BAYESIAN SOFTWARE PREDICTION MODELS. Volume I.
An Imperfect Debugging Model for Reliability and other Quantitative Measures of Software Systems

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In this report a stochastic model for software failure phenomena is developed for the case when the errors are not corrected with certainty. Expressions for several quantities of interest are derived to establish quantitative measures for software performance assessment. Approximations for large-scale software systems using a gamma distribution are also discussed. Numerical examples are used to illustrate the computations and usefulness of various quantities. Volume V will be published at a later date.
Block 7.

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</table>
EVALUATION

The necessity for more complex software systems in such areas as command and control and avionics has led to the desire for better methods for predicting software errors to ensure that software produced is of higher quality and of lower cost. This desire has been expressed in numerous industry and Government sponsored conferences, as well as in documents such as the Joint Commanders' Software Reliability Working Group Report (Nov 1975). As a result, numerous efforts have been initiated to develop and validate mathematical models for predicting such quantities as the number of remaining errors in a software package, the time to achieve a desired reliability level, and a measure of the software reliability. However, early efforts have not produced models with the desired accuracy of prediction and with the necessary confidence limits for general model usage.

This effort was initiated in response to this need for developing better and more accurate software error prediction models and fits into the goals of RADC TPO No. 5, Software Cost Reduction (formerly RADC TPO No. 11, Software Sciences Technology), in the subthrust of Software Quality (Software Modeling). This report summarizes the development of a mathematical model for predicting quantities such as the expected number of remaining errors, achieved reliability, and time to detect and correct a specified number of errors that assumes a software error is not corrected at a given time with probability 1 (i.e., imperfect debugging). The importance of this development is that it represents the first attempt to develop software error prediction models that incorporate imperfect debugging, and thus more closely reflect the actual software error detection and correction process.

The theory and equations developed under this effort will lead to much needed predictive measures for use by software managers in more accurately tracking software development projects in terms of test time needed to achieve given reliability and error objectives. In addition, the associated confidence limits and other related statistical quantities developed under this effort will insure more widespread use of these modeling techniques. Finally, the predictive measures and equations developed under this effort will be applicable to current Air Force software development projects and thus help to produce the high quality, low cost software needed for today's systems.

ALAN N. SUKERT
Project Engineer
1. INTRODUCTION

A considerable emphasis has been placed in recent years on the study of software error phenomena with the objective of developing analytical models which can be used to obtain quantitative measures for software performance. Most of these studies assume an exponential distribution for times between software errors with a failure rate that depends on the number of remaining errors, see for example, [3, 6, 8, 9, 10, 11, 13, 15, 18].

A key assumption made in most of these studies is that the errors are removed with certainty, when detected. However, as pointed out in Miyamoto [7] and Thayer et al. [15], in practice errors are not always corrected when detected. The existing models do not provide a solution for such situations. The purpose of this report, then, is to develop an analytical model for software error phenomenon when the errors are not removed/corrected with certainty, i.e., for the case of imperfect debugging. The model is developed in Section 2 and expressions for the following quantities of interest are derived in Section 3:

(i) Distribution of time to a completely debugged system.
(ii) Distribution of time to a specified number of remaining errors.
(iii) Distribution of number of remaining errors.
(iv) Expected number of errors detected by time $t$.

The distribution of time between software failure is obtained in Section 4 and approximate solutions using a gamma distribution are discussed in Section 5. Numerical examples are used to illustrate the computations and usefulness of various quantities.
2. MODEL DEVELOPMENT

The following assumptions are made for developing the model.

(i) The error causing a software failure, when detected, is corrected with probability \( p \) (\( 0 \leq p \leq 1 \)), while with probability \( q \) with \( p + q = 1 \) we fail to completely remove it. Thus, \( q \) is the probability of imperfect debugging.

(ii) Errors in the software package are independent of each other and have a constant occurrence rate \( \lambda \).

(iii) The probability of two or more errors occurring simultaneously is negligible.

(iv) The time to remove an error is considered to be negligible in this model.

(v) No new errors are introduced during the debugging process.

(vi) At most one error is removed at correction time.

Let \( X(t) \) denote the number of errors remaining in the package at time \( t \). We will use this random variable to describe the state of the error process at time \( t \). Further, let \( N \) be the number of errors at the beginning of the debugging phase, i.e., \( X(0) = N \).

Suppose that there are \( i \) errors in the package at some time. Then from assumption (i), we note that after the occurrence of the next failure

\[
X(t) = \begin{cases} 
    i-1 & \text{with probability } p \\
    i & \text{with probability } q 
\end{cases} \quad (2.1)
\]

In other words, if we were to observe the \( X(t) \) process at times of software failures, then its behavior is governed by equation (2.1). The transition probabilities \( P_{ij} \) from state \( i \) to state \( j \),
\[ \begin{array}{cccccccccc}
0 & 1 & 2 & \cdots & i-2 & i-1 & i & \cdots & N-2 & N-1 & N \\
0 & 1 & 0 & 0 & 0 & \cdots & 0 & 0 & \cdots & 0 & 0 \\
1 & p & q & 0 & 0 & \cdots & 0 & 0 & \cdots & 0 & 0 \\
2 & 0 & p & q & 0 & \cdots & 0 & 0 & \cdots & 0 & 0 \\
\vdots & \vdots & \vdots & \vdots & \vdots & \vdots & \vdots & \vdots & \vdots & \vdots & \vdots \\
N & 0 & 0 & \cdots & 0 & \cdots & 0 & \cdots & p & q & 0 \\
N & 0 & 0 & \cdots & 0 & \cdots & 0 & \cdots & 0 & p & q \\
\end{array} \]

A diagrammatic representation of transitions between states corresponding to equation (2.2) is given in Figure 2.1.

Now, assumptions (i) and (ii) imply that the times between successive software failures (error occurrences) follow an exponential distribution. Suppose at some time \( t = \tau, x(\tau) = i \), \( i = 0, 1, \ldots, N \). Then the probability density function (pdf) \( f_1(t) \) of the time to next failure, \( T_1 \), is given by the distribution of the first order statistic of \( i \) exponential distributions each with parameter \( \lambda \), i.e.,

\[ f_i(t) = \binom{i}{1} \lambda e^{-\lambda t} (e^{-\lambda t})^{i-1} \]

or

\[ f_i(t) = i\lambda e^{-i\lambda t} \]

and the cumulative distribution function (cdf) is given by

\[ F_i(t) = 1 - e^{-i\lambda t} \].
Figure 2.1  A Diagrammatic Representation of Transitions Between States of $X(t)$
We note that even though the stochastic process $X(t)$ makes transitions from state to state in accordance with equation (2.2), the times spent in various states are random and are given by equation (2.3). Hence $\{X(t), t \geq 0\}$ forms a semi-Markov process. A typical realization of this process is shown in Figure 2.2. It should be pointed out that in our formulation the process $X(t)$ undergoes both real and virtual transitions. This means that after an attempt to remove an error the state of $X(t)$ may change or may remain unchanged. In Figure 2.2 real transitions occur at states $N, N-2$ and $i$ while a virtual transition occurs at state $N-1$.

Let $Q_{ij}(t)$ denote the one step transition probability that after making a transition into state $i$, the process $X(t)$ next makes a transition into state $j$ by time $t$. In other words if a software package has $i$ remaining errors at time zero, then $Q_{ij}(t)$ represents the probability that the next failure, resulting in $j$ remaining errors, will be by time $t$. Hence, for $i, j = 0, 1, 2, \ldots, N$, we can write

$$Q_{ij}(t) = \int_0^t P[X(u) = j, T_i = u | X(0) = i] \cdot du.$$  

Since the events $\{X(u)=j\}$ and $\{T_i=u\}$ are independent, we get

$$Q_{ij}(t) = \int_0^t P[X(u)=j | X(0)=i] \cdot P[T_i=u | X(0)=i] \cdot du$$

$$= \int_0^t P_{ij} \cdot P[T_i=u | X(0)=i] \cdot du$$

$$= P_{ij} \int_0^t \lambda \cdot e^{-i \lambda u} \cdot du$$

$$= P_{ij} \cdot \left[1 - e^{-i \lambda t}\right]$$

\[5\]
Figure 2.2  A Typical Realization of the X(t) Process
or \[ Q_{ij}(t) = P_{ij} \cdot F_i(t) \] (2.5)

for \( i, j = 0, 1, 2, \ldots, N \).

It is obvious that \( Q_{ij}(t) \) must satisfy

\[ Q_{ij}(t) \geq 0, \quad i, j = 0, 1, 2, \ldots, N, \quad t \geq 0 \]

and

\[ \sum_{j=0}^{N} Q_{ij}(t) = p_{i+j} = 1, \quad i = 0, 1, \ldots, N. \]

The probabilities \( Q_{ij}(t) \) are obtained by multiplying the probabilities \( P_{ij} \) from (2.2) and \( F_i(t) \) from (2.4). Thus, for example,

\[ Q_{N,N-1}(t) = P_{N,N-1} \cdot F_N(t) \]

or

\[ Q_{N,N-1}(t) = p(1 - e^{-N\lambda t}). \]

Proceeding similarly for all \( i, j \) we get \( Q_{ij}(t) \) as shown in Equation (2.6) on the following page.

For known parameters \( N, p \) and \( \lambda \), the probabilities \( Q_{ij}(t) \) are obtained from Equation (2.6). This equation represents the basic model that will be used in the following sections for obtaining the various quantities of interest for the software error phenomenon.
\[ \{ Q_{ij}(t) \} = \begin{bmatrix}
0 & 1 & 2 & \ldots & N-2 & N-1 & N \\
0 & 1 & 0 & 0 & \ldots & 0 & 0 & 0 \\
p_{F_1}(t) & q_{F_1}(t) & 0 & \ldots & 0 & 0 & 0 \\
0 & p_{F_2}(t) & q_{F_2}(t) & \ddots & \vdots & \vdots & \vdots \\
0 & 0 & \ddots & \ddots & p_{F_{N-1}}(t) & q_{F_{N-1}}(t) & 0 \\
0 & 0 & \ddots & 0 & p_{F_N}(t) & q_{F_N}(t) & \end{bmatrix} \quad (2.6) \]
3. DERIVATION OF VARIOUS QUANTITIES OF INTEREST

3.1 Distribution of Time to a Completely Debugged Software System

Suppose \( i \) is the number of errors remaining in a software system at some time during the debugging process. Let \( g_{i,0}(t) \) and \( G_{i,0}(t) \) denote the pdf and cdf, respectively, of the first passage time from \( i \) to 0. In other words \( g_{i,0}(t) \) and \( G_{i,0}(t) \) represent, respectively, the pdf and cdf of the time required to obtain a completely debugged software system when the initial number of errors is \( i \).

Recall that at time zero, \( X(0) = N \) and at the time of the next failure

\[
X(t) = \begin{cases} 
N-1 & \text{with probability } p \\
N & \text{with probability } q
\end{cases}
\]  

(3.1)

as shown in Figure 2.1. Now, from the definition of \( Q_{i,j}(t) \), the probability of going from \( N \) to \( N-1 \) errors in time \([u, u+du]\) is \( dQ_{N,N-1}(u) \). Then the process \( X(t) \) restarts with \( (N-1) \) remaining errors at time \( u \) and the cdf of the first passage time is \( G_{N-1,0}(t-u) \). For the case of perfect debugging the cdf of the first passage time is

\[
\int_0^t G_{N-1,0}(t-u) \cdot dQ_{N,N-1}(u) = Q_{N,N-1} \ast G_{N-1,0}(t),
\]

(3.2)

where \( \ast \) denotes convolution.

Similarly, if the debugging at the first error occurrence is imperfect, the cdf of the first passage time is
Since the events depicted in Equations (3.2) and (3.3) are mutually exclusive, we get the renewal equation

\[ G_{N,0}(t) = Q_{N,N-1} * G_{N-1,0}(t) + Q_{N,N} * G_{N,0}(t) \]  

(3.4)

In general, we get the renewal equation

\[ G_{i,0}(t) = Q_{i,i-1} * G_{i-1,0}(t) + Q_{i,i} * G_{i,0}(t) \]  

(3.5)

for \( i = 1, 2, \ldots, N \)

where \( G_{0,0}(t) = 1 \).

We use Laplace-Stieltjes (L-S) transforms to solve renewal equations (3.5), where the L-S of \( G_{i,0}(t) \) is defined as:

\[ \tilde{G}_{i,0}(s) = \int_0^\infty e^{-st} G_{i,0}(t) \text{d}t \]  

(3.6)

From (3.5) we get

\[ \tilde{G}_{i,0}(s) = \tilde{G}_{i,i-1}(s) \tilde{G}_{i-1,0}(s) + \tilde{Q}_{i,i}(s) \tilde{G}_{i,0}(s), \quad i = 1, 2, \ldots, N \]  

(3.7)

where

\[ \tilde{Q}_{i,i-1}(s) = \frac{ip\lambda}{s + i\lambda}, \]  

(3.8)

and

\[ \tilde{Q}_{i,i}(s) = \frac{iq\lambda}{s + i\lambda}. \]  

(3.9)

Solving (3.7) recursively, we get the L-S transform of \( G_{N,0}(t) \) as

\[ \tilde{G}_{N,0}(s) = \prod_{i=1}^{N} \tilde{G}_{i,0}(s), \]  

(3.10)
\[ G_{N,0}(s) = \sum_{j=1}^{N} \frac{j\lambda}{s + j\lambda} = \sum_{j=1}^{N} C_{N,j} \frac{j\lambda}{s + j\lambda} \]  

where

\[ C_{N,j} = \binom{N}{j}(-1)^{j-1} \]

By taking the inverse of \( G_{N,0}(s) \), the cdf of the first passage time from \( N \) to 0 is:

\[ G_{N,0}(t) = \sum_{j=1}^{N} C_{N,j} \left(1 - e^{-j\lambda t}\right). \]

The pdf of the first passage time from \( N \) to 0 is given by

\[ g_{N,0}(t) = \sum_{j=1}^{N} C_{N,j} \cdot j\lambda \cdot e^{-j\lambda t}. \]

To illustrate the above result let us consider a software system with \( N = 10, \lambda = 0.02 \) and \( p = 0.8 \). Then

\[ G_{10,0}(t) = \sum_{j=1}^{10} \binom{10}{j}(-1)^{j-1} \cdot (1 - e^{-j\lambda t}(.02)t). \]

The values of this function for various \( t \) are plotted in Figure 3.1. From this plot we note that the probability of getting an error free system by 275 time units is 0.9 and by 500 units is 1.0. Such a plot is useful for calculating the time required to get an error free system with a desired probability.

Similar plots for values of \( p = .85, .90, .95 \) and 1.0 are also shown in Figure 3.1. As would be expected the cdf for a larger \( p \) dominates that for a smaller \( p \). In other words the better the debugger, the faster is the process of debugging.
Figure 3.1  CDF of Time to a Completely Debugged Software System
3.2 Distribution of Time to a Specified Number of Remaining Errors

In many instances a completely debugged software is not cost effective and we may be willing to tolerate a certain number of remaining errors, say \( n_0 \), which will ensure some desired reliability. The distribution of time to \( n_0 \) is then of interest.

Using an approach similar to that of Section 3.1 we get the renewal equation

\[
G_{i,n_0}(t) = Q_{i-1}G_{i-1,n_0}(t) + Q_iG_{i,n_0}(t),
\]

for \( i = n_0+1, \ldots, N \) \hspace{1cm} (3.14)

where \( G_{n_0,n_0}(t) = 1 \).

Then the L-S transform of \( G_{N,n_0}(t) \) is given by

\[
\mathcal{G}_{N,n_0}(s) = \prod_{j=n_0+1}^{n-n_0} \frac{s+j\lambda}{s+(n_0+j)\lambda} = \sum_{j=1}^{n-n_0} B_{N,j,n_0} \frac{(n_0+j)p\lambda}{s+(n_0+j)p\lambda}
\]

where

\[
B_{N,j,n_0} = \frac{N!}{n_0^! j! (N-n_0-j)!} (-1)^j \frac{j}{n_0+j}.
\]

The cdf is obtained by taking the inverse L-S transform of \( \mathcal{G}_{N,n_0}(s) \),

\[
G_{N,n_0}(t) = \sum_{j=1}^{n-n_0} B_{N,j,n_0} \left\{ 1 - e^{-(n_0+j)p\lambda t} \right\},
\]

and the pdf is

\[
g_{N,n_0}(t) = \sum_{j=1}^{n-n_0} B_{N,j,n_0} (n_0+j)p\lambda e^{-(n_0+j)p\lambda t}
\]

13
To see the nature of the pdf and cdf, let us consider the case when $N=10$, $\lambda = 0.02$, and $p = 0.9$. These are shown in Figures 3.2 and 3.3, respectively, for various values of $n_0$. The plots are self explanatory.

Now let a random variable $T_{N,n_0}$ denote the first passage time from $N$ to $n_0$ errors. Then, from (3.18) we can obtain the $i^{th}$ moment of $T_{N,n_0}$ as

$$E[T_{N,n_0}^i] = \sum_{j=1}^{N-n_0} B_{N,j,n_0} \frac{r(i+1)}{(n_0+j)p_i^j}. \tag{3.19}$$

From (3.19), the mean and variance are

$$ET_{N,n_0} = \sum_{j=1}^{N-n_0} B_{N,j,n_0}/(n_0+j)p_i \tag{3.20}$$

$$\text{var}(T_{N,n_0}) = ET_{N,n_0}^2 - (ET_{N,n_0})^2 \tag{3.21}$$

The values for mean and variance of first passage time for various $n_0$ are given in Table 3.1, where $N=10$, $p=0.9$ and $\lambda = 0.02$. 

Figure 3.2  PDF of Time to a Specified Number ($n_0$) of Remaining Errors
Figure 3.3  CDF of Time to a Specified Number ($n_0$) of Remaining Errors
<table>
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<th>Mean</th>
<th>Variance</th>
<th>Standard Deviation</th>
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<td>5.56</td>
<td>30.86</td>
<td>5.56</td>
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<tr>
<td>8</td>
<td>11.73</td>
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<td>8.30</td>
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<td>7</td>
<td>18.67</td>
<td>117.19</td>
<td>10.83</td>
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<td>6</td>
<td>26.61</td>
<td>180.18</td>
<td>13.42</td>
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<td>79.39</td>
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<td>30.42</td>
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<td>1696.81</td>
<td>41.19</td>
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<tr>
<td>0</td>
<td>162.72</td>
<td>4783.23</td>
<td>69.16</td>
</tr>
</tbody>
</table>

Table 3.1

Mean and Variance of First Passage Time for Various $n_o$

($N = 10, p = 0.9, \lambda = 0.02$)
3.3 Distribution of Number of Remaining Errors

First, we develop the expressions for the distribution of the number of remaining errors after a specified time period, \( t \). Then, the expected number of remaining errors at time \( t \) is obtained.

Let \( P_{N,n_0}(t) \) represent the probability that there are \( n_0 \) errors remaining in a software package at time \( t \), given that there are \( N \) errors at the beginning of debugging, i.e.,

\[
P_{N,n_0}(t) = P[X(t) = n_0 | X(0) = N]
\]  

(3.22)

which is the so-called state occupancy probability. Conditioning on the next failure and following an approach similar to that of Section 3.1, we get the following renewal equation.

\[
P_{n_0, n_0}(t) = e^{-n_0 \lambda t} + Q_{n_0, n_0} * P_{n_0, n_0}(t), \quad n_0 \leq N.
\]  

(3.23)

Conditioning on the first passage time, we get

\[
P_{N,n_0}(t) = P_{n_0, n_0} * G_{N,n_0}(t), \quad n_0 < N.
\]  

(3.24)

By taking the L-S transform of \( P_{n_0, n_0}(t) \) and rearranging, we get

\[
\mathcal{F}_{n_0,n_0}(s) = \frac{s}{s + n_0 \rho \lambda} = 1 - \frac{n_0 \rho \lambda}{s + n_0 \rho \lambda}
\]  

(3.25)

Substituting (3.25) into the L-S transform of \( P_{N,n_0}(t) \), we obtain

\[
\mathcal{F}_{N,n_0}(s) = \tilde{G}_{N,n_0}(s) - \frac{n_0 \rho \lambda}{s + n_0 \rho \lambda} \tilde{G}_{N,n_0}(s)
\]

\[= \tilde{G}_{N,n_0}(s) - \tilde{G}_{N,n_0-1}(s). \]  

(3.26)
By taking the inverse L-S transform of \( \hat{P}_{N,n_0}(s) \) we get

\[
P_{N,n_0}(t) = G_{N,n_0}(t) - G_{N,n_0-1}(t), \quad n_0 = 0, 1, 2, \ldots, N \tag{3.27}
\]

where

\[
G_{N,N}(t) = 1, \\
G_{N,-1}(t) = 0.
\]

Figure 3.4 shows \( P_{N,n_0}(t) \) for various \( n_0 \), where \( N = 10, \ p = 0.9, \) and \( \lambda = 0.02 \). From this figure we can see how the distribution of the number of remaining errors changes with time.

Now, we obtain the expected number of remaining errors in the software at time \( t \) as follows:

\[
E[X(t) | X(0) = N] = \sum_{n_0 = 0}^{N} n_0 P_{N,n_0}(t)
\]

\[
= \sum_{n_0 = 0}^{N} n_0 \{G_{N,n_0}(t) - G_{N,n_0-1}(t)\}
\]

\[
= \sum_{n_0 = 0}^{N} \{1 - G_{N,n_0}(t)\}
\]

Now, using the expression in (3.17), we get

\[
E[X(t) | X(0) = N] = Ne^{-\lambda t}. \tag{3.28}
\]

Figure 3.5 shows the expected number of remaining errors at time \( t \) for various \( p \), where \( N = 10, \) and \( \lambda = 0.02 \). As can be seen, software errors can be eliminated faster if larger values of \( p \) are chosen. In other words, a good debugger can eliminate software errors fast. For example, for \( n_0 = 1 \) a debugger with \( p = 1 \) requires
Figure 3.4 Probability Distributions of Number of Remaining Errors, $n_0$, at Time $t$
Figure 3.5  Expected Number of Remaining Errors versus Time $t$

$N = 10$
$\lambda = 0.02$

$\begin{align*}
N & = 10 \\
\lambda & = 0.02
\end{align*}$

$\begin{align*}
p & = 0.80 \\
p & = 0.85 \\
p & = 1.00 \\
p & = 0.95 \\
p & = 0.90
\end{align*}$
debugging time $t = 118$, and the debugger with $p = 0.8$ requires $t = 148$. The difference between the two debuggers is 30 in the sense of expectation.
3.4 Expected Number of Errors Detected by Time \( t \)

We introduce a new random variable \( N(t) \) which denotes the total number of errors detected by time \( t \). The process \([N(t), t \geq 0]\) is called a counting process. We are interested in obtaining the expression for the expected number of errors detected, \( \mu_N(t) \), during the debugging period, \( t \), when the initial number of errors is \( N \), i.e.

\[
\mu_N(t) = E[N(t) | X(0) = N]
\]

which is called a Markov renewal function. By conditioning on the next software failure, we obtain the renewal equations.

\[
\mu_j(t) = F_j(t) + p\mu_{j-1}F_j(t) + q\mu_jF_j(t), \quad j = 1, 2, \ldots, N
\]

where \( \mu_0(t) = 0 \).

Using the L-S transforms of \( \mu_j(t) \), \( j = 1, 2, \ldots, N \), we get

\[
\tilde{\mu}_N(s) = \frac{1}{p} \sum_{k=1}^{N} \sum_{j=k}^{N} \frac{\lambda^j}{s + \lambda^j} = \frac{1}{p} \sum_{k=1}^{N} \tilde{G}_{N,k-1}(s).
\]

The expression for \( \mu_N(t) \) in terms of the first passage time distribution is then given by

\[
\mu_N(t) = \frac{1}{p} \sum_{k=1}^{N} G_{N,k-1}(t) = \frac{N}{p} (1 - e^{-\lambda t}).
\]

Note that if we let \( t \rightarrow \infty \) we have

\[
\mu_N(\infty) = \frac{N}{p}
\]

which is the expected number of software errors detected by the end of debugging.

Figure 3.6 shows the expected number of errors detected by time \( t \), \( \mu_N(t) \), for various \( N \) when \( p = 0.9 \) and \( \lambda = 0.02 \).
Figure 3.6 Expected Number of Errors Detected by Time $t$ for Various $N$
Let us now consider the case when the detected errors are separated as new errors and errors which were not corrected due to imperfect debugging. Let \( N_1(t) \) be a random variable which denotes the total number of imperfect debugging errors by time \( t \). Then we can show that

\[
D_{\text{N}}(t) = q M_{\text{N}}(t),
\]

(3.34)

where

\[
D_{\text{N}}(t) = E[N_1(t) | X(0) = N].
\]

Note that \( D_{\text{N}}(\cdot) = q \frac{N}{p} \).

Plots of \( M_{\text{N}}(t) \) and \( D_{\text{N}}(t) \) for the case when \( N=10 \), \( p=0.9 \) and \( \lambda=0.02 \) are shown in Figure 3.7.
Figure 3.7 Plots of $M_N(t)$ and $D_N(t)$
4. DISTRIBUTION OF TIME BETWEEN SOFTWARE FAILURES

In the previous section we studied the stochastic behavior of the number of errors in the software system during the debugging period. In this section we investigate the distribution of the time between software failures and study the problem of reliability growth. From Section 2 recall that the random variable $T_i$ denotes the time to next failure when the number of remaining errors is $i$ and $F_i(t)$ is the cdf of $T_i$. Let $X_k$ denote the time between the $(k-1)^{st}$ and $k^{th}$ software failures and $\Phi_k(x)$ be the cdf of $X_k$. Note that $X_k$ does depend on the number of remaining errors at the $(k-1)^{st}$ failure but this number is not explicitly known. Further, let $\eta_k$, a r.v., denote the number of remaining errors between the $(k-1)^{st}$ and $k^{th}$ software failures. Then, from Section 2 we have

\begin{equation}
\eta_1 = N \tag{4.1}
\end{equation}

\begin{equation}
\Phi_1(x) = F_N(x) \tag{4.2}
\end{equation}

and

\begin{equation}
\Phi_2(x) = pF_{N-1}(x) + qF_N(x). \tag{4.3}
\end{equation}

In general, we have

\begin{equation}
\Phi_k(x) = p(X_k \leq x) = \sum_{i=N-(k-1)}^{N} p(X_k \leq x | \eta_k = i) p(\eta_k = i) \tag{4.4}
\end{equation}

or
This is called a mixture of exponential distributions with binomial mixing portions. As proved in Barlow and Proschan [1], \( \phi_k(x) \) is a decreasing failure rate (DFR) distribution. The reliability function at the kth stage, i.e., between (k-1)st and kth failure, is given by

\[
R_k(x) = P(X_k > x) = 1 - \phi_k(x) = \sum_{j=0}^{k-1} \binom{k-1}{j} p^{k-j-1} q_j \overline{F}_{N-(k-j-1)}(x) \quad (4.6)
\]

where

\[
\overline{F}_N(x) = 1 - F_N(x) = e^{-N\lambda x} \quad (4.7)
\]

Also the corresponding failure rate is given by

\[
r_k(x) = \phi_k(x)/R_k(x) \quad (4.8)
\]

where \( \phi_k(x) \) is the p.d.f. of \( X_k \). The behavior of \( R_k(x) \) with respect to \( k \) is of interest. To study this behavior we have the following theorem.

**Theorem:** The reliability function \( R_k(x) \) is increasing in \( k \) for any time \( x > 0 \), i.e.
\[ R_k(x) < R_{k+1}(x), \quad k = 1, 2, \ldots \] (4.9)

**Proof:** It suffices to show that
\[ \nabla R_k(x) = R_{k+1}(x) - R_k(x) \] (4.10)

is positive for \( x > 0 \). Then we have
\[ \nabla R_k(x) = \sum_{j=0}^{k-1} \binom{k-1}{j} p^{k-j} q^j \left( \overline{F}_{N-(k-j)}(x) - \overline{F}_{N-(k-j-1)}(x) \right). \] (4.11)

It holds that for \( x > 0, j = 0, 1, 2, \ldots \)
\[ \overline{F}_{N-(k-1)}(x) > \overline{F}_{N-(k-j-1)}(x). \] (4.12)

Hence we get
\[ \nabla R_k(x) > 0 \quad \text{for} \quad x > 0. \] (4.13)

Q.E.D.

The reliability growth curves are shown in Figure 4.1, where \( N = 10, p = 0.9 \) and \( \lambda = 0.02 \). The p.d.f.'s and the failure rates of \( X_k \) are shown in Figures 4.2 and 4.3, respectively.

Note that the number of software errors remaining at the time between \((k-1)\)th and \(k\)th software failures is \( N-(k-1) \), where the random variable \( I \) is distributed as a binomial with \((k-1,q)\).
Therefore, the expected number of software errors remaining is given by \( N-p(k-1) \). This observation will be useful in constructing a likelihood function to estimate unknown parameters.
Figure 4.1 Reliability Growth Curves
Figure 4.2  PDF of Time Between Software Failures
Figure 4.3 Failure Rates of Time Between Software Failures
5. GAMMA APPROXIMATION FOR A LARGE-SCALE SOFTWARE SYSTEM

In Section 3 we obtained the quantities of interest, e.g. state occupancy probability and renewal function, in terms of first passage time distribution. Once we have computed $G_{N,n_0}(t)$, we can easily obtain $P_{N,n_0}(t)$ and $M_{N}(t)$. However, it should be noted that the computation of $G_{N,n_0}(t)$ is almost impossible for a large-scale software system because of the difficulty in computing the coefficient, $B_{N,j,n}$. Through numerical study we have found that the computations become very messy and almost impossible for $N-n_0 \geq 20$. In this section we study methods for obtaining approximate solutions for these quantities.

Of prime interest is the approximation of first passage time distribution by using a Gamma distribution. From a study of the pdf's of first passage times in Figure 3.2, we feel that these distributions might be approximated by Gamma distributions. We use the method of moments to obtain estimates of the parameters of a Gamma distribution corresponding to $G_{N,n_0}(t)$. In order to do that, we first discuss how to obtain the moments of $G_{N,n_0}(t)$ without computing the coefficient $B_{N,j,n}$. Let $T_{N,n_0}$ be a random variable which denotes the first passage time from $N$ to $n_0$. The random variable of holding time at state $N$, denoted by $W_N$, has an exponential distribution with parameter $N\lambda$. Therefore, we have

$$\mu_N = E(T_N) = \frac{1}{N\lambda} \quad (5.1)$$

$$\text{Var}(T_N) = E(T_N - \mu_N)^2 = 1/(N\lambda)^2. \quad (5.2)$$

The following recursive equations are easily obtained:
Solving (5.3) recursively, we get

$$T_{N, n_0} = \frac{1}{p} \sum_{j=n_0+1}^{N} T_j.$$ \hspace{1cm} (5.4)

Then, we have

$$\mu_{N, n_0} = E T_{N, n_0} = \frac{1}{p} \sum_{j=n_0+1}^{N} E T_j = \frac{1}{p} \sum_{j=n_0+1}^{N} 1/j \lambda.$$ \hspace{1cm} (5.5)

and

$$\text{Var}(T_{N, n_0}) = \frac{1}{p} \sum_{j=n_0+1}^{N} \text{Var}(T_j) = \frac{1}{p} \sum_{j=n_0+1}^{N} 1/(j \lambda)^2.$$ \hspace{1cm} (5.6)

These are identical to the ones obtained in Section 3.2. Suppose the Gamma distribution corresponding to $G_{N, n_0}(t)$ has a shape parameter $\alpha$ and a scale parameter $\beta$, so the mean and variance are given by $\alpha/\beta$ and $\alpha/\beta^2$, respectively. Then the parameters $\alpha$ and $\beta$ can be estimated by using the method of moments, i.e.,

$$\frac{1}{p} \sum_{j=n_0+1}^{N} 1/(j \lambda) = \alpha/\beta,$$ \hspace{1cm} (5.7)

$$\frac{1}{p^2} \sum_{j=n_0+1}^{N} 1/(j \lambda)^2 = \alpha/\beta^2.$$ \hspace{1cm} (5.8)

Therefore, we have
\begin{align*}
\hat{\beta} &= \frac{\sum_{j=n_0+1}^{N} 1/(j\lambda)}{\sum_{j=n_0+1}^{N} 1/(j\lambda)^2} \quad (5.9) \\
\hat{\sigma} &= \left[ \sum_{j=n_0+1}^{N} 1/(j\lambda) \right]^2 / \sum_{j=n_0+1}^{N} 1/(j\lambda)^2. \quad (5.10)
\end{align*}

Numerical examples for various \( n_0 \) are given in Table 5.1, where \( N=100 \), \( p=0.9 \) and \( \lambda=0.02 \). We also compute the relative losses for third and fourth moments around the mean to see how good the approximations are. Since the third and fourth moments around the mean of a Gamma distribution with parameters \( \alpha \) and \( \beta \) are given by \( 2\alpha/\beta^3 \) and \( 9\alpha/\beta^4 \), respectively, we define the relative losses for third and fourth moments around the mean as

\begin{align*}
\frac{|E(T_{N,n_0} - \mu_{N,n_0})^3 - 2\alpha/\beta^3|}{E(T_{N,n_0} - \mu_{N,n_0})^3} \quad (5.11)
\end{align*}

and

\begin{align*}
\frac{|E(T_{N,n_0} - \mu_{N,n_0})^4 - 9\alpha/\beta^4|}{E(T_{N,n_0} - \mu_{N,n_0})^4} \quad (5.12)
\end{align*}

respectively, where

\begin{align*}
E(T_{N,n_0} - \mu_{N,n_0})^3 &= \frac{1}{p^3} \sum_{j=n_0+1}^{N} E(T_j - \mu_j)^3 = \frac{2}{p^3} \sum_{j=n_0+1}^{N} 1/(j\lambda)^3 \quad (5.13)
\end{align*}

and
Table 5.1
Gamma Approximation for First Passage Time Distributions
(N = 100, p = 0.9, λ = 0.02)

<table>
<thead>
<tr>
<th>n₀</th>
<th>Mean</th>
<th>Variance</th>
<th>( \hat{\alpha} )</th>
<th>( \hat{\beta} )</th>
<th>3rd Moment</th>
<th>4th Moment</th>
</tr>
</thead>
<tbody>
<tr>
<td>10</td>
<td>125.47</td>
<td>263.01</td>
<td>59.85</td>
<td>0.477</td>
<td>28.15</td>
<td>4.21</td>
</tr>
<tr>
<td>15</td>
<td>103.84</td>
<td>168.34</td>
<td>64.05</td>
<td>0.617</td>
<td>21.58</td>
<td>2.70</td>
</tr>
<tr>
<td>20</td>
<td>88.31</td>
<td>119.82</td>
<td>65.09</td>
<td>0.737</td>
<td>16.81</td>
<td>1.92</td>
</tr>
<tr>
<td>25</td>
<td>76.19</td>
<td>90.31</td>
<td>64.28</td>
<td>0.844</td>
<td>13.19</td>
<td>1.44</td>
</tr>
<tr>
<td>30</td>
<td>66.24</td>
<td>70.47</td>
<td>62.27</td>
<td>0.940</td>
<td>10.37</td>
<td>1.12</td>
</tr>
<tr>
<td>35</td>
<td>57.81</td>
<td>56.23</td>
<td>59.44</td>
<td>1.028</td>
<td>8.14</td>
<td>0.89</td>
</tr>
<tr>
<td>40</td>
<td>50.49</td>
<td>45.49</td>
<td>56.04</td>
<td>1.110</td>
<td>6.35</td>
<td>0.72</td>
</tr>
<tr>
<td>45</td>
<td>44.02</td>
<td>37.12</td>
<td>52.21</td>
<td>1.186</td>
<td>4.92</td>
<td>0.58</td>
</tr>
<tr>
<td>50</td>
<td>38.23</td>
<td>30.40</td>
<td>48.07</td>
<td>1.257</td>
<td>3.77</td>
<td>0.47</td>
</tr>
<tr>
<td>55</td>
<td>32.99</td>
<td>24.90</td>
<td>43.70</td>
<td>1.324</td>
<td>2.84</td>
<td>0.39</td>
</tr>
<tr>
<td>60</td>
<td>28.19</td>
<td>20.30</td>
<td>39.15</td>
<td>1.389</td>
<td>2.09</td>
<td>0.31</td>
</tr>
<tr>
<td>70</td>
<td>19.70</td>
<td>13.07</td>
<td>29.69</td>
<td>1.507</td>
<td>1.03</td>
<td>0.20</td>
</tr>
<tr>
<td>80</td>
<td>12.33</td>
<td>7.63</td>
<td>19.92</td>
<td>1.616</td>
<td>0.41</td>
<td>0.11</td>
</tr>
<tr>
<td>90</td>
<td>5.82</td>
<td>3.39</td>
<td>9.99</td>
<td>1.716</td>
<td>0.09</td>
<td>0.05</td>
</tr>
</tbody>
</table>
Figure 5.1 shows the relative losses for third and fourth moments around the mean with $N$, where $p = 0.9$, $\lambda = 0.02$ and $n_0 = 0.2N$. As we see in this figure, the maximum relative losses for third and fourth moments around the mean are about 17% and 10%, respectively. This means that the Gamma approximation of first passage time distributions for large-scale software systems is reasonably good.

Plots of first passage time using Gamma approximation for $N=100$, $p=0.9$ and $\lambda=0.02$ are given in Figure 5.2 for $n_0=0,1,2,3,5$, and 9. Also, plots of state occupancy probabilities using this approximation are given in Figure 5.3.
Figure 5.1  Relative Loss for the Third and Fourth Moments
Figure 5.2 Plots of First Passage Times Using Gamma Approximation
Figure 5.3 State Occupancy Probabilities
Using Gamma Approximations
6. CONCLUDING REMARKS

An imperfect debugging model (IDM) for software systems was developed in this report. Various quantities of interest were derived in terms of the first passage time distribution of the underlying semi-Markov process. Computations for and usefulness of these quantities were illustrated via numerical examples. An approximation method for obtaining these quantities for large-scale software system was also presented.

It should be pointed out that most of the models reported in the literature, for example the models in [3], [6], [9], [10], and [13], are special cases of IDM.
7. SELECTED REFERENCES


