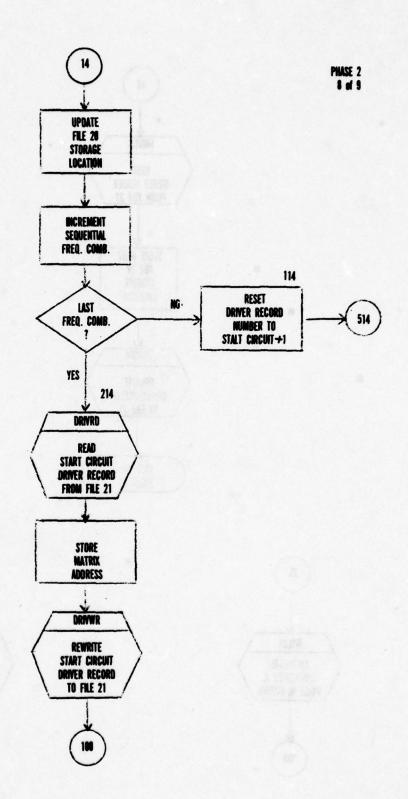


PHASE 2 6 of 9

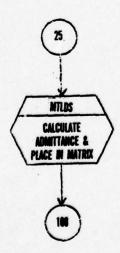














NAME · PHASE2

TYPE · SUBROUTINE

GENERAL PURPOSE:

Construct admittance matrices

VARIABLES:

P = Complex admittance or admittance matrix
element

FREQ = Frequency values

FRTPL = Frequency number table

IMTBC = Impedance codes

JRECS = Intermediate storage for driver file record

LOCPE = Starting address of admittance matrix on file 20

LPE = Length of admittance matrix in words

NCOMP = Length of linear or nonlinear component data
record

NF = Total number of defined frequencies

NNF = Total number of frequency combinations

S = Angular frequency

SCOMB = Sequential frequency combination

SCODE = Frequency combination code

SUPROUTINES CALLED:

DATARD DATAWR DRIVRD DRIVWR IMLOC

MTLDS MTNDS PHASE3 PTIME SSCOD2

TIMOUT - ZIADD

CALLING PROCRAMS .

PHASE1

DESCRIPTION:

Subroutine PHASE 2 is the principal routine of NCAP's second computational phase. Its function is to construct and store the admittance matrices used in the circuit analysis. Because admittances are dependent upon frequency, one admittance matrix is constructed for each possible combination of the circuit's input frequencies.

An admittance matrix is a square array of complex numbers. The order (dimension) of an admittance matrix is determined by the number of nodes in the circuit. Each complex matrix element occupies two decimal storage locations, one for its real part and one for its imaginary part. The matrices are stored by columns on disk file 2%. Every complex element of an admittance matrix is individually addressable. Successive admittance matrices are stored contiguously, in the order they are created, as one data record associated with the START CIRCUIT (MODE=13) driver record.

The construction of admittance matrices in subroutine PHASE2 is controlled by scanning the driver file $(2^{\rm NF}-1)$ times; once for each possible combination of the NF input frequencies. The START

CIRCUIT driver record initiates the construction of the admittance matrix. During its construction in PHASE2 the matrix remains disk-resident. Whenever an admittance is to be placed in the matrix, the appropriate element is read from file 20, updated, and rewritten to its original location.

As the driver file scan proceeds, the admittances from circuit elements are placed in the matrix as their driver records are encountered. The END CIRCUIT driver record (MODE=14) terminates the matrix construction for a particular frequency combination and restarts the driver file scan for the next combination.

The first driver file scan begins by initializing the driver file record number to DRVREC=1. A driver record is read from file 21 by subroutine DRIVRD and subroutine PHASE2 responds to the MODE of the record as follows:

MODE=1 Driver File Header

The next available file 20 data storage location is set to DATREC=MISC(1) and the total number of input frequencies for the circuit is taken as NF=MISC(2). The number of possible frequency combinations is calculated according to

NNF=2**NF-1

The 20-word frequency table is read from file 20 by subroutine DATARD to the common data buffer BUFF. The frequency values are transferred from BUFF(1)-(10) to the decimal vector FREQ and the

frequency numbers are transferred from BUFF(11)-(20) to the integer vector FRTBL. Control then returns to the driver file scan.

MODE=13 Start Circuit

The order of the admittance matrix is defined by ORDER= MISC(4) and the sequential frequency combination is initialized to SCOMB=1. The disk address of the first admittance matrix is taken as LOCPE=DATREC and the length of each admittance matrix (in words) is calculated according to:

LPE=ORDER*ORDER*2

The admittance matrix storage area is cleared, one element at a time, by writing the complex number D=(0., 0.) to the LPE/2 successive file 20 locations beginning at DATREC. The frequency combination code, SCODE, and angular frequency, S, are calculated by subroutine SSCOD2 and control returns to the driver file scan.

MODE=3 Generator

The generator's impedance combination codes are read from file 20 by subroutine DATARD and transferred from the common data buffer BUFF to the integer vector IMTBG. The function IMLOC calculates the location of the admittance value corresponding to the frequency combination SCODE in the generator data record. If

function IMLOC returns an admittance location I.LE.0, there is no admittance for the present frequency combination, and control returns to the driver file scan without further processing.

Otherwise the disk address of the admittance is calculated and subroutine DATARD reads the complex admittance value from file 20 to the common data buffer BUFF. Subroutine ZIADD is called to place the generator's admittance in the matrix and control returns to the driver file scan.

MODE=4 Linear Components

MODE=5 Nonlinear Components

The number of words in each component description is defined as I4=4 for linear components or I4=13 for nonlinear components. The component data record is read from file 20 by subroutine DATARD to the common data buffer BUFF and the length of the data record is set by NCOMP=MISC(1).

For each component (I=1, NCOMP, I4) in the data record, the component type M is taken from the first word of its I4-word description, the nodes of connection K and L are taken from the second and third words, and the component value AA is taken from the fourth word. The component's admittance, a function of the input value and angular frequency S, is calculated and stored in D according to:

Resistor (M=1):

D(1)=1./(AA+1.E-7)

. Primasoon D(2)=0. Tuodiku niga akii saviji aki ba anii da

Capacitor (M=2):

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D(2)=S*AA

Inductor (M=3):

D(1) = 0.

D(2) = -1./(S*AA+1.E-7)

where the factor 1.E-7 prevents division by zero. ZIADD is then called to place the admittance in the matrix. After every component in the data record has been processed, control returns to the driver file scan.

MODE=6 Vacuum Diode

MODE=10 Semiconductor Diode

The relative location of the device admittance within the data record is set by N=13 for the vacuum diode or N=30 for the semiconductor diode. The admittance is read from file 20 by subroutine DATARD and stored in the common data buffer BUFF. After the imaginary part of the admittance, BUFF(2), is multiplied by the angular frequency S, subroutine ZIADD is called

to place the admittance in the matrix. Control then returns to the driver file scan.

MODE=7 Vacuum Triode

MODE=8 Vacuum Pentode

MODE=9 Bipolar Junction Transistor

MODE=20 Junction Field Effect Transistor

The device submatrix is read from file 20 by subroutine DATARD and stored by columns in the common data buffer BUFF. The submatrix is the nodal representation of the device's linear parts. For a device with N nodes, the submatrix is an NxN square array of complex elements. The admittance matrix is the nodal representation of the circuit's linear parts, where successive nodes are represented by rows and columns of like number. The nodes of the device model are numbered sequentially, beginning with the user-assigned base node MISC(3), and are represented in the admittance matrix by the NxN square subset of elements bounded by the MISC(3)rd row and column.

The device's admittances are placed in the matrix by adding the submatrix to the admittance matrix as follows: The first column of the submatrix is added to the MISC(3)rd column of the admittance matrix beginning at (MISC(3), MISC(3)). Then the second column of the submatrix is added to the (MISC(3)+1)st column of the admittance matrix beginning with (MISC(3), MISC(3)+1), and so on until the Nth column of the submatrix has been added to the (MISC(3)+N-1)st column of the admittance matrix.

The appropriate matrix elements are accessed by calculating their relative locations with respect to the origin of the admittance matrix. Matrix elements are complex numbers which occupy two storage locations each, and are stored and operated upon by columns. The locations of the NxN matrix elements bounded by the MISC(3)rd row and column are represented by PHASE2 variables as follows:

entapostos est usua	MISC(3)	MISC(3)+1	willisabus annelo	MISC(3)+N-1
MISC(3)	11 add	11+13	adea tene	12 11969
MISC (3)+1	11+2	11+13+2	W ID by 18	12+2
Labor tout Of	k kiriisa 90	nevolube od5 .		av of complex
ovi e sovocii	eloua (e)	agg wasanibari		70 motosfosect
yodq , ,, sedmo	n edil lo s	naulas bas esol		es are dipases
MISC(3)+N-1	11+14	12+13+14	as lebos (12+14

where: " The desdua sample and tens vo kitting concattang est

I4=2*N-1

I3=2*ORDER

I1= (MISC(3)-1) * (ORDER+1) *2+1

I2=I1+ORDER*(I4-1)

N=Order of submatrix

ORDER=Order of admittance matrix

MISC(3)=Base node of connection

In the PHASE2 code, the index I addresses the column origins 3-46R

and ranges from Il to I2 in increments of I3. Within a column, the index J addresses successive row elements and ranges from I to I+I4 in increments of 2. The index K addresses the submatrix elements stored by columns in the data buffer BUFF. The algorithm for adding a device submatrix to the admittance matrix is summarizes as follows:

STEP 1: Initialization

a) Define I4, a constant depending on the number of nodes in the device model as follows:

MODE	DEVICE	NO. NODES	14
7	Triode	empara 3 milos de	5
8	Pentode	4	7
9	Transistor	4	7
20	JFET	4	7

- b) Initialize submatrix index K= -1
- c) Calculate II, I2, and I3 for admittance matrix column origins
- d) Initialize column origin I=Il

STEP 2: Operate on Ith Column

a) Set end of Ith column J2=I+I4.
Initialize row address J=I

- b) Increment submatrix index K=K+2
- Calculate absolute address of matrix element

 L=DATREC-1+J. Read element from file 20 to D and

 add submatrix element by:

D(1) = D(1) + BUFF(K)

D(2) = D(2) + BUFF(K+1) *S

Write updated matrix element to file 20.

d) Increment row address J=J+2. If J.LE.J2, go to Step 2b. Otherwise go to Step 3.

STEP 3: Increment & Test Column Origin

- a) Increment column origin I=I+I3
- b) If I.LE.I2, go to Step 2. Otherwise algorithm complete, return to driver file scan.

MODE=25 Linear Dependent Source

MODE=26 Nonlinear Dependent Source

The admittances for the linear and nonlinear dependent sources are calculated and placed in the admittance matrix by subroutines MTLDS and MTNDS respectively. Control then returns to the driver file scan.

MODE=14 End Circuit

The next available file 20 storage location is calculated by DATREC=DATREC+LPE and the sequential frequency combination SCOMB is incremented. If the last admittance matrix has not been constructed (SCOMB.LE.NNF), the driver file record number is reset for the next driver file scan by DRVREC=MISC(1)+1. The admittance matrix storage area is cleared, the frequency combination code and angular frequency are calculated by subroutine SSCOD2, and control returns to the driver file scan.

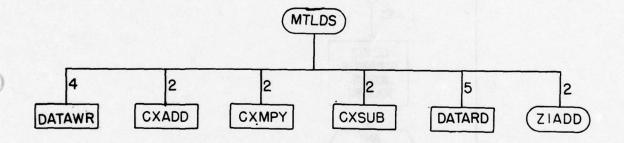
If the last admittance matrix has been constructed (SCOMB.GT.NNF), the matrix data record is attached to the start circuit driver record as follows: The present driver record number is saved in JRECS, and the driver record number is reset to the start circuit by DRVREC=MISC(1). The start circuit driver record is read from file 21 by subroutine DRIVRD, and after saving the matrix data location in MISC(5), the record is rewritten by subroutine DRIVWR. The driver record number is restored from JRECS and control returns to the driver file scan.

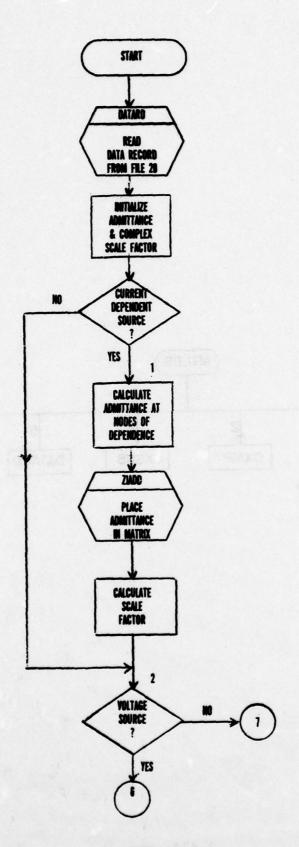
MODE=16 End

The driver file header is read by subroutine DRIVRD, updated to reflect the next available file 20 storage location, and rewritten by subroutine DRIVWR. Program control then transfers to subroutine PHASE3.

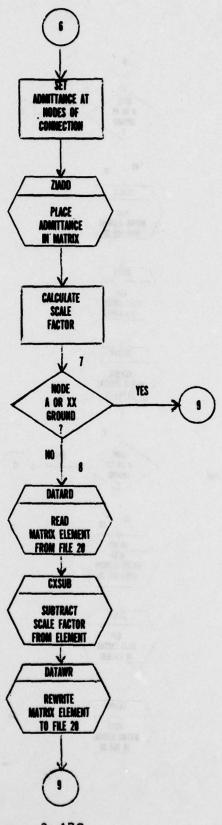
For other values of MODE, control returns directly to the driver file scan without further processing by subroutine PHASE2.

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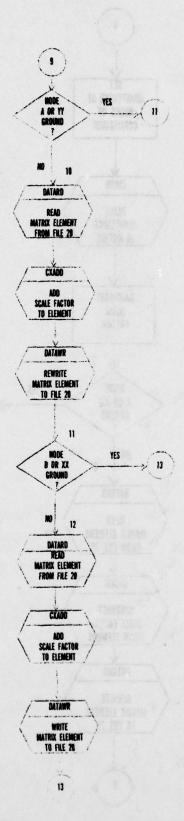


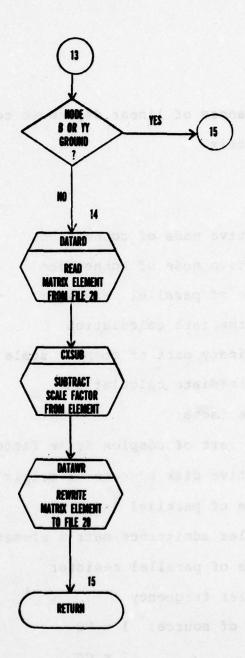
MTLDS 1 of 4



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3-47C





NAME: MTLDS

TYPE: SUBROUTINE

GENERAL PURPOSE .

Places admittances of linear dependent source in admittance matrix

VARIABLES:

A = Positive node of connection

B = Negative node of connection

C = Value of parallel capacitor

CC = Intermediate calculation

CI = Imaginary part of complex scale factor

CL = Intermediate calculation

CONST = Scale factor

CR = Real part of complex scale factor

I = Relative disk address of matrix element

L = Value of parallel inductor

PE = Complex admittance matrix element

R = Value of parallel resistor

S = Angular frequency

SOURCE = Type of source; 1 = VC

2 = CC

3 = VV

4 = CV

X = Intermediate calculation

XX = Positive node of dependence

YY = Negative node of dependence

SUBROUTINES CALLED:

CXADD CXMPY CXSUB DATARD DATAWR

ZIADD

CALLING PROGRAMS:

PHASE2

DESCRIPTION:

Subroutine MTLDS calculates the admittances for the linear dependent source and places them in the admittance matrix. The linear dependent source data record is read from file 20 by subroutine DATARD and stored in the common data buffer BUFF. The complex admittance is initialized to X=(1.,1.) and the scale factor is transferred from the data buffer to CONST.

Then for current dependent sources (SOURCE=2 or 4), the admittance at the nodes of dependence is derived from the values of the parallel components R, C, and L as follows:

X(1)=1./R for $R\neq\emptyset$

X(2) = (S*C) - (1./S*L) for $L \neq \emptyset$

where S is the angular frequency transmitted to subroutine MTLDS as an argument. The admittance at the nodes of dependence (XX and YY) is placed in the admittance matrix by subroutine ZIADD. Subroutine CXMPY multiplies the scale factor by the admittance

and the result is stored in CONST.

For current sources (SOURCE=3 or 4), the admittance at the nodes of connection (A and B) is defined as X=(1.0E10., 0.) and placed in the admittance matrix by subroutine ZIADD. Subroutine CXMPY multiplies the scale factor by the admittance and the result is stored in CONST.

Then the calculated scale factor CONST is placed in the admittance matrix by subtracting it from the matrix elements (A,XX) and (B,YY) and adding it the elements (A,YY) and (B,XX). If neither node A nor XX is ground $(A*XX\neq\emptyset)$, the location of the element (A,XX) relative to the origin of the admittance matrix is calculated by the function LOCATE according to:

I = LOCATE(A, XX) = 2*((XX-1)*ORDER+A)-1

where ORDER is the number of elements in each column of the matrix transmitted to subroutine MTLDS through global common. The absolute address is computed by I=DATREC-1+I and the element is read from file 20 by subroutine DATARD and stored in PE. The complex scale factor is subtracted from PE by subroutine CXSUB and the updated matrix element is rewritten to file 20 by subroutine DATAWR.

If neither node A nor YY is ground $(A*YY\neq\emptyset)$, the location of the element (A,YY) relative to the origin of the admittance matrix is calculated by the function LOCATE according to:

I = LOCATE(A, YY) = 2*((YY-1)*ORDER+A)-1

The absolute address is computed by I=DATREC-1+I and the element is read from file 20 by subroutine DATARD and stored in PE. The complex scale factor is added to PE by subroutine CXADD and the

updated matrix element is rewritten to file 20 by subroutine DATAWR.

If neither node B nor XX is ground $(B*XX\neq\emptyset)$, the location of the element (B,XX) relative to the origin of the admittance matrix is calculated by the function LOCATE according to:

I = LOCATE(B,XX) = 2*((XX-1)*ORDER+B)-1

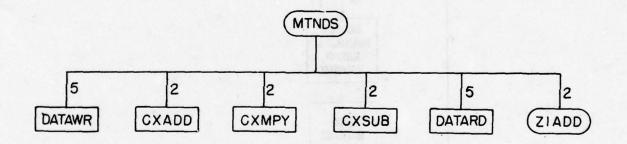
The absolute address is computed by I=DATREC-1+I and the element is read from file 20 by subroutine DATARD and stored in PE. The complex scale factor is added to PE by subroutine CXADD and the updated matrix element is rewritten to file 20 by subroutine DATAWR.

If neither node B nor YY is ground (B*YY \neq 0), the location of the element (B, YY) relative to the origin of the admittance matrix is calculated by the function LOCATE according to:

I = LOCATE(P, YY) = 2*((YY-1)*ORDER+B)-1

The absolute address is computed by I=DATREC-1+I and the element is read from file 20 by subroutine DATARD and stored in PE. The complex scale factor is subtracted from PE by subroutine CXSUB and the updated matrix element is rewritten to file 20 by subroutine DATAWR.

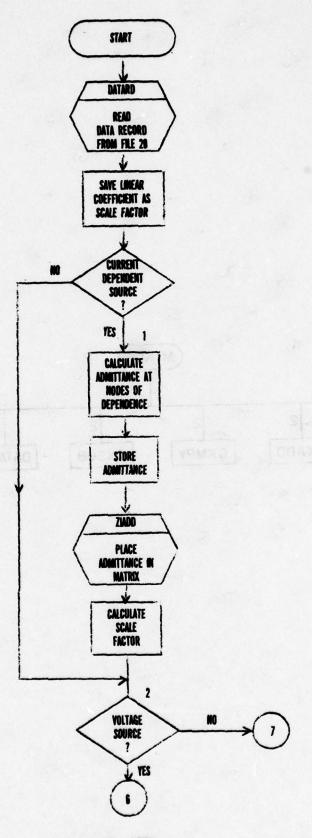
After the admittances have been properly placed in the matrix, subroutine MTLDS returns to the calling program.



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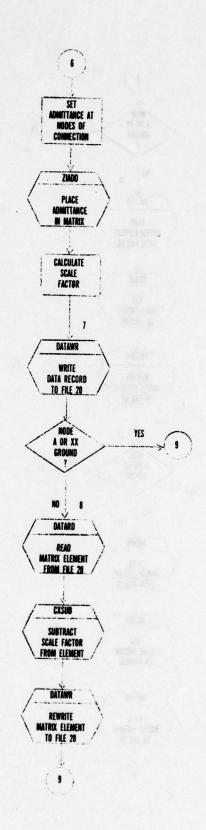
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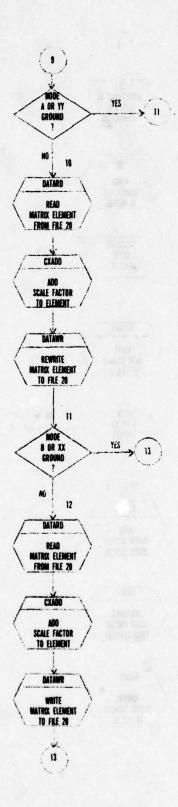
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NONLINEAR CIRCUIT ANALYSIS PROGRAM (NCAP) DOCUMENTATION. VOLUME--ETC(U)
SEP 79 J B VALENTE , S STRATAKOS
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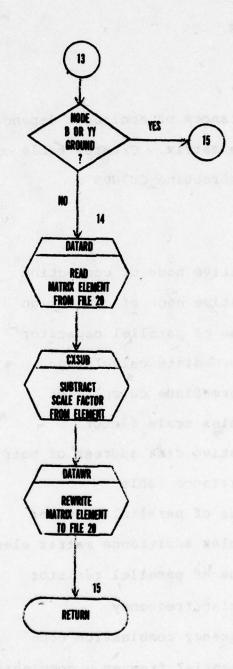


MTNOS 1 of 4

3-48B







NAME: MTNDS

TYPE: SUBROUTINE

GENERAL PURPOSE .

Places admittances of nonlinear dependent source in admittance matrix. Creates table of admittances for use by Subroutine CURNDS

VARIABLES:

A = Positive node of connection

B = Negative node of connection

C = Value of parallel capacitor

CC = Intermediate calculation

CL = Intermediate calculation

CONST = Complex scale factor

I = Relative disk address of matrix element

IMPS = Admittance table

L = Value of parallel inductor

PE = Complex admittance matrix element

P = Value of parallel resistor

S = Angular frequency

SCODE = Frequency combination code

SCOMB = Sequential frequency combination

SCURCE = Type of source 1 = VC, 2 = CC, 3 = VV,

4 = CV

X = Intermediate calculation

XX = Positive node of dependence

YY = Negative node of dependence

SUBROUTINES CALLED:

CXADD CXMPY CXSUB DATARD DATAWR

ZIADD

CALLINC PROGRAMS:

PHASE2

DESCRIPTION:

Subroutine MTNDS calculates the admittances for the nonlinear dependent source and places them in the admittance matrix. The nonlinear dependent source data record is read from file 20 by subroutine DATARD and stored in the common data buffer PUFF. The complex admittance is initialized to X=(1.,1.) and the scale factor is transferred from the data buffer to CONST.

Then for current dependent sources (SOURCE= 2 or 4), the admittance at the nodes of dependence is derived from the values of the parallel components R, C, and L as follows:

X(1)=1./R for $R\neq \emptyset$

X(2) = (S*C) - (1./S*L) for $L \neq \emptyset$

where S is the angular frequency transmitted to subroutine MTNDS as an argument. The complex admittance is stored in the nonlinear dependent source data record at IMPS(1,SCOMB) and IMPS (2,SCOMB) where SCOMB is the sequential frequency combination

transmitted to subroutine MTNDS as an argument.

The admittance at the nodes of dependence (XX and YY) is placed in the admittance matrix by subroutine ZIADD. Subroutine CXMPY multiplies the scale factor by the admittance and the result is stored in CONST.

For current sources (SOURCE= 3 or 4), the admittance at the nodes of connection (A and B) is defined as X=(1.0E10.,0.) and placed in the admittance matrix by subroutine ZFADD. Subroutine CXMPY multiplies the scale factor by the admittance and the result is stored in CONST.

Then the nonlinear dependent source data record is rewritten to file 20 by subroutine DATAWR and the scale factor is placed in the admittance matrix by subtracting it from the matrix elements (A,XX) and (B,YY) and adding it to the elements (A,YY) and (B,XX).

If neither node A nor XX is ground $(A*XX\neq\emptyset)$, the location of the element (A,XX) relative to the origin of the admittance matrix is calculated by the function LOCATE according to:

I= LOCATE(A,XX) = 2*((XX-1)*ORDER+A)-1

where ORDER is the number of elements in each column of the matrix transmitted to subroutine MTNDS through global common. The absolute address is computed by I=DATREC-1+I and the element is read from file 20 by subroutine DATARD and stored in PE. The complex scale factor is subtracted from PE by subroutine CXSUP and the updated matrix element is rewritten to file 20 by subroutine DATAWR.

If neither node A nor YY is ground (A*YY #0), the location of

the element (A,YY) relative to the origin of the admittance matrix is calculated by the function LOCATE according to:

I = LOCATE(A, YY) = 2*((YY-1)*ORDER+A)-1

The absolute address is computed by I=DATREC-1+I and the element is read from file 20 by subroutine DATARD and stored in PE. The complex scale factor is added to PE by subroutine CXADD and the updated matrix element is rewritten to file 20 by subroutine DATAWR.

If neither node E nor XX is ground $(E*XX\neq\emptyset)$, the location of the element (E,XX) relative to the origin of the admittance matrix is calculated by the function LOCATE according to:

I = LOCATE(B,XX) = 2*((XX-1)*ORDER+B)-1

The absolute address is computed by I=DATREC-1+I and the element is read from file 20 by subroutine DATARD and stored in PE. The complex scale factor is added to PE by subroutine CXADD and the updated matrix element is rewritten to file 20 by subroutine DATAWR.

If neither node B nor YY is ground $(E*YY\neq\emptyset)$, the location of the element (B,YY) relative to the origin of the admittance matrix is calculated by the function LOCATE according to:

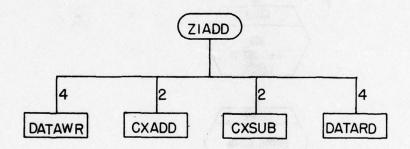
I = LOCATE(B, YY) = 2*((YY-1)*ORDER+B)-1

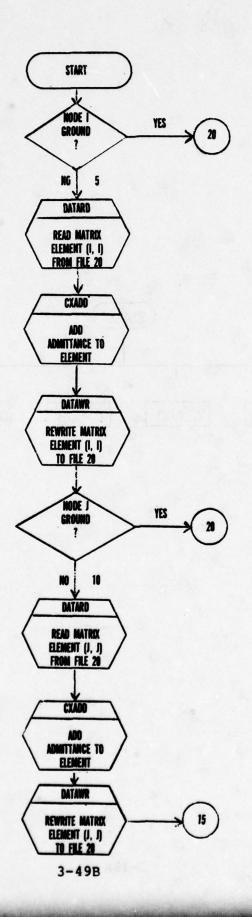
The absolute address is computed by I=DATREC-1+I and the element is read from file 20 by subroutine DATARD and stored in PE. The complex scale factor is subtracted from PE by subroutine CXSUB and the updated matrix element is rewritten to file 20 by subroutine DATAWR.

After the admittances have been properly placed in the

matrix, subroutine MTNDS returns to the calling program.

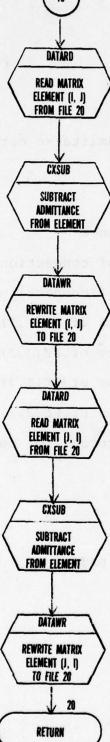
updated matrix element is rewritten to file in the sub-routing





ZWDD 1 of 2





NAME . ZIADD

TYPE · SUEROUTINE

GENERAL PURPOSE:

Places admittances in admittance matrix

VARIAPLES:

A = Complex admittance

I = Positive node of connection

J = Negative node of connection

K = Relative address of PE(I,I) on disk file 29

L = Relative address of PE(J,J) on disk file 20

M = Relative address of PE(I,J) on disk file 20

N = Pelative address of PE(J,I) on disk file 20

PE = Intermediate storage for complex matrix element

SUPROUTINES CALLED:

CXADD CXSUE DATARD DATAWR

CALLING PROGRAMS:

MTLDS MTNDS PHASE2

DESCRIPTION:

Subroutine ZIADD places the complex admittance for a circuit element in the admittance matrix. Arguments transmitted to the subroutine are the complex admittance A and the numbers of the positive and negative nodes of connection, I and J. No arguments are returned to the calling program.

The admittance is placed in the matrix by adding it to the matrix elements (I,I) and (J,J) and subtracting it from the elements (I,J) and (J,I). If the positive node of connection is ground $(I=\emptyset)$, subroutine ZIADD returns to the calling program without performing any operations.

Otherwise the location of the element (I,I) relative to the the origin of the admittance matrix is calculated according to:

Y = 1 + 2*(I-1)*(OFDER+1)

where OPDEP is the number of elements in each matrix column transmitted to the subroutine in global common. The absolute address is computed as LOC(K)=DATREC-1+K and the element is read from file 2% by subroutine DATARD and stored in PE. The complex admittance A is added to PF by subroutine CXADD and the updated matrix element is rewritten to file 2% by subroutine DATAWR.

If the negative node of connection is ground (J=0), subroutine ZIADD returns to the calling program without performing further operations. Otherwise the location of the element (J,J) relative to the origin of the admittance matrix is calculated according to

L = 1 + 2* (J-1)* (ORDER+1)

The absolute address is computed as LOC(L)=DATREC-1+L and the element is read from file 20 by subroutine DATARD and stored in PE. The complex admittance A is added to PE by subroutine CXADD and the updated matrix element is rewritten to file 20 by subroutine DATAWR.

Next the location of the element (I,J) relative to the origin of the admittance matrix is calculated according to:

M= K+(J-I) *OFDER*2

The absolute address is computed as LOC(M)=DATREC-1+M and the element is read from file 20 by subroutine DATARD and stored in PE. The complex admittance A is subtracted from PE by subroutine CXSUB and the updated matrix element is rewritten to file 20 by subroutine DATAWR.

Finally the location of the element (J,I) relative to the origin of the admittance matrix is calculated according to:

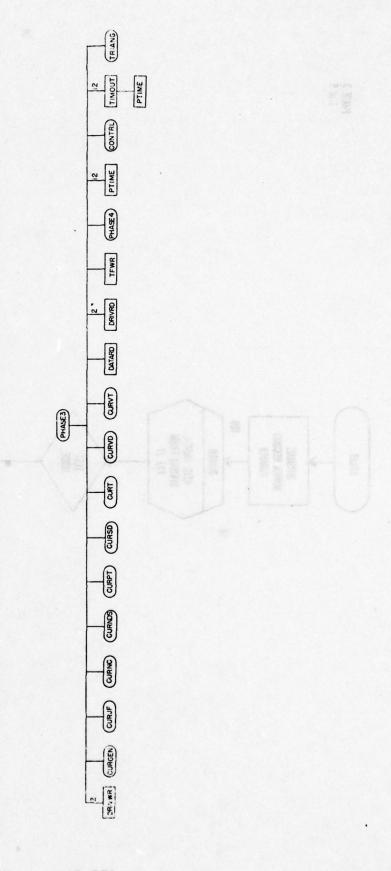
N = K + (J - I) * 2

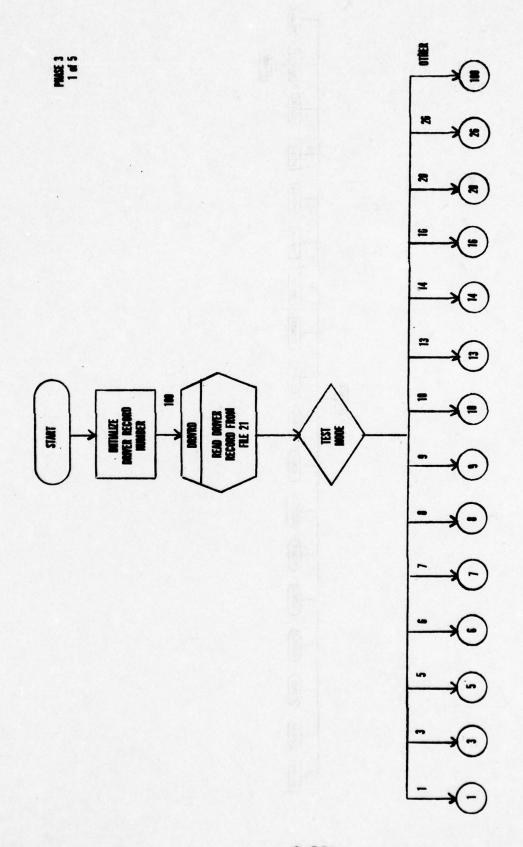
The absolute address is computed as LCC(N)=DATREC-1+N and the element is read from file 20 by subroutine DATARD and stored in PE. The complex admittance A is subtracted from PE by subroutine CXSUB and the updated matrix element is rewritten to file 20 by subroutine DATAWR.

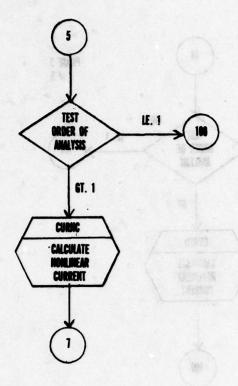
After the admittance has been properly placed in the matrix, subroutine ZIADD returns to the calling program.

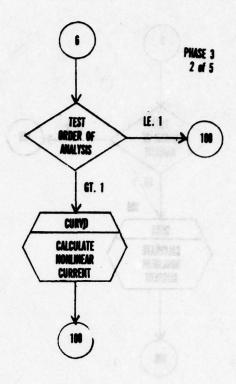
PHASE3

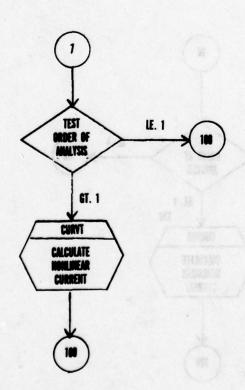
BETA CONTRL CROSS CURGEN CURJF CURNC CURNDS CURPT CURSD CURT CURVD CURVT FNCTN FRPRM LOCTF QFN SSCODE TRIANG UPFRQ UPSEQ

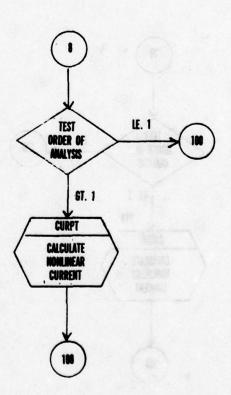


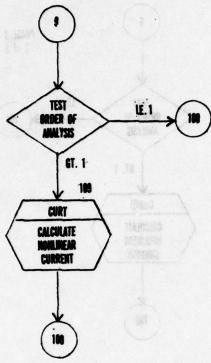


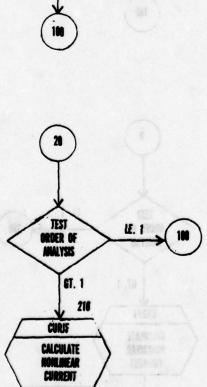


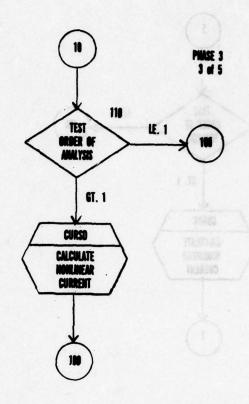


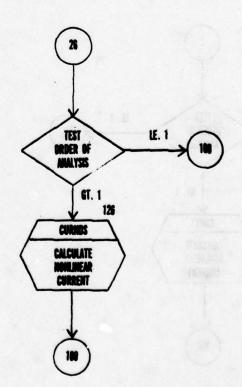


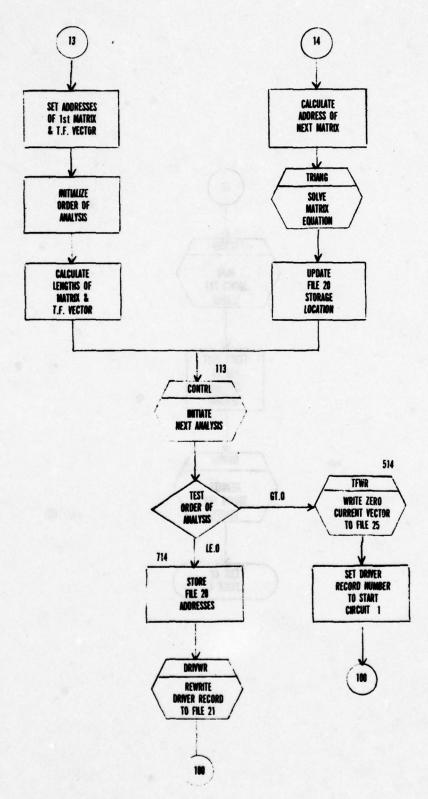














NAME: PHASE3

TYPE: SUPROUTINE

GENERAL PURPOSE .

Calculates current vectors. Computes linear and nonlinear transfer functions for each possible frequency combination by solving the matrix equation

(ADMITTANCE MATRIX) X (TRANSFER FUNCTION) = (CURRENT VECTOR)

VARIABLES:

CUR = Intermediate storage for current vector

element

FREQ = Frequency values

FRTBL = Frequency number table;

Subscript = Sequential frequency number

Value = User-assigned frequency number

LOCPE = Address of 1st admittance matrix on file 20

LOCT = Address of 1st T.F. vector on file 20

LPE = Length of each admittance matrix in words

LT = Length of each T.F. vector in words

MXCRD = Maximum order of analysis

NORDR = Order of analysis

S = Angular frequency

SCODE = Frequency combination code

SCOMB = Seguential frequency combination

SUBROUTINES CALLED:

CONTRL	CURGEN	CURJF	CURNC	CURNDS
CURPT	CURSD	CURT	CURVD	CURVT
DATARD	DRIVRD	DRIVWR	PHASE4	PTIME
TFWR	ттмошт	TRIANG		

CALLING PROGRAMS:

PHASE2

DESCRIPTION:

Subroutine PHASE3 is the principal routine of the last computational phase of the NCAP circuit analysis. It controls the creation of current vectors and calculates and stores the transfer functions for each possible combination of the circuit's input frequencies.

The current vector is a single-dimension array of complex elements, one for each node in the circuit. The nonlinear currents calculated in PHASE3 are stored in the vector so that the current at node I is represented by the Ith complex vector element. The current vector is stored on disk file 25. Every complex vector element occupies two decimal storage locations and is individually addressable.

The transfer functions at each node of the circuit are calculated for every possible combination of the circuit's input frequencies. The transfer functions for a single frequency

combination are stored in a vector of complex elements, one for each node in the circuit, so that the transfer function at node I is represented by the Ith complex vector element. The transfer function vectors are stored on disk file 20. Every complex vector element occupies two decimal storage locations and is individually addrescable. Successive transfer function vectors are stored contiguously in the order they are calculated in PHASE3, as one data record associated with the End Circuit driver record.

The construction of current vectors and calculation of transfer functions in subroutine PHASE3 is controlled by scanning the driver file (2^{MXORD} -1) times, once for each possible combination of the MXORD input frequencies. The Start Circuit driver record initiates the construction of a current vector. During the construction in PHASE3 the current vector remains disk-resident. When a current is to be placed in the vector, the appropriate element is read from file 25, updated, and rewritten to its original location.

As the driver file scan proceeds, the currents generated by circuit elements are calculated and placed in the vector as their driver records are encountered. The End Circuit driver record terminates the current vector construction for a particular frequency combination, calculates and stores the transfer functions for that frequency combination, and restarts the driver file scan for the next combination.

The first driver file scan begins by initializing the driver file record number to DRVREC=1. A driver record is read from

file 21 by subroutine DRIVRD, and subroutine PHASE2 responds to the MODE of the record as follows:

MODE=1 Driver File Header

The next available file 20 data storage location is set by DATREC=MISC(1) and the total number of input frequencies for the circuit is taken as MXORD=MISC(2). The 20-word frequency table is read from file 20 by subroutine DATARD and stored in the common data buffer BUFF. The frequency values are transferred from BUFF(1)-(10) to the decimal vector FREQ and the frequency numbers are transferred from BUFF(11)-(20) to the integer vector FRTBL. Control then returns to the driver file scan.

MODE=13 Start Circuit

The disk address of the first admittance matrix on file 20 is set to LOCPE=MISC(5), the disk address of the first transfer function vector on file 20 is taken as LOCT=DATREC, and the number of nodes in the circuit is set by ORDER=MISC(4). The order of analysis is initialized to NORDR=0. The length of each transfer function vector (in words) is calculated by LT=2*ORDER and the length of each admittance matrix (in words) is calculated by LPE=LT*ORDER.

Subroutine CONTRL is called to initiate the PHASE3

calculations for the first frequency combination. It updates the order of analysis, forms the frequency combination, and creates and stores a table of selected lower order transfer function vectors for use by the current generator subroutines. The call to subroutine CONTRL and the remainder of the initialization is performed by the END CIRCUIT code, and is described in detail under MODE=14. After the initialization is complete, subroutine PHASE3 resumes the driver file scan.

MODE=3 Cenerator

For linear frequency combinations (NORDR=1), subroutine CUPGEN is called to compute the generator current and place it in the current vector. Control then returns to the driver file scan. For higher orders of analysis (NORDR#1), subroutine PHASE3 returns to the driver file scan without further processing.

MCDE=5 Nonlinear Components

For linear frequency combinations (NORDR=1), subroutine PHASE3 returns to the driver file scan without further processing. For higher orders of analysis (NORDR.GT.1), subroutine CURNC is called to compute the nonlinear currents and place them in the current vector. Control then returns to the driver file scan.

MODE=6 Vacuum Diode

For linear frequency combinations (NORDR=1), subroutine PHASE3 returns to the driver file scan without further processing. For higher orders of analysis (NORDR.GT.1), subroutine CURVD is called to compute the nonlinear current and place it in the current vector. Control then returns to the driver file scan.

MODE=7 Vacuum Triode

For linear frequency combinations (NORDF=1), subroutine PHASE3 returns to the driver file scan without further processing. For higher orders of analysis (NORDR.GT.1), subroutine CURVT is called to compute the nonlinear current and place it in the current vector. Control then returns to the driver file scan.

MODE=8 Vacuum Pentode

For linear frequency combinations (NORDR=1), subroutine PHASE3 returns to the driver file scan without further

processing. For higher orders of analysis (NORDR.GT.1), subroutine CURPT is called to compute the nonlinear currents and place them in the current vector. Control then returns to the driver file scan.

MCDE=9 Bipolar Junction Transistor

For linear frequency combinations (NORDR=1), subroutine PHASE3 returns to the driver file scan without further processing. For higher orders of analysis (NORDR.GT.1), subroutine CURT is called to compute the nonlinear currents and place them in the current vector. Control then returns to the driver file scan.

MODE=10 Semiconductor Diode

For linear frequency combinations (NORDR=1), subroutine PHASE3 returns to the driver file scan without further processing. For higher orders of analysis (NORDR.GT.1), subroutine CURSD is called to compute the nonlinear current and place it in the current vector. Control then returns to the driver file scan.

MODE=20 Junction Field Effect Transistor

PHASE3 returns to the driver file scan without further processing. For higher orders of analysis (NORDR.GT.1), subroutine CURJF is called to compute the nonlinear currents and place them in the current vector. Control then returns to the driver file scan.

MODE=26 Nonlinear Dependent Source

For linear frequency combinations (NORDR=1), subroutine PHASE3 returns to the driver file scan without further processing. For higher orders of analysis (NCRDR.CT.1), subroutine CUPNDS is called to compute the nonlinear current and place it in the current vector. Control then returns to the driver file scan.

MODE=14 End Circuit

The transfer function vector is calculated by subroutine TRIANC and stored on file 20. The next available file 20 storage location is updated to DATREC=DATREC+LT and subroutine CONTRL is called to initiate the calculations for the next frequency combination. It updates the order of analysis, forms the next frequency combination, and creates and stores a table of selected

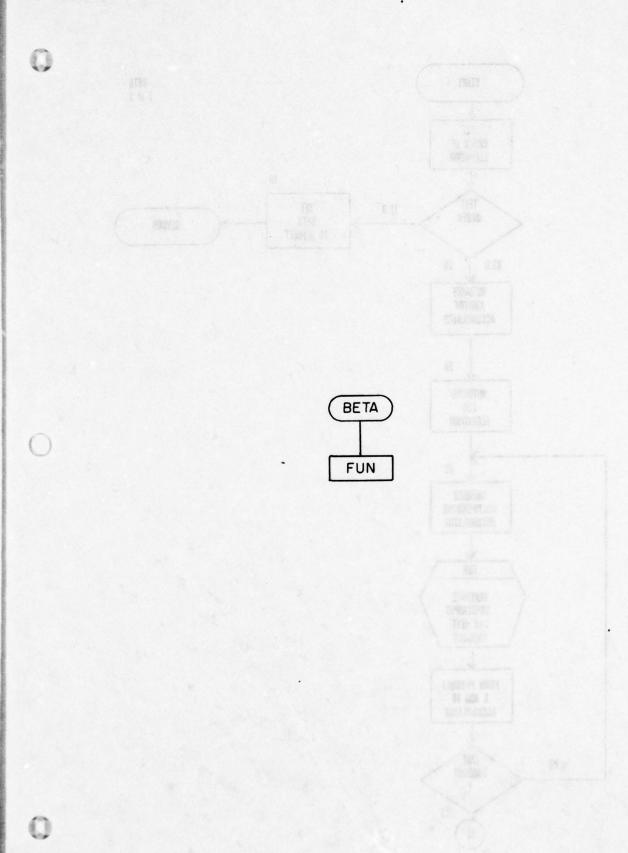
lower order transfer function vectors for use by the current generator subroutines. After returning from subroutine CONTRL, if the last transfer function vector has not been computed (NORDR.GT.0), the current vector storage area is cleared, one element at a time as follows: subroutine TFWR writes the complex number CUP = (0.0.0) to the first LT/2 vector elements on file 25. The driver file record number is reset for the next driver file scan by DRVREC = MISC(1) + 1 and control returns to the driver file scan.

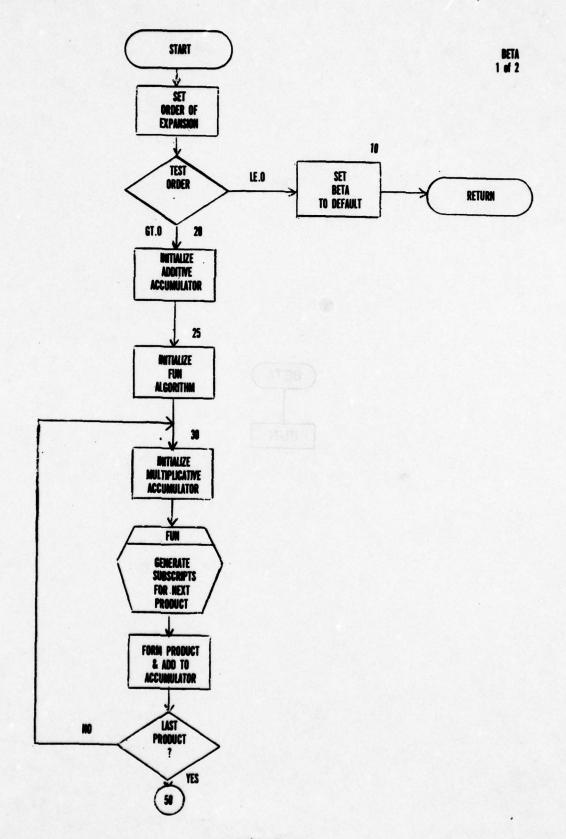
If the last transfer function vector has been calculated (NORDP=-1), the transfer function data record is attached to the End Circuit driver record by MISC(2)=LOCT. The next available file 20 storage location is saved in MISC(3) and the End Circuit record is rewritten to the driver file by subroutine DRIVWR. Control then returns to the driver file scan.

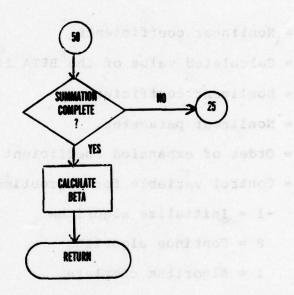
MODE=16 End

The driver file header is read by subroutine DRIVRD, updated to reflect the next available file 20 storage location, and rewritten by subroutine DRIVWR. Program control then transfers to subroutine PHASE4.

For other values of MODE, control returns to the driver file scan without further processing by subroutine PHASE3.







NAME: BETA

TYPE: SUBROUTINE

GENERAL PURPOSE .

Calculates the function B for use by subroutine CURT M

VARIABLES:

ALPHS = Nonlinear coefficients

B = Calculated value of the BETA function

HM = Nonlinear coefficients

ICHMO = Nonlinear parameter

II = Order of expansion coefficient

ISW = Control variable for subroutine FUN

-1 = Initialize algorithm

0 = Continue algorithm

1 = Algorithm complete

K = Index for outer summation of current

calculation

KI = Subscripts calculated in subroutine FUN

KS = Nonlinear coefficients

M = Index for inner summation of current

calculation

PAK = Additive accumulator

PK = Multiplicative accumulator

SNCS = Nonlinear parameter

SUBROUTINES CALLED:

PIIN

CALLING PROGRAMS.

CHRT

DESCRIPTION:

Subroutine BETA calculates the function B_M for use in calculating the nonlinear current for the bipolar junction transistor. Arguments transmitted to the subroutine are: K, the index for the inner summation in the nonlinear current calculation: M, the index for the outer summation in the nonlinear current calculation; ALPHS, HM, and KS, nonlinear coefficients calculated in Phase 1; ICHMØ and SNCS, nonlinear parameters calculated in Phase 1. The calculated function B_M^K is returned to the calling program in B.

If K=M, the function B_M^K is calculated according to B=ICHMØ*HM(M) and subroutine BETA returns to the calling program. Otherwise the function B_M^K is defined as:

$$B = SNCS * HM(M) * \sum_{T=1}^{K-M} ALPHS(I) * F_{I}^{K-M} (KS)$$

where each F_{I}^{K-M} (KS), a sum of products of nonlinear coefficients

KS, is computed by repeated calls to subroutine FUN under the control of the variable ISW. For ISW=-1 the algorithm of subroutine FUN is initialized and each call returns a set of integers in the vector KI which constitute the subscripts of KS's to be multiplied together to form one additive term of the expression. Upon return from subroutine FUN, if ISW=0 additional products remain to be calculated, and subroutine FUN is called again. When the expression is complete, SWTCH is returned as 1 and the algorithm stops.

The complete procedure for calculating the function B_M^{K} in subroutine BETA is summarized as follows:

STEP 1: Initialization

- a) Clear additive accumulator PAK=0.
- b) Initialize index for summation I=1
- c) Initialize FUN algorithm by ISW=-1

STEP 2: Calculate One Additive Term ALPHS(I)*F K-M(KS)

- a) Initialize multiplicative accumulator PK=1.
- b) Call subroutine FUN to generate subscripts KI(J), J=1,2,...I for one additive term.
- c) Form the product

$$PK = \begin{array}{c} I \\ \pi \\ J=1 \end{array} KS(KI(J))$$

d) Form product ALPHS(I) *PK and add to accumulator PAK.

STEP 3. Test ISW

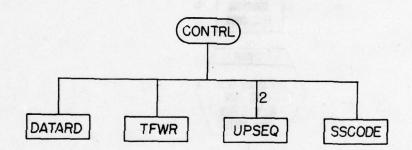
- a) If ISW=0, more products remain to be calculated. Go to Step 2.
- b) Otherwise one additive term complete in PK, go to Step 4.

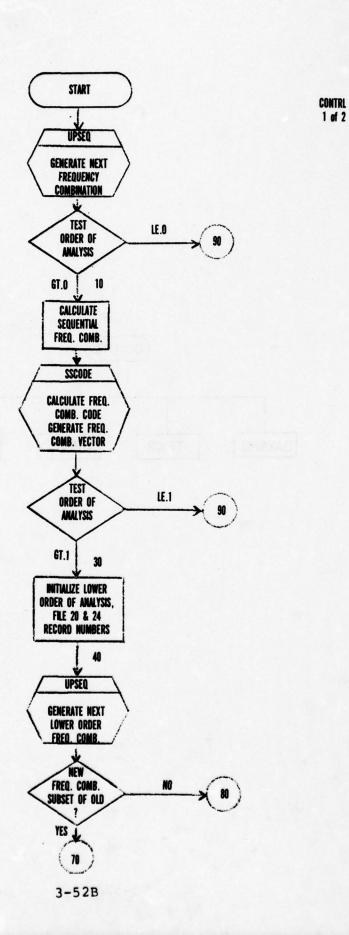
STEP 4: Update and Test Summation Index I

- a) Increment I=I+1
- b) If (I.LE.(K-M)), go to Step lc.
- c) Otherwise summation complete in PAK, go to Step 5.

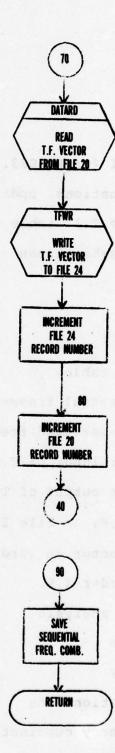
STEP 5. Complete Calculation of BK

- a) Compute B=SNCS*HM(M)*PAK.
- b) Return to calling program with B_{M}^{K} in B.









NAME: CONTRL

TYPE: SUPROUTINE

GENERAL PURPOSE:

Provides frequency control for PHASE3.

Cenerates frequency combinations, updates order of analysis, creates T.F. table for use by current generator subroutines

VARIAPLES:

FREQ = Frequency values

FRTPL = Frequency number table:

Subscript = Sequential frequency number

Value = User-assigned frequency number

K = Record number for input of T.F. from file 20

LL = Record number for output of T.F. to file 24

LOCT = Address of 1st T.F. on file 20

LT = Length of T.F. vector in words

MORDR = Order of lower-order T.F.

MXORD = Maximum order of analysis

NORDR = Order of analysis

S = Angular frequency

SCODE = Frequency combination code

SCOME = Sequential frequency combination

SEQ = Frequency combination vector

SEC1 = Lower-order frequency combination vector

T = Intermediate storage for T.F. element

SUBPOUTINES CALLED.

DATARD SSCORE TEWR UPSEQ

CALLING PROCRAMS:

PHASE3

DESCRIPTION:

Subroutine CONTRL initiates the PHASE3 calculations, of the current vector and transfer functions for a particular frequency combination by forming that frequency combination and creating a table of selected lower order transfer function vectors from which the circuit's nonlinear currents are derived.

Arguments transmitted to CONTRL from subroutine PHASE3 are: MXORD, the number of input frequencies for the circuit (maximum order of analysis): FREQ, the input frequency values: FRTBL, the frequency number table: LOCT, the disk address of the first transfer function vector on file 20: LT, the length of each transfer function vector in words: NORDR, the number of frequencies in the previous frequency combination (order of analysis): and SEQ, the previous frequency combination vector. Arguments returned to subroutine PHASE3 are: SCODE, the new frequency combination code: NOPDR, the number of frequencies in the new frequency combination (updated order of analysis): S, the calculated angular frequency for the new frequency

combination and SEQ, the new frequency combination vector.

Upon entering subroutine CONTRL, the next frequency combination is generated by subroutine UPSEQ. If the last possible frequency combination has already been generated (NORDR.LE.0), subroutine CONTRL returns to the calling program without further processing. Otherwise NORDR (the updated order of analysis) represents the number of frequencies in the new combination which is transmitted to subroutine CONTRL in the integer vector SEQ, where the NORDR frequencies in the combination are stored according to increasing sequential frequency numbers in the first NORDR elements of SEQ.

Next the sequential frequency combination is derived from the frequency combination vector according to:

$$\begin{array}{c} \text{NORDR} \\ \text{SCOMB} = \sum_{T=1}^{\infty} \text{SEQ}(I) \end{array}$$

Subroutine SSCODE is called to calculate the frequency combination code, SCCDE, and the angular frequency, S.

Pecause linear currents are not derived from previous transfer functions, no transfer function vector table is created for first order frequency combinations. Therefore for NOPDR=1, subroutine CONTRL stores the sequential frequency combination in SEQ(11) and returns to the calling program without further processing. For higher orders of analysis (NCRDR.GT.1), the remainder of subroutine CONTRL is devoted to creating and storing

the transfer function vector table for the new frequency combination.

In the NCAP circuit analysis, the nonlinear currents for a particular frequency combination are derived from previously calculated transfer functions. In particular, an nth order nonlinear current evaluated at the frequency combination (f1,f2,...fn) is derived from the 2ⁿ-2 lower order transfer function vectors from frequency combinations which are subsets of (f1,f2,...fn). For example in a circuit with four input frequencies, the third order nonlinear currents at the frequency combination (f1,f3,f4) are derived from the 2³-2=6 transfer function vectors from all the first and second order combinations of f1, f3, and f4:

T(f1)

T(f3)

T(f4)

T(f1,f3)

T(f1,f4)

T(f3,f4)

The transfer functions vectors T(f2), T(f1,f2), T(f2,f3), T(f2,f4) do not contribute to the nonlinear current since they correspond to frequency combinations that are not subsets of (f1,f3,f4).

In subroutine CONTRL the lower-order transfer function vectors required for the nonlinear current calculations at the new frequency combination are selected from file 20 and placed in a table on file 24 for use by the nonlinear current generator

subroutines of Phase 3. The selection process involves recreating every lower-order frequency combination while maintaining the file 20 address of the corresponding transfer function vector. When a subset of the new frequency combination is encountered, the appropriate transfer function vector is read from file 20 and appended to the table on file 24. This procedure guarantees that only those transfer function vectors required for the new frequency combination are included in the table.

The lower order frequency combinations are generated by repeated calls to subroutine UPSEQ. The order of the combination is maintained in MORDR and each frequency combination is returned to subroutine CONTRL in the first MORDR elements of the integer vector SEQ1.

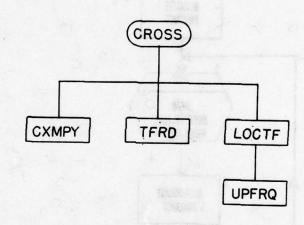
The creation of the transfer function vector table in subroutine CONTRL begins by initializing MORDR=0. The file 20 address is initialized to K=LOCT and the file 24 address is set to LL=1. Subroutine UPSEQ is called to generate a frequency combination of order MORDR, which is returned in SFQ1.

If the last combination has not been generated (MORDR.LT NORDR), subroutine CONTRL determines whether the frequency combination in SEQ1 is a subset of the combination in SEQ. If every SEQ1(I), I=1,MORDR equals some SEQ(J), J=1,NORDR, then the corresponding transfer function vector stored at K is read from file 20 by subroutine DATARD and appended to the table on file 24 by subroutine TFWR. Then the file 24 address is updated to LL=LL+LT, the file 20 address is updated to K=K+LT, and the next

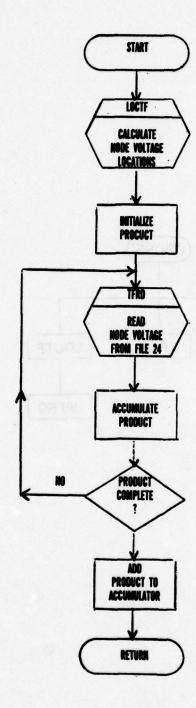
frequency combination is generated.

If SEQ1 is not a subset of SEQ, the read/write operations are by-passed and only the file 20 address is updated before generating the next frequency combination.

This process continues until MORDR=NORDR, indicating that all lower order combinations have been generated and the transfer function vector table is complete. Subroutine CONTRL then saves the sequential frequency combination in SEQ(11) and returns to the calling program.



-



NAME: CROSS

TYPE: SUBROUTINE

GENERAL PURPOSE:

Accesses voltages in T.F. vector table,
computes T.F. product for one frequency
permutation, adds product to accumulator
for symmetrization by subroutine FNCTN

VARIABLES:

FNK = Subscripts calculated in subroutine FUN

FREQC = Permuted frequency vector

II = Subscript for permuted frequency vector

K = Number of voltages in T.F. product

LCCTF = Origin of T.F. vector on file 24

N = Order of nonlinear current being calculated

NODE = Base node of circuit element

C = Node pattern of T.F. product

SUBS = Locations of voltages in T.F. vector

TS = Intermediate storage for voltage

YY = Sum of T.F. products

ZZ = T.F. product

SUPPOUTINES CALLED:

CXMPY LOCTF TFRD

CALLING PROGRAMS:

CURPT FNCTN

DESCRIPTION:

Subroutine CROSS accesses node voltages in the modified transfer function table, computes the voltage product for a single frequency permutation, and adds the product to an accumulator for use by subroutine FNCTN. Arguments transmitted to the subroutine are: FNK, the orders of the voltages in the transfer function product: Q, the node numbers (relative to the base node of the device) of the voltages in the transfer function product: FREQC, the frequency permutation vector: N, the order of the nonlinear current being calculated; K, the number of voltages in the transfer function product: and YY, the accumulated sum of transfer function products. The updated sum YY is returned as an argument to the calling program.

Subroutine CROSS begins by calculating the storage locations for each of the K voltages in the transfer function product, saving the locations in SUBS(I) for I=1,2,...K. The voltages are stored in the modified transfer function table on file 24. For an Nth order nonlinear curent calculation, the table contains 2^N -2 transfer function vectors. Each vector occupies one record and contains ORDER complex elements. The location of a particular node voltage in the transfer function table is derived by:

VOLTAGE LOCATION = T.F. ORIGIN + NODE DISPLACEMENT

First the function subprogram LOCTF calculates the record number of the transfer function vector containing the desired voltage. Then the transfer function origin is defined as:

TRANSFER FUNCTION ORIGIN = (LOCTF-1) *ORDER

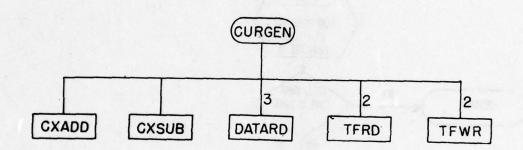
Then using O(I), the location of the Ith node voltage relative to the base node of the circuit element, and NODE, the number of the base node, the node displacement is defined as:

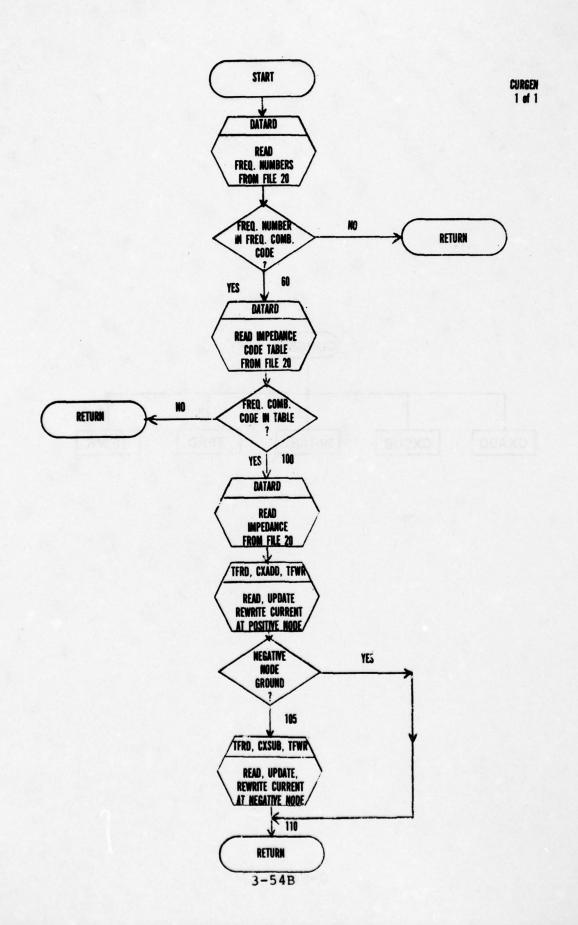
NODE DISPLACEMENT = Q(I) + NODE - 1

The complete expression for the location of the Ith voltage in the transfer function product is:

SUBS(I) = ((LOCTF-1)*ORDER)+Q(I)+NODE-1

After all of the K voltage locations have been calculated, the transfer function product is formed. A voltage is read from file 24 by subroutine TFRD and stored in TS. The product is accumulated in ZZ by subroutine CXMPY. When all K voltages have been accessed and multiplied, the resulting product is added to the accumulator YY and subroutine CROSS returns to the calling program.





NAME: CURGEN

TYPE: SURROUTINE

CENERAL PURPOSE:

Calculates linear current for generator and places it in current vector

VARIABLES:

A = Linear current (impedance)

CUR = Current vector element

FNC = User-assigned frequency numbers

I = Relative location of impedance on file 20

IMTPL = Impedance code table

MFREC = Number of frequencies

MIMPS = Number of impedances

NCDE = Positive node of connection

MCDEP = Negative node of connection

SCODE = Frequency combination code

SUPROUTINES CALLED.

CXACD CXSUE PATARD TFRD TFWR

CALLING PROGRAMS.

PHASE3

DESCRIPTION:

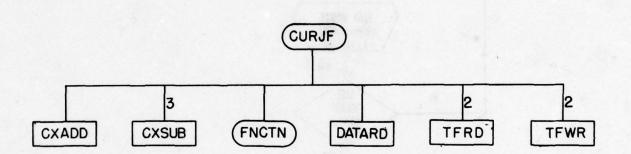
Subroutine CURCEN calculates the current for the generator at a given frequency combination and places it in the current vector. The frequency combination code SCODE, for which the current is to be calculated, is transmitted to subroutine CURCEN as an argument. The generator is an independent voltage source which generates linear current defined as the generator's impedance at the given frequency combination. The impedances are stored in the generator data record together with a table of corresponding frequency combination codes IMTPL where IMTPL(I) is the frequency combination for the Ith impedance in the data record. The number of input frequencies and impedances for the generator, NFREQ and NIMPS, and the nodes of connection, NCDE and NODEB, are stored in the generator driver record in global common.

First the user-assigned frequency numbers are read from the generator data record on file 20 by subroutine DATARD and stored in the integer vector FNC. If the frequency number encoded in SCODE is not among the frequencies FNO defined for the generator, subroutine CURGEN returns to the calling program without further processing.

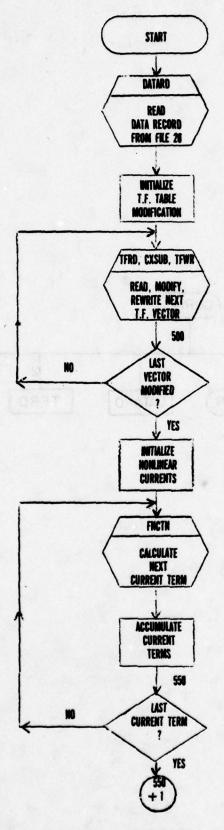
from the generator data record on file 20 by subroutine DATARD and stored in IMTEL. Then SCODE is compared against every frequency combination code in the table, IMTEL(I), I=1,2,...NIMPS. If no match is found, subroutine CURGEN returns to the calling program. If a match is found, the table position

I is used to locate the desired impedance in the data record. The impedance is then read from file 20 by subroutine DATARD and stored in A.

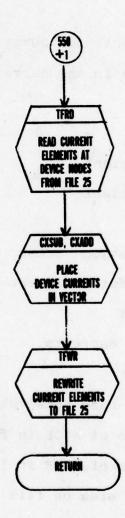
The current is placed in the current vector on file 25 by adding A to the current at the positive node of connection, NODE, and subtracting A from the current at the negative node of connection, NODER. The NODE current, located at I=2*NODE-1, is read from file 25 by subroutine TFRD and stored in CUR. A is added to CUR by subroutine CXADD and the updated current is rewritten to its original location by subroutine FTWR. If the negative node of connection is ground (NODER=0), subroutine CURGEN returns to the calling program. Otherwise the NODER current, located at I=2*NODER-1, is read from file 25 by subroutine TFRD and stored in CUR. A is subtracted from CUR by subroutine CXSUE, the updated current is rewritten to its original location, and subroutine CURCEN returns to the calling program.







3-55B



NAME: CURJF

TYPE: SUPROUTINE

GENERAL PURPOSE:

Calculates nonlinear currents for Junction Field Effect
Transistor and places them in the current vector

VARIABLES:

A = Nonlinear coefficient

C = Nonlinear coefficient

CGD = Input value

FREQ = Frequency combination vector

GG = T.F. expression for nonlinear current

HG = Nonlinear current

HH = Sum of nonlinear currents

HL = Nonlinear current

I = Record number of base node voltage (current)

Il = Relative location of NODE in first T.F.

12 = Relative location of NODE in last T.F.

I3 = Location of work area on file 24

K = Index for summation of current calculation

E = Record number for file 24 work area

IT = Length of each T.F. vector in words

NODE = Base node of connection

NORDR = Order of analysis

- S = Angular frequency
- X = Base node voltage (current)
- X1 = (Base+1) node voltage (current)
- x2 = (Base+2) node voltage (current)

SUBROUTINES CALLED:

CXADD CXSUB DATARD FNCTN

TFRD TFWR

CALLING PROGRAMS .

PHASE3

DESCRIPTION .

Subroutine CURJF calculates the nonlinear currents for the junction field effect transistor and places them in the current vector. Arguments transmitted to the subroutine are: NORDR, the order of the nonlinear currents to be calculated and S, the angular frequency. The nonlinear currents are functions of the device coefficients calculated in Phase 1 and the node voltages derived from the lower order transfer functions stored on file 24. The base node of connection, NODE, is stored in the JFET driver record in global common. The remaining nodes are numbered sequentially by the system as NODE+1, NODE+2, and NODE+3.

First the JFET data record is read from file 20 by subroutine DATARD and stored in the common data buffer BUFF. The node voltages for the nonlinear current calculations are computed by replacing the transfer function at NODE+2 with the difference

between the transfer functions at NODE and NODE+2 for every vector in the table. This "modified transfer function table" serves as the basis for the nonlinear current calculations.

The transfer function table consists of 2 NORDR-2 contiguous vectors beginning at record number 1 on file 24. Each complex vector element is the transfer function at some node in the circuit and occupies two decimal storage locations. The transfer functions are stored in the vectors according to increasing node numbers. Because the JFET nodes are numbered sequentially, the corresponding transfer functions occupy successive elements in each vector. This method of storage allows the transfer functions to be accessed and modified by the following iterative procedure:

STEP 1. Initialization

- a) Length of each T.F. vector: LT=2*ORDER
- b) Location of NODE in first vector: I1=2*NODE-1
- c) Location of NODE in last vector: I2=I1+(2 NORDR-3)*LT
- d) Origin of scratch area on file 24: L = 2 NORDR * LT
- e) Set index at first vector: I=Il

STEP 2. Access and Modify Next Vector

a) Read transfer functions at sequential device nodes

beginning at I, store in X, X1, and X2.

- b) Write device voltages to scratch area at L, increment L=L+2.
- c) Calculate X2=X-X2
- d) Write updated transfer functions to file 24 begin-

STEP 3: Increment and Test Index

- a) Increment I=I+LT
- b) If (I.LE.I2) more vectors remain to be modified. Go to Step 2.
- c) Otherwise modified transfer function table complete.

The nonlinear currents of order NORDR for the JFET are defined as:

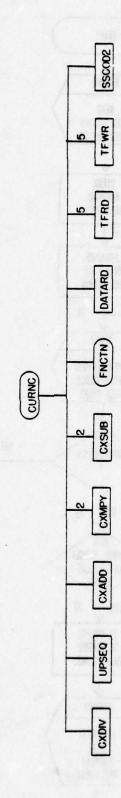
NORDR
HG =
$$\Sigma$$
 A(K) * F_K^{NORDR} (NODE + 3)
K=2

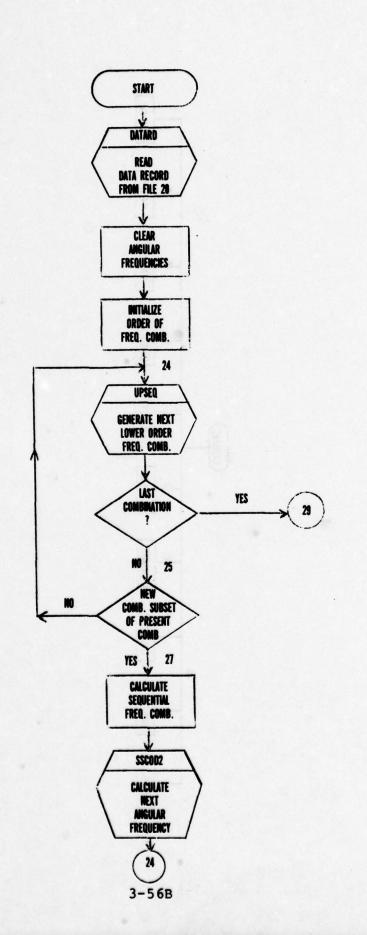
NORDR
HL =
$$\Sigma$$
 S * C(K) * F_K^{NORDR} (NODE + 3)

Subroutine CURJF initializes HG and HL to zero, then each NORDR symmetric transfer function expression F_K (NODE+3) for K=2,3,... NORDR is computed by subroutine FNCTN, multiplied by the appropriate coefficients A(K) or S*C(K), and accumulated in HG and HL.

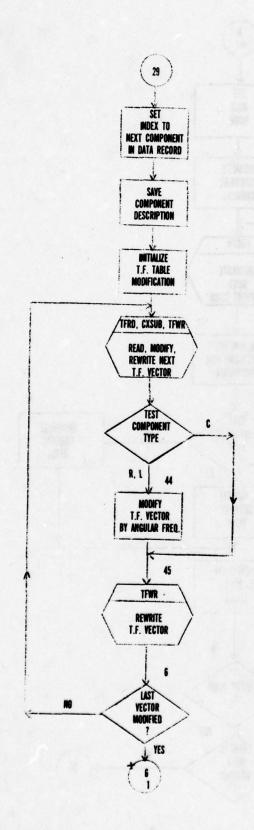
Then the transfer function table is restored by reading the transfer functions saved in the scratch area in file 24 and rewriting them to the original locations in the table. The procedure is similar to that employed for modifying the transfer function table.

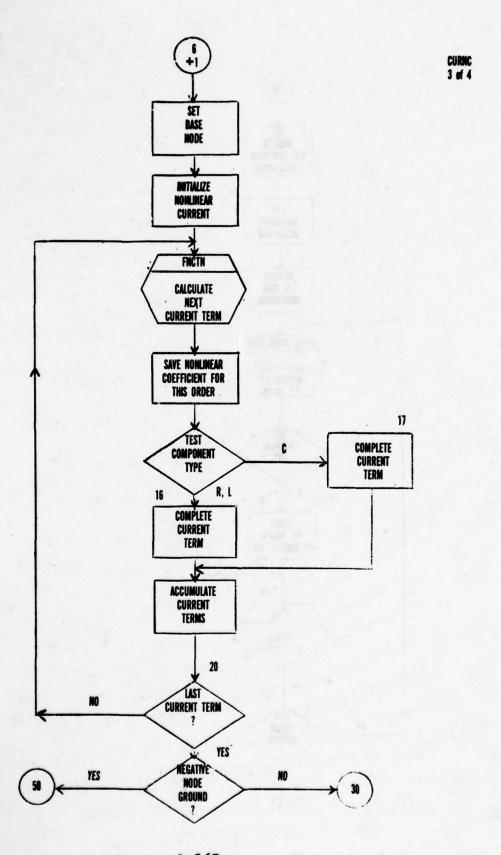
The nonlinear currents are placed in the current vector on file 25 by subtracting HL from the current at NODE, subtracting HG from the current at NODE+1, and adding HG and HL to the current at NODE+2. The current elements for the sequential JFET nodes, beginning at I=2*NODE-1, are read from file 25 by subroutine TFRD and stored in X, X1, and X2. HL is subtracted from X and HG is subtracted from X1 by subroutine CXSUB. HG and HL are added to X2 by subroutine CXADD. The updated current elements are rewritten to their original locations by subroutine TFWR and subroutine CURJF returns to the calling program.

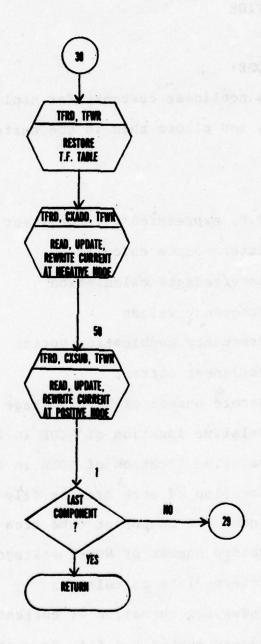




CURNC 1 of 4







NAME: CURNC

TYPE: SUBROUTINE

GENERAL PURPOSE:

Calculates nonlinear currents for nonlinear components and places them in the current vector

VARIABLES:

A = T.F. expression ror nonlinear current

E = Intermediate calculation

CONST = Intermediate calculation

FREQ = Frequency values

FRQ = Frequency combination vector

HK = Nonlinear current

I = Record number of NODE voltage (current)

Il = Relative location of NODE in first T.F.

12 = Relative location of MODE in last T.F.

I3 = Location of work area on file 24

I4 = Index for component data area

J = Record number of NODEE voltage

J2PI = Intermediate calculation

F = Index for summation of current calculation

I = Pecord number for file 24 work area

LT = Length of each T.F. vector in words

MCCMP = Length of nonlinear component data record

NODE = Positive node of connection

NODEB = Negative node of connection

NORDR = Order of analysis

RCL = Component type; 1 = Resistor

The total wash nonement and a 2 = Capacitor as areas well-units as

a a saladausus 3 = Inductor sawoi da aidad a mada

S = Angular frequency

SP = Frequency product

X = NODE voltage (current)

Y = NODEB voltage (current)

SUPROUTINES CALLED:

CXADD CXDIV CXMPY CXSUB DATARD

FNCTN SSCOD2 TFRD TFWR UPSEQ

CALLING PROGRAMS.

PHASE3

DESCRIPTION: The State of the S

Subroutine CURNC calculates the nonlinear currents for nonlinear components and places them in the current vector.

Arguments transmitted to the subroutine are: NORDR, the order of the nonlinear current to be calculated; FREQ, the input frequency values for the circuit; FRTBL, the frequency number table: SEQ, the frequency combination vector: and S, the angular frequency. The nonlinear current is a function of the device coefficients stored in the nonlinear components data

record, the node voltages derived from the lower order transfer functions stored on file 24, and, in the case of the nonlinear inductor, the angular frequency at the lower orders of analysis.

First the nonlinear components data record is read from file

20 by subroutine DATARD and stored in the common data buffer

BUFF. Then a table of lower order angular frequencies is created

for each vector in the transfer function table. The method is to

recreate every lower order frequency combination which is a

subset of the present frequency combination vector SEC, calculate

the angular frequency for that combination, and store it in the

vector FPTAB. The frequency combinations are recreated in the

same order as the transfer function vectors were selected and

stored in the transfer function table by subroutine CONTRL, so

that the angular frequency stored at FPTAP(I) corresponds to the

Ith transfer function vector on file 24.

The recreation of lower order frequency combinations is performed by repeated calls to subroutine UPSEQ under the control of the integer variable MORDR, which maintains the order of the combination being formed. Subroutine SSCOD2 calculates the angular frequency and appends it to the FRTAB vector. The complete procedure for creating the angular frequency table is summarized as follows:

STEP 1. Initialization

a) Initialize FRTAP subscript to LL=0.

- b) Initialize maximum order of frequency combination required to MXCRD=NCRDR.
- c) Initialize number of frequencies in lower order combination to MCPDR=0.

STEP 2 Generate Mext Frequency Combination

- a) Call subroutine UPSEC to generate next frequency combination vector in SEC1.
- b) If (MCRDP.IT.MCRDP), go to Step 3.
- c) Ctherwise angular frequency table complete.

STEP 3 Test for SECl a Subset of SEC

- a) If every SEQ1(I) for I=1,2,...MORDR equals some SEQ(J) for J=1,2,...MOPDR, go to Step 4.
- b) Otherwise go to Step 2.

STEP 4: Calculate Angular Frequency and Append to Table

a) Calculate sequential frequency combination by:

$$\begin{array}{c} \text{MORDR} \\ \text{SCOMB} = \sum_{\mathbf{I}=1}^{N} \text{SEQl}(\mathbf{I}) \end{array}$$

- b) Increment FRTAB index to LL=LL+1.
- c) Call subroutine SSCOD2 to calculate angular frequency and append to table at FRTAB(LL).
- d) Go to Step 2.

For each nonlinear component description in the data record, subroutine CURNC transfers the component type to RCL and the positive and negative nodes of connection to NODE and NODEB respectively. Then the node voltages for the nonlinear current calculations are computed by replacing the transfer function at NODE with the difference between the transfer functions at NODE and NODEB for every vector in the transfer function table. For nonlinear inductors, the updated transfer function at NODE is divided by the corresponding angular frequency from the table FRTAB. This "modified transfer function table" serves as the basis of the nonlinear current calculations. If the negative node of connection is ground (NODEP=0) for any component in the data record, the modification of transfer functions for that component is by-passed.

The transfer function table consists of 2 NORDR -2 contiguous vectors beginning at record number 1 on file 24. Each complex vector element is the transfer function at some node in the circuit and occupies two decimal storage locations. The transfer functions are stored in the vectors according to increasing node numbers. This method of storage allows the transfer functions to

be accessed and modified by the following iterative procedure:

STEP 1. Initialization

- a) Length of each T.F. vector: LT=2*ORDER
- b) Location of NCDE in first vector: Il=2*NODE-1
- c) Location of NODE in last vector: I2=I1+(2 NORDR_3)*LT
- d) Origin of scratch area on file 24: L=2 *LT
- e) Subscript for FRTAB vector: M=0
- f) Set index at first vector: I=Il

STEP 2: Access and Modify Next Vector

- a) Increment FRTAB subscript M=M+1
- b) Read transfer function from I, store in X.
- c) Write X to scratch area at L, increment L=L+2.
- d) If NODEB=0, go to Step 2h.
- e) Otherwise locate NODEB in vector: J=I+(NODEB-NODE) *2
- f) Read transfer function from J, store in Y.
- g) Calculate X=X-Y.
- h) If component is inductor (RCL=3), calculate X=X/FRTAB(M)
 Otherwise go to Step 3i.
- i) Write X to file 24 at I.

STEP 3: Increment and Test Index

- a) Increment I=I+LT for next vector
- b) If (I.LE.I2), more vectors remain to be modified. Go to Step 2.
- c) Otherwise modified transfer function table complete.

The nonlinear current of order NORDR for the nonlinear resistor (RCL=1) is defined as:

HK =
$$\Sigma$$
 CONST * F_K^{NORDR} (NODE)

For the nonlinear capacitor (RCL=2):

NORDR
$$HK = \sum_{K=2}^{NORDR} CONST * S/K * F_K^{NORDR} (NODE)$$

For the nonlinear inductor (RCL=3):

where CONST is the Kth nonlinear coefficient taken from the

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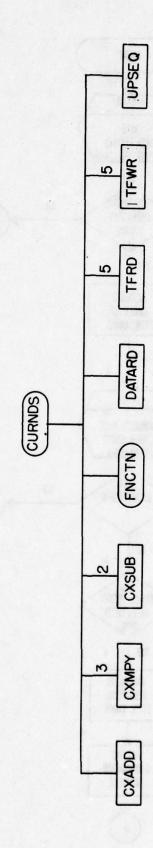
component description in BUFF.

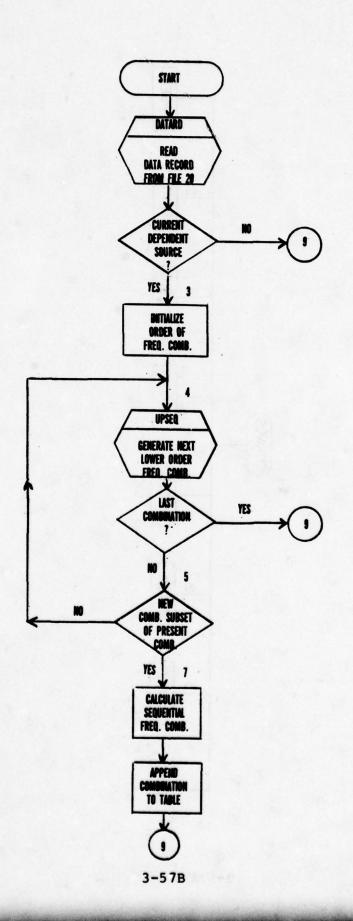
Subroutine CURNC initializes HK to zero, then each symmetric transfer function expression $F_K^{NORDR}(\overline{NODE})$ for K=2,3...NORDR is computed by subroutine FNCTN. After the intermediate calculation for CCNST has been performed depending on the component type, the product is formed and accumulated in HK.

Then the transfer function table is restored by reading the transfer functions saved in the scratch area on file 24 and rewriting them to their original locations in the table. The procedure is similar to that employed for modifying the transfer function table. If the negative node of connection is ground (NCDEP=0), the restore procedure is by-passed.

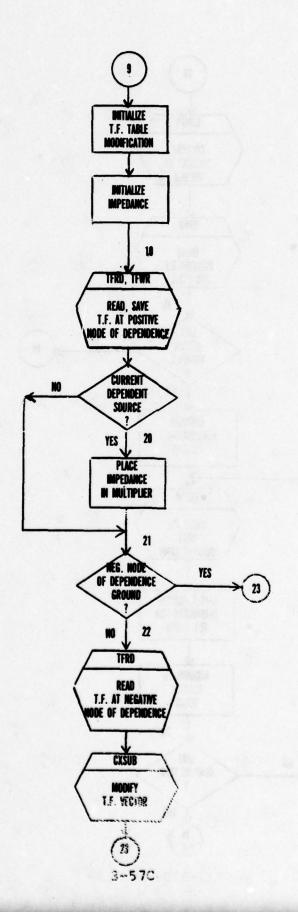
The nonlinear current is placed in the current vector on file 25 by adding it to the current at NODEP (for NODEB #0) and subtracting it from the current at NODE. The NODEB current, located at I=2*NCDEE-1, is read from file 25 by subroutine TFRD and stored in X. HK is added to X by subroutine CXADD and the updated NODEP current is rewritten to its original location by subroutine TFWR. The procedure is repeated for the NODE current, except that the NODE current is located at I=2*NODE-1 and HK is subtracted from X by subroutine CXSUE.

After the nonlinear currents have been calculated and placed in the current vector for every component in the data record, subroutine CURNC returns to the calling program.

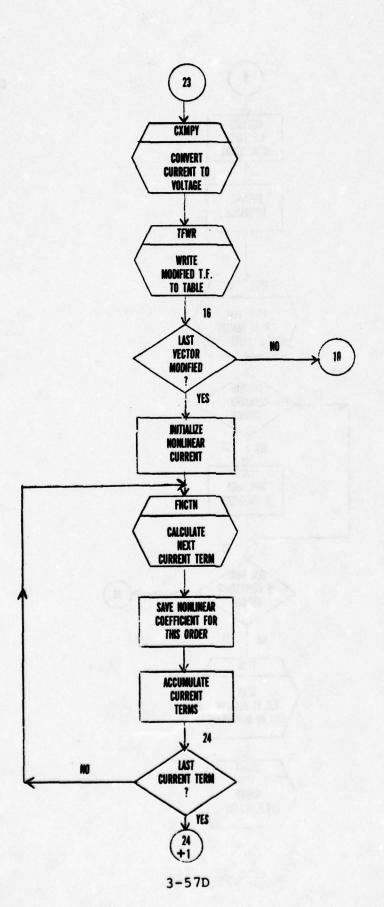




CURNOS 1 of 4

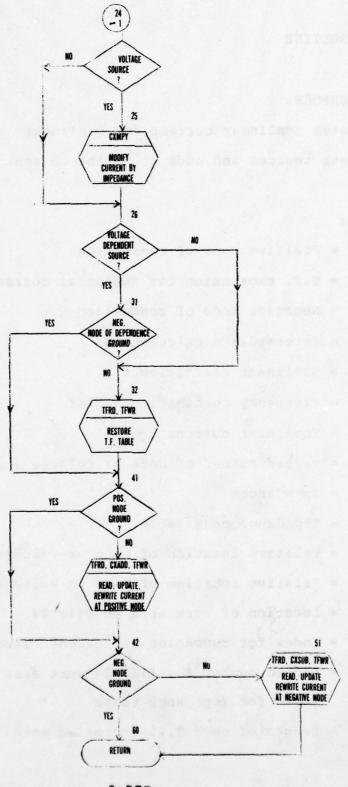


CURNOS 2 of 4



CURNOS 3 of 4

0



3-57E

NAME: CURNDS

TYPE: SUBROUTINE

CENERAL PURPOSE:

Calculates nonlinear current for nonlinear dependent sources and adds it to the current vector

VARIABLES:

A = Positive node of connection

AA = T.F. expression for nonlinear current

B = Negative node of connection

CC = Intermediate calculation

COEF = Nonlinear coefficients

FREQ = Frequency combination vector

HK = Nonlinear current

I = Record number of node xx voltage (current)

IMPS = Impedances

IMTBL = Impedance code table

Il = Relative location of node xx voltage in first T.F.

12 = Relative location of node xx voltage in last T.F.

I3 = Location of work area on file 24

K = Index for summation of current calculation

L = Record number for file 24 work area and
index for impedance table

LT = Length of each T.F. vector in words

M = Index for impedance code table

MORDR = Order of impedance

MXORD = Maximum order of analysis

NORDR = Order of analysis

S = Angular frequency

SCOMP = Sequential frequency combination

SEQ = Frequency combination vector

SEQ1 = Lower-order frequency combination vector

SOURCE = Type of source: 1 = VC

2 = CC

3 = VV

4 = CV

X = Node XX voltage (current)

XX = Positive node of dependence

Y = Node YY voltage (current)

YY = Negative node of dependence

Z = Impedance value

SUBROUTINES CALLED:

CXADD CXMPY CXSUB DATARD FNCTN

TFRD TFWP UPSEC

CALLING PROCRAMS:

PEASE3

DESCRIPTION .

Subroutine CUPNDS calculates the nonlinear currents for

nonlinear dependent sources and places them in the current vector. Arguments transmitted to the subroutine are: NORDR, the order of the nonlinear current to be calculated; S, the angular frequency: and SEO, the frequency combination vector. The nonlinear current is a function of the device coefficients stored in the nonlinear dependent source data record and the node voltages derived from the lower order transfer functions stored on file 24. For current dependent sources, the nonlinear current also depends on the impedance at the nodes of dependence, while for voltage sources, the nonlinear current depends on the impedance at the nodes of connection. The positive and negative nodes of dependence (XX and YY), the positive and negative nodes of connection (A and B), and the type of SCURCE are stored in the nonlinear dependent source driver record in global common.

First the nonlinear dependent source data record is read from file 20 by subroutine DATARD and stored in the common data buffer EUFF. Then for current dependent sources only (SOURCE=2 or 4), a table of lower order sequential frequency combinations is created for each vector in the transfer function table. The method is to recreate every lower order frequency combination which is a subset of the frequency combination vector SEC, calculate the sequential frequency combination, and store it in the vector IMTEL. The frequency combinations are recreated in the same order as the transfer function vectors were selected and stored in the transfer function table by subroutine CONTPL, so that the sequential frequency combination IMTPL(I) corresponds to the Ith transfer function vector on file 24.

The recreation of lower order frequency combinations is performed by repeated calls to subroutine UPSEQ under the control of the integer variable MORDR, which maintains the order of the combination being formed. The complete procedure for creating the sequential frequency combination table is summarized as follows:

STEP 1. Initialization

- a) Initialize IMTEL index to L=0.
- b) Initialize maximum order of frequency combination required to MXORD=NORDR.
- c) Initialize number of frequencies in lower order combination to MORDR=0.

STEP 2: Generate Next Frequency Combination

- a) Call subroutine UPSEC to generate next frequency combination vector in SEC1.
- b) If (MORDR.LT.NORDR), go to Step 3.
- c) Otherwise sequential frequency combination table complete.

STEP 3. Test for SEQ1 a Subset of SEQ

- a) If every SEQ1(I) for I=1,2,...MORDR equals some SEQ(J) for J=1,2,...NORDR, go to Step 4.
- b) Otherwise go to Step 2.

STEP 4: Calculate Sequential Frequency Combination and Append to Table

a) Calculate sequential frequency combination by:

$$\begin{array}{ccc} & \text{MORDR} \\ \text{SCOMB} & = & \Sigma & \text{SEQ1(I)} \\ & & \text{I=1} \end{array}$$

- b) Increment IMTBL index to L=L+1.
- c) Append sequential frequency combination to table by IMTBL(L)=SCOMB.
- d) Go to Step 2.

Then the node voltages for the nonlinear current calculations are computed by replacing the transfer function at node XX with the difference between the transfer functions at nodes XX and YY for every vector in the transfer function table. For current dependent sources (SOURCE=2 or 4), the updated transfer function vector at node XX is divided by the corresponding impedance IMPS(IMTBL(M)) from the nonlinear

dependent source data record. This "modified transfer function table" serves as the basis for the nonlinear current calculations. If the negative node of dependence is ground (YY=0), the subtraction of transfer functions is by-passed.

The transfer function table consists of 2NORDR-2 contiguous vectors beginning at record number 1 on file 24. Each complex vector element is the transfer function at some node in the circuit and occupies two decimal storage locations. The transfer functions are stored in the vectors according to increasing node numbers. This method of storage allows the transfer functions to be accessed and modified by the following iterative procedure:

STEP 1: Initialization

- a) Subscript for IMTBL vector: M=0.
- b) Location of node XX in first vector: I1=2*XX-1.
- c) Length of each T.F. vector: LT=2*ORDER
- d) Location of node XX in last vector I2=I1+(2 NORDR-3)*LT.
- e) Origin of scratch area on file 24: $L=2^{NORDR} *LT$.
- f) Impedance value: Z=(1., 0.)
- g) Set index at first vector: I=Il.

STEP 2: Access and Modify Next Vector

- a) Read transfer function from I, store in X.
- b) Increment IMTBL subscript M=M+1
- c) Write X to scratch area at L, increment L=L+2.
- d) For current dependent sources (SOURCE=2 or 4), store
 impedance value Z=IMPS(IMTBL(M)).
- e) If node YY=0, go to Step 2i.
- f) Otherwise locate node YY in vector: J=I+(YY-XX)*2.
- g) Pead transfer function from J, store in Y.
- h) Calculate X=X-Y
- i) Calculate X=X/Z
- j) Write X to file 24 at I.

STEP 3: Increment and Test Index

- a) Increment I=I+LT
- b) If (I.LE.I2), more vectors remain to be modified. Co to Step 2.
- c) Otherwise modified transfer function table complete.

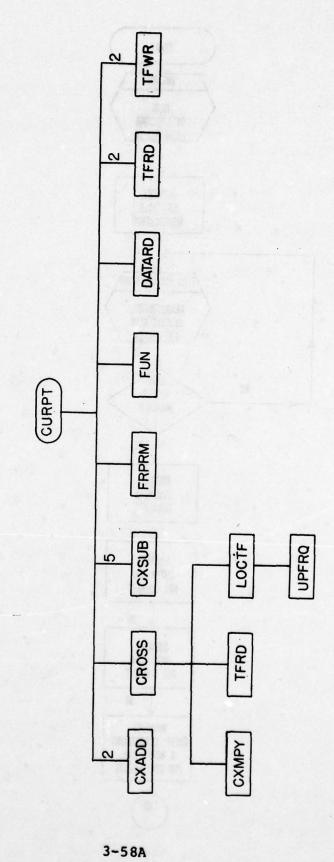
The nonlinear current of order NCRDR for the nonlinear dependent sources is defined as:

$$\begin{array}{lll}
\text{NORDR} \\
\text{HK} &= \sum_{K=2} & \text{COEF}(K) * F_{K}^{\text{NORDR}} & \overline{(XX)}
\end{array}$$

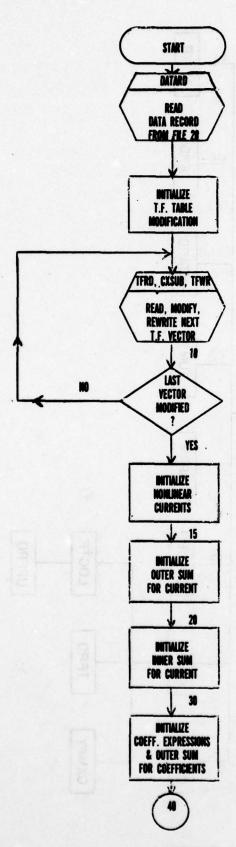
Subroutine CURNDS initializes HK to zero, then each symmetric transfer function expression FNORDR (\overline{XX}) for K=2,3,...NORDR is computed by subroutine FNCTN, multiplied by the appropriate coefficient COEF(K), and accumulated in HK. For voltage sources only (SOURCE=3 or 4), HK is divided by the impedance at the nodes of connection, Z = (1.0E10., 0.).

Then the transfer function table is restored by reading the transfer functions saved in the scratch area on file 24 and rewriting them to their original locations in the table. The procedure is similar to that employed for modifying the transfer function table. If the negative node of dependence is ground (YY=0), the restore procedure is by-passed.

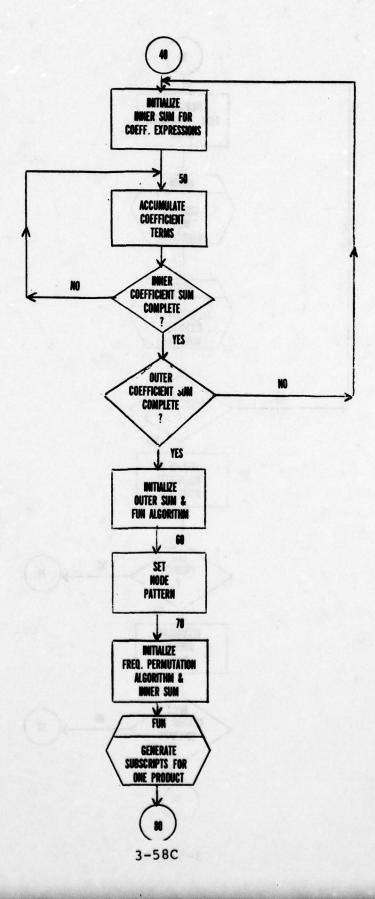
The nonlinear current is placed in the current vector on file 25 by adding it to the current at the positive node of connection A (for A≠0) and subtracting it from the negative node of connection B (for B≠0). The node A current, located at I=2*A-1, is read from file 25 by subroutine TFRD and stored in X. HK is added to X by subroutine CXADD and the updated node A current is rewritten to its original location by subroutine TFWR. The procedure is repeated for the node B current, except that node P is located at I=2*B-1 and HK is subtracted from X by subroutine CXSUE. Subroutine CURNDS then returns to the calling program.



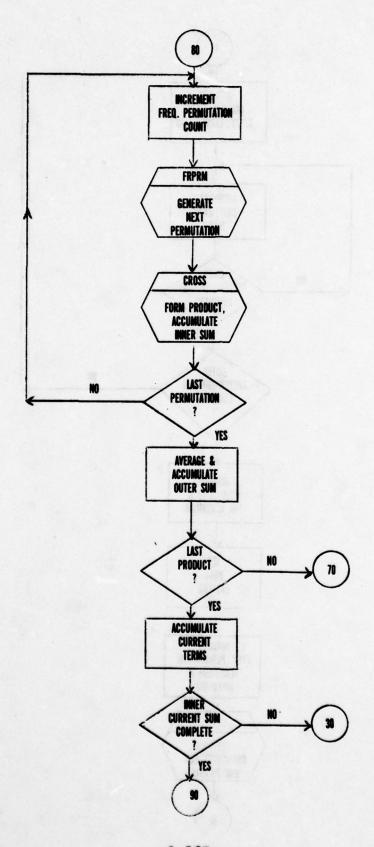


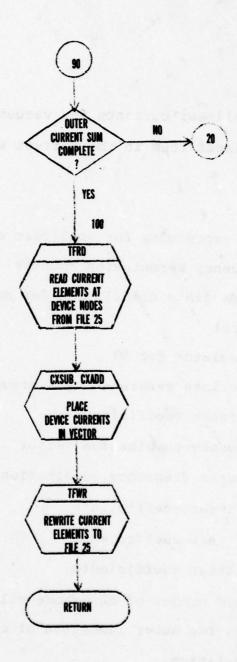












SYRACUSE UNIV N Y
NONLINEAR CIRCUIT ANALYSIS PROGRAM (NCAP) DOCUMENTATION. VOLUME--ETC(U)
SEP 79 J B VALENTE , S STRATAKOS F30602-79-C-0011 AD-A076 317 UNCLASSIFIED RADC-TR-79-245-VOL-3 NL 6 OF 8 AD 46317 1. * NAME: CURPT

TYPE: SUBROUTINE

GENERAL PURPOSE:

Calculates nonlinear currents for vacuum pentode and places them in the current vector

VARIABLES.

A = T.F. expression for nonlinear current

CNT = Frequency permutation counter

COEF = Intermediate calculation for nonlinear current

FACT = Accumulator for N!

FNK = Subscripts generated by subroutine FUN

FP = Nonlinear coefficients

FREQ = Frequency combination vector

FREQC = Permuted frequency combination vector

F2 = Nonlinear coefficients

G = Nonlinear coefficients

H = Nonlinear coefficients

I = Record number of base node voltage (current):
 index for outer summation of current
 calculation

IHIGH = Pelative location of NOTE in last T.F.
 vector

ILCW = Relative location of NODE in first T.F.

vector

INC = Binomial coefficient for I,N

IP = Nonlinear current

ISWF = Control variable for subroutine FRPRM

ISWT = Control variable for subroutine FUN

12 = Nonlinear current

13 = Location of work area on file 24

J = Index for inner summation of current calculation

JNIC = Trinomial coefficient for J, N-I

L = Index for QQ vector. Record number for file
24 work area.

LT = Length of each T.F. vector in words

Ll = Subscript for coefficient calculation

L2 = Subscript for coefficient calculation

L3 = Subscript for coefficient calculation

MU = Input value

N = Order of T.F. expression

NCDE = Base node of connection

NORDR = Order of analysis

PROD = Sum of T.F. products returned by subroutine
CROSS

CQ = Node pattern of T.F. expression

P = Index for inner summation of coefficient expression

PJC = Pinomial coefficient for R,J

- S = Index for outer summation of coefficient expression
- SIC = Binomial coefficient for S,I
- SP = Coefficient expression for IP
- S2 = Coefficient expression for 12
- TERM = Intermediate calculation for coefficient expression
- X = Base node voltage (current)
- XX = Angular frequency
- x1 = (Base+1) node voltage (current)
- x2 = (Base+2) node voltage (current)
- x3 = (Base+3) node voltage (current)

SUBROUTINES CALLED:

CROSS CXADD CXSUB DATARD

FRPRM FUN TFRD TFWR

CALLING PROGRAMS:

PHASE3

DESCRIPTION:

Subroutine CURPT calculates the nonlinear currents for the vacuum pentode and places them in the current vector. Arguments transmitted to the subroutine are: NORDR, the order of the nonlinear current to be calculated and XX, the angular frequency. The nonlinear currents are functions of the device parameters and coefficients calculated in Phase 1 and the node voltages derived

from the lower order transfer functions stored on file 24. The base node of connection, NODE, is stored in the vacuum pentode driver record in global common. The remaining nodes are numbered sequentially by the system as NODE+1, NODE+2, and NODE+3.

First the vacuum pentode data record is read from file 20 by subroutine DATARD and stored in the common data buffer BUFF. Then the node voltages for the nonlinear current calculations are computed for each vector in the transfer function table as follows: the transfer function at NODE is replaced with the difference between the transfer functions at NODE and NODE+3; the transfer function at NODE+1 is replaced with the difference between the transfer functions at NODE+1 and NODE+3; and the transfer function at NODE+2 is replaced with the difference between the transfer functions at NODE+2 and NODE+3. This "modified transfer function table" serves as the basis for the nonlinear current calculations.

The transfer function table consists of 2 -2 contiguous vectors beginning at record number 1 on file 24. Each complex vector element is the transfer function at some node in the circuit and occupies two decimal storage locations. The transfer functions are stored in the vector according to increasing node numbers. Because the vacuum pentode nodes are numbered sequentially, the corresponding transfer functions occupy successive elements in each vector. This method of storage allows the transfer functions to be accessed and modified by the following iterative procedure

STEP 1: Initialization

- a) Location of NODE in first vector: ILOW=2*NODE-1.
- b) Length of each T.F. vector: LT=2*ORDER.
 - c) Location of NODE in last vector:

 IHIGH = ILOW + (2 NORDR -3) * LT
 - d) Origin of scratch area on file 24: I3=2NORDR*LT
- e) Set index to first vector: I=ILOW.

STEP 2: Access and Modify Next Vector

- a) Read transfer functions for sequential device nodes beginning at I, store in X, X1, X2, and X3.
- b) Write device voltages to scratch area at L, increment L = L + 2.
 - c) Calculate X=X-X3, X1=X1-X3, and X2=X2-X3.
 - d) Write updated transfer functions to file 24 beginning at I.

STEP 3: Increment and Test Index

- a) Increment I=I+LT.
- b) If (I.LE.IHIGH), more vectors remain to be modified.
 Go to Step 2.
- c) Otherwise modified transfer function table complete.

The nonlinear currents of order NORDR for the vacuum pentode are defined as:

$$12 = \sum_{N=2}^{NORDR} \sum_{I=0}^{N} \sum_{J=0}^{N-I} \frac{1}{N!} {N \choose I} {N-I \choose J} S2_{N,I,J} F_{N}^{NORDR} Q[\overline{NODE,NODE+1,NODE+2}]$$

where S2_{N,I,J}is:

$$S2_{N,I,J} = \sum_{\substack{S = MAX \\ (I=1,0)}}^{I} \sum_{R=0}^{J} \binom{I}{S} \binom{J}{R} G(I-S+1) H(J-R+S+1) MU^{R-J} F2(N-I-J+1,R+1)$$

and .

$$IP = \begin{array}{cccc} & NORDR & N & N-I \\ & \Sigma & \Sigma & \Sigma & \Sigma \\ & N=2 & I=0 & J=0 & \overline{N}! & \begin{pmatrix} N \\ I \end{pmatrix} \begin{pmatrix} N-I \\ J \end{pmatrix} & SP_{N,I,J} & F_N^{NORDR} & Q[\overline{NODE,NODE+1,NODE+2}] \\ \end{array}$$

where SPN,I,Jis: The state of t

$$SP_{N,I,J} = \sum_{S=}^{I} \sum_{R=0}^{J} \binom{I}{S} \binom{J}{R} G(I-S+1) H(J-R+S+1)MJ^{R-J} FP(N-I-J+1,R+1) MAX(I-1,0)$$

Because the nonlinear currents for the vacuum pentode are functions of the voltages at three nodes, the Q function of subroutine QFN cannot be used. Subroutine CURPT contains its own Q function to determine the pattern of node voltages to be multiplied together to form one term of the symmetric transfer function expression F_N O(NODE,NODE+1,NODE+2). Subroutine CURVT apportions the relative node numbers to the integer vector QQ so that each product in the symmetric transfer function expression contains exactly N lower order voltages where the first I are from the base NODE, the next J are from NODE+2, and the remaining N-(I+J) are from NODE+1.

$$OQ = (1, 1, ... 1, 3, 3, ... 3, 2, 2, ... 2)$$
 I
 J
 $N - (I + J)$

Voltages Voltages

Subroutine FNCTN, which is used to calculate the terms of the symmetric transfer function expressions for NCAP's other nonlinear devices, calls subroutine QFN and therefore cannot be be used directly for the vacuum pentode. However, the algorithm of subroutine FNCTN does apply to the pentode's nonlinear currents and is implemented directly as a part of subroutine CURPT. The algorithm uses subroutine FUN to determine the orders of the voltages in each product, subroutine FRPRM to permute the frequency combination, and subroutine CROSS to access the voltages in the modified transfer function table and to form the

voltage product for a single frequency permutation. The complete algorithm used by subroutine CURPT to calculate the nonlinear currents of order NORDR is summarized as follows:

STEP 1: Initialize Summation Over N

- a) Clear nonlinear currents I2=0. and IP=0.
- b) Initialize N! accumulator to FACT=1
- c) Set index at N=2

STFP 2. Initialize Summation Over I

- a) Accumulate N! in FACT
- b) Set index at I=0
- c) Initialize accumulator for $\binom{N}{I}$ to INC=1

STEP 3. Initialize Summation Over J

- a) Set index at J=0
- b) Initialize accumulator for $\binom{N-1}{J}$ to JNIC=1

STEP 4: Initialize Algorithm for F_N

- a) Calculate coefficient expressions S2 and SP
- b) Initialize sum of symmetric voltage products to A=0.
- c) Initialize FUN algorithm by ISWT=-1
- d) Establish node pattern in QQ vector

STEP 5. Initialize for one Additive Term of F_N^{NORDR}

- a) Initialize FRPRM algorithm by ISWF=-1
- b) Clear frequency permutation count to CNT=0.
- c) Initialize symmetric accumulator to PROD=0.
- d) Call subroutine FUN to determine pattern of voltage orders, return orders in FNK vector.

STEP 6: Accumulate Products over all Permutations

- a) Increment permutation count CNT=CNT+1.
- b) Call subroutine FRPPM to create frequency permutation, return in FREQC vector
- c) Call subroutine CROSS to calculate voltage product for given permutation, add to PROD
- d) If (ISWF=0), more permutations remain: go to Step 6a.

Otherwise permutations complete. Go to Step 7.

STEP 7: Test for F_N at I, J Complete

a) Average over all permutations, add symmetric product to A:

$$A = A + PROD/CNT$$

- b) If (ISWT=0), more symmetric products remain to be calculated. Go to Step 5.
- c) Otherwise F_N complete. Accumulate current terms according to:

$$12 = 12 + \frac{1}{N!} {N \choose I} {N-I \choose J} * S2 * A$$

$$IP = IP + \frac{1}{N!} {\binom{N}{I}} {\binom{N-I}{J}} \qquad * SP * A$$

STEP 8: Increment and Test Summation Over J

- a) Accumulate $\binom{N-1}{J}$ in JNIC
- b) Increment J=J+1
- c) If (J.LE.(N-I)), go to Step 4. Otherwise summation over J complete. Go to Step 9.

STEP 9: Increment and Test Summation over I

- a) Accumulate $\binom{N}{I}$ in INC
- b) Increment I=I+1
- c) IF (I.LE.N), go to Step 3. Otherwise summation over I complete. Go to Step 10.

STEP 10: Increment and Test Summation over N

- a) Increment N=N+1
- b) If (N.LE.NORDR), go to STEP 2. Otherwise nonlinear currents complete.

The coefficient expressions S2 and SP are also calculated by an iterative process, summarized as follows:

STEP 1: Initialize Summation over S

- a) Set index at S=MAX(I-1,0)
- b) Initialize accumulator for (S) to SIC=S+1
- c) Clear coefficient expressions S2=0. and SP=0.

STFP 2. Initialize summation over R

- a) Set index at R=0
- b) Initialize accumulator for $\binom{J}{R}$ to RJC=1

STEP 3: Accumulate Coefficient Terms

a) Calculate

0

S2 = S2 +
$$\binom{I}{S}\binom{J}{R}$$
 G(I-S+1) H(J-R+S+1)MU^{R-J} F2(N-I-J+1,R+1)

b) Calculate

$$SP = SP + {1 \choose S} {J \choose R} G(I-S+1) H(J-R+S+1) MU^{R-J} FP(N-I-J+1,R+1)$$

STEP 4: Increment and Test Summation over R

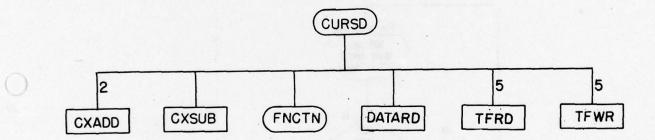
- a) Accumulate $\binom{J}{R}$ in RJC
- b) Increment R=R+1
- c) If (R.LE.J), go to Step 3. Otherwise summation over R complete, go to Step 5.

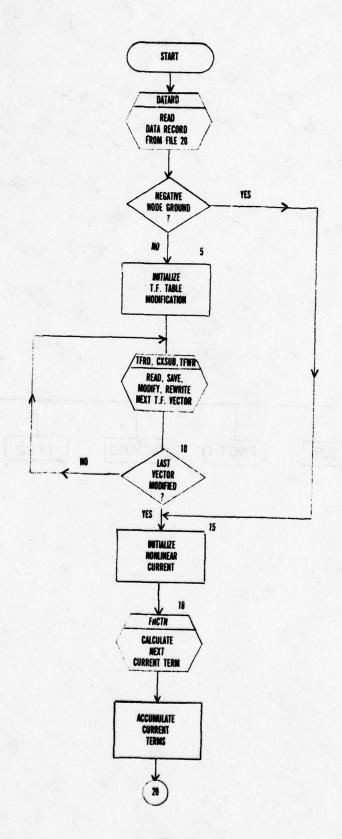
STEP 5: Increment and Test Summation over S

- a) Accumulate $\binom{I}{S}$ in SIC
- b) Increment S=S+1
- c) If (S.LE.I), go to Step 2. Otherwise coefficient expression complete.

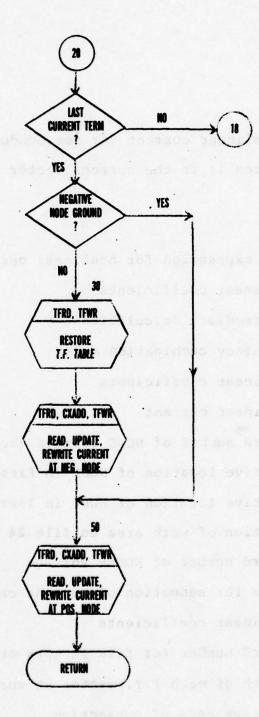
Then the transfer function table is restored by reading the transfer functions saved in the scratch area on file 24 and rewriting them to the original locations in the table. The procedure is similar to that employed for modifying the transfer function table.

The nonlinear currents are placed in the current vector on file 25 by subtracting I2 from the current at NODE+1, subtracting IP from the current at NODE+2, and adding I2 and IP to the current at NODE+3. The current elements for the sequential vacuum pentode nodes, beginning at I=2*NODE-1, are read from file 25 by subroutine TFRD and stored in X, X1, X2, and X3. I2 and IP are subtracted from X1 and X2 respectively by subroutine CXSUE. I2 and IP are added to X3 by subroutine CXADD. The updated current vector elements are rewritten to their original locations by subroutine TFWR and subroutine CURPT returns to the calling program.





CURSD 1 of 2



NAME: CURSD

TYPE: SUBROUTINE

GENERAL PURPOSE:

Calculates nonlinear current for semiconductor diode and places it in the current vector

VARIABLES .

A = T.F. expression for nonlinear current

C2 = Nonlinear coefficients

D = Intermediate calculation

FREQ = Frequency combination vector

GAMMA = Nonlinear coefficients

HKC = Nonlinear current

I = Record number of NODE voltage (current)

Il = Relative location of NODE in first T.F.

12 = Relative location of NODE in last T.F.

I3 = Location of work area on file 24

J = Record number of NODEP voltage

K = Index for summation of current calculation

KS = Nonlinear coefficients

L = Record number for file 24 work area

LT = Length of each T.F. vector in words

NODE = Positive node of connection

NODER = Negative node of connection

NORDR = Order of analysis

S = Angular frequency = 300 mm and a second second

X = NODE voltage (current)

Y = NODEP voltage (current)

SUBROUTINES CALLED.

CXADD CXSUP DATARD FNCTN

TFRD TE FTWR 100 38 BOLDONE SETENSIS 381 381 38 38816 SOUTH

CALLING PROGRAMS:

PHASE3

DESCRIPTION:

Subroutine CURSD calculates the nonlinear current for the semiconductor diode and places it in the current vector. Arguments transmitted to the subroutine are: NORDR, the order of the nonlinear current to be calculated, and S, the angular frequency. The nonlinear current is a function of the device coefficients and parameters calculated in Phase 1 and the node voltages derived from the lower order transfer functions stored on file 24. The positive and negative nodes of connection, NOD and NODER, are stored in the semiconductor diode driver record in global common.

te accepted and model of by the following interactive crocedure

First the semiconductor diode data record is read from file 20 by subroutine DATARD and stored in the common data buffer PUFF. Then the node voltages for the nonlinear current calculations are computed by replacing the transfer function at

NODE with the difference between the transfer functions at NODE and NODEE for every vector in the table. This "modified transfer function table" serves as the basis of the nonlinear current calculations. If the negative node of connection is ground (NODEE =0), modification of the transfer function table is by-passed.

The transfer function table consists of 2 -2 contiguous vectors beginning at record number 1 on file 24. Each complex vector element is the transfer function at some node in the circuit and occupies two decimal storage locations. The transfer functions are stored in the vectors according to increasing node numbers. This method of storage allows the transfer functions to be accessed and modified by the following iterative procedure:

STEP 1: Initialization

- a) Length of each T.F. vector: LT=2*CFDEP
- b) Location of NODE in first vector: I1=2*NCDE-1
- c) Location of MCDF in last vector: I2=I1+(2NORDR -3)*IT
- d) Crigin of scratch area on file 24: $I=2^{NORDR} *I.T$
- e) Set index at first vector: I=Il

STEP 2: Access and Modify Next Vector

- a) Pead transfer function from I, store in X.
 - b) Write X to scratch area at L, increment L=L+2

- c) Locate NODEB in vector: J=I+(NODEB-NODE) *2
 - d) Read transfer function from J, store in Y.
- e) Calculate X=X-Y and write to file 24 at I.

STEP 3: Increment and Test Index

- a) Increment I=I+LT for next vector.
- b) If (I.LE.I2), more vectors remain to be modified. Go to Step 2.
- c) Otherwise modified transfer function table complete.

The nonlinear current of order NORDR for the semiconductor diode is defined as:

HKG =
$$\sum_{K=2}^{NORDR}$$
 (KS(K) * D) * F_K^{NORDR} (NODE)

where:

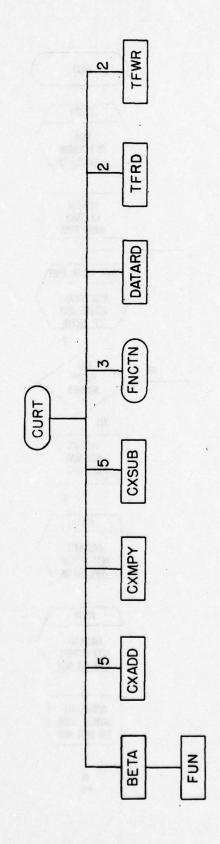
$$D=S*(GAMMA(K)+C2(2)*KS(K-1)/K)$$

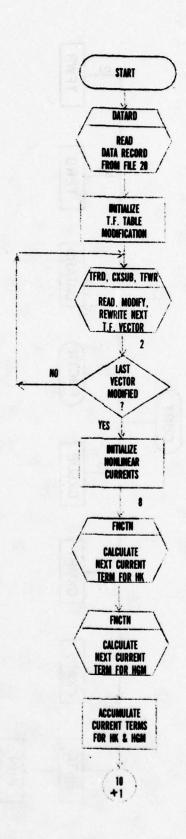
Subroutine CURSD initializes HKG to zero, then each symmetric transfer function expression F_K^{NORDR} (NODE) for K=2,3,... NORDR is computed by subroutine FNCTN. After the intermediate

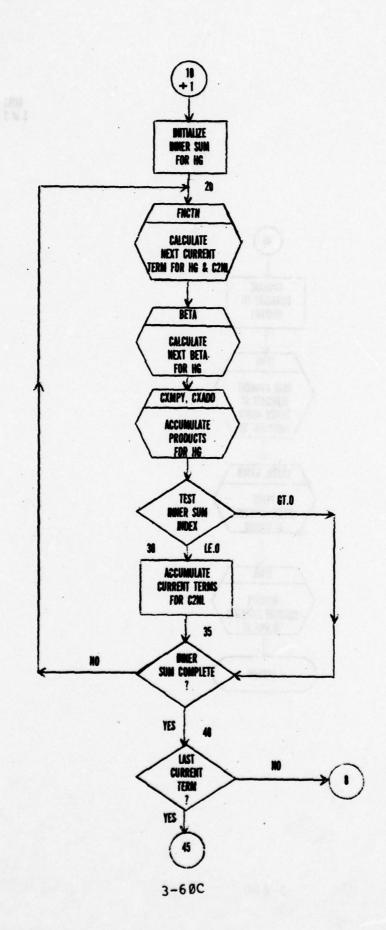
calculation is performed for D, the product is formed and accumulated in HKG.

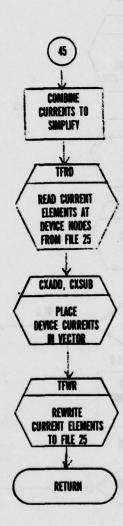
Then the transfer function table is restored by reading the transfer functions saved in the scratch area on file 24 and rewriting them to their original locations in the table. The procedure is similar to that employed for modifying the transfer function table. If the negative node of connection is ground (NODEB=0), the restore procedure is by-passed.

The nonlinear current is placed in the current vector on file 25 by adding it to both the current at NODEB (for NODEB #0) and the current at NODE. The NODEB current, located at I=2*NODEB-1, is read from file 25 by subroutine TFRD and stored in X. HKG is added to X by subroutine CXADD and the updated NODEB current is rewritten to its original location by subroutine TFWR. The procedure is repeated for the NODE current located at I=2*NODE-1, and subroutine CURSD returns to the calling program.









NAME: CURT

TYPE: SUBROUTINE

GENERAL PURPOSE:

Calculates nonlinear currents for bipolar junction transistor and places them in current vector

VARIABLES:

A = T.F. expression for nonlinear current

ALPHA = Nonlinear coefficients

P = Intermediate calculation for HG

C = Intermediate calculation for HG

C2NL = Nonlinear current

D = Intermediate calculation for C2NL

FREQ = Frequency combination vector

GAMMA = Nonlinear coefficients

HC = Nonlinear coefficients

HG = Nonlinear current

HGM = Nonlinear current

HK = Nonlinear current

HM = Nonlinear coefficients

ICHMO = Nonlinear parameter

I = Record number of NODE voltage (current)

Il = Relative location of NODE in first T.F.

12 = Relative location of NODE in last T.F.

I3 = Location of work area on file 24

K = Index for outer summation of current
calculation

L = Record number for file. 24 work area

KS = Nonlinear coefficients

LT = Length of each T.F. vector in words

M = Index for inner summation of current calculation

NODE = Base node of connection

NORDR = Order of analysis

S = Angular frequency

SCO = Nonlinear parameter

SIGN = Sign of HGM current terms

SNCS = Nonlinear parameter

X = Base node voltage (current)

Xl = (Base+1) voltage (current)

X2 = (Base+2) voltage (current)

X3 = (Pase+3) voltage (current)

SUBPOUTINES CALLED.

BETA CXADD CXMPY CXSUE

DATARD FNCTN TFRD

TFWR

CALLING PROGRAMS:

PHASE3

DESCRIPTION :

Subroutine CURT calculates the nonlinear currents for the bipolar junction transistor and places them in the current vector. Arguments transmitted to the subroutine are: NORDR, the order of the nonlinear currents to be calculated, and S, the angular frequency. The nonlinear currents are functions of the device parameters and coefficients calculated in Phase 1 and the node voltages derived from the lower order transfer functions stored on file 24. The base node of connection, NODE, is stored in the transistor driver record in global common. The remaining nodes are sequentially numbered by the system as NODE+1, NODE+2, and NODE+3.

First the transistor data record is read from file 20 by subroutine DATARD and stored in the common data buffer BUFF. Then the node voltages for the nonlinear current calculations are computed for each vector in the transfer function table as follows: The transfer function at NODE is replaced by the difference between the transfer functions at NODE+2 and NODE; the transfer function at NODE+2 is replaced by the difference between the transfer functions at NODE+2 and NODE+1: and the transfer function at NODE+1 is replaced by the difference between the transfer functions at NODE+1 and NODE+3. This "modified transfer function table" serves as the basis for the nonlinear current calculations.

The transfer function table consists of 2^{NORDR} -2 contiguous vectors beginning at record number 1 on file 24. Each complex vector element is the transfer function at some node in the

circuit and occupies two decimal storage locations. The transfer functions are stored in the vectors according to increasing node numbers. Because the transistor nodes are numbered sequentially, the corresponding transfer functions occupy successive elements in each vector. This method of storage allows the transfer functions to be accessed and modified by the following iterative procedure:

STEP 1: Initialization

- a) Length of each T.F. vector: LT=2*ORDER
- b) Location of NODE in first vector: I1=2*NODE-1
- c) Location of NODE in last vector: I2=I1+(2NORDR -3)*LT
 - d) Origin of scratch area on file 24: L = 2 NORDR * LT
 - e) Set index at first vector: I=Il

STEP 2: Access and Modify Next Vector

- a) Read transfer functions for sequential device nodes beginning at I, store in X, X1, X2, and X3
- b) Write device voltages to scratch area at L, increment L=L+2
- c) Calculate X=X2-X, X2=X2-X1, and X1=X1-X3
- d) Write updated transfer function to file 24 beginning at I

STEP 3: Increment and Test Index

- a) Increment I=I+LT
- b) If (I.LE.I2), more vectors remain to be modified. Go to Step 2.
 - c) Otherwise modified transfer function table complete.

The nonlinear currents of order NORDR for the bipolar junction transistor are defined as:

$$HK = \sum_{K=2}^{NORDR} KS(K) * F_{K}^{NORDR} (NODE + 2)$$

HGM =
$$\sum_{K=2}^{NORDR}$$
 S * GAMMA (K) * F_K^{NORDR} (NODE + 3)

$$HG = \sum_{K=2}^{NORDR} \sum_{M=0}^{K} B_{M}^{K} \star F_{K}^{NORDR} Q_{M}^{K} (NODE + 2, NODE + 1)$$

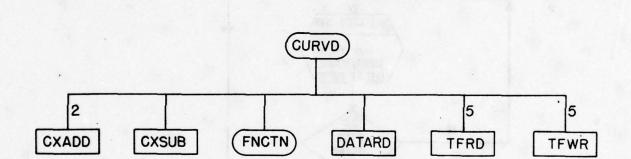
C2NL =
$$\sum_{K=2}^{NORDR}$$
 D * F_K^{NORDR} (NODE + 1)

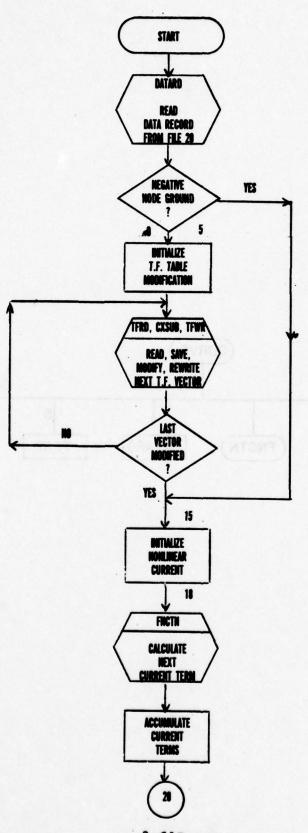
where D=S*HC(2)*KS(K-1)/K. Subroutine CURT initializes HK, HGM, HG, and C2NL to zero. Then the symmetric transfer function NORDR expressions F_K (NODE+2) and F_K (NODE+3) for K=2,3,...NORDR are computed by subroutine FNCTN, multiplied by the appropriate coefficients KS(K) or S*GAMMA(K), and accumulated in HK and HGM respectively. Then the symmetric transfer function expression (NODE+2, NODE+1) for K=2,3,... NORDR and $M=\emptyset,1,...K$ is computed by subroutine FNCTN. Subroutine BETA is called to compute BM and the product is formed and accumulated in HG. The symmetric transfer function expression FK (NODE+1) is identical to the expression for HG evaluated at M=0. intermediate calculation for D is performed and the product is formed and accumulated in C2NL.

Then the transfer function table is restored by reading the transfer functions saved in the scratch area on file 24 and rewriting them to the original locations in the table. The procedure is similar to that employed for modifying the transfer function table.

The nonlinear currents are placed in the current vector on file 25 by adding HG and HGM to the current at NODE+1, adding HK and C2NL to the current at NODE+3, subtracting HK and C2NL from the current at NODE+1, and subtracting HG and HCM from the current at NODE+2. The current elements for the sequential transistor nodes, beginning at I=2*NODE-1, are read from file 25 by subroutine TFRD and stored in X, X1, X2, and X3. HG and HGM are added to X1 and HK and C2NL added to X3 by subroutine CXADD.

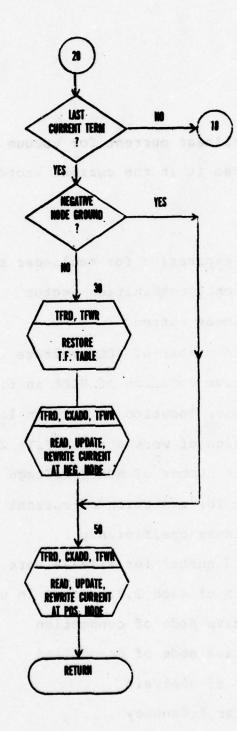
HK and C2NL are subtracted from X1 and HG and HGM are subtracted from X2 by subroutine CXSUB. The updated current vector elements are rewritten to their original locations by subroutine TFWR and subroutine CURT returns to the calling program.





CURVO 1 of 2

3-61B



NAME: CURVD

TYPE · SUPROUTINE

GENERAL PURPOSE:

Calculates nonlinear current for vacuum diode and places it in the current vector

VARIABLES:

A = T.F. expression for nonlinear current

FREQ = Frequency combination vector

PK = Nonlinear current

I = Record number of NCDE voltage (current)

Il = Pelative location of NCDE in first T.F.

I2 = Relative location of NODE in last T.F.

I3 = Location of work area on file 24

J = Record number of NODEP voltage

K = Index for summation of current calculation

KS = Monlinear coefficients

I = Record number for file 24 work area

LT = Length of each T.F. vector in words

NODE = Positive node of connection

NODEE = Negative node of connection

NORDR = Order of analysis

S = Angular frequency

Y = NODE voltage (current)

Y = NODEB voltage (current)

SUBROUTINES CALLED:

CXADD CXSUB DATARD FNCTN TFRD
TFWR

CALLING PROGRAMS:

PHASE3'

DESCRIPTION:

Subroutine CURVD calculates the nonlinear current for the vacuum diode and places it in the current vector. Arguments transmitted to the subroutine are: NORDE, the order of the nonlinear current to be calculated, and S, the angular frequency. The nonlinear current is a function of the device coefficients calculated in Phase 1 and the node voltage derived from the lower order transfer functions stored on file 24. The positive and negative nodes of connection NODE and NODEE are stored in the vacuum diode driver record in global common.

First the vacuum diode data record is read from file 20 by subroutine DATARD and stored in the common data buffer BUFF. Then the node voltages for the nonlinear current calculations are computed by replacing the transfer function at NODE with the difference between the transfer functions at NODE and NODEF for every vector in the table. This "modified transfer function table" serves as the basis for the nonlinear current calculations. If the negative node of connection is ground

(NODER=0), modification of the transfer function table is by-passed.

The transfer function table consists of 2 -2 contiguous vectors beginning at record number 1 on file 24. Each complex vector element is the transfer function at some node in the circuit and occupies two decimal storage locations. The transfer functions are stored in the vectors according to increasing node numbers. This method of storage allows the transfer functions to be accessed and modified by the following iterative procedure:

STEP 1: Initialization

- a) Length of each T.F. vector: LT=2*ORDER
- b) Location of NCDE in first vector: I1=2*NODE-1
- c) Location of NODE in last vector: I2=I1+(2 -3)*LT
- d) Origin of scratch area on file 24: L=2 *LT
- e) Set index at first vector: I=Il

STEP 2: Access and Modify Next Vector

- a) Read transfer function from I, store in X.
 - b) Write X to scratch area at L, increment L=L+2.
 - c) Locate NCDEP in vector: J=I+(NODEB-NODE)*2.
 - d) Read transfer function from J, store in Y.
 - e) Calculate X=X-Y and write to file 24 at I.

STEP 3: Increment and Test index

- a) Increment I=I+LT for next vector.
- b) If (I.LE.I2), more vectors remain to be modified.

 Go to Step 2.
 - c) Otherwise modified transfer function table complete.

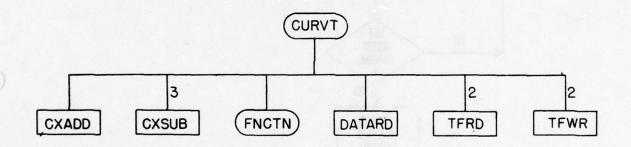
The nonlinear current of order NORDR for the vacuum diode is defined as:

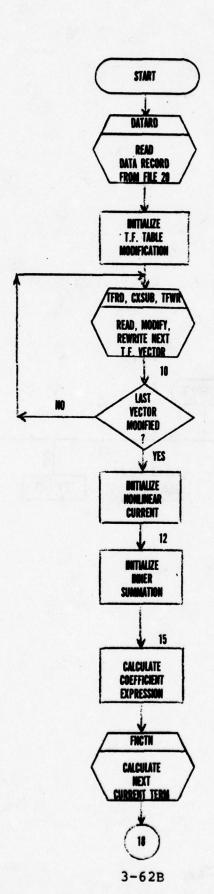
$$HK = \sum_{K=2}^{NORDR} KS(K) * F_{K}^{NORDR} (\overline{NODE})$$

Subroutine CURVD initializes HK to zero, then each symmetric transfer function expression F_K^{NORDR} (NODE) for K=2,3,...NORDR is computed by subroutine FNCTN, multiplied by the appropriate coefficient KS(K), and accumulated in HK.

Then the transfer function table is restored by reading the transfer functions saved in the scratch area on file 24 and rewriting them to their original locations in the table. The procedure is similar to that employed for modifying the transfer function table. If the negative node of connection is ground (NODEB=C), the restore procedure is by-passed.

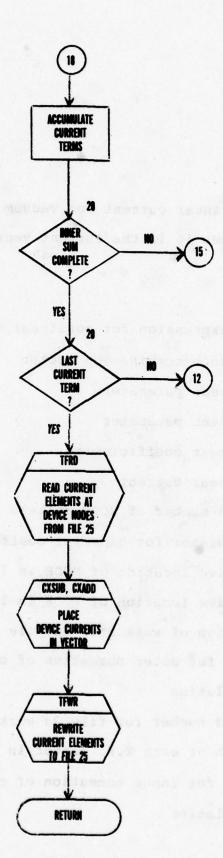
The nonlinear current is placed in the current vector on file 25 by subtracting it from the current at NODEB (for NODEB #0) and adding it to the current at NODE. The NODEB current, located at I=2*NODEP-1, is read from file 25 by subroutine TFRD and stored in X. HK is subtracted from X by subroutine CXSUB and the updated NCDEB current is rewritten to its original location by TFWF. The procedure is repeated for the NODE current, except that NODE is located at I=2*NODE-1 and HK is added to X by subroutine CXADD. Subroutine CURVD then returns to the calling program.





CURYT 1 of 2

V



CURVT 2 of 2

3-62C

NAME: CURVT

TYPE: SUBROUTINE

GENERAL PURPOSE:

Calculates nonlinear current for vacuum triode and places it in the current vector

VARIABLES:

A = T.F. expression for nonlinear current

FREQ = Frequency combination vector

G = Nonlinear parameter

GP = Nonlinear parameter

H = Nonlinear coefficients

HG = Nonlinear current

I = Record number of NODE voltage (current).

IFACT = Accumulator for binomial coefficient

Il = Relative location of NODE in first T.F.

12 = Relative location of NODE in last T.F.

I3 = Location of work area on file 24

K = Index for outer summation of current calculation

L = Record number for file 24 work area

LT = Length of each T.F. vector in words

M = Index for inner summation of current calculation

MU = Input value

NOTE = Base node of connection

NORDR = Order of analysis

Q = Intermediate calculation

S = Angular frequency

X = Base node voltage (current)

X1 = (Base+1) voltage (current)

x2 = (Base+2) voltage (current)

SUBROUTINES CALLED:

CXADD CXSUB DATARD FNCTN

TFRD TFWR

CALLING PROGRAMS:

PHASE3

DESCRIPTION .

Subroutine CURVT calculates the nonlinear current for the vacuum triode and places it in the current vector. Arguments transmitted to the subroutine are: NORDR, the order of the nonlinear current to be calculated, and S, the angular frequency. The nonlinear current is a function of the device parameters and coefficients calculated in Phase 1 and the node voltages derived from the lower order transfer functions stored on file 24. The base node of connection, NODE, is stored in the vacuum triode driver record in global common. The remaining nodes are sequentially numbered by the system as NODE+1 and NODE+2.

First the vacuum triode data record is read from file 20 by subroutine DATARD and stored in the common data buffer BUFF. Then the node voltages for the nonlinear current calculations are computed as follows: the transfer function at NODE is replaced by the difference between the transfer functions at NODE and NODE+2 for every vector in the table and the transfer function at NODE+1 is replaced by the difference between the transfer functions at NODE+1 and NODE+2 divided by the input parameter MU for every vector in the table. This "modified transfer function table" serves as the basis for the nonlinear current calculations.

The transfer function table consists of 2 -2 contiguous vectors beginning at record number 1 on file 24. Each complex vector element is the transfer function at some node in the circuit and occupies two decimal storage locations. The transfer functions are stored in the vectors according to increasing node numbers. Because the vacuum triode nodes are numbered sequentially, the corresponding transfer functions occupy successive elements in each vector. This method of storage allows the transfer functions to be accessed and modified by the following iterative procedure:

STEP 1: Initialization

- a) Length of each T.F. vector: LT=2*ORDER
- p) Location of MODE in first vector · Il=2*NODE-1

- c) Location of NODE in last vector: I2=I1+(2NORDR -3)*LT
- d) Origin of scratch area on file 24: L = "NORDR * LT
- e) Set index at first vector. I=Il

STEP 2: Access and Modify Next Vector

- a). Read transfer functions for sequential device nodes beginning at I, store in X, X1, and X2.
- b) Write device voltages to scratch area at L, increment L=L+2
- c) Calculate X=X-X2 and X1=(X1-X2)/MU
- d) Write updated transfer functions to file 24 beginning at I

STEP 3: Increment and Test Index

- a) Increment I=I+LT
- b) If (I.LE.LT), more vectors remain to be modified. Go to Step 2.
 - c) Otherwise modified transfer function table complete.

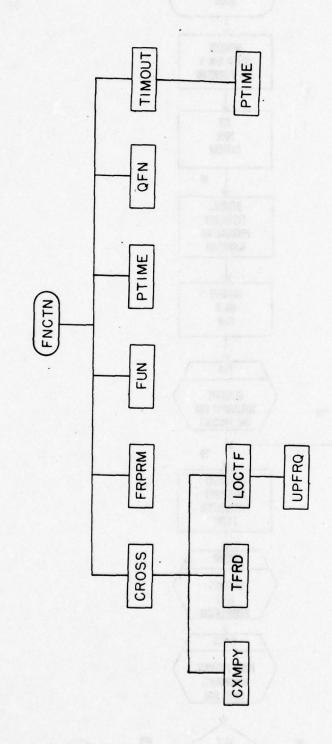
The nonlinear current of order NCPER for the vacuum triode is defined as:

$$HG = \sum_{K=2}^{NORDR} \sum_{M=0}^{K} D_{M,K} * F_{K}^{NORDR} Q_{M}^{K} [NODE + 1, NODE + 2]$$

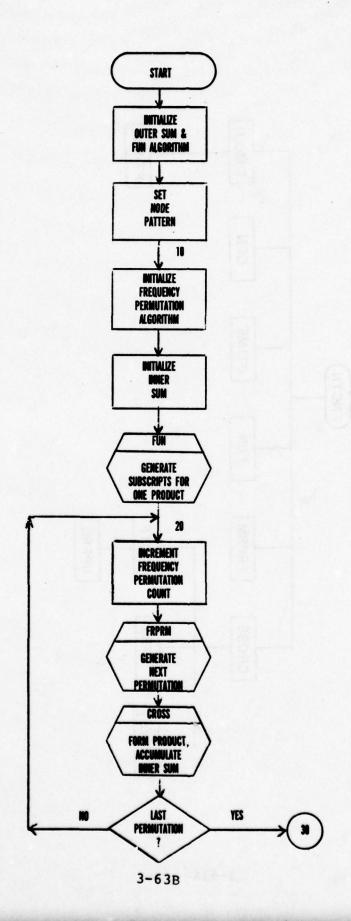
where $D_{M,K} = (M)$ (K-M)*GP*H(K)/K+G*H(K+1). Subroutine CURVT initializes HG to zero, then each symmetric transfer function expression F_{K}^{NORDR} Q_{M}^{K} [NODE+1, NODE+2] for K=2,3,... NORDR and $M=\emptyset$,1...K is computed by subroutine FNCTN. After the intermediate calculations are performed for $D_{M,K}$, the binomial coefficient $\binom{K}{M}$ is accumulated in IFACT, and the product is formed and accumulated in HG.

Then the transfer function table is restored by reading the transfer functions saved in the scratch area on file 24 and rewriting them to the original locations in the table. The procedure is similar to that employed for modifying the transfer function table.

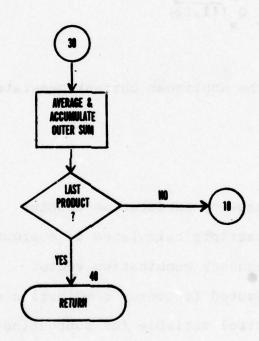
The nonlinear current is placed in the current vector on file 25 by subtracting it from the current at NODE+1 and adding it to the current at NODE+2. The current elements for the sequential vacuum triode nodes, beginning at I=2*NODE-1, are read from file 25 by subroutine TFRD and stored in X, X1, and X2. HG is subtracted from X1 by subroutine CXSUB and added to X2 by subroutine CXADD. The updated current elements are rewritten to their original locations by subroutine TFWR and subroutine CURVT returns to the calling program.



10



FINCTIN 1 of 2



NAME: FNCTN

TYPE: SUBROUTINE

GENERAL PURPOSE:

Calculates the symmetrized transfer function expression

for use by the nonlinear current generator subroutines

VARIABLES:

CNT = Frequency permutation counter

FNK = Subscripts calculated in subroutine FUN

FREQ = Frequency combination vector

FREQC = Permuted frequency combination vector

ISWF = Control variable for subroutine FRPRM

ISWT = Control variable for subroutine FUN

K = Length of FNK vector

M = Number of nodes I2 in node pattern

N = Order of T.F. expression

Q = Node pattern of T.F. expression

U = Location of nonlinearity relative to base node

V = Location of dependence relative to base node

XX = Calculated symmetrized T.F. expression

YY = Sum of T.F. products calculated in subroutine CROSS

SUBROUTINES CALLED:

CROSS FRPRM FUN PTIME QFN
TIMOUT

CALLING PROGRAMS:

CURVT

CURJF CURNC CURNDS CURSD CURT

DESCRIPTION :

CURVD

For given N, K and M subroutine FNCTN calculates the expression F_K^N [Q $_M^K$ (U,V)], which is a sum of symmetric products of voltages from the modified transfer function table used in calculating nonlinear currents during Phase 3.

A nonlinear current is a function of the voltage across the nonlinearity as well as the voltage at some other node in the circuit. A nonlinear current of order N can be expressed as a power series in voltage at two nodes, U and V:

$$I_{N} (f_{1}, f_{2}, ... f_{N}) = \sum_{K=2}^{N} \sum_{M=0}^{K} KS_{K,M} F_{K}^{N} [\overline{Q_{M}^{K} (U, V)}]$$

The KS's are functions of the Taylor series coefficients calculated in Phase 1. F_K^N [Q_M^K (U, V)] is a sum of products of K lower order voltages from the modified transfer function table. Each product has the following general form:

The function Q_M^K (U, V) determines the pattern of node voltages U and V so that each transfer function product contains K voltages, where the first K-M voltages are from node U and the remaining M voltages are from node V. The function F_K^N determines the order P_i of each transfer function in the product so that P_i + P_i +... P_i =N. The overbar depicts the symmetry of the product, achieved by averaging over the N! possible permutations of the N input frequencies.

For example, a third order nonlinear current is given by:

$$13(f_{1},f_{2},f_{3}) = \sum_{K=2}^{3} \sum_{M=0}^{K} KS_{K,M} F_{K}^{3} [\overline{Q_{M}^{K}(u,v)}]$$

$$= \sum_{K=2}^{3} KS_{K,0} F_{K}^{3} [\overline{Q_{0}^{K}(u,v)}] + KS_{K,1} F_{K}^{3} [\overline{Q_{1}^{K}(u,v)}] + \dots$$

$$+ KS_{K,K} F_{K}^{3} [\overline{Q_{K}^{K}(u,v)}]$$

Rather than expanding the expression completely, consider the current component at K=2:

$$F_{2,0}^{3} F_{2}^{3} [Q_{0}^{2} (u,v)] + KS_{2,1} F_{2}^{3} [Q_{1}^{2} (u,v)] + KS_{2,2} F_{2}^{3} [Q_{2}^{2} (u,v)]$$

Each F $_2^3$ represents a sum of products of first and second order voltages. The nodes from which the voltages are derived are determined by the Q function where Q $_0^2$ represents the node pattern UU, Q_1^2 represents VU, and Q_2^2 represents VV. Together the two functions produce the following additive terms:

$$F_{2}^{3} \overline{[Q_{0}^{2} (U,V)]} = \overline{U_{2}U_{1} + U_{1}U_{2}}$$

$$F_{2}^{3} \overline{[Q_{1}^{2} (U,V)]} = \overline{V_{2}U_{1} + V_{1}U_{2}}$$

$$F_{2}^{3} \overline{[Q_{2}^{2} (U,V)]} = \overline{V_{2}V_{1} + V_{1}V_{2}}$$

The symmetry of each product, depicted by the overbar, is achieved by averaging over all the 3!=6 possible permutations of the three input frequencies. Let the six permutations be

represented by:

Then the symmetric product U2U1 is

$$\begin{split} &\frac{1}{6} \left[\text{U}(\textbf{f}_{1}, \ \textbf{f}_{2}) \text{U}(\textbf{f}_{3}) \ + \ \text{U}(\textbf{f}_{1}, \ \textbf{f}_{3}) \text{U}(\textbf{f}_{2}) \ + \ \text{U}(\textbf{f}_{2}, \ \textbf{f}_{1}) \text{U}(\textbf{f}_{3}) \ + \\ & \text{U}(\textbf{f}_{2}, \ \textbf{f}_{3}) \text{U}(\textbf{f}_{1}) \ + \ \text{U}(\textbf{f}_{3}, \ \textbf{f}_{1}) \text{U}(\textbf{f}_{2}) \ + \ \text{U}(\textbf{f}_{3}, \ \textbf{f}_{2}) \text{U}(\textbf{f}_{1}) \right] \end{split}$$

The symmetric product U1U2 is:

$$\frac{1}{6} \left[U(f_1)U(f_2,f_3) + U(f_1)U(f_3,f_2) + U(f_2)U(f_1,f_3) + U(f_2)U(f_3,f_1) + U(f_3)U(f_1,f_2) + U(f_3)U(f_2,f_1) \right]$$

The symmetric product V U is:

$$\frac{1}{6} \left[V(f_1, f_2) U(f_3) + V(f_1, f_3) U(f_2) + V(f_2, f_1) U(f_3) + V(f_2, f_2) U(f_3) + V(f_3, f_2) U(f_3$$

The symmetric product $\overline{V_1U_2}$ is:

$$\frac{1}{6} \left[V(f_1) U(f_2, f_3) + V(f_1) U(f_3, f_2) + V(f_2) U(f_1, f_3) + V(f_2) U(f_3, f_3) + V(f_2) U(f_3, f_3) + V(f_3) U(f_3, f_3$$

The symmetric product $v_2^{V_1}$ is:

$$\frac{1}{6} \left[v(f_1, f_2) v(f_3) + v(f_1, f_3) v(f_2) + v(f_2, f_1) v(f_3) + v(f_2, f_3) v(f_1) + v(f_3, f_1) v(f_2) + v(f_3, f_2) v(f_1) \right]$$

The symmetric product V_1V_2 is:

$$\frac{1}{6} \left[v(f_{1}) v(f_{2}, f_{3}) + v(f_{1}) v(f_{3}, f_{2}) + v(f_{2}) v(f_{1}, f_{3}) + v(f_{2}) v(f_{3}, f_{3}) + v(f_{3}) v(f_{3}, f_{3}) + v(f_{3})$$

The nonlinear current component at K=3 is derived in a similar fashion:

$$KS_{3'0} F_{3}^{3} [Q_{0}^{3} (U, V)] + KS_{3'1} F_{3}^{3} [Q_{1}^{3} (U, V)] + KS_{3'2} F_{3}^{3} [Q_{3}^{3} (U, V)]$$

where each F represents a product of three first order

voltages. The nodes from which the voltages are derived are determined by the Q function where Q_0^3 represents the node pattern UUU, Q_1^3 represents VUU, Q_2^3 represents VVU, and Q_3^3 represents VVV. Together the two functions produce the following additive terms:

$$F_{3}^{3} = \overline{[Q_{0}^{3} (U, V)]} = \overline{U_{1}U_{1}U_{1}}$$

$$F_{3}^{3} = \overline{[Q_{1}^{3} (U, V)]} = \overline{V_{1}U_{1}U_{1}}$$

$$F_{3}^{3} = \overline{[Q_{2}^{3} (U, V)]} = \overline{V_{1}V_{1}U_{1}}$$

$$F_{3}^{3} = \overline{[Q_{3}^{3} (U, V)]} = \overline{V_{1}V_{1}V_{1}}$$

Some nonlinear currents are functions of the voltage at a single node. In this case M= 0 and the current expression reduces to:

$$I_N (f_1, f_2, ... f_N) = \sum_{K=2}^N KS_K F_K^N (V)$$

Subroutine FNCTN calculates one additive term F_K^N [C_M^K (U,V)] of the nonlinear current expression for given N, K, and M. Arguments transmitted to the subroutine are: N, the order of the current being calculated: K, the number of voltages in each transfer function product. M, the number of voltages from node V

in each product: U and V, the nodes from which the voltages are derived (relative to the base node of the nonlinear device) and FREC, the frequency combination vector. The calculated term is returned to the calling program in XX.

Subroutine FNCTN uses subroutine QFN to determine the node pattern returned in the integer vector Q, subroutine FUN to determine the orders of the transfer functions returned in the vector FNK, and subroutine FRPRM to produce the frequency permutations returned in the integer vector FREQC. Subroutine CRCSS is called to access the voltages in the modified transfer function table and to form the voltage product for a single frequency permutation.

The algorithm used by subroutine FNCTN to calculate the term $F \times \{O \times \{U, V\}\}$ is summarized as follows:

STEP 1.

- a) Clear expression to XX=0.
- b) Initialize FUN algorithm by ISWT=-1
- c) Call subroutine QFN to generate Q_M^K (U, V): return node pattern in Q vector.

STEP 2.

a) Initialize FRPRM algorithm by ISWF=-1

- b) Clear frequency permutation count CNT=0.
- c) Initialize symmetric product YY=0.
- d) Call subroutine FUN to generate orders of transfer functions for one term of F_K^N ; return orders in FNK vector.

refurned in the integer vector . C. sabrestine !

STEP 3:1 Degrater and larged relenant and to large out to

- a) Increment permutation count CNT=CNT+1.
- b) Call subroutine FRPRM to create a frequency permutation: return permutation in FREQC vector.
 - c) Call subroutine CROSS to calculate product for this frequency permutation and add product to YY
 - d) If more permutations (ISWF=0), go to Step 3a.
 - e) If permutations complete (ISWF=1), go to Step 4.

STEP 4:

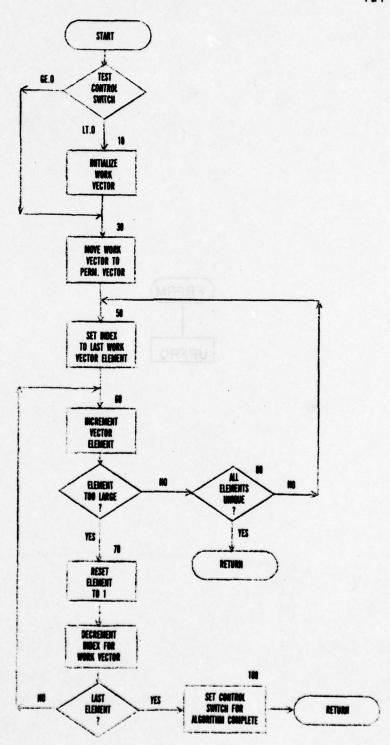
a) Calculate symmetric product and add to expression:

XX=XX+YY/CNT

- b) If more F_{K}^{N} terms (ISWF=0), go to Step 2b.
- c) If F_K^N complete (ISWF=1), exit with calculated expression in XX.



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NAME: FRPRM

TYPE: SUBROUTINE

GENERAL PURPOSE .

Cenerates a frequency permutation for use by subroutine FNCTN

VARIATLES .

FREC = The integers 1 thru 10

FPEOC = Permuted frequency combination vector

ISW = Control switch: -1 = Initialize algorithm

Ø = Continue algorithm

M.T. the number of trechescies to my permuted

1 = Algorithm complete

M = Number of frequencies to be permuted

SEC = Work vector

SUPROUTINES CALLED.

NONE

CALLINC PROCRAMS.

CUPPT FNCTN

DESCRIPTION:

Subroutine FRPPM generates frequency permutations for use by subroutine FNCTN. Arguments transmitted to the subroutine are:

N, the number of frequencies to be permuted; FREQ, the sequential frequency numbers to be permuted; and ISW, a control variable described below. Arguments returned to the calling program are: FREQC, the permuted frequency vector and ISW, the updated control variable.

In order to generate every possible permutation of N frequency numbers, subroutine FRPRM must be invoked N! times. These iterative calls to subroutine FPPRM are controlled by the integer variable ISW, which is initialized to -1 by the calling program, tested and updated in FRPRM, then retested in the calling program. ISW=-1 initializes the algorithm of subroutine. FRPRM by creating the first permutation of order N in the work vector SEQ. When ISW=0, subroutine FPPRM transfers the permutation from the work vector to FPPCC and then generates the next permutation in SEQ, saving it in the labelled common area PRMFL. ISW is set to 1 when the last permutation has been generated.

The frequency permutations generated by subroutine FRPRM are stored in the first N elements of the integer vector SEQ. Each element of the vector is numerically equal to the sequential frequency number it represents. The largest frequency that any vector element can contain is N. Subroutine FRPRM ignores permutations in which a frequency number appears more than once.

Frequency permutations are generated recursively by subroutine FPFF. The first permutation of M frequencies is defined to be the ordered set of integers (1, 2,...) which is generated in the work vector by

SEC(I)=I for I=1, 2,...N

Fach successive permutation of order N is derived from the previous SFC vector according to the following algorithm:

STTP 1 Initialize Algorithm at SEQ(N)

a) Set SFC subscript to J=M.

STEP 2: Update and Test SEQ(J) .

- a) Define SEC(J)=SEC(J)+1
- b) If (SEC(J).GT.N), this value too large for SEC(J). Go to Step 3.
- c) If (SEQ(J).LE.P), this value O.K. for SEQ(J). Go to Step 4.

STEP 3 Redefine SFC(J), Update and Test J

- a) Set SEQ(J)=1
- b) Decrement SEQ subscript J=J-1
- c) If (J.CT.0), go to Step 2. Otherwise algorithm complete.

Step 4: Test SEQ for Uniqueness

- a) If SEQ(I) = SEQ(J) for any I and J $(I \neq J)$, go to Step 1.
 - b) Otherwise algorithm complete.
- A tabulation of all frequency permutations for every order up to N=6 (the largest order allowed in NCAP) is presented below:

n = 2

			FR	EOC		
NO.	1	2	3	4	5	6
1	1	2	0	0	0	0
2	2	1	0	0	0	0

n = 3

NO.	1	2	FR 3	EOC 4	5	6
1	1	2	3	0	0	0
2	1	3	2	0	0	0
3	2	1	3	0	0	0
4	2	3	1	0	0	0
5	3	1	2	0	0	0
5	. 3	2	1	0	0	0

n = 4

NO.	1	2	FR 3	EQC 4	5	6
1 2 3 4 5 6 7 8 9 10 11 12 13 14 15 16 17 18 19 20 21 22 23 24	1 1 1 1 2 2 2 2 2 2 3 3 3 3 3 3 4 4 4 4 4 4 4 4	2 2 3 3 4 4 1 1 2 2 4 4 1 1 2 2 3 3 3	3 4 2 4 2 3 3 4 1 4 1 3 2 4 1 4 1 2 2 3 1 3 1 2	4 3 4 2 3 2 4 3 4 1 3 1 4 2 4 1 2 1 3 2 3 1 2 1	5 0000000000000000000000000000000000000	6 0000000000000000000000000000000000000

NO.	1.	2	FR 3	EQC 4	5	6		NO.	1	2	FR 3	EQC 4	5	6
12345678901123415678901123456789011234567890112345678901123456789011234567890112345678901123456789011234567890112345678901123456789000112345678900112345678900112345678900112345678900112345678900112345678900011234567890001123456789000112345678900001123456789000000000000000000000000000000000000		22222333333444444555555	3344552244552233552233443344554455335533442244554455	45353445252435252334242345353445-5-435-5-334-4-345252445-5-4	54534354524253523243424245345545545-4-535-3-434-3-545242545-4-	000000000000000000000000000000000000000		61 62 63 64 65 66 67 77 77 77 77 77 77 77 77 77 77 77	333333333333333444444444444444444444444	444445555555	1122551-22442233551-33551-22551-22332233441-33441-22441-2233	25-5-224-4-235252335-5-325-5-223-3-234242334-4-324-4-223-3-2	525-2-424-2-535232535-3-525-2-323-2-434232434-3-424-2-323-2-	000000000000000000000000000000000000000

NO.	1	2	FR 3	EQC 4	5	6		NO.	1	2	FR 3	EQC 4	5	6
12345678901123456789012345678901233456789012344567890123445678900		222222222222222222222222222222222222222	3333334444445555555666666666222222224444455555566666622222222	445566335566334466334455445566225566224466224455335566225566	5646455636354636344535345646455600054600044505045086355600005	050454656353646343545343656454656252646242545242656353656252		61 62 63 64 65 66 66 67 77 77 77 77 77 77 77 77 77 77		444444444455555555555555555555555555555	555555566666000000000000000000000000000	2233662233553344662244662233662233443344552244455223355223344	362623352523463634462624362624362623342423453534452524352524352523342423	636232535252646346646242636232434343453545345545242535232434232

0

n = 6 (continued)

No.	1	2	FR 3	EQC 4	5	6		NO.	1	2	FR.	EQC 4	5	6
121 122 122 123 124 125 127 129 120 121 123 123 124 125 127 129 127 129 129 129 129 129 129 129 129 129 129	222222222222222222222222222222222222222	111111111111111111111111111111111111111	3333334444445555556666661111114444445555556666666111113333333	455663355663344663344554455665566446644553355665566	564645563635463634453534564645566-6-6-6-445-5-456363556-6-5	050454050353040343545343050454050151040141545141056353050151		181 182 183 184 185 187 189 199 199 199 199 199 199 199 199 199	~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~	444444444445555555555555555555555555555	5555556666661111113333333444444666666611111333333344444455555555	1-3366-133553344661-44661-3366-133443344551-44551133551-3344	361613351513463634461614361613341413453534451514351513341413	636131535131646343646141636131434131545343545141535131434131

n = 6 (continued)

NO.	1	2	FF 3	EQC 4	5	6	NO.	1	2	FREQ 3 4	C 5	6
241 242 243 2445 245 245 245 245 245 245 245 245 24		111111111111111111111111111111111111111	22222244444555555566666611111144444555555666666111111222222	445566225566224466224455445566115566114466114455225566115566	564645562625462624452524564645561615461614451514562625561615	656454656252646242545242656454656151646141545141656252656151	301 302 303 304 305 307 307 310 311 311 311 311 311 311 311 311 311		444444444445555555555555555555555555555	5555555666666111111112222222444444446666666611111112222222444555112224445551122244455555555	261612251512462624461614261612241412452524451514251512241412	626121525121646242646141626121424121545242545141525121424121

n = 6 (continued)

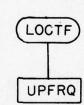
NO.	.1	2	FR 3	EQC 4	5	6		NO.	1	2	FR 3	EQC 4	5	6
361 362 363 363 366 367 367 377 377 377 377 377	444444444444444444444444444444444444444		222223333355555555566666611111333335555555666666111111222222	3355662255662233662233553355665566336633552255665566	56363556262536262335252356363556-6-536-6-335-5-356262556-6-5	656353656252636232535232656353656-5-636-3-535-3-656252656-5-		4223424 423424 4234424 4234424 4234424 42344334 43344334	444444444444444444444444444444444444444	33333333333333555555555555555555555555	555555666661	1-22662255223366336622662233223355335522552233	26-6-225-5-236262336-6-326-6-223-3-235252335-5-325-5-223-3-2	626121525121636232636131626121323121535232535131525121323121

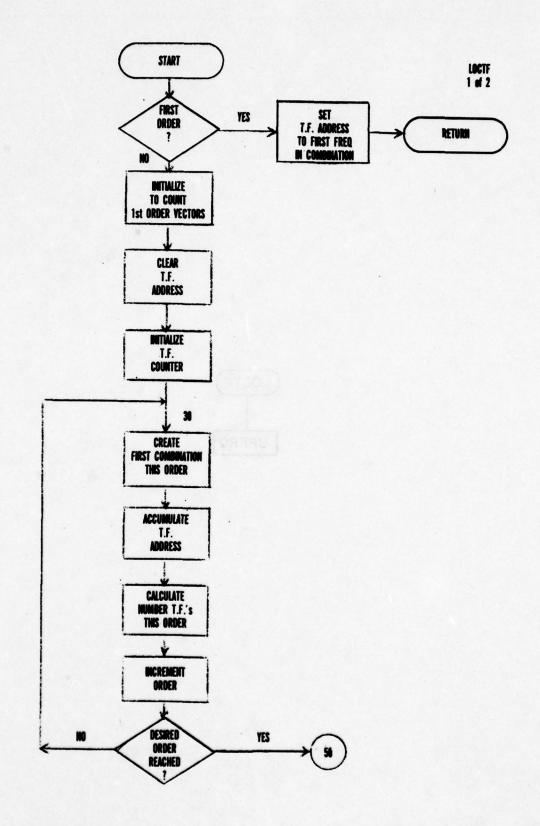
n = 6 (continued)

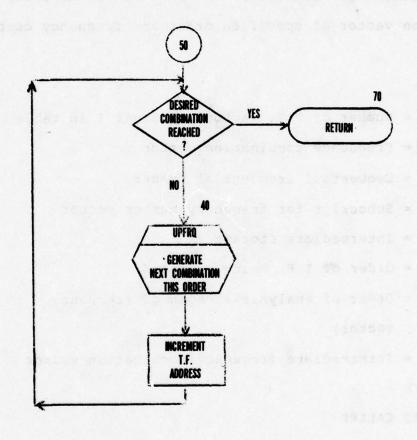
NO.	1	2	FR 3	EQC 4	5	6	No.	1	2	FR 3	EQC 4	5	6
4823 4844 4854 4874 4894 4994 4995 5000	50556555555555555555555555555555555555		22222233333344444466666661111113333333344444466666661111112222222	334466224466223366223344334466114466113366113344224466114466	463634462624362623342423463634461614361613341413462624461614	646343646242636232434232646343646141636131434131646242646141	54123 5423 5445 5555 5555 5555 5555 5555 55	50555555555555555555555555555555555555	333333333333334444444444444444444444444	444446666661111112222222333336666666111112222223333333444444	112266112244223366113366112266112233223344113344112244112233	261612241412362623361613261612231312342423341413241412231312	626121424121636232636131626121323121434232434131424121323121

n = 6 (continued)

		ole	FR	EQC 4						FR	EQC 4		
NO.	₃ I	2	3	4	5	6	NO.	1	2	3	4	5	6
6012366066606666666666666666666666666666	000000000000000000000000000000000000000		22222233333344444455555555	3344552244552233552233443344554455335533442244554455	45353445252435252334242345353445-5-435-5-334-4-345252445-5-4	545343545242535232434232545343545-4-535-3-434-3-545242545-4-	661 662 663 664 665 666 667 670 671 672 673 674 675 676 681 683 684 685 686 687 689 690 691 702 703 704 705 707 710 711 711 711 711 711 712	000000000000000000000000000000000000000	33333333333334444444444444444444444444	4444445555555	1-2255-12244223355-13355-122551-2233223344-13344-12244-12233	25-5-224-4-235252335-5-325-5-223-3-234242334-4-324-4-223-3-2	525121424121535232535131525121323121434232434131424121323121







NAME · LOCTF

TYPE: FUNCTION

GENERAL PURPOSE:

Calculates file 24 record number of a modified transfer function vector of specified order and frequency combination

VARIABLES:

C = Number of T.F. vectors of order I in table

FRQS = Frequency combination vector

I = Sequential frequential number

J = Subscript for frequency number vector

I. = Intermediate storage

M = Order of T.F. being accessed

E = Order of analysis (length of frequency

vector)

SEC = Intermediate frequency combination vector

SUBROUTINES CALLED:

UPFRC

CALLING PROGRAMS .

CFOSS

PESCRIPTION.

Function LOCTF calculates the file 24 address of the modified transfer function vector of a specified order and frequency combination for use in computing nonlinear currents in Phase 3. Arguments transmitted to the function are: N, the order of the nonlinear current being computed: M, the order of the transfer function vector being accessed; and FRQS, the frequency combination of the transfer function vector being accessed. The calculated address is returned to the calling program in LOCTF.

During the calculation of Nth order nonlinear currents, the transfer function table contains 2 n -2 lower-order transfer function vectors. These vectors are stored according to increasing frequency combinations in the order they were created in Phase 3. The address of a particular vector is a function of its order and frequency combination, and can be calculated by finding the address of the first vector of like order and adding a displacement factor based on the frequency combination of the desired vector.

There are exactly $\binom{N}{I}$ transfer function vectors of each order I=1, 2,...M in the table. Therefore the first vector of order M is located at

LOCIF =
$$1 + \sum_{I=1}^{M-1} {N \choose I}$$

The frequency combination corresponding to this vector is the ordered set of sequential frequency numbers (1,2,...M). This frequency combination is created in the first M elements of the integer vector SEQ. Each element of SEQ is numerically equal to the frequency number it represents and the frequencies are stored in the vector so that they decrease numerically as the subscript of SEQ increases.

The displacement factor is added to LOCTF by successively updating the frequency combination while at the same time incrementing LOCTF. Subroutine UPFRC updates the frequency combination. When the calculated frequency combination matches FRCS, the frequency combination of the desired transfer function vector, then LOCTF contains the desired address.

For first order transfer function vectors, the desired address is equal to the sequential frequency number. The algorithm for calculating the address of higher order vectors is summarized as follows:

STEP 1. Initialization

- a) Set sequential frequency number to I=1
- b) Initialize SEQ subscript to J=M
- c) Clear address by LOCTF=0
- d) Initialize number of vectors of order I to C=1

STEP 2. Locate First Vector of Order I, Generate Corresponding Frequency Combination in SEQ

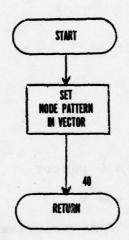
- a) Store sequential frequency number SEQ(J)=I
- b) Update address by LOCTF=LOCTF+C
- c) Calculate number of vectors of order I by C=C*L/I
- d) Increment order to I=I+1
- e) Decrement SEQ subscript to J=J-1
- f) If (J.GT.0), frequency combination not complete. Goto Step 2a.
- g) Otherwise LOCTF contains address of first transfer function vector of order M. Corresponding frequency combination is complete in SEQ. Go to Step 3.

STEP 3 Compare Desired Frequency Combination to SEQ

- a) If each FRQS element equal some SEQ element, algorithm complete. Exit with desired address in LOCTF.
- b) Otherwise go to Step 4.

STEP 4: Update Frequency Combination and Address

- a) Call subroutine UPFRQ to generate next frequency combination
- b) Increment address to LOCTF=LOCTF+1
- c) Go to Step 3



NAME: QFN

TYPE: SUBROUTINE

CENERAL PURPOSE .

Calculates the function

V (U,V)

for use by subroutine FNCTN

VARIABLES:

K = Length of Q vector

M = Number of nodes V in Q vector

C = Pattern of nodes U and V

U = Location of nonlinearity relative to base

node

V = Location of dependence relative to base node

SUBROUTINES CALLED.

NONE

CALLING PROGRAMS:

FNCTN

DESCRIPTION:

Subroutine QFN calculates the function C_M^K (U, V) which determines the pattern of node voltages for each transfer function product computed in subroutine FNCTN. Arguments transmitted to the subroutine are: K, the number of voltages in the transfer function product; M, the number of voltages from node V in the product: U and V, the nodes from which the voltages are derived (relative to the base node of the nonlinear device). The node pattern is returned to the calling program in the integer vector C.

The function Q_M^K (U, V) insures that each transfer function product calculated by subroutine FNCTN contains K voltages, where the first K-M voltages are derived from node U and the remaining M voltages are derived from node V. The node pattern is stored in the first K elements of the integer vector Q such that the first K-M elements are equal to U and the remaining M elements are equal to V.

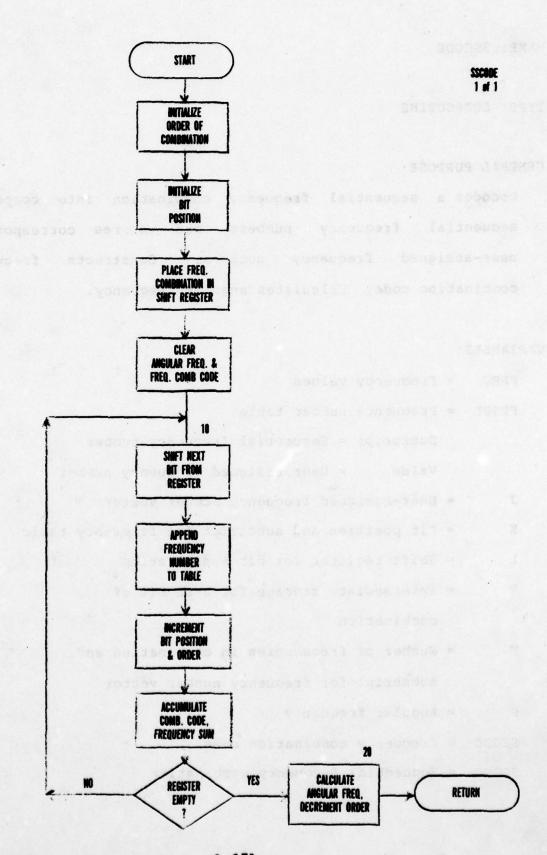
The integers U and V represent the locations of the voltages in each transfer function relative to the base node of the nonlinear device. The base node is represented by 1 in the C

vector, base+1 is represented by 2, and so on.

As an example, consider a nonlinear current which is a function of the voltages at the base node and base node+1. Then U=1, V=2, and the complete Q K M (U, V) function up to sixth order (the largest order allowed in NCAP) is tabulated as follows:

ants de la	K=2	K=3	K=4
M=0.	$Q_0^2 = (1,1)$	$Q_0^3 = (1,1,1)$	$Q_0^4 = (1,1,1,1)$
M=1	$Q_1^2 = (2,1)$	$Q_1^3 = (2,1,1)$	$Q_1^4 = (2,1,1,1)$
M=2	$Q_2^2 = (2,2)$	$Q_2^3 = (2,2,1)$	$Q_2^4 = (2, 2/1, 1)$
M=3		$0\frac{3}{3} = (2,2,2)$	$Q_3^4 = (2,2,2,1)$
M=4		n of the sales of	$Q_4^4 = (2,2,2,2)$
M=5		n visj baktispå av	pagaskay 148 sessi
M=6		nieta ud stur, mo	l boviled eas do ad
		ampedant eat to be	seld I detail a

	K=5		K=6
M=0	$Q_0^5 = (1,1,1,1,1)$		$Q_0^6 = (1,1,1,1,1,1)$
M=1	$Q_1^5 = (2,1,1,1,1)$	V	$Q_1^6 = (2,1,1,1,1,1)$
M=2	$Q_2^5 = (2,2,1,1,1)$	And the second	$Q_2^6 = (2,2,1,1,1,1)$
M=3	$0_3^5 = (2,2,2,1,1)$		$Q_3^6 = (2,2,2,1,1,1)$
M=4	$Q_4^5 = (2,2,2,2,1)$	n dela	$Q_4^6 = (2,2,2,2,1,1)$
M=5	$Q_5^5 = (2,2,2,2,2)$		$Q_5^6 = (2,2,2,2,2,1)$
M=6	Low said to appropriate the day		$Q_6^6 = (2,2,2,2,2,2)$



NAME: SSCODE

TYPE: SUBROUTINE

GENERAL PURPOSE:

Decodes a sequential frequency combination into component sequential frequency numbers and stores corresponding user-assigned frequency numbers. Constructs frequency combination code. Calculates angular frequency.

VARIABLES:

FREQ = Frequency values

FRTPL = Frequency number table;

Subscript = Sequential frequency number

Value = User assigned frequency number

J = User-assigned frequency number vector

K = Pit position and subscript for frequency table

L = Shift register for bit manipulation

M = Intermediate storage for k-th bit of combination

N = Number of frequencies in combination and subscript for frequency number vector

S = Angular frequency

SCODE = Frequency combination code

SCOME = Sequential frequency combination

SUBROUTINES CALLED:

NONE restricted parketing for couring or little tider today

CALLING PROGRAMS:

CONTRL

DESCRIPTION:

Subroutine SSCODE decodes a sequential frequency combination into its component sequential frequency numbers, stores the corresponding user-assigned frequency numbers, constructs the frequency combination code, and calculates the angular frequency. Arguments transmitted to subroutine SSCODE are: the sequential frequency combination, SCOMB; the frequency number table, FRTBL; and the frequency values FREQ. Arguments returned to the calling program are: the user-assigned frequency numbers in the combination, J; the number of frequencies in the combination, N; the calculated angular frequency, S; and the frequency combination code, SCODE.

The relationship between the sequential frequency combination and the frequency combination code is that if the Kth sequential frequency is in the combination (i.e., the Kth bit of SCOMB is 1), then the user-assigned frequency number FRTBL(K) is to be included in the code (i.e., the FRTBL(K)th bit of SCODE is 1). Similarly the Kth frequency value, FREQ(K), is included in the angular frequency, S, if the Kth sequential frequency appears in the frequency combination.

The method employed by subroutine SSCODE is to extract each

sequential frequency number from SCOMB and, using the frequency number table FRTBL, to combine the corresponding user-assigned frequency numbers and values into the frequency combination code and angular frequency. The sequential frequency numbers are extracted from the combination by using a simulated shift register L, which is initially loaded with the sequential frequency combination SCOMB. A frequency number is extracted from SCOMB by shifting the right-most bit from the register into the integer variable M. At the same time, the integer K maintains a count on the bit position being examined.

If the Kth bit is set (M=1), then the Kth user-assigned frequency number is saved by J(N) = FRTBL(K), the FRTBL(K)th bit of the frequency combination code is set to 1, and FREQ(K) is added to the angular frequency. The calculations are complete when the shift register is empty.

The complete algorithm of subroutine SSCODE is summarized as follows.

STEP 1: Initialization

- a) Frequency count: N=1
- b) Bit position counter: K=0
- c) Load shift register: L=SCOMB
- d) Angular frequency: S=0
- e) Frequency combination code · SCCDE=0

STEP 2: Extract and Store Kth bit of Sequential Combination

- a) Increment bit position count K=K+1
- b) Shift right-most bit from register and store in M:
 M=L
 L=L/2

..

M=M-L*2

c) Save user-assigned frequency number J(N)=FRTBL(K) and increment frequency count by N=N+M

STEP 3: Map Kth Bit to SCODE and S

a) Set FRTEL(K)th bit of frequency combination code according to value of M:

SCODE = SCODE + 2 ** (FRTBL(K) -1) * M

b) Unclude Kth frequency value in angular frequency according to value of M:

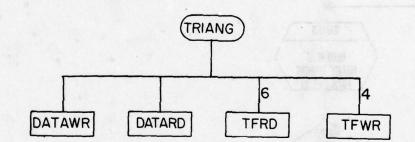
S = S + FREQ(K) * FLOAT(M)

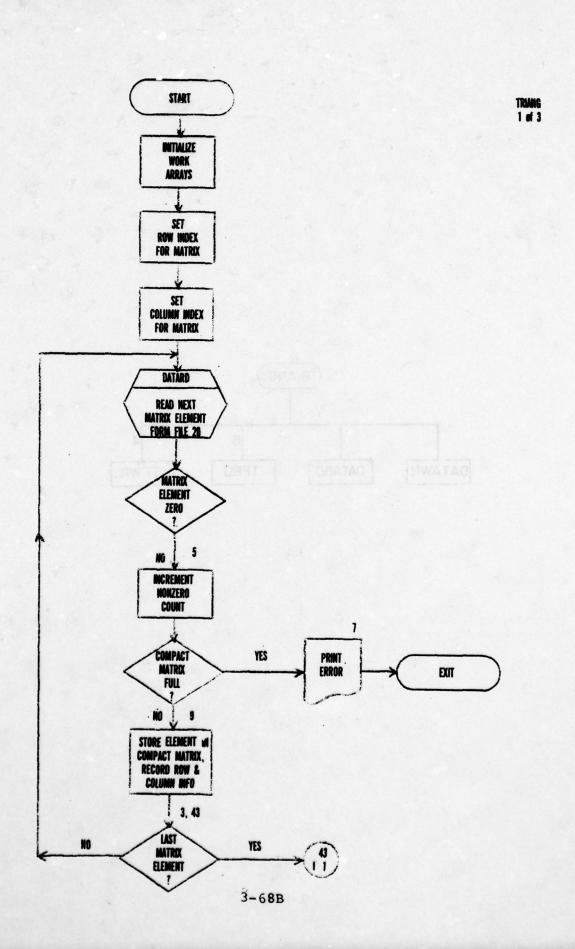
STEP 4: Test Shift Register

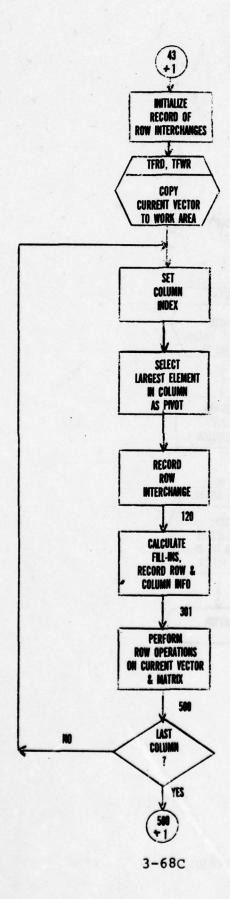
- a) If L=0 (shift register empty), frequency combination code is complete. Go to Step 5.
- b) Otherwise go to Step 2.

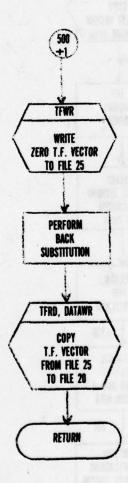
STEP 5: Complete Angular Frequency Calculation

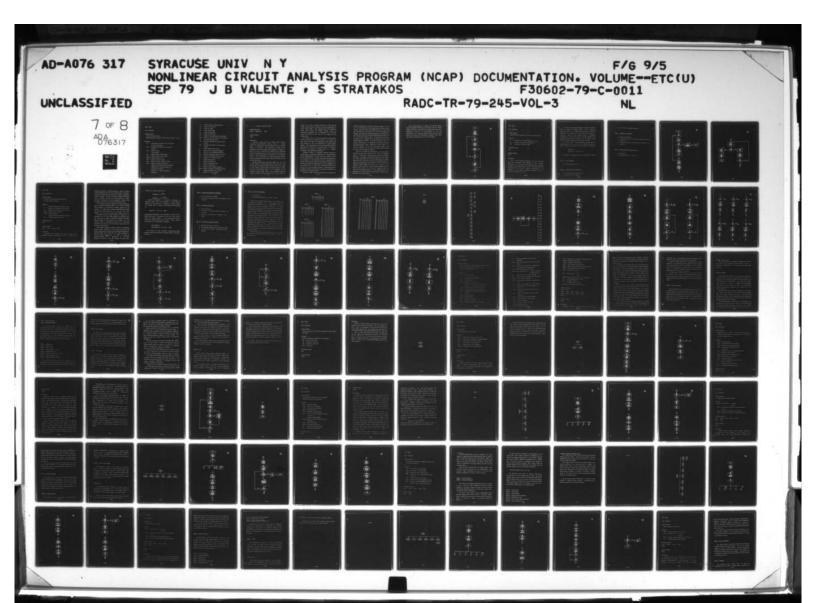
- a) Multiply angular frequency by 2
- b) Decrement frequency count N=N-1
- c) Exit











NAME: TRIANG

TYPE: SUBROUTINE

GENERAL PURPOSE:

Solves the matrix equation

[Admittance Matrix] * [Transfer Function Vector] = [Current Vector]

without continuition woll a

VARIABLES:

AM = Intermediate calculation for row operation
and back substitution

BIGX = Intermediate storage for pivot element selection

CCUR = Complex current

vector element

CUR = Complex current vector element

CUR1 = Complex current vector element

CXZERO = Complex constant $(\emptyset.,\emptyset.)$

DIM = Length of U and JU arrays.

IBIG = Pointer for pivot element selection

II = Column index for triangularization

III = Index for IHOW array

IP = Pivot row number

IPT = Pivot element number

IS = Index for IHOW array

ISS = Index for IHOW array

IT = Number of row operations to be performed

ITOT = Number of fill-ins created by row operation

IU = Row identification vector

JIP = Relative disk address of current vector element

JNR = Relative disk address of current vector element

JU = Column identification vector

Kl = Pointer for row undergoing operation

K2 = Pointer for row undergoing operation

LOC = Address of admittance matrix on file 20

LPE = Length of admittance matrix in words

LT = Length of T.F. vector in words

Ll = Pointer for pivot row

L2 = Pointer for pivot row

N = Order of admittance matrix

NN1 = Index for IHOW array

NR = Number of row undergoing operation

NZERO = Number of nonzeroes in admittance matrix

PE = Complex admittance matrix element

T = Complex T.F. vector element

SUBROUTINES CALLED:

DATARD DATAWR TFRD TFWR

CALLING PROGRAMS:

PHASE3

DESCRIPTION:

Subroutine TRIANG solves the complex matrix equation [Admittance Matrix]*[Transfer Function Vector] = [Current Vector]. The mathematical technique employed is that of Stoner's Method as described in "The Application of Sparse Matrix Techniques to Logisitcs Related Linear Programming Problems" by Barton and Lloyd. The method is to collapse the full admittance matrix into compacted form so that only the non-zero elements are stored in core and manipulated in the subroutine. The compacted matrix is then transformed into upper triangularized form by Gaussian elimination and the desired transfer functions are found by back substitution.

Arguments transmitted to subroutine TRIANG are: LOC, the starting address of the admittance matrix on disk file 20; LPE, the length of the admittance matrix in words: and LT, the length of the transfer function vector (and current vector) in words. The order of the admittance matrix, N, is transmitted to the subroutine through global common. The admittance matrix is stored by columns on disk file 20, while the current vector is

stored in the first LT words of disk file 25. The calculated transfer function vector is written to disk file 20 beginning at record number DATREC. Each complex element of the admittance matrix, current vector, and transfer function vector is individually addressable and is represented by two decimal numbers, the first for its real part and the second for its imaginary part.

The program code of subroutine TRIANG was taken directly from the main program and subroutine TRIANG of the Barton and Lloyd paper with some minor variations in input technique and pivot strategy as well as modifications to handle complex numbers. This discussion will focus on the changes required to adapt Barton and Lloyd's work to the NCAP program. The reader is directed to the previously referenced paper for a complete description of the mathematical techniques.

Subroutine TRIANG can accommodate admittance matrices of varying orders by adjusting the dimensions of U, the storage area for the compacted matrix, JU and IU, which are indexing vectors for the compacted matrix, and IHOW, which records the sequence of row operations for the elimination process. Experience with a variety of NCAP circuits has shown that sufficient storage can be allocated by setting the dimensions of U and JU (represented by the variable DIM) at seven times the number of nodes in the circuit and of IU at the number of nodes+1. Although the size of the IHOW array cannot be predicted as accurately, a dimension of IHOW(2*DIM) will generally suffice.

Subroutine TRIANG begins by collapsing the full admittance

matrix to compacted form, storing the non-zero elements by rows in U and recording the row and column indices in the IU and JU vectors. First the nonzero count NZERO and the row count IROW are initialized to zero and the U, IU, JU, and IHOW arrays are cleared. The elements of the admittance matrix are read one at a time from disk file 20 by subroutine DATARD under the control of the row and column indices I and J. If the element is not zero, the nonzero count is incremented and the element is placed in the compacted matrix at U(NZERO). In the event that the U array is overrun (NZERO.GT.DIM), an appropriate error message is printed and program execution stops. Otherwise the column index J is stored at JU(NZERO). If this is the first nonzero element in the Ith row (I.GT.IROW), the row count is incremented and the origin of the Ith row is recorded at IU(IROW).

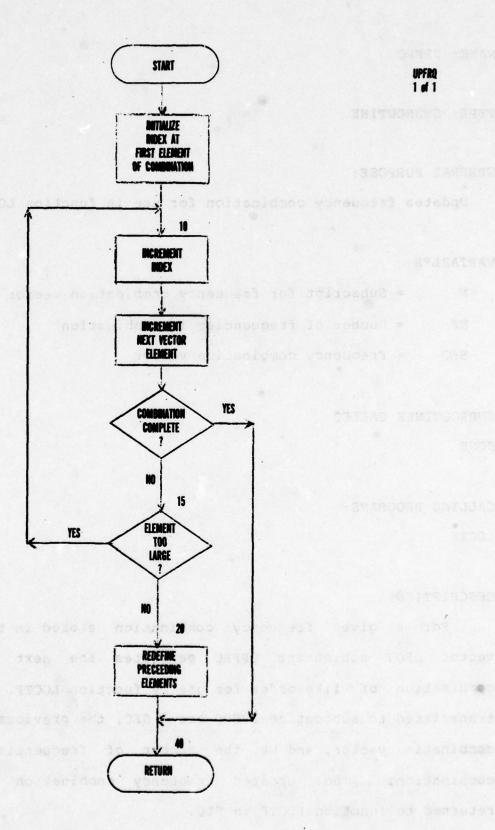
After the full matiix has been compacted, the current vector is copied to a work area on disk file 25. Then the compacted matrix is transformed to upper triangularized form by Gaussian elimination. The program code follows directly from the Barton and Lloyd paper except that the column element with the largest magnitude (sum of the squares of the real and imaginary parts) is used as the pivot element for that column. The row in which that element is located becomes the pivot row for a series of row operations which eliminate all the non-zero elements below the pivot in that column. As fill-ins are generated they are inserted in the compacted matrix U. At the completion of all row operations for a given column, the same sequence of operations is performed on the current vector.

When triangularization is complete, the transfer function vector storage area on disk file 25 is cleared and the desired transfer functions are derived by back substitution and written temporarily to file 25. At the completion of the back substitution, the resulting transfer function vector is copied from file 25 to file 20 for accessing by subsequent NCAP phases and subroutine TRIANG returns to the calling program.

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NAME: UPFRO

TYPE: SUBROUTINE

CENERAL PURPOSE:

Updates frequency combination for use in function LOCTF

VARIABLES:

Y = Subscript for frequency combination vector

NF = Number of frequencies in combination

SEQ = Frequency combination vector

SUBROUTINES CALLED:

MONE

CALLING PROGRAMS:

LOCTE

DESCRIPTION:

For a given frequency combination stored in the integer vector SEQ, subroutine UPFRC generates the next frequency combination of like order for use by function LCCTF. Arguments transmitted to subroutine UPFRC are: SEC, the previous frequency combination vector, and NF, the number of frequencies in the combination. The updated frequency combination vector is returned to function LCCTF in SEC.

The frequency combinations generated by subroutine UPFRQ are based on sequential frequency numbers. The NF sequential frequency numbers in a combination are stored in the first NF elements of the integer vector SEQ, where each element is numerically equal to the frequency number it represents. Frequencies are stored in the vector so that they decrease as the subscript of SEQ increases. Therefore for any vector element SEQ(N), the largest frequency it can contain is NF-N+1.

The method used in subroutine UPFRC is to find the first element SEC(N) for N=1,2...NF such that (SEC(N)+1).LE.(NF-N+1). Having isolated this element, the new combination is formed by:

K=SEQ(N)+1SEQ(I)=K+N-I for I=1,2...N

The algorithm of subroutine UPFRC is summarized as follows:

STEP 1: Initialization

a) Set SEQ subscript at N=0

STEP 2 Increment and Test Subscript

- a) Increment SEC subscript by N=N+1
- b) If (M.LF.NF), go to Step 3.

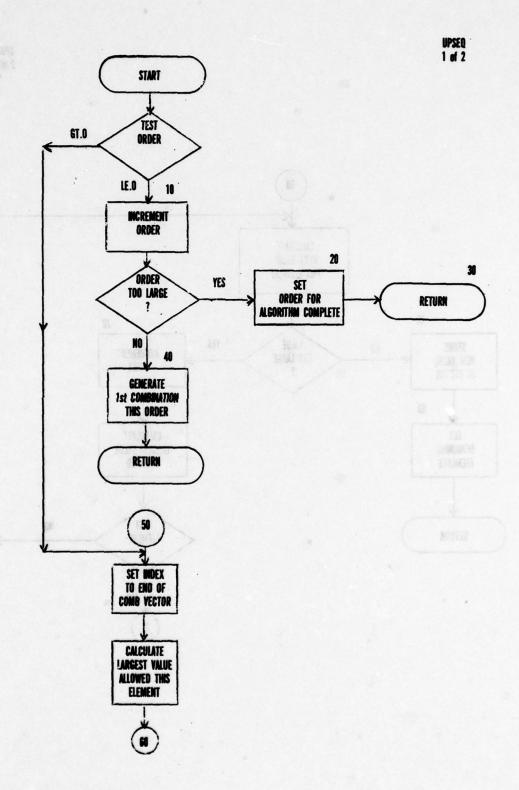
c) If (N. GT.NF), algorithm complete.

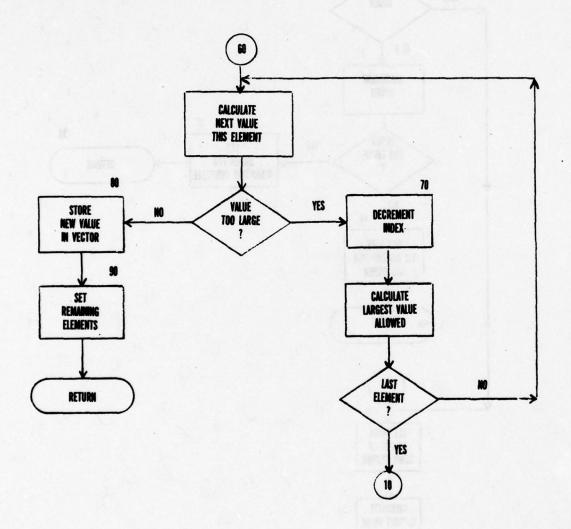
STEP 3: Update and Test SEC(N)

- a) Define SEQ(N)=SEQ(N)+1
- b) If SEC(N).GT.(NF-N+1), this value too large for SEC(N).
 Co to Step 2.
- c) If SEC(N).LE.(NF-N+1), this value O.K. for SEC(N). Co to Step 4.

STEP 4: Redefine First N Vector Elements

- a) Set K=SEO(N).
- b) Define SEC(I)=K+N-I for I=1,2...N.





NAME: UPSEQ

TYPE: SUBROUTINE

GENERAL PURPOSE:

Generates frequency combination vector and maintains order of analysis

VARIAPLES:

I = Index for frequency combination vector

M = Largest valid frequency code for SEC(I)

MM = Largest valid frequency code in vector

MXORD = Maximum order of analysis

CRDER = Order of analysis

(length of frequency combination vector)

SEQ = Frequency combination vector

SUPROUTINES CALLED.

NONE as offereness essential vest dand or recovered at the cost exe

CALLING PROGRAMS:

CONTRL CUPIC CURIDS PRSET

DESCRIPTION ·

Subroutine UPSEQ maintains the order of analysis and generates frequency combinations for use in Phases 3 and 4.

Arguments transmitted to subroutine UPSEQ are: MXORD, the number of input frequencies for the circuit (maximum order of analysis); the number of frequencies in the previous frequency combination (order of analysis); and SEQ, the previous frequency combination vector. Arguments returned to the calling program are: ORDER, the number of frequencies in the new combination (updated order of analysis) and SEQ, the new frequency combination vector.

The frequency combinations generated by subroutine UPSEQ are based on sequential frequency numbers. For a circuit with MXORD input frequencies, the sequential frequency numbers are the integers 1,2,...MXORD and the frequency combinations are ordered sets of these integers taken one at a time for first order (ORDER=1), two at a time for second order (ORDER=2), and so on.

The ORDER sequential frequency numbers in a combination are represented by the first ORDER elements of the integer vector SEQ. Each SEQ(I) is a power of 2 derived by mapping a sequential frequency number to the bit position of like number (i.e. the frequency number N is represented by $SEQ(I) = 2^{N-1}$). Frequencies are stored in the vector so that they increase numerically as the subscript of SEQ increases.

This method of encoding and storing frequency combinations imposes certain limitations on the numeric values of the SEQ elements which provide the basis for the algorithm implemented in subroutine UPSEQ. First, since each vector element represents a frequency number in the range 1 to MXORD, no element may numerically exceed 2 Furthermore, since the vector elements are stored so that they increase numerically, each

element has a largest possible value:

SEQ(ORDER) .LE. 2 MXORD-1

SEQ(ORDER-1) .LE. 2 MXORD-2 ,etc.

Frequency combinations are generated recursively by subroutine UPSEQ. The first combination of a given ORDER is defined to be the ordered set of sequential frequency numbers (1,2...ORDER) which is mapped to the elements of SEQ according to:

SEQ(I)=2 I-1 for I=1,2,...ORDER

Each successive combination of the same ORDER is derived from the previous SEQ vector. The method is to find the first element SEQ(I) for I=ORDER, ORDER-1,...l such that SEQ(I)*2 does not exceed the largest possible value for SEQ(I). Having isolated this element, the new combination is formed by:

SEQ(I)=SEQ(I)*2 SEQ(J)=SEQ(J-1)*2 for J=I+1,...ORDER

The algorithm by which successive combinations of ORDER frequency numbers are generated by subroutine UPSEQ is summarized as follows:

STEP 1: Initialize Algorithm at SEQ(ORDER)

- a) Set SEQ subscript at I=ORDER.
- .b) Define largest possible value for SEQ(ORDER) as M=2** (MXCRD-1).

STEP 2: Update and Test SEC(I)

- a) Define J=SEQ(I)*2.
- b) If (J.GT.M), this value too large for SEQ(I). Go to Step 3.
- c) If (J.LE.M), this value O.K. for SEQ(I). Go to Step 4.

previous SEO vector. The metacd is to find the first glement

STEP 3: Move Search to Next SEC Element

- a) Decrement SEQ subscript by I=I-1.
- b) Define largest possible value for SEQ(I) as M=M/2.
- c) If (I.GT.6), go to Step 2. Otherwise this ORDER complete.

STEP 4. Define New SEQ Vector

- a) Set SEQ(I)=J.
- b) Define SEQ(J)=SEQ(J-1)*2 for J=I+1,...ORDER.

Each call to subroutine UPSEC results in the generation of one frequency combination. For a circuit with MXORD input frequencies, subroutine UPSEC must be invoked 2 MXORD -1 times to generate every possible frequency combination. The integer variable CRDER, which primarily defines the number of frequencies in a particular combination, also controls the iterative calls to subroutine UPSEQ. OFDER is initialized to zero by the calling program, tested and appropriately updated by subroutine UPSEC. then retested by the calling program. For ORDER=0,1,2...MXCRD, both the calling program and subroutine UPSEQ are alerted that additional combinations remain to be generated. After all combinations for a particular CPDER have been generated, CPDER is incremented and tested in UPSEQ. If OFFER.LE.MYOFF, the first combination for the new CFDEE is created. When CFDEE exceed MXCRD, subroutine UPSFO returns the value CRDER=-1 which signals the calling program that all possible combinations for all orders have been cenerated.

I tabulation of all frequency combinations for every order of analysis up to MYOFF=6 (the largest order allowed in DCAP), is presented below:

MXORD = 2

SEQ							
NO.	adi	2	3	4	5	6	
1.	1	0	0	0	0	0	
2							
3	2	1	0	0	0	0	

one frequency coubligation. For a circuit with tyord input

MXORD = 3	the count of dang These and

	SEQ								
No.		2	3	4	5	6			
1	1	0	0	0	0	0			
2	2	0	0	0	0	0			
3	3	0	0	0	0	0			
4	. 2	1	0	0	0	0			
5	3	1	0	0	0	0			
6	3	2	0	0	0	0			
7	3	2	.1	0	0.	0.			

MXORD = 4

	SEQ									
NO.	1	2	3	EQ 4	5	6				
1 2 3 4 5 6 7 8 9 10 11 12 13 14 15	1	0	0	0.	0	000000000000000000000000000000000000000				
2	2	0	0	0	0	0				
3	3	0	0	0	0	0				
4	4	0	0	0	0	0				
5	2	1	0	0	0	0				
6	3	1	0	0	0	0				
7	4	1	0	0	0	0				
8	3	2.	0	0	0	0				
9	4	2	0	0	0	0				
10	4	3	0	0	0	0				
11	3	2	1	0	0	0				
12	4	2	1	0	0	0				
13	4	3	1	0	.0	0				
14	1 2 3 4 2 3 4 4 3 4 4 4 4 4	000011122322333	00000000011122	000000000000000000000000000000000000000	000000000000000000000000000000000000000	0				
15	4	3	2	1	0	0				

MXORD = 5

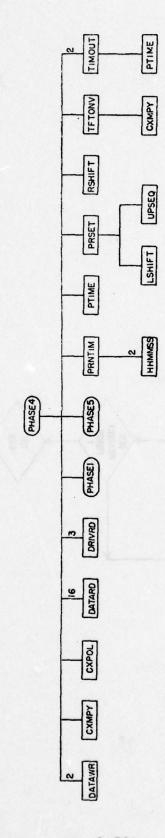
		0.000						
				S	EQ			
	NO.	d L	2	3	4	5	6	
	1 2	1 2	0	000	00	00	00	
	3	3.	0	0	0	0	0	
oppilately upo	1 2 3 4 5 6 7 8 9	5 2	0	000	0 0	0	0	
. 35.79613 87	7 8	3	1	0000	0000000	0	0	
	9	5	1 2	0	0	0	0	
	11 12 13 14 15 16 17 18	5	2	00000	0	0.0	0	
	13	5	3	0	0	0	0	
	15 16	5	4 2	0	0	0	0	
	17	5	2	1	0	0	0	
	19	4 5	3	110	000	0	0	
	21 22 23 24	5 4 5 5	3 3 4	1 2 2 2	0000	0000	0000	
	25 26	5	4 3	3 2	0	0	0	
	19 20 21 22 23 24 25 26 27 28 29 30 31	1234523453454553454554555455555	22233422233443344444	222322333	1 1 2 2	0 0 0 0 1	00000	

MXORD = 6

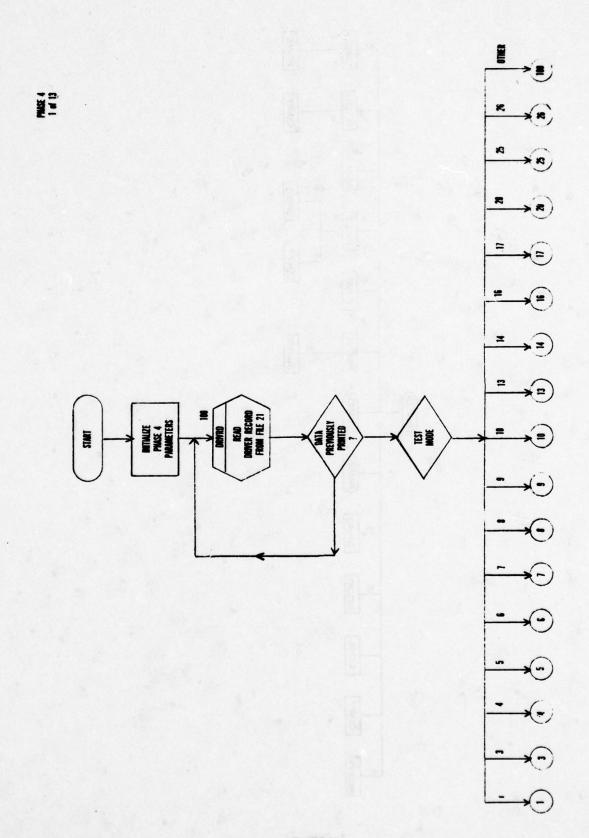
йо•	1	2	3	EQ 4	5	6	NO.	1	2	3	EQ 4	5	6
1 2 3 4 5 6 7 8 9 10 1 12 13 14 15 16 17 18 19 20 1 22 3 4 25 26 27 28 9 30 31 32	12345623456345645656634564565664	0000001111122223333445222233334453	000000000000000000000000000000000000000	000000000000000000000000000000000000000	000000000000000000000000000000000000000	000000000000000000000000000000000000000	33 34 35 36 37 38 39 40 41 42 44 45 46 47 48 49 50 50 50 50 50 50 60 60 60 60 60 60 60 60 60 60 60 60 60	5 6 5 6 6 5 6 6 6 4 5 6 5 6 6 6 5 6 6 6 6	33445445533344544554455544555555	22222333422222233344333344444	0 0 0 0 0 0 0 0 0 1 1 1 1 1 1 1 1 1 2 2 2 2	000000000000000000000000000000000000000	900000000000000000000000000000000000000

PHASE 4

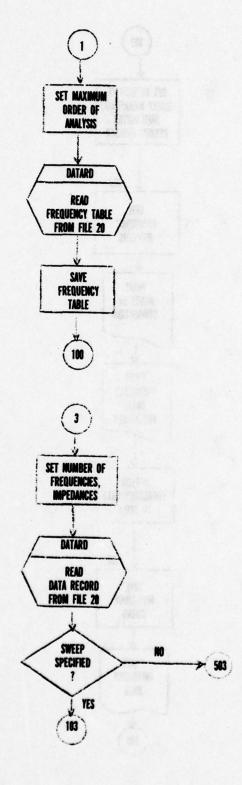
HHMMSS PRNTIM PRSET TFTONV



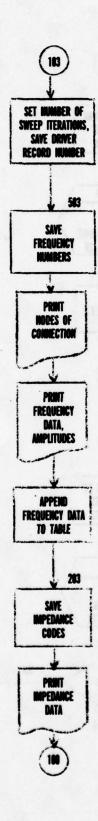
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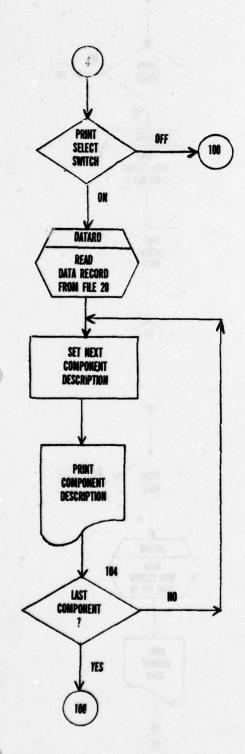


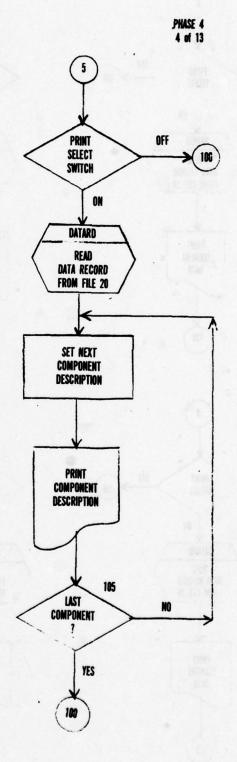


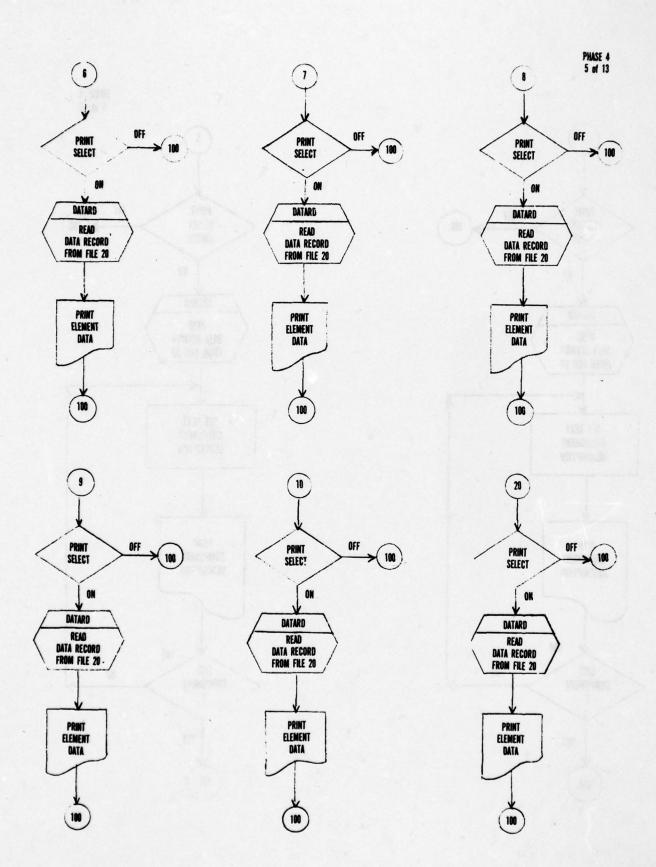


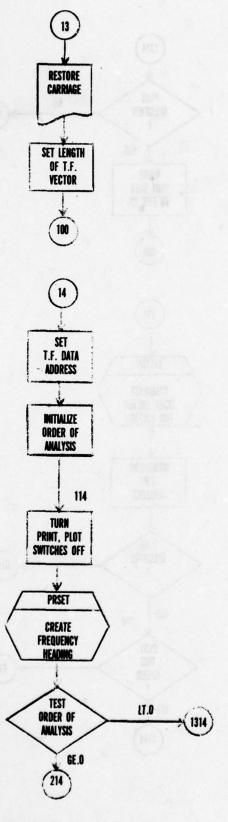


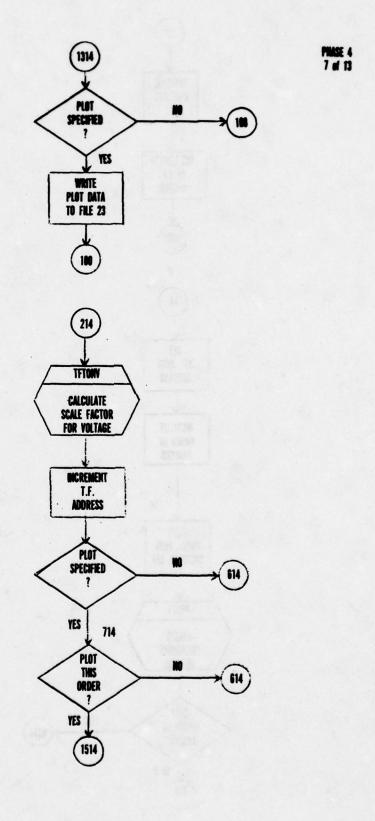


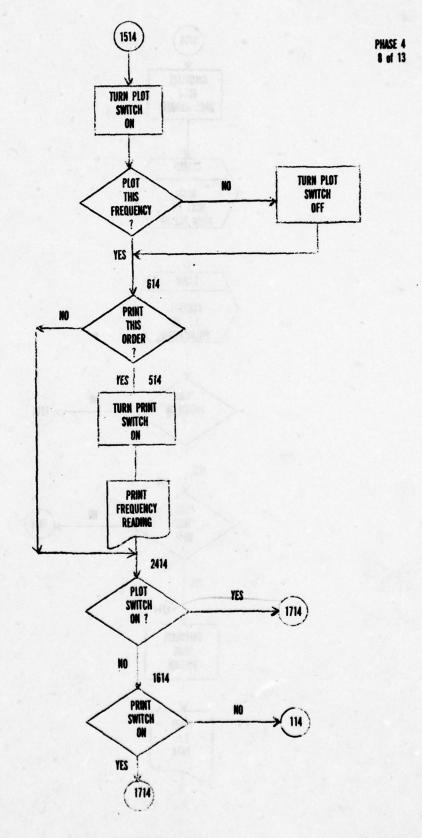


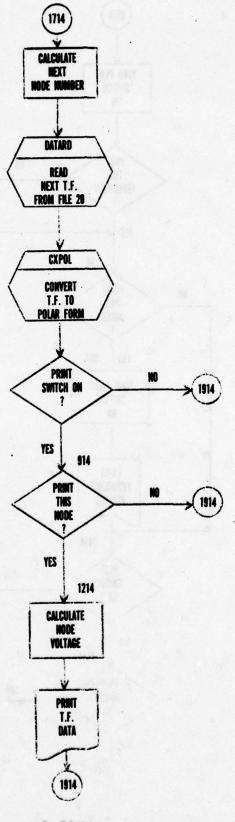




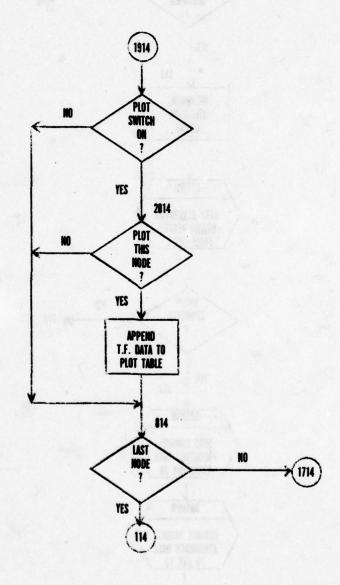


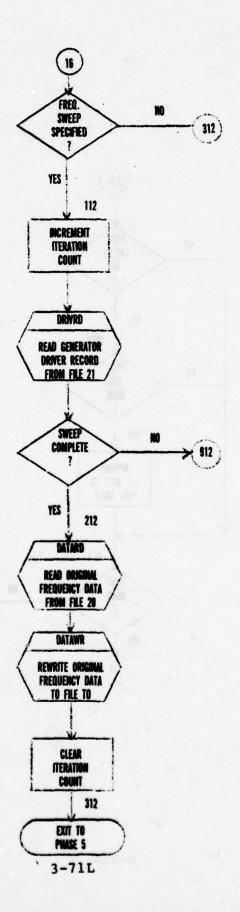




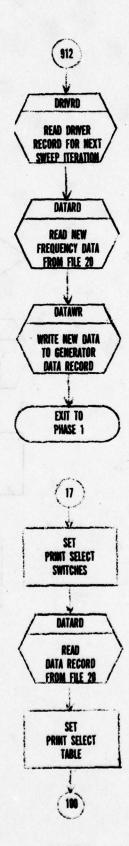


PHASE 4 9 of 13



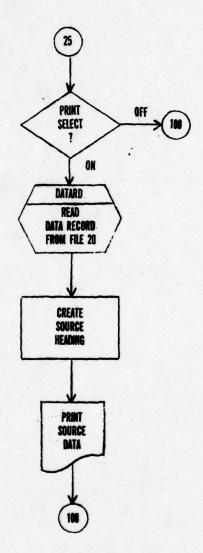


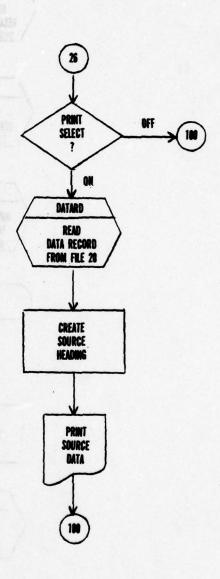
PHASE 4 11 of 13



3-71M

PHASE 4 13 of 13





MAME: PFASE4

TYPE: SUPECUTINE

CENEFAL PUPPOSE

Prints input data, transfer functions and node voltages for the circuit analysis. Controls frequency sweeping. Tabulates plot data.

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VICIATION

""FL = /mplitude values

F = Sur of frequency values

TTC = User-assigned frequency numbers

FFFC = Frequency values

FPTPI = Frequency number table

Subscript = Sequential Stequency number

Value = User-assigned fraguency number

CFRFC = Frequency values for single concrator

ICCUP = Linear and nonlinear component type

IFP = User-assigned number of frequency to be

rlotted

IMPC = Impedance value

IMTTC = Impedance code

IMCD = Fodes to be plotted

IOPD = Orders to be plotted

ICUT = Logical unit number of line printer

JADD = Address of frequency data for sweep restoration

JRFCC = Intermediate storage for driver file record
number

vFECC = Left-justified alphanumeric representation
 of user-assigned frequency number

ICCT = File 20 address of T.F. vector

IT = Length of each T.F. vector in words

"AXIT = Total number of frequency sweep iteration's

"MCDF = "aximum order of analysis

MCCMP = Length of linear/nonlinear component data
area

MIPP = Mumber of nodes specified in print select

PF = Total number of defined frequencies in circuit

TPPEC = Number of frequencies defined for given cenerator

WIMPS = Number of impedances defined for given generator

MIT = Frequency sweet iteration counter

MCDAS = Modes specified in print select

MCFEF = Order of analysis

"PLOT = Furber of plot specifications

MVAPPL = Arplitude table

TVMIC = Angle of node voltage

MVEC = Number of T.F. values for ordinate of plot

NVFREC = Frequency value table

NVMAC = Magnitude of node voltage

NVCLT = Node voltage (rectangular form)

ORDPR = Number of orders specified in print select

ORDRS = Crders specified in print select

PLSW = Plot switch = Off

 $\mathbf{i} = \mathbf{c} \mathbf{n}$

PRSW = Print switch: 0 = Off

1 = Cn

SCFACT = Scale factor for T.F. to M.V. calculation

SFC = Frequency combination vector

SFT = Frequency heading

T = Transfer function element

VECTOR = T.F. values for ordinate of plot

SUPROUTINES CALLED

CXI:PY CXPCL

DATAFD DATAWR DRIVED PHASE1 PHASE5

PRNTIM PRSET PTIME PSEIFT TETONV

TIMOUT TIMOUT

CALLING PROGRAMS

PEASE3

DESCRIPTION .

Subroutine FMASE4 is the primary routine of NCAP's output

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phase. It prints the output from a circuit analysis, tabulates plot data for use in a subsequent phase, and controls frequency sweeping. These functions are performed under the control of a driver file scan.

The standard printed output consists of all input values for each circuit element, as well as all scaled nonlinear transfer functions and node voltages resulting from a circuit analysis. The printing of circuit element data is controlled by the print select CV/OFF switch, which is initialized ON at the beginning of PEASF4 and subsequently reset whenever a print select driver record is processed. When a circuit element record is encountered in the driver file scan, the print select switch is interrogated. If the switch is ON, the circuit element data is printed if the switch is OFF, the printing of element data is suppressed.

The End Circuit driver record initiates printing of the transfer functions and node voltages under the control of the print select NCDE and CRDEP specifications. If no print selection is specified in the input stream, the complete transfer function vector and the corresponding node voltages are printed, one node per line, for each frequency combination at which the circuit was analyzed. The transfer functions and node voltages are output in both rectangular and log polar form in the order they were calculated by Phase 3.

If node and/or order print selection is specified, a table of selected nodes and orders is created during the processing of the print select driver record. When the End Circuit is

encountered, only those transfer functions corresponding to the nodes and orders in the print select table are printed.

Subroutine PHASE4 begins with initialization of phase parameters. The frequency counter is cleared by NF=0, the print select ON/OFF switch is turned ON by LSTOFF=0, and the logical unit number for the printer is stored in IOUT. The driver record number is initialized to DRVFEC=1 and the print select node and order counters, NDPR and CEDPR are cleared.

A driver record is read from file 21 by subroutine DFIVED. If MISC(7)=1, the printing of data associated with the driver record is by-passed, and control returns to the driver file scan. Otherwise subroutine PHASE4 responds to the MODE of the record as follows:

MCDE=1 Driver File Header

The total number of input frequencies for the circuit is set by MXCRD=MISC(2) and the frequency table is read from file 2% by subroutine DATARD and stored in the common data buffer BUFF. The frequency values are transferred from BUFF(1)-(10) to the decimal vector FPEC and the frequency numbers are transferred from BUFF(11)-(20) to the integer vector FPTPL. Control then returns to the driver file scan.

MODE=13 Start Circuit

The printer carriage is restored in preparation for output, the length of each transfer function vector in words is calculated as LT=2*MISC(4), and control returns to the driver file scan.

MODE=3 Cenerator

The generator data record is read from file 20 by subroutine DATARE and stored in the common data buffer BUFF. The number of frequencies and impedances associated with the generator are set by WFREQ=FISC(1) and NIMPS=MISC(2). The number of frequency sweep iterations is defined as NIT=MISC(6) and the generator driver record number is saved in JRECO. The user-assigned frequency numbers are transferred from BUFF(31)-(40) to the integer vector FN0.

The positive and negative nodes of connection are printed from MISC(3) and MISC(4) and the frequency number, frequency value, and complex amplitude for each of the NFREQ input frequencies are printed. The frequency values and amplitudes are appended to the tables NVFPEQ and NVAMP to be used in converting transfer functions to node voltages.

For each impedance in the data record, the impedance code is transferred from the data buffer to IMTPC, and the impedance code and complex admittance are printed. After all of the impedance data has been output, control returns to the driver file scan.

MODE=4 Linear Components

MCDE=5 Nonlinear Components

If the print select switch is off (LSTCFF.NF.S), the printing of component data is by-passed and control returns to the driver file scan. Otherwise the component data record is read from file 20 by subroutine DATARD and stored in the common data buffer PUFF. The length of the component data record is set by FCCMP=MISC(1). For each component in the data record, its type, nodes of connection, and value (or coefficient set) is printed. Control then returns to the driver file scan.

MCDE=6 Vacuum Diode

MCDE=7 Vacuum Triode

MCDE=8 Vacuum Pentode

MCDF=9 Biploar Junction Transistor

MODF=10 Semiconductor Diode

MCDF=20 Junction Field Effect Transistor

MCDF=25 Linear Dependent Source

MCDE=26 Nonlinear Dependent Source

If the print select switch is off (LSTOFF.NT.3), the printing of the device data is by-passed and control returns to the driver file scan. Otherwise the device data record is read

from file 20 by subroutine DATAPD and stored in the common data buffer PUFF. The device data is printed according to a predetermined format and control returns to the driver file scan.

MODE = 17 Print Select

The number of orders and nodes specified in the print select function are stored by CREPE=MISC(1) and NEPR=MISC(2). The print select CN/OFF switch is set by LSTOFF=MISC(3) and the print select data record is read from file 2% by subroutine DATARD and stored in the common data buffer BUFF. The orders to be printed are transferred from BUFF(1)-(10) to the integer vector ORDRS and the nodes to be printed are transferred from BUFF(11)-(20) to the integer vector MCDES. Control returns to the driver file scan.

MCDE=14 End Circuit

The printing of transfer functions and node voltages is initiated by the End Circuit driver record. The transfer functions are read from file 20, converted to node voltages, and printed in the order they were calculated in Phase 3. The transfer functions are accessed by recreating every frequency combination at which the circuit was analyzed while maintaining the file 20 address of the corresponding transfer function vector.

The End Circuit processing begins by initializing the address of the transfer function vector to LOCT=MISC (2)-LT. The order of analysis is initialized to NORDR=0 and the print and plot switches are turned off by PRSW=0 and PLSW=0.

Subroutine PRSET is called to generate the next frequency combination and to form the alphanumeric frequency heading. If the last frequency combination has not been generated (NORDR.GT. 0), subroutine TFTONV is called to calculate the scale factor used in converting transfer functions to node voltages. The address of the next transfer function vector is calculated by LOCT=LOCT+LT.

When plot functions are included in the NCAP input (NPLOT. GT.0), the plot switch is turned on (PLSW=1) if the order of analysis NORDR is among the orders to be plotted. The plot switch is subsequently turned off (PLSW=0) if the frequency number is not among the frequencies to be plotted.

When print select order specifications are included in the NCAP input (ORDPR.GT.0), the print switch is turned on (PRSW=1) if the order of analysis NORDR is among the orders to be printed.

The frequency heading which contains the order of analysis, frequency value, and frequency combination is printed. If the print and plot switches are both off, further processing for the present frequency combination is by-passed, and the next frequency combination is generated.

Otherwise for each transfer function in the vector I=1, LT, 2 the corresponding node number is calculated by K=I/2+1 and the complex vector element is read from file 20 by subroutine DATARD and stored in T. The transfer function is converted to log polar form by subroutine CXPOL and stored in POLAR.

When print select node specifications are included in the NCAP input (NDPR.CT.0), the printing of the transfer function is by-passed if the node K is not among the nodes in the print select table. Otherwise the transfer function is converted to a node voltage and the node number, transfer function (in both rectangular and polar form), and node voltage are printed.

If the plot switch is on and K is among the nodes to be plotted, the transfer function is appended to the plot 'table at VECTOR(J) for J=1, NPLOT.

After every frequency combination has been generated and the printing of transfer functions is complete (NORDR.LT.0), control returns to the driver file scan.

MODE=16 End Fant was choir striams above soles soles and

The End driver record controls frequency sweeping. If frequency sweeping is not specified for the circuit (MAXIT.LT.0), program control passes to subroutine PHASES. Otherwise the frequency sweep iteration count is incremented by NIT=NIT+1 and the generator driver record is read from file 21 by subroutine PPIVPD.

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If the frequency sweep is complete (NIT.CT.MAXIT), the criginal frequency data is read from file 20 by subroutine PATAFF and restored to the generator data record by subroutine

DATAWR. The frequency sweep iteration count is reset by NIT=0 and program control passes to subroutine PHASE5.

If the frequency sweep is not complete (NIT.LE.MAXIT), the next sweep iteration is initialized as follows: the address of the generator data record is saved in JADD and the driver record for the next sweep iteration is read from file 21 by subroutine DPIVPD. The sweep data is read from file 20 by subroutine DATARD and written to the generator data record by subroutine DATAWR. A new circuit analysis is initiated by calling subroutine PHASFI.

For other values of MODE, subroutine PHASE4 returns to the driver file scan without further processing.

NAME: HHMMSS

TYPE: SUBROUTINE

GENERAL PURPOSE:

Converts clock time for NCAP timing routines to hours, minutes, and seconds.

of Time. Minutes and cadculated by multiplying the fractional

VARIABLES: Au Medadal edd politicaeu has

ATIME = Intermediate storage for time calculation

HRS = Hours of clock time

MIN = Minutes of clock time

SEC = Seconds of clock time

TIME = Clock time

SUBROUTINES CALLED:

NONE

CALLING PROGRAMS:

PRNTIM

DESCRIPTION:

Subroutine HHMMSS converts clock time for the NCAP timing routines to hours, minutes, and seconds. The time, expressed in hours, tenths, and hundredths, is transmitted to the subroutine in the decimal argument TIME. The converted time is returned to the calling program in the integer arguments HRS, MIN, and SEC.

Hours of clock time are derived by taking the integer part of TIME. Minutes are calculated by multiplying the fractional part of TIME by 60., and retaining the integer part of the result. Seconds are then found by multiplying the hundredths portion of TIME by 60., and subroutine HHMMSS returns to the calling program.



MAN'E: PPNTIM

TYPE · SUPROUTINE

CENERAL PURPOSE

Prints output from timing routines

VARIAPLES.

ICUT = Logical unit number of line printer .

LIST = Alphanumeric designator of program segements

SUM1 = Total clock time for one circuit

SUM2 = Total clock time for NCAP job

TIME = Elapsed clock time per segment

TTOTAL = Accumulated clock time per segment

SUPRCUTINES CALLED:

UHMMES

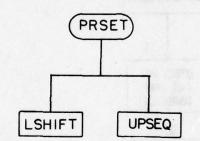
CALLING PROCEAMS

PRASE4

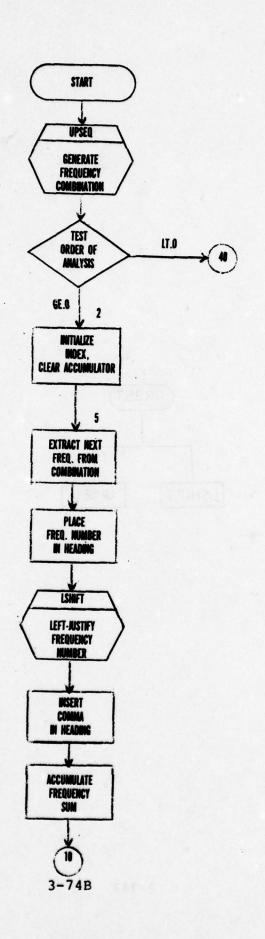
DESCRIPTION .

Subjoutine PRNTIM prints the output from the NCAP timing routines. Elapsed and accumulated times for each program segment, TIME and TTOTAL, are transmitted to the subroutine in the latelled common area TIMEFL.

First the total time for the past circuit analysis and total time for the NCAP job are accumulated in SUM1 and SUM2. Then for each program segment executed during the previous circuit analysis (TIME(I).CT.0), the elapsed and accumulated segment times are converted to hours, minutes, and seconds by subroutine HEMMSS and printed in tabular form. After each segment's time has been printed, the TIME array is zeroed and subroutine PRNTIM returns to the calling program.

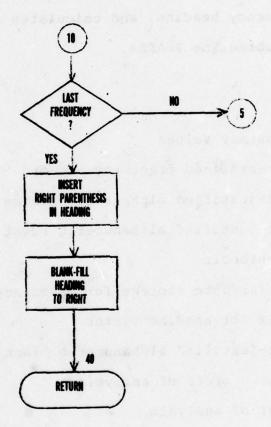


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PR SET 1 of 2

1



MAME: PRSET

TYPE · SUPROUTINE

CEMERAL PURPOSE

Controls order of analysis, forms frequency combinations, creates frequency heading, and calculates sum of frequencies for use by subroutine PHASE4.

VARIABLES:

FREC = Frequency values

FRTEL = User-assigned frequency numbers

ICMMA = Left-justified alphanumeric comma

IRPTR = Left-justified alphanumeric right

parenthesis

ISFQ = Intermediate storage for frequency number

L = Index for heading vector

LFLANK = Left-justified alphanumeric blank

MXORD = Maximum order of analysis

MORDR = Crder of analysis

MUF = Index for frequency number and value

s = Sum of frequencies

SEQ = Frequency combination vector

SET = Alphanumeric frequency heading

SUPPOUTINES CALLED

LSFIFT UPSEC

CALLING PROGRAMS:

PHASE4

DESCRIPTION:

Subroutine PRSFT forms the frequency combinations and creates the alphanumeric frequency heading for use by subroutine PMASEA. Arguments transmitted to the subroutine from the calling program are: FRFQ, the circuit's input frequency values; FRTFL, the frequency number table: MXOPD, the number of input frequencies for the circuit (maximum order of analysis): NORDR, the number of frequencies in the last frequency combination (order of analysis): and SFQ, the previous frequency combination vector. Arguments returned to subroutine PHASEA are: NORDR, the number of frequencies in the new frequency combination (updated order of analysis): S, the calculated frequency sum for the new frequency combination: SET, the alphanumeric frequency heading: and SEQ, the updated frequency combination vector.

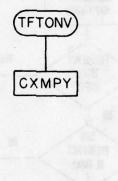
Upon entering subroutine PRSET, the next frequency combination is generated by subroutine UPSEQ. If the last possible frequency combination has already been generated (MORDP.IF.C), subroutine PRSET returns to the calling program without further processing. Otherwise NORDP represents the number of frequencies in the new combination, which is stored according to increasing sequential frequency numbers in the integer vector SEC.

The frequency sum is initialized to S=0. and the index for the frequency heading is set at L=1. The frequency heading is created by converting each sequential frequency number in the frequency combination to its user-assigned number. The user-assigned number is then converted to left-justified alphanumeric format and placed in the heading.

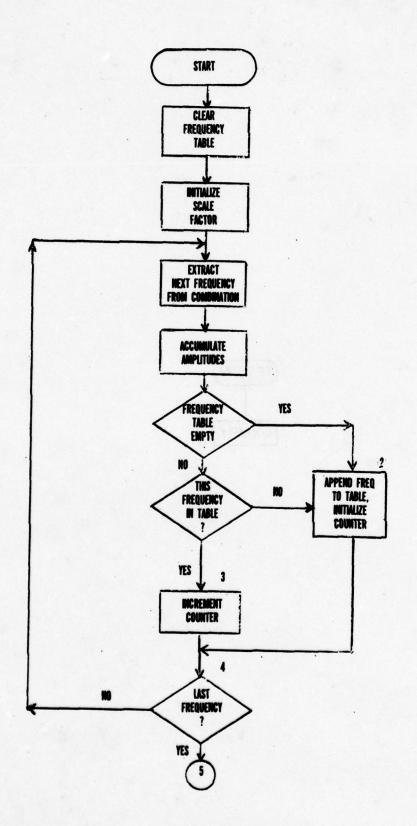
First the frequency combination vector is decoded as follows: the vector element SEC(I) is stored in ISEQ. Then ISEQ is successively shifted right one bit at a time while a count, NUM, is maintained on the number of shifts performed. When ISEQ=1, then NUM contains the sequential frequency number represented by SEQ(I).

The corresponding user-assigned frequency number is taken from the frequency table as FRTEL(NUM), stored in the frequency heading at SET(L), and converted to left-justified alphanumeric format by subroutine LSHIFT. A comma is inserted after the frequency number in the heading by SET(L+1)=ICMMA and the heading index is updated to L=L+2. Then the frequency value FREC(NUM) is added to the frequency sum.

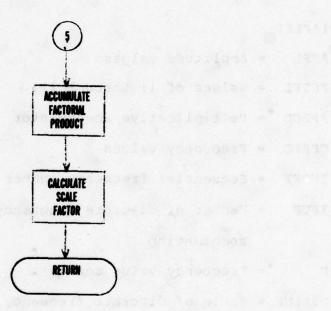
After every frequency in the combination has been placed in the heading, a right parenthesis is inserted after the last frequency number by SET(L-1)=IRPTR, the heading is blank-filled to the right, and subroutine PPSET returns to the calling program.



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TFTONY 1 of 2



NAME TFTONV

TYPE · SUBFCUTINE

CENERAL PURPOSE:

Calculates scale factor for use in converting transfer functions to node voltages

VARIATIES:

AMPL = Amplitude values

FCTFL = Values of 1! through 10!

FPRCD = Multiplicative accumulator

CFFEC = Frequency values

INDEX = Sequential frequency number

= Frequency value counters

MVALUE = Table of discrete frequency values in combination

MOPDE = Order of analysis

SCFACT = Calculated scale factor

SEC = Frequency combination vector

TFTF = Intermediate calculation

CULBOUTINES CALLED:

CXI'PY

CALLING PROGRAMS:

PHASE4

DESCRIPTION: VALUE OF THE PROPERTY OF THE PROP

Subroutine TFTCNV calculates the scale factor for use in converting transfer functions to node voltages for use by subroutine PFASE4. The arguments transmitted to subroutine TFTCNV from the calling program are. NCRDR, the number of frequencies in the frequency combination (order of analysis) GFREQ, the circuit's input frequency values: AMPI, the circuit's input amplitude values: and SEQ, the frequency combination vector. The calculated scale factor is returned to the calling program in SCFACT.

For a given frequency combination, the scale factor is a function of the amplitude at each frequency in that combination, the order of analysis, and the number of times each discrete frequency value appears in the combination. Subroutine TPTONV creates a table of the discrete frequency values in the combination and counts the number of times each value is used in the frequency combination. The discrete frequency values and the corresponding counters are stored in the vectors NVALUE and M so that M(I) counts the number of times MVALUE(I) appears in the combination. The length of the MVALUE and M vectors is accumulated in IPTP.

Subroutine TETCHV begins by clearing the MVALUE and "vectors and setting IPTR=F. The complex scale factor is

initialized to SCFACT=(1., 0.). Then for each element in the frequency combination vector, SEQ(1) through SEQ(NCPDR), the sequential frequency number represented by SEQ(I) is decoded by the function IND and stored in INDEX. Subroutine CXMPY is called to multiply the accumulated scale factor by the amplitude AMP(INDEX), storing the product in SCFACT.

If no frequency values have been placed in the table (IPTF =0), the table length is updated to IPTF=IPTF+1, the value CFREQ(INDEX) is appended to the table at FVALUE(IPTF), and the frequency counter M(IPTF) is initialized to 1.

If there are frequency values in the table (IPTR.CT.0), the value CFREC(INDEX) is compared to each MVALUE. If no match is found the new value is appended to the table as above. If a match is found, the corresponding counter M(J) is incremented.

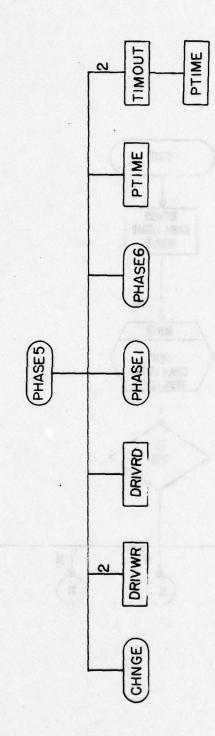
After all NORDF frequencies in the SEC vector have been processed, the values of M(I)! for every frequency counter are multiplied together and stored in FPFCD. After the intermediate calculation:

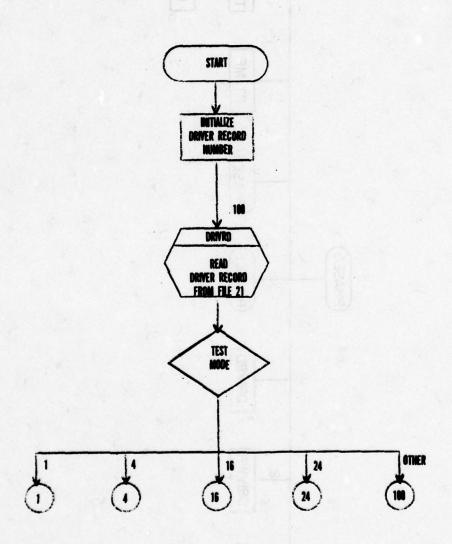
TEMP=NCRDR! / FPROD*2 NORDR-1

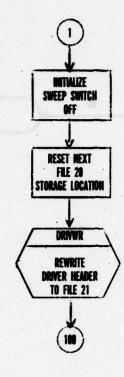
has been performed, the complex scale factor follows directly by multiplying the real and imaginary parts of SCFACT by TEMP.

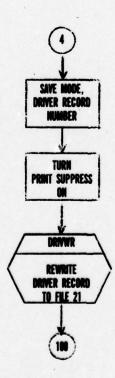
PHASE5

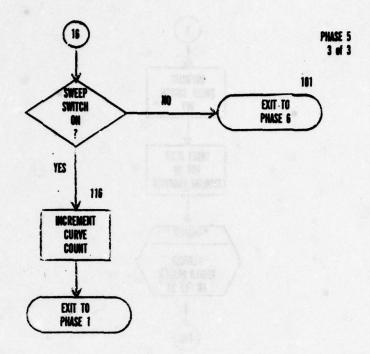
CHNGE













NAME: PHASE 5

TYPE - SUPROUTINE

CENERAL PURPOSE:

Controls linear component sweeping

VARIABLES:

IDONE = Sweep switch: 0 = No sweep

1 = Sweep

JRFCC = Intermediate storage for driver file record

number

MDESV = Intermediate storage for device mode

MCURV = Mumber of curves defined by circuit analysis

SUBROUTINES CALLED:

CHNGE DRIVED DRIVER PHASEL PHASE6

PTIME TIMOUT

CALLING PROGRAMS:

PEASE4

DESCRIPTION -

Subroutine PFASE5 controls linear component sweeping by initiating a new circuit analysis for each iteration of a component sweep in the circuit description. The method is to

scan the driver file for sweep functions which are represented by a MODE=4 (linear component) driver record followed by a MODE= 24 (linear component sweep) driver record. For each such function, the component values for the next iteration are brought forward from the sweep data record to replace the previous linear component data, and a new circuit analysis is performed beginning with Phase 1.

The driver file scan begins by initializing the driver file record number to DRVREC=1. A driver record is read from file 21 by subroutine DRIVRD, and subroutine PHASES responds to the MODE of the record as follows:

MODE=1 Driver File Header

The sweep switch is turned off by IDONE=0, indicating that no sweep functions have been processed. The file 20 storage address is reset by MISC(1)=MISC(4), allowing the admittance matrices and transfer function vectors for the next circuit analysis to be written over those from the last analysis. The updated driver file header is rewritten by subroutine DRIVWR and control returns to the driver file scan.

MODE=4 Linear Components

The device mode and driver record number are saved by

MDESV=4 and JRECO=DRVREC-1 in preparation for a possible sweep. MISC(7) is set to 1 to suppress printing of the old component values in PHASE4. The updated linear component driver record is rewritten by subroutine DRIVWR and control returns to the driver file scan.

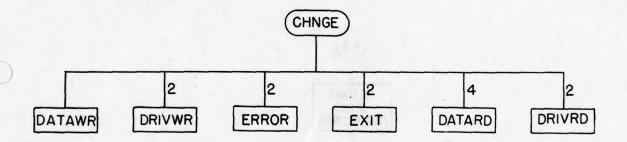
MODE=24 Linear Element Sweep

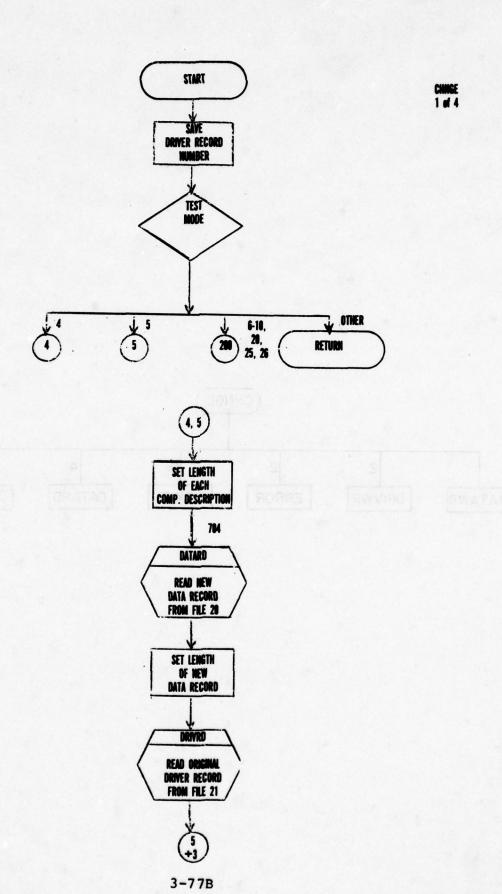
The sweep switch is turned on by IDONE=1 and subroutine CFNGE is called to replace the previous component data with the data values for the new sweep iteration. The saved device mode is changed to MDESV=24, causing any remaining sweep records associated with the present linear component definition to be by-passed in the driver scan (see subroutine CHNCF), and control returns to the driver file scan.

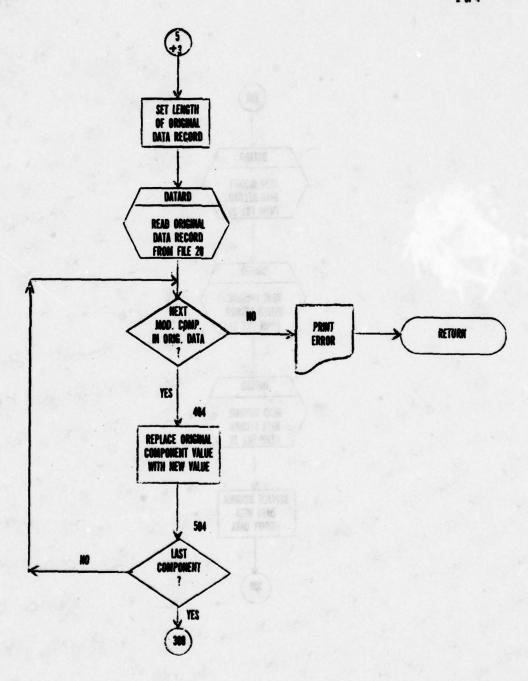
VCDE=16 Fnd

If the sweep switch is on (IDONE=1), the curve count is incremented by NCURV=NCURV+1 and a new circuit analysis is initiated by calling subroutine PMASE1. If the sweep switch is off (IDONE=6), program control passes to subroutine PMASE6.

For other values of MODE, control returns to the driver file scan without further processing by subroutine PHASES.











MAME: CHNGE

TYPE: SUPPOUTINE

GENERAL PURPOSE:

Performs data replacement for linear component sweep and device modification

VARIABLES.

PUFF = I/C buffer for modified data

DUFFC = I/C buffer for original data

JRCSV = Address of modify driver record

JPFCC = Address of driver record of device being

rodified

MDESV = Mode of device being modified

MCMP = Length of modified component data area

NCMPO = Length of original component data area

!'W = Length of each component description

SUPECUTINES CALLED

DATABL DATAWR DRIVED DRIVER ERROR

CALLING PROCEAMS.

PHASES PHASE6

DESCRIPTION :

Subroutine CHNGE performs the data replacement for linear component sweeping and device modification performed in Phases 5 and 6. Arguments transmitted to the subroutine are: MDESV, the mode of the device being modified and JRECO, the driver record number of the device being modified. The modify or sweep driver record is transmitted to subroutine CHNGE in the first ten words of global common.

First the driver record number of the sweep or modify record from which the data is derived is set by JRCSV=DRVREC-1. Then subroutine CHNGE interrogates the mode of the device being modified, responding to the value of MDESV as follows:

component driver record is rewritten by

MDESV=4 Linear Components

MDESV=5 Nonlinear Components

The length of each component description in the data record is set to NW=4 for linear components or NW=13 for nonlinear components. The sweep or modify data record is read from file 20 by subroutine DATARD and stored in the common data buffer BUFF. Its length is saved by NCMP=MISC(1).

The driver record number is reset to DRVREC=JRECO and the original component driver record is read by subroutine DRIVRD. Then the component data record is read from file 20 by subroutine DATARD and stored in BUFFO. Its length is saved by NCMPO=MISC(1).

The type and nodes of connection for each component in the modify data record (BUFF(I), BUFF(I+1), and BUFF(I+2) for I=1, NCMP,NW) are compared against those in the original component data record (BUFFO(J), BUFFO(J+1) and BUFFO(J+2) for J=1, NCMPO, NW). If no match is found, an appropriate error message is printed and the program terminates. When a match is found, the original component value or coefficients are replaced with the modified component value or coefficients according to:

BUFFO(J-1+K)=BUFF(I-1+K) for K=4, NW

The component driver record is updated by MISC(7)=0 to enable printing of the new component data in PHASE4. The updated component driver record is rewritten by subroutine DRIVWR and the updated component data record is written to file 20 by subroutine DATAWR. The sweep or modify driver record is disabled by writing a no-operation (MODE=2) record in its place in the driver file and subroutine CHNGE returns to the calling program.

MDESV=6 Vacuum Diode

MDESV=7 Vacuum Triode

MDESV=8 Vacuum Pentode

MDESV=9 Bipolar Junction Transistor

MDESV=10 Semiconductor Diode

MDESV=20 Junction Field Effect Transistor

MDESV=25 Linear Dependent Source

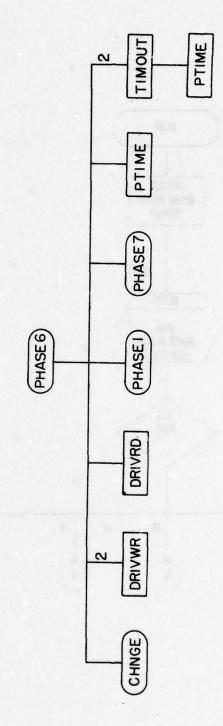
MDESV=26 Nonlinear Dependent Source

The modify data record is read from file 20 by subroutine DATARD and stored in BUFFO. The driver record number is reset by DRVREC=JRECO and the original device driver record is read from file 20 by subroutine DRIVRD.

The component driver record is them updated by MISC(7)=0 to enable the printing of the new device data in PHASE4. The updated device driver record is rewritten by subroutine DRIVWR and the modify data record is written from BUFFO to the original device data record on file 20 by subroutine DATAWR. The modify driver record is disabled by writing a no-operation (MODE=2) record in its place in the driver file and subroutine CHNGE returns to the calling program.

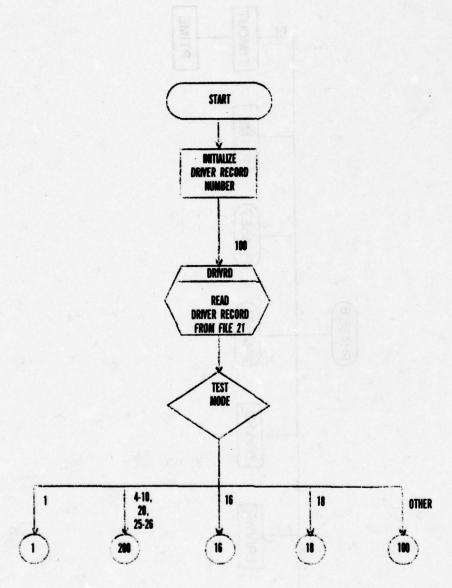
Any other value of MDESV encountered by subroutine CHNGE indicates a misplaced modify function. An appropriate error message is printed and the program terminates.

PHASE 6



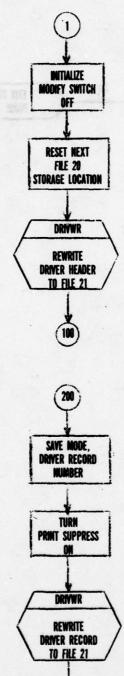
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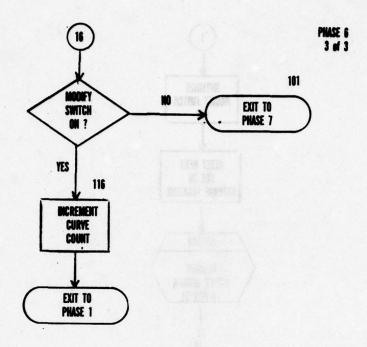


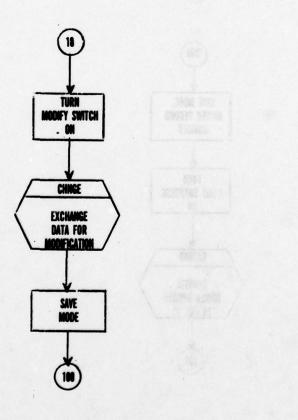




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MAME: PPASE6

TYPE: SUFFCUTINE AND THE STATE OF THE STATE

CENERAL PURPOSE: Designation and analysis of the sound of

Controls device mudification

VARIAPLES

IDCND = Modify switch: 2 = No modify

1 = Modify

ewollor as brober and lo

JEECC = Intermediate storage for driver file record

number

MTESV = Intermediate storage for device mode

MCUFV = Number of curves defined by circuit analysis

SUFFOUTIFFS CALLER:

CUNCE PRIVED DRIVER PRASE1 PHASE7

PTIME TIMEUT

CALLINC PROCEAMS

DEASES.

PESCRIPTION .

Cubroutine PHASTE controls linear and nonlinear element redifications (except for generator modifications). It scans the driver file for modify functions which are represented by a

device driver record (MODE=4,5,6,7,8,9,10,20,25,26) followed by a modify record (MODE=18). For each such function, the modify data record is brought forward to replace the original device data, and a new circuit analysis is performed beginning with Phase 1.

The driver file scan begins by initializing the driver file record number to DFVEC=1. A driver record is read from file 21 by subroutine DRIVED, and subroutine PMASE6 responds to the MODE of the record as follows:

MODE=1 Driver File Header

The modify switch is turned off by IDCNE=5, indicating that no modify functions have been processed. The file 20 storage address is reset by MISC(1)=MISC(4), allowing the admittance matrices and transfer function vectors for the next circuit analysis to be written over those from the last analysis. The updated driver file header is rewritten by subroutine DRIVWR and control returns to the driver file scan.

MCDE=4 Linear Components

MCDF=5 Honlinear Components

MCDE=6 Vacuum Diode

MCDE=8 Vacuum Pentode

MODE=9 Pipolar Junction Transistor

MODE=10 Semiconductor Diode

MCDF=20 Junction Field Effect Transistor

MODE=25 Linear Dependent Source

MCDE=26 Monlinear Dependent Source

The device mode and driver record number are saved by MIDSV=MCDE and JRECO=DRVFEC-1 in preparation for a possible modification. MISC(7) is set to 1 to suppress printing of the old device values in PMASE4. The updated device driver record is rewritten by subroutine DRIVWE and control returns to the driver file scan.

MODE=18 Modify

The modify switch is turned on by IDCNE=1 and subroutine CPNCE is called to replace the original device data with the modify data. The saved device mode is changed to MDESV=18, causing any remaining modify records associated with the present device to be by-passed in the driver scan (see subroutine CHNCE), and control returns to the driver file scan.

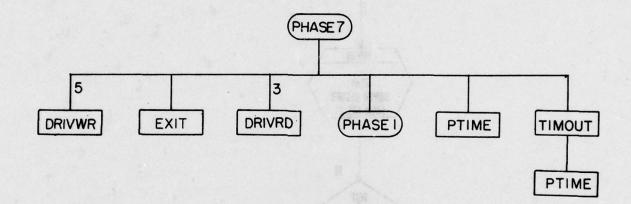
"OFF=16 End

If the modify switch is on (IDONE=1), the curve count is incremented by NCURV=NCURV+1 and a new circuit analysis is initiated by calling subroutine PHASE1. If the modify switch is

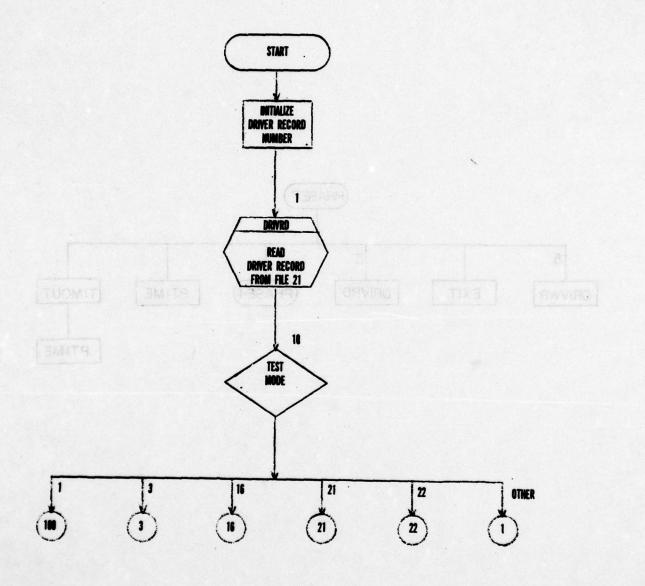
off (IDCNE=0), program control passes to subroutine PHASE7.

For other values of MODE, control returns to the driver file scan without further processing by subroutine PHASE6.

PHASE7

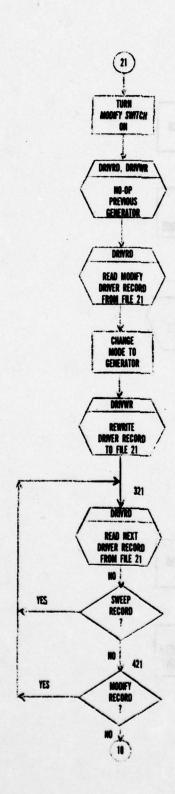


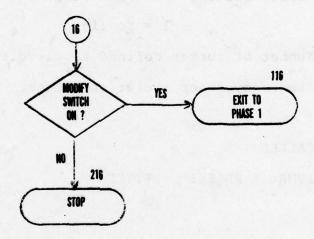
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NAME: PHASE7

TYPE: SUPROUTINE

GENERAL PURPOSE :

Controls generator modification

VARIABLES

JPECO := Address of generator driver record

MODEY = Modify switch; 0 = No modify

1 = Modify

NCURV = Number of curves defined by circuit analysis

FIT = Frequency sweep iteration counter

SUPRCUTINES CALLED

DRIVED DRIVER PHASEL PTIME

TIMCUT

CALLING PROGRAMS:

PHASE 6

DESCRIPTION .

Subroutine PHASE7, the only routine of NCAP's final phase, controls generator modification. It scans the driver file for generator modifications which are represented by a MCDE=3 (generator) driver record followed by a MCDE=21 (generator)

modification) driver record. Unlike other device modifications, generator modification does not involve data replacement. Instead, the original generator description is disabled by writing a no-operation (MODE=2) record in its place in the driver file, and the modified generator is activated by changing its MODE from 21 to 3.

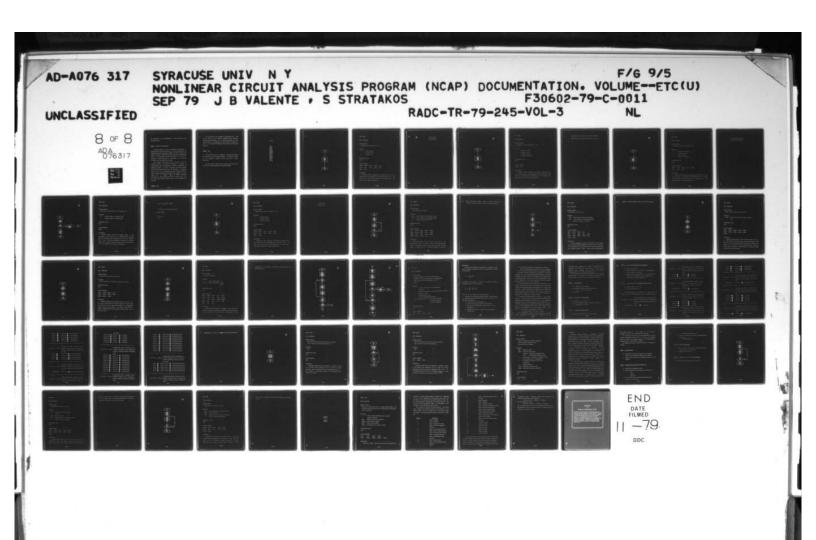
The driver file scan begins by initializing the driver file record number to DRVREC=1. A driver record is read from file 21 by subroutine DRIVRD, and subroutine PHASE7 responds to the MODE of the record as follows:

MODE=1 Driver File Header

The modify switch is turned off by MODFY=0, indicating that no modify functions have been processed. The file 20 storage address is reset by MISC(1)=MISC(4), allowing the admittance matrices and transfer functions for the next circuit analysis to be written over those from the last analysis. The updated driver file header is rewritten by subroutine DRIVWR and control returns to the driver file scan.

MODE=3 Generator

The generator driver record number is saved by JRECO=DRVREC-1 and the number of frequency sweep iterations for



the generator is set by NIT=MISC(6)+1. Control returns to the driver file scan.

MODE=21 Generator Modification

The modify switch is turned on by MODFY=1, indicating that a modify function has been encountered. The original generator is disabled by writing no-operation (MODE=2) records in place of the generator driver record and any sweep driver records associated with it. The no-operation records are written to file 21 by subroutine DRIVWR beginning at DRVREC=JRECO and proceeding through the next successive NIT records.

The generator modification driver record is read by subroutine DRIVRD. The modified generator is activated by setting MODE=3, and the updated driver record is rewritten by subroutine DRIVWR. In order to by-pass any sweep or modify driver records associated with the new generator, a secondary driver file scan is initiated by calling subroutine DRIVRD to read the record immediately following the new generator in the driver file. If MODE=12 (sweep) or MODE=21 (generator modification), control remains within the secondary scan. For any other value of MODE, control returns to the standard driver file scan.

The curve count is set to NCURV=1 in preparation for a new set of curves for the modified frequency values. MISC(5)=0 causes new plot parameters to be calculated during the next circuit analysis. The updated plot driver record is rewritten to file 21 by subroutine DRIVWR and control returns to the driver file scan.

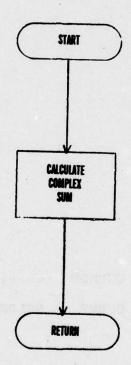
MODE=16 End

If the modify switch is on (MODFY=1), a new circuit analysis is initiated by calling subroutine PHASE1. If the modify switch is off (MODFY=0), program execution ends with a normal termination.

For other values of MODE, control returns to the driver file scan without further processing by subroutine PHASE7.

UTILITY

CXADD CXDIV CXMPY CXPOL CXSUE DATARD DATAWR DRIVRD DRIVWR ERROR FUN LSHIFT RSHIFT SSCOD2 TFRD TFWR TIMOUT



NAME: CXADD

TYPE · SUFFCUTINE

CENERAL PURPOSE .

Calculates the complex sum C = A + F

VARIABLES

A = Complex addend

Complex addend

C = Complex sum

SUPROUTINES CALLED.

MONE

CALLINC PROCRAMS:

CURGEN CURJF CURNC CURNDS CURPT

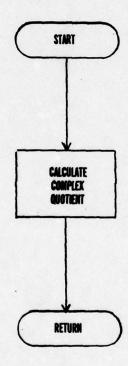
CUPSD CURT CURVD CURVT MILDS

MTNDS ZIADD

TESCRIPTION ·

Subroutine CXADE calculates the complex sum C=A+P. Each complex number is represented by two decimal numbers, the first for its real part and the second for its imaginary part. The complex sum is defined according to.

$$C(2) = A(2) + P(2)$$



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MAME . CXDIV

TYPE · SUPPOUTINE

CEMERAL PURPOSE:

Calculates the complex quotient C = P/B

VA LES:

> = Complex dividend

P = Complex divisor

C = Complex guotient

P = Intermediate storage

SUBFOUTINES CALLED.

MONE

CALLING PROCPATE

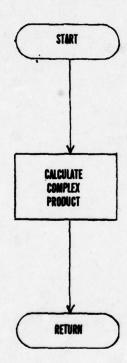
CUENC CENIE

PESCRIPTION .

Fubroutine CXDIV calculates the complex quotient C=A/B. Fach complex number is represented by two decimal numbers, the first for its real part and the second for its imaginary part. The complex cuotient is defined according to

C(1) = (A(1) *B(1) + A(2) *B(2)) / ((B(1) *B(1) + B(2) *B(2))

C(2) = (A(2) *B(1) - A(1) *B(2)) / ((B(1) *B(1) + B(2) *B(2))



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MAME: CXMPY

TYPF · SUPROUTINE

CENEFAL PURPOSF .

Calculates the complex product C = A*B

VARIABLES:

A = Complex multiplicand

E = Complex multiplier

C = Complex product

P = Intermediate storage

SUPPOUTINES CALLED-

MONE

CALLING PROGRAMS

CPOSS CUPNC CUPNDS CURT. MTLDS

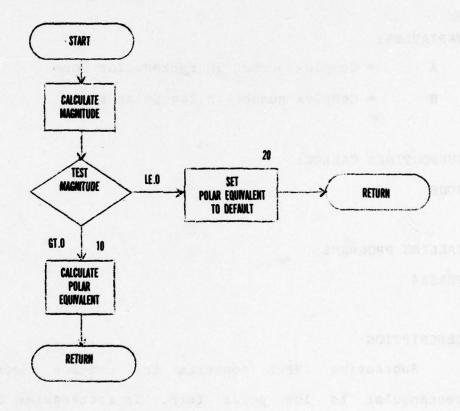
MTNDS PHASE4 TFTONV

PESCRIPTION .

Subroutine CXMPY calculates the complex product C=A*B. Each complex number is represented by two decimal numbers, the first for its real part and the second for its imaginary part. The complex product is defined according to:

C(1) = A(1) * B(1) - A(2) * P(2)

C(2) = A(2) * E(1) + A(1) * E(2)



NAME: CXPOL

TYPE: SUBROUTINE

GENERAL PURPOSE:

Converts the complex number A to log polar form B

VARIABLES:

A = Complex number in rectangular form

B = Complex number in log polar form

SUBROUTINES CALLED:

NONE

CALLING PROGRAMS:

PHASE 4

DESCRIPTION .

Subroutine CXPOL converts the complex number A from rectangular to log polar form. In rectangular form, the real part of the complex number is stored in A(1) and the imaginary part in A(2). The polar coordinates are returned to the calling program in B, where the radius is stored in B(1) and the angle in B(2). The conversion is carried out according to:

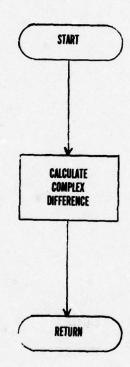
 $E(1) = 20 \log_{10} \sqrt{A(1)^2 + A(2)^2}$

P(2) = Arctan (A(2)/A(1))*57.2957795131

or if
$$\sqrt{A(1)^2 + A(2)^2} = c$$

F(1)=-1.0F20

P(2) = G.



NAME · CXSUB

TYPE: SUPROUTINE

GENERAL PURPOSE:

Calculates the complex difference C = A-D

VARIABLES .

A . = Complex minuend

D = Complex subtrahend

C = Complex difference

SUPROUTINES CALLED:

MONE

CALLING PROCRAMS:

CURCEL CURJF CURNC CURNES CURPT

CURSD CURT CURVE CURVT MTLDS

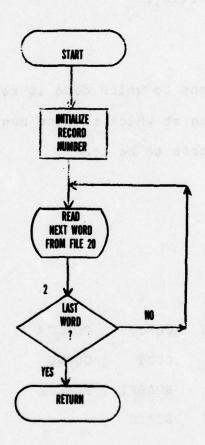
MTNDS ZIADD

PESCPIPTION .

Subroutine CXSUF calculates the complex difference C=A-B. Fach complex number is represented by two decimal numbers, the first for its real part and the second for its imaginary part. The complex difference is defined according to:

$$C(1) = A(1) - F(1)$$

$$C(2) = h(2) - P(2)$$



MAME: DATARD

TYPE: SUPRCUTINE

CENERAL BURPOSE -

Reads data from disk file 20

VARIABLES.

AREA = Core locations to which data is read

IPTP = Disk location at which reading begins

M = Number of words to be read

SUBROUTINES CALLED

NONE

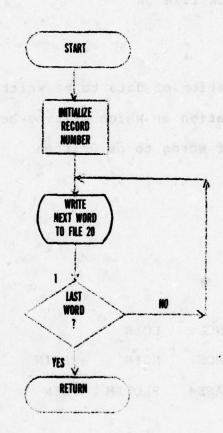
CALLING PROGRAMS:

CHNGE	CONTRL	CURCEN	CURJF	CURNC
CUFNDS	CUPPT	CURSD	CURT	CURVD
MODGEN	MTLDS	MTNDS	PHASE1	PHASE2
PHASE3	PEASE4	TRIANC	ZIADD	

PESCRIPTION .

Subroutine DATARD reads data from disk file 28. The transmission begins from record number IPTR and continues one word at a time until N words have been read from file 26 and stored in AREA(1)-AREA(N). The data is assumed to be decimal,

and is read without format. After the data transmission is complete, subroutine DATARD returns to the calling program.



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NAME: DATAWR

TYPE: SUBROUTINE

GENERAL PURPOSE:

Writes data to disk file 20

VARIABLES:

AREA = Core location of data to be written

IPTR = Disk location at which writing begins

N = Number of words to be written

SUBROUTINES CALLED:

NONE

CALLING PROGRAMS:

CHNGE DATOUT GENOUT LCIN

LDSIN MTLDS MTNDS NCIN NDSIN

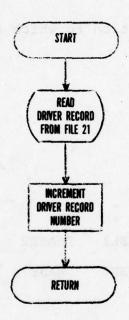
PHASE1 PHASE2 PHASE4 PLOTIN PSIN

TRIANG ZIADD

DESCRIPTION .

Subroutine DATAWR writes data to disk file 20. The decimal values stored at AREA(1)-AREA(N) are written to file 20 one word at a time beginning at record number IPTR. The write operations are performed without format. After the data transmission is

complete, subroutine DATAWR returns to the calling program.



NAME: DRIVED

TYPE: SUBROUTINE

GENERAL PURPOSE:

Reads a driver record from disk file 21

VARIABLES:

LCMMN = Starting core location to which record is
 read

SUBROUTINES CALLED:

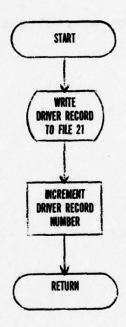
NONE

CALLING PROGRAMS:

CHNGE GENOUT MODGEN PHASE1 PHASE2
PHASE3 PHASE4 PHASE5 PHASE6 PHASE7
PRNDRV

DESCRIPTION:

Subroutine DRIVRD reads a driver record from disk file 21. The 10-word driver record stored at record number DRVREC is read from file 21 by an unformatted random access read statement and stored in the first ten words of global common. The driver record number is updated to DRVREC=DRVREC+1 and subroutine DRIVRD returns to the calling program.



NAME: DRIVWR

TYPE: SUBROUTINE

GENERAL PURPOSE .

Writes a driver record to disk file 21

VARIABLES:

LCMMN = Starting core location of record to be written

SUBROUTINES CALLED.

NONE

CALLING PROGRAMS:

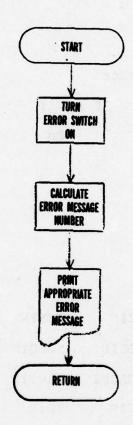
CHNGE DRVOUT GENOUT LCIN

PHASE1 PHASE2 PHASE3 PHASE5

PHASE6 PHASE7 PLOTIN

DESCRIPTION :

Subroutine DRIVWR writes a driver record to disk file 21. The driver record stored in the first ten words of global common is written to file 21 at record number DRVREC by an unformatted random access write statement. The driver record number is updated to DRVREC=DRVREC+1 and subroutine DRIVWR returns to the calling program.



MAME - EFFCR

TYPE - SUBROUTINE

CENEFAL PUPPOSE:

Prints error messages

VARIABLES .

J = Error message number

JADT = Input error switch · 0 = Off

1 = On

SUPROUTINES CALLED.

MONE

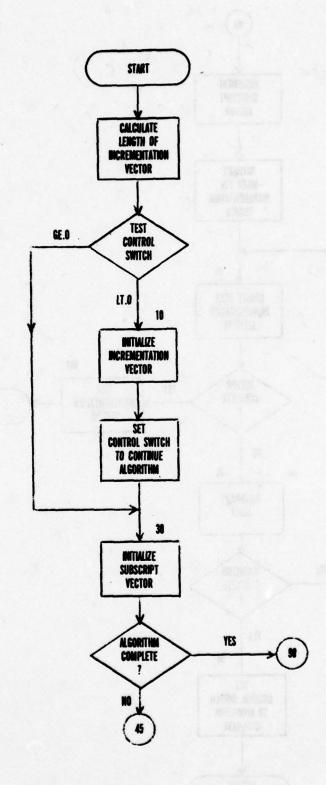
CALLING PROCRAMS -

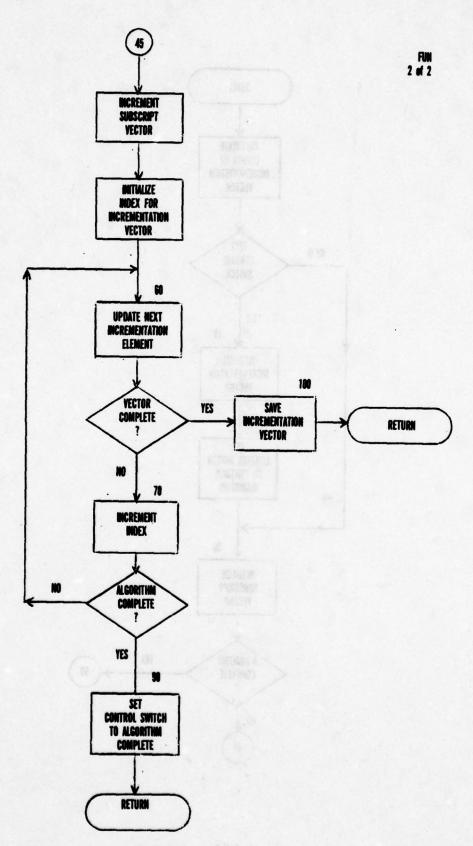
CENCE	PATOUT	DECIF	ECIN	GENIN
CENOUT	JFETIN	LCIN	LDSIN	MODGIN
rCIr	MDSIM	MODEIN	PUASE1	PLOTIN
PFCSS	PSIN	PTIN	SCIN	SDIN
TPNIN	VDIN	VTIN		

PESCRIPTION

Subroutine ERPOP prints error messages. First the error switch is turned on by JAPT=1, the appropriate error message is printed according to the error number I transmitted to the

subroutine as an argument, and subroutine ERROR returns to the calling program.





3-90B

MAHE - FUN

TYPE - SUPPOUTINE

CENERAL PURPOSE.

Cenerates subscripts of terms to be multiplied together to form one additive term of the polynomial power series expansion coefficient of order I, degree J.

VARIAPLES :

J = Degree of coefficient (number of subscripts returned)

= Subscript vector

I = Incrementation vector

"M = Length of incrementation vector

SWTCH = Control switch; -1 = Initialize algorithm

M = Continue algorithm

1 = Algorithm complete

SUPPOUTINES CALLED.

PONE

CALLING PROGRAMS -

PETA CPT CURPT FNCTH

DESCRIPTION:

Subroutine FUN generates (indirectly) the polynomial power series expansion coefficient $\mathbf{F}_{\mathbf{J}}^{\mathbf{I}}$ (A), the coefficient of $\mathbf{X}^{\mathbf{I}}$ when the polynomial:

$$\mathbf{P} = \sum_{n=1}^{\infty} \mathbf{A}_n \mathbf{x}^n$$

is raised to the Jth power. In NCAP it is convenient to denote a polynomial raised to the Jth power by:

$$P^{J} = \sum_{I=J}^{\infty} F_{J}^{I} (A) x^{I}$$

where F_J^I (A) is a function of the coefficients A.

Although F $_{\bf J}^{\bf I}$ (A) cannot be formally defined mathematically, it possesses certain properties which serve as the basis for its algorithmic evaluation.

- 1) $F_{J}^{I}(A)$ is a sum of products of A.
- 2) Each product contains exactly J coefficients A.
- The sum of the subscripts in each product is I, the exponent of X.
- 4) Each subscript within a product lies in the range 1 through I-J+1.

- 5) Each possible ordered set of subscripts of length J and sum I appears once and only once in the expansion.
- 6) The number of additive terms in the expansion is $\binom{I-1}{J-1}$.

Subroutine FUN does not generate the expansion coefficient directly, but rather through repeated calls produces a sequence of integer vectors K, each of which constitutes the subscripts of coefficients to be multiplied out to form a single additive term of the expansion. The calling program then forms the products and adds them to form the correct expansion.

The iterative calls to subroutine FUN are controlled by the integer variable SWTCH which is initialized to -1 by the calling program and subsequently tested and updated in FUN. Upon entry into FUN, if SWTCH is negative, the algorithm is initialized and the first subscript vector created. SWTCH is set to +1 when the last subscript vector has been formed and the algorithm is complete. For intermediate iterations SWTCH=C.

In subroutine FUN the subscripts are stored in the integer vector K. For F_J^I (A) each call to subroutine FUN creates the J subscripts which constitute one additive term of the expansion returning them to the calling program in K(1) through K(J). The subscripts are generated by applying an incrementation operation to the initial subscript vector $K=(1,1,\ldots 1)$ of length J. This operation is controlled by an incrementation vector L(M), $M=1,2,\ldots,(I-J)$, where K incremented by L results in:

$$Y(L(M))=K(L(M))+1$$
 for $M=1,2,...,(I-J)$

The incrementation vectors are formed as follows. For the first J iterations, L(1) ranges from 1 through J and the

remaining vector elements L(2)...L(I-J)=1. For the next (J-1) iterations, L(1) ranges from 2 through J, L(2)=2, and the remaining vector elements L(3)...L(I-J)=1. This process continues until all (I-J) vector elements equal J.

The algorithm of subroutine FUN generates in K all possible ordered sets of J integers in the range 1 through (I-J+1) such that K(1)+K(2)+...+K(J)=I and can be summarized as follows:

STEP 1. Initialization

- a) If (SWTCH.GE.0) go to Step 2.
- b) Otherwise initialize incrementation vector:
 I(N)=1 for N=1,2,...,(I-J)
- c) Set SWTCH=0

STEP 2: Initialize Subscript Vector

- a) Set K(N)=1 for N=1,2,...,J
 - h) If I=J, algorithm complete; exit with subscripts inF. Otherwise go to Step 3.

STEP 3. Increment Subscript Vector

a) Set K(L(M))=K(L(M))+1 for M=1.2,...,(I-J)

STEP 4. Form Incrementation For Next Iteration

- a) Initialize index N=1
- b) Upcate L(N)=L(N)+1. If (L(N).LE.J), incrementation vector complete go to Step 5.
- c) Otherwise increment index N=N+1. If (N.LE.(I-J)), go go to Step 4b. Otherwise algorithm complete: set SFTCH=1 and exit with subscripts in K.

STEP 5. Save Incrementation Vector for Next Iteration

- a) Define L(N) = L(N) for M=1, N
- b) Exit with subscripts in R.

A complete tabulation of F $_{\bf J}^{\bf I}$ (A) up to 6th order (the maximum order of analysis allowed in NCAP) is presented below:

Incrementation vector L contains $(I-J)=\emptyset$ For F $_2^2$ (A): I=2,J=2 elements, subscript vector K contains J=2 elements, expansion contains one additive term.

L=Inapplicable K=(1,1) \Rightarrow A(1)*A(1)

Incrementation vector L contains (I-J)=1For F $\frac{3}{2}$ (A): I=3,J=2 element, subscript vector K contains J=2 clements, expansion contains two additive terms.

$$I = (1)$$
 \Rightarrow $K = (2,1)$ \Rightarrow $\Lambda(2) * \Lambda(1)$
 $I = (2)$ \Rightarrow $\Lambda(1) * \Lambda(2)$

For F $\frac{3}{3}$ (A): I=3,J=3 elements, subscript vector K contains

J=3 elements, expansion contains one additive term.

I=Inapplicable F=(1,1,1) \blacktriangleright F(1)*A(1)*A(1)

Incrementation vector L contains (I-J)=2For $\Gamma = \frac{4}{2}$ (A) I=4,J=2 elements, subscript vector K contains J=2 elements, expansion contains three additive terms.

$$L=(1,1)$$
 \Rightarrow $K=(3,1)$ \Rightarrow $\Lambda(3)*\Lambda(1)$
 $L=(2,1)$ \Rightarrow $M=(2,2)$ \Rightarrow $\Lambda(2)*\Lambda(2)$
 $M=(2,2)$ \Rightarrow $M=(1,2)$ \Rightarrow $M=(1)*\Lambda(3)$

Incrementation vector L contains (I-J)=1

For F 3 (A) I=4.J=3 element, subscript vector X contains J=3

elements, expansion contains three additive terms

$$I = (1)$$
 $K = (2,1,1)$
 $A(2) *A(1) *A(1)$
 $E = (2)$
 $K = (1,2,1)$
 $A(1) *A(2) *A(1)$
 $E = (3)$
 $A(1) *A(1) *A(2)$
 $A(1) *A(1) *A(2)$

- Incrementation vector L contains $(I-J)=\emptyset$ For F $_4^4$ (A): I=4.J=4 elements, subscript vector K contains J=4 elements, expansion contains one additive term.
 - L=Inapplicable $K = (1,1,1,1) \implies A(1) *A(1) *A(1) *A(1)$
- Incrementation vector L contains (I-J)=3

 For F 2 (1): I=5.J=2 elements, subscript vector K contains J=2 elements, expansion contains 4 additive terms.

$$\Gamma = (1,1,1)$$
 \Rightarrow $\Gamma = (4,1)$ \Rightarrow $\Lambda(4) *\Lambda(1)$
 $\Gamma = (2,1,1)$ \Rightarrow $\Gamma = (3,2)$ \Rightarrow $\Lambda(3) *\Lambda(2)$
 $\Gamma = (2,2,1)$ \Rightarrow $\Gamma = (2,2,2)$ \Rightarrow $\Gamma = (1,4)$ \Rightarrow $\Gamma = (1,4)$ \Rightarrow $\Gamma = (1,4)$ \Rightarrow $\Gamma = (1,4)$

- Incrementation vector I contains (I-J)=2 For F $_3^5$ (A) I=5,J=3 elements, subscript vector X contains J=3 elements, expansion contains six additive terms.
 - L=(1,1) \Rightarrow Y=(3,1,1) \Rightarrow $\Lambda(3)*\Lambda(1)*\Lambda(1)$

L=(2,1)
$$\Rightarrow$$
 K=(2,2,1) \Rightarrow A(2)*A(2)*A(1)
L=(3,1) \Rightarrow K=(2,1,2) \Rightarrow A(2)*A(1)*A(2)
L=(2,2) \Rightarrow K=(1,3,1) \Rightarrow A(1)*A(3)*A(1)
L=(3,2) \Rightarrow K=(1,2,2) \Rightarrow A(1)*A(2)*A(2)
L=(3,3) \Rightarrow K=(1,1,3) \Rightarrow A(1)*A(1)*A(3)

Incrementation vector I contains (I-J)=1For F $_4^5$ (A): I=5,J=4 elements, subscript vector K contains J=4 elements, expansion contains four additive terms.

$$L = (1) \qquad \qquad K = (2,1,1,1) \qquad \qquad \Lambda(2) * \Lambda(1) * \Lambda(1) * \Lambda(1)$$

$$L = (2) \qquad \qquad M = (1,2,1,1) \qquad \qquad \Lambda(1) * \Lambda(2) * \Lambda(1) * \Lambda(1)$$

$$L = (3) \qquad \qquad K = (1,1,2,1) \qquad \qquad \Lambda(1) * \Lambda(1) * \Lambda(2) * \Lambda(1)$$

$$L = (4) \qquad \qquad K = (1,1,1,2) \qquad \qquad \Lambda(1) * \Lambda(1) * \Lambda(1) * \Lambda(2)$$

Incrementation vector L contains (I-J)=0For F $_5^5$ (A): I=5,J=5 elements, subscript vector K contains J=5 elements, expansion contains one additive term.

L=INAPPLICABLE $K=(1,1,1,1,1) \implies A(1)*A(1)*A(1)*A(1)*A(1)$

Incrementation vector L contains (I-J)=4For F₂ (A) · I=6,J=2 elements, subscript vector K contains J=2 elements, expansion contains five additive terms.

L=(1,1,1,1)
$$\blacktriangleright$$
 K=(5,1) \blacktriangleright A(5)*A(1)
L=(2,1,1,1) \blacktriangleright K=(4,2) \blacktriangleright A(4)*A(2)
L=(2,2,1,1) \blacktriangleright K=(3,3) \blacktriangleright A(3)*A(3)
L=(2,2,2,1) \blacktriangleright K=(2,4) \blacktriangleright A(2)*A(4)
L=(2,2,2,2) \blacktriangleright K=(1,5) \blacktriangleright A(1)*A(5)

Incrementation vector L contains (I-J)=3

For F 3 (A): I=6,J=3 elements, subscript vector K contains J=3

elements, expansion contains ten additive terms.

L=(1,1,1)	•	K = (4,1,1)	•	A(4)*A(1)*A(1)
L=(2,1,1)	•	K=(3,2,1)	•	A(3)*A(2)*A(1)
L=(3,1,1)	•	K=(3,1,2)	•	A(3)*A(1)*A(2)
L=(2,2,1)	•	K = (2,3,1)	•	A(2)*A(3)*A(1)
L=(3,2,1)	•	K=(2,2,2)	•	A(2)*A(2)*A(2)
L=(3,3,1)		K=(2,1,3)	•	A(2)*A(1)*A(3)
L=(2,2,2)	•	K = (1, 4, 1)	•	A(1)*A(4)*A(1)
L=(3,2,2)	-	K=(1,3,2)	•	A(1)*A(3)*A(2)
L=(3,3,2)	•	K = (1, 2, 3)	•	A(1)*A(2)*A(3)
L=(3,3,3)	•	K = (1, 1, 4)	•	A(1) *A(1) *A(4)

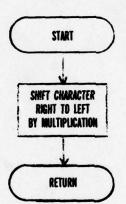
Incrementation vector L contains (I-J)=2 For F $_4^6$ (A): I=6,J=4 elements, subscript vector K contains J=4 elements, expansion contains ten additive terms.

Incrementation vector L contains (I-J)=1 For F $_5^6$ (A): I=6,J=5 element, subscript vector K contains J=5 elements, expansion contains five additive terms.

L=(1)
$$\Rightarrow$$
 K=(2,1,1,1,1,) \Rightarrow A(2)*A(1)*A(1)*A(1)*A(1)
L=(2) \Rightarrow K=(1,2,1,1,1,) \Rightarrow A(1)*A(2)*A(1)*A(1)*A(1)
L=(3) \Rightarrow K=(1,1,2,1,1,) \Rightarrow A(1)*A(1)*A(2)*A(1)*A(1)
L=(4) \Rightarrow K=(1,1,1,2,1,1) \Rightarrow A(1)*A(1)*A(1)*A(2)*A(1)
L=(5) \Rightarrow K=(1,1,1,1,2,1) \Rightarrow A(1)*A(1)*A(1)*A(2)*A(2)

Incrementation vector L contains (I-J)=0 For F $_6^6$ (A): I=6,J=6 elements, subscript vector K contains J=6 elements, expansion contains one additive term.

L=Inapplicable K=(1,1,1,1,1,1) A(1)*A(1)*A(1)*A(1)*A(1)*A(1)



NAME: LSHIFT

TYPE: SUBROUTINE

GENERAL PURPOSE:

Shifts the character stored in the least significant bits of ICHAR to the most significant bits

VARIABLES:

NONE

SUBROUTINES CALLED:

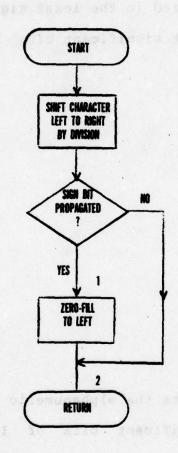
NONE

CALLING PROGRAMS .

PRSET

DESCRIPTION .

Subroutine LSHIFT shifts the alphanumeric character stored in the NECHAR least significant bits of ICHAR to the most significant bits by multiplication by an appropriate power of 2. The resulting left-justified character is automatically zero filled to the right.



NAME: RSHIFT

TYPE: SUBROUTINE

GENERAL PURPOSE:

Shifts the character stored in the least significant bits of ICHAR to the most significant bits

VARIABLES .

NONE

SUBROUTINES CALLED:

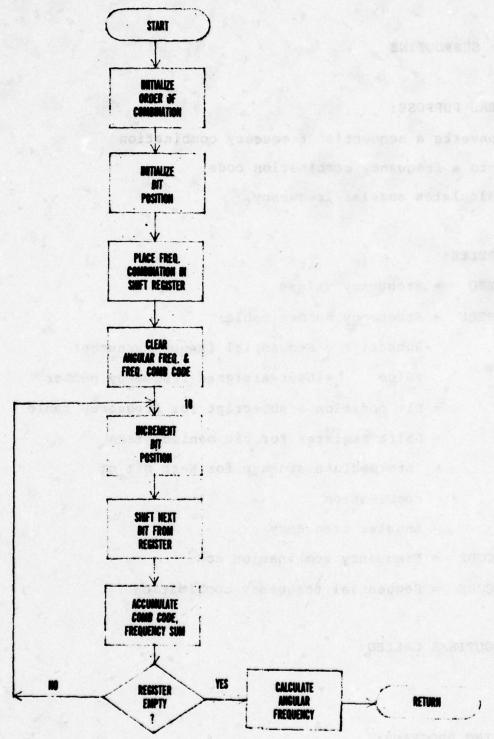
NONE

CALLING PROGRAMS .

MAIN PHASE4 SHIFT

DESCRIPTION:

Subroutine RSHIFT shifts the alphanumeric character stored in the NBCHAR most significant bits of ICHAR to the least significant bits by division by an appropriate power of 2. If the operation causes the sign bit to propagate (ICHAR.LT.0), then the word is zero filled to the left by adding 2 NBCHAR.



0

NAME: SSCOD2

TYPE: SUBROUTINE

GENERAL PURPOSE:

Converts a sequential frequency combination into a frequency combination code.

Calculates angular frequency.

VARIABLES:

FREQ = Frequency values

FRTBL = Frequency number table;

Subscript = Sequential frequency number

Value = User-assigned frequency number

K = Bit position & subscript for frequency table

L = Shift register for bit manipulation

M = Intermediate storage for K-th bit of combination

S = Angular frequency

SCODE = Frequency combination code

SCOMB = Sequential frequency combination

SUBROUTINES CALLED:

NONE

CALLING PROGRAMS:

CUPNC PFASE2

DESCRIPTION:

Subroutine SSCOD2 converts a sequential frequency combination to a frequency combination code and calculates the angular frequency for that combination. Arguments transmitted to the subroutine are: the sequential frequency combination, SCOME; the frequency number table, FRTPL and the frequency values, FREC. Arguments returned to the calling program are SCODE and S, the calculated frequency combination code and angular frequency.

The relationship between the sequential frequency combination and the combination code is that if the Kth sequential frequency is in the combination (i.e., the Kth bit of SCOPP is 1), then the user-assigned frequency number FRTPL(K) is included in the code (i.e., the FPTPL(K)th bit of SCODE is 1). Similarly the Kth frequency value is included in the angular frequency, S, if the Kth sequential frequency appears in the frequency combination.

The method employed by subroutine SSCCD2 is to extract each secuential frequency number from SCCME and, using the frequency number table FRTBL, to combine the corresponding user-assigned frequency numbers and values into the frequency combination code and angular frequency. The sequential frequency numbers are extracted from the combination by using a simulated shift register L, which is initially loaded with the sequential frequency combination SCCME. A frequency number is extracted from SCCME by shifting the right-most bit from the register into

the integer variable M. At the same time, the integer K maintains a count on the bit position being examined.

If the Kth bit is set (M=1), then the FRTBL(K)th bit of the frequency combination code is set to 1 and the Kth frequency value is added to the angular frequency. The calculations are complete when the shift register is empty.

The complete algorithm of subroutine SSCOD2 is summarized as follows:

STEP1: Initialization

- a) Initialize bit position counter K=0 and load shift register with sequential combination by L=SCOMB.
- b) Clear angular frequency and combination code by S=0. and SCODE=0

STEP 2: Extract Kth Bit and Map to SCODE and S

- a) Increment bit position by K=K+l
- b) Shift right-most bit from register and store in M:

M=L

L=L/2

M=M-(L*2)

c) Set PTRBL(K)th bit of frequency combination code according to value of M:

SCODE=SCODE+2**(FTREL(K)-1)*M

d) Include Kth frequency in angular frequency according to value of \mathbb{M}^{\perp}

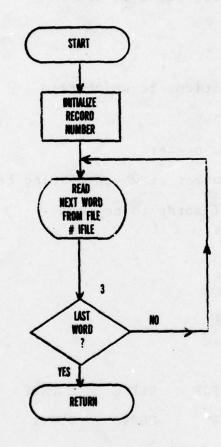
S=S*FREQ(K)*FLOAT(M)

STEP 3 Test Shift Register

- e) If L=0 (shift register empty), frequency combination code is complete, go to Step 4.
- b) Otherwise go to Step 2.

STEF 4. Complete Angular Frequency Calculation

a) "ultiply angular frequency by 2π and exit



NAME: TFRD

TYPE: SUDRCUTINE

CENERAL PURPOSE .

Reads data from disk files 24 and 25

VARIACLES:

AREA = Core locations to which data

is read

IFILE = Disk file number

IPTP = Record number at which reading begins

M = Number of words to be read

SUBPOUTINES CALLED :

MCNE

CALLING PROGRAMS.

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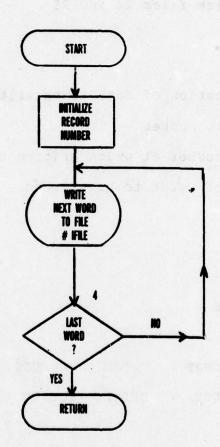
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TRIANC

DESCRIPTION .

Subroutine TFRD reads data from disk file IFILE. The transmission begins from record number IPTR and continues one word at a time until N words have been read from file IFILE and

stored in AREA(1)-AREA(N). The data is assumed to be decimal and is read without format. After the data transmission is complete subroutine TFRD returns to the calling program.



NAME: TEWR

TYPE: SUPROUTINE

GENEPAL PUPPOSE

Writes data to disk files 24 and 25

VARIAPLES:

AREA -= Core location of data to be written

IFILE = Disk file number

IPTE = Record number at which writing begins

F = Number of words to be written

SUBFOUTINES CALLED.

MONE

CALLING PROCRAMS ...

CONTRL CURCEN CUPJF CURNC CURNDS

CURPT CUPSD CURT CURVD CURVT

PHASE3 TRIANC

DESCRIPTION -

Subroutine TFWP writes data to disk file IFILE. The decimal values stored at AREA(1)-AREA(N) are written to file IFILE one word at a time beginning at record number IPTE. The write operations are performed without format. After the data

transmission is complete, subroutine TFWR returns to the calling program.

THOMIT - SHA

Calculates sispeed clock time for a single pregram segment and acquestioners that sign of the segment

TIMOUT

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CALLING PROGRAMS

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NAME: TIMOUT

TYPE: SUBROUTINE

GENERAL PURPOSE:

Calculates elapsed clock time for a single program segment and accumulates total clock time for all executions of that segment within an NCAP job

VARIABLES:

I = Index for TIME and TTOTAL arrays; program

segment identifier

TIME = Elapsed clock time per segment

TSTART = Start time per segment

TSTOP = Stop time per segment

TTOTAL = Accumulated clock time per segment

SUBROUTINES CALLED:

PTIME

CALLING PROGRAMS:

ENDIN FNCTN PHASE1 PHASE2

PHASE3 PHASE4 PHASE5 PHASE6 PHASE7

DESCRIPTION .

Subroutine TIMOUT calculates the clock time elapsed during

execution of a single program segment (a portion of a subprogram or one or more subprograms) and accumulates total clock time consumed for all executions of that segment during an NCAP job. The program segment is identified by the integer argument I, and the start time for the segment, TSTART, is transmitted to the subroutine in the labelled common area TIMEFL.

NCAP allows the timing of 30 program segments. This number may be increased or decreased by adjusting the dimension of the TSTART, TSTOP, TIME, and TTOTAL arrays in the labelled common area TIMEFL. The program segments presently selected for timing are:

Segment	De	esc	cription
1	Phase	Ø	Complete
2	Phase	1	Complete
3	Phase	1	Generator Routines
4	Phase	1	Nonlinear Components
	Routin	nes	£\$
5	Phase	1	Vacuum Diode Routines
6	Phase	1	Vacuum Triode Routines
7	Phase	1	Vacuum Pentode Routines
8	Phase	1	Transistor Routines
the clock Carre	Phase	1	Semiconductor Diode
	Routin	nes	rnemejaje šilajuteka Jul
10	Phase	1	JFET Routines
11 and a way	Phase	1	Linear Dependent Source
	Routin	es	ser no local famol on me
	3-96C		

12 Phase 1 Nonlinear Dependent S	Source
And the Lagrant Res at the cost Pout ines to 19019019	
13 Phase 2 Complete	
14 Phase 3 Complete	
15 Phase 3 Cenerator Routines	for Boats
Phase 3 Nonlinear Components	Routines
17 Phase 3 Vacuum Diode Routines	ADW .
18 Phase 3 Vacuum Triode Routine	es ed year
19 Phase 3 Vacuum Pentode Routin	nes TRATET
20 Phase 3 Transistor Routines	
Phase 3 Semiconductor Diode F	Routines
Phase 3 JFET Routines	
Phase 3 Nonlinear Dependent S	Source
Poutines	
24 Subroutine CONTRL	
25 Santauos voissanes Subroutine TRIANG	
26 Subroutine FNCTN	
27 Phase 4 Complete	
28 Alaman at the Market Phase 5 Complete	
29 Phase 6 Complete	
30 Phase 7 Complete	

Each program segment is timed by addressing the clock before its first executable statement to acquire its TSTART time and by calling subroutine TIMOUT after its last executable statement to accuire its TSTCP time. The method of addressing the clock is a machine dependent function performed in the present NCAP version

by subroutine PTIME. Subroutine TIMOUT then calculates the elapsed time for each program segment as:

TIME(I) = TIME(I) + (TSTOP(I) - TSTAPT(I))

and the accumulated time as:

TTOTAL(I) = TTOTAL(I) + (TSTOP(I) - TSTAFT(I))

The resulting times are saved in the labelled common area TIMEFL for subsequent output by subroutine PENTIM, and subroutine TIMCUT returns to the calling program.

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