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A COMPUTATIONAL ALGORITHM FOR THE
EIGENVECTORS OF A SINGULAR MATRIX

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I. INTRODUCTION

In this paper the eigenvalues and the eigenvectors of a singular matrix, A , are investigated.¹ Present techniques for determining the eigenvalues of a matrix require that the matrix be nonsingular or that some iterative procedure such as root searching techniques be applied to the characteristic equation. Here, a computational (non-iterative type) algorithm is given to construct a (smaller) nonsingular matrix, Q , that has exactly the same nonzero eigenvalues as A and only these eigenvalues. Thus, a standard technique such as the QR algorithm² can be applied to the matrix Q to evaluate the eigenvalues and eigenvectors of Q .

The algorithm given in this paper yields the eigenvectors of the nonzero eigenvalues of A , which are obtained from the eigenvectors of Q , and the rank of the matrix A . The eigenvectors associated with the zero eigenvalues of A are an immediate consequence of the algorithm.

II. MOTIVATION OF THE ALGORITHM

Let $A = (a_{ij})$ be an $n \times n$ matrix and U an $n \times n$ nonsingular matrix. Define $W = U^{-1}AU$. Let λ and X be an eigenvalue and an eigenvector, respectively, of A , then

$$\begin{aligned} AX &= \lambda X \\ U^{-1}A(UU^{-1})X &= \lambda U^{-1}X \\ W(U^{-1}X) &= \lambda(U^{-1}X), \end{aligned} \tag{1}$$

that is, λ and $U^{-1}X$ are an eigenvalue and eigenvector of W . Conversely, if μ and Z are an eigenvalue and an eigenvector of W , then μ and UZ are an eigenvalue and an eigenvector of A .

Suppose A is of rank r and there exists a U such that W can be partitioned into the following form:

¹The algorithm given in this paper was developed by the author in response to a problem posed by H. McCoy, TRASANA, to obtain the eigenvalues and eigenvectors of a special singular matrix.

²Wilkinson, J. H., The Algebraic Eigenvalue Problem, Clarendon Press, Oxford, 1965.

$$W = U^{-1}AU = \begin{pmatrix} Q_r & O_1 \\ B & O_2 \end{pmatrix}, \quad (2)$$

where Q_r is an $r \times r$ matrix, O_1 and O_2 are $r \times (n-r)$ and $(n-r) \times (n-r)$ zero matrices, respectively, and B is an $(n-r) \times r$ matrix. Let $\lambda \neq 0$ be an eigenvalue of W , and of A , and Z be the associated eigenvector. If the vector Z , as a matrix, is partitioned into Z_r , a $r \times 1$ matrix, and Y , a $(n-r) \times 1$ matrix, then we can write the following matrix equation

$$\begin{pmatrix} Q_r & O_1 \\ B & O_2 \end{pmatrix} \begin{pmatrix} Z_r \\ Y \end{pmatrix} = \lambda \begin{pmatrix} Z_r \\ Y \end{pmatrix} \quad (3)$$

as

$$Q_r Z_r = \lambda Z_r \quad (4)$$

and

$$BZ_r = \lambda Y, \quad (5)$$

that is, λ and Z_r are an eigenvalue and eigenvector of the matrix Q_r .

Thus, the nonzero eigenvalues of W , therefore for A , are the eigenvalues (nonzero) of Q_r . The associated eigenvectors of A are given by

$$X = UZ = U \begin{pmatrix} Z_r \\ Y \end{pmatrix} = U \begin{pmatrix} Z_r \\ \lambda^{-1} BZ_r \end{pmatrix}, \quad (6)$$

which are determined from the eigenvectors of Q_r .

Conversely, if $\lambda \neq 0$ is an eigenvalue of Q_r and Z_r the eigenvector, then it follows that X , which is defined by (6), is an eigenvector of A and λ is the eigenvalue.

The eigenvectors of A associated with the zero eigenvalues are immediate. Define the column vectors

$$E_i = \begin{pmatrix} \delta_{1,r+i} \\ \delta_{2,r+i} \\ \vdots \\ \delta_{n,r+i} \end{pmatrix}, \quad i = 1, 2, \dots, n-r, \quad (7)$$

where δ_{k1} is the Kronecker delta function. Then $WE_i = 0$, that is, E_i , $i = 1, 2, \dots, n-r$, are the associated eigenvectors for the zero eigenvalues of W . Thus, UE_i , $i = 1, 2, \dots, n-r$, are the eigenvectors of the zero eigenvalues of A ; but these eigenvectors are just the column vectors represented by the last $n-r$ columns of the matrix U .

Since the rank of A is r , it follows that $\{UE_i\}$ is the total set of (distinct) eigenvectors associated with the zero eigenvalue. If Q_r is singular, this implies that the characteristic equation has a zero of multiplicity greater than $n-r$; but no greater collection of eigenvectors. The process would then be applied to Q_r to determine a matrix of smaller order that has the same nonzero eigenvalues as A .

III. EXISTENCE OF THE MATRIX U

We will now show the existence of the matrix U . Let A_i , $i = 1, 2, \dots, n$, denote the column vectors represented by the columns of the matrix A . The definition and notation for the inner product of any two vectors is given by

$$(A_i, A_j) = \sum_{k=1}^n a_{ki} \bar{a}_{kj}, \quad (8)$$

where \bar{a} denotes the complex conjugate of a .

We will now apply the Gram-Schmidt orthogonalization process (see [3]) to the set of vectors $\{A_i\}$. If A_1 is the zero vector, then interchange A_1 and A_n . Let U_1 be the elementary matrix that interchanges column one and column n , when post matrix multiplication is applied (this is the identity matrix with the first and n th columns interchanged). If, after this interchange, A_1 is zero, then interchange A_1 and A_{n-1} , where U_2 is the appropriate elementary matrix. Continue until A_1 is not the zero vector. Define $V_1 = A_1$ and

³ Berberian, S. K., Introduction to Hilbert Space, Oxford University Press, New York, 1961.

$$V_2 = A_2 - \alpha_{21}V_1, \quad (9)$$

where $\alpha_{21} = \frac{(A_2, V_1)}{(V_1, V_1)}.$

If V_2 is the zero vector, interchange A_2 and A_k , where A_{k+1} was the last vector interchanged; and U_j the appropriate elementary matrix. Define U_{j+1} to be the elementary matrix that adds $-\alpha_{21}$ times the second column to column k . This is given by appending to the identity matrix $-\alpha_{21}$ in row 2 and column k . Construct a new V_2 . If V_2 is not the zero vector, then continue in the construction until a zero vector is generated, that is, define

$$V_i = A_i - \sum_{j=1}^{i-1} \alpha_{ij} V_j, \quad (10)$$

where $\alpha_{ij} = \frac{(A_i, V_j)}{(V_j, V_j)}, j = 1, \dots, i-1$. If V_i is the zero vector, interchange A_i with the last vector not interchanged, say A_m , and assign an appropriate elementary matrix U_k . Since $V_j, j = 1, \dots, i-1$, can be written in terms of A_1, \dots, A_j , then A_i can be written in the following form:

$$A_i = \beta_{i1}A_1 + \dots + \beta_{i,i-1}A_{i-1}. \quad (11)$$

Furthermore, let $U_{k+j}, j = 1, \dots, i-1$, be the elementary matrices which adds $-\beta_{ij}$ times column j to column m . Continue this construction for V_i until A_{i+1} was the last vector interchanged.

Define U to be the product of the elementary matrices generated above, that is,

$$U = U_1 U_2 \dots U_t. \quad (12)$$

Then AU has the form:

$$AU = \begin{pmatrix} Q & 0_1 \\ R & 0_2 \end{pmatrix}. \quad (13)$$

For example, if V_i is the zero vector, then

$$A_i = \sum_{j=1}^{i-1} \beta_{ij} A_j . \quad (14)$$

The product of the associated elementary matrices, say U_k, \dots, U_{k+i-1} , would interchange column i with column m and subtract β_{ij} times column j , $j = 1, \dots, i-1$, from column m . This would result in a zero column in column m .

If U_p interchanges column i with column j , then premultiplication by U_p^{-1} would interchange row i with row j ; and if U_q would add $-\beta$ times column i to column j , U_q^{-1} would add β times row j to row i . Thus,

$$U^{-1} = U_t^{-1} \dots U_2^{-1} U_1^{-1} \quad (15)$$

would result in row operations. Therefore, $U^{-1}AU$ would have the same form as (13), that is,

$$U^{-1}AU = \begin{pmatrix} Q_r & 0_1 \\ B & 0_2 \end{pmatrix} . \quad (16)$$

The resulting collection of nonzero vectors, A_1, \dots, A_r , under the Gram-Schmidt process, is the largest collection of independent vectors represented by the columns of A ; therefore, r is the rank of A .

IV. CONSTRUCTION OF THE MATRIX U

In this section we will develop a computationally feasible algorithm for the construction of the matrices U , Q_r and B . It should be noted that the actual construction of U^{-1} will not be required to obtain the final result.

To start the construction, set $U = (u_{ij}) = (\delta_{ij})$, the identity matrix. If there exists a vector V_i equal to the zero vector, then from (14)

$$A_i = \sum_{j=1}^{i-1} \beta_{ij} A_j . \quad (18)$$

Assume A_{m+1} was the last column interchanged. If $u_{ii} = 1$, set $u_{ii} = 0$, $u_{im} = 1$, $u_{mi} = 1$ and $u_{mm} = 0$. If $u_{ii} = 0$, set $u_{m+1,i} = 0$, $u_{m+1,m} = 1$, $u_{mi} = 1$ and $u_{mm} = 0$. This interchanges column i and column m .

Simultaneously, replace column i with column m in matrix A . Note that column m need not be replaced by column i , for this column is assumed to be the zero vector. In order to accomplish this, add $-\beta_{ij}$,

$j = 1, \dots, i-1$, to column m in the U matrix, where the i th row is determined from the one and only one nonzero element in the j th column of the U matrix. This nonzero element is unity. This construction is continued for i until $i+1$ was the last column interchanged. Thus, A has been reduced to the form of (13).

In order to obtain Q_r and B , similar operations must be done on rows of the reduced matrix A . An accounting must be kept on the column operations, for the row operations must be done in the same order, that is, if A has been postmultiplied by $U_1 U_2 \dots U_t$, then A must be pre-multiplied by $U_t^{-1} \dots U_2^{-1} U_1^{-1}$. Note that the row operations only involve the first r columns, for the remaining $n-r$ columns are assumed to be zero vectors. Therefore, if a column operation involved interchanging column ℓ with column m , then the row operation would interchange row ℓ with row m . Similarly, if a column operation adds $-\beta$ times column ℓ to column m , then the row operation would add β times row m to row ℓ .

In order to construct β_{ij} , the vectors, V_i , must be generated, from which β_{ij} can be obtained recursively. From (10)

$$V_i = A_i - \sum_{k=1}^{i-1} \alpha_{ik} V_k , \quad (19)$$

where $\alpha_{ik} = \frac{(A_i, V_k)}{(V_k, V_k)}$. Suppose

$$V_k = A_k - \sum_{j=1}^{k-1} \beta_{kj} A_j , \quad k = 2, \dots, i-1, \quad (20)$$

where $V_1 = A_1$. Then from (19) and (20)

$$\begin{aligned}
V_i &= A_i - \alpha_{i1}A_1 - \sum_{k=2}^{i-1} \alpha_{ik}(A_k - \sum_{j=1}^{k-1} \beta_{kj}A_j) \\
&= A_i - \alpha_{i1}A_1 - \sum_{k=2}^{i-1} \alpha_{ik}A_k + \sum_{k=2}^{i-1} \sum_{j=1}^{k-1} \alpha_{ik}\beta_{kj}A_j \\
&= A_i - \sum_{j=1}^{i-1} \alpha_{ij}A_j + \sum_{j=1}^{i-2} \left(\sum_{k=j+1}^{i-1} \alpha_{ik}\beta_{kj} \right) A_j \\
&= A_i - \alpha_{i,i-1}A_{i-1} - \sum_{j=1}^{i-2} \left(\alpha_{ij} - \sum_{k=j+1}^{i-1} \alpha_{ik}\beta_{kj} \right) A_j. \quad (21)
\end{aligned}$$

Therefore,

$$\beta_{ij} = \begin{cases} \alpha_{ij} - \sum_{k=j+1}^{i-2} \alpha_{ik}\beta_{kj} & , j = 1, \dots, i-2 \\ \alpha_{ij} & , j = i-1 \end{cases} \quad (22)$$

V. SOME ILLUSTRATIVE EXAMPLES

A. Example 1.

Let

$$A = \begin{bmatrix} 1 & 2 & 0 & 0 \\ 1 & 1 & 1 & 0 \\ 2 & 3 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix}.$$

Then

$$V_1 = \begin{bmatrix} 1 \\ 1 \\ 2 \\ 0 \end{bmatrix},$$

which is not the zero vector.

$$V_2 = A_2 - \alpha_{21}V_1 = \begin{bmatrix} 2 \\ 1 \\ 3 \\ 0 \end{bmatrix} - \frac{3}{2} \begin{bmatrix} 1 \\ 1 \\ 2 \\ 0 \end{bmatrix} = \begin{bmatrix} .5 \\ -.5 \\ 0 \\ 0 \end{bmatrix},$$

where

$$\alpha_{21} = \frac{(A_2, V_1)}{(V_1, V_1)} = \frac{2 + 1 + 6 + 0}{1 + 1 + 4 + 0} = \frac{3}{2}$$

and $\beta_{21} = \alpha_{21}$

$$V_3 = A_3 - \alpha_{31}V_1 - \alpha_{32}V_2 = \begin{bmatrix} 0 \\ 1 \\ 1 \\ 0 \end{bmatrix} - \frac{1}{2} \begin{bmatrix} 1 \\ 1 \\ 2 \\ 0 \end{bmatrix} - (-1) \begin{bmatrix} .5 \\ -.5 \\ 0 \\ 0 \end{bmatrix} = \begin{bmatrix} 0 \\ 0 \\ 0 \\ 0 \end{bmatrix},$$

where

$$\alpha_{31} = \frac{(A_3, V_1)}{(V_1, V_1)} = \frac{0 + 1 + 2 + 0}{6} = \frac{1}{2}$$

$$\alpha_{32} = \frac{(A_3, V_2)}{(V_2, V_2)} = \frac{0 + 0 + (-.5) + 0}{.25 + .25 + 0 + 0} = -1$$

and

$$\beta_{31} = \alpha_{31} - \alpha_{32}\beta_{21} = 2$$

$$\beta_{32} = \alpha_{32} = -1.$$

Now V_3 is the zero vector. Therefore, replace column 3 with column 4 in A . We will assume column 4 is the zero column. For in theory, $-\beta_{31}$ times column 1 and $-\beta_{32}$ times column 2 are added to column 4. Hence, the reduced A matrix is

$$AU = \begin{bmatrix} 1 & 2 & 0 & 0 \\ 1 & 1 & 0 & 0 \\ 2 & 3 & 0 & 0 \\ 0 & 0 & 1 & 0 \end{bmatrix}$$

If the same column operations are applied to U, which was initially set to the identity matrix, then

$$U = \begin{bmatrix} 1 & 0 & 0 & -2 \\ 0 & 1 & 0 & 1 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{bmatrix}.$$

Since V_3 was the zero vector, A_3 is renamed (interchange columns). A new V_3 is generated, namely,

$$V_3 = \begin{bmatrix} 0 \\ 0 \\ 0 \\ 1 \end{bmatrix} - 0 \cdot V_1 - 0 \cdot V_2 = \begin{bmatrix} 0 \\ 0 \\ 0 \\ 1 \end{bmatrix}$$

This "new" V_3 is not the zero vector. Hence, column operations terminate. The row operations are the following:

1. Interchange row 3 and row 4.
2. β_{31} times row 4 added to row 1.
3. β_{32} times row 4 added to row 2.

Thus,

$$U^{-1}AU = \begin{bmatrix} 5 & 8 & 0 & 0 \\ -1 & -2 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 2 & 3 & 0 & 0 \end{bmatrix}$$

The rank of A is 3. The eigenvector associated with the zero eigenvalue of A is the fourth column of U. Since

$$Q_r = \begin{bmatrix} 5 & 8 & 0 \\ -1 & -2 & 0 \\ 0 & 0 & 1 \end{bmatrix}$$

is nonsingular, no further reduction need be applied. The eigenvalues of Q_r are $\lambda_1 = 1$, $\lambda_2 = \frac{1}{2}(3 + \sqrt{17})$ and $\lambda_3 = \frac{1}{2}(3 - \sqrt{17})$. The respective eigenvectors are the following:

$$Z_1 = \begin{bmatrix} 0 \\ 0 \\ 1 \end{bmatrix}, \quad Z_2 = \frac{1}{2} \begin{bmatrix} -7-\sqrt{17} \\ 2 \\ 0 \end{bmatrix}, \quad Z_3 = \frac{1}{2} \begin{bmatrix} -7+\sqrt{17} \\ 2 \\ 0 \end{bmatrix}.$$

To obtain the eigenvectors for the nonzero eigenvalues of A, (6) is applied, where

$$B = \begin{bmatrix} 2 & 3 & 0 \end{bmatrix}.$$

Thus,

$$X_1 = \begin{bmatrix} 0 \\ 0 \\ 0 \\ 1 \end{bmatrix}, \quad X_2 = \frac{1}{4} \begin{bmatrix} -4 \\ -1-\sqrt{17} \\ -5-\sqrt{17} \\ 0 \end{bmatrix}, \quad X_3 = \frac{1}{4} \begin{bmatrix} -4 \\ -1+\sqrt{17} \\ -5+\sqrt{17} \\ 0 \end{bmatrix}$$

and the eigenvector for the zero eigenvalue is

$$X_4 = \begin{bmatrix} -2 \\ 1 \\ 1 \\ 0 \end{bmatrix},$$

the last column of U.

B. Example 2.

Let

$$A = \begin{bmatrix} 3 & -2 & -1 & 3 \\ 1 & 0 & 0 & 1 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \end{bmatrix}$$

It is easy to see that A_1 , A_2 and A_3 are independent vectors (also V_1 , V_2 and V_3) and $A_4 = A_1$. Thus, the only column operation is to add -1 times column 1 to column 4. Therefore,

$$Q_T = \begin{bmatrix} 3 & -2 & 0 \\ 1 & 0 & 0 \\ 0 & 1 & 0 \end{bmatrix}$$

and

$$B = [0 \ 0 \ 1],$$

where

$$U = \begin{bmatrix} 1 & 0 & 0 & -1 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix}.$$

If we apply the algorithm again to Q'_r , without intermediate calculations, we have

$$Q'_r = \begin{bmatrix} 3 & -2 \\ 1 & 0 \end{bmatrix},$$

$$B' = [0 \ 1]$$

and U' is the 3×3 identity matrix, where the prime denotes the partitioning of Q_r . The two eigenvalues of Q'_r are $\lambda_1 = 2$ and $\lambda_2 = 1$, and the associated eigenvectors are

$$z'_1 = \begin{bmatrix} 2 \\ 1 \end{bmatrix}$$

and

$$z'_2 = \begin{bmatrix} 1 \\ 1 \end{bmatrix}.$$

The respective eigenvectors of Q_r are

$$z_1 = U' \begin{bmatrix} z'_1 \\ \lambda_1^{-1} B' z'_1 \end{bmatrix} = \frac{1}{2} \begin{bmatrix} 4 \\ 2 \\ 1 \end{bmatrix}$$

and

$$z_2 = U' \begin{bmatrix} z'_2 \\ \lambda_2^{-1} B' z'_2 \end{bmatrix} = \begin{bmatrix} 1 \\ 1 \\ 1 \end{bmatrix}.$$

Finally, the eigenvectors of the nonzero eigenvalues of A are

$$x_1 = U \begin{bmatrix} z_1 \\ \lambda_1^{-1} B z_1 \end{bmatrix} = \frac{1}{4} \begin{bmatrix} 7 \\ 4 \\ 2 \\ 1 \end{bmatrix}$$

and

$$x_2 = U \begin{bmatrix} z_2 \\ \lambda_2^{-1} B z_2 \end{bmatrix} = \begin{bmatrix} 0 \\ 1 \\ 1 \\ 1 \end{bmatrix}.$$

The eigenvector associated with the zero eigenvalue is

$$x_3 = \begin{bmatrix} -1 \\ 0 \\ 0 \\ 1 \end{bmatrix},$$

that is, the fourth column of U . Since the rank of A is three, this is the only eigenvector associated with the zero eigenvalue.

VI. CONCLUSION

In order to apply this algorithm to a computer code, an a priori decision must be made for a "zero vector", due to machine round-off. In the author's program, which generates the matrices U , A_r and B , a vector, V , is the zero vector if

$$(V, V) < \epsilon,$$

where ϵ is some preassigned value.

An alternate application of the procedure could eliminate some computer code "bookkeeping". Since matrix multiplication is associative, a column operation, U_p , could be immediately followed by the row operation U_p^{-1} . This would, of course, increase the arithmetical operations, since the knowledge of the number of zero columns could not be totally incorporated.

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