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NOTES IN THE THEORY OF  
DYNAMIC PROGRAMMING - IV:  
A VARIATIONAL PROBLEM  
WITH CONSTRAINTS

by

Richard Bellman

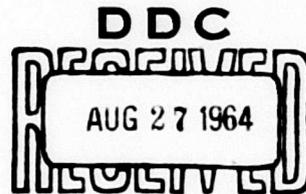
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SUMMARY

✓ The purpose of this paper is to illustrate how the functional equation technique of the theory of dynamic programming may be employed to treat a class of variational problems with constraints. ( ) ↖

NOTES IN THE THEORY OF DYNAMIC PROGRAMMING - IV:  
A VARIATIONAL PROBLEM WITH CONSTRAINTS

by

Richard Bellman

§1. Introduction

We have, in another place, [5], considered the problem of maximizing the functional

$$(1.1) \quad J(y) = \int_0^T F(x,y)dt$$

subject to the constraints

$$(1.2) \quad \begin{aligned} \text{a.} \quad & dx/dt = G(x,y), \quad x(0) = c, \\ \text{b.} \quad & 0 \leq y \leq x, \end{aligned}$$

using the classical techniques of the calculus of variations, with modifications imposed by the constraint (1.2b).

We now wish to treat the problem using the approach of the theory of dynamic programming. Expositions of the techniques we shall employ may be found in [1], [2] or [3], and we refer to these works for more detailed discussion.

The basic idea is to consider a variational problem as a multi-stage decision process of continuous type. The emphasis will now be upon determining  $y(0)$  as a function of  $c$  and  $T$  above, rather than determining  $y$  as a function of  $t$  for  $0 \leq t \leq T$ .

We shall show how we may obtain a functional equation for

$$U(c,T) = \text{Max}_y J(y),$$

and how a limiting form of this functional equation may be used to determine the structure of the extremal curve given certain simple structural properties of  $F(x,y)$  and  $G(x,y)$ .

In order to establish these results rigorously, we consider first the finite analogue of the variational problem above, namely, the problem of maximizing

$$(1.3) \quad J(\{y_k\}) = \sum_{k=0}^N F(x_k, y_k),$$

where the  $y_k$  are subject to the constraints

$$(1.4) \quad \begin{aligned} \text{a.} \quad & x_{k+1} - x_k = G(x_k, y_k), \quad k = 0, 1, 2, \dots, N-1, \\ \text{b.} \quad & 0 \leq y_k \leq x_k. \end{aligned}$$

This we again attack by means of the functional equation approach, employing a technique we have used repeatedly, cf. [1], [2], [3]. This approach is particularly suitable for a machine computation of the problem.

## §2. The Basic Functional Equation

Let us make the assumption that the solution to the variational problem described in (1.1) and (1.2) exists and that  $y(t)$  is a continuous function of  $t$  in  $[0, T]$  for every  $T \geq 0$ .

We then define the function

$$(2.1) \quad U(c, T) = \underset{y}{\text{Max}} J(y).$$

Let us now proceed formally to obtain a partial differential equation for  $U$  on the assumption that  $U$  has continuous partials with respect to  $c$  and  $T$ .

We have, along an extremal,

$$(2.2) \quad U(c, S+T) = \int_0^S F(x, y) dt + \int_S^{S+T} F(x, y) dt.$$

Hence, employing the "principle of optimality", the section of the function  $y(t)$  in  $[0, S]$  is determined by the equation

$$(2.3) \quad U(c, S+T) = \underset{y[0, S]}{\text{Max}} \left[ \int_0^S F(x, y) dt + \int_S^{S+T} F(x, y) dt \right].$$

Consider the second integral over  $[S, S+T]$ . At  $t = S$ , the value of  $x$  is  $x(S)$  as determined by the differential equation of (1.2a). Since the problem is independent of the starting time, we see that

$$(2.4) \quad \int_S^{S+T} F(x, y) dt = U(c(S), T).$$

Hence, (2.3) becomes

$$(2.5) \quad U(c, S+T) = \text{Max}_{y[0, S]} \left[ \int_0^S F(x, y) dt + U(c(S), T) \right].$$

If we now let  $S \rightarrow +0$  and assume the continuity of all functions appearing, we obtain in the limit the partial differential equation

$$(2.6) \quad U_T = \text{Max}_{0 \leq v \leq c} [ F(c, v) + G(c, v)U_c ],$$

where we set  $v = y(0)$  for typographical convenience. The condition  $0 \leq v \leq c$  is the consequence of the restriction  $0 \leq y \leq x$ .

### §3. Heuristic Considerations

Let us now see what we can deduce from (2.6) under the assumption that  $U_c$  is a positive, increasing function of  $T$ . This is a very natural condition to expect to be satisfied in a number of problems of engineering and economic origin.

We first of all impose the conditions that  $F(x, y)$  and  $G(x, y)$  be concave functions of  $y$  for all  $x \geq y \geq 0$ . It follows that the function

$$(3.1) \quad K(v) = F(c, v) + G(c, v)U_c$$

for any fixed values of  $c$  and  $U_c$  will be a concave function and possess a single maximum, which may be at  $v = 0$ ,  $v = c$ , or in between.

If we assume that  $U_c$  is a continuous function of  $T$ , as we have every right to expect, we obtain the important result that

as  $T$  varies, a region where  $v = 0$  can never abut a region where  $v = c$ . There must always be a transition region where  $0 < v < c$ , a region we call the Euler region since the Euler equation must be valid there, a region of free variation.

The critical value of  $v$  is the point where  $K'(v) = 0$ , or

$$(3.2) \quad F_v + G_v U_c = 0.$$

Let us make the further assumption that  $G_v > 0$  for all  $v$  in the range  $0 \leq v \leq c$  for any  $c > 0$ . If we assume that  $U_c \rightarrow \infty$  as  $T \rightarrow \infty$ , again a physically plausible result in many of these problems, it follows that, for sufficiently large  $T$ ,  $K'(v) = F_v + G_v U_c$  will be positive for  $0 \leq v \leq c$ . The maximum will then occur at  $v = c$ .

It is easy to show that  $U_c \rightarrow 0$  as  $T \rightarrow 0$ . If we assume that  $F_v < 0$  for all  $0 \leq v \leq c$ , as is the case in the case  $F(x,y) = x - y$  treated below, we see that  $K'(v) < 0$  for small  $T$  and hence the maximum is at  $v = 0$ .

Under the assumption that  $U_c$  is positive and monotone increasing in  $T$ , we see that the solution will be the following:

$$(3.3) \quad \begin{array}{ll} v = 0, & 0 \leq c \leq T_1 \\ 0 < v < c, & T_1 < T < T_2 \\ v = c, & T_2 \leq T, \end{array}$$

where the transition points  $T_1$  and  $T_2$ , in general, will be functions of  $c$ , and may degenerate, i.e.  $T_1 = 0$ ,  $T_1 = T_2$  or  $T_2 = T$ .

We see then the utility of the functional equation as a heuristic device for obtaining information concerning the structure of the solution.

To prove these results, we have two avenues open to us. We may first of all use the method of successive approximations on (2.6), obtaining a sequence of functions  $U_n(c, T)$  in the following way. A choice of  $v_0(c, T)$  subject to the inequality  $0 \leq v_0(c, T) \leq c$ , yields a function  $U_0(c, T)$  determined by the partial differential equation

$$U_{0T} = F(c, v_0) + G(c, v_0)U_{0c},$$

$$U_0(c, 0) = 0 \text{ for } c \geq 0.$$

Having obtained  $U_0$ , we choose a new  $v_1(c, T)$  as the function yielding the maximum of

$$K(v, U_0) = F(c, v) + G(c, v)U_{0c}$$

subject to  $0 \leq v \leq c$ . The choice of  $v_1$  yields in turn a function  $U_1$  as the solution of

$$U_{1T} = K(v_1, U_1)$$

$$U_1(c, 0) = 0 \text{ for } c \geq 0.$$

We now continue in this fashion.

We shall in a subsequent paper, in a series of which [4] is the first, consider in detail the existence and uniqueness of solutions of functional equations of the form of (2.6) and the convergence of the successive approximations above.

Here we shall follow a different path and consider the discrete multi-stage process corresponding to the variational problem. As we noted above, a derivation of these results employing classical techniques is contained in [5].

#### §4. Finite Version

Replacing the integral by a Riemann sum and the differential equation by a difference equation, we are led in a natural way to consider the problem of determining the maximum of

$$(4.1) \quad J(\{y_k\}) = \sum_{k=0}^N F(x_k, y_k),$$

where  $x_k$  and  $y_k$  are related by the equations

$$(4.2) \quad \begin{aligned} x_{k+1} - x_k &= G(x_k, y_k), \quad k = 0, 1, 2, \dots, N-1, \\ x_0 &= c, \end{aligned}$$

and  $y_k$  is restricted by the condition,

$$(4.3) \quad 0 \leq y_k \leq x_k, \quad k = 0, 1, 2, \dots, N.$$

We assume that  $F$  and  $G$  are continuous functions of  $x$  and  $y$ , and set

$$(4.4) \quad U_N(c) = \text{Max}_{\{y_k\}} J(\{y_k\}) .$$

The basic functional equation is now

$$(4.5) \quad U_{N+1}(c) = \text{Max}_{0 \leq v \leq c} [F(c, v) + U_N(c + G(c, v))], \quad N = 0, 1, 2, \dots,$$

using the notation  $v = y_0$ , with

$$(4.6) \quad U_0(c) = \text{Max}_{0 \leq v \leq c} F(c, v).$$

In order to simplify the analysis we shall consider only the case where

$$(4.7) \quad F(c, v) = c - v ,$$

a case of some interest in itself.

It corresponds to a multi-stage allocation process where we have a single resource, measured by  $x_k$  at time  $k$ . A certain quantity of this resource,  $y_k$ , may be used to increase  $x_k$  yielding

$$(4.8) \quad x_{k+1} = x_k + G(x_k, y_k)$$

at time  $k+1$ . The profit on the other hand will be measured by  $x_k - y_k$ . The problem is to determine the allocation policy which maximizes the total  $N$ -stage profit, cf. [1], [2], for a discussion of similar problems.

### §5. A Simple Case

Let us begin with a simple case, where the analysis will not overshadow the ideas. The discrete version is the problem of maximizing

$$(5.1) \quad J(\{y_k\}) = \sum_{k=0}^N (x_k - y_k),$$

where

$$(5.2) \quad (a) \quad x_{k+1} = x_k + b(y_k)$$

$$(b) \quad 0 \leq y_k \leq x_k, \quad k = 0, 1, 2, \dots, N.$$

Setting

$$U_N(c) = \text{Max}_{\{y_k\}} J(\{y_k\}),$$

we clearly have the functional equation

$$(5.3) \quad U_{N+1}(c) = \text{Max}_{0 \leq v \leq c} [c - v + U_N(c + b(v))], \quad N = 1, 2, \dots$$

$$U_0(c) = c.$$

Our aim is to determine the structure of the optimal policy under appropriate assumptions concerning  $b(y)$ . We shall assume

$$(5.4) \quad (a) \quad b(0) = 0, \quad b'(0) = \infty$$

$$(b) \quad b'(y) > 0, \quad b'(y) \rightarrow 0 \text{ as } y \rightarrow \infty$$

$$(c) \quad b''(y) < 0.$$

A simple function satisfying these conditions is  $b(y) = y^{1/2}$ .

We shall show that for each  $N$ , the optimal first allocation  $v_N(c)$  has the following form as a function of  $c$ ,

$$(5.5) \quad v_N(c) = c, \quad 0 \leq c \leq c_N$$

$$0 < v_N(c) < c, \quad c_N < c$$

where  $\{c_N\}$  is a sequence we shall determine inductively below. In this case, there is only one transition point.

The proof will be inductive. Let us begin with the case  $N=1$ .

We have

$$(5.6) \quad U_1(c) = \text{Max}_{0 \leq v \leq c} [c-v + U_0(c+b(v))] = \text{Max}_{0 \leq v \leq c} [G_0(v,c)] .$$

If an internal maximum occurs, it occurs at a point where

$$\frac{\partial G_0}{\partial v} = 0 .$$

This equation is

$$(5.7) \quad \frac{1}{b'(v)} = U_0'(c+b(v)) .$$

Since  $U_0(c) = c$ ,  $U_0'(c) = 1$ , equation (5.7) by virtue of the assumptions concerning  $b$  has precisely one positive root. If this root is less than  $c$ , there is an internal maximum. The critical value of  $c$  is then given by the root of

$$(5.8) \quad \frac{1}{b'(c)} = 1;$$

call this value  $c_1$ .

If  $c < c_1$ , the maximum in (5.6) will occur at  $v=c$ . If  $c > c_1$ , there will be an internal maximum at  $v_1$ .

Hence, we have

$$(5.9) \quad U_1(c) = U_0(c+b(c)), \text{ for } c \leq c_1, v = c, \\ = c - v_1 + U_0(c+b(v_1)), \text{ for } c > c_1, v = v_1 < c.$$

The function  $U_1(c)$  is clearly differentiable for  $c < c_1$  and for  $c > c_1$ . We have

$$(5.10) \quad U_1'(c) = (1+b'(c)) U_0'(c+b(c)) = 1 + b'(c), \quad c < c_1, \\ = 1 + U_0'(c+b(c)) = 2, \quad c > c_1.$$

At  $c=c_1$ ,  $b'(c) = 1$ , and we have equality. Hence,  $U_1(c)$  has a continuous derivative. Furthermore,  $U_1'(c) > U_0'(c)$  for  $c \geq 0$ .

Let us now investigate convexity. We have

$$(5.11) \quad U_1''(c) = b''(c), \quad c < c_1 \\ = 0, \quad c > c_1$$

Hence  $U_1''(c) \leq 0$  for all  $c > 0$ .

We now turn to the case  $N = 2$ .

$N = 2$ . We have

$$(5.12) \quad U_2(c) = \text{Max}_{0 \leq v \leq c} [c - v + U_1(c+b(v))].$$

If an internal maximum occurs, it occurs at the point where

$$(5.13) \quad \frac{1}{b'(v)} = U_1'(c+b(v)).$$

Since  $U_1''(c+b(v)) \leq 0$  and  $(\frac{1}{b'(v)})' > 0$ , there is precisely one root of this equation, which we call  $v_2 = v_2(c)$ . Since  $U_1' > U_0'$ , it follows that  $v_2(c) > v_1(c)$ .

As above, the critical value of  $c$  is the root of

$$(5.14) \quad \frac{1}{b'(c)} = U_1'(c+b(c)).$$

Call this root  $c_2$ . Since  $U_1' > U_0'$  and  $(1/b')' > 0$ ,  $c_2 > c_1$ .

We thus have

$$(5.15) \quad U_2(c) = U_1(c+b(c)), \quad \text{for } c \leq c_2 \\ = c - v_2 + U_1(c+b(v_2)), \quad \text{for } c \geq c_2 .$$

Furthermore, we have

$$(5.16) \quad U_2'(c) = [1+b'(c)]U_1'(c+b(c)), \quad \text{for } c < c_2 \\ = 1+U_1'(c+b(v_2)) + \frac{dv_2}{dc} [-1+b'(v_2)U_1'(c+b(v_2))] \\ = 1+U_1'(c+b(v_2)) = 1+1/b'(v_2), \quad \text{for } c > c_2 .$$

At  $c=c_2$ ,  $U_2'(c)$  is continuous, recalling (5.14).

Let us now examine the convexity of  $U_2(c)$ .

$$(5.17) \quad U_2''(c) = b''(c)U_1'(c+b(c)) + [1+b'(c)]^2 U_1''(c+b(c)), \quad \text{for } c < c_2 \\ = U_1''(c+b(v_2)) [1+b'(v_2)\frac{dv_2}{dc}], \quad \text{for } c > c_2 .$$

Using (5.13),

$$(5.18) \quad -\frac{b''(v_2)}{[b'(v_2)]^2} \frac{dv_2}{dc} = U_1''(c+b(v_2)) [1+b'(v_2)\frac{dv_2}{dc}] ,$$

or

$$(5.19) \quad \frac{dv_2}{dc} [b'(v_2)U_1''(c+b(v_2)) + \frac{b''(v_2)}{[b'(v_2)]^2}] = -U_1''(c+b(v_2)).$$

Hence,  $\frac{dv_2}{dc} < 0$ . Returning to (5.18) this yields  $1+b'(v_2)\frac{dv_2}{dc} > 0$ . Using this in (5.17),  $U_2''(c) < 0$  for  $c > c_2$ . Since (5.17) shows that  $U_2''(c) < 0$  for  $0 \leq c < c_2$ , we see that  $U_2''(c) < 0$  for all  $c \geq 0$ , but is not continuous at  $c = c_2$ .

Now to the final step that

$$(5.20) \quad U_2'(c) > U_1'(c) .$$

Having established this, we have all the ingredients of an inductive proof.

To establish this inequality, we must consider three distinct intervals  $[0, c_1]$ ,  $[c_1, c_2]$ , and  $[c_2, \infty]$ .

In  $[0, c_1]$ , we have

$$(5.21) \quad U_2'(c) = [1 + b'(c)] U_1'(c+b(c)) > [1 + b'(c)] U_0'(c+b(c)) = U_1'(c) .$$

In  $[c_2, \infty]$ , we have

$$(5.22) \quad U_2'(c) = 1 + 1/b'(v_2) > 1 + 1/b'(v_1) = U_1'(c) ,$$

Since  $v_2 > v_1$ . The remaining interval to consider is  $[c_1, c_2]$ . In  $[c_1, c_2]$  we have

$$(5.23) \quad U_1'(c) = 1 + 1/b'(v_1)$$

$$U_2'(c) = [1 + b'(c)] U_1'(c + b(c)) .$$

Since  $v_2 = c$  in  $[c_1, c_2]$ ,

$$(5.24) \quad -1 + b'(v)U'_1(c + b(v)) \geq 0$$

for  $0 \leq v \leq c_2$ . In particular,

$$(5.25) \quad b'(c)U'_1(c + b(c)) \geq 1.$$

We now wish to show that

$$(5.26) \quad U'_1(c + b(c)) \geq 1/b'(v_1)$$

for  $c_1 \leq c \leq c_2$ . Since  $1/b'(v)$  is increasing and  $v_1(c) \leq c$ ,  $1/b'(v_1) \leq 1/b'(c)$ . By (5.25) it follows that, for  $c \leq c_2$ ,

$$(5.27) \quad U'_1(c + b(c)) \geq 1/b'(c) \geq 1/b'(v_1) .$$

We now have all the material required for an inductive proof of the following

(5.28) Theorem. For each  $N$  there exists a function  $v_N(c)$  with the following properties:

- (a)  $v_N(c)$  is monotone decreasing as  $c$  increases;
- (b)  $v_{N+1}(c) > v_N(c)$ ,  $N = 1, 2, \dots$
- (c) There is a unique solution of  $v_N(c) = c$  which we call  $c_N$ , and  $c_{N+1} > c_N$ ;
- (d) For  $0 \leq c \leq c_N$ , we have  $U_N(c) = U_{N-1}(c + b(c))$ ;
- (e) For  $c_N \leq c$ , we have  $U_N(c) = c - v_N(c) + U_{N-1}[c + b(v_N(c))]$ ;
- (f)  $U'_N(c) \geq U'_{N-1}(c)$ ,  $N = 1, 2, \dots$ ,  $c \geq 0$ ;

$$(g) \quad U_N^*(c) \leq 0, \quad N = 1, 2, \dots, \quad c \geq 0 \quad .$$

§6. A More General Problem

Let us now consider the problem of maximizing

$$(1) \quad J(y) = \int_0^T (x-y) dt$$

subject to the relation

$$(2) \quad \frac{dx}{dt} = G(x, y), \quad x(0) = c \quad ,$$

and the constraint

$$(3) \quad 0 \leq y \leq x \quad .$$

As above, we begin by considering the discrete version of the problem, where

$$(4) \quad J[\{y_k\}] = \sum_{k=0}^N (x_k - y_k) \quad ,$$

and

$$(5) \quad (a) \quad x_{k+1} = x_k + G(x_k, y_k), \quad k = 0, 1, 2, \dots, N-1,$$

$$(b) \quad 0 \leq y_k \leq x_k \quad .$$

Setting

$$(6) \quad U_N(c) = \text{Max}_y J[\{y_k\}]$$

we clearly have

$$(7) U_1(c) = c ,$$

$$U_{N+1}(c) = \text{Max}_{0 \leq v \leq c} [ c-v + U_N(c+G(c,v)) ] , \quad N = 1,2,\dots$$

Our aim is to determine the structure of the optimal policy under appropriate assumptions concerning  $G(c,v)$ . This is equivalent to determining the structure of  $v = v_N(c)$  as a function of  $c$  and  $N$ .

We shall assume that  $G(c,v)$  satisfies the following conditions

- (8) (a)  $G_v(c,v) > 0$ ,  $G_v(c,v) \rightarrow 0$  as  $v \rightarrow \infty$ ,  $G_v(c,v) \rightarrow \infty$  as  $v \rightarrow 0$ , uniformly in  $c$ .
- (b)  $G_c > 0$ ,  $G_{cv} \leq 0$ .
- (c)  $[ 1/G_v(c,v) ]_{v=c}$  is monotone increasing in  $c$ .
- (d)  $r^2 G_{cc} + 2rs G_{cv} + s^2 G_{vv}$  is a negative definite form.
- (e)  $[ 1 + G_c(c,v) ] / G_v(c,v)$  is monotone increasing in  $v$ .\*

As above, we shall employ an inductive approach.

### §7. The Case $N=2$

We have

$$(1) U_2(c) = \text{Max}_{0 \leq v \leq c} [ c-v + U_1(c+G(c,v)) ]$$

Let us set

$$(2) F_1(v,c) = c-v + U_1(c+G(c,v)) .$$

If an internal maximum occurs in (1), it occurs at a point where

\* In the limit of the continuous case, this requires that  $1/G_v(c,v)$  be monotone increasing in  $v$ , which is a consequence of (d).

$$(3) \quad \frac{\partial F_1}{\partial v} = 1 \quad ,$$

or

$$(4) \quad 1/G_v(c,v) = U'_1(c+G(c,v)) = 1$$

By virtue of our assumptions concerning  $G$ , this equation has precisely one root for any value of  $c$ . If this root is less than  $c$ , there is an internal maximum.

The critical value of  $c$  is then given by the root of

$$(5) \quad [1/G_v(c,v)]_{v=c} = 1 \quad .$$

Let us call this value  $c_1$ . Our assumption ensures that this equation has precisely one root.

If  $c < c_1$ , there will be no internal maximum and the maximum will occur at  $v = c$ . If  $c > c_1$ , there will be an internal maximum at a value  $v < c$  which we call  $v_1 = v_1(c)$ .

Hence we have

$$(6) \quad U_2(c) = U_1(c+G(c,c)) \quad , \quad c \leq c_1 \quad ,$$

$$= c - v_1 + U_1(c+G(c,v_1)) \quad , \quad c \geq c_1 \quad ,$$

The function  $f_2(c)$  is clearly differentiable for  $c < c_1$  and  $c > c_1$ . We have

$$(7) \quad U'_2(c) = 1 + (G_c(c,c) + [G_v(c,v)]_{v=c}) \quad , \quad c < c_1$$

$$= 1 + (1 + G_c(c,v_1)) \quad , \quad c > c_1 \quad .$$

For further use, we note that (4) shows that we may write

$$(8) \quad U'_2(c) = 1 + (1 + G_c(c,v_1))/G_v(c,v_1) \quad , \quad c > c_1 \quad .$$

At  $c = c_1$ ,  $v_1(c) = c$  and  $[G_v(c, v)]_{v=c} = 1$ . Hence we have equality of right and left hand derivatives of  $U_2(c)$  at  $c = c_1$ . Thus  $U_2'(c)$  is continuous for all values of  $c$ .

Since  $U_1'(c) = 1$ , (7) shows that

$$(9) \quad U_2'(c) > U_1'(c)$$

for all  $c \geq 0$ .

We now wish to show that  $U_2'(c)$  is monotone increasing, which is to say that  $U_2(c)$  is concave. To accomplish we shall prove a useful lemma.

### 58. A Lemma on Concavity

Let us prove

Lemma 1. Let  $\phi(x, y)$  be a concave function of  $x, y$  for  $x, y \geq 0$ ,

$$(1) \quad \phi(\lambda x_1 + (1-\lambda)x_2, \lambda y_1 + (1-\lambda)y_2) \geq \lambda \phi(x_1, y_1) + (1-\lambda)\phi(x_2, y_2)$$

for  $0 \leq \lambda \leq 1$ .

Then

$$(2) \quad f(x) = \text{Max}_{0 \leq y \leq x} \phi(x, y)$$

is concave.

Proof: If  $y = x$ , we have  $f(x) = \phi(x, x)$ , clearly concave. If  $y = 0$ ,  $f(x) = \phi(x, 0)$ , also clearly concave. If  $0 < y < x$ , then  $y$  is determined by

$$(3) \quad \phi_y(x, y) = 0 \quad .$$

Then

$$(4) \quad f'(x) = \phi_x + \phi_y \frac{dy}{dx} = \phi_x \quad .$$

Thus

$$(5) \quad f''(x) = \phi_{xx} + \phi_{xy} \frac{dy}{dx} \\ = \frac{\phi_{xx}\phi_{yy} - \phi_{xy}^2}{\phi_{yy}} \quad ,$$

using (3) which yields  $\phi_{xy} + \phi_{yy} \frac{dy}{dx} = 0 \quad .$

The condition that  $\phi(x,y)$  be concave is equivalent to

$$(6) \quad r^2\phi_{xx} + 2rs\phi_{xy} + s^2\phi_{yy}$$

being negative definite. Hence the right side of the equation in (5) is negative.

Let us now show that

$$(7) \quad F(c,v) = c-v + U(G(c,v))$$

is concave if  $U$  is concave and monotone increasing and  $G$  is concave.

We have

$$(8) \quad F_{cc} = G_{cc}U'(G) + G_c^2U''(G)$$

$$F_{vv} = G_{vv}U'(G) + G_v^2U''(G)$$

$$F_{cv} = G_{cv}U'(G) + G_cG_vU''(G)$$

From this it is clear that  $F(c,v)$  is concave.

Combining these results we see that  $F_2(c)$  is concave.

### 59. The Case N=3

We have

$$\begin{aligned} (1) \quad U_3(c) &= \text{Max}_{0 \leq v \leq c} [c-v+U_2(c+G(c,v))] \\ &= \text{Max}_{0 \leq v \leq c} [F_2(v,c)] \end{aligned}$$

If an internal maximum occurs, it occurs at a point where

$$(2) \quad 1/G_v(c,v) = U_2'(c+G(c,v))$$

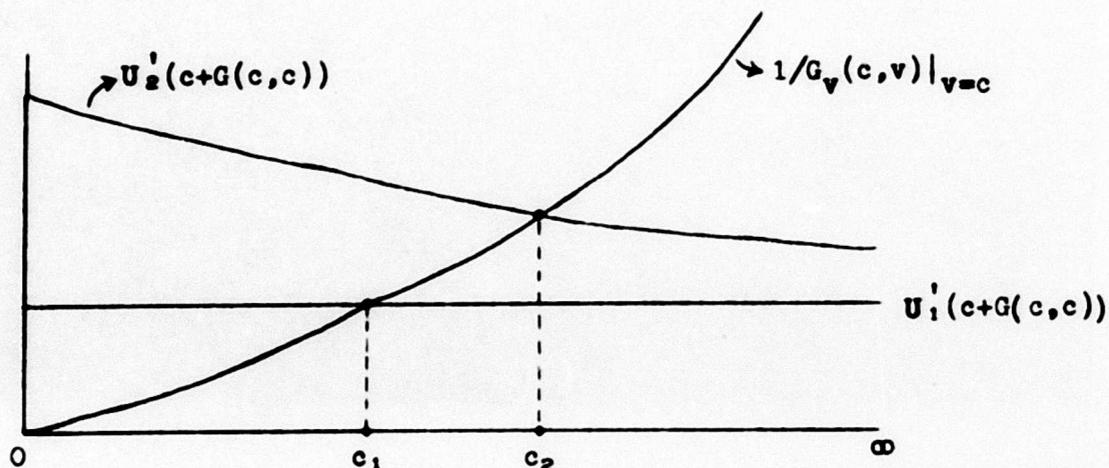
Since  $1/G_v$  is monotone increasing in  $v$ , and  $U_2'(c+G(c,v))$  is monotone decreasing in  $v$ , there is exactly one root of this equation which we call  $v_2 = v_2(c)$ . Since  $U_2' > U_1'$  it is clear that  $v_2(c) > v_1(c)$ .

The critical value of  $c$  is the root of

$$(3) \quad [1/G_v(c,v)]_{v=c} = U_2'(c+G(c,c)) .$$

Call this root  $c_2$ . It is clear that  $c_2 > c_1$ .

From (2) we see that  $dv/dc$  is negative.



Referring to the figure above it is clear that  $c_2 > c_1$ . Hence

$$(4) \quad U_3(c) = U_2(c+G(c,c)) \quad , \quad c < c_2 \\ = c-v_2+U_2(c+G(c,v_2)) \quad , \quad c > c_2 \quad .$$

Thus

$$(5) \quad U_3(c) = \{1 + [G_c(c,c) + [G_v(c,v)]_{v=c}]\} U'_2(c+G(c,c)), \quad c < c_2 \\ = 1 + (1+G_c(c,v_2))U'_2(c+G(c,v_2)), \quad c > c_2 \\ = 1 + \frac{1+G_c(c,v_2)}{G_v(c,v_2)}$$

The concavity of  $U_3(c)$  follows as before.

That  $U_3(c) > U_2(c)$  is clear for  $c > c_2$  and  $0 \leq c \leq c_1$ . It remains to prove that the inequality holds in  $[c_1, c_2]$ . After having established this, we have all the material for an inductive proof of the structure of  $v_N(c)$ .

In  $[c_1, c_2]$  we have

$$(6) \quad U'_2(c) = 1 + (1+G_c(c,v_1))U'_1(c+G(c,v_1)) \\ U_3(c) = [1 + (G_c(c,c) + [G_v(c,v)]_{v=c})] U'_2(c+G(c,c))$$

The fact that  $v_2 = c$  in  $[c_1, c_2]$  implies that

$$(7) \quad -1 + G_v U'_2(c+G) \geq 0$$

for  $0 \leq v \leq c$  and  $c_1 \leq c \leq c_2$ . Hence

$$(8) \quad U_3(c) \geq \frac{[1 + G_c(c,c) + [G_v(c,v)]_{v=c}]}{[G_v(c,v)]_{v=c}}$$

$$\geq 1 + \frac{1 + G_c(c, c)}{[G_v(c, v)]_{v=c}}$$

On the other hand

$$(9) \quad U_2^1(c) = 1 + \frac{1 + G_c(c, v_1)}{G_v(c, v_1)}$$

Since  $v_1(c) \leq c$  for  $c_1 \leq c < \infty$ , and  $[1 + G_c(c, v)]/G_v(c, v)$  a function of  $v$  is monotone increasing, by assumption, we see that  $U_3^1(c) > U_2^1(c)$ . We now have all the material required for an inductive proof.

Summarizing our results we see that we have a monotone increasing sequence,  $\{c_N\}$ , possessing the property that

$$(10) \quad \begin{aligned} v_N = c, \quad 0 \leq c \leq c_N \\ = v_N(c) < c, \quad c_N < c < \infty. \end{aligned}$$

Each function  $U_N^1(c)$  is monotone increasing in  $c$  and concave and the sequence  $\{U_N^1(c)\}$  is monotone increasing in  $N$ .

This property carries over in the limit as the discrete process goes over to the continuous and enables us to use the functional equation to determine the properties of the solution. It may be proved by a straightforward argument that the limit of the discrete case is indeed the continuous. We shall omit the proof here.

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