

Computer-Aided Network Design by Optimization in the Frequency Domain

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ABSTRACT

The filter design problem is considered as an optimization problem. An iterative search technique is employed to adjust the variable network element values to approximate some desired network response, with a minimum of error. Explicit constraints are employed to ensure physical realizability. The design process uses a combination of a modified version of Calahan's network analysis program with a direct search method of minimization developed by Hooke and Jeeves. The result is a procedure which utilizes the circuit designer's experience and knowledge to set up the problem but relieves him of the tedious labor now performed by the high-speed digital computer.

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I. INTRODUCTION

A. COMPUTER-AIDED NETWORK DESIGN

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Mathematical programming techniques have found wide use in operations research, econom.cs, and other related fields. However, it has only been in recent years that such techniques have gained acceptance as tools for the design and avaluation of electronic circuits. The development of several general network-analysis programs has made computer-aided network design quite attractive. What is computer-aided network design? The circuit operation is first analyzed by means of a computer. It is then modified and analyzed again until the desired result is achieved--a trial-and-error procedure. Naturally the more experienced the engineer, the fewer the trials before a satisfactory design is realized.

The engineer today has a variety of analysis programs which may suit his needs in the design of networks. Some of the more well-known ones are: NET-1, ECAP, SCEPTRE, NASAP, CIRCUS, LISA, PANE, CAIAHAN, and CORNAP. Programs such as these have offered great assistance to the engineer in the analysis and design of networks. Although the obvious advantages in saving of time and tedious labor are quite apparent, there are certain features that would be desirable and perhaps possible in future programs of the type mentioned. Some of these features may include:

1. A graphical output on remote terminals which will allow the engineer to check his results and make on the spot changes as necessary.

2. Automatic means for improving the circuit design; i.e., some optimization technique to obtain "best" element values.

As valuable an aid as the computer is, a significant part of the design procedure will still require the engineering judgement of the designer. The cost for relieving the engineer of all the tedious calculations required for analysis is not an inexpensive one. The engineer must use his knowledge to specify the network topology, the response desired, constraints on element values, error criteria, reasonable initial element values, and other information which only he can provide.

B. USE OF OPTIMIZATION TECHNIQUES IN COMPUTE A-AIDED NETWORK DESIGN

The network designer is basically confronted with the problem of designing a circuit to meet some prescribed performance requirements. The design may be accomplished in one of many ways. If the requirements are such that an existing synthesis technique will provide the answer, the problem is essentially solved, and a satisfactory solution is obtained. In some caces a perfectly good design may be achieved in the laboratory by physically wiring the circuit on the "bread-board" and experimentally determining the "best" element values for the design.

There are classical synthesis techniques which provide a step-bystep design procedure, resulting in the circuit configuration and element values [1]. However, there are some design problems which may not be amenable to solution by any of the known synthesis techniques. The designer may be given a requirement in the form of a table of values or a graph of the response desired. Such a requirement cannot be satisfied by the classical synthesis techniques. If the circuit contains a large number of variable elements, design by a trial-and-error process in the laboratory is also highly unfeasible. Again, as with the analysis, the high-speed digital computer has offered an alternative approach to

the problem. We can now use some optimization technique to find element values in a given design configuration which yield a solution nearest to the prescribed performance requirement. The optimization technique iteratively adjusts the elecent values until the requirement is approximated as closely as possible.

Although synthesis techniques are available for the design of standard high- and low-pass filters, they do not take into account any constraints on the network configuration or element values. Problems of this nature would certainly be amenable to solution by an optimization technique. Networks whose transfer functions are extremely complex comprise another class of problems which could be solved by optimization. Optimization may also be used to obtain models for active devices. An optimization scheme could well be used in the design of matched filters. There are countless other examples, but suffice it to say that a combination of a good network analysis program and an efficient optimization program is certainly an excellent application of computer-aided network design.

C. OPTIMIZATION TECHNIQUES

In any optimization procedure two requirements must be satisfied. First, there must be some means to determine the behavior or performance of the system for any set of parameter values. Second, a performance measure must be selected which is a numerical measure of the behavior of the system. The optimization is basically a matter of minimizing the performance measure, which is a function of the parameters. If we think of this in geometric terms, the points in the parameter space represent different circuit element values and any change in the element

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value will result in movement to another point in the space. The performance measure is defined on this parameter space and requires an additional dimension if it is to be represented geometrically.

Minimization techniques generally do not yield the global minimum. What is found is a local minimum, but by changing the starting values of the parameters, it can be determined whether the local minimum is also the global minimum. If the minimization procedure converges to different values of the performance measure, the smallest value of the performance measure is then selected as the global minimum. The different minimization techniques may be classified by the method which is utilized to find a local minimum. They may be generally classified in the three following categories:

1. Direct search methods: those which do not compute the partial derivatives of the performance measure with respect to the parameters, but use only the value of the performance measure.

2. Gradient methods: those which require the calculation of the first partial derivatives of the performance measure with respect to the parameters.

3. Second-Order methods: those requiring higher-order partial derivatives.

No attempt will be made here to discuss the various methods under each category. An excellent discussion of the methods can be found in Ref. [2].

D. THE GENERAL NATURE OF THE PROBLEM

Earlier it was stated that for problems which cannot be solved by existing synthesis techniques or by experimenting with the circuit, an optimization technique may be feasible as an alternative solution.

Problems which are amenable to solution by optimization will generally be stated as follows: "Given a particular network with a fixed number of variable elements, adjust these elements until the response of the network minimizes some preassigned criterion". The key words in this general statement are "particular network with a fixed number of variable elements". For a particular desired response there may be several network configurations which will yield comparable results. The job of the designer then is to choose among the configurations he tries, and to select the best design which satisfies the requirements. One of the desirable features the engineer would like in future programs for computer-aided design is the ability of the program to also produce the network configuration as well as the element values for the optimal network for a given response. With the existing programs the engineer starts with a particular network configuration and a fixed number of variable elements are adjusted until the performance meets some preassigned criterion.

Now that the type of design problem is defined, the important features of the optimization process may be studied. Figure I-1 gives the essential elements which should be a part of the optimization technique for the circuit design problems.

The choice of the network configuration and the initial element values is the task of the engineer before the actual optimization process begins. The features to be discussed now are the evaluation of the response, the performance measure and the decision-making process.

The first thing the optimization program must be capable of doing is to evaluate the response from a given set of element values. This

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Fig. I-1 The Optimization Process

must be done each time the element values are adjusted. The analysis programs mentioned above all have this capability.

Once the response has been evaluated, it must be compared with the desired response. The measure of behavior is the performance measure. It is impossible to choose a single criterion and call it the universal performance measure. The nature of the problem may determine what performance measure is to be used. The experienced circuit designer will generally have in mind what performance measure is best for a given situation. The performance measure chosen must remain the same throughout

the design procedure. Some problems will dictate what performance measure is to be used, but in many cases the choice is purely subjective [3].

There are several typical forms of the performance measure the designer may use in the optimization process. The simplest form would be

$$J(\underline{p}) = \sum_{i=1}^{N} |R_A(f_i) - R_D(f_i)|$$

where $J(\underline{p})$ is the performance measure, a function of the parameter values \underline{p} , the terms $\underline{R}_A(\underline{f}_i)$ and $\underline{K}_D(\underline{f}_i)$ represent the actual and desired frequency responses, respectively, and N is the total number of points. This form of the performance measure indicates that only the magnitude of the difference is of interest, positive and negative deviations having equal weight. Another performance measure is

$$J(\underline{p}) = \sum_{i=1}^{N} \left[R_{A}(f_{i}) - R_{D}(f_{i}) \right]^{n}$$

where n is some even integer. When the difference between the actual and desired responses are small, say less than 1.0, then the value of this performance measure will decrease with increasing values of the exponent, n. These are just two examples of what may be used for performance measures. By defining the performance measure in some way such as mentioned above, the optimization procedure is one in which a search is conducted to find the minimum of the performance measure.

The search for this minimum generally results in finding a local minimum as mentioned above. Only if it is known that the function is unimodal will the local minimum be the global minimum. Otherwise, it is necessary to conduct a systematic search of the entire parameter

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space in order to locate the global minimum. For problems that have more than three or four variable parameters this may not be feasible. However, if the search is started with new initial values of the elements, and the function value converges to the same value, then one can be relatively certain that a global minimum has been located. By finding a minimum, whether it be local or global, a perfectly acceptable solution may be obtained. In the final analysis, it is the engineer's decision whether the final network configuration behaves in a satisfactory manner or whether further investigation is necessary to locate a "better" minimum.

This final aspect, the decision of the engineer, is perhaps the most critical item in the optimization. He must weigh the cost of further exploration to find a smaller minimum against the solution he already has. A great deal depends on what the circuit requirements are. The tolerances on the element values may be such that the procedure may have to be repeated again and again. On the other hand, there may be very weak restrictions on the element values so long as the response matches the desired response within say 0.1%. The computer relieves the engineer of the tedious work involved in any optimization procedure, but he is still responsible for making the knowledgable decisions to use the computer most advantageously.

II. THE OPTIMIZATION PROGRAM

A. ANALYSIS PROGRAM

As mentioned in Chapter I, the optimization procedures will include an analysis program and a minimization program. Among the various programs available, the CALAHAN and ECAP programs were considered as likely choices. Both programs were subjects of study in a course in computer-aided design, taught by Professor S. G. Chan, offered at the Naval Postgraduate School. The CALAHAN program was the final choice since it provides for a graphical output of the frequency response, an output not available with ECAP. Presumably the engineer who is designing by optimization will either choose a program which he has used successfully, or he will write a program to suit his needs.

1. A Linear Network Analysis Program

The CAIAHAN program is a general-purpose program designed for the analysis of linear electrical networks [4]. Input data to the program consists of the number of nodes in the circuit, the number of passive elements, the number of active elements, the input and output node numbers, and the type of output desired. A list of element values must be provided as well as a range of values of frequency (time) over which the frequency (transient) response is to be calculated. Outputs from the program are the coefficients of the specified network function, the poles and zeros, frequency and/or transient responses. The program is designed so that the user need only provide the required data cards, to obtain the desired output.

2. Modification of CALAHAN for Use in the Optimization Program

In order to incorporate CALAHAN into the optimization program, it was necessary to make some modifications to the original CALAHAN program. Before this was done, the original version was run a considerable number of times to yield frequency responses of circuits for which the actual responses were known. From the closed-form expression of the voltage transfer function of the Butterworth filter [5], the theoretical frequency responses were obtained for various orders of this type of filter. Using values of the normalized Butterworth filters [6], frequency responses were calculated by CALAHAN for different orders of the filter. The responses calculated by CALAHAN were almost exactly the same as the theoretical responses. The procedure was also repeated with the Chebyshev filter and similarly good results were obtained.

Since the goal of the optimization is to determine a set of element values for a particular network configuration whose frequency response is to match a given response as closely as possible, only the portion of CAIAHAN that calculates the frequency response is needed. The main program from the original CAIAHAN was reduced until only the portions involving the frequency response remained. Several subroutines that are not essential to the calculation of the frequency response were removed.

B. THE MINIMIZATION PROGRAM

The minimization program used in conjunction with CALAHAN to form the optimization program is a direct search technique [7]. This category of minimization techniques requires only the calculation of

the value of the performance measure, the calculation of derivatives not being a requirement. A gradient technique used in the design of filters by optimization is the subject of the Naval Postgraduate School thesis written by Major Charles A. Henry, USMC. Results obtained using the two methods are discussed and compared in Chapter IV.

1. Direct Search

Direct search may be basically described as a sequential examination of trial solutions which involves the comparison of the trial solution with the "best" solution obtained up to that time and a method for determining what the next trial solution will be [7]. Among the various types of problems which can be solved by direct search are solution of system of equations, curve-fitting problems, solution of integral equations, and minimizing (or maximizing) functions with or without constraints on the variables. The application of direct search methods to the solution of problems of the types mentioned above is basically the same regardless of the type of problem.

A space of P points, representing the solution space, must be defined. There must be some means to determine that one point is "better" than another. Presumably there is a "best" solution P^* in the solution space. Direct search is then accomplished in the following manner: A point B_1 , designated the first base point, is arbitrarily selected in the space. A second point, P_2 , is then selected and compared with B_1 . If P_2 is "better" than B_1 then P_2 becomes B_2 , the second base point. However if P_2 is not "better" than B_1 , then B_1 remains the base point. The process continues with each new point selected and compared with the current base point. The technique for selecting new trial points is determined by various conditions which

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arise as a function of results of trials made. The technique to be used in the minimization program is pattern search.

2. The Specific Technique--Pattern Search

Pattern search is a direct search technique for finding the minimum of a function F(p) of the variables $p = (p_1, p_2, p_3, \dots, p_N)^T$. The argument p is varied until a minimum value of F(p) is obtained. The successive values of p represent points in an N-dimensional space.

The operation of the pattern-search routine will now be described. First. a few definitions will be of aid in the ensuing discussion. The procedure of going from one point to another point is termed a move. If the value of F(p) decreases, then the move is a <u>success</u>; if the function F(p) does not decrease, then the move is a <u>failure</u>. Pattern search makes two types of moves. The <u>explore</u> move is to acquire knowledge about the behavior of the function F(p). The second type of move is the <u>pattern</u> move which utilizes the information gained from the explore moves to accomplish the actual minimization of the function by moving in the direction of an established pattern. The point from which a pattern move is made is known as a <u>base point</u>. Basically the pattern-search procedure is movement from base point to base point.

The explore move provides the information which indicates a probable direction for a successful move. A pattern is thus established. The pattern move from a given base point duplicates the combined moves from a previous base point if the direction of the pattern is unchanged. This process continues as long as the moves are successful, the step lengths increasing in magnitude. The result of each pattern move then is either a success or a failure. If the

Each explore move is carried out in the following manner: a single coordinate of the point is varied by either increasing or decreasing the coordinate by some fixed amount and seeing if the move is a success. If a success occurs, the new coordinate value is used; otherwise the old coordinate is retained. For each coordinate, these explore moves are made until the final point, as a result of all the explore moves, becomes the new base point.

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If, on the other hand, the pattern move is a failure, the search continues by retreating to the base point and starting over again with new explore moves until a new pattern is established.

The pattern-search technique can be better understood with the aid of a simple example. Figure II-1 serves as an illustration of what has been discussed in the previous paragraphs. A two-dimensional parameter space is shown with equal-cost contours represented by F_1 , F_2 , ..., F_8 ; where $F_k > F_{k+1}$. The argument of the function F is $p = (p_1, p_2)^T$. The point B_1 is selected as the first base point and explore moves are conducted from this point. First the p_1 coordinate is stepped in its positive and negative directions; a success is achieved when the step is negative. Next the p_2 coordinate is tested and it is determined that a positive step yields a success. The explore move produces a new base point B_2 . The most probable direction of a success is in the direction of the line segment $\overline{B_1B_2}$; therefore, the patteru is established and the pattern move results in TP_1 , a temporary point. Each pattern move is followed by a series of explore moves. If the result of the explore moves is a success, then the pattern move is termed a



success. The point resulting from the successful vplore moves becomes the new base point. If, however, the explore moves are failures, the function value at the temporary point is calculated and if this value is greater than the function value at the previous base point, the old base point becomes the new base point and explore moves are tried again. At TP, explore trials are made and a successful move is made to B, which becomes the new base point. This indicates that the pattern move to TP₁ was a success. By the same line of reasoning, the pattern move to TP_2 is a success. The pattern move to TP_3 is a failure since all explore moves from this point are failures and the value of the function at TP₃ is greater than the value at E_4 ; therefore B_4 becomes B_5 , the new base point, and the explore moves are tried again. The region within the F_8 contour is enlarged by a factor of five and shown in Fig. II-2. The point B_6 is a successful explore but it should be noted that there is no change in the p_2 coordinate from B_5 . A change in the p_2 coordinate would yield a failure. The pattern move to TP4 is a failure, so B₆ now becomes B₇ and explore moves are made. Perturbations of both p_1 and p_2 in the original step size do not produce any successful explore moves. It is, therefore, necessary to reduce the step size by some fixed amount. The step reduction factor in this case is 0.2, which means that the original step size has now been recuced by a factor of five. Once again explore trials are made, now with the new step size, and a success is achieved at B_g . The pattern move to TP_g is a failure, so B_8 becomes B_q . A successful explore move is schieved at B_{10} but the pattern move to TP_6 is a failure and B_{10} becomes B_{11} . This process of reducing the step size and then making the pattern moves continues until the difference between two consecutive steps is less than



some prescribed amount. If this criterion is a very small number then the step size will be sufficiently small to ensure that the minimum bas been closely approximated. Care must be taken in the choice of both the step size and the step reduction factor. Too large a reduction factor will result in a slowdown of the search procedure. If an imitial step size is too large, the minimum may be missed altogether.

The direct search procedure described above is termed pattern search since the minimization is basically performed by the pattern moves. Although the explore moves provide some reduction, their main purpose is to provide information for the improvement of the pattern move. The pattern-search program used in the optimization program is a program written by R. Halleary¹, with some modifications to include constraints on the independent variables.

C. THE OPTIMIZATION PROGRAM -- A COMBINATION

In sections A and B the individual programs in the optimization program were discussed in some detail. A brief description was given of the modifications to CALAHAN to accommodate the particular problems to be considered. How are CALAHAN and DIRECT together to be implemented into one program to be used in the design of networks by optimization? The answer to this question is the subject of this section.

The basic type of filter design problem which will be solved by the optimization program is one in which a particular frequency response is given and the objective is to design a filter which approximates the response as closely as possible. In the next chapter the exact problem

¹Subroutine DIRECT, Naval Postgraduate School Computer Facility.

problem formulation and specific example problems will be discussed in detail, but the above problem description is adequate for a general discussion of the optimization program. After the particular circuit has been selected, a choice of initial element values must be made, the engineer's knowledge and experience playing a vital role in the choice. Other information which must be supplied to the program includes the following: the number of frequency points to be matched, the values of the desired frequency response at the points to be matched, the explicit constraints on the element values, the step size, the step reduction factor, and the termination criterion for the minimization.

The initial element values serve as coordinates of the first base point for the pattern search routine, called DIRECT. An external function subprogram then calculates the value of the function to be minimized by DIRECT. The exact form of this function may differ for different problems, but in all cases it is a comparison between the actual frequency response and the desired response. This calculation is performed as part of the function subprogram utilizing the modified version of CAIAHAN. Once the minimum has been determined, the element values producing the minimum are supplied to the analysis program and the actual frequency response is calculated. This process may be repeated until the overall design satisfies all of the requirements.

III. IMPLEMENTATION OF THE OPTIMIZATION PROGRAM

The use of the optimization program is dictated by the requirements for the design. An optimization technique should be used only when classical synthesis methods and experimental methods are either impossible or unfeasible. The purpose of this chapter is to discuss the use of the optimization program in circuit design. The first section of the chapter is a general discussion of the features of the program. In the concluding section several examples are given to illustrate the use of the program.

A. PROGRAM FEATURES

1. The Input Data

For input data, the optimization program requires a topological description of the network, a list of element values, a range of frequencies over which the desired and actual responses are to be matched, a description of the desired response, the number of varying and non-varying elements, and a list of the constraints on the varying elements. The following information required by DIRECT must also be included in the input data: the step size, the step reduction factor, the termination criterion, and the maximum allowable number of evaluations. Figure III-lis a flow chart showing the sequence and coding of the input data cards.

2. The Output

The output from the optimization program consists of two parts. The first part is a result of the minimization and includes the value of the function at convergence, and the optimum values of the variable

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Fig. III-1 Coding Flow Chart for Optimization Program

Next Card(s) Active Elements Node No. of Active Element - Node No. of Controlling Branch Co1 / 12345678 + 012345678 + 012345678 90 Lyalue of G or M Transconductance G or Mutual Inductance M FORMAT(4(12,1X),A1,1X,F10.0)

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Sext Card(s) Desired Prequency Response PORMAT(P15.7)

Fig. III-1 (continued)

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Fig. III-1 (continued)

elements. The second part of the output is a result of the analysis program. With the optimum element values calculated by DIRECT, the frequency response is calculated. The output is in tabular form as well as in a graphical form. The circuit designer merely compares the values of the calculated response with those of the desired response. If the design requirements are satisfied, the element values calculated in the first part of the output are the element values of the design.

3. Accuracy of the Optimization Program

In all examples used in testing the program, the performance measure to be minimized was of the form,

$$J(\underline{p}) = \sum_{i=1}^{N} \left[R_{A} (f_{i}) - R_{D} (f_{i}) \right]^{n}$$

where $J(\underline{p})$ is a function of the network elements \underline{p} . \mathbb{R}_{A} (f_{i}) is the actual frequency response at the ith comparison point and \mathbb{R}_{D} (f_{i}) is the desired response at the same point of comparison. N is the total number of frequency points to be compared and n is some even integer. Theoretically, the minimum of this performance measure is zero if a perfect match of frequency responses is achieved. In practice, however, a zero output is rarely, if ever, achieved. The measure of accuracy is determined by the function value at exit from DIRECT; the smaller the function value, the closer the actual response approaches the desired response.

The accuracy of the output is basically dependent upon the choice of the termination criterion for DIRECT. The pattern search ends when the difference between consecutive step sizes falls below this pre-selected termination criterion. A small criterion will result in a small function value, consequently a closer approximation to the

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desired response. The program allows the user to specify the termination criterion as part of the input. Table III-1 shows a comparison of execution times, and function values, for a normalized fourth-order Butterworth filter, as a function of the step size, the step reduction factor, and the termination criterion. The performance measure used was

$$J(\underline{p}) = \sum_{i=1}^{21} \left[R_A(f_i) - R_D(f_i) \right]^2.$$

The function values for a termination criterion of 10^{-4} differ by a factor of 100 from those for a termination criterion of 10^{-6} ; whereas the difference between function values for termination criteria 10^{-6} and 10^{-9} is insignificant. In this case there is no particular advantage in the choice of a termination criterion less than 10^{-6} , since the function values only change slightly but the execution times are longer. A comparison between the desired response, for the frequency range specified, and the largest and smallest function values is given in Table III-2.

4. Execution Time

The execution time for the program is dependent on seve: 11 factors which will be discussed in this section. The initial choice of element values will certainly affect the execution time; if the initial guess is a poor choice the program may take an inordinate amount of time if it converges at all to a minimum. Convergence to a minimum also may be quite slow for circuits with a large number of elements. The only solution to this problem is to choose a simpler circuit configuration which may yield a response within acceptable tolerances. In the pattern search, the execution time is a function of the termination criterion. The choice of the termination criterion is a compromise between speed and accuracy; one is sacrificed for the other. If more

TABLE	III-1

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<u>Trial</u>	Step <u>Size</u>	Step Red. Factor	Termination Criterion	Execution Time (Sec)	Function Value x 10 ⁸
A	0.05	0.25	10 ⁻⁴	50.38	203.67
В	0.05	0.25	10-6	68,50	0.31191
C	C.05	0.25	10 ⁻⁹	78.01	0.23612
σ	0.05	0.125	10 ⁻⁴	54,93	105,01
E	0.05	0.125	10 ⁻⁶	70.18	0.11424
F	0 .0 5	0.125	10 ⁻⁹	71.77	0.11424
G	0.1	0.25	10-4	46.37	386,91
H	0.1	0.25	10 ⁻⁶	70.90	0.55285
I	0.1	0.25	10 ⁻⁹	?8.97	0.54994
J	0.1	0.125	10 ⁻⁴	49.09	60.625
K	0.1	0.125	10 ⁻⁶	57.52	7.753
L	0.1	0.125	10 ⁻⁹	63.03	7.723
M	0.5	0.25	10 ⁻⁴	68.27	16.730
N	0.5	0.25	10 ⁻⁶	87.31	0.14581
0	0.5	0.25	10 ⁻⁹	90.89	0.14581
Р	0.5	0.125	10 ⁻⁴	85.22	284.77
Q	0.5	0.125	10 ⁻⁶	104.78	0.12187
R	0.5	0.125	10 ⁻⁹	110.14	0.12187

TABLE III-2

Desired Response	Trial G <u>Response</u>	Trial F <u>Response</u>
- 0.1042320	- 0.1032002	- 0.1042283
- 0.2204427	- 0.2196893	- 0.2204416
- 0.4314420	- 0.4310329	- 0.4314427
0.7850979	- 0.7850114	- 0.7851012
- 1.3305276	- 1.3306713	- 1.3305340
- 2.1019602	- 2.1021585	- 2.1019697
- 3.1032674	- 3.1033316	- 3.1032581
- 4.3047100	- 4.3045549	- 4.3047056
- 5.6547211	- 5.6543004	- 5.6547127
- 7.0972899	- 7.0966568	- 7.0972862
- 8.5844102	- 8.5836172	- 8.5844040
-10.0806382	-10.0797758	-10.0806456
-11.5623696	-11.5614929	-11.5623751
-13.0151248	-13.0142879	-13.0151329
-14.4307494	-14.4300060	-14.4307604
-15.8052081	-15.8045635	-15.8052015
-17.1370394	-17.1365509	-17.1370392
-18.4263366	-18.4259949	-18.4263458
-19.6740983	-19.6738892	-19.6740875
-20.8818219	-20.8817902	-20.8818054
-22.0512536	-22.0513916	-22.0512390

accurate results are required then the execution time is necessarily longer. Table III-1 shows the effects of different step reduction factors and termination criteria. A further comparison of execution time as a function of the number of points compared is made for the normalized fourth-order Butterworth filter. The results of this comparison are shown in Table III-3.

TABLE ITI-3

No. Points	Execution Time (Sec)	Function Value
50	147.06	0.8699953×10^{-8}
40	112.77	0.4037205×10^{-7}
30	83.93	0.3040103×10^{-8}
20	71.77	0.1142364×10^{-8}

B. DESIGN XXAMPLES

To illustrate the use of the optimization program, several examples of filter design will be discussed in this section. In all of the examples, the desired frequency response is in the form of a table of values. These values are to be matched as closely as possible by the circuit configuration selected. In general, design specifications are not quite as stringent as this. A more likely specification would be to design a maximally flat filter in a pass band whose cut-off frequencies are at f_1 and f_2 and with a dropoff of a specified number of db per octave; however, to illustrate the capability of the program, point-by-point comparisons will be made. Example 1

Problem: Find the optimum element values for the filter configuration shown in Fig. III-2, whose frequency response from 0.15Hz to 0.24Hz most closely approximates the 5th-order Butterworth response over the same range of frequencies.



Fig. III-2 Circuit for Example 1

Constraints on element values

 $0.5 \le L_1 \le 1.75$ $1.0 \le C_2 \le 2.5$ $1.0 \le L_3 \le 1.6 \quad \mathbb{R} = 1.0$ $0.4 \le C_4 \le 1.0$ $0.1 \le L_5 \le 0.75$

Solution:

Trial 1 Initial guess: $L_1=1.0$, $C_2=2.0$, $L_3=1.5$, $C_4=.8$, $L_5=.5$ At exit from program: $L_1=1.29$, $C_2=2.12$, $L_3=1.19$, $C_4=.888$, $L_5=.161$ Function value = .2435 x 10⁻⁶

Trial 2 Initial guess: $L_1=1.6$, $C_2=2.0$, $L_3=1.5$, $C_4=.8$, $L_5=.5$ At exit from program: $L_1=1.7$, $C_2=1.5$, $L_3=1.51$, $C_4=.875$, $L_5=.372$ Function value = .6473 x 10⁻⁷

Trial 3 Initial guess: $L_1 = 1.6$, $C_2 = 2.0$, $L_3 = 1.5$, $C_4 = .9$, $L_5 = .5$ At exact from program: $L_1 = 1.54$, $C_2 = 1.7$, $L_3 = 1.38$, $C_4 = .895$, $L_5 = .307$ Function value = .5150 x 10⁻¹³

Trial 4 Initial guess: $L_1=1.5$, $C_2=2.0$, $L_3=1.5$, $C_4=.9$, $L_5=.3$ At exit from program: $L_1=1.43$, $C_2=1.87$, $L_3=1.29$, $C_4=.903$, $L_5=.253$ Function value = .1384 x 10⁻⁷

Trial 5 Initial guess: $L_1 = 1.6$, $C_2 = 1.8$, $L_3 = 1.5$, $C_4 = .9$, $L_5 = .4$ At exit from program: $L_1 = 1.54$, $C_2 = 1.7$, $L_3 = 1.38$, $C_4 = .895$, $L_5 = .307$ Function value = .429 x 10⁻¹³

The optimum element values are those values calculated in trials 3 and 5. Comparison of the trial frequency responses with the Butterworth response is shown in Table III-4.

Discussion--In this problem only ten points were compared, if more accurate results are desired more points should be compared. The element values for trials 3 and 5 are very close to the values for the fifth order normalized Butterworth filter.

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TABLE III-4

req.	Butterworth	Trial 1	Trial 2	Trial 3	Trial 4	<u>Trial 5</u>
. 15	- 1.9116645	- 1.8978281	- 1.9217176	- 1.9116898	- 1.9047451	- 1.9116688
.16	- 3.1264136	- 3,1432962	- 3.1149416	- 3.1270609	- 3,1350603	- 3,1270981
.17	- 4.6734886	- 4.6764345	- 4.6712990	- 4.6733904	- 4.6751757	- 4.6734352
.18	- 6.4580720	- 6.4457130	- 6,4668531	- 6.4577456	- 6.4521284	- 6.4577742
.19	- 8.3755182	- 8.3625135	- 8.3847713	- 8.3752613	- 8,3690329	- 8.3752661
.20	-10.3421541	-10.3396997	-10.3439083	-10.3421764	-10.3406610	-10.3421516
.21	-12.3032792	-12.3126831	-12.2965260	-12.3035507	-12.3075638	-12.3335259
.22	-14.2275041	-14.2418127	-14.2172155	- 14 . 227 8433	-14.2342997	-14.2279185
.23	-16.0987370	-16.1062012	-16,0933380	-16.0988770	-1.6.1024323	-16.1988617
.24	-17.9099567	-17.8969879	-17.9191589	-17.9096069	-17.9040222	-17.4096222
Example 2

Problem: Design a filter which approximates the straight-line characteristic shown in Fig. III-4.

Solution: Select the circuit configuration by determining the slope of the straight line after cutoff. The slope is approximately 24db per octave. Each 6db/octave represents one order of a low-pass filter; therefore the circuit to be used for the design is a fourth-order lowpass filter as shown in Fig. 111-3.



Fig. III-3 Circuit for Example 2

Results:

The first initial guess: $L_1=2.0$, $C_2=2.0$, $L_3=2.0$, $C_4=2.0$ At exit from program: $L_1=1.44$, $C_2=1.62$, $L_3=1.10$, $C_4=0.379$ The second initial guess: $L_1=1.5$, $C_2=1.5$, $L_3=1.0$, $C_4=0.5$ At exit from program: $L_1=1.44$, $C_2=1.63$, $L_3=1.10$, $C_4=0.381$

A plot of the actual and desired responses is shown in Fig. III-4.





Example 3

Problem: Design a filter that has a Gaussian distribution response in the frequency range from OHz to 4Hz.

Solution: The first step in the solution is to change the Gaussian response from a voltage ratio to db for use in the optimization program. Table III-5 contains the values for a 21-point comparison. The response is plotted in Fig. III-5.

The first trial design was a ninth-order low-pass filter. The resulting response is plotted in Fig. III-5 showing rather marked deviations from the desired response. At the lower frequencies the deviations are much greater.

The second trial design was a modified fifth-order low-pass filter. The response is plotted in Fig. III-5. There is a slight ' improvement in the approximation; however the deviations at some points are quite large.

Discussion: In this problem, only low-pass filter configurations were considered. Both design responses deviated considerably from the desired responses. This points out a limitation of the optimization program. The success of the optimization technique is dependent upon the circuit configuration selected. In this example, presumably there is a better circuit configuration which would approximate the desired response with less deviation.

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TABLE	III-5
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Freq.	Desired Gain	Design 1 Gain	Design 2 Cain
0.	0.	0.	0
0.2	- 0,174	0 389	- 1 296
0.4	- 0 693	1 927	- 1.270
0.4	- 0.095	1.027	- 2.880
0.0	- 1.563	1.172	- 3.000
0.8	- 2.779	- 3.918	- 1.969
1.0	- 4.341	- 7.501	- 3.687
1.2	- 0.253	- 8.378	- 8.179
1.4	- 8.512	- 6.531	-11.557
1.6	-11.119	-10.663	-13.150
1.8	-14.067	-17.770	-12.961
2.0	-17.368	-22.048	-14.552
2.2	-21.012	-23.421	-21.813
2.4	-25.005	-18.813	-28.598
2.6	-29.345	-22.974	-34.113
2.8	-34.067	-36.294	-38.755
3.0	-39.172	-44.084	-42.794
3.2	-44.437	-50.013	-46.388
3.4	-50.458	-54.899	-49.640
3.6	-56.478	-59.072	-52.618
3.8	-61.938	-62.697	-55.371
4.0	-70.458	-65.865	-57.933



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Example 4

Problem: Design a simple bandpass filter to match the desired frequency response of Table III-6.

Solution: A simple third-order low-pass filter is transformed into a bandpass filter by frequency transformation. The resultant circuit is shown in Fig. III-6.



Fig. III-6 Bandpass Filter Design

Results:

The initial guess: $L_1=0.5$, $C_2=1.2$, $C_3=1.2$, $L_4=0.5$, $L_5=0.75$, $C_6=0.75$ Exit from program: $L_1=0.278$, $C_2=0.968$, $C_3=1.07$, $L_4=0.231$, $L_5=0.512$, $C_6=0.5$ Results are tabulated in Table III-6.

TABLE III-6

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Freq.	Desired Gain	Actual Cain
0.12	-10,6043911	-10 6046772
0.14	7.0007050	-10.00-0/72
0.24	- /.800/050	- 7.8000298
0.16	- 6.5953960	- 6.5954 628
0.18	- 6.1859341	- 6.1866503
0.20	- 6.0634375	- 6.0643120
6.22	- 6.0301752	- 6.0309124
0.24	- 6.0222807	- 6.0227900
0.26	- 6.0207853	- 6.0211134
0.28	- 6.0205956	- 6.0207987
0.30	- 6.0205870	- 6.0207443
0.32	- 6.0205832	- 6.0207691
0.34	- 6.0205908	- 6.0208607
0.36	- 6.0206118	- 6.0210257
0.38	- 6.0207939	- 6.0213528
0.40	- 6.0214853	- 6.0221939
0.42	- 6.0233574	- 6.0242023
0.44	- 6.0275412	- 6.0284853
0.46	- 6.0356464	- 6.0366526
0.48	- 6.0498857	- 6.0509090
0.50	- 6.0731249	- 6.0740948
0.52	- 6.1088524	- 6.1096964
0.54	- 6.1611710	- 6.1618414
0.65	- 6.2347078	- 6.2351713

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TABLE III-6 (continued)

Freq.	Desired Gain	Actual Gain
0.58	- 6.3344698	- 6.3346691
0.60	- 6.4655190	- 6,4654360
0.62	- 6.6326761	- 6.632.3280
0.64	- 6.8401451	- 6.8395615
0.66	- 7.0910606	- 7.0902948
0.68	- 7.3872252	- 7.3863807
0.70	- 7.7288332	- 7.7280092
0.72	- 8.1145258	- 8.1138344
0.74	- 8.5414162	- 8.5409737
0.76	- 9.0054874	- 9.0054045
0.78	- 9.5018768	- 9.5022497
0.80	-10.0253372	-10.0262442
0.82	-10.5705452	-10.5720444

Example 5

Problem: Determine the element values of a fourth-order low-pass filter whose frequency response approximates the response of a fourthorder Butterworth filter using:

(a) ideal elements

- (b) inductances with nominal resistance of 0.01 ohrs
- (c) inductances with nominal resistance of 0.5 ohms

Solution: The circuit selected is the same as in Fig. III-3 but there are series resistors with the inductances when non-ideal elements are considered. The initial guess for the elements is the same for all three situations. All parameters for DIRECT remain the same. The frequency range is from 0.1 Hz to 0.3 Hz, comparing 21 points. The zesults are shown in Table III-7. Figure III-7 is a comparison of the frequency responses for circuits with ideal and non-ideal elements.

Discussion: For nominal resistances of 0.01 ohms the element values did not change much from the values of the ideal elements since the resistances are so small. When the resistance is of the same order of magnitude as the inductance then the final element values differ considerably from the ideal element values. Also, with non-ideal elements the frequency response as shown in Fig. III-7 is attenuated at the lowfrequency end.

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Fig. III-7 Response of Ideal and Non-ideal Circuits

TABLE III-7

Ideal Element Values	Non-ideal Fl	ement Values
$R_s = 0$	$R_{s} = 0.01$	$R_{s} = 0.5$
$L_1 = 1.53$	$L_1 = 1.51$	$L_1 = 1.11$
C ₂ = 1.58	$C_2 = 1.59$	$c_2 = 1.31$
$L_3 = 1.08$	$L_3 = 1.09$	$L_3 = 1.53$
$C_4 = 0.383$	$C_4 = 0.386$	$C_4 = 0.215$

IV. SUMMARY AND CONCLUSIONS

The subject of computer-aided design by optimization techniques, although only one facet of computer-aided design, is in itself quite a diverse field. There is a large variety of optimization methous which he effectively employed in network design. The main reason for using an optimization technique instead of a classical synthesis technique in circuit des_gn is that classical techniques cannot satisfy all possible design specifications. A specification such as matching the response of a circuit to some desired response given by a table of values or a graph cannot be realized by classical techniques. Constraints on circuit element values generally cannot be accommodated by classical methods. Such design specifications which cannot be realized by classical methods can often be satisfied by optimization techniques.

A. SUMMARY

Chapter I is an introductory chapter presenting a general discussion of computer-aided design and application of optimization techniques in computer-aided design. The three basic categories of optimization techniques are described and the general nature of the problem is presented.

In Chapter II the optimization program is described. The optimization program used 's a combination of the linear network analysis program by Calahan and the pattern-search technique for minimization of a function of several variables. Modifications to the original analysis program were made in order to incorporate it into the

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optimization program. Basically this was a matter of reducing the size of CALAHAN, since only the portion perraining to frequency response was required. The specific method of pattern search, DIRECT, in conjunction with the molified version of CALAHAN constitute the optimization program.

The specific details regarding the implementation of the optimization program are included in Chapter III. Instructions for coding of input data cards are shown as a coding flowchart in Fig. III-1. The factors affecting the accuracy of the program are shown by the data o: Table III-1. A comparison of the accuracy of the program for the worst and best approximations over a series of trial runs is shown in Table III-1 and Table III-3. Five lesign examples are provided at the end of the chapter.

B. CAPABILITIES AND LIMITATIONS OF THE PROGRAM

For the designs attempted, results indicated that the optimization program is highly accurate and relatively fast. A comparative study between the gradient-projection method described in the Naval Postgraduate School thesis by Major C. A. Henry, and the pattern-search method was conducted to determine the relative accuracy and speed of the two methods. Examples 2, 4, and 5(a) in Chapter III were selected for the comparisons. The results are shown in Table IV-1. Very accurate results can be achieved, as shown by Example 4 in Chapter III. On the other hand, results may deviate considerably from what is desired, as illustrated in Example 3 in Chapter III. A high degree of accuracy can be achieved if the circuit configuration chosen is the proper one for the desired response. At present there is no known

optimization program that automatically alters the configuration of the network to yield an optimum solution.

TABLE IV-1

Design Problem	Method	Function Value	Execution Time(sec)
Straight-Line	Gradient Projection	1.355	132
Approx.	Pattern Search	1.651	50
Fifth-Order	Grac. wit Projection	0.458x10 ⁻²	133
Butterworth	Pattern Search	0.204x10 ⁻⁵	49
Band Pass	Gradient Projection	9.564x10 ⁻⁴	328
	Pattern Search	0.220x10 ⁻⁵	184

One of the main limitations of the optimization program is that an excessive execution time is required for circuits with more than 12 or 13 elements. The reason for this is that CALAHAN finds the tree for the network each time the elements are perturbed. This is not necessary since the circuit configuration remains the same throughout the optimization process; however no attempt was made to alter this.

The total memory requirements for the program are approximately 110 K bytes. This may or may not present a problem depending upon the computer system available to the circuit designer.

C. FUTURE REFINEMENTS

Possible areas in which the program may be improved or implemented are:

(1) Modification of the tree-finding process so that the tree is found only once for each circuit configuration.

(2) Use of the program to optimize active networks.

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(3) Development of a means to "grow" elements; i.e. development of a technique that will change the circuit configuration. In this manner the circuit configuration as well as optimum element values would be calculated. STATES AND IN COMPANY

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С С C C C C A NETWORK OPTIMIZATION PROGRAM

CCCCCCCCCCCCCC A COMPINATION OF THE CALAHAN LINEAR NETWORK ANALYSIS PROGRAM AND THE DIRECT SEARCH MINIMIZATION PROGRAM FOR THE SOLUTION OF CIRCUIT DESIGN PROBLEMS MAIN PROGRAM ExTERNAL FE DIMENSION MP(:00,3),ML(5C,5),ELT(100),MAP(20,5),ELTA 1(20),VAL(100),VALA(20),C(50),G(50),H(50),VII(60),VI2 2(60),V(6C),Z(6C),V2I(60),V22(60),VALL(50),ZZ(60,2),D 3(1CC),R2(100),PP(60,2),BU(15),BL(15) COMMON VAL,UMGMIN,CMGMAX,Y,R2,D,ELT,ELTA,VALA,V11,V12, 1VALL,Y2,Y22,Z,Z,PP,LIN,NOM.JP,JZ,KEY1,ND,NPL,NN,JI, 2KI,JO,KG,NAL,KEY2,MP,MAP,JW,NVAR,KEY3,NRES REAL IHC / 4HC 10 CLOCK=ITIME(0)**01 READ(5,11,END=26) LABEL 11 FORMAT(2CA4) WRITE (6,12) LABEL 12 FORMAT (1H1,10X,2CA4) READ(5,13)NPL,NAL,NN,JI,KI,JO,KO,(MP(J,1),MP(J,2),ELT 1(J),VAL(J),J=1,NPL) 13 FCRMAT(7/,10X,19HCIRCUIT INPUT DATA MRITE(6,14) 14 FCRMAT(7/,10X,19HCIRCUIT INPUT DATA MRITE(6,15) NPL,NAL,NN,JI,KI,JO,KO,(MP(J,1),MP(J,2), 14 ELT(J),VAL(J),J=1,NPL) 15 FCRMAT(7/,10X,19HCIRCUIT INPUT DATA MRITE(6,16) NPL,NAL,NN,JI,KI,JO,KO,(MP(J,1),MP(J,2), 16 LT(J),VAL(J),J=1,NPL) 15 FCRMAT(7/,10X,19HCIRCUIT INPUT DATA MRITE(6,16) NPL,NAL,NN,JI,KI,JO,KO,(MP(J,1),MP(J,2), 16 LT(J),VAL(J),J=1,NPL) 15 FCRMAT(7/,10X,19HCIRCUIT INPUT DATA MRITE(6,16) NPL,NAL,NN,JI,KI,JO,KO,(MP(J,1),MP(J,2), 16 LT(J),VAL(J),J=1,NPL) 15 FCRMAT(9X,7(12,1X)/(9X,12,1X,12,1X,A1,1X,F18.9)) 16 ACTIVE ELEMENTS,READ AND PRINT 17 F(NAL)16,19,16 16 AREAD (5,17) ((MAP(J,1),J=1,4),ELTA(J),VALA(J),J=1,NAL) 17 FCRMAT(412,1X),A1,1X,F10.0) MRITE(6,16)(MAP(J,1),I=1,4),ELTA(J),VALA(J),J=1,NAL) 18 NCRMAT(9X,4(12,1X),A1,1X,F18.9) 19 KEY1=1 NVAR=0 NVAR=0 NVAL=C PAAD(5,2C)LIN,NOM,OMGMIN,OMGMAX EXTERNAL FE C C NVAL=C READ(5,2C)LIN,NOM,OMGMIN,OMGMAX FORMAT(I1,1X,I3/2F10.C) ND INCLUDES THE STARTING POINT 20 ND INCLODES THE STARTING FOLT NC=NCM+1 D(I) IS THE DESIRED FREQUENCY RESPONSE READ(5,21)(D(I),I=1,ND) FCRMAT(F15.7) 21 JW=1 NRES IS THE NUMBER OF CONSTANT ELEMENTS NFX IS THE NUMBER OF VARIABLE ELEMENTS READ(5,22) NRES,NEX 22 FCRMAT(2I5) READ(5,23)DEL,RHO,DEC,MAXEV 23 FORMAT(3FI5.7,I5) BU(1) IS THE UPPER BOUND HL(1) IS THE LOWER BOUND READ(5,24)(BU(1),BL(1),I=1,NFX) 24 FCRMAT(2FI5.7) NPM=NPL-NRES CALL DIRECT(VAL,NPM,SPSI,DEL,RHO,DEC,FE,KON,MAXEV,-1, 1BU,BL) Jh=1 CALL DIRECT (VAL, NPM, SPSI, DEL, RHO, DEC, FE, KON, MAXEV, -IBU, BL) KEY1=2 CALL FRECO(LIN, NOM, OMGMIN, OMGMAX, JP, JZ, Y, Z, KEY1, R2) CLOCK=ITIME(0)*.01-CLOCK

SUBROUTINE DIRECT

TC LOCATE A MINIMUM OF A FUNCTION, S, OF K VARIABLES BY THE METHOD OF DIRECT SEARCH (HOOKE AND JEEVES)

DESCRIPTION OF PARAMETERS

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PSI IS THE VECTOR OF K INDEPENDENT VARIABLES. IT IS INITIALLY FILLED BY USER WITH FIRST GUESS OF SOLUTION AT EXIT FROM DIRECT IT CONTAINS BEST VALUES ATTAINED. K IS THE NO. OF INDEPENDENT VARIABLES OF THE FUNCTION, S, TO BE MINIMIZED SPSI AT EXIT FROM DIRECT CONTAINS SMALLEST S(PSI) ATTAINED DELCAP IS THE INITIAL STEP LENGTH DELCAP IS THE STEP REDUCTION FACTOR SUGGESTED VALUES ARE .125 OR .25 DELLC IS THE STEP REDUCTION FACTOR SUGGESTED VALUES ARE .125 OR .25 DELLC IS THE TERMINATION CRITERION WHEN THE CURRENT STEP SIZE IS LESS THAN DELLC THE SEARCH IS ENDED. S IS THE NAME OF THE EXTERNAL FUNCTION,S(PHI).TO BE MINIMIZED. A FUNCTION SUBPROGRAM OF THE SAME NAME MUST BE SUPPLIED BY THE USER KCNVRG IS AN INDICATOR TESTED UPON EXIT FROM DIRECT. KCNVRG GREATER THAN ZERO THEN THIS NUMBER IS THE NUMBER OF EVALUATIONS OF THE FUNCTION, MAXEY IS THE MAX. NO. OF EVALUATIONS USER ALLOWS TO FIND THE MINIMUM. KA IS AN INDICATOR USED TO OBTAIN OUTPUT KN=-1 OUTPUT OF FUNCTION VALUE AND VARIABLES IS MADE AT ORIGIN, AFTER EACH EXPLOPE MOVE, AFTER EACH PATTERN MUSER OF EVALUATIONS OF THE FUNCTION. KA IS AN INDICATOR USED TO OBTAIN OUTPUT KN=-1 OUTPUT OF FUNCTION VALUE AND VARIABLES IS MADE AT ORIGIN, AFTER EACH EXPLOPE MOVE, AFTER EACH PATTERN MCVE, AND AT EXIT. KN=C, ND OUTPUT BY OIRECT XN=1, SAME AS FOR -1 EXCEPT EXPLORE MOVES ARE OMMITED. SUBROUTINE DIRECT (X,K,SPSI,DELCAP,RHO,DELLC,S,KONVRG,

SUBROUTINE DIRECT (X,K,SPSI,DELCAP,RHO,DELLC,S,KONVRG, 1MAXEV,KN,BU,BL) CIMENSION (50),PSI(15),PHI(15),SLC(15),X(15),BU(15) 1,BL(15) INTEGER EVAL DO 100 I=1,K

100 PSI(1) = X(1)С IF(K.GT.15) GO TO 50 IF(K) 50,50,4 4 IF(DELCAP) 50,50,5 5 IF(RHD) 50,50,6 6 IF(RHD.GE.1.) GO TO 50 IF(DELLC) 50,50,7 7 MAXEVL = MAXEV IF(MAXEVL) 8,8,9 8 MAXEVL = 500 С 9 DC 60 I=1,K 0 SLC(I) = DELCAP SPSI = S(PSI) EVAL = 1 60 С IF(KN) 61,1,61 61 WRITE (6,63) DELCAP, RHO, DELLC, MAXEVL, KM, (I,I=1,K) 63 FOPMAT (14HDIRECT SEARCH,2X,8HDELCAP=,E15.6,2X,5HRHO 1 =,E15.6,2X,7HDELLC =,E15.6,2X,8HMAXEVL =,I8,2X,5H KN 2=,I3//8HC MOVE ,15H FUNCTION VALUE,3X,3X,I2,6HST VAR, 34X,3X,I2,6HND VAR,4X, 3X,I2,6HRD VAR,4X,3(3X,I2,6HTH 4 VAR,4X)/ 26X,6(3X,12,6HTH VAR,4X)/26X,6(3X,I2,6HTH VA 3R,4X)) WRITE (6,62) SPSI, (PCI(I),I=1,K) 62 FCPMAT(8HCORIGIN,E15.7,3X,6E15.6 /(26X,6E15.6)) С 1 SS = SPSI DC 10 I=1,K 10 PHI(I)= PSI(I) ASSIGN 11 TU IBK GO TO 40 С 11 IF(KN) 12,13,13 12 WRI E (6,14) SS,(PHI(I),I=1,K) 14 FORMAT(8HCEXPLORE,E15.7,3X,6E15.6 /(26X,6E15.6)) С 13 IF(SS.GE.SPSI) GO TO 3 2 IF (EVAL.GE.MAXEVL) GO TO 51 С DO 20 I=1,K IF(SLC(I)) 21,50,22 21 IF(PHI(I).GT.PSI(I)) SLC(GC TO 23 22 IF(PHI(I).LT.PSI(I)) SLC(23 THET = PSI(I) PSI(I) = PHI(I) PHI(I) = 2.*PHI(I) - THET PHI(I)=AMAX1(PHI(I),BU(I)) PHI(I)=AMAX1(PHI(I),BL(I)) 20 CCNTINUE SLC(I) = -SLC(I)SLC(I) = -SLC(I)20 CCNTINUE С SPSI = SS SPHI=S(PHI) SS=SPHI EVAL = EVAL +? ASSIGN 25 TO I&K 40 DC 41 I=1,K THET = PHI(I) SLCI = SLC(I) PHI(I) = THET + SLCI PHI(I)=AMINI(PHI(I),BU(I)) PHI(I)=AMAX1(PHI(I),BL(I)) SPHI = S(PHI) EVAL = EVAL +1 IF(SPHI.LT.SS) GO TO 42 PH:(I) = THET - SLCI PHI(I)=AMAX1(PHI(I),BL(I)) PHI(I)=AMIN1(PHI(I),PU(I)) SPHI=S(PHI) С

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C	42 44 41	EVAL IF(S SLC(SS=C GC T PHI(CONT GD T	= PH1 I)= PH1 0 4 I)= INL 0 1	EV I-G I-S I I I I I I B K	AL E•S LCI ET	+1 S)	GO 25)	T	34	- 4												
c c	25 27 29 28	IF(K WRIT FORM IF(S	N) E(AT (S•(27 6, 8H	28 29) PA SPS	•2 5 TT	7 5,(ERN GD	РН • Е: Т(I(] 15. D]	;;;	I =] 3X (L, K , 6E	15	•6	1	(26	sx,	6E1	5.	6)]	ł	
C	26 3	UU 2 IF(A CONT IF(D)-(• D		(1 C))). 	GT	•0•	5+ 52	•AB	S (:	SL	C ()	())) (50	TO	2	
C	30 50	DO 3 SLC(GO T KCNY	RG	[=] = _	-1	*	SL	C (1)													
	51 52 53 55 56 54	GD T KONV GD T KONV IF(K WRIT FORM I(26X RETU END	RG RG RG N) E(AT (RN	53 53 55 55 115	0 E VA 56) 56) HOK •6)	L • 5 0 N	5 KON VRG	VR =	G.S.	; P S	I. 8H	(PS E)		1)	• I •	=1 E1	K) 5.7	',3)	6 و ۲	E15	•6/	,
CCCCC							FU	NC	TIC)N	FE											
		FUNC DIME 1.MP(1 2).C(3ZZ(6	T I NS 0C 50	CN 1 ON 3) 1 y 2 1 ,	FE(VA • ML 11(PP(X1 15 60	100 0,5 1,Y ,21	;; ;; 12	X() EL (6(), 00 ¥2		100 14P), (2 ,Y	Y(0, 22	60 5) (6)), 2 , EL	(6) TA VA	0), (20	R2(). 501	100 /AL/	liżo
	2	COMM 1VALL 2KI,J NPM= DO 2 VAL	ION ,Y IO,I NPI	VA 21, KO, L-N =1, =X	L C Y22 NAL RES NP)MG ,Z ,K	MIN •ZZ EY2	, P	MG1 P,1 P,1	4A) LIN 4AF	(,Y ,N),J	•R2 DM W •	2,0 JP VV A),E ,J \R,	LT Z, KE	FE KE Y3	LT4 71, , NF	ND ND RES	ALA NP	Y) L,I	L1 • Y VN • J	/12, II,
	46	CALL IMAP, CALL IMAP, IMAP, IF(1) IF(2) JP=N JZ=N	EL EL (1) (1) (1) (1) (1) (1) (1) (1) (1) (1)	UPU TAPU J=16 J=16 J=16 J=16 J=16 J=16 J=16 J=16	L (N VAL VAL VAL VAL VAL	IPL A, IPL A, 6,6	¥11 •NA ¥21	L , , Y , Y	NN 12 NN 22		,K LL ,K	, 1 , Ji , Ji	JO , JO , J , N	KO IVA KO IVA	,3 R,2 R,		EY2 Y31 EY2 Y31	2 , M Y ; ! 2 , M 2 , I	P,E P,E NZ,		VAL VAL 2)	- 7
	8 10 12 14 16	GU 1 DO 1 Y(K) DO 1 Z(K) CONT CONT DO 2 JJ=J	0 =Y 2 =Z IN IN 0 P-	10=1+ (K#1+ UE=1 10=1+	, NF 1) 1) , NZ)																

IF (Y (JJ))18,20,18 19 JP=JP-J+1 GO TO 22 20 CCNTINUE 22 DO 26 J=1,60 JJ=JZ-J+1 IF (Z'JJ)24,26,24 24 JZ=J/-J+1 GC TJ 28 26 CONTINUE C CALCULATE ZEROS 28 CALL MULLER(Y,JP,ZZ) C CALCULATE POLES CALL MULLER(Z,JZ,PP) C CALCULATE FREQUENCY RESPONSE CALL FREQQ(LIN,NOM,OMGMIN,OMGMAX,JP,JZ,Y,Z,KEY1,R2) FE=0. DG 1 I=1,ND FE=FE+(RZ(I)-D(I))**2 RETURN END

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THE MODIFIED VERSION OF CALAHAN The Subroutines are Listed Alphabetically for Convenience	SUBROUTINE ASBS (LN,C,G,G,H,X,NP,MG,I,KEY1,JN,JP,II,NPL,NAL) IF (I-2)101 c (50);G(50),H(50),X(60),2(60),MG(30,5),Y(3,50) IF (I-2)101 100,101 GG TO(2:3:3,4),KEV1 JM=1	00=2 60 TO 5 VP=2#LN-3 JM=2	JP=3 60 T0 5 NP=2#LN+1 JM=0	JP=1 JN=LN-JM N=M6(JP3) re(ND-3)30.30.102	F(NP)31,31,32°C X(1)=0. NP=1	RETURN X(1)=0. VD=1	RETURN DG 130 J=1.NP X(J)*0. Z(1)=1.	K1#1 K2#1 DO 1 J#1,J K#UM+J N#M6(K,3)	00 116 K*1 3 00 117 1=K1,K2r Y(K,I)=0.
	m	4	N	0					
00000000000	100			يب بب	30	31	102 130		117

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CONTINUE IF(C(N))81,90,81 N1=2 N2=2 N2=2 N2=2	Y(3, K+2)=2(K)*C(N) IF(G(N))86, 87, 86 IF(G(N))82, 83, 82 N2#1	DO 111 K=K1,K2 Y(2,K+1)=Z(K)*G(N) F(H(N))85,34,85 IF(H(N))88,89,88 N2=0 N1=0	DO 112 K=K1.K2 Y(1.K)=Z(K)+H(N) X1=K1+N1 X2=K7+N2	DC 113 K=K1,K2 Z(K)=Y(1,K)+Y(2,K)+Y(3,K) CONTINUE	IF (11) 1 5 5, 1 33, 1 3 5 00 131 J=K1, K2 X(J)=-Z(J) DET 1000	DD 134 J=K1.K2 X(J)=Z(J) RETURN END	SUBROUTINE DPOLRT (XCUF,COF,M,ROOTR,ROOTI,IER) PURPOSE COMPUTES THE REAL AND COMPLEX RNOTS OF A REAL POLYNGMIAL USAGE USAGE USAGE CALL POLRT(XCOF,COF,M,ROOTR,RNOTI,IER) DESCPIPTION OF PARAMETERS DESCPIPTION OF PARAMETERS DESCPIPTION OF PARAMETERS COF -WORKING VECTOR OF LENGTH M+1 COF -WORKING VECTOR OF LENGTH M+1 COF -WORKING VECTOR OF LENGTH M+1
116 81	110 80 82	883341 883341 883341	112844	113	132 131	110 10 10 10 10 10 10 10 10 10 10 10 10	000000000000000000000000000000000000000

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RESULTS THE DCUBLE PRECISION VERSION OF THIS SUBROUTINE MUST ALSO CONTAIN DOUBLE PRECISION FORTRAN FUNCTIONS. ABS IN STATEMENTS 78 AND 122 MUST BE CHANGED TO DABS. • 500 INTERATIONS ш METHOD NEWTON-RAPHSON ITERATIVE TECHNIQUE. THE FINAL ITERATIONS ON EACH ROUT ARE PERFORMED USING THE ORIGINAL POLYNOMIAL RATHER THAN THE REDUCED POLYNOMIAL TO AVOID ACCUMULATED ERRORS IN THE REDUCED POLYNOMIAL. TH DOUBLE PRECISION XCOF,COF,RODTR,RODII,X0,Y0,X,Y,XPR,YPR,UX,UV,V, Y1,X1,U,XT2,YT2,SUMSQ,DX,DY,TEMP,ALPHA **R00TS** REMOVED FROM DOUBLE PRECISION STATEMENTS ROUTINES USED IN CONJUNCTION WITH THIS PRECISION R ROOTR-RESULTANT VECTOR OF LENGTH M CONTAINING REAL R ROOTI-RESULTANT VECTOR OF LENGTH M CONTAINING REAL R CORRESPONDING IMAGINARY ROOTS OF THE POLYNOMIA IER -ERROR CODE WHERE IER -ERROR CODE WHERE IER -ERROR OF MAN ONE IER - M GREATER THAN 36 IER - W ABLE TO DE TERMINE ROOT WITH 500 INTER IER - W ABLE TO DE TERMINE ROOT WITH 500 INTER IER - HIGH ORDER COEFFICIENT IS ZERD REMARKS LIMITED TO 36TH ORDER POLYNOMIAL OK LESS. FLOATING POINT OVERFLOW MAY OCCUR FOR HIGH ORDER POLYNOMIALS BUT WILL NOT AFFECT THE ACCURACY OF THE •••• LE PRECISION VERSION OF THIS ROUTINE IS MN I SHOULD BE REMOVED FROM THE DOUBLE F WHICH FOLLOWS. 4 DIMENS.JN XCOF(1), COF(1), RODTR(1), RODTI(1) ALSO BE • CIN COLUMN STATEMENT WH THE C MUST A APPEARING IN ROUTINE.

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INCREMENT INITIAL VALUES AND COUNTER SET X AND Y TO CURRENT VALUE ALRO INITIAL VALUE COUNTER IFIT=0
N=M
IER=0
IF(xCOF(N+1))10,25,10
IF(N) I5,15,32 SET ERRUR CODE TO 1 4 SET ERROR CODE TO 2 SET ERROR CONE TO SET INITIAL VALUES D IER=2 GG TO 20 S IF(N-36) 35,35,30 NXX=N+1 NXX=N+1 NXX=1 +1 KJ1= N+1 KJ1= N+1 KJ1= N+1 COF(MT)=KGP(L) 45 X0=. C0500101 Y0=C.01000101 X0=-10.0+Y0 YC=-10.0+X IER=4 60 T0 20 15 IER=1 20 RETURN SQ X#X0 10 25 2 Nin Nin 30 **4**0 ပပပ ပပပ ပပပ

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TEMP=XCUF(MI) CCCF(L)=TEMP CCF(L)=TEMP CCF(L)=TEMP N=KP=N N=N N=N N=N N=N N=N N=N N=N	Y#YPK IFIT=0 ALPHA=X+X SUMSG=X#X+Y#Y N=N-2	GU TO 140 XX=C*O XX=XX+1 XX=XX+1 XX=XX+1 XX=XX+1 XX=0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0	ALTHATA N=N-1 CCF(2)=COF(2)+ALPHA*COF(1; CCF(2)=COF(2)+ALPHA*COF(1; CCF(1)]=COF(L+1)+ALPHA*COF(L)-SUMSQ*COF(L~1) COT1(N2)=Y ROOTR(N2)=X N2=N2+1 N2=N2+1 1F(SUMSQ) 160,165,160	Y=-Y SUMSQ=0.0 GO TO 155 IF(N) 20,20,45 END	SUBROUTINE FREQ (NA.NB.A.B.G.W) DIMENSION A(60), AE(30), AD(30), AE(30), AD(30), B(30) CALL PARTS (NA.A.M1.AE.N1.AO.M11.AE1.N11.AO1), B(60 CALL PARTS (NB.B.M2.BE,N2.800.ML2.BE1.N12.BO1) CALL PARTS (NB.B.M2.BE,N2.800.ML2.BE1.N12.BO1) CONNEW*SUM (N11AC.W) CODNEW*SUM (N11AO.W) CODINESUM (N11AO.W) CODINESUM (N11AO.W)
105	1200 1200 1200	130 135	ユユユ つちつち	160 165	33

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9(1 IN' NOM, DMGMIN, DMGMAXINA NB'A, B, KEYI R2) , B(60), G(4), RI(100), R2(100), R3(100), R4(100), R(60) 010+0000*EV10)/80770M-(EVN*0001N+000N*EV1N)/70P ABLE FOR ./.20X, W = .1PEIS.6, RESPONSE DATA P+BOTTOM) - 1.0E+65) 1,1,2 28318 GMAX-DMGMIN//SL+DMGMIN =4, 3429448#ALDG(G(2)) =6(4) =6(3) (X<u>/</u>X) KEQ (NA, NB, A, B, G, WW) 0000+ ĂXPTS) 86.87.87 899189 X-W187,87,1 # 00 **X** H = (E V d W PORMA PORMA I U U 330×888 1341000 81111222 Ш SUBI u 4 ũũ -UNZ 3 Ъ. ΰ 11209 22 87 80 N ပပ

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Bart Month Propagation and a

• "X, IOHPHASE(DEG))// ,3X,10HPHASE(DEG))//

 WRITE (6:301)
 CUMIINUATION OF ABOVE TABLE AND GRAPH*)

 I2X, INDELAY (SEC)
 XX: VAGAIN (DB)
 XX. IOHPHASE(DEG)

 I2X, INDELAY (SEC)
 XX: V)
 4X, 9HGAIN (DB)
 XX. IOHPHASE(DEG)

 WRITE(6)
 BI
 XX: V)
 4X, 9HGAIN (DB)
 XX. IOHPHASE(DEG)

 I2X, INDELAY (SEC)
 XX: V)
 4X, 9HGAIN (DB)
 XX. IOHPHASE(DEG)

 WRITE(6)
 BI
 XX: V)
 4X, 9HGAIN (DB)
 XX. IOHPHASE(DEG)

 WRITE(6)
 BI
 XX: V)
 4X, 9HGAIN (DB)
 XX. IOHPHASE(DEG)

 WRITE(5)
 BI
 XX: V)
 XX: V)
 XX. IOHPHASE(DEG)

 PHASE (DEG) . DELAY (SEC)" PHASE (DEG) . OHMS . . / . 32 X . GAIN (DB) .. 141,5%, "CONTINUATION OF ABOVE TABLE AND GRAPH") 141,5%, "CONTINUATION OF ABOVE TABLE AND GRAPH") 233010/ 1000000 /031 //2010 /001 /2010 /0000 FORMAT (2(F20.7,F15.7,F10.3)) 60 10 33 WRITE (619) (R1(J),R2(J),P3(J),R4(J),J=1,KL) FORMAT (2(F20.7,F15.7,F10.3,F15.7)) Y-AXIS Y-AXIS Y-AXIS Y-AXIS Y-AXIS SOHMH9 . X4. .4X,9HOHMS .33) (R1(J).82(J).83(J).J=1.KL) .60.1) CALL PLOT(R1.R2,KL) .200) //.iox..plot of above table: 2x.ix-axis freq (H2).) 2(f20.7,f15.7,f10.3)) // IOX.PLOT OF ABOVE TABLE: Axis freq (HZ)) Eq. 1) call plot(R1,R3,Kl) 6,202) ABLES ABLES ABLES 6.300) 6.333) 27.2(9x,10HFREQ (HZ) 6,300) 6,302) //.2(9x,10Hfreq (HZ) , , F15, , 612 (81, R2, KL) , 612 (81, R2 313) KEYL (6,205) KRITE RELATION RELATI 204 FO' MAT FORMAT 205 FORMAT 203 FOPEAN CHOHOHO LANCHO LINCHOLIC 200 FOR 1189 5 113 300 202 60 24 128 302 114 303 301 93 æ നപ ပပ

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ELETS IHC / 4HC / IHR/4HR / IHG /4HG / List of elements and node numbers in both current and voltage 7,10X,4HTHE ,12,50HTH ELEMENT VALUE HAS BEEN REPLACED BY P. MAP, ELT, ELTA, VAL, VALA, NPL, NAL, NE)) ELT 100) • VALA(20) • Č(50) • G(50) • AS PRECEDING ELEMENTS, PLACE IN PERMANENT LIST OF MPL+NAL E PARALLEL RLC ELEMENTS INTO SINGLE TOPOLOGICAL ELEMENT 120 K=1,NLL ELEMENTS (BY NODE NUMBER) [2] [2] GO TO 197 FREQ (HZ) 1 (5005) 1200) F (OMGMAX-W) 94, 94, 95 7.32X, 'X-AXIS NEN1 ž 불물 DOUIA 120 N 196 **46** 197 40 300 100 ပပ ں G

ČÁS PRECEDING ELEMENT,ADD VALUE TO THAT OF PRECEDING)*K CTIVE ELEMENT NODE NUMBERS TO PERMANENT LIST ELEMENT AND STORE VALUE Rigo to 106 FUNCTION IGN(NN,LN,MG) DIMENSION MG(30,5),ME(20,2),MM(100) K1=1 F (ELT(J) .NE. IHR) GD T0 116 #1./VAL(J)+6(K) 0 109 60 70 11 JI ... NE. IHC) GD TO 108 LT(J) .NE. IHG) LA(J)+G(NM) 1. /V AL (J) +H(JK) VAL (J)+H(K) (WN)H+17) V] = MAP (] . 1) = MAP (] . 2) = WAP (J , B) VAL - KK 5 C E E E E DCCCCCC DLLLLA XXXXXX 115 108 10 186 106 115 100 118 121 105 111 111

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TWC REAL+8 ARRAYS OF SIZE (N+1) MUST BE FURNISHED FOR THE COEFFICIENTS CR (REAL PART) AND CI (IMAG PART) ••• ••• SUBROUTINE MAKPOL (N, ROOTR, ROOTI, CR, CI) RFAL*8 ROOTR(1), ROGTI(1), CR(1), CI(1) R C C0015 C0015 C0055 C0655 20CC RRRR PARTS PARTS PARTS PARTS 1 1 + UMBER OF ROOTS ARRAY OF REAL ARRAY OF OF REAL AGC ACTR(I) RETURN <u>8</u>5 00 0018 ***** ***** ***** ***** DABS (R001 DABS (R001 DABS (C 1 1 I=1,K **~~**** 22 N Ī CR 001 Z N. α วซิพี

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COEFF IN ASCENDING ORDER ORDER OF POLNOMIAL PLUS ONE * REAL ROOT • CORRESPONDING IMAG PART OF ROOT DIMENSION ZRO(60), COE(60), Z(60, 2) SUBROUTINE MULLER(ZRO,N8,Z) ¿ZÓX,4HNONE 1000,1001 74-N1)19,37,19 1NUE 0.8 COE(1) 9,7,9 J=1,NB = ŽRO (K) i o<u>oo</u> , 00 ZRO NB Z(1.1) Z(1.2) R=AXE I=AX R=1 C Ħ Ŧ 60709 BET2R BET2R ۵ A A A 200 C1000 11000 1001 60 51 12 17 1003

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667 10 99 Akriittfer Akriittfer Akriittfer Akriittfer 60 10 99 60 10 99 Akritter Artifer Artifer Altif	M=3 TEMR=COE(1) DC 100 0 TE1=TEMR#AXR-TEMI#AXI TE1=TEMR#AXR-TEMR#AXI TEMR= TEMR= TEMR= TEMR= TEMR= TEMR= TEMR=	TENTA 102,103,102 TEMI=AXR-2(1,1) TEMI=AXR-2(1,1) TEITEMITEMI+TEM2+TEM2 TEITEMIEMEFEM1+TEM2+TEM2 TEIETEMEFEM1+TEM2) TE28 TE28(TEMEFEM1+TEM1+TEM2) /TE10	TEMI#(TFMI#TEMI~TEMR#TEM2)/TEIA TEMR#TE2 Co To(11,12,13,15,33,34),M Return END	SUBROUTINE OPT(NML, NEL, KEYL) DIMENSION NML(50,5) NS=1 S=1 S=2 NS=2 NS=3 S=2 NS=3 S=2 NS=3 S=2 NS=3 S=2 NS=3 S=2 NS=3 S=2 NS=3 S=2 S=2 S=2 S=2 S=2 S=2 S=2 S=2 S=2 S=2
9 9 9	99 100	102	101 00 031 01 00	4° 10' 40

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X 2 2 2 2 2 4 2 4 2 4 2 4 2 4 2 4 2 4 2 4 2 4 2 4 2 4 2 4 2 4 2 4 2 2 4 2 2 2 2 2 2 2 2 2 2 2 2 2	DD 109 J=NS,NEL 1F(J-NS) 108,109,108	KK=J-I DO I21 K=NS,KK IF(NML(J,K1)-NML(K,K1))1121,103,1121	IF(NML(J,K)) - NML(K,K2) (121,122,121) IF(NML(J,K2)-NML(K,K1))121,104,121 DO 100 JJ#1,5 NIFNML(J,JJ)	LL=J-K-1 IF(LL)100,100,107 D0 101 L=1,LL	Ll=J-L+I NML (Ll,JJ)=NML (Ll-1,JJ) NML (K+1,JJ)=N1		IF(K1-1)2,1,2 Return End	SURROUTINE PARTS (NA,A,MK,AE,NK,AO,MIK,AE1,NIK,AO1) DIMENSION A(60),AE(30),AO(30,AE1(30),AOI(30) Iti MK=1	NK=0 M1K=0 AF(1)=0(1)	ÎF(ÑÁ~])3,3,1 I=I+1 NK=NK+1	AO(NK)=A(I) DUMMY=I-I ACI(NX)=DIMMY=A(I)	TF(N N + 1 - 10)FE	[+]=]
		108		107	101	121				-		2	
", IOX, GRAPHS' DF LESS THAN FOUR POINTS ARE NOT PLOTTED), Y (NN), RANGE (4) ANGE (1), XMAX), (RANGE (2), XMIN), (RANGE (3), YMAX), E (4), YMIN) GO TO 200 GUAGE INPUT & FIND MAX & MIN FOR X & Y; CALL UTPLOT X - THE X-AXIS COORDINATE Y - THE Y-AXIS COORDINATE NN- THE NUMBER OF POINTS TO BE PLOTTED TPLDT (X, Y, NN, RANGE, 1) (6, 101) YMAX, XYMAX, YMIN, XYMIN SUBROUTINE PLOT(X, Y, NN) NN Max1 6,6,2 HIN) 3,3,7 -YMAX1 8,8,4 -YHIN) 5,5,1 DUMMY*A(I) • 3• 1 DIMENSION X(NN), EQUIVALENCE (RAN fourvalence (Range) IF (NN .GE 4) WRITE (6,201) FORMAT (10()).1(TF (N WRITEN FORMAT (1001) RETURN RETURN KMANAT (110) KMINET (100)
 YMAX=1
 YCO

 YMAX=1
 FECO

 YMAN
 FECO

 DO
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NTS NTS /•10X | Max #=•1PE12•4•• AT X=•E12•4•20X••MIN Y=• E12•4•• AT X=• • # > X.X="IPE12.4," AT Y="E12.4,20X,"MIN X=" E12.4," AT (37), RODTR(36), RODTI(36), DZR(36), DZI(36), NO ROOT CHECK. . . //) FORM TWO NEW POLYNOMIALS FROM THE TWO SETS OF RODTS CALCULATE THE ERROR CRITERIA The set of roots to be printed out and use when expanded vield the most nearly (compared (abs values) to the original CÓF, COF, M, ROOTR, ROOTI, IER)) 60 TO 20 100) XMAX, YXMAX, XMIN, YXMIN 10X. OLRT (XCOF,COF,M,ROOTR,ROOTI,I) • EQ. 0) GO TO 20 6,200) (//, 0 DPOLRT UNABLE TO FACTOR. , COF ERR1 ++05 (M, DZR, CZI, COFI (M, ROOTR, ROOTI, 2(60,2) 0011 (1) 1100 (1) 1100 =1 , NB MAKPOL ** EETURN END FORMAT (/ WRITE C FORMAT RETURN D0 21 D2R(1) D21(1) CALL P CALL P ACC ACC SUBF ERR] 22 100 200 101 0 0000 2 21 ပပ 000000

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FFML(J.5)-JUPI117.116.117 FFML(J.5)-FKD1115.119.115 FFML(J.5)-FKD1115.120.115 INSERT LAST ELEMENT IN PLACE OF REMOVED ELEMENT DOD 105 KT SKI 5 DOD 105 KT SKI 5 DOT 105 KT SKI 5	CONTINUE GO TOLIJS, 5,6) KEVI INSERT FEST ELEMENT ACROSS INPUT IN CURRENT GRAPH,ACROSS DUTPUT IN VOLTAGE GRAPH ML(1.3) #98 ML(1.5) #40 ML(1.5) #40 ML(1.5) #40 ML(1.5) #40 ML(1.5) #40	MILICZERT MILICZ		FUNCTION SUM #0.0 SUM #0.0 SUM #0.0 SUM #0.0 SUM #0.0 SUM #0.0 SUM #0.0 SUM #1.0 SUM #2.0 SUM
	·UU~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~	0 0	ح	********
J	60	с С	υu	(VE VT - ¥C)

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TALL PART

5'TO'ADÓ ONE NEW NODE AT A TIME UNTIL TED FOR (SEE TRETST) N. KYOUT . K 10) X-LN)7,10,10 VINUE K-X,19,8,9 05 21 J*1, NEL 40(J)=0 811=1N 05 30, J=410, NEL + 2 2 2 2 3 X##6(1, X1) X##6(1, X1) MM(B)#1 XM(X)#1 XM(X)#1 XM(X)#1 XM(X)#1 E NSION S TEST NODES хх) ххун С 410 11 NX B RE TURN END 07141 01220 പ്പം (XX 110 22 2 32 30 3

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فتمكوهم فالاستعمام والمحافظة والمحافظة والمنافعة المحافظة والمحافظة والمحافظة والمحافظ المراجع ومعالاتها والمسارك والمنا

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.1X,11HT0P0 4H+ / Lready been calculated in Bilinear Form, ion 5015) SUCEST 100 MAP (20.5), ELTA 20 % 464 4 20 LE TOPOLOGICAL ELEMENT L'VALA, NPL, NAL, NE) I3.103.104) KEY2 IB€R,CÖRRESPONOING TOPOLOGICAL ELEMENT NUMBER NAL, NN, JI, KI, JU, KU, KEYI, KEYI, KEY2, MP, ELT, VAL, VAL, KW, NVAR, KEY3, Y, NP, JP) . NML, J1, KI, JD, KD) IOS) (J,MP(J,3),J=1,NPL) ,IX,19HELEMENT ASSUCIATION ,/,IX,8HPHYSICAL X,25HTREES,2-TREES,OR 3-TREES [J.MAP(J.5), J.1, NAL) IS ACROSS INPUT NVVI 133,134,133 5 J=1 200 *VALL (KW)*Y1 (J)/VALL (KW-L) 48 16 J=1 NP =VALL (KW-1)*Y1(J)/VALL(KW) -48 1X,1HY, 12) 81,106 (S · N ž R-NP P NO # 106 è 105 82 103 40 MM 200 31 130 80

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BEEN FORMED, REMOVE FATAL ELEMENT FROM LIST LE TREE-ADMITTANCE PRODUCT SBS([n,C,G,H,X,NP,MG,[,KEVI,JN,JP,I[,NPL,NAL) SET OF ELEMENTS TO BE TESTED AS TREE TOTIOD Network is active,test for sign of tree 0_50 Etwork IS Active.test voltage graph All1233,235,233 URRENT GRAPH KVIN, KY OUT, LL, MG) 9,19 870,871,671 YINTO ITCN IN, LN. KG 0 899 K=NEL+2-NN N 200 0 12] 0# 2 -. ž 20 50 233 141 22 <u>الم</u> 274 800 850 070 671 672 872 うううう こでのの ろうろし

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ATED IN BILÌNEAR FORM,TEST FUR VARIABLE ELEMENT 146 TEST ELEMENTS HAVE BEEN PERMUTED ((HV, I2, 2X))) PRODUCI [NVAR,3]]147,44,147 K 23)-MAP(NV,5))148,44,148 EXAMINE THE ORIGINAL IN THIS AREA K#NG(L) D0 77 J=1,5 MG(L,J)=ML(K,J) G0 T0 100 Permute Next Element IN LIST Sol Sul 48,48 50,50 . 201 wa HALKK . NY αÑ + × 12 1 80 90 51 9 52²⁶⁷ * ^ 68 00 ئ 50 265 5 428 347 147 148 44 4440 440 44 400 4 410 46 5 1 0000 J

427,48 100,100 12,NELTER FOR CONNECTED GRAPH 953,953 ASSEMBLE AND SHIFT COEFFICIENTS 1,555 1,555	TST NN, LN, K VIN, KYOUT, LL, MG) AUE 64 TRYING TO ADD ONE NEW NODE AT A TIME UNTIL BEEN ACCOUNTED FOR ONE NEW NODE AT A TIME UNTIL 5
$\begin{array}{c} \label{eq:constraint} \\ \label{eq:constraint} \\$	ALLE SUBRAUTIN ALLE SUBRAUTIN
10 10 10 10 10 10 10 10 10 10	

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IS PRESENT.	щ
DO 7 J=2. N=MGG(J.*KI) N=MGG(J.*KI) N=MGG(J.*KI) N=MGG(J.*KI) N=MGG(J.*KI) N=MGG(N=NEW N=NCONFERST PASS.CIRCUT IFF(NMMN)) IFF(NMN)) IFF(NMN	SUBROUTINE UTPLOT (X, YSCALE(5); YSCALE(7)) DIMENSION CRID(61,101); XSCALE(6); YSCALE(7) DIMENSION X (1), YSCALE(6); YSCALE(7) INTEGER*2 ZERO/IHOV, IYORG(101) INTEGER*2 CGRID, BLANK, YCHAR, DOT INTEGER*2 CGRID, BLANK, YCHAR, DOT CRID IS THE MATRIX USED TO PLOT THE POINTS CRID IS THE MATRIX USED TO PLOT THE POINTS OD & I=1, 101 NG FR =0 XMAX=RANGE(1) XMAX=RANGE(2) YMAX=RANGA YMAX=RANGA YMAX=RANGA YMAX=RANGA YMAX=RANGA Y
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PLACING PDINTS IN THEIR PROPER GRID POSITIONS AXIS , IF NECESSARY XRANGE+1.5 DO 70 I=1 NDATA KKZ IPTX=60.#1 YMAX-7(I)/YRANGE+1.5 .KK2 205,205,220 - XMIN) 203, 210, 210 ŽMAX1215,215,212 30,30 XAXIS=60.**MAX/YRANGE+1.5 0 60 1*1.101 RID(IXAXIS,1)=DOT BLANKING DUT MATRIX-(GRID) ŻERO **Ă**XĪS)=00T AND Y PLOTTING X AND Y XRANGE=XMAX-XMIN YRANGE=YMAX-YMIN 101 = 101 ERR+1 **CONTINUE** ERR-1 301 0 220 205 30 212 215 **\$** 333 60 444 222 J ပ ပပ ပပပ ပပပ

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[PE10.3,5(10X,E10.3)/15X,2H**,10(10H+********),3H+** N WITH GRAPH IYORG XSCALE(6),XSCALE(4),XSCALE(3),XSCALE(2), xSCALE(6), xSCALE(5), XSCALE(4), XSCALE(3), XSCALE(2), 01 [K=1,61 [91,91,92 E [6,18] YSCALE(II),(GRID(IK,IX),IX=1,101),YSCALE(II) MT[3X,1PE10.3,4H + ,101A1,4H + ,E10.3)) 102 IK •NE• IXAXIS) GO TO 192 IK •NE• IXAXIS) GO TO 192 E •6.400) (GRID(IK,IX)•IX=1,101) AT (ДХ,440.00.3X,1H±,IX,101A1,2H *,3X,440.00) AT (ДХ,440.00.3X,1H±,IX,101A1,2H *,3X,440.00) 14) . 10(10H+********),3H+**/, 1P 1X,610.31) 0RG N BER OF POINTS OUT OF RANGE TG15,19, (CRID(IK,IX',IX=1,101) TG15X,1H*,1X,101A1,1X,1H*) TO GRID(IPTX, IPTV) = XCHAR CCMPUTE PROPER SCALE NUMBERS SCALE(I-1)-YINCR W XSCALE(I-I)-XINCR 500,1001 -10/101,103,103 ž X INCR-XRAN び し し て お ち じ し く 1 Ē 22 F DR 1 92 1 92 1 92 161 101 1001 2001 8000 80 81 400 401 ~ 6 ບບບ υu

LIST OF REFERENCES

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ABSTRACT					
The filter design problem is const	idered as an o	optimizatio	on problem. An		
iterative search technique is employed	d to adjust th	e variable	e network ele-		
ment values to approximate some desire	ed network res	ponse, wit	ch a minimum of		
error. Explicit constraints are emplo	oyed to ensure	physical	realizability.		
The design process uses a combination	of a modified	version o	of Calahan's		
network analysis program with a direct	t search metho	d of minim	nization developed		
by Hooke and Jeeves. The result is a	procedure whi	ch utilize	& the circuit		
designer's experience and knowledge to	o set up the p	roblem but	relieves him		
of the tedious labor now performed by	the high-ence	d dioital	computer		
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