

Further results on constructions of generalized bent Boolean functions

Fengrong ZHANG¹, Shixiong XIA^{1*}, Pantelimon STĂNICĂ² & Yu ZHOU³

¹*School of Computer Science and Technology, China University of Mining and Technology, Xuzhou 221116, China;*

²*Naval Postgraduate School, Applied Mathematics Department, Monterey, CA 93943, USA;*

³*Science and Technology on Communication Security Laboratory, Chengdu 610041, PR China*

Received ; accepted

Citation Zhang F R, Xia S X, Stănică P, et al. Further results on constructions of generalized bent Boolean functions. *Sci China Inf Sci*, 2016, 59(3): *****, doi: *****

Dear editor,

Boolean bent functions were introduced by Rothaus in 1976 as an interesting combinatorial object with the important property of having optimal nonlinearity [1]. Since bent functions have many applications in sequence design, cryptography and algebraic coding, they have been extensively studied during the last thirty years [2, 3]. Over the past decades, based on bent functions, several constructions of highly nonlinear balanced functions were presented [4, 5].

In recent years several researchers have proposed generalizations of Boolean functions [6–9] and studied the effect of the Walsh-Hadamard transform on these classes. In [6], Schmidt presented the connection between words in multi-code code-division multiple access (MC-CDMA) systems and generalized bent functions from \mathbb{Z}_2^m to \mathbb{Z}_4 , and considered functions from \mathbb{Z}_2^n to \mathbb{Z}_q from the viewpoint of cyclic codes over rings. Later, Solé and Tokareva [7] called these functions from \mathbb{Z}_2^n to \mathbb{Z}_q generalized Boolean functions and presented the direct links between Boolean bent functions and generalized bent functions. More recently, Stănică et al. [9] investigated the properties of generalized bent functions and presented several constructions of such generalized bent functions

for both n even and n odd. They characterized a class of generalized bent functions symmetric with respect to two variables and generalized bent functions defined on \mathbb{Z}_2^n in \mathbb{Z}_8 . However, is there a technique that provides generalized bent functions symmetric with respect to m variables, where m is even? Additionally, in [9, Example 20, 21] the authors provided an explicit construction only for the even case. These give us a motivation to identify those generalized bent functions.

Let us denote the set of integers, real numbers and complex numbers by \mathbb{Z}, \mathbb{R} and \mathbb{C} , respectively and let the ring of integers modulo r be denoted by \mathbb{Z}_r . We denote the addition over \mathbb{Z}, \mathbb{R} and \mathbb{C} by ‘+’. Moreover, addition modulo q ($\neq 2$) is also denoted by ‘+’ and it is understood from the context. Let \mathbb{Z}_2^n be the n -dimensional vector space over \mathbb{Z}_2 . We denote the addition over \mathbb{Z}_2^n and \mathbb{Z}_2 by ‘ \oplus ’. Letting $\omega = (\omega_1, \dots, \omega_n)$ and $\mathbf{x} = (x_1, \dots, x_n) \in \mathbb{Z}_2^n$, we define the inner (or scalar) product by $\omega \cdot \mathbf{x} = \omega_1 x_1 \oplus \dots \oplus \omega_n x_n$. If $z = a + bi \in \mathbb{C}, a, b \in \mathbb{R}$, then $|z| = \sqrt{a^2 + b^2}$ denotes the absolute value of z , where $i^2 = -1$. We denote the vectors $(0, 0, \dots, 0) \in \mathbb{Z}_2^n$ by $\mathbf{0}_n$.

A function from \mathbb{Z}_2^n to \mathbb{Z}_q ($q \geq 2$ a positive integer) is called a *generalized Boolean function* in n variables [7]. Let \mathcal{GB}_n^q be the set of all n -variable

* Corresponding author (email: xiasx@cumt.edu.cn)

The authors declare that they have no conflict of interest.

generalized Boolean functions from \mathbb{Z}_2^n to \mathbb{Z}_q . If $q = 2$, we obtain the classical Boolean functions in n variables, whose set will be denoted by \mathcal{B}_n . The Hamming weight $\text{wt}(\mathbf{u})$ of a vector $\mathbf{u} \in \mathbb{Z}_2^n$ is the weight (number of 1's) of the binary string.

The (generalized) Walsh-Hadamard transform of $f \in \mathcal{GB}_n^q$ is the complex valued function over \mathbb{Z}_2^n which is defined by $\mathcal{H}_f(\boldsymbol{\omega}) = \sum_{\mathbf{x} \in \mathbb{Z}_2^n} \zeta^{f(\mathbf{x})} (-1)^{\boldsymbol{\omega} \cdot \mathbf{x}}$ where $\zeta (= e^{2\pi i/q})$ is the complex q -primitive root of unity. When $q = 2$, we obtain the Walsh transform of $f \in \mathcal{B}_n$, which will be denoted by \mathcal{W}_f .

A generalized Boolean function $f \in \mathcal{GB}_n^q$ is called *generalized bent* (or *gbent*, for short) if and only if $|\mathcal{H}_f(\boldsymbol{\omega})| = 2^{n/2}$ for all $\boldsymbol{\omega} \in \mathbb{Z}_2^n$. Note that when $q = 2$, Boolean bent functions exists only if the number n of variables is even. For $q > 2$, if f is a gbent function in n variables, it does not follow that n must be even. Such functions for $q = 4$ were investigated by Schmidt [6], Solé and Tokareva [7], Stănică, Martinsen, Gangopadhyay, and Singh [9], etc.

The sum $\mathcal{C}_{f,g}(\mathbf{u}) = \sum_{\mathbf{x} \in \mathbb{Z}_2^n} \zeta^{f(\mathbf{x})-g(\mathbf{x} \oplus \mathbf{u})}$ is the *crosscorrelation* of f and g at $\mathbf{u} \in \mathbb{Z}_2^n$. The *autocorrelation* of $f \in \mathcal{GB}_n^q$ at $\mathbf{u} \in \mathbb{Z}_2^n$ is $\mathcal{C}_{f,g}(\mathbf{u})$ above, which is denoted by $\mathcal{C}_f(\mathbf{u})$.

Lemma 1. Let $f \in \mathcal{GB}_n^q$. Then f is a gbent function if and only if

$$\mathcal{C}_f(\mathbf{u}) = \begin{cases} 2^n, & \text{if } \mathbf{u} = \mathbf{0}_n, \\ 0, & \text{if } \mathbf{u} \neq \mathbf{0}_n. \end{cases}$$

By using Lemma 1, we can prove the following theorem.

Theorem 1. Let n be a positive integer and m, q be even positive integers. Let $f \in \mathcal{GB}_n^q$ be gbent. Let $f + \frac{q}{2}g_i \in \mathcal{GB}_n^q$ be gbent, where $i = 0, 1$. Let $\mathbf{y} = (\mathbf{y}', \mathbf{y}'')$, $\mathbf{y}' = (y_1, y_2, \dots, y_{m/2})$, $\mathbf{y}'' = (y_{m/2+1}, y_{m/2+2} \dots y_m)$ and $\vartheta(\mathbf{y}) = \mathbf{y}' \cdot \mathbf{y}''$. Let $\mathbf{c} \in \mathbb{Z}_2^m$ and $\text{wt}(\mathbf{c})$ be even. Then the function $h \in \mathcal{GB}_n^q$, defined by

$$h(\mathbf{x}, \mathbf{y}) = f(\mathbf{x}) + \frac{q}{2}(\mathbf{c} \cdot \mathbf{y})g_{\mathbf{c} \cdot \mathbf{y}}(\mathbf{x}) + \frac{q}{2}\vartheta(\mathbf{y}) \quad (1)$$

is a gbent function in $n + m$ variables.

In Table 1, we compare our approach to other methods [9, 10] in terms of the form of gbent functions.

In what follows, we first provide some notations.

If $f \in \mathcal{B}_n$ is bent, then the dual function \tilde{f} of f , defined on \mathbb{Z}_2^n by $\mathcal{W}_{\tilde{f}}(\boldsymbol{\omega}) = 2^{n/2}(-1)^{\tilde{f}(\boldsymbol{\omega})}$ is also bent and it is known that $\tilde{\tilde{f}} = f$.

Lemma 2. For every $\mathbf{a}, \mathbf{b} \in \mathbb{Z}_2^n$ and for every bent function f , the dual of the function $f(\mathbf{x} \oplus \mathbf{b}) \oplus \mathbf{a} \cdot \mathbf{x}$ equals $\tilde{f}(\mathbf{x} \oplus \mathbf{a}) \oplus \mathbf{b} \cdot (\mathbf{x} \oplus \mathbf{a})$.

The original Maiorana-McFarland's (M - M) class of bent functions is the set of all the (bent) Boolean functions on $\mathbb{Z}_2^{2n} = \{(\mathbf{x}, \mathbf{y}) | \mathbf{x}, \mathbf{y} \in \mathbb{Z}_2^n\}$ of the form

$$f(\mathbf{x}, \mathbf{y}) = \mathbf{x} \cdot \phi(\mathbf{y}) \oplus g(\mathbf{y}), \quad (2)$$

where ϕ is any permutation of \mathbb{Z}_2^n and $g \in \mathcal{B}_n$.

Let $f \in \mathcal{B}_n$. If there exists an even integer $0 \leq r \leq n$, such that $\|\{\boldsymbol{\omega} | \mathcal{W}_f(\boldsymbol{\omega}) \neq 0, \boldsymbol{\omega} \in \mathbb{F}_2^n\}\| = 2^r$, where $\|\cdot\|$ denotes the size (cardinality) of a set, and $(\mathcal{W}_f(\boldsymbol{\omega}))^2$ equals 2^{2n-r} or 0, for every $\boldsymbol{\omega} \in \mathbb{F}_2^n$, then f is called an r -order plateaued function in n variables. If f is a $2\lceil(n-2)/2\rceil$ -order plateaued function in n variables, then f is also called a *semibent* function.

Let $f \in \mathcal{GB}_n^8$ be as

$$f(\mathbf{x}) = v_0(\mathbf{x}) + v_1(\mathbf{x}) \cdot 2 + v_2(\mathbf{x}) \cdot 2^2, \quad (3)$$

where $v_i(\mathbf{x}) \in \mathcal{B}_n, i = 0, 1, 2$.

In [9, Theorem 19], Stănică et al. presented a sufficient and necessary condition for a function f as in (3) to be gbent.

Theorem 2 ([9]). Let $f \in \mathcal{GB}_n^8$ be as in (3). The following are true:

(i) If n is even, then f is gbent if and only if $v_2, v_0 \oplus v_2, v_1 \oplus v_2, v_0 \oplus v_1 \oplus v_2$ are all bent, and $\mathcal{W}_{v_0 \oplus v_2}(\mathbf{u})\mathcal{W}_{v_1 \oplus v_2}(\mathbf{u}) = \mathcal{W}_{v_2}(\mathbf{u})\mathcal{W}_{v_0 \oplus v_1 \oplus v_2}(\mathbf{u})$ for all $\mathbf{u} \in \mathbb{Z}_2^n$;

(ii) If n is odd, then f is gbent if and only if $v_2, v_0 \oplus v_2, v_1 \oplus v_2, v_0 \oplus v_1 \oplus v_2$ are all semibent, and $\mathcal{W}_{v_0 \oplus v_2}(\mathbf{u}) = \mathcal{W}_{v_2}(\mathbf{u}) = 0$ and $|\mathcal{W}_{v_1 \oplus v_2}(\mathbf{u})| = |\mathcal{W}_{v_0 \oplus v_1 \oplus v_2}(\mathbf{u})| = 2^{\frac{n+1}{2}}$; or, $|\mathcal{W}_{v_0 \oplus v_2}(\mathbf{u})| = |\mathcal{W}_{v_2}(\mathbf{u})| = 2^{\frac{n+1}{2}}$ and $\mathcal{W}_{v_1 \oplus v_2}(\mathbf{u}) = \mathcal{W}_{v_0 \oplus v_1 \oplus v_2}(\mathbf{u}) = 0$, for all $\mathbf{u} \in \mathbb{Z}_2^n$.

From the above theorem, we know that the sufficient conditions that a function f as in (3) is gbent are abstract. Hence, we provide some sufficient conditions for a function f as in (3) to be gbent.

Theorem 3. Let n be an even integer, $v_0, v_1, v_2 \in \mathcal{B}_n$ and $f \in \mathcal{GB}_n^8$ be as in (3). The following v_0, v_1, v_2 satisfy the sufficient conditions of Theorem 2 for the even case.

(i) Let v_0, v_1, v_2 be bent functions and $v_2, v_0 \oplus v_2, v_1 \oplus v_2, v_0 \oplus v_1 \oplus v_2$ be all bent, and $(v_0 \oplus v_2)(\mathbf{x}) = \tilde{v}_0(\mathbf{x}) \oplus \tilde{v}_2(\mathbf{x}), (v_1 \oplus v_2)(\mathbf{x}) = \tilde{v}_1(\mathbf{x}) \oplus \tilde{v}_2(\mathbf{x}), (v_0 \oplus v_1 \oplus v_2)(\mathbf{x}) = \tilde{v}_0(\mathbf{x}) \oplus \tilde{v}_1(\mathbf{x}) \oplus \tilde{v}_2(\mathbf{x})$.

(ii) Let $v_2 \in \mathcal{B}_n$ be a bent function, $v_0 = v_1$ and $v_0 \oplus v_2$ be bent.

(iii) Let $v_0(\mathbf{x}) = \mathbf{a}_0 \cdot \mathbf{x}$ and $v_1(\mathbf{x}) = \mathbf{a}_1 \cdot \mathbf{x}$ respectively, be two linear functions. Let $v_2 \in \mathcal{B}_n$ be a bent function, and $\tilde{v}_2(\mathbf{x}) \oplus \tilde{v}_2(\mathbf{x} \oplus \mathbf{a}_0) \oplus \tilde{v}_2(\mathbf{x} \oplus \mathbf{a}_1) \oplus \tilde{v}_2(\mathbf{x} \oplus \mathbf{a}_0 \oplus \mathbf{a}_1) = 0$.

Table 1 Form of gbent functions comparison

Number of variables	q	From	Resource
$n + 2$	2	$h(\mathbf{x}, \mathbf{y}) = f(\mathbf{x}) \oplus (y_1 \oplus y_2)g(\mathbf{x}) \oplus y_1y_2$	Ref. [10]
$n + 2$	Even integer	$h(\mathbf{x}, \mathbf{y}) = f(\mathbf{x}) + (y_1 \oplus y_2)g(\mathbf{x}) + \frac{q}{2}y_1y_2$	Ref. [9]
$n + m$	Even integer	$h(\mathbf{x}, \mathbf{y}) = f(\mathbf{x}) + \frac{q}{2}(\mathbf{c} \cdot \mathbf{y})g_{\mathbf{c} \cdot \mathbf{y}}(\mathbf{x}) + \frac{q}{2}\vartheta(\mathbf{y})$	New

(iv) Let $v_0(\mathbf{x}) = \mathbf{a}_0 \cdot \mathbf{x}$, be a linear function. Let $v_2 \in \mathcal{B}_n$ be a bent function, $v_1 \oplus v_2$ be bent and $\widetilde{v_2}(\mathbf{x}) \oplus \widetilde{v_2}(\mathbf{x} \oplus \mathbf{a}_0) \oplus (v_1 \oplus v_2)(\mathbf{x}) \oplus (v_1 \oplus v_2)(\mathbf{x} \oplus \mathbf{a}_0) = 0$.

We now discuss the case when n is odd. Let n be a positive odd integer and $g_1, g_2 \in \mathcal{B}_n$. We say that g_1 and g_2 are *complementary semibent functions* in n variables if they are semibent (that is, $(n - 1)$ -order plateaued) functions and satisfy the property that $\mathcal{W}_{g_1}(\omega) = 0$ if and only if $\mathcal{W}_{g_2}(\omega) \neq 0$.

Lemma 3. Let n be an even integer and $f \in \mathcal{B}_n$. Then f is bent if and only if the two functions on \mathbb{Z}_2^{n-1} , $f(x_1, \dots, x_{j-1}, 0, x_{j+1}, \dots, x_n)$ and $f(x_1, \dots, x_{j-1}, 1, x_{j+1}, \dots, x_n)$, are complementary semibent functions on \mathbb{Z}_2^{n-1} , where $j = 1, \dots, n$.

Theorem 4. Let k, n be two integers and $n = 2k - 1$. Let $\varphi = (\varphi_1, \dots, \varphi_k), \phi = (\phi_1, \dots, \phi_k)$ be Boolean maps from \mathbb{Z}_2^k to \mathbb{Z}_2^k such that both ϕ and $\phi \oplus \varphi = (\phi_1 \oplus \varphi_1, \dots, \phi_k \oplus \varphi_k)$ are permutations on \mathbb{Z}_2^k . Set $\Delta_j = \{\phi(\mathbf{y}) | \mathbf{y} \in \mathbb{Z}_2^{j-1} \times \{0\} \times \mathbb{Z}_2^{k-j}\}$, $\mathbf{y}_\epsilon^{(j)} = (y_1, \dots, y_{j-1}, \epsilon, y_{j+1}, \dots, y_k)$, where $\epsilon \in \mathbb{Z}_2, j = 1, 2, \dots, k$. Let $f \in \mathcal{GB}_n^8$ be as in (3), and let $v_0(\mathbf{x}, \mathbf{y}_0^{(j)}) = \mathbf{a}_0 \cdot \mathbf{x} \oplus \varphi(\mathbf{y}_0^{(j)}) \cdot \mathbf{x}$, $v_1(\mathbf{x}) = (\phi(\mathbf{y}_0^{(j)}) \oplus \phi(\mathbf{y}_1^{(j)})) \cdot \mathbf{x} \oplus g(\mathbf{y}_0^{(j)}) \oplus g(\mathbf{y}_1^{(j)})$ and $v_2(\mathbf{x}) = \phi(\mathbf{y}_0^{(j)}) \cdot \mathbf{x} \oplus g(\mathbf{y}_0^{(j)})$, where $\mathbf{a}_0 \in \mathbb{Z}_2^k$. If there exists one positive integer $\varrho (\leq k)$ such that

$$\{(\phi \oplus \varphi)(\mathbf{y}) | \mathbf{y} \in \mathbb{Z}_2^{\varrho-1} \times \{0\} \times \mathbb{Z}_2^{k-\varrho}\} = \Delta_\varrho \quad (4)$$

(if $\mathbf{a}_0 \neq \mathbf{0}_k$ we further require Δ_ϱ to be a linear subspace of \mathbb{Z}_2^k and $\mathbf{a}_0 \in \Delta_\varrho$), then v_0, v_1, v_2 satisfy the conditions of Theorem 2 for the odd case, that is, f is gbent.

Acknowledgements This work was supported by National Natural Science Foundation of China (Grant Nos. 61303263, 61309034), and in part by Fundamen-

tal Research Funds for the Central Universities (Grant No. 2015XKMS086), and in part by China Postdoctoral Science Foundation Funded Project (Grant No. 2015T80600).

Supporting information The supporting information is available online at info.scichina.com and link.springer.com. The supporting materials are published as submitted, without typesetting or editing. The responsibility for scientific accuracy and content remains entirely with the authors.

References

- Rothaus O S. On ‘bent’ functions. *J Combin Theory A*, 1976, 20: 300–305
- Carlet C. Two new classes of bent functions. In: *Proceedings of Workshop on the Theory and Application of Cryptographic Techniques*, Lofthus, 1994. 77–101
- Zhang F, Wei Y, Pasalic E. Constructions of bent-negabent functions and their relation to the completed Maiorana-McFarland class. *IEEE Trans Inf Theory*, 2015, 61: 1496–1506
- Zhang W G, Xiao G Z. Constructions of almost optimal resilient Boolean functions on large even number of variables. *IEEE Trans Inf Theory*, 2009, 55: 5822–5831
- Zhang W G, Jiang F Q, Tang D. Construction of highly nonlinear resilient Boolean functions satisfying strict avalanche criterion. *Sci China Inf Sci*, 2014, 57: 049101
- Schmidt K U. Quaternary constant-amplitude codes for multicode CDMA. *IEEE Trans Inf Theory*, 2009, 55: 1824–1832
- Solé P, Tokareva N. Connections between quaternary and binary bent functions. <http://eprint.iacr.org/2009/544.pdf>
- Stănică P, Gangopadhyay S, Chaturvedi A, et al. Nega-Hadamard transform, bent and negabent functions. In: *Proceedings of 6th International Conference on Sequences and Their Applications*, Paris, 2010. 359–372
- Stănică P, Martinsen T, Gangopadhyay S, et al. Bent and generalized bent Boolean functions. *Des Codes Cryptogr*, 2013, 69: 77–94
- Zhao Y, Li H L. On bent functions with some symmetric properties. *Discret Appl Math*, 2006, 154: 2537–2543