SIMULATION PROGRAMMING USING SIMSCRIPT II

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The simulation described in this paper has been designed to illustrate as much of SIMSCRIPT II as possible in a natural, problem-oriented setting. While the program uses all of the language's features that are important in simulation studies, it does not use every SIMSCRIPT feature. A complete description is contained in P. J. Kiviat and R. Villanueva, The SIMSCRIPT II Programming Language, R-460-PR, The RAND Corporation, September 1968.

Even though the features illustrated are not exhaustive, the example may still seem forced and artificial. This should not be surprising, as it is a rare program that requires the full facilities of any complex programming language. The example is an elaboration of the job shop model of Chapter 3 of the SIMSCRIPT II report.**

The plan of the paper is as follows: the first section describes a model of a system in general terms, presents some problems the model has been designed to study, and places the rest of the paper in perspective. The next section contains a listing of the complete simulation program followed by a set of typical data cards. The last section works through the program section by section and, occasionally, where it is warranted, statement by statement, explaining the syntax and semantics of the statements.

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THE SYSTEM

The system under study is shown abstractly in Fig. 1. It is a shop containing N production centers, each containing \( M \) identical machines, and finished goods inventory storage area. The shop produces \( P \) standard products for local sale and distribution, and variations of standard products for local and export distributors. Each product ordered travels through the shop, undergoing processing at production centers according to standard routings, production times and expediting procedures.

Fig. 1 -- System Under Investigation

Each production center has an in-process inventory area where products in process are stored if they cannot be worked on when they arrive. To minimize the value of in-process inventory, the production rules of the shop remove partially completed products from production center queues according to a "high value first" rule.

Table 1 shows the entity-attribute-set model of the shop and its product line. These descriptors explain the shop's static structure. Permanent entities are used for production centers and product descriptions. Temporary entities are used for jobs and for job processing specifications.

The shop operates roughly as follows. When orders for standard products come into the shop, a standard production sequence is copied from an order book onto a job's production routing tag. The job is sent to the first production center, where it is worked on if a
machine is free. If no machines are available, the job is put in a work-in-process queue until a machine becomes available. When a job finishes processing, its routing tag is examined, and the job is either sent on to another processing center or to the finished goods inventory.

The shop's dynamic structure is captured in two events -- SALE and END.OF.PROCESS. Two other events, WEEKLY.REPORT and END.OF.SIMULATION, print periodic system performance data and stop the simulation. SALE is set up so it can be triggered both internally and externally. When triggered internally, SALE represents a local sale of a standard product. When triggered externally, SALE represents either a local or export sale of a special order. Two external event data tapes are provided to supply special order information.

In SALE, jobs are assigned to machines and the system state changes to reflect such assignments. When a job is assigned to a machine, an END.OF.PROCESS event is scheduled to terminate the processing, make the machine available for another job, and pass the job on for additional processing or for shipment.

The simulation model is designed to determine the number of machines needed at each production center to provide "adequate" customer service. To study the effects of varying the number of machines in each center, a TALLY statement looks at the length of time jobs spend in the shop, and an ACCUMULATE statement looks at the waiting lines that build up at the production centers. Some number of machines will be chosen that balances the cost of degraded customer service with the costs of additional machines.

The program listing that follows has been written and annotated to make it as readable as possible. Statements that are not clear from the program itself are discussed in the section following the listing.


<table>
<thead>
<tr>
<th>Entity</th>
<th>Set or Attribute</th>
<th>Comment</th>
</tr>
</thead>
<tbody>
<tr>
<td>PRODUCTION.CENTER</td>
<td>NUMBER.IDLE</td>
<td>The number of idle machines in a production center</td>
</tr>
<tr>
<td></td>
<td>QUEUE</td>
<td>Each production center owns a collection of in-process products.</td>
</tr>
<tr>
<td>PRODUCT</td>
<td>SALES,FREQUENCY</td>
<td>Characterizes the frequency with which orders for a standard product arrive at the shop.</td>
</tr>
<tr>
<td></td>
<td>NAME</td>
<td>Identifies a product</td>
</tr>
<tr>
<td></td>
<td>STRUCTURE</td>
<td>Each product has a list of standard operations that have to be performed to produce it.</td>
</tr>
<tr>
<td>JOB</td>
<td>VALUE</td>
<td>The dollar value of a job</td>
</tr>
<tr>
<td></td>
<td>DUE,DATE</td>
<td>The time a job is promised to a customer</td>
</tr>
<tr>
<td></td>
<td>ARRIVAL,TIME</td>
<td>The time a job is ordered</td>
</tr>
<tr>
<td></td>
<td>EXPEDITE,FACTOR</td>
<td>The degree to which a job's processing can be speeded up at a production center.</td>
</tr>
<tr>
<td></td>
<td>ROUTING</td>
<td>A list of production centers through which a job has to be processed</td>
</tr>
<tr>
<td></td>
<td>FINISHED,Goods.INVENTORY</td>
<td>Jobs can be placed in finished goods inventory awaiting shipment</td>
</tr>
<tr>
<td>THE SYSTEM</td>
<td>FINISHED,Goods.INVENTORY</td>
<td>Jobs are placed in finished goods inventory if finished before their due date.</td>
</tr>
<tr>
<td>OPERATION</td>
<td>MACHINE.DESTINED</td>
<td>The production center at which a job must be processed.</td>
</tr>
<tr>
<td></td>
<td>CODE</td>
<td>A number representing a particular processing task.</td>
</tr>
<tr>
<td></td>
<td>PROCESS,TIME</td>
<td>The time it takes to perform a processing task</td>
</tr>
<tr>
<td></td>
<td>STRUCTURE</td>
<td>The standard parts list on which different operations appear</td>
</tr>
<tr>
<td></td>
<td>ROUTING</td>
<td>The processing operations required for a particular job</td>
</tr>
</tbody>
</table>
SAMPLE SIMSCRIPT II SIMULATION PROGRAM
A SHOP SIMULATION

PREAMBLE
NORMALLY MODE IS INTEGER AND DIMENSION IS 0

PERMANENT ENTITIES----
EVERY PRODUCT HAS A SALES FREQUENCY AND A NAME AND OWNS A STRUCTURE
DEFINE SALES,FREQUENCY AS A REAL RANDOM LINEAR VARIABLE
DEFINE NAME AS AN ALPHA VARIABLE
EVERY PRODUCT, PRODUCT HAS A PRODUCT SALES(*/2)
EVERY PRODUCTION CENTER HAS A (MAX, IN, QUEUE(*1/2), MAX, QUEUE(*2/2)) IN ARRAY 1,
A (NUM(*1/2), NUM(*2/2)) IN ARRAY 2, A WSM, A PSUM, A NUMBER, IDLE
AND OWN'S A QUEUE
DEFINE NUMBER, IDLE AS A VARIABLE MONITORED ON THE LEFT

TEMPORARY ENTITIES----
EVERY JOB HAS A VALUE IN WORD 2, A DUE DATE, AN ARRIVAL TIME,
AN EXPEDITED, FACTOR FUNCTION, MAY BELONG TO A QUEUE, OWNS A ROUTING
AND MAY BELONG TO THE WAITING SET
DEFINE EXPEDITED, FACTOR AS A REAL FUNCTION
DEFINE VALUE, DUE DATE AND ARRIVAL TIME AS REAL VARIABLES
DEFINE ROUTING AS A FIFO SET WITHOUT P AND N ATTRIBUTES
DEFINE QUEUE AS A SET RANKED BY HIGH VALUE
EVERY OPERATION HAS A CODE(*1/2) AND MACHINE, DESTINEC(*2/2) ) IN WORD 1
AND A PROCESS, Time AND BELONGS TO A STRUCTURE AND A ROUTING
DEFINE STRUCTURE AS A SET RANKED BY LOW CODE WITHOUT P ATTRIBUTE
AND WITHOUT A ROUTING
DEFINE PROCESS, TIME AS A REAL VARIABLE

EVENT NOTICES INCLUDE WEEKLY REPORT
EVERY SALE HAS A PRODUCT, TYPE, A PRICE AND A PRIORITY
DEFINE PRICE AS A REAL VARIABLE
EVERY ENDC, CF, PROCESS HAS AN ITEM AND A PRODUCER

BREAK SALE TIES BY HIGH PRICE THEN BY LOW PRIORITY
EXTERNAL EVENTS ARE ENDC, CF, SIMULATION AND SALE
EXTERNAL EVENT UNITS ARE LOCAL, SALES AND EXPORT, SALES
PRIORITY ORDER IS ENDC, CF, PROCESS, SALE, WEEKLY REPORT AND ENDC, CF, SIMULATION

BEGIN FILING AND REMOVING FROM QUEUE CALL QUEUE, CHECK
BEFORE DESTROYING JOB, CALL STAY, TIME

DEFINE STAY AS A REAL DUMMY VARIABLE
TALLY AVG, STAY AS THE WEEKLY MEAN, VAR, STAY AS THE WEEKLY VARIANCE, SUM, STAY AS THE WEEKLY SUM, SUM, SQUARES, STAY AS THE WEEKLY SUM, CF, SQUARES, AND NUM, STAY AS THE WEEKLY NUMBER OF STAY
ACCUMULATE WSL, AS THE WEEKLY SUM, WNUM AS THE WEEKLY NUMBER, AVG, QUEUE AS THE WEEKLY MEAN, MAX, QUEUE AS THE WEEKLY MAXIMUM AND PRI(10) TO MAX QUEUE AS THE WEEKLY HISTOGRAM OF N, QUEUE
ACCUMULATE WSL, AS THE MONTHLY SUM, WNUM AS THE MONTHLY NUMBER, AVG, QUEUE AS THE MONTHLY MEAN, MAX, QUEUE AS THE MONTHLY MAXIMUM OF N, QUEUE

THE SYSTEM OWNS A FINISHED GOODS INVENTORY
DEFINE FINISHED, GOODS, INVENTORY AS A SET RANKED BY DATE
DEFINE LOCAL TO MEAN DEFINE 1, K, L, M AND N AS SAVI, INTEGER VARIABLES
DEFINE WEEK TO MEAN *MOURS, V01 MOURS
DEFINE PRIORITY, FREQUENCY AS A 2-DIMENSIONAL ARRAY
DEFINE TITLE AS A TEXT VARIABLE
DEFINE WEEK.COUNTER AND TAPEJFLAG AS INTEGER VARIABLES
DEFINE AVERAGE AS A REAL FUNCTION WITH 1 ARGUMENT

END

PAIK

*INITIALIZE* PERFORM INITIALIZATION
LET BETWEEN.V="TRACE" START SIMULATION

** PERFORM NEXT EXPERIMENT **
FOR EACH JOB IN FINISHED.GOODS.INVENTORY, DO
  REMOVE THE JOB FROM FINISHED.GOODS.INVENTORY
  DESTROY THE JOB
LCOP

FOR EACH PRODUCTION.CENTER,
  FOR EACH JOB IN QUEUE, DO
    FOR EACH OPERATION IN ROUTING, DO
      REMOVE THE OPERATION FROM ROUTING
      DESTROY THE OPERATION
    LCOP
    REMOVE THE JOB FROM QUEUE
    DESTROY THE JOB
  LCOP

FOR EACH PRODUCT, CC
  FOR EACH OPERATION IN STRUCTURE, DO
    REMOVE THE OPERATION FROM THE STRUCTURE
    DESTROY THE OPERATION
  LCOP

ALSO FOR EACH RANDOM.E IN SALES.FREQUENCY, DO
  REMOVE THE RANDOM.E FROM SALES.FREQUENCY
  DESTROY THE RANDOM.E
LCOP

RELEASE NAME, F.STRUCTURE, L.STRUCTURE, N.STRUCTURE, F.SALES.FREQUENCY,
  PRODUCT.SALES, NUMBER.IDLE, F.QUEUE, L.QUEUE, N.QUEUE, MAX.IN.QUEUE,
  MAX.QUECLE, WSUP, NSUP, WNUP, MNUP
RELEASE PRIORITY.FREQUENCY
ERASE TITLE
RESET TOTALS OF STAY
FOR EACH PRODUCTION.CENTER, RESET TOTALS OF N.QUEUE

LET WEEK.COUNTER=0
WRITE AS "******/
LET TAPEJFLAG=0

** RELEASE EXTERNAL EVENTS IN NEXT EXPERIMENT
REWIN: LOCAL.SALES AND EXPORT.SALES

GC INITIALIZE
STOP   END
HCOTINE FOR INITIALIZATION

LOCAL
DEFINE PB TO MEAN PRIORITY.FREQUENCY
DEFINE SF TO MEAN SALES.FREQUENCY
DEFINE CHECK AS AN ALPHA VARIABLE
LET ECF.V=1
INPUT TITLE
IF ECF.V=2, PRINT 1 LINE AS FOLLOWS
ENC OF DATA HIT
STOP ELSE

READ A.PRODUCTION.CENTER
CREATE EVERY PRODUCTION.CENTER
FOR EACH PRODUCTION.CENTER, READ NUMBER.IDLE

READ A.PRODUCT
CREATE EVERY PRODUCT
RESERVE PRIORITY.FREQUENCY(*) AS A.PRODUCT BY *
FOR EACH PRODUCT, CC
READ NAME
READ SALES.FREQUENCY
RESERVE PRIORITY.FREQUENCY(PRODUCT,*1) AS PRODUCT
FOR I=1 TO PRODUCT, READ PRIORITY.FREQUENCY(PRODUCT,I)
UNTIL MODE IS ALPHA, DO THIS....
CREATE AN OPERATION
FILE THE OPERATION IN STRUCTURE
READ CODE, MACHINE.DESTINED AND PROCESS.TIME

LOOP
SKIP 1 FIELD
CAUSE A SALE IN SF HOURS
LET PRODUCT.TYPE=PRODUCT
LET PRICE=PRODUCT*RANDOM.F(I)
LET PRIORITY=PF(PRODUCT, TRUNC.F(PRICE)+1)

LCCP

READ LOCAL.SALES, EXPORT.SALES AND SAVE.TAPE

READ MCNTY, DAY AND YEAR CALL ORIGIN.R(MONTH, DAY AND YEAR)

READ CHECK IF CHECK EQUALS "OK", CALL REPORT RETURN
OTHERWISE PRINT 1 LINE AS FOLLOWS
EITHER TOO MUCH DATA OR DATA HAS BEEN READ INCORRECTLY
STOP ENC

EVENT SALE(PRODUCT, PRICE, PRIORITY) SAVING THE EVENT NOTICE
DEFINE SF TO MEAN SALES.FREQUENCY
LOCAL
IF SALE IS EXTERNAL, READ PRODUCT, PRICE AND PRIORITY AS II 30, I 5, O(10,3), I 5
REGARDLESS ADD 1 TO PRODUCT.SALES(PRODUCT, TRUNC.F(PRICE)+1)
CREATE A JOB
LET VALUE=PRICE
LET DUE.DATE=TIME.V + PRICE + PRIORITY
LET ARRIVAL.DATE=TIME.V
IF SALE IS INTERNAL,
FOR EACH PIECE OF STRUCTURE, FILE PIECE IN ROUTING GO TO JOB
** PROCESS SPECIAL ORDERS
Cotherwise until mode is Alpha, do the following...
read n
for each piece in structure with code(piece) = n,
find the first case, if none go to loop
file piece in routing
loop
loop
job
new attend.tc.job
if sale is external, destroy the sale return
cotherwise
schedule the sale(product, product*random.f(1), priority, frequency(product, trunc.f(price)+1)) in sf hours
return
end

routine to attend.tc.job
lit production.center=machine.destined(f.routing(job))
if number.idle is positive,
subtract 1 from number.idle perform allocation
return
cotherwise
file job in queue return
end

routine for allocation
r@evi tpf first operation from this routing
schedule an end.of.process given job and production.center in
price*time*expedite.factor.hours
return end

routine expedite.factor
if time.v is greater than due.date return with 0.5 else
return with min.fit(due.date-time.v)/process.time,1)
end

lit lit of.process given job and production.center
if routing is empty, if due.date <= time.v,
destroy this job go to pc
else file this job in finished.goods.inventory go to pc
cotherwise call attend.tc.job
*pl*
if queue is empty,
acc 1 to number.idle return
else remove the first job from queue
perform allocation return
end

event for weekly.report saving the event notice
write this weekly.report in 1 week
all 1 1 week.count
new mيلي
write weekly totals of stay
FOR EACH PRODUCTION.CENTER, RESET WEEKLY TOTALS OF N.QUEUE IF MOD(F(WEEK.COUNTER,4),0) = 0, FOR EACH PRODUCTION.CENTER, RESET MONTHLY TOTALS OF N.QUEUE ELSE RETURN END

EVENT FOR END.OF.SIMULATION
FOR I = 1 TO EVENTS.V, FOR EACH NOTICE IN EV.S(I), DO
REMOVE THE NOTICE FROM EV.S(I)
DESTROY THE NOTICE

ROUTINE FOR QUEUE.CHECK GIVEN ENTITY AND I LOCAL
IF I.LE.1.LE.N.PRODUCTION.CENTER RETURN
OTHERWISE PRINT 1 LINE WITH I AS FOLLOWS
STOPPED TRYING TO REFERENCE QUEUE(* *) TRACE STOP END

ROUTINE FOR STAY.TIME GIVEN JOB
LET STAY.TIME.V = ARRIVAL.TIME(JOB)
RETURN END

ROUTINE TO TRACE
LOCAL
IF FINISHED.GOODS.INVENTORY IS EMPTY, GO AROUND
ELSE FOR EACH JOB IN FINISHED.GOODS.INVENTORY UNTIL DUE.DATE > TIME.V, DO
REMOVE THE JOB FROM FINISHED.GOODS.INVENTORY DESTROY THE JOB

"AROUND"
GC TO END.CF.PROCESS, SALE, WEEKLY.REPORT AND END.OF.SIMULATION PER EVENT.V
"ENC.CF.PROCESS" WRITE ITEM, PRODUCER, TIME.V AS "ITEM", I 5, "STOPPED",
"PROCESSING ON MACHINE", I 5, "AT TIME="@10.3), /
RETURN
"SALE" WRITE TIME.V, PRODUCT.TYPE, PRICE AND PRIORITY AS "SALE OF TYPE", I 3, 
"PRODUCT AT TIME="@10.3), " FOR S", @16.2), "PRIORITY="@1.3, 
/ RETURN
"WEEKLY.REPORT"
"ENC.CF.SIMULATION"
RETURN END

LEFT HCUTINE NUMBER,ICLE(M)
DEFINE J AS A SAVED, 2-DIMENSIONAL ARRAY
DEFINE K AS A SAVED, 1-DIMENSIONAL ARRAY
ENTER WITH N
IF TAPE.FLAG=0, LET TAPE.FLAG=1
RELEASE J AND K
RESERVE K(0) AS N.PRODUCTION.CENTER
RESERVE J(0,0) AS 100 BY N.PRODUCTION.CENTER
REGARDLESS ADD 1 TO K(M)
LET J(K(M),M)+N
IF K(M) = 100, WRITE M AS 1 3 USING SAVE.TAPE
FOR I=1 TO 100, WRITE J(I,M) AS 1 3 USING SAVE.TAPE
WRITE AS / USING SAVE.TAPE
REGARDLESS MOVE FROM N
RETURN END

RCLTIME TO REPORT
LOCAL
IF TIME=0,
START NEW PAGE "AND" OUTPUT TITLE
SKIP 2 OUTPUT LINES
PRINT 3 LINES AS FOLLOWS
NAME
SALES FREQUENCY PRODUCTION SEQUENCE
PROB. VALUE CODE CNTR TIME
FOR EACH PRODUCT, DO
LET I=F.SALES.FREQUENCY
LET J=F.STRUCTURE
PRINT 1 LINE WITH NAME, PROB.A(I), RVALUE.A(I), CODE(J),
MACHINE.DESTINED(I) AND PROCESS.TIME(I) THUS
*000 000 000 000
IF I=0, LET I=F.SALES.FREQUENCY(I) ELSE
IF J=0, LET J=F.STRUCTURE(I) ELSE
IF I=0 AND J=0, GO TO "LOOP" ELSE
IF I=0 AND J=0, PRINT 1 LINE WITH PROB.A(I) AND RVALUE.A(I) THUS
000 000
ELSE IF I=0 AND J=0, PRINT 1 LINE WITH CODE(I), MACHINE.DESTINED(I),
AND PROCESS.TIME(I) THUS
000 000 000 000
ELSE IF I=0=J, PRINT 1 LINE WITH PROB.A(I), RVALUE.A(I), CODE(I),
MACHINE.DESTINED(I) AND PROCESS.TIME(I) THUS
000 000 000 000
*LOOP* LOOP
SKIP 2 OUTPUT LINES
PRINT 2 LINES AS FOLLOWS
NAME
SALES FREQUENCY PRODUCTION SEQUENCE
PROB. VALUE CODE CNTR TIME
FOR EACH PRODUCTION.CENTER, PRINT 1 LINE WITH PRODUCTION.CENTER AND
NUMBER.IDLE THUS
000 000
SKIP 2 OUTPUT LINES
PRINT 2 LINES AS FOLLOWS
INITIAL EVENTS
EVENT TYPE TIME
FOR I=1 TO EVENTS.V, FOR EACH J IN EV.S(I),
PRINT 1 LINE WITH I AND TIME.A(I) THUS
000 000
REGARDLESS....
START NEW PAGE
PRINT 1 LINE AS FOLLOWS
PRINT 3 LINES LIKE THIS.... WEEklY REPORT
PRINT 2 LINES WITH AVG.STAY AND VAR.STAY AS FOLLOWS

JOB STAY STATISTICS ARE:
AVG.STAY = 0.90
VARIANCE = 0.00

SKIP 3 OUTPUT LINES
BEGIN REPORT
BEGIN HEADING
PRINT 2 LINES AS FOLLOWS

PRODUCTION CENTER QUEUING REPORT
CNTR AVG. QUEUE MAX. QUEUE

ENC **HEADING
FOR EACH PRODUCTION.CENTER, PRINT 1 LINE WITH PRODUCTION.CENTER,
AVG.QUEUE AND MAX.QUEUE THUS

** 0.90

ENC ** REPORT
PRINT 1 LINE WITH AVERAGE(AVG.QUEUE*) LIKE THIS
OVERALL AVERAGE QUEUE LENGTH OF ALL QUEUES IS 0.00

SKIP 3 OUTPUT LINES
FOR EACH PRODUCTION.CENTER, DO
BEGIN REPORT PRINTING FOR 1 TO 25 IN GROUPS OF 5
BEGIN HEADING
PRINT 1 LINE WITH PRODUCTION.CENTER LIKE THIS
HISTOGRAM OF QUEUE LENGTH FOR PRODUCTION CENTER **
ENC ** HEADING
PRINT 1 LINE WITH A GROUP OF FREQUENCY*2) FIELDS THUS

ENC ** REPORT

LOOP
IF MCC.TYPE.WEEK.COUNTER*41==0, RETURN
OTHERWISE START NEW PAGE
PRINT 1 LINE AS FOLLOWS
MONTHLY REPORT

SKIP 2 OUTPUT LINES
SKIP 3 OUTPUT LINES
BEGIN REPORT
BEGIN HEADING
PRINT 2 LINES AS FOLLOWS

PRODUCTION CENTER QUEUING REPORT
CNTR AVG. QUEUE MAX. QUEUE

ENC **HEADING
FOR EACH PRODUCTION.CENTER, PRINT 1 LINE WITH PRODUCTION.CENTER,
AVG.QUEUE AND MAX.QUEUE THUS

** 0.00

ENC ** REPORT
PRINT 1 LINE WITH AVERAGE(AVG.QUEUE*) LIKE THIS
OVERALL AVERAGE QUEUE LENGTH OF ALL QUEUES IS 0.00
RETURN

ENC

ROUTINE FOR AVERAGE GIVEN ARRAY
LOCAL
DEFINE ARRAY AS A 1-DIMENSIONAL ARRAY
FOR J=1 TO DIP.F.1ARRAY*, COMPUTE M AS THE MEAN OF ARRAY1J
RETURN WITH M
ENC
SAMPLE DATA FOR SEVERAL JOB SHOP EXPERIMENTS USING THE LEVEL K, SECTION 5.09 JOB SHOP SIMULATION PROGRAM. TITLES, AS SHOWN, CAN EXTEND OVER SEVERAL CARDS AND ARE ENDED BY A MARK-V CHARACTER *.

5 10 10 9 9 9

<table>
<thead>
<tr>
<th>TOP</th>
<th>YCYC</th>
<th>COLL</th>
</tr>
</thead>
<tbody>
<tr>
<td>0.50</td>
<td>0.10</td>
<td>0.10</td>
</tr>
<tr>
<td>0.75</td>
<td>0.28</td>
<td>0.28</td>
</tr>
<tr>
<td>20.0</td>
<td>0.56</td>
<td>0.56</td>
</tr>
<tr>
<td>1.00</td>
<td>1.00</td>
<td>1.00</td>
</tr>
</tbody>
</table>

THE FOLLOWING CARDS ARE SAMPLES OF THE DATA CARDS PUNCHED FOR ONE OF THE TWO EXTERNAL EVENTS TAPES. THIS IS JUST A SAMPLE FROM THE TAPE.

<table>
<thead>
<tr>
<th>SALT</th>
<th>ENC.LF.SIMULATION</th>
</tr>
</thead>
<tbody>
<tr>
<td>7/2/68</td>
<td>9/1/69</td>
</tr>
<tr>
<td>12 12</td>
<td>12 12</td>
</tr>
<tr>
<td>1.00</td>
<td>1.00</td>
</tr>
<tr>
<td>1.00</td>
<td>1.00</td>
</tr>
</tbody>
</table>

* THIS IS A TITLE CARD FOR THE SECOND SIMULATION EXPERIMENT OF THE SERIES. DATA CARDS FOR THIS EXPERIMENT WILL HAVE THE SAME FORMAT AS THOSE OF THE PREVIOUS EXPERIMENT AND WILL END WITH AN "OK" CARD *.

THE FOLLOWING CARDS ARE SAMPLES OF THE DATA CARDS PUNCHED FOR ONE OF THE TWO EXTERNAL EVENTS TAPES. THIS IS JUST A SAMPLE FROM THE TAPE.
COMMENTS ON THE SIMULATION PROGRAM

The program is arranged functionally and is discussed as it appears. The order of the program is preamble, main routine, initialization, events and routines of the simulation model, routines for monitoring, debugging and analysis.

Preamble

The preamble is divided into seven sections: permanent entities, temporary entities, event notices, event control, debugging, analysis and miscellaneous declarations. Most simulation programs can be organized this way.

One compound and two simple permanent entities are declared. The special features of each are these:

PRODUCT has a RANDOM attribute and an ALPHA attribute, each requiring definition in a DEFINE statement.

PRODUCT.SALES, the single attribute of the compound entity PRODUCT, is intrapacked to conserve storage space.

PRODUCTION.CENTER has two pairs of attributes that are packed in the same array, and one attribute that is monitored. The packed attributes use field-packing, equivalence and array specification. The monitored attribute requires an additional DEFINE statement.

PRODUCT.SALES could just as easily have been defined as a global array or as a two-dimensional system attribute. As a global array though, it could not have been packed; as a system attribute it would not be eligible for implied subscripting.

Two temporary entities are declared. The special features of each follows:

JOB has one attribute placed in a specific word in its entity record, and has a function attribute. The function attribute requires a DEFINE statement to declare its mode.

Two sets in which a JOB participates have their implied properties modified by DEFINE statements.

Two attributes of OPERATION are packed in the first word of each entity record.
A set to which an OPERATION belongs has its removal routines deleted by a DEFINE statement.

Three event notices are declared. The special features of each are these:

WEEKLY.REPORT has no attributes, and neither owns nor belongs to sets other than the standard one defined for all event notices.

One event, SALE, breaks ties among competing event notices through a BREAK TIES declaration. The other internal events break ties, if they occur, on a first-come, first-served basis.

Two external event classes, END.OF.SIMULATION and SALE, are declared. Two input devices are declared as suppliers of external event triggers. The priority order of the four event classes is declared in a PRIORITY statement.

Two BEFORE statements are used. Each states that a certain routine is to be called before a specified action takes place. The arguments of these routines are not stated, but implied.

One TALLY and two ACCUMULATE statements are used. The special features of each follow:

The TALLY statement compiles statistics for a DUMMY variable, which is declared in a separate DEFINE statement.

All of the statistical counters used in the TALLY and ACCUMULATE statements are defined so they can be released. If they were not named, they would be given local names such as A.1 and A.2 by the SIMSCRIPT II compiler.

FREQ is a two-dimensional array. The first dimension is an entity index. The variable for which it accumulates a histogram is an attribute of PRODUCTION.CENTER. The second dimension is the histogram index and is an integer between 1 and 26.

The remaining statements are:

Declare a system owned set.
Use DEFINE TO MEAN statements to create shorthand notation.
Declare four global variables: a two-dimensional array, two INTEGER variables, and a TEXT variable.
Declare a function and specify the number of arguments it must always have.

**Main Program**

The main routine has three functions: it initializes the model so simulation can start, it transfers control to the timing routine when initialization is complete, and it resets the entire system to an "empty" condition at the end of a simulation run so another experiment can begin.

Initialization takes place in the routine INITIALIZATION. After initialization, the SUBPROGRAM system variable BETWEEN.V is set to the routine name 'TRACE', indicating that this routine is to be called before each event is executed. Simulation begins at the START SIMULATION statement that removes the first event from the events list and transfers control to it. Simulation proceeds until an END.OF. SIMULATION event occurs. This event, aside from its obvious task of reporting the results of the simulation experiment, empties the events list sets. When END.OF.SIMULATION returns control to the timing routine, the lack of scheduled events causes control to pass to the statement after START.SIMULATION. In many simulations this will be STOP. In this example, it is the first of many statements that release and destroy all permanent and temporary entities. After these statements have been executed, all the memory structures set up by the previous experiment have been erased.

To perform this erasure, the system set FINISHED.GOODS.INVENTORY and the sets owned by the permanent and temporary entities are emptied and their members destroyed. Finally, all attributes of permanent entities are released. Special features to notice are these:

- The PRODUCTION.CENTER loop in which operations owned by jobs owned by production centers are successively removed and destroyed.
- The RANDOM variable SALES.FREQUENCY is treated as a set when it is emptied.
- All permanent entity attributes, including set pointers and statistical accumulators, are released.
In many programs, so extensive a reinitialization process is not necessary. For example, it is usually sufficient to zero out all attribute values and empty all sets. This example has been written to illustrate what seems to be the worst case. When single simulation runs are made and no reinitialization is necessary, the initialization routine can be released and its space regained for array and entity storage. The following routine shows how this is done.

Add to the preamble:

```
DEFINE INITIALIZATION AS A RELEASABLE ROUTINE
```

Use this routine:

```
MAIN
PERFORM INITIALIZATION
RELEASE INITIALIZATION
START SIMULATION
STOP END
```

**Program Initialization**

INITIALIZATION starts with some declarations. The first takes advantage of the DEFINE TO MEAN statement of the preamble to define some local INTEGER variables I, J, K, L, M and N. The next two statements are local DEFINE TO MEAN declarations that create shorthand notations for two lengthy variable names. The last declaration describes a local ALPHA variable that verifies whether or not all input data have been read.

Since a mistake may have been made in setting up a simulation run, EOF>V is set to 1 to give the program control over the actions taken when the end of the input data file is reached. If an end-of-file is encountered when reading TITLE, EOF>V is set to 2 and this fact is picked up in the following IF statement. A sequence of simulation experiments can also be stopped this way. When all the data for a sequence of runs are exhausted, these statements will stop the program.
The INPUT statement reads characters from the current input unit READ,V until an asterisk, the MARK,V default symbol, is reached. A typical simulation TITLE card might be:

SIMULATION RUN NO. 1 JOB SHOP WITH 10 CENTERS *

If some symbol other than * is to be used as a TEXT terminator, a statement such as LET MARK,V="?" is put at the head of INITIALIZATION.

A value that is the number of production centers is then read and used to reserve the arrays that hold the attributes of PRODUCTION.CENTER. This value is also used to read in the number of machines in each production center (which are all idle when simulation begins.)

A similar process then takes place for PRODUCT. It is more complex in that a richer variety of data structures are associated with PRODUCT than with PRODUCTION.CENTER. The data structures are the following:

- PRIORITY,FREQUENCY—a "ragged table" whose rows are "sized" dynamically.
- SALES,FREQUENCY—a RANDOM variable whose sampling data is read in a standard format.
- STRUCTURE—a set with OPERATIONS as members.

Also, an initial local SALE for a standard product must be scheduled for each product type. In scheduling these sales, the PRICE of each SALE is a random variable between 0 and the product type, e.g., a type 3 product can be sold for between $0 and $3, and the PRIORITY assigned to a sale can be determined by a random draw from the PRIORITY,FREQUENCY table.

At the end of initialization the numbers of the input devices for the LOCAL.SALES and EXPORT.SALES external event units are read. This allows devices to be changed between simulation runs. Finally, the ORIGIN,R routine is invoked to set the simulation calendar so that calendar dates can be used on the external event tapes.

If the last data field read is not the character string OK, the run terminates with an error message.
Events and Routines of the Simulation Model

The event SALE is written to react properly to both internal and external event triggers. The event creates a job, determines its routing through the shop, and starts it in its processing sequence. If the sale is internal, a new order is scheduled for the same product some time in the future.

The EVENT statement defines SALE as an event routine with three input arguments. It also declares that when a SALE event notice is selected as the next event, the first three programmer defined attributes of SALE are to be assigned to the local variables, PRODUCT, PRICE and PRIORITY, and that the event notice is not to be destroyed. An important point to note here is that PRODUCT, PRICE and PRIORITY are local variables; they are different from the variables defined as attributes of the event notice in the preamble.

The first two statements are local declarations that we have seen before. The third statement is the one that allows SALE to be used with both internal and external event triggers. It says: if SALE has occurred externally, read three data items, otherwise do not read any data. Regardless of how the event was triggered, values get assigned to PRODUCT, PRICE and PRIORITY.

The next statement adds 1 to an element of PRODUCT.SALES, counting the number of times particular products are sold at different prices.

The next section of code creates a JOB entity and assigns values to its attributes. The JOB is the entity that will flow through the shop, and will represent the sale from now on. If SALE is triggered internally, it is a sale for a standard product, and the standard sequence of operations for producing that product is transferred from STRUCTURE(PRODUCT) to ROUTING(JOB). Notice that implied subscripts are used in the program for both STRUCTURE and ROUTING. If the SALE is triggered externally, it represents a possible special order. As special order operations are subsets of operations that produce standard products, data are read that select a subset of operations for producing the order and store it in the JOB routine set.
With this, the task of SALE is almost completed. The now completely specified JOB is started through processing by the routine ATTEND.TO.JOB. After this routine deals with the job, it returns to SALE where arrangements are made for the next SALE. If the SALE just processed was triggered externally, the event notice is destroyed and control returned to the timing mechanism. The next external event will be selected automatically by the timing mechanism from the external events tapes. If the SALE was triggered internally, the event notice is reused to schedule another sale for this same product, with a different price and priority. The time at which this sale is to occur is determined by a random sample from the RANDOM variable SALES.FREQUENCY(PRODUCT); again implied subscripting is used.

ATTEND.TO.JOB uses the routing of the current JOB to select the production center in which the first operation is to be performed. The first statement in ATTEND.TO.JOB is important as it illustrates several basic operations in SIMSCRIPT II programming. First, it illustrates the use of a set pointer to select a set member; F.ROUTING(JOB) picks out an entity identification number. This identification number is then used with an attribute of the entity type it represents (OPERATION) to determine a value; MACHINE.DESTINED(R.ROUT.NG(JOB)) is the production center number the first operation requires. This number is assigned to PRODUCTION.CENTER so that implied subscripting can be used later on in the program.

After the production center is determined, NUMBER.IDLE(PRODUCTION.CENTER) is examined to determine whether a machine is available in this production center to process the job. If a machine is available, it is allocated to the job by first subtracting 1 from the idle machine counter and then calling the routine ALLOCATION. If a machine is not available, the job is filed in a QUEUE owned by the production center. When a machine becomes free at some later date, the QUEUE will be examined and a job removed for processing. All the time a job spends in the shop in excess of its actual processing time is spent in queues belonging to different production centers and in finished goods inventory.
ALLOCATE knows that it is given a certain job because the
identification number of this job is in the global variable
JOB. It removes the first operation from the ROUTING of this job,
and schedules an END.OF.PROCESS event. The job identification
and the production center identification are assigned to the first
two attributes of this event notice. The event is scheduled for a
standard time, PROCESS.TIME(OPERATION), modified by a factor
that depends upon the current time and the time the job was
promised to the customer. Note the use of implied subscripts in
communicating between the routines ALLOCATION and EXPEDIT.FACTOR,
and within the routines themselves.

EXPEDIT.FACTOR looks at the DUE.DATE of the current job and
compares it with the current simulation time. If the job is late,
EXPEDIT.FACTOR returns a value of 0.5 to shorten the processing time
of the operation. If the job is not due, a value between 0 and 1 is computed,
which depends upon the time remaining before the job will be late
and the processing time of the current operation. Again
implied subscripts are employed.

The event END.OF.PROCESS does two things, it takes care of a
job that has just finished processing at a production center and it
takes care of the production center. First it deals with the job.
If the ROUTING of the job is empty, all its operations have been
completed and it can pass from the shop. If the job is finished on
time, or is late, its entity record is destroyed. If it is early,
it is filed in FINISHED.GOODS.INVENTORY where it stays until its
DUE.DATE. If the job is not completed, the routine ATTEND.TO.JOB
assigns it to its next operation. Note the use of the REMOVE state-
ment in ALLOCATION that makes this assignment automatic.

It is important at this point to understand the flow of control
between events and routines. The reader is advised to make up some
data, or use the data at the end of the program to trace several
jobs and their flow through the shop.

After disposing of the job, END.OF.PROCESS deals with the produc-
tion center. If no jobs are awaiting processing in the production
center QUEUE, the machine just released is returned to the idle state. If jobs are waiting, one is selected according to the queue's priority rule and processing is started on it.

This completes the routines and events that are directly involved in the shop simulation. The remaining routines and events deal with preparing reports, stopping the simulation, collecting data and monitoring the model.

Events and Routines for Monitoring, Debugging and Analysis

The event WEEKLY.REPORT occurs periodically. It keeps track of the number of simulated weeks that have gone by, resets counters that are used for collecting periodic statistical information, calls on a report routine and reschedules itself. The feature of interest in this routine is the use of the word WEEK in the RESCHEDULE statement. This word was defined to mean *HOURS.V*7 HOURS in the program preamble; the RESCHEDULE statement is therefore compiled as RESCHEDULE THIS WEEKLY.REPORT IN 1*HOURS.V*7 HOURS. When declarations such as this are made, care must be taken that the defined word is not used inadvertently in another context, e.g. there can be no variable, routine, label or set named WEEK in this program.

The event END.OF.SIMULATION is triggered externally and has as its main purpose the termination of a simulation run. It does this by emptying the events sets so control will pass to the statement after START SIMULATION when END.OF.SIMULATION returns control. In addition to stopping the simulation, END.OF.SIMULATION calls REPORT and lists an array.

The next two routines, QUEUE.CHECK and STAY.TIME, are associated with the BEFORE statements of the preamble. QUEUE.CHECK is called before filing or removing is done in any QUEUE; STAY.TIME is called before any JOB is destroyed. The sole purpose of QUEUE.CHECK is error checking; its code is straightforward. STAY.TIME has a different purpose. It computes the time a job spends in the shop, assigning this time to a variable that will have a TALLY operation. Note that STAY is used to compute statistics through
its TALLY computations, but as it is not used anywhere else in the program, it is declared DUMMY and given no storage location.

TRACE is more complicated. Triggered by a call from BETWEEN.V before every event, it is used to trace events and to release jobs from finished goods inventory when the simulation clock reaches their due date. The first part of the routine deals with this task. The code makes use of the fact that jobs are ranked in FINISHED.GOODS.INVENTORY in order of their DUE.DATE.

The trace section of the program uses EVENT.V to branch to a different output statement according to the type of event that has been selected to occur next. These output statements print different items of information about each event type. The program could as easily be written to take actions on the event types, such as turning off the trace by setting BETWEEN.V=0 when TIME.V reaches a certain value or a special kind of event occurs.

Routine NUMBER.IDLE is a left-handed routine that implements the monitoring of the attribute NUMBER.IDLE. It has several unusual features. One reason for defining NUMBER.IDLE as a left-hand monitored variable is to save values of the number of machines idle over time for later processing, without putting the code to do this in the simulation routines. To remove this feature from the program at some later date, one need only remove the preamble card that states that NUMBER.IDLE is monitored and the routine NUMBER.IDLE, and recompile the program. No changes need be made to any other routines.

The program uses two SAVED local arrays to collect and write on tape successive values of NUMBER.IDLE for each production center. A global variable TAPE.FLAG is used to tell the routine when initialization of the SAVED local variables is required; TAPE.FLAG is set to zero at the start of every simulation experiment. The routine demonstrates SAVED values, local arrays, a left-handed function, subscripted subscripts, and monitored variables.

The last routines deal with reports of system activity during a simulation experiment. They print out the parameters of the experiment and the measurements made during the experiment. They illustrate the report generating facilities of SIMSCRIPT II as well as the COMPUTE statement.