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## ELEMENTARY PROOF OF THE WIELANDT-HOFFMAN THEOREM

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#### J. H. Wilkinson

<u>Abstract</u>: An elementary proof is given of the Wielandt-Hoffman Theorem for normal matrices and of a generalization of this theorem. The proof makes no direct appeal to results from linearprogramming theory.

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#### 1. Introduction

In [2] Wielandt and Hoffman proved a theorem on the eigenvalues of normal matrices which is of considerable importance in the error analysis of eigenvalue algorithms based on the use of unitary transformations [4,5]. Their proof was very elegant and was based on the use of linear programming techniques. In [5] Wilkinson gave an elementary proof in the case when the matrices are Hermitian, which was based on an earlier proof due to Givens [1]. This proof did not extend easily to the general case. Here we give an elementary proof for the general case which applies immediately to a generalization of the Wielandt-Hoffman theorem due to Kahan [3]. Not surprisingly the proof involves techniques which are familiar in the area of linear programming but no direct appeal is made to results from that field.

#### 2. The Basic Theorem

The proof depends on a theorem which is not directly concerned with normal matrices. Before stating this theorem we give two definitions.

<u>DEFINITION 1</u>. The set of n elements  $a_{1,i_1}, a_{2,i_2}, \dots, a_{n,i_n}$  of an nxn matrix A is called a <u>diagonal</u> of A if  $i_1, i_2, \dots, i_n$  is a permutation of the integers  $1, 2, \dots, n$ . If  $i_j = j$   $(j = 1, \dots, n)$ then we have the <u>principal</u> diagonal.

<u>DEFINITION 2</u>. A matrix X is called a <u>doubly stochastic matrix</u> if  $x_{ij} \ge 0$  and  $\sum_{i=1}^{n} x_{ij} = \sum_{i=1}^{n} x_{ji} = 1$  (j = 1,...,n) i.e., all row and column runs are unity.

<u>THEOREM 1</u>. If P is a real matrix such that the sum of the elements on the principal diagonal is not greater than the sum of the elements on any other diagonal, and X is any doubly stochastic matrix, then  $S(X) \equiv \sum \sum p_{ij} x_{ij}$  is a minimum when  $\lambda \in I$ .

<u>Proof</u>. The minimum is attained, possibly for many different X. Let us choose X to be a minimizing doubly stochastic matrix having the maximum number of zero off-diagonal elements. We shall show that all its off-diagonals must be zero. For suppose that this is not true. Let  $x_{i_1,i_2}$  be a non-zero off-diagonal. Then  $x_{i_2,i_2} < 1$  and hence there is a non-zero element  $x_{i_2,i_3}$  (say) in row  $i_2$ . If  $i_3 \neq i_2$ then similarly there is a non-zero element  $x_{i_3,i_4}$  in row  $i_3$ . Continue in this way until we reach an  $x_{i_{m-1},i_m}$  for which  $i_m$  equals some earlier  $i_k$ . Let x be the smallest of the <u>positive</u> elements

 $x_{i_k,i_{k+1},x_{i_{k+1},i_{k+2}},\dots,x_{i_{m-1},i_k}}$ 

Construct a matrix Y such that

$$y_{i_{s},i_{s}} = x_{i_{s},i_{s}} + x$$
,  $s = k,k+1,...,m-1$  (2.1)

$$y_{i_{s},i_{s+1}} = x_{i_{s},i_{s+1}} - x$$
,  $s = k,k+1,...,m-1$  (2.2)

$$y_{i,j} = x_{i,j}$$
 otherwise. (2.3)

Then Y is clearly a doubly stochastic matrix and

$$\sum \sum p_{ij} y_{ij} - \sum \sum p_{ij} x_{ij} = x \left[ \sum_{s=k}^{m-1} p_{i_s, i_s} - \sum_{s=k}^{m-1} p_{i_s, i_{s+1}} \right] . \quad (2.4)$$

The expression in brackets cannot be positive since otherwise by replacing the elements  $p_{i_s,i_s}$  in the principal diagonal by the elements  $p_{i_s,i_{s+1}}$  we could obtain a smaller diagonal sum. Hence

$$\sum \sum \mathbf{p}_{i,j} \mathbf{y}_{i,j} \leq \sum \sum \mathbf{p}_{i,j} \mathbf{x}_{i,j}$$

But Y is clearly a doubly stochastic matrix and it has at least one more off-diagonal zero than X, contradicting the hypothesis. Hence all off-diagonal elements of X must be zero, i.e., X = I.

An exactly analogous theorem holds when the principal diagonal has the maximum sum.

#### 3. The Wielandt-Hoffman Theorem

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<u>THEOREM 2</u>. If A and B are normal matrices and C = A - B, and if  $\mathbf{a}_i$  and  $\mathbf{b}_i$  are the eigenvalues of A and B arranged so that  $\sum_{i=1}^{n} |\mathbf{a}_i - \mathbf{b}_i|^2$  is a minimum for all possible orderings, then

$$\sum_{i=1}^{n} |\mathbf{a}_{i} - \mathbf{b}_{i}|^{2} \leq ||\mathbf{C}||_{F}^{2} \cdot (||\mathbf{C}||_{F} = \text{the Frobenius norm of C}) \quad (3.1)$$

<u>Proof.</u> Since A and B are normal there exist unitary  ${\tt Q}_1$  and  ${\tt Q}_2$  such that

$$A = Q_1 \operatorname{diag}(a_1)Q_1^H$$
,  $B = Q_2 \operatorname{diag}(b_1)Q_2^H$ . (3.2)

(Note then we are free to prescribe the ordering of the  $a_i$  and  $b_i$ and we choose the ordering which gives  $\sum |a_i - b_i|^2$  a minimum value. Hence

$$A - B = Q_1 \operatorname{diag}(a_1)Q_1^H - Q_2 \operatorname{diag}(b_1)Q_2^H = C$$
 (3.3)

giving

$$\operatorname{diag}(a_{1})Q_{1}^{H}Q_{2} - Q_{1}^{H}Q_{2} \operatorname{diag}(b_{1}) = Q_{1}^{H}CQ_{2} \quad . \tag{3.4}$$

Writing  $Q = Q_1^H Q_2$ , a unitary matrix, we have

$$\|\operatorname{diag}(a_{i}) Q - Q \operatorname{diag}(b_{i})\|_{F}^{2} = \|C\|_{F}^{2}$$
 (3.5)

since the Frobenius norm is unitarily invariant. Hence

$$\sum \sum |a_{i} - b_{j}|^{2} |q_{ij}|^{2} = ||C||_{F}^{2} . \qquad (3.6)$$

Now the matrix P with  $p_{ij} = |a_i - b_j|^2$  is real and from the ordering of the  $a_i$  and  $b_i$  its principal diagonal is minimal. Further, since Q is unitary, the matrix Z with  $z_{ij} = |q_{ij}|^2$  is a doubly stochastic matrix. Hence by Theorem 1 and equation (3.6)

$$\sum_{i=1}^{n} |\mathbf{a}_{i} - \mathbf{b}_{i}|^{2} \leq \sum \sum |\mathbf{a}_{i} - \mathbf{b}_{j}|^{2} |\mathbf{q}_{ij}|^{2} = \|\mathbf{c}\|_{F}^{2}$$
(3.7)

and the result is proved.

When A and B are Hermitian, the  $a_i$  and  $b_i$  are real, and it is easy to prove that the orderings  $a_1 \ge a_2 \ge \cdots \ge a_n$ ,  $b_1 \ge b_2 \ge \cdots \ge b_n$ give the minimal value. In fact, returning to Theorem 1 in the case when  $p_{ij} = (a_i - b_j)^2$  with  $a_i$  and  $b_i$  real and monotonically ordered, the proof is much simpler. For if X has a non-zero off diagonal element

in row 1 or column 1 it must have at least one such in both. Suppose  $x_{lr}$  and  $x_{sr}$  are non-zero and x is the smaller. If we increase  $x_{ll}$  and  $x_{sr}$  by x and diminish  $x_{lr}$  and  $x_{sl}$  by x the sum is changed by

$$x[(a_{1} - b_{1})^{2} + (a_{s} - b_{r})^{2} - (a_{1} - b_{r})^{2} - (a_{s} - b_{1})]^{2} = x(a_{1} - a_{s})(b_{r} - b_{1}) \leq 0 \quad (3.3)$$

Hence continuing in this way the minimizing X has no non-zero off-diagonal elements in row 1 or column 1, and continuing again the minimizing X is I. (Notice we do not even have to show that for this P, the principal diagonal is minimal; this emerges from the proof.)

#### 4. Generalization of the Wielandt-Hoffman Theorem

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A generalization of the Wielandt-Hoffman Theorem which is of practical importance is the following.

THEOREM 3. If X is an nxr matrix with orthonormal columns, A is an nxn normal matrix, B is an rxr normal matrix and R an nxr matrix is defined by

$$AX - XB = R , \qquad (4.1)$$

if the eigenvalues  $\mathbf{a}_{i}$  (i = 1,...,n) of A and  $\mathbf{b}_{i}$  (i = 1,...,r) of B are ordered so that  $\sum_{i=1}^{r} |\mathbf{a}_{i} - \mathbf{b}_{i}|^{2}$  is a minimum, then

$$\sum_{i=1}^{r} \left| \mathbf{a}_{i} - \mathbf{b}_{i} \right|^{2} \leq \left\| \mathbf{R} \right\|_{\mathbf{F}}^{2} \qquad (4.2)$$

A weaker result with  $||R||_{F}^{2}$  replaced by  $2^{1/2} ||R||_{F}^{2}$  was given by Wilkinson in [5] and the result itself by Kahan [3].

Notice we are interested only in the selection and ordering of the relevant r of the  $a_i$  to be associated with the  $b_i$ . Writing

$$A = Q_1 \operatorname{diag}(a_1)Q_1^H$$
,  $B = Q_2 \operatorname{diag}(b_1)Q_2^H$  (4.3)

with the prescribed ordering of the  $a_i$  and  $b_i$ , we have

$$\|\operatorname{diag}(\mathbf{a}_{i}) Q - Q \operatorname{diag}(\mathbf{b}_{i})\|_{F}^{2} = \|Q_{1}^{H} R Q_{2}\|_{F}^{2} = \|R\|_{F}^{2}$$
 (4.4)

where Q is an nxr matrix with ortho-normal columns. Hence

$$\sum_{j=1}^{r} \sum_{i=1}^{n} |a_{i} - b_{j}|^{2} |q_{ij}|^{2} = ||R||_{F}^{2} .$$
 (4.5)

Let Y = [Q | Z] be an nxn unitary matrix given by the completion of Q; then if

$$p_{ij} = |a_i - b_j|^2$$
  $(j \le r)$ ,  $p_{ij} = 0$   $(j > r)$ . (4.6)

$$\sum_{i=1}^{n} \sum_{j=1}^{n} p_{ij} |y_{ij}|^{2} = \sum_{j=1}^{r} \sum_{i=1}^{n} |a_{i} - b_{j}|^{2} |q_{ij}|^{2}$$
(4.7)

and from the definition of the ordering of the  $a_i$  and  $b_i$ , the diagonal of P is minimal. Hence by Theorem 1 and Equation (4.5)

$$\sum_{i=1}^{n} p_{ii} = \sum_{i=1}^{r} |a_{i} - b_{i}|^{2} \le \sum_{j=1}^{r} \sum_{i=1}^{n} |a_{i} - b_{j}| |q_{ij}|^{2} = ||R||_{F}^{2}.$$
(4.8)

This theorem is of practical value when r orthonormal approximate eigenvectors  $x_1, \ldots, x_r$  are known corresponding to alleged eigenvalues  $\mu_1, \ldots, \mu_r$ . If

$$Ax_{i} - \mu_{i}x_{i} = r_{i}$$
 (i = 1,...,r) (4.9)

Then

$$AX - X \operatorname{diag}(\mu_{\star}) = R \tag{4.10}$$

with an obvious notation, and  $diag(\mu_1)$  is the matrix B of Theorem 3. This then states that there exist r eigenvalues  $a_1, \ldots, a_r$  of A such that

$$\sum_{i=1}^{r} (a_{i} - \mu_{i})^{2} = ||R||_{F}^{2} \qquad . \qquad (4.11)$$

Notice that the  $\mu_i$  can include multiple or pathologically chic eigenvalues. The result is well known when r = 1 and the Wielandt-Hoffman theorem corresponds to the case r = n. We observe that by using less than r of the alleged eigenvectors we can obtain results of the type (4.11) corresponding to any  $s (\leq r)$  of the approximate eigenvalues.

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