Working Paper No. 123

AD 659989

PROPER EFFICIENCY AND THE THEORY OF VECTOR MAXIMIZATION

by

ARTHUR M. GEOFFRION

August, 1967

CCT 24 1957

WESTERN MANAGEMENT SCIENCE INSTITUTE

rill, die 1

University of California, Los Angeles

18 (E) (G) (

The same of

24

University of California

Los Angeles

Western Management Science Institute

Working Paper No. 123

PROPER EFFICIENCY AND THE THEORY OF VECTOR MAXIMIZATION

bу

Arthur M. Geoffrion

August 1967

This working paper should be regarded as preliminary in nature, and subject to change before publication in the open literature. It should not be quoted without prior consent of the author. Comments are cordially invited.

This work was jointly sponsored by the United States Air Force under Project RAND, and by the Western Management Science Institute under grants from the National Science Foundation and the Office of Naval Research. It is a pleasure to acknowledge the helpful comments of Allen Klinger. SUMMARY

The concept of <u>efficiency</u> in problems with multiple criterion functions--sometimes under an alias such as "admissibility" or "Pareto optimality"--has long played an important role in economics, game theory, statistical decision theory, and in all optimal decision problems with noncomparable criteria. Here we propose a slightly restricted definition of efficiency that eliminates efficient points of a certain anomalous nature. This new definition, which we call <u>proper</u> efficiency, is related in spirit to the notion of "proper" efficiency introduced by Kuhn and Tucker in their celebrated paper of 1950; but the present definition avoids certain drawbacks inherent in the earlier one. A comprehensive theory of vector maximization is constructed using the new definition, with and without various constraint qualification, convexity, and differentiability assumptions. The theory includes as a special case the standard theory of nonlinear programming.

I. INTRODUCTION

Given a vector-valued criterion function $f(x) = (f_1(x), \dots, f_p(x))$ and a set $X \subseteq R^n$ of feasible points, the Vector Maximum Problem

(VMP) V-MAX f(x) subject to x « X

is the problem of finding all points that are <u>efficient</u>: x° is said to be efficient if $x^{\circ} \in X$ and there exists no other feasible point x such that $f(x) \ge f(x^{\circ})$ but $f(x) \ne f(x^{\circ})$. The concept of efficiency--sometimes under an alies such as "admissibility," "maximality," noninferiority," or "Pareto optimality"--has long played an important role in economics, game theory, statistical decision theory, and in all optimal decision problems with noncomparable criteria.

In this study we propose a slightly restricted definition of efficiency that (a) eliminates efficient points of a certain anomalous type; and (b) lends itself to more satisfactory characterization (see Theorem 2 below, and Sec. II). We shall call this new definition proper efficiency, although Kuhn and Tucker [7] have previously used the same term. Their intent appears to have been much the same as ours but, as we shall see, the present definition is of greater generality and seems to be somewhat more natural.

PROPER EFFICIENCY

<u>Definition:</u> x^{0} is said to be a proper efficient solution of (VMP) if it is efficient and if there exists a scalar M > 0 such that, for each i, $f_{i}(x) > f_{i}(x^{0})$ and $x \in X$ implies

$$\frac{f_{i}(x) - f_{i}(x^{0})}{f_{i}(x^{0}) - f_{i}(x)} \leq H$$

for some j such that $f_j(x) < f_j(x^0)$.

-1-

An efficient point that is not properly efficient is said to be <u>improperly</u> efficient. Thus for x^0 to be improperly efficient means that to every scalar M > 0 (no matter how large) there is a point x & X and an i such that $f_i(x) > f_i(x^0)$ and

$$\frac{f_{i}(x) - f_{i}(x^{0})}{f_{i}(x^{0}) - f_{i}(x)} > M$$

for all j such that $f_j(x) < f_j(x^0)$. If we take a sequence $\langle M^{\vee} \rangle \to \infty$ and remember that there is but a finite number of criteria, we see that for some criterion i_0 , the marginal gain in f_i can be made arbitrarily large relative to each of the marginal losses incurred by other criteria. Assuming that the decisionmaker's desire for f_i is not satiated, x^0 certainly seems undesirable. An example of improper efficiency is given in Sec. III.

CHARACTERIZATION

A matter of great interest, both computationally and theoretically, is the relation of the Vector Maximum Problem to the following scalar maximum problem:

$$(P_{\lambda}) \qquad MAX \quad \sum_{i=1}^{p} \lambda_{i}f_{i}(x) \text{ subject to } x \in X,$$

where the λ_i are nonnegative parameters often normalized according to $\sum_{i=1}^{p} \lambda_i = 1$. The fundamental results characterizing proper vector maxima in terms of the solutions of (P_{λ}) are given in Theorems 1 and 2.

Theorem 1. Let $\lambda_i > 0$ (i = 1, ..., p) be fixed. If x^0 is optimal in (P_{λ}) , then x^0 is properly efficient in (VMP)

<u>Proof</u>: It is easy to show by contradiction that x° is efficient. Let $\overline{\lambda} = \max_{i,j} \{\lambda_j/\lambda_i\}$. We shall show that x° is <u>properly</u> efficient in (VMP) i,j with $M = p\overline{\lambda}$; that is, we shall show that for each i, $f_i(x) \ge f_i(x^{\circ})$ and $x \in X$ implies $f_i(x) - f_i(x^{\circ}) \le p\overline{\lambda}$ ($f_j(x^{\circ}) - f_j(x)$) for some $j \in J_X$, where $J_X \triangleq \{1 \le j \le p: f_j(x) < f_j(x^{\circ})\}$. Suppose to the contrary that, for some i and $x \in X$, we have $f_i(x) \ge f_i(x^{\circ})$ and

$$f_{i}(x) - f_{i}(x^{0}) > p\bar{\lambda} (f_{j}(x^{0}) - f_{j}(x)), \forall j \in J_{x}.$$

Summing the latter inequalities and dividing the result by the number q of elements of J_x , we obtain

$$(f_{i}(x) - f_{i}(x^{o})) > \frac{p}{q} \sum_{j \in J_{x}} \overline{\lambda}(f_{j}(x^{o}) - f_{j}(x)).$$

Using p/q > 1 and the definitions of $\tilde{\lambda}$ and J_{χ} , in that order, it follows that

$$(f_{i}(x) - f_{i}(x^{\circ})) > \sum_{j \neq i} (\lambda_{j} / \lambda_{i}) (f_{j}(x^{\circ}) - f_{j}(x))$$

Multiplying through by λ_i and rearranging, we see that this contradicts the optimality of x^0 in (P_λ) . Hence x^0 must be properly efficient.

> <u>Theorem 2</u>. Let X be a convex set, and let the f_i be concave on X. Then x^o is properly efficient in (VMP) if and only if x^o is optimal in (P_λ) for some λ with strictly positive components.

<u>Proof</u>: The "if" part of the theorem is provided by Theorem 1. If x° is properly efficient, then there exists a scalar M > 0 such that for each i (i = 1, ..., p) the system

$$f_{i}(x) > f_{i}(x^{o})$$

 $f_{i}(x) + M f_{j}(x) > f_{i}(x^{o}) + M f_{j}(x^{o}), all j \neq i$

admits no solution in X. By a fundamental property of concave functions [2, p. 62], for the ith system there exist $\lambda_j^i \ge 0$ (j = 1, ..., p) with p Σ $\lambda_j^i = 1$ such that j=1

$$\lambda_{i}^{i} f_{i}(\mathbf{x}) + \sum_{j \neq i} \lambda_{j}^{i} (f_{i}(\mathbf{x}) + \mathbf{H} f_{j}(\mathbf{x})) \leq \lambda_{i}^{i} f_{i}(\mathbf{x}^{0}) + \sum_{j \neq i} \lambda_{j}^{i} (f_{i}(\mathbf{x}^{0}) + \mathbf{H} f_{j}(\mathbf{x}^{0}))$$

for all x & X. Summing over i yields, after some simplification and rearrangement,

$$\sum_{j=1}^{\mathbf{p}} (1 + \mathbf{M} \sum_{i \neq j} \lambda_{j}^{i}) f_{j}(\mathbf{x}) \leq \sum_{j=1}^{\mathbf{p}} (1 + \mathbf{M} \sum_{i \neq j} \lambda_{j}^{i}) f_{j}(\mathbf{x}^{o})$$

for all x & X. This completes the proof.

ā.

Thus from a computational viewpoint, finding proper efficient solutions is reduced to a parametric programming problem; (P_{λ}) yields only properly efficient solutions as λ varies over

$$\Lambda^{+} \Delta \{ \lambda \in \mathbb{R}^{P} : \text{ all } \lambda_{i} > 0 \text{ and } \sum_{i=1}^{P} \lambda_{i} = 1 \},$$

A more complete characterization theory for the Proper Vector Maximum Problem is developed in the next section. It provides, for example, necessary conditions for a proper vector maximum in the absence of concavity.

* In this regard see, for example, Charnes and Cooper [3, Ch. 9], Markowitz [8], and Geoffrion [4].

II. THEORY

We shall give the theory of the Proper Vector Maximum Problem in terms of the relationships between the following six problems. In problems 3, 4, and 5. X is taken to be of the form $X = \{x; g(x) \ge 0\}$, where $g(x) = (g_1(x), \dots, g_m(x))$. In problems 3 and 4, the differentiability of all functions is presumed.

<u>Problem 1</u> - Find a point $\hat{\mathbf{x}}$ that is a proper efficient solution of (VMP).

<u>Problem ?</u> - Find a point $\bar{\mathbf{x}}$ that is a locally proper efficient solution of (VMP).

<u>Problem 3</u> - Find a feasible point \bar{x} such that none of the p systems^{**} (i = 1, ..., p)

 $\nabla_{\mathbf{x}} \mathbf{f}_{\mathbf{i}}(\mathbf{x}) \cdot \mathbf{u} \ge 0$

 $\nabla_{\mathbf{x}} t_{\mathbf{j}}(\mathbf{x}) \cdot \mathbf{u} \neq 0$, all $\mathbf{j} \neq \mathbf{i}$

$$\nabla_{\mathbf{x}} g_{\mathbf{j}}(\mathbf{x}) \cdot \mathbf{u} \ge 0$$
, all $\mathbf{j} \ni g_{\mathbf{j}}(\mathbf{x}) = 0$

has a solution u in Rⁿ.

Froblem 4 - Find a feasible point $\bar{\mathbf{x}}$, a point $\bar{\mathbf{y}} \ge 0$ in \mathbb{R}^m , and a point $\bar{\lambda} \in \Lambda^+$ such that $\bar{\mathbf{y}} \ge (\bar{\mathbf{x}}) = 0$ and

$$\nabla_{\mathbf{x}} [\overline{\mathbf{x}} \cdot \mathbf{f}(\overline{\mathbf{x}}) + \overline{\mathbf{y}} \ \mathbf{g}(\overline{\mathbf{x}})] = 0$$

<u>Problem 5</u> - Find a feasible point \bar{x} , a point $\bar{y} \ge 0$ in \mathbb{R}^m , and a point $\bar{\lambda} \in \Lambda^+$ such that $\bar{y} \cdot g(\bar{x}) = 0$ and \bar{x} achieves the unconstrained maximum of $\bar{\lambda} \cdot f(x) + \bar{y} \cdot g(x)$.

<u>Problem 6</u> - Find a point $\bar{\mathbf{x}}$ and a point $\bar{\mathbf{\lambda}} \in \Lambda^+$ such that $\bar{\mathbf{x}}$ is optimal in $(P_{\bar{\mathbf{x}}})$.

** $\nabla_{\varphi}(\bar{x})$ represents the gradient vector of the function φ evaluated at $x = \bar{x}$.

 $[\]tilde{x}$ is said to be a locally proper efficient solution of (VMP) if it is properly efficient in N- $f \in X$, where N- is some (open convex) neighborhood of \tilde{x} .

Problem 1 is the central problem of interest. Problem 2 is its "local" equivalent, and problem 3 is the local problem in terms of directional derivatives. Problem 4 represents the generalized Lagrange multiplier or Kuhn-Tucker conditions in differential form associated with problem 1. Problem 5 is precisely equivalent to the following saddle-point problem:

> Find a point \bar{x} , a point $\bar{y} \ge 0$ in \mathbb{R}^m , and a point $\bar{\lambda} \in \Lambda^+$ such that the pair (\bar{x}, \bar{y}) is a saddle-point subject to $y \ge 0$ of the function $F(x,y) = \bar{\lambda} \cdot f(x) + y \cdot g(x)$; i.e., such that $F(\bar{x},y) \le F(\bar{x}, \bar{y}) \le F(x, \bar{y})$ for all $x \in \mathbb{R}^m$ and $y \ge 0$ in \mathbb{R}^m .

Problem 5 is also of interest for its own sake. Problem 6 is just (P_1)

In stating the relations between these problems, we shall use the A_1, \cdots notation j \longrightarrow k, which is to be understood as follows. Let (u,v) be the unknowns of problem j and (u,w) the unknowns of problem k. Then this notation is to be read. "If (\bar{u},\bar{v}) solves problem j, and if assumptions A_1, \cdots hold, then there exists \bar{w} such that (\bar{u},\bar{w}) solves problem k." Or, somewhat more loosely, "Under assumptions A_1 , \ldots , every solution of problem j is also a solution of problem k." The assumptions which will be used at one time or another are:

Assumption C: All functions are <u>concave</u> on $\mathbf{E}^{\mathbf{n}}$.

<u>Assumption D:</u> All functions are continuously differentiable on E^n .

Assumption Q1: The following constraint <u>qualifica-</u> tion holds: there exists a feasible point \hat{x} such that $g_1(\hat{x}) > 0$ for g_j nonlinear.

Assumption Q₂: The Kuhn-Tucker constraint <u>qualifica-</u> tion holds [7, p. 483].

-6-

We are now in a position to state the relationships between the six problems.



For example, the Comprehensive Theorem asserts $(1 \rightarrow 2)$ that every proper efficient solution of (VMP) is a locally proper efficient solution of (VMP), and $(1 \stackrel{C}{=} 2)$ that the converse is true under Assumption C. It also asserts $(5 \rightarrow 6)$ that if $(\bar{x}, \bar{y}, \bar{\lambda})$ solves problem 5, then $(\bar{x}, \bar{\lambda})$ solves problem 6; and $(5 \stackrel{C,Q_1}{\longleftarrow} 6)$ that if $(\bar{x}, \bar{\lambda})$ solves problem 6, then there exists a point $\bar{y} \in \mathbb{R}^m$ such that $(\bar{x}, \bar{y}, \bar{\lambda})$ solves problem 5.

Because of its length, we give the proof in Appendix A.

The Comprehensive Theorem is actually many theorems in one. Its significance is that it gives, under various assumptions, necessary and/or sufficient conditions for proper efficiency. In order to be explicit, we state the most important of these conditions as three simple corollaries of the Comprehensive Theorem. Corollary 1 asserts that under Assumptions D and Q_2 , the conditions of problem 4 are necessary first order conditions for proper efficiency. Corollary 2 characterizes problem 1 as being equivalent (in the appropriate sense) to problems 2, 5, and 6 under Assumptions C and Q_1 . Corollary 3 asserts

-7-

that all six problems are equivalent under C, D, and either Q_1 or Q_2 .

<u>COR 1</u> - If Assumptions D and Q hold, then problem $2 \rightarrow$ problem 4.

<u>COR 2</u> - If Assumptions C and Q₁ hold, these problem $1 \stackrel{\frown}{\rightarrow}$ problem 2 $\stackrel{\frown}{\rightarrow}$ problem 5 $\stackrel{\frown}{\rightarrow}$ problem 6.

<u>COR 3</u> - If Assumptions C, D, and either Q_1 or Q_2 hold, then problem 1 $\xrightarrow{\sim}$ problem j for j = 2, ..., 6.

The Comprehensive Theorem subsumes, of course, the cases in which there are no constraints or only equality constraints. Again for the sake of explicitness, we shall state the main results for these cases in Appendix B.

It is of interest to note that in the special case all of the f_i are identical or p = 1, the notion of proper efficiency coincides with the notion of a constrained maximum, so that the results of the Comprehensive Theorem reduce to well-known counterparts in the theory of nonlinear programming.

LII. DISCUSSION

We turn now to further discussion of the notion of proper efficiency.

Just how slight a restriction proper efficiency is over efficiency can perhaps be better appreciated in the light of the following. Denote the set of all efficient (properly efficient) points by X^2 ($X^2_{pr.}$), and the image in R^p of X^2 under f by $f[X^2]$. If the f_i are continuous and concave on the closed convex set X, then $f[X^2_{pr.}] \subseteq f[X^2] \subseteq \overline{f[X^2_{pr.}]}$, where the bar denotes closure. This result is a consequence of Theorem 2 and a result due to Arrow, Barankin and Blackwell [1]. Thus under the given conditions, which are almost always satisfied in concave programming, the outcome of any improperly efficient point is always the limit of the outcomes of some sequence of properly efficient points.

COMPARISON WITH THE DEFINITION OF KUHN AND TUCKER

The notion of "proper" efficiency introduced by Kuhn and Tucker applies only when assumptions D and Q_2 hold. Under these assumptions, x° is said to be "properly" efficient if it is efficient and if it solves problem 3. Let us denote the problem of finding such a "properly" efficient point as $(X^{\geq}, 3)$. Then the results obtained by Kuhn and Tucker are^{**} (in the presence of D and Q_2):

-9-

^{*}If S is a closed convex set in R^n , then the set of efficient points of S contains the subset of points of S for which there is a supporting hyperplane whose normal has all positive components, and is contained in the closure of the last mentioned set.

Each of these assertions can be obtained as an immediate corollary of the Comprehensive Theorem.



To justify excluding efficient solutions that are not "proper," Kuhn and Tucker give an example with p = 2 in which such a solution admits a first-order gain in one criterion at the expense of but a second-order loss in the other. Indeed, every "improperly" efficient solution poses an equally objectionable anomaly.⁴ The converse, however, is not true--not every anomalous efficient point is "improper" in the sense of Kuhn and Tucker, as the following example shows. Put n = 1, m = 1, p = 2, g(x) = x, $f_1(x) = x^2$, $f_2(x) = -x^3$, $x^\circ = 0$. Assumptions D and Q_2 hold, and x° is "properly" efficient, but for x positive and sufficiently small the gain in f_1 can be made arbitrarily large with respect to the loss in f_2 (the gain-to-loss ratio is 1/x for $x \ge 0$).

Since $1 \xrightarrow{D,Q_2} 3$ (see Comprehensive Theorem), the set of points "properly" efficient in the sense of Kuhn and Tucker contains all those properly efficient in the present sense. The above example (in which x° is improperly efficient in the sense of Sec. I) shows that the containment can be strict.

To summarize, the advantages of the present definition of proper efficiency over that of Kuhn and Tucker seem to be that it excludes all of a precise class of anomalies, and that it applies even in the absence of Assumption D or Q_2 .

-10-

For an explicit proof see **Elinger** [6]; his proof seems to require the locus of $\mathbf{g}(t)$ in the definition of Q₂ to be linear, but this restriction can be removed (cf. the proof of 2 ______3 in Appendix A).

CONCLUSION

We began with the premise that, in optimization problems with multiple criteria, it is natural to restrict attention to efficient decisions that are properly so--in the sense that at least one potential marginal gain-to-loss ratio must be bounded. We then obtained, in Theorems 1 and 2, basic characterization results for proper efficiency in terms of the scalar parametric problem (P_{λ}) . These results were extended in the Comprehensive Theorem to include the relationships with four other intimately related problem formulations, with and without various constraint qualifications, differentiability and convexity assumptions. The result is a coherent theory of the Proper Vector Maximum Problem which generalizes the well-known Kuhn-Tucker theory for nonlinear programming. This theory seems more satisfactory than that possible using either the usual definition of efficiency or the closely related definition of "proper" efficiency proposed by Kuhn and Tucker.

-11-

Appendix A PROOF OF THE COMPREHENSIVE THEOREM

A. 1 _____6. This is a restatement of Theorems 1 and 2 (with λ normalized).

B. $6 \xrightarrow{C,Q_1}{5} \xrightarrow{D} 4$. These assertions are all known results from the theory of nonlinear programming applied to (P_{λ}) .

 $6 \xrightarrow{C,Q_1} 5$ is a consequence of a slightly more general form of the Farkas-**Highsughi** Theorem [2, p. 67].

5 6 is easily verified directly.

 $5 \xrightarrow{D} 4$ occurs because the gradient of a continuously differentiable function must vanish at an unconstrained extremum.

4 C 5 occurs because a concave function (which $\bar{\lambda} \cdot f(x) + \bar{y} \cdot g(x)$ must be, since $\bar{\lambda} \ge 0$ and $\bar{y} \ge 0$) achieves an unconstrained supremum at any point for which its gradient vanishes.

C. 1 2. 1 2 is trivial.

Let \bar{x} be a locally proper efficient solution in the neighborhood $N_{\bar{x}}$. Under Assumption C, Theorem 2 tells us that \bar{x} maximizes $\bar{\lambda} \cdot f(x)$ on $N_{\bar{x}} \cap X$ for some $\bar{\lambda} \in \Lambda^+$. Again from Assumption C, \bar{x} must maximize $\bar{\lambda} \cdot f(x)$ over X, and so by Theorem 1 \bar{x} must be properly efficient. Thus $2 - \frac{C}{2} - 1$.

D. $3 \longrightarrow 4$. This result can essentially be found in [7, Theorems 4 and 5].

-13-

equation by u, we readily see by contradiction that \bar{x} must be a solution of Problem 3.

To see 3 \rightarrow 4, let \overline{x} be a solution of Problem 3 and apply the Farkas-Minkowski Theorem in turn to each of the p systems. As a result, there must exist numbers $w_1^L \ge 0$ and $x_1^i \ge 0$, such that, for i = 1, ..., p,

$$\nabla_{\mathbf{x}} \mathbf{f}_{\mathbf{i}}(\mathbf{\ddot{x}}) + \sum_{\mathbf{j} \neq \mathbf{i}} \mathbf{\nabla}_{\mathbf{x}} \mathbf{f}_{\mathbf{j}}(\mathbf{\ddot{x}}) + \sum_{\mathbf{j} \neq \mathbf{j}} \mathbf{\nabla}_{\mathbf{x}} \mathbf{g}_{\mathbf{j}}(\mathbf{\ddot{x}}) = 0.$$

Summing over i yields

$$\begin{array}{l} \mathbf{p} \\ \Sigma (\mathbf{1} + \Sigma \mathbf{w}_{\mathbf{j}}^{\mathbf{i}}) \nabla_{\mathbf{x}} \mathbf{f}_{\mathbf{j}}(\bar{\mathbf{x}}) + \Sigma (\Sigma z_{\mathbf{j}}^{\mathbf{i}}) \nabla_{\mathbf{x}} \mathbf{g}_{\mathbf{j}}(\bar{\mathbf{x}}) = 0, \\ \mathbf{j} = \mathbf{1} \qquad \mathbf{i} \neq \mathbf{j} \qquad \mathbf{j} \in \mathbf{J} \quad \mathbf{i} = \mathbf{1} \end{array}$$

Put
$$\hat{\lambda}_{i} = (1 + \sum_{i \neq j} w_{j}^{i}), \ \hat{y}_{j} = (\sum_{i=1}^{p} z_{j}^{i}) \text{ for } j \in J$$

and
$$\hat{y}_{j} = 0$$
 for $j \notin J$. Clearly $\overline{x}, \overline{\lambda}_{j} = \hat{\lambda}_{j} / (\sum_{l=1}^{p} \hat{\lambda}_{l}),$

and

 $\overline{y}_{j} = \hat{y}_{j} / (\sum_{i=1}^{p} \hat{\lambda}_{i})$ solves Problem 4. E. 2 $\xrightarrow{D, Q_2}$ 3. We have previously shown 3 \xrightarrow{Q} 4 C \xrightarrow{C} 5 \xrightarrow{C} 6 \xrightarrow{C} 1 \xrightarrow{C} 2; hence 3 \xrightarrow{C} 2. To C complete the proof of the Comprehensive Theorem it remains only to

show $2 \xrightarrow{D, Q2} 3$.

Let \bar{x} be a locally proper efficient solution of (VMP), and let Assumptions D and Q_2 hold. Suppose, contrary to what we desire to show, that \overline{x} is not a solution of Problem 3. Then one of the p systems, say the first, has a solution: there exists \overline{u} such that

$$\nabla_{\mathbf{x}} \mathbf{f}_{1} (\overline{\mathbf{x}}) \cdot \overline{\mathbf{u}} > 0$$

$$\nabla_{\mathbf{x}} \mathbf{f}_{j} (\overline{\mathbf{x}}) \cdot \overline{\mathbf{u}} \ge 0, \ \mathbf{j} = 2, \dots, \mathbf{p}$$

$$\nabla_{\mathbf{x}} \mathbf{g}_{1} (\overline{\mathbf{x}}) \cdot \overline{\mathbf{u}} \ge 0, \ \mathbf{all} \ \mathbf{j} \ge \mathbf{g}_{1} (\overline{\mathbf{x}}) = 0.$$

By Assumption Q_2 there exists a continuously differentiable arc $\hat{\mathbf{x}}(t)$, $0 \le t \le 1$, contained in the feasible region, with $\hat{\mathbf{x}}(0) = \overline{\mathbf{x}}$ and some positive scalar α such that $\left(\frac{d \hat{\mathbf{x}}_1(0)}{dt}, \dots, \frac{d \hat{\mathbf{x}}_n(0)}{dt}\right) = \alpha \overline{\mathbf{u}}$. Consider the functions $f_1(\hat{\mathbf{x}}(t))$. From Taylor's Theorem we have:

$$f_{i}(\hat{x}(t)) = f_{i}(\hat{x}(0)) + t \frac{d f_{i}(\hat{x}(t_{i}))}{dt}$$

$$= f_{i}(\overline{x}) + t \sum_{j=1}^{n} \frac{\partial f_{i}(x)}{\partial x_{j}} \left| \frac{d \hat{x}_{i}(t)}{dt} \right| t_{i}$$

where t_i is some scalar between 0 and t. Denote the summation in the last term by $s_i(t)$, so that $f_i(\hat{x}(t)) = f_i(\bar{x}) + t s_i(t)$. Evidently $s_i(0) = \alpha \nabla_x f_i(\bar{x}) \cdot \bar{u}$ and $s_i(t)$ is continuous (from the right) at t = 0. Now for t sufficiently near 0, $\hat{x}(t)$ is in the neighborhood within which \bar{x} is properly efficient. Consider a sequence $< t^{\vee} > \rightarrow 0$, where $t^{\vee} > 0$. By taking a subsequence, if necessary, we may assume that the set $\{j: f_j(\hat{x}(t^{\vee})) < f_j(\bar{x})\}$ is constant for all $\nu - call$ it J⁻. We know that $< s_j(t^{\vee}) > \rightarrow \alpha \nabla_x f_j(\bar{x}) \cdot \bar{u} \ge 0$, all $j \in J^-$. But $s_j(t^{\vee}) < 0$ by definition for all ν and $j \in J^-$, and so $< s_j(t^{\vee}) > \rightarrow 0$ for all $j \in J^-$. Furthermore, $< s_1(t^{\vee}) > \rightarrow \alpha \nabla_x f_1(\bar{x}) \cdot u \ge 0$.

-15-

$$\langle \frac{f_{1}(\hat{\mathbf{x}}(t^{\vee})) - f_{1}(\bar{\mathbf{x}})}{f_{j}(\bar{\mathbf{x}}) - f_{j}(\hat{\mathbf{x}}(t^{\vee}))} \rangle , j \in J^{-},$$

which can be written

$$\langle \frac{t^{\vee} s_{1}(t^{\vee})}{-t^{\vee} s_{1}(t^{\vee})} \rangle$$
, j ϵJ^{-} ,

•

all diverge to + =. But this contradicts the local proper efficiency of \overline{x} , and so \overline{x} must indeed be a solution of Problem 3.

Appendix B

NO CONSTRAINTS AND EQUALITY CONSTRAINTS

NO CONSTRAINTS

Here we consider the case in which X is an open set in \mathbb{R}^n (perhaps the whole of \mathbb{R}^n). Corollary 4 gives necessary, and Cor. 5 sufficient conditions for a locally proper efficient (1.p.e.) solution.

<u>COR 4</u> - Let the f_i be continuously differentiable on X.

If \mathbf{x}° is l.p.e., then $\nabla_{\mathbf{x}}[\lambda \cdot f(\mathbf{x}^{\circ})] = 0$ for som $\lambda \in \Lambda^{+}$. <u>Proof</u>: With $\mathbf{w} = 0$ and $\mathbf{x}^{\circ} \in \mathbf{X}$, \mathbf{Q}_{2} becomes superfluous, and the

<u>COR 5</u> - Let the f_i be twice continuously differentiable on an open set $X \subseteq \mathbb{R}^n$. If $x^o \in X$ satisfies $\nabla_x[\lambda \cdot f(x)] = 0$ for some $\lambda \in \Lambda^+$, and the Hessian $\nabla_x^2 [\lambda \cdot f(x)]$ is negative definite, then x^o is l.p.e.

<u>Proof</u>: The assumptions imply that $\chi \cdot f(x)$ is strictly concave on some (convex) open neighborhood N_xo of x^O. Hence x^O maximizes this function on N_xo, and so by Theorem 1 x^O must be l.p.e.

It is clear from the proof of Cor. 5 that the hypothesis " f_i twice continuously differentiable and $\nabla_x^2[\lambda^*f(x^*)]$ negative definite" can be weakened to " f_i continuously differentiable and $\lambda^*f(x)$ concave on some neighborhood of x^0 ."

EQUALITY CONSTRAINTS

Here we consider the case $X = \{x:g_j(x) = 0, j = 1, ..., m\}$. The Comprehensive Theorem subsumes this case if we write X as $\{x:g_j(x) \ge 0, j = 1, ..., m \text{ and } \frac{-\Sigma}{i=1}g_i(x) \ge 0\}$. Assumption Q_i is satisfied if and only if all constraints are linear; and the directions u of concern in Q_2 are those for which $\nabla_{\mathbf{x}} \mathbf{g}_{\dagger}(\mathbf{x}) \cdot \mathbf{u} = 0$, $\mathbf{j} = 1, ..., m$.

Corollary 6 is a Lagrange Multiplier Theorem, and Cor. 7 examines the linear constraints case.

<u>COR 6</u> - Let the f_i and g_j be continuously differentiable on some neighborhood of x° , and let Q_2 hold at x° . If x° is l.p.e., then

$$\nabla_{\mathbf{x}} [\lambda \cdot f(\mathbf{x}^{\circ}) + \mu \cdot g(\mathbf{x}^{\circ})] = 0$$

for some $\lambda \in \Lambda^+$ and $\mu \in \mathbb{R}^m$.

 $\mu_{j} = \overline{y}_{j} - \overline{y}_{m+1}.$

<u>COR 7</u> - Let the g_j be linear, and the f_i concave. Then each of the following conditions is necessary and sufficient for x° to be properly efficient:

- (i) x° maximizes $\lambda \cdot f(x)$ subject to g(x) = 0for some $\lambda \in \Lambda^+$;
- (ii) x° is feasible, and maximizes $\lambda \cdot f(x) + \mu \cdot g(x)$ over all x for some $\lambda \in \Lambda^+$ and $\mu \in \mathbb{R}^m$;
- (111) there exists $\mu^{\circ} \in \mathbb{R}^{m}$ such that (x°, μ°) is a saddlepoint of the function $F(x, \mu) = \lambda^{\circ} \cdot f(x) + \mu^{\circ}g(x)$ for some $\lambda^{\circ} \in \Lambda^{+}$; i.e., $F(x^{\circ}, \mu) \leq F(x^{\circ}, \mu^{\circ}) \leq F(x, \mu^{\circ})$ for all $x \in \mathbb{R}^{n}$ and $\mu \in \mathbb{R}^{m}$.

If, in addition, the f_i are continuously differentiable, then a fourth equivalent condition is:

(iv) x^o satisfies



for some $\lambda \in \Lambda^+$ and $\mu \in \mathbb{R}^m$.

<u>Proof</u>: Directly from the Comprehensive Theorem.

REFERENCES

- Arrow, K. J., E. W. Barankin, and D. Blackwell, "Admissible Points of Convex Sets," in R. W. Kuhn and A. W. Tucker (eds.), <u>Contribu-</u> <u>tions to the Theory of Games</u>, Princeton University Press. 1953, pp. 87-91.
- Berge, C., and A. Gouila-Houri, <u>Programming, Games, and Transportation</u> <u>Networks</u>, John Wiley and Sons, New York, 1965.
- Charnes, A., and W. W. Cooper, <u>Management Models and Industrial</u> <u>Applications of Linear Programming</u>, Vol. 1, John Wiley and Sons, New York, 1961.
- Geoffrion, A. M., "Strictly Concave Parametric Programming, Parts I and II," <u>Management Science</u>, Vols. 13, Nos. 3 and 5, November 1966 and January 1967, pp. 244-253 and 359-370.
- Karlin, S., <u>Mathematical Methods and Theory in Games</u>, <u>Programming</u>, <u>and Economics</u>, Vol. I, Addison-Wesley, Reading, Massachusetts, 1959.
- 6. Klinger, A., "Improper Solutions of the Vector Maximum Problem," Operations Research, Vol. 15, No. 3, May-June 1967, pp. 570-572.
- Kuhn, H. W., and A. W. Tucker, "Nonlinear Programming," <u>Proc. Second</u> <u>Berkeley Symposium on Mathematical Statistics and Probability</u>, University of California Press, Berkeley, California, 1950, pp. 481-492.
- Markowitz, H., "The Optimization of a Quadratic Function Subject to Linear Constraints," <u>Naval Research Logistics Quarterly</u>, Vol. 3, Nos. 1 and 2, March and June 1955, pp. 111-133.

DOCUM	ENT CONTROL DATA - R&D	D Isred when	the overall report is classified)		
ORIGINATIN & ACTIVITY (Corporate author)		2. REPO	RT SECURITY CLASSIFICATION		
Western Management Science Institute		Unclassified			
University of California at	: Los Angeles	25 GROUI	P		
PROPER EFFICIENCY AND THE T	THEORY OF VECTOR MA	XIMIZ	ATION		
DESCRIPTIVE NOTES (Type of report and inclusive	dales)				
WORKING Paper AUTHOR(S) (Lest name first name initial)	······································				
Geoffrion, Arthur M.					
REPORT DATE	74 TOTAL NO OF PA	GFS	75 NO OF REFS		
August 1967	21		8		
Nonr 233(75)	9. ORIGINATOR'S REI	PORTNUM	BER(S)		
L PROJECT NO	Working Pa	aper No	0. 123		
τ	SO OTHER REPORT N	O(S) (Any	other numbers that may be assigned		
		this report)			
	this report)				
J O AVAILABILITY LIMITATION NOTICES Distribution of this documen			Linge Institute		
J O AVAILABILITY LIMITATION NOTICES Distribution of this documen is unlimited.			Distance Institute		
 ³ ³ ABSTRACT The concept of eff functionssometimes under a optimality"has long played theory, statistical decision lems with noncomparable crit definition of efficiency tha anomalous nature. This new is related in spirit to the oy Kuhn and Tucker in their definition, with and wit comprehensive theory of vect new definition, with and wit convexity, and differentiabit a special case the standard 	Los A Loci, Los A Los A Loci, Los A Los A Los Los A Los A Los A Los A Los A Los Los A Los A	Any ACTION Any ACTION admission in each opose a cient p we can define of 1950 in the s const craint The far pros	h multiple criteria ibility" or "Pareta conomics, game imal decision prob- a slightly restric points of a certain i proper efficient ciency introduced 0; but the present earlier one. A tructed using the qualification, theory includes as gramming.		

Security Classification

14 KEY WORDS		LINK A		LINK B		LINK C	
		WT	ROLE	WT	ROLE	WT	
Pareto optimality Nonlinear programming Concave programming Multiple criteria Vector maximization							
~							

INSTRUCTIONS

1. ORIGINATING ACTIVITY: Enter the name and address of the contractor, subcontractor, grantee, Department of Defense activity or other organization (corporate author) issuing the report.

2a. REPORT SECURITY CLASSIFICATION: Enter the overall security classification of the report. Indicate whether "Restricted Data" is included. Marking is to be in accordance with appropriate security regulations.

2b. GROUP: Automatic downgrading is specified in DoD Directive 5200.10 and Armed Forces Industrial Manual. Enter the group number. Also, when applicable, show that optional markings have been used for Group 3 and Group 4 as authorized.

3 REPORT TITLE: Enter the complete report title in all capital letters. Titles in all cases should be unclassified. If a meaningful title cannot be selected without classification, show title classification in all capitals in parenthesis immediately following the title.

4. DESCRIPTIVE NOTES: If appropriate, enter the type of report, e.g., interim, progress, summary, annual, or final. Give the inclusive dates when a specific reporting period is covered.

5. AUTHOR(S): Enter the name(s) of author(s) as shown on or in the report. Enter last name, first name, middle initial. If military, show rank and branch of service. The name of the principal author is an absolute minimum requirement.

ti. REPORT DATE. Enter the date of the report as day, munth, year, or month, year. If more than one date appears in the report, use date of publication.

7.6. TOTAL NUMBER OF PAGES: The total page count should follow normal pagination procedures, i.e., enter the number of pages containing information.

76 NUMBER OF REFERENCES. Enter the total number of references cited in the report.

8.7 CONTRACT OR GRANT NUMBER: If appropriate, enter the applicable number of the contract or grant under which the report was written.

86. 8. 8. BROJECT NUMBER: Enter the appropriate additary department identification, such as project number, subproject number, system numbers, task number, etc.

9a. ORIGINATOR'S REPORT NUMBER(S): Enter the official report number by which the document will be identified and controlled by the originating activity. This number must be unique to this report.

9b. OTHER REPORT NUMBER(S): If the report has been assigned any other report numbers (either b): the originator or by the sponsor), also enter this number(s).

10. AVAILABILITY 'LIMITATION NOTICES: Enter any limitations on further dissemination of the report, other than those

imposed by security classification, using standard statements such as:

- (1) "Qualified requesters may obtain copies of this report from DDC."
- (2) "Foreign announcement and dissemination of this report by DDC is not authorized."
- (3) "U. S. Government agencies may obtain copies of this report directly from DDC. Other qualified DDC users shall request through

(4) "U.S. military agencies may obtain copies of this report directly from DDC. Other qualified users shall request through

(5) "All distribution of this report is controlled. Qualified DDC users shall request through

If the report has been turnished to the Office of Technical Services, Department of Commerce, for sale to the public, indicate this fact and enter the price, if known.

11. SUPPLEMENTARY NOTES. Use for additional explanatory notes.

12. SPONSORING MILITARY ACTIVITY: Enter the name of the departmental project office or laboratory sponsoring (paying for) the research and development. Include address.

13 ABSTRACT. Enter an abstract giving a brief and factual summary of the document indicative of the report, even though it may also appear elsewhere in the body of the technical report. If additional space is required, a continuation sheet shall be attached

It is highly desirable that the abstract of classified reports be unclassified. Each paragraph of the abstract shall end with an indication of the military security classification of the information in the paragraph represented as TS = (S) + (C) or (U)

There is no limitation on the length of the abstract. However, the suggested length is from 150 to 225 words

14 KEY WORDS: Key words are technically meaningful terms or short phrases that characterize a report and may be used as index entries for cataloging the report. Key words must be selected so that no security classification is required. Identifiers, such as equipment model designation, trade name, military project code name, geographic location, may be used as key words but will be followed by an indication of technical context. The assignment of links, roles, and weights is optional

Security Classification