Software Rejuvenation For Secure Tracking Control Of **Cyber-Physical Systems**

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Motivation

Software rejuvenation (SR) protects cyberphysical systems (CSPs) against cyber attacks on the run time code by **periodically** refreshing the system with an uncorrupted software image.

During SR, the system is in **open loop**, then this mechanism of protection may imply severe issues from the control perspective, such as stability and the inability to complete a mission (tracking performances).

Goal

To propose a **secure tracking control scheme** based on **software rejuvenation** for nonlinear and linear systems and provide the general conditions that guarantee the property of **safety** and liveness.

Introduction

Software Engineering

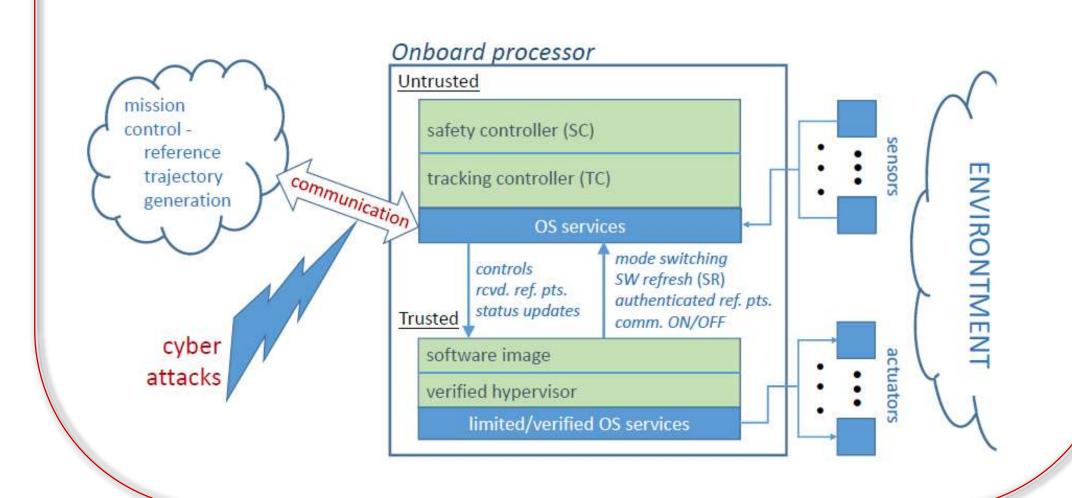
'...Software Rejuvenation is a periodic preemptive rollback of continuously running applications to prevent failures in the future." [Huang et al, 1995]

- I reboot;
- I restart the application from a clean internal state.

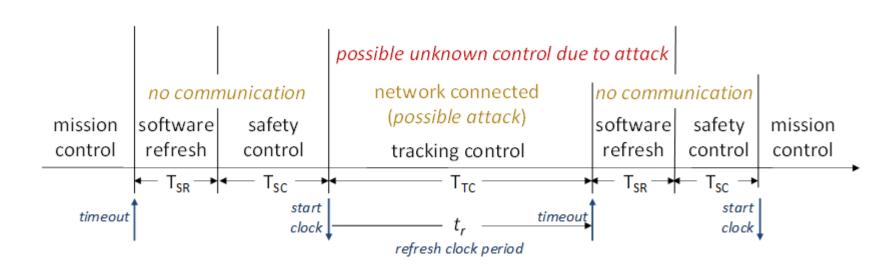
Control System

- I Fault Tolerant Control;
- I Secure Control of CPS [Abdi et al 2018, Arroyo et al 2017]

Attack Model and Architecture

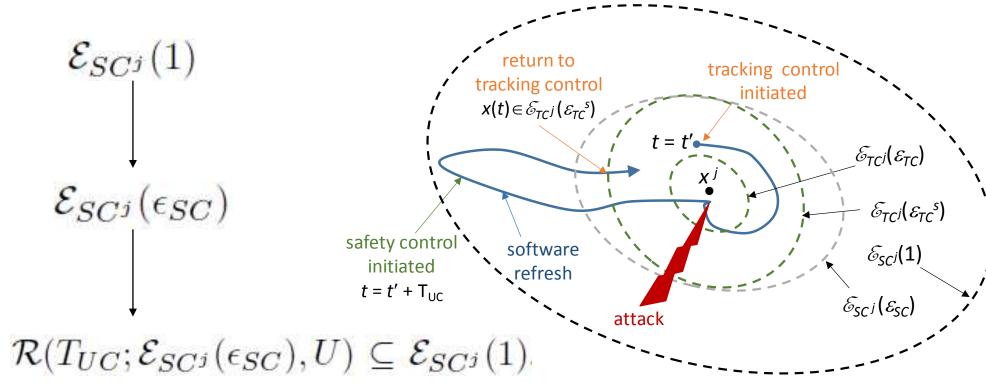


How SR works



The refresh clock period has to guarantee that for any control input, the system cannot leave the safety set.

Safety Control



Lyapunov Functions and Invariant Sets

 $\dot{x} = f_{\varphi}(x) \triangleq f(x, \varphi(x)),$ Controlled system:

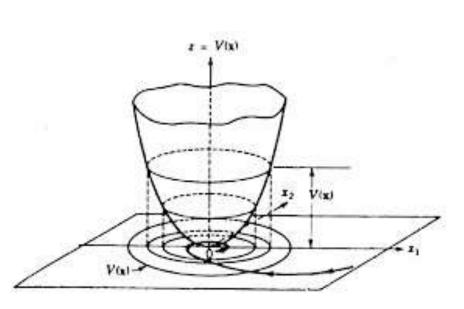
 $\begin{aligned}
V_{\varphi}: \mathbb{R}^n \to \mathbb{R} \\
V_{\varphi}(x_{eq}) &= 0
\end{aligned} \dot{V}_{\varphi}(x) = \frac{\partial V}{\partial x} \cdot f_{\varphi}(x) < 0.$ Lyapunov Function:

 $\mathcal{E}_{\varphi}(\epsilon) = \{x \in \mathcal{N}_{V_{\varphi}}(x_{eq}) | V_{\varphi}(x) \le \epsilon \}$ Lyapunov level Set:

Positively Invariant Set:

 $\forall t > 0, \ \mathcal{R}(t; \mathcal{E}_{\varphi}(\epsilon), \varphi) \subseteq \mathcal{E}_{\varphi}(\epsilon).$

 $0 < \epsilon' < \epsilon, \, \mathcal{E}_{\varphi}(\epsilon') \subset \mathcal{E}_{\varphi}(\epsilon)$



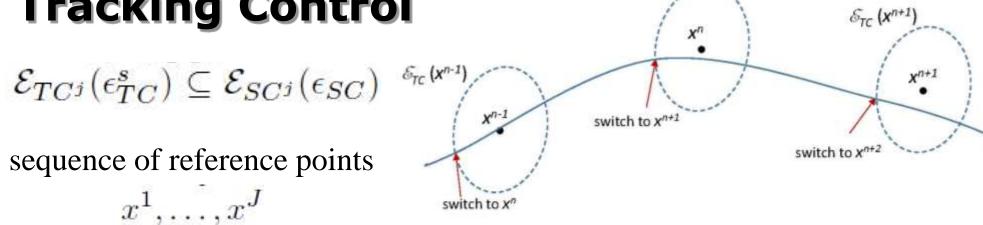
Two Fundamental Propositions

Proposition 2.1: Given system (1) with stabilizing controller φ for equilibrium state $(x_{eq}, \varphi(x_{eq}))$ and Lyapunov function $V_{\varphi}(x)$ as defined above, given $\epsilon > 0$ for any $\epsilon < \epsilon' \le 1 \ \exists \ \gamma > 0 \ \ni \forall \ t \ge (\epsilon' - \epsilon) \gamma^{-1},$

$$\mathcal{R}(t; \mathcal{E}_{\varphi}(\epsilon'), \varphi) \subseteq \mathcal{E}_{\varphi}(\epsilon). \tag{10}$$

Proposition 2.2: For any $U \subseteq \mathcal{U}$ and any $0 < \epsilon < \epsilon' \le 1$, $\exists T_U > 0 \ni \mathcal{R}(t; \mathcal{E}_{\varphi}(\epsilon), U) \subseteq \mathcal{E}_{\varphi}(\epsilon') \ \forall \ t < T_U.$

Tracking Control



 $\mathcal{E}_{TC^{j}}(\epsilon_{TC}) \subset \mathcal{E}_{TC^{j+1}}(\epsilon'), \quad \exists t^{j+1} > 0 \ni \forall \ t \geq t_{j+1},$ $\mathcal{R}(t; \mathcal{E}_{TC^j}(\epsilon_{TC}), TC^{j+1}) \subset \mathcal{E}_{TC^{j+1}}(\epsilon_{TC}).$

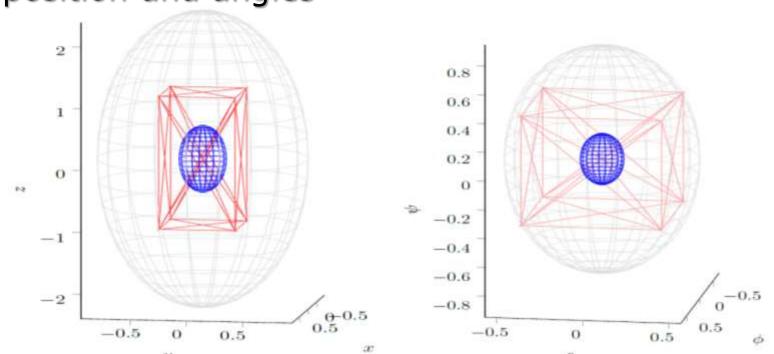
Safety and Liveness

- safety: when $x^j \to x^{j+1}$, the system has to be in $\mathcal{E}_{TC^{j+1}}(\epsilon_{TC}^s);$
- liveness: in presence of software refresh, the tracking controller has to drive the system to $\mathcal{E}_{TC^{j}}(\epsilon_{TC})$.

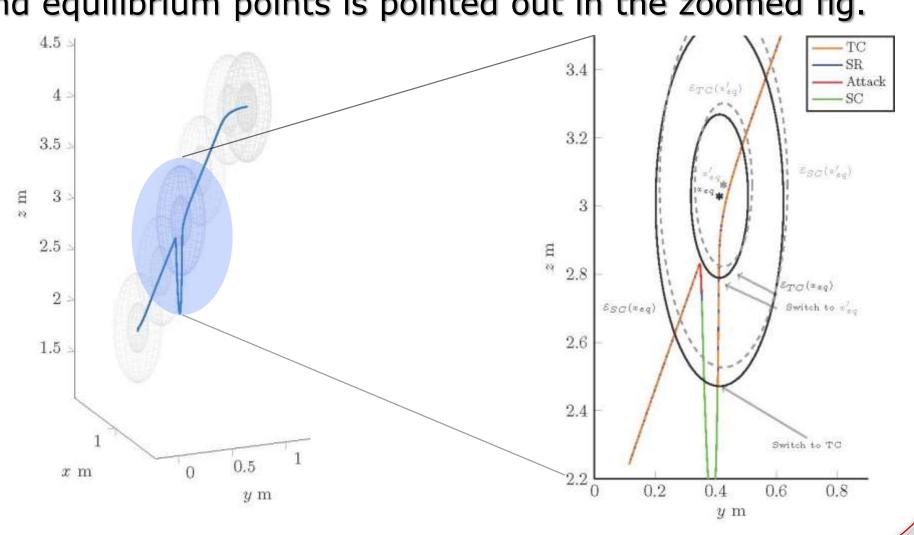
Example: quadrotor

- 6 DOF quadrotor using the PX4 jMAVSim quadrotor simulator
- Linearized model





Simulation results in presence of an attack. Details about the switching between the serveral controllers and equilibrium points is pointed out in the zoomed fig.



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